



Amortized Rejection Sampling in Universal Probabilistic Programming



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TL;DR

- Rejection sampling loops are widely used in implementing probabilistic models e.g. sampling from constrained distributions.
- Unbounded loops (including rejection sampling) add major complexities to the inference task.
- We address the problem of efficient amortized importance-sampling-based inference, in particular, Inference Compilation (IC) [4] in such models.
- We show naive application of IC can produce estimators with unbounded variance.
- We propose an unbiased estimator for handling rejection sampling loop, prove its finite variance and implement it in pyprob [1; 2] in a way that requires minimal modifications to user code.

| 1: $x \sim p(x)$ 2: 3: for $k \in \mathbb{N}^+$ do 4: $z^k \sim p(z x)$ 5: 6: 7: if $c(x, z^k)$ then 8: $z = z^k$ 9: break 10: observe $(y, p(y z, x))$ | $x \sim q(x y)$ $w \leftarrow \frac{p(x)}{q(x y)}$ $\mathbf{for} \ k \in \mathbb{N}^+ \mathbf{do}$ $\mathbf{z}^k \sim q(\mathbf{z} x,y)$ $w^k \leftarrow \frac{p(\mathbf{z}^k x)}{q(\mathbf{z}^k x,y)}$ $w \leftarrow w \ w^k$ $\mathbf{if} \ c(x,\mathbf{z}^k) \ \mathbf{then}$ $\mathbf{z} = \mathbf{z}^k$ \mathbf{break} $w \leftarrow wp(y \mathbf{z},x)$ |
|-----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|-----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| (a) Original program 1: $x \sim p(x)$ 2: 3: $\boldsymbol{z} \sim p(\boldsymbol{z} x, c(x, z))$ 4: 5: observe $(y, p(y \boldsymbol{z}, x))$ (c) Equivalent to above | (b) Inference compilation $x \sim q(x y)$ $w \leftarrow \frac{p(x)}{q(x y)}$ $z \sim q(z x, y, c(x, z))$ $w \leftarrow w \frac{p(z x, c(x, z))}{q(z x, y, c(x, z))}$ $w \leftarrow w p(y z, x)$ (d) Our IS estimator |

IC weighting

$$w_{IC} = \frac{p(x)}{q(x|y)} p(y|x,z) \prod_{k=1}^{L} w^k$$

Theorem: Under some mild conditions if the following holds then the variance of w_{IC} is infinite.

$$\mathbb{E}_{\boldsymbol{z} \sim q(\boldsymbol{z}|\boldsymbol{x}, \boldsymbol{y})} \left[\frac{p(\boldsymbol{z}|\boldsymbol{x})^2}{q(\boldsymbol{z}|\boldsymbol{x}, \boldsymbol{y})^2} (1 - p(\boldsymbol{A}|\boldsymbol{x}, \boldsymbol{z})) \right] \geq 1$$

where A is the event of c(x, z) being satisfied.

Collapsed weighting

$$w_C = \frac{p(x)}{q(x|y)} \frac{p(\boldsymbol{z}|x, A)}{q(\boldsymbol{z}|x, y, A)} p(y|x, \boldsymbol{z})$$

- $\mathbb{E}\left[w_{IC}\right] = \mathbb{E}\left[w_{C}\right]$ but this weighting scheme does not introduce infinite variance to the estimator.
- Unfortunately, we cannot directly compute w_C

Our method (ARS)

$$w_C = \frac{p(x)}{q(x|y)} \frac{p(\boldsymbol{z}|x)}{q(\boldsymbol{z}|x,y)} p(y|x,\boldsymbol{z}) \frac{q(A|x,y)}{p(A|x)}$$

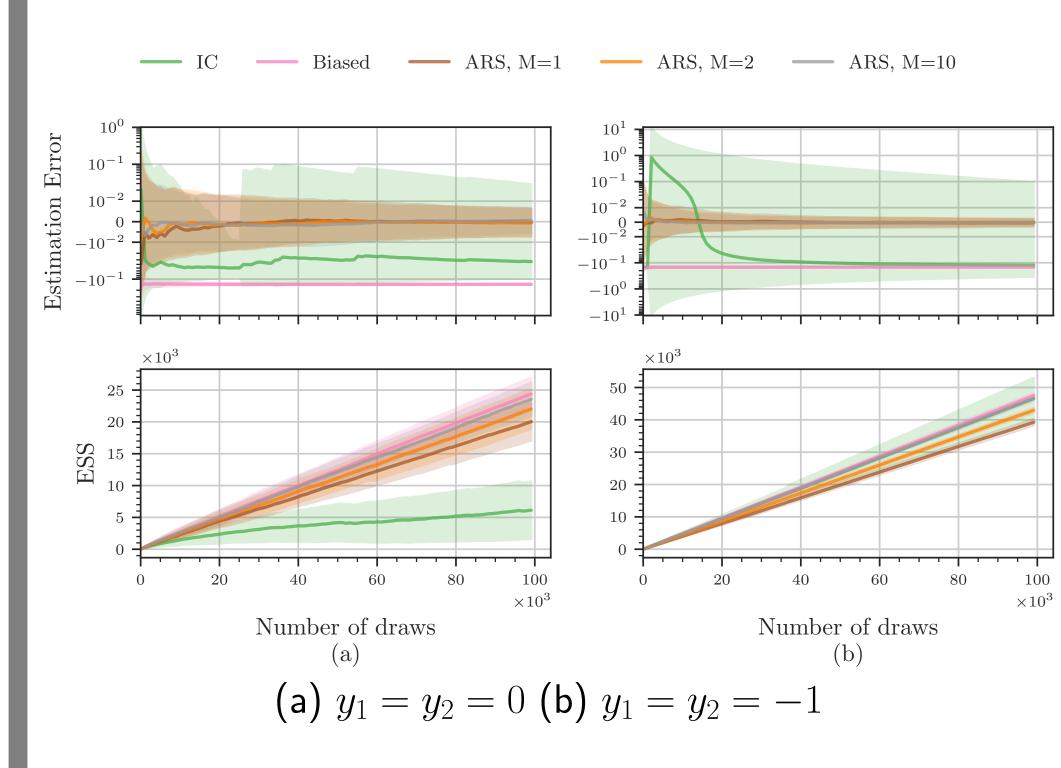
- q(A|x,y) is a Bernoulli distribution. We get an unbiased estimate of its parameter via Monte Carlo.
- The number of samples p(z|x) until the first acceptance follows a geometric distribution with mean $\frac{1}{p(A|x)}$.
- We get an unbiased Monte Carlo estimates of q(A|x,y) and $\frac{1}{p(A|x)}$.

Marsaglia Experiment

Marsaglia polar method [5] is a pseudo-random number sampling method for generating samples from a Normal distribution.

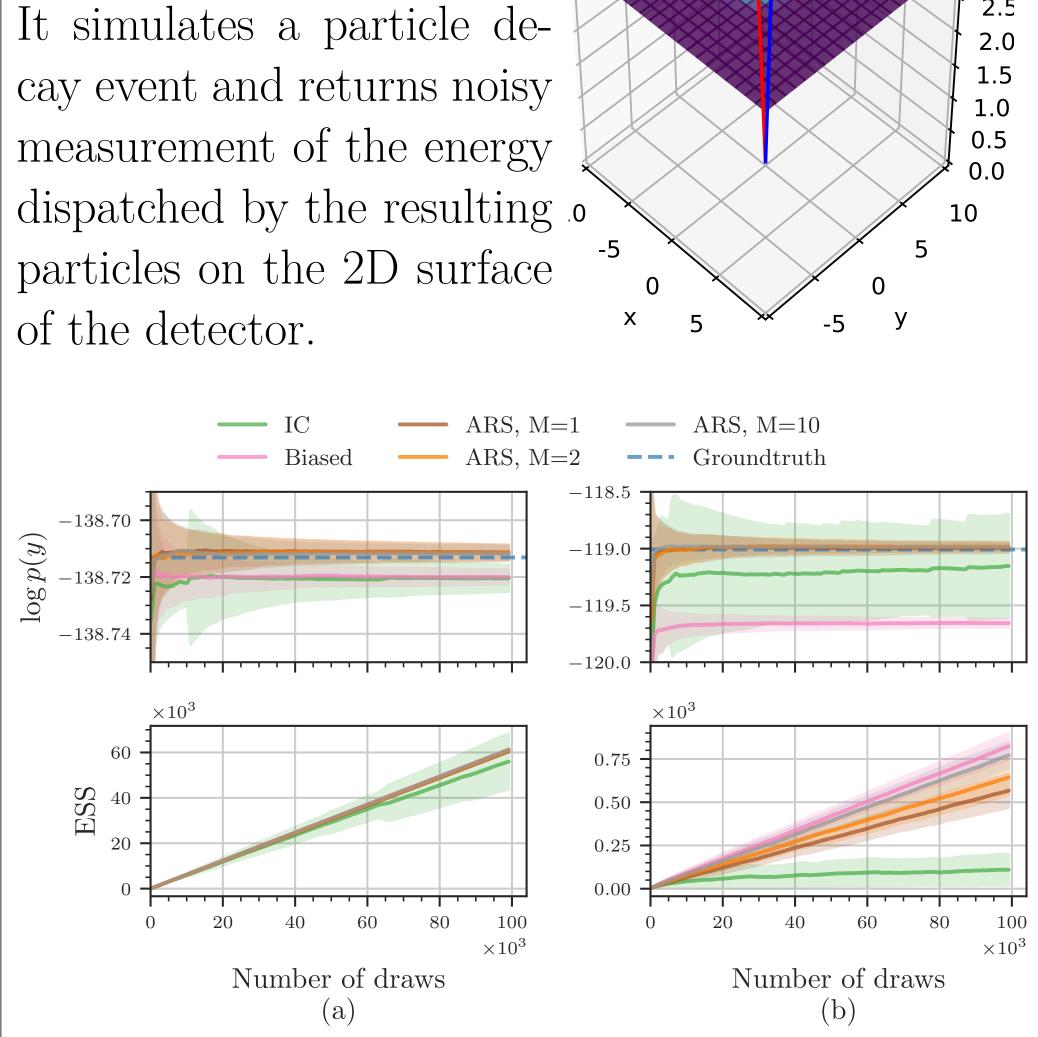
$$\begin{cases} (a,b) \sim \text{Unit circle} \\ \mu = a \sqrt{\frac{-2\log(a^2+b^2)}{a^2+b^2}} \end{cases} \equiv \mu \sim \mathcal{N}(0,1)$$
 $y_i \sim \mathcal{N}(mu,\sigma^2), i \in \{1,2\}$

We target estimating $p(y_1, y_2)$.



Mini-SHERPA Experiment

Mini-SHERPA is a simplified model of high-energy reactions of particles [3]. It simulates a particle demeasurement of the energy of the detector.



Algorithm

| 1: | $I \setminus I \cup J$ |
|----|------------------------------------|
| 2: | $w \leftarrow \frac{p(x)}{q(x y)}$ |
| 3: | for $k \in \mathbb{N}^+$ do |

3:
$$\mathbf{for} \ k \in \mathbb{N}^+ \mathbf{do}$$

4: $\mathbf{z}^k \sim q(\mathbf{z}|x,y)$

5: **if**
$$c(x, \mathbf{z}^k)$$
 then

$$oldsymbol{z}-oldsymbol{z}$$
 break

8:
$$w \leftarrow w \frac{p(\boldsymbol{z}|x)}{q(\boldsymbol{z}|x,y)}$$

9: $K \leftarrow 0$

- 10: **for** $i \in 1, ... N$ **do**
- 11: $\boldsymbol{z}_i' \leftarrow q(\boldsymbol{z}|x,y)$ 12: $K \leftarrow K + c(\boldsymbol{z}, x)$
- 13: **for** $j \in 1, ... M$ **do** 14: for $l \in \mathbb{N}^+$ do $oldsymbol{z}_{i,l}'' \leftarrow q(oldsymbol{z}|x,y)$ if $c(x, \boldsymbol{z}_{i,l}'')$ then $T_j \leftarrow l$ 19: $T \leftarrow \frac{1}{M} \sum_{j=1}^{M} T_j$ 20: $w \leftarrow w \frac{KT}{N}$

21: $w \leftarrow w p(y|\boldsymbol{z}, x)$

Implementation

We introduce two new functions to tag the beginning and end of rejection sampling loops.

| Original | Annotated |
|--------------------------|--------------------------|
| $x = sample(P_x)$ | $x = sample(P_x)$ |
| while True: | while True: |
| | rs_start() |
| $z = sample(P_z(x))$ | $z = sample(P_z(x))$ |
| if $c(x, z)$: | if c(x, z): |
| | rs_end () |
| break | break |
| observe($P_y(x,z)$, y) | observe($P_y(x,z)$, y) |
| return x, z | return x, z |
| | |

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