Notebook

August 6, 2024

```
[]: from typing import List from collections import defaultdict
```

Simple Computation

```
[]: | # https://leetcode.com/problems/maximum-star-sum-of-a-graph/description/
     \# find maximum star sum of a graph with atmost k edges neighbors to it
     # pick onlyy positive values from the graph
     # find all possible stars and then find the maximum sum of the star
     import collections
     class Solution:
         # just find the center of the graph which has (n-1) edges
         def findCenter(self, e: List[List[int]]) -> int:
             return e[0][0] if e[0][0] == e[1][0] or e[0][0] == e[1][1] else e[0][1]
         def maxStarSum(self, vals: List[int], edges: List[List[int]], k: int) -> __
      ⇔int:
             g = collections.defaultdict(set)
             for u, v in edges:
                 if vals[v] > 0 : g[u].add(v)
                 if vals[u] > 0 : g[v].add(u)
             stars = []
             for idx , val in enumerate(vals):
                 vv = [vals[j] for j in g[idx]]
                 vv.sort(reverse=True)
                 stars.append(val + sum(vv[:k]))
             return max(stars)
```

```
[]: # https://leetcode.com/problems/maximal-network-rank/description/
# count the indegree of each pair of nodes and then find the maximum rank
class Solution:
```

```
def maximalNetworkRank(self, n: int, roads: List[List[int]]) -> int:
    g = defaultdict(set)
    for u, v in roads:
        g[u].add(v)
        g[v].add(u)

ans = 0
    for i in range(n):
        for j in range(i+1, n):
            ans = max(ans, (len(g[i]) + len(g[j]) - (i in g[j])))
    return ans
```

\mathbf{DFS}

```
[]: # https://leetcode.com/problems/reconstruct-itinerary/description/
     # start from JFK and reconstruct the itinerary
     class Solution:
         def findItinerary(self, tickets: List[List[str]]) -> List[str]:
             # Construct the graph
             self.flights = defaultdict(list)
             self.itinerary = []
             for ticket in tickets:
                 self.flights[ticket[0]].append(ticket[1])
             # Ensure lexicographical order
             for key in self.flights:
                 self.flights[key].sort(reverse=True)
             self.visit("JFK")
             return self.itinerary[::-1]
         def visit(self, airport):
             while self.flights[airport]:
                 next_airport = self.flights[airport].pop()
                 self.visit(next_airport)
             self.itinerary.append(airport)
```

```
[]: #### https://leetcode.com/problems/word-search/description/

def exist(self, board, word):
    if not board:
        return False
    for i in range(len(board)):
        for j in range(len(board[0])):
            if self.dfs(board, i, j, word):
                return True
    return False
```

```
# check whether can find word, start at (i,j) position
def dfs(self, board, i, j, word):
    if len(word) == 0: # all the characters are checked
        return True
   if i<0 or i>=len(board) or j<0 or j>=len(board[0]) or word[0]!=board[i][j]:
       return False
   tmp = board[i][j] # first character is found, check the remaining part
   board[i][j] = "#" # avoid visit agian
    # check whether can find "word" along one direction
   res = self.dfs(board, i+1, j, word[1:]) or self.dfs(board, i-1, j, word[1:
 →]) \
   or self.dfs(board, i, j+1, word[1:]) or self.dfs(board, i, j-1, word[1:])
   board[i][j] = tmp
   return res
#### https://leetcode.com/problems/word-search-ii/description/
class TrieNode:
   def init (self):
        self.children = defaultdict(TrieNode)
        self.word = None
   def addWord(self, word):
       cur = self
        for c in word:
            cur = cur.children[c]
       cur.word = word
class Solution:
   def findWords(self, board: List[List[str]], words: List[str]) -> List[str]:
       m, n = len(board), len(board[0])
       DIR = [0, 1, 0, -1, 0]
       trieNode = TrieNode()
       ans = []
        for word in words:
           trieNode.addWord(word)
        def dfs(r, c, cur):
            if r < 0 or r == m or c < 0 or c == n or board[r][c] not in cur.
 ⇔children: return
            orgChar = board[r][c]
            cur = cur.children[orgChar]
            board[r][c] = '#' # Mark as visited
            if cur.word != None:
                ans.append(cur.word)
```

```
cur.word = None # Avoid duplication!
for i in range(4): dfs(r + DIR[i], c + DIR[i + 1], cur)
board[r][c] = orgChar # Restore to org state

for r in range(m):
    for c in range(n):
        dfs(r, c, trieNode)
return ans
```

```
[]: # https://leetcode.com/problems/
      -reorder-routes-to-make-all-paths-lead-to-the-city-zero/description/
     # reorder the routes to make all paths lead to city O
     # we make graph with edges in two directions from u to v and v to u and mark_\sqcup
      →the opposite direction as negative since we dont need to change it
     # Time complexity: O(n) where n is the number of cities / nodes
     # Space complexity: O(n)
     class Solution:
         def minReorder(self, n: int, connections: List[List[int]]) -> int:
             def dfs(graph, visited, node):
                 change = 0
                 visited[node] = True
                 for next_node in graph[node]:
                     if not visited[abs(next_node)]:
                         change += dfs(graph, visited, abs(next_node)) + (1 if⊔
      →next_node > 0 else 0)
                 return change
             graph = [[] for _ in range(n)]
             for c in connections:
                 graph[c[0]].append(c[1])
                 graph[c[1]].append(-c[0])
             visited = [False] * n
             return dfs(graph, visited, 0)
```

0.0.1 BFS

```
dist = \{\}
q = deque()
q.append((0, 0))
seen = set()
while q:
    u, d = q.popleft()
    if u in seen: continue
    seen.add(u)
    dist[u] = d
    for v in g[u]:
        q.append((v, d+1))
ans = 0
for index in range(1, len(patience)):
    resendInterval = patience[index]
    shutOffTime = dist[index] * 2
    lastSecond = shutOffTime - 1
    print(resendInterval, shutOffTime)
    lastResentTime = (lastSecond // resendInterval) * resendInterval
    lastPacketTime = lastResentTime + shutOffTime
    ans = max(ans, lastPacketTime)
return ans+1
```

```
[]: # https://github.com/doocs/leetcode/blob/main/solution/0300-0399/0317.

Shortest%20Distance%20from%20All%20Buildings/README_EN.md

### maintain cnt and dist for each cell and then find the minimum distance cell

"""

You are given an m x n grid grid of values 0, 1, or 2, where:

each 0 marks an empty land that you can pass by freely,
each 1 marks a building that you cannot pass through, and
each 2 marks an obstacle that you cannot pass through.
You want to build a house on an empty land that reaches all buildings in the
shortest total travel distance. You can only move up, down, left, and right.

Return the shortest travel distance for such a house. If it is not possible to
build such a house according to the above rules, return -1.

The total travel distance is the sum of the distances between the houses of the
friends and the meeting point.
```

```
The distance is calculated using Manhattan Distance, where distance(p1, p2) = \Box
\Rightarrow /p2.x - p1.x/ + /p2.y - p1.y/
11 11 11
from collections import deque
class Solution:
    def shortestDistance(self, grid: List[List[int]]) -> int:
        m, n = len(grid), len(grid[0])
        q = deque()
        total = 0
        cnt = [[0] * n for _ in range(m)]
        dist = [[0] * n for _ in range(m)]
        for i in range(m):
            for j in range(n):
                if grid[i][j] == 1:
                    total += 1
                    q.append((i, j))
                     visited = [[False] * n for _ in range(m)]
                    visited[i][j] = True
                     level = 0
                     while q:
                         level += 1
                         for _ in range(len(q)):
                             x, y = q.popleft()
                             for a, b in [[0, 1], [1, 0], [0, -1], [-1, 0]]:
                                 nx, ny = x + a, y + b
                                 if 0 \le nx \le m and 0 \le ny \le n and not_{\sqcup}
 →visited[nx][ny] and grid[nx][ny] == 0:
                                     visited[nx][ny] = True
                                     cnt[nx][ny] += 1
                                     dist[nx][ny] += level
                                     q.append((nx, ny))
        res = float('inf')
        for i in range(m):
            for j in range(n):
                if grid[i][j] == 0 and cnt[i][j] == total:
                    res = min(res, dist[i][j])
        return res if res != float('inf') else -1
grid = [[1,0,2,0,1],[0,0,0,0,0],[0,0,1,0,0]]
s = Solution()
print(s.shortestDistance(grid)) # 7
```

```
[]: # https://leetcode.com/problems/01-matrix/description/
     class Solution:
         def updateMatrix(self, mat: List[List[int]]) -> List[List[int]]:
             m, n = len(mat), len(mat[0])
             DIR = [0, 1, 0, -1, 0]
             q = deque([])
             for r in range(m):
                 for c in range(n):
                     if mat[r][c] == 0:
                         q.append((r, c))
                     else:
                         mat[r][c] = -1 # Marked as not processed yet!
             while q:
                 r, c = q.popleft()
                 for i in range(4):
                     nr, nc = r + DIR[i], c + DIR[i + 1]
                     if nr < 0 or nr == m or nc < 0 or nc == n or mat[nr][nc] != -1:
      ⇔continue
                     mat[nr][nc] = mat[r][c] + 1
                     q.append((nr, nc))
             return mat
```

Kahn's Algorithm (Topological Sort)

```
[]:
```

```
[]: # https://leetcode.com/problems/build-a-matrix-with-conditions/description/
     # fill the matrix with the given row and column conditions
     # topological sort - find row and column order and then fill the matrix
     Input: k = 3, rowConditions = [[1,2],[3,2]], colConditions = [[2,1],[3,2]]
     Output: [[3,0,0],[0,0,1],[0,2,0]]
     11 11 11
     class Solution:
         def buildMatrix(self, k: int, rowConditions: List[List[int]], colConditions:

    List[List[int]]) → List[List[int]]:

             def topo_sort(A):
                 deg, graph, order = defaultdict(int), defaultdict(list), []
                 for x, y in A:
                     graph[x].append(y)
                     deg[y] += 1
                 q = deque([i for i in range(1, k+1) if deg[i] == 0])
                 while q:
```

```
x = q.popleft()
        order.append(x)
        for y in graph[x]:
            deg[y] = 1
            if deg[y] == 0: q.append(y)
    return order
left right order = topo sort(colConditions)
above_below_order = topo_sort(rowConditions)
if len(left_right_order) < k or len(above_below_order) < k : return []</pre>
val_to_row = {x: i for i,x in enumerate(above_below_order)}
val_to_col = {x: i for i,x in enumerate(left_right_order)}
ans = [[0] * k for _ in range(k)]
for num in range(1, k+1):
    r, c = val_to_row[num], val_to_col[num]
    ans[r][c] = num
return ans
```

Dijkstra Algorithm (Priority Queue)

[]:

```
[]: # https://leetcode.com/problems/

-find-the-city-with-the-smallest-number-of-neighbors-at-a-threshold-distance/
-description

# given undirected graph with n cities and distance threshold, find the city_u

-with smallest number of neighbors at a threshold distance
```

```
# use dijkstra to find distance from each city to all other cities and then
 →find the number of cities within the threshold distance
# Time complexity: O(n^2 \log n)
class Solution:
    def findTheCity(self, n: int, edges: List[List[int]], distanceThreshold:
 ⇒int) -> int:
        graph = defaultdict(list)
        for u, v, w in edges:
            graph[u].append((v, w))
            graph[v].append((u, w))
        rechable = [0] * n
        def dijkstra(node):
            dist = [float('inf')] * n
            dist[node] = 0
            minHeap = []
            heappush(minHeap, (0, node))
            while minHeap:
                curr_dist, curr_node = heappop(minHeap)
                for next_node, w in graph[curr_node]:
                    if w + curr_dist < dist[next_node]:</pre>
                        dist[next_node] = w + curr_dist
                        heappush(minHeap, (dist[next_node], next_node))
            return dist
        ans node, ans reachable = -1, n
        for i in range(n):
            dist = dijkstra(i)
            dlen = len([x for idx, x in enumerate(dist) if x <=__</pre>
 ⇔distanceThreshold and idx != i])
            if dlen <= ans_reachable:</pre>
                ans_reachable = dlen
                ans_node = i
        return ans_node
```

```
class Solution:
    def minCost(self, grid: List[List[int]]) -> int:
        m, n = len(grid), len(grid[0])
        vis =[[False]*n for _ in range(m)]
        q = deque()
        q.append((0, 0, 0))
        dirs = [[0, 1], [0, -1], [1, 0], [-1, 0]]
        while q:
            x, y, cost = q.popleft()
            if vis[x][y]: continue
            vis[x][y] = True
            if x == m-1 and y == n-1: return cost
            for i in range(len(dirs)):
                dx, dy = dirs[i][0], dirs[i][1]
                nx, ny = x+dx, y+dy
                if nx < 0 or nx >= m or ny < 0 or ny >= n: continue
                if i+1 == grid[x][y]:
                    q.appendleft((nx, ny, cost))
                else:
                    q.append((nx, ny, cost+1))
        return -1
```

```
[]: # https://leetcode.com/problems/path-with-minimum-effort/description/
     # start from 0,0 and reach m-1, n-1 with minimum effort
     # A route's effort is the maximum absolute difference in heights between two !!
      ⇔consecutive cells of the route.
     ,, ,, ,,
     Input: heights = [[1,2,2],[3,8,2],[5,3,5]]
     Output: 2 / Explanation: The route of [1,3,5,3,5] has a maximum absolute \Box
      \hookrightarrow difference of 2 in consecutive cells.
     ,,,,,,
     # standard dijkstra algorithm to solve this problem
     # Time complexity - O(ElogV) = O(M*N log M*N)
     # Space complexity - O(M*N)
     class Solution:
         def minimumEffortPath(self, heights: List[List[int]]) -> int:
             m, n = len(heights), len(heights[0])
             dist = [[inf] * n for _ in range(m)]
             dist[0][0] = 0
             minHeap = [(0, 0, 0)] # (dist, row, col)
             dirs = [[1, 0], [0, 1], [-1, 0], [0, -1]]
             while minHeap:
```

```
d, x, y = heappop(minHeap)
                  if d > dist[x][y]: continue # this is an outdated version \rightarrow skip_{\sqcup}
      \rightarrow it
                 if x == m-1 and y == n-1: return d
                 for dx, dy in dirs:
                     nx, ny = x + dx, y + dy
                      if 0 \le nx \le m and 0 \le ny \le n:
                          newDist = max(d, abs(heights[nx][ny]-heights[x][y]))
                          if dist[nx][ny] > newDist:
                              dist[nx][ny] = newDist
                              heappush(minHeap, (dist[nx][ny], nx, ny))
     # Binary search + DFS solution
[]: # https://leetcode.com/problems/swim-in-rising-water/description/
     # we have to react end of the qrid(m-1, n-1) with minimum time can swim to
      →equal or less than the value of the grid neighbors
     # at time t, we can swim to the cell with value \leftarrow t
     # can be solved used (binary search and dfs) / Priority queue
     # Priority queue solution
     # ans is max(qrid[0][0], qrid[m-1][n-1]) and then do standard priority queue
      ⇔dijkstra algorithm
     class Solution:
         def swimInWater(self, grid: List[List[int]]) -> int:
             m, n = len(grid), len(grid[0])
             pq = []
             vis = [[False]*n for _ in range(m)]
             vis[0][0] = True
             ans = \max(\text{grid}[0][0], \text{grid}[m-1][n-1])
             heappush(pq, (ans, 0, 0))
             dirs = [[1, 0], [0, 1], [-1, 0], [0, -1]]
             while pq:
                 cost, x, y = heappop(pq)
                 ans = max(ans, cost)
                 if x == m-1 and y == n-1: return ans
                 for dx, dy in dirs:
                     nx, ny = dx + x, dy + y
                      if nx == m-1 and ny == n-1: return ans
                      if nx < 0 or nx >= m or ny < 0 or ny >= n or vis[nx][ny]:
                          continue
                      heappush(pq, (grid[nx][ny], nx, ny))
                      vis[nx][ny] = True
```

return -1

```
# Binary search and (dfs or bfs) solution
class Solution:
    def swimInWater(self, grid: List[List[int]]) -> int:
        n = len(grid)
        1, h = grid[0][0], n*n-1
        dirs = [[1, 0], [0, 1], [-1, 0], [0, -1]]
        def bfs(m):
            q = [(0, 0)]
            seen = set((0, 0))
            for x, y in q:
                if grid[x][y] <= m:</pre>
                    if x == y == n-1: return True
                    for dx, dy in dirs:
                        nx, ny = x+dx, y+dy
                         if 0 \le nx \le n and 0 \le ny \le n and (nx, ny) not in seen:
                             seen.add((nx, ny))
                             q.append((nx, ny))
            return False
        def dfs(vis, x, y, m):
            vis[x][y] = True
            for dx, dy in dirs:
                nx, ny = dx+x, dy+y
                if 0 \le nx \le n and 0 \le ny \le n and not vis[nx][ny] and
 ⇒grid[nx][ny] <=m :
                    if nx == n-1 and ny == n-1: return True
                    if dfs(vis, nx, ny, m): return True
            return False
        def isValid(m):
            vis = [[False]*n for _ in range(n)]
            # return dfs(vis, 0, 0, m)
            return bfs(m)
        while 1 < h:
            m = 1 + (h - 1) // 2
            if isValid(m):
                h = m
            else:
                1 = m+1
        return 1
```

[]: # https://leetcode.com/problems/find-the-safest-path-in-a-grid/description/

```
# given a grid with two types of cells, 0 and 1 --> 0 is empty cell and 1 is _{\sqcup}
 ⇔obstacle (thief)
# first check if start (0,0) and end (m-1, n-1) is reachable
# do bfs from all thief cells and calculate the distance to all empty cells
# use dijkstra algorithm to find the minimum distance from start to end cell
class Solution:
    def maximumSafenessFactor(self, grid: List[List[int]]) -> int:
        n = len(grid)
        if grid[0][0] or grid[n-1][n-1]: return 0
        score = [[inf] * n for _ in range(n)]
        dirs = [[-1, 0], [0, -1], [0, 1], [1, 0]]
        def bfs():
            q = deque()
            for i in range(n):
                for j in range(n):
                    if grid[i][j]:
                        score[i][j] = 0
                        q.append((i, j))
            while q:
                x, y = q.popleft()
                s = score[x][y]
                for dx, dy in dirs:
                    nx, ny = x + dx, y + dy
                    if 0 \le nx \le n and 0 \le ny \le n and score[nx][ny] > s + 1:
                        score[nx][ny] = s + 1
                        q.append((nx, ny))
        bfs()
        vis = [[False] * n for _ in range(n)]
        pq = [(-score[0][0], 0, 0)]
        while pq:
            safe, x, y = heappop(pq)
            safe = -safe
            if x == n-1 and y == n-1: return safe
            vis[x][y] = True
            for dx, dy in dirs:
                nx, ny = dx + x, dy + y
                if 0 \le nx \le n and 0 \le ny \le n and not vis[nx][ny]:
                    s = min(score[nx][ny], safe)
                    heappush(pq, (-s, nx, ny))
                    vis[nx][ny] = True
        return -1
```

[]: # https://github.com/doocs/leetcode/blob/main/solution/1100-1199/1102.

Path%20With%20Maximum%20Minimum%20Value/README_EN.md

```
import heapq
     def maximumMinimumPath(grid):
         # Define the directions for moving in 4 cardinal directions
         directions = [(0, 1), (1, 0), (0, -1), (-1, 0)]
         m, n = len(grid), len(grid[0])
         # Max-heap (we invert the values to use heapq as a max-heap)
         \max_{\text{heap}} = [(-\text{grid}[0][0], 0, 0)]
         # To track the maximum minimum value path to each cell
         max_min_values = [[-float('inf')] * n for _ in range(m)]
         max_min_values[0][0] = grid[0][0]
         while max_heap:
             current_min_val, x, y = heapq.heappop(max_heap)
             current_min_val = -current_min_val
             # If we reach the bottom-right corner, return the value
             if x == m - 1 and y == n - 1:
                 return current_min_val
             for dx, dy in directions:
                 nx, ny = x + dx, y + dy
                 if 0 \le nx \le m and 0 \le ny \le n:
                     new_min_val = min(current_min_val, grid[nx][ny])
                     if new_min_val > max_min_values[nx][ny]:
                         max_min_values[nx][ny] = new_min_val
                         heapq.heappush(max_heap, (-new_min_val, nx, ny))
         \# If there is no valid path (although per problem statement, there should \sqcup
      →always be one)
         return -1
     # Example usage:
     grid = [[5, 4, 5], [1, 2, 6], [7, 4, 6]]
     print(maximumMinimumPath(grid)) # Output: 4
     grid = [[3,4,6,3,4],[0,2,1,1,7],[8,8,3,2,7],[3,2,4,9,8],[4,1,2,0,0],[4,6,5,4,3]]
     print(maximumMinimumPath(grid)) # Output: 3
[]: | # https://leetcode.com/problems/network-delay-time/description/
     # will send a signal from node k to all other nodes in the network
```

dijkstra algorithm can be used to solve this problem

```
class Solution:
         def networkDelayTime(self, times: List[List[int]], n: int, k: int) -> int:
             graph = defaultdict(list)
             dist = \{\}
             minHeap = [(0, k)]
             for u, v, w in times:
                 graph[u].append((v, w))
             while minHeap:
                 time, node = heappop(minHeap)
                 if node not in dist:
                     dist[node] = time
                     for v, w in graph[node]:
                         heapq.heappush(minHeap, (time+w, v))
             return max(dist.values()) if len(dist) == n else -1
[]: # https://leetcode.com/problems/
      →minimum-weighted-subgraph-with-the-required-paths/description/
     # we need to find the shortest path from src1 and src2 to dst
     # To solve with we have intermediate node (IM) we try to go from src1 to IM and
      ⇒then from IM to dst
     # standard\ dijkstra\ algorithm\ with\ priority\ queue\ but\ need\ three\ times\ call_{\sqcup}
      →dijkstra(graph, src1), dijkstra(graph, src2), dijkstra(graph, dst)
     from heapq import heappop, heappush
     class Solution:
         def minimumWeight(self, n: int, edges: List[List[int]], src1: int, src2:⊔
      →int, dest: int) -> int:
             def dijkstra(g, src):
                 distances = defaultdict(lambda : inf)
                 heap = [(0, src)]
                 while heap:
                     dist, node = heappop(heap)
                     if node in distances: continue
                     distances[node] = dist
```

newdist = distances[node] + g[node][neigh]

for neigh in g[node].keys():

if neigh in distances: continue

heappush(heap, (newdist, neigh))

```
return distances
      graph = defaultdict(dict)
      rev_graph = defaultdict(dict)
      for u, v, w in edges:
           graph[u][v] = w if v not in graph[u] else min(w, graph[u][v])
           rev_graph[v][u] = w if u not in rev_graph[v] else min(w,__
→rev_graph[v][u])
      src1_distances = dijkstra(graph, src1)
      src2_distances = dijkstra(graph, src2)
      dest_distances = dijkstra(rev_graph, dest)
      res = inf
      for node in range(n):
           local_res = src1_distances[node] + src2_distances[node] +

dest_distances[node]

           res = min(local_res, res)
      return res if res != inf else -1
```

Union Find

```
[]: | # https://github.com/doocs/leetcode/blob/main/solution/1100-1199/1102.
      →Path%20With%20Maximum%20Minimum%20Value/README_EN.md
     # we need to find the maximum minimum value path from (0,0) to (m-1, n-1)
     # we can solve it using union find / dijkstra algorithm
     # we can sort the values and then start from the largest value and then find \Box
     ⇔the path to reach the destination
     # we do union find until parent of 0 and parent of m*n-1 are same
     from typing import List
     class UnionFind:
         __slots__ = ("p", "size")
         def __init__(self, n):
             self.p = list(range(n))
             self.size = [1] * n
         def find(self, x: int) -> int:
             if self.p[x] != x:
                 self.p[x] = self.find(self.p[x])
             return self.p[x]
         def union(self, a: int, b: int) -> bool:
```

```
pa, pb = self.find(a), self.find(b)
        if pa == pb:
            return False
        if self.size[pa] > self.size[pb]:
            self.p[pb] = pa
            self.size[pa] += self.size[pb]
        else:
            self.p[pa] = pb
            self.size[pb] += self.size[pa]
        return True
class Solution:
    def maximumMinimumPath(self, grid: List[List[int]]) -> int:
        m, n = len(grid), len(grid[0])
        uf = UnionFind(m * n)
        q = [(v, i, j) for i, row in enumerate(grid) for j, v in enumerate(row)]
        print(q)
        q.sort()
        ans = 0
        vis = set()
        dirs = [[1, 0], [0, 1], [-1, 0], [0, -1]]
        while uf.find(0) != uf.find(m * n - 1):
            v, i, j = q.pop() # pop the largest value, pop gets the last element
            print(v, i, j)
            ans = v
            vis.add((i, j))
            for a, b in dirs:
                x, y = i + a, j + b
                if (x, y) in vis:
                    uf.union(x * n + y, i * n + j)
        return ans
grid = [[3,4,6,3,4],[0,2,1,1,7],[8,8,3,2,7],[3,2,4,9,8],[4,1,2,0,0],[4,6,5,4,3]]
print(f'ans is {s.maximumMinimumPath(grid)}')
```

Minimum Spanning Tree

- Prim Algorithm Priority Queue
- Kruskal Algorithm Union Find

```
[]: # https://leetcode.com/problems/min-cost-to-connect-all-points/description/ # we need to find the minimum cost to connect all points - minimum spanning tree # given points and manhattan distance between them is the cost (abs(x1-x2) +_{\square} +abs(y1-y2))
```

```
# Prim's algorithm - standard minimum spanning tree algorithm (Priority queue)
# Time complexity - O(n^2 \log n) , due to priority queue operations
# Space complexity - O(n)
class Solution:
    def minCostConnectPoints(self, points: List[List[int]]) -> int:
        def manhattanDistance(p1, p2):
            return abs(p1[0]-p2[0]) + abs(p1[1]-p2[1])
        n = len(points)
        visited = [False] * n
        dist = \{\}
        minHeap = [(0, 0)]
        mst_weight = 0
        while minHeap:
            w, u = heappop(minHeap)
            if visited[u] or dist.get(u, inf) < w:</pre>
                continue
            visited[u] = True
            mst_weight += w
            for v in range(n):
                if not visited[v]:
                    new_dist = manhattanDistance(points[u], points[v])
                    if new_dist < dist.get(v, inf):</pre>
                        dist[v] = inf
                        heappush(minHeap, (new_dist, v))
        return mst_weight
# Kruskal's algorithm - standard minimum spanning tree algorithm (Union find)
# Time complexity - O(n^2 \log n) , due to priority queue operations / edges sort
# Space complexity - O(n)
class UnionFind:
    def __init__(self, n):
        self.parent = list(range(n))
        self.rank = [0] * n
    def find(self, u):
        if self.parent[u] == u: return u
        self.parent[u] = self.find(self.parent[u])
        return self.parent[u]
    def union(self, u, v):
        pu, pv = self.find(u), self.find(v)
        if pu == pv: return False
        if self.rank[pu] > self.rank[pv]:
            pu, pv = pv, pu
```

```
self.parent[pu] = pv
        if self.rank[pu] == self.rank[pv]:
            self.rank[pv] += 1
        return True
class Solution:
    def minCostConnectPoints(self, points: List[List[int]]) -> int:
        def manhattanDistance(p1, p2):
            return abs(p1[0]-p2[0]) + abs(p1[1]-p2[1])
        n = len(points)
        uf = UnionFind(n)
        edges = []
        for i in range(n):
            for j in range(i+1, n):
                distance = manhattanDistance(points[i], points[j])
                heappush(edges, (distance, i, j))
                #edges.append((distance, i, j))
        #edges.sort()
        mst_weight, mst_edges = 0, 0
        while edges:
            # w, u, v = edges.pop(0)
            w, u, v = heappop(edges)
            if uf.union(u, v):
                mst_weight += w
                mst edges += 1
                if mst_edges == n-1: break
        return mst_weight
```

```
[]: # https://leetcode.com/problems/second-minimum-time-to-reach-destination/
      ⇔description/
     # since we need to find the second minimum time to reach the destination, well
      ⇔need to maintain two distances
     # since weight of the edge is time and it is same for all edges we can use BFS_{\sqcup}
      ⇔instead of dijkstra algorithm
     # since red signal happens when time \% change == 0, we can update the new cost
      ⇔based on the red signal
     class Solution:
         def secondMinimum(self, n: int, edges: List[List[int]], time: int, change:
      ⇒int) -> int:
             dist, dist2 = [float('inf')]*(n+1), [float('inf')]*(n+1)
             graph = collections.defaultdict(list)
             for u, v in edges:
                 graph[u].append(v)
                 graph[v].append(u)
```

```
dist[1] = 0
Q = deque()
Q.append((0, 1))
while Q:
    cost, node = Q.popleft()
    for nei in graph[node]:
        new_cost = cost + time
        # update red signal
        if (cost // change) % 2 == 1:
            new_cost += change - (cost % change)
        # update two distance
        if dist[nei] > new_cost:
            dist2[nei], dist[nei] = dist[nei], new_cost
            Q.append((new_cost, nei))
        elif new_cost > dist[nei] and new_cost < dist2[nei]:</pre>
            dist2[nei] = new_cost
            Q.append((new_cost, nei))
return dist2[n]
```

 $[\]: \ \# \ https://leetcode.com/problems/path-with-maximum-probability/solutions/731767/\\ \circ java-python-3-2-codes-bellman-ford-and-dijkstra-s-algorithm-w-brief-explanation-and-analysing and the second s$

Flyod Warshall Algorithm

```
[]: # https://leetcode.com/problems/
      {\small \hookrightarrow} find-the-city-with-the-smallest-number-of-neighbors-at-a-threshold-distance/
     # given undirected graph with n cities and distance threshold, find the city,
      with smallest number of neighbors at a threshold distance
     # flyod warshall algorithm to find the shortest distance between all pairs of \Box
      \hookrightarrow cities
     # Time complexity - O(n^3)
     class Solution:
         def findTheCity(self, n: int, edges: List[List[int]], distanceThreshold:
      ⇒int) -> int:
              dist = [[math.inf] * n for _ in range(n)]
             for i in range(n):
                  dist[i][i] = 0
             for u, v, w in edges:
                  dist[u][v] = dist[v][u] = w
              # Floy-Warshall's algorithm
             for k in range(n):
                  for i in range(n):
```

```
for j in range(n):
    if dist[i][j] > dist[i][k] + dist[k][j]:
        dist[i][j] = dist[i][k] + dist[k][j]

cnt = [0] * n

for i in range(n):
    for j in range(n):
        if i == j: continue
        if dist[i][j] <= distanceThreshold:
            cnt[j] += 1

ans = 0

for i in range(n):
    if cnt[i] <= cnt[ans]:
        ans = i

return ans</pre>
```

```
[]: | # https://leetcode.com/problems/minimum-cost-to-convert-string-i/description/
     # use floyd warshall algorithm to find the shortest path between all pairs of \Box
      \rightarrownodes
     # dijkstra algorithm can also be used to solve this problem but leads to TLE
     from typing import List
     from collections import defaultdict
     import sys
     class Solution:
         def minimumCost(self, source: str, target: str, original: List[str], u
      →changed: List[str], cost: List[int]) -> int:
             # Create a set of all characters in original and changed lists
             all_chars = set(original + changed)
             # Initialize the distance dictionary with infinity
             dist = defaultdict(lambda: defaultdict(lambda: float('inf')))
             # Distance from a node to itself is zero
             for char in all_chars:
                 dist[char][char] = 0
             # Fill initial distances based on the input lists
             for orig, chng, c in zip(original, changed, cost):
                 dist[orig][chng] = min(dist[orig][chng], c)
             # Floyd-Warshall algorithm to compute shortest paths between all pairsu
      ⇔of nodes
             for k in all_chars:
                 for i in all_chars:
```

```
for j in all_chars:
    if dist[i][k] < float('inf') and dist[k][j] < float('inf'):
        dist[i][j] = min(dist[i][j], dist[i][k] + dist[k][j])

total_cost = 0

# Compute the total cost by looking up the precomputed shortest paths
for s, t in zip(source, target):
    if s == t:
        continue
    if dist[s][t] == float('inf'):
        return -1
        total_cost += dist[s][t]

return total_cost

# https://leetcode.com/problems/minimum-cost-to-convert-string-ii/description/</pre>
```