Notebook

August 2, 2024

```
[]: from typing import List from collections import defaultdict
```

Simple Computation

```
[]: | # https://leetcode.com/problems/maximum-star-sum-of-a-graph/description/
     \# find maximum star sum of a graph with atmost k edges neighbors to it
     # pick onlyy positive values from the graph
     # find all possible stars and then find the maximum sum of the star
     import collections
     class Solution:
         # just find the center of the graph which has (n-1) edges
         def findCenter(self, e: List[List[int]]) -> int:
             return e[0][0] if e[0][0] == e[1][0] or e[0][0] == e[1][1] else e[0][1]
         def maxStarSum(self, vals: List[int], edges: List[List[int]], k: int) -> __
      ⇔int:
             g = collections.defaultdict(set)
             for u, v in edges:
                 if vals[v] > 0 : g[u].add(v)
                 if vals[u] > 0 : g[v].add(u)
             stars = []
             for idx , val in enumerate(vals):
                 vv = [vals[j] for j in g[idx]]
                 vv.sort(reverse=True)
                 stars.append(val + sum(vv[:k]))
             return max(stars)
```

\mathbf{DFS}

```
[]: # https://leetcode.com/problems/reconstruct-itinerary/description/
# start from JFK and reconstruct the itinerary
```

```
class Solution:
    def findItinerary(self, tickets: List[List[str]]) -> List[str]:
        # Construct the graph
        self.flights = defaultdict(list)
        self.itinerary = []
        for ticket in tickets:
            self.flights[ticket[0]].append(ticket[1])
        # Ensure lexicographical order
        for key in self.flights:
            self.flights[key].sort(reverse=True)
        self.visit("JFK")
        return self.itinerary[::-1]
    def visit(self, airport):
        while self.flights[airport]:
            next_airport = self.flights[airport].pop()
            self.visit(next_airport)
        self.itinerary.append(airport)
```

```
[]: #### https://leetcode.com/problems/word-search/description/
     def exist(self, board, word):
        if not board:
            return False
        for i in range(len(board)):
            for j in range(len(board[0])):
                 if self.dfs(board, i, j, word):
                    return True
        return False
     # check whether can find word, start at (i,j) position
     def dfs(self, board, i, j, word):
        if len(word) == 0: # all the characters are checked
             return True
        if i<0 or i>=len(board) or j<0 or j>=len(board[0]) or word[0]!=board[i][j]:
            return False
        tmp = board[i][j] # first character is found, check the remaining part
        board[i][j] = "#" # avoid visit agian
        # check whether can find "word" along one direction
        res = self.dfs(board, i+1, j, word[1:]) or self.dfs(board, i-1, j, word[1:
         or self.dfs(board, i, j+1, word[1:]) or self.dfs(board, i, j-1, word[1:])
        board[i][j] = tmp
        return res
```

```
#### https://leetcode.com/problems/word-search-ii/description/
     class TrieNode:
         def __init__(self):
             self.children = defaultdict(TrieNode)
             self.word = None
         def addWord(self, word):
             cur = self
             for c in word:
                 cur = cur.children[c]
             cur.word = word
     class Solution:
         def findWords(self, board: List[List[str]], words: List[str]) -> List[str]:
             m, n = len(board), len(board[0])
             DIR = [0, 1, 0, -1, 0]
             trieNode = TrieNode()
             ans = \Pi
             for word in words:
                 trieNode.addWord(word)
             def dfs(r, c, cur):
                 if r < 0 or r == m or c < 0 or c == n or board[r][c] not in cur.
      ⇔children: return
                 orgChar = board[r][c]
                 cur = cur.children[orgChar]
                 board[r][c] = '#' # Mark as visited
                 if cur.word != None:
                     ans.append(cur.word)
                     cur.word = None # Avoid duplication!
                 for i in range(4): dfs(r + DIR[i], c + DIR[i + 1], cur)
                 board[r][c] = orgChar # Restore to org state
             for r in range(m):
                 for c in range(n):
                     dfs(r, c, trieNode)
             return ans
[]: # https://leetcode.com/problems/
      \rightarrowreorder-routes-to-make-all-paths-lead-to-the-city-zero/description/
     # reorder the routes to make all paths lead to city O
     # we make graph with edges in two directions from u to v and v to u and mark_\sqcup
      sthe opposite direction as negative since we dont need to change it
```

Time complexity: O(n) where n is the number of cities / nodes

```
# Space complexity: O(n)
class Solution:
    def minReorder(self, n: int, connections: List[List[int]]) -> int:
        def dfs(graph, visited, node):
            change = 0
            visited[node] = True
            for next_node in graph[node]:
                if not visited[abs(next_node)]:
                    change += dfs(graph, visited, abs(next_node)) + (1 if_
 →next_node > 0 else 0)
            return change
        graph = [[] for _ in range(n)]
        for c in connections:
            graph[c[0]].append(c[1])
            graph[c[1]].append(-c[0])
        visited = [False] * n
        return dfs(graph, visited, 0)
```

0.0.1 BFS

```
[]: # https://leetcode.com/problems/the-time-when-the-network-becomes-idle/
      ⇔description/
     class Solution:
         def networkBecomesIdle(self, edges: List[List[int]], patience: List[int])_u
      →-> int:
             g = defaultdict(list)
             for u, v in edges:
                 g[u].append(v)
                 g[v].append(u)
             dist = \{\}
             q = deque()
             q.append((0, 0))
             seen = set()
             while q:
                 u, d = q.popleft()
                 if u in seen: continue
                 seen.add(u)
                 dist[u] = d
                 for v in g[u]:
                     q.append((v, d+1))
             ans = 0
             for index in range(1, len(patience)):
```

```
resendInterval = patience[index]
shutOffTime = dist[index] * 2
lastSecond = shutOffTime - 1
print(resendInterval, shutOffTime)
lastResentTime = (lastSecond // resendInterval) * resendInterval
lastPacketTime = lastResentTime + shutOffTime
ans = max(ans, lastPacketTime)
return ans+1
```

```
[]: | # https://qithub.com/doocs/leetcode/blob/main/solution/0300-0399/0317.
      →Shortest%20Distance%20from%20All%20Buildings/README_EN.md
     ### maintain cnt and dist for each cell and then find the minimum distance cell
     You are given an m x n grid grid of values 0, 1, or 2, where:
     each O marks an empty land that you can pass by freely,
     each 1 marks a building that you cannot pass through, and
     each 2 marks an obstacle that you cannot pass through.
     You want to build a house on an empty land that reaches all buildings in the \sqcup
      shortest total travel distance. You can only move up, down, left, and right.
     Return the shortest travel distance for such a house. If it is not possible to_{\sqcup}
      \hookrightarrowbuild such a house according to the above rules, return -1.
     The total travel distance is the sum of the distances between the houses of the \Box
      ⇔friends and the meeting point.
     The distance is calculated using Manhattan Distance, where distance(p1, p2) = 1
      \Rightarrow /p2.x - p1.x/ + /p2.y - p1.y/
     11 11 11
     from collections import deque
     class Solution:
         def shortestDistance(self, grid: List[List[int]]) -> int:
             m, n = len(grid), len(grid[0])
             q = deque()
             total = 0
             cnt = [[0] * n for _ in range(m)]
             dist = [[0] * n for _ in range(m)]
             for i in range(m):
                 for j in range(n):
                      if grid[i][j] == 1:
                          total += 1
```

```
q.append((i, j))
                         visited = [[False] * n for _ in range(m)]
                         visited[i][j] = True
                         level = 0
                         while q:
                             level += 1
                             for _ in range(len(q)):
                                  x, y = q.popleft()
                                  for a, b in [[0, 1], [1, 0], [0, -1], [-1, 0]]:
                                      nx, ny = x + a, y + b
                                      if 0 \le nx \le m and 0 \le ny \le n and not_{\sqcup}
      →visited[nx][ny] and grid[nx][ny] == 0:
                                          visited[nx][ny] = True
                                          cnt[nx][ny] += 1
                                          dist[nx][ny] += level
                                          q.append((nx, ny))
             res = float('inf')
             for i in range(m):
                 for j in range(n):
                     if grid[i][j] == 0 and cnt[i][j] == total:
                         res = min(res, dist[i][j])
             return res if res != float('inf') else -1
     grid = [[1,0,2,0,1],[0,0,0,0,0],[0,0,1,0,0]]
     s = Solution()
     print(s.shortestDistance(grid)) # 7
[]: # https://leetcode.com/problems/01-matrix/description/
     class Solution:
         def updateMatrix(self, mat: List[List[int]]) -> List[List[int]]:
             m, n = len(mat), len(mat[0])
             DIR = [0, 1, 0, -1, 0]
             q = deque([])
             for r in range(m):
                 for c in range(n):
                     if mat[r][c] == 0:
                         q.append((r, c))
                     else:
                         mat[r][c] = -1 # Marked as not processed yet!
             while q:
                 r, c = q.popleft()
                 for i in range(4):
```

```
nr, nc = r + DIR[i], c + DIR[i + 1]
    if nr < 0 or nr == m or nc < 0 or nc == n or mat[nr][nc] != -1:
continue

mat[nr][nc] = mat[r][c] + 1
    q.append((nr, nc))
return mat</pre>
```

Kahn's Algorithm (Topological Sort)

```
[]:
[]: # https://leetcode.com/problems/build-a-matrix-with-conditions/description/
     # fill the matrix with the given row and column conditions
     # topological sort - find row and column order and then fill the matrix
     11 11 11
     Input: k = 3, rowConditions = [[1,2],[3,2]], colConditions = [[2,1],[3,2]]
     Output: [[3,0,0],[0,0,1],[0,2,0]]
     n n n
     class Solution:
         def buildMatrix(self, k: int, rowConditions: List[List[int]], colConditions:

    List[List[int]]) → List[List[int]]:

             def topo sort(A):
                 deg, graph, order = defaultdict(int), defaultdict(list), []
                 for x, y in A:
                     graph[x].append(y)
                     deg[y] += 1
                 q = deque([i for i in range(1, k+1) if deg[i] == 0])
                 while q:
                     x = q.popleft()
                     order.append(x)
                     for y in graph[x]:
                         deg[y] = 1
                         if deg[y] == 0: q.append(y)
                 return order
             left_right_order = topo_sort(colConditions)
             above_below_order = topo_sort(rowConditions)
             if len(left_right_order) < k or len(above_below_order) < k : return []</pre>
             val_to_row = {x: i for i,x in enumerate(above_below_order)}
             val_to_col = {x: i for i,x in enumerate(left_right_order)}
             ans = [[0] * k for _ in range(k)]
             for num in range(1, k+1):
                 r, c = val_to_row[num], val_to_col[num]
```

```
ans[r][c] = num
return ans
```

Dijkstra Algorithm (Priority Queue)

[]:

```
[]: # https://leetcode.com/problems/
      \rightarrow find-the-city-with-the-smallest-number-of-neighbors-at-a-threshold-distance/
     # given undirected graph with n cities and distance threshold, find the city_{\sqcup}
      with smallest number of neighbors at a threshold distance
     # use dijkstra to find distance from each city to all other cities and then
      sfind the number of cities within the threshold distance
     # Time complexity: O(n^2 \log n)
     class Solution:
         def findTheCity(self, n: int, edges: List[List[int]], distanceThreshold:
      →int) -> int:
             graph = defaultdict(list)
             for u, v, w in edges:
                 graph[u].append((v, w))
                 graph[v].append((u, w))
             rechable = [0] * n
             def dijkstra(node):
                 dist = [float('inf')] * n
                 dist[node] = 0
                 minHeap = []
                 heappush(minHeap, (0, node))
```

```
while minHeap:
               curr_dist, curr_node = heappop(minHeap)
               for next_node, w in graph[curr_node]:
                    if w + curr_dist < dist[next_node]:</pre>
                        dist[next_node] = w + curr_dist
                        heappush(minHeap, (dist[next_node], next_node))
           return dist
       ans node, ans reachable = -1, n
       for i in range(n):
           dist = dijkstra(i)
           dlen = len([x for idx, x in enumerate(dist) if x <=__</pre>
⇒distanceThreshold and idx != i])
           if dlen <= ans_reachable:</pre>
               ans_reachable = dlen
               ans_node = i
       return ans_node
```

```
[]: # https://leetcode.com/problems/
      \rightarrow minimum-cost-to-make-at-least-one-valid-path-in-a-grid/description/
     # Can solve using simple queue or priority queue
     # when the grid value matches the direction of the path, then the cost is O_{\sqcup}
      ⇔else 1
     # Time complexity is O(m*n) and space complexity is O(m*n)
     class Solution:
         def minCost(self, grid: List[List[int]]) -> int:
             m, n = len(grid), len(grid[0])
             vis =[[False]*n for _ in range(m)]
             q = deque()
             q.append((0, 0, 0))
             dirs = [[0, 1], [0, -1], [1, 0], [-1, 0]]
             while q:
                 x, y, cost = q.popleft()
                 if vis[x][y]: continue
                 vis[x][y] = True
                 if x == m-1 and y == n-1: return cost
                 for i in range(len(dirs)):
                     dx, dy = dirs[i][0], dirs[i][1]
                     nx, ny = x+dx, y+dy
                     if nx < 0 or nx >= m or ny < 0 or ny >= n: continue
                     if i+1 == grid[x][y]:
                          q.appendleft((nx, ny, cost))
```

```
return -1
[]: # https://leetcode.com/problems/path-with-minimum-effort/description/
     # start from 0,0 and reach m-1,n-1 with minimum effort
     # A route's effort is the maximum absolute difference in heights between two !!
     ⇔consecutive cells of the route.
     Input: heights = [[1,2,2],[3,8,2],[5,3,5]]
     Output: 2 / Explanation: The route of [1,3,5,3,5] has a maximum absolute \Box
      ⇔difference of 2 in consecutive cells.
     # standard dijkstra algorithm to solve this problem
     # Time complexity - O(EloqV) = O(M*N loq M*N)
     # Space complexity - O(M*N)
     class Solution:
         def minimumEffortPath(self, heights: List[List[int]]) -> int:
             m, n = len(heights), len(heights[0])
             dist = [[inf] * n for _ in range(m)]
             dist[0][0] = 0
             minHeap = [(0, 0, 0)] # (dist, row, col)
             dirs = [[1, 0], [0, 1], [-1, 0], [0, -1]]
             while minHeap:
                 d, x, y = heappop(minHeap)
                 if d > dist[x][y]: continue # this is an outdated version -> skip_
      \hookrightarrow it
                 if x == m-1 and y == n-1: return d
                 for dx, dy in dirs:
                     nx, ny = x + dx, y + dy
                     if 0 \le nx \le m and 0 \le ny \le n:
                         newDist = max(d, abs(heights[nx][ny]-heights[x][y]))
                         if dist[nx][ny] > newDist:
                             dist[nx][ny] = newDist
                             heappush(minHeap, (dist[nx][ny], nx, ny))
     # Binary search + DFS solution
[]: # https://leetcode.com/problems/swim-in-rising-water/description/
```

else:

q.append((nx, ny, cost+1))

we have to react end of the grid(m-1, n-1) with minimum time can swim to

⇔equal or less than the value of the grid neighbors

```
# at time t, we can swim to the cell with value \leftarrow t
# can be solved used (binary search and dfs) / Priority queue
# Priority queue solution
# ans is max(grid[0][0], grid[m-1][n-1]) and then do standard priority queue
⇔dijkstra algorithm
class Solution:
    def swimInWater(self, grid: List[List[int]]) -> int:
        m, n = len(grid), len(grid[0])
        pq = []
        vis = [[False]*n for _ in range(m)]
        vis[0][0] = True
        ans = \max(\text{grid}[0][0], \text{grid}[m-1][n-1])
        heappush(pq, (ans, 0, 0))
        dirs = [[1, 0], [0, 1], [-1, 0], [0, -1]]
        while pq:
            cost, x, y = heappop(pq)
            ans = max(ans, cost)
            if x == m-1 and y == n-1: return ans
            for dx, dy in dirs:
                nx, ny = dx + x, dy + y
                if nx == m-1 and ny == n-1: return ans
                if nx < 0 or nx >= m or ny < 0 or ny >= n or vis[nx][ny]:
                    continue
                heappush(pq, (grid[nx][ny], nx, ny))
                vis[nx][ny] = True
        return -1
# Binary search and (dfs or bfs) solution
class Solution:
    def swimInWater(self, grid: List[List[int]]) -> int:
        n = len(grid)
        1, h = grid[0][0], n*n-1
        dirs = [[1, 0], [0, 1], [-1, 0], [0, -1]]
        def bfs(m):
            q = [(0, 0)]
            seen = set((0, 0))
            for x, y in q:
                if grid[x][y] <= m:</pre>
                    if x == y == n-1: return True
                    for dx, dy in dirs:
                        nx, ny = x+dx, y+dy
                        if 0 <= nx < n and 0 <= ny < n and (nx, ny) not in seen:
```

```
seen.add((nx, ny))
                           q.append((nx, ny))
           return False
      def dfs(vis, x, y, m):
           vis[x][y] = True
           for dx, dy in dirs:
               nx, ny = dx+x, dy+y
               if 0 \le nx \le n and 0 \le ny \le n and not vis[nx][ny] and

grid[nx][ny] <=m :</pre>
                   if nx == n-1 and ny == n-1: return True
                   if dfs(vis, nx, ny, m): return True
           return False
      def isValid(m):
           vis = [[False]*n for _ in range(n)]
           # return dfs(vis, 0, 0, m)
           return bfs(m)
      while 1 < h:
          m = 1 + (h - 1) // 2
           if isValid(m):
              h = m
           else:
               1 = m+1
      return 1
```

```
[]: | # https://leetcode.com/problems/find-the-safest-path-in-a-grid/description/
     # given a grid with two types of cells, 0 and 1 --> 0 is empty cell and 1 is _{f L}
      ⇔obstacle (thief)
     # first check if start (0,0) and end (m-1, n-1) is reachable
     # do bfs from all thief cells and calculate the distance to all empty cells
     # use dijkstra algorithm to find the minimum distance from start to end cell
     class Solution:
         def maximumSafenessFactor(self, grid: List[List[int]]) -> int:
             n = len(grid)
             if grid[0][0] or grid[n-1][n-1]: return 0
             score = [[inf] * n for _ in range(n)]
             dirs = [[-1, 0], [0, -1], [0, 1], [1, 0]]
             def bfs():
                 q = deque()
                 for i in range(n):
                     for j in range(n):
                         if grid[i][j]:
                             score[i][j] = 0
                             q.append((i, j))
```

```
while q:
        x, y = q.popleft()
        s = score[x][y]
        for dx, dy in dirs:
            nx, ny = x + dx, y + dy
            if 0 \le nx \le n and 0 \le ny \le n and score[nx][ny] > s + 1:
                score[nx][ny] = s + 1
                q.append((nx, ny))
bfs()
vis = [[False] * n for _ in range(n)]
pq = [(-score[0][0], 0, 0)]
while pq:
    safe, x, y = heappop(pq)
    safe = -safe
    if x == n-1 and y == n-1: return safe
    vis[x][y] = True
    for dx, dy in dirs:
        nx, ny = dx + x, dy + y
        if 0 \le nx \le n and 0 \le ny \le n and not vis[nx][ny]:
            s = min(score[nx][ny], safe)
            heappush(pq, (-s, nx, ny))
            vis[nx][ny] = True
return -1
```

```
# If we reach the bottom-right corner, return the value
             if x == m - 1 and y == n - 1:
                 return current_min_val
             for dx, dy in directions:
                 nx, ny = x + dx, y + dy
                 if 0 \le nx \le m and 0 \le ny \le n:
                     new_min_val = min(current_min_val, grid[nx][ny])
                     if new_min_val > max_min_values[nx][ny]:
                         max_min_values[nx][ny] = new_min_val
                         heapq.heappush(max_heap, (-new_min_val, nx, ny))
         \# If there is no valid path (although per problem statement, there should \sqcup
      ⇔always be one)
         return -1
     # Example usage:
     grid = [[5, 4, 5], [1, 2, 6], [7, 4, 6]]
     print(maximumMinimumPath(grid)) # Output: 4
     grid = [[3,4,6,3,4],[0,2,1,1,7],[8,8,3,2,7],[3,2,4,9,8],[4,1,2,0,0],[4,6,5,4,3]]
     print(maximumMinimumPath(grid)) # Output: 3
[]: | # https://leetcode.com/problems/network-delay-time/description/
     # will send a signal from node k to all other nodes in the network
     # dijkstra algorithm can be used to solve this problem
     class Solution:
         def networkDelayTime(self, times: List[List[int]], n: int, k: int) -> int:
             graph = defaultdict(list)
             dist = \{\}
             minHeap = [(0, k)]
             for u, v, w in times:
                 graph[u].append((v, w))
             while minHeap:
                 time, node = heappop(minHeap)
                 if node not in dist:
                     dist[node] = time
                     for v, w in graph[node]:
                         heapq.heappush(minHeap, (time+w, v))
             return max(dist.values()) if len(dist) == n else -1
```

```
[]: # https://leetcode.com/problems/
      →minimum-weighted-subgraph-with-the-required-paths/description/
     # we need to find the shortest path from src1 and src2 to dst
     # To solve with we have intermediate node (IM) we try to go from src1 to IM and
      →then from IM to dst
     # standard dijkstra algorithm with priority queue but need three times call_{\mathsf{L}}
      -dijkstra(graph, src1), dijkstra(graph, src2), dijkstra(graph, dst)
     from heapq import heappop, heappush
     class Solution:
         def minimumWeight(self, n: int, edges: List[List[int]], src1: int, src2:
      →int, dest: int) -> int:
             def dijkstra(g, src):
                 distances = defaultdict(lambda : inf)
                 heap = [(0, src)]
                 while heap:
                     dist, node = heappop(heap)
                     if node in distances: continue
                     distances[node] = dist
                     for neigh in g[node].keys():
                         if neigh in distances: continue
                         newdist = distances[node] + g[node][neigh]
                         heappush(heap, (newdist, neigh))
                 return distances
             graph = defaultdict(dict)
             rev_graph = defaultdict(dict)
             for u, v, w in edges:
                 graph[u][v] = w if v not in graph[u] else min(w, graph[u][v])
                 rev_graph[v][u] = w if u not in rev_graph[v] else min(w, u
      →rev_graph[v][u])
             src1_distances = dijkstra(graph, src1)
             src2_distances = dijkstra(graph, src2)
             dest_distances = dijkstra(rev_graph, dest)
             res = inf
             for node in range(n):
                 local_res = src1_distances[node] + src2_distances[node] +__
      →dest_distances[node]
                 res = min(local_res, res)
```

```
return res if res != inf else -1
```

Union Find

```
[]: | # https://qithub.com/doocs/leetcode/blob/main/solution/1100-1199/1102.
      →Path%20With%20Maximum%20Minimum%20Value/README_EN.md
     # we need to find the maximum minimum value path from (0,0) to (m-1, n-1)
     # we can solve it using union find / dijkstra algorithm
     # we can sort the values and then start from the largest value and then find_{\mathsf{L}}
      ⇔the path to reach the destination
     # we do union find until parent of 0 and parent of m*n-1 are same
     from typing import List
     class UnionFind:
         __slots__ = ("p", "size")
         def __init__(self, n):
             self.p = list(range(n))
             self.size = [1] * n
         def find(self, x: int) -> int:
             if self.p[x] != x:
                 self.p[x] = self.find(self.p[x])
             return self.p[x]
         def union(self, a: int, b: int) -> bool:
             pa, pb = self.find(a), self.find(b)
             if pa == pb:
                 return False
             if self.size[pa] > self.size[pb]:
                 self.p[pb] = pa
                 self.size[pa] += self.size[pb]
             else:
                 self.p[pa] = pb
                 self.size[pb] += self.size[pa]
             return True
     class Solution:
         def maximumMinimumPath(self, grid: List[List[int]]) -> int:
             m, n = len(grid), len(grid[0])
             uf = UnionFind(m * n)
             q = [(v, i, j) for i, row in enumerate(grid) for j, v in enumerate(row)]
             print(q)
```

```
q.sort()
        ans = 0
        vis = set()
        dirs = [[1, 0], [0, 1], [-1, 0], [0, -1]]
        while uf.find(0) != uf.find(m * n - 1):
            v, i, j = q.pop() # pop the largest value, pop gets the last element
            print(v, i, j)
            ans = v
            vis.add((i, j))
            for a, b in dirs:
                x, y = i + a, j + b
                if (x, y) in vis:
                    uf.union(x * n + y, i * n + j)
        return ans
grid = [[3,4,6,3,4],[0,2,1,1,7],[8,8,3,2,7],[3,2,4,9,8],[4,1,2,0,0],[4,6,5,4,3]]
s = Solution()
print(f'ans is {s.maximumMinimumPath(grid)}')
```

Minimum Spanning Tree

- Prim Algorithm Priority Queue
- Kruskal Algorithm Union Find

```
[]: | # https://leetcode.com/problems/min-cost-to-connect-all-points/description/
     # we need to find the minimum cost to connect all points - minimum spanning tree
     # given points and manhattan distance between them is the cost (abs(x1-x2) +
      \hookrightarrow abs(y1-y2))
     # Prim's algorithm - standard minimum spanning tree algorithm (Priority queue)
     # Time complexity - O(n^2 \log n) , due to priority queue operations
     # Space complexity - O(n)
     class Solution:
         def minCostConnectPoints(self, points: List[List[int]]) -> int:
             def manhattanDistance(p1, p2):
                 return abs(p1[0]-p2[0]) + abs(p1[1]-p2[1])
             n = len(points)
             visited = [False] * n
             dist = \{\}
             minHeap = [(0, 0)]
             mst_weight = 0
             while minHeap:
                 w, u = heappop(minHeap)
                 if visited[u] or dist.get(u, inf) < w:</pre>
                      continue
                 visited[u] = True
```

```
mst_weight += w
            for v in range(n):
                if not visited[v]:
                    new_dist = manhattanDistance(points[u], points[v])
                    if new_dist < dist.get(v, inf):</pre>
                        dist[v] = inf
                        heappush(minHeap, (new_dist, v))
        return mst_weight
# Kruskal's algorithm - standard minimum spanning tree algorithm (Union find)
# Time complexity - O(n^2 \log n) , due to priority queue operations / edges sort
# Space complexity - O(n)
class UnionFind:
    def __init__(self, n):
        self.parent = list(range(n))
        self.rank = [0] * n
    def find(self, u):
        if self.parent[u] == u: return u
        self.parent[u] = self.find(self.parent[u])
        return self.parent[u]
    def union(self, u, v):
        pu, pv = self.find(u), self.find(v)
        if pu == pv: return False
        if self.rank[pu] > self.rank[pv]:
            pu, pv = pv, pu
        self.parent[pu] = pv
        if self.rank[pu] == self.rank[pv]:
            self.rank[pv] += 1
        return True
class Solution:
    def minCostConnectPoints(self, points: List[List[int]]) -> int:
        def manhattanDistance(p1, p2):
            return abs(p1[0]-p2[0]) + abs(p1[1]-p2[1])
        n = len(points)
        uf = UnionFind(n)
        edges = []
        for i in range(n):
            for j in range(i+1, n):
                distance = manhattanDistance(points[i], points[j])
                heappush(edges, (distance, i, j))
                #edges.append((distance, i, j))
```

```
#edges.sort()
mst_weight, mst_edges = 0, 0
while edges:
    # w, u, v = edges.pop(0)
    w, u, v = heappop(edges)
    if uf.union(u, v):
        mst_weight += w
        mst_edges += 1
        if mst_edges == n-1: break
return mst_weight
```

```
[]: # https://leetcode.com/problems/second-minimum-time-to-reach-destination/
      ⇔description/
     # since we need to find the second minimum time to reach the destination, we \Box
      →need to maintain two distances
     # since weight of the edge is time and it is same for all edges we can use BFS_{\sqcup}
     ⇔instead of dijkstra algorithm
     # since red signal happens when time \% change == 0, we can update the new cost
      ⇔based on the red signal
     class Solution:
         def secondMinimum(self, n: int, edges: List[List[int]], time: int, change:
      ⇒int) -> int:
             dist, dist2 = [float('inf')]*(n+1), [float('inf')]*(n+1)
             graph = collections.defaultdict(list)
             for u, v in edges:
                 graph[u].append(v)
                 graph[v].append(u)
             dist[1] = 0
             Q = deque()
             Q.append((0, 1))
             while Q:
                 cost, node = Q.popleft()
                 for nei in graph[node]:
                     new_cost = cost + time
                     # update red signal
                     if (cost // change) % 2 == 1:
                         new_cost += change - (cost % change)
                     # update two distance
                     if dist[nei] > new_cost:
                         dist2[nei], dist[nei] = dist[nei], new_cost
                         Q.append((new_cost, nei))
                     elif new_cost > dist[nei] and new_cost < dist2[nei]:</pre>
                         dist2[nei] = new_cost
                         Q.append((new_cost, nei))
             return dist2[n]
```

 $[\]: \ \# \ https://leetcode.com/problems/path-with-maximum-probability/solutions/731767/\\ \hookrightarrow java-python-3-2-codes-bellman-ford-and-dijkstra-s-algorithm-w-brief-explanation-and-analysing states of the problems of the pro$

Flyod Warshall Algorithm

```
[]: # https://leetcode.com/problems/
      \rightarrow find-the-city-with-the-smallest-number-of-neighbors-at-a-threshold-distance/
     # given undirected graph with n cities and distance threshold, find the city_{\sqcup}
      with smallest number of neighbors at a threshold distance
     # flyod warshall algorithm to find the shortest distance between all pairs of \Box
      \hookrightarrow cities
     # Time complexity - O(n^3)
     class Solution:
         def findTheCity(self, n: int, edges: List[List[int]], distanceThreshold:
      →int) -> int:
             dist = [[math.inf] * n for _ in range(n)]
             for i in range(n):
                  dist[i][i] = 0
             for u, v, w in edges:
                  dist[u][v] = dist[v][u] = w
              # Floy-Warshall's algorithm
             for k in range(n):
                  for i in range(n):
                      for j in range(n):
                          if dist[i][j] > dist[i][k] + dist[k][j]:
                              dist[i][j] = dist[i][k] + dist[k][j]
             cnt = [0] * n
             for i in range(n):
                  for j in range(n):
                      if i == j: continue
                      if dist[i][j] <= distanceThreshold:</pre>
                          cnt[j] += 1
             ans = 0
             for i in range(n):
                  if cnt[i] <= cnt[ans]:</pre>
                      ans = i
             return ans
```

[]: | # https://leetcode.com/problems/minimum-cost-to-convert-string-i/description/

```
# use floyd warshall algorithm to find the shortest path between all pairs of \Box
 \hookrightarrownodes
# dijkstra algorithm can also be used to solve this problem but leads to TLE
from typing import List
from collections import defaultdict
import sys
class Solution:
    def minimumCost(self, source: str, target: str, original: List[str], u
 ⇔changed: List[str], cost: List[int]) -> int:
        # Create a set of all characters in original and changed lists
        all_chars = set(original + changed)
        # Initialize the distance dictionary with infinity
        dist = defaultdict(lambda: defaultdict(lambda: float('inf')))
        # Distance from a node to itself is zero
        for char in all_chars:
            dist[char][char] = 0
        # Fill initial distances based on the input lists
        for orig, chng, c in zip(original, changed, cost):
            dist[orig][chng] = min(dist[orig][chng], c)
        # Floyd-Warshall algorithm to compute shortest paths between all pairs
 ⇔of nodes
        for k in all_chars:
            for i in all_chars:
                for j in all_chars:
                    if dist[i][k] < float('inf') and dist[k][j] < float('inf'):</pre>
                        dist[i][j] = min(dist[i][j], dist[i][k] + dist[k][j])
        total_cost = 0
        # Compute the total cost by looking up the precomputed shortest paths
        for s, t in zip(source, target):
            if s == t:
                continue
            if dist[s][t] == float('inf'):
                return -1
            total cost += dist[s][t]
        return total_cost
# https://leetcode.com/problems/minimum-cost-to-convert-string-ii/description/
```