

# Week 4 Lecture 10

Theory

# What's in this lecture?

- Reasoning about the Efficiency of Algorithms

# Sorting

- Previous lectures: explored bubble sort, insertion sort, and merge sort
- Today: quantifying running time of algorithms
- For any algorithm, we seek to describe a **recurrence relation** that talks about its running time in terms of component operations

# Recap: Merge Sort

- Input: array of size  $l$ 
  - Return the array itself (already sorted)
- Input: an array of size  $N$ 
  - Split into 2 sub-arrays of size  $N/2$
  - Recursively sort sub-arrays
  - Merge two sorted sub-arrays

# Talking about Time

- $T(I)$  means: “operations necessary for merge\_sort of array of size  $I$ ”
  - Thus,  $T(I) = I$
- $T(N)$  means: “operations necessary for merge\_sort of array of size  $N$ ”
  - $T(N) = T(\text{split}(N)) + T(N/2) + T(N/2) + T(\text{merge}(N/2, N/2))$

# Digging In

- $T(N) = T(\text{split}(N)) + T(N/2) + T(N/2) + T(\text{merge}(N/2, N/2))$
- $T(\text{split}(N)) = N * \text{copy from original array}$
- $T(N/2) = \text{time of recursive merge\_sort on array of size } N/2$
- $T(\text{merge}(N/2, N/2)) = N * \text{copy from original to destination array}$

# Simplifying Terms

- $T(N) = T(\text{split}(N)) + T(N/2) + T(N/2) + T(\text{merge}(N/2, N/2))$
- $T(N) = N + T(N/2) + T(N/2) + N$
- $T(N) = 2N + 2T(N/2)$

# Question

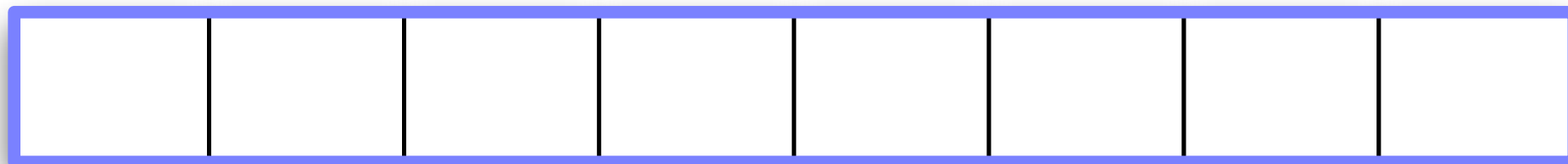
- How many times can something be split in half before it is less than or equal to 1?



# Answer

- The number of times is less than  $\text{ceil}(\lg(N))$
- Thus, merge sort does at most  $\text{ceil}(\lg(N))$  recursive calls

# Recursive Calls



$T(8)$



$2T(4)$



$4T(2)$



$8T(1)$

# Working it Out

- $$\begin{aligned} T(8) &= 8 + 2T(4) + 8 \\ &= 16 + 2(4 + 2T(2) + 4) \\ &= 32 + 4T(2) \\ &= 32 + 4(2 + 2T(1) + 2) \\ &= 48 + 8T(1) \\ &= 48 + 8 \\ &= 56 \end{aligned}$$

# Formalizing

- For merge sort, we have:  
$$T(N) = 2T(N/2) + 2N$$
- Using the master theorem (not covered here), the running time of merge sort as  $N$  approaches infinity is proportional to the function:  $N * \lg(N)$
- Intuitively, there are  $\lg N$  levels in the recursion tree, and the total work in each level is  $N$

# Take-Aways

- It is possible to analyze and compare algorithms based on recurrence relations
- Calculating concrete recurrences is beyond the scope of this class
- For many algorithms and data structures, the running times of operations are well known

# Exercises

- Read wikipedia entries on Merge Sort, Insertion Sort, and Bubble Sort
- Re-read CLR sections which talk about running-time analysis of Merge Sort and Insertion Sort in a best-attempt to understand them