



# DEADLOCK DETECTION

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## Project Overview

### Objective:

Implement a program to detect and resolve deadlocks in a Resource Allocation Graph (RAG).

### Key Features:

- Deadlock detection using cycle detection.
- Deadlock resolution via process termination.
- Visual representation of RAG.
- Reset and priority simulation features.

# Resource Allocation Graph (RAG)

What is RAG?

A directed graph where:

Nodes represent processes (P) and resources (R).

Edges indicate resource allocation or requests.

Example Components:

Processes: P0, P1, P2

Resources: R0, R1, R2

Key Concept:

Allocation Edge: Resource assigned to a process (e.g.,  $R0 \rightarrow P0$ ).

Request Edge: Process requesting a resource (e.g.,  $P0 \rightarrow R1$ ).

## Key Code Components

### 1.Graph Functions:

- initGraph: Initialize the RAG with a given number of nodes.
- addEdge: Add edges between nodes (process-resource or resource-process).
- removeEdge: Remove edges when processes release resources.
- displayRAG: Visualizes the RAG structure.

### 2.Deadlock Detection Functions:

- detectCycle: Uses Depth-First Search (DFS) to find cycles.
- detectCycleUtil: Helper function for cycle detection.

### 3.Deadlock Resolution Functions:

- resolveDeadlock: Identifies processes in the cycle and terminates one to resolve deadlock.

## RAG Creation Example

```
initGraph(&graph, totalNodes);  
addEdge(&graph, 0, 3); // P0 → R3  
addEdge(&graph, 1, 4); // P1 → R4  
addEdge(&graph, 2, 5); // P2 → R5  
addEdge(&graph, 3, 1); // R3 → P1  
addEdge(&graph, 4, 2); // R4 → P2  
addEdge(&graph, 5, 0); // R5 → P0
```

## Resulting Graph Visualization:

```
graph TD  
P0 <--> R3  
P1 <--> R4  
P2 <--> R5  
P0 -.-> R0  
P1 -.-> R3  
P2 -.-> R4
```

P0 ---> R3 <--- P1  
|  
|  
V  
R0 <--- P2 ---> R5

# Semaphore:

Wait and signal  
functions

```
/**
 * Function: wait
 * -----
 * Simulates the request of a resource by a process.
 * Adds an edge from the process to the resource in the resource allocation graph,
 * indicating that the process is waiting for the resource.
 */
void wait(int process, int resource) {
    // Add an edge from the process to the resource in the resource graph.
    // This represents that the process is now waiting for the resource.
    addEdge(&resourceGraph, process, resource);
    printf("Process P%d is requesting Resource R%d.\n", process, resource);

    // Check if adding this edge introduces a deadlock in the system.
    if (isDeadlock(&resourceGraph)) {
        // Notify that a deadlock has been detected.
        // This warns the user/system about potential problems due to this request.
        printf("Warning: Deadlock detected. Process P%d's request may cause issues.\n", process);
    }
}

/**
 * Function: signal
 * -----
 * Simulates the release of a resource by a process.
 * Removes the edge from the process to the resource in the resource allocation graph,
 * indicating that the resource is no longer held or needed by the process.
 */
void signal(int process, int resource) {
    // Remove the edge from the process to the resource in the resource graph.
    // This signifies that the process has released the resource.
}
```

## Deadlock Detection :

### Code :

```
// Function to check if there is a deadlock in the Resource Allocation Graph (RAG).
int isDeadlock(Graph *g) { // Iterate through all nodes in the graph to detect cycles.
    for (int i = 0; i < g->numNodes; i++) { // If a cycle is detected starting from
        node i, a deadlock exists.
            if (detectCycle(g, i))
                return 1; }
    return 0; // No deadlock found.}
```

### Logic :

1. Uses detectcycle to find cycles in the graph
2. If cycle exists deadlock is confirmed.

# Deadlock Resolution

- **Process:** Select a process involved in the deadlock (e.g., P0).
- Remove edges associated with the terminated process.
- Free up resources to resolve the deadlock

```
function to resolve deadlock by terminating a process and releasing its resources.
resolveDeadlock(Graph *g) {
    printf("Attempting to resolve deadlock...\n");
    for (int i = 0; i < g->numNodes; i++) {
        // Check if a cycle (deadlock) exists starting from node `i`.
        if (detectCycle(g, i)) {
            printf("Process P%d is involved in deadlock.\n", i);
            printf("Terminating Process P%d to resolve deadlock.\n", i);

            // Release all resources held by Process P[i].
            for (int j = 0; j < g->numNodes; j++) {
                removeEdge(g, i, j);
            }

            printf("Deadlock resolved.\n");
            return; // Exit after resolving the deadlock.
        }
    }

    printf("No deadlock to resolve.\n");
}
```



# Priority Simulation

```
setPriority(0, 1); // Set priority of P0 to 1  
setPriority(1, 2); // Set priority of P1 to 2  
setPriority(2, 3); // Set priority of P2 to 3
```

**Purpose:** Allows processes to have priorities to influence scheduling or deadlock resolution strategies.

# OUTPUT

```
[vamsigudipudi@Vamsis-MacBook-Air OS % gcc -o deadlock_detection main.c graph.c semaphore.c deadlock_detection.  
[  
vamsigudipudi@Vamsis-MacBook-Air OS % ./deadlock_detection  
  
Setting priority of Process P0 to 1  
Setting priority of Process P1 to 2  
Setting priority of Process P2 to 3  
Simulating resource allocation scenario...  
Resource Allocation Graph created.  
Resource Allocation Graph:  
Node 0 -> Node 3  
Node 1 -> Node 4  
Node 2 -> Node 5  
Node 3 -> Node 1  
Node 4 -> Node 2  
Node 5 -> Node 0  
Deadlock detected after allocations.  
Attempting to resolve deadlock...  
Process P0 is involved in deadlock.  
Terminating Process P0 to resolve deadlock.  
Deadlock resolved.  
Resetting Resource Allocation Graph...  
Graph reset successfully.  
vamsigudipudi@Vamsis-MacBook-Air OS % $
```

Here a RAG is drawn, a deadlock is detected , and the deadlock is resolved

# Before and After Resolution

- Before Deadlock Resolution:

- $P0 \rightarrow R3 \leftarrow P1$

V	V

$R0 \leftarrow P2 \rightarrow R5$

- After Deadlock Resolution:

- $P1 \rightarrow R4 \leftarrow P2$

V	V

$R1 \leftarrow P2 \rightarrow R5$

# Conclusion :

## Key Takeaways:

- Deadlock detection is crucial for resource allocation in concurrent systems.
  - The RAG model simplifies detection and resolution.
  - Simulations allow testing under various scenarios.



Thank You