

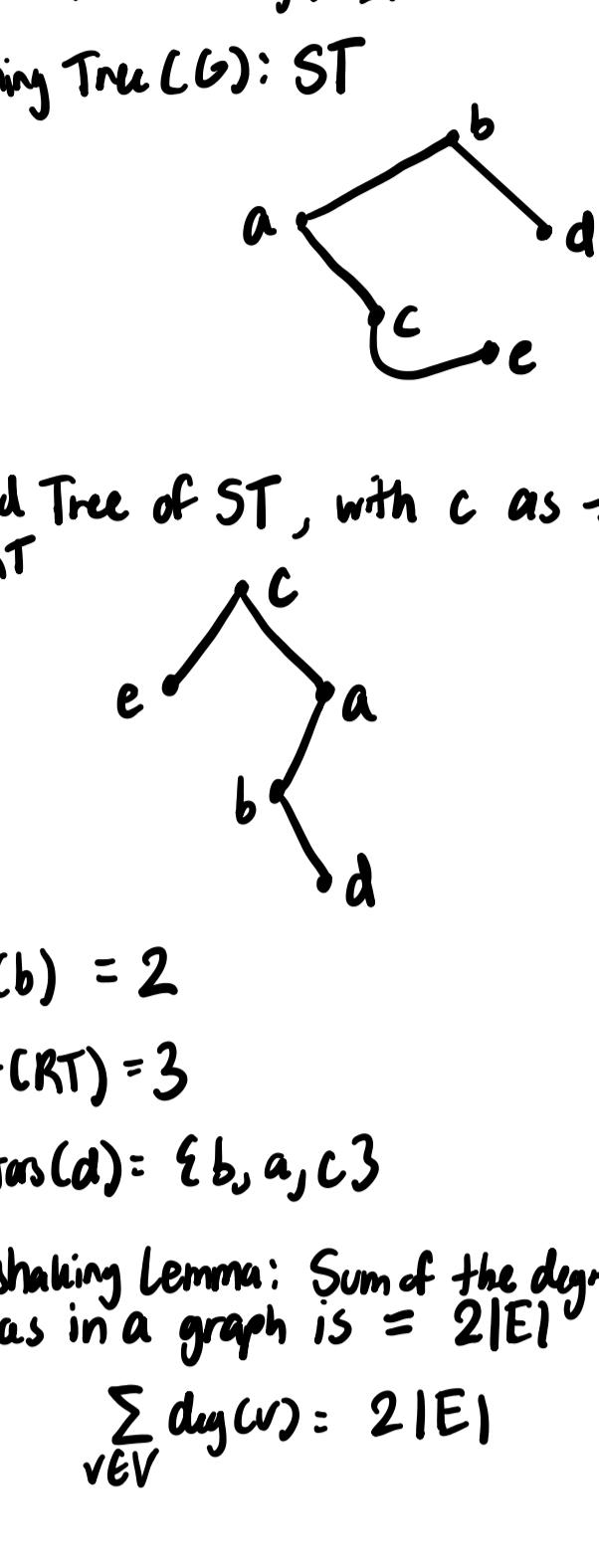
### Announcements

- Midterm #2 on 11/21 (Thursday) in Phil B152
  - Graphs, Trees, Spanning Trees, Hamiltonian Cycles, Eulerian Circuits, Directed Graphs, Graph Coloring
- My office hours are after class (midterm review)
- TA application for Discrete Math I (to be released soon...)

### Midterm Review (Content)

#### Undirected Graphs:

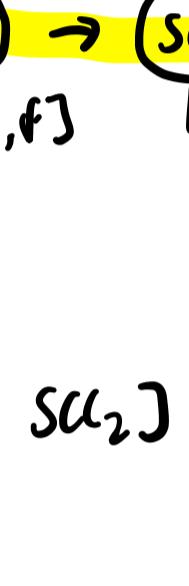
$G(V, E)$



$$V = \{a, b, c, d, e\}$$

$$E = [\{a, c\}, \{c, e\}, \{b, d\}, \{c, d\}, \{a, b\}]$$

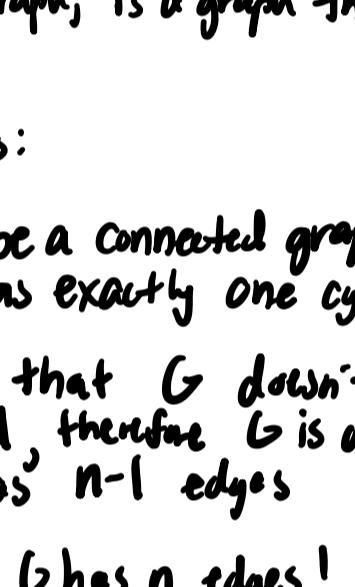
- $N(c) = \{a, d, e\}$
- $\delta(G) = 1$
- $\Delta(G) = 3$
- Walk( $G$ ) (example):  $a \rightarrow c \rightarrow d \rightarrow b \rightarrow d \rightarrow c \rightarrow e$
- Path( $G$ ) (example):  $a \rightarrow c \rightarrow d \rightarrow b$
- Cycle( $G$ ):  $a \rightarrow b \rightarrow d \rightarrow c \rightarrow a$
- Subgraph( $G$ ) (example):



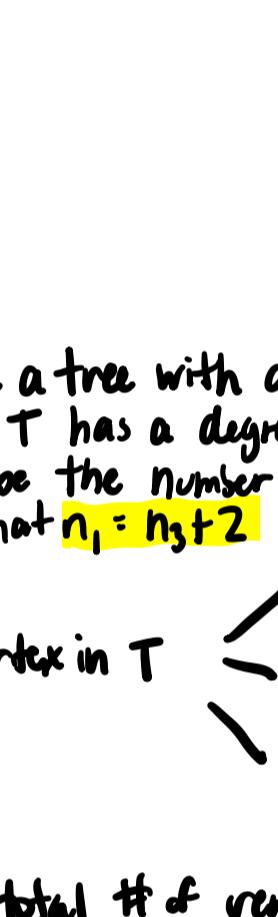
- |Connected Components| = 1, aka the graph is connected

•  $G$  is a tree. Why? It has a cycle

- Spanning Tree( $G$ ): ST



- Rooted Tree of ST, with  $c$  as the root



$$\text{level}(b) = 2$$

$$\text{height}(RT) = 3$$

$$\text{ancestors}(d) = \{b, a, c\}$$

- Handshaking Lemma: Sum of the degrees of all the vertices in a graph is  $= 2|E|$

$$\sum_{v \in V} \deg(v) = 2|E|$$

- Every graph with  $n$  vertices and  $m$  edges has at least  $n-m$  connected components

- Properties of a tree:
  - acyclic and connected
  - connected and  $n-1$  edges
  - minimally connected
  - acyclic, but adding a new edge creates a cycle
  - Every pair of vertices are linked by a unique path

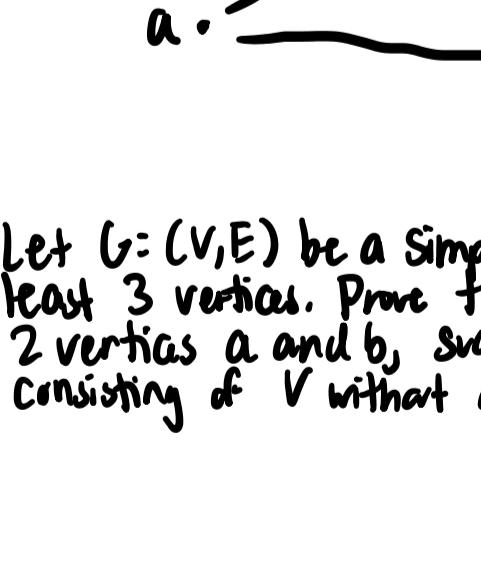
- Every connected graph  $G = (V, E)$  contains a spanning tree

- A hamiltonian cycle in  $G$ , is a cycle in which every vertex of  $G$  appears exactly once

- An eulerian circuit is a closed walk in which every edge appears exactly once

- A connected graph is Eulerian (has an Eulerian circuit) iff every vertex has even degree

#### Directed Graphs : $G$



$$\text{Outdegree}(b) = 2$$

$$\text{indegree}(d) = 2$$

$$\text{total degree}(b) = 3 = \text{outdegree}(b) + \text{indegree}(b)$$

$$\text{sources}(G) = []$$

$$\text{sinks}(G) = [c]$$

- SCCs ( $G$ ):  $\text{SCC}_1 \rightarrow \text{SCC}_2$

$$[a, b, d, e, f] \uparrow [c]$$

- $G_{\text{SCC}}$  is a DAG

$$\tau(G_{\text{SCC}}) = [\text{SCC}_1, \text{SCC}_2]$$

- Sum of outdegree of all nodes

$$= \sum \text{outdegree of all nodes}$$

- First vertex in a topo-sort is a source, last is a sink

#### Graph Coloring:

- A graph is  $k$ -colorable if each vertex can be colored using one of the  $k$  colors so that adjacent vertices are colored using different colors

- The chromatic number of a graph,  $\chi(G)$ , is the smallest value of  $k$  for which  $G$  is  $k$ -colorable.

- A bipartite graph, is a graph that is 2-colorable

#### Practice Questions:

- Let  $G = (V, E)$  be a connected graph such that  $|E| = |V|$ . Prove that  $G$  has exactly one cycle.

- Assume F.P.O.C. that  $G$  doesn't have a cycle and is connected, therefore  $G$  is a tree and thus  $G$  has  $n-1$  edges

Contradiction!  $G$  has  $n$  edges!

- Every connected graph has a spanning tree with  $n$  vertices and  $n-1$  edges.

Add back the missing edge,  $\rightarrow$  creates exactly 1 cycle.

- Let  $T$  be a tree with at least 2 vertices, and no vertex in  $T$  has a degree which is larger than 3 ( $\Delta(T) \leq 3$ ). Let  $n_i$  be the number of vertices of degree exactly  $i$ , prove that  $n_1 = n_3 + 2$

Every vertex in  $T$   $\begin{cases} \text{Vertices of degree 1} \\ \text{Vertices of degree 2} \\ \text{Vertices of degree 3} \end{cases}$

let the total # of vertices in  $T$  be  $n$

$$n = n_1 + n_2 + n_3 \quad (1)$$

By the handshaking lemma:

$$2(n-1) = n_1(1) + n_2(2) + 3(n_3) \quad (2)$$

$$2(n-1) = 2n - 2$$

$$2n_1 + 2n_2 + 2n_3 - 2 = n_1 + 2n_2 + 3n_3$$

$$n_1 = n_3 + 2$$

- How many edges are required to make a graph with 3 connected components, connected?

2



- Consider a directed graph,  $G$ , with no self-loops and no cycles of length 2 (no parallel edges).

Define an undirected graph  $G' = (V, E')$  where  $G'$  has the same vertices as  $G$ , and moreover in  $G'$ , we have an edge  $\{x, y\}$  if  $x \rightarrow y$  OR  $y \rightarrow x$  are in  $G$ .

- Prove/disprove: If  $G$  is strongly connected, then  $G'$  is connected

Prove: W.T.S. that  $G'$  is connected and therefore for any pair of vertices  $a$  and  $b$ , I can get from  $a \rightarrow b$  in  $G'$  and a path from  $b \rightarrow a$ .

Because  $G$  is strongly connected, there is a path from  $a \rightarrow b$  in  $G$  and a path from  $b \rightarrow a$ .

Consider this first path:  $a \rightarrow b$

$$a \rightarrow x_1 \rightarrow \dots \rightarrow x_n \rightarrow b$$

Undirected ✓

- Prove or disprove: If  $G$  is a DAG, then  $G'$  is acyclic



- Let  $(V, E)$  be a simple connected graph with at least 3 vertices. Prove that there exists at least 2 vertices  $a$  and  $b$ , such that a subgraph of  $G$  consisting of  $V$  without  $a$  and  $b$  is connected

- An edge  $\{a, b\}$  is a chord of the cycle  $C$  in an undirected simple graph  $G = (V, E)$  if  $a$  and  $b$  are vertices in the cycle but the edge  $\{a, b\} \notin E$  is not an edge of the cycle



$\{a, b\}$

- Prove if  $\delta(G) \geq 3$ , then  $G$  contains a cycle with a chord.

Let P:  $v_1, v_2, \dots, v_n$  be a maximal path in  $G$



$\{v_d, v_n\}$

$\{v_d, v_n\}$  and  $\{v_b, v_n\}$  exist

$$v_d \rightarrow v_n \rightarrow v_{n-1} \rightarrow \dots \rightarrow v_b \rightarrow v_n$$

Contradiction!  $\rightarrow$   $\{v_d, v_n\}$  and  $\{v_b, v_n\}$  don't exist

- Final Announcement:

- Very likely, this is my final class this semester!

• Thanks for being a fantastic 1st class at GW, and don't hesitate to reach out (email / LinkedIn)

- I think I'll be teaching Discrete I next semester, and looking for TAs

- Official applications are soon... let's chat

- Summer 2025 - Research Endeavor?

- Class Social ~ 12/13? Next Semester?

- Good luck on the exam, and the final exam (cumulative)