

Announcements

- Homework 5 due October 1
- Midterm #1 is October 8
 - please arrive ON TIME
 - if you need any accommodations or need to take a make-up please e-mail us TODAY
 - review session in lab next week
 - Questions?

① $2^{n+1} \in O(2^n)??$

$$2^{n+1} \leq c \cdot 2^n$$

$$2^{n+1} = 2 \cdot 2^n$$

We need to show that

$$\exists b, B > 0 \text{ s.t. } 0 \leq |f(n)| \leq g(n) \cdot B \quad \forall n \geq b$$

$$0 \leq 2^{n+1} \leq 2^n \cdot B$$

E.g. $b=1; B=2$

$$0 \leq 2 \cdot 2^n \leq 2 \cdot 2^n \quad \forall n \geq 1$$

② Is $2^{2^n} \in O(2^n)??$ Assume it is.

No! What would it mean for it to be $O(2^n)$.

$$\exists b, B \text{ s.t. } 2^{2^n} \leq B \cdot 2^n \quad \forall n \geq b \quad (*)$$

$$2^{2^n} = 2^n \cdot 2^n \leq B \cdot 2^n$$

Divide by 2^n

$$\Rightarrow 2^n \leq B \quad \forall n \geq b$$

We can express (*) as a conditional statement

$$\exists b, B \text{ s.t. } n \geq b \Rightarrow 2^{2^n} \leq B \cdot 2^n$$

$$\Rightarrow \log_2(2^{2^n}) \leq \log_2(B \cdot 2^n)$$

$$\Rightarrow 2n \leq \log_2 B + n$$

$$\Rightarrow n \leq \log_2 B$$

Consider $n' > \max(b, \log_2 B)$.

This is : $n' \geq b$ but also $n' \geq \log_2 B$

OCTOBER 1, 2024

(2) from worksheet

$$\boxed{f \in O(g)} \Rightarrow kf(n) \in O(g(n)) \text{ for } k \in \mathbb{R} \quad k \neq 0$$

$$f \in O(g)$$

$$\Rightarrow \exists b, B > 0 \text{ s.t. } 0 \leq |f(n)| \leq B \cdot g(n) \quad \forall n \geq b$$

We want to show:

$$\exists b', B' > 0 \text{ s.t. } 0 \leq |k \cdot f(n)| \leq B' \cdot g(n) \quad \forall n \geq b'$$

$$\Rightarrow \exists b, B > 0 \text{ s.t. } 0 \leq |k| \cdot |f(n)| \leq |k| \cdot B \cdot g(n) \quad \forall n \geq b$$

Because $|k| > 0$, the direction of the inequality is preserved.

$$\text{Define } B' = |k| \cdot B$$

$$B' > 0$$

So there exists a $B' = |k| \cdot B > 0$ and $b' = b > 0$

$$\text{s.t. } 0 \leq |k \cdot f(n)| \leq B' \cdot g(n) \quad \forall n \geq b'$$

By the def of Big-Oh, $kf \in O(g)$

(3) Is big-Oh reflexive?

$$f(n) \in O(f(n))$$

$$\exists b, B > 0 \text{ s.t. } 0 \leq |f(n)| \leq B \cdot g(n) \quad \forall n \geq b$$

$$B=1; b=1$$

$$f(n) = 1 \times f(n) \quad \forall n \geq 1$$

$$|f(n)| \leq 1 \times f(n) \quad \forall n \geq 1$$

(c) Transitive?

$$\boxed{f(n) \in O(g(n)) \wedge g(n) \in O(h(n))} \Rightarrow f(n) \in O(h(n))$$

$$\Rightarrow \exists b, B > 0 \text{ s.t. } 0 \leq f(n) \leq B \cdot g(n) \quad \forall n \geq b$$

and

$$\exists b', B' > 0 \text{ s.t. } 0 \leq g(n) \leq B' \cdot h(n) \quad \forall n \geq b'$$

[Big-Oh def]

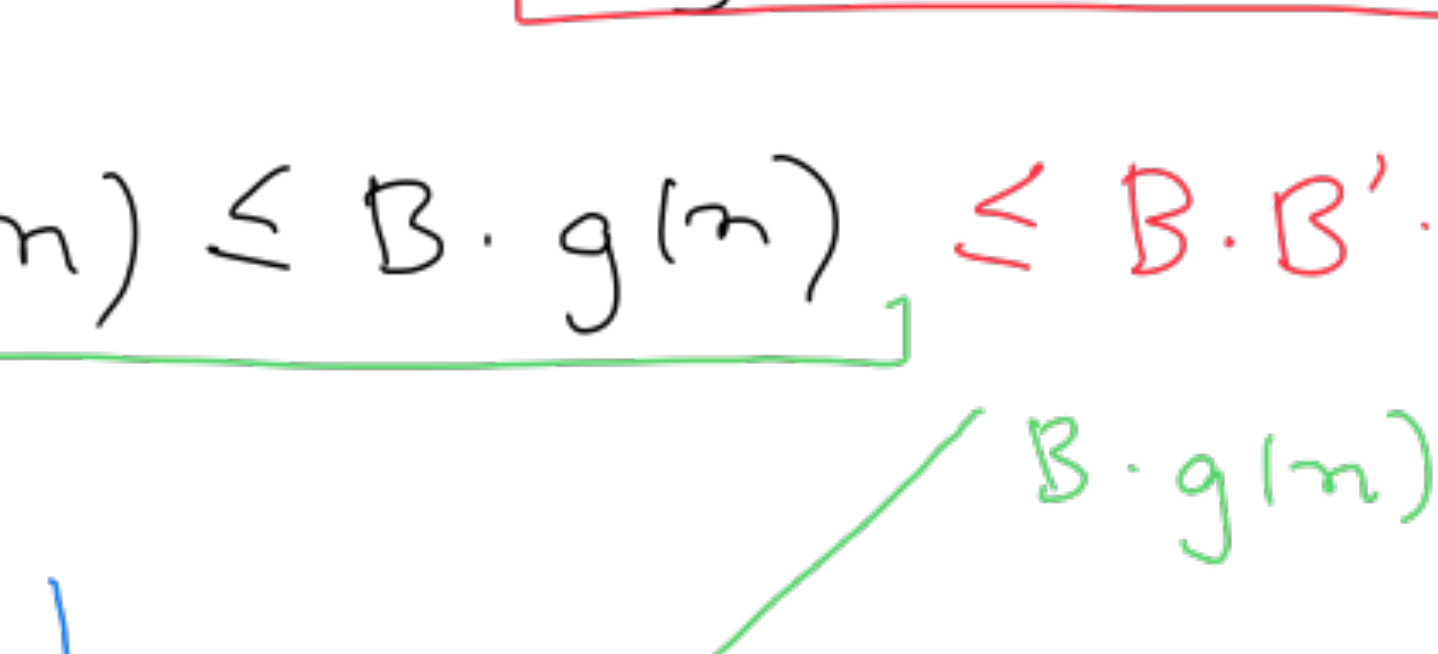
We want to show:

$$\exists b'', B'' > 0 \text{ s.t. } 0 \leq f(n) \leq B'' \cdot h(n) \quad \forall n \geq b''$$

(or that $f \in O(h)$)

$$\Rightarrow \exists b', B' > 0 \text{ s.t. } 0 \leq B \cdot g(n) \leq B \cdot B' \cdot h(n) \quad \forall n \geq b'$$

$$\Rightarrow 0 \leq f(n) \leq B \cdot g(n) \leq B \cdot B' \cdot h(n) \quad \forall n \geq \max(b, b')$$



↓
Because this is true when n is bigger than both b and b'

$$\Rightarrow B'' = B \cdot B' > 0 \quad b/c \quad B > 0 \text{ and } B' > 0 \quad (\text{by def})$$

$$b'' = \max(b, b') > 0 \quad b/c \quad b > 0 \text{ and } b' > 0 \quad (\text{by def})$$

$$\exists b'', B'' > 0 \text{ s.t. } 0 \leq f(n) \leq B'' \cdot h(n) \quad \forall n \geq b''$$

$$\Rightarrow f \in O(h) \quad \checkmark$$

(4) Show that for any $k, d \geq 0$, $n^k \in O(n^d) \Leftrightarrow d \geq k$

(a) First let's show that $n^k \in O(n^d) \Rightarrow d \geq k$ for $k, d \geq 0$

Let's do a proof by contradiction.

Assume $d < k$

$$n^k \in O(n^d)$$

$$\Rightarrow \exists b, B > 0 \text{ s.t. } 0 \leq n^k \leq B \cdot n^d \quad \forall n \geq b$$

$$\Rightarrow \exists b, B > 0 \text{ s.t. } 0 \leq n^{k-d} \leq B \quad \forall n \geq b$$

$$\Rightarrow \exists b, B > 0 \text{ s.t. } 0 \leq n \leq B^{\frac{1}{k-d}} \quad \forall n \geq b$$

(raising everything $\wedge \frac{1}{k-d}$)

$$\text{Consider } n > \max(b, B^{\frac{1}{k-d}})$$

$$n > b, \text{ so the inequality } 0 \leq n \leq B^{\frac{1}{k-d}}$$

$$\text{should hold, but } n > B^{\frac{1}{k-d}}$$

\Rightarrow contradiction!

(b) Now let's show $d \geq k \Rightarrow n^k \in O(n^d)$

We'll want to show that:

$$\exists b, B > 0 \text{ s.t. } n^k \leq B \cdot n^d \quad \forall n \geq b$$

$$d \geq k$$

$$\Rightarrow n^d \geq n^k \quad \forall n \geq 1$$

$$\Rightarrow n^k \leq 1 \cdot n^d \quad \forall n \geq 1$$

$$\text{so } \exists b, B=1 \text{ s.t. } n^k \leq B \cdot n^d \quad \forall n \geq b$$

$$\Rightarrow n^k \in O(n^d)$$

(6) Let $f, g : \mathbb{N} \rightarrow \mathbb{R}^+$ show that $f \in O(g) \Leftrightarrow g \in \Omega(f)$

To show that $g \in \Omega(f)$:

$$\exists a, A > 0 \text{ s.t. } A f(n) \leq g(n) \quad \forall n \geq a \quad (1)$$

$$f \in O(g)$$

$$\Rightarrow \exists b, B > 0 \text{ s.t. } f(n) \leq B \cdot g(n) \quad \forall n \geq b \quad (2)$$

(we want to make (2) look like (1))

$$\Rightarrow \exists b, B > 0 \text{ s.t. } \frac{1}{B} f(n) \leq g(n) \quad \forall n \geq b$$

$$\exists A = \frac{1}{B} > 0 \text{ and } a = b > 0 \text{ s.t.}$$

$$A f(n) \leq g(n) \quad \forall n \geq a$$