

ECE 1125 Data Structures and Algorithms – Fall 2015 – Exam 1 Solution

1. (a) Rank of complexities from least complex to most complex:

$\Theta(1)$, $\Theta(\log(N))$, $\Theta(N)$, $\Theta(N \cdot \log(N))$, $\Theta(N^2)$, $\Theta(N^3)$, $\Theta(N!)$.

1. (b) Operations on a stack (assumes top starts at 0):

```
pop (returns error, top = 0)
pop (returns error, top = 0)
push 45 (returns success, top = 1, A = [45])
push 67 (returns success, top = 2, A = [45, 67])
pop (returns success, copies out 67, top = 1, A = [45])
pop (returns success, copies out 45, top = 0, A = [])
pop (returns error)
push 88 (returns success, top = 1, A = [88])
push 11 (returns success, top = 2, A = [88, 11])
pop (returns success, copies out 11, top = 1, A = [88])
```

1. (c) Operations on a queue of size 5:

```
Initial state: {h=0, t=0, A[X, X, X, X, X]}
deq (returns error, state unchanged)
deq (returns error, state unchanged)
enq 5 (returns success, st: {h=0, t=1, A[5, X, X, X, X]})
enq 4 (returns success, st: {h=0, t=2, A[5, 4, X, X, X]})
enq 3 (returns success, st: {h=0, t=3, A[5, 4, 3, X, X]})
deq (returns success, copies out 5, st:{h=1, t=3, A[5, 4, 3, X, X]})
enq 2 (returns success, state: {h=1, t=4, A[5, 4, 3, 2, X]})
enq 1 (returns success, state: {h=1, t=0, A[5, 4, 3, 2, 1]})
deq (returns success, copies out 4, st:{h=2, t=0, A[5, 4, 3, 2, 1]})
deq (returns success, copies out 3, st:{h=3, t=0, A[5, 4, 3, 2, 1]})
enq 0 (returns success, state: {h=3, t=1, A[0, 4, 3, 2, 1]})
enq -1 (returns success, state: {h=3, t=2, A[0, -1, 3, 2, 1]})
```

2. (a) Declaration and instantiation of doubly-linked list to hold double-precision floating point numbers:

```
struct list_node {
    double data;
    struct list_node *prev;
    struct list_node *next;
};
struct list_node *head = NULL;
```

2. (b) Function in C:

```
void print_reverse_below_mean(struct list_node *head)
{
    int count = 0;
    double mean = 0.0;
    struct list_node *n = head;
    while (n) {
        mean += n->data;
        count++;
        n = n->next;
    }

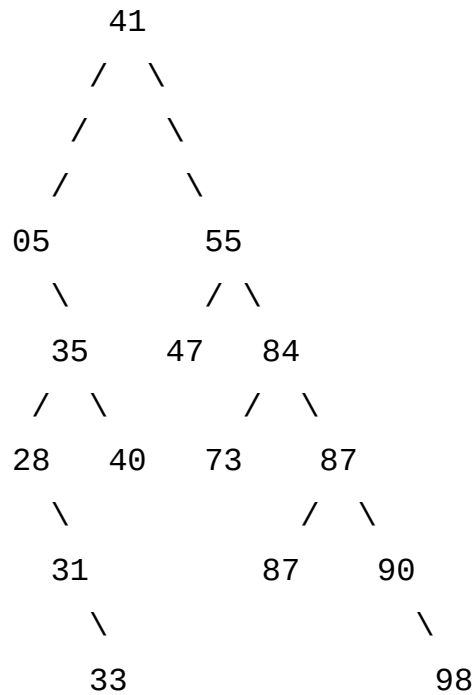
    if (count > 0) {
        mean /= count;
    }

    n = head;
    while (n && n->next) {
        n = n->next;
    }

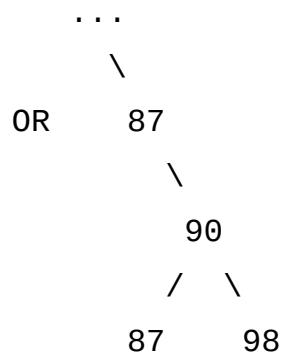
    while (n) {
        if (n->data < mean) {
            printf("%f\n", n->data);
        }
        n = n->prev;
    }
}
```

2. (c) Time Complexity: $\Theta(N)$. Space Complexity: $\Theta(1)$.

3. (a) Many solutions are possible. Binary search tree when inserted in the given order:

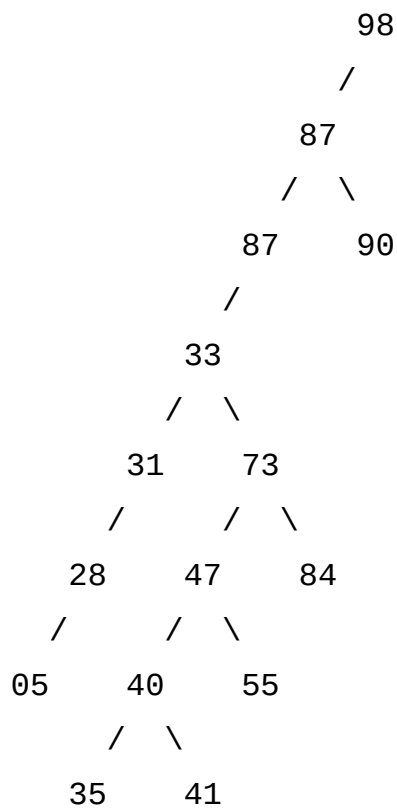


If Left \leq Parent < Right

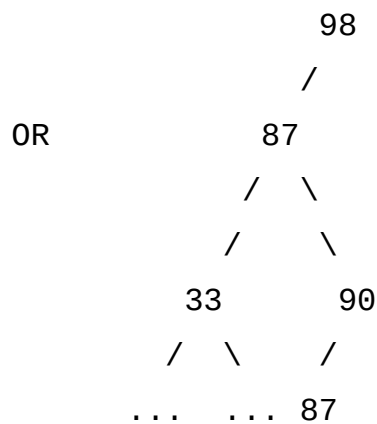


If Left < Parent \leq Right

Next is shown inserting in reverse order:

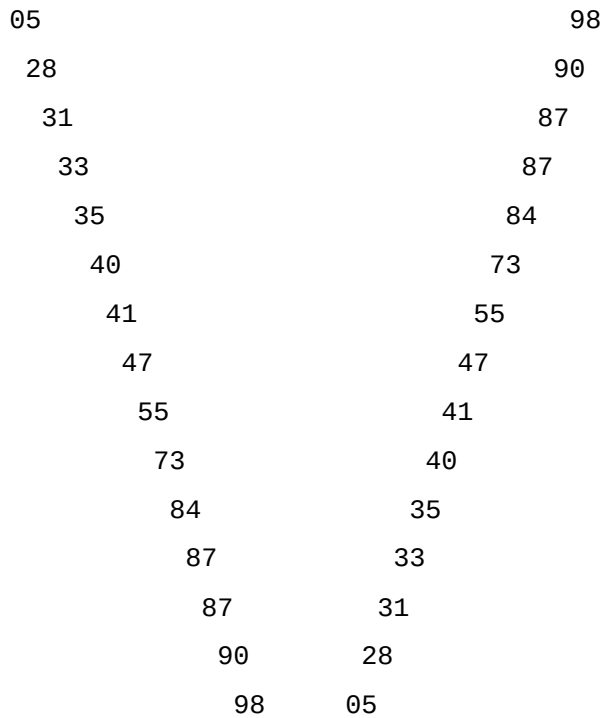


If Left \leq Parent < Right

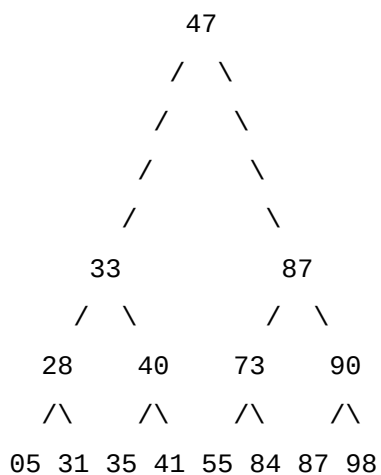


If Left < Parent \leq Right

3. (b) Two worst case binary search trees:



3. (c) Perfectly balanced tree (only Left < Parent <= Right is possible for perfect):



4. Recursive function for printing a BST in reverse order:

```

void print_reverse(struct tree_node *n)
{
    if (n) {
        print_reverse(n->right);
        printf("%d\n", n->data);
        print_reverse(n->left);
    }
}
  
```

5. (a) Approximate with $2^{10} \approx 10^3$ and solve for k:

$$2^k - 1 = 10^{11}$$

$$2^k \approx 10^{11}$$

$$\log_2(2^k) \approx \log_2(10^{11})$$

$$k \approx \log_2((10^3)^3 * 100)$$

$$k \approx \log_2((2^{10})^3 * 100)$$

$$k \approx 30 + \log_2(100)$$

The maximum depth must be an integer, so use the ceiling of this:

$$\text{ceil}(k) = 30 + 7 = 37$$

(b) The maximum number of comparisons to search is the maximum depth, 37.

(c) $10\text{ns} * 37 = 370\text{ ns}$.

(d) $10\text{ns} * 10^{11}\text{ comparisons} = 10 * 10^{(11-9)} = 10^3\text{ seconds}$.

6. (a) Time Complexity: $\Theta(\sqrt{N})$. Space Complexity: $\Theta(1)$.

(b) Each loop iteration performs 1 floating-point multiply, 1 floating-point comparison, and 1 floating-point add, so takes 4ns.

The loop runs until $y*y == N$, which occurs when $y = \sqrt{80000} = \sqrt{8} * 100 \approx 283$. Each whole number step takes 10^9 iterations of the loop, so the total time is:

$$283 * 10^9 * 4\text{ns} \approx 283 * 4\text{ seconds} = 1132\text{ seconds. (Approximations are ok.)}$$

(c) A faster binary-search based algorithm. (This is the basic idea in pseudo-code, this will not work for all inputs on a real machine without some modification.)

```
double mysqrt(double N)
{
    double lo = 0.0, hi = N, mid;
    for (;;) {
        mid = (lo + hi) / 2.;
        if (mid * mid == N) {
            return mid;
        } else if (mid * mid < N) {
            lo = mid;
        } else {
            hi = mid;
        }
    }
}

/* Note: A real robust version would break if mid repeats itself. */
```

6. (d) Time Complexity: $\Theta(\log_2(N))$. Space Complexity: $\Theta(1)$.