

Number Theory and the RSA Cryptosystem

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Why number theory?



Euclidean Algorithm and Linear Equation Theorem

Linear Equation Theorem

Can always find a pair (x, y) such that $ax + by = \gcd(a, b)$

Given: $22x + 60y = \gcd(22, 60)$

$$60 = 2 \times 22 + 16$$

$$22 = 1 \times 16 + 6$$

$$16 = 2 \times 6 + 4$$

$$6 = 1 \times 4 + 2$$

$$4 = 2 \times 2 + 0$$

Congruences

We say that " a is congruent to $b \pmod{m}$ "

$$a \equiv b \pmod{m}$$

Example

- $7 \equiv 2 \pmod{5}$
- $47 \equiv 35 \pmod{6}$

Fermat's little theorem

Theorem (Fermat's little theorem)

For a prime p and integer a , we have $a^{p-1} \equiv 1 \pmod{p}$

Example

■ $3^6 \equiv 1 \pmod{7}$

$x \pmod{7}$	1	2	3	4	5	6
$3x \pmod{7}$	3	6	2	5	1	4

■
$$\underbrace{(3 \cdot 1)(3 \cdot 2)(3 \cdot 3)(3 \cdot 4)(3 \cdot 5)(3 \cdot 6)}_{\text{numbers in second row}} \equiv \underbrace{1 \cdot 2 \cdot 3 \cdot 4 \cdot 5 \cdot 6}_{\text{numbers in first row}} \pmod{7}.$$

Euler's ϕ function and Euler's formula

Euler's ϕ function

The number of integers between 1 and m that are relatively prime to m .

Euler's formula

If $\gcd(a, m) = 1$, then $a^{\phi(m)} \equiv 1 \pmod{m}$

Powers modulo m

Solves really large powers $(\text{mod } m) : 5^{100000000000} (\text{mod } 12826832)$

Example

■ $7^{327} (\text{mod } 853)$

$$\begin{aligned}7^1 &\equiv 7 \equiv 7 \pmod{853} \\7^2 &\equiv (7^1)^2 \equiv 7^2 \equiv 49 \equiv 49 \pmod{853} \\7^4 &\equiv (7^2)^2 \equiv 49^2 \equiv 2401 \equiv 695 \pmod{853} \\7^8 &\equiv (7^4)^2 \equiv 695^2 \equiv 483025 \equiv 227 \pmod{853} \\7^{16} &\equiv (7^8)^2 \equiv 227^2 \equiv 51529 \equiv 349 \pmod{853} \\7^{32} &\equiv (7^{16})^2 \equiv 349^2 \equiv 121801 \equiv 675 \pmod{853} \\7^{64} &\equiv (7^{32})^2 \equiv 675^2 \equiv 455625 \equiv 123 \pmod{853} \\7^{128} &\equiv (7^{64})^2 \equiv 123^2 \equiv 15129 \equiv 628 \pmod{853} \\7^{256} &\equiv (7^{128})^2 \equiv 628^2 \equiv 394384 \equiv 298 \pmod{853}\end{aligned}$$

■

$$\begin{aligned}7^{327} &= 7^{256+64+4+2+1} \\&= 7^{256} \cdot 7^{64} \cdot 7^4 \cdot 7^2 \cdot 7^1 \\&\equiv 298 \cdot 123 \cdot 695 \cdot 49 \cdot 7 \pmod{853}.\end{aligned}$$

■

k^{th} roots modulo m

$$x^k \equiv b \pmod{m} \implies x = ??$$

Method

φ function $\implies ku - \varphi v = 1$ with Euclidean algorithm \implies successive squaring

Example

- $x^{131} \equiv 758 \pmod{1073}$
- $\varphi(1073) = 28 \times 36 = 1008$
- $131u - 1008v = 1 \implies 131 \times 731 - 1008 \times 95 = 1$
- Notice: $x^{131^{731}} = x^{131 \times 731} = x^{1 + 1008 \times 95} = x \times x^{1008 \times 95}$
- Euler's : $x^{1008} \equiv 1 \pmod{1073}$
- $x^{131^{731}} \equiv x \pmod{1073}$
- $x \equiv x^{131^{731}} \equiv 758^{731} \pmod{1073} \equiv 905 \pmod{1073}$

Key Generation and Distribution

- 1 Choose two distinct prime numbers p and q
- 2 Compute $m = p \times q$
- 3 Compute $\phi(m) = (p - 1) \times (q - 1)$ since p and q are prime
- 4 Choose a number k that is relatively prime to $\phi(m)$

Private

- p, q - the two primes
- $\phi(m)$

Public

- k
- m

This is called a public-key cryptosystem.

RSA Encryption

Now...anyone who wants to send us a message uses the values of m and k to encode in the following manner:

- 1 Convert message into a string of digits:

A	B	C	D	E	F	G	H	I	J	K	L	M
11	12	13	14	15	16	17	18	19	20	21	22	23

N	O	P	Q	R	S	T	U	V	W	X	Y	Z
24	25	26	27	28	29	30	31	32	33	34	35	36

- 2 Break the string of digits into numbers less than m
- 3 Use successive squaring to compute $a^k \pmod{m}$
- 4 Example: $p = 73$, $q = 97$, $m = 7081$, $k = 347$

H	E	L	L	O
1815		2222		25
$1815^{347} \pmod{7081}$		$2222^{347} \pmod{7081}$		$25^{347} \pmod{7081}$
2212		6844		6593

RSA Decryption

Decryption requires using the k^{th} roots modulo m method which requires finding $\phi(m)$, easy for us if we know the factors p and q since $\phi(m) = (p - 1)(q - 1)$:

2212		6844		6593
$x^{347}(\text{mod } 7081)$		$x^{347}(\text{mod } 7081)$		$x^{347}(\text{mod } 7081)$
1815		2222		25
H	E	L	L	O



Code!

5301435910, 4794709296

$m = 6022651441$

$k = 57737$

References

-  Silver, Joseph, H.
A Friendly Introduction to Number Theory.
Pearson, 2012.
-  <https://www.math.brown.edu/johsilve/frint.html>