

## Introduction

### 1.1 Graph Isomorphism

**Definition 1.1:** An undirected graph  $G$  is a set of vertices,  $V(G)$ , and a function  $g : V(G) \rightarrow P(V(G))$ , defined by  $g(v) = N(v)$ , where  $N(v)$  is the set of vertices that are adjacent to  $v$  and  $P(V(G))$  is the power-set of  $V(G)$ .

**Remark:** We call  $N(v)$  the *neighborhood* of  $v$ .

**Remark:** Since the graphs we are concerned with are undirected, for any graph  $G$  and any vertices  $u, v \in V(G)$ ,  $u \in N(v)$  if and only if  $v \in N(u)$ .

**Definition 1.2:** Two graphs,  $G$  and  $G'$ , are isomorphic, denoted  $G \cong G'$ , if there exists a bijection  $f : V(G) \rightarrow V(G')$  such that for every  $u, v \in V(G)$ ,  $u \in N(v)$  if and only if  $f(u) \in N(f(v))$ .

# The String Isomorphism Problem

## 2.1 A graph as a string

Following Luks [1982], a graph can be represented as a string using the following procedure. Let  $\Omega = [1, \dots, n]$ . Let  $G_1$  and  $G_2$  be undirected graphs on  $n$  vertices, labelled using  $\Omega$ . Let  $\delta$  be the indicator function,  $\delta : \binom{V(G)}{2} \rightarrow \{0, 1\}$ , defined by

$$\delta(x, y) = \begin{cases} 1, & x \in N(y) \\ 0, & \text{o/w} \end{cases},$$

where  $\binom{V(G)}{2}$  denotes the set of all unordered pairs in  $V(G)$  for some undirected graph  $G$  (Babai 2018). To create a string from delta that represents  $G$ , we must order the pre-image of  $\delta$ . We choose a straightforward ordering. We have

$$S = \delta(1, 2) \dots \delta(1, n) \delta(2, 3) \dots \delta(2, n) \dots \delta(k, k+1) \dots \delta(k, n) \dots \delta(n-1, n).$$

Notice that  $S$  has a total length of  $\binom{n}{2}$ .



Figure 2.1: Two non-isomorphic graphs on 4 vertices.

**Example 2.1:** Consider graphs  $G_1$  and  $G_2$  in Figure 2.1. Since  $G_1$  has a vertex with degree 2 while  $G_2$  does not, we know immediately that the two are non-isomorphic. The string representations of  $G$  and  $G'$  are both of length  $\binom{4}{2} = 6$ . Applying the ordering used

above, the string representation of  $G_1$  is

$$001100,$$

while that of  $G_2$  is

$$101000.$$

Notice that the string representation of graph is only unique after an order has been established on  $\binom{V(G)}{2}$ . We also have

$$\text{Sym}(V(G)) = \{(), (12), (13), (14), (23), (24), (34), (12)(34), (13)(24), (14)(23)\}.$$

**Definition 2.3:** Let  $G$  be an undirected graph on  $n$  vertices. We define  $\text{Sym}(V(G))$  as the group of possible permutations of  $V(G)$ .

**Definition 2.4:** Let  $G$  be an undirected graph on  $n$  ordered vertices. Let  $\delta$  be the indicator function for  $G$  while  $S$  is the string representation of  $G$  produced by  $\delta$  and the order on  $\binom{V(G)}{2}$ .

Then

$$\text{Sym}(V(G))^{(2)} = \{\{\sigma a, \sigma b\} : a, b \in V(G), \sigma \in \text{Sym}(V(G))\}.$$

When does  $\delta(\{\sigma a, \sigma b\}) = \delta(\{a, b\})$  come into play (i.e. valid vs invalid)?

**Definition 2.5:** Let  $S_1$  and  $S_2$  be string representations of undirected graphs  $G_1$  and  $G_2$ , respectively. Then  $S_1$  and  $S_2$  are  $\text{Sym}(V(G_1))^{(2)}$ -isomorphic, denoted  $S_1 \cong_{\text{Sym}(V(G_1))^{(2)}} S_2$ , if there exists  $\bar{\sigma} \in \text{Sym}(V(G_1))^{(2)}$  such that  $S_1^{\bar{\sigma}} = S_2$ .

**Example 2.2:** Consider graphs  $G$  and  $G'$  in Fig. 2.1. Let the order of the string be 12, 13, 14, 23, 24, 34. To show that  $S \cong_{\text{Sym}(V(G))^{(2)}} S'$ , let  $\sigma$  be the permutation of  $S'$  that is induced by (14)(23) applied to  $V(G')$ . Then  $\sigma \in \text{Sym}(V(G))^{(2)}$ . QUESTION: What is  $\sigma$  in this case? Hence, by Definition 2.3,  $S \cong_{\text{Sym}(V(G))^{(2)}} S'$ .



Figure 2.2: Undirected, isomorphic graphs  $G$  and  $G'$

**Lemma 2.1:** Two graphs are isomorphic if and only if their string representations are  $\text{Sym}(V(G))^{(2)}$ -isomorphic.

*Proof.*  $\rightarrow$

Let  $S$  and  $S'$  be string representations of undirected graphs  $G$  and  $G'$ , respectively. Suppose  $G \cong G'$ . Without loss of generality, suppose  $G$  and  $G'$  each have  $n$  vertices. By Definition 1.2, there exists a bijection  $f : V(G) \rightarrow V(G')$  such that for every  $u, v \in V(G)$ ,  $u \in N(v)$  if and only if  $f(u) \in N(f(v))$ . This means that for every vertex  $u \in G$ , there is a vertex  $f(u) \in G'$  such that  $f(N(u)) = N(f(u))$ . Recall that  $S$  and  $S'$  both contain  $n$  binary  $n$ -tuples that describe the neighborhood of each vertex in graphs  $G$  and  $G'$ , respectively. We want to show that there exists  $\sigma \in \text{Sym}(V(G))^{(2)}$  such that  $S^\sigma = S'$ . I'm stuck here. I'm having trouble figuring out how to deal with elements of  $\text{Sym}(V(G))^{(2)}$ .

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