

Combinatorial Games: Open Problems and Computational Approaches

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Hood College

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Game Solutions and Winning Strategies





Two Players (Left/Right)





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- No Chance or Randomness







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(Finite) Combinatorial Games





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Let's play a game!





Rules for Chomp





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 Players select a square in remaining in the board, and remove all squares above and to the right of it.



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- Player forced to take the bottom left square loses.



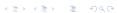


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Goal: Find solutions to games.





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- **Ultra-Weak:** There exists a strategy for the first (or second) player to guarantee a win, given optimal play on both sides.
- **Weak:** There is an algorithm to *describe* the optimal moves to secure a win or draw for the first (or second) player from the initial position of the game.
- **Strong:** There is an algorithm to secure a win for a player from *any* game position, even if sub-optimal moves were made in previous play.



Poset Games

Definition of Posets

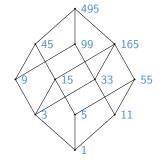
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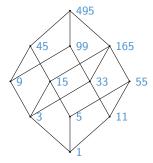
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Poset Games

Given a non-empty poset P, players alternate turns selecting an element $x \in P$ and removing x and all elements y with $x \leq y$. The loser takes the last element remaining in P.

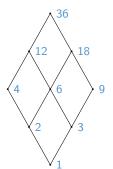


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Game of Divisors

Given a positive integer n, the game of divisors on n is the poset game of the divisor poset D_n .

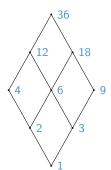




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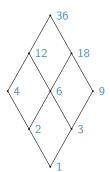


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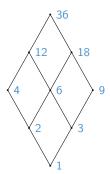
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What about a weak or strong solution?





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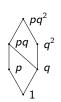












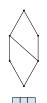
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One board of game for each antichain of the poset.























Sprague-Grundy Function for Games

Sprague-Grundy Function

Let \mathcal{B}_L be the set of boards in game G for which a player has no permissible moves (losing boards.)

- Set SG(B) = 0 for each board $B \in B_L$.
- Define $SG(B) = \max(\mathcal{M}_B)$, recursively for each $B \in \mathcal{B}$.



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Game Solutions



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 - Solve $D(p_1p_2\cdots p_k)$.





MAA MD-VA-DC Section Meeting: Fall 2015

MAA Section Meeting, Fall 2015

Thanks for listening!

Code from this talk is available at Github and on my website at:

https://github.com/

gwynwhieldon/PosetGames

http://cs.hood.edu/~whieldon

