

# CIS 700: “algorithms for Big Data”

## Lecture 5: Dimension Reduction

Slides at <http://grigory.us/big-data-class.html>

**Grigory Yaroslavtsev**

<http://grigory.us>



# Today

- Dimensionality reduction
  - AMS as dimensionality reduction
  - Johnson-Lindenstrauss transform

# $L_p$ -norm Estimation

- Stream:  $m$  updates  $(x_i, \Delta_i) \in [n] \times \mathbb{R}$  that define vector  $f$  where  $f_j = \sum_{i: x_i=j} \Delta_i$ .
- **Example:** For  $n = 4$

$$\langle (1,3), (3, 0.5), (1,2), (2, -2), (2,1), (1, -1), (4,1) \rangle$$
$$f = (4, -1, 0.5, 1)$$

- $L_p$ -norm:  $\|f\|_p = (\sum_i |f_i|^p)^{\frac{1}{p}}$

# $L_p$ -norm Estimation

- $L_p$ -norm:  $\|f\|_p = (\sum_i |f|^p)^{\frac{1}{p}}$
- Two lectures ago:
  - $\|f\|_0 = F_0$ -moment
  - $\|f\|_2^2 = F_2$ -moment (via AMS sketching)
- Space:  $O\left(\frac{\log n}{\epsilon^2} \log \frac{1}{\delta}\right)$
- Technique: linear sketches
  - $\|f\|_0$ :  $\sum_{i \in S} f_i$  for random sets  $S$
  - $\|f\|_2^2$ :  $\sum_i \sigma_i f_i$  for random signs  $\sigma_i$

# AMS as dimensionality reduction

- Maintain a “linear sketch” vector

$$\mathbf{Z} = (Z_1, \dots, Z_k) = Rf$$

$$Z_i = \sum_{j \in [n]} \sigma_{ij} f_j, \text{ where } \sigma_{ij} \in_R \{-1, 1\}$$

- Estimator  $\mathbf{Y}$  for  $\|f\|_2^2$ :

$$\frac{1}{k} \sum_{i=1}^k Z_i^2 = \frac{\|Rf\|_2^2}{k}$$

- “Dimensionality reduction”:  $x \rightarrow Rx$ , “heavy” tail

$$\Pr \left[ \left| \mathbf{Y} - \|f\|_2^2 \right| \geq c \left( \frac{2}{k} \right)^{\frac{1}{2}} \|f\|_2^2 \right] \leq \frac{1}{c^2}$$

# Normal Distribution

- Normal distribution  $N(0,1)$ 
  - Range:  $(-\infty, +\infty)$
  - Density:  $\mu(x) = (2\pi)^{-\frac{1}{2}} e^{-\frac{x^2}{2}}$
  - Mean = 0, Variance = 1
- Basic facts:
  - If  $X$  and  $Y$  are independent r.v. with normal distribution then  $X + Y$  has normal distribution
  - $Var[cX] = c^2 Var[X]$
  - If  $X, Y$  are independent, then  $Var[X + Y] = Var[X] + Var[Y]$

# Johnson-Lindenstrauss Transform

- Instead of  $\pm 1$  let  $\sigma_i$  be i.i.d. random variables from normal distribution  $N(0,1)$

$$Z = \sum_i \sigma_i f_i$$

- We still have  $\mathbb{E}[Z^2] = \sum_i f_i^2 = \|f\|_2^2$  because:
  - $\mathbb{E}[\sigma_i]\mathbb{E}[\sigma_j] = 0$ ;  $\mathbb{E}[\sigma_i^2] = \text{“variance of } \sigma_i \text{”} = 1$
- Define  $\mathbf{Z} = (Z_1, \dots, Z_k)$  and define:

$$\|\mathbf{Z}\|_2^2 = \sum_j Z_j^2 \quad \left( \mathbb{E}[\|\mathbf{Z}\|_2^2] = k \|f\|_2^2 \right)$$

- **JL Lemma:** There exists  $C > 0$  s.t. for small enough  $\epsilon > 0$ :

$$\Pr \left[ \left| \|\mathbf{Z}\|_2^2 - k \|f\|_2^2 \right| > \epsilon k \|f\|_2^2 \right] \leq \exp(-C \epsilon^2 k)$$

# Proof of JL Lemma

- **JL Lemma:**  $\exists C > 0$  s.t. for small enough  $\epsilon > 0$ :  
$$\Pr \left[ \left| \|\mathbf{Z}\|_2^2 - k \|f\|_2^2 \right| > \epsilon k \|f\|_2^2 \right] \leq \exp(-C \epsilon^2 k)$$
- Assume  $\|f\|_2^2 = 1$ .
- We have  $\mathbf{Z}_i = \sum_j \sigma_{ij} f_j$  and  $\mathbf{Z} = (\mathbf{Z}_1, \dots, \mathbf{Z}_k)$   
$$\mathbb{E} \left[ \|\mathbf{Z}\|_2^2 \right] = k \|f\|_2^2 = k$$
- **Alternative form of JL Lemma:**  
$$\Pr \left[ \|\mathbf{Z}\|_2^2 > k(1 + \epsilon)^2 \right] \leq \exp(-\epsilon^2 k + O(k \epsilon^3))$$



# Proof of JL Lemma

- Alternative form of JL Lemma:

$$\Pr \left[ \|\mathbf{Z}\|_2^2 > k(1 + \epsilon)^2 \right] \leq \exp(-\epsilon^2 k + O(k \epsilon^3))$$

- Let  $Y = \|\mathbf{Z}\|_2^2$  and  $\alpha = k(1 + \epsilon)^2$
- For every  $s > 0$  we have:

$$\Pr[Y > \alpha] = \Pr[e^{sY} > e^{s\alpha}]$$

- By Markov and independence of  $\mathbf{Z}'_i$ s:

$$\Pr[e^{sY} > e^{s\alpha}] \leq \frac{\mathbb{E}[e^{sY}]}{e^{s\alpha}} = e^{-s\alpha} \mathbb{E} \left[ e^{s \sum_i Z_i^2} \right] = e^{-s\alpha} \prod_{i=1}^k \mathbb{E} \left[ e^{s Z_i^2} \right]$$

- We have  $Z_i \in N(0,1)$ , hence:

$$\mathbb{E} \left[ e^{s Z_i^2} \right] = (2\pi)^{-\frac{1}{2}} \int_{-\infty}^{\infty} e^{s t^2} e^{-\frac{t^2}{2}} dt = \frac{1}{\sqrt{1 - 2s}}$$

# Proof of JL Lemma

- Alternative form of JL Lemma:

$$\Pr \left[ \|\mathbf{Z}\|_2^2 > k(1 + \epsilon)^2 \right] \leq \exp(-\epsilon^2 k + O(k \epsilon^3))$$

- For every  $\mathbf{s} > 0$  we have:

$$\Pr[\mathbf{Y} > \alpha] \leq e^{-\mathbf{s}\alpha} \prod_{i=1}^k \mathbb{E} \left[ e^{\mathbf{s}Z_i^2} \right] = e^{-\mathbf{s}\alpha} (1 - 2\mathbf{s})^{-\frac{k}{2}}$$

- Let  $\mathbf{s} = \frac{1}{2} \left( 1 - \frac{k}{\alpha} \right)$  and recall that  $\alpha = k(1 + \epsilon)^2$
- A calculation finishes the proof:

$$\Pr[\mathbf{Y} > \alpha] \leq \exp(-\epsilon^2 k + O(k \epsilon^3))$$

# Johnson-Lindenstrauss Transform

- Single vector:  $k = O\left(\frac{\log \frac{1}{\delta}}{\epsilon^2}\right)$ 
  - Tight:  $k = \Omega\left(\frac{\log \frac{1}{\delta}}{\epsilon^2}\right)$  [Woodruff'10]
- $n$  vectors simultaneously:  $k = O\left(\frac{\log n \log \frac{1}{\delta}}{\epsilon^2}\right)$ 
  - Tight:  $k = \Omega\left(\frac{\log n \log \frac{1}{\delta}}{\epsilon^2}\right)$  [Molinaro, Woodruff, Y. '13]
- Distances between  $n$  vectors =  $O(n^2)$  vectors:
$$k = O\left(\frac{\log n \log \frac{1}{\delta}}{\epsilon^2}\right)$$