

# CS 497: Cybersecurity

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# Computer Networks

# Circuit and Packet Switching

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## • Circuit switching

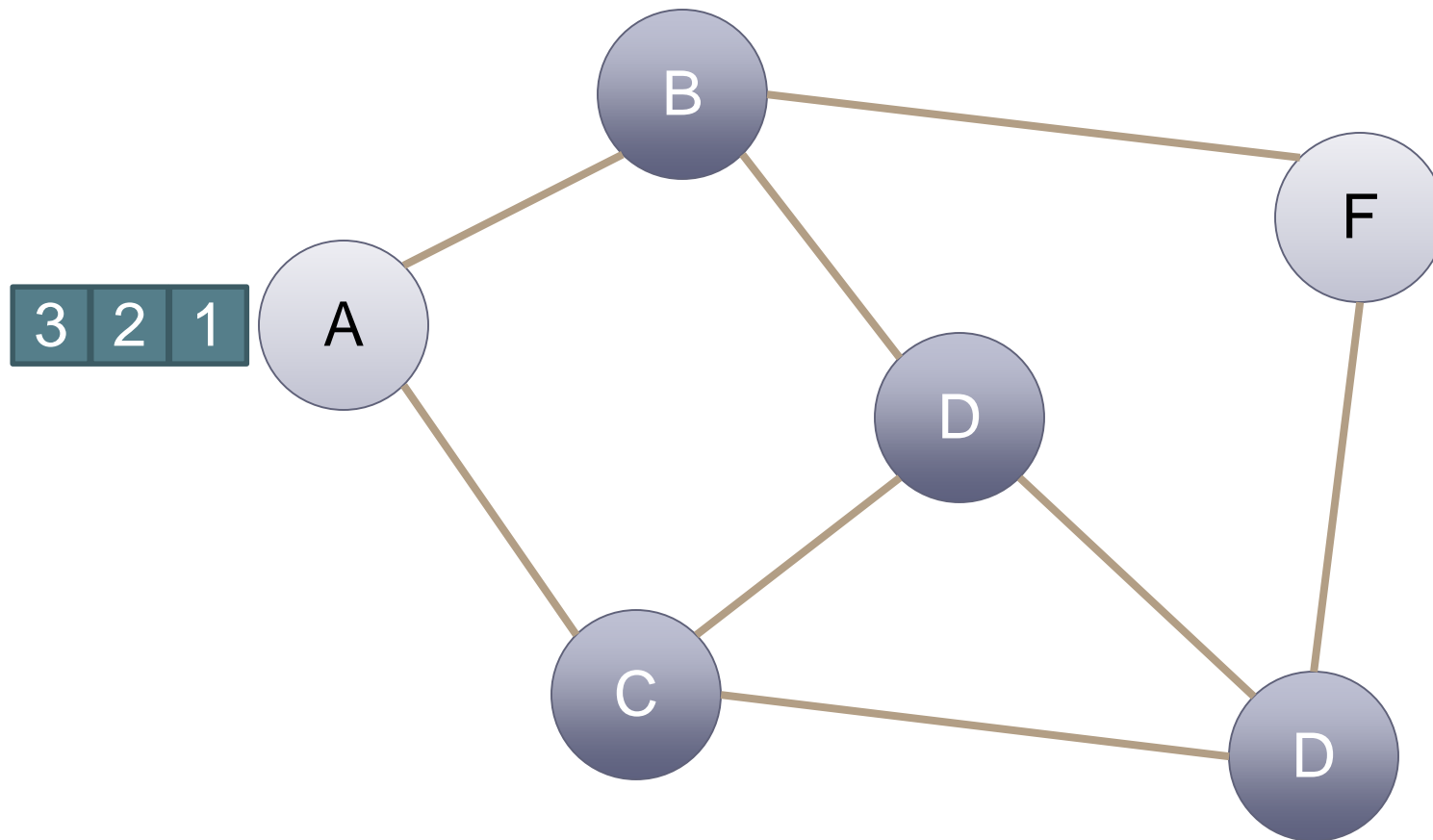
- Legacy phone network
- Single route through sequence of hardware devices established when two nodes start communication
- Data sent along route
- Route maintained until communication ends

## • Packet switching

- Internet
- Data split into packets
- Packets transported independently through network
- Each packet handled on a best efforts basis
- Packets may follow different routes

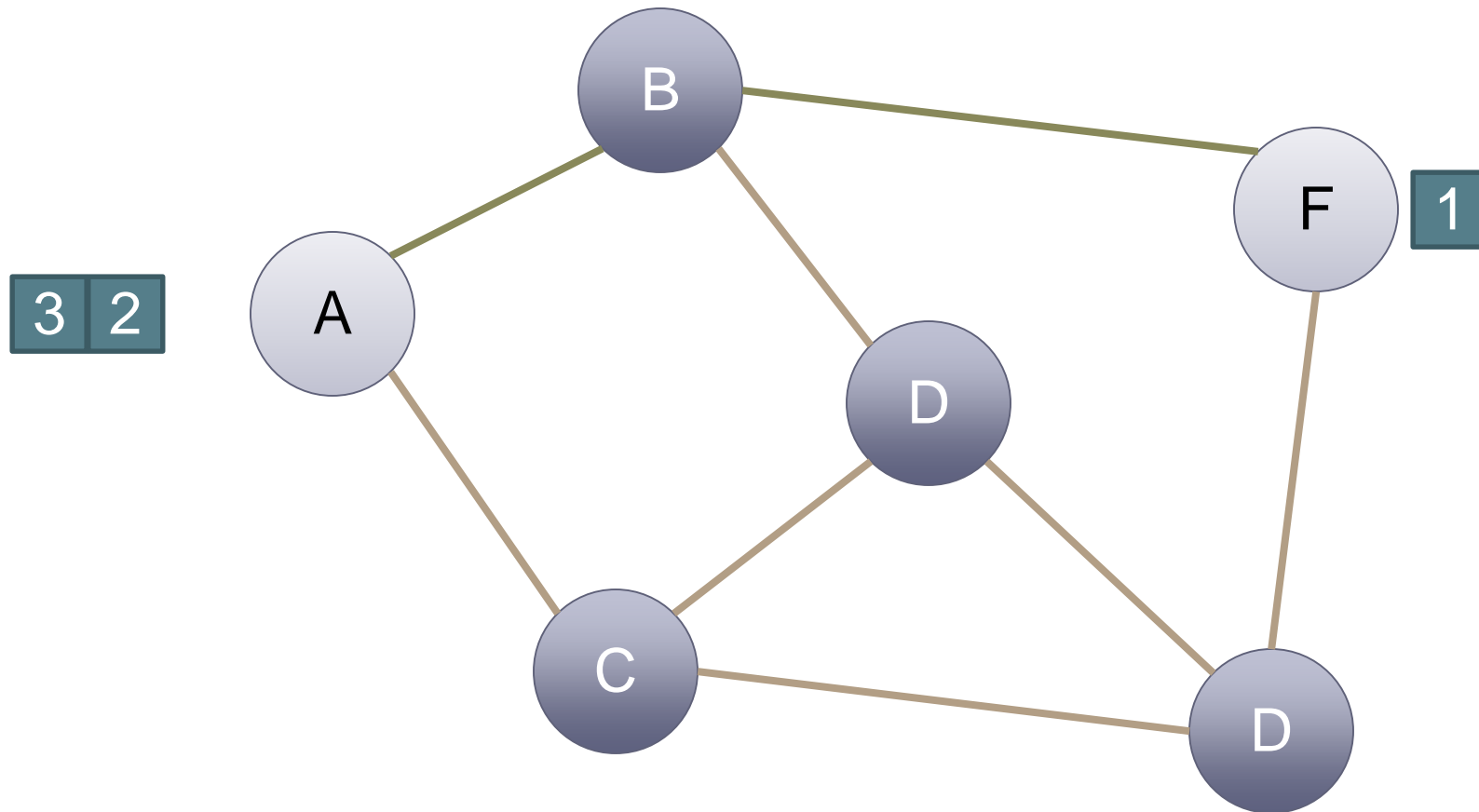
# Packet Switching

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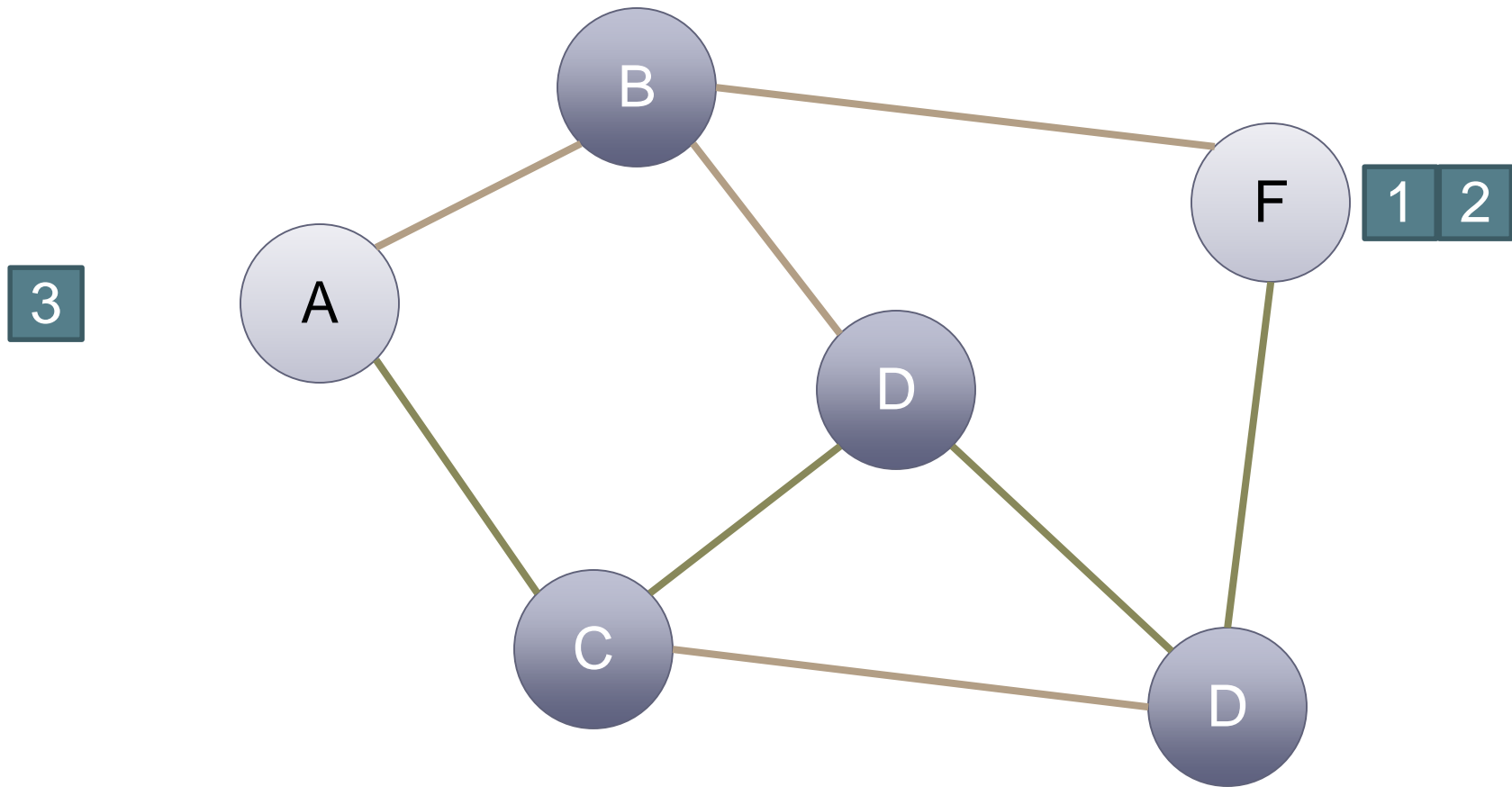
# Packet Switching

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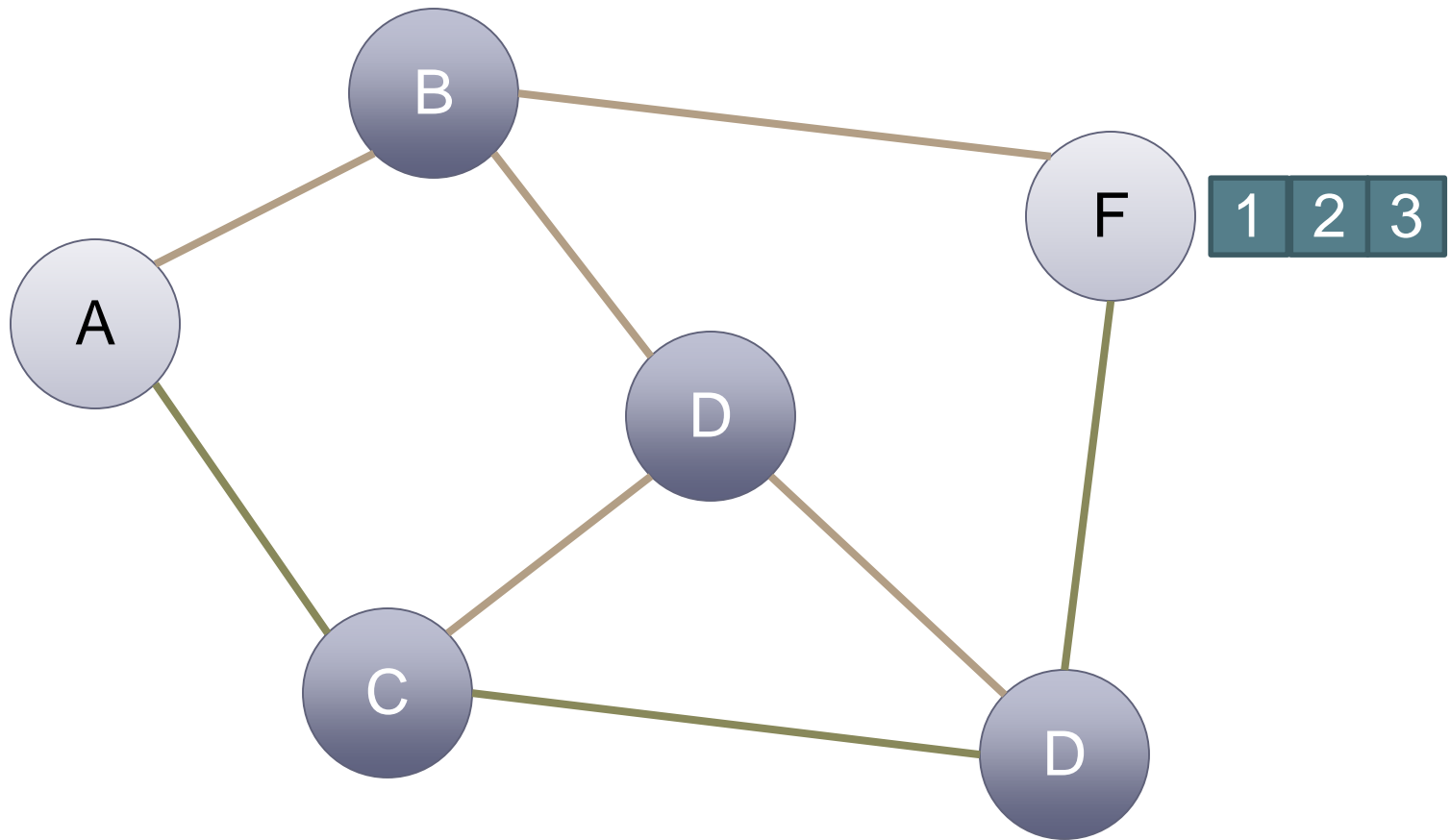
# Packet Switching

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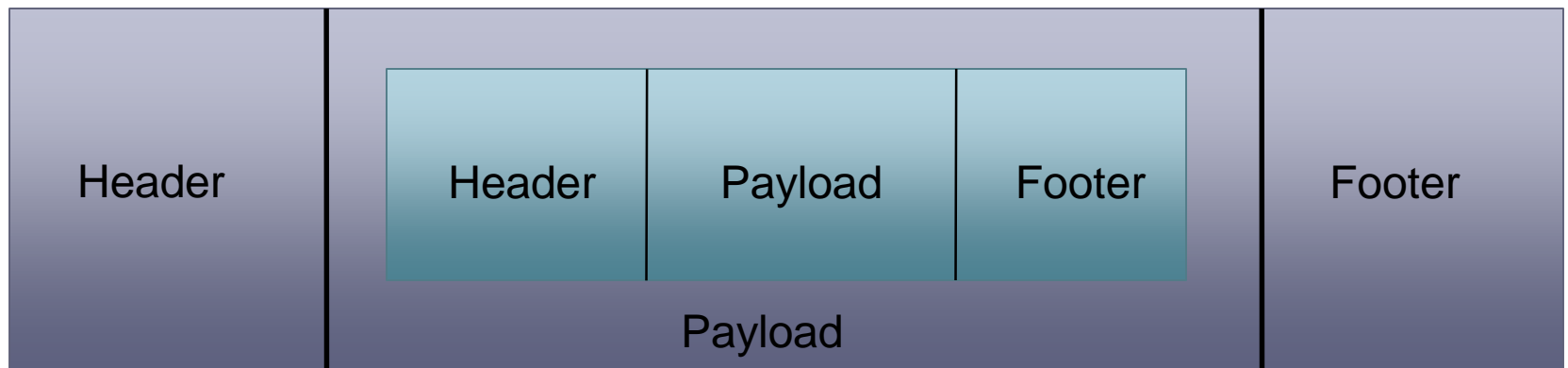
# Packet Switching

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# Encapsulation

- **A packet typically consists of**
  - Control information for addressing the packet: header and footer
  - Data: payload
- **A network protocol N1 can use the services of another network protocol N2**
  - A packet p1 of N1 is encapsulated into a packet p2 of N2
  - The payload of p2 is p1
  - The control information of p2 is derived from that of p1



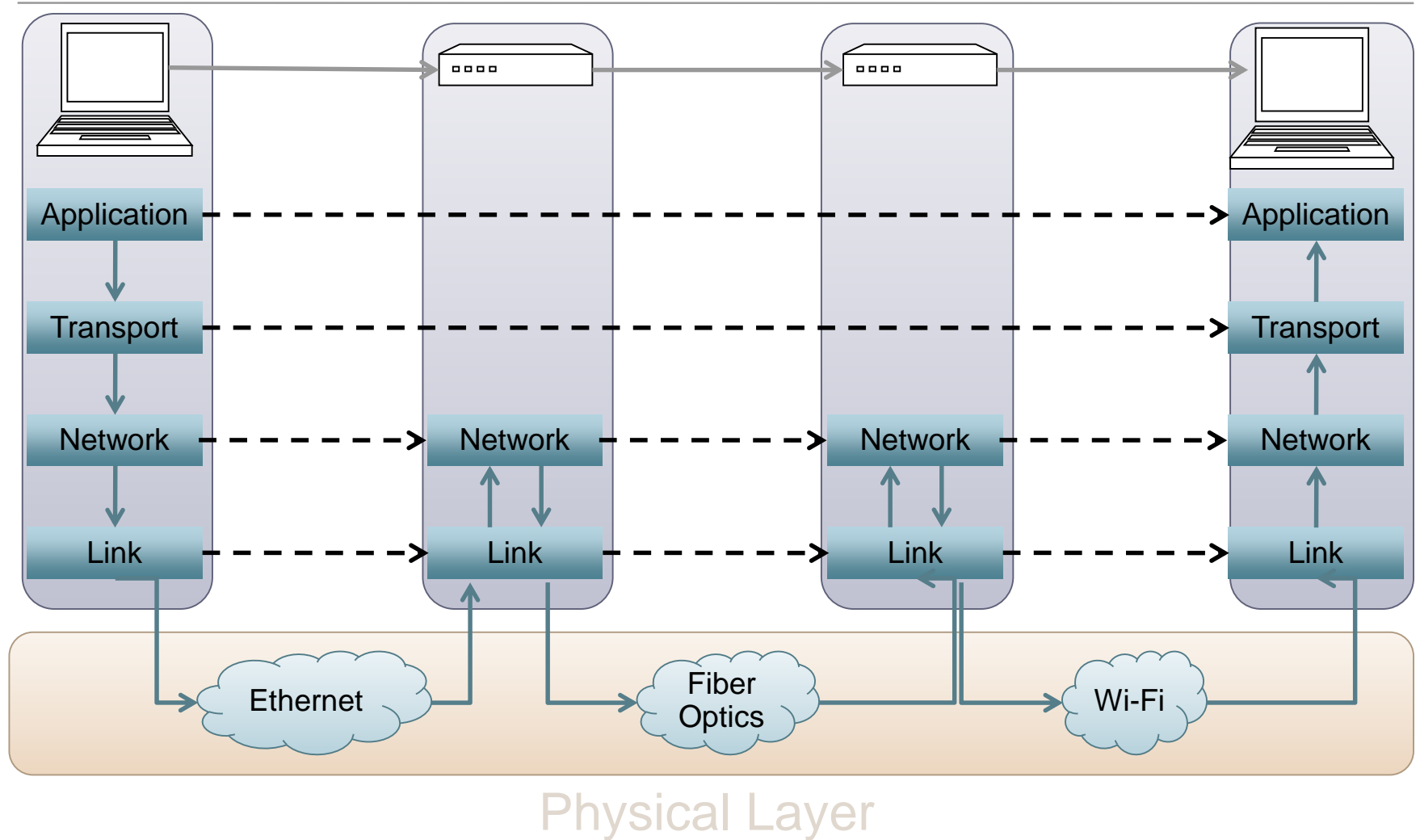


# Network Layers

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- **Network models typically use a stack of layers**
  - Higher layers use the services of lower layers via encapsulation
  - A layer can be implemented in hardware or software
  - The bottommost layer must be in hardware
- **A network device may implement several layers**
- **A communication channel between two nodes is established for each layer**
  - Actual channel at the bottom layer
  - Virtual channel at higher layers

# Internet Layers



# Intermediate Layers

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- **Link layer**

- Local area network: Ethernet, WiFi, optical fiber
- 48-bit media access control (MAC) addresses
- Packets called frames

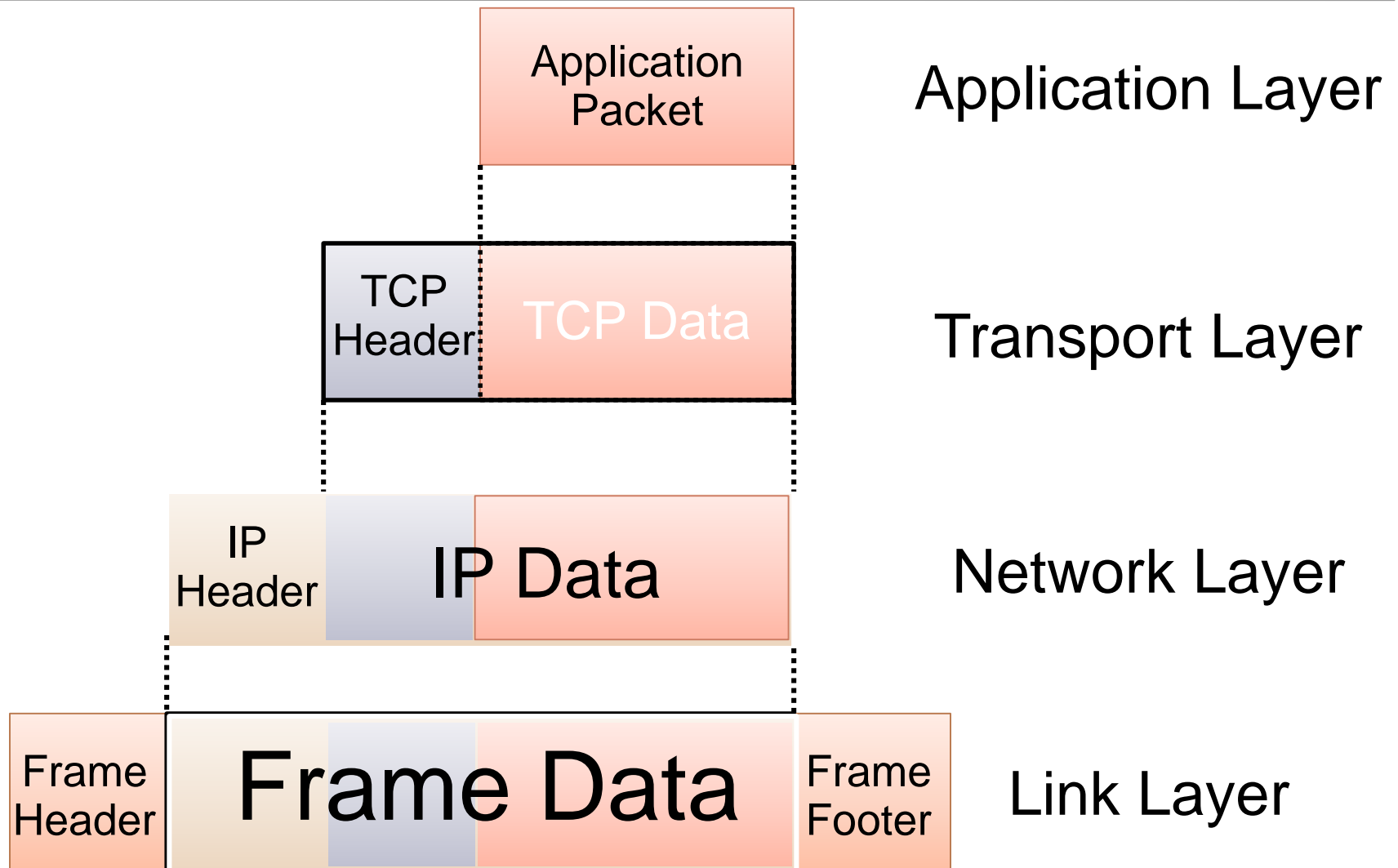
- **Network layer**

- Internet-wide communication
- Best efforts
- 32-bit internet protocol (IP) addresses in IPv4
- 128-bit IP addresses in IPv6

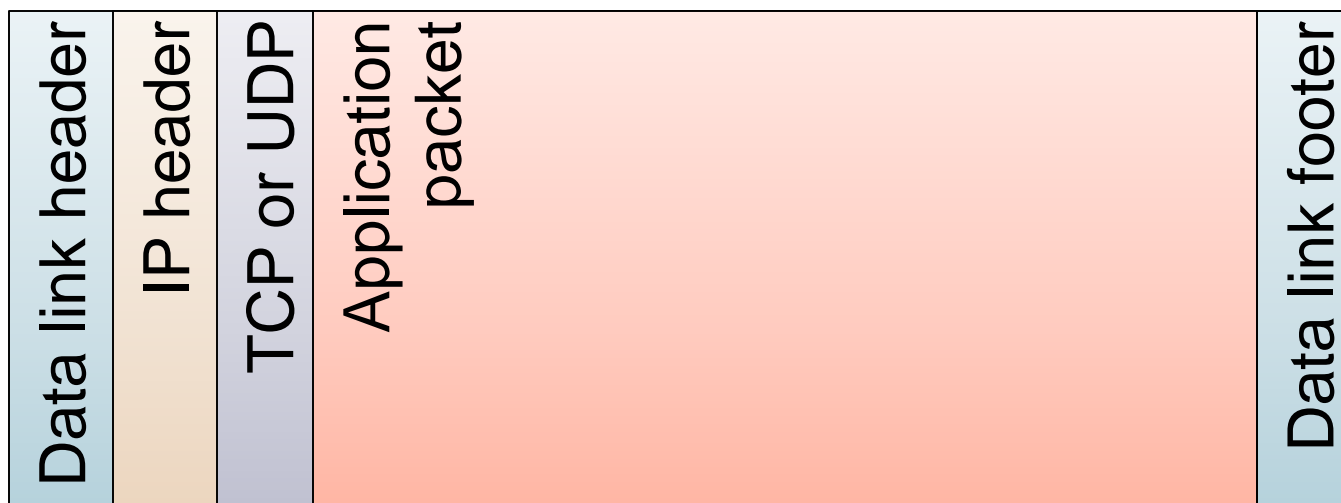
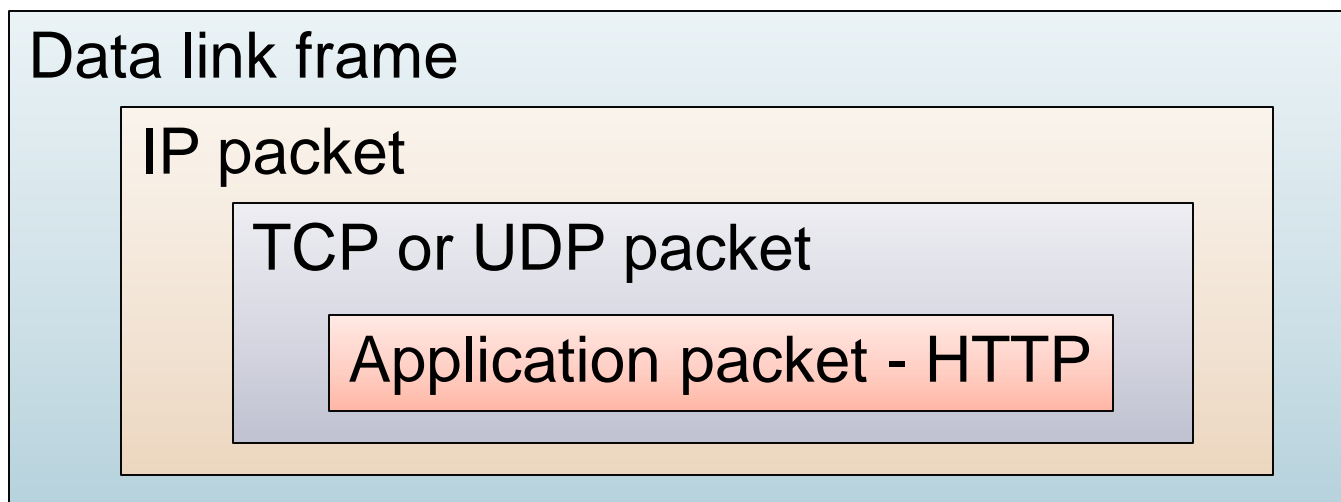
- **Transport layer**

- 16-bit addresses (ports) for classes of applications
- Connection-oriented transmission layer protocol (TCP)
- Connectionless user datagram protocol (UDP)

# Internet Packet Encapsulation

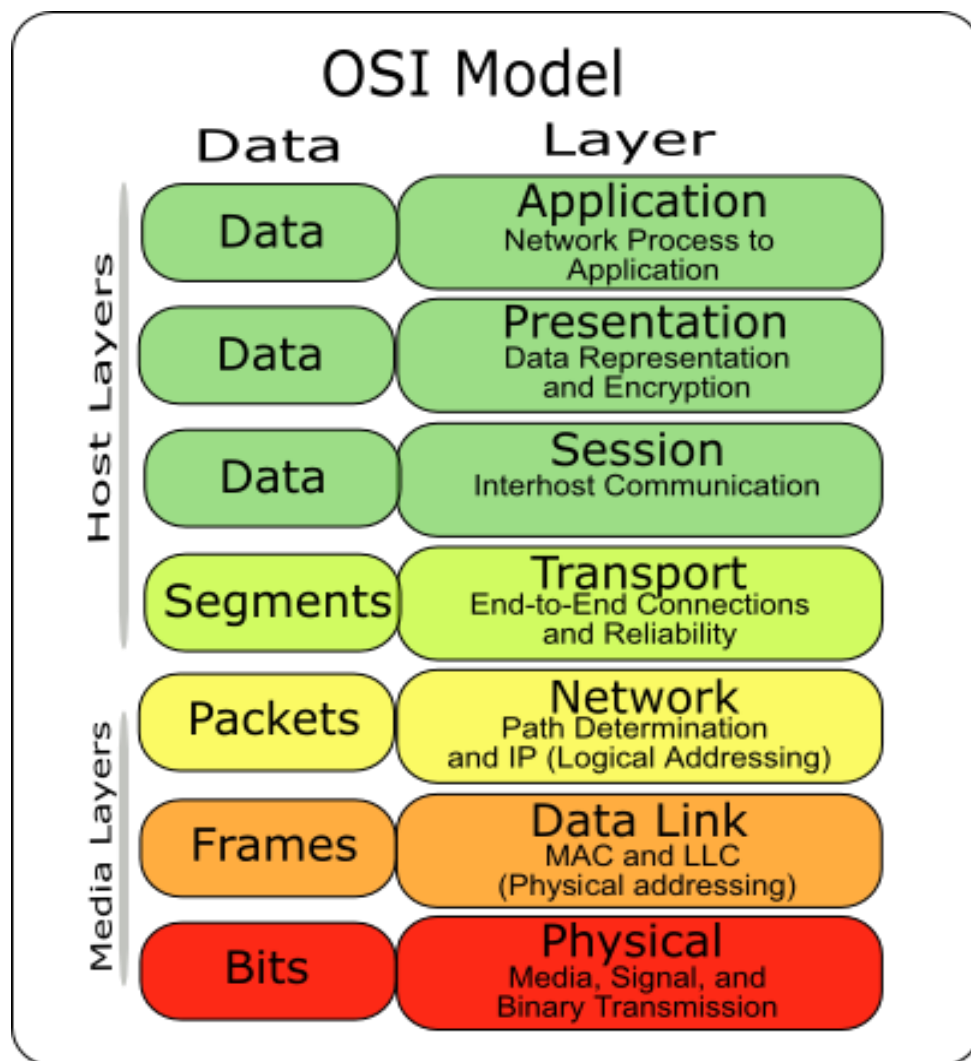


# Internet Packet Encapsulation



# The OSI Model

- **The OSI (Open System Interconnect) Reference Model is a network model consisting of seven layers**
- **Created in 1983, OSI is promoted by the International Standard Organization (ISO)**



# Network Interfaces

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- **Network interface: device connecting a computer to a network**
  - Ethernet card
  - WiFi adapter
- **A computer may have multiple network interfaces**
- **Packets transmitted between network interfaces**
- **Most local area networks, (including Ethernet and WiFi) broadcast frames**
- **In regular mode, each network interface gets the frames intended for it**
- **Traffic sniffing can be accomplished by configuring the network interface to read all frames (promiscuous mode)**

# MAC Addresses

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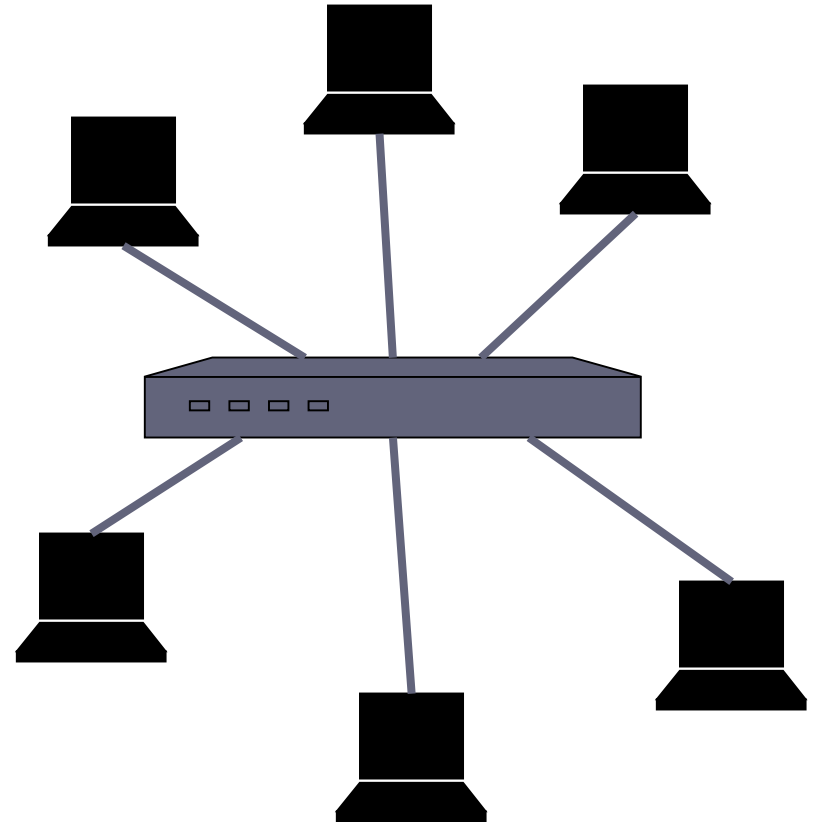
- **Most network interfaces come with a predefined MAC address**
- **A MAC address is a 48-bit number usually represented in hex**
  - E.g., 00-1A-92-D4-BF-86
- **The first three octets of any MAC address are IEEE-assigned Organizationally Unique Identifiers**
  - E.g., Cisco 00-1A-A1, D-Link 00-1B-11, ASUSTek 00-1A-92
- **The next three can be assigned by organizations as they please, with uniqueness being the only constraint**
- **Organizations can utilize MAC addresses to identify computers on their network**
- **MAC address can be reconfigured by network interface driver software**



# Switch

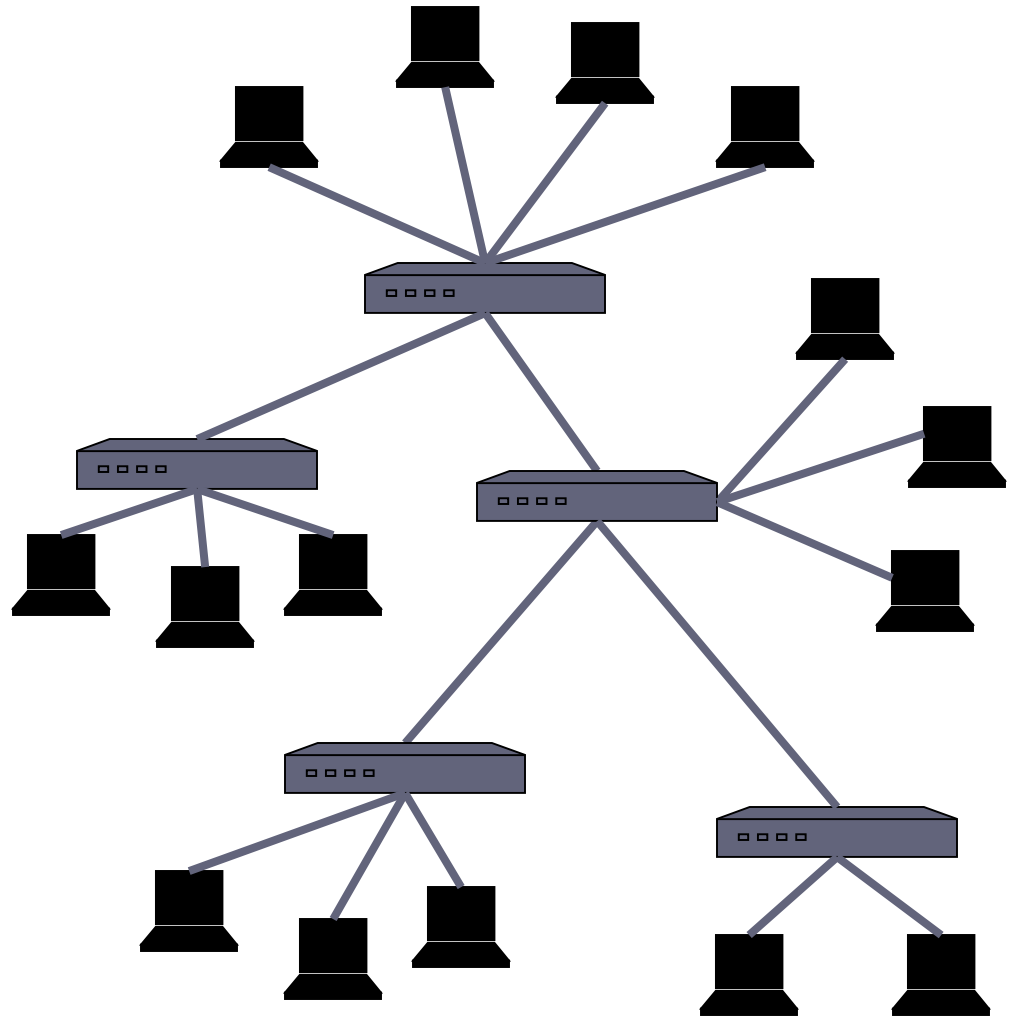
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- **A switch is a common network device**
  - Operates at the link layer
  - Has multiple ports, each connected to a computer
- **Operation of a switch**
  - Learn the MAC address of each computer connected to it
  - Forward frames only to the destination computer



# Combining Switches

- Switches can be arranged into a **tree**
- Each port learns the MAC addresses of the machines in the segment (subtree) connected to it
- Fragments to unknown MAC addresses are broadcast
- Frames to MAC addresses in the same segment as the sender are ignored



# MAC Address Filtering

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- **A switch can be configured to provide service only to machines with specific MAC addresses**
- **Allowed MAC addresses need to be registered with a network administrator**
- **A MAC spoofing attack impersonates another machine**
  - Find out MAC address of target machine
  - Reconfigure MAC address of rogue machine
  - Turn off or unplug target machine
- **Countermeasures**
  - Block port of switch when machine is turned off or unplugged
  - Disable duplicate MAC addresses

# Viewing and Changing MAC Addresses

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- **Viewing the MAC addresses of the interfaces of a machine**
  - Linux: `ifconfig`
  - Windows: `ipconfig /all`
- **Changing a MAC address in Linux**
  - Stop the networking service: `/etc/init.d/networking stop`
  - Change the MAC address: `ifconfig eth0 hw ether <MAC-address>`
  - Start the networking service: `/etc/init.d/networking start`
- **Changing a MAC address in Windows**
  - Open the Network Connections applet
  - Access the properties for the network interface
  - Click “Configure ...”
  - In the advanced tab, change the network address to the desired value
- **Changing a MAC address requires administrator privileges**

# ARP

- The address resolution protocol (ARP) connects the network layer to the data layer by converting IP addresses to MAC addresses
- ARP works by broadcasting requests and caching responses for future use
- The protocol begins with a computer broadcasting a message of the form  
who has <IP address1> tell <IP address2>
- When the machine with <IP address1> or an ARP server receives this message, it broadcasts the response  
<IP address1> is <MAC address>
- The requestor's IP address <IP address2> is contained in the link header
- The Linux and Windows command `arp - a` displays the ARP table

Internet Address	Physical Address	Type
128.148.31.1	00-00-0c-07-ac-00	dynamic
128.148.31.15	00-0c-76-b2-d7-1d	dynamic
128.148.31.71	00-0c-76-b2-d0-d2	dynamic
128.148.31.75	00-0c-76-b2-d7-1d	dynamic
128.148.31.102	00-22-0c-a3-e4-00	dynamic
128.148.31.137	00-1d-92-b6-f1-a9	dynamic

# ARP Spoofing

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- **The ARP table is updated whenever an ARP response is received**
- **Requests are not tracked**
- **ARP announcements are not authenticated**
- **Machines trust each other**
- **A rogue machine can spoof other machines**

# ARP Poisoning (ARP Spoofing)

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- **According to the standard, almost all ARP implementations are stateless**
- **An arp cache updates every time that it receives an arp reply... even if it did not send any arp request!**
- **It is possible to “poison” an arp cache by sending gratuitous arp replies**
- **Using static entries solves the problem but it is almost impossible to manage!**

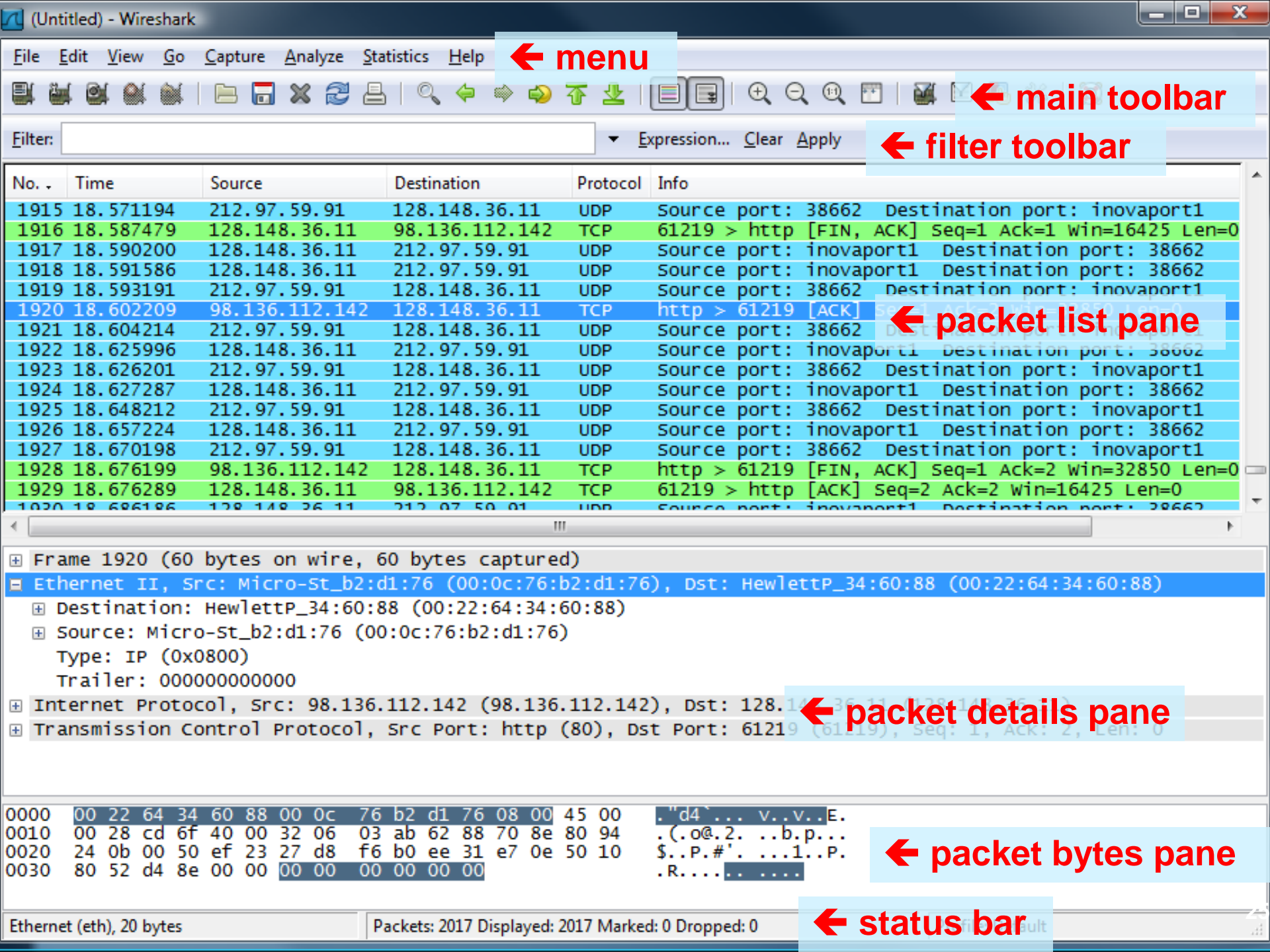


# Wireshark

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- **Wireshark is a packet sniffer and protocol analyzer**
  - Captures and analyzes frames
  - Supports plugins
- **Usually required to run with administrator privileges**
- **Setting the network interface in promiscuous mode captures traffic across the entire LAN segment and not just frames addressed to the machine**
- **Freely available on [www.wireshark.org](http://www.wireshark.org)**





← menu

← main toolbar

← filter toolbar

No. -	Time	Source	Destination	Protocol	Info
1915	18.571194	212.97.59.91	128.148.36.11	UDP	Source port: 38662 Destination port: inovaport1
1916	18.587479	128.148.36.11	98.136.112.142	TCP	61219 > http [FIN, ACK] Seq=1 Ack=1 win=16425 Len=0
1917	18.590200	128.148.36.11	212.97.59.91	UDP	Source port: inovaport1 Destination port: 38662
1918	18.591586	128.148.36.11	212.97.59.91	UDP	Source port: inovaport1 Destination port: 38662
1919	18.593191	212.97.59.91	128.148.36.11	UDP	Source port: 38662 Destination port: inovaport1
1920	18.602209	98.136.112.142	128.148.36.11	TCP	http > 61219 [ACK] Seq=1 Ack=2 win=32850 Len=0
1921	18.604214	212.97.59.91	128.148.36.11	UDP	Source port: 38662 Destination port: inovaport1
1922	18.625996	128.148.36.11	212.97.59.91	UDP	Source port: inovaport1 Destination port: 38662
1923	18.626201	212.97.59.91	128.148.36.11	UDP	Source port: 38662 Destination port: inovaport1
1924	18.627287	128.148.36.11	212.97.59.91	UDP	Source port: inovaport1 Destination port: 38662
1925	18.648212	212.97.59.91	128.148.36.11	UDP	Source port: 38662 Destination port: inovaport1
1926	18.657224	128.148.36.11	212.97.59.91	UDP	Source port: inovaport1 Destination port: 38662
1927	18.670198	212.97.59.91	128.148.36.11	UDP	Source port: 38662 Destination port: inovaport1
1928	18.676199	98.136.112.142	128.148.36.11	TCP	http > 61219 [FIN, ACK] Seq=1 Ack=2 win=32850 Len=0
1929	18.676289	128.148.36.11	98.136.112.142	TCP	61219 > http [ACK] Seq=2 Ack=2 win=16425 Len=0
1930	18.686186	128.148.36.11	212.97.59.91	UDP	Source port: inovaport1 Destination port: 38662

← packet list pane

+ Frame 1920 (60 bytes on wire, 60 bytes captured)  
- Ethernet II, Src: Micro-St\_b2:d1:76 (00:0c:76:b2:d1:76), Dst: HewlettP\_34:60:88 (00:22:64:34:60:88)  
+ Destination: HewlettP\_34:60:88 (00:22:64:34:60:88)  
+ Source: Micro-St\_b2:d1:76 (00:0c:76:b2:d1:76)  
Type: IP (0x0800)  
Trailer: 000000000000  
+ Internet Protocol, Src: 98.136.112.142 (98.136.112.142), Dst: 128.148.36.11 (128.148.36.11)  
+ Transmission Control Protocol, Src Port: http (80), Dst Port: 61219 (61219), Seq: 1, ACK: 2, Len: 0

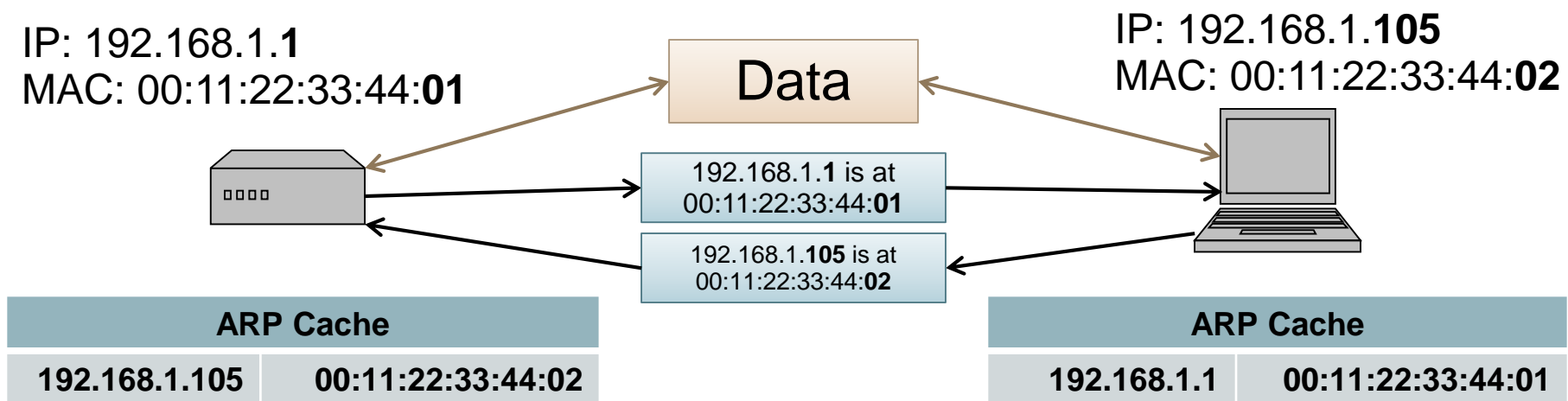
← packet details pane

0000 00 22 64 34 60 88 00 0c 76 b2 d1 76 08 00 45 00 . "d4`... v..v...E.  
0010 00 28 cd 6f 40 00 32 06 03 ab 62 88 70 8e 80 94 . (.o@.2. ..b.p...  
0020 24 0b 00 50 ef 23 27 d8 f6 b0 ee 31 e7 0e 50 10 \$. .P.#'. ...1..P.  
0030 80 52 d4 8e 00 00 00 00 00 00 00 00 00 00 .R.... .....

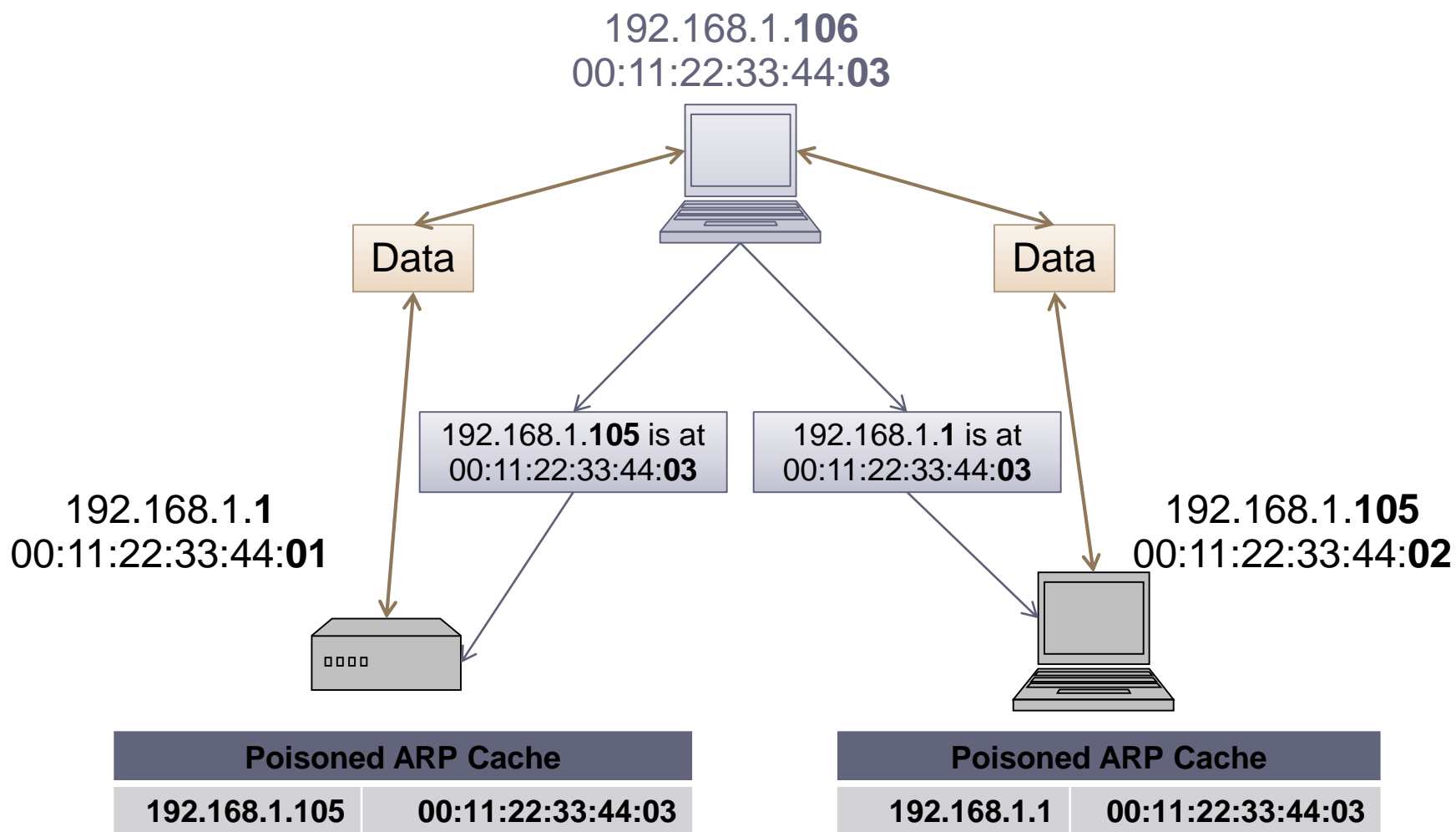
← packet bytes pane

← status bar

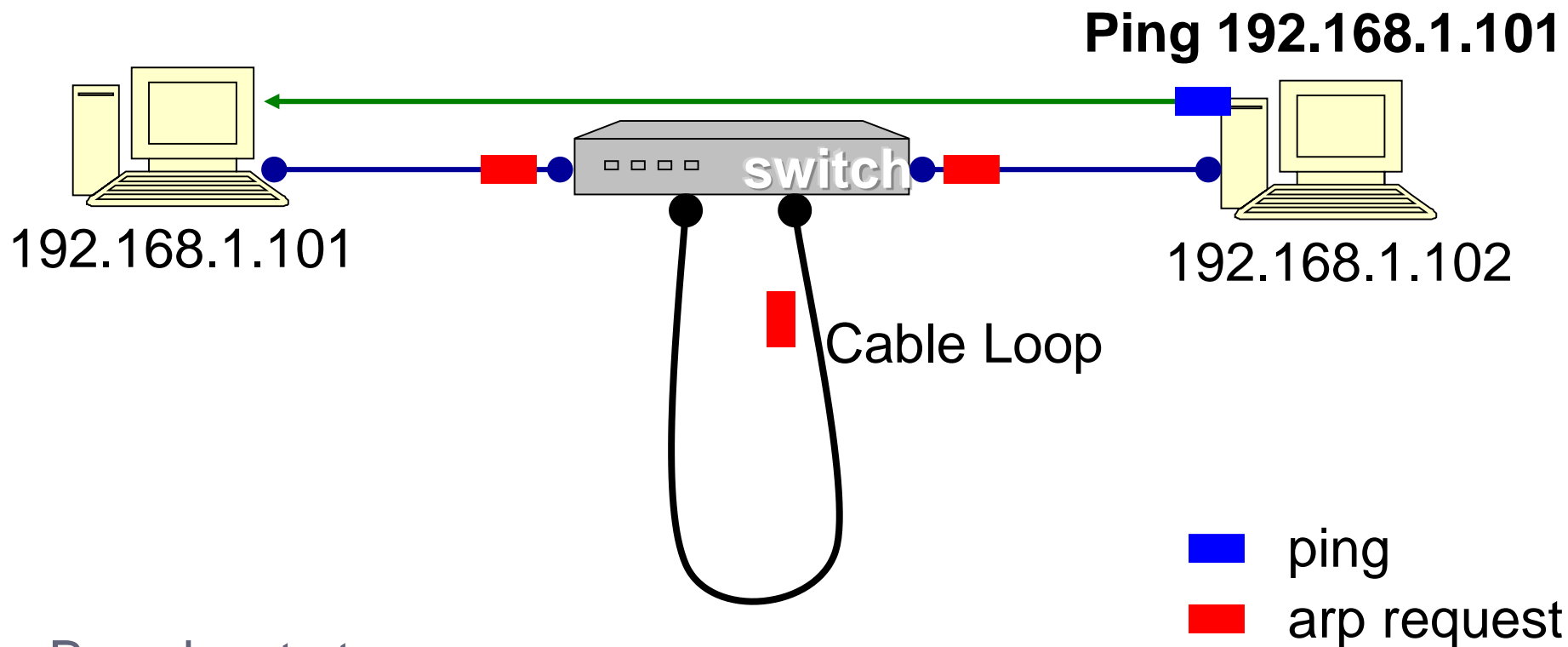
# ARP Caches



# Poisoned ARP Caches



# Network DoS using ARP



Broadcast storm

How can it be avoided?