

Asynchronous Fault Diagnosis of Discrete-Event Systems With Partially Observable Outputs

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APPENDIX

The proofs of the following lemmas are omitted because they follow directly from the definitions of $\hat{D}_{oc}^{l_o}$ and $\hat{D}_{dc}^{l_o}$.

Lemma 1. For any path $p = q_{-1} \xrightarrow{l_o^0} q_0 \xrightarrow{l_o^1} q_1 \cdots q_{k-1} \xrightarrow{l_o^k} q_k \cdots q_n \xrightarrow{l_o^{n+1}} q_k$ in $\hat{D}_{oc}^{l_o}$ ending with a cycle, The following hold:

- 1) $\kappa(q_j) = \kappa(q_r)$ for any $j, r \in [k, n]$;
- 2) There exists a path $\bar{p} = \bar{q}_0 \xrightarrow{\sigma_1} \bar{q}_1 \xrightarrow{\sigma_2} \bar{q}_2 \cdots \bar{q}_{b-1} \xrightarrow{\sigma_b} \bar{q}_b \cdots \bar{q}_c \xrightarrow{\sigma_c} \bar{q}_b$ in \mathbf{G} such that $O(\lambda(\bar{q}_0 \bar{q}_1 \cdots \bar{q}_c \bar{q}_b)) = l_o^0 l_o^1 \cdots l_o^n l_o^k$ and $Q_p \subseteq Q_{\bar{p}}$, where Q_p and $Q_{\bar{p}}$ are the set of states composed of paths p and \bar{p} , respectively.

In a cycle, all states correspond to the same condition, which is either normal or F_i ($i \in [1, m]$), owing to the assumption that all failure modes are permanent. We know that $\hat{D}_{oc}^{l_o}$ is defined on the basis of $\hat{D}_{dm}^{l_o}$ and $\hat{D}_{dm}^{l_o}$ derives from \mathbf{G} . Thus, there exists a transition path \bar{p} in \mathbf{G} corresponding to the indication path p in $\hat{D}_{oc}^{l_o}$.

Lemma 2. For any path $p = (q_{-1}, q_{-1}) \xrightarrow{l_o^0} (q_0^1, q_0^2) \xrightarrow{l_o^1} (q_1^1, q_1^2) \cdots (q_{k-1}^1, q_{k-1}^2) \xrightarrow{l_o^k} (q_k^1, q_k^2) \cdots (q_n^1, q_n^2) \xrightarrow{l_o^{n+1}} (q_k^1, q_k^2)$ ($k < n$) in $\hat{D}_{dc}^{l_o}$ ending with a cycle, the following hold:

- 1) There exist a path p_{noc}^1 in \hat{D}_{noc} and a path p_{foc}^2 in \hat{D}_{foc} ending with cycles, namely, $p_{noc}^1 = q_{-1} \xrightarrow{l_o^0} q_0^1 \xrightarrow{l_o^1} q_1^1 \cdots q_{k-1}^1 \xrightarrow{l_o^k} q_k^1 \cdots q_n^1 \xrightarrow{l_o^{n+1}} q_k^1$ and $p_{foc}^2 = q_{-1} \xrightarrow{l_o^0} q_0^2 \xrightarrow{l_o^1} q_1^2 \cdots q_{k-1}^2 \xrightarrow{l_o^k} q_k^2 \cdots q_n^2 \xrightarrow{l_o^{n+1}} q_k^2$.
- 2) $\kappa(q_j^1) = \kappa(q_r^1) = N$ and $\kappa(q_j^2) = \kappa(q_r^2)$ for any $j, r \in [k, n]$.

Based on Lemmas 1 and 2, the proof of Theorem 1 is shown as follows.

Proof. (only if): Suppose \mathbf{G} is F_i -asynchronously diagnosable w.r.t. $l_o^0 \in \Lambda_o$, but there exists a cycle cl in $\hat{D}_{dc}^{l_o}$, $cl = (q_k^1, q_k^2) \xrightarrow{l_o^{k+1}} (q_{k+1}^1, q_{k+1}^2) \cdots (q_n^1, q_n^2) \xrightarrow{l_o^{n+1}} (q_k^1, q_k^2)$ ($k < n$) such that $\kappa(q_j^2) = F_i$ ($i \in [1, m], j \in [k, n]$).

Since $\hat{D}_{dc}^{l_o}$ is reachable, there exists a path p in $\hat{D}_{dc}^{l_o}$ ending with the cycle cl , i.e., $p = (q_{-1}, q_{-1}) \xrightarrow{l_o^0} (q_0^1, q_0^2) \xrightarrow{l_o^1} (q_1^1, q_1^2) \cdots (q_{k-1}^1, q_{k-1}^2) \xrightarrow{l_o^k} (q_k^1, q_k^2) \cdots (q_n^1, q_n^2) \xrightarrow{l_o^{n+1}} (q_k^1, q_k^2)$

(q_k^1, q_k^2). Based on Lemma 2, we know that there exist one path p_{noc}^1 in \hat{D}_{noc} and one path p_{foc}^2 in \hat{D}_{foc} ending with cycles, namely, $p_{noc}^1 = q_{-1} \xrightarrow{l_o^0} q_0^1 \xrightarrow{l_o^1} q_1^1 \cdots q_{k-1}^1 \xrightarrow{l_o^k} q_k^1 \cdots q_n^1 \xrightarrow{l_o^{n+1}} q_k^1$ and $p_{foc}^2 = q_{-1} \xrightarrow{l_o^0} q_0^2 \xrightarrow{l_o^1} q_1^2 \cdots q_{k-1}^2 \xrightarrow{l_o^k} q_k^2 \cdots q_n^2 \xrightarrow{l_o^{n+1}} q_k^2$ with $\kappa(q_j^1) = \kappa(q_r^1) = 0$ and $\kappa(q_j^2) = \kappa(q_r^2) = i$ for any $j, r \in [k, n]$. Further from Lemma 1, we know that for the path p_{foc}^2 there exists a path $\bar{p}^2 = \bar{q}_0^2 \xrightarrow{\sigma_1^2} \bar{q}_1^2 \xrightarrow{\sigma_2^2} \bar{q}_2^2 \cdots \bar{q}_{b-1}^2 \xrightarrow{\sigma_b^2} \bar{q}_b^2 \cdots \bar{q}_c^2 \xrightarrow{\sigma_c^2} \bar{q}_b^2$ in \mathbf{G} such that $O(\lambda(\bar{q}_0^2 \bar{q}_1^2 \cdots \bar{q}_c^2 \bar{q}_b^2)) = l_o^0 l_o^1 \cdots l_o^n l_o^k$ and $Q_{p_{foc}^2} \subseteq Q_{\bar{p}^2}$.

Let $\mathcal{Q}_e(q, n)$ denote the set of state estimations calculated after the occurrence of n events from state q . Suppose the state estimation $x_{0e} \in \mathcal{Q}_e(\bar{q}_s^2, 0)$ ($s \in [1, b]$), where \bar{q}_s^2 ($\kappa(\bar{q}_s^2) = F_i$) is the first faulty state in the path \bar{p}^2 , is given. When the system evolves along the path \bar{p}^2 , there exists a state estimation $x_{e'}^{n'} \in \mathcal{Q}_e(\bar{q}_s^2, n')$ ($n' = d + m_2 * k_2$, $d = b - s$, and $m_2 = c - b + 1$ is the length of $\bar{q}_b^2 \cdots \bar{q}_c^2$) such that $q_k^1 \in x_{e'}^{n'}$ and $q_k^2 \in x_{e'}^{n'}$ for any nonnegative integer k_2 . Then $\mathbf{D}(x_{e'}^{n'}) = -1$. Since \mathbf{G} is asynchronously diagnosable w.r.t. l_o^0 , there exists an integer N_i such that for any $x_e^n \in \mathcal{Q}(q_s^2, n)$, $x_e^n \subseteq Q_{F_i}$ holds for $n \geq N_i$. We choose an integer k_2 such that $n' \geq N_i$. Then we have that $x_{e'}^{n'} \subseteq Q_{F_i}$, i.e., $\mathbf{D}(x_{e'}^{n'}) = i$, which leads to a contraction. So the necessity holds.

(if): Suppose for every cycle $cl = (q_k^1, q_k^2) \xrightarrow{l_o^{k+1}} (q_{k+1}^1, q_{k+1}^2) \cdots (q_n^1, q_n^2) \xrightarrow{l_o^{n+1}} (q_k^1, q_k^2)$ ($k < n$) in $\hat{D}_{dc}^{l_o}$, we have $\kappa(q_j^1) = \kappa(q_j^2) = N$ ($j \in [k, n]$). From the second clause of Lemma 2, we can infer that for any $q^d = (q^1, q^2)$ in $\hat{D}_{dc}^{l_o}$, q^d is not contained in a loop if $\kappa(q^1) \neq \kappa(q^2)$, which further implies that for any state sequence $q_1^d q_2^d \cdots q_k^d$ in $\hat{D}_{dc}^{l_o}$ with $q_r^d = (q_r^1, q_r^2)$ ($r \in [1, k]$) if $\kappa(q_r^1) \neq \kappa(q_r^2)$, then the length of this state sequence is finite, which is less than the number of states in $\hat{D}_{dc}^{l_o}$.

Suppose the state estimation $x_{0e} \in \mathcal{Q}_e(q_{F_i}, 0)$ ($i \in [1, m]$) is given when the system first reaches the fault state q_{F_i} ($\kappa(q_{F_i}) = F_i$). For any $x_e^n \in \mathcal{Q}_e(q_{F_i}, n)$ with $n > |\hat{Q}_d| \times (|\mathcal{Q}| - 1)$, we claim that $\mathbf{D}(x_e^n) = i$. After n transitions from state q_{F_i} , we have the state sequence $s = q_{F_i} q_{F_i}^1 \cdots q_{F_i}^n$ with the observed output sequence $O(\lambda(s))$ ($n' = |O(\lambda(s))|$). From above, for any state $q^d \in \hat{Q}_d$ that can be reached from (q_{-1}, q_{-1}) in $\hat{D}_{dc}^{l_o}$, we have that for any state sequence starting from q^d , a state $\hat{q}^d = (\hat{q}^1, \hat{q}^2) \in \hat{Q}_d$ with $\kappa(\hat{q}^1) = \kappa(\hat{q}^2)$ can be reached within $|\hat{Q}_d| - 1$ transitions. This implies that for any $n' > |\hat{Q}_d| - 1$ and $x_e^{n'} \in \mathcal{Q}_e(q_{F_i}, n')$, we have that $\mathbf{D}(x_e^{n'}) = i$. From the assumption in Remark 2, each observed output can be followed by at most $|\mathcal{Q}| - 1$ unobserved outputs. It follows that for the above state sequence s , $n \leq (n' + 1) \times (|\mathcal{Q}| - 1)$, i.e., $n' \geq n / (|\mathcal{Q}| - 1) - 1$. So if $n > |\hat{Q}_d| \times (|\mathcal{Q}| - 1)$,

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83 then $n' \geq n/(|Q| - 1) - 1 > |\hat{Q}_d| - 1$, establishing our
 84 claim. Note that we have assumed implicitly that $|Q| > 1$;
 85 otherwise if $|Q| = 1$, then from the assumption of no path
 86 cycles, no transition labeled by a failure event exists, so that
 87 the system is trivially diagnosable. Based on Definition 9, we
 88 can conclude that \mathbf{G} is diagnosable w.r.t, l_o^0 . So the sufficiency
 89 also holds. \square