

# Asynchronous Fault Diagnosis of Discrete-Event Systems With Partially Observable Outputs

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## APPENDIX

The proofs of the following lemmas are omitted because they follow directly from the definitions of  $\hat{\mathcal{D}}_{oc}$  and  $\hat{\mathcal{D}}_{dc}$ .

**Lemma 1.** *For any path  $p = \{(q_{-1}, l_0, q_0), (q_0, l_1, q_1), \dots, (q_{k-1}, l_k, q_k), \dots, (q_n, l_k, q_k)\}$  in  $\hat{\mathcal{D}}_{oc}$  ending with a cycle, The following hold:*

- 1)  $\kappa(q_j) = \kappa(q_r)$  for any  $j, r \in [k, n]$ ;
- 2) *There exists a path  $\bar{p} = \{(\bar{q}_0, \sigma_1, \bar{q}_1), (\bar{q}_1, \sigma_2, \bar{q}_2), \dots, (\bar{q}_{b-1}, \sigma_b, \bar{q}_b), \dots, (\bar{q}_c, \sigma_c, \bar{q}_b)\}$  in  $\mathbf{G}$  such that  $O(\lambda(\bar{q}_0\bar{q}_1 \dots \bar{q}_c\bar{q}_b)) = l_0l_1 \dots l_nl_k$  and  $Q_p \subseteq Q_{\bar{p}}$ , where  $Q_p$  and  $Q_{\bar{p}}$  are the set of states composed of paths  $p$  and  $\bar{p}$ , respectively.*

In a cycle, all states correspond to the same condition, which is either normal or  $F_i$  ( $i \in [1, m]$ ), owing to the assumption that all failure modes are permanent. We know that  $\hat{\mathcal{D}}_{oc}$  is defined on the basis of  $\hat{\mathcal{D}}_{ds}$  and  $\hat{\mathcal{D}}_{ds}$  derives from  $\mathbf{G}$ . Thus, there exists a transition path  $\bar{p}$  in  $\mathbf{G}$  corresponding to the indication path  $p$  in  $\hat{\mathcal{D}}_{oc}$ .

**Lemma 2.** *For any path  $p = \{((q_{-1}, q_{-1}), l_0, (q_0^1, q_0^2)), ((q_0^1, q_0^2), l_1, (q_1^1, q_1^2)), \dots, ((q_{k-1}^1, q_{k-1}^2), l_k, (q_k^1, q_k^2)), \dots, ((q_n^1, q_n^2), l_k, (q_k^1, q_k^2))\}$  ( $k < n$ ) in  $\hat{\mathcal{D}}_{dc}$  ending with a cycle, the following hold:*

- 1) *There exist a path  $p_{noc}^1$  in  $\hat{\mathcal{D}}_{noc}$  and a path  $p_{foc}^2$  in  $\hat{\mathcal{D}}_{foc}$  ending with cycles, namely  $p_{noc}^1 = \{(q_{-1}, l_0, q_0^1), (q_0^1, l_1, q_1^1), \dots, (q_{k-1}^1, l_k, q_k^1), \dots, (q_n^1, l_k, q_k^1)\}$  and  $p_{foc}^2 = \{(q_{-1}, l_0, q_0^2), (q_0^2, l_1, q_1^2), \dots, (q_{k-1}^2, l_k, q_k^2), \dots, (q_n^2, l_k, q_k^2)\}$ ;*
- 2)  $\kappa(q_j^1) = \kappa(q_r^1) = 0$  and  $\kappa(q_j^2) = \kappa(q_r^2)$  for any  $j, r \in [k, n]$ .

◇

Based on Lemmas 1 and 2, the proof of Theorem 1 is shown as follows.

*Proof.* (only if): Suppose  $\mathbf{G}$  is asynchronously diagnosable w.r.t.  $l_0 \in \Lambda_o$ , but there exists a cycle  $cl$  in  $\hat{\mathcal{D}}_{dc}$ ,  $cl = \{((q_k^1, q_k^2), l_{k+1}, (q_{k+1}^1, q_{k+1}^2)), \dots, ((q_n^1, q_n^2), l_k, ((q_k^1, q_k^2)))\}$  ( $k < n$ ) such that  $\kappa(q_j^2) = i \in [1, m]$  ( $j \in [k, n]$ ).

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Since  $\hat{\mathcal{D}}_{dc}$  is reachable, there exists a path  $p$  in  $\hat{\mathcal{D}}_{dc}$  ending with the cycle  $cl$ , i.e.,  $p = \{((q_{-1}, q_{-1}), l_0, (q_0^1, q_0^2)), ((q_0^1, q_0^2), l_1, (q_1^1, q_1^2)), \dots, ((q_{k-1}^1, q_{k-1}^2), l_k, (q_k^1, q_k^2)), \dots, ((q_n^1, q_n^2), l_k, (q_k^1, q_k^2))\}$ . Based on Lemma 2, we know that there exist one path  $p_{noc}^1$  in  $\hat{\mathcal{D}}_{noc}$  and one path  $p_{foc}^2$  in  $\hat{\mathcal{D}}_{foc}$  ending with cycles, namely  $p_{noc}^1 = \{(q_{-1}, l_0, q_0^1), (q_0^1, l_1, q_1^1), \dots, (q_{k-1}^1, l_k, q_k^1), \dots, (q_n^1, l_k, q_k^1)\}$  and  $p_{foc}^2 = \{(q_{-1}, l_0, q_0^2), (q_0^2, l_1, q_1^2), \dots, (q_{k-1}^2, l_k, q_k^2), \dots, (q_n^2, l_k, q_k^2)\}$  and  $\kappa(q_j^1) = \kappa(q_r^1) = 0$  and  $\kappa(q_j^2) = \kappa(q_r^2) = i$  for any  $j, r \in [k, n]$ . Further from Lemma 1, we know that for the path  $p_{noc}^1$  there exists a path  $\bar{p}^2 = \{(\bar{q}_0^2, \sigma_1^2, \bar{q}_1^2), (\bar{q}_1^2, \sigma_2^2, \bar{q}_2^2), \dots, (\bar{q}_{b-1}^2, \sigma_b^2, \bar{q}_b^2), \dots, (\bar{q}_c^2, \sigma_c^2, \bar{q}_b^2)\}$  in  $\mathbf{G}$  such that  $O(\lambda(\bar{q}_0^2\bar{q}_1^2 \dots \bar{q}_c^2\bar{q}_b^2)) = l_0l_1 \dots l_nl_k$  and  $Q_{p_{foc}^2} \subseteq Q_{\bar{p}^2}$ .

Let  $\mathcal{Q}_e(q, n)$  denote the set of state estimations calculated after the occurrence of  $n$  events from state  $q$ . Suppose the state estimation  $x_{0e} \in \mathcal{Q}_e(\bar{q}_s^2, 0)$  ( $s \in [1, b]$ ), where  $\bar{q}_s^2$  ( $\kappa(\bar{q}_s^2) = i$ ) is the first fault state in the path  $\bar{p}^2$ , is given. When the system evolves along the path  $\bar{p}^2$ , there exists a state estimation  $x_e^{n'} \in \mathcal{Q}_e(\bar{q}_s^2, n')$  ( $n' = d + m_2 * k_2$ ,  $d = b - s$ , and  $m_2 = c - b + 1$  is the length of  $\bar{q}_b^1 \dots \bar{q}_c^1$ ) such that  $q_k^1 \in x_e^{n'}$  and  $q_k^2 \in x_e^{n'}$  for any nonnegative integer  $k_2$ . Then  $\mathbf{D}(x_e^{n'}) = -1$ . Since  $\mathbf{G}$  is asynchronously diagnosable w.r.t.  $l_0$ , there exists an integer  $N_i$  such that for any  $x_e^n \in \mathcal{Q}_e(q_s^2, n)$ ,  $x_e^n \subseteq Q_{F_i}$  holds for  $n \geq N_i$ . We choose an integer  $k_2$  such that  $n' \geq N_i$ . Then we have that  $x_e^{n'} \subseteq Q_{F_i}$ , i.e.,  $\mathbf{D}(x_e^{n'}) = i$ , which leads to a contraction. So the necessity holds.

(if): Suppose for every cycle  $cl = \{((q_k^1, q_k^2), l_{k+1}, (q_{k+1}^1, q_{k+1}^2)), \dots, ((q_n^1, q_n^2), l_k, ((q_k^1, q_k^2)))\}$  ( $k < n$ ) in  $\hat{\mathcal{D}}_{dc}$ , we have  $\kappa(q_j^1) = \kappa(q_j^2) = 0$  ( $j \in [k, n]$ ). From the second clause of Lemma 2, we can infer that for any  $q^d = (q^1, q^2)$  in  $\hat{\mathcal{D}}_{dc}$ ,  $q^d$  is not contained in a loop if  $\kappa(q^1) \neq \kappa(q^2)$ , which further implies that for any state sequence  $(q^d, q^d, \dots, q^d)$  in  $\hat{\mathcal{D}}_{dc}$  with  $q_r^d = (q_r^1, q_r^2)$  ( $r \in [1, k]$ ) if  $\kappa(q_r^1) \neq \kappa(q_r^2)$ , then the length of this state sequence is finite, which is less than the number of states in  $\hat{\mathcal{D}}_{dc}$ .

Suppose the state estimation  $x_{0e} \in \mathcal{Q}_e(q_{F_i}, 0)$  ( $i \in [1, m]$ ) is given when the system first reach the fault state  $q_{F_i}$  ( $\kappa(q_{F_i}) = i$ ). For any  $x_e^n \in \mathcal{Q}_e(q_{F_i}, n)$  with  $n > |\hat{Q}_d| \times (|\hat{Q}| - 1)$ , we claim that  $\mathbf{D}(x_e^n) = i$ . After  $n$  transitions from state  $q_{F_i}$ , we have the state sequence  $s = q_{F_i}q_{F_i}^1 \dots q_{F_i}^n$  with the observed output sequence  $O(\lambda(s))$  ( $n' = |O(\lambda(s))|$ ). From above, for any state  $q^d \in \hat{Q}_d$  that can be reached from  $(q_{-1}, q_{-1})$  in  $\hat{\mathcal{D}}_{dc}$ , we have that for any state sequence starting from  $q^d$ , a state  $\hat{q}^d = (\hat{q}^1, \hat{q}^2) \in \hat{Q}_d$  with  $\kappa(\hat{q}^1) = \kappa(\hat{q}^2)$  can be reached within  $|\hat{Q}_d| - 1$  transitions. This implies that for any  $n' > |\hat{Q}_d| - 1$  and  $x_e^{n'} \in \mathcal{Q}_e(q_{F_i}, n')$ , we have that  $\mathbf{D}(x_e^{n'}) = i$ . From the assumption in Remark 1, each observed output can be followed by at most  $|\hat{Q}| - 1$  unobserved outputs. It follows that for the above state sequence  $s$ ,  $n \leq (n' + 1) \times (|\hat{Q}| - 1)$ ,

86 i.e.,  $n' \geq n/(|Q| - 1) - 1$ . So if  $n > |\hat{Q}_d| \times (|Q| - 1)$ ,  
 87 then  $n' \geq n/(|Q| - 1) - 1 > |\hat{Q}_d| - 1$ , establishing our  
 88 claim. Note that we have assumed implicitly that  $|Q| > 1$ ;  
 89 otherwise if  $|Q| = 1$ , then from the assumption of no path  
 90 cycles, no transition labeled by a failure event exists, so that  
 91 the system is trivially diagnosable. Based on Definition 9, we  
 92 can conclude that  $\mathbf{G}$  is diagnosable w.r.t,  $l_0$ . So the sufficiency  
 93 also holds.  $\square$