

Assignment 02

1 $f(n) = n^2 + f(n-1) + n$
 $f(1) = 2$

$O(n^3)$ because there are 3 for loops nested within each other.

2 $n^2 + 3n^3 = f(n)$ $c = 3$ $3n^3 = c \cdot g(n)$

$$3n^3 + n^2 \geq 3n^3 \quad \forall n \geq 0$$

$$n^2 \geq 0$$

$$\therefore f(n) \geq c \cdot g(n)$$

$\therefore c \cdot g(n)$ is a lower bound on $f(n)$

$$\therefore f(n) \in \Omega(n^3)$$

$n \geq 0$ because

$$3n^3 + n^2 = 0$$

$$\hookrightarrow n^2 = 0$$

$$3n + 1 = 0$$

$$\hookrightarrow n = -\frac{1}{3}$$

$$n^2 + 3n^3 = f(n) \quad g(n) = n^3 \quad c = 4$$

$$3n^3 + n^2 \leq 4n^3$$

$$\frac{3n^3 + n^2}{n^3} \leq 4$$

$$3 + \frac{1}{n} \leq 4$$

$$\frac{1}{n} \leq 1 \quad \forall n$$

$$\therefore f(n) \leq c \cdot g(n) \quad \forall n \geq 0$$

$\therefore c \cdot g(n)$ is an upper bound on $f(n)$

$$\therefore f(n) \in O(n^3)$$

$$\therefore f(n) \in \Theta(n^3)$$

3 For $2^{n+1} + 2^n$, every integer constant yields

$$c_1 2^n \geq 2^{n+1} \quad \text{which satisfies } O(2^n)$$

if $c_2 \leq 2$, then $c_2 2^n \leq 2^{n+1}$ satisfying $\Omega(2^n)$
($c_2 2^n = 2^{n+1}$ if $c_2 \in \mathbb{N}$)

So, to generalize:

$a^{n+1} \in \Theta(a^n)$ because

$$\begin{aligned} b \cdot a^n &\geq a^{n+1} \quad \text{iff} \quad b \geq a \rightarrow O(a^n) \\ c \cdot a^n &\leq a^{n+1} \quad \text{iff} \quad c \leq a \rightarrow \Omega(a^n) \end{aligned}$$

$$b \frac{a^n}{a^n} \geq \frac{a^n}{a^n} \cdot a$$

$$c \frac{a^n}{a^n} \leq \frac{a^n}{a^n} \cdot a$$

Because of this, $a^{n+1} \in \Theta(a^n)$

4

n^2

For each node in the graph, we check for an edge between it and all other nodes \leftarrow itself. Worst case, all nodes are checked and the graph is complete.

5

See excel workbooks in repo

5. The recorded timings indicate that the `.sort()` function in `<algorithm>` of the C++ standard library has $O(n \lg(n))$. This is because $n \lg(n)$ is an upper bound on the graph of the recorded timings of the sort method; for $f(n)$ = record timings and $g(n) = n \lg(n)$, for every value of n , $f(n) < g(n)$. Because of this, $f(n) < c_1 * g(n)$ for all $c_1 \geq 1$, proving the time complexity is $O(n \lg(n))$. This means that an estimation of the worst runtimes could be much closer to the values recorded for $n \lg(n)$, and my trial was much lower than that scenario.