Assignment 5

1. Arbitrage.py

```
def weight alias(x: float):
  return {'weight': -log(x)}
edges = [("USD", "GBP", weight_alias(0.75)),
    ("GBP", "AUD", weight alias(2)),
    ("AUD", "USD", weight_alias(0.7))]
def main():
 #Set up Graph
 G = nx.DiGraph()
 for currency in ["USD", "GBP", "AUD"]:
    G.add_node(currency)
  G.add_edges_from(edges)
                                                           Algorithm Begins Here:
 try:
                                                                O(V * E): V=# of vertices, #=# of edges
    path = nx.find_negative_cycle(G,"USD")
  except nx.exception.NetworkXError:
    print("There is no profitable exchange path for USD.")
                                                                  C1
  print("\nPath: " + str(path[1:]))
                                                                                  C1
  profit = nx.path_weight(G, path, 'weight')
                                                                                 O(V2)
  profit = pow(e,-profit)
                                                                                  C2
  #Round to 2 decimal places because USD
                                                                                  C3
  profit = round(profit, 2)
  print("$" + str(profit))
  return
main()
                                                             = O(V*E) + C1 + O(V2) + C2 + C3
                                                                          €O(n)
```

- 2. Prim's algorithm visits nodes multiple times to get the minimum distance/lowest cost/lowest weight whereas Kruskal's algorithm traverses each node once. Because of this behavior, Prim's algorithm is ideal for graphs with many nodes with smaller weights (dense graphs) and Kruskal's is ideal for graphs with greater weights and potentially less nodes (sparse graphs).
- 3. See prim_kruskal.py
- 4. The algorithm for making change operates on the assumption that we have an infinite amount of all coins. If we have 0 nickels, dimes, or quarters (or any other coins of greater value), the algorithm can still work by giving enough pennies equal to the amount covered by the missing coin(s). However, if we have 0 pennies the algorithm will fail. This is because n (the value to obtain) is found by adding one coin (d_i) to the amount $n d_i$ until the condition $n \ge d_i$ fails.

4.1. Suppose we need to make 9 cents but have no pennies. On the first iteration, we would add a nickel, and the remaining change is four cents. There is no coin available that would satisfy the condition 4 >= (coin value).