```
# Importing base classes

from priority_queue_base import PriorityQueueBase  # PriorityQueueBase defines Priority Queue ADT (abstract data type)
from positional_list import PositionalList  # PostionalList is an implementation of a doubly linked list
from exceptions import Empty
```

Unsorted Priority Queue

In our first concrete implementation of a priority queue, we store entries within an <code>unsorted list</code>. Our <code>UnsortedPriorityQueue</code> class is given in the following, inheriting from the <code>PriorityQueueBase</code> class introduced in <code>priority_queue_base.py</code>. For internal storage, key-value pairs are represented as composites, using instances of the inherited <code>Item</code> class. These items are stored within a <code>PositionalList</code>, identified asthe <code>__data</code> member of our class. We assume that the positional list is implemented with a doubly-linked list, as in <code>positional list.py</code>, so that all operations of that ADT execute in O(1) time.

We begin with an empty list when a new priority queue is constructed. At all times, the size of the list equals the number of key-value pairs currently stored in the priority queue. For this reason, our priority queue __len__ method simply returns the length of the internal data list. By the design of our PriorityQueueBase class, we inherit a concrete implementation of the is empty method that relies on a call to our __len__ method.

Each time a key-value pair is added to the priority queue, via the add method, we create a new __Item composite for the given key and value, and add that item to the end of the list. Such an implementation takes O(1) time.

The remaining challenge is that when min or remove min is called, we must locate the item with minimum key. Because the items are not sorted, we must inspect all entries to find one with a minimum key. For convenience, we define a nonpublic find min utility that returns the *position* of an item with minimum key. Knowledge of the position allows the remove min method to invoke the delete method on the positional list. The min method simply uses the position to retrieve the item when preparing a key-value tuple to return. Due to the loop for finding the minimum key, both min and remove min methods run in O(n) time, where n is the number of entries in the priority queue.

A summary of the running times for the UnsortedPriorityQueue class is given as:

Operation	Running Time	
len()	O(1)	
is_empty()	O(1)	
add(o)	O(1)	
o = min()	O(n)	
o = remove_min()	O(n)	
space	O(n)	

In [2]:

```
class UnsortedPriorityQueue (PriorityQueueBase): # base class defines Item
    """A min-oriented priority queue implemented with an unsorted list."""
                ----- nonpublic behavior -----
   def find min(self):
       """Return Position of item with minimum key."""
       if self.is empty():
                                                      # is empty inherited from base class
           raise Empty('Priority queue is empty')
       small = self._data.first()
       walk = self. data.after(small)
       while walk is not None:
           if walk.element() < small.element():</pre>
               small = walk
           walk = self. data.after(walk)
       return small
                     ----- public behaviors ------
        init (self):
       """Create a new empty Priority Queue."""
```

```
self. data = PositionalList()
    def len (self):
        """Return the number of items in the priority queue."""
        return len(self._data)
    def add(self, key, value):
        """Add a key-value pair."""
        self. data.add last(self. Item(key, value))
    def min(self):
        """Return but do not remove (k,v) tuple with minimum key.
        Raise Empty exception if empty.
        p = self._find_min()
        item = p.element()
        return (item._key, item._value)
    def remove min(self):
        """Remove and return (k,v) tuple with minimum key.
        Raise Empty exception if empty.
        p = self. find min()
        item = self._data.delete(p)
        return (item._key, item._value)
    # Custom function to display the entire elements
    def list elems(self):
        listed elem = []
        if self.is empty():
                                                           # is empty inherited from base class
            return listed elem
        walk = self. data.first()
        while walk is not None:
            item = walk.element()
            listed elem.append((item. key, item. value))
            walk = self._data.after(walk)
        return listed elem
In [3]:
# For testing purpose, assign a random names
import random
names = ["Alfa", "Bravo", "Charlie", "Delta", "Echo", "Foxtrot", "Golf", "Hotel", "India", "Juliett
", "Kilo", "Lima", "Mike", "November", "Oscar", "Papa", "Quebec", "Romeo", "Sierra", "Tango", "Unif orm", "Victor", "Whiskey", "X-ray", "Yankee", "Zulu"]
Q = UnsortedPriorityQueue()
for name in random.sample(names, 5):
   key = random.randint(0, 100)
   print(key, name)
    Q.add(key, name)
47 Victor
68 Kilo
25 Zulu
56 Yankee
39 Sierra
In [4]:
Q.list_elems()
Out[4]:
[(47, 'Victor'), (68, 'Kilo'), (25, 'Zulu'), (56, 'Yankee'), (39, 'Sierra')]
In [5]:
print(Q.remove min())
print(Q.remove min())
```

```
print(Q.remove_min())
print(Q.remove_min())
print(Q.remove_min())

(25, 'Zulu')
(39, 'Sierra')
(47, 'Victor')
(56, 'Yankee')
(68, 'Kilo')
```

Sorted Priority Queue

An alternative implementation of a priority queue uses a positional list, yet maintaining entries sorted by nondecreasing keys. This ensures that the first element of the list is an entry with the smallest key.

Our SortedPriorityQueue class is given in the following code. The implementation of min and remove_min are rather straightforward given knowledge that the first element of a list has a minimum key. We rely on the first method of the positional list to find the position of the first item, and the delete method to remove the entry from the list. Assuming that the list is implemented with a doubly linked list, operations min and remove min take O(1) time.

This benefit comes at a cost, however, for method add now requires that we scan the list to find the appropriate position to insert the new item. Our implementation starts at the end of the list, walking backward until the new key is smaller than an existing item; in the worst case, it progresses until reaching the front of the list. Therefore, the add method takes O(n) worst-case time, where n is the number of entries in the priority queue at the time the method is executed. In summary, when using a sorted list to implement a priority queue, insertion runs in linear time, whereas finding and removing the minimum can be done in constant time.

This table compares the running times of the methods of a priority queue realized by means of a sorted and unsorted list, respectively. We see an interesting trade-off when we use a list to implement the priority queue ADT. An unsorted list supports fast insertions but slow queries and deletions, whereas a sorted list allows fast queries and deletions, but slow insertions.

Operation	Unsorted List	Sorted List
len()	O(1)	O(1)
<pre>is_empty()</pre>	O(1)	O(1)
add(o)	O(1)	O(n)
o = min()	O(n)	O(1)
<pre>o = remove_min()</pre>	O(n)	O(1)
space	O(n)	O(n)

In [6]:

```
class SortedPriorityQueue(PriorityQueueBase): # base class defines Item
    """A min-oriented priority queue implemented with a sorted list."""
                       ----- public behaviors -----
        init__(self):
        """Create a new empty Priority Queue."""
        self. data = PositionalList()
         <u>len</u> (self):
        """Return the number of items in the priority queue."""
        return len(self. data)
    def add(self, key, value):
        """Add a key-value pair."""
        newest = self. Item(key, value)
                                                                # make new item instance
        walk = self. data.last()
                                            # walk backward looking for smaller key
        while walk is not None and newest < walk.element():</pre>
           walk = self._data.before(walk)
        if walk is None:
            self._data.add first(newest)
                                                                    # new key is smallest
                                                             # newest goes after walk
           self._data.add_after(walk, newest)
        """Return but do not remove (k,v) tuple with minimum key.
        Raise Empty exception if empty.
```

```
if self.is empty():
           raise Empty('Priority queue is empty.')
        p = self._data.first()
        item = p.element()
        return (item. key, item. value)
    def remove min(self):
        """Remove and return (k,v) tuple with minimum key.
        Raise Empty exception if empty.
        if self.is empty():
           raise Empty('Priority queue is empty.')
        item = self._data.delete(self._data.first())
        return (item._key, item._value)
    # Custom function to display the entire elements
    def list elems(self):
        listed elem = []
        if self.is_empty():
                                                         # is empty inherited from base class
            return listed elem
        walk = self._data.first()
        while walk is not None:
            item = walk.element()
            listed_elem.append((item._key, item._value))
            walk = self. data.after(walk)
        return listed_elem
In [7]:
Q2 = SortedPriorityQueue()
for name in random.sample(names, 5):
   key = random.randint(0, 100)
   print(key, name)
   Q2.add(key, name)
70 Juliett
77 Sierra
21 November
28 Charlie
63 Golf
In [8]:
Q2.list_elems()
Out[8]:
[(21, 'November'),
 (28, 'Charlie'),
 (63, 'Golf'),
 (70, 'Juliett'),
 (77, 'Sierra')]
In [9]:
print(Q2.remove min())
print(Q2.remove min())
print(Q2.remove min())
print(Q2.remove_min())
print(Q2.remove_min())
(21, 'November')
(28, 'Charlie')
(63, 'Golf')
(70, 'Juliett')
(77, 'Sierra')
```

Heap-based priority Queue

We provide a Python implementation of a heap-based priority queue in the following code. We use an array-based representation, maintaining a Python list of item composites. Although we do not formally use the binary tree ADT, the code includes nonpublic utility functions that compute the level numbering of a parent or child of another. This allows us to describe the rest of our algorithms using tree-like terminology of *parent*, *left*, and *right*. However, the relevant variables are integer indexes (not "position" objects). We use recursion to implement the repetition in the upheap and downheap utilities.

The table displays the running time of the methods in different priority queue implementations.

Operation	Unsorted List	Sorted List	Heap-based
len()	O(1)	O(1)	O(1)
is_empty()	O(1)	O(1)	O(1)
add(o)	O(1)	O(n)	$O(\$\log n\$)$
o = min()	O(n)	O(1)	O(1)
<pre>o = remove_min()</pre>	O(n)	O(1)	$O(\$\log n\$)$
space	O(n)	O(n)	O(n)

In [10]:

```
# Copyright 2013, Michael H. Goldwasser
# Developed for use with the book:
        Data Structures and Algorithms in Python
        Michael T. Goodrich, Roberto Tamassia, and Michael H. Goldwasser
        John Wiley & Sons, 2013
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# You should have received a copy of the GNU General Public License
# along with this program. If not, see <a href="http://www.gnu.org/licenses/">http://www.gnu.org/licenses/</a>.
from priority queue base import PriorityQueueBase
from exceptions import Empty
class HeapPriorityQueue (PriorityQueueBase): # base class defines _Item
    """A min-oriented priority queue implemented with a binary heap."""
    #----- nonpublic behaviors ------
   def parent(self, j):
       return (j-1) // 2
   def left(self, j):
       return 2*j + 1
   def right(self, j):
       return 2*j + 2
   def has left(self, j):
                                                      # index beyond end of list?
        return self. left(j) < len(self. data)</pre>
   def has right(self, j):
       return self._right(j) < len(self._data)</pre>
                                                      # index beyond end of list?
   def swap(self, i, j):
        """Swap the elements at indices i and j of array."""
       self. data[i], self. data[j] = self. data[j], self. data[i]
   def _upheap(self, j):
       parent = self. parent(j)
       if j > 0 and self. data[j] < self. data[parent]:</pre>
           self._swap(j, parent)
```

```
self. upheap (parent)
                                                    # recur at position of parent
   def downheap(self, j):
       if self._has_left(j):
           left = self. left(j)
           small child = left
                                                            # although right may be smaller
           if self. has right(j):
               right = self._right(j)
                if self._data[right] < self._data[left]:</pre>
                   small child = right
            if self._data[small_child] < self._data[j]:</pre>
               self._swap(j, small_child)
               self. downheap(small child)
                                                   # recur at position of small child
                         ----- public behaviors ------
   def __init__(self):
    """Create a new empty Priority Queue."""
        init (self):
       self. data = []
         _len__(self):
        """Return the number of items in the priority queue."""
       return len(self._data)
   def add(self, key, value):
       """Add a key-value pair to the priority queue."""
       self._data.append(self._Item(key, value))
       self._upheap(len(self._data) - 1)
                                                                  # upheap newly added position
   def min(self):
       """Return but do not remove (k,v) tuple with minimum key.
       Raise Empty exception if empty.
       if self.is empty():
           raise Empty('Priority queue is empty.')
       item = self._data[0]
       return (item._key, item._value)
   def remove min(self):
        """Remove and return (k,v) tuple with minimum key.
       Raise Empty exception if empty.
       if self.is empty():
          raise Empty('Priority queue is empty.')
       self. swap(0, len(self. data) - 1)
                                                                # put minimum item at the end
       item = self._data.pop()
                                                                            # and remove it from the
list;
       self. downheap(0)
                                                                                   # then fix new roc
       return (item._key, item._value)
   def display_heap(self, j=0): # j: position
       if len(self) == 0:
           return
       height = int((j+1)/2)
       print('+-'*height+str(self. data[j]))
       if self._has_left(j):
           left = self. left(j)
           self.display heap(left)
       if self. has right(j):
           right = self._right(j)
           self.display_heap(right)
```

In [11]:

```
Qh = HeapPriorityQueue()

for name in random.sample(names, 5):
    key = random.randint(0, 100)
    print("Inserted:", key, name)
    Qh.add(key, name)
    Qh.display_heap()
    print()
```

```
Inserted: 48 Papa
(48, Papa)
Inserted: 54 Tango
(48, Papa)
+- (54, Tango)
Inserted: 77 Charlie
(48, Papa)
+-(54, Tango)
+-(77, Charlie)
Inserted: 79 Lima
(48, Papa)
+- (54, Tango)
+-+-(79, Lima)
+-(77, Charlie)
Inserted: 91 Oscar
(48, Papa)
+-(54, Tango)
+-+-(79, Lima)
+-+-(91, Oscar)
+-(77, Charlie)
In [12]:
for i in range(5):
    print("Removed:", Qh.remove_min())
    Qh.display_heap()
    print()
Removed: (48, 'Papa')
(54, Tango)
+-(79, Lima)
+-+-(91, Oscar)
+-(77, Charlie)
Removed: (54, 'Tango')
(77, Charlie)
+-(79, Lima)
+-(91,0scar)
Removed: (77, 'Charlie')
(79,Lima)
+- (91, Oscar)
Removed: (79, 'Lima')
(91,0scar)
Removed: (91, 'Oscar')
```

Sorting with priority queue

As our first application of priority queues, we demonstrate how they can be used to sort a collection C of comparable elements. That is, we can produce a sequence of elements of C in increasing order (or at least in nondecreasing order if there are duplicates). The algorithm is quite simple—we insert all elements into an initially empty priority queue, and then we repeatedly call remove min to retrieve the elements in nondecreasing order.

```
In [13]:
```

```
def pq_sort(elements, PQ):
    n = len(elements)
    sorted = []
    for elem in elements:
        PQ.add(elem, elem)

for j in range(n):
        (k, v) = PQ.remove_min()
```

```
return sorted

In [14]:
import random
import time

unsorted_list = []
for i in range(1000):
    unsorted_list.append(random.randrange(1000))

#print(unsorted_list)
```

In [15]:

```
# Selection Sort

start_time = time.time()
sorted_list = pq_sort(unsorted_list, SortedPriorityQueue())
end_time = time.time()
sorted_time = end_time - start_time

#print(sorted_list)
```

In [16]:

```
# Insertion Sort

start_time = time.time()
sorted_list = pq_sort(unsorted_list, UnsortedPriorityQueue())
end_time = time.time()
unsorted_time = end_time - start_time

#print(sorted_list)
```

In [17]:

In [18]:

```
print(sorted_time, unsorted_time, heapsort_time)
```

 $0.6700878143310547\ 1.3490607738494873\ 0.03217196464538574$

Python's heapq module

heapq module provides a convenient Heap implementation in Python. https://docs.python.org/3.8/library/heapq.html

In [19]:

```
from heapq import heappush, heappop, heapify

def heapsort(iterable):
    h = []
    for value in iterable:
        heappush(h, value)
    return [heappop(h) for i in range(len(h))]
```

```
In [20]:
print(heapsort([1, 3, 5, 7, 9, 2, 4, 6, 8, 0]))
[0, 1, 2, 3, 4, 5, 6, 7, 8, 9]
In [21]:
items = random.sample(range(100), 10)
print(items)
heapify(items) # This runs in linear time => in-place and bottom-up construction!
print(items)
[63, 9, 34, 8, 52, 16, 57, 84, 12, 75]
[8, 9, 16, 12, 52, 34, 57, 84, 63, 75]
In [22]:
h = []
heappush(h, (5, 'write code'))
heappush(h, (7, 'release product'))
heappush(h, (1, 'write spec'))
heappush(h, (3, 'create tests'))
In [23]:
print (heappop (h) )
print(heappop(h))
print(heappop(h))
print(heappop(h))
(1, 'write spec')
(3, 'create tests')
(5, 'write code')
(7, 'release product')
In [ ]:
```