

Data Structure_Assignment#2

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#A4 - Linked List (5+1 points)

#A4-1(1 point)

singly_linked_list의 경우, x와 y를 swap하는 과정에서 x의 previous node와 y의 previous node를 모두 알아야 한다. 이 과정에서 각각의 previous node를 찾기 위해 $O(n)$ time이 소요된다. 반면 doubly_linked_list의 경우 x와 y의 previous 노트 탐색을 $O(1)$ time에 수행 가능하다.(모든 node가 previous_node_field와 next_node_field를 갖고 있기 때문) 따라서 singly_linked_list의 작업 소요 시간이 더 길다.

#A4-2(1 point)

```
from positional_list import PositionalList
class personal_PL_2(PositionalList):
    def __reversed__(self):
        cursor = self.last()
        while cursor is not None:
            yield cursor.element()
            cursor = self.before(cursor)
```

#A4-3(1 point)

```
from positional_list import PositionalList
class personal_PL(PositionalList):
    def add_last(self, elem):
        if self.is_empty() == True: #리스트가 비어있으면
            return self.add_first(elem) #header 노드 뒤에 elem 삽입
        else:
            last = self.last()
            return self.add_after(last, elem)

    def add_before(self, pos, elem):
        if self.is_empty() == True:
            return self.add_first(elem)
        else:
            before_node = self.before(pos) #pos 이전 node를 before_node에 할당
            if before_node == None:
                return self.add_first(elem)
            else:
                return self.add_after(before_node, elem) #before_node의 다음 node
에 elem 삽입
```

```
#A4-4(1 point)
from positional_list import PositionalList
class personal_PL_4(PositionalList):
    def max(self):
        k = 0
        iterator = iter(self)
        for i in iterator:
            if type(i) != int or type(i) != float:
                raise TypeError('Check your element type')
            else:
                if k < i:
                    k = i
        return k
```

```
#A4-5(1 point)
from positional_list import PositionalList
class personal_PL_5(PositionalList):
    def find(self, elem):
        iterator = iter(self)
        for i in iterator:
            if elem == i:
                return self._make_position(i)
            else:
                return None
```

```
#A4-Bonus(1 point)
from positional_list import PositionalList
class personal_PL_B(PositionalList):
    walk = self.first()
    def find(self, elem):
        if self.is_empty():
            raise EmptyError('The list is empty')
        else:
            if self.first().element() == elem:
                return self._make_position(walk.element())
            else:
                walk = self.after(walk)
                if walk.element() == elem:
                    return self._make_position(walk)
                else:
                    return self.find(self.element())
```

L의 capacity를 n이라 하면, 위 함수는 L에서 elem을 탐색하기 위해 n번의 재귀호출이 필요하다. 이는 새로 생성되는 space가 n번 생성되는 뜻과 상통하고, 따라서 공간복잡도는 $O(n)$ 이 된다.

#A5 - Trees (7 points)

#A5-1(1 point)

1. /user/rt/courses/
2. /user/rt/courses/, cs016/, homeworks/, programs/, cs252/, projects/, papers/, /demos/
3. cs016/ has 9 decendants.
4. It has only 1 ancestor.(//user/rt/courses/)

5. Its siblings are grades/ and programs/.
6. papers/ and demos/
7. 4 (*root의 height == 1로 가정하였습니다.*)
8. 5 (위의 과정을 따름)

#A5-2 (1 point)

```

E
/\
X N
/\
A U
/\
M F

```

#A5-3 (2 points)

```

def root():
    for node in T:
        if f(node) == 0:
            return node

def is_root(p):
    return f(p) == 0

def parent(p):
    if f(p) % 2 == 0:
        for node in T:
            return T[(f(p))/2-1]
    elif f(p) % 2 == 1:
        for node in T:
            return T[(f(p)-1)/2]

def left(p):
    for node in T:
        if node == parent(T[2*f(p)+1]):
            return node
    else:
        return None

def right(p):
    for node in T:
        if node == parent(T[2*f(p)+2]):
            return node
    else:
        return None

def is_leaf(p):
    if left(p) != None or right(p) != None:
        False
    else:
        True

```

```
#A5-4 (3 points)
from linked_binary_tree import LinkedBinaryTree
class personal_LBP(LinkedBinaryTree):
    def _delete_subtree(self, p):
        if self.num_children(p) != 0:
            for child in self.children(p):
                self._delete_subtree(child)
            return self._delete(child)
        else:
            return self._delete(p)
```

running time의 경우, $O(n)$ time이 소요된다.

#A6 - Priority Queues, Heaps (8+2 points)

#A6-1 (1 point)

heap에서의 remove_min 실행의 시간복잡도는 $O(\log(n))$ 이다. 그런데, 문제에서 removing smallest elements의 횟수를 $\log(n)$ 이라 지정하였으므로, 총 시간복잡도는 $O((\log(n))^2)$ 이다.

#A6-2 (1 point)

1. remove_min()
In PQ: [(1, D), (4, B), (5, A), (7, F)] ---> (1, D)
2. remove_min()
In PQ: [(3, j), (4, B), (5, A), (6, L), (7, F)] ---> (3, j)
3. remove_min()
In PQ: [(4, B), (5, A), (6, L), (7, F)] ---> (4, B)
4. remove_min()
In PQ: [(5, A), (6, L), (7, F), (8, G)] ---> (5, A)
5. remove_min()
In PQ: [(2, H), (6, L), (7, F), (8, G)] ---> (2, H)
6. remove_min()
In PQ: [(6, L), (7, F), (8, G)] ---> (6, L)

Result: (1, D), (3, j), (4, B), (5, A), (2, H), (6, L)

#A6-3 (1 point)

- Priority Queue
future event를 추가하고, time stamp가 가장 낮은 것부터 extract 해야한다. 따라서 key값의 우선 순위를 비교하여 pop하는 priority queue의 구조와 흡사하다.

#A6-4 (1 point)

Input: Sequence: (22, 15, 36, 44, 10, 3, 9, 13, 29, 25) Queue: ()

Phase1

a. S: (15, 36, 44, 10, 3, 9, 13, 29, 25) Q: (22)

b. S: (36, 44, 10, 3, 9, 13, 29, 25) Q: (22, 15)

...

j. S: () Q: (22, 15, 36, 44, 10, 3, 9, 13, 29, 25)

Phase2

a. S: (3)	Q: (22, 15, 36, 44, 10, 9, 13, 29, 25)
b. S: (3, 9)	Q: (22, 15, 36, 44, 10, 13, 29, 25)
c. S: (3, 9, 10)	Q: (22, 15, 36, 44, 13, 29, 25)
d. S: (3, 9, 10, 13)	Q: (22, 15, 36, 44, 29, 25)
e. S: (3, 9, 10, 13, 15)	Q: (22, 36, 44, 29, 25)
f. S: (3, 9, 10, 13, 15, 22)	Q: (36, 44, 29, 25)
g. S: (3, 9, 10, 13, 15, 22, 25)	Q: (36, 44, 29)
h. S: (3, 9, 10, 13, 15, 22, 25, 29)	Q: (36, 44)
i. S: (3, 9, 10, 13, 15, 22, 25, 29, 36)	Q: (44)
j. S: (3, 9, 10, 13, 15, 22, 25, 29, 36, 44)	Q: ()

#A6-5 (1 point)

Input: Sequence: (22, 15, 36, 44, 10, 3, 9, 13, 29, 25) Priority Queue: ()

Phase1

a. S: (15, 36, 44, 10, 3, 9, 13, 29, 25)	Q: (22)
b. S: (36, 44, 10, 3, 9, 13, 29, 25)	Q: (15, 22)
c. S: (44, 10, 3, 9, 13, 29, 25)	Q: (15, 22, 36)
d. S: (10, 3, 9, 13, 29, 25)	Q: (15, 22, 36, 44)
e. S: (3, 9, 13, 29, 25)	Q: (10, 15, 22, 36, 44)
f. S: (9, 13, 29, 25)	Q: (3, 10, 15, 22, 36, 44)
g. S: (13, 29, 25)	Q: (3, 9, 10, 15, 22, 36, 44)
h. S: (29, 25)	Q: (3, 9, 10, 13, 15, 22, 36, 44)
i. S: (25)	Q: (3, 9, 10, 13, 15, 22, 29, 36, 44)
j. S: ()	Q: (3, 9, 10, 13, 15, 22, 25, 29, 35, 44)

Phase2

a. S: (3)	Q: (9, 10, 13, 15, 22, 25, 29, 35, 44)
b. S: (3, 9)	Q: (10, 13, 15, 22, 25, 29, 35, 44)
...	
j. S: (3, 9, 10, 13, 15, 22, 25, 29, 35, 44)	Q: ()

#A6-6 (3 points)

```
from heap_priority_queue import HeapPriorityQueue
class heapq(HeapPriorityQueue):
    def heappushpop(self, elem):
        self.add(elem[0], elem[1]) #elem의 key, value쌍 입력
        return self.remove_min() #downheap의 경우 remove_min에서 자동으로 수행

    def heapreplace(self, elem):
        show = self.remove_min() #조건에 맞추어 새로운 element가 add되기 전 최솟값을
        self.add(elem[0], elem[1])
        return show
```

#A6-Bonus (2 points)

```
from heap_priority_queue import HeapPriorityQueue
class personal_HPQ(HeapPriorityQueue):
    def upheap(self, i): #upheap과 downheap을 문제 조건에 맞게 재구성
```

```

parent = self._parent(i)
if i > 0 and self._data[i] > self._data[parent]:
    self._swap(i, parent)
    self._upheap(parent)

def downheap(self, j):
    if self._has_left(j):
        left = self.left(j)
        big_child = left
        if self._has_right(j):
            right = self._right(j)
            if self._data[right] > self._data[left]:
                big_child = right
        if self._data[big_child] > self._data[j]:
            self.swap(j, big_child)
            self.downheap(big_child)

def heapsort(a):
    heap = personal_HPQ()
    for key in a:
        heap.add(key, 'value')
    for i in range(len(a)):
        return self.remove_min()

```