

Minimal Proof System of Matching Logic

Formal Systems Laboratory

November 13, 2017

The following inference rules are shown to constitute a sound and complete proof system of matching logic. We conjecture that all rules marked with “*” are derivable from the rest plus the axiom (DEFINEDNESS). In other words, we conjecture that all rules that are *not* marked with “*” constitute a sound and complete proof system of matching logic.

PROPOSITIONAL ₁	$\varphi_1 \rightarrow (\varphi_2 \rightarrow \varphi_1)$
PROPOSITIONAL ₂	$(\varphi_1 \rightarrow (\varphi_2 \rightarrow \varphi_3)) \rightarrow ((\varphi_1 \rightarrow \varphi_2) \rightarrow (\varphi_1 \rightarrow \varphi_3))$
PROPOSITIONAL ₃	$(\neg\varphi_1 \rightarrow \neg\varphi_2) \rightarrow (\varphi_2 \rightarrow \varphi_1)$
VARIABLE SUBSTITUTION	$\forall x.\varphi \rightarrow \varphi[y/x]$
\forall	$\forall x.(\varphi_1 \rightarrow \varphi_2) \rightarrow (\varphi_1 \rightarrow \forall x.\varphi_2)$ if x does not occur free in φ_1
K	$\sigma(\dots, \varphi_1 \vee \varphi_2, \dots) \leftrightarrow \sigma(\dots, \varphi_1, \dots) \vee \sigma(\dots, \varphi_2, \dots)$
BARCAN	$\sigma(\dots, \exists x.\varphi, \dots) \leftrightarrow \exists x.\sigma(\dots, \varphi, \dots)$ if x does not occur free in the left-hand side of the double implication
*MEMBERSHIP SYMBOL	$x \in \sigma(\dots, \varphi_i, \dots) \rightarrow \exists y.y \in \varphi_i \wedge x \in \sigma(\dots, y, \dots)$
*MEMBERSHIP \neg	$x \in \neg\varphi \rightarrow \neg(x \in \varphi)$
*MEMBERSHIP \exists	$x \in \forall y.\varphi \rightarrow \forall y.(x \in \varphi)$ where x is distinct from y
*DEFINEDNESS	$[x]$
N	From φ deduce $\neg\sigma(\neg\varphi)$
MODUS PONENS	From φ_1 and $\varphi_1 \rightarrow \varphi_2$ deduce φ_2
UNIVERSAL GENERALIZATION	From φ deduce $\forall x.\varphi$
EQUATIONAL SUBSTITUTION	From $\varphi_1 \leftrightarrow \varphi_2$ deduce $\varphi_3[\varphi_1/x] \leftrightarrow \varphi_3[\varphi_2/x]$
FRAMING	From $\varphi_1 \rightarrow \varphi_2$ deduce $\sigma(\varphi_1) \rightarrow \sigma(\varphi_2)$
*INTRODUCTION _{\in}	From φ deduce $x \in \varphi$ If x does not occur free in φ
*ELIMINATION _{\in}	From $x \in \varphi$ deduce φ If x does not occur free in φ