

The Semantics of K

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1 Matching Logic

Let us recall the basic grammar of matching logic from [?]. Assume a matching logic *signature* (S, Σ) , and let Var_s be a countable set of *variables* of sort s . For simplicity, here we assume that the sets of *sorts* S and of *symbols* Σ are finite. We partition Σ in sets of symbols $\Sigma_{s_1 \dots s_n, s}$ of *arity* $s_1 \dots s_n, s$, where $s_1, \dots, s_n, s \in S$. Then *patterns* of sort $s \in S$ are generated by the following grammar:

Add references.

$$\begin{aligned} \varphi_s ::= & x:s \quad \text{where } x \in Var \\ & | \varphi_s \wedge \varphi_s \\ & | \neg \varphi_s \\ & | \exists x:s'. \varphi_s \quad \text{where } x \in N \text{ and } s' \in S \\ & | \sigma(\varphi_{s_1}, \dots, \varphi_{s_n}) \quad \text{where } \sigma \in \Sigma \text{ has } n \text{ arguments, and } \dots \end{aligned}$$

The grammar above only defines the syntax of (well-formed) patterns of sort s . It says nothing about their semantics. For example, patterns $x:s \wedge y:s$ and $y:s \wedge x:s$ are distinct elements in the language of the grammar, in spite of them being semantically/provably equal in matching logic.

For notational convenience, we take the liberty to use mix-fix syntax for operators in Σ , parentheses for grouping, and omit variable sorts when understood. For example, if $Nat \in S$ and $-, +, -, * \in \Sigma_{Nat \times Nat, Nat}$ then we may write $(x + y) * z$ instead of $-, * \in \Sigma_{Nat \times Nat, Nat}$ then we may write $(x + y) * z$ instead of $-, * \in \Sigma_{Nat \times Nat, Nat}$.

A matching logic *theory* is a triple (S, Σ, A) where (S, Σ) is a signature and A is a set of patterns called *axioms*. Like in many logics, sets of patterns may

I think we also need to talk about: other logical connectives as derived, free variables, capture-free substitution, equality. Add more as we need them.

be presented as *schemas* making use of meta-variables ranging over patterns, sometimes constrained to subsets of patterns using side conditions. For example:

$\varphi[\varphi_1/x] \wedge (\varphi_1 = \varphi_2) \rightarrow \varphi[\varphi_2/x]$ where φ is any pattern and φ_1, φ_2 are any patterns of same sort as x

$(\lambda x.\varphi)\varphi' = \varphi[\varphi'/x]$ where φ, φ' are *syntactic patterns*, that is, ones formed only with variables and symbols

$\varphi_1 + \varphi_2 = \varphi_1 +_{\text{Nat}} \varphi_2$ where φ, φ' are *ground* syntactic patterns of sort *Nat*, that is, patterns built only with symbols **zero** and **succ**

$(\varphi_1 \rightarrow \varphi_2) \rightarrow (\varphi[\varphi_1/x] \rightarrow \varphi[\varphi_2/x])$ where φ is a *positive context in x* , that is, a pattern containing only one occurrence of x with no negation (\neg) on the path to x , and where φ_1, φ_2 are any patterns having the same sort

One of the major goals of this paper is to propose a formal language and an implementation, that allows us to write such pattern schemas.

2 A Calculus of Matching Logic

In this section, we propose a calculus of matching logic as a matching logic theory.

Xiaohong: Defining the meta-level in matching logic seems to be a good idea ...

Many people have developed calculi for mathematical reasoning. A calculus of logics is often called a *logical framework*. I prefer to speak of a *meta-logic* and its *object-logic*.

By L. Paulson, *The Foundation of a Generic Theorem Prover*

In this proposal, the *object-logic* refers to matching logic, and we propose to use matching logic itself as the *meta-logic*¹. The calculus of matching logic, denoted as $K = (S_K, \Sigma_K, A_K)$, is a matching logic theory where S_K is a set of sorts, Σ_K is a set of symbols, and A_K is a set of axioms. In S_K , there are sorts for abstract syntax trees (ASTs) of matching logic patterns and proof system infrastructure, beside built-in sorts.

$$S_K = \{ \text{Bool}, \text{String}, \text{Int}, \dots, \\ \text{Pattern}, \text{Sort}, \text{Symbol}, \text{Signature}, \text{Theory}, \\ \text{Goal}, \text{InferenceRule}, \text{Tactic}, \text{Validation}, \text{Tactical}, \\ \text{StringList}, \text{PatternList}, \text{SymbolList}, \dots \}.$$

¹Coq and Isabelle use fragments of higher-order logic as their meta-logics.

In Σ_K , there are constructors of ASTs and the proof system infrastructure, constructors of lists (on various sorts), and many other symbols and built-in symbols, which we all conclude in the next figure, in which we feel free to use all kinds of conventional notations, overloading, and syntactic sugar to make our lives easier.

$true, false: \rightarrow Bool$
 $and, or, implies: Bool \times Bool \rightarrow Bool$
 $\#sort: String \rightarrow Sort$
 $\#variable: String \times Sort \rightarrow Pattern$
 $\#and, \#or, \#implies, \#iff,$
 $\#equals, \#contains: Pattern \times Pattern \rightarrow Pattern$
 $\#not: Pattern \rightarrow Pattern$
 $\#exists, \#forall: String \times Sort \times Pattern \rightarrow Pattern$
 $\#application: Symbol \times PatternList \rightarrow Pattern$
 $\#top, \#bottom: \rightarrow Pattern$
 $wellFormed: Pattern \rightarrow Bool$
 $getSort \in \Sigma_{Pattern, Sort}$
 $replaceAll: Pattern \times Pattern \times Pattern \rightarrow Pattern$
 $deducible: Pattern \rightarrow Bool$
 \dots

2.1 Faithfulness

² It remains a question whether the calculus of matching logic M *faithfully* captures the reasoning in matching logic. Before we define faithfulness, let us introduce some useful notations. Suppose $T = (S, \Sigma, A)$ is a matching logic theory. The meta-theory of T , denoted as $\#T$, is a matching logic theory obtained by adding

$$deducible(\#\varphi) = true \quad \text{for all } \varphi \in A$$

as axioms to the meta-logic theory M , where the sharp $\#_-$ is the lifting function that maps object-patterns to their meta-representations (ASTs) in M . For simplicity, we denote $deducible(\#\varphi) = true$ as $\llbracket \varphi \rrbracket$, a pattern of the sort *Pattern*.

The following definition is inspired by Paulson's paper (Definition 1).

Definition 1. Let T be a matching logic theory and $\#T$ is its meta-theory. Tet $\varphi_1, \dots, \varphi_n$ and ψ be patterns of T . Then say

- $\#T$ is *sound* for T , if for every $\#T$ -proof of $\llbracket \psi \rrbracket$ from $\llbracket \varphi_1 \rrbracket, \dots, \llbracket \varphi_n \rrbracket$, there is a T -proof of ψ from $\varphi_1, \dots, \varphi_n$.

²It is strongly recommend that readers of this proposal read L. Paulson's *The Foundation of a Generic Theorem Prover*, especially Section 2, 3, and 4.

- $\#T$ is *complete* for T , if for every T -proof of ψ from $\varphi_1, \dots, \varphi_n$, there is a $\#T$ -proof of $\llbracket \psi \rrbracket$ from $\llbracket \varphi_1 \rrbracket, \dots, \llbracket \varphi_n \rrbracket$.
- $\#T$ is *faithful* for T if it is both sound and complete.

As a result of $\#T$ being faithful for T , for any pattern φ , φ is deducible in T iff $\llbracket \varphi \rrbracket$ is deducible in $\#T$.

Theorem 2. *For any matching logic theory T , its meta-theory $\#T$ is faithful for T .*

```
fth M-LAMBDA is
  extending M .

  ---- the following syntactic sugar is just for readability.

  ops app lambda0 : -> Symbol .
  eq app = #symbol("app", ("Exp", "Exp"), "Exp") .
  eq lambda0 = #symbol("lambda0", ("Exp", "Exp"), "Exp") .

  op lambda_._ : String K -> K .
  eq lambda X:VariableId . E:K
    = #exists(X:VariableId, "Exp",
      #application(lambda0, (#variable(X:VariableId, "Exp"), E:K))) .
  op _[_] : K K -> K .
  eq E1:K[E2:K]
    = #application(app, (E1:K, E2:K)) .

  ---- side conditions checker

  op isLTerm : K -> Bool .

  eq isLTerm(#variable(X:VariableId, "Exp")) = true .

  eq isLTerm(E1:K[E2:K])
    = isLTerm(E1:K) and isLTerm(E2:K) .

  eq isLTerm(lambda(X:VariableId, E:K))
    = isLTerm(E:K) .

  ---- the (Beta) axiom

  cmb #equal((lambda X:VariableId . E1:K)[E2:K],
    replace(...)) : Theorem
    if isLTerm(E1:K) and isLTerm(E2:K) .
endth
```

3 The Kore Language

We have shown a calculus of matching logic M in which we can specify everything about matching logic theories, for example, whether a pattern is well-formed, what sort a patten has, which patterns are deducible, free variables, fresh variables generation, substitution and replacement, alpha-renaming, etc. This meta-theory provides a universe of (meta-representations of) patterns and the proof system of matching logic. On the other hand, it is usually easier to work in the object-level rather than the meta-level. Even if all reasoning in the object-level theory T can be faithfully lifted to and conducted in the meta-theory $\#T$, it does not mean one should always do so.

3.1 Syntax of Kore

We propose the next Kore syntax which allows the use of meta-variables at object-level. This is a syntactic sugar that lets users write compact axioms at object-level.

```
// Namespaces for sorts, variables, metavariables,
// symbols, and Kore modules.
SortId      = String
VariableId  = String
MetaVariableId = String
SymbolId    = String
ModuleId    = String

Sort        = SortId

Variable    = VariableId:SortId
MetaVariable = MetaVariableId::SortId

Pattern     = Variable | MetaVariable
             | \and(Pattern, Pattern)
             | \not(Pattern)
             | \exists(Variable, Pattern)
             | SymbolId(List{Pattern})

Sentence    = import ModuleId
             | syntax SortId
             | syntax SortId ::= SymbolId(List{SortId})
             | axiom Pattern

Sentences   = Sentence | Sentences Sentences

Module      = module ModuleId
             Sentences
             endmodule
```

3.2 Semantics of Kore

The semantics of Kore is given as a translation from Kore definitions to meta-theories. The main principle is that every object-level constructors are translated to ground terms (ASTs) while meta-variables are translated to variables in the meta-theories. In other words, the sentence **axiom** φ when φ is an object-pattern is a shorthand of **axiom** $\llbracket \varphi \rrbracket$. Since we now allow φ containing meta-variables, the notation $\llbracket \varphi \rrbracket$ now means that $musc(\varphi) \rightarrow (deducible(\# \varphi) = true)$, where the lifting function $\#_-$ and the meta-variable sort constraint $musc$ are defined in Algorithm 1 and 2, respectively.

Algorithm 1: Lifting Function $\#_-$

Input: An object-pattern φ
Output: The meta-representation of φ

```

1 if  $\varphi$  is  $x:s$  then
2   |  $\#variable(x, \#sort(s))$ 
3 else if  $\varphi$  is  $x::s$  then
4   |  $x:Pattern \wedge \#sort(s) \in getSort(x:Pattern)$ 
5 else if  $\varphi$  is  $\varphi_1 \wedge \varphi_2$  then
6   |  $\#and(\# \varphi_1, \# \varphi_2)$ 
7 else if  $\varphi$  is  $\neg \varphi_1$  then
8   |  $\#not(\# \varphi_1)$ 
9 else if  $\varphi$  is  $\exists x:s. \varphi$  then
10  |  $\#exists(x, \#sort(s), \# \varphi)$ 
11 else if  $\varphi$  is  $\sigma(\varphi_1, \dots, \varphi_n)$  then
12  |  $\#application(\#symbol(\sigma, \dots), \# \varphi_1, \dots, \# \varphi_n)$ 

```

Algorithm 2: Meta-Variable Sort Constraint Collection $musc$

Input: An object-pattern φ
Output: The meta-variable sort constraint of φ

```

1 Collect in set  $W$  all meta-variables appearing in  $\varphi$ ;
2 Let  $C = \emptyset$ ;
3 foreach  $x::s \in W$  do
4   |  $C = C \cup (\#sort(s) \in getSort(x:Pattern))$ 
5 Return  $\bigwedge C$ ;

```

3.3 Examples of Kore

The following examples may need some change.

The **BOOL** module.

```

module BOOL
  syntax Bool
  syntax Bool ::= trueBool | falseBool | notBool(Bool)
                | andBool(Bool, Bool) | orBool(Bool, Bool)
  axiom \or(trueBool, falseBool)
  axiom \exists(X:Bool, \equal(X:Bool, trueBool))
  axiom \exists(X:Bool, \equal(X:Bool, notBool(Y:Bool)))
  axiom ... ..
endmodule

```

The NAT module.

```

module NAT
  import BOOL
  syntax Nat
  syntax Nat ::= zero | succ(Nat) | plus(Nat, Nat)
                | mult(Nat, Nat)
  syntax Bool ::= gt(Nat, Nat) | gte(Nat, Nat)
  axiom \equal(mult(X:Nat, zero), zero)
  axiom \equal(mult(X::Nat, Y::Nat), mult(Y::Nat, X::Nat))
  axiom ... ..
  // hook to a model (not finished)
  hook-sort Nat "external-model Naturals"
  hook-symbol zero \value("0", Nat)
  hook-symbol succ \value("succ", Nat)
  ... ..
endmodule

```

The LAMBDA module.

```

module LAMBDA
  import M      ---- import the meta-logic theory M
                ---- we can rename it to K ...
  syntax Exp
  syntax Exp ::= app(Exp, Exp)
                | lambda0(Exp, Exp)
  syntax M.Bool ::= isLTerm(K)      ---- I think the Bool here is the
                                     ---- one in the meta-logic M, not
                                     ---- not one in the Kore module
                                     ---- BOOL
  axiom \equals(app(\exists(X:Exp, lambda0(X:Exp, E::Exp)), E'::Exp),
                \replace(E::Exp, E'::Exp, X:Exp))
  requires \isLTerm(E::Exp) M.and \isLTerm(E'::Exp) ---- requires is a sugar

  axiom isLTerm(X:Exp) = M.true
  axiom isLTerm(app(E::Exp, E'::Exp)) = isLTerm(E::Exp) M.and isLTerm(E'::Exp)

```

```

...
endmodule

```

3.4 Lambda Calculus

```

module LAMBDA
  syntax Exp
  syntax Exp ::= lambda0(Exp, Exp)
                | app(Exp, Exp)
  syntax Bool ::= isLTerm(Pattern)
  axiom isLTerm(X:Exp) = true
  axiom isLTerm(\exists(X:Exp, lambda0(X:Exp, E:Exp)))
    = isLTerm(E:Exp)
  axiom isLTerm(app(E:Exp, E':Exp))
    = isLTerm(E:Exp) andBool isLTerm(E':Exp)
  axiom app(\exists(X:Exp, lambda0(X:Exp, E:Exp)), E':Exp)
    = E:Exp[E':Exp / X:Exp]
  requires isLTerm(E:Exp) andBool isLTerm(E':Exp)
endmodule

```

Many discussions in the next section (Sec.4 Object-level and Meta-level) should be moved to Section 3 as examples. Sec.5 Binders and Sec.6 Contexts should also move to a subsection of Sec.3 as applications and examples.

4 Object-level and Meta-level

It is an aspect of life in mathematical logics to distinguish the *object-level* and *meta-level* concepts. In matching logic, we put more emphasize and care on metavariables and their range, that is, the set of patterns that they stand for. It turns out that having metavariables that range over all well-formed patterns will lead us to inconsistency theories immediately. As an example, consider the (β) axiom in the matching logic theory LAMBDA of lambda calculus:

$$(\lambda x.e)[e'] = e[e'/x].$$

If we do not put any restriction on the range of metavariables e and e' , we have an inconsistency issue as the following reasoning shows:

$$\perp \stackrel{(N)}{=} (\lambda x.\top)[\perp] \stackrel{(\beta)}{=} \top[\perp/x] = \top.$$

Therefore, in matching logic, one should explicitly specify the range of metavariables whenever he uses them.

Definition 3 (Restricted metavariables). Let φ be a metavariable of sort $s \in S$. The range of φ is a set of patterns of sort s . We write $\varphi :: R$ if the range of φ is $R \subseteq \text{Patterns}_s$.

Remark 4 (Metavariables in first-order logic). In first-order logic, one often uses metavariables in axiom schemata, but the inconsistency issue does not arise. This is because in first-order logic, we do not need to distinguish metavariables for terms from logic variables, thanks to the next (Substitution) rule:

$$\forall x.\varphi(x) \rightarrow \varphi(t).$$

The predicate metavariables are not a problem because there are no object level symbols on top of them.

I don't get the point of predicate metavariables.

Variables and metavariables for variables For any matching logic theory $\mathsf{T} = (S, \Sigma, A)$, it comes for each sort $s \in S$ a countably infinite set V_s of variables. We use $x : s, y : s, z : s, \dots$ for variables in V_s , and omit their sorts when that is clear from the contexts. Different sorts have disjoint sets of variables, so $\text{Var}_s \cap \text{Var}_{s'} = \emptyset$ if $s \neq s'$.

Proposition 5. *Let A be a set of axioms and $\bar{A} = \forall A$ be the universal quantification closure of A , then for any pattern φ , $A \vdash \varphi$ iff $\bar{A} \vdash \varphi$.*

Remark 6 (Free variables in axioms). The free variables appearing in the axioms of a theory can be regarded as implicitly universal quantified, because a theory and its universal quantification closure are equal.

Example 7.

$$\begin{aligned} A_1 &= \{\text{mult}(x, 0) = 0\} \\ A_2 &= \{\forall x.\text{mult}(x, 0) = 0\} \\ A_3 &= \{\forall y.\text{mult}(y, 0) = 0\} \\ A_4 &= \{\text{mult}(x, 0) = 0\} \\ A_5 &= \{\forall x.\text{mult}(x, 0) = 0\} \\ A_6 &= \{\forall y.\text{mult}(y, 0) = 0\} \\ A_7 &= \{\text{mult}(x, 0) = 0, \text{mult}(y, 0) = 0, \text{mult}(z, 0) = 0, \dots\} \\ A_8 &= \{\forall x.\text{mult}(x, 0) = 0, \forall y.\text{mult}(y, 0) = 0, \forall z.\text{mult}(z, 0) = 0, \dots\} \end{aligned}$$

All the eight theories are equal. Theories A_4, A_5, A_6 are finite representations of theories A_7, A_8, A_8 respectively.

Remark 8. There is no need to have metavariables for variables in the Kore language, because (1) if they are used as bound variables, then replacing them with any (matching logic) variables will result in the same theories, thanks to alpha-renaming; and (2) if they are used as free variables, then it makes no difference to consider the universal quantification closure of them and we get to the case (1).

~~Given said that, there are cases when metavariables for variables make sense. In those cases we often want our metavariables to range~~

~~over all variables of all sorts, in order to make our Kore definitions compact.~~ No, we do not need metavariables over variables. I was thinking of the definedness symbols. We might want to write only one axiom schema of $[x]$ instead many $[x:s]_s^{s'}$'s, but we cannot do that unless we allow polymorphic and overloaded symbols in Kore definitions.

Patterns and metavariables for patterns It is in practice more common to use metavariables that range over all patterns. One typical example is axiom schemata. For example, $\vdash \varphi \rightarrow \varphi$ in which φ is the metavariable that ranges all well-formed patterns.

There has been an argument on whether metavariables for patterns should be sorted or not. Here are some observations. Firstly, since all symbols are decorated and not overloaded, in most cases, the sort of a metavariable for patterns can be inferred from its context. Secondly, the only counterexample against the first point that I can think of is when they appear alone, which is not an interesting case anyway. Thirdly, we do want the least amount of reasoning and inferring in using Kore definitions, so it breaks nothing if not helping things to have metavariables for patterns carrying their sorts.

Example 9.

$$\begin{aligned} A_1 &= \{\text{merge}(h1, h2) = \text{merge}(h2, h1)\} \\ A_2 &= \{\forall h1 \forall h2. \text{merge}(h1, h2) = \text{merge}(h2, h1)\} \\ A_3 &= \{\text{merge}(\varphi, \psi) = \text{merge}(\psi, \varphi)\} \end{aligned}$$

All three theories are equal. It is easier to see that fact from a model theoretic point of view, since all theories require that the interpretation of `merge` is commutative and nothing more. On the other hand, it is not straightforward to obtain that conclusion from a proof theoretic point of view. For example, to deduce $\text{merge}(\text{list}(\text{one}, \text{cons}(\text{two}, \text{epsilon})), \text{top}) = \text{merge}(\text{top}, \text{list}(\text{one}, \text{cons}(\text{two}, \text{epsilon})))$ needs only one step in A_3 , but will need a lot more in either A_1 or A_2 , because one cannot simply substitute any patterns for universal quantified variables in matching logic.

5 Binders

In matching logic there is a unified representation of binders. We will be using the theory of lambda calculus `LAMBDA` as an example in this section. Recall that the syntax for untyped lambda calculus is

$$\Lambda ::= V \mid \lambda V. \Lambda \mid \Lambda \Lambda$$

where V is a countably infinite set of atomic λ -terms, a.k.a. variables in lambda calculus. The set of all λ -terms, denoted as Λ , is the smallest set satisfying the above grammar.

The matching logic theory **LAMBDA** has one sort **Exp** for lambda expressions. It also has in its signature a binary symbol **lambda**₀ that builds a λ -terms, and a binary symbol **app** for lambda applications. To mimic the binding behavior of λ in lambda calculus, we define syntactic sugar $\lambda x.e = \exists x.\text{lambda}_0(x, e)$ and $e_1 e_2 = \text{app}(e_1, e_2)$ in theory **LAMBDA**. Notice that by defining λ as a syntactic sugar using the existential quantifier $\exists x$, we get alpha-renaming for free. The β -reduction is captured by the next axiom:

$$(\lambda x.e)e' = e[e'/x] \quad , \text{ where } e \text{ and } e' \text{ are metavariables for } \lambda\text{-terms.}$$

Two important observations are made about the (β) axiom. Firstly, e and e' cannot be replaced by logic variables, because λ -terms in matching logic are (often) not functional patterns. Secondly, metavariables e and e' cannot range over all patterns of sort **Exp**, but only those which are (syntactic sugar of) λ -terms. Allowing e and e' to range over all patterns of **Exp** will quickly lead to an inconsistent theory, because of the next contradiction:

$$\perp \stackrel{(N)}{=} (\lambda x.\top)[\perp] \stackrel{(\beta)}{=} \top.$$

Therefore, when defining the lambda calculus, we need a way

Theorem 10 (Consistency). *Consider a theory of a binder α , with a sort S and two binary symbols α_0 and $- \cdot -$. Define $\alpha x.e$ as syntactic sugar of $\exists x.\alpha(x, e)$ where x is a variable and e is a pattern. Define α -terms be patterns satisfying the next grammar*

$$T_\alpha ::= V_s \mid \alpha x.T_\alpha \mid T_\alpha T_\alpha.$$

If a theory contains only axioms of the form $e = e'$ where e and e' are α -terms, then the theory is consistent.

Proof. The final model M exists, in which the carrier set is a singleton set, and the two symbols are interpreted as the total function over the singleton set. One can then prove that all α -terms interpret to the total set, so all axioms hold in the final model. \square

Corollary 11. *The theory **LAMBDA** is consistent.*

Definition 12 (Common ranges of metavariables).

- Full range Pattern_s ;
- Syntactic terms range (variables plus symbols without logic connectives);
- Ground syntactic terms range (symbols only);
- Variable range Var_s (metavariables for variables).

Remark 13. Syntactic terms (and ground syntactic terms) are purely defined syntactically and not equal to terms or functional patterns. When all symbols are functional symbols, the set of syntactic terms equals the set of terms, and both of them are included in the set of all functional patterns.

Remark 14. We need to design a syntax for specifying ranges of metavariables in the Kore language.

Remark 15. We have not proved that matching logic is a conservative extension of untyped lambda calculus, which bothers me a lot. I will remain skeptical about everything we do in this section until we prove that conservative extension result.

The benefit of such a unified theory of binders and binding structures in matching logic is more of theoretical interest. In practice (K backends), one will never want to implement the lambda calculus by desugaring $\lambda x.e$ as $\exists x.\lambda_0(x, \varphi)$ but rather dealing with $\lambda x.\varphi$ directly.

Example 16 (Lambda calculus in Kore).

```
module LAMBDA
  import BOOL
  import META-LEVEL

  syntax Exp
  syntax Exp ::= app(Exp, Exp)
                | lambda0(Exp, Exp)

  axiom \implies(true = andBool(#isLTerm(#up(E:Exp)),
                                #isLTerm(#up(E':Exp))),
                app(\exists(x:Exp, lambda0(x:Exp, E:Exp)), E':Exp)
                  = E:Exp(E':Exp / x:Exp)

  // Q1: what is substitution?
  // Q2: we know #up is not a part of the logic, so what does
  //      it mean?

  syntax Bool ::= #isLTerm(Pattern)
  axiom #isLTerm(#variable(x:Name, s:Sort)) = true
  axiom #isLTerm(#application(
    #symbol(#name("app"), #appendSortList(...), #sort("Exp"))),
    #appendPatternList(#PatternListAsPattern(#P),
                      #PatternListAsPattern(#P')))))
  = andBool(#isLTerm(#P), #isLTerm(#P'))
  ...
endmodule
```

Rewriting logic

6 Contexts

Introduce a binder γ together with its application symbol which we write as $_{-}[-]$. Binding variables of the binder γ are often written as \square , but in this proposal and hopefully in future work we will use regular variables x, y, z, \dots instead of \square , in order to show that there is nothing special about contexts but simply a

theory in matching logic. Patterns of the form $\gamma x.\varphi$ are often called *contexts*, denoted by metavariables C, C_0, C_1, \dots . Patterns of the form $\varphi[\psi]$ are often called *applications*.

Definition 17. The context $\gamma x.x$ is called the identity context, denoted as l . Identity context has the axiom schema $\mathsf{l}[\varphi] = \varphi$ where φ is any pattern.

Example 18. $\mathsf{l}[\mathsf{l}] = \mathsf{l}$.

Definition 19. Let $\sigma \in \Sigma_{s_1 \dots s_n, s}$ is an n -arity symbol. We say σ is *active* on its i th argument ($1 \leq i \leq n$), if

$$\sigma(\varphi_1, \dots, C[\varphi_i], \dots, \varphi_n) = (\gamma x.\sigma(\varphi_1, \dots, C[x], \dots, \varphi_n))[\varphi_i],$$

where $\varphi_1, \dots, \varphi_n$, and C are any patterns. Orienting the equation from the left to the right is often called *heating*, while orienting it from the right to the left is called *cooling*.

Example 20. Assume the next theory of IMP.

$$\begin{aligned} A = \{ & \mathsf{ite}(C[\varphi], \psi_1, \psi_2) = (\gamma x.\mathsf{ite}(C[x], \psi_1, \psi_2))[\varphi], \\ & \mathsf{while}(C[\varphi], \psi) = (\gamma x.\mathsf{while}(C[x], \psi))[\varphi], \\ & \mathsf{seq}(C[\varphi], \psi) = (\gamma x.\mathsf{seq}(C[x], \psi))[\varphi], \\ & C[\mathsf{ite}(\mathsf{tt}, \psi_1, \psi_2)] \Rightarrow C[\psi_1], \\ & C[\mathsf{ite}(\mathsf{ff}, \psi_1, \psi_2)] \Rightarrow C[\psi_2], \\ & C[\mathsf{while}(\varphi, \psi)] \Rightarrow C[\mathsf{ite}(\varphi, \mathsf{seq}(\psi, \mathsf{while}(\varphi, \psi)), \mathsf{skip})], \\ & C[\mathsf{seq}(\mathsf{skip}, \psi)] \Rightarrow C[\psi] \}. \end{aligned}$$

We can simply require that ψ_1, ψ_2, ψ_3 , and C are any patterns. That will allow us to do any reasoning that we need, but will that lead to inconsistency?

Example 20(a).

$$\begin{aligned} \mathsf{seq}(\mathsf{skip}, \mathsf{skip}) &= \mathsf{l}[\mathsf{seq}(\mathsf{skip}, \mathsf{skip})] \\ &\Rightarrow \mathsf{l}[\mathsf{skip}] \\ &= \mathsf{skip}. \end{aligned}$$

Example 20(b).

$$\begin{aligned} \mathsf{seq}(\mathsf{ite}(\mathsf{tt}, \psi_1, \psi_2), \psi_3) &= \mathsf{seq}(\mathsf{l}[\mathsf{ite}(\mathsf{tt}, \psi_1, \psi_2)], \psi_3) \\ &= (\gamma x.\mathsf{seq}(\mathsf{l}[x], \psi_3))[\mathsf{ite}(\mathsf{tt}, \psi_1, \psi_2)] \\ &\Rightarrow (\gamma x.\mathsf{seq}(\mathsf{l}[x], \psi_3))[\psi_1] \\ &= \mathsf{seq}(\mathsf{l}[\psi_1], \psi_3) \\ &= \mathsf{seq}(\psi_1, \psi_3). \end{aligned}$$

Example 21. Consider the following theory written in the Kore language:

```

module IMP
  import ...
  syntax AExp
  syntax AExp ::= plusAExp(AExp, AExp)
  syntax AExp ::= minusAExp(AExp, AExp)
  syntax AExp ::= AExpAsNat(Nat)
  syntax BExp
  syntax BExp ::= geBExp(AExp, AExp)
  syntax BExp ::= BExpAsBool(Bool)
  syntax Pgm
  syntax Pgm ::= skip()
  syntax Pgm ::= seq(Pgm Pgm)
  syntax Pgm ::=
  syntax Heap
  syntax Cfg
endmodule

```

Example 22. Following the above example, extend A with the next axioms:

$$\begin{aligned}
&\{C[x][\text{mapsto}(x, v)] \Rightarrow C[v][\text{mapsto}(x, v)], \\
&C[\text{asgn}(x, v)][\text{mapsto}(x, v')] \Rightarrow C[\text{skip}][\text{mapsto}(x, v)], \\
&C[\text{asgn}(x, v)][\varphi] \Rightarrow C[\text{skip}][\text{merge}(\varphi, \text{mapsto}(x, v))]\}
\end{aligned}$$

The above example is meant to show the loopup rule, but it does not work because the third axiom is incorrect. Instead of simply writing φ , we should say that φ does not assign any value to x . One solution (that is used in the current K backend) is to introduce a strategy language and to extend theories with strategies.

Example 23. Suppose f and g are binary symbols who are active on their first argument. Suppose a, b are constants, and x is a variable. Let \square_1 and \square_2 be two hole variables. Define two contexts $C_1 = \gamma\square_1.f(\square_1, a)$ and $C_2 = \gamma\square_2.g(\square_2, b)$.

Because f is active on the first argument,

$$\begin{aligned}
C_1[\varphi] &= (\gamma\square_1.f(\square_1, a))[\varphi] \\
&= (\gamma\square_1.f(\text{!}[\square_1], a))[\varphi] \\
&= f(\text{!}[\varphi], a) \\
&= f(\varphi, a), \text{ for any pattern } \varphi.
\end{aligned}$$

And for the same reason, $C_2[\varphi] = g(\varphi, b)$. Then we have

$$\begin{aligned}
C_1[C_2[x]] &= C_1[f(x, a)] \\
&= g(f(x, a), b).
\end{aligned}$$

On the other hand,

$$\begin{aligned} g(f(x, a), b) &= g(C_1[x], b) \\ &= (\gamma\Box.g(C_1[\Box], b))[x] \\ &= (\gamma\Box.g(f(\Box, a), b))[x]. \end{aligned}$$

Therefore, the context $\gamma\Box.g(f(\Box, a), b)$ is often called the *composition* of C_1 and C_2 , denoted as $C_1 \circ C_2$.

Example 24. Suppose f is a binary symbol with all its two arguments active. Suppose C_1 and C_2 are two contexts and a, b are constants. Then easily we get

$$\begin{aligned} f(C_1[a], C_2[b]) &= (\gamma\Box_2.f(C_1[a], C_2[\Box_2]))[b] \\ &= (\gamma\Box_2.((\gamma\Box_1.f(C_1[\Box_1], C_2[\Box_2]))[a]))[b]. \end{aligned}$$

What happens above is similar to *curring* a function that takes two arguments. It says that there exists a context C_a , related with C_1, C_2, f and a of course, such that $C_a[b]$ returns $f(C_1[a], C_2[b])$. The context C_a has a binding hole \Box_2 , and a body that itself is another context C'_a applied to a . In other words, there exists C_a and C'_a such that

- $f(C_1[a], C_2[b]) = C_a[b]$,
- $C_a = \gamma\Box_2.(C'_a[a])$,
- $C'_a = \gamma\Box_1.f(C_1[\Box_1], C_2[\Box_2])$.

A natural question is whether there is a context C such that $C[a][b] = f(C_1[a], C_2[b])$.

Proposition 25. $C_1[C_2[\varphi]] = C[\varphi]$, where $C = \gamma\Box.C_1[C_2[\Box]]$.

6.0.1 Normal forms

In this section, we consider *decomposition* of patterns. A decomposition of a pattern P is a pair $\langle C, R \rangle$ such that $C[R] = P$. Let us now consider patterns that do not have logical connectives.

Fixed points

7 Appendix: The First Kore Language

The next grammar is the firstly-proposed Kore language at [here](#).

Definition = Attributes
Set{Module}

Module = module ModuleName

```

Set{Sentence}
endmodule
Attributes

Sentence = import ModuleName Attributes
| syntax Sort Attributes                // sort declarations
| syntax Sort ::= Symbol(List{Sort}) Attributes // symbol declarations
| rule Pattern Attributes
| axiom Pattern Attributes

Attributes = [ List{Pattern} ]

Pattern = Variable
| Symbol(List{Pattern})                // symbol applications
| Symbol(Value)                        // domain values
| \top()
| \bottom()
| \and(Pattern, Pattern)
| \or(Pattern, Pattern)
| \not(Pattern)
| \implies(Pattern, Pattern)
| \exists(Variable, Pattern)
| \forall(Variable, Pattern)
| \next(Pattern)
| \rewrite(Pattern, Pattern)
| \equals(Pattern, Pattern)

Variable = Name:Sort                  // variables

ModuleName = RegEx1
Sort       = RegEx2
Name       = RegEx2
Symbol     = RegEx2
Value      = RegEx3

RegEx1 == [A-Z][A-Z0-9-]*
RegEx2 == [a-zA-Z0-9.@#%~_-]+ | ' [^']* '
RegEx3 == <Strings> // Java-style string literals, enclosed in quotes

```

In the grammar above, `List{X}` is a special non-terminal corresponding to possibly empty comma-separated lists of `X` words (trivial to define in any syntax formalism). `Set{X}`, on the other hand, is a special non-terminal corresponding to possibly empty space-separated sets of `X` words. Syntactically, there is no difference between the two (except for the separator), but Kore tools may choose to implement them differently.

7.1 Builtin theories

```
module BOOL
syntax Bool
syntax Bool ::= true | false | notBool(Bool)
              | andBool(Bool, Bool) | orBool(Bool, Bool)

// axioms for functional symbols
axiom \exists(T:Bool, \equals(T:Bool, true))
axiom \exists(T:Bool, \equals(T:Bool, false))
axiom \exists(T:Bool, \equals(T:Bool, \notBool(X:Bool)))
axiom \exists(T:Bool, \equals(T:Bool, andBool(X:Bool, Y:Bool)))
axiom \exists(T:Bool, \equals(T:Bool, orBool(X:Bool, Y:Bool)))

// axioms for commutativity
axiom \equals(andBool(X:Bool, Y:Bool), andBool(Y:Bool, X:Bool))
axiom \equals(orBool(X:Bool, Y:Bool), orBool(Y:Bool, X:Bool))

// the no-junk axiom for constructors
axiom \or(true, false)

axiom \equals(notBool(true), false)
axiom \equals(notBool(false), true)
axiom \equals(andBool(true, T:Bool), T:Bool)
axiom \equals(andBool(false, T:Bool), false)
axiom \equals(orBool(true, T:Bool), true)
axiom \equals(orBool(false, T:Bool), T:Bool)
endmodule

module META-LEVEL
syntax
endmodule

module LAMBDA
syntax Exp
syntax Exp ::= lambda0(Exp, Exp) | app(Exp, Exp)

endmodule
```