Concurrency Problems | DBMS

► Database Management System

Transactions in DBMS-

Before you go through this article, make sure that you have gone through the previous article on **Transactions in DBMS**.

We have discussed-

- A transaction is a set of logically related operations.
- A transaction goes through different states throughout its life cycle.
- ACID Properties are followed by each transaction to ensure the consistency of database.

In this article, we will discuss about concurrency problems of transactions.

Concurrency Problems in DBMS-

- · When multiple transactions execute concurrently in an uncontrolled or unrestricted manner, then it might lead to several problems.
- Such problems are called as concurrency problems.

The concurrency problems are-



- 1. Dirty Read Problem
- 2. Unrepeatable Read Problem
- 3. Lost Update Problem
- 4. Phantom Read Problem

1. Dirty Read Problem-

Reading the data written by an uncommitted transaction is called as dirty read.

This read is called as dirty read because-

- There is always a chance that the uncommitted transaction might roll back later.
- Thus, uncommitted transaction might make other transactions read a value that does not even exist.
- This leads to inconsistency of the database.

NOTE-

- Dirty read does not lead to inconsistency always.
- It becomes problematic only when the uncommitted transaction fails and roll backs later due to some reason.

Example-

Transaction T1	Transaction T2	
R (A) W (A)		
	R (A) W (A)	// Dirty Read
	Commit	
Failure		

Here,

- 1. T1 reads the value of A.
- 2. T1 updates the value of A in the buffer.
- 3. T2 reads the value of A from the buffer.
- 4. T2 writes the updated the value of A.
- 5. T2 commits.
- 6. T1 fails in later stages and rolls back.

In this example,

- $\bullet~$ T2 reads the dirty value of A written by the uncommitted transaction T1.
- T1 fails in later stages and roll backs.
- Thus, the value that T2 read now stands to be incorrect.
- Therefore, database becomes inconsistent.

2. Unrepeatable Read Problem-

This problem occurs when a transaction gets to read unrepeated i.e. different values of the same variable in its different read operations even when it has not updated its value.

Example-

Transaction T1	Transaction T2	_
R (X)		
	R (X)	
W (X)		
	R (X) // Ur	repeated Read

Here,

- 1. T1 reads the value of X (= 10 say).
- 2. T2 reads the value of X (= 10).
- 3. T1 updates the value of X (from 10 to 15 say) in the buffer.
- 4. T2 again reads the value of X (but = 15).

In this example,

- T2 gets to read a different value of X in its second reading.
- T2 wonders how the value of X got changed because according to it, it is running in isolation.

3. Lost Update Problem-

This problem occurs when multiple transactions execute concurrently and updates from one or more transactions get lost.

Example-

Transaction T1	Transaction T2
R (A)	
W (A)	W (A) Commit
Commit	

Here,

- 1. T1 reads the value of A (= 10 say).
- 2. T2 updates the value to A (= 15 say) in the buffer.
- 3. T2 does blind write A = 25 (write without read) in the buffer.
- 4. T2 commits.
- 5. When T1 commits, it writes A = 25 in the database.

In this example,

- T1 writes the over written value of X in the database.
- Thus, update from T1 gets lost.

NOTE-

- This problem occurs whenever there is a write-write conflict.
- In write-write conflict, there are two writes one by each transaction on the same data item without any read in the middle.

4. Phantom Read Problem-

This problem occurs when a transaction reads some variable from the buffer and when it reads the same variable later, it finds that the variable does not exist.

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Example-

Transaction T1	Transaction T2
R (X)	
	R (X)
Delete (X)	
	Read (X)

Here,

- 1. T1 reads X.
- 2. T2 reads X.
- 3. T1 deletes X.
- 4. T2 tries reading X but does not find it.

In this example,

- T2 finds that there does not exist any variable X when it tries reading X again.
- T2 wonders who deleted the variable X because according to it, it is running in isolation.

Avoiding Concurrency Problems-

- To ensure consistency of the database, it is very important to prevent the occurrence of above problems.
- Concurrency Control Protocols help to prevent the occurrence of above problems and maintain the consistency of the database.

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