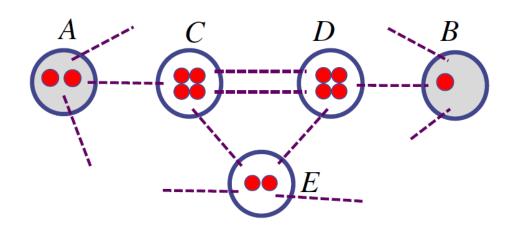
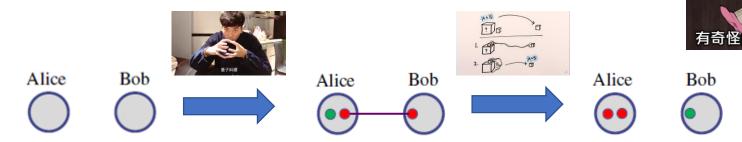
Data Structures Programming Project #1

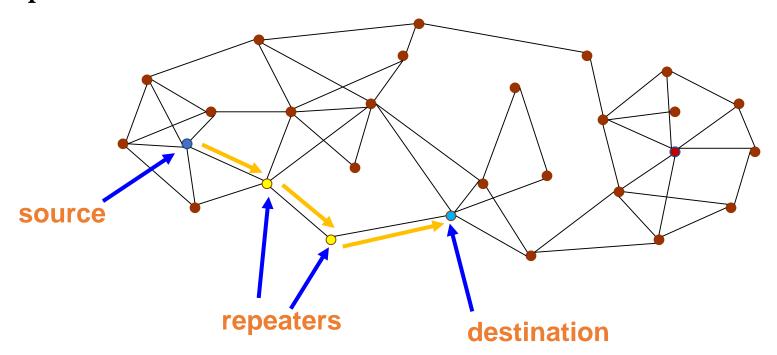
- A quantum network:
- Nodes has a limited number of quantum memories
- Nodes are interconnected with a limited number of quantum channels (e.g., optical fiber)



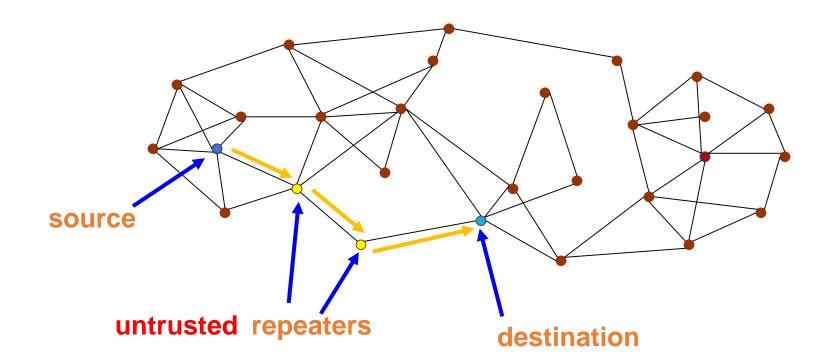
- Entangling (building an entangled link):
 Create an entangled pair between two nodes
- Precondition:
 Two nodes each with a quantum memory are interconnected with a quantum channel
- However, the success probability is $p_c = e^{-\alpha \cdot L}$, where α is a constant and L is the distance between the two nodes



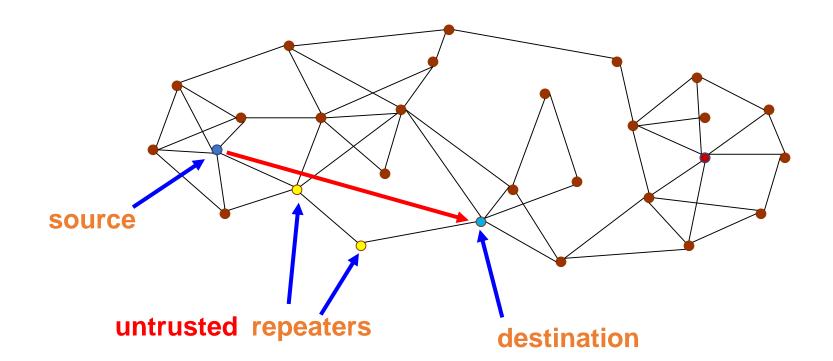
- The two nodes may be distant from each other
- Classical networks:
 Repeaters use store and forward to transmit packets from a source to a destination



- However, the data qubit may visit untrusted repeaters
- It could be destroyed, peeked at, or faked

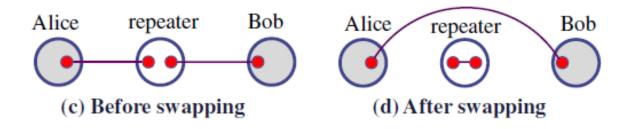


- Can we send the data qubits without letting repeaters know?
- Yes, via Entanglement Swapping

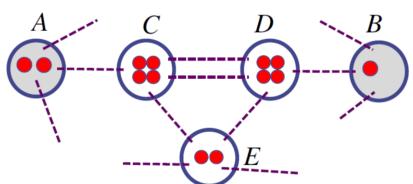


Entanglement Swapping

- Alice has a data qubit for Bob
- Entangling: build the links between Alice and the repeater and between Bob and the repeater
- Swapping: build a long-distance entanglement
- However, the success probability is *q*



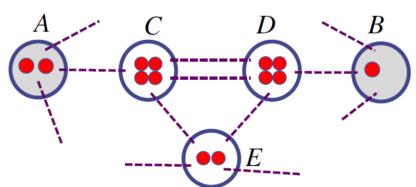
- Given:
- A quantum network with limited resources
- Multiple source-destination (SD) requests
- Constants α and β
- Goal: minimize the expected completion time
- Constraints:
- Find paths for the SD requests over time
- Limited channels
- Limited quantum memories



- Given:
- A quantum network with limited resources
- Multiple source-destination (SD) requests
- Constants α and β



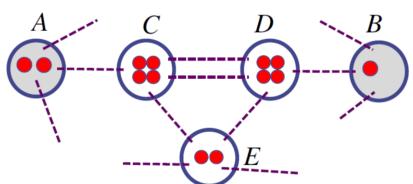
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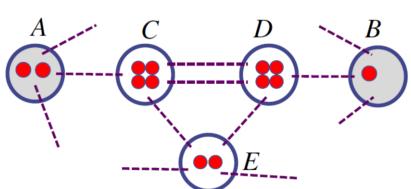


- Constraints:
- Find paths for the SD requests over time
- Limited channels
- Limited quantum memories





- Given:
- A quantum network with limited resources
- Multiple source-destination (SD) requests
- Goal: maximize the number of accepted SD requests
- Constraints:
- Find a path with sufficient resource for each accepted SD request
- Limited channels
- Limited quantum memories



Bad News

- The problem is still NP-hard
- We may not always find the optimal solution in polynomial time
- Alternatively, we aim at a near-optimal solution

- Input:
 - A node-weighted edge-weighted network G = (V, E)
 - Multiple SD requests
- Procedure:
 - Accept or reject each SD request
 - Find a path with sufficient resource for each accepted one
- Output:
 - The accepted SD requests and their paths
- The grade is proportional to the number of accepted SD requests

- Input:
 - A node-weighted edge-weighted network G = (V, E)
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 - A node-weighted edge-weighted network G = (V, E)
 - Multiple SD requests
- Procedure:
 - Accept or reject each SD request
 - Find a path with sufficient resource for each accepted one
- Output:
 - The accepted SD requests and their paths
- The grade is proportional to the number of accepted SD requests
- We have a competition (see next page)

The Competition

- The grade is proportional to # accepted SD requests
- Basic: 75 (deadline)
 - A baseline solution (see the following pages)
- Performance ranking (decided after the deadline)
 - [0%, 50%) (bottom): +0
 - [50%, 75%): + 5
 - \bullet [75%, 90%): + 9
 - [90%, 95%): + 12
 - [95%, 100%] (top): + 15
- Homework assistant (superb deadline)
 - +10

The Competition

The grade

• Basic: 75

A baselii

Performa

• [0%, 50]

• [50%, 75]

• [75%, 90

• [90%, 9!

• [95%, 10

Homework

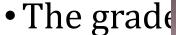
• +10



SD requests

deadline)

The Competition



- Basic: 75
 - A baselii

Performa

- [0%, 50]
- [50%, 75]
- [75%, 90
- [90%, 9!
- [95%, 10
- Homework
 - +10



相信你們在做完作業以後



The Baseline Algorithm

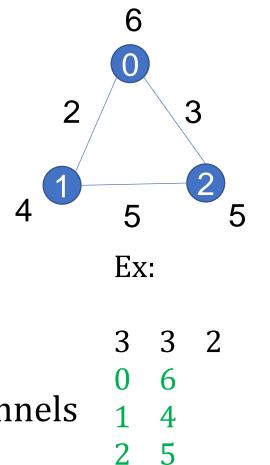
- Sequentially find the shortest path for each input SD request
- If there is a tie, select the one that visits the node with a smaller ID first
 - 10 -> 4 -> 11 -> 15 -> 12 vs 10 -> 4 -> 9 -> 17 -> 12
 - Select 10 -> 4 -> 9 -> 17 -> 12
- Examine whether the shortest path has sufficient resources to accommodate the SD request
 - If yes, then accept the SD request
 - Otherwise, reject the SD request
- Repeat the above actions until all SD requests are examined

Input Sample: use scanf

Format:

#Nodes #Links #Req
NodeID #QuantumMemories
...
LinkID LinkEnd1 LinkEnd2 #Channels

...
ReqID ReqSrc ReqDst



2 1 2 5

Output Sample: use printf

Format:

```
#AccReq
ReqID ReqSrc Rep1 Rep2... ReqDst
...
2
0 0 2
1 1 2
```

Note

- Superb deadline: 9/27 Tue (adjust?)
- Deadline: 10/4 Tue (adjust?)
- Pass the test of our online judge platform
- Submit your code to E-course2
- Demonstrate your code remotely with TA
- C Source code (i.e., only .c)
- Show a good programming style