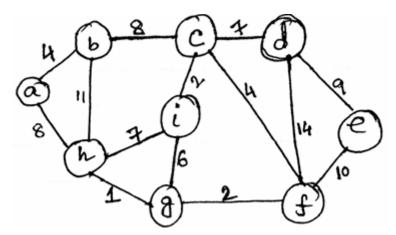
PYQ of (Analysis and Design of Algorithms) 5 Mark

2019

(a) Write Prims algorithm. Construct a minimum spanning tree for the following graph using prim's algorithm.



What is the minimum weight of your spanning tree? Is there any other MST with same weight? 3+4+1+2=10

b) Compare between breadth first search and depth first search. Write pseudo code for BFS. Explain your pseudo-code with suitable example. 2+5+3=10

2021

- a) Describe insertion sort with a suitable example considering at least 5 inputs.
- b) Write the merge sort algorithm. Explain the algorithm with suitable example. 5+5.

2022

- a. State some distinguishing features of insertion sort. Write insertion sort algorithm and trace its execution sequence with a suitable example.
- b. Write KMP algorithm for string processing and trace its execution sequence with a suitable example.

2022-23

- a) Write quick sort algorithm and use it to sort the file 2, 7, 9, 1, 11, 6, 5, 43, 12, 10
- b) What is precondition for applying binary search algorithm? Write binary search algorithm and determine its time complexity.

ANSWERS

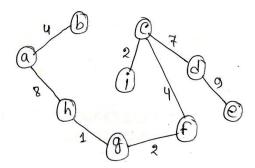
2019

a) Write Prim's algorithm. Construct a minimum spanning tree for the following graph using Prim's algorithm. What is the minimum weight of your spanning tree? Is there any other MST with the same weight?

Prim's Algorithm:

- 1. Initialize a set MST to store the vertices included in the MST.
- 2. Choose an arbitrary starting vertex and add it to MST.
- 3. While MST does not include all vertices:
 - Find the edge with the minimum weight that connects a vertex in MST to a vertex outside
 MST
 - o Add the selected edge and the vertex to MST.
- 4. Repeat until all vertices are included in MST.

MST for the following Graph:



Total Weight of MST:

Is there any other MST with the same weight?

Yes. Prim's algorithm is **not unique** in path selection — another MST with **different edges but same total weight (37)** may exist, depending on the choice of edges when weights are equal. For example, starting from a different vertex or choosing a different edge of equal weight can still yield the same total weight MST.

b) Compare between breadth-first search and depth-first search. Write pseudo code for BFS. Explain your pseudo-code with a suitable example.

Comparison:

Feature	Breadth-First Search (BFS)	Depth-First Search (DFS)
Traversal Method	Level by level	Deep into a branch before backtracking

Feature	Breadth-First Search (BFS)	Depth-First Search (DFS)		
Data Structure Used	d Queue	Stack (or recursion)		
Completeness	Finds shortest path in unweighted graph	s May not find shortest path		
Space Complexity	O(V)	O(V)		
Time Complexity	O(V + E)	O(V + E)		
BFS Pseudo Code:				
pseudo				
function BFS(graph, start):				
create a queue Q				
create a set visited				
enqueue start into Q				
add start to visited				
while Q is not empty:				
node = dequeue from Q				
process node				
for each neighb	or of node:			
if neighbor n	ot in visited:			
enqueue neighbor into Q				
add neighbor to visited				
Explanation:				

- Initialize a queue and a visited set.
- Start from the given node, mark it as visited, and enqueue it.
- While the queue is not empty:
 - o Dequeue a node, process it.
 - o For each unvisited neighbor, enqueue it and mark it as visited.

Example:

For the graph:

A---B---C | | D-----E

Starting from A:

- Queue: [A]
- Dequeue A, process A, enqueue B, D.
- Queue: [B, D]
- Dequeue B, process B, enqueue C.
- Queue: [D, C]
- Dequeue D, process D, enqueue E.

- Queue: [C, E]
- Dequeue C, process C.
- Queue: [E]
- Dequeue E, process E.

Visited Order: A, B, D, C, E

2021

a) Describe insertion sort with a suitable example considering at least 5 inputs.

Insertion Sort Algorithm:

- 1. Start from the second element. Compare it with the elements before it.
- 2. Shift all elements greater than the current element to the right.
- 3. Insert the current element into its correct position.
- 4. Repeat for all elements.

Example:

Consider the array: [5, 2, 9, 1, 5]

- Start with 2:
 - o Compare with 5, shift 5 to the right.
 - Insert 2 at the beginning.
 - o Array: [2, 5, 9, 1, 5]
- Next, 9:
 - No change needed.
 - o Array: [2, 5, 9, 1, 5]
- Next, 1:
 - Compare with 9, shift 9 to the right.
 - Compare with 5, shift 5 to the right.
 - o Compare with 2, shift 2 to the right.
 - o Insert 1 at the beginning.
 - o Array: [1, 2, 5, 9, 5]
- Next, 5:
 - o Compare with 9, shift 9 to the right.
 - o Insert 5.
 - o Array: [1, 2, 5, 5, 9]

b) Write the merge sort algorithm. Explain the algorithm with a suitable example.

Merge Sort Algorithm:

- 1. Divide the array into two halves.
- 2. Recursively sort each half.
- 3. Merge the sorted halves.

Example:

Consider the array: [38, 27, 43, 3, 9, 82, 10]

- Divide: [38, 27, 43, 3] and [9, 82, 10]
- Sort each half:
 - \circ [38, 27, 43, 3] \rightarrow [3, 27, 38, 43]
 - \circ [9, 82, 10] \rightarrow [9, 10, 82]
- Merge:
 - o [3, 27, 38, 43] and [9, 10, 82]
 - Merge to: [3, 9, 10, 27, 38, 43, 82]

Final Sorted Array: [3, 9, 10, 27, 38, 43, 82]

2022

a) State some distinguishing features of insertion sort. Write insertion sort algorithm and trace its execution sequence with a suitable example.

Distinguishing Features of Insertion Sort:

- Simple and intuitive.
- Efficient for small datasets or nearly sorted data.
- Adaptive: performs better on nearly sorted arrays.
- Stable: maintains the relative order of equal elements.
- In-place: requires only a constant amount of additional memory.

Insertion Sort Algorithm:

- 1. Start from the second element.
- 2. Compare it with the elements before it.
- 3. Shift all elements greater than the current element to the right.
- 4. Insert the current element into its correct position.
- 5. Repeat for all elements.

Example Array: [12, 11, 13, 5, 6]

- Step 1: 11 is compared with 12 → shift 12 right → Insert 11
 Array: [11, 12, 13, 5, 6]
- Step 2: 13 > 12 → no shift needed
 Array: [11, 12, 13, 5, 6]
- Step 3: 5 is compared with 13, 12, and 11 → shift them right → insert 5
 Array: [5, 11, 12, 13, 6]
- Step 4: 6 compared with 13, 12, and 11 → shift them → insert 6
 Array: [5, 6, 11, 12, 13]

Final Sorted Array: [5, 6, 11, 12, 13]

b) Write KMP algorithm for string processing and trace its execution sequence with a suitable example.

KMP Algorithm (Knuth-Morris-Pratt):

- Used for efficient pattern matching in a string.
- Avoids rechecking characters after a mismatch using a prefix table (LPS array).

Steps:

- 1. Preprocess the pattern to create the LPS (Longest Prefix Suffix) array.
- 2. Start matching the pattern from left to right with the text.
- 3. On mismatch, use LPS array to avoid rechecking previous characters.

Example:

Text = "ABABDABACDABABCABAB"
Pattern = "ABABCABAB"

LPS Array for Pattern:

Index 0 1 2 3 4 5 6 7 8

Value 0 0 1 2 0 1 2 3 4

Matching Process:

- Start comparing text and pattern.
- After a few shifts, match found at index 10 in text.

Output: Pattern found at index 10.

a) Write quick sort algorithm and use it to sort the file 2, 7, 9, 1, 11, 6, 5, 43, 12, 10

Quick Sort Algorithm:

- 1. Choose a pivot.
- 2. Partition array into two halves:
 - o Elements less than pivot go left.
 - o Elements greater go right.
- 3. Recursively sort both halves.

Example: Sort [2, 7, 9, 1, 11, 6, 5, 43, 12, 10]

Choose pivot: 10

• Partition:

Left: [2, 7, 9, 1, 6, 5] Right: [11, 43, 12]

• Recursively sort left:

o Pivot: 5 → Left: [2, 1], Right: [7, 9, 6]

 \circ Sort $[2, 1] \rightarrow [1, 2]$

 \circ Sort [7, 9, 6] → Pivot 6 → [7, 9] → Final [6, 7, 9]

o Combine: [1, 2, 5, 6, 7, 9]

Sort right [11, 43, 12]:

o Pivot: 12 → Left: [11], Right: [43] → Combine: [11, 12, 43]

Final Sorted Array: [1, 2, 5, 6, 7, 9, 10, 11, 12, 43]

b) What is precondition for applying binary search algorithm? Write binary search algorithm and determine its time complexity.

Precondition:

• The array **must be sorted** in ascending or descending order.

Binary Search Algorithm:

pseudo

BinarySearch(arr, key):

low = 0

high = length(arr) - 1

```
while low <= high:
  mid = (low + high) / 2
  if arr[mid] == key:
    return mid
  else if arr[mid] < key:
    low = mid + 1
  else:
    high = mid - 1
return -1 // not found
Time Complexity:
```

Best Case: O(1) (key found at middle)

Worst Case: O(log n)

• Average Case: O(log n)