

第一题

200P 9/12

$$1. \text{对于实矩阵 } (c_{ij}) = \sum_{k=1}^n a_{ik} b_{kj}$$

计算这一行元素的加法/减法是 $n-1$. 计算需要的乘法是 n .
 总加加法是 $n^2(n-1) = n^3$ 乘法是 n^3 . 都是 $O(n^3)$ 并且没有 $O(n^3)$ 的项.

$$\text{对于复矩阵 } (c_{ij}) = \sum_{k=1}^n (Re(a_{ik}) + Im(a_{ik})) (Re(b_{kj}) + Im(b_{kj}))$$

$$= (Re(a_{ik}) \cdot Re(b_{kj}) - Im(a_{ik}) Im(b_{kj})) i$$

计算这一行元素的加法/减法是 $4n-2$. 乘法是 $4n$.

总加加法/减法是 $(4n-2)n^2 \approx 4n^3$. 乘法是 $4n^3$, 故有 $(4n-2)n^2 \approx 4n^3$

故用计算实矩阵运算方法计算复矩阵运算, 而且加法/减法的乘法次数
 几近相同.

第二题

方法一

1. 方法二: 利用LU分解. 对于一般对称矩阵 A , 有单位下三角阵

$$\text{使得 } A = \tilde{L}\tilde{U}, \text{ 其中 } \tilde{U} = D^{-1}U, \text{ 且 } A = \tilde{L}D\tilde{U}, A^T = \tilde{U}^T D^{-1} \tilde{L}^T$$

$$\text{其中 } D = \text{diam}(u_{11}, u_{22}, \dots, u_{nn})$$

$$\text{于是 } \tilde{L}\tilde{D}\tilde{U} = \tilde{U}^T D^{-1} \tilde{L}^T, \text{ 且 } D^{-1} \tilde{U}^T \tilde{L}^T = \tilde{U}^T \tilde{U}^{-1}$$

右边是上三角阵左边是下三角阵. 于是有 $\tilde{U}^T \tilde{U}^{-1} = I$. $\tilde{U} = \tilde{U}^T$

$$\text{进而 } A = \tilde{L}D\tilde{U}$$

算法的目标是求出 \tilde{U} 的各元素, 以及 U 的上三角阵元素.

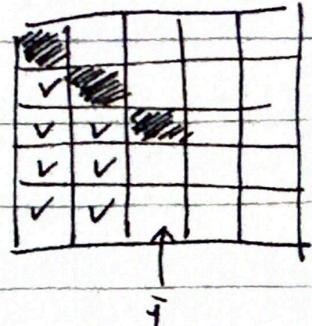
方法二

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方法二：利用待求解法。已知 $A = LDL^T$, 且 $L = \begin{bmatrix} l_{11} & & \\ & l_{22} & \\ & & \ddots \\ & & & l_{nn} \end{bmatrix}$ 。
有 $a_{ij} = \sum_{p=1}^i l_{ip}l_{pj}$ ($1 \leq j \leq i \leq n$)。

计算 $a_{11} \rightarrow$ 计算 L 的第一列 $\rightarrow l_{12} \rightarrow$
 $\text{[第二行]} \rightarrow \dots \rightarrow l_{nn}$.



文件 T2.py

```
import numpy as np
from math import sqrt

# Method One: get Cholesky factorization through LU factorization
def cholesky_1(A):
    n = A.shape[0]
    for i in range(n):
        for j in range(i+1, n):
            factor = A[j, i] / A[i, i]
            A[j, i] = factor
            A[j, i+1:] = A[j, i+1:] - factor * A[i, i+1:]
    diag_elements = np.diag(A)
    diag_matrix = np.diag(np.sqrt(diag_elements))

    L_tilde = np.tril(A, k = -1)
    np.fill_diagonal(L_tilde, 1)
    L = np.dot(L_tilde, diag_matrix)
    return L
```

```

# Method Two: get Cholesky factorization through undetermined coefficient
def cholesky_2(A):
    n = A.shape[0]
    for i in range(n):
        A[i, i] = sqrt(A[i, i])
        A[i+1:, i] = A[i+1:, i] / A[i, i]
        for j in range(i+1, n):
            A[j:, j] = A[j:, j] - A[j, i] * A[j:, i]
    L = np.tril(A)
    return L

def check_factorization(A, fac_fun):
    L = fac_fun(A)
    A_result = np.dot(L, L.T)
    print(f"{fac_fun.__name__} gets \n{A_result}")
    return A_result

# test
A = np.array([[1,2,3],
              [2,7,6],
              [3,6,10]], dtype = np.float64)

A_copy = np.copy(A)

check_factorization(A, cholesky_1)
check_factorization(A_copy, cholesky_2)

```

第三题

3. 设 $A = [a_{ij}]$ 是 $n \times n$ 的矩阵, 经一次高斯消元后为 $\begin{bmatrix} a_{11} & a_{12}^T \\ 0 & A_2 \end{bmatrix}$

为方便表示, 设 $A_2 = [\tilde{a}_{ik}]$. 矩阵的角标沿用 A 中角标.

例如 在实际上是 A 的 a_{22} 经一次高斯消元后的值. $A_2 - \{j\}$ -列的元素是 \tilde{a}_{2j} .

$$|\tilde{a}_{ii}| = \left| a_{ii} - \frac{a_{i1}}{a_{11}} \cdot a_{12} \right| = |a_{ii}| - \frac{|a_{i1}|}{|a_{11}|} \cdot |a_{12}| \quad ①$$

$$|\tilde{a}_{ik}| = \left| a_{ik} - \frac{a_{i1}}{a_{11}} \cdot a_{1k} \right| \leq |a_{ik}| + \frac{|a_{i1}|}{|a_{11}|} \cdot |a_{1k}| \quad (2 \leq k \leq n, k \neq i) \quad ②$$

$$\text{左边 } ① - ② : \left(|a_{ii}| - \sum_{\substack{k=2 \\ k \neq i}}^n |a_{ik}| \right) - \frac{|a_{i1}|}{|a_{11}|} \left(\sum_{\substack{k=2 \\ k \neq i}}^n |a_{ik}| + |a_{12}| \right).$$

$$> |a_{ii}| - \frac{|a_{i1}|}{|a_{11}|} \left(\sum_{\substack{k=2 \\ k \neq i}}^n (|a_{ik}| + |a_{1k}|) \right)$$

$$\text{由 } |a_{ii}| > \sum_{\substack{k=2 \\ k \neq i}}^n |a_{ik}| + |a_{12}| = \sum_{\substack{k=1 \\ k \neq i}}^n |a_{ik}|$$

$$\text{可得 } ① - \sum_{\substack{k=2 \\ k \neq i}}^n |a_{ik}| > 0 \quad \text{且 } |\tilde{a}_{ii}| > |\tilde{a}_{ik}|, k \neq i \text{ 得证.} \quad (i=1, \dots, n)$$

第四题

文件 T4.py

```
import numpy as np
import time
import matplotlib.pyplot as plt

def partial_pivoting_gaussian_elimination(A):
    copy_A = np.copy(A)
    n = copy_A.shape[0]

    for i in range(n):
        max_row_i = np.argmax(np.abs(A[i:, i])) + i
        if max_row_i != i:
            A[[i, max_row_i], :] = A[[max_row_i, i], :]

        for j in range(i+1, n):
            factor = A[j, i] / A[i, i]
            A[j, i:] = A[j, i:] - factor * A[i, i]

    return copy_A

def complete_pivoting_gaussian_elimination(A):
    copy_A = np.copy(A)
    n = copy_A.shape[0]

    for i in range(n):

        max_row_i = np.argmax(np.abs(A[i:, i])) + i
        if max_row_i != i:
            A[[i, max_row_i], :] = A[[max_row_i, i], :]

        max_col_i = np.argmax(np.abs(A[i, i:])) + i
        if max_col_i != i:
            A[:, [max_col_i, i]] = A[:, [i, max_col_i]]

        for j in range(i+1, n):
            factor = A[j, i] / A[i, i]
            A[j, i:] = A[j, i:] - factor * A[i, i:]
```

```

return copy_A

def generate_test_matrix(n):
    A = np.random.randn(n,n)
    return A


def measure_execution_time(n, strategy):
    A = generate_test_matrix(n)
    start_time = time.time()
    strategy(A)
    end_time = time.time()
    return end_time - start_time


def plot_log_log(strategy):
    dimention = np.arange(1000, 2001, 10)
    times = []

    for n in dimention:
        exec_time = measure_execution_time(n, strategy)
        times.append(exec_time)

    log_dimentions = np.log10(dimention)
    log_times = np.log10(times)

    plt.figure(figsize=(8, 6))
    plt.plot(log_dimentions, log_times, marker='o', linestyle='--')
    plt.xlabel('log(matrix dimension)', fontsize=12)
    plt.ylabel('log(execution time)', fontsize=12)
    plt.title(f'Log(matrix dimension)-log(execution time) for {strategy.__name__}')
    plt.grid(True, which='both', ls='--')
    plt.show(block=False)

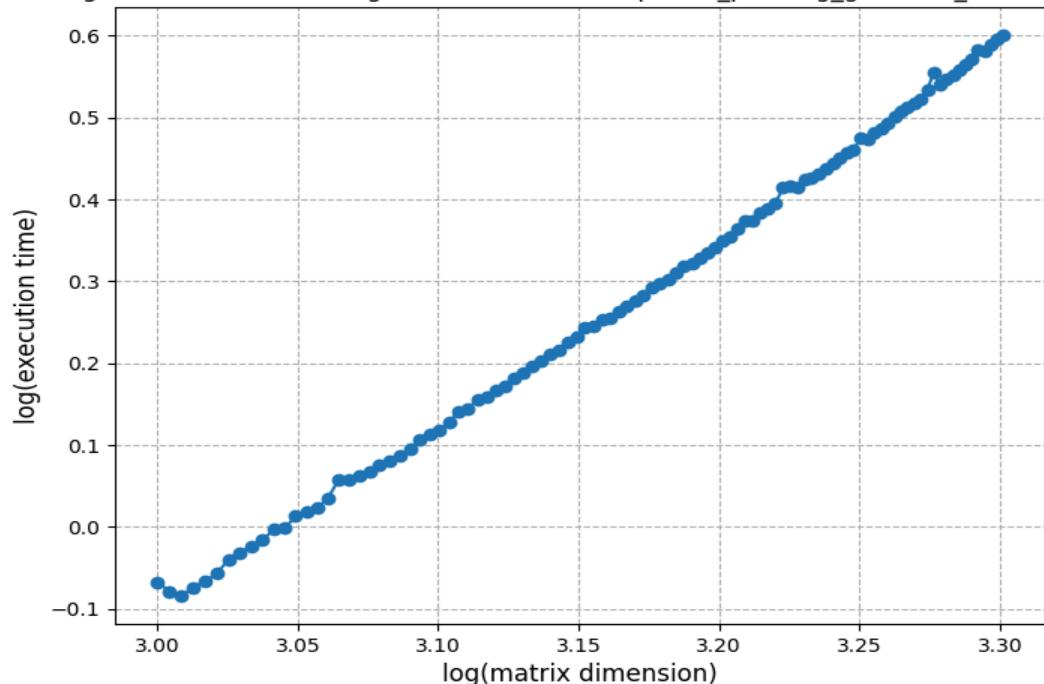
plot_log_log(partial_pivoting_gaussian_elimination)
plot_log_log(complete_pivoting_gaussian_elimination)
plt.show()

```

绘图结果

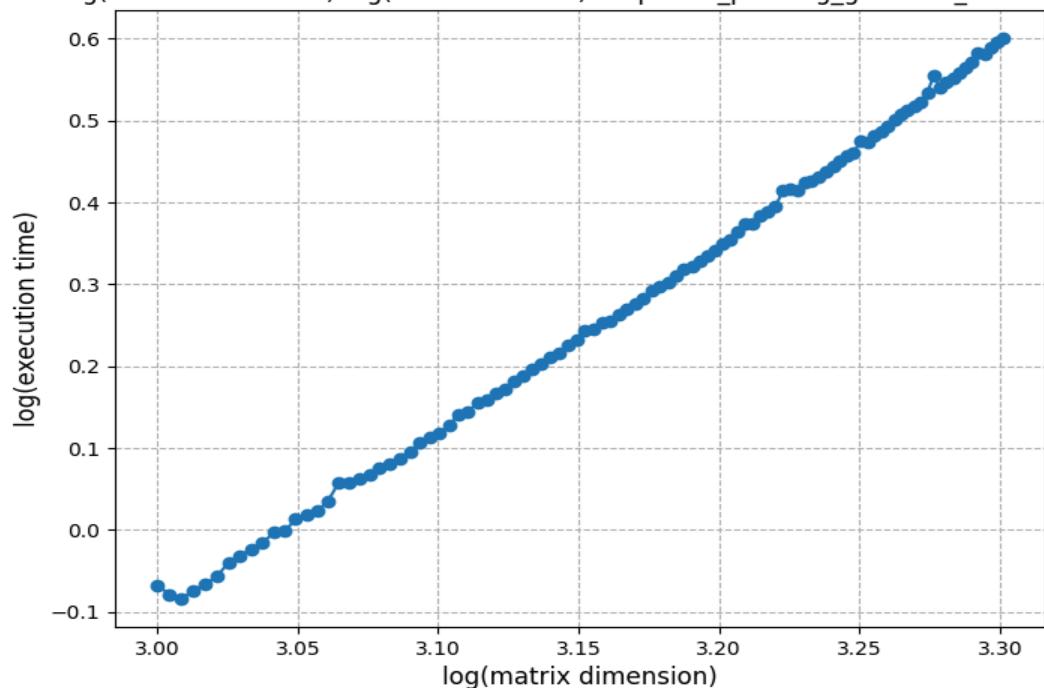
部分主元高斯消元

Log(matrix dimension)-log(execution time) for partial_pivoting_gaussian_elimination



全主元高斯消元

Log(matrix dimension)-log(execution time) for partial_pivoting_gaussian_elimination



第五题

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$$\text{设 } y_k = Ax, A = [a_{ij}], x = [x_1, x_2, \dots, x_n]^T$$

$$n \quad a_{ij} = [a_{1j}, a_{2j}, \dots, a_{nj}]^T$$

$$\text{则有 } y_k = \sum_{i=1}^m x_i a_{ik}, k=1, \dots, m$$

$$f_l(y_k) = \left\{ \begin{array}{l} [x_1 a_{1k}(1+\epsilon_{1k}) + x_2 a_{2k}(1+\epsilon_{2k})] (1+f_{1k}) + x_3 a_{3k}(1+\epsilon_{3k}) \end{array} \right\} (1+f_{2k}) + \dots \\ = \sum_{i=1}^n x_i a_{ik}(1+\epsilon_{ik}) \prod_{j=i+1}^{n-1} (1+f_{jk}) \quad - f_{0k}=0 \\ \approx \sum_{i=1}^n x_i a_{ik}(1+\epsilon_{ik}) \quad \text{其中 } 1+\epsilon_{ik} = (1+\delta_{ik}) \prod_{j=i+1}^{n-1} (1+f_{jk}).$$

目标是估计 ϵ_{ik} 大小, 记 u 是机器精度满足 $f_l(x) = x(1+f), f_l \leq u$.

$$\text{由于 } |r_{ik}|, |f_{jk}| \leq u, \text{ 故 对 } i=1, (1-u)^n \leq 1+\epsilon_{ik} \leq (1+u)^n \quad ①$$

$$\text{对 } i=2, \dots, n \quad (1-u)^{n-i+2} \leq 1+\epsilon_{ik} \leq (1+u)^{n-i+2} \quad ②$$

$$\text{对 } f_l(y_k) = \sum_{i=1}^n x_i a_{ik}(1+\epsilon_{ik}), a_{ik} \in A \text{ 的作用, } a_{ik}\epsilon_{ik} \text{ 是 } \epsilon_{ik} \text{ 的倍数.}$$

$$\text{记 } \beta = [\beta_1, \beta_2, \dots, \beta_n], \beta_i = [b_{1i}, b_{2i}, \dots, b_{ni}]$$

$$\text{故 } a_{ik}\epsilon_{ik} = b_{ik} \text{ 是 } \beta \text{ 中第 } k \text{ 行第 } i \text{ 列的元素. 由 } ①② \text{ 获得以下估计}$$

$$\Rightarrow a_{ik} \geq 0 \text{ 且. } b_{ik} \text{ 为: 对 } i=1, [(1-u)^n - 1] a_{ik} \leq b_{ik} \leq [(1+u)^n - 1] a_{ik}$$

$$\text{对 } i \geq 2, [(1-u)^{n-i+2} - 1] a_{ik} \leq b_{ik} \leq [(1+u)^{n-i+2} - 1] a_{ik}$$

$$\Rightarrow a_{ik} \leq 0 \text{ 且, } b_{ik} \text{ 为: 对 } i=1, [(1+u)^n - 1] a_{ik} \leq b_{ik} \leq [(1-u)^n - 1] a_{ik}$$

$$\text{对 } i \geq 2, [(1+u)^{n-i+2} - 1] a_{ik} \leq b_{ik} \leq [(1-u)^{n-i+2} - 1] a_{ik}$$

$$\text{该数系中: } F = \{0\} \cup \{f, f = \pm 1, d_1 d_2 \dots d_n \times 2^J, 0 \leq d_i < 2, 1 \leq J \leq u\} \cup \{\pm \infty\}$$

假定 $x > 0$, $\exists \alpha$ 满足 $2^\alpha \leq x < 2^{\alpha+1}$, 此时有:

$$|f_l(x) - x| \leq \frac{1}{2} 2^{\alpha+1} = \frac{1}{2} 2^\alpha \cdot 2^{-1} \leq \frac{1}{2} x \cdot 2^{-1} - \frac{|f_l(x) - x|}{x \cdot 2^{-1}} \cdot (2+1)$$

此时可以令 $u = 2^{-(\alpha+1)}$, 取浮点数的尾数位数为 1.