Mathematics and numerics for data assimilation and state estimation – Lecture 11





Summer semester 2020

Overview

Bayesian inversion and optimization

2 Entropy and Kullback-Leibler divergence

Summary of lecture 10

■ Weak convergence of distributions $\mathbb{P}_k \Rightarrow \mathbb{P}$.

Bayesian inversion in the linear-Gaussian setting

$$Y = AU + \eta$$
, π_U, π_η Gaussian pdfs.

■ Consistency of posterior $\pi(u|y)$ in small noise limit when η "disappears", when A is overdetermined, determined and underdetermined.

Overview

Bayesian inversion and optimization

2 Entropy and Kullback-Leibler divergence

Problem setting

$$Y = G(U) + \eta \tag{1}$$

with $G: \mathbb{R}^d \to \mathbb{R}^k$, $\eta \sim \pi_{\eta}$, $U \sim \pi_U$ and $\eta \perp U$.

For an observation Y = y, we obtained

$$\pi(u|y) \propto \pi_{\eta}(y - Au)\pi_{U}(u)$$

And in the linear-Gaussian setting

$$\pi(u|y) \propto \exp\left(-\frac{1}{2}|y - G(u)|_{\Gamma}^2 - \frac{1}{2}|u - \hat{m}|_{\hat{C}}^2\right) = \exp(-\mathsf{J}(u))$$

where, decomposing into loss and regularization terms,

$$\mathsf{L}(u) := -\log(\pi_{\eta}(y - G(u))) \quad \text{and} \quad \mathsf{R}(u) := -\log(\pi_{U}(u))$$
 and
$$\underbrace{\mathsf{J}(u)}_{\mathsf{Objective fcn}} := \mathsf{L}(u) + \mathsf{R}(u) \tag{2}$$

Assuming π_{η} , $\pi_{U} > 0$, we extend the notation (2) to general settings:

$$\pi(u|y) \propto \pi_{\eta}(y - Au)\pi_{U}(u) = \exp(-\mathsf{J}(u)) = \exp(-\mathsf{L}(u) - \mathsf{R}(u)).$$

MAP estimators and Tikhonov regularization

Maximizing the posterior is equivalent to minimizing the objective function:

$$u_{MAP}[\pi(\cdot|y)] = \arg\max_{u \in \mathbb{R}^d} \pi(u|y) =$$
 $= \arg\min_{u \in \mathbb{R}^d} J(u)$

■ In Gaussian setting, with $U|Y = y \sim N(m, C)$ and $U \sim N(0, \lambda^{-1}I)$,

$$u_{MAP}=m=rg\min_{u\in\mathbb{R}^d}rac{1}{2}|y-G(u)|_\Gamma^2+rac{\lambda}{2}|u|^2.$$

■ This corresponds to Tikhonov regularization. Unique, closed form solution in linear setting G(u) = Au.

Laplace-distributed prior and LASSO regression

 Alternatively, consider the prior with iid Laplace-distributed components

$$\pi_U(u) = \prod_{i=1}^d \pi_{U_i}(u_i) \propto \prod_{i=1}^d e^{-\lambda |u_i|} = e^{-\lambda |u|_1}$$

where

$$|u|_p := \Big(\sum_{i=1}^d |u_i|^p\Big)^{1/p}, \qquad p > 0.$$

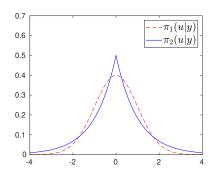
This yields

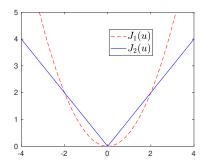
$$\mathsf{R}(u) \propto \lambda |u|_1$$
 and $u_{MAP} = \arg\min_{u \in \mathbb{R}^d} \frac{1}{2} |y - G(u)|_\Gamma^2 + \lambda |u|_1$

which corresponds to lasso (least absolute shrinkage and slection operator) regression.

Generally, lasso has no closed-form solution, but a solution is typically attainable. It tends to produce more sparse solutions than Tikhonov.

Posterior setting with R \gg L and regularizers so that approximately $\pi_1(u|y) \propto \exp(-|u|^2/2)$ and $\pi_2(u|y) \propto \exp(-|u|_1)$.





Attainability of u_{MAP}

Theorem 1

Assume that the objective fcn $J: \mathbb{R}^d \to \mathbb{R}$ is bounded from below, continuous and that $J(u) \to \infty$ as $|u| \to \infty$. Then J attains its infimum, which implies that

$$u_{MAP}[\pi(\cdot|y)]$$
 is attained for $\pi(u|y) \propto \exp(-J(u))$.

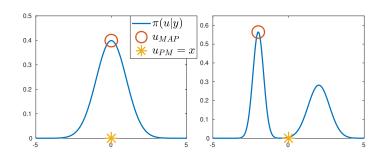
Sufficient conditions for attainable u_{MAP} :

- lacksquare $G \in C(\mathbb{R}^d, \mathbb{R}^k)$ and $\eta \sim N(0, \Gamma)$,
- R(u) = $\lambda |u|_p^p$ for any $\lambda, p > 0$ (as this implies J(u) $\to \infty$ as $|u| \to \infty$).

Examples of the MAP performing poorly

- "All happy families are alike; each unhappy family is unhappy in its own way."
 Leo Tolstoy, in Anna Karenina
- Paraphrasing: "All unimodal densities are alike; each multimodal density is multimodal in its own way"

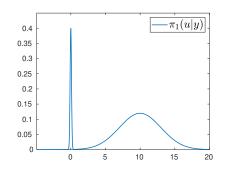
In Lecture 7 we already saw that u_{MAP} can be of limited value for bimodal densities:

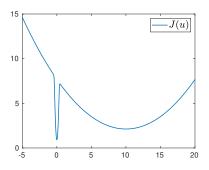


Slab-spike figure

For

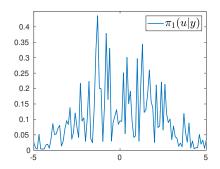
$$\pi(u|y) = \frac{\exp(-|u|^2/0.02) + 0.3\exp(-|u-10|^2/18)}{\sqrt{2\pi}}$$

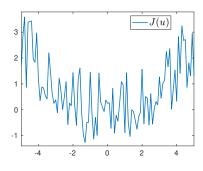




Low-regularity objective function

```
normalF = @(x) (x).^2/10;
objective = normalF(x)+1.5*(1-2*rand(size(x)));
posterior = exp(-objective);
posterior = exp(-objective)/(trapz(posterior)*dx);
```





And low-regularity in higher dimensions ...



Figure: Photo by Michel Royon / Wikimedia Commons

Overview

1 Bayesian inversion and optimization

2 Entropy and Kullback-Leibler divergence

Low-rank approximations of posteriors

- We have seen that one-parameter/vector compression of a posterior, like MAP or posterior mean, may provide little information.
- Natural next step: Extend the compressed representations of posteriors to best fitting in a class of candidate densities:

$$p^* = \arg\inf_{p \in \mathcal{A}} d(p, \pi(\cdot|y))$$

for some $d:\mathcal{M}\times\mathcal{M}\to[0,\infty)$

Here we will restrict ourselves to

$$\mathcal{A} = \{ p = PDF(N(\mu, C)) \mid \mu \in \mathbb{R}^d \text{ and } C \in \mathbb{R}^{d \times d} \text{ and pos definite} \}$$

which can be viewed as a two-parameter (two-moment) compression of a posterior.

Kullback-Leibler divergence

Definition 2 (K-L divergence)

■ For positive discrete measures: Let

$$\mathcal{P}_+ = \{ \text{Probability measures on } A \mid \mathbb{P}(x), \mathbb{Q}(x) > 0 \text{ for all } x \in A \}.$$

For all $\mathbb{P}, \mathbb{Q} \in \mathcal{P}_+$,

$$d_{\mathcal{KL}}(\mathbb{P}||\mathbb{Q}) := \sum_{x \in A} \log \left(\frac{\mathbb{P}(x)}{\mathbb{Q}(x)} \right) \mathbb{P}(x).$$

For positive pdfs on \mathbb{R}^d : Let

For all $\pi, p \in \mathcal{M}_+$

$$d_{\mathcal{KL}}(\pi||p) := \int_{\mathbb{R}^d} \log\left(\frac{\pi(x)}{p(x)}\right) \pi(x) dx = \mathbb{E}^{\pi}\left[\log\left(\frac{\pi}{p}\right)\right]$$

 $\mathcal{M}_+ := \{ \pi \in \mathcal{M} \mid \pi(x) > 0 \mid \forall x \in \mathbb{R}^d \}.$

Properties of the K-L divergence

For all $\pi, p \in \mathcal{M}_+$, it holds that $d_{KL}(\pi||p) \in [0, \infty]$ (similar result holds for prob measures).

Example of infinite K-L divergence:

$$p(x) \propto e^{-|x|}, \quad \pi \propto (1+|x|)^{-2}, \quad x \in \mathbb{R}$$

Then

$$d_{KL}(\pi||p) = \int_{\mathbb{R}} \log\left(\frac{\pi(x)}{p(x)}\right) \pi(x) dx$$

$$= C \int_{\mathbb{R}} \left(\log(\pi(x)) - \log(p(x))\right) \pi(x) dx$$

$$= C \int_{\mathbb{R}} \frac{-2\log((1+|x|)) + |x|}{(1+|x|)^2} \pi(x) dx$$

$$= \infty.$$

Properties of the K-L divergence

 d_{KL} is not a metric; neither does it saitisfy the triangle inequality nor is it symmetric in its arguments.

Example: Let
$$A = \{1, 2, 3\}$$
 and $\mathbb{P}(1) = \mathbb{P}(2) = \mathbb{P}(3) = 1/3$ and $\mathbb{Q}(1) = 1/2$, $\mathbb{Q}(2) = 1/3$, $\mathbb{Q}(3) = 1/6$. Then

$$d_{KL}(\mathbb{P}||\mathbb{Q}) = \sum_{x_i \in A} \log \left(\frac{\mathbb{P}(x_i)}{\mathbb{Q}(x_i)}\right) \mathbb{P}(x_i)$$
$$= \frac{\log(2/3) + \log(1) + \log(2)}{3} \approx 0.0959$$

while

$$d_{\mathit{KL}}(\mathbb{Q}||\mathbb{P}) = \frac{3\log(3/2) + 2\log(1) + \log(1/2)}{6} \approx 0.0872$$

Properties of the K-L divergence

K-L divergence has natural applications in information theory and thermodynamics.

■ In Bayesian inference, for a prior π_U and a posterior $\pi(\cdot|y)$, $d_{KL}(\pi(\cdot|y), \pi_U)$ is a measure of the information gain of replacing the prior by the posterior.

■ The logarithm base in the definition of K-L divergence is flexible; use what is most suitable for the application (here, log denotes the natural logarithm).

Lemma 3 (Lower bounds for K-L divergence, (SST 4.2))

For any $pi, p \in \mathcal{M}_+$ it holds that

$$d_H(\pi,p)^2 \leq rac{1}{2} d_{KL}(\pi||p)$$
 and $d_{TV}(\pi,p)^2 \leq d_{KL}(\pi||p)$.

Proof of first inequality:

$$d_H(\pi,p)^2 = \frac{1}{2} \int_{\mathbb{R}^d} (\sqrt{\pi} - \sqrt{p})^2 dx$$

$$= \int_{\mathbb{R}^d} \left(1 - \sqrt{\frac{p}{\pi}}\right) \pi \ dx \le -\frac{1}{2} \int_{\mathbb{R}^d} \log \left(\frac{p}{\pi}\right) \pi \ dx = \frac{1}{2} d_{\mathsf{KL}}(\pi||p).$$

where we used that

$$1-\sqrt{x} \leq -rac{1}{2}\log(x) \quad orall x \in [0,\infty].$$

Comments

■ Second inequality follows from $d_{TV}(\pi, p) \leq \sqrt{2} d_H(\pi, p)$.

■ The lemma implies that K-L divergence is point/density separating: For all $\pi, p \in \mathcal{M}_+$,

$$d_{KL}(\pi||p) \geq 0$$

and

$$d_{KL}(\pi||p) = 0 \iff p = \pi.$$

(Similar for measures.)

Entropy in information theory

Suppose you want to transmit a very long text encoded in some alphabet, e.g., $A = \{a, b, c, d, e\}$,

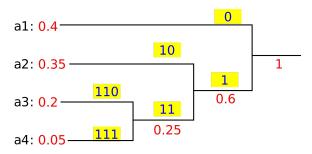
TEXT= "abbedeeeedcaababecbddaeedeccabe..."

and that

- the data-transmission problem can to good approximation be viewed as transmitting a sequence iid characters drawn with relative frequencies $\mathbb{P}(a)$, $\mathbb{P}(b)$ etc.
- you want to send the text over a digital communication channel with alphabet $\{0,1\}$. Hence, each letter in your original alphabet must be replaced with a codeword, e.g. $a=101,\ b=111$, and your want the digitally encoded text to be as short as possible.
- Core idea: assign shortest codeword to most frequent letter in the text, second shortest codeword to ... (then there is a subtle issue with uniqueness/reversibility of encoding).

Huffman encoding

Input alphabet: $A = \{a1, a2, a3, a4\}.$



Letter frequency: $\mathbb{P}(a1) = 0.4$, $\mathbb{P}(a2) = 0.35$ etc

Digital codewords: a1 = 0, a2 = 10, etc

NB! A shorter encoding is possible: a1 = 0, a2 = 1, a3 = 10 and a4 = 11 but this encoding is, unlike Huffman's, not uniquely reversible, since it is not injective when applied to strings:

$$a4 \mapsto 11 \quad a2a2 \mapsto 11$$

Shannon's approach

Shannon relates the text-frequency of a letter to the information content:

Definition 4 (Information content of a character)

For an event/character a which occurs with probability $\mathbb{P}(a)$ we define its information content by

$$I(a) := -\log_2(\mathbb{P}(a))$$

Idealized motivation: if there are $1/\mathbb{P}(a)$ many independent events, each occurring with probability $\mathbb{P}(a)$, how many bits do I need to distinguish all these events when encoded in $\{0,1\}$?

Example Alphabet $A = \{a, b, c, d, e\}$ with uniform letter probability 1/5. Then at least $-\lceil \log_2(1/5) \rceil = 3$ bits are needed to distinguish the letters/events.

Shannon entropy

Generalization: Information content straightforwardly generalizes from a character to any text string ${\cal B}$

$$I(B) := -\log_2(\mathbb{P}(B))$$

where we recall that letter sequences, e.g., B = abeba, are assumed to consist of iid characters,

$$\mathbb{P}(abeba) = \mathbb{P}(a)\mathbb{P}(b)\mathbb{P}(e)\mathbb{P}(b)\mathbb{P}(a)$$

Lemma 5 (Information content of independent events)

Let B and C denote two independent events (i.e., text strings), then the information content of B and C is additive

$$I(BC) := I(B) + I(C)$$

Verification for two-character sequence: Consider basic events B=a and C=b. Then

$$I(ab) = -\log_2(\mathbb{P}(ab)) = -\log_2(\mathbb{P}(a)\mathbb{P}(b)) = I(a) + I(b)$$

Shannon entropy

Question: Given a text encoded in the alphabet $A = \{a_1, \ldots, a_n\}$ with relative frequencies $\{\mathbb{P}(a_k)\}_k$, and a digital encoding representing the letter a_k by $I(a_k)$ bits (we allow fractional-bit encoding in this thought experiment) then if the original text consists of $N \gg 1$ characters, how long does the digitally encoded text become?

Answer:

$$N \times \text{mean num of bits for single } A\text{-character} = N \sum_{k=1}^{n} I(a_k) \mathbb{P}(a_k)$$

Introducing the information content rv

$$\mathit{I}_{\mathbb{P}}(\mathit{a}) := -\log_2(\mathbb{P}(\mathit{a})), \quad (\mathit{I}_{\mathbb{P}} : \mathit{A} \to [0, \infty], \text{ and } \mathbb{P}_{\mathit{I}_{\mathbb{P}}}(\mathit{I}_{\mathbb{P}}(\mathit{a})) = \mathbb{P}(\mathit{a})),$$

we may associate the above with the expected information content/Shannon entropy

$$\mathbb{E}^{\mathbb{P}}[I_{\mathbb{P}}] = \sum_{k=1}^{n} I_{\mathbb{P}}(a_k) \mathbb{P}(a_k) = -\sum_{k=1}^{n} \log_2(\mathbb{P}(a_k)) \mathbb{P}(a_k)$$

Comparison of encoding methods

Assume that a text encoded in $A = \{a_1, \ldots, a_n\}$ has true relative frequencies $\{\mathbb{P}(a_k)\}$, but that

- lacksquare you only have an approximation of the relative frequencies $\{\mathbb{Q}(a_k)\}$
- and that given \mathbb{Q} , your encoding in $\{0,1\}$ is optimal, meaning it uses $I_{\mathbb{Q}}(a_k) = -\log_2(\mathbb{Q}(a_k))$ bits to encode the letter a_k .

K-L divergence is a comparison of efficiency $\mathbb{Q}\text{-}$ vs $\mathbb{P}\text{-encoding:}$

[mean
$$\mathbb Q$$
-bits in encoded A -char] — [mean $\mathbb P$ -bits in encoded A -char]
$$=\sum_{k=0}^n (I_{\mathbb Q}(a_k)-I_{\mathbb P}(a_k))\mathbb P(a_k)$$

$$=\sum_{k=1}^{n}(\log_2(\mathbb{P}(a_k))-\log_2(\mathbb{Q}(a_k))\mathbb{P}(a_k)$$

$$=\sum_{k=1}^n \log_2\Big(rac{\mathbb{P}(a_k)}{\mathbb{Q}(a_k)}\Big)\mathbb{P}(a_k)=d_{KL}(\mathbb{P}||\mathbb{Q})$$

Best encoding in a set

Given a collection of encodings, a natural task is to find the most efficient one:

$$\mathbb{Q}^* = \arg\min_{\mathbb{Q} \in \mathcal{A}} d_{\mathit{KL}}(\mathbb{P}||\mathbb{Q}).$$

Example: Let $A = \{a, b, c, d, e\}$ and $\mathbb{P}(a) = \mathbb{P}(b) = \ldots = \mathbb{P}(d) = 1/5$, and $A = \{\mathbb{Q}_1, \mathbb{Q}_2\}$ with

$$\mathbb{Q}_1(a) = \mathbb{Q}_1(b) = \mathbb{Q}_1(c) = \mathbb{Q}_1(d) = 2^{-4}, \quad \mathbb{Q}_1(e) = 3/4$$

and

$$\mathbb{Q}_2(a) = \mathbb{Q}_2(b) = \mathbb{Q}_2(c) = \mathbb{Q}_2(d) = 2^{-5}, \quad \mathbb{Q}_2(e) = 7/8.$$

Result: $\mathbb{Q}^* = \mathbb{Q}_1$ as

$$d_{\mathit{KL}}(\mathbb{P}||\mathbb{Q}_1) = \frac{4\log_2(16/5) + \log_2(4/15)}{5} \approx 0.9611$$

and

$$d_{KL}(\mathbb{P}||\mathbb{Q}_2) = \frac{4\log_2(32/5) + \log_2(8/35)}{5} \approx 1.7166$$

Connecting information theory and random variables

For discrete distributions $\mathbb P$ and $\mathbb Q$ on A we defined the information content rv

$$I_{\mathbb{P}}(a) = -\log(\mathbb{P}(a)), \quad I_{\mathbb{Q}}(a) = -\log(\mathbb{Q}(a))$$

and the K-L divergence from $\mathbb Q$ to $\mathbb P$ takes the form

$$d_{\mathit{KL}}(\mathbb{P}||\mathbb{Q}) = \mathbb{E}^{\mathbb{P}}[I_{\mathbb{Q}} - I_{\mathbb{P}}] = \sum_{a \in A} \log \left(\frac{\mathbb{P}(a)}{\mathbb{Q}(a)}\right) \mathbb{P}(a)$$

For continuous rv X,Y with densities $\pi_X,\pi_Y\in\mathcal{M}_+$, we define the information content as

$$I_{\pi_X}(x) = -\log(\pi_X(x)), \quad I_{\pi_Y}(x) = -\log(\pi_Y(x))$$

and

$$d_{\mathsf{KL}}(\pi_X||\pi_Y) = \mathbb{E}^{\pi_X}[I_{\pi_Y} - I_{\pi_X}] = \int_{\mathbb{R}^d} \log\left(\frac{\pi_X(x)}{\pi_Y(x)}\right) \pi_X(x) \, dx$$

Expected information gain Bayesian inversion

For the additive Gaussian inverse problem

$$Y = G(U) + \eta \tag{3}$$

with $\pi_{\eta}, \pi_{U} \in \mathcal{M}_{+}$ and $U \perp \eta$, the posterior is also a strictly positive pdf

$$\pi(u|y) = \frac{\exp(-\mathsf{L}(u))\pi_U(u)}{Z}.\tag{4}$$

Then

$$d_{\mathsf{KL}}(\pi(\cdot|y)||\pi_U) = \mathbb{E}^{\pi(\cdot|y)}[I_{\pi_U} - I_{\pi(\cdot|y)}]$$

is a measure of the information gained by revising the prior π_U into the posterior $d_{KL}(\pi(\cdot|y)||\pi_U)$

Interpretation: wrt $\pi(\cdot|y)$, $I_{\pi(\cdot|y)}$ yields the minimum expected information content, so, as we already know,

$$\mathbb{E}^{\pi(\cdot|y)}[I_{\pi_U}-I_{\pi(\cdot|y)}]\geq 0.$$

Variational formulation of Bayes theorem

Theorem 6 (SST Thm 4.9)

For the inverse problem (3) it holds that

$$\pi(\cdot|y) = \arg\min_{p \in \mathcal{M}_+} d_{\mathit{KL}}(p||\pi_U) + \mathbb{E}^p[L(u)]$$

Verification: Recalling that $\pi(\cdot|y) = \frac{\exp(-L(u))\pi_U(u)}{Z}$,

$$d_{KL}(p||\pi(\cdot|y)) = \int_{\mathbb{R}^d} \log\left(\frac{p\pi_U}{\pi(x|y)\pi_U}\right) p(x) dx$$

$$= \int_{\mathbb{R}^d} \log\left(\frac{pZ \exp(L(u))}{\pi_U}\right) p(x) dx$$

$$= \int_{\mathbb{R}^d} \log\left(\frac{p}{\pi_U}\right) + L(u) p(x) dx + \log(Z)$$

$$= d_{KL}(p||\pi_U) + \mathbb{E}^p[L] + \log(Z)$$

and

$$\pi(\cdot|y) = \arg\min_{p \in \mathcal{M}_+} d_{\mathit{KL}}(p||\pi(\cdot|y)).$$

Best Gaussian fit and K-L divergence

Consider again the posterior obtained from the inverse problem (3),

$$\pi(u|y) = \frac{\exp(-\mathsf{L}(u))\pi_U(u)}{Z}.$$
 (5)

(6)

Theorem 7

Assume that L is non-negative, continuous, and globally bounded from above and that $U \sim N(0, \lambda^{-1}I)$ for some $\gamma > 0$. Then there exists at least one pdf p in

$$\mathcal{A} := \{ \rho = PDF(N(\mu, C)) \mid \mu \in \mathbb{R}^d \text{ and } C \in \mathbb{R}^{d \times d} \text{ and pos definite} \}.$$

which satisfies the best-Gaussian-fit-of-posterior condition

$$d_{\mathsf{KL}}(p||\pi(\cdot|y)) = \inf_{\rho \in A} d_{\mathsf{KL}}(\rho||\pi(\cdot|y))$$

Essential fitting idea:

make
$$\log\left(\frac{p(x)}{\pi(x|y)}\right)$$
 small i.e., $\frac{p}{\pi(\cdot|y)}\lessapprox 1.$

Ideas in proof

For $p_{\mu,C} = PDF(N(\mu,C))$ it is possible to show that for

$$I(\mu, C) := d_{KL}(p_{\mu, C}||\pi(\cdot|y))$$

it holds that

$$I(0,I) < \infty, \quad \lim_{\|\mu\| \to \infty} I(\mu,C) = \infty$$

and

$$\lim_{trace(C)\to 0} I(\mu, C) = \lim_{trace(C)\to \infty} I(\mu, C) = \infty.$$

Consequently, there exists R > r > 0 s.t.

$$\arg\inf_{p\in\mathcal{A}}d_{\mathit{KL}}(p||\pi)\in ilde{\mathcal{A}}_{r,R}$$

where

$$ilde{\mathcal{A}}_{r,R} = \{p_{\mu,\mathcal{C}} \in \mathcal{A} \mid |\mu| < R, \quad \text{and} \quad r < trace(\mathcal{C}) < R\}.$$

Best Gaussian fit by moment matching

One may also fit p to π by minimizing $d_{\mathit{KL}}(\pi(\cdot|y)||p)$

Theorem 8 (SST Thm 4.5)

Let $\pi(\cdot|y)$ denote the posterior density of the inverse problem (3). If $\bar{\mu} = \mathbb{E}^{\pi(\cdot|y)}[u]$ is finite and $\bar{C} = \mathbb{E}^{\pi(\cdot|y)}[(u-\bar{\mu})(u-\bar{\mu})^T]$ is finite and positive definite then

$$p_{\bar{\mu},\bar{C}} = \arg\inf_{p \in \mathcal{A}} d_{KL}(\pi||p),$$

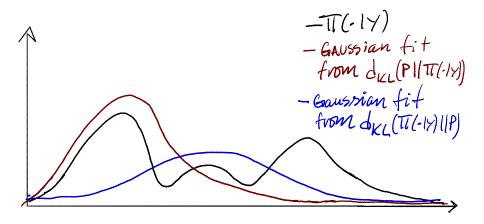
and the minimizer $p_{\bar{\mu},\bar{C}}$ is unique.

Essential fitting idea:

make
$$\log(\frac{\pi(x|y)}{p(x)})$$
 small, i.e., $\frac{\pi(\cdot|y)}{p} \lessapprox 1$.

Comparison of the fitting approaches

- For $d_{KL}(p||\pi(\cdot|y))$: make $\frac{p}{\pi(\cdot|y)} \lesssim 1$
- For $d_{KL}(\pi(\cdot|y)||p)$: make $\frac{\pi(\cdot|y)}{p} \lesssim 1$



Next time

discrete time continuous state-space Markov chains

Markov chain Monte Carlo methods

■ introduction to smoothing and filtering in continuous state-space