COMPUTER ORGANIZATION & ARCHITECTURE

Course Instructor: Dr. Shafina LAB SESSION: 03 DATE: February 17, 2023

Title: A QUICK GUIDE OF MIPS ARCHITECTURE

LAB SESSION 3

Computer Organization & Architecture

MIPS Architecture and Assembly Language Overview

Data Types and Literals

Data types:

- Instructions are all 32 bits
- byte(8 bits), halfword (2 bytes), word (4 bytes)
- a character requires 1 byte of storage
- an integer requires 1 word (4 bytes) of storage

Literals:

- numbers entered as is. <u>e.g.</u> 4
- characters enclosed in single quotes. <u>e.g.</u> 'b'
- strings enclosed in double quotes. <u>e.g.</u> "A string"

Registers

- 32 general-purpose registers
- register preceded by \$ in assembly language instruction two formats for addressing:
 - o using register number e.g. \$0 through \$31
 - o using equivalent names <u>e.g.</u> \$t1, \$sp
- special registers Lo and Hi used to store result of multiplication and division
 - not directly addressable; contents accessed with special instruction mfhi ("move from Hi") and mflo ("move from Lo")
- stack grows from high memory to low memory

Program Structure

- just plain text file with data declarations, program code (name of file should end in suffix .s to be used with QTSPIM simulator)
- data declaration section followed by program code section

Data Declarations

- placed in section of program identified with assembler directive .data
- declares variable names used in program; storage allocated in main memory (RAM)

Code

- placed in section of text identified with assembler directive .text
- contains program code (instructions)
- starting point for code execution given label main:
- ending point of main code should use exit system call (see below under System Calls)

Comments

- anything following # on a line# This stuff would be considered a comment
- Template for a MIPS assembly language program:

```
# Comment giving name of program and description
of function
# Template.s
# Bare-bones outline of MIPS assembly language
program

.
data # variable declarations
follow this line
# ...

. text # instructions follow this
line

main: # indicates start of code
(first instruction to execute)
# ...

# End of program, leave a blank line afterwards
to make SPIM happy
```

Data Declarations

format for declarations:

```
name: storage_type value(s)
```

- create storage for variable of specified type with given name and specified value
- value(s) usually gives initial value(s); for storage type .space, gives number of spaces to be allocated

Note: labels always followed by colon (:)

Load / Store Instructions

- RAM access only allowed with load and store instructions
- all other instructions use register operands

load:

```
lw register destination, RAM source
```

#copy word (4 bytes) at source RAM location to destination register.

```
lb register_destination, RAM_source
```

#copy byte at source RAM location to low-order byte of destination register, # and sign-extend to higher-order bytes

store word:

```
sw register_source, RAM_destination
```

#store word in source register into RAM destination

```
sb register source, RAM destination
```

#store byte (low-order) in source register into RAM destination

load immediate:

li register_destination, value #load immediate value into destination register

Indirect and Based Addressing

• Used only with load and store instructions

load address:

```
la $t0, var1
```

• copy RAM address of var1 (presumably a label defined in the program) into register \$t0

indirect addressing:

```
lw $t2, ($t0)
```

load word at RAM address contained in \$t0 into \$t2

```
sw $t2, ($t0)
```

• store word in register \$t2 into RAM at address contained in \$t0

based or indexed addressing:

```
lw $t2, 4($t0)
```

- load word at RAM address (\$t0+4) into register \$t2
- "4" gives offset from address in register \$t0

```
sw $t2, -12($t0)
```

- store word in register \$t2 into RAM at address (\$t0 12)
- negative offsets are fine

Note: based addressing is especially useful for:

- arrays; access elements as offset from base address
- stacks; easy to access elements at offset from stack pointer or frame pointer

```
example
```

Arithmetic Instructions

- most use 3 operands
- all operands are registers; no RAM or indirect addressing

• operand size is word (4 bytes)

```
$t0,$t1,$t2
                     # $t0 = $t1 + $t2; add as signed (2's
complement) integers
       $t2,$t3,$t4
                     # $t2 = $t3 Đ $t4
sub
addi
       $t2,$t3, 5
                   # $t2 = $t3 + 5; "add immediate" (no sub
immediate)
addu $t1,$t6,$t7
                  # $t1 = $t6 + $t7;
                                         add as unsigned integers
subu $t1,$t6,$t7 # $t1 = $t6 + $t7; subtract as unsigned
integers
                    # multiply 32-bit quantities in $t3 and $t4, and
mult
      $t3,$t4
store 64-bit
# result in special registers Lo and Hi: (Hi,Lo) = $t3 * $t4
div
       $t5,$t6
                     # Lo = $t5 / $t6 (integer quotient)
                     # Hi = $t5 mod $t6 (remainder)
mfhi $t0
                     # move quantity in special register Hi to $t0:
#$t0 = Hi
mflo $t1
                     # move quantity in special register Lo to $t1:
#$t1 = Lo
                     # used to get at result of product or quotient
move $t2,$t3 # $t2 = $t3
```

Control Structures

Branches

• comparison for conditional branches is built into instruction

<u>Jumps</u>

```
j target # unconditional jump to program label target
jr $t3 # jump to address contained in $t3 ("jump register")
```

Subroutine Calls

subroutine call: "jump and link" instruction

```
jal sub label # "jump and link"
```

- copy program counter (return address) to register \$ra (return address register)
- jump to program statement at sub_label

subroutine return: "jump register" instruction

```
jr $ra # "jump register"
```

• jump to return address in \$ra (stored by jal instruction)

Note: return address stored in register \$ra; if subroutine will call other subroutines, or is recursive, return address should be copied from \$ra onto stack to preserve it, since jal always places return address in this register and hence will overwrite previous value

System Calls and I/O (SPIM Simulator)

- used to read or print values or strings from input/output window, and indicate program end
- use **syscall** operating system routine call
- first supply appropriate values in registers \$v0 and \$a0-\$a1
- result value (if any) returned in register \$v0

The following table lists the possible **syscall** services.

Service	Code in \$v0	Arguments	Results
print_int	1	\$a0 = integer to be printed	
print_float	2	\$f12 = float to be printed	
print_double	3	\$f12 = double to be printed	

print_string	4	\$a0 = address of string in memory	
read_int	5		integer returned in \$v0
read_float	6		float returned in \$v0
read_double	7		double returned in \$v0
read_string	8	\$a0 = memory address of string input buffer \$a1 = length of string buffer (n)	
sbrk	9	\$a0 = amount	address in \$v0
exit	10		

- The print_string service expects the address to start a nullterminated character string. The directive .asciiz creates a nullterminated character string.
- The read_int, read_float and read_double services read an entire line of input up to and including the newline character.
- The read_string service has the same semantices as the UNIX library routine fgets.
 - It reads up to n-1 characters into a buffer and terminates the string with a null character.
 - If fewer than n-1 characters are in the current line, it reads up to and including the newline and terminates the string with a null character.
- The sbrk service returns the address to a block of memory containing n additional bytes. This would be used for dynamic memory allocation.
- The exit service stops a program from running.

```
Print out integer value contained in register $t2
e.g.
li
       $v0, 1 # load appropriate system call code into register $v0;
               # code for printing integer is 1
       $a0, $t2
                      # move integer to be printed into $a0: $a0 = $t2
move
                      # call operating system to perform operation
syscall
      Read integer value, store in RAM location with label int value
(presumably declared in data section)
li
    $v0, 5
                              # load appropriate system call code into
register $v0;
                              # code for reading integer is 5
syscall
                      # call operating system to perform operation
```

e.g. Print out string (useful for prompts)

.data

string1 .asciiz "Print this.\n" # declaration for string

variable,

.asciiz directive makes

string null terminated

.text

main: li \$v0, 4 # load appropriate system call code into register \$v0;

code for printing string

is 4

la \$a0, string1 # load address of string

to be printed into \$a0

syscall # call operating system to

perform print operation

 $\underline{\text{e.g.}}$ To indicate end of program, use \mathbf{exit} system call; thus last lines of program should be:

li \$v0, 10\$ # system call code for exit = 10 syscall # call operating sys