

# **Compiler Design Assignment**

**on Semantic Analysis**

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Question 1):- what is an attribute? Explain the various types of Attributes and their uses.

Attributes grammar is a special form of context-free grammar where some additional information are appended to one or more of its non-terminals in order to provide context. Sensitive information values, such as integer, float, character, string and expressions.

There are two types of Attributes:-

① Inherited:- An attribute is said to be inherited attribute if its parse tree node value is determined by the attributes value at parent or siblings node. The production must have non-terminal as a symbol in its body. An inherited attribute at Node  $N$  defined only in terms of attributes value  $N$ 's parent,  $N$  itself and  $N$ 's siblings.

② Synthesized:-

A synthesized attribute for a Nonterminal  $A$  at parse tree node  $N$  is defined by a semantic rule associated with the production at  $N$ . Note that the production must have  $A$  as its head. A synthesized attribute at Node  $N$  is defined only in terms of attribute values at the children of  $N$  and at itself.

Production

$S \rightarrow E$

$E \rightarrow E_1 \text{ OR } T$

$E \rightarrow E_1 \text{ AND } T$

$E \rightarrow T$

$T \rightarrow id$

Semantic Rules

check  $(E.val)$

$E.val = E_1.val \text{ 'OR' } T.val$

$E.val = E_1.val \text{ 'AND' } T.val$

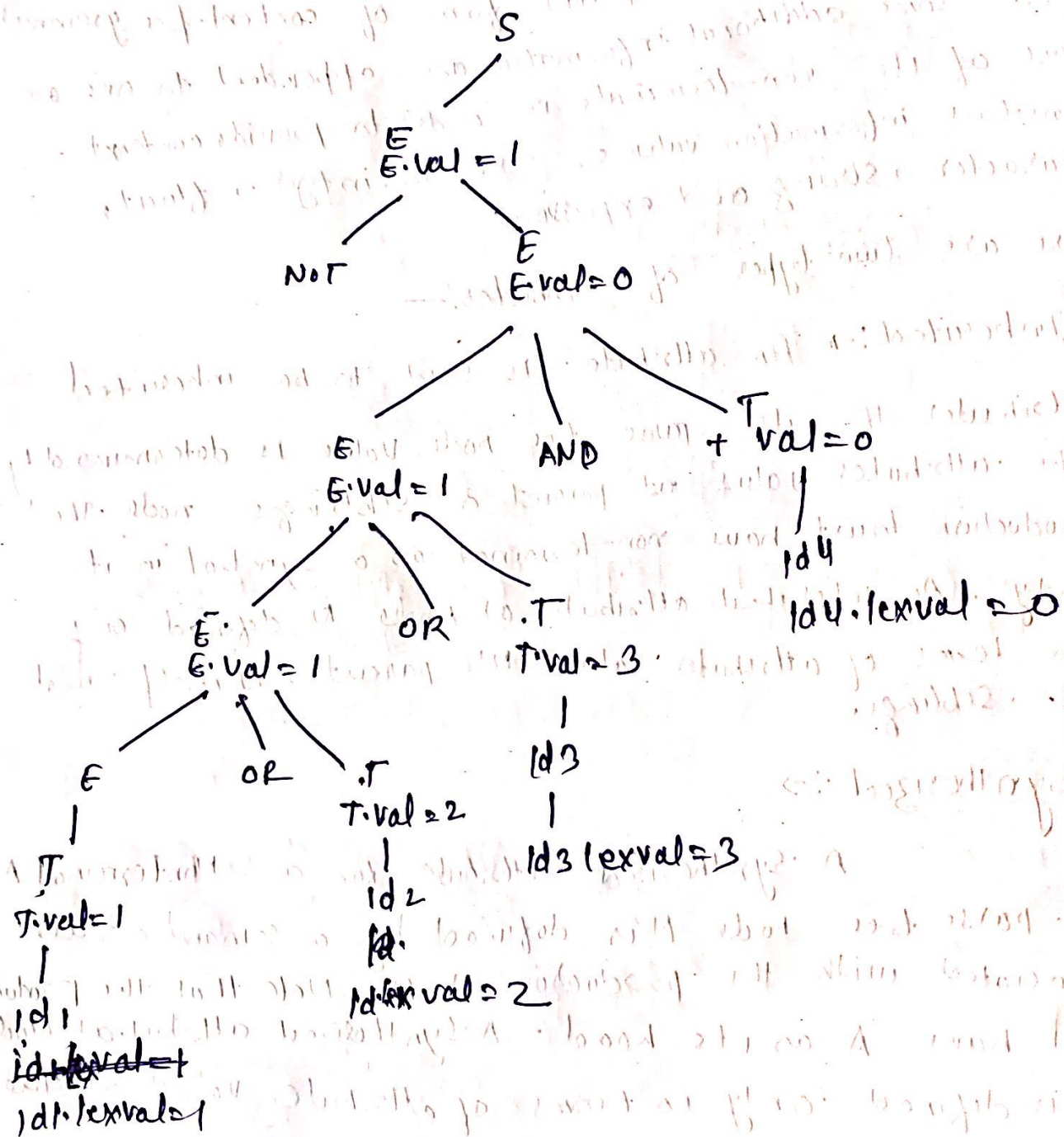
$E.val = T.val$

$T.val = id.lexval$



$E \rightarrow \text{NOT } E \text{ implies } E.\text{val} = \sim E.\text{val}$

Example:- NOT (id1 OR id2 OR id3 AND id4)  
 NOT (1 OR 2 OR 3 AND 0)



Q2:- Design the syntax Directed Definition for the constructing syntax tree for the boolean expressions use suitable examples to create a tree for any boolean expression from their selected CFG.

$E \rightarrow E_1 \cdot OR T$

$E \rightarrow E_1 AND T$

$E \rightarrow NOT E$

$E \rightarrow T$

$T \rightarrow id$

$T \rightarrow num$

$E.nptr := mknode("OR", E.nptr,$

$T.nptr)$

$mknode("AND", E.nptr, T.nptr)$

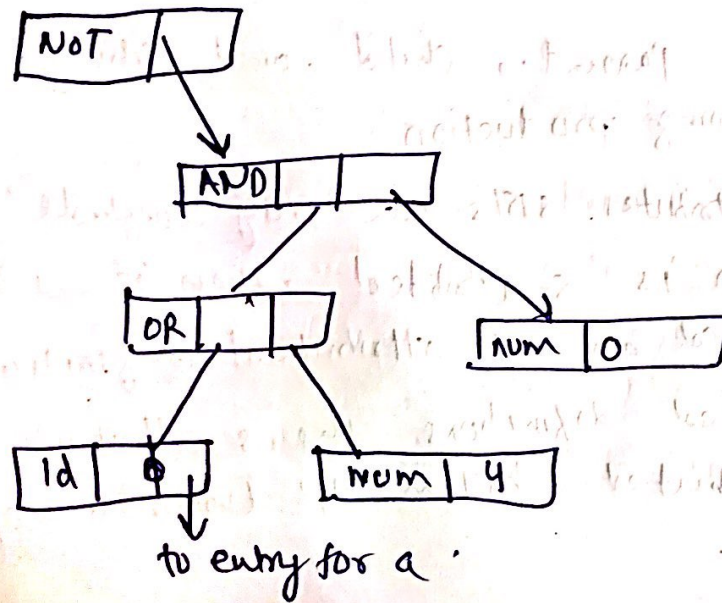
$mknode("NOT", E.nptr)$

$T.nptr$

$mkleaf(id, id.lexval)$

$mkleaf(num, num.val)$

NOT (a OR 4 AND 0)

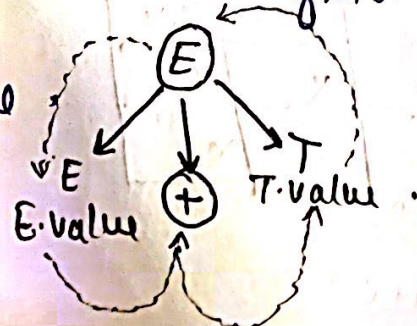


Question 3:- Difference b/w S Attributes Definition & L attributes Definition Also explain how Semantic rule are.

Synthesized attributes it is called S- attributed SDT.

These attributes are evaluated using s- attributed SDTs that have their semantic actions written after the production.

$$E.val = E.value + T.value$$





Attributes in S-attributed SDTs are evaluated in bottom up parsing as values of the pattern node depend upon the values of the child nodes.

L-attributed SDT:— This form of SDT uses inherited attributes with ~~inherited~~ restriction of not taking value from right siblings.

In L-attributed SDTs, a non-terminal can get values from its parent, child and sibling nodes as in the following production.

L-attributes SDTs we may conclude that if a definition is S-attributed, then it is also L-attributed as L-attributed definition encloses S-attributed definitions the S-attributed definition are also L-attributed bottom up Evaluation of S-attributed definition.

① maintain the values of the synthesized attributes of the grammar symbols into a parallel stack.

② The evaluation of values of the attributes are done during reduction operation.

Z	Z.Z
Y	Y.Y
X	X.X
X	X

IF  $A \rightarrow XY Z$

After  
Reduction

A	A.A
X	X.X



Q:-4:- Along with mentioning appropriate rules and annotated parse tree perform type checking on the given expression :-

if (a != b)

{ if (i = j)

{

a = b + c;

}

}

// Note a, b, c, i and j are of integer

// type:

$S \rightarrow \text{if}(E) \{ S1 \}$

$S.type = (\text{if } E.type = \text{Boolean, then } S1.type \text{ else type-error})$

$S \rightarrow \text{id} \mid E$

$S \rightarrow S1 : S2$

$S.type = (\text{if id.type} = E.type \text{ then void else type-error})$

$E \rightarrow E1 + E2$

$S.type = (\text{if } S1.type = \text{void and } S2.type = \text{void then void else type-error})$

$E \rightarrow F \mid = E$

$E \rightarrow F == E$

$E \rightarrow F$

$F \rightarrow \text{id}$

$E.type = (\text{if } E1.type = E2.type \text{ then } E1.type \text{ else type-error})$

$F \rightarrow \text{num}$

$E.type = \text{if } (E1.type = E2.type \text{ then Boolean else type-error})$

$E.type = (\text{if } E1.type = F.type \text{ then Boolean else type-error})$

f. type = integer

