

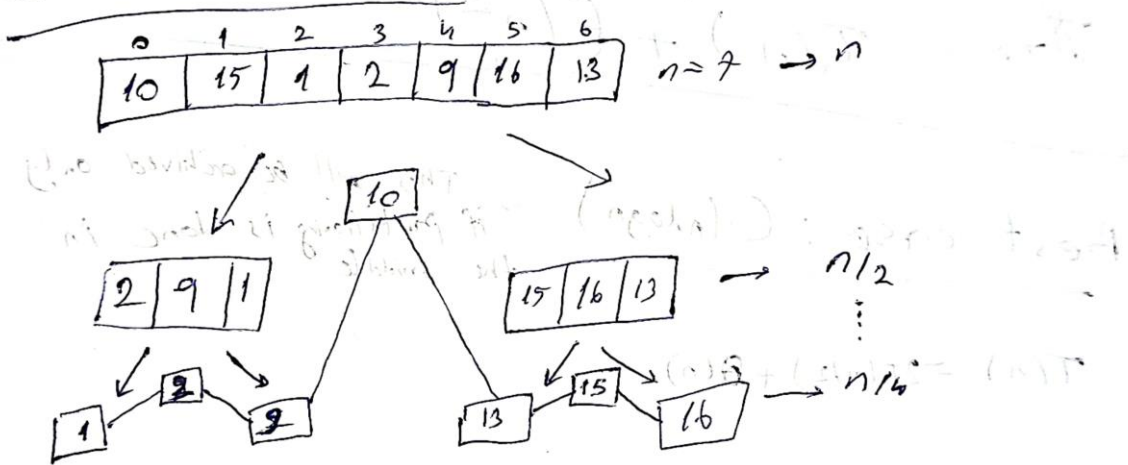
```
100luk:
mean merge sort: 0.044
mean quick sort: 0.029
mean new sort: 0.032

1000luk:
mean merge sort: 0.094
mean quick sort: 0.062
mean new sort: 0.473

10000luk:
mean merge sort: 0.5103333333333333
mean quick sort: 0.4783333333333333
mean new sort: 30.826
```

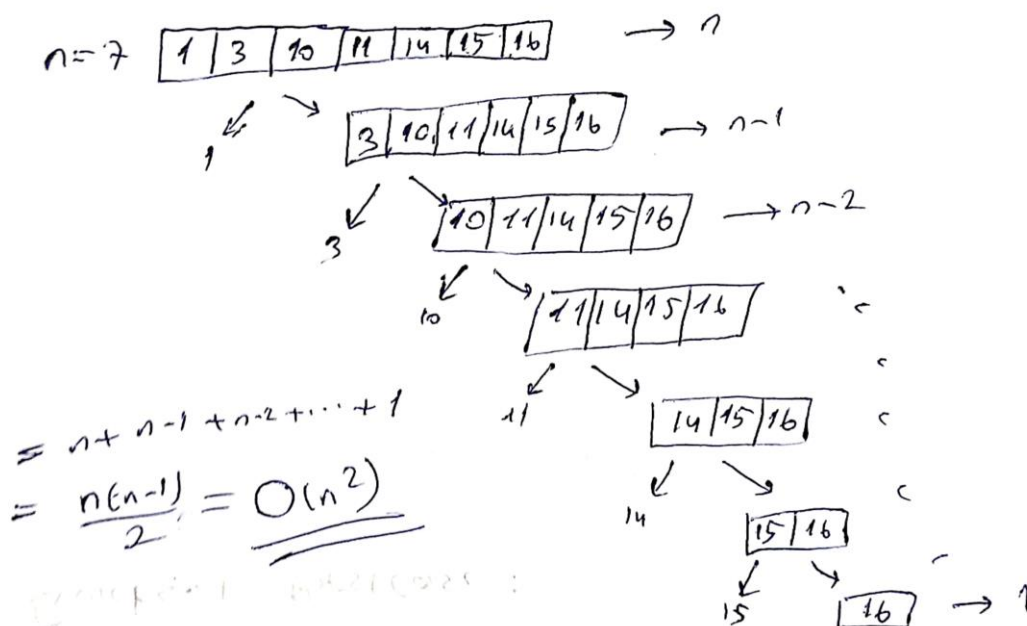
QuickSort:

QuickSort Bestcase.



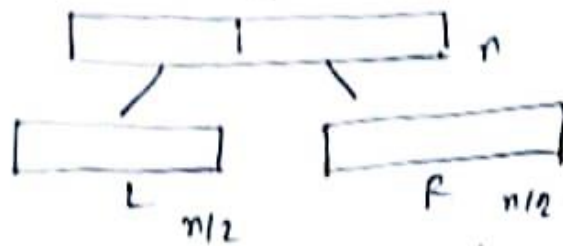
$$\frac{n}{2^k} \approx 1 \quad n = 2^k \quad = \underline{\underline{O(n \cdot \log n)}}$$

QuickSort worst case: Array is already sorted



MergeSort:

Time Complexity for Merge Sort



$$\begin{aligned}
 T(n) &= 2T(n/2) + n \\
 &= 2 \left[2T(n/4) + \frac{n}{2} \right] + n \\
 &= 2^2 T\left(\frac{n}{4}\right) + n + n \\
 &= 2^2 T\left(\frac{n}{2^2}\right) + 2n \\
 &= 2^2 \left[2T\left(\frac{n}{8}\right) + \frac{n}{4} \right] + 2n \\
 &= 2^3 T\left(\frac{n}{2^3}\right) + n + 2n \\
 &= 2^3 T\left(\frac{n}{2^3}\right) + 3n \\
 &= 2^k T\left(\frac{n}{2^k}\right) + nk
 \end{aligned}$$

$T(1)$
 $= O(1)$
 $= 1$

$$\frac{n}{2^3} = 1$$

$$\frac{n}{2^k} = 1$$

$$n = 2^k$$

$$\log_2 n = k \log_2 2$$

$$\log_2 n = k$$

$$k = \log_2 n$$

$$= 2^{\log_2 n} T(1) + n \log_2 n$$

$$= 2^{\log_2 n} \cdot 1 + n \log_2 n$$

$$= n + n \log_2 n$$

$$= O(n \log_2 n) \text{ worst case}$$

$$O(n^2)$$

newSort:

```

*/
public static void min_max_finder_helper(Integer[] nums, int head, int tail, MinMaxDuals p) {
    // if the array contains only one element
    if (head == tail) {
        if (nums[p.max] < nums[head]) {
            p.max = head;
        }
        if (nums[p.min] > nums[tail]) {
            p.min = tail;
        }
        return;
    }
    // if the array contains only two elements
    if (tail - head == 1) {
        if (nums[head] < nums[tail]) {
            if (nums[p.min] > nums[head]) {
                p.min = head;
            }
            if (nums[p.max] < nums[tail]) {
                p.max = tail;
            }
        } else {
            if (nums[p.min] > nums[tail]) {
                p.min = tail;
            }
            if (nums[p.max] < nums[head]) {
                p.max = head;
            }
        }
        return;
    }
    //middle element
    int mid = (head + tail) / 2;
    //left subarray
    min_max_finder_helper(nums, head, mid - 1, p);
    //right subarray
    min_max_finder_helper(nums, mid, tail, p);
}

```

$T(1)$
 $T(n) = 2T(n/2)$
 $T(n/2) = 2T(n/4)$
 $T(n) = 2^2 T(n/2^2)$
 \vdots k times
 $T(n) = 2^k \cdot T(n/2^k)$
 $k = \log(n)$
 $\rightarrow O(2^{\log n})$

$\rightarrow T(n/2)$
 $\rightarrow T(n/2)$

```

0
10 public static Integer[] newSort(Integer[] arr, int head, int tail) {
2     if (head > tail)
3         return arr;
4
5     MinMaxDuals mm = new MinMaxDuals();
6     mm = min_max_finder(arr, head, tail);
7
8     swap(arr, head, mm.min);
9     swap(arr, tail, mm.max);
10
11     return newSort(arr, head + 1, tail - 1);
12
13 }
14
15 /**

```

$\rightarrow \Theta(1)$
 $\rightarrow O(2^{\log n})$
 $T(n) = O(n \cdot 2^{\log n})$
 $\rightarrow \Theta(1)$
 $\rightarrow T(n-2)$