min = 2 (h-1) +1 - h (In the complete binary tree, there my be only one child)

max = 2 h(n-1)+1 - (2h-1). In the complete binary tree, there my be at most oneless then the perfect way tree)

Avarge = 
$$\frac{mh+mnx}{2}$$
  
 $= (2h(h-1)+1-h)+(2h(h-1)+1-(2h-1).h)$   
 $= 2h(h-4)-15$ 

$$\frac{1-\sqrt{n}}{n} (n) = T(n/2) + 1$$

$$T(n/2) = T(n/2^2) + 1$$

$$T(n) = T(n/2^{2}) + 1 + 1$$

$$T(n/2^{2}) = T(n/2^{3}) + 1$$

$$T(n) = T(n/2^3) + 1 + 1 + 1$$

$$T(n) = T(n/2k) + k$$

$$k = \log n$$

$$T(n) = 1 + \log n = \log n$$

1-C) yes, There some rast riching:

Theorem 1: let T be a binary free. For every k > 0, there are no more than 2k nodes in level k

Theorem 2: let T be as binary free wit  $\lambda$  levels, then T has no more than  $2^{\lambda}-1$  nodes.

N: nodes

L: nom of Leaves

Can determine the other two.

let The a nonempty, full binary tree;

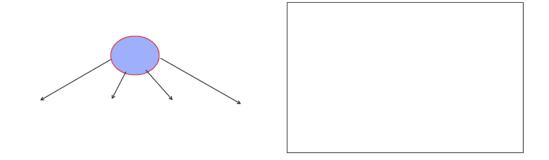
-) If T was I internal nodes, the number of leaves is

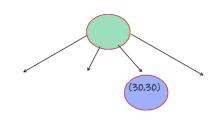
$$L = I + 1$$

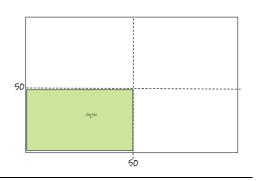
$$N = 2I + 1$$

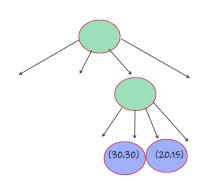
 $\Rightarrow$  If T has a total of N modes:  $5 I = \frac{N-1}{2}$   $L = \frac{N+1}{2}$ 

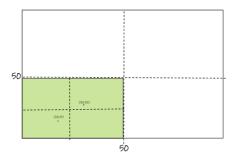
-> If T has L leaves -> N=2L-1 -> I=L-1

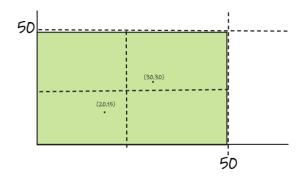




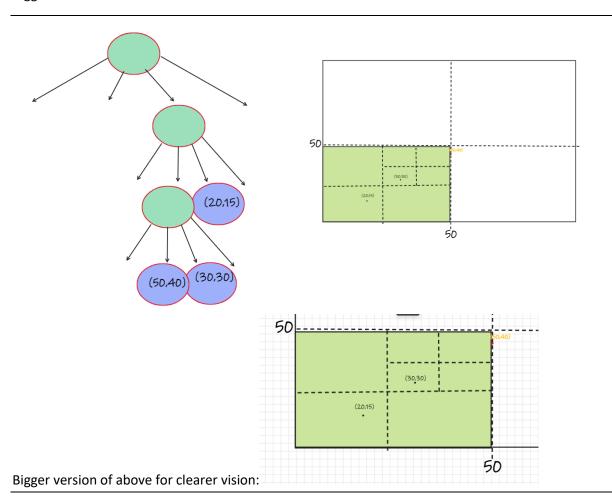


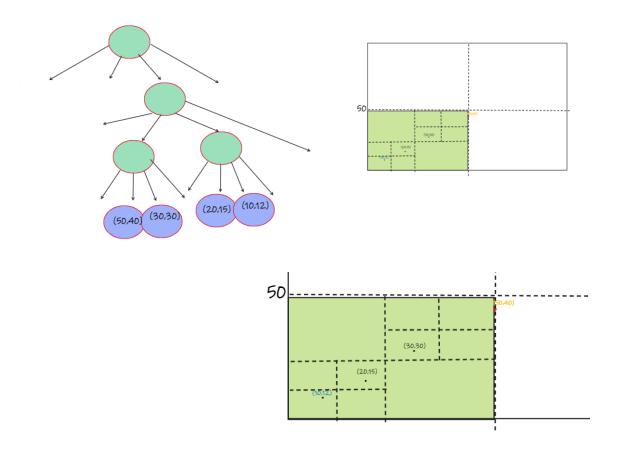




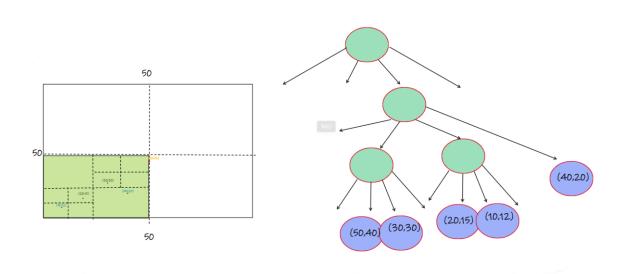


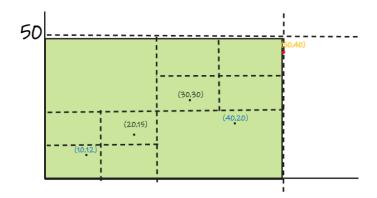
Bigger version of above for clearer vision:



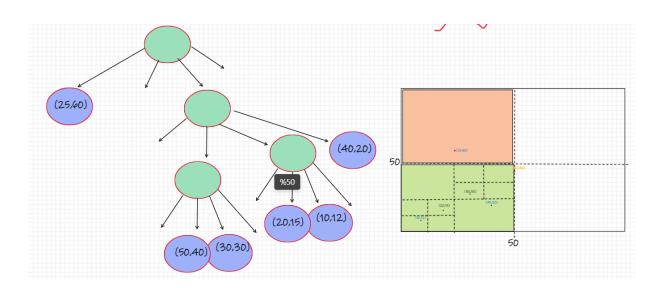


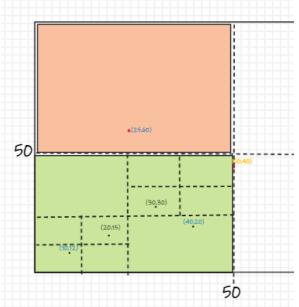
Bigger version of above for clear vision:



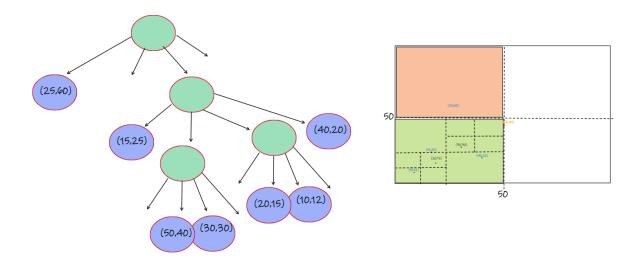


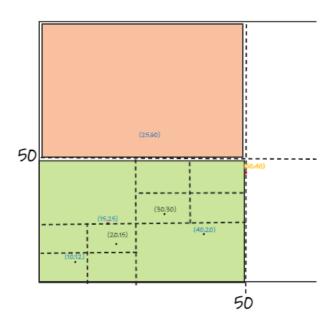
Bigger version of above for clear vision:





Bigger version of above for clear vision:





Bigger version of above for clear vision:

Note: the points that in exactly boundary of the sub square assumed in child of left-top square