Proposition - Part II

Course on Discrete Mathematics for GATE 2023

Sanchit Jain • Lesson 14 • Sept 16, 2022

Educator highlights

- Works at **Knowledge Gate**
- Studied at **Delhi Technological University**
- Have qualified gate Have experience of more than 8 years Have a YouTube channel with 5 lakh follower
- Lives in **Ghaziabad, Uttar Pradesh, India**
- Unacademy Educator since **19th July, 2019**
- **844,388 live minutes** taught in last 30 days
- Knows **Punjabi, Hinglish, Hindi and English**

Sanchit Jain

Follow

Legend in GATE - CS & IT

I am passionate for teaching computer science, having experience of more than 10 years. I teach all computer science subjects for GATE.

69M Watch mins | **3M** Watch mins (last 30 days) | **30K** Followers | **3K** Dedications

भाई में है दम
10 Subject ख़तम

Months Wise Calendar for GATE-2023 (Unacademy Plus) (Sanchit Jain)

Month	Morning Class (6:30am – 8:00am)	Evening Class (6:30pm – 8:30pm)
Feb-2022 Mar-2022 Apr-2022 May-2022 Jun-2022 Jul-2022 Aug-2022	CN(16-Feb to 26-March) (54 Hours) TOC (6-Apr to 26-May) (58 Hours) Compiler (1-June to 20-June) (34 Hours) (DS & Programming) (22-June to 23-July) (56 Hours) Algo (28-July to 25-Aug) (38 Hours)	DM (2-Feb to 2-Mar) (50 Hours) DBMS (16-Mar to 2-May) (50 Hours) DE (4-May to 30-May) (40 Hours) COA (8-June to 19-July) (38 Hours) OS (21-July to 27-Aug) (58 Hours)
Sep-2022 Oct-2022 Nov-2022 Dec-2022 Jan-2023 Feb-2023	**DM (1-Sept to 10-Oct) (50 Hours)** DBMS (13-Oct to 21-Nov) (50 Hours) DE (24-Nov to 26-Dec) (40 Hours) COA (29-Dec to 30-Jan) (38 Hours)	**CN (01-Sept to 12-Oct) (54 Hours)** TOC (13-Oct to 30-Nov) (58 Hours) Compiler (1-Dec to 23-Dec) (34 Hours) DS & C (22-Dec to 2-Feb) (56 Hours)

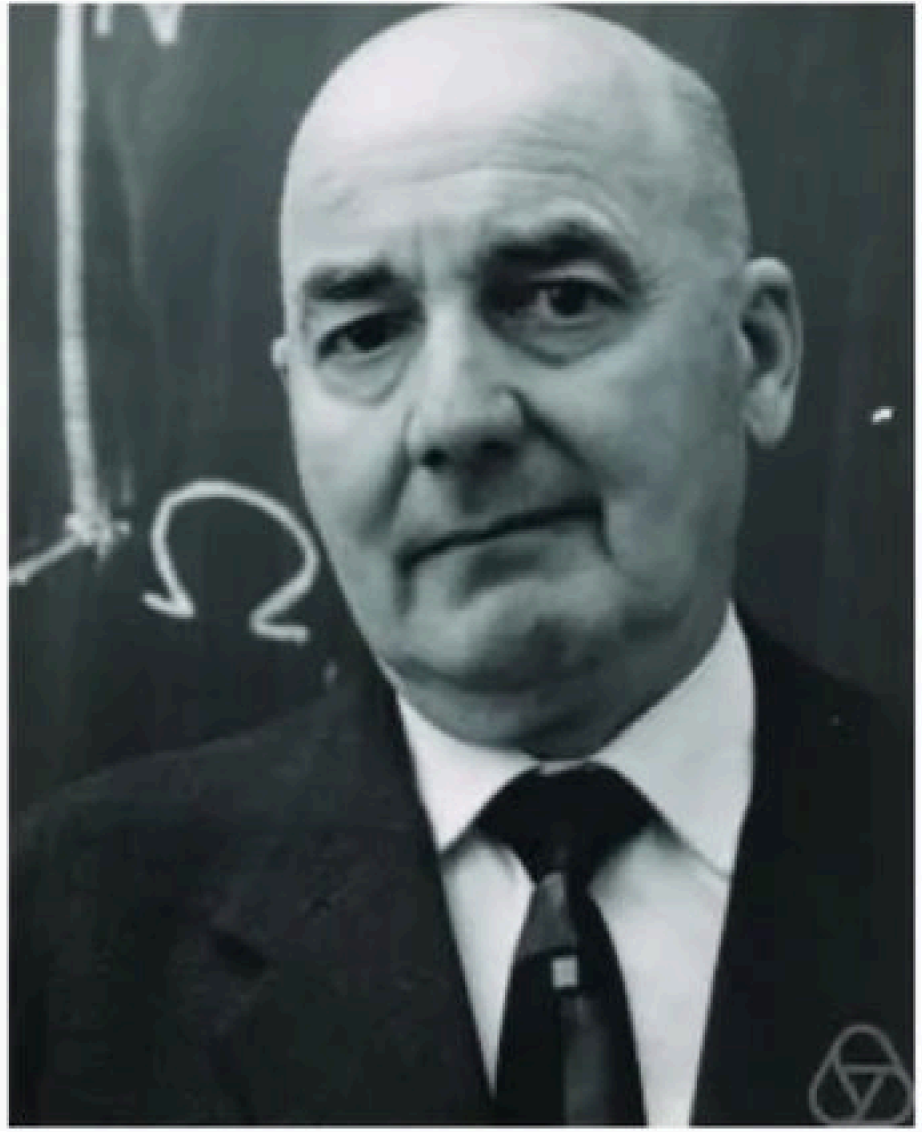
Conversion of POSET into a Hasse Diagram

- If we want to study Partial order relation further then it will be better to convert it into more convenient notation so that it can be studied easily.

- This graphical representation is called Hasse Diagram

- In order theory, a Hasse diagram is a type of mathematical diagram used to represent a finite partially ordered set, in the form of a drawing of its transitive reduction.

- The diagrams are named after Helmut Hasse (1898–1979)



Steps to convert partial order relation into hasse diagram

1- Draw a vertex for each element in the Set

2- If (a, b) ϵ R then draw an edge from a to b

3- Remove all Reflexive and Transitive edges

4- Remove the direction of edges and arrange them in the increasing order of heights.

Q Consider a Partial order relation and convert it into hasse diagram?

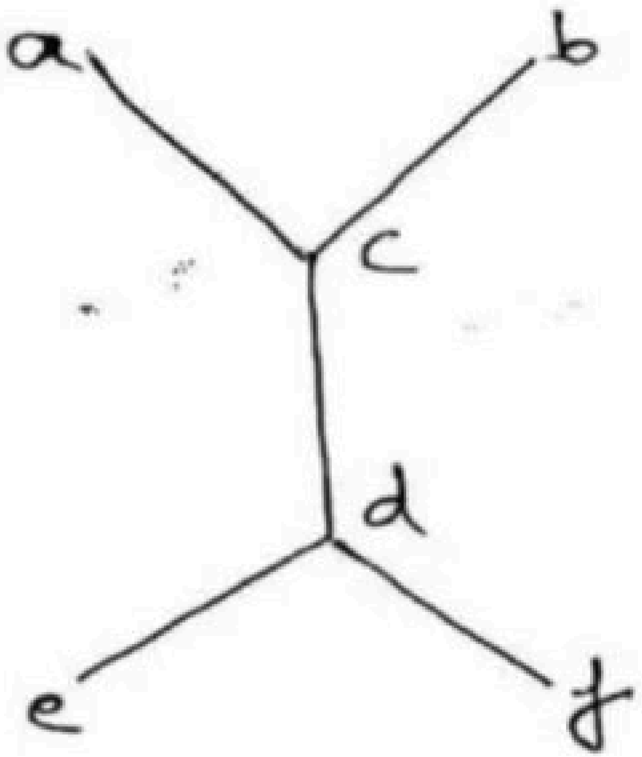
R = {(1,1), (1,2), (1,4), (1,8), (2,2), (2,4), (2,8), (4,4), (4,8), (8,8)}

Q Consider a Partial order relation and convert it into hasse diagram?

R = {(1,1), (1,2), (1,3), (1,6), (2,2), (2,6), (3,3), (3,6), (6,6)}

Break

Q Study the following hasse diagram and find which of the following are valid?



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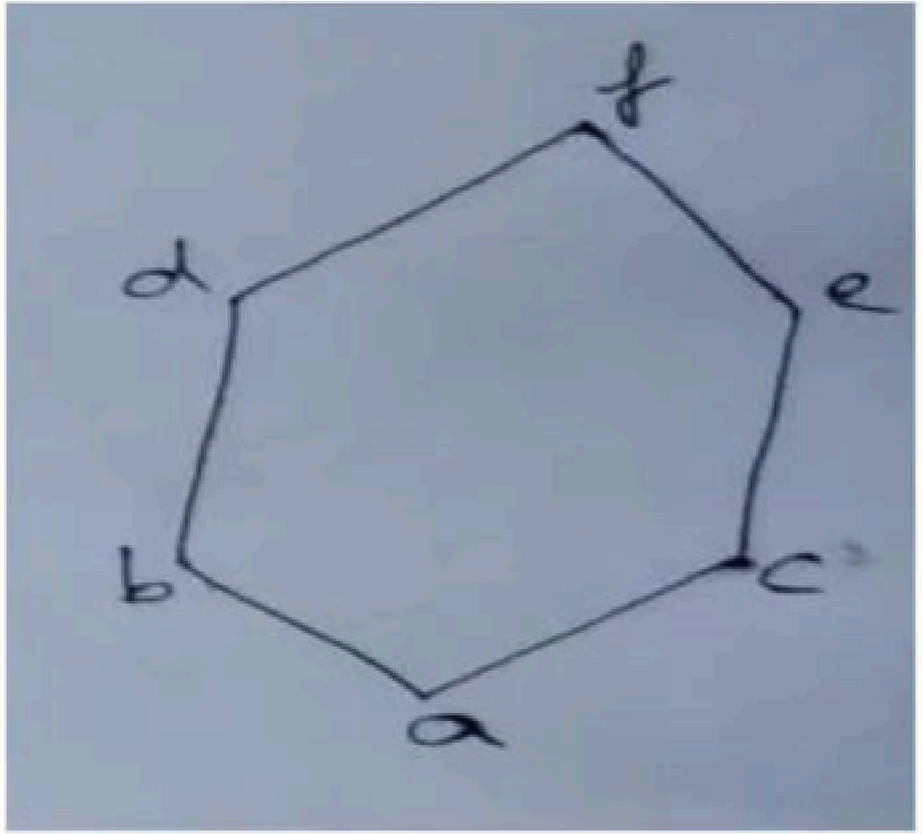

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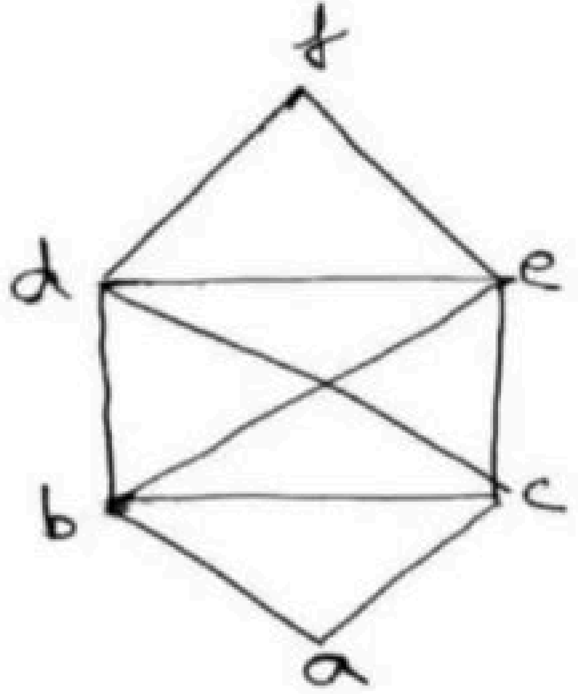

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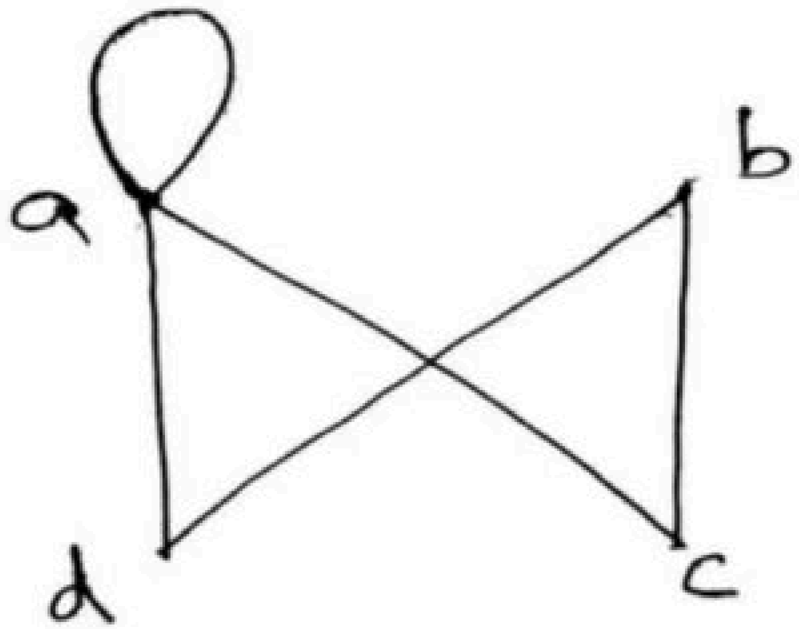

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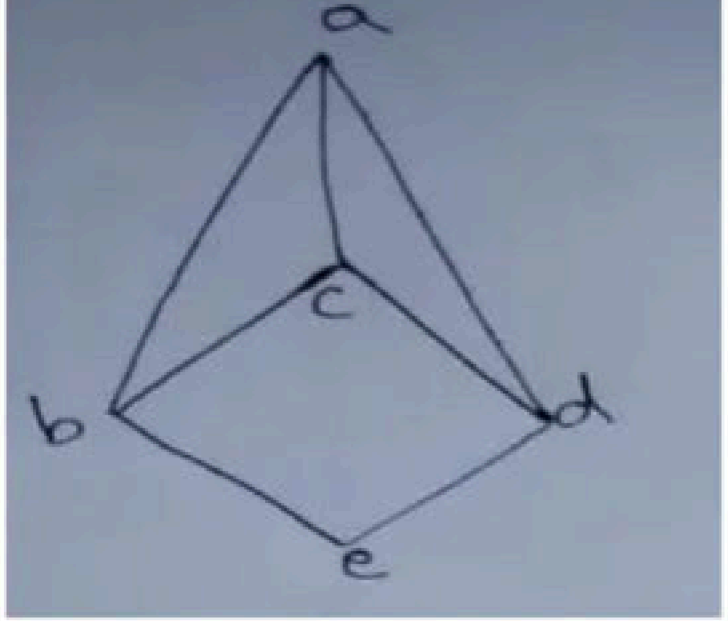

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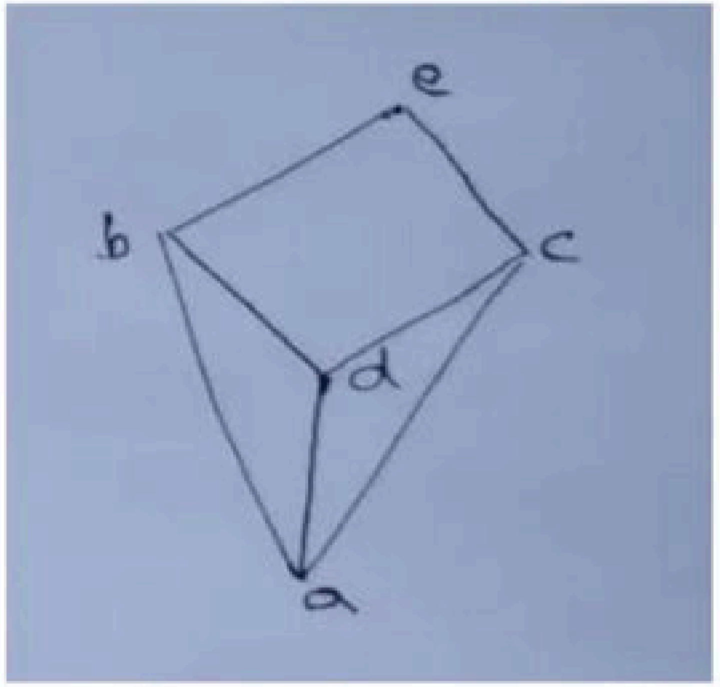

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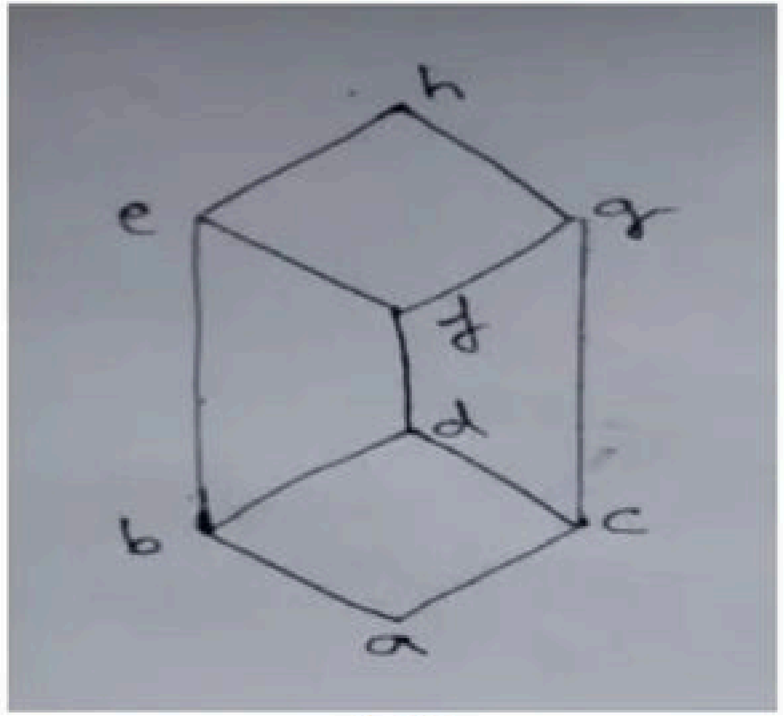

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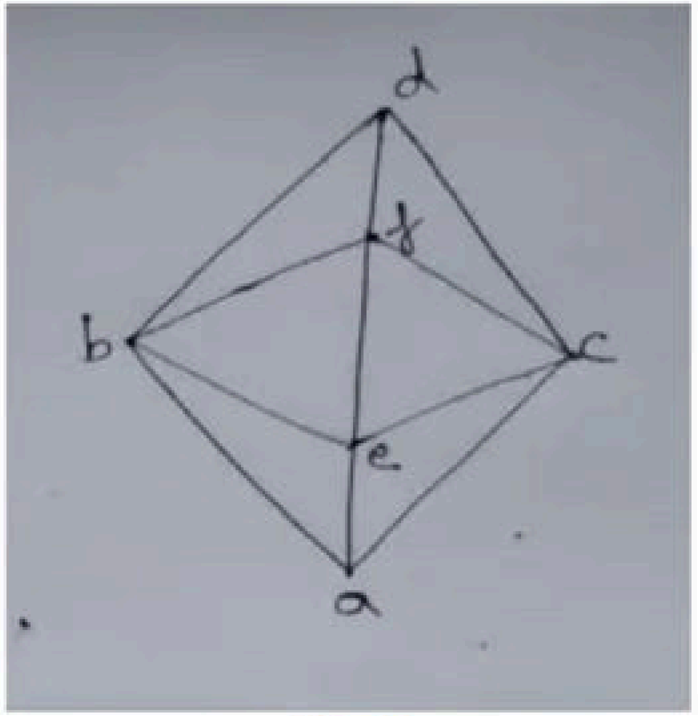

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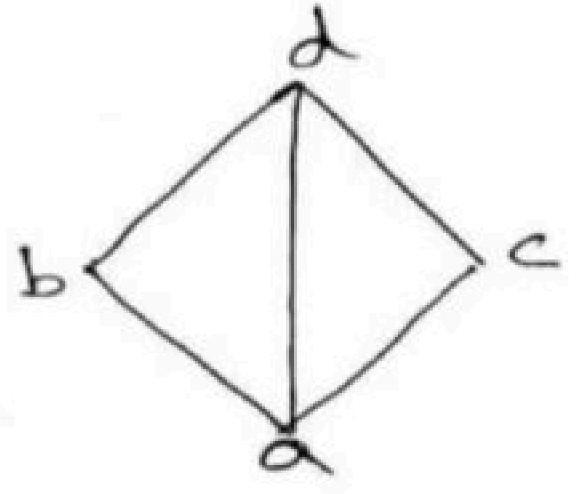

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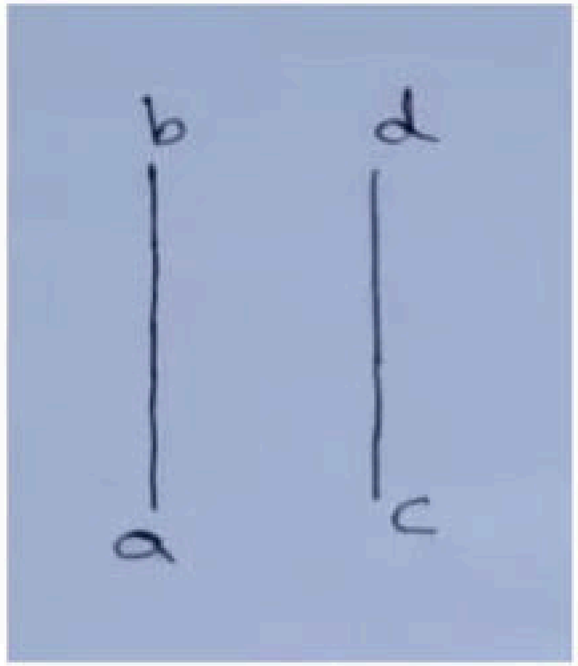

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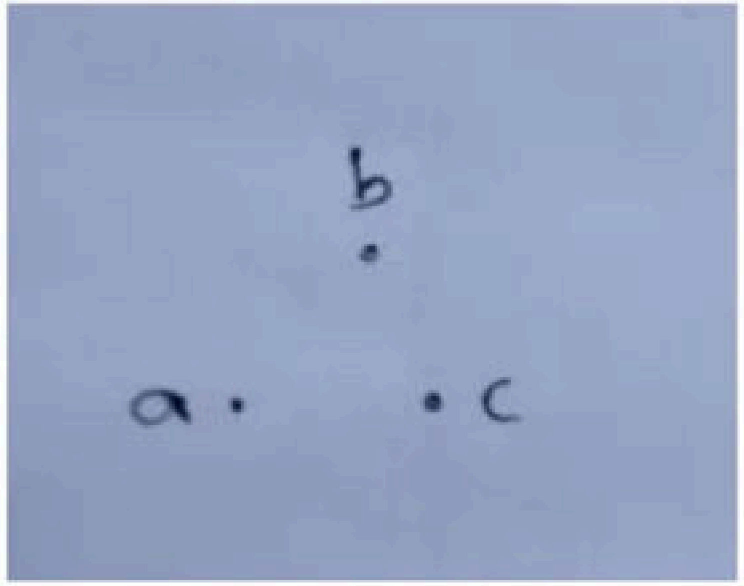


(10)

Conclusion

- We can not have a horizontal edge in a hasse diagram
- We can not have a reflexive and transitive edge in Hasse Diagram

Q Let X = {2,3,6,12,24}, Let ≤ be the partial order defined by X ≤ Y if x divides y. Number of edges as in the Hasse diagram of (X, ≤) is. **(GATE-1996) (1 Marks)**

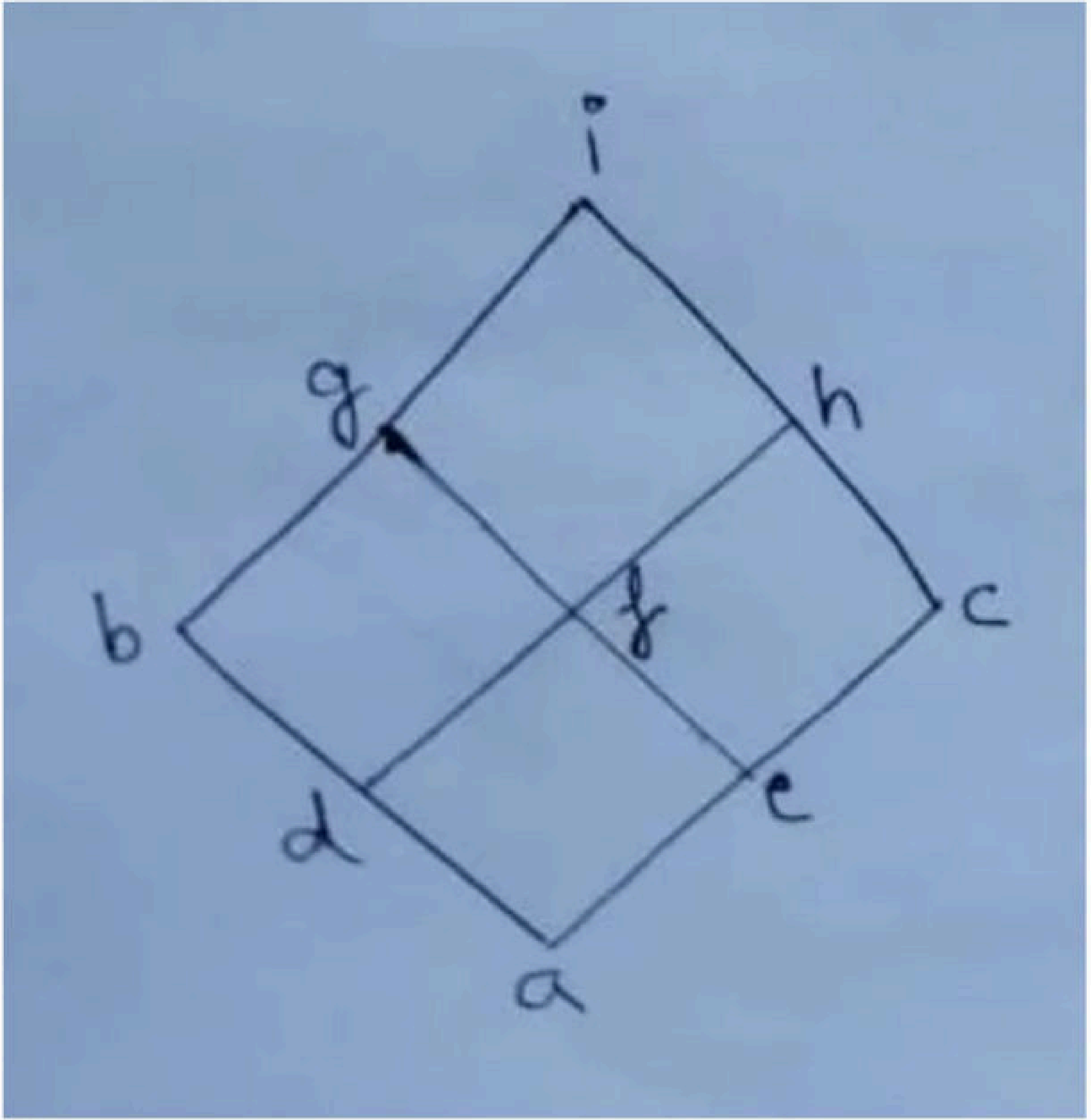
(a) 3 **(b)** 4 **(c)** 9 **(d)** None of the above

Break

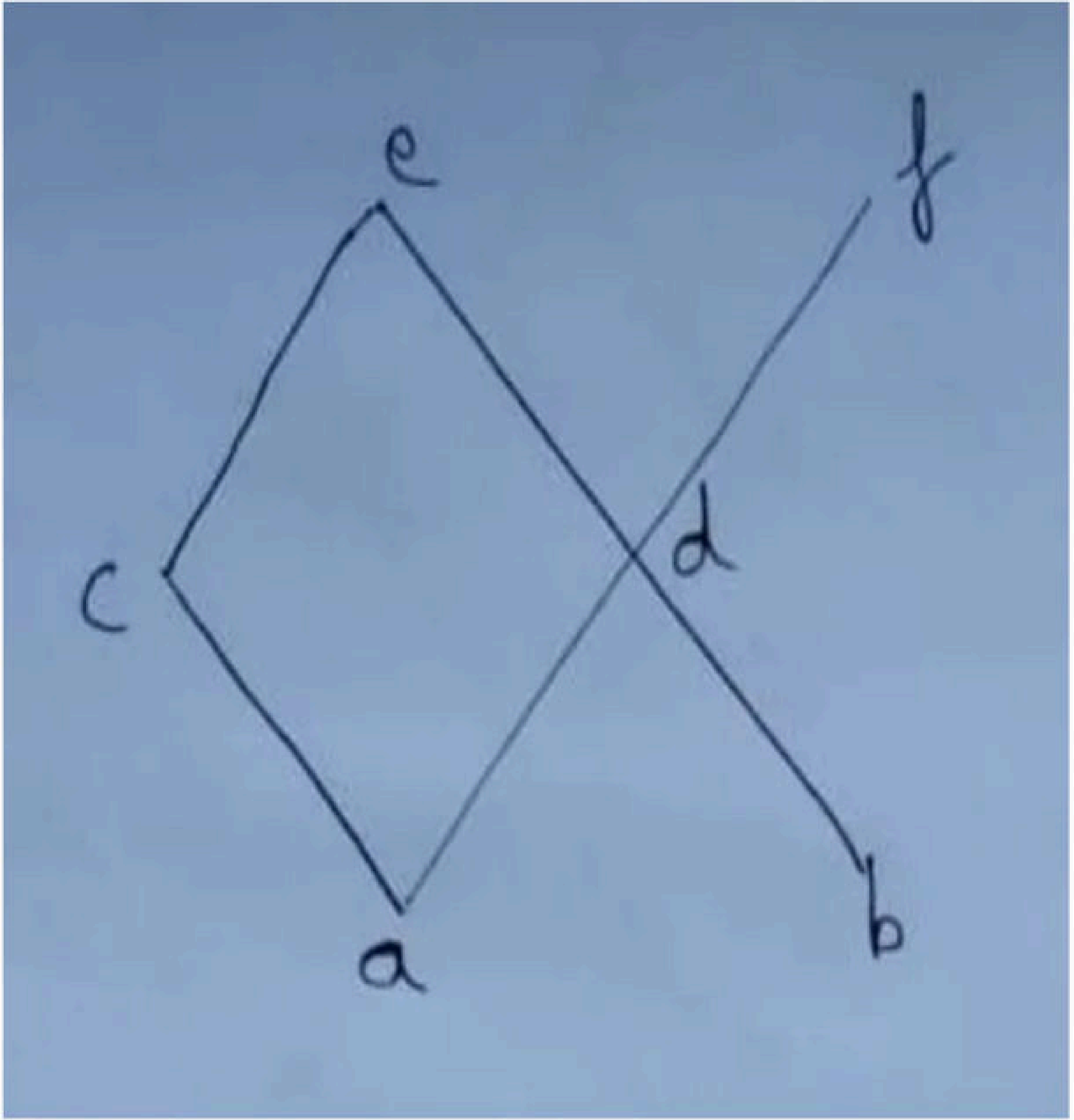
Elements of a Poset

1. **Maximal Element: -** An element is said to be maximal if it is not related to any other element in the Partial order relation.

2. **Minimal Element: -** An element is said to be minimal if no other element is related to it in the Partial order relation.



	Elements
Minimal	
Least	
Maximal	
Greatest	



	Elements
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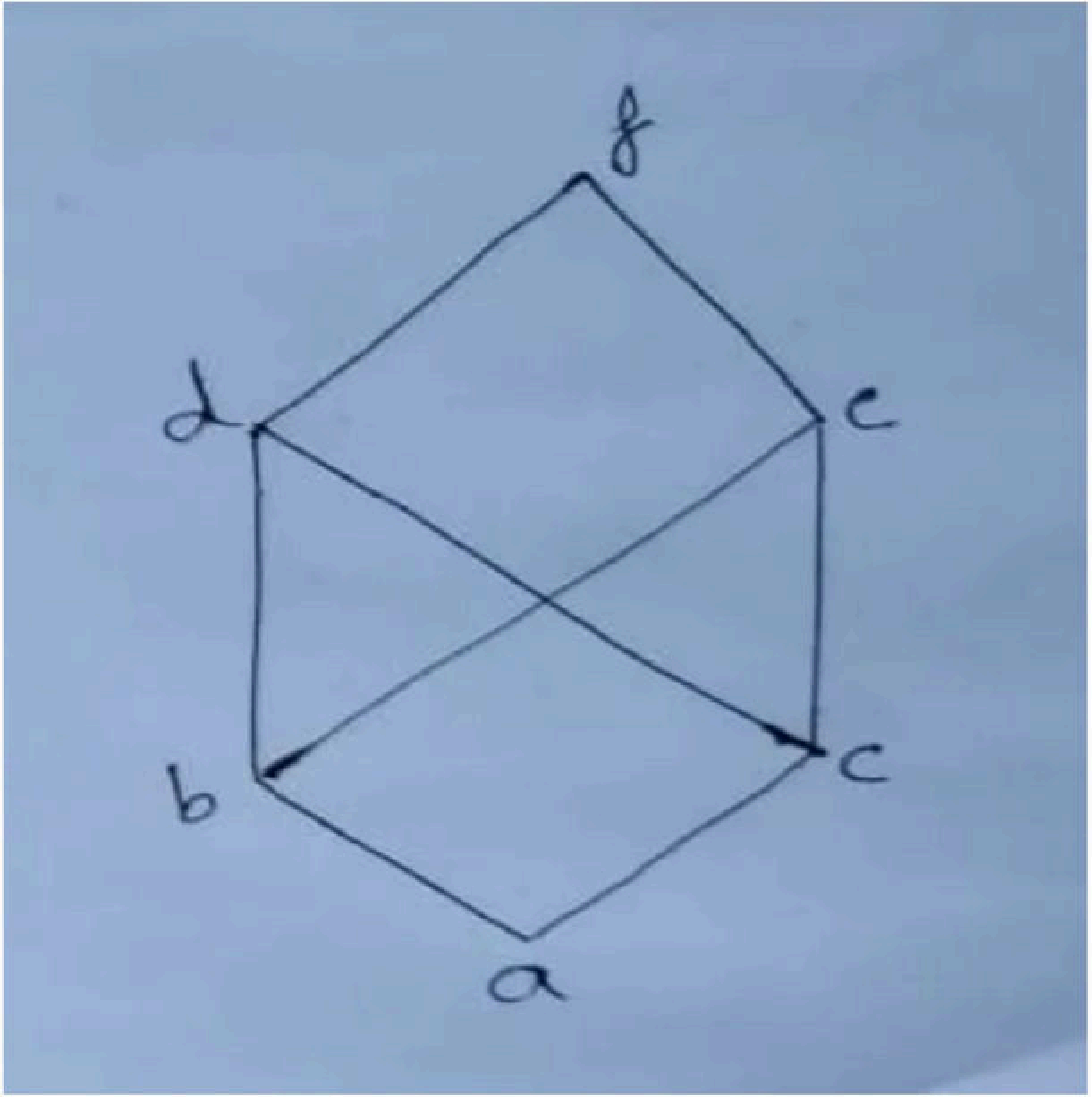
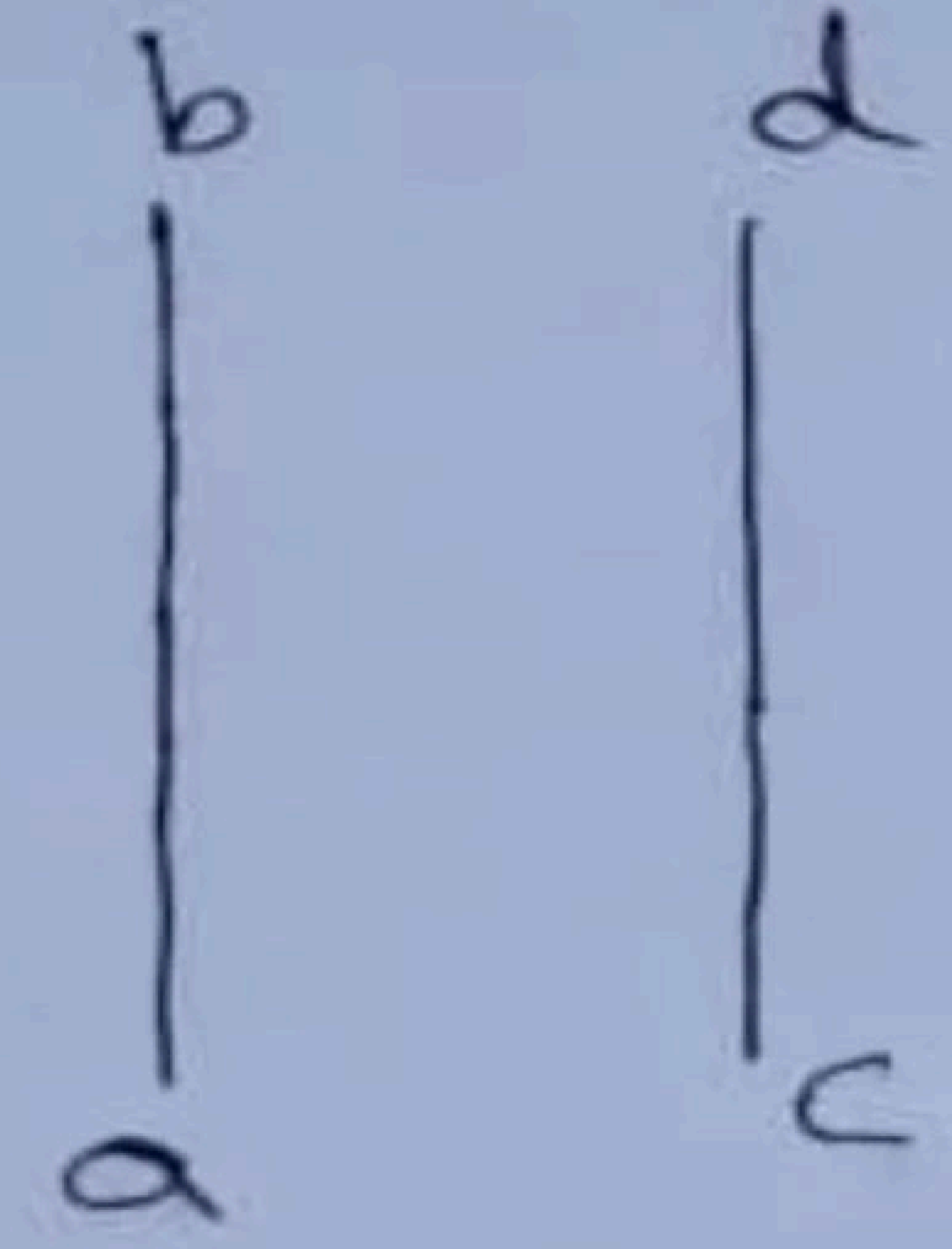
Least	
Maximal	
Greatest	



	Elements
Minimal	
Least	
Maximal	
Greatest	



	Elements
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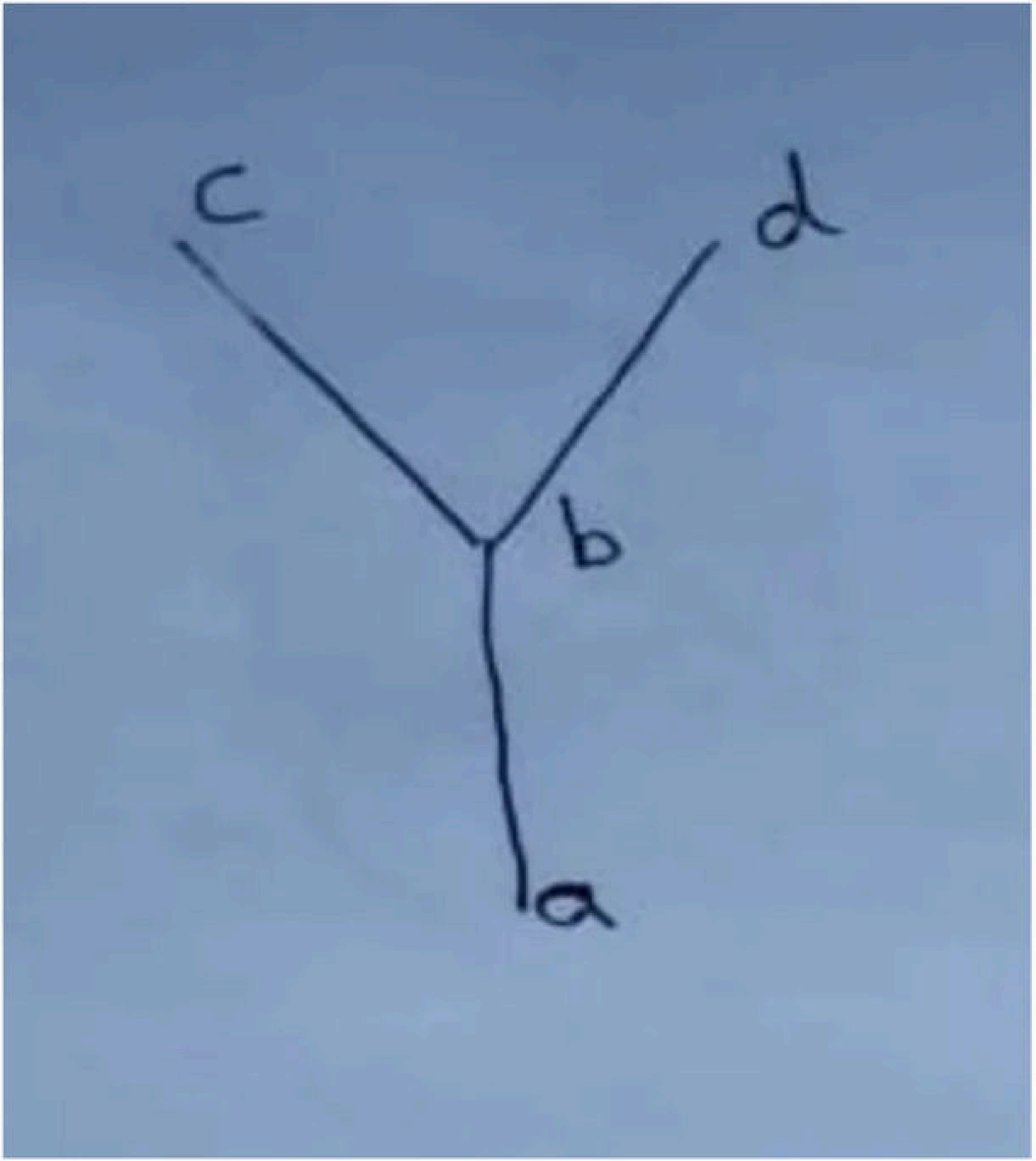
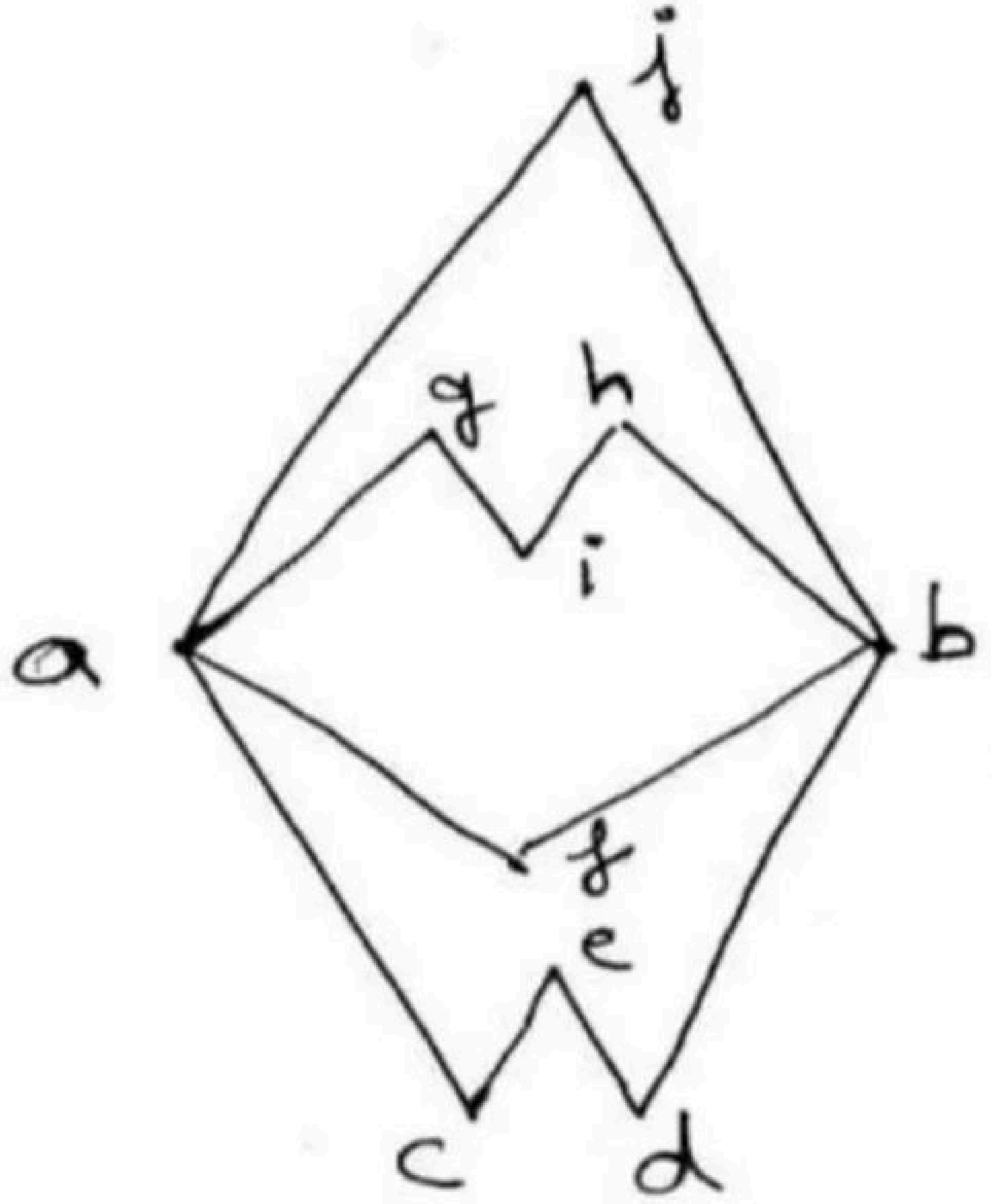
Least	
Maximal	
Greatest	



	Elements
Minimal	
Least	
Maximal	
Greatest	



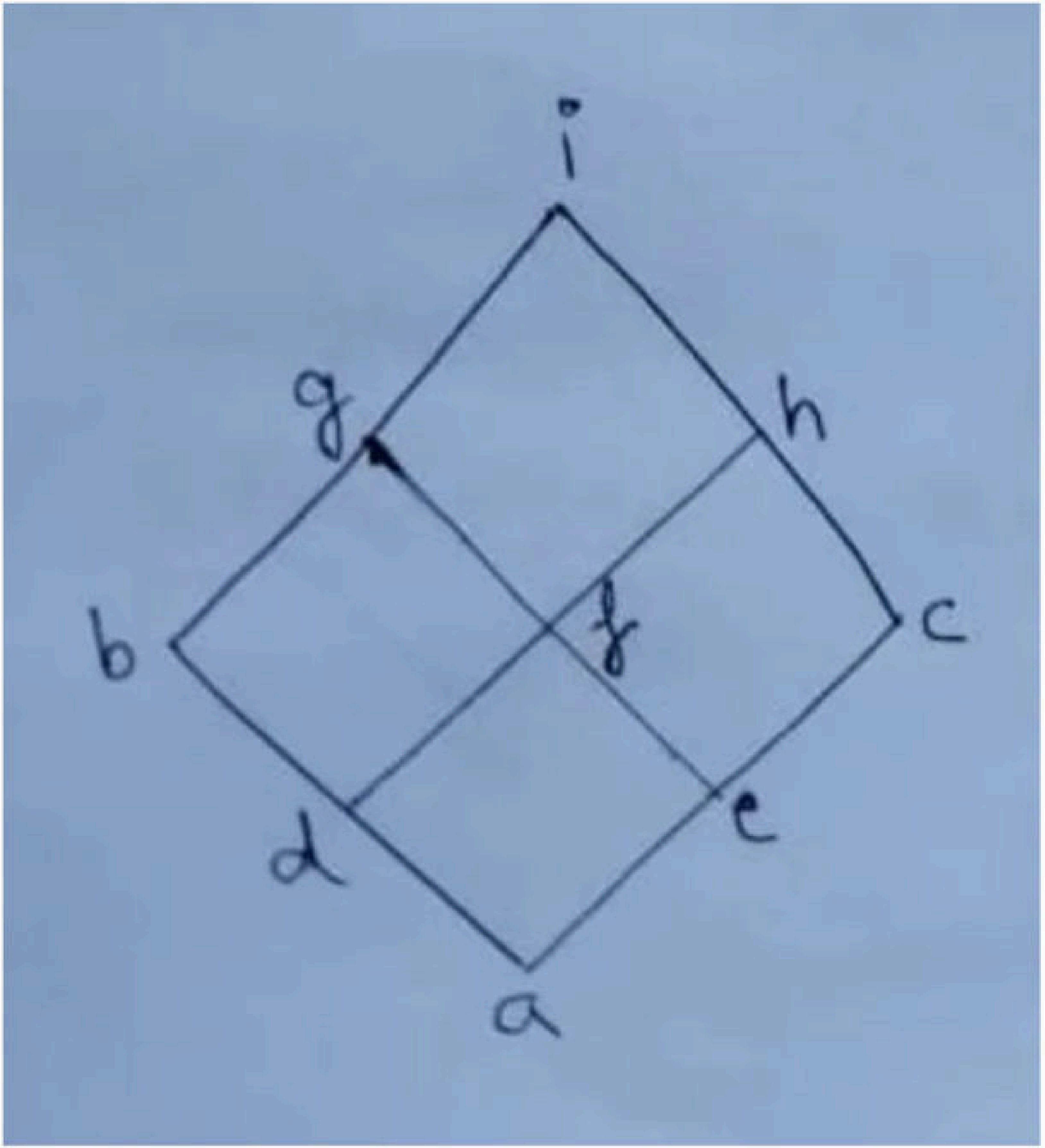
	Elements
Minimal	
Least	
Maximal	
Greatest	

- Every hasse diagram will have at least one Maximal and Minimal element(one or more).

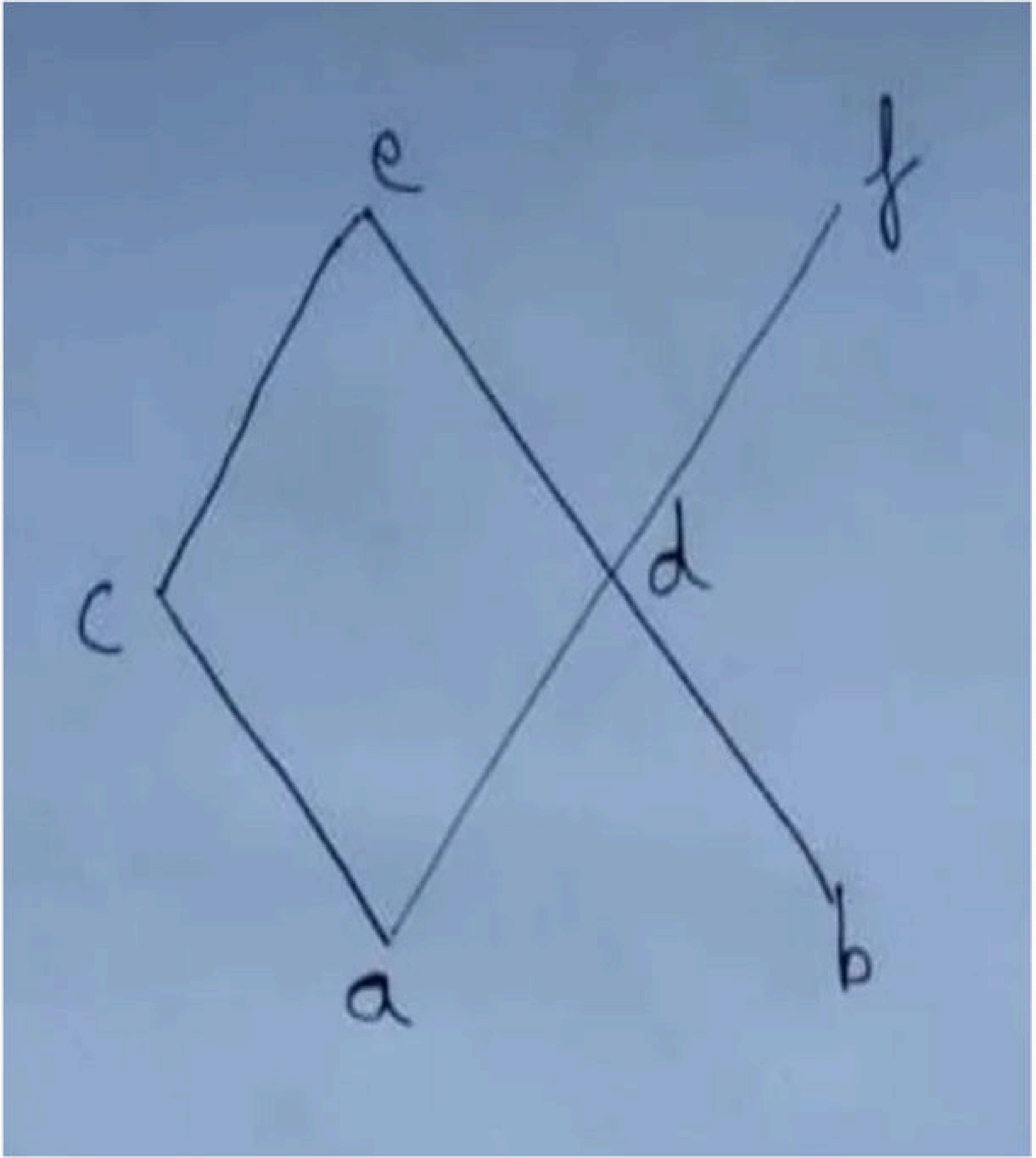
Break

- **<u>Greatest Element</u>: -** An element is said to be Maximum/Greatest if it is not related to any other element but every element is related to it in Partial order relation. Or if a hasse diagram has only one Maximal element then it will also be Maximum/Greatest element.

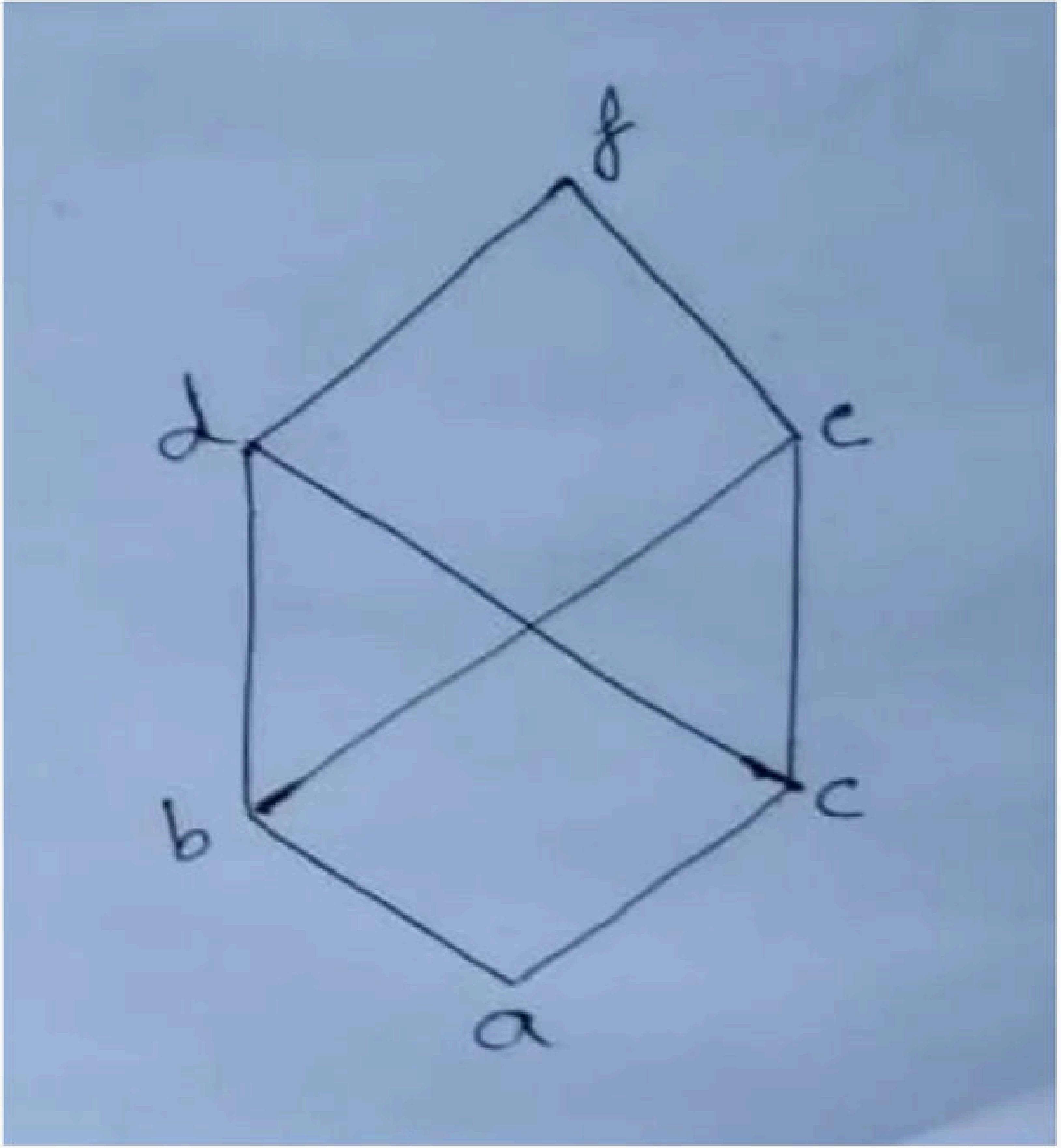
- **<u>Least Element</u>: -** An element is said to be Minimum/Least if no other element is related to it but it is related to every element Partial order relation. Or if a hasse diagram has only one Minimal element then it will also be Minimum/Least element.



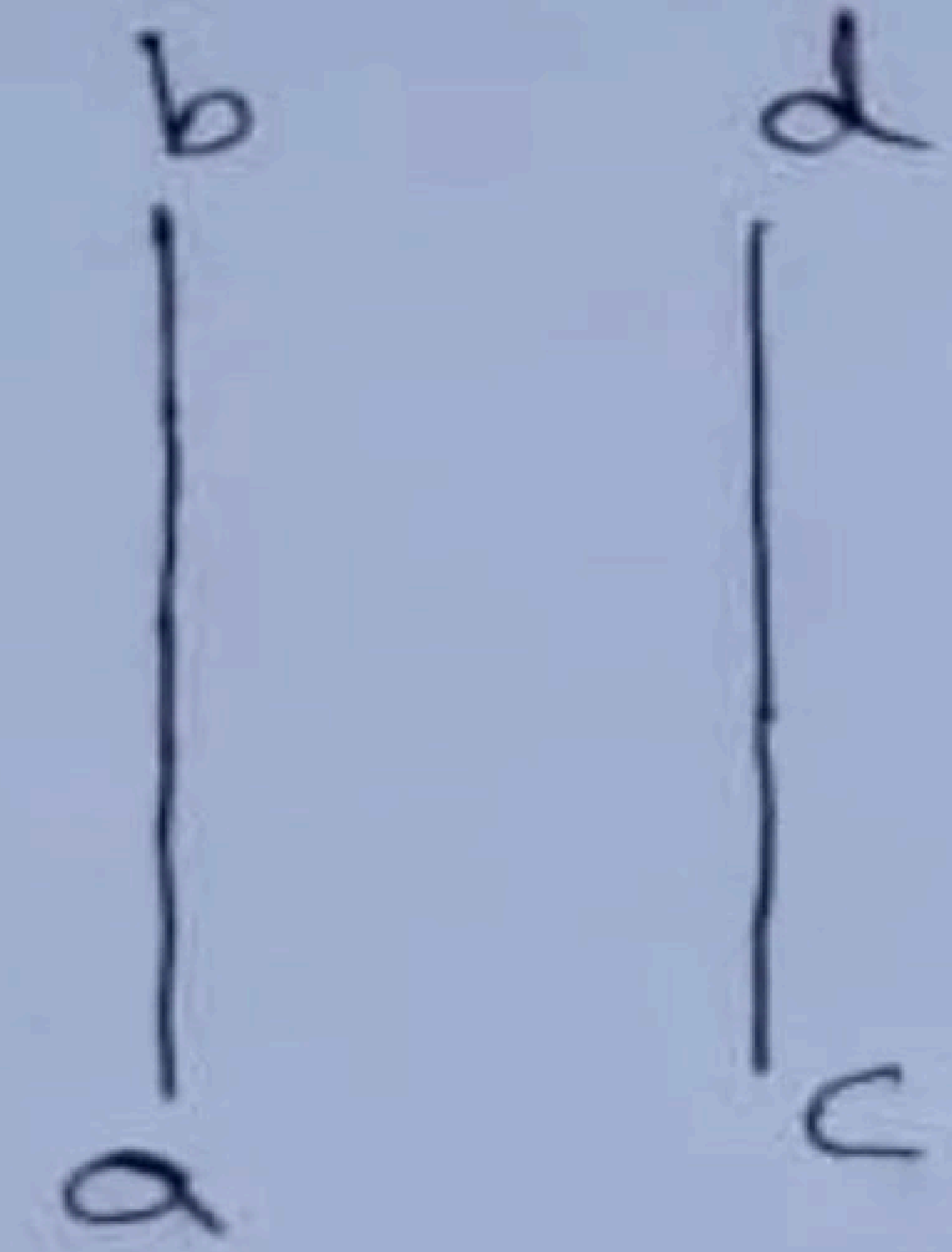
	Elements
Minimal	a
Least	
Maximal	i
Greatest	



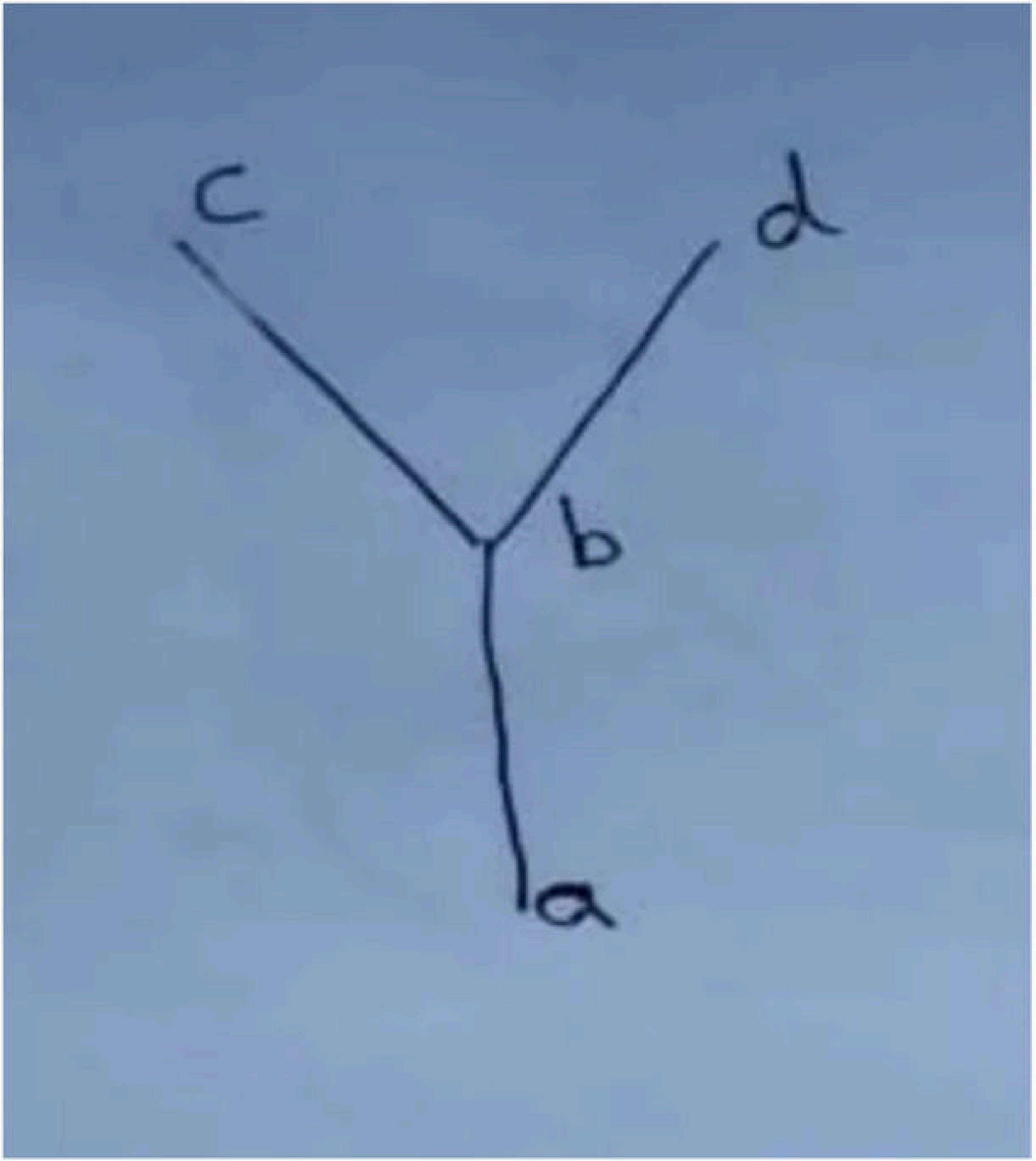
	Elements
Minimal	a, b
Least	
Maximal	e, f
Greatest	



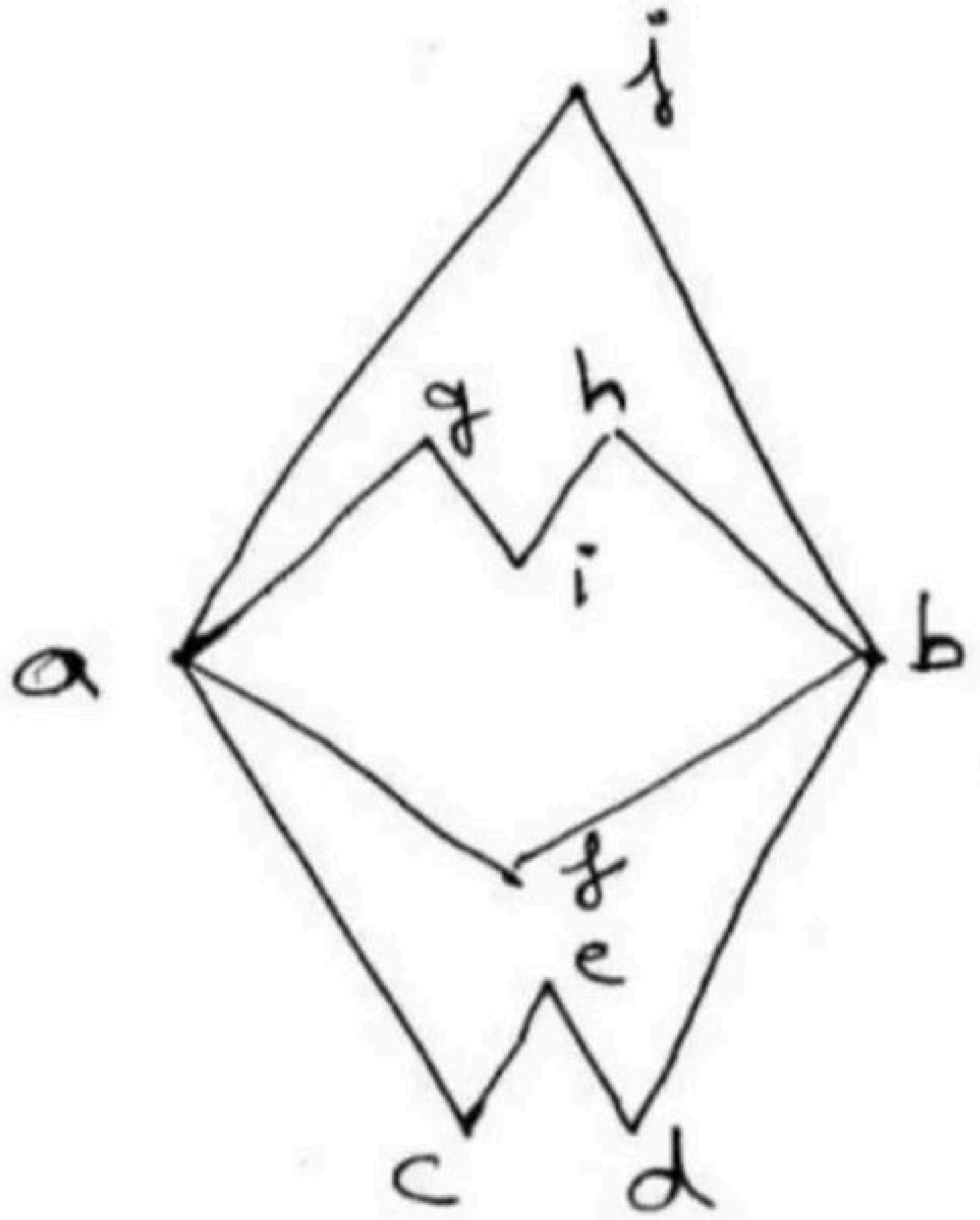
	Elements
Minimal	a
Least	
Maximal	f
Greatest	



	Elements
Minimal	a, c
Least	
Maximal	b, d
Greatest	



	Elements
Minimal	a
Least	
Maximal	c, d
Greatest	



	Elements
Minimal	c, d, f, i
Least	
Maximal	j, g, h, e
Greatest	

1. Every hasse diagram will have at most one Greatest and Least element(zero or one) (T/F)

2. Every Greatest element is also Maximal (T/F)

3. Every Least element is also Minimal (T/F)

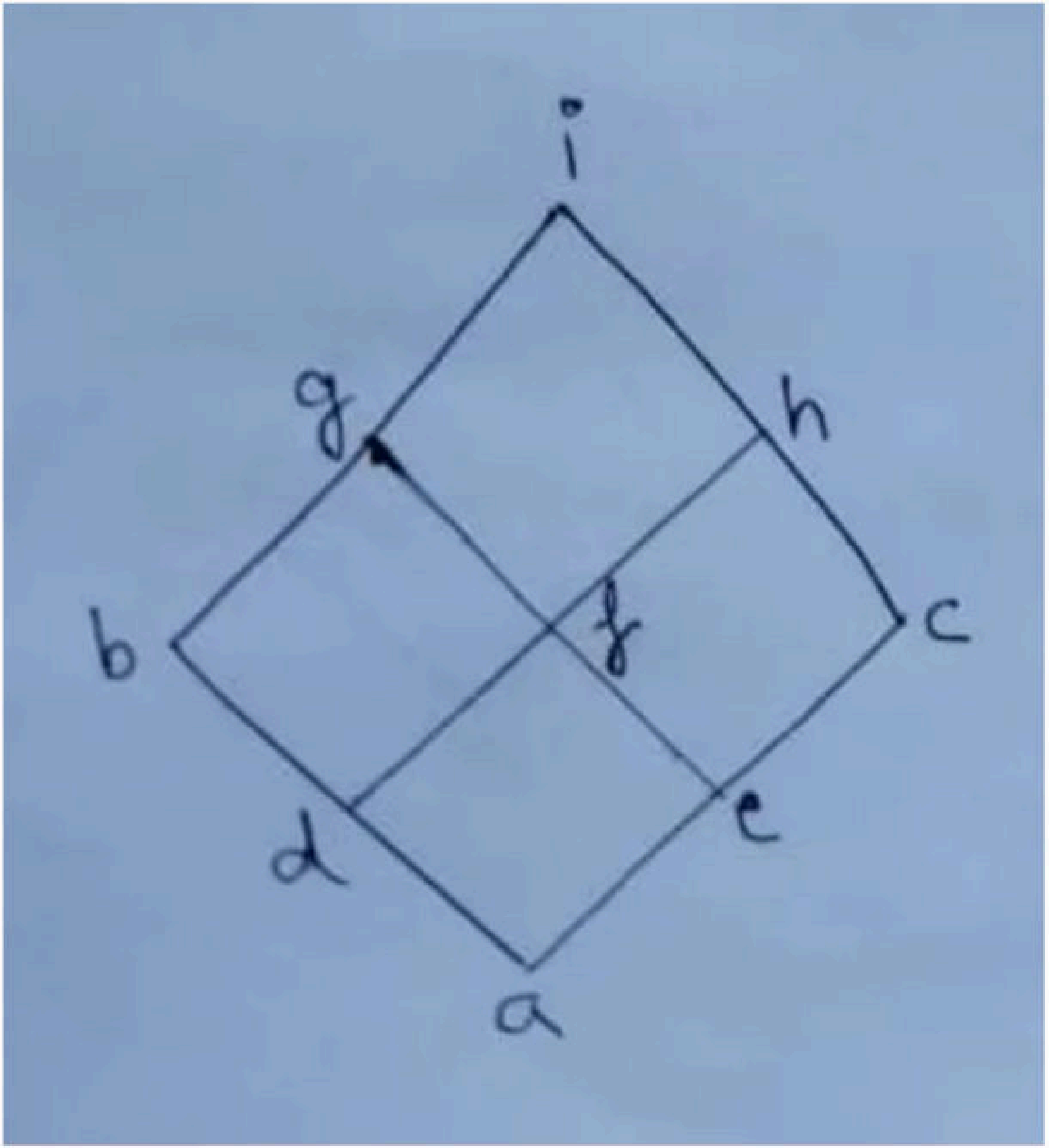
4. If there is only one Maximal element then it is called Greatest (T/F)

5. If there is only one Minimal element then it is called Least (T/F)

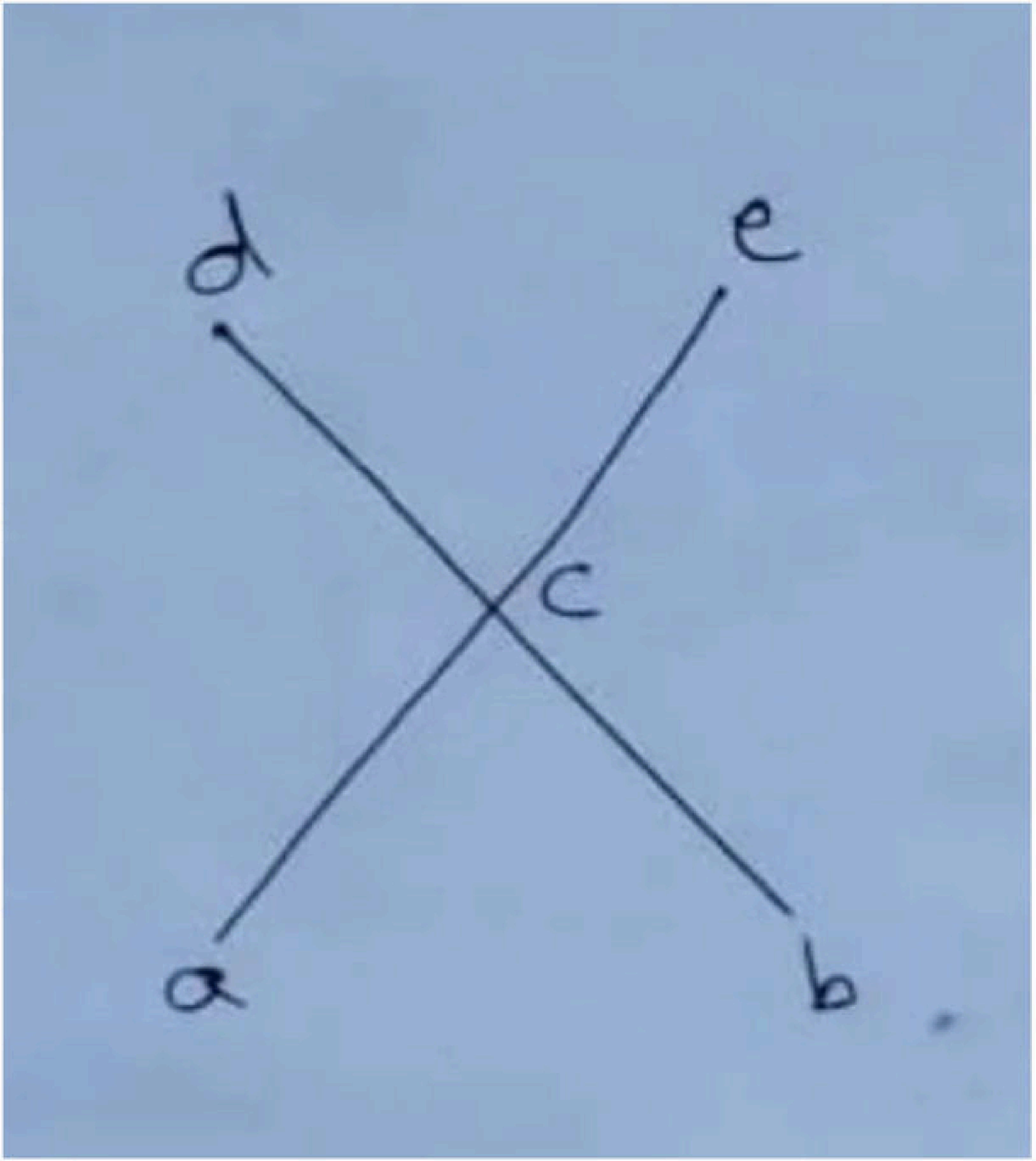
Break

- **Upper Bound: -** Upper bound of a subset B with respect to set A, will contain all those element to which all the elements of B is related.

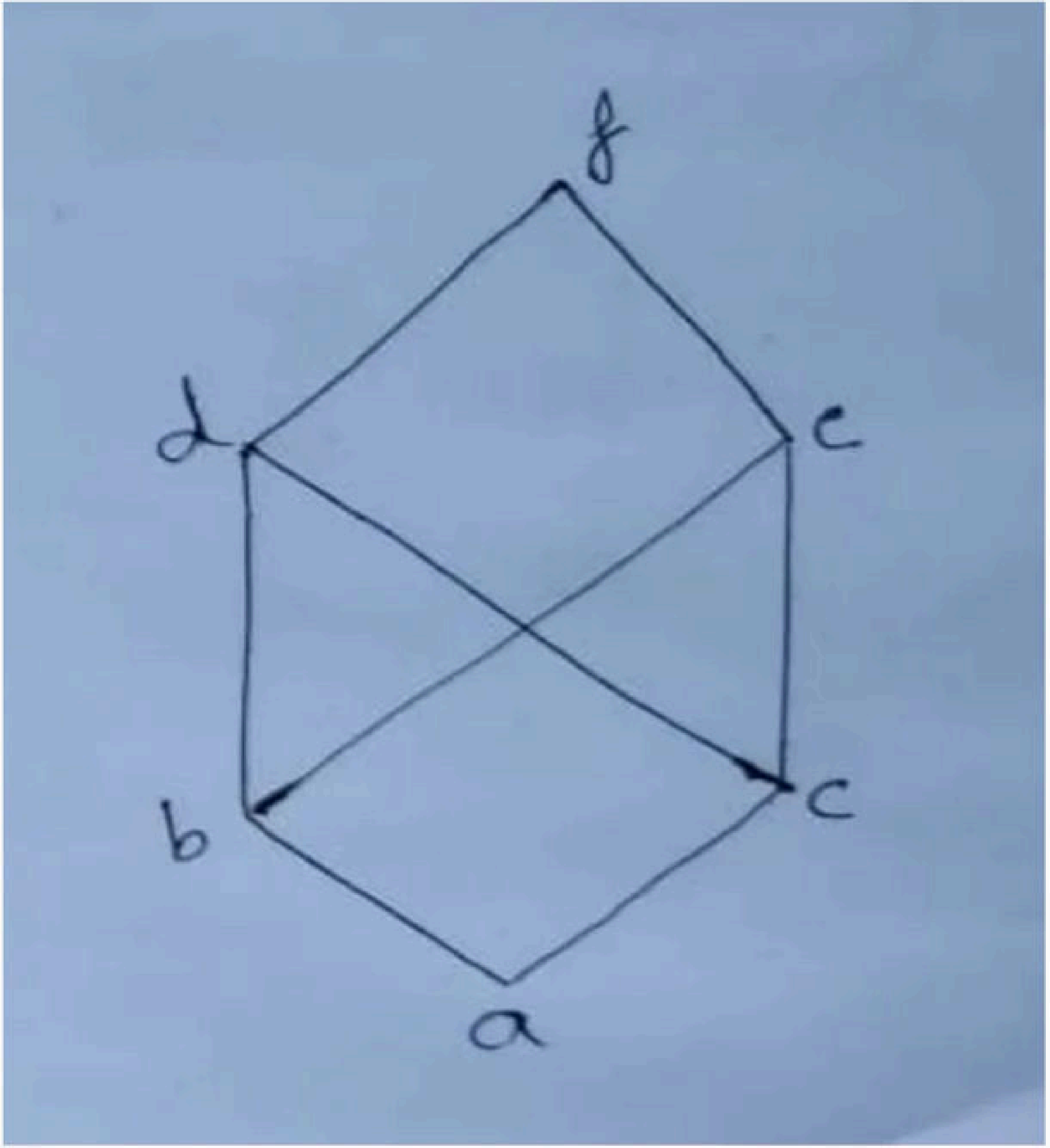
- **Lower Bound: -** lower bound of a subset B with respect to A, will contain all those elements which are related to every element of B.



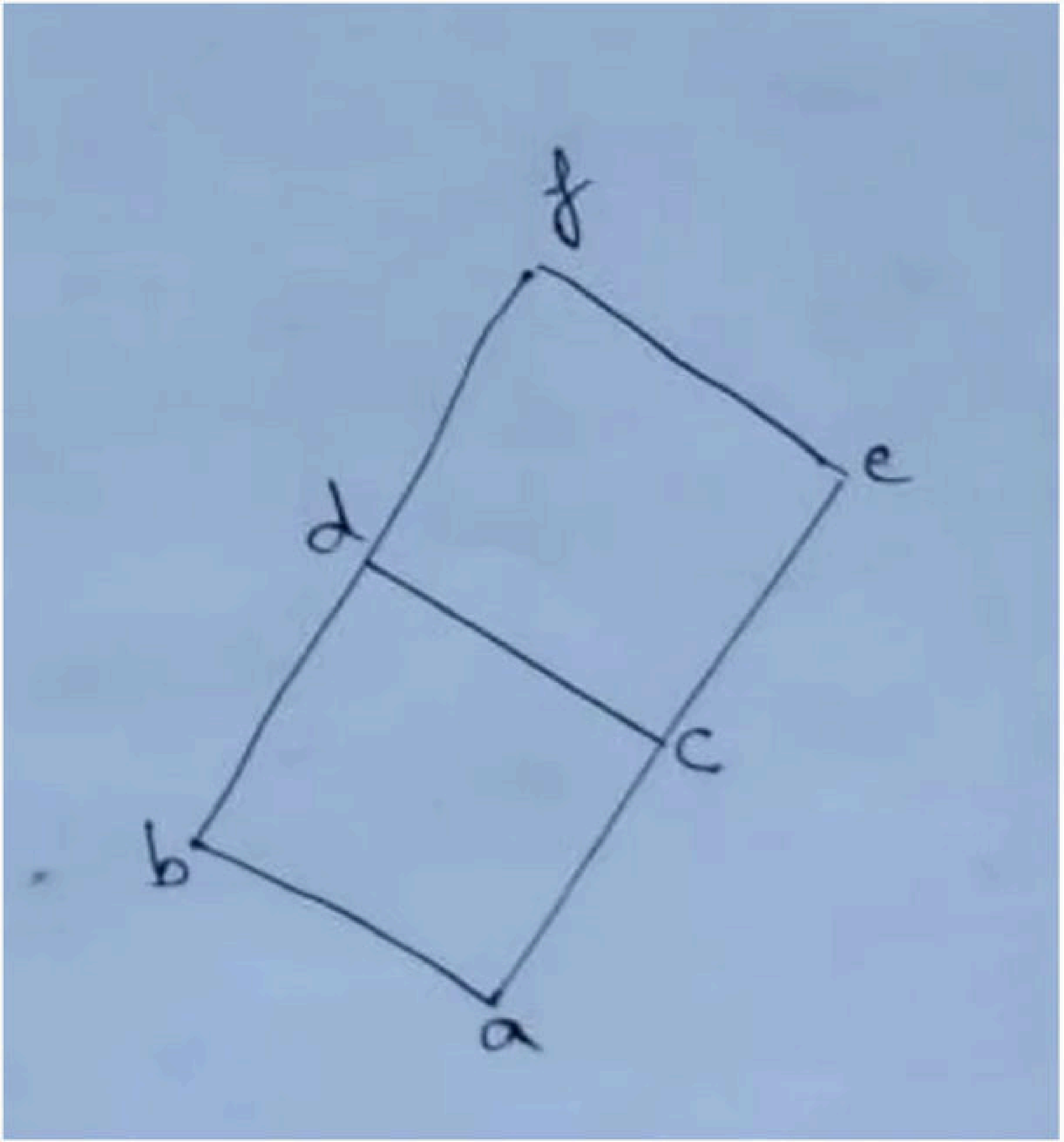
Elements	B = {g, f, e, c}	B = {d, f}
Upper Bound		
Lower Bound		



Elements	B = {d, e}	B = {c}
Upper Bound		
Lower Bound		



Elements	B = {d, e}	B = {b, c}
Upper Bound		
Lower Bound		



Elements	B = {d, c}	B = {b, e}
Upper Bound		
Lower Bound		

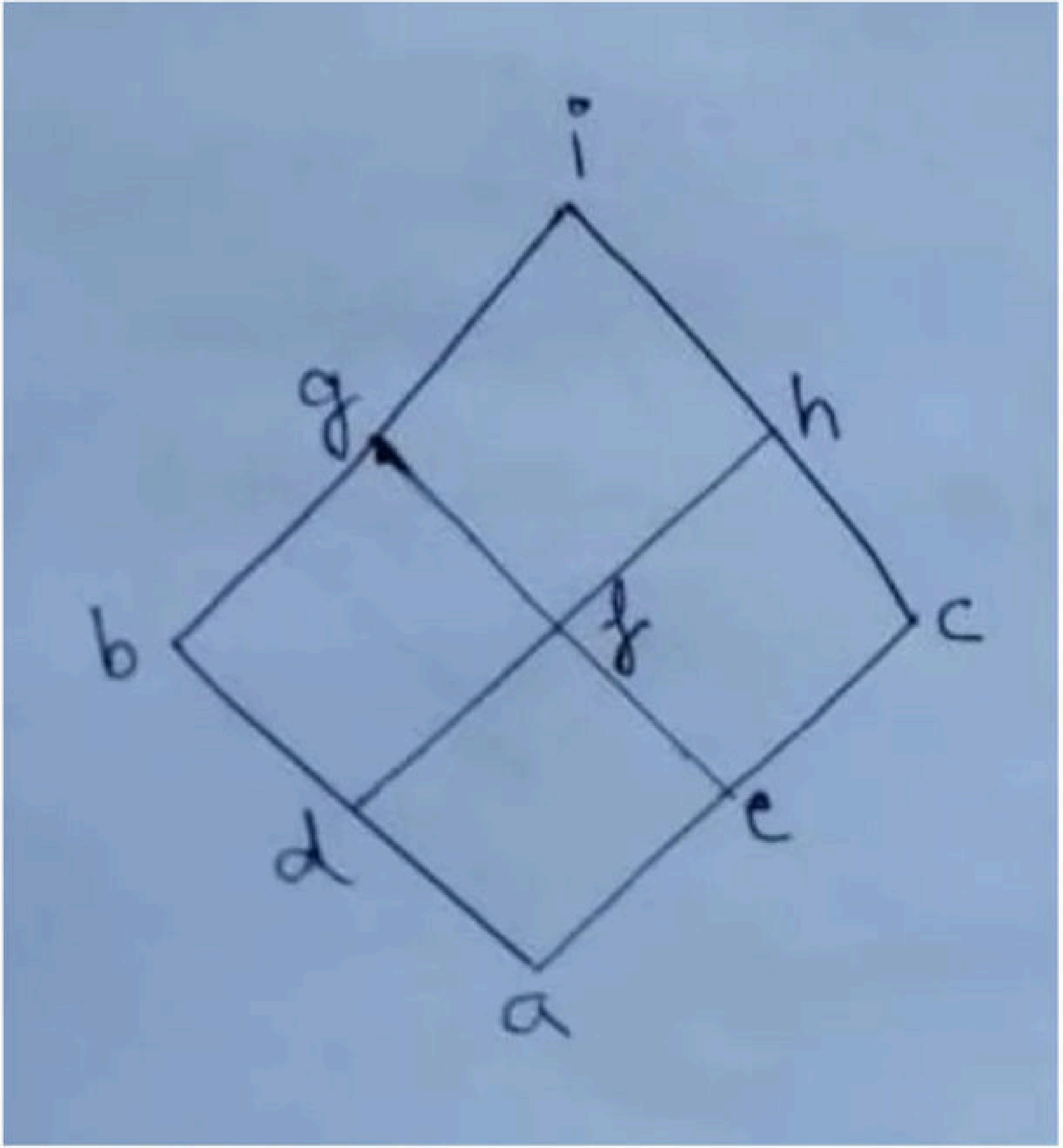
Break

Least Upper Bound / LUB / Join / Supremum / $\vee$

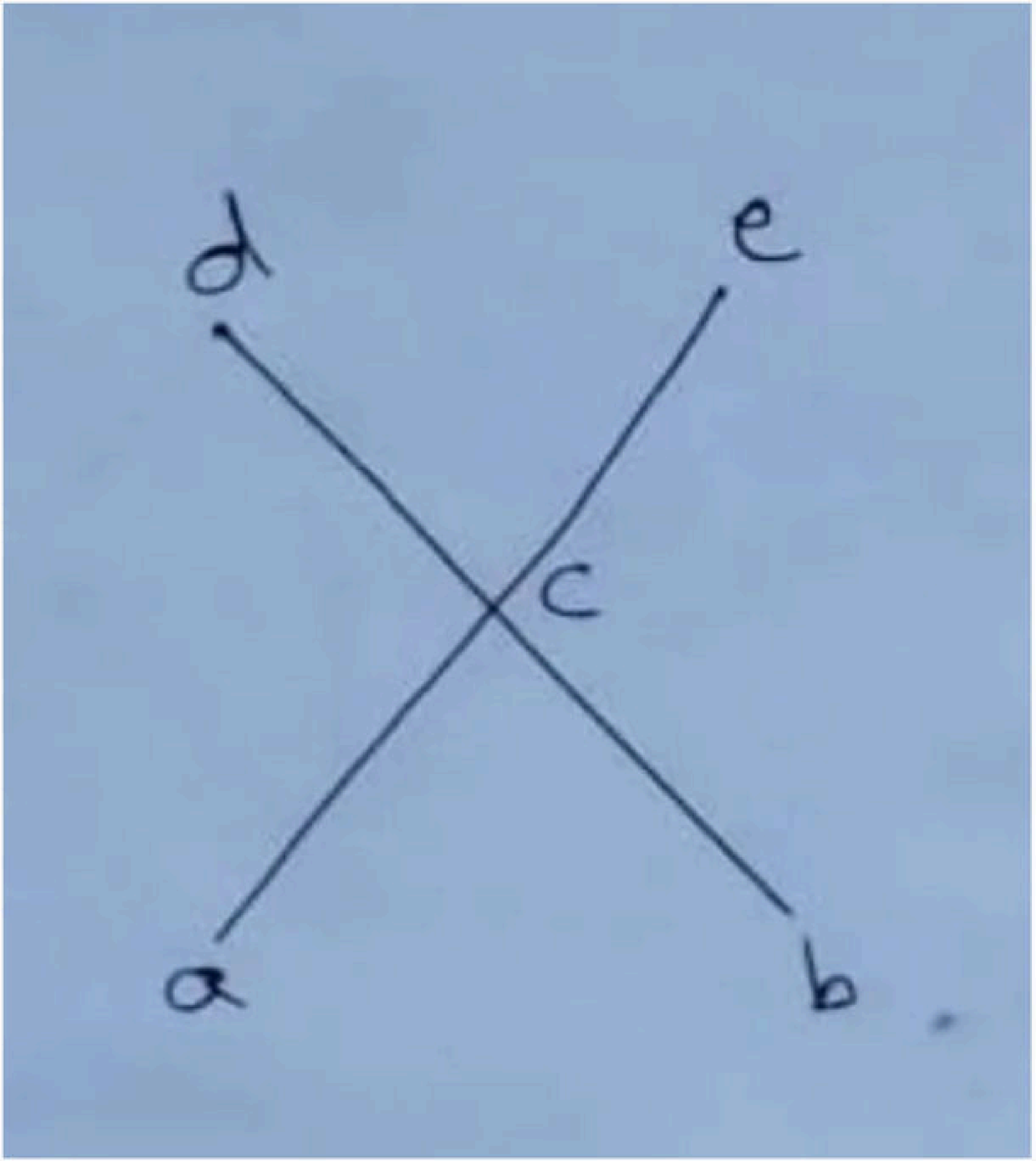
Least value in the upper bound

Greatest Lower Bound / GLB / Meet / Infimum / ∧

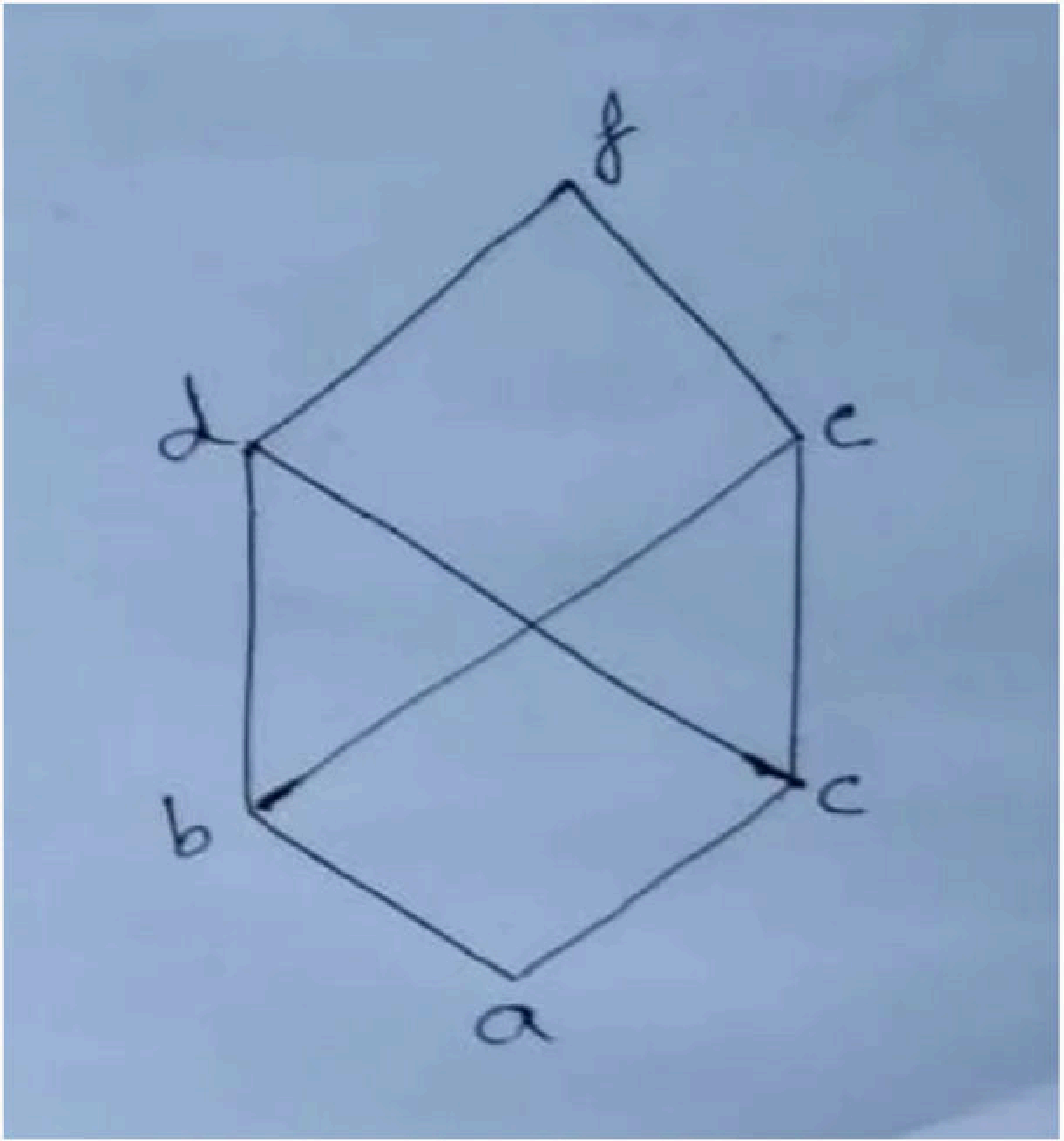
Greatest value in the lower bound



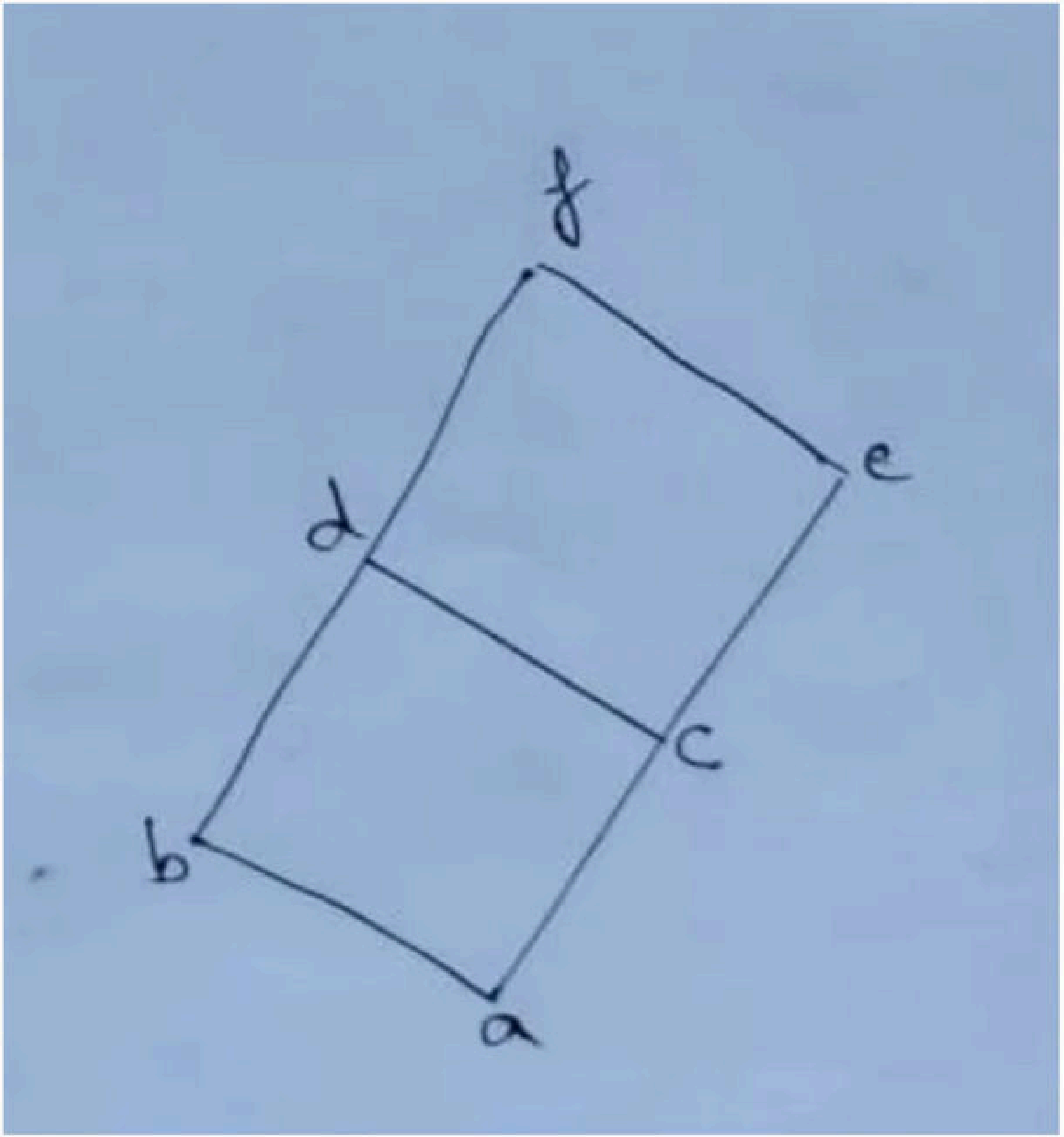
Elements	B = {g, f, e, c}	B = {d, f}
Upper Bound	{i}	{i, g, h, f}
Least Upper Bound		
Lower Bound	{a, e}	{a, d}
Greatest Lower Bound		



Elements	B = {d, e}	B = {c}
Upper Bound	{}	{d, e, c}
Least Upper Bound		
Lower Bound	{a, b, c}	{a, b, c}
Greatest Lower Bound		



Elements	B = {d, e}	B = {b, c}
Upper Bound	{f}	{d, e, f}
Least Upper Bound		
Lower Bound	{a, b, c}	{a}
Greatest Lower Bound		



Elements	B = {d, c}	B = {b, e}
Upper Bound	{f, d}	{f}
Least Upper Bound		
Lower Bound	{a, c}	{a}
Greatest Lower Bound		

Q Consider the Poset ({3,5,9,15,24,45}, /). Which of the following is correct for the given Poset? **(NET-JUNE-2019)**

a) There exist a greatest element and a least element

b) There exist a greatest element but not a least element

c) There exist a least element but not a greatest element

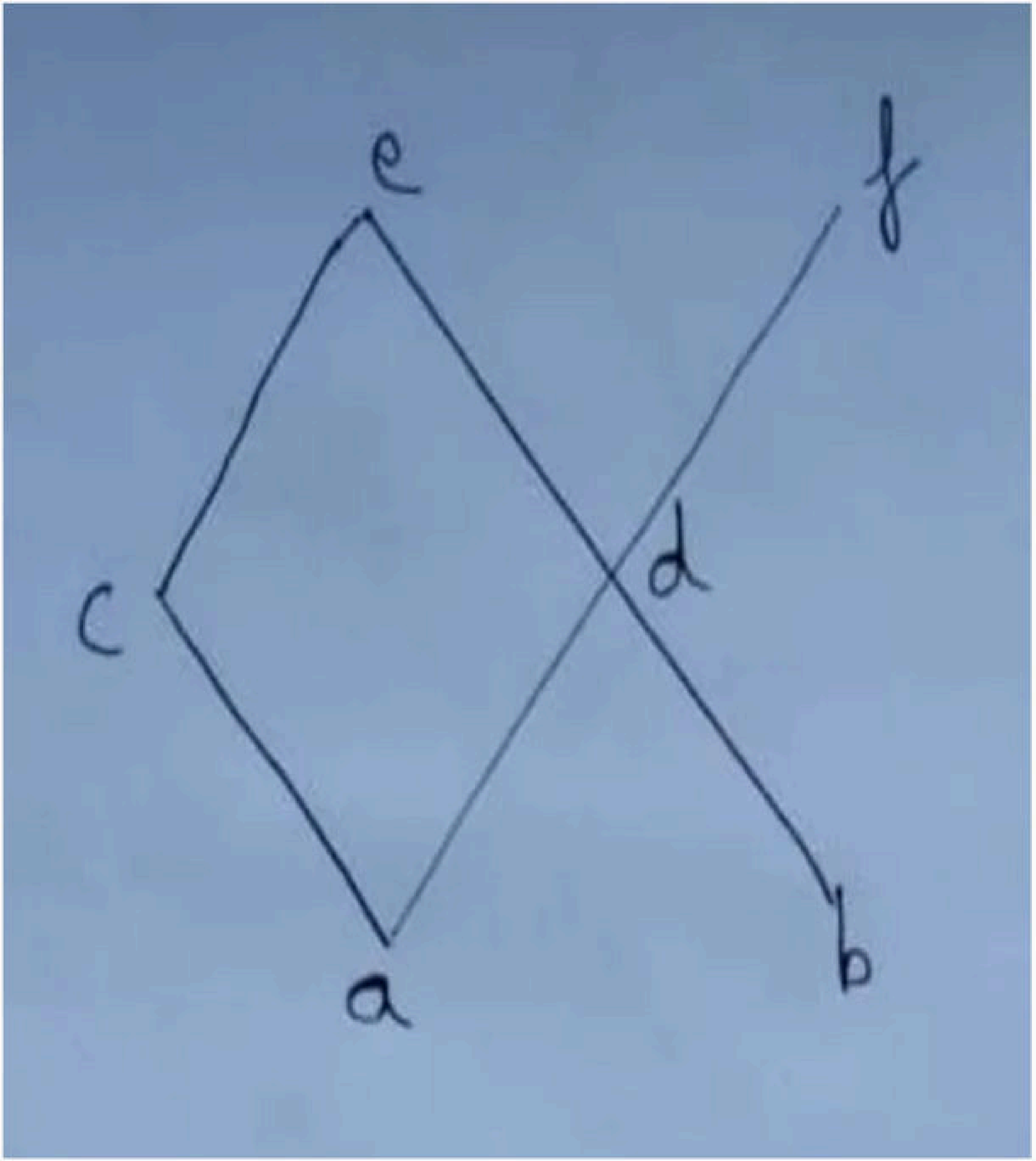
d) There does not exist a greatest element and a least element

Break

Join Semi Lattice :- A hasse diagram/Partial order relation is called Join Semi Lattice if for every elements their exists a Join.

<u>Meet Semi Lattice</u> :- A hasse diagram/Partial order relation is called Meet Semi Lattice if for every elements their exists a Meet.

<u>Lattice</u> :- A hasse diagram/Partial order relation is called Lattice if their exist a Join and Meet for every pair of element. Or A hasse diagram/Partial order relation is called Lattice if it is both Join Semi Lattice and Meet Semi Lattice.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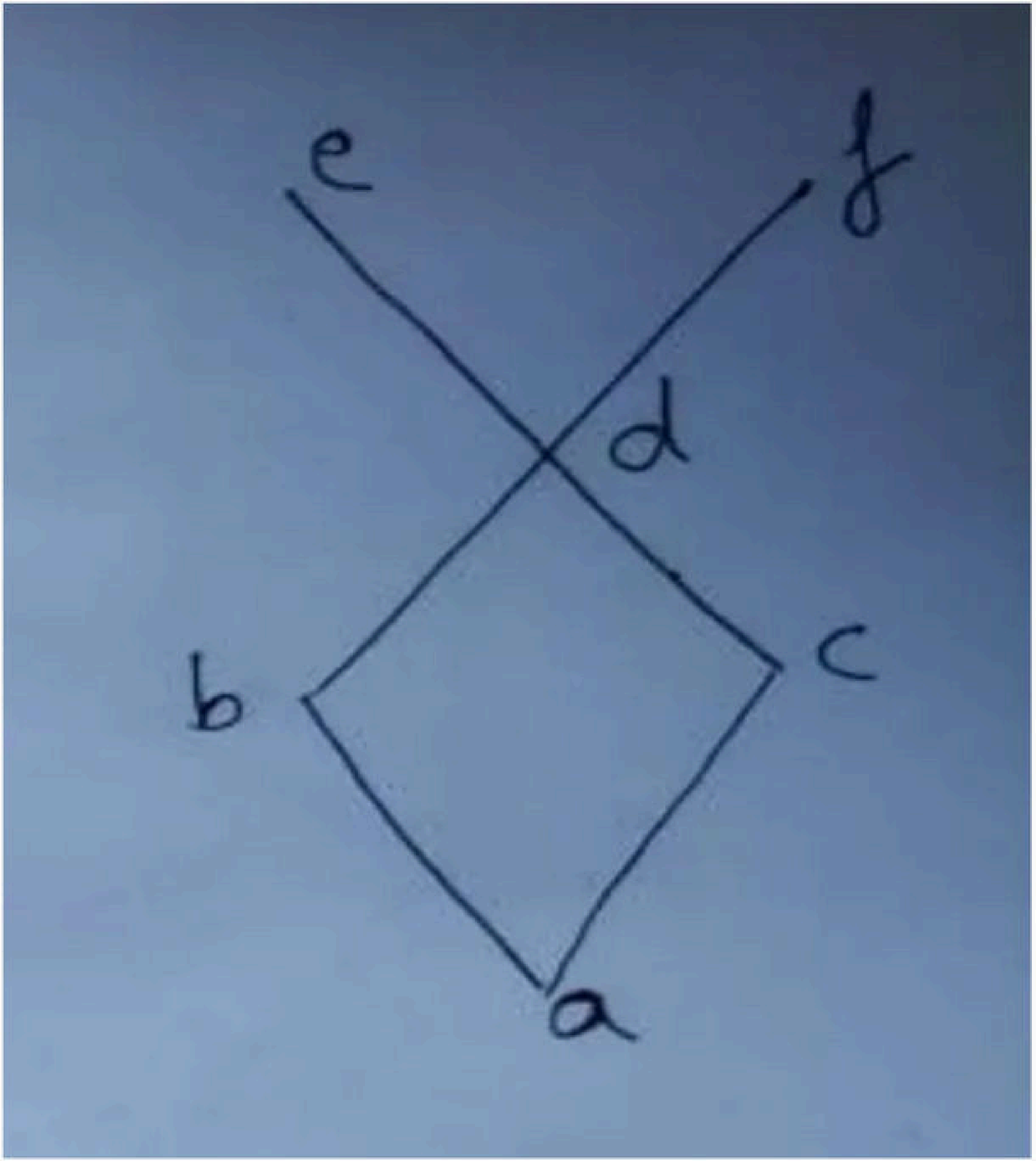
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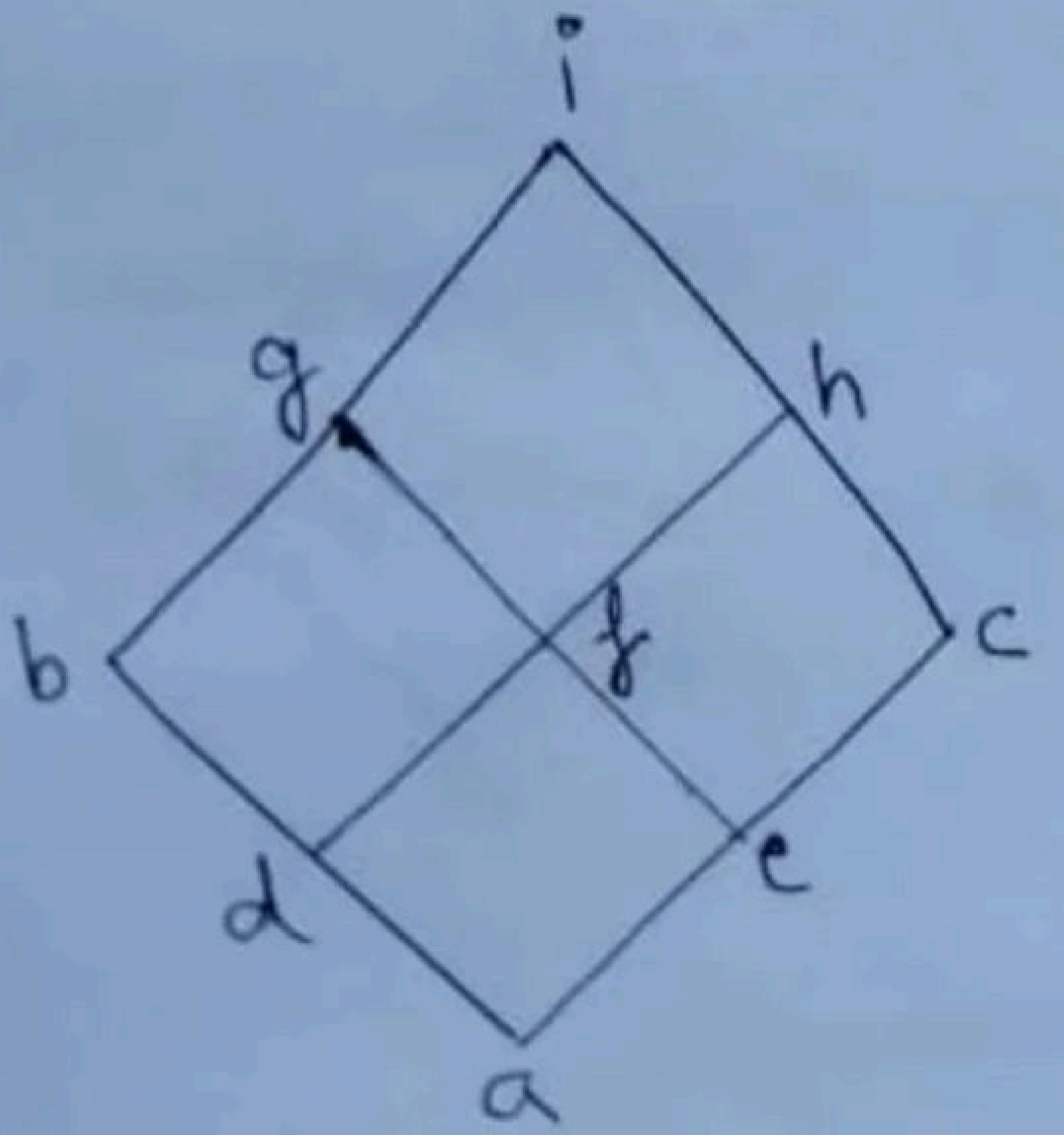
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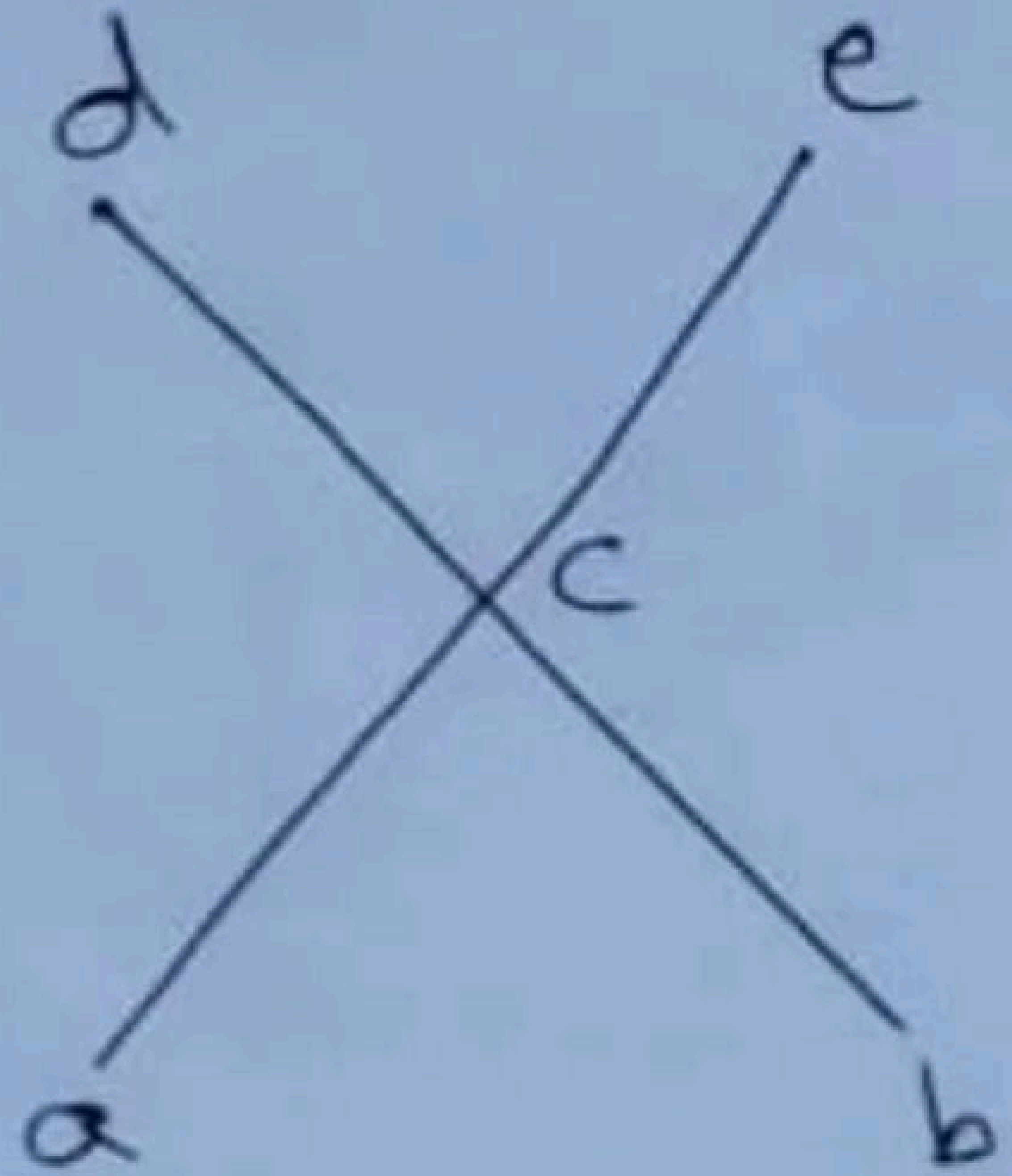
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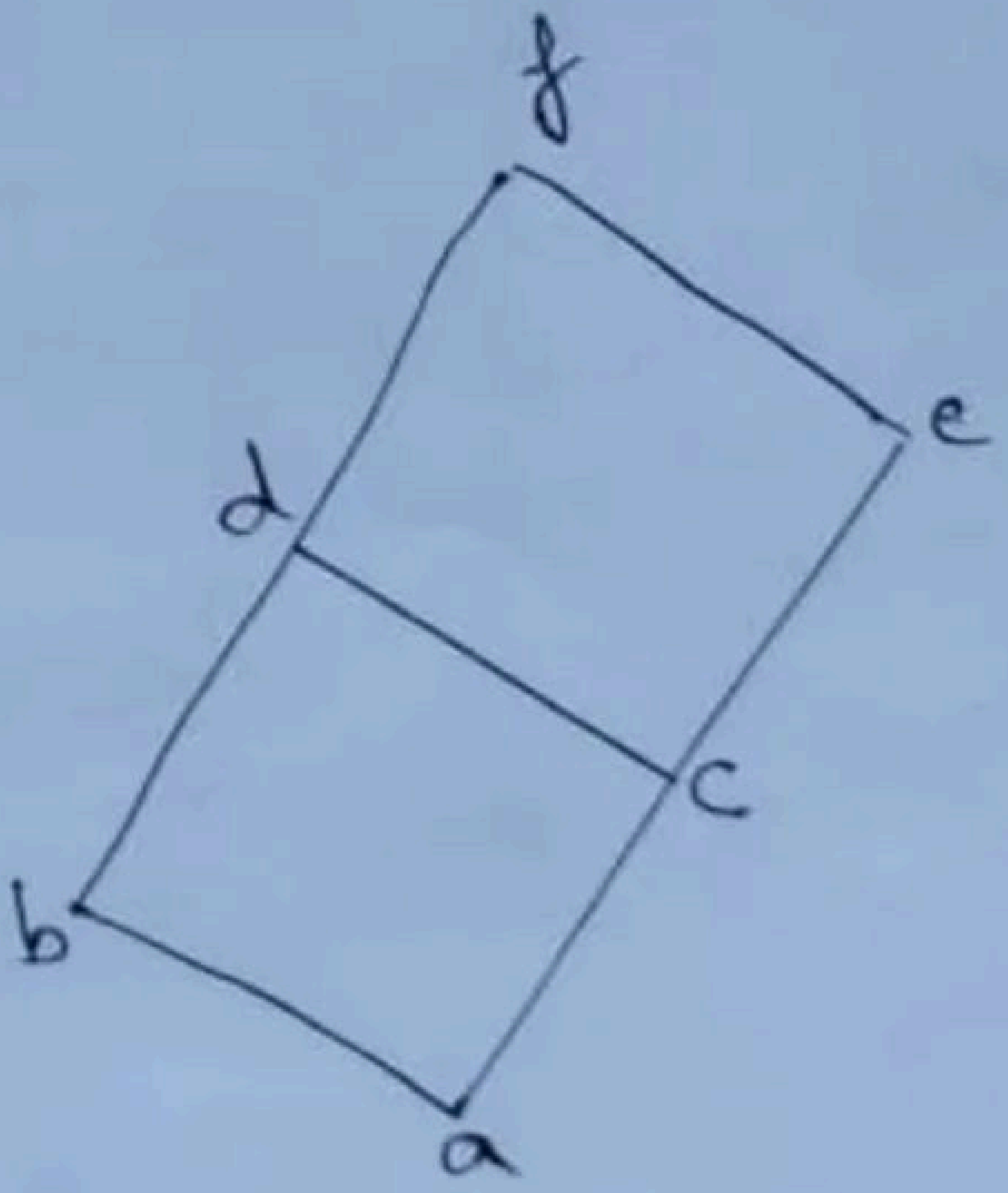
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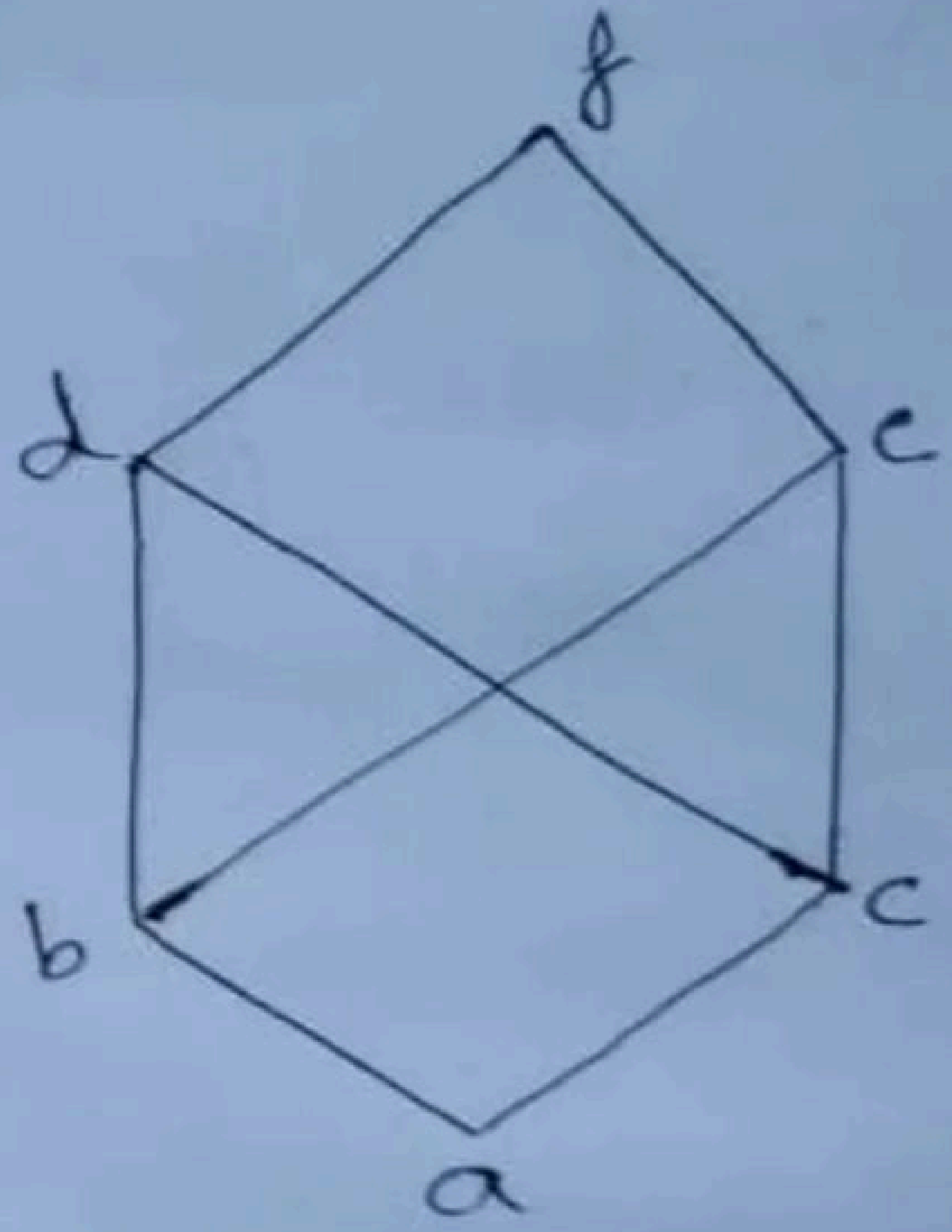
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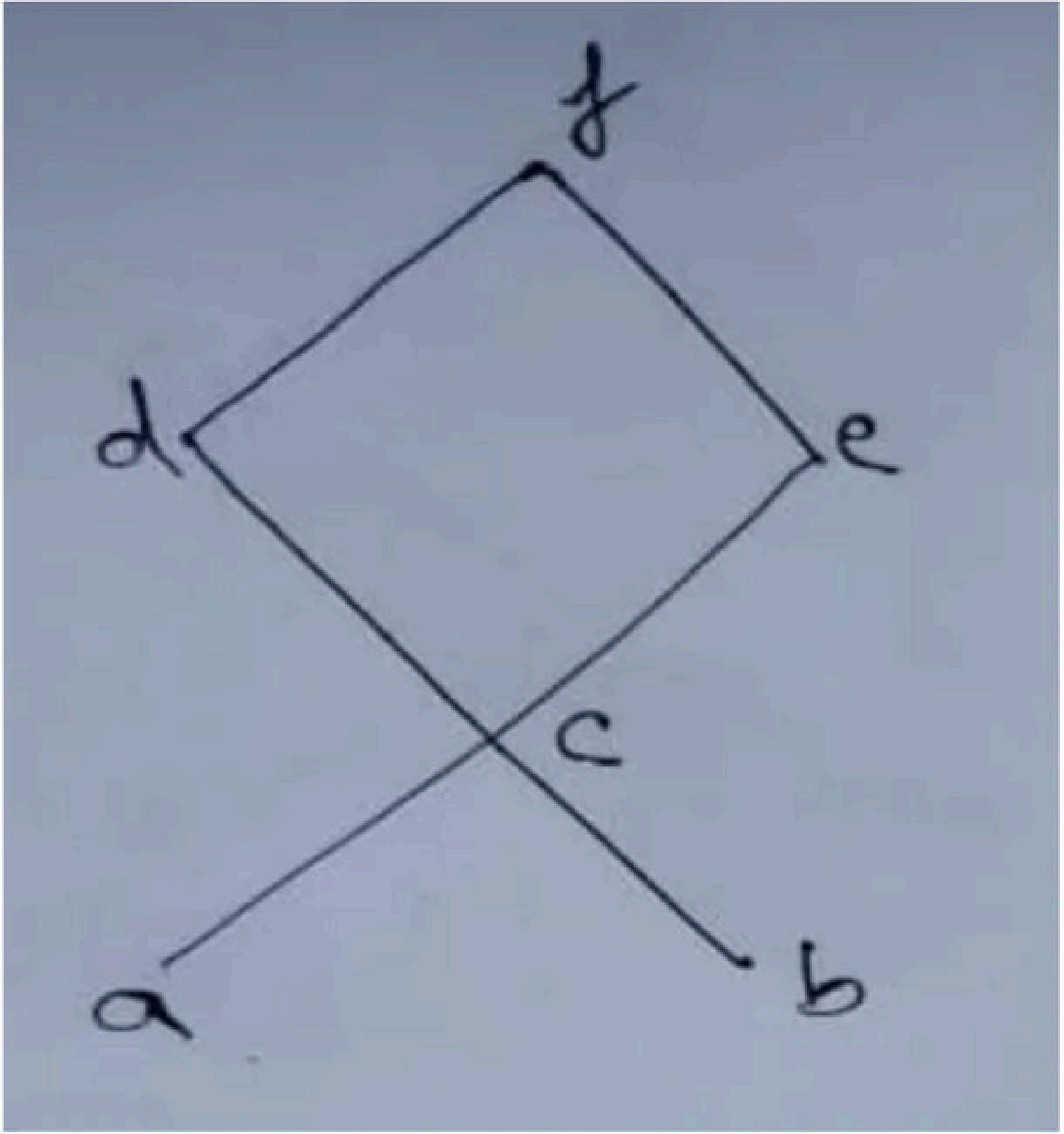
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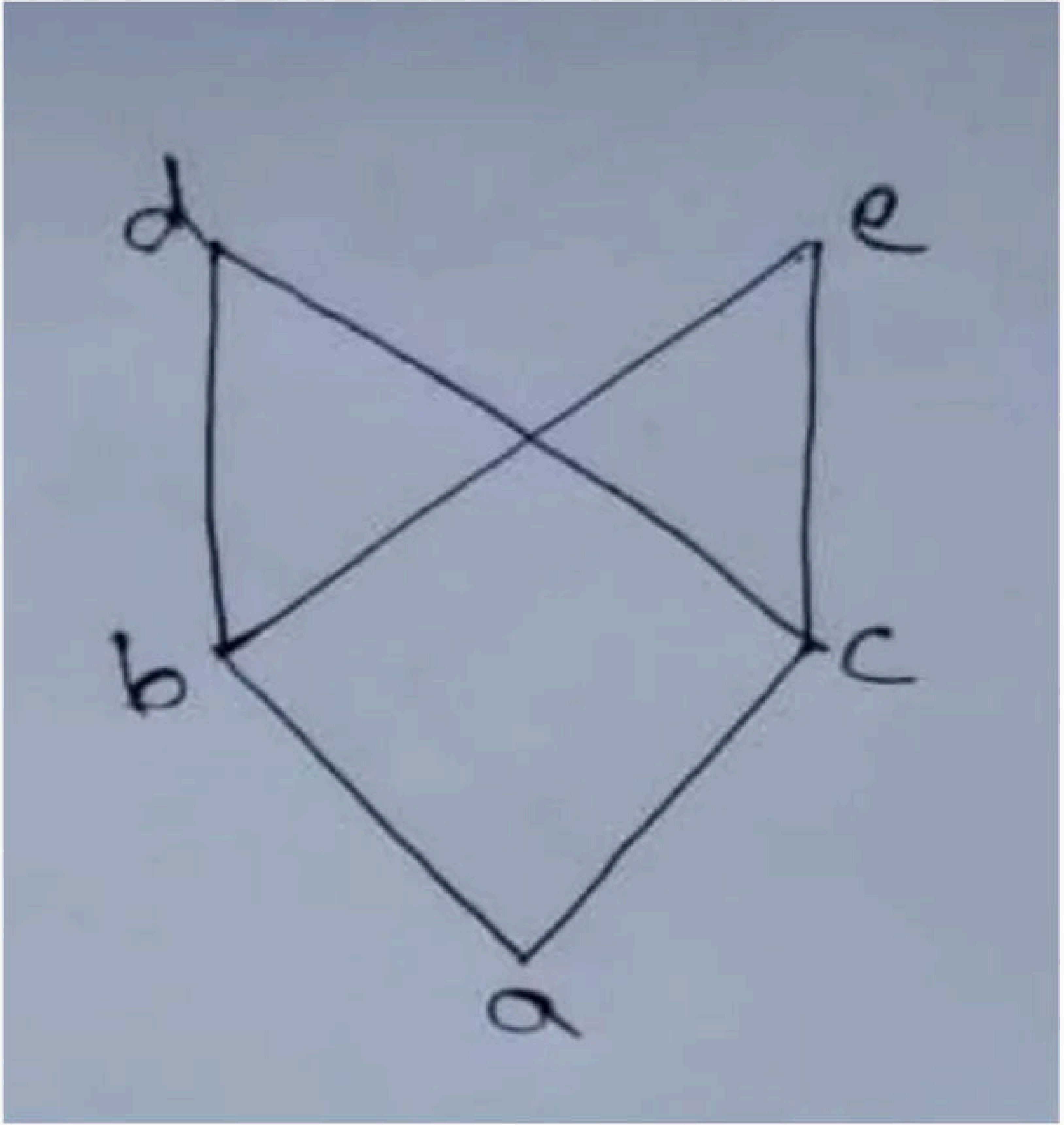
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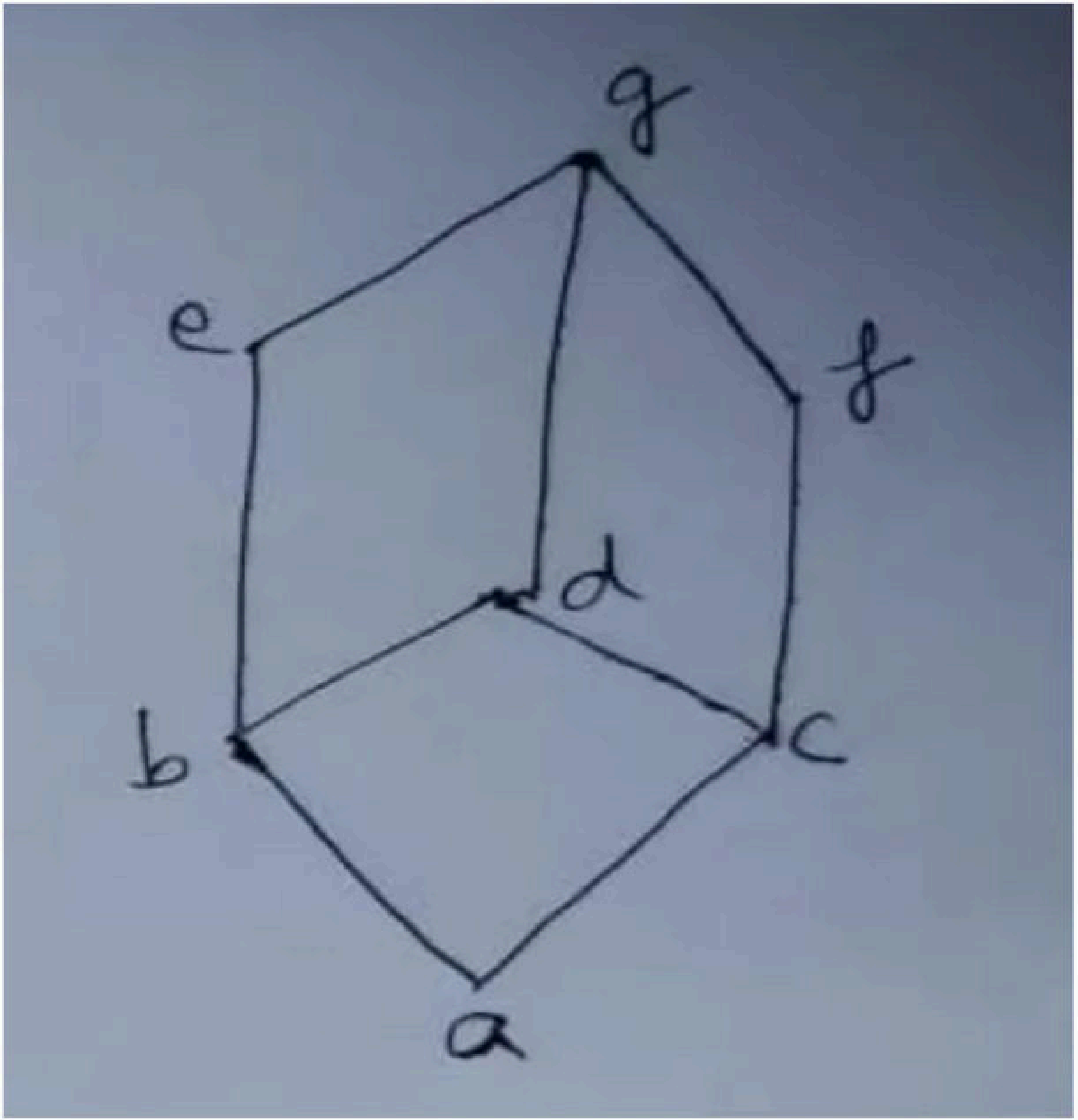
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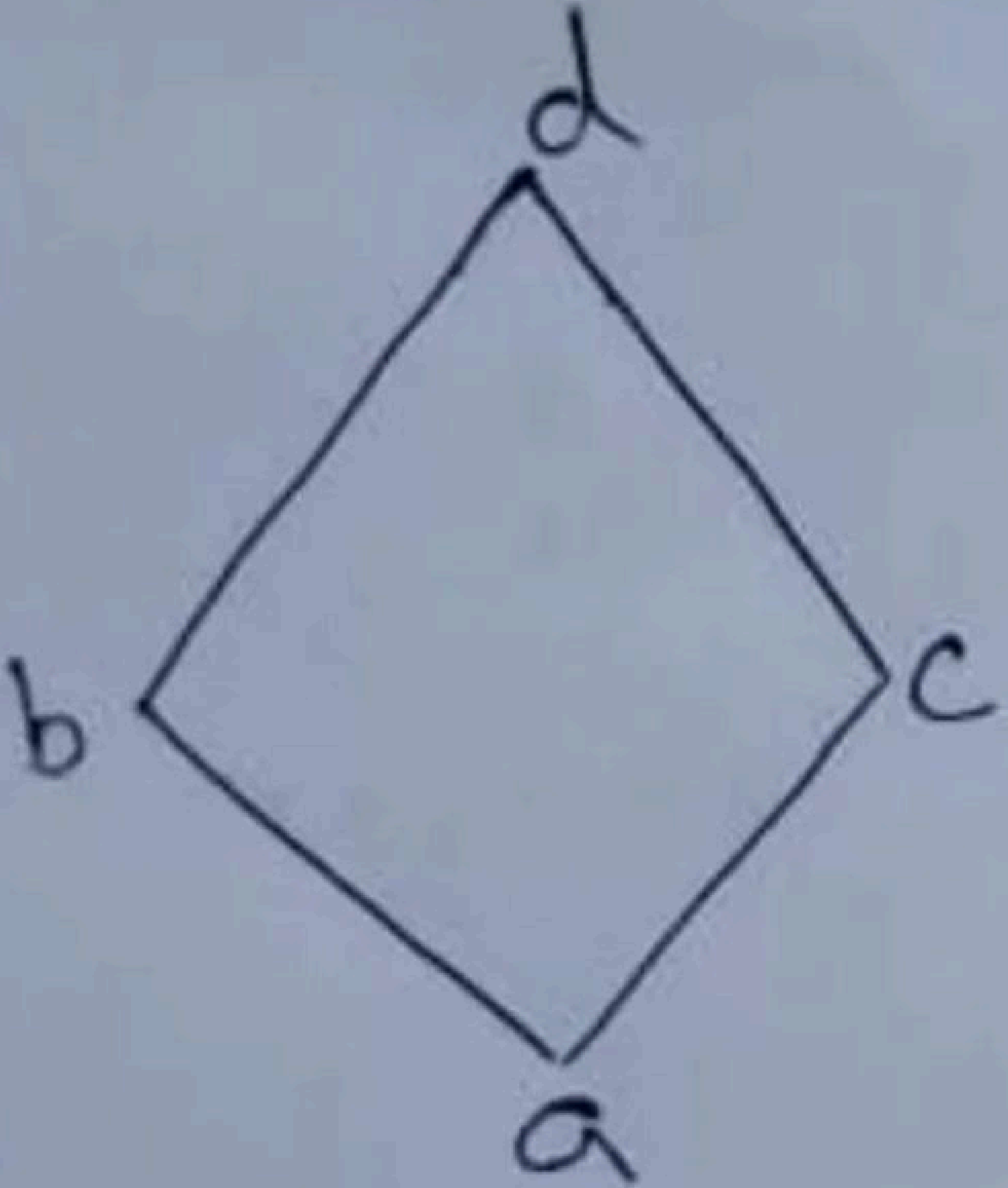
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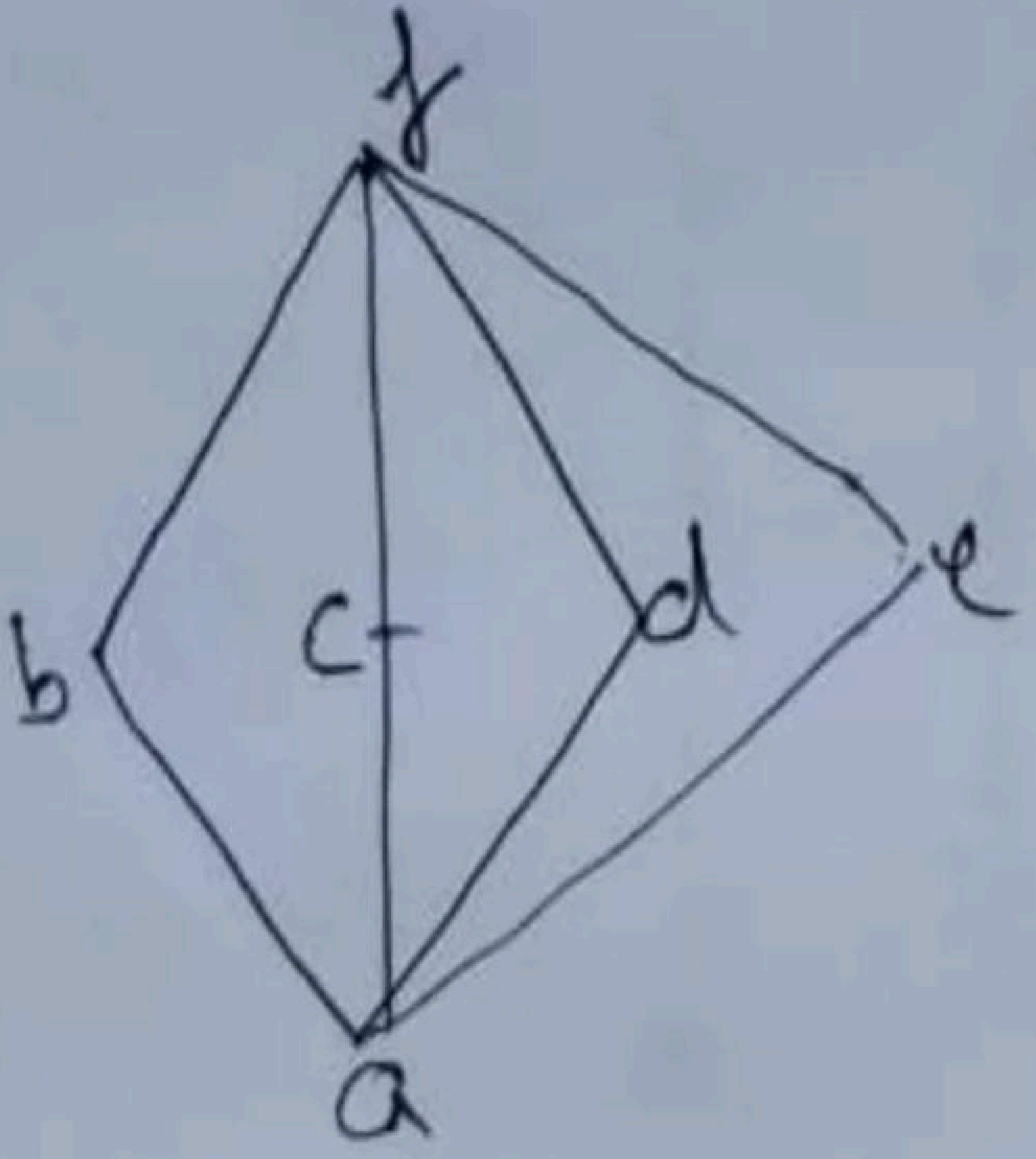
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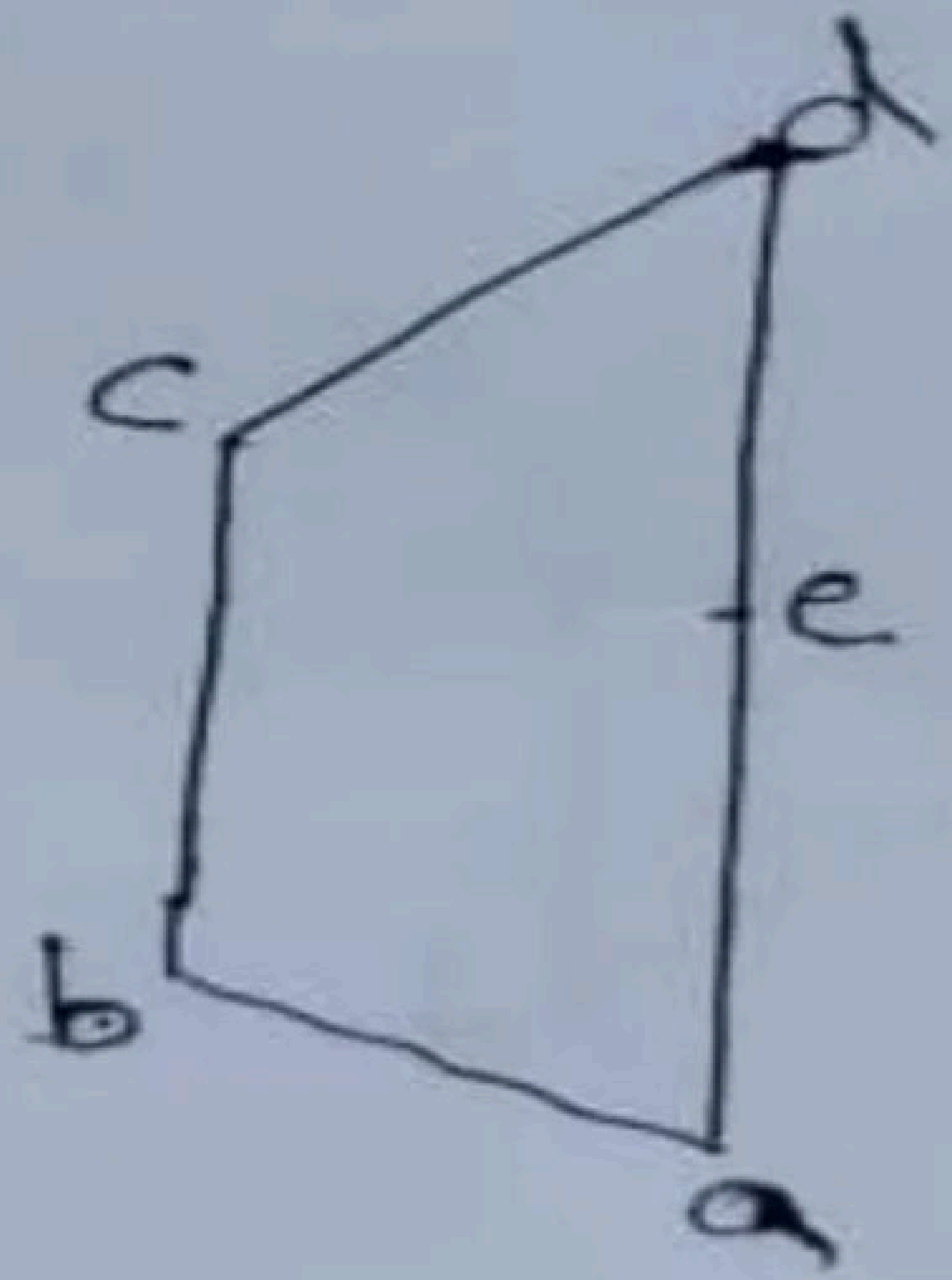
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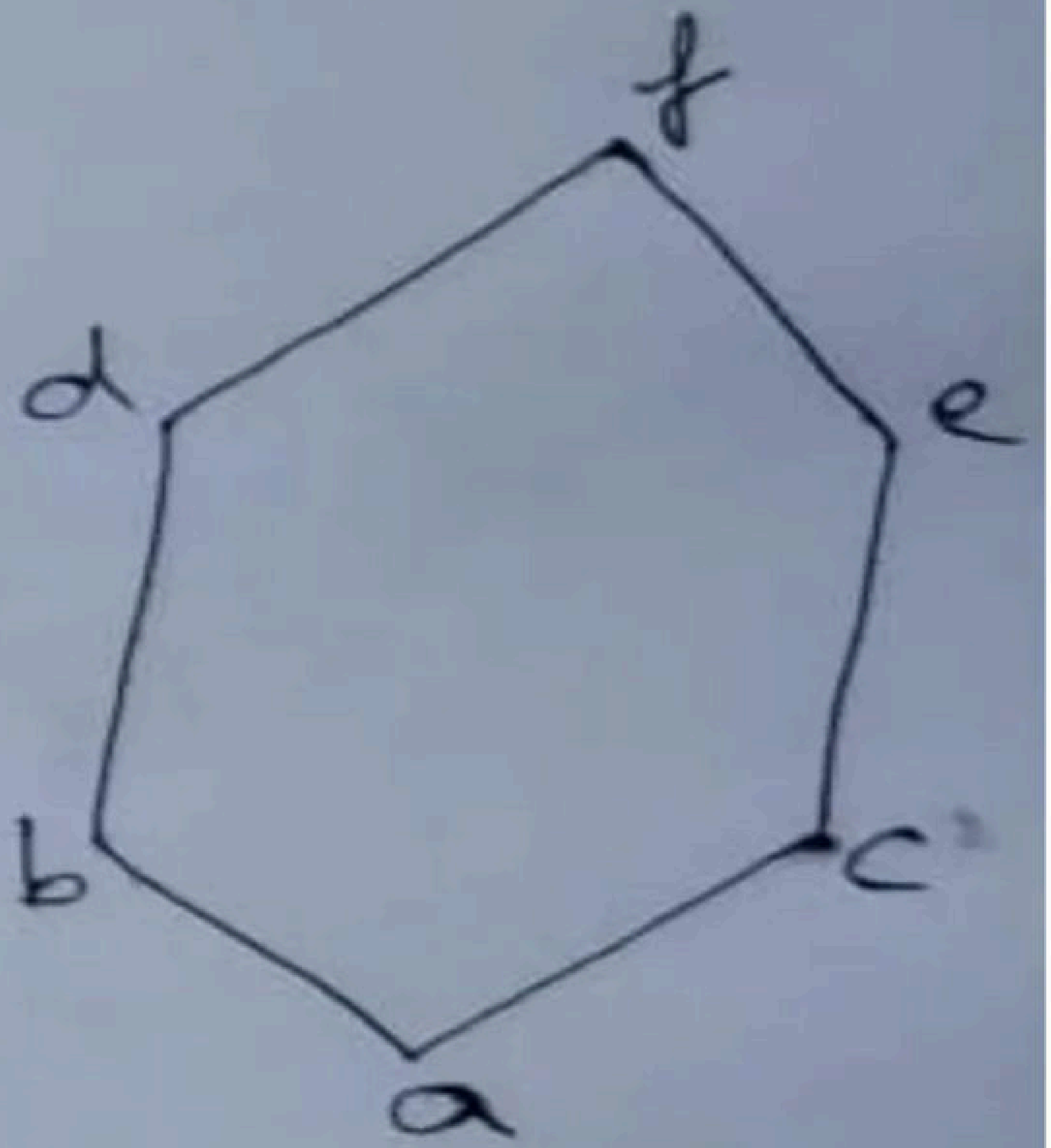
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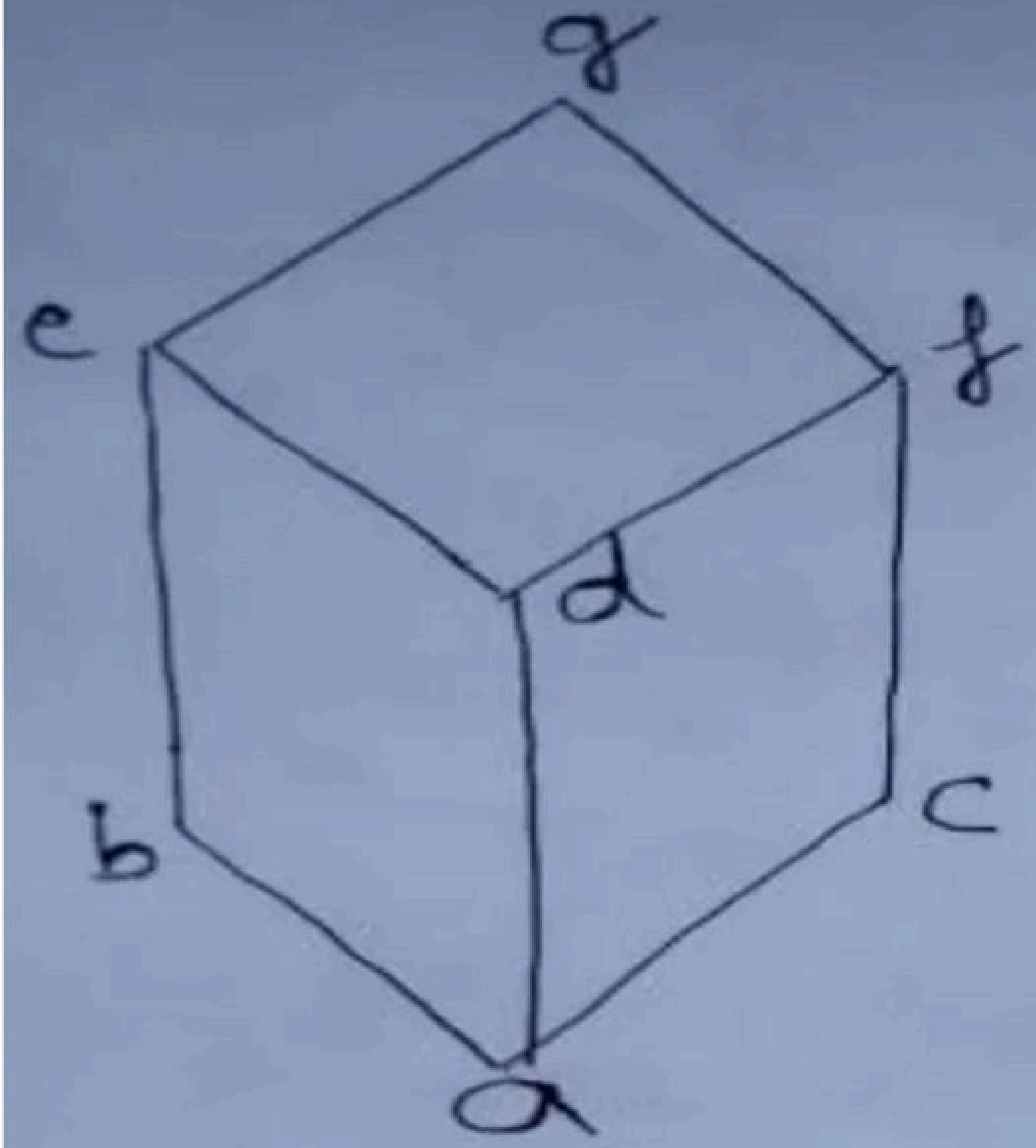
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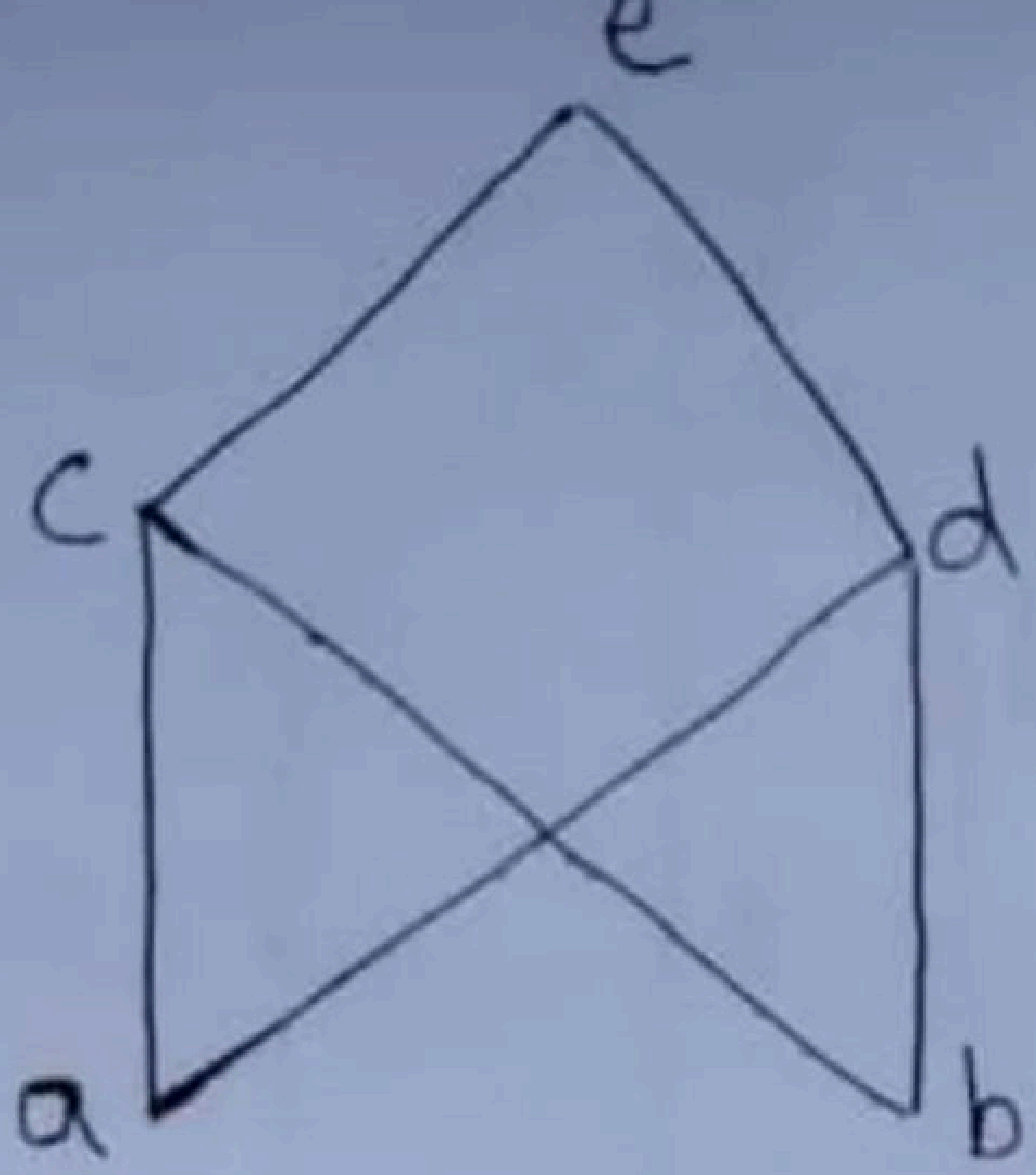
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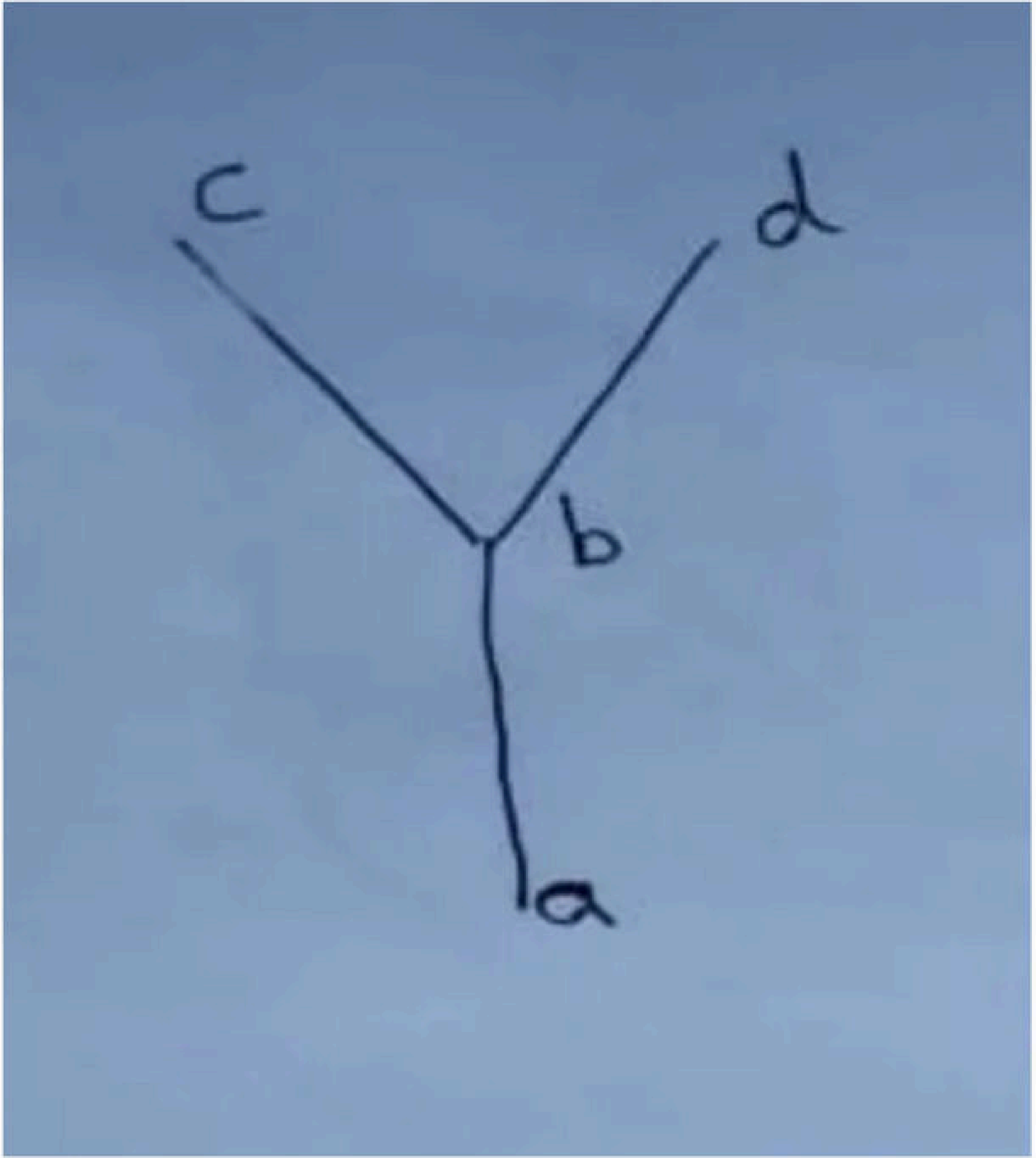
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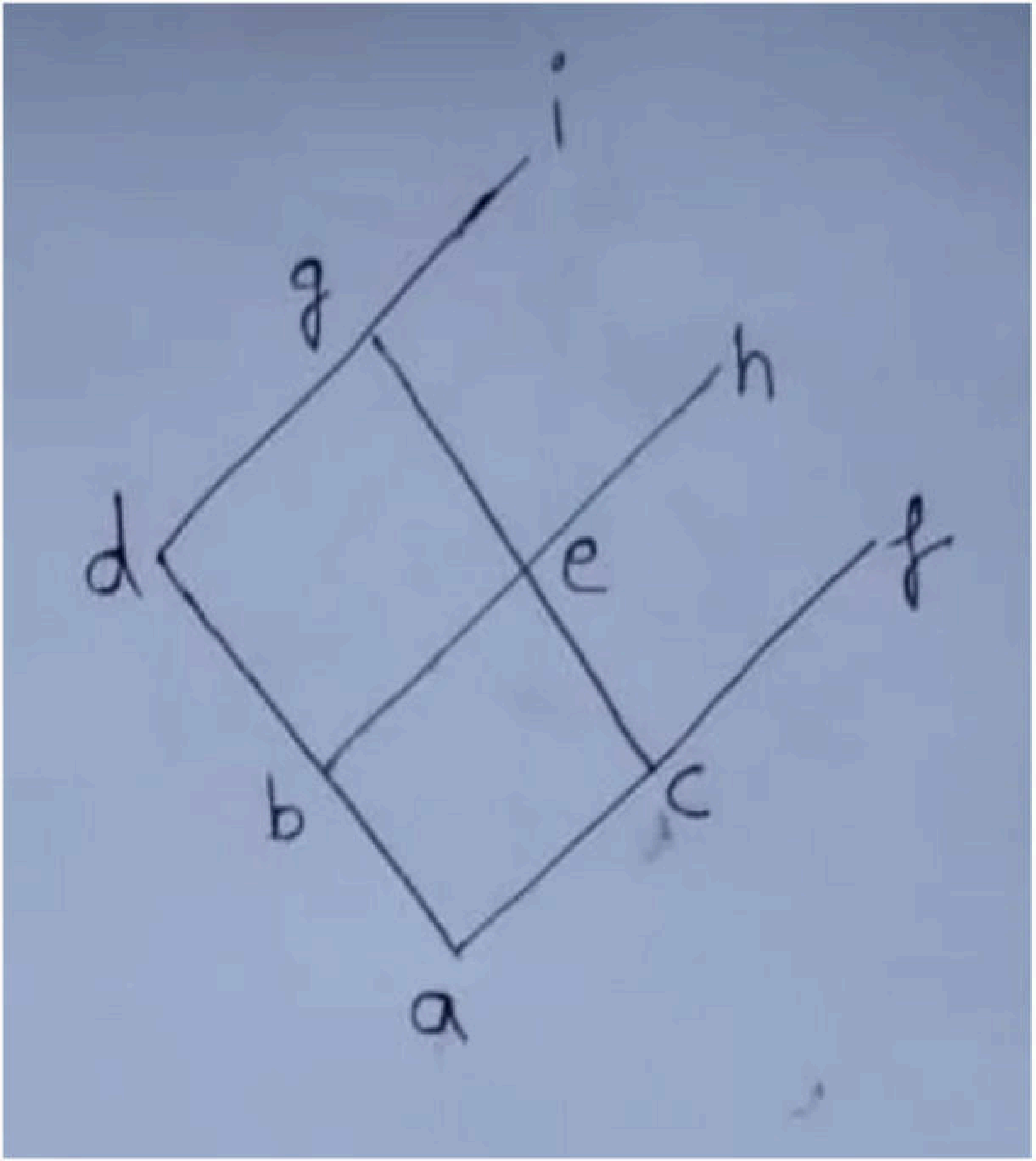
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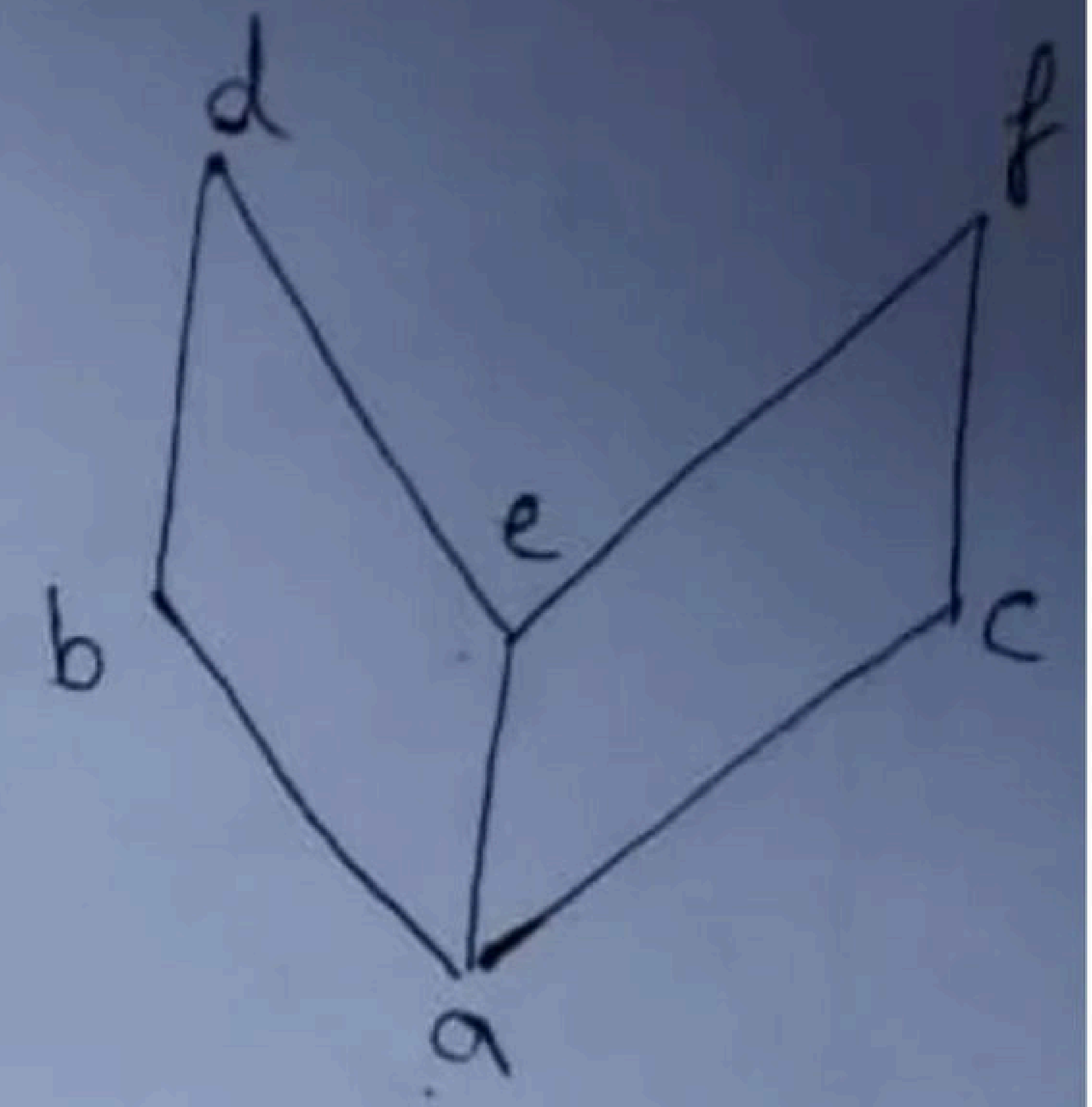
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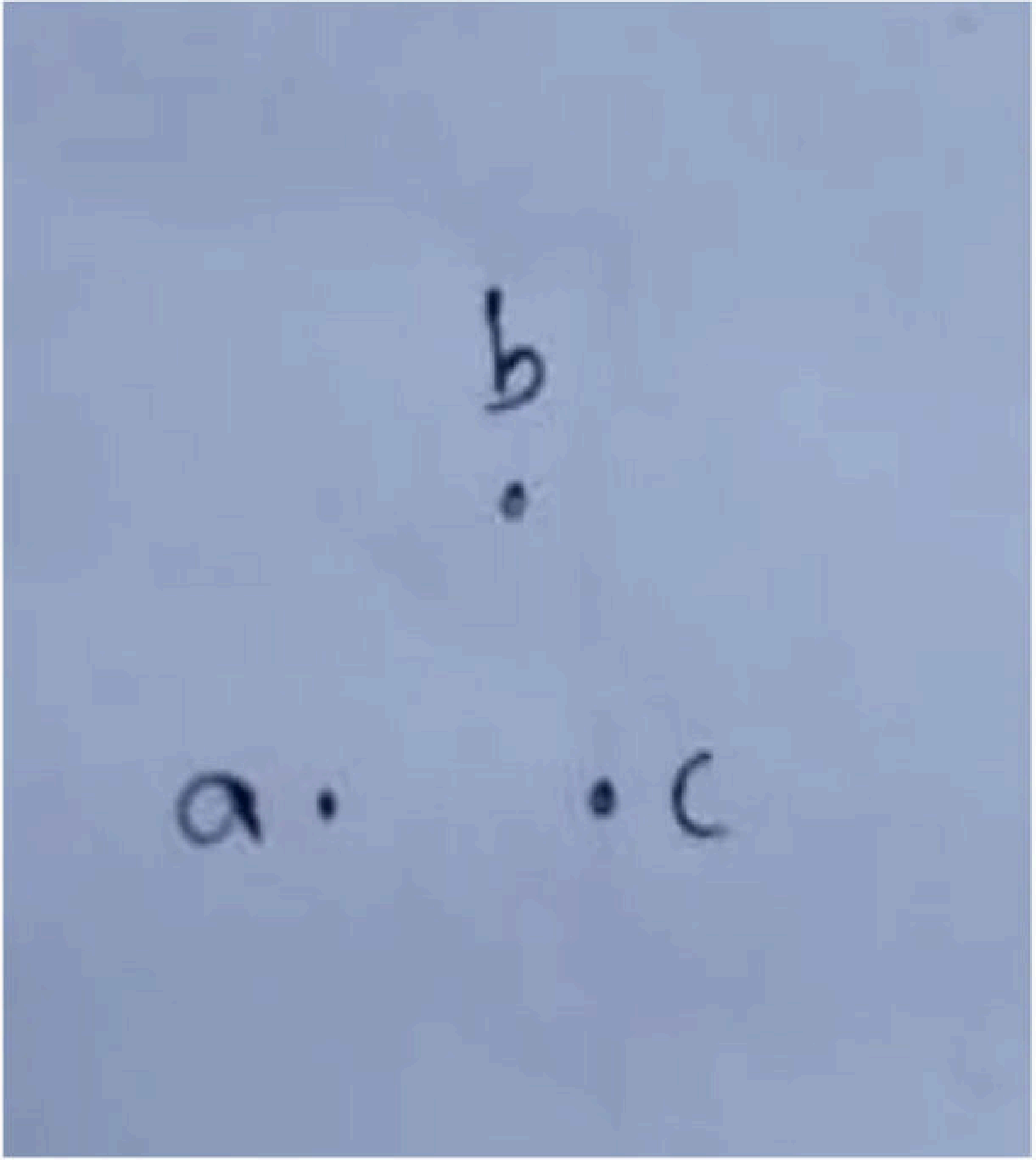
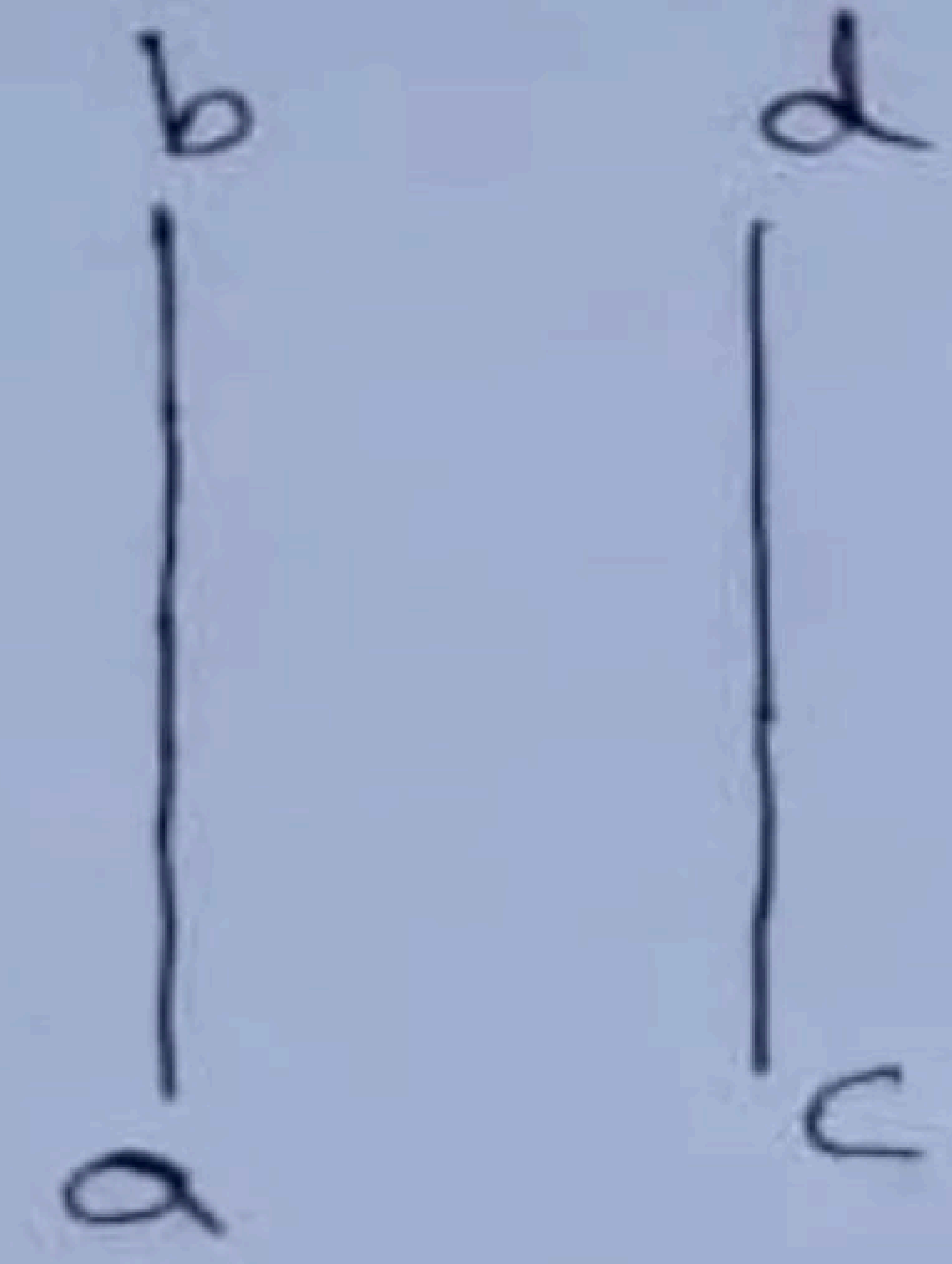

b
a
c

a


b
d
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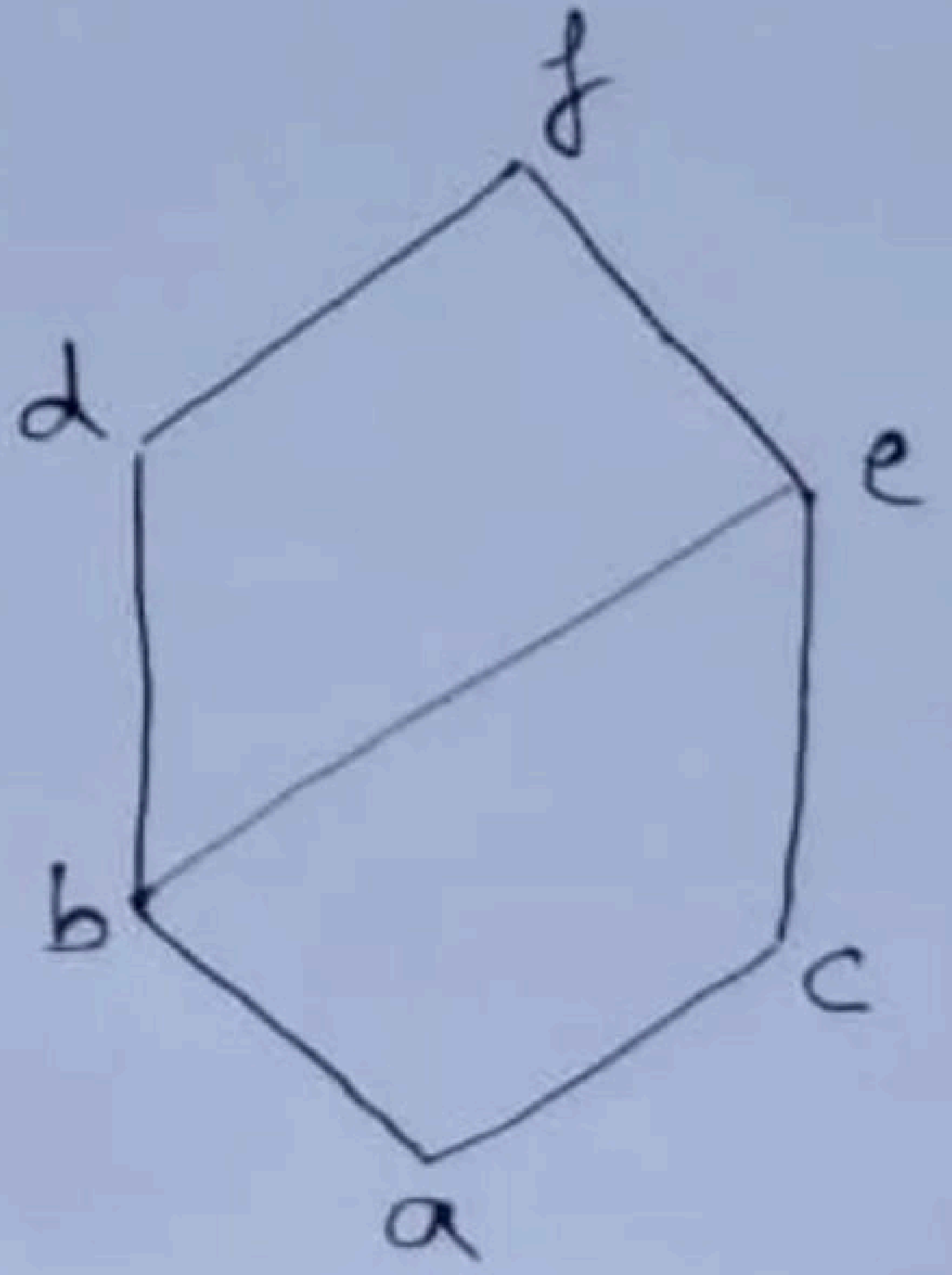
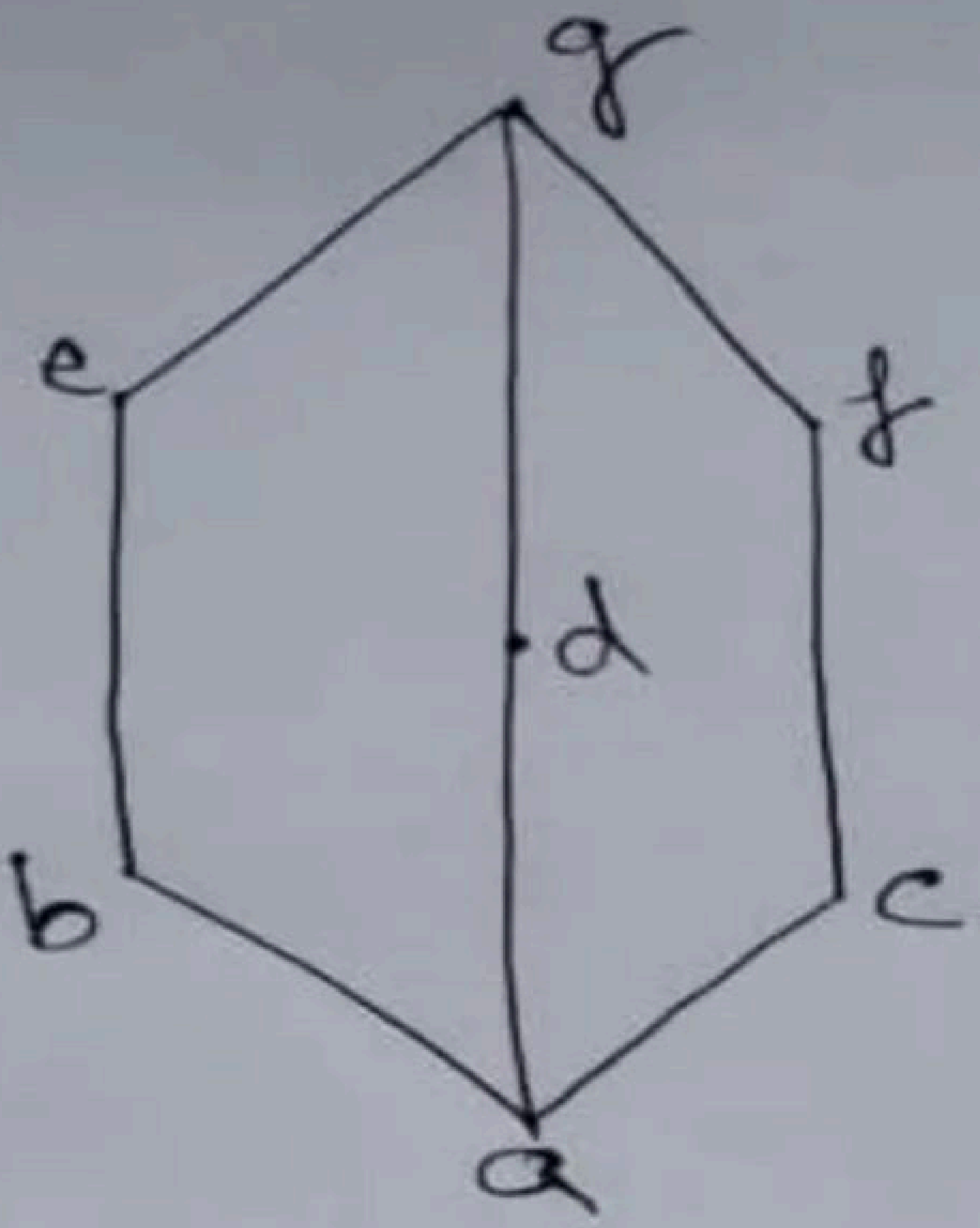


f
d
e
b
c
a


g
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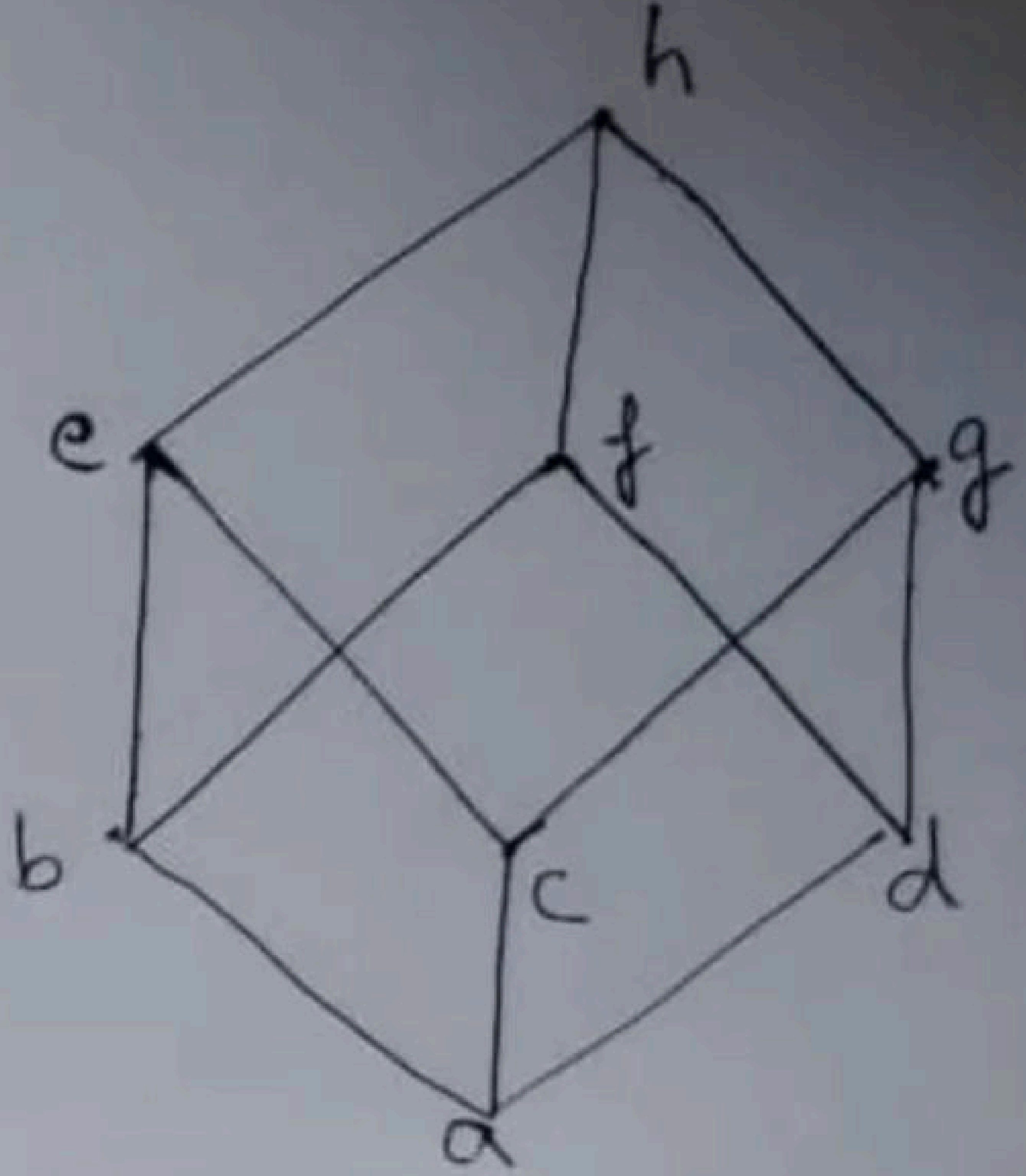
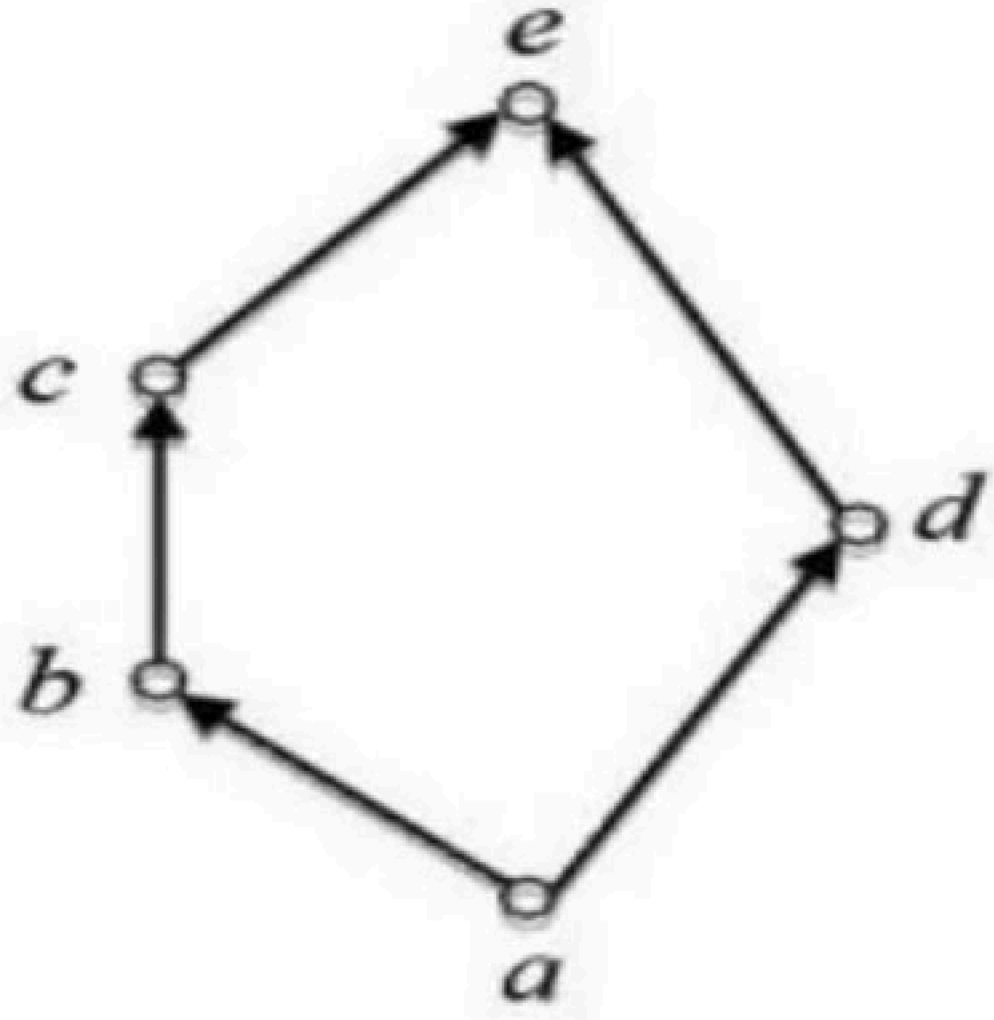


h
e
f
g
b
c
d
a

Break

Q Consider the set **X**={a, b, c, d, e} under partial ordering **R**={(a, a), (a, b), (a, c), (a, d), (a, e), (b, b), (b, c), (b, e), (c, c), (c, e), (d, d), (d, e), (e, e)}. The Hasse diagram of the partial order (X, R) is shown below. The minimum number of ordered pairs that need to be added to R to make (X, R) a lattice is ______ **(GATE-2017) (1 Marks)**



Q A partially ordered set is said to be a lattice if every two elements in the set have **(NET-Dec-2010)**

a) a unique least upper bound

b) a unique greatest lower bound

c) both (A) and (B)

d) none of the above

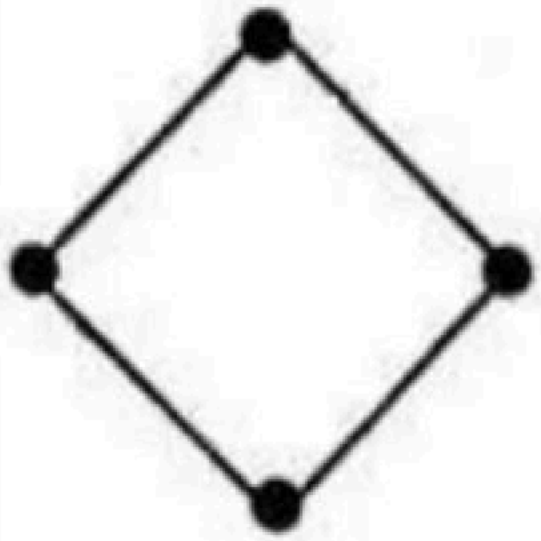
Q Consider the following Hasse diagrams
Which all of the above represent a lattice? **(GATE-2008) (2 Marks)**

(A) (i) and (iv) only

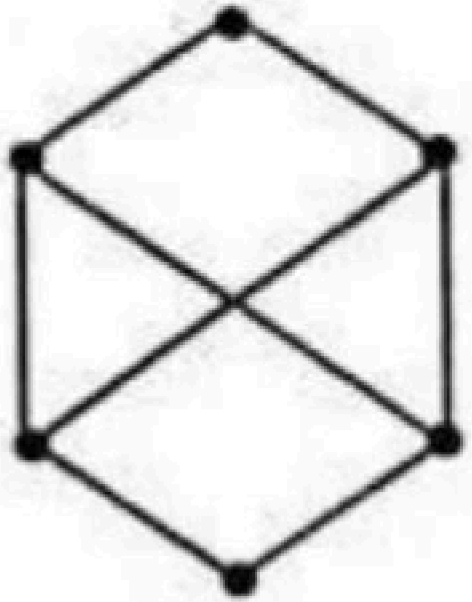
(B) (ii) and (iii) only

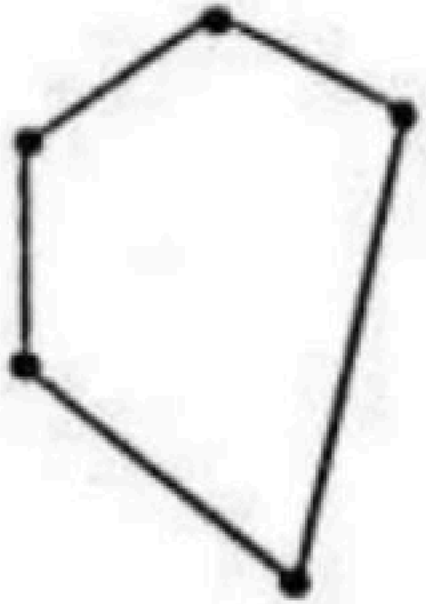
(C) (iii) only

(D) (i), (ii) and (iv) only


(i)


(ii)


(iii)


(iv)

Q the inclusion of which of the following set into S = {{1,2}, {1,2,3}, {1,3,5}, {1,2,4}, {1,2,3,4,5}} is necessary and sufficient to make S a complete lattice under the partial order defined by set containment? **(GATE-2004) (2 Marks)**

a) {1}

b) {1}, {2,3}

c) {1}, {1,3}

d) {1}, {1,3}, {1,2,3,4}, {1,2,3,5}

Break

Boolean algebra

- **Unbounded Lattice :-** If a lattice has infinite of elements then it is called Unbounded Lattice.

- **<u>Bounded Lattice </u>:-** If a lattice has finite number of elements then it is called Bounded lattice, there will be upper and lower bound in lattice.

- **Complement of an element in a Lattice :-** If two elements a and a^c, are complement of each other, then the following equations must always holds good.

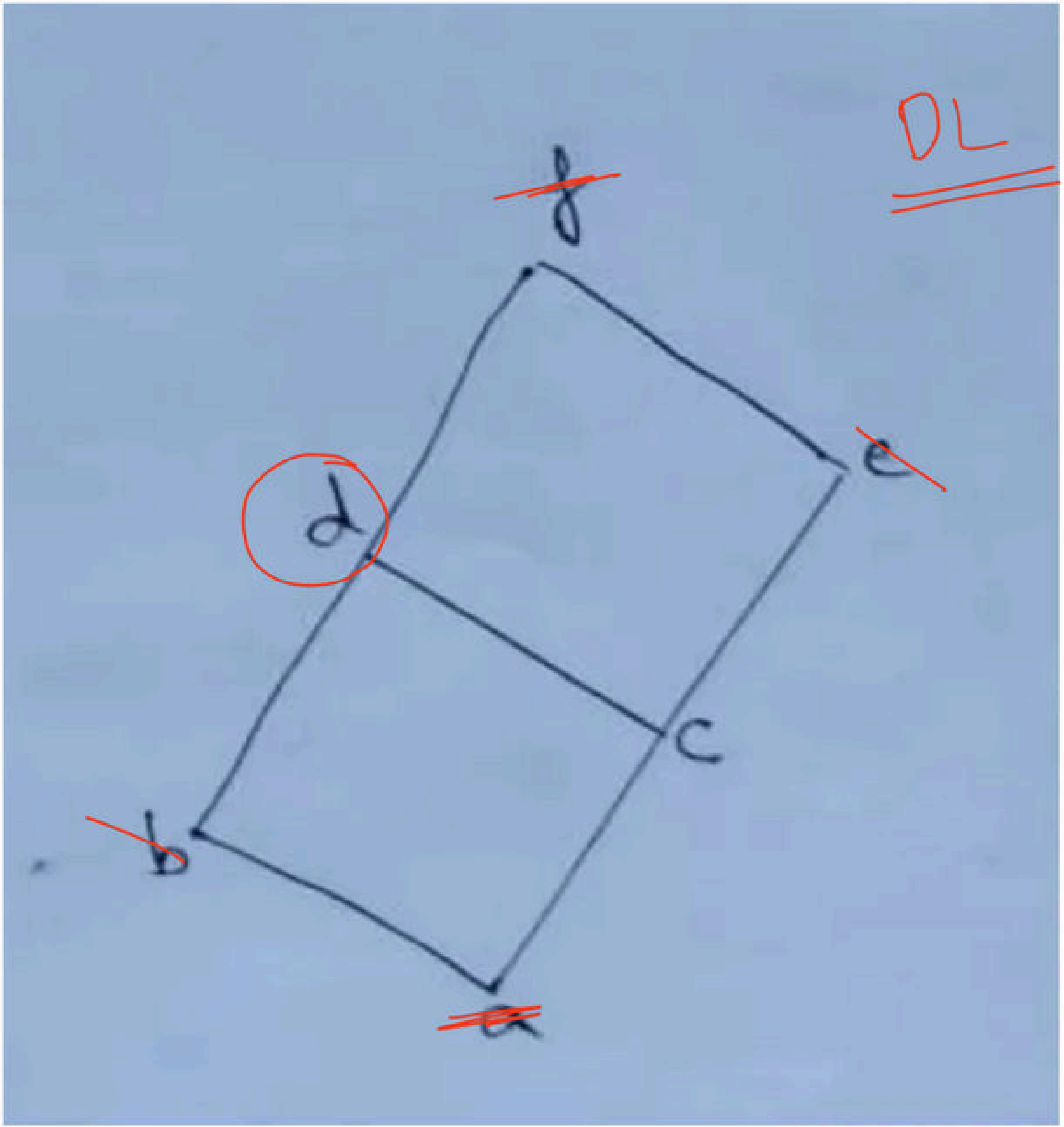
$a \vee a^c$ = Upper bound of lattice

$a \wedge a^c$ = Lower bound of lattice

- **<u>Distributive Lattice</u> :-** A lattice is said to be distributed lattice. if for every element their exist at most one completemt(zero or one).

- **<u>Complement Lattice</u> :-** A Lattice is said to be Complement lattice. if for every element their exist at least one complement(one or more).

- **Boolean Algebra :-** A Lattice is said to be Boolean Algebra, if for every element their exist exactly one complement. Or if a lattice is both complemented and distributed then it is called Boolean Algebra.



$A \cup A^c = U$ [U.B]

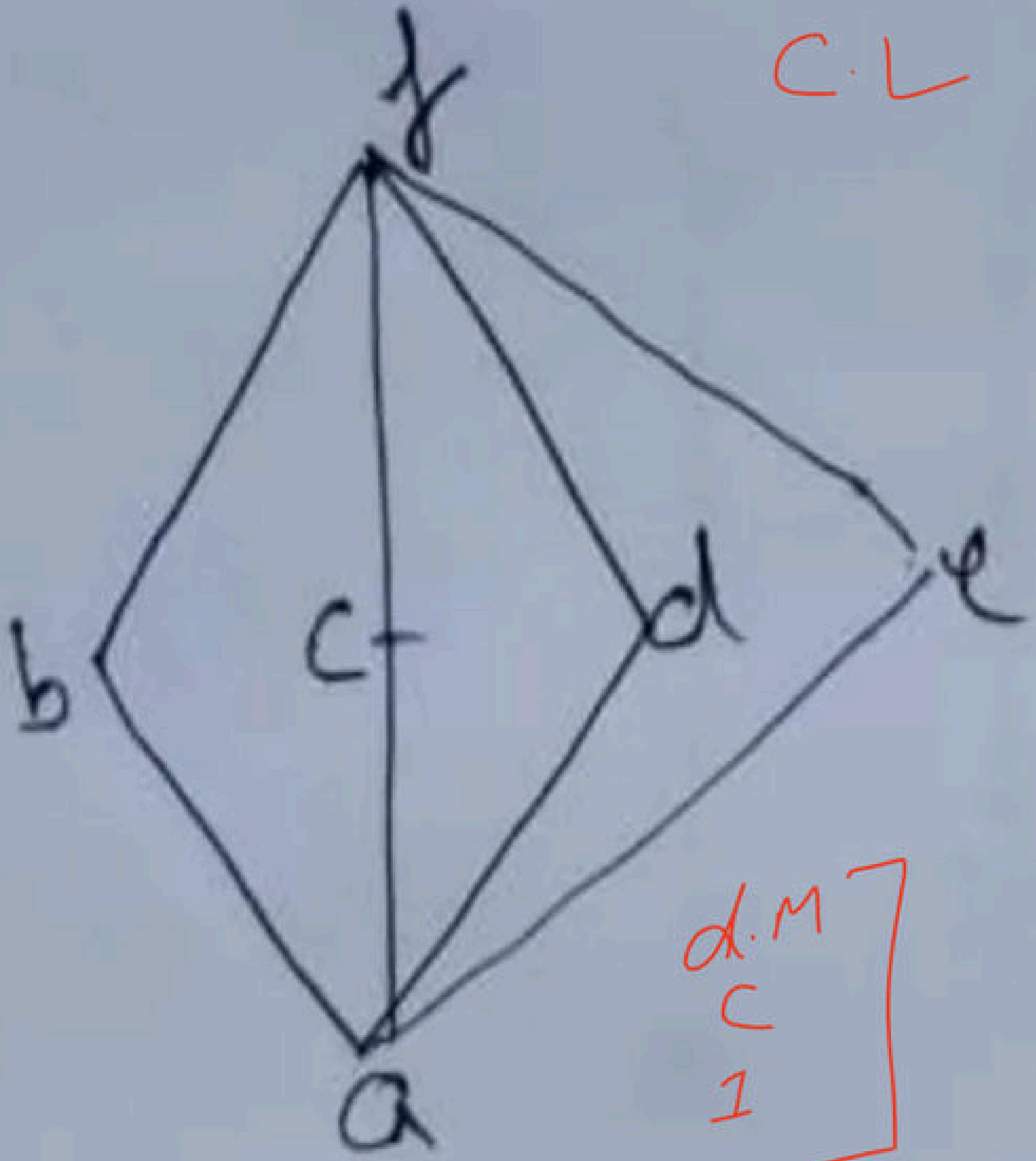
$A \cap A^c = \phi$ [L.B]

$b \vee e = UB \quad f$

$b \wedge e = L.B = a$

$d \vee c = d$

$d \wedge c = c$



C.L

d.M

C

1

$$c \wedge c \wedge c \wedge c \wedge c \,--\, c = c$$

$$c \vee c \vee c \vee c \,---\, \vee c = c$$

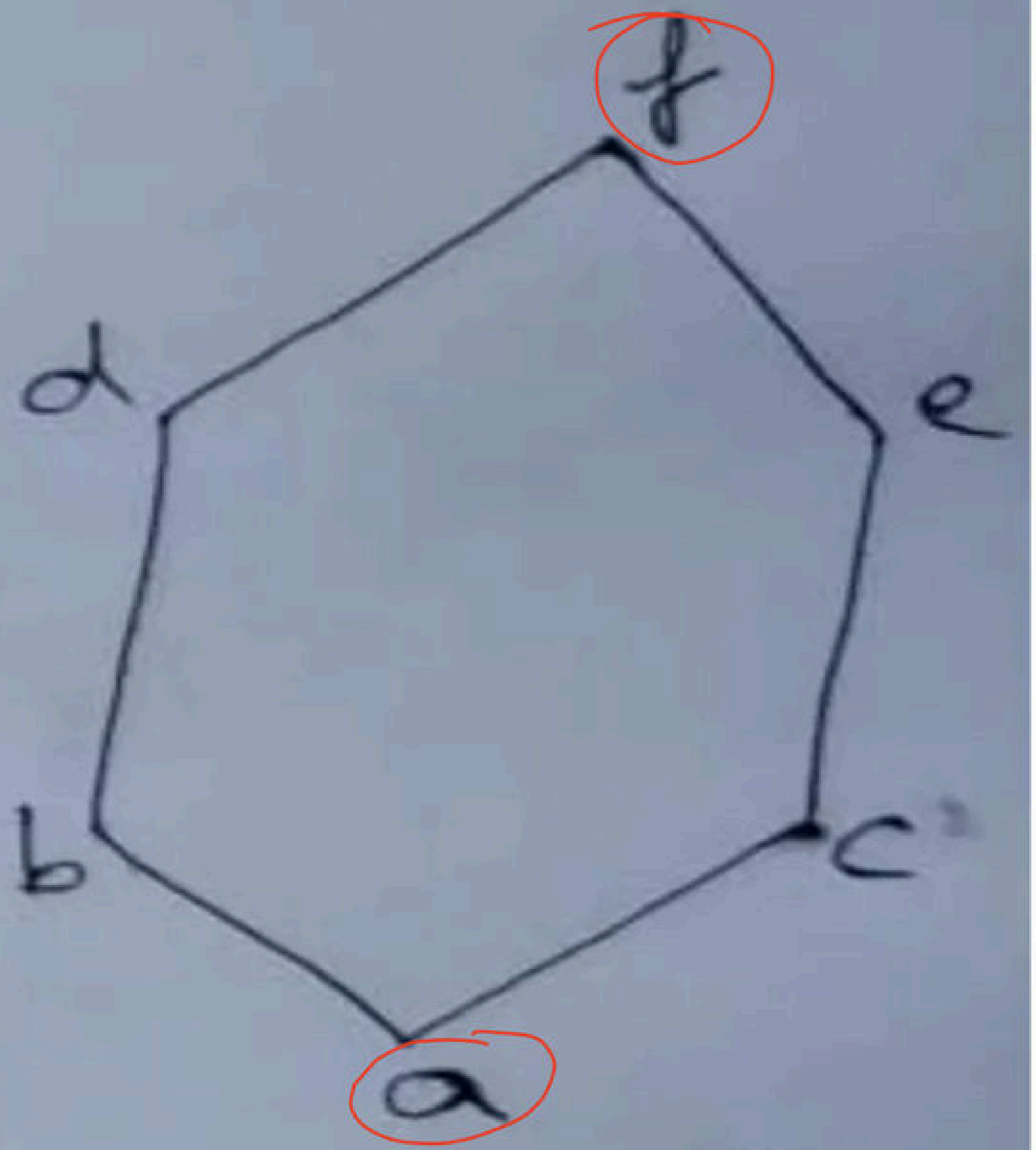
$$b \vee c = c \vee b$$

$$(\underbrace{\cancel{b \vee c}}_{\cancel{f}}) \vee d = \cancel{b \vee (c \vee d)}$$

$$f \qquad\qquad f \; \cancel{f}$$

$$\cancel{b} \vee \underbrace{\cancel{(c \wedge d)}}_{\cancel{a}} = \underbrace{\cancel{(b \vee c)}}_{\cancel{f}} \wedge \underbrace{\cancel{(b \vee d)}}_{\cancel{f}}$$

$$b \neq f$$



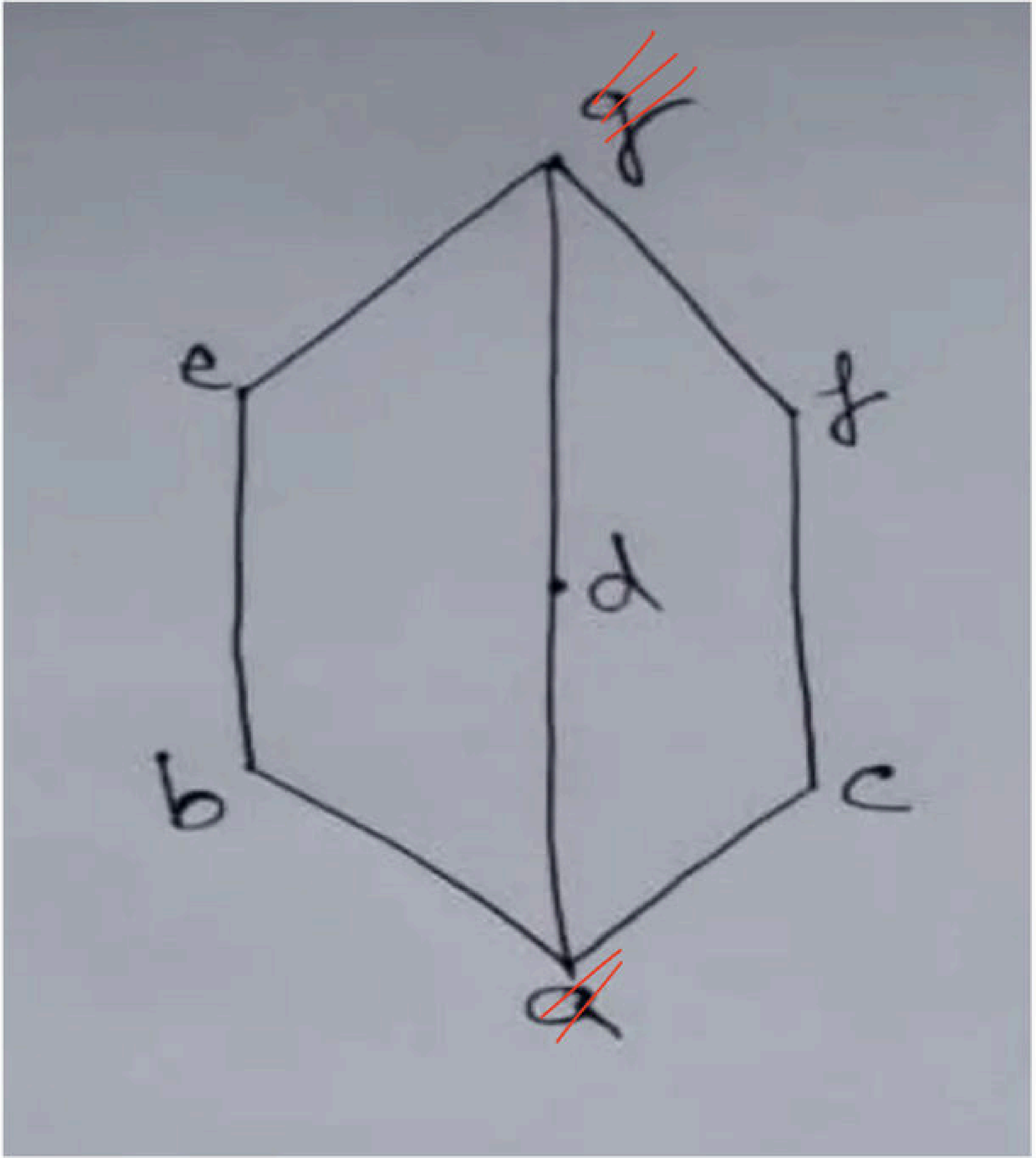
$b' = c, e$

$d' = c, e$

$e' = b, d$

$c' = b, d$

C.L



$e' = d\,f\,c$

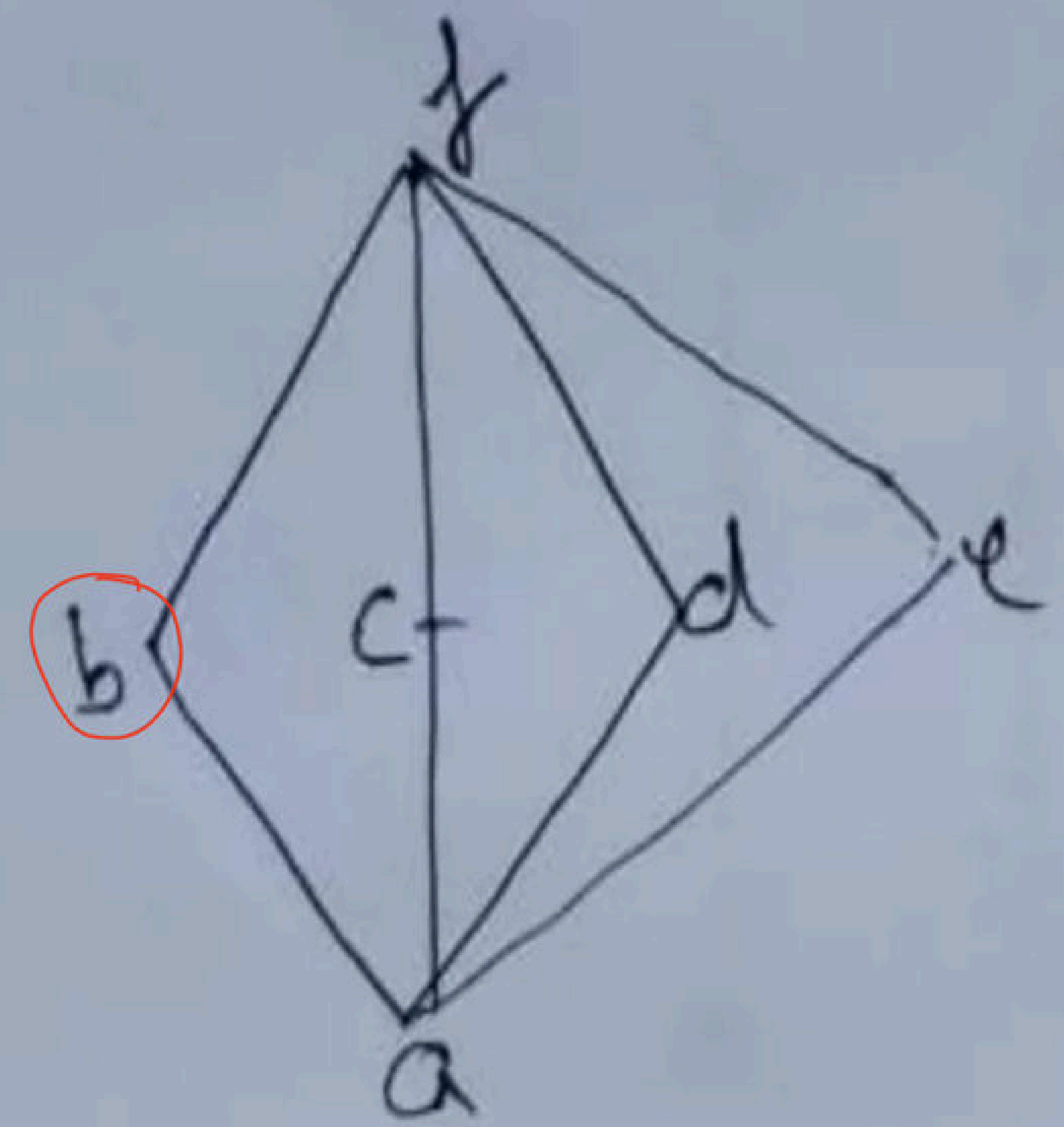
$b' = d\,f\,c$

$f' = e\,b\,d$

$c' = e\,b\,d$

$d' = b\,e\,f\,c$

C.L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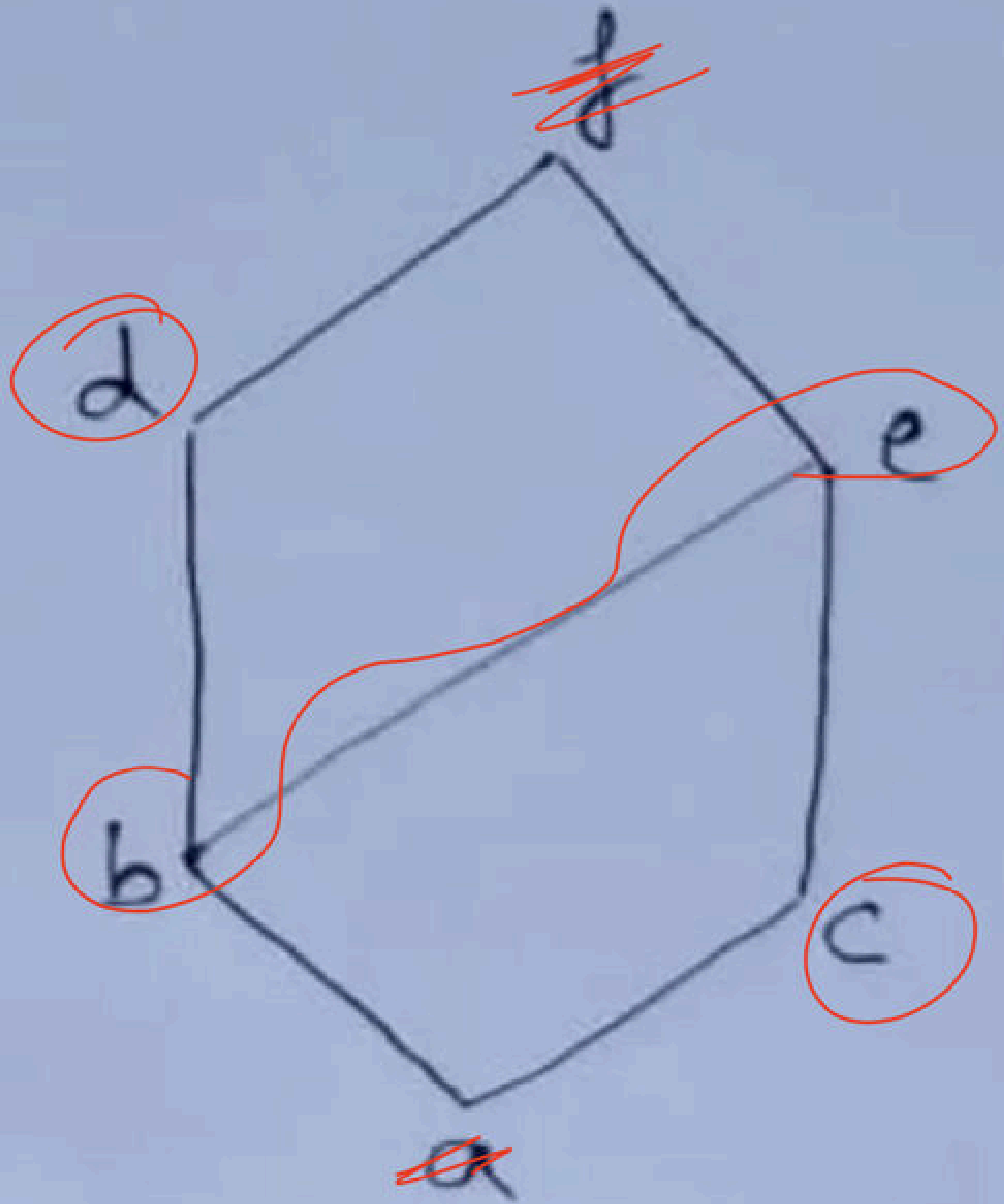
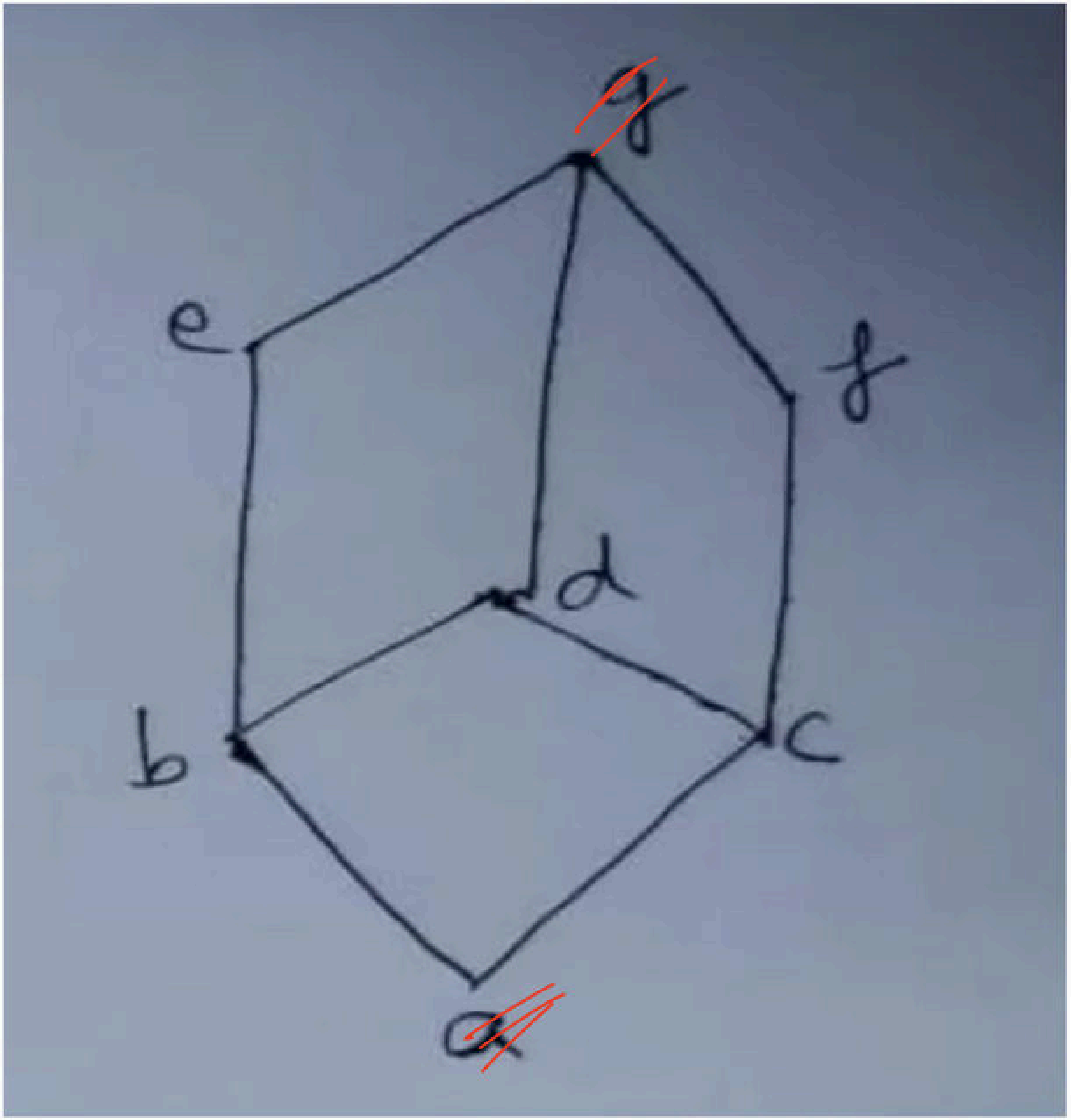


d. L



$\boxed{a \leftrightarrow g}$

$e' = f, c$ — $\cancel{D.L}$

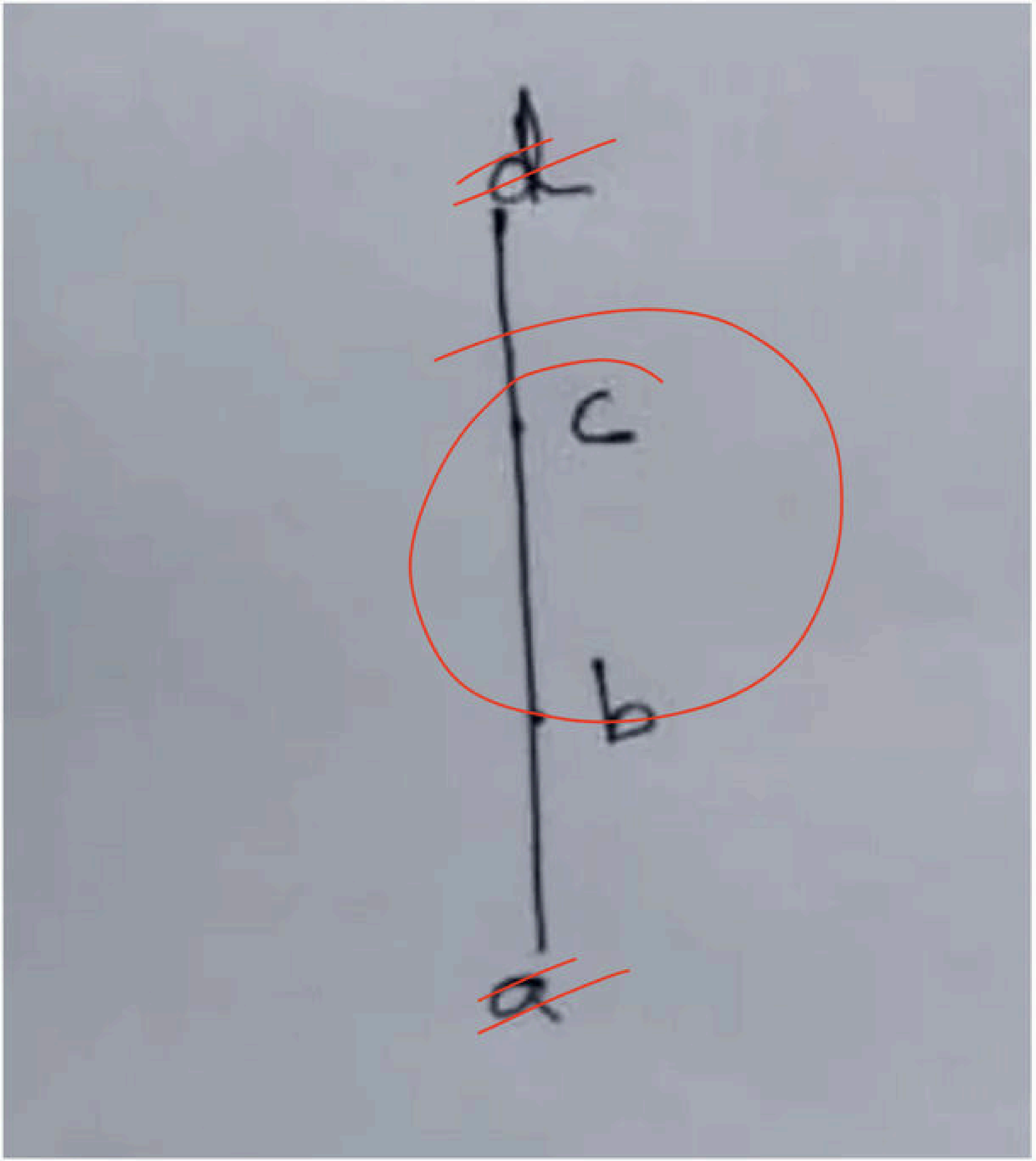
$f' = e, b$

$b' = f$

$c' = e$

$d' = ?$ — $\cancel{C.L}$

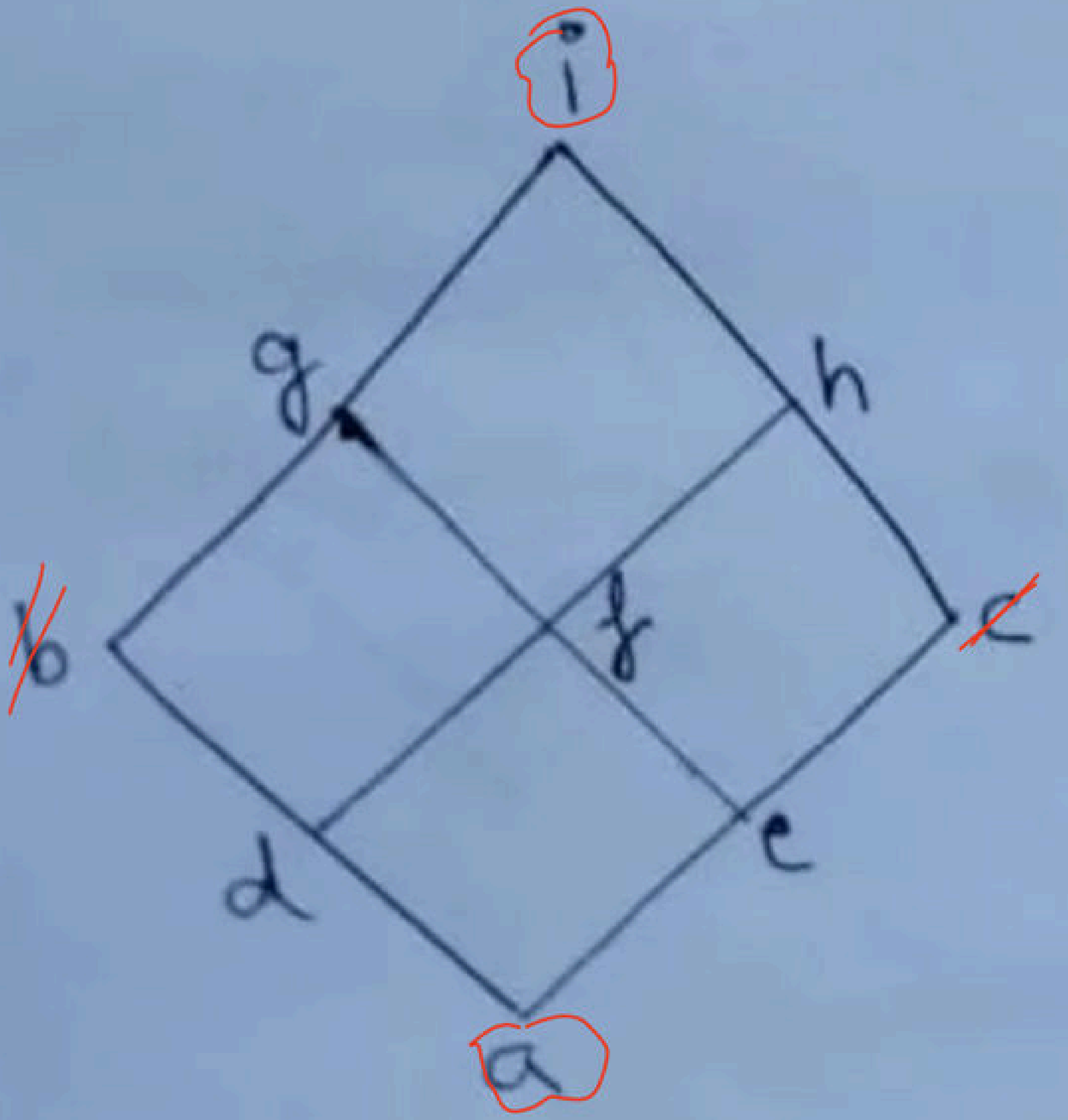
Break



$c' = ?$

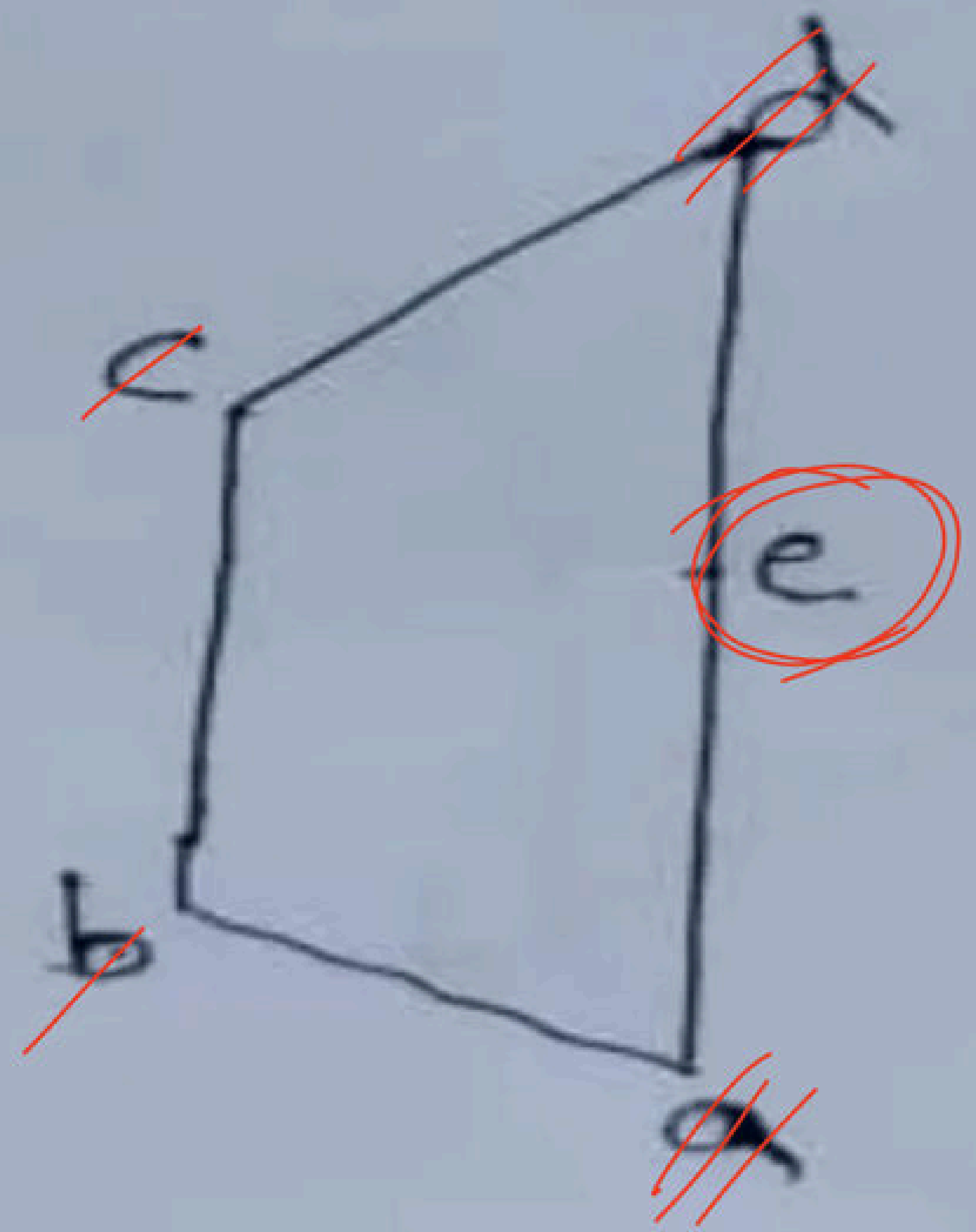
$b' = ?$

D.L

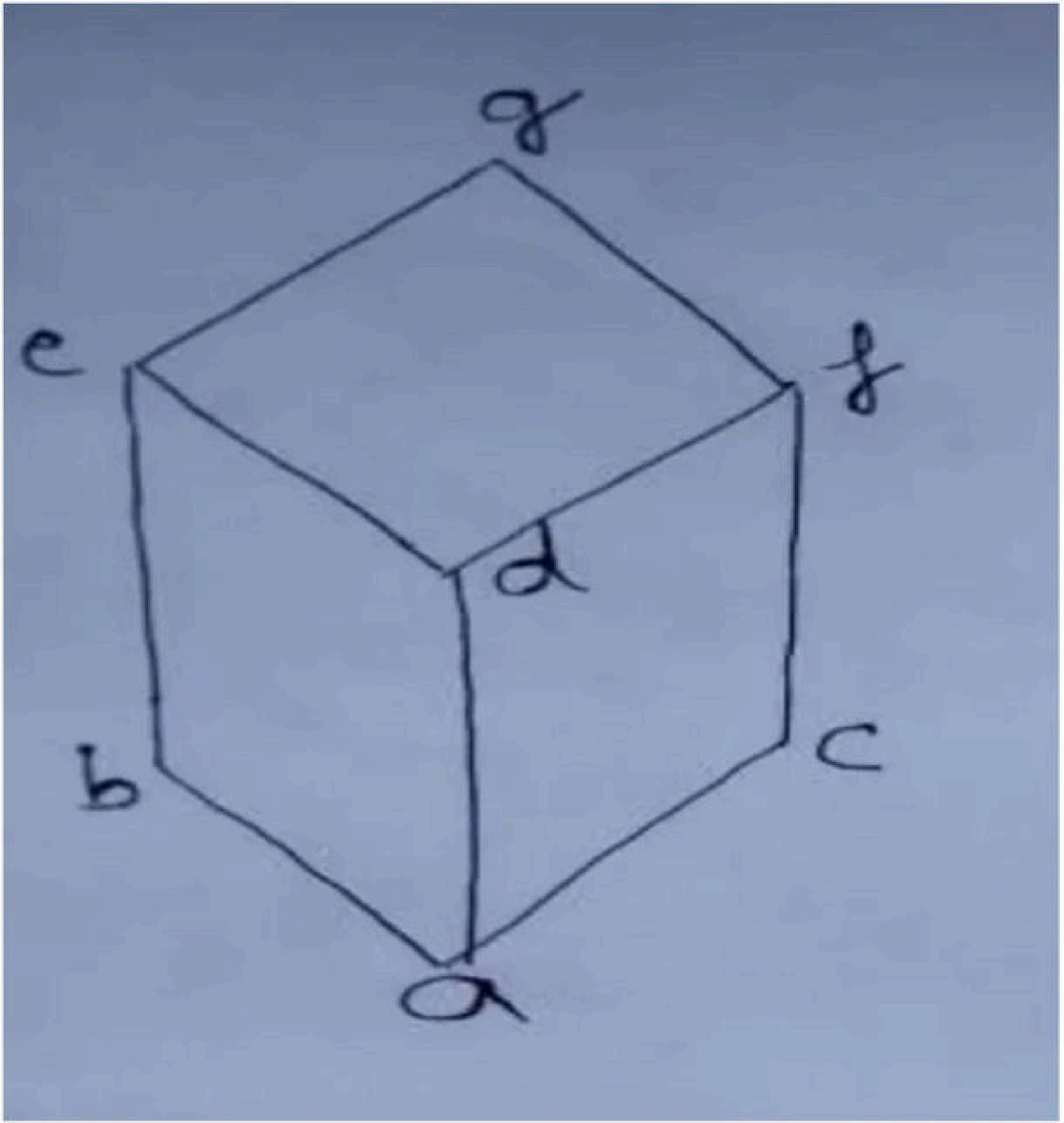


D.L

$g' h' f' d' e' = ?$



e: c, b


g
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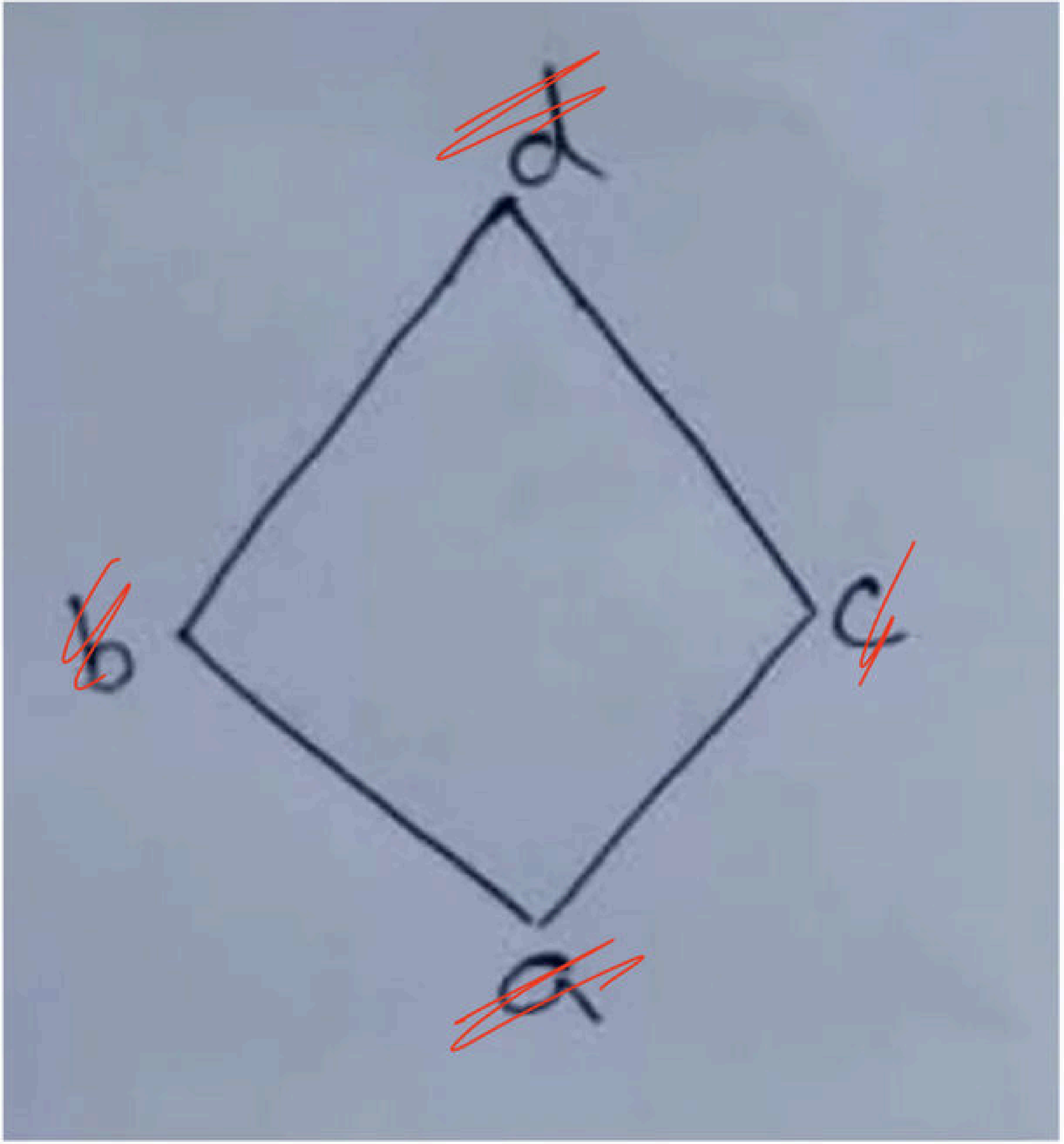
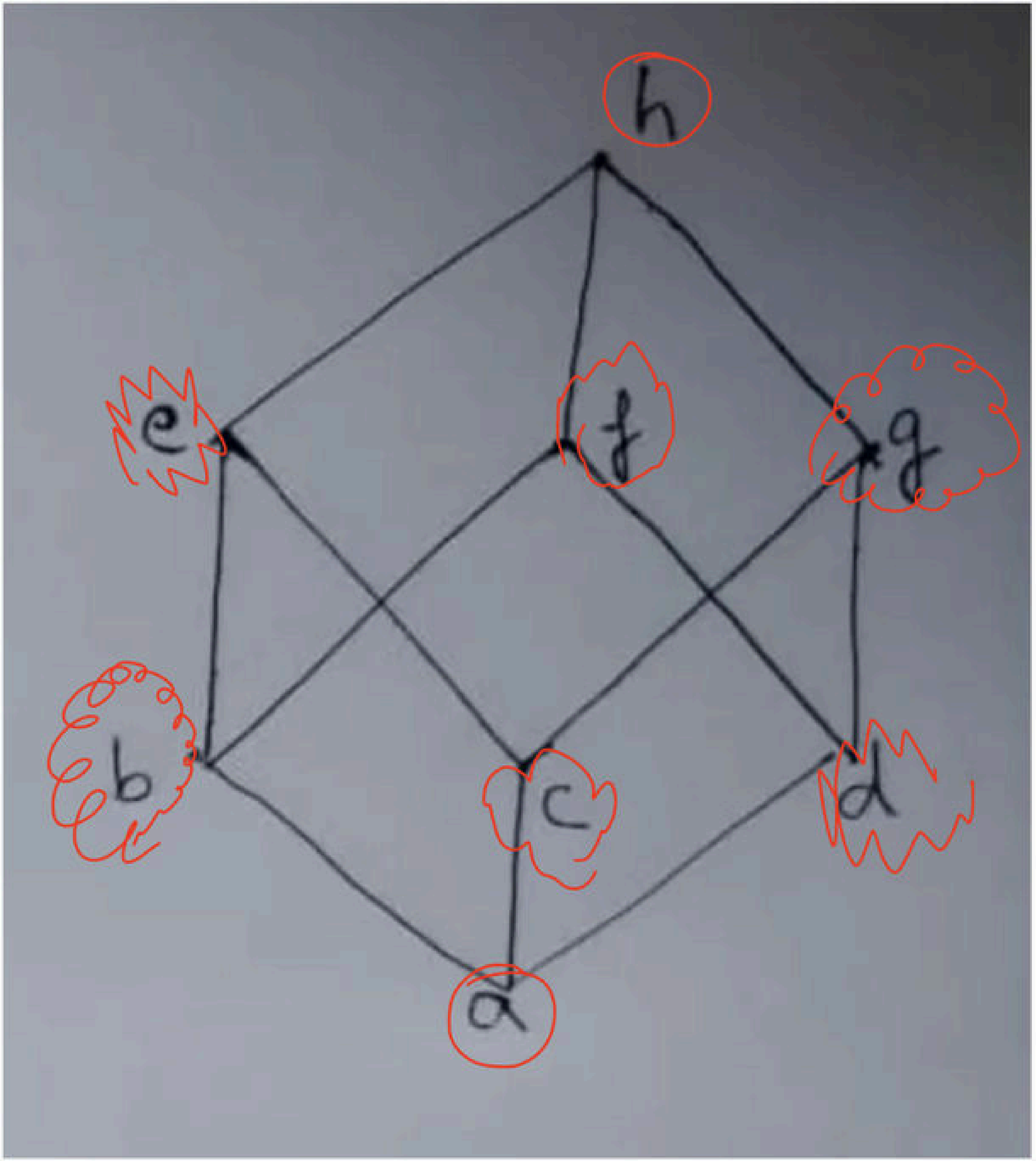


BA



h h ehg

f c

Break

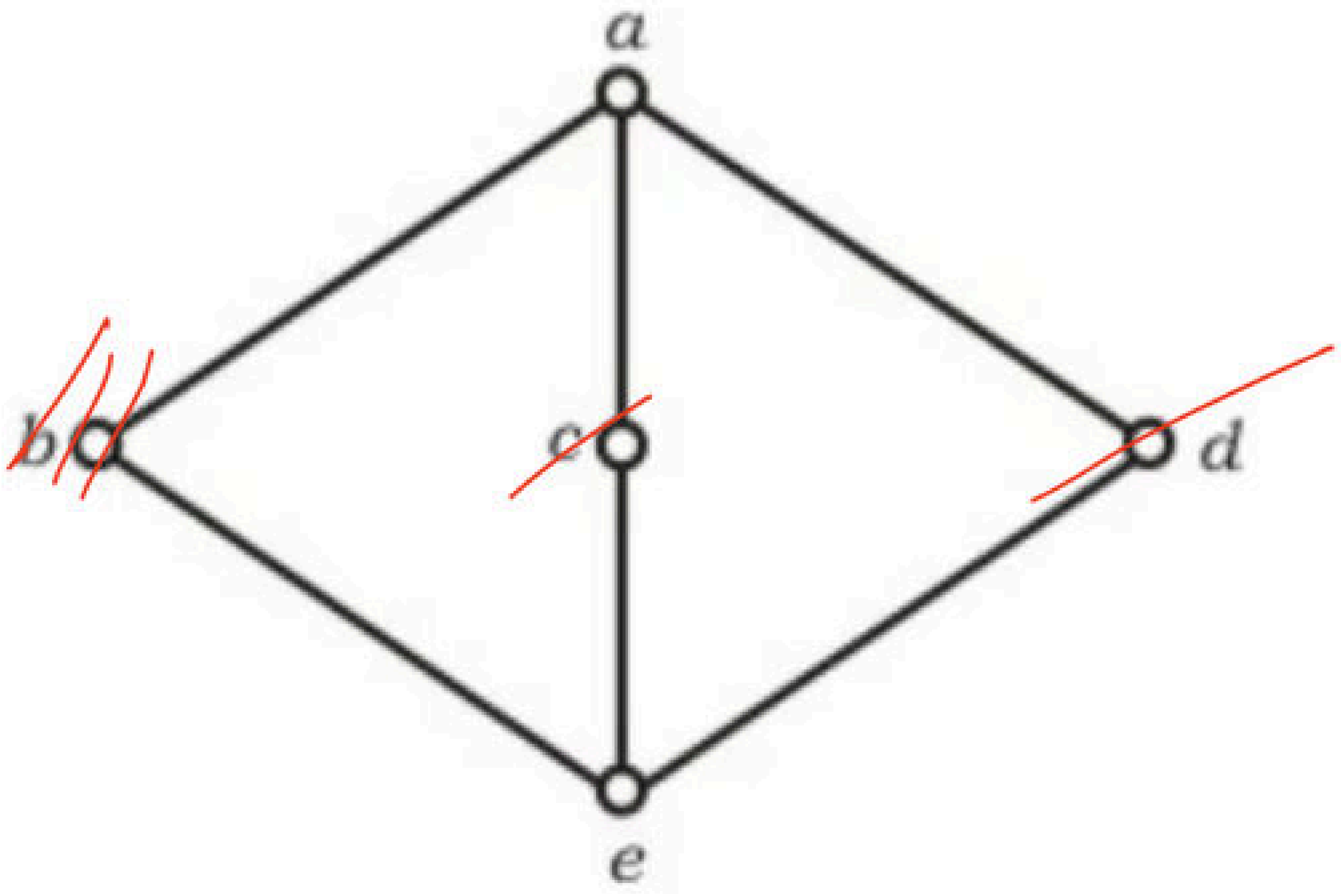
Q The following is the Hasse diagram of the Poset [{a, b, c, d, e}, ≤]
The Poset is **(GATE-2005) (1 Marks)**

(A) ~~not a lattice~~ – 1

(B) a lattice but not a distributive lattice – 66

(C) a distributive lattice but not a Boolean algebra – 23

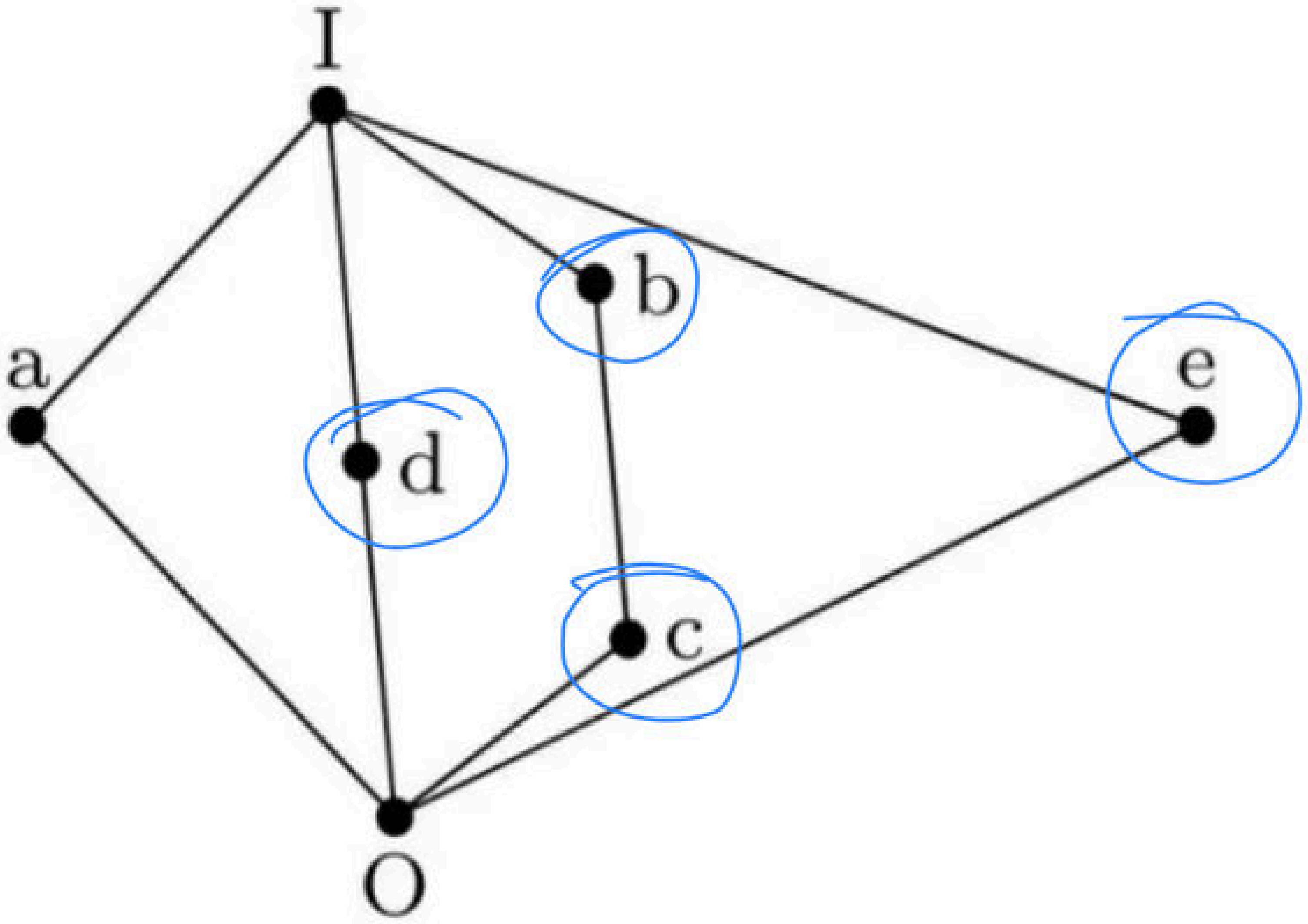
(D) a ~~Boolean algebra~~ – 10



C.L

Q The complement(s) of the element 'a' in the lattice shown in below figure is (are) ____4

(GATE-1988) (2 Marks)



Q Find which of the following is a lattice and Boolean Algebra?

(1) $[D_{10}, /]$

(2) $[D_{12}, /]$

(3) $[D_{30}, /]$

(4) $[D_{45}, /]$

(5) $[D_{64}, /]$

(6) $[D_{81}, /]$

(7) $[D_{91}, /]$

(8) $[D_{110}, /]$

Q Find which of the following is a lattice and Boolean Algebra?

(1) [{1,2,3,4,6,9}, /]

(2) [{2,3,4,6,12}, /]

(3) [{1,2,3,5,30}, /]

(4) [{1,2,3,6,9,18}, /]

(5) [{2,3,4,9,12,18}, /]

(6) [R, <=]

(7) [P(A), $\sqsubseteq$], A = {1,2,3}

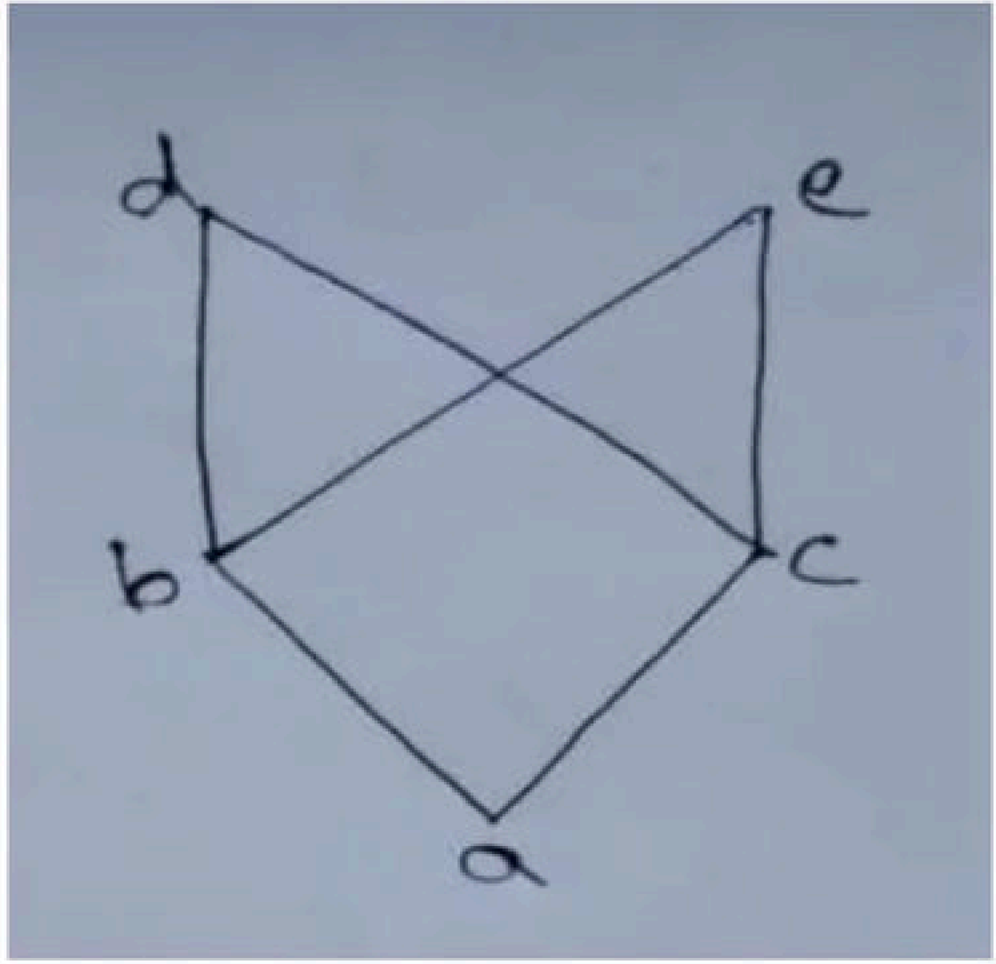
Q Consider the following hasse diagram, find which of the following is true?

a) it is a lattice

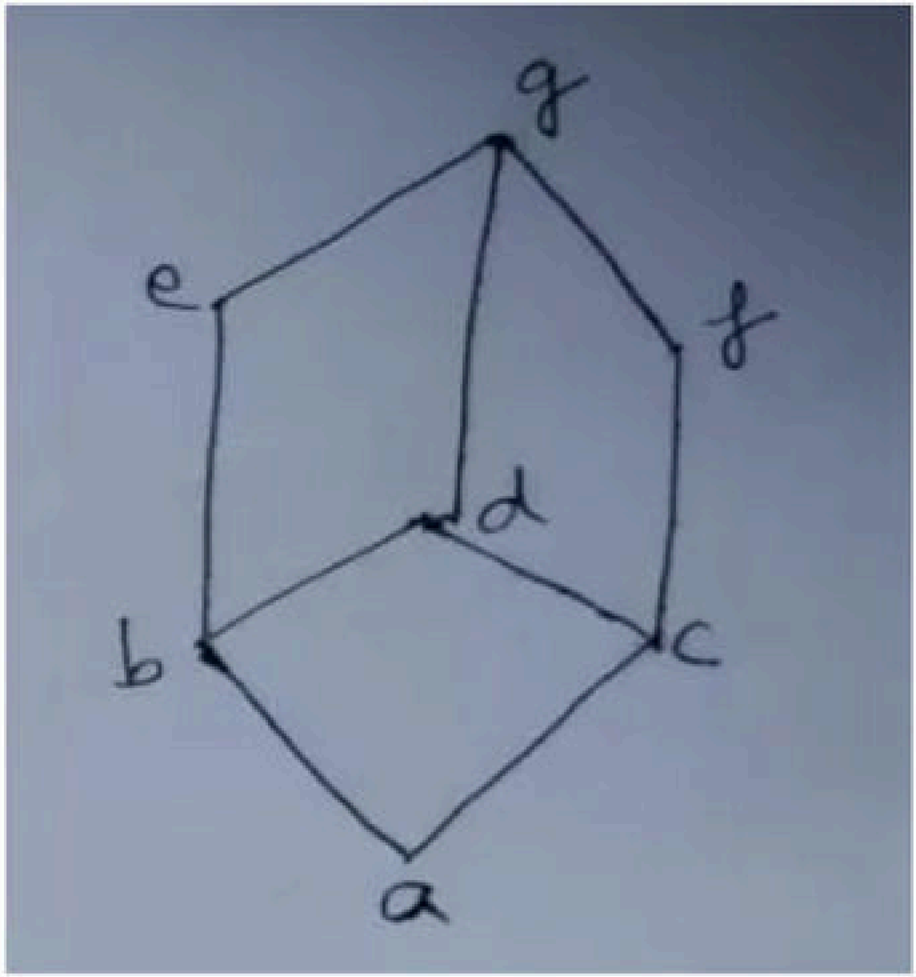
b) subset {a, b, c , d} is a lattice

c) subset {b, c, d, e} is a lattice

d) subset {a, b, c, e} is a lattice



Q Consider the following hasse diagram, find which of the following is true?



a) subset {a, b, c, g} is a lattice

b) subset {a, b, f, g} is a lattice

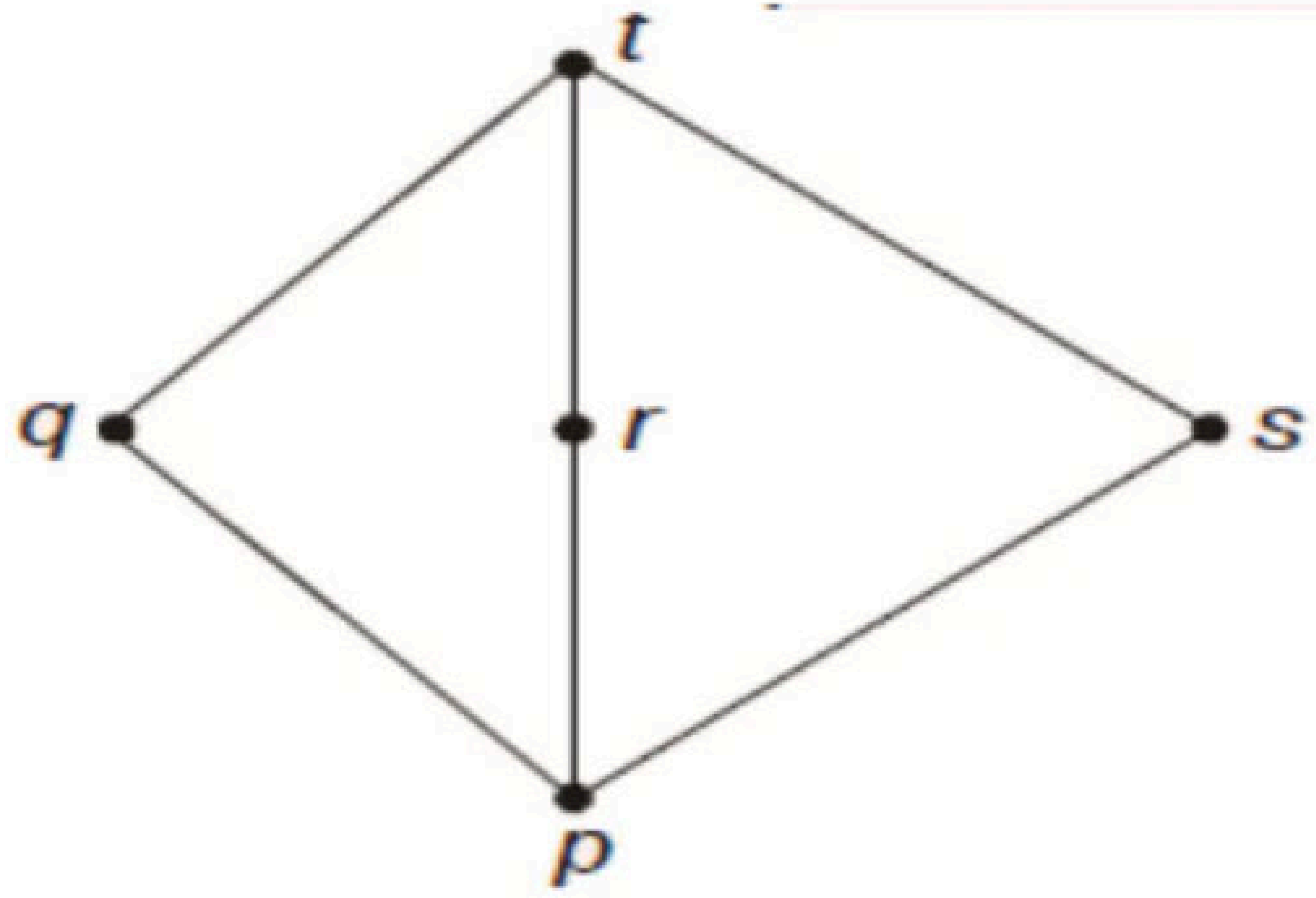
c) subset {a, d, e, g} is a lattice

d) subset {a, c, e, g} is a lattice

Break

Q Suppose L = {p, q, r, s, t} is a lattice represented by the following Hasse diagram: For any x, y ∈ L, not necessarily distinct, $x \vee y$ and $x \wedge y$ are join and meet of x, y respectively. Let $L^3 = \{(x, y, z): x, y, z \in L\}$ be the set of all ordered triplets of the elements of L. Let P_r be the probability that an element $(x, y, z) \in L^3$ chosen equiprobably satisfies $x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z)$. Then **(GATE-2015) (2 Marks)**

(A) Pr = 0 **(B)** Pr = 1 **(C)** 0 < Pr ≤ 1/5 **(D)** 1/5 < Pr < 1



Q Consider the set S = {a, b, c, d}.
Consider the following 4 partitions $\pi_1, \pi_2, \pi_3, \pi_4$ on
S: π_1={~~abcd~~}, π_2={~~ab~~, ~~cd~~}, π_3={~~abc~~, ~~d~~}, π_4={~~a~~, ~~b~~, ~~c~~, ~~d~~}
Let $\prec$ be the partial order on the set of partitions $S'=\{\pi_1,\pi_2,\pi_3,\pi_4\}$ defined as follows: $\pi_i \prec \pi_j$ if and only if π_i refines π_j. The Poset diagram for $(S', \prec)$ is: **(GATE-2007) (2 Marks)**