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# Recursive Sorting

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### Lecture Topics

• Recursive Bubble Sort

Merge Sort

Quicksort

#### Using Recursion to perform a Bubble Sort

• You've seen the bubble sort (and its variants) implemented using an iterative algorithm.

 Since any problem that can be solved iteratively can be solved recursively, we can design a recursive replacement for the iterative bubble sort.

### Bubble Sort (Iterative Algorithm)

```
void bubbleSort(int a[], int length) {
    for(int i = 0; i < length; i++) {
        for(int j = 1; j < length; j++) {
            if(a[j-1] > a[j]) {
                int temp = a[j-1];
                      a[j-1] = a[j];
                      a[j] = temp;
            }
        }
    }
}
```

#### Using Recursion to perform a Bubble Sort

- For an ascending sort, the first pass will move the largest value to the end of the array (length 1).
  - The next pass will move the second largest value to index length 2.
  - The next pass will move the third largest value to index length 3.
  - And so on...

- The last pass of the iterative algorithm will ensure the smallest value is placed in index 0.
  - At this point, the smallest value is guaranteed to already be in index 0.

#### Using Recursion to perform a Bubble Sort

- The base case is
  - When the algorithm sorts for index 0.
  - Since all other values will have already been sorted in the array, it implies the value at index 0 is already in the correct position.
- The recursive case is for all sorting other indexes.

### Bubble Sort (Recursive Algorithm)

```
void bubbleSort(int a[], int length) {
   if(length == 1) {
       return;
   for(int i = 0; i < length-1; i++) {
       if(a[i] > a[i+1]) {
           int temp = a[i+1];
           a[i+1] = a[i];
           a[i] = temp;
   bubbleSort(a, length-1);
```

#### Using Recursion to perform a Bubble Sort

• This recursive bubble sort algorithm, like the iterative version, also performs in polynomial time.

- There are no advantages (speed/complexity) by using an a recursive bubble sort over an iterative bubble sort.
  - This is example is here mainly to show you the two different implementations (recursive vs iterative).

#### Recursive Sorting Algorithms

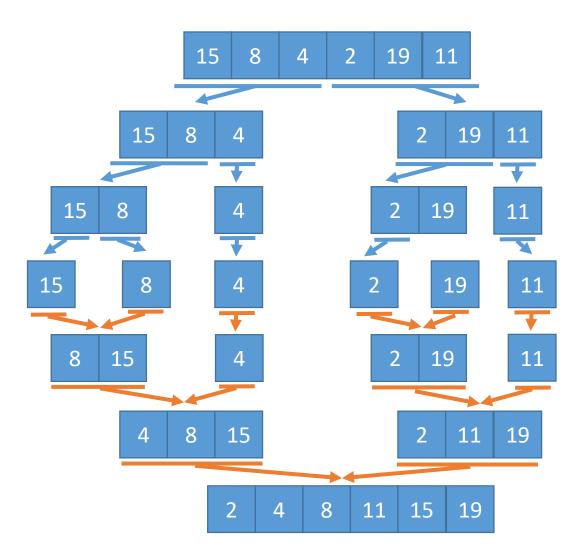
- We'll now see a pair of comparative sorting algorithms, Merge Sort and Quicksort, that are typically implemented using recursion.
  - They can be done iteratively, but the code is much easier to read when implemented recursively.
- Both algorithms take a "divide and conquer" approach to sorting, much like the how the binary search algorithm performs its searches.

#### Merge Sort Algorithm

- In the Merge Sort algorithm, the array is repeatedly (recursively) split in half until it reaches halves that only contain one element.
  - At this point, the lowest depth has been reached.
- Then, working backwards, it sorts/merges the smaller arrays back together

### Merge Sort Algorithm

- The image on the right gives the basic idea of how it works.
  - It divides up the array (Blue lines)
  - Then merges the array back together (Orange lines)
- Each merge involves two, sorted arrays.
  - Since the arrays to merge are in order, merging them is not computationally difficult.



#### Merge Sort Algorithm

- The C++ functions are a bit too long to put here in their entirety.
  - See the Sample Code provided.
- The next slides explain, very briefly, what is needed.
  - Two functions:
    - One function for the algorithm
    - Another function that handles the merging process

### Merge Sort

```
mergesort(array, left, right):

If left boundary < right boundary:

Find the middle, m

mergesort(array, left, m)

mergesort(array, m+1, right)

merge(array, left, m, right)
```

#### Merge Sort

#### Quicksort Algorithm

- In the Quicksort algorithm, the array is repeatedly (recursively) split into two smaller partitions, until the partitions only contain one element.
  - At this point, the lowest depth has been reached.
- The algorithm chooses a value in each partition, called the **pivot**.
  - One of the two partitions will contain any values less than the pivot.
  - The other partition will contain any values greater than the pivot.
- The process repeats recursively until partitions of length 1 are reached.
  - At which point, the array will have been sorted through the pivot processes.

### Quicksort Algorithm

- There are a few ways of selecting the pivot:
  - Always use the last element.
  - Always use the middle element.
  - Always use the first element.
  - Always use a randomly chosen element.
- The Sample Code provided uses the middle element as the selected pivot value.
- The next slides explain, very briefly, what is needed.

#### Quicksort

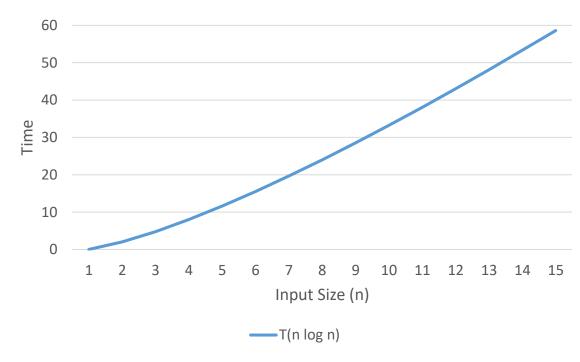
```
quicksort(array, start, end) :
    Select the pivot (the value in the middle)
    Partition the array
    quicksort(lower half)
    quicksort(upper half)
```

#### Merge Sort and Quicksort

- Both algorithms use the divide and conquer process like a binary search.
  - Which we already determined performs in logarithmic time.
- Both algorithms, at one point or another, will have partitions with a length of 1.
  - The algorithms will perform the logarithmic divide and conquer operations for as many elements that exist in the array.
- Their time complexity is a mix of logarithmic and linear time.

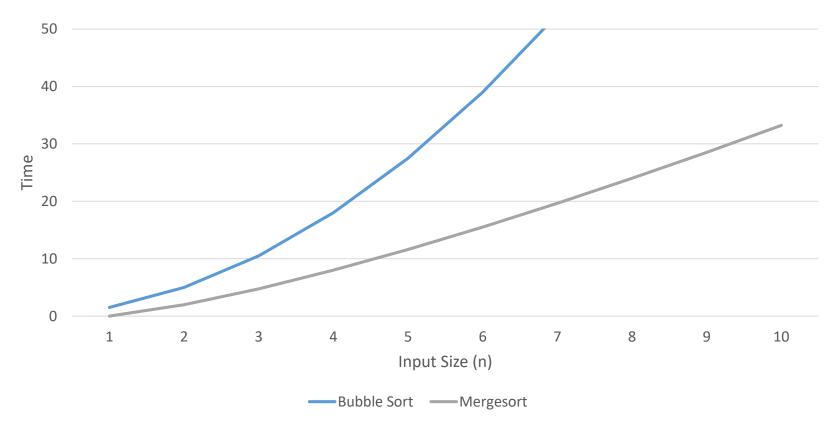
#### Quasi-Linear Time

- Quasi-Linear (or Log-Linear) time is when an algorithm executes nnumber of operations, where each operation performs in logarithmic time.
  - $T(n) = n \log_2 n$
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  - $T(4) = 4 \log_2 4$
  - $T(8) = 8 \log_2 8$
  - $T(12) = 12 \log_2 12$



### Bubble Sort vs Merge Sort

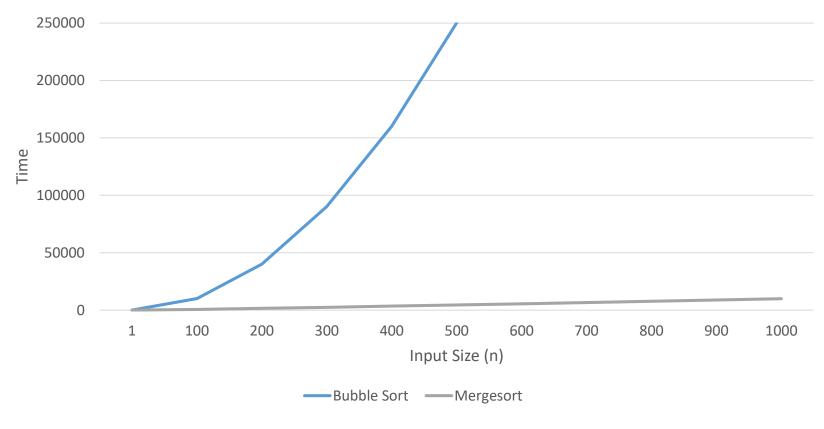
• Max array length: 10



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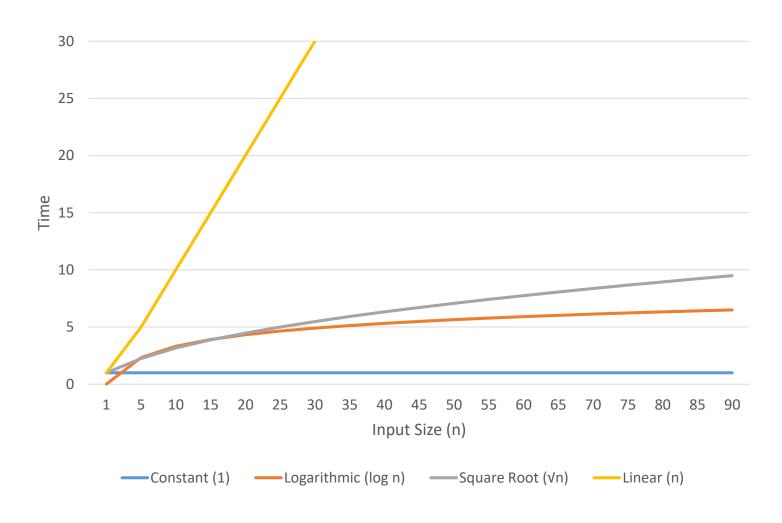
#### Bubble Sort vs Merge Sort

• Max array length: 1000



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# Complexities



## Complexities

