Revised: 3/6/2020

Graphs I

Michael C. Hackett
Computer Science Department

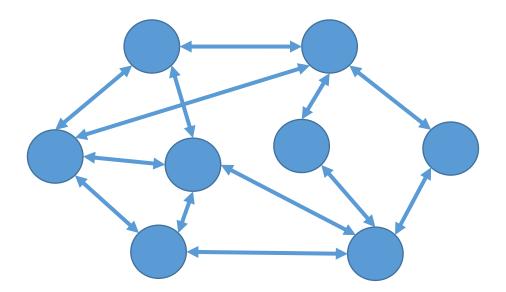
Community College of Philadelphia

Lecture Topics

- Graph Terminology
- Implementing Graphs
 - Adjacency Lists
 - Adding an edge
 - Removing an edge
 - Checking if an edge exists
 - Adjacency Matrices
 - Adding an edge
 - Removing an edge
 - Checking if an edge exists
 - Complexity Comparison

- Directed Graphs
 - Adding an edge
 - Removing an edge
 - Checking if an edge exists
 - Complexity Comparison

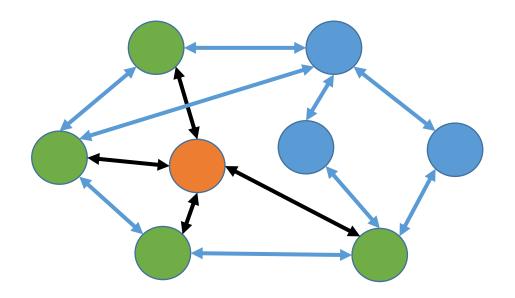
• A **graph** is a non-linear data structure, where each node is connected by one or more edges.



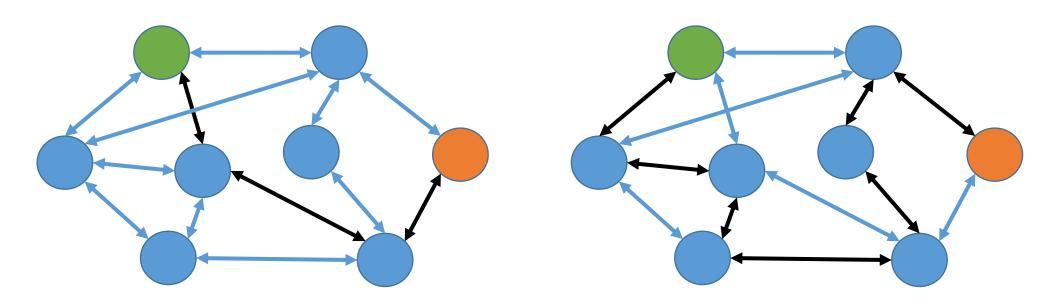
• There is no "root" node.

- Each node has one or more edges.
 - A node could connect to every other node, or connect to only one.
- No parent/child relationships

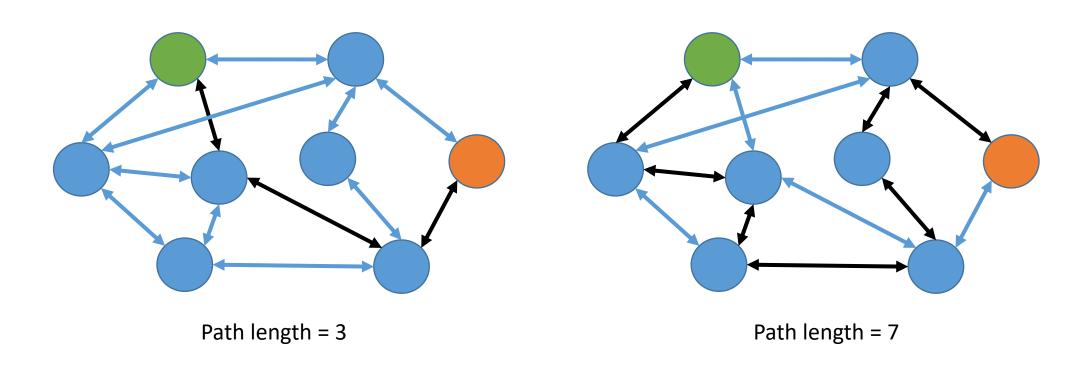
- Two nodes are **adjacent** if they are connected by an edge.
 - "Neighbors"
 - The green nodes below are adjacent to the orange node



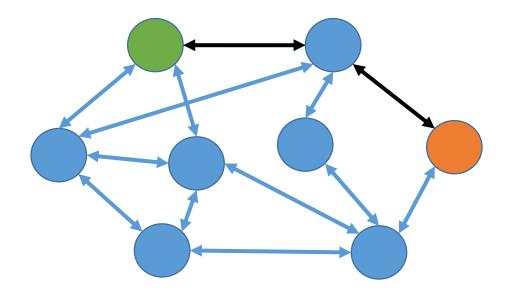
- A path is the sequence of edges between two nodes in the structure.
 - The green node below is the starting node and the orange node is the ending node; Black edges are a possible path.
 - More than one path may exist.



• The path length is the number of edges in a path.

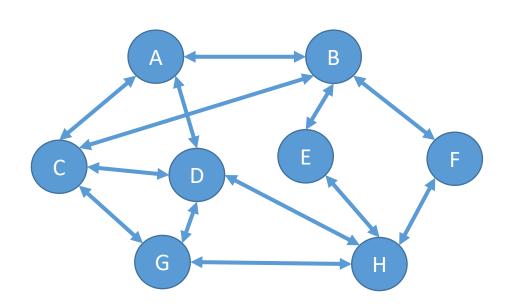


- The **distance** between two nodes is the <u>shortest</u> path length between them.
 - The distance between the two nodes in this example graph is 2.



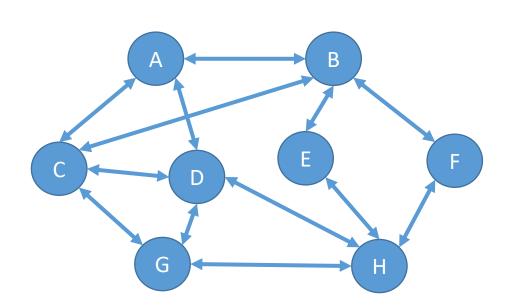
Adjacency Lists

- One way to implement a graph is using an adjacency list.
 - Each node corresponds to a List of adjacent nodes



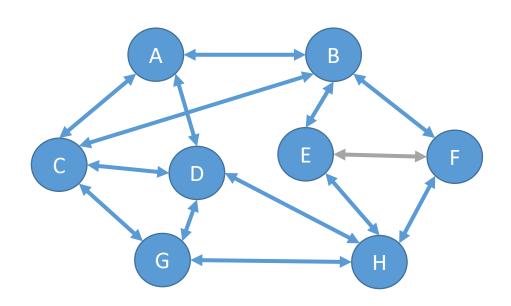
Node	Adjacent Nodes					
А	B, C, D					
В	A, C, E, F					
С	A, B, D, G					
D	A, C, G, H					
E	В, Н					
F	В, Н					
G	C, D, H					
Н	D, E, F, G					

 To add an edge between two nodes, each node is placed in the other's adjacency list



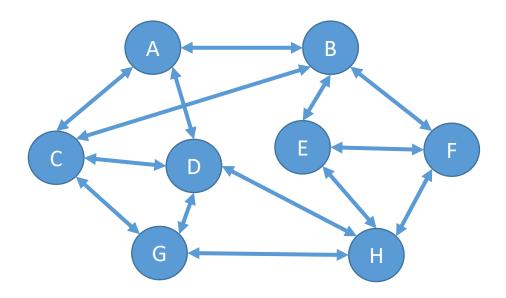
Node	Adjacent Nodes				
Α	B, C, D				
В	A, C, E, F				
С	A, B, D, G				
D	A, C, G, H				
E	В, Н				
F	В, Н				
G	C, D, H				
Н	D, E, F, G				

Adding an edge between E and F



Node	Adjacent Nodes				
А	B, C, D				
В	A, C, E, F				
С	A, B, D, G				
D	A, C, G, H				
E	В, Н				
F	В, Н				
G	C, D, H				
Н	D, E, F, G				

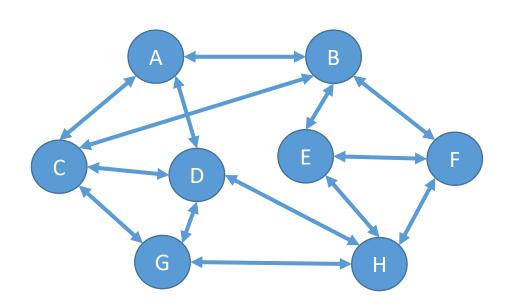
- E is added to F's adjacency list
- F is added to E's adjacency list



Node	Adjacent Nodes			
А	B, C, D			
В	A, C, E, F			
С	A, B, D, G			
D	A, C, G, H			
E	В, Н, F			
F	В, Н, Е			
G	C, D, H			
Н	D, E, F, G			

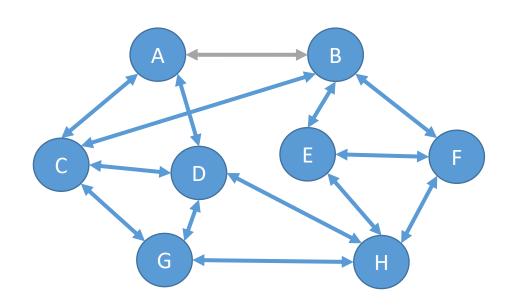
- Time Complexity (Adding an edge): O(1)
 - Adding an edge simply appends the node/edge to a node's adjacency list
 - Adding/appending to a List takes constant time

 To remove an edge between two nodes, each node is removed from the other's adjacency list



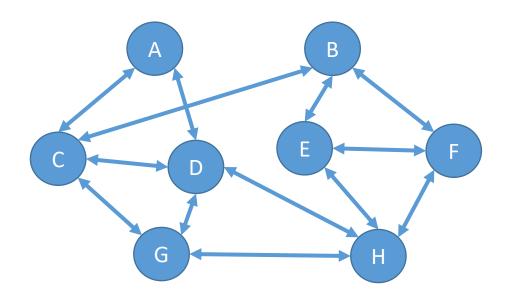
Node	Adjacent Nodes			
Α	B, C, D			
В	A, C, E, F			
С	A, B, D, G			
D	A, C, G, H			
E	B, H, F			
F	В, Н, Е			
G	C, D, H			
Н	D, E, F, G			

Removing the edge between A and B



Node	Adjacent Nodes			
А	B, C, D			
В	A, C, E, F			
С	A, B, D, G			
D	A, C, G, H			
E	B, H, F			
F	В, Н, Е			
G	C, D, H			
Н	D, E, F, G			

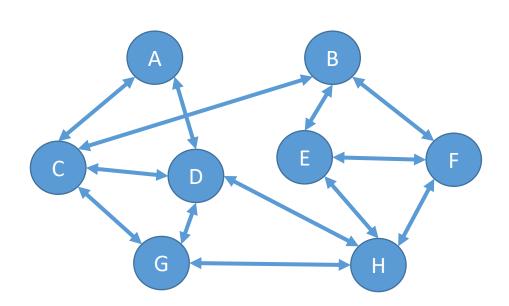
- B is removed from A's adjacency list
- A is removed from B's adjacency list



Node	Adjacent Nodes				
А	B , C, D				
В	A, C, E, F				
С	A, B, D, G				
D	A, C, G, H				
E	B, H, F				
F	В, Н, Е				
G	C, D, H				
Н	D, E, F, G				

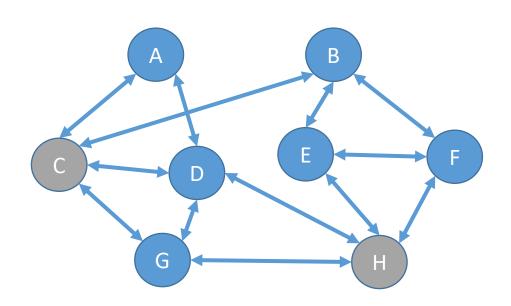
- Time Complexity (Removing an edge): O(e₁ + e₂)
 - e_1 = Number of nodes/edges in the first node's adjacency list
 - e_2 = Number of nodes/edges in the second node's adjacency list
 - The first node's adjacency list is searched (linearly) to find and remove the second node
 - The second node's adjacency list is searched (linearly) to find and remove the first node

 To check if an edge exists between two nodes, the node's adjacency list is checked for the other node



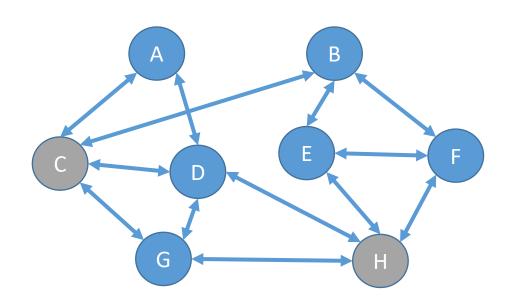
Node	Adjacent Nodes				
Α	C, D				
В	C, E, F				
С	A, B, D, G				
D	A, C, G, H				
E	B, H, F				
F	В, Н, Е				
G	C, D, H				
Н	D, E, F, G				

Checking for an edge between C and H



Node	Adjacent Nodes				
Α	C, D				
В	C, E, F				
С	A, B, D, G				
D	A, C, G, H				
E	B, H, F				
F	В, Н, Е				
G	C, D, H				
Н	D, E, F, G				

- H is not in C's adjacency list
 - Alternatively could have checked H's adjacency list for C
 - Does not imply a path doesn't exist

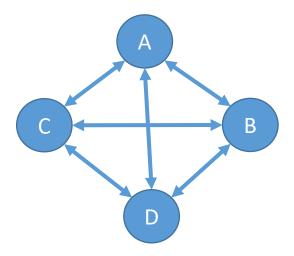


Node	Adjacent Nodes				
Α	C, D				
В	C, E, F				
С	A, B, D, G				
D	A, C, G, H				
E	В, Н, F				
F	В, Н, Е				
G	C, D, H				
Н	D, E, F, G				

- Time Complexity (Checking for an edge): O(e)
 - e = Number of nodes/edges in the node's adjacency list
 - Determining one node's adjacency to a second node is verified by iterating through the adjacent nodes in either node's adjacency list.

Adjacency Lists – Space Complexity

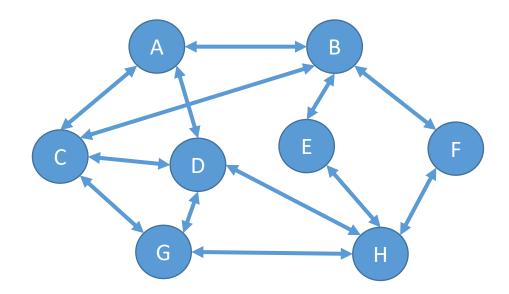
- (Worst Case) Space Complexity: O(V+E)
 - V = Total number of vertices (Nodes)
 - E = Total number of edges (Adjacent Nodes)



Node	Adjacent Nodes
А	B, C, D
В	A, C, D
С	A, B, D
D	A, B, C

Adjacency Matrices

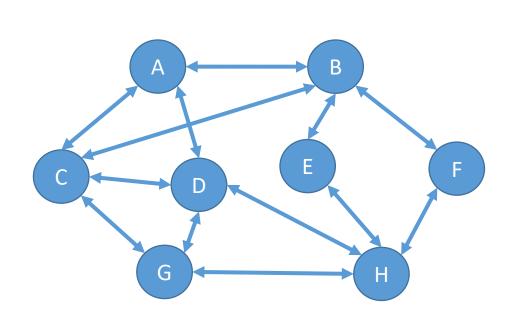
- Another way to implement a graph is using an adjacency matrix.
 - 2-D Array
 - 1 if an edge exists, 0 if not



	A	В	С	D	Е	F	G	Н
A		1	1	1				
В	1		1		1	1		
С	1	1		1			1	
D	1		1				1	1
Е		1						1
F		1						1
G			1	1				1
Н				1	1	1	1	

Adjacency Matrices – Adding an Edge

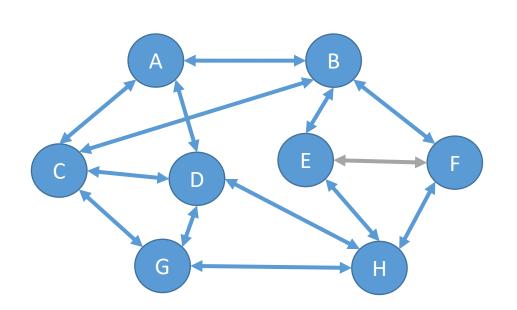
• To add an edge between two nodes, a 1 is placed in the matrix



	A	В	С	D	Ε	F	G	Н
Α		1	1	1				
В	1		1		1	1		
С	1	1		1			1	
D	1		1				1	1
E		1						1
F		1						1
G			1	1				1
Н				1	1	1	1	

Adjacency Matrices – Adding an Edge

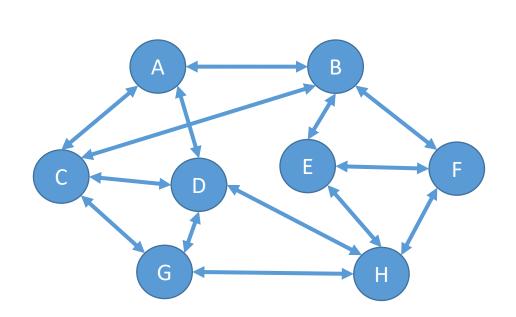
Adding an edge between E and F



	A	В	С	D	E	F	G	Н
Α		1	1	1				
В	1		1		1	1		
С	1	1		1			1	
D	1		1				1	1
E		1						1
F		1						1
G			1	1				1
Н				1	1	1	1	

Adjacency Matrices – Adding an Edge

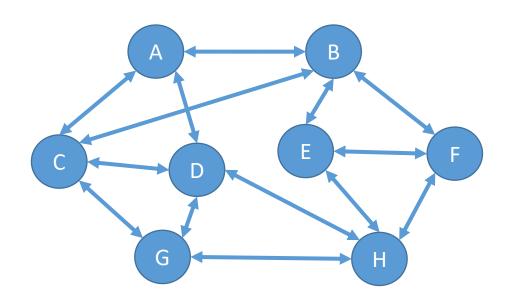
A 1 is placed in the corresponding locations in the matrix



	A	В	С	D	Е	F	G	Н
Α		1	1	1				
В	1		1		1	1		
С	1	1		1			1	
D	1		1				1	1
E		1				1		1
F		1			1			1
G			1	1				1
Н				1	1	1	1	

Adjacency Matrices – Removing an Edge

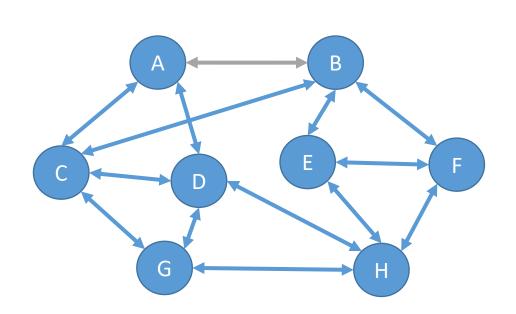
- To remove an edge between two nodes, a 0 is placed in the corresponding locations in the matrix
 - In this example, blank cells represent 0



	A	В	С	D	E	F	G	Н
Α		1	1	1				
В	1		1		1	1		
С	1	1		1			1	
D	1		1				1	1
E		1				1		1
F		1			1			1
G			1	1				1
Н				1	1	1	1	

Adjacency Matrices – Removing an Edge

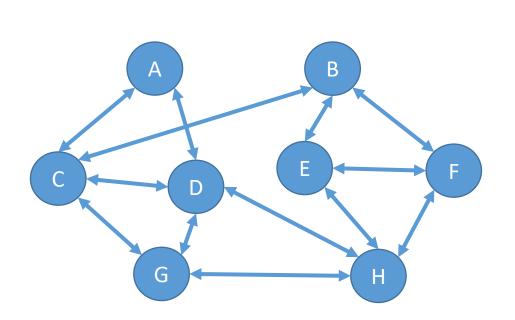
Removing the edge between A and B



	A	В	С	D	Е	F	G	Н
Α		1	1	1				
В	1		1		1	1		
С	1	1		1			1	
D	1		1				1	1
Е		1				1		1
F		1			1			1
G			1	1				1
Н				1	1	1	1	

Adjacency Matrices – Removing an Edge

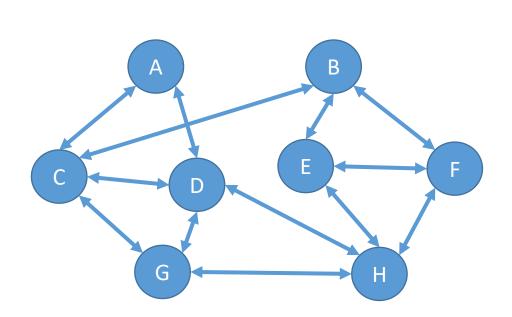
A 0 is placed in the corresponding locations in the matrix



	A	В	С	D	Е	F	G	Н
Α		0	1	1				
В	0		1		1	1		
С	1	1		1			1	
D	1		1				1	1
Е		1				1		1
F		1			1			1
G			1	1				1
Н				1	1	1	1	

Adjacency Matrices – Checking for an edge

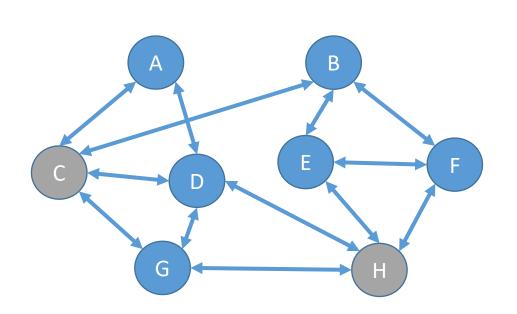
 To check if an edge exists between two nodes, that location in the matrix is checked for a 1



	A	В	С	D	E	F	G	Н
Α			1	1				
В			1		1	1		
С	1	1		1			1	
D	1		1				1	1
E		1				1		1
F		1			1			1
G			1	1				1
Н				1	1	1	1	

Adjacency Matrices – Checking for an edge

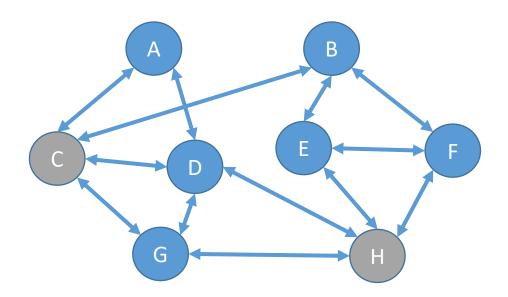
Checking if C is adjacent to H



	A	В	С	D	E	F	G	Н
Α			1	1				
В			1		1	1		
С	1	1		1			1	
D	1		1				1	1
Е		1				1		1
F		1			1			1
G			1	1				1
Н				1	1	1	1	

Adjacency Matrices – Checking for an edge

- There is not a 1 in [C][H]
 - Alternatively could have checked [H][C]
 - Does not imply a path doesn't exist



	A	В	С	D	E	F	G	Н
Α			1	1				
В			1		1	1		
С	1	1		1			1	0
D	1		1				1	1
Е		1				1		1
F		1			1			1
G			1	1				1
Н			0	1	1	1	1	

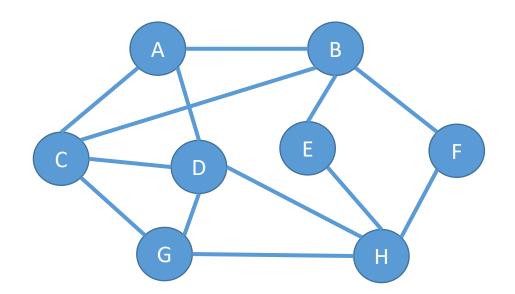
Adjacency Matrices

- Time Complexity (Adding an edge): **O(1)**
- Time Complexity (Removing an edge): **O(1)**
- Time Complexity (Checking for an edge): O(1)

Array access takes constant time

Adjacency Matrices

- Space Complexity: always O(V²)
 - V = Number of vertices



	A	В	С	D	Е	F	G	Н
A		1	1	1				
В	1		1		1	1		
С	1	1		1			1	
D	1		1				1	1
Е		1						1
F		1						1
G			1	1				1
Н				1	1	1	1	

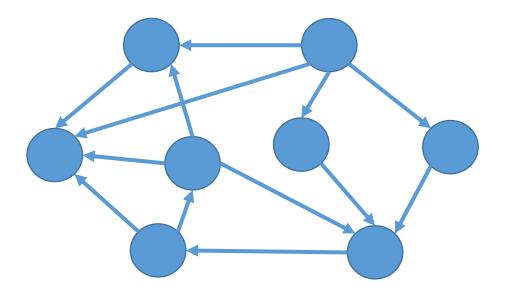
Complexity Comparison

Graph Implementation	Add Edge	Remove Edge	Check Edge	Space
Adjacency List	O(1)	$O(e_1 + e_2)$	O(e)	O(V + E)
Adjacency Matrix	O(1)	O(1)	O(1)	$O(V^2)$

Green - Constant Orange - Linear Red - Polynomial

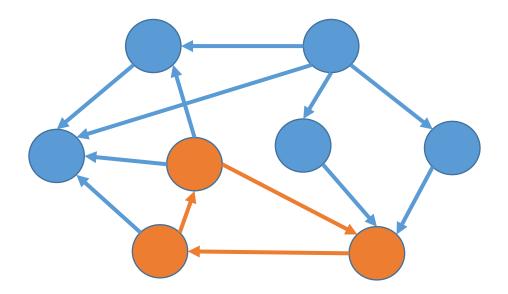
Directed Graphs

- A directed graph (or digraph) is a graph where each node is connected via <u>one-way</u> edges.
 - Previously, we were using a bi-directional graph



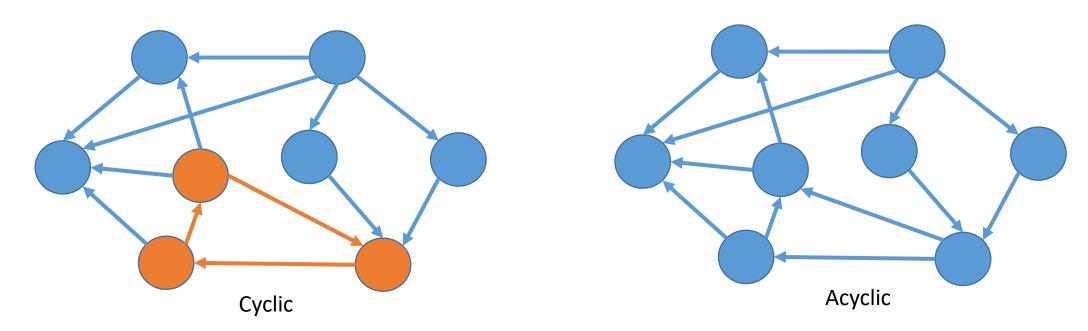
Directed Graphs

• A **cycle** is a path that can begin and end at the same node in a directed graph.



Directed Graphs

- A graph that contains a cycle is cyclic
 - Bi-directional graphs are inherently cyclic
- A graph that does not contain a cycle is acyclic



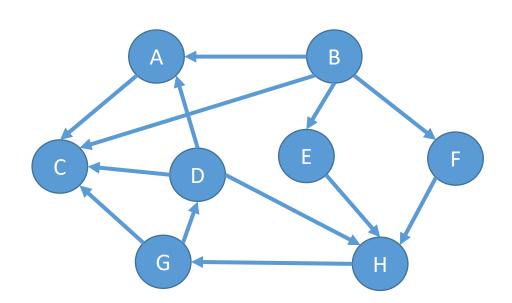
Directed Graphs

 Digraphs can be implemented using adjacency lists or adjacency matrices

 Processes for adding, removing, and checking edges in digraphs are mostly the same as for bi-directional graphs

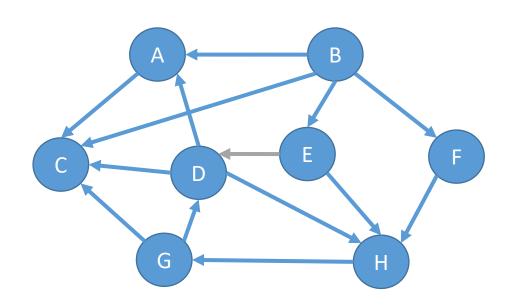
The remaining slides only show processes for a digraph using an adjacency list

 To add an edge between two nodes, the end node is placed in the starting node's adjacency list



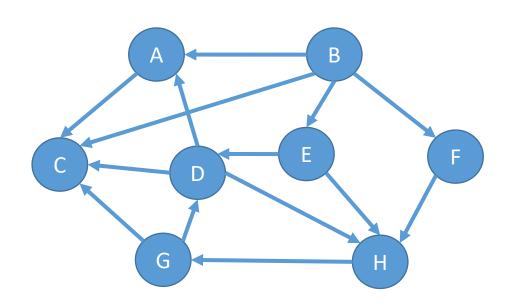
Node	Adjacent Nodes
А	С
В	A, C, E, F
С	
D	A, C, H
E	Н
F	Н
G	C, D
Н	G

Adding an edge from E to D



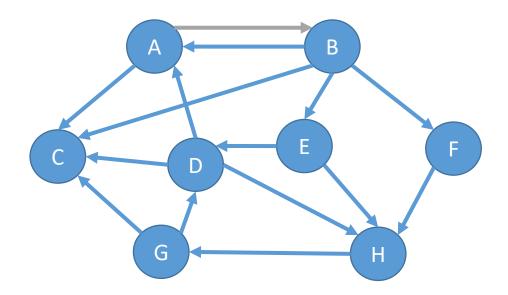
Node	Adjacent Nodes			
А	С			
В	A, C, E, F			
С				
D	A, C, H			
E	Н			
F	Н			
G	C, D			
Н	G			

• D is added to E's adjacency list



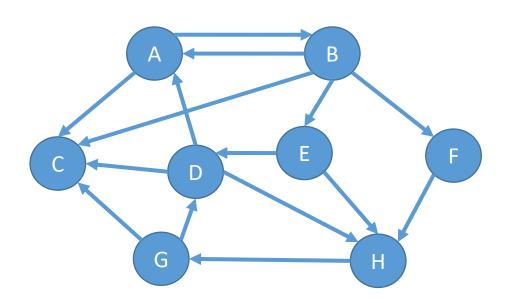
Node	Adjacent Nodes
А	С
В	A, C, E, F
С	
D	A, C, H
E	H, D
F	Н
G	C, D
Н	G

- Digraphs can have bi-directional edges
 - Adding an edge from A to B



Node	Adjacent Nodes
А	С
В	A, C, E, F
С	
D	A, C, H
E	H, D
F	Н
G	C, D
Н	G

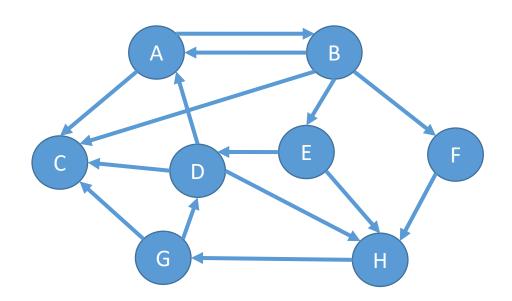
• B is added to A's adjacency list



Node	Adjacent Nodes
А	С, В
В	A, C, E, F
С	
D	A, C, H
E	H, D
F	Н
G	C, D
Н	G

Directed Graphs – Removing an Edge

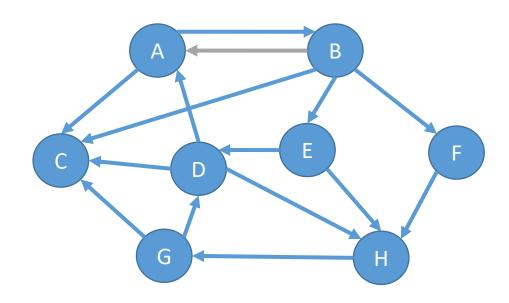
 To remove an edge, the end node is removed from the starting node's adjacency list



Node	Adjacent Nodes
А	С, В
В	A, C, E, F
С	
D	A, C, H
E	H, D
F	Н
G	C, D
Н	G

Directed Graphs – Removing an Edge

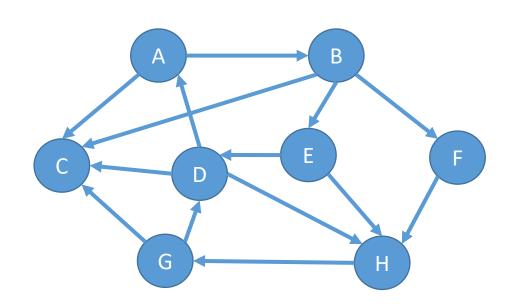
Removing the edge from B to A



Node	Adjacent Nodes
А	С, В
В	A, C, E, F
С	
D	A, C, H
E	H, D
F	Н
G	C, D
Н	G

Directed Graphs – Removing an Edge

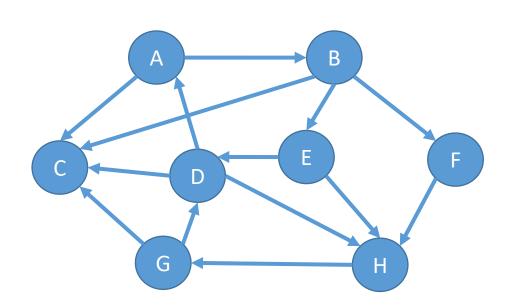
A is removed from B's adjacency list



Node	Adjacent Nodes
А	С, В
В	C, E, F
С	
D	A, C, H
Е	H, D
F	Н
G	C, D
Н	G

Directed Graphs – Checking for an edge

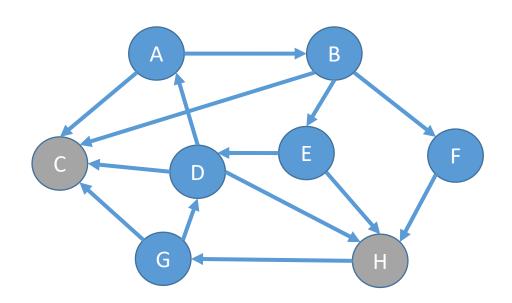
 To check if an edge exists between two nodes, the starting node's adjacency list is checked for the other node



Node	Adjacent Nodes
А	С, В
В	C, E, F
С	
D	A, C, H
E	H, D
F	Н
G	C, D
Н	G

Directed Graphs – Checking for an edge

Checking if H is adjacent to C



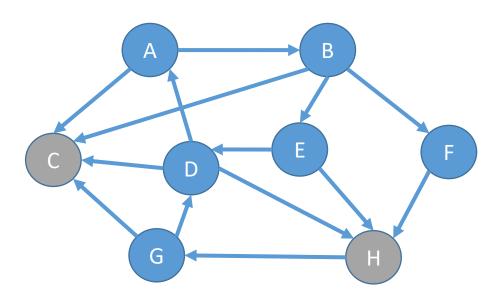
Node	Adjacent Nodes
Α	С, В
В	C, E, F
С	
D	A, C, H
E	H, D
F	Н
G	C, D
Н	G

Directed Graphs – Checking for an edge

• C does not exist in H's adjacency list

• Since this is a digraph, we could not alternatively check to see if H is in C's

adjacency list



Node	Adjacent Nodes
Α	С, В
В	C, E, F
С	
D	A, C, H
E	H, D
F	Н
G	C, D
Н	G

Directed Graphs – Complexity Comparison

Graph Implementation	Add Edge	Remove Edge	Check Edge	Space
Adjacency List	O(1)	$O(e_1 + e_2)$	O(e)	O(V + E)
Adjacency Matrix	O(1)	O(1)	O(1)	$O(V^2)$
Adjacency List (Digraph)	O(1)	O(e)	O(e)	O(V + E)
Adjacency Matrix (Digraph)	O(1)	O(1)	O(1)	$O(V^2)$

Green - Constant Orange - Linear Red - Polynomial