

# Graphs II

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# Lecture Topics

- Graph Traversals
  - Depth-First
  - Breadth-First
- Finding distance with breadth-first
- Weighted Graphs
  - Dijkstra's Algorithm

# Graph Traversal

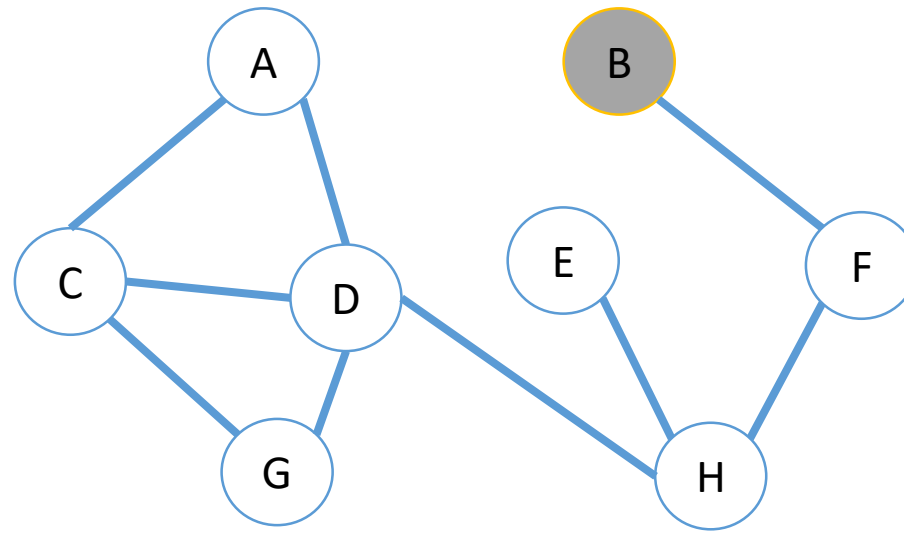
- Graphs can be traversed using depth-first or breadth-first traversals.
- Similar to traversing a tree, but we can begin at any node in the graph.
- Depth-first traversal uses a *stack*
- Breadth-first traversal uses a *queue*

# Depth-First Traversal

- Depth-first traversals work by coloring each node
  - White: Node has not been seen previously
  - Gray: Node has been seen previously
  - Black: Node has been visited and we are done with it
- This prevents getting stuck in a cycle, if one existed
  - The traversal is the same for bi-directional graphs or digraphs

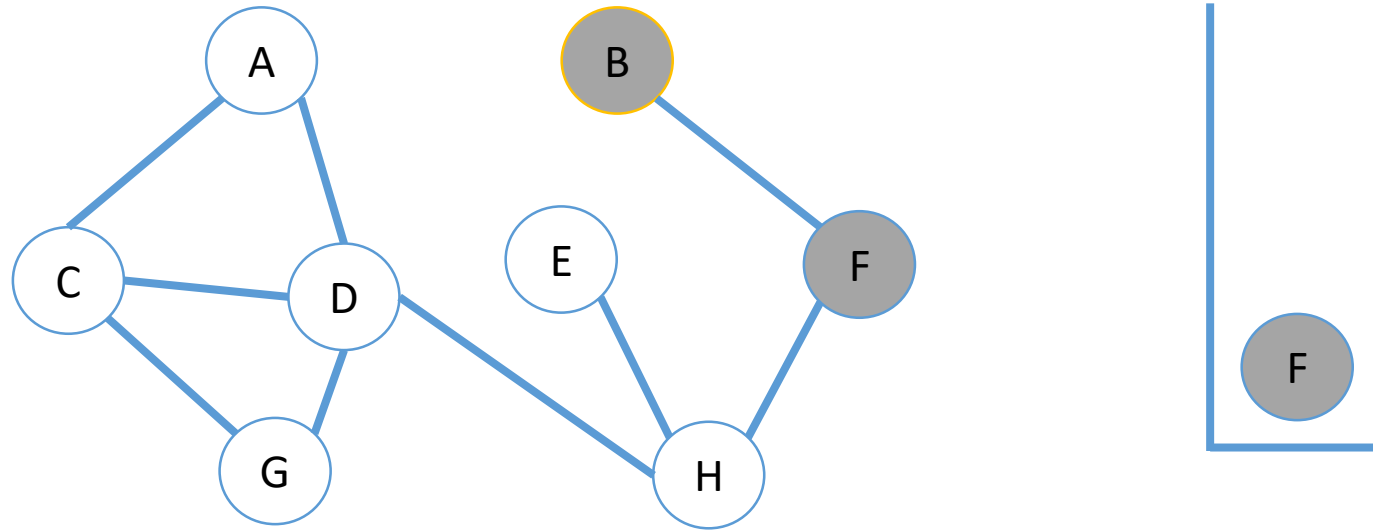
# Depth-First Traversal

- Starting at B
- Order visited: N/A



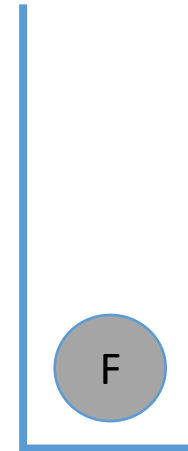
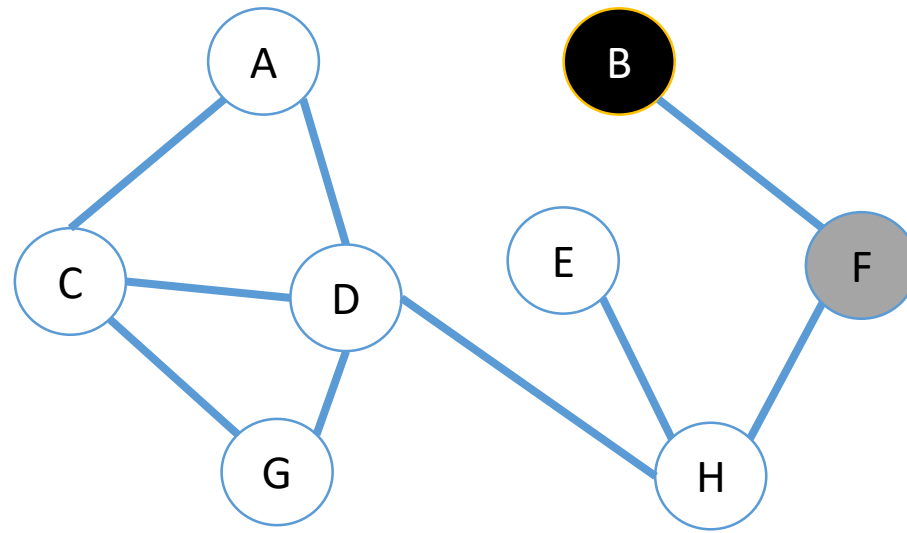
# Depth-First Traversal

- Adds its adjacent white node(s) to the stack (turning the white node(s) gray)
- Order visited: N/A



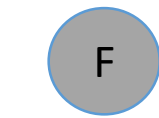
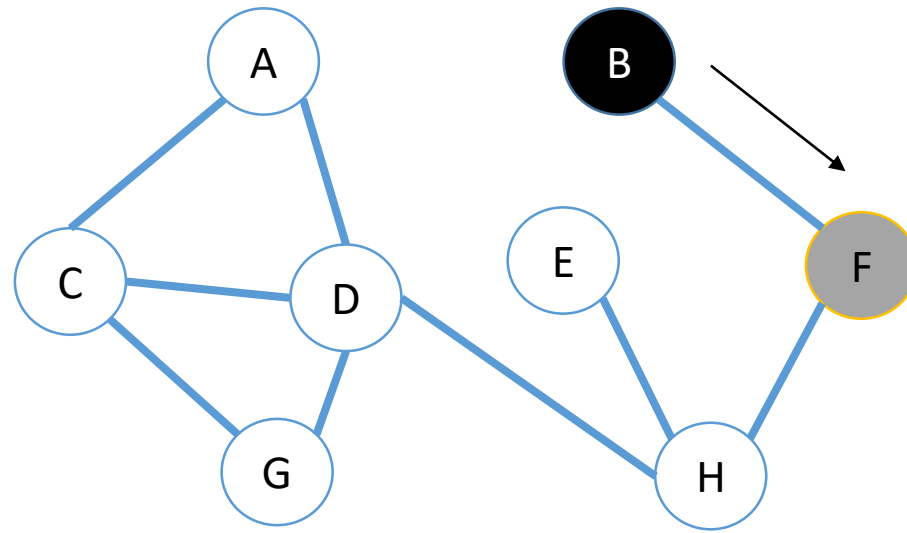
# Depth-First Traversal

- B was visited
  - Changed to black
- Order visited: B



# Depth-First Traversal

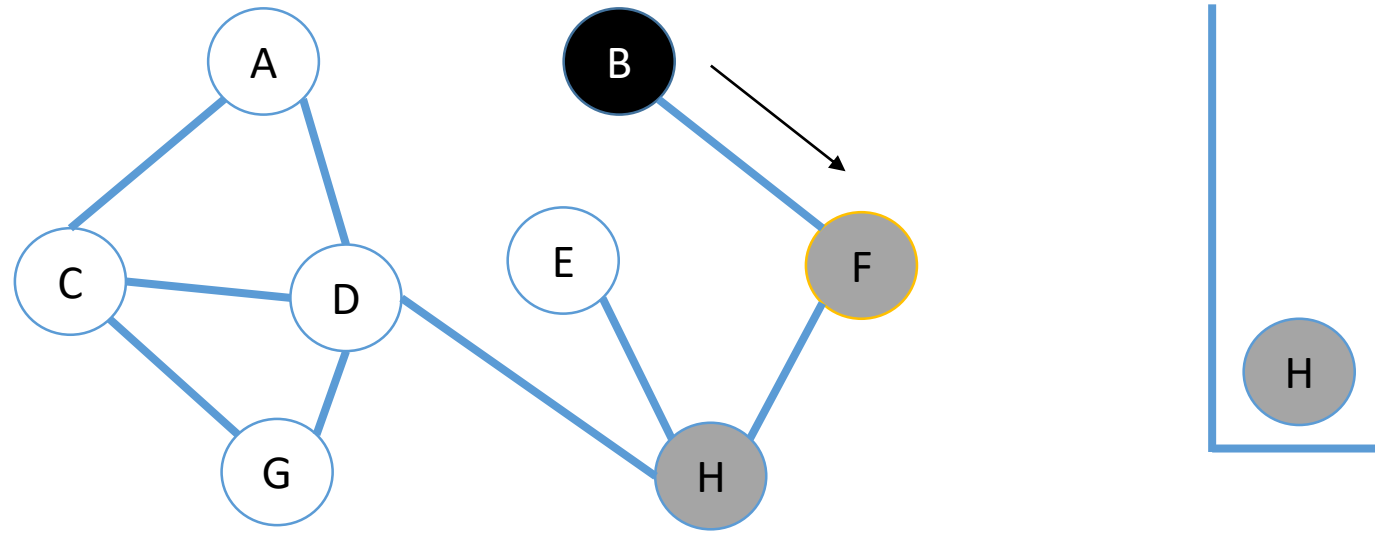
- Next node popped from stack
- Order visited: B





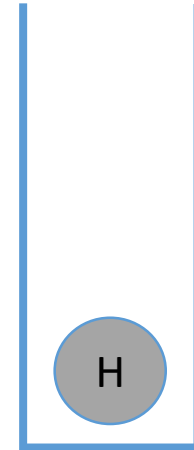
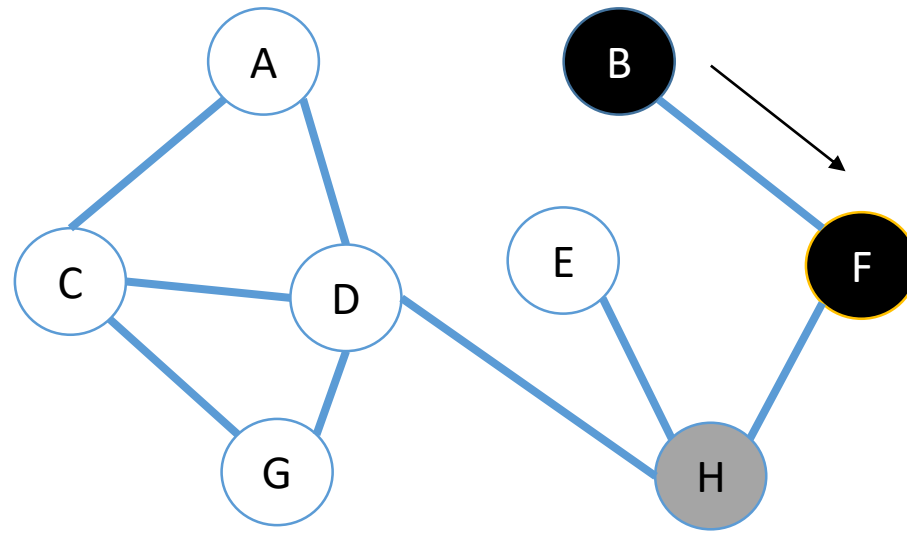
# Depth-First Traversal

- Adds its adjacent white nodes to the stack (turning them gray)
- Order visited: B



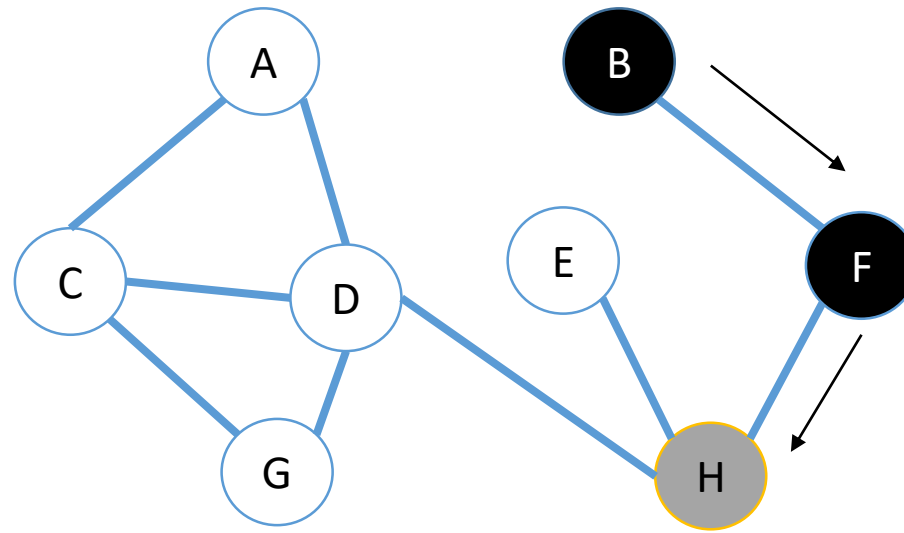
# Depth-First Traversal

- F was visited
- Order visited: B, F



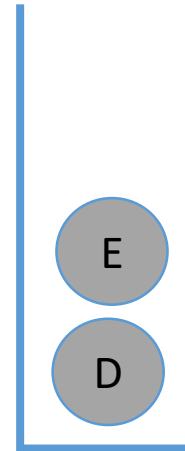
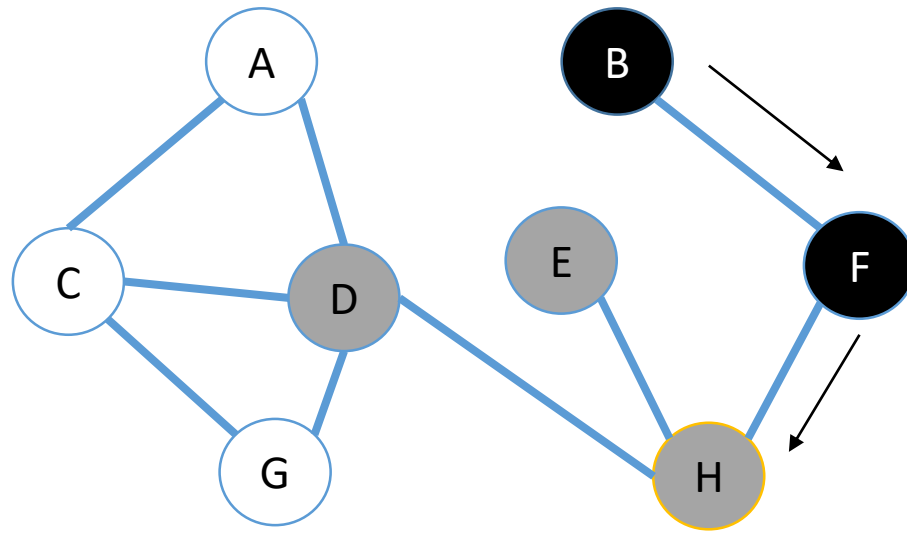
# Depth-First Traversal

- Next node popped from stack
- Order visited: B, F



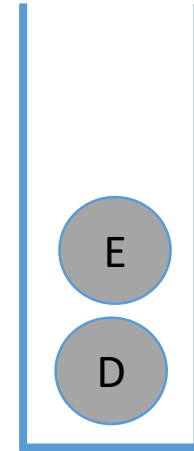
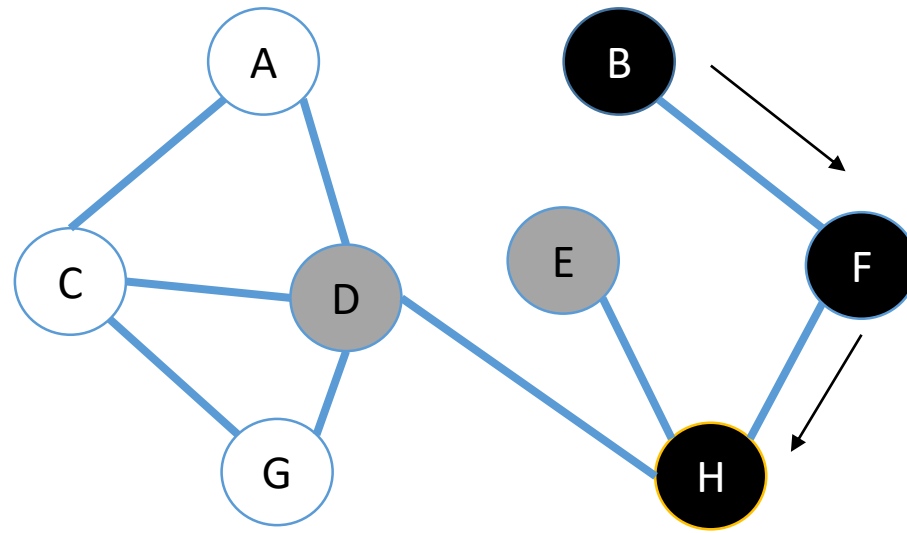
# Depth-First Traversal

- Adds its adjacent white nodes to the stack
- Order visited: B, F



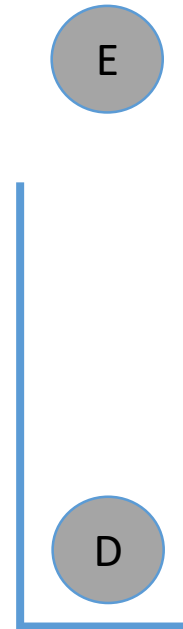
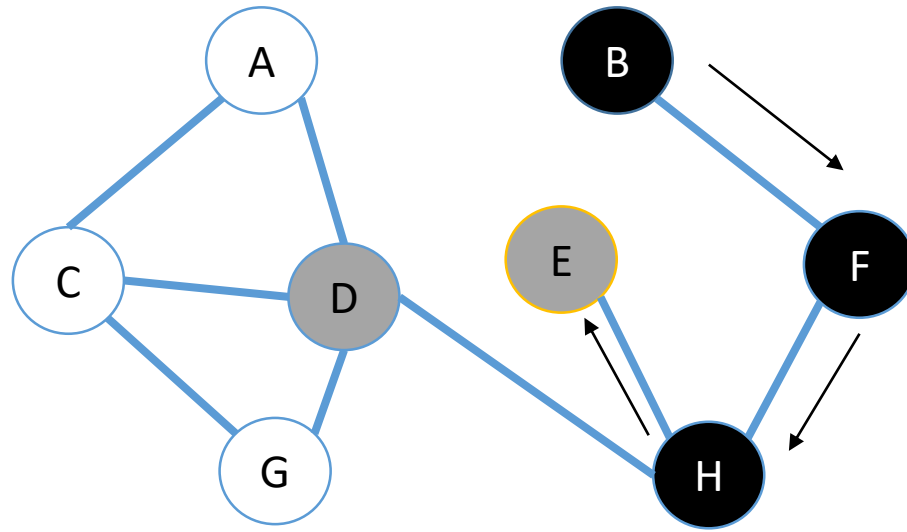
# Depth-First Traversal

- H was visited
- Order visited: B, F, H



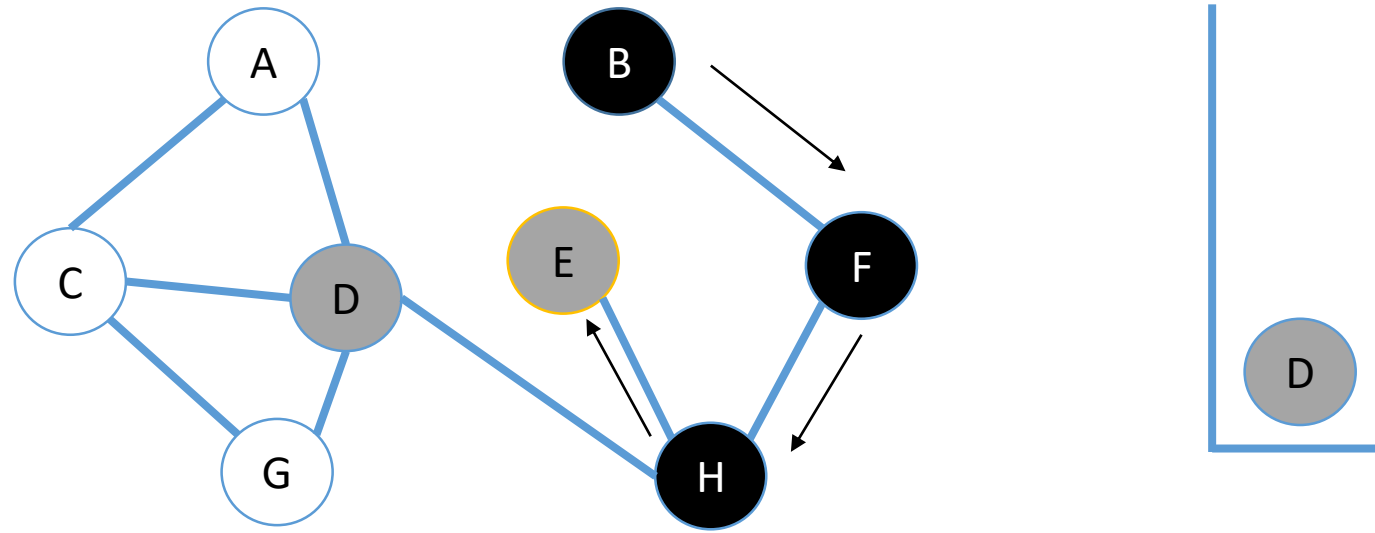
# Depth-First Traversal

- Next node popped from stack
- Order visited: B, F, H



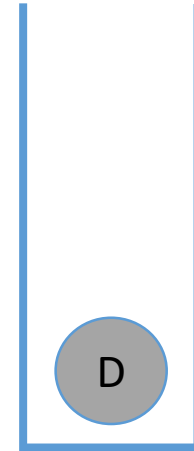
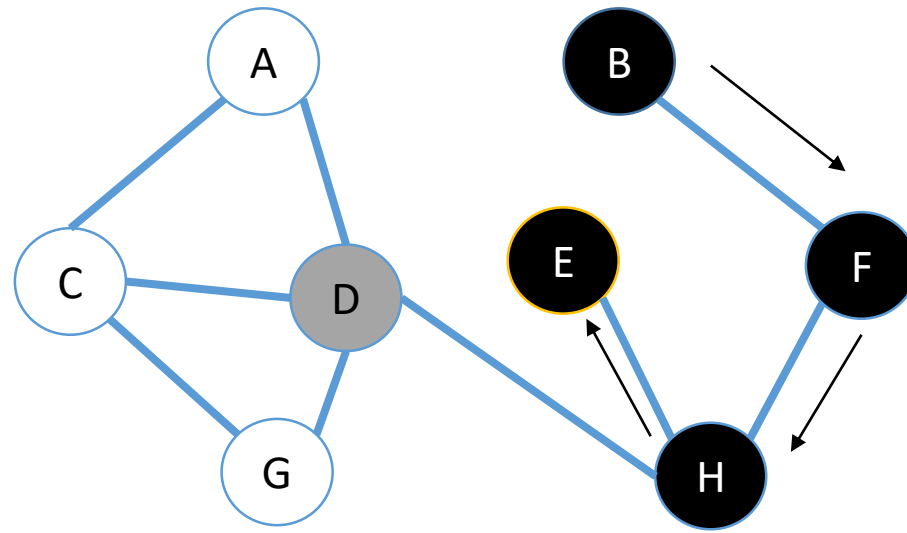
# Depth-First Traversal

- Add its adjacent white nodes to the stack (there are none)
- Order visited: B, F, H



# Depth-First Traversal

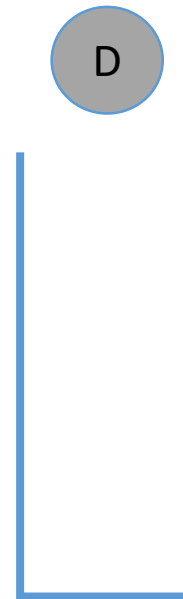
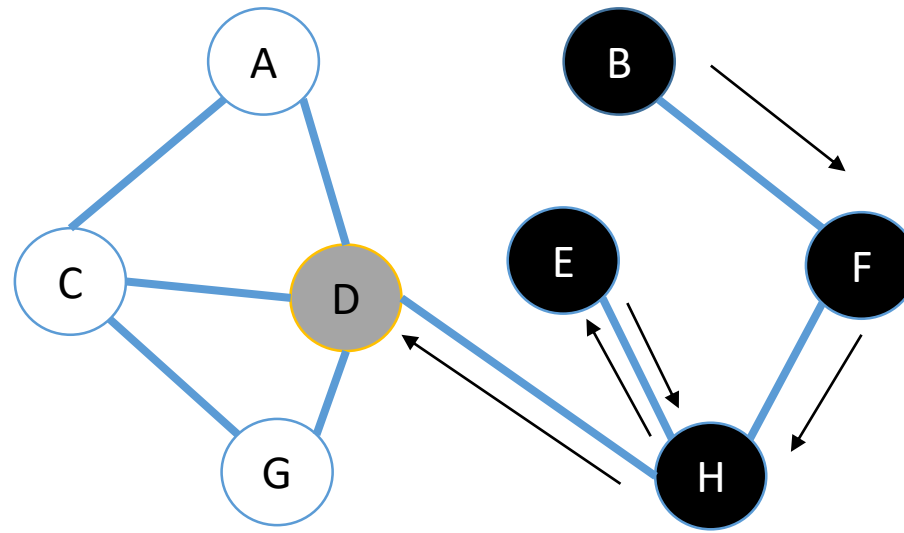
- E was visited
- Order visited: B, F, H, E





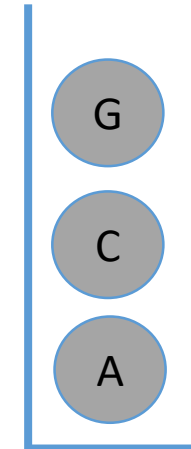
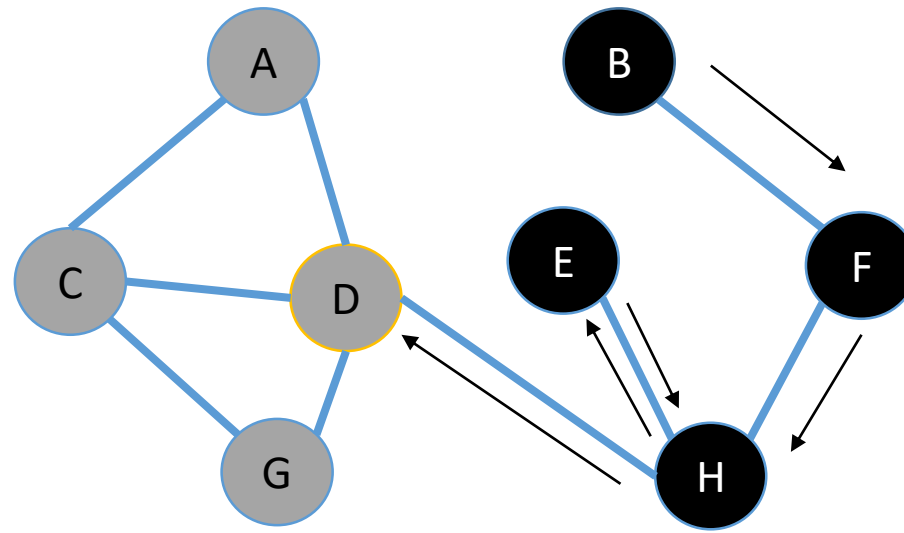
# Depth-First Traversal

- Next node popped from stack
- Order visited: B, F, H, E



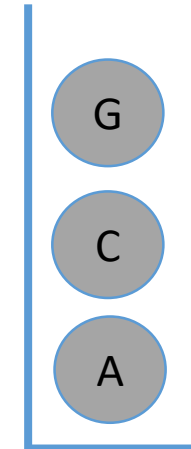
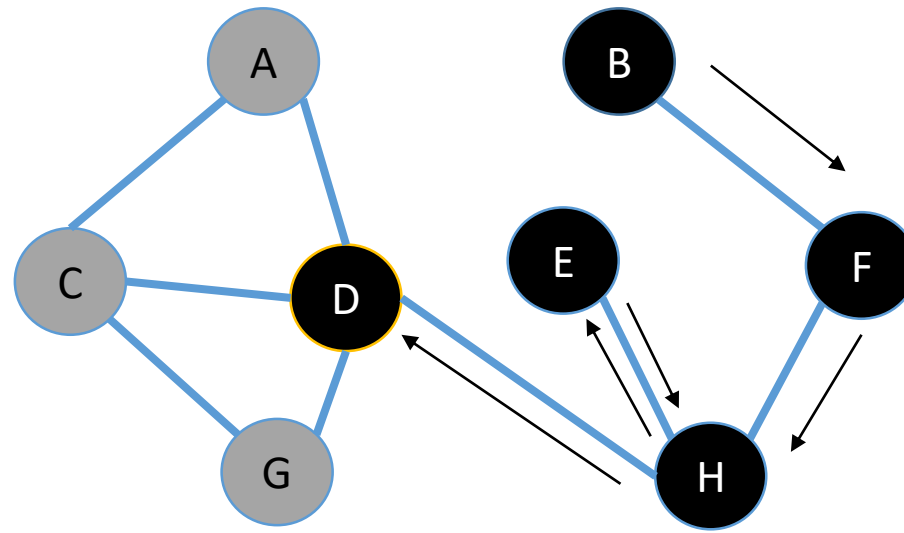
# Depth-First Traversal

- Add its adjacent white nodes to the stack
- Order visited: B, F, H, E



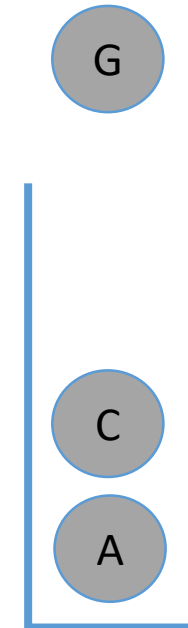
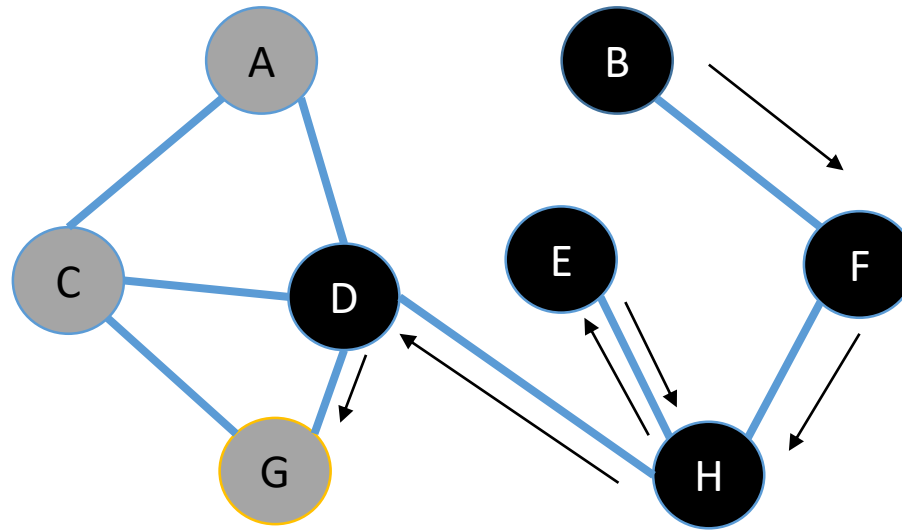
# Depth-First Traversal

- D was visited
- Order visited: B, F, H, E, D



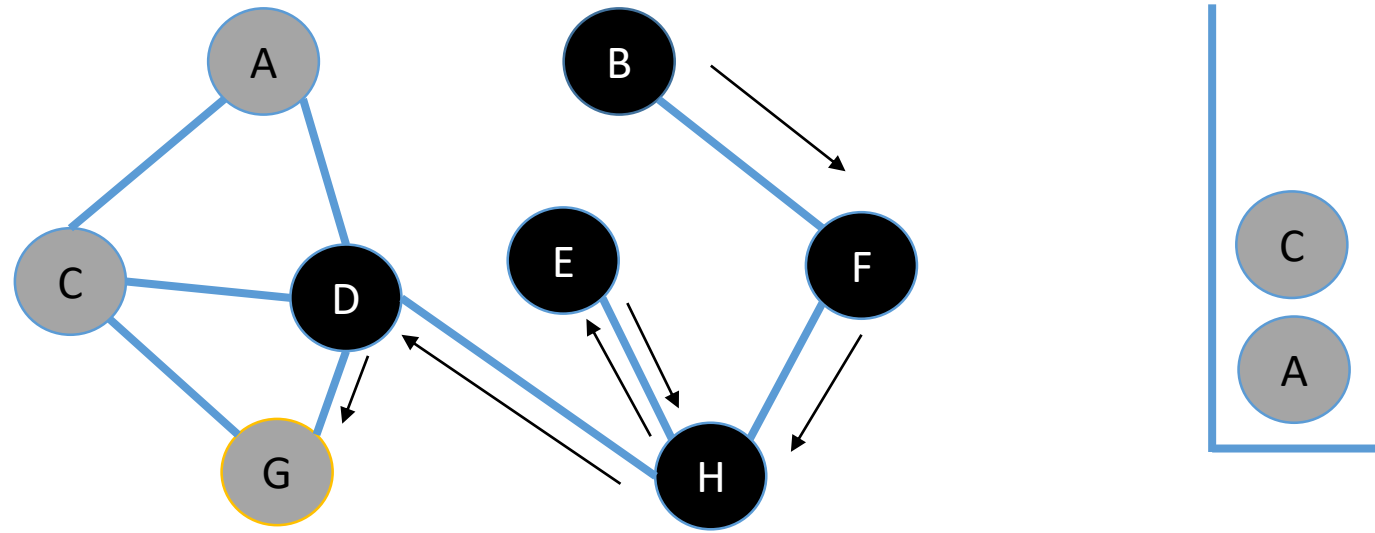
# Depth-First Traversal

- Next node popped from the stack
- Order visited: B, F, H, E, D



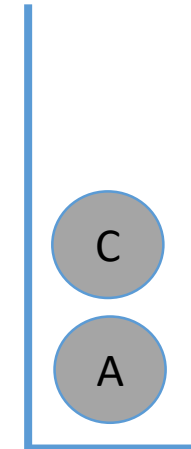
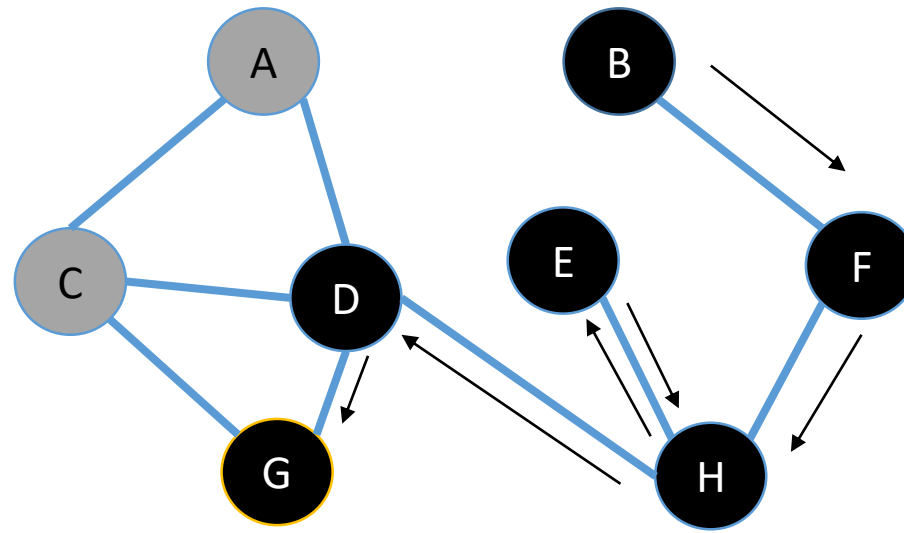
# Depth-First Traversal

- Add its adjacent white nodes to the stack (there are none)
- Order visited: B, F, H, E, D



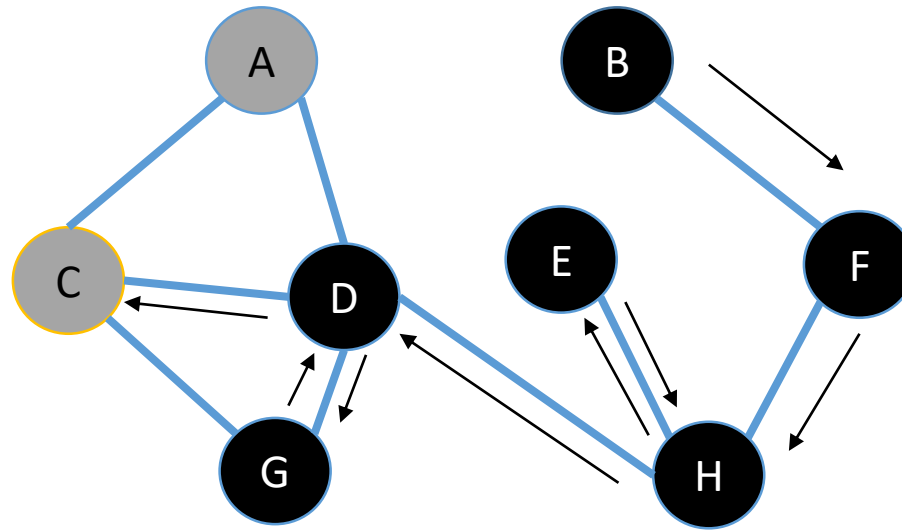
# Depth-First Traversal

- G has been visited
- Order visited: B, F, H, E, D, G



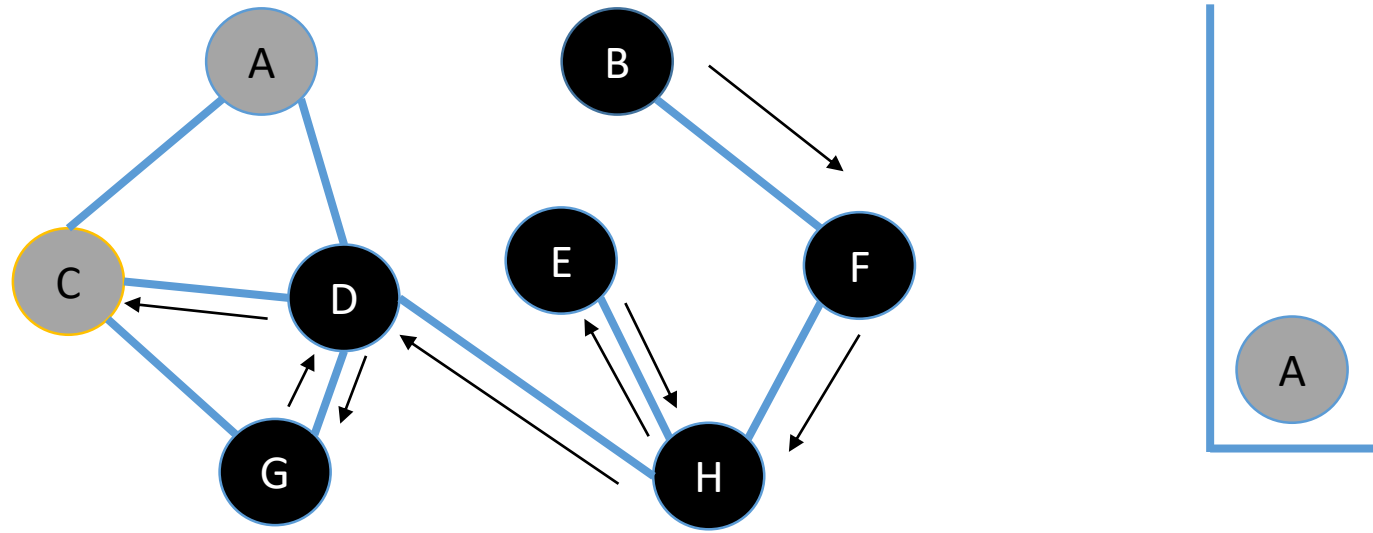
# Depth-First Traversal

- Next node popped from the stack
- Order visited: B, F, H, E, D, G



# Depth-First Traversal

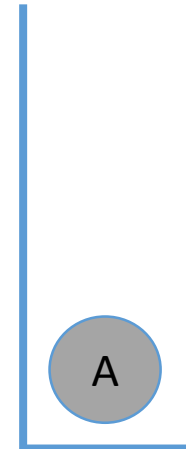
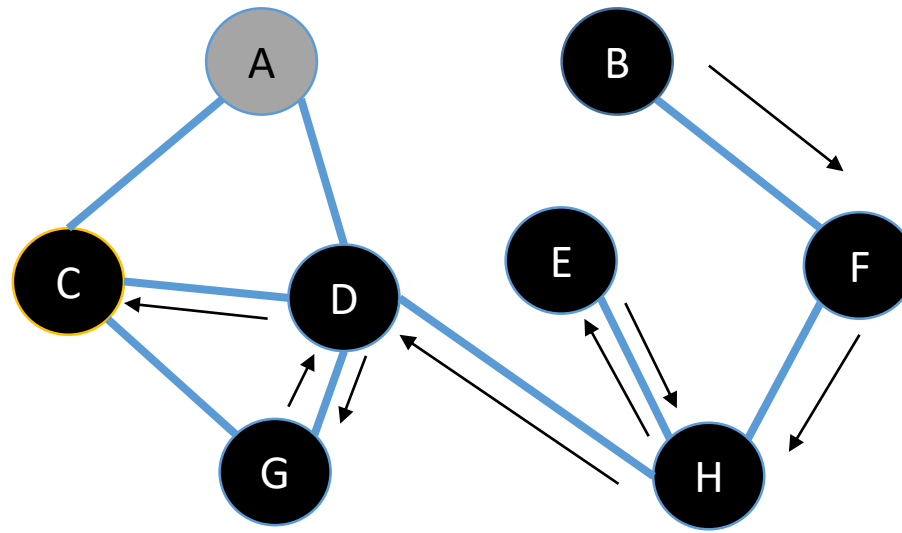
- Add its adjacent white nodes to the stack (there are none)
- Order visited: B, F, H, E, D, G





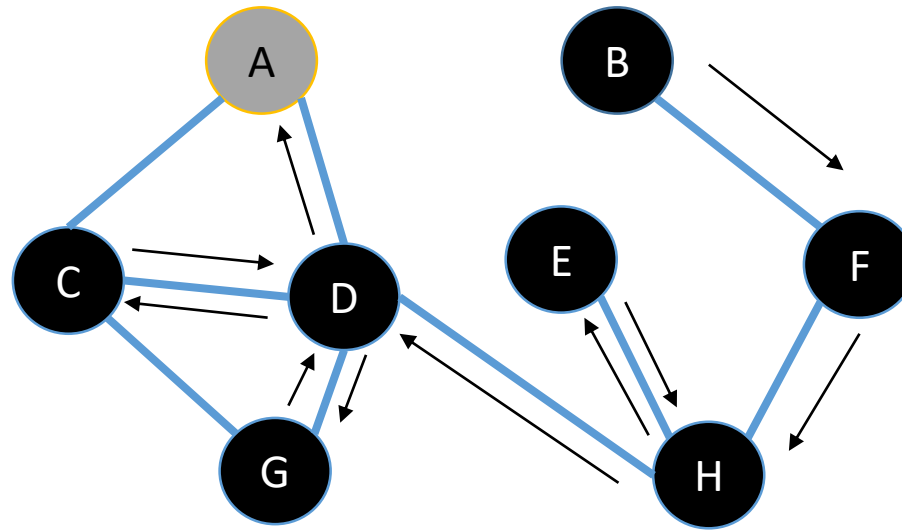
# Depth-First Traversal

- C has been visited
- Order visited: B, F, H, E, D, G, C



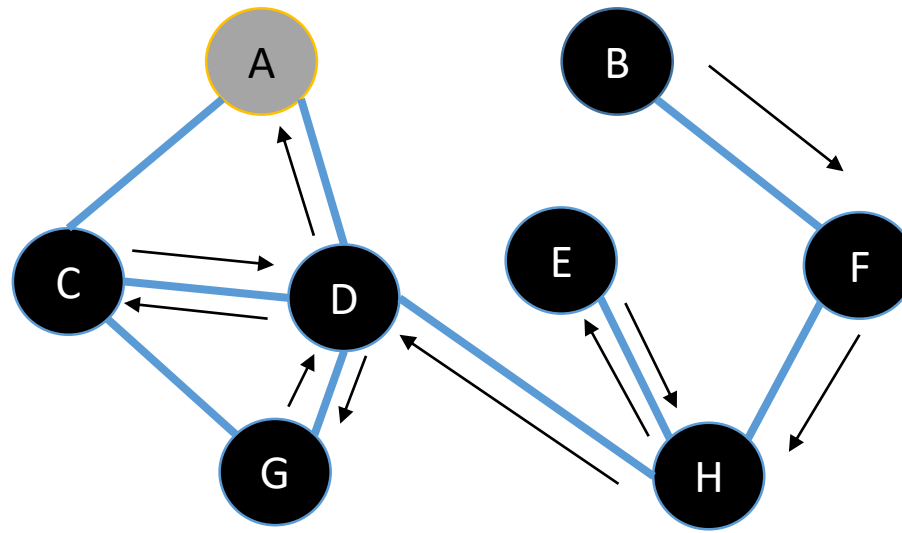
# Depth-First Traversal

- Next node is popped from the stack
- Order visited: B, F, H, E, D, G, C



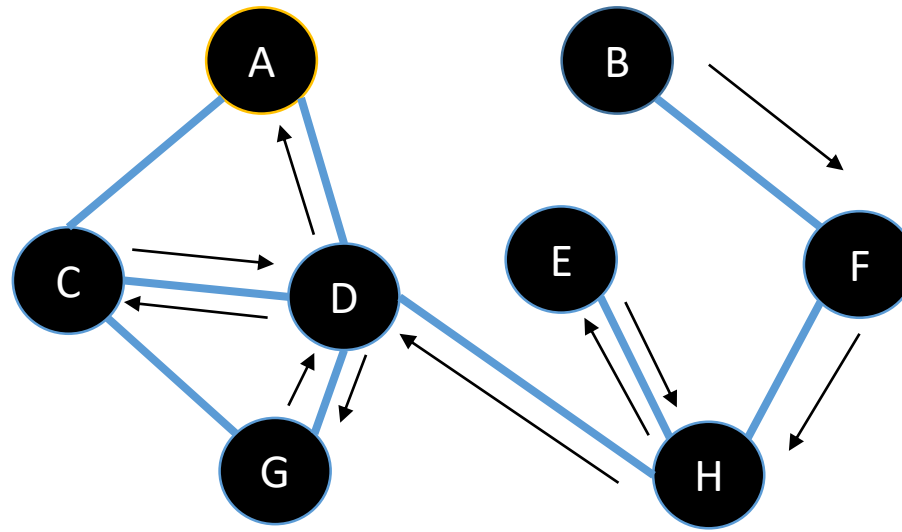
# Depth-First Traversal

- Add its adjacent white nodes to the stack (there are none)
- Order visited: B, F, H, E, D, G, C



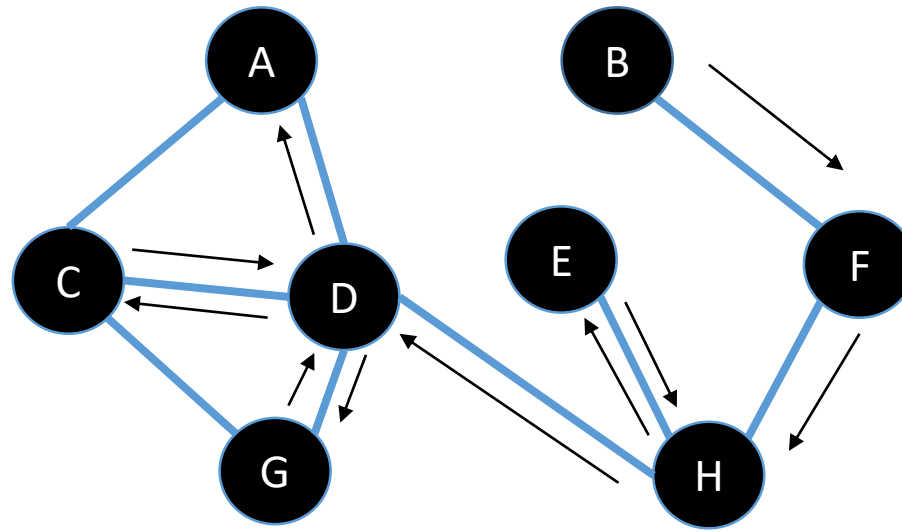
# Depth-First Traversal

- A has been visited
- Order visited: B, F, H, E, D, G, C, A



# Depth-First Traversal

- Next node is popped from the stack (there are none)
  - Traversal is complete when the stack is empty
- Order visited: **B, F, H, E, D, G, C, A**

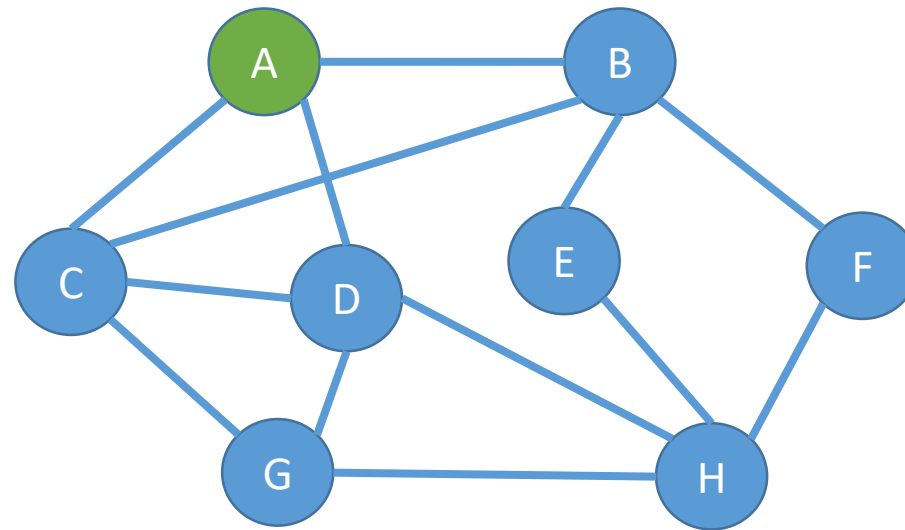


# Breadth-First Traversal

- Breadth-first traversal of a graph is similar to the breadth-first traversal of a tree
- Uses a queue to manage the next nodes to visit
- Maintains an array of “seen” nodes to prevent getting stuck in a cycle, if one existed
  - The traversal process is the same for bi-directional graphs or digraphs

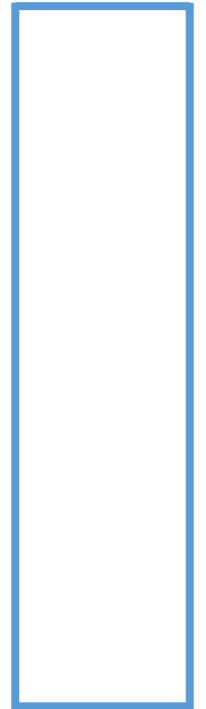
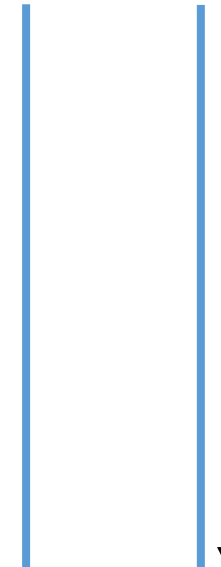
# Breadth-First Traversal

- Starting at A
- Order visited: N/A



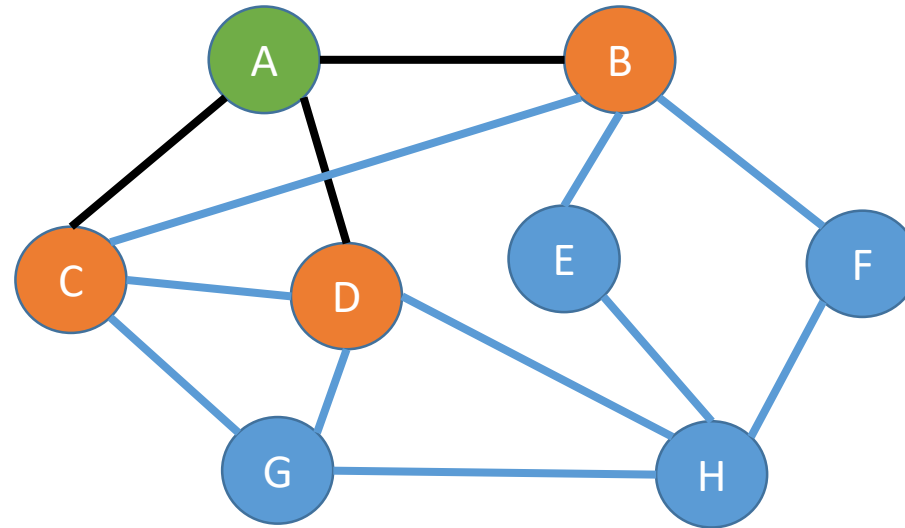
Queue

Nodes Seen

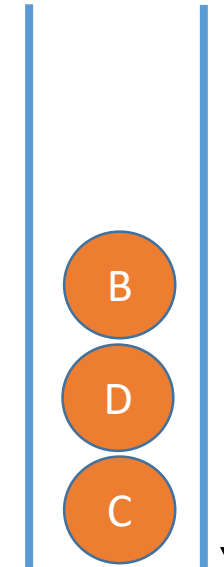


# Breadth-First Traversal

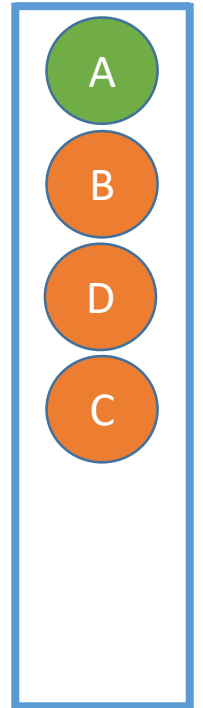
- A's unseen neighbors added to queue
  - A and unseen neighbors are added to nodes seen
- Order visited: A



Queue



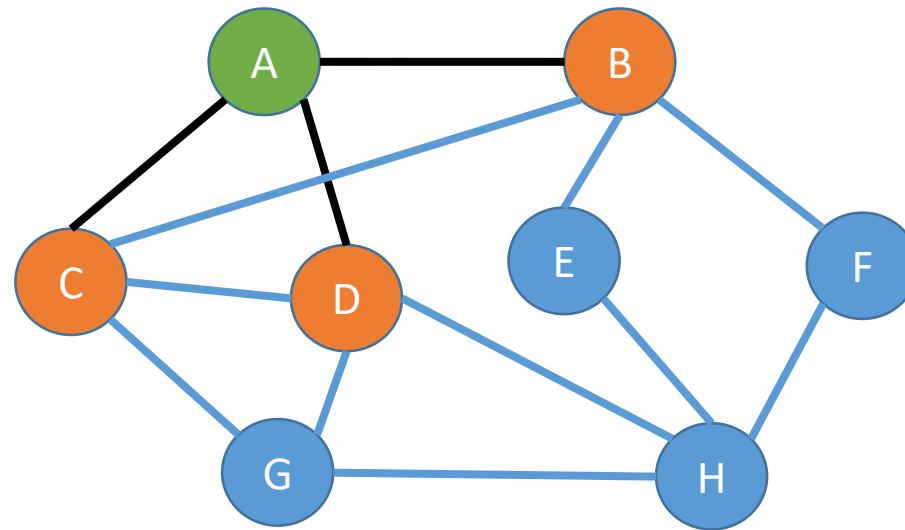
Nodes Seen



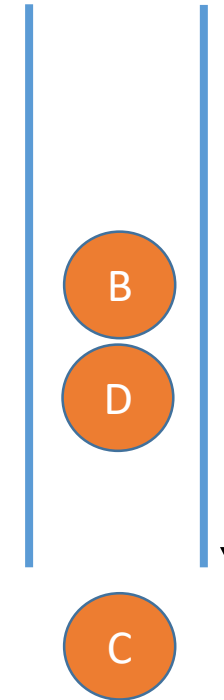


# Breadth-First Traversal

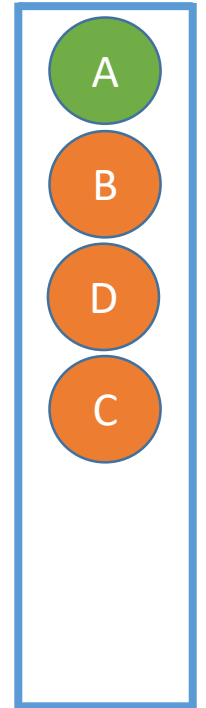
- C popped from the queue
- Order visited: A



Queue

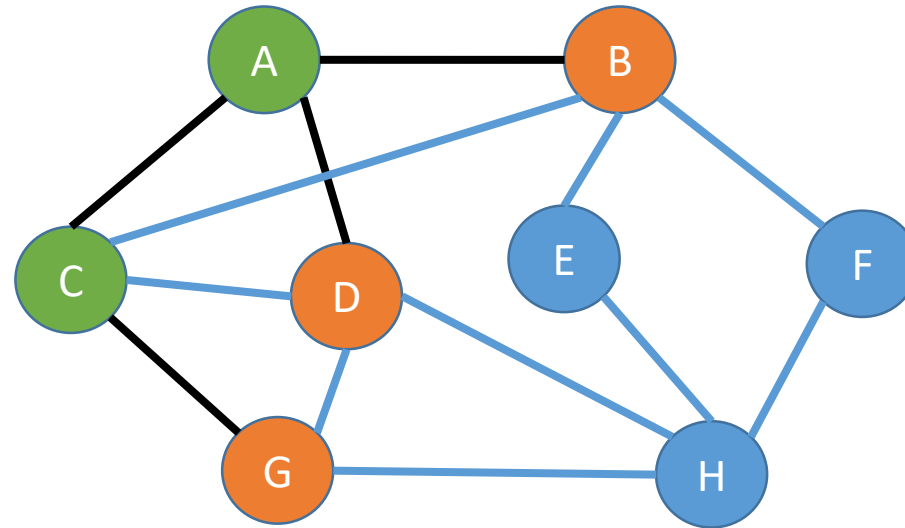


Nodes Seen

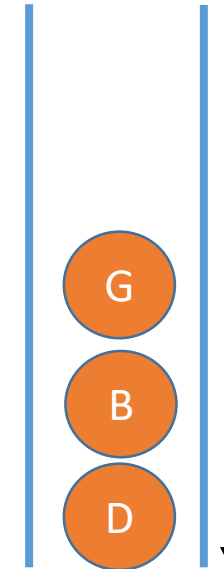


# Breadth-First Traversal

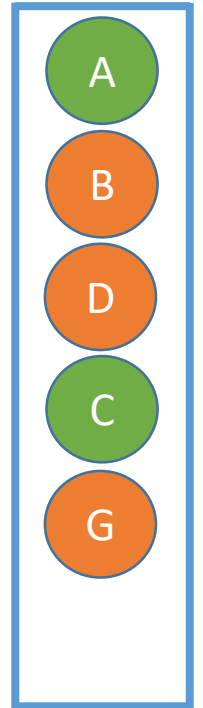
- C's unseen neighbors added to queue
  - Unseen neighbors are added to nodes seen
- Order visited: A, C



Queue

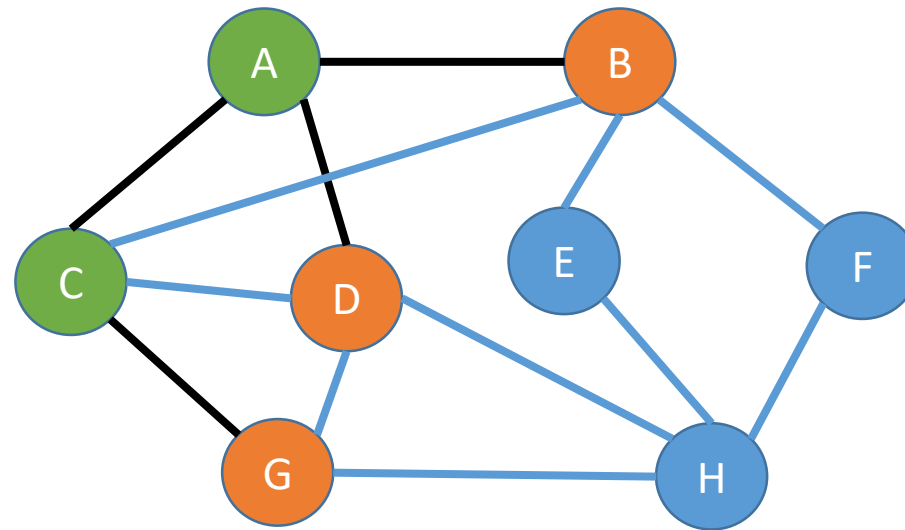


Nodes Seen

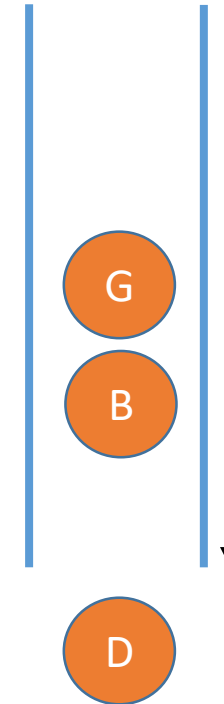


# Breadth-First Traversal

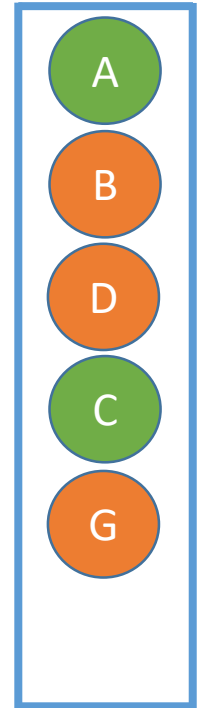
- D popped from the queue
- Order visited: A, C



Queue

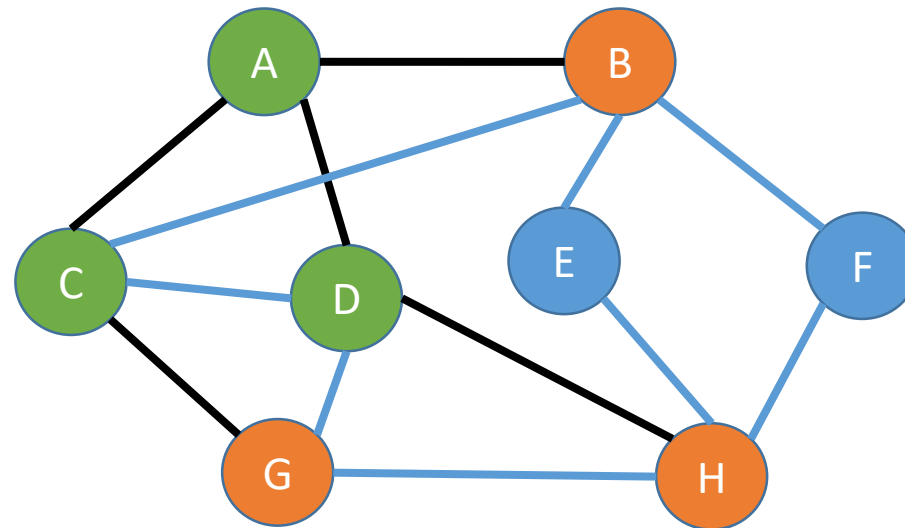


Nodes Seen

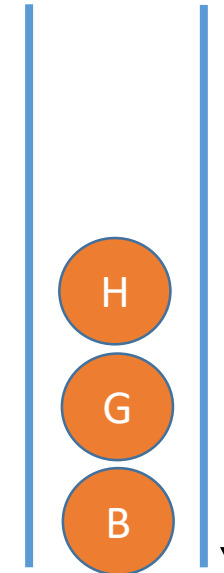


# Breadth-First Traversal

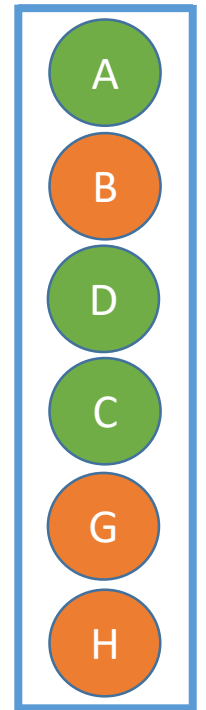
- D's unseen neighbors added to queue
  - Unseen neighbors are added to nodes seen
- Order visited: A, C, D



Queue

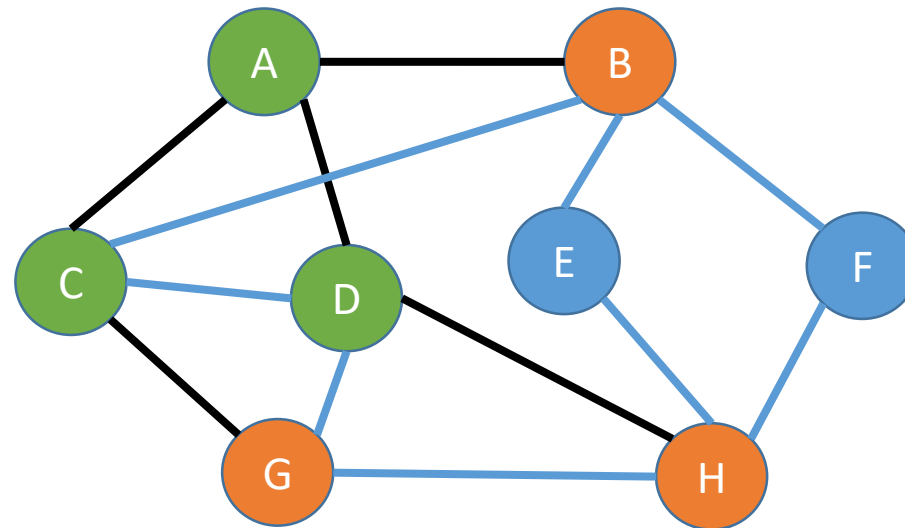


Nodes Seen

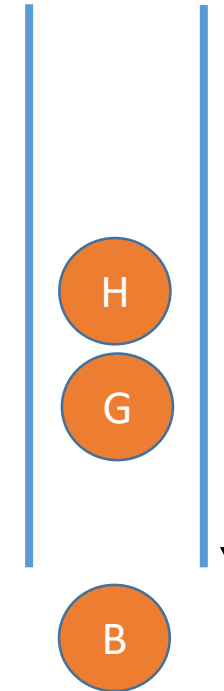


# Breadth-First Traversal

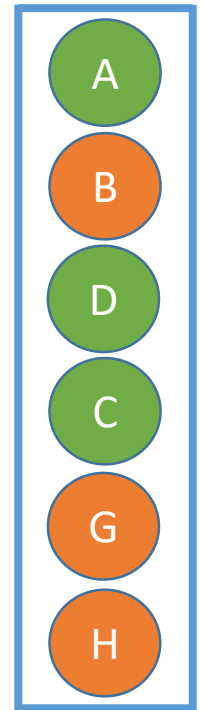
- B popped from the queue
- Order visited: A, C, D



Queue

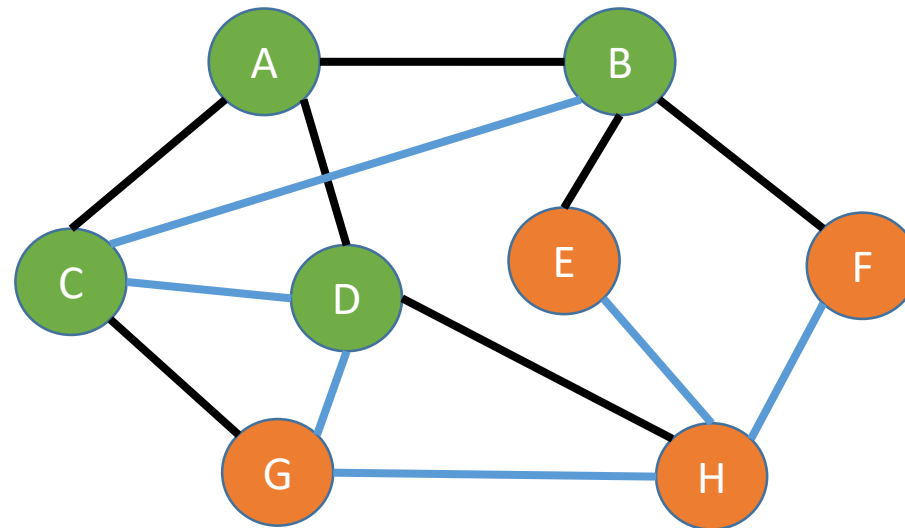


Nodes Seen

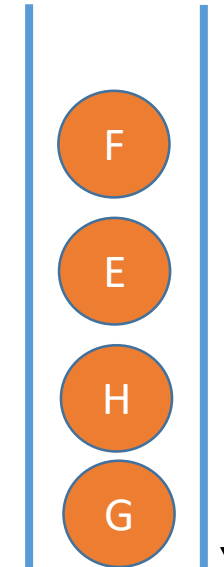


# Breadth-First Traversal

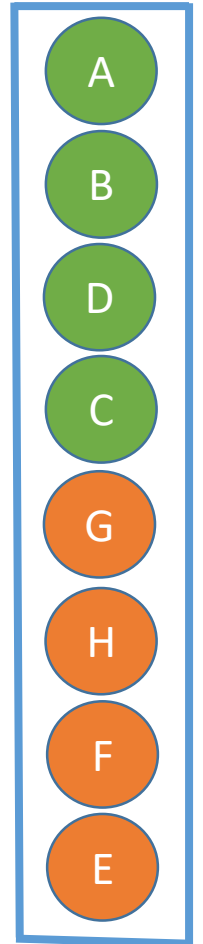
- B's unseen neighbors added to queue
  - Unseen neighbors are added to nodes seen
- Order visited: A, C, D, B



Queue

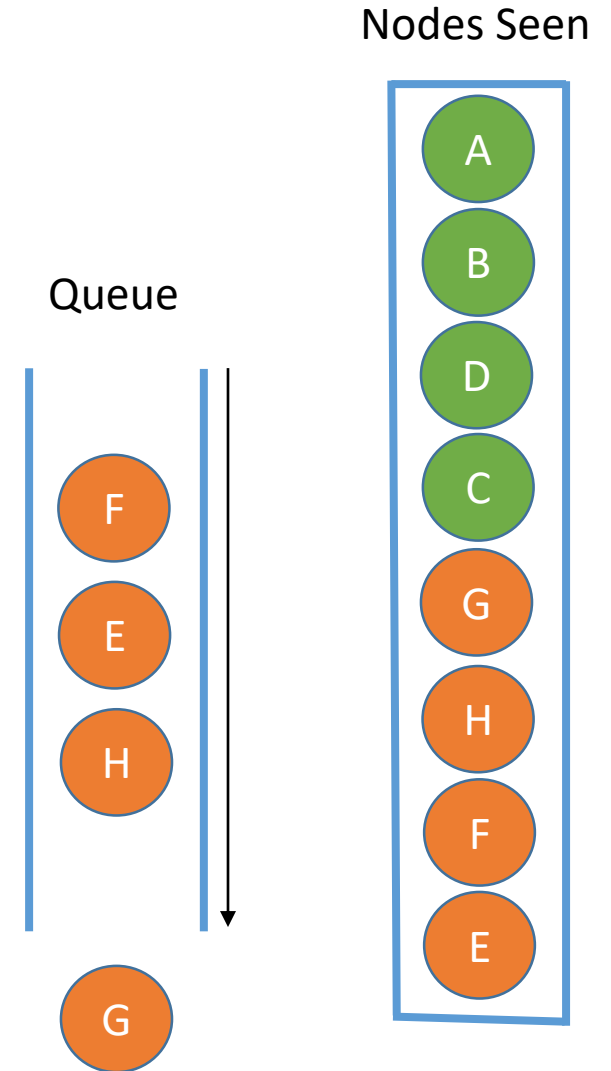
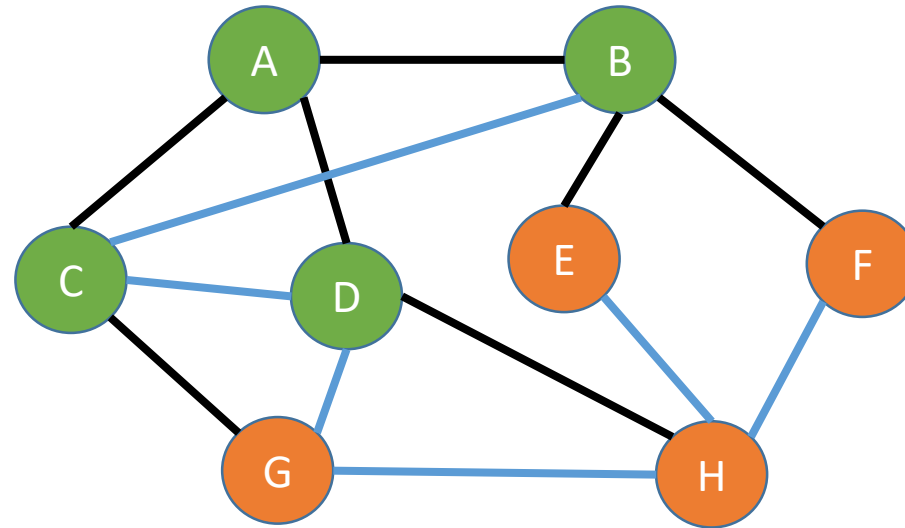


Nodes Seen



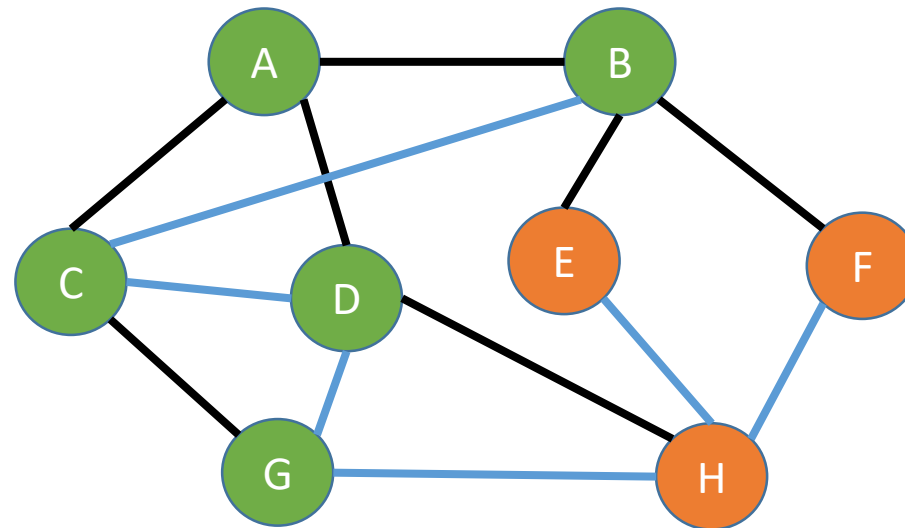
# Breadth-First Traversal

- G popped from the queue
- Order visited: A, C, D, B

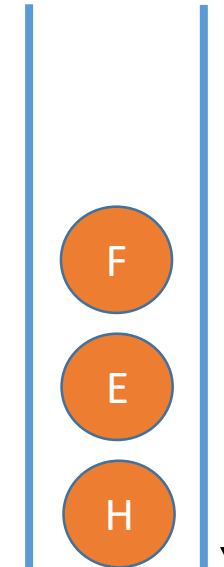


# Breadth-First Traversal

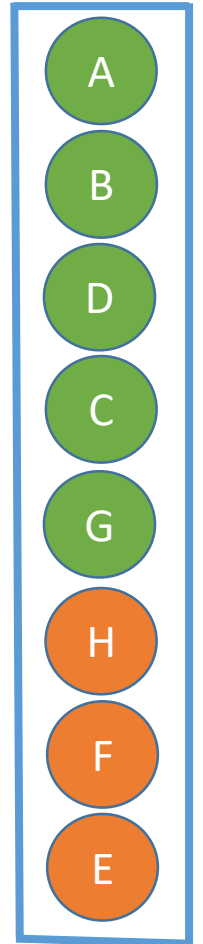
- G's unseen neighbors added to queue (none)
  - Unseen neighbors are added to nodes seen
- Order visited: A, C, D, B, G



Queue



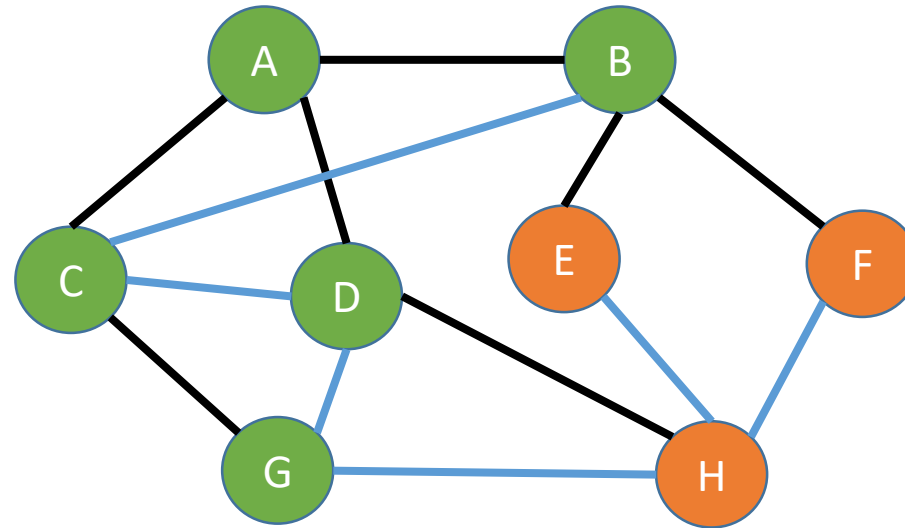
Nodes Seen



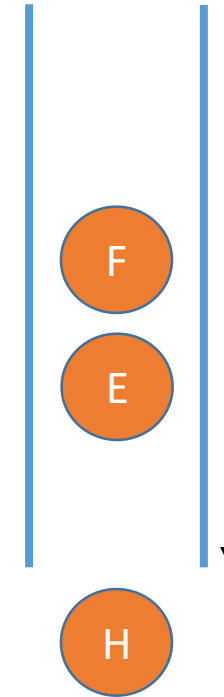


# Breadth-First Traversal

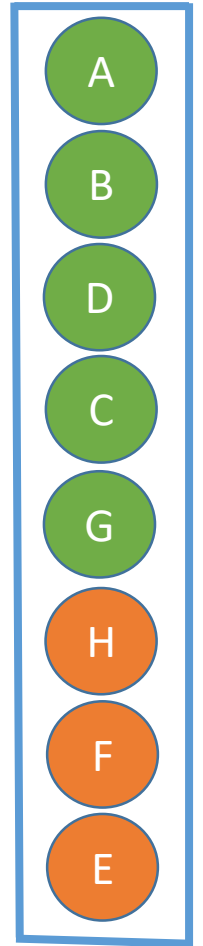
- H popped from the queue
- Order visited: A, C, D, B, G



Queue

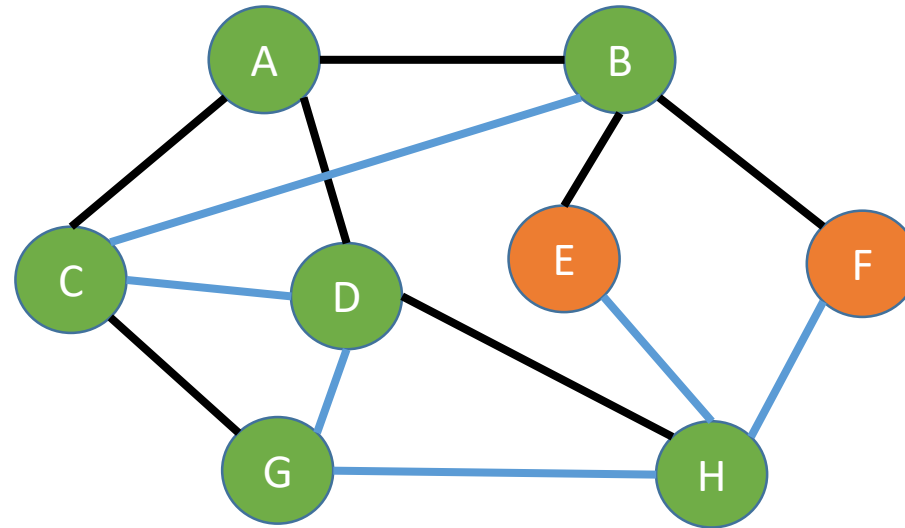


Nodes Seen

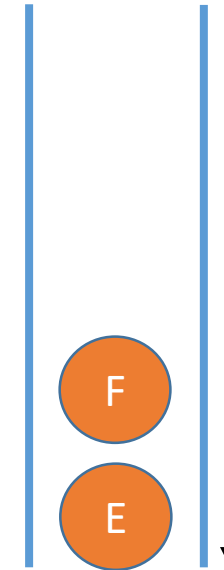


# Breadth-First Traversal

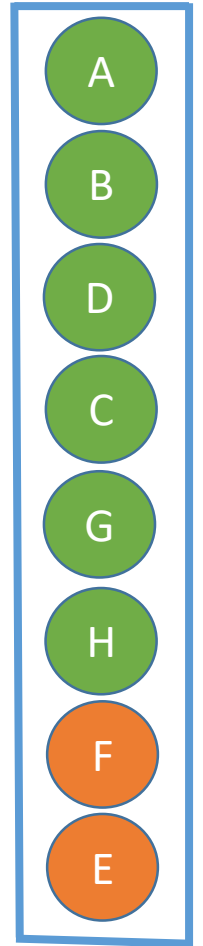
- H's unseen neighbors added to queue (none)
  - Unseen neighbors are added to nodes seen
- Order visited: A, C, D, B, G, H



Queue

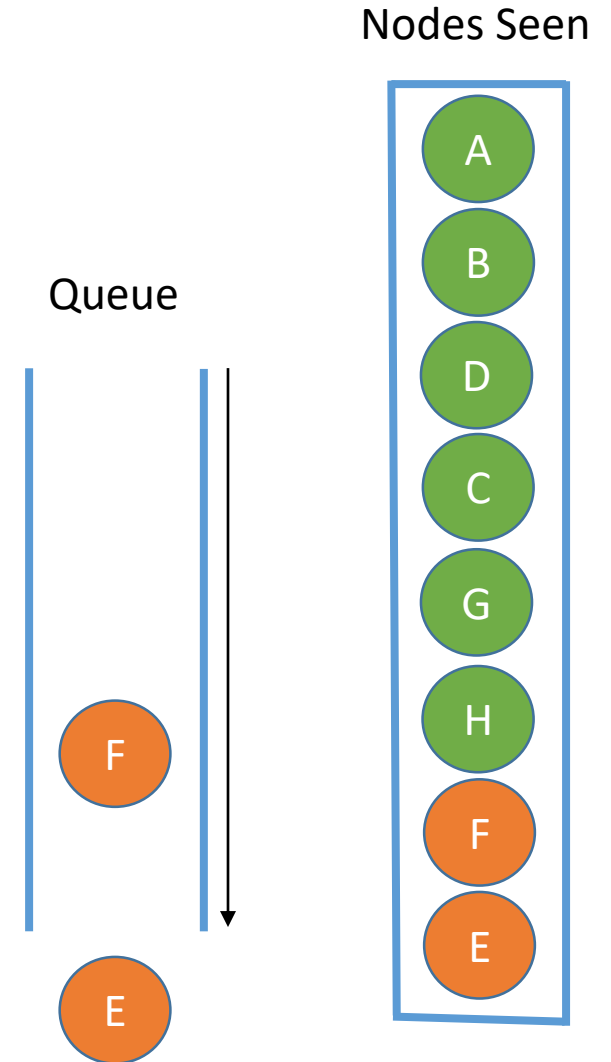
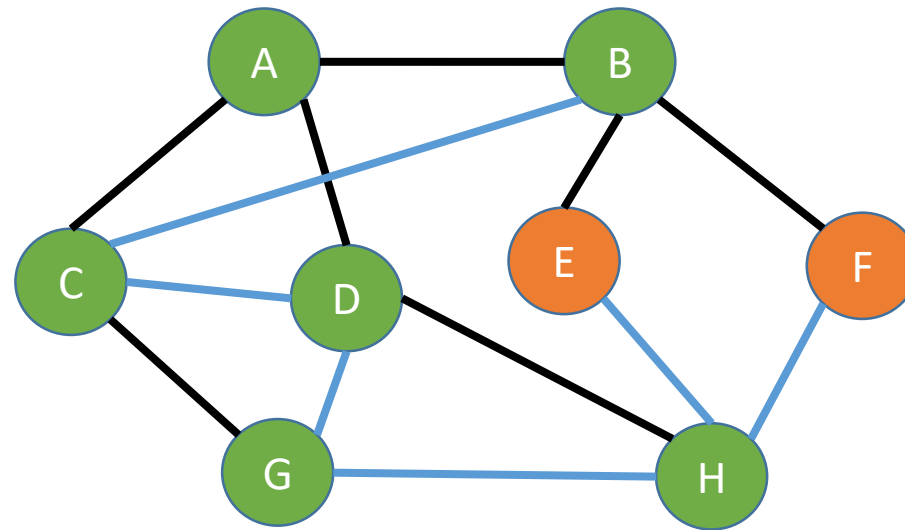


Nodes Seen



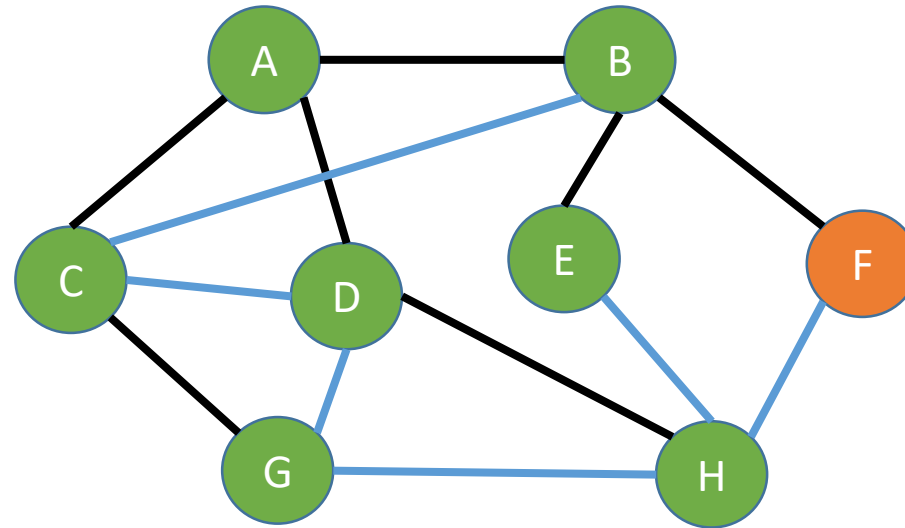
# Breadth-First Traversal

- E popped from the queue
- Order visited: A, C, D, B, G, H

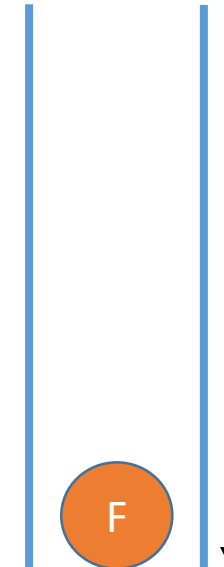


# Breadth-First Traversal

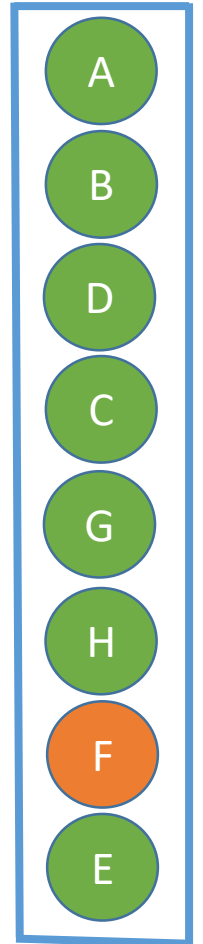
- E's unseen neighbors added to queue (none)
  - Unseen neighbors are added to nodes seen
- Order visited: A, C, D, B, G, H, E



Queue

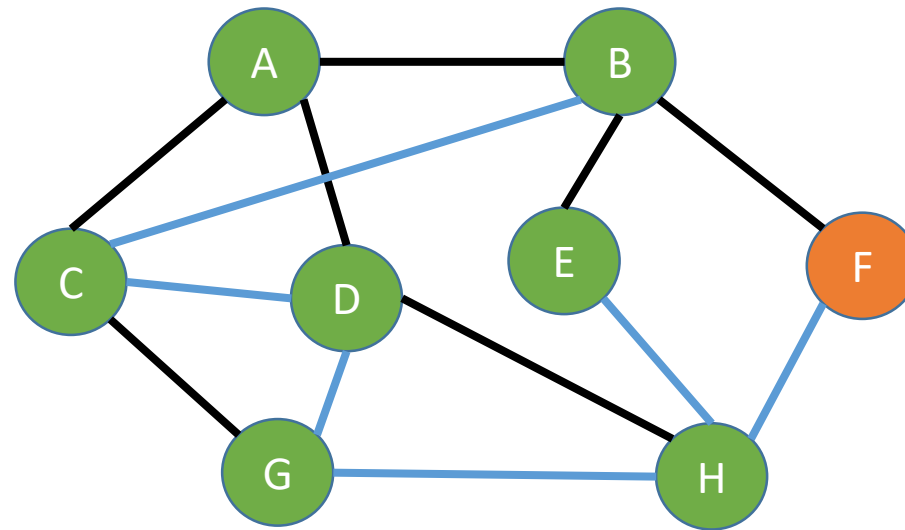


Nodes Seen

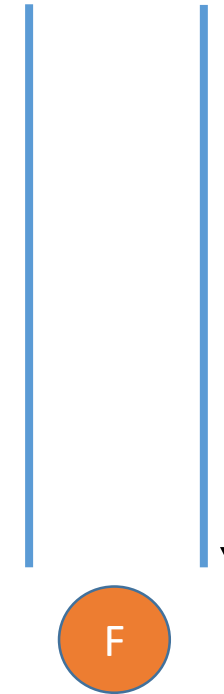


# Breadth-First Traversal

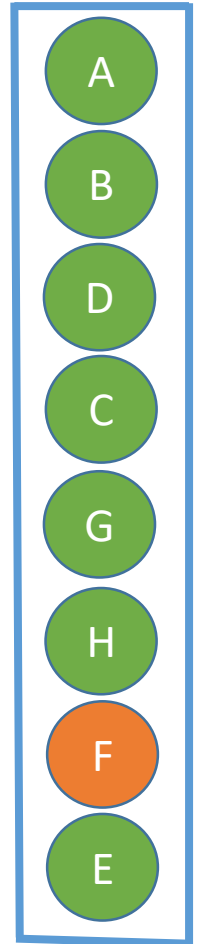
- F popped from queue
- Order visited: A, C, D, B, G, H, E



Queue

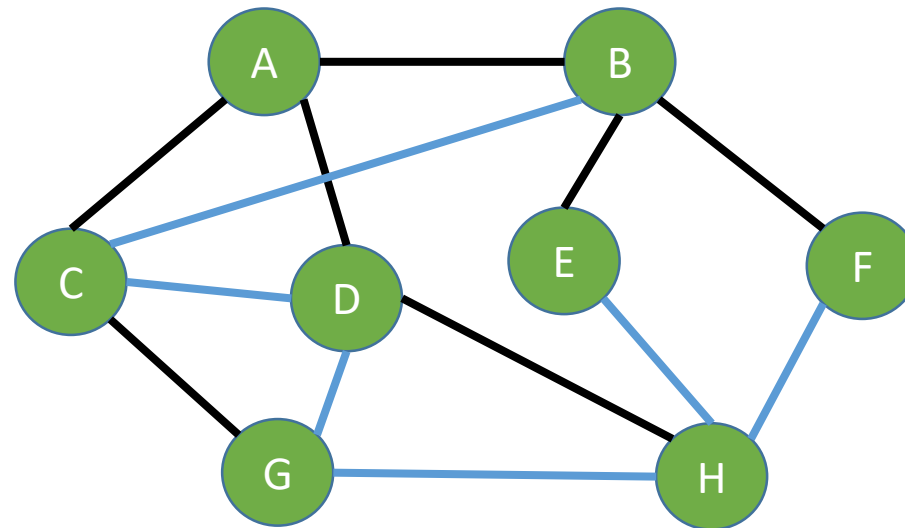


Nodes Seen

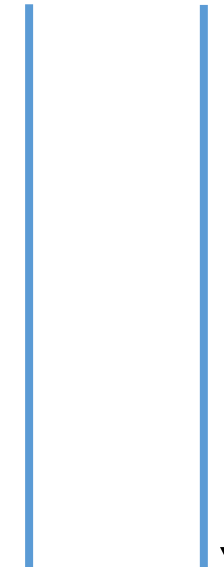


# Breadth-First Traversal

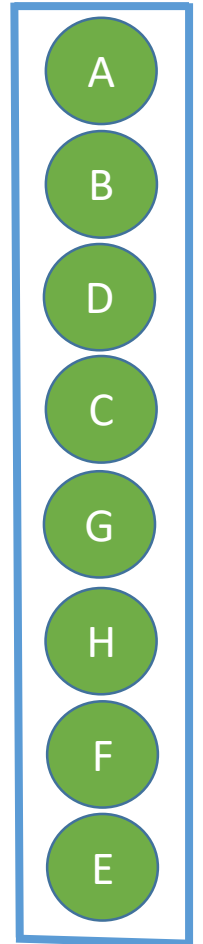
- F's unseen neighbors added to queue (none)
  - Unseen neighbors are added to nodes seen
- Order visited: A, C, D, B, G, H, E, F



Queue

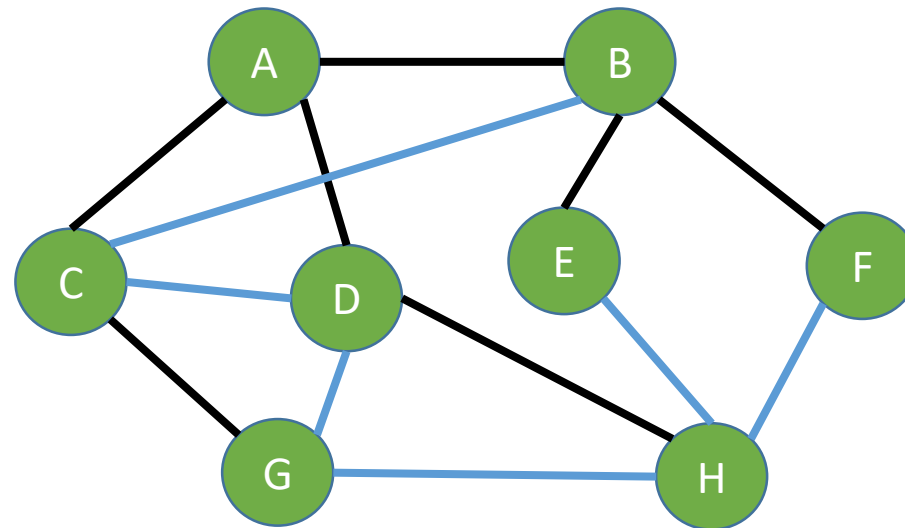


Nodes Seen



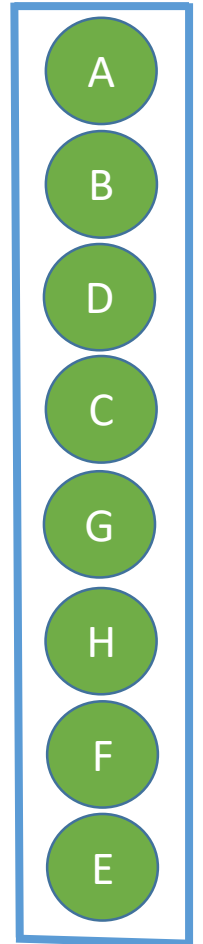
# Breadth-First Traversal

- Queue is empty
  - Traversal is complete
- Order visited: **A, C, D, B, G, H, E, F**



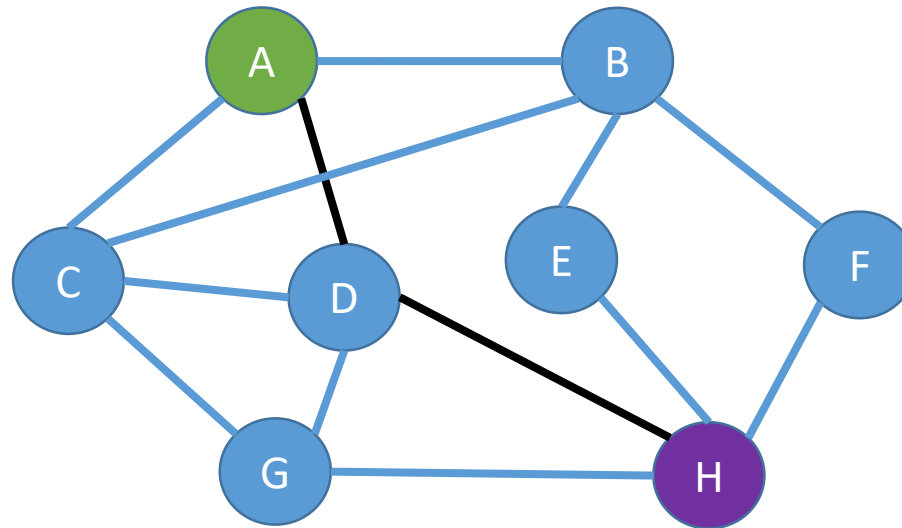
Queue

Nodes Seen



# Finding Distance with Breadth-First

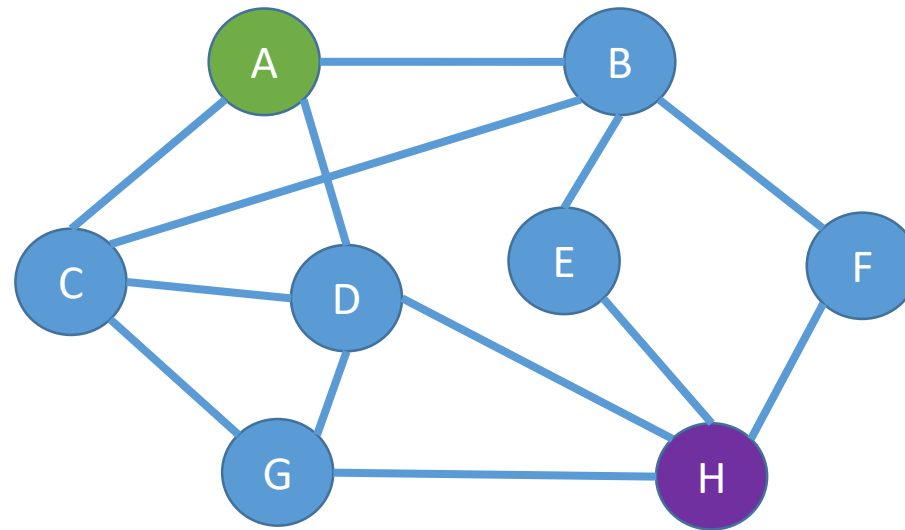
- Based on the last example, its easy to see a breadth-first traversal can be used to find the distance (shortest path) between two nodes
- As an example, we'll use a breadth-first traversal to find the distance between A and H





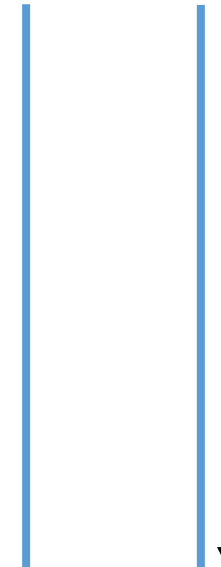
# Finding Distance with Breadth-First

- Starting at A
- Order visited: N/A



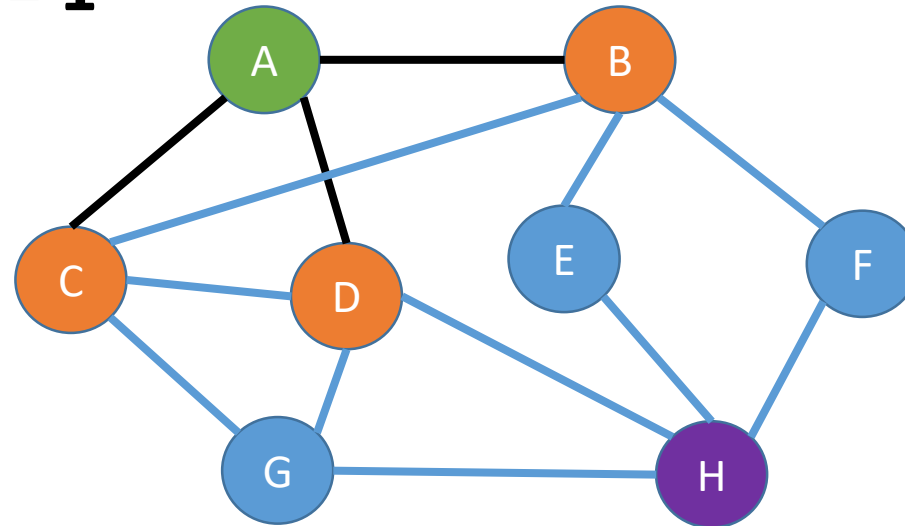
Queue

Nodes Seen

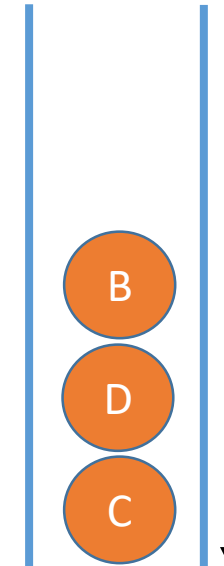


# Finding Distance with Breadth-First

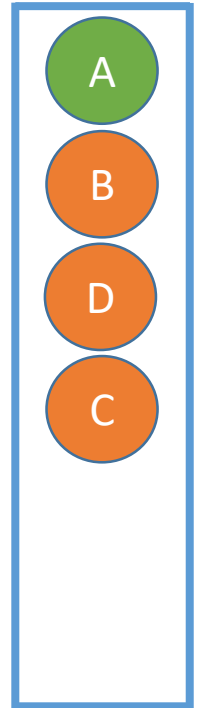
- A's unseen neighbors added to queue
  - A and unseen neighbors are added to nodes seen
- Order visited: A
- **DEEPEST LEVEL = 1**



Queue

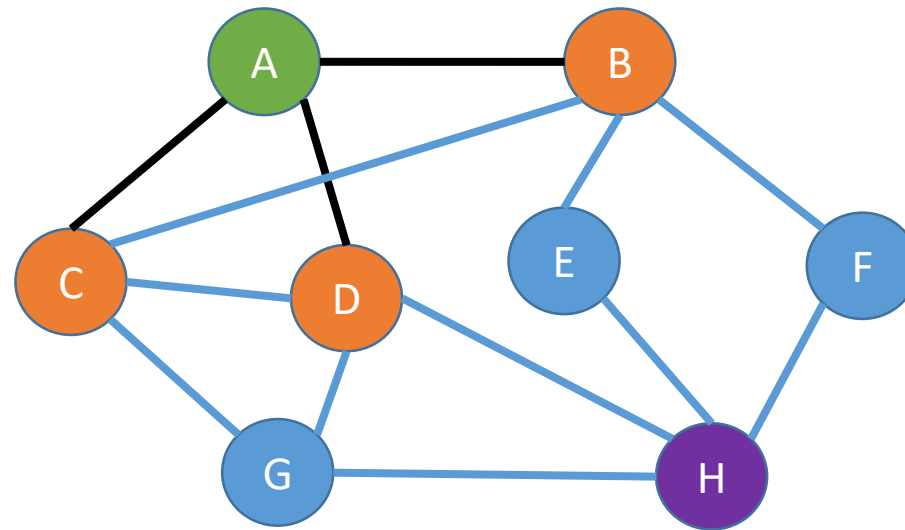


Nodes Seen

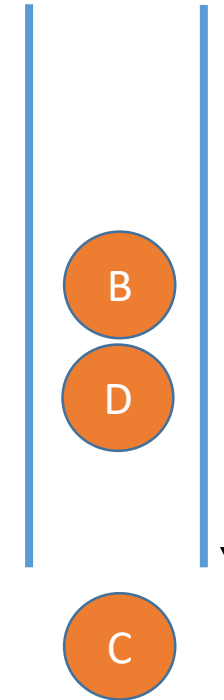


# Finding Distance with Breadth-First

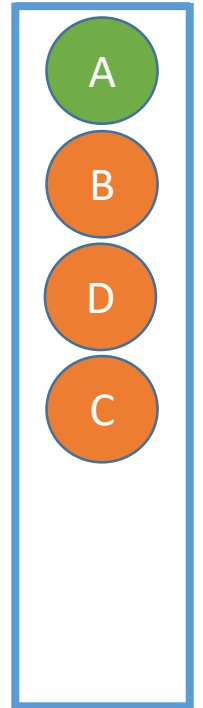
- C popped from the queue
- Order visited: A
- **DEEPEST LEVEL = 1**



Queue

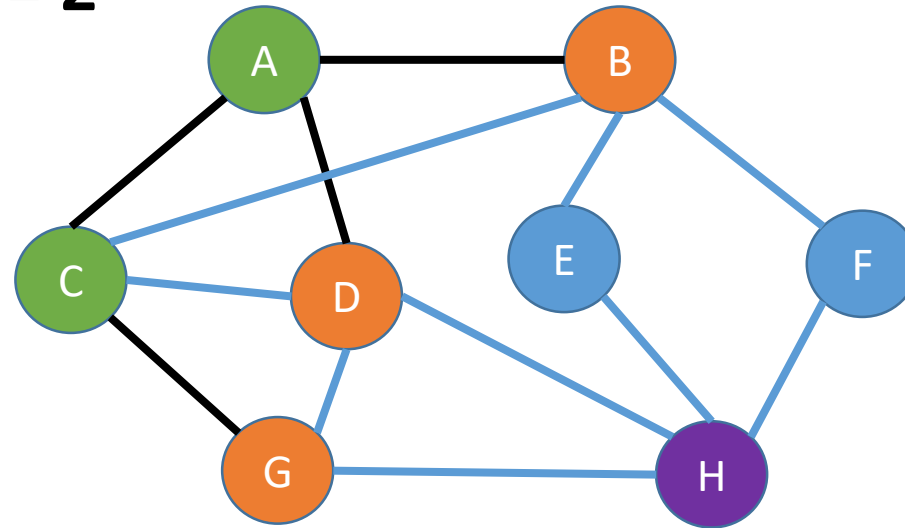


Nodes Seen

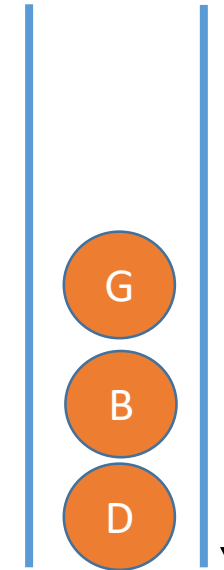


# Finding Distance with Breadth-First

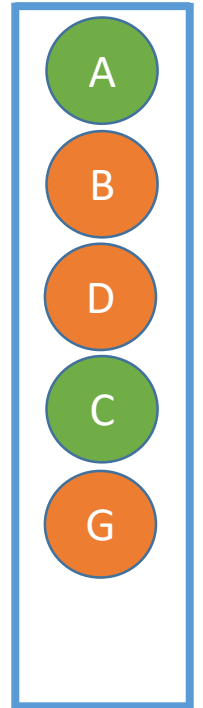
- C's unseen neighbors added to queue
  - Unseen neighbors are added to nodes seen
- Order visited: A, C
- **DEEPEST LEVEL = 2**



Queue

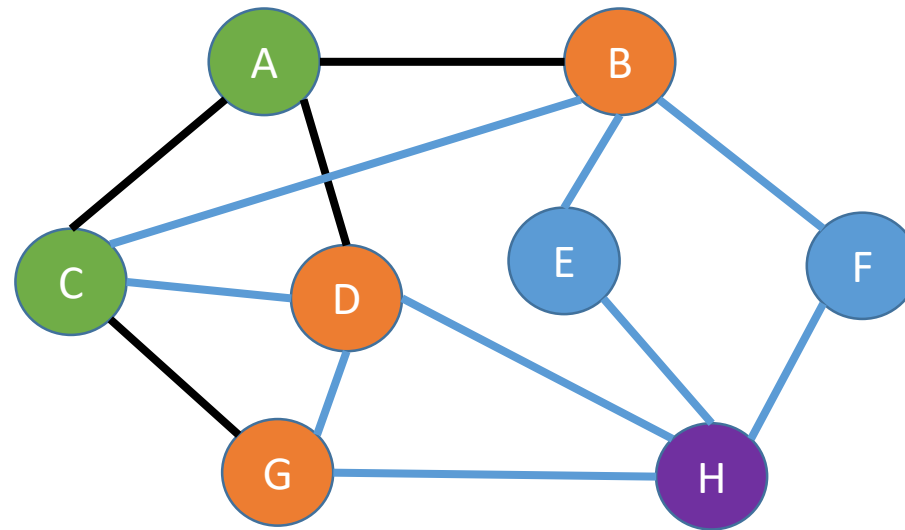


Nodes Seen

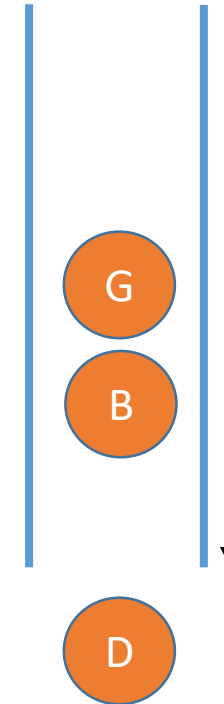


# Finding Distance with Breadth-First

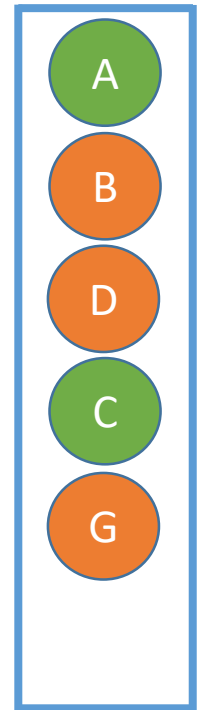
- D popped from the queue
- Order visited: A, C
- **DEEPEST LEVEL = 2**



Queue

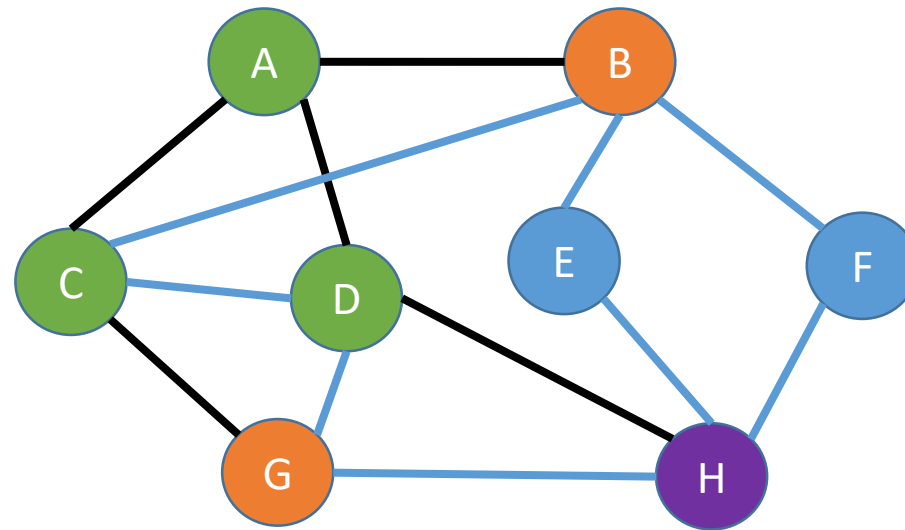


Nodes Seen

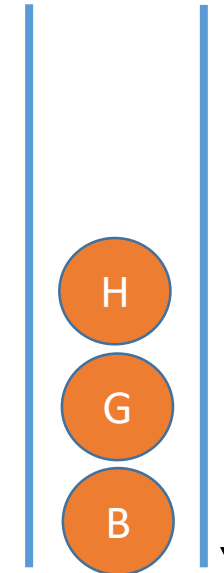


# Finding Distance with Breadth-First

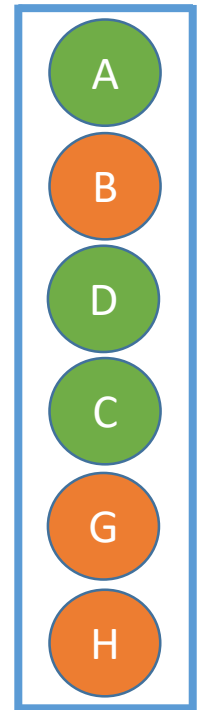
- D's unseen neighbors added to queue
  - H was reached
  - Traversal stops
- **DEEPEST LEVEL = 2**
  - DISTANCE = 2



Queue

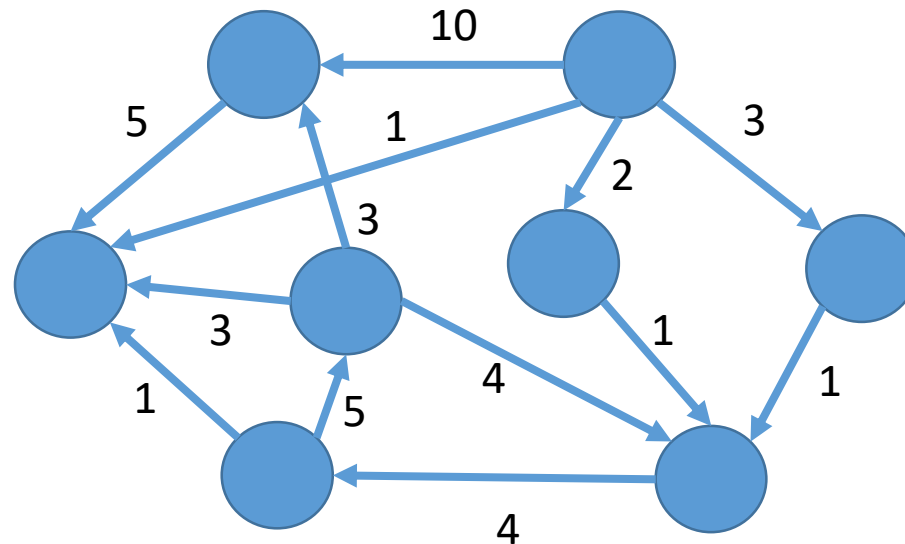


Nodes Seen



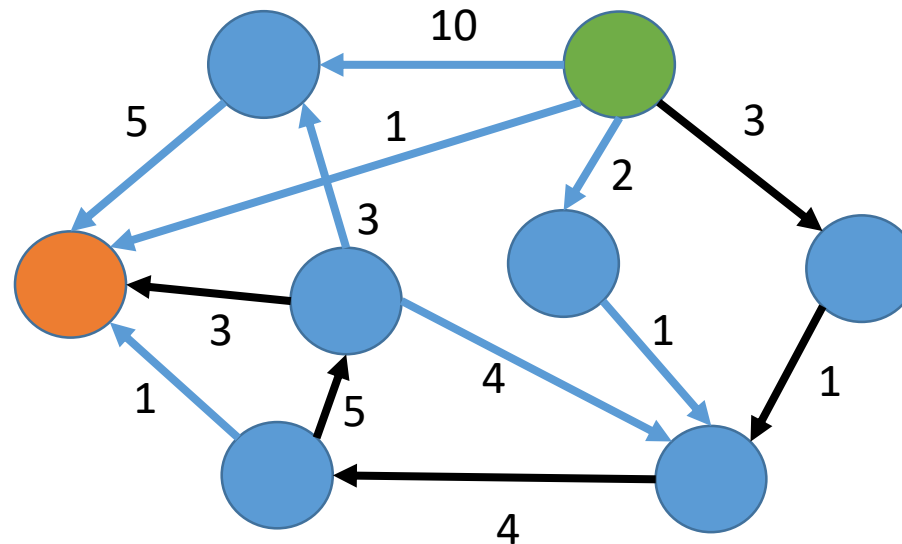
# Weighted Graphs

- A **weighted graph** is a graph where each edge has a weight or cost.
  - Weighted graphs can be undirected or directed



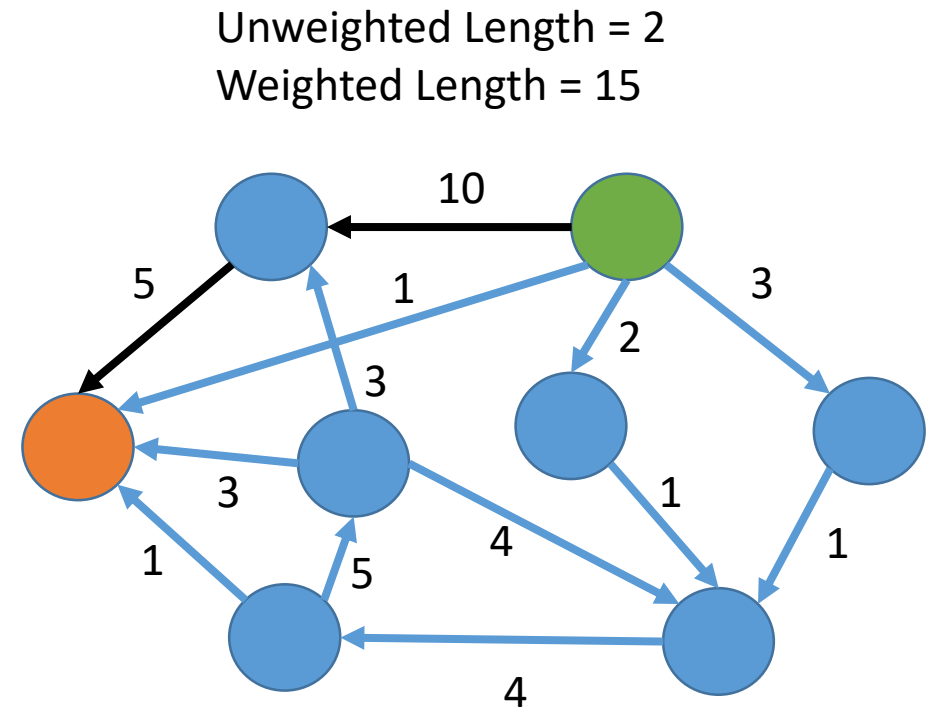
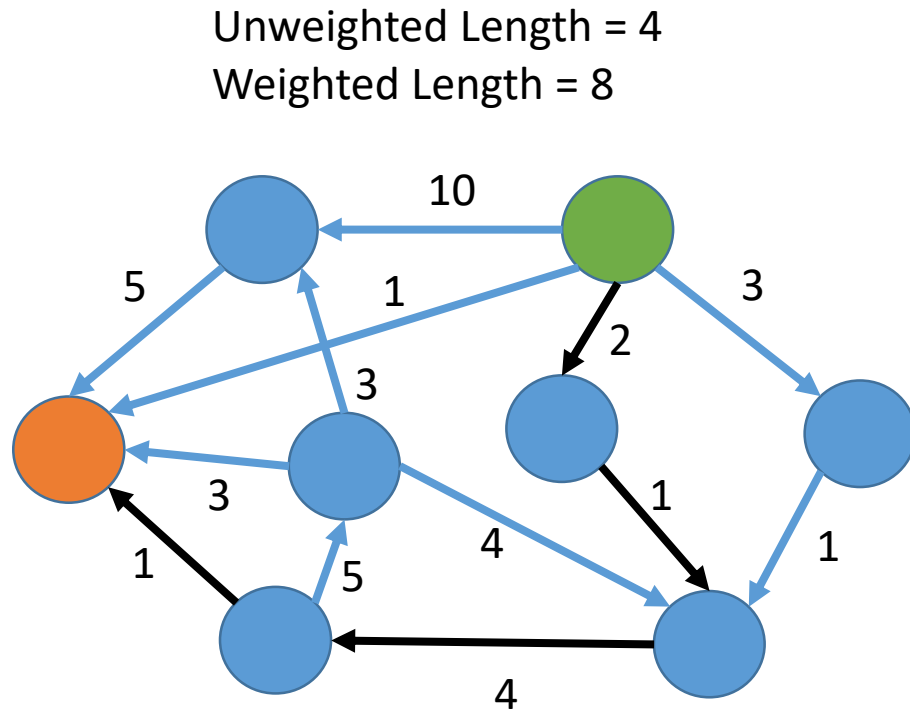
# Weighted Graphs

- The path length of a weighted graph is the sum of the edge costs.
  - $3+1+4+5+3 = \mathbf{16}$





# Weighted Graphs



- The first path is less costly than the second path, despite it being twice as long

# Weighted Graphs

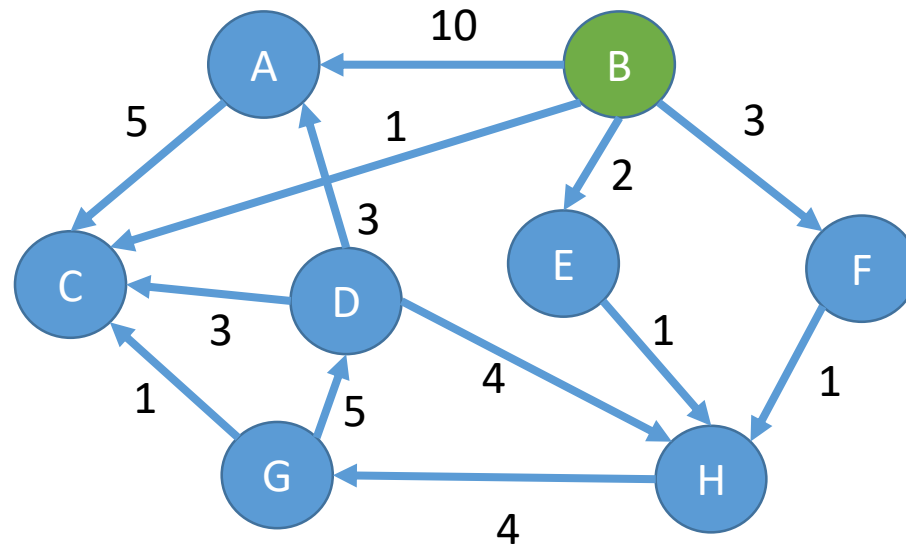
- A breadth-first traversal would not be useful for finding the path with the least cost
- Other algorithms are used to find the path with the least cost between two nodes
- The most well known is Dijkstra's Algorithm

# Dijkstra's Algorithm

- This algorithm finds the shortest path between one node and **every other node** in a graph.
- *For each vertex*
  - The algorithm determines the vertex's distance (shortest/least costly path) from the starting vertex
  - The algorithm determines the vertex's predecessor pointer- the previous vertex with the shortest (or least costly) path from the starting vertex
- Can be used on:
  - Bi-directional and digraphs
  - Weighted graphs and unweighted graphs

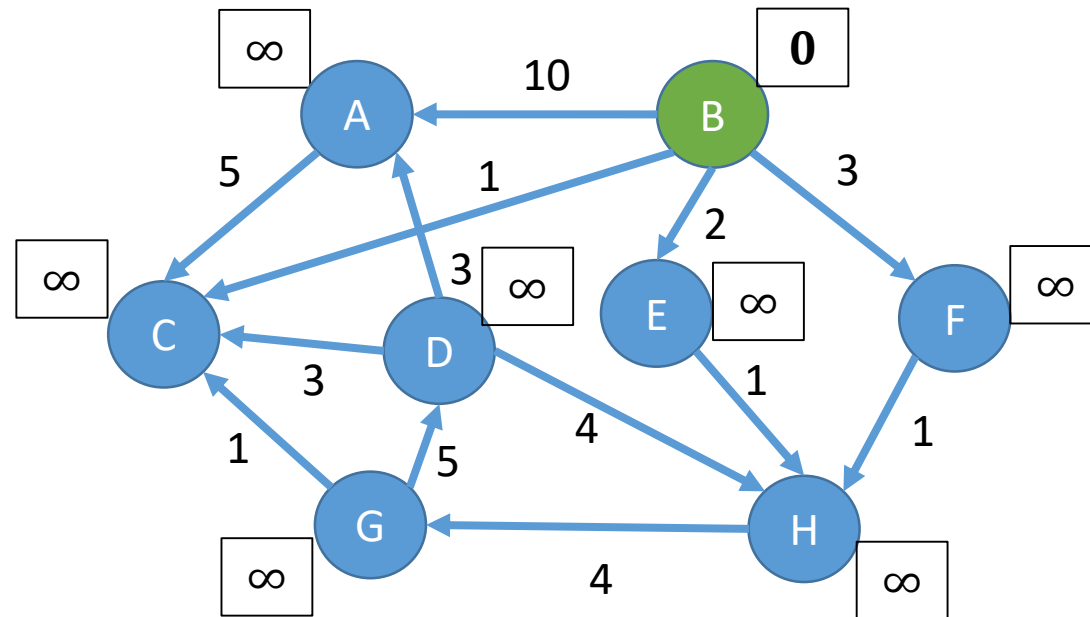
# Dijkstra's Algorithm

- We can start with any node, but we'll start with node B since a path exists from B to all other nodes



# Dijkstra's Algorithm

- We'll remember the cost from each node back to node B.
- B has a cost of 0; All other nodes are assumed to have a cost of infinity
  - Ensures the path found will be less than that

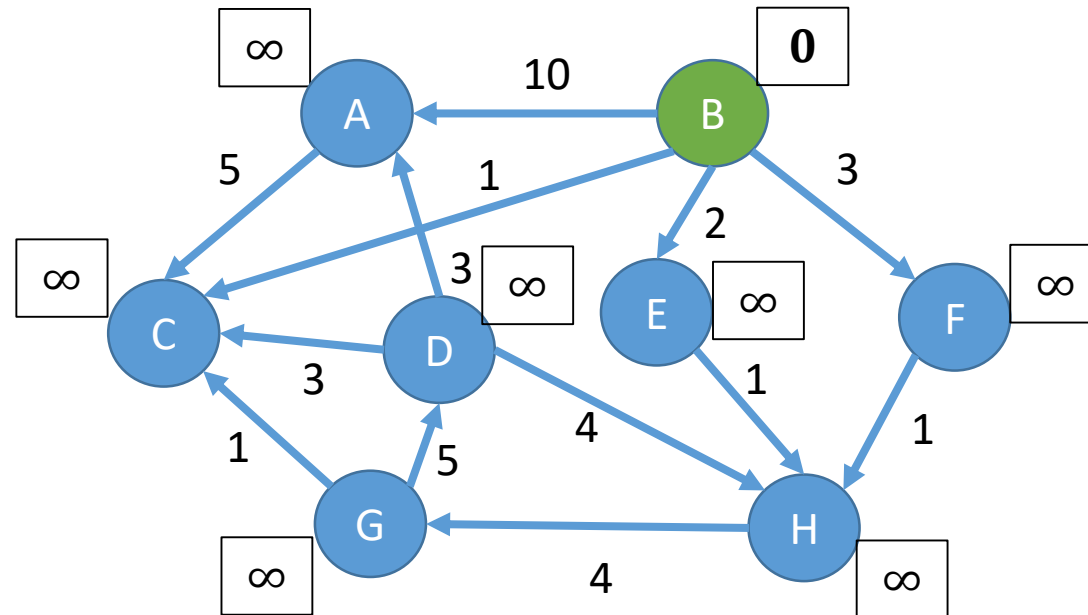
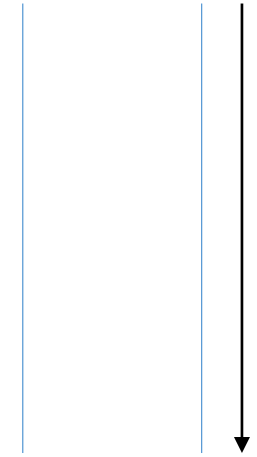


A cost =  $\infty$   
B cost = 0  
C cost =  $\infty$   
D cost =  $\infty$   
E cost =  $\infty$   
F cost =  $\infty$   
G cost =  $\infty$   
H cost =  $\infty$

# Dijkstra's Algorithm

- A priority queue or min-heap is used to prioritize nodes with lower costs to be visited first

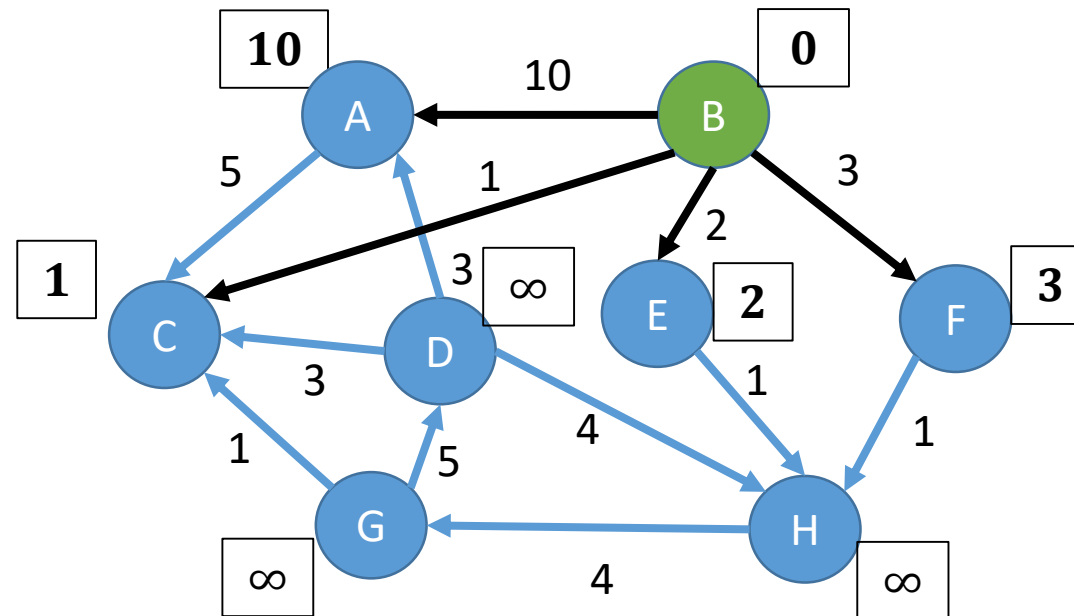
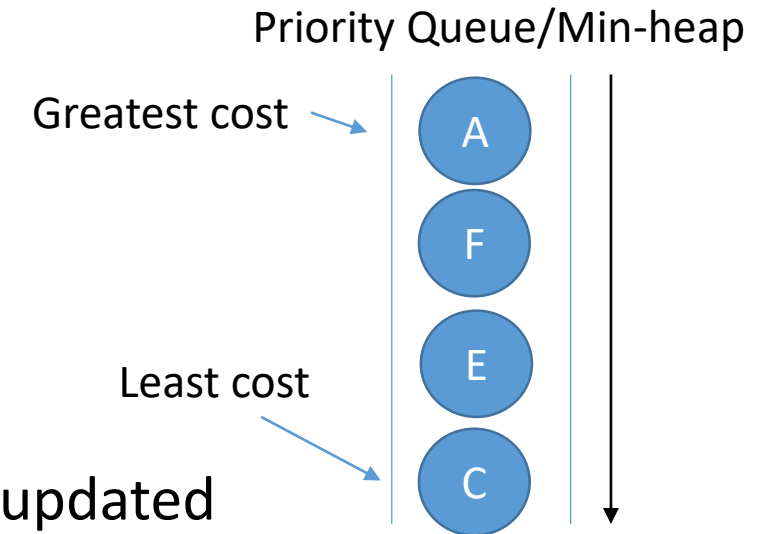
Priority Queue/Min-heap



A cost =  $\infty$   
B cost = 0  
C cost =  $\infty$   
D cost =  $\infty$   
E cost =  $\infty$   
F cost =  $\infty$   
G cost =  $\infty$   
H cost =  $\infty$

# Dijkstra's Algorithm

- We'll now look at the nodes adjacent to B.
- Cost of each is B's cost + cost of the edge
  - All of which are less than infinity so the costs are updated

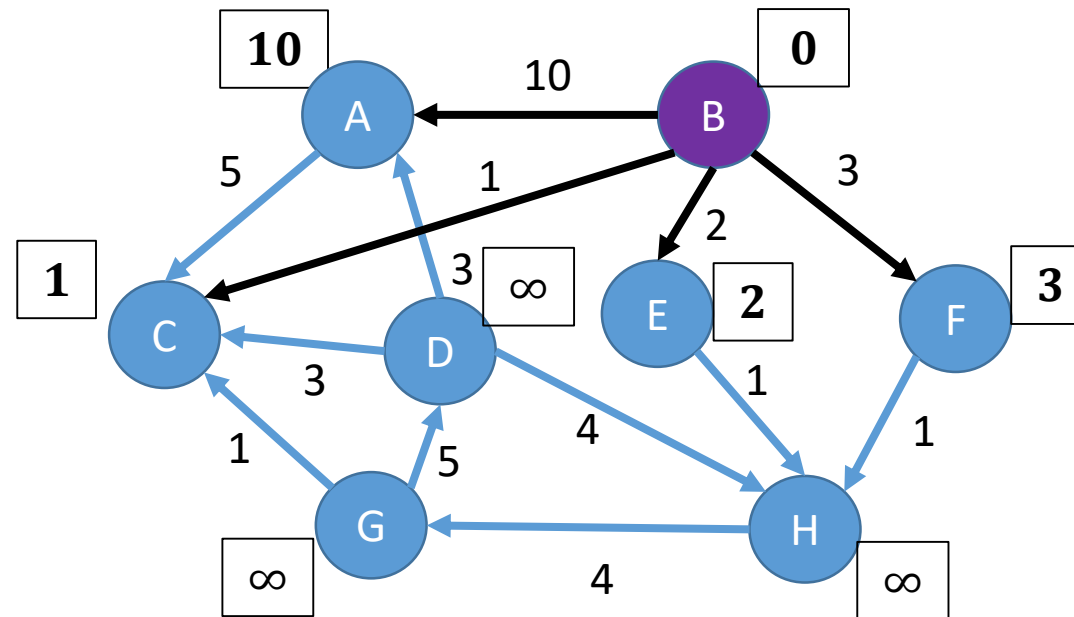
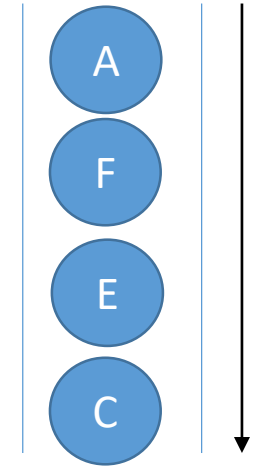


A cost =  $0 + 10 = 10$   
B cost = 0  
C cost =  $0 + 1 = 1$   
D cost =  $\infty$   
E cost =  $0 + 2 = 2$   
F cost =  $0 + 3 = 3$   
G cost =  $\infty$   
H cost =  $\infty$

# Dijkstra's Algorithm

- We are finished with B

Priority Queue/Min-heap

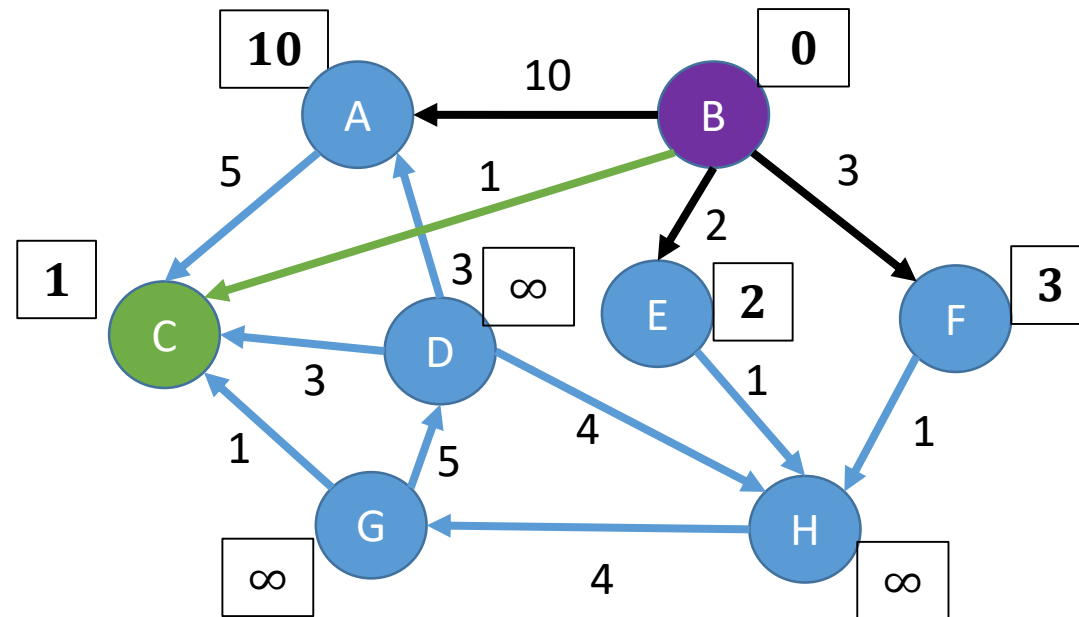


A cost = 10  
B cost = 0  
C cost = 1  
D cost =  $\infty$   
E cost = 2  
F cost = 3  
G cost =  $\infty$   
H cost =  $\infty$

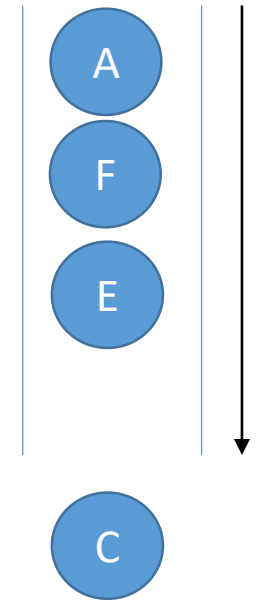


# Dijkstra's Algorithm

- We move on to the next node in the priority queue



Priority Queue/Min-heap

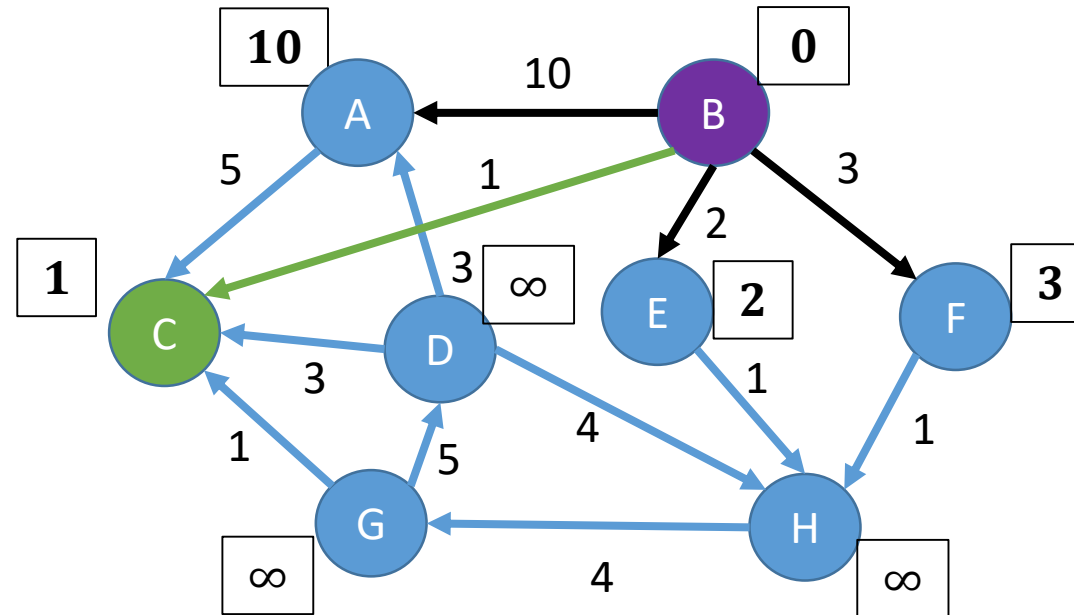
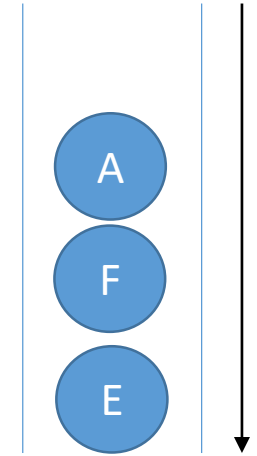


A cost = 10  
B cost = 0  
C cost = 1  
D cost =  $\infty$   
E cost = 2  
F cost = 3  
G cost =  $\infty$   
H cost =  $\infty$

# Dijkstra's Algorithm

- We'll now look at the nodes adjacent to C.
  - There are none

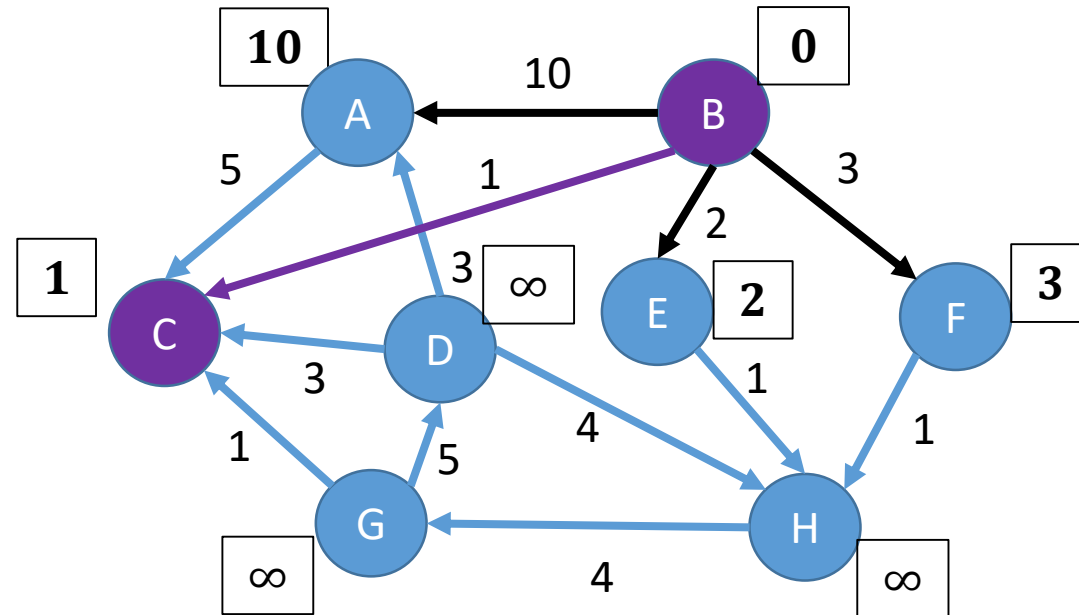
Priority Queue/Min-heap



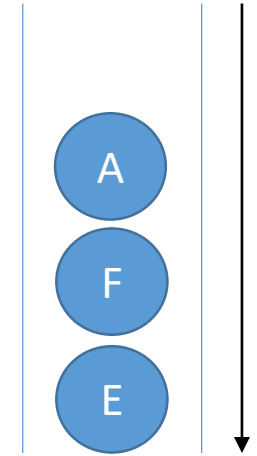
A cost = 10  
B cost = 0  
C cost = 1  
D cost =  $\infty$   
E cost = 2  
F cost = 3  
G cost =  $\infty$   
H cost =  $\infty$

# Dijkstra's Algorithm

- We are finished with C



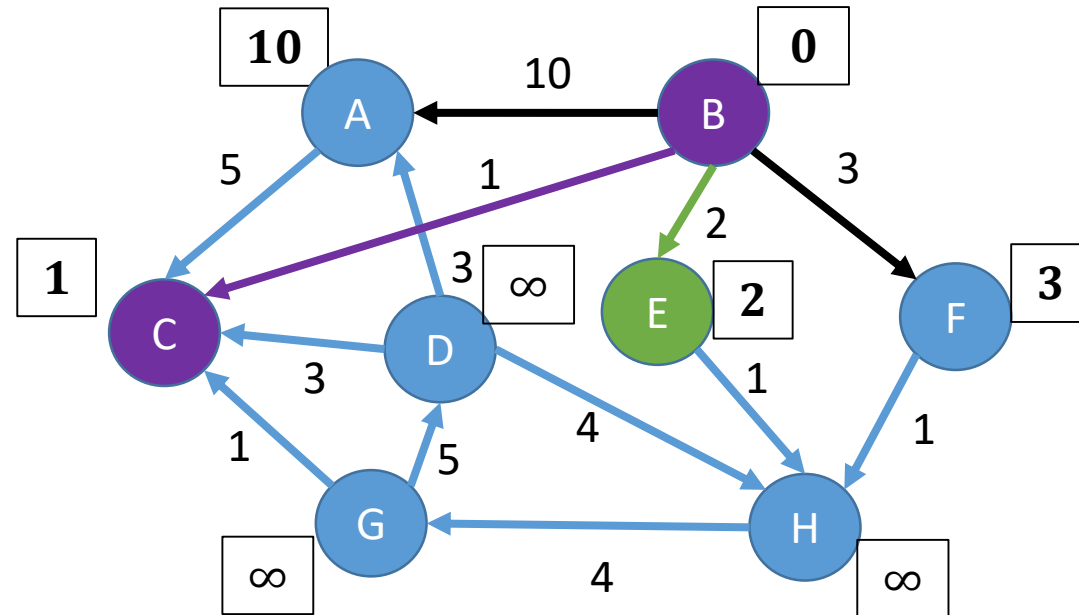
Priority Queue/Min-heap



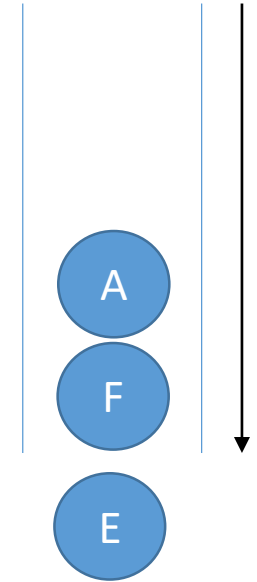
A cost = 10  
B cost = 0  
C cost = 1  
D cost =  $\infty$   
E cost = 2  
F cost = 3  
G cost =  $\infty$   
H cost =  $\infty$

# Dijkstra's Algorithm

- We move on to the next node in the priority queue



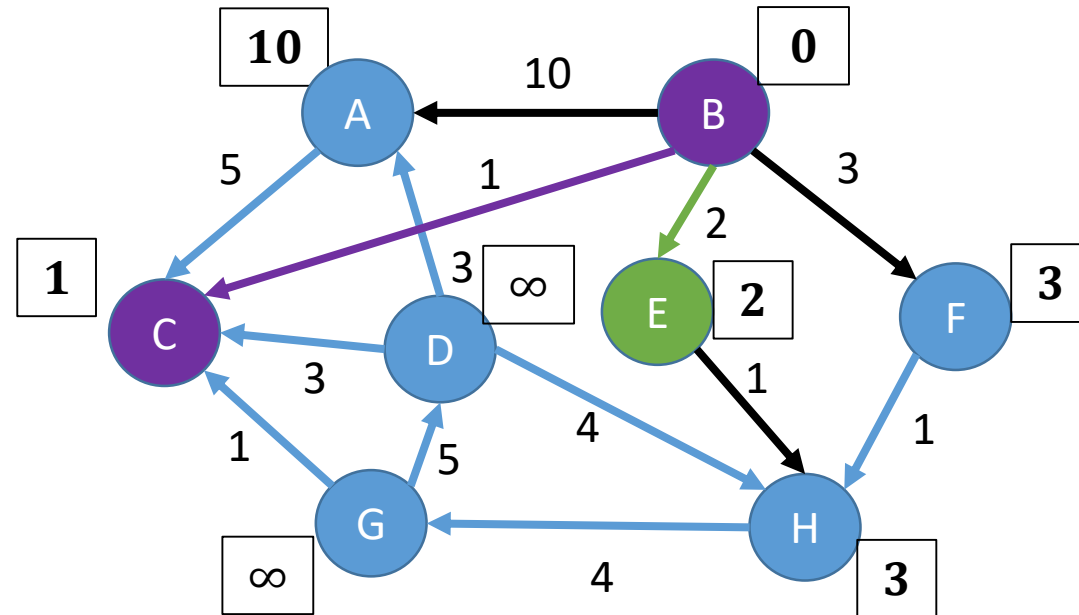
Priority Queue/Min-heap



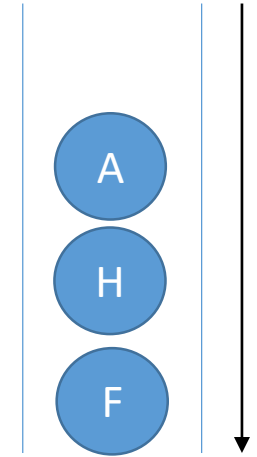
A cost = 10  
B cost = 0  
C cost = 1  
D cost =  $\infty$   
E cost = 2  
F cost = 3  
G cost =  $\infty$   
H cost =  $\infty$

# Dijkstra's Algorithm

- We'll now look at the nodes adjacent to E.
  - Cost of each is E's cost + cost of the edge



Priority Queue/Min-heap



A cost = 10

B cost = 0

C cost = 1

D cost =  $\infty$

E cost = 2

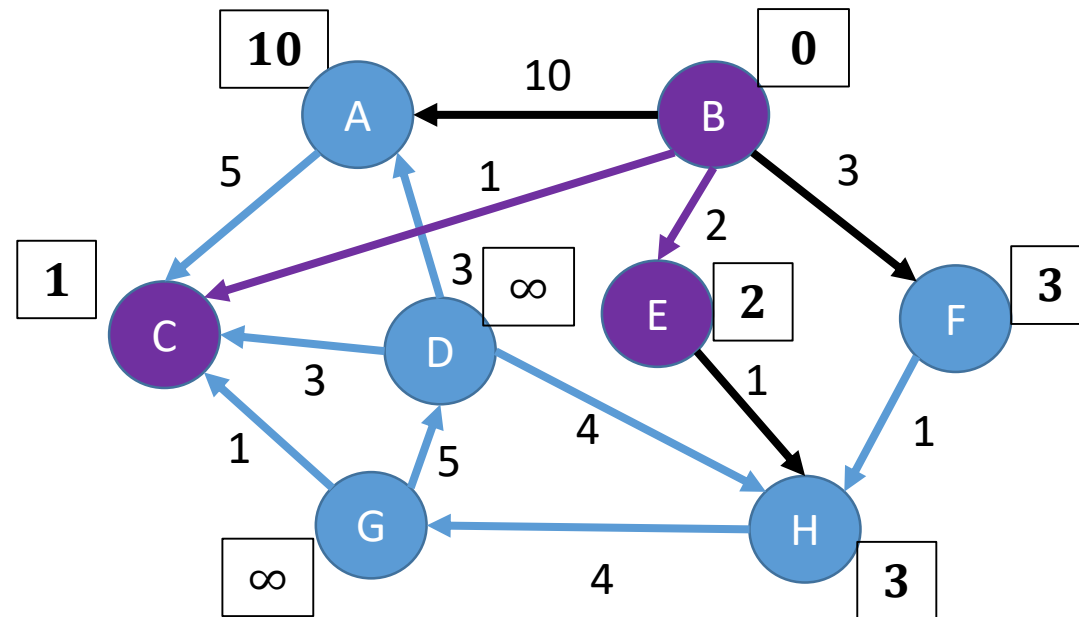
F cost = 3

G cost =  $\infty$

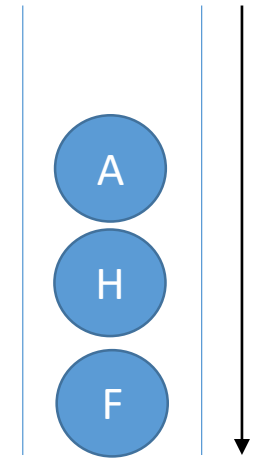
H cost = 2 + 1 = 3

# Dijkstra's Algorithm

- We are finished with E.



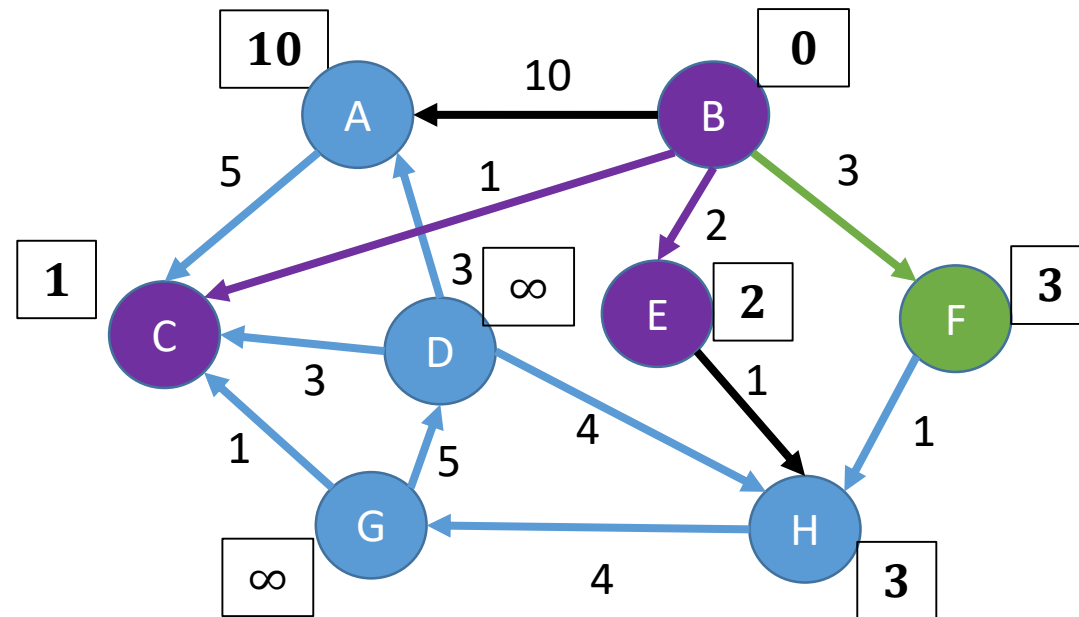
Priority Queue/Min-heap



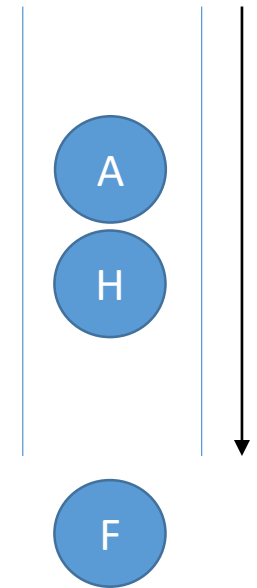
A cost = 10  
B cost = 0  
C cost = 1  
D cost =  $\infty$   
E cost = 2  
F cost = 3  
G cost =  $\infty$   
H cost = 3

# Dijkstra's Algorithm

- We move on to the next node in the priority queue



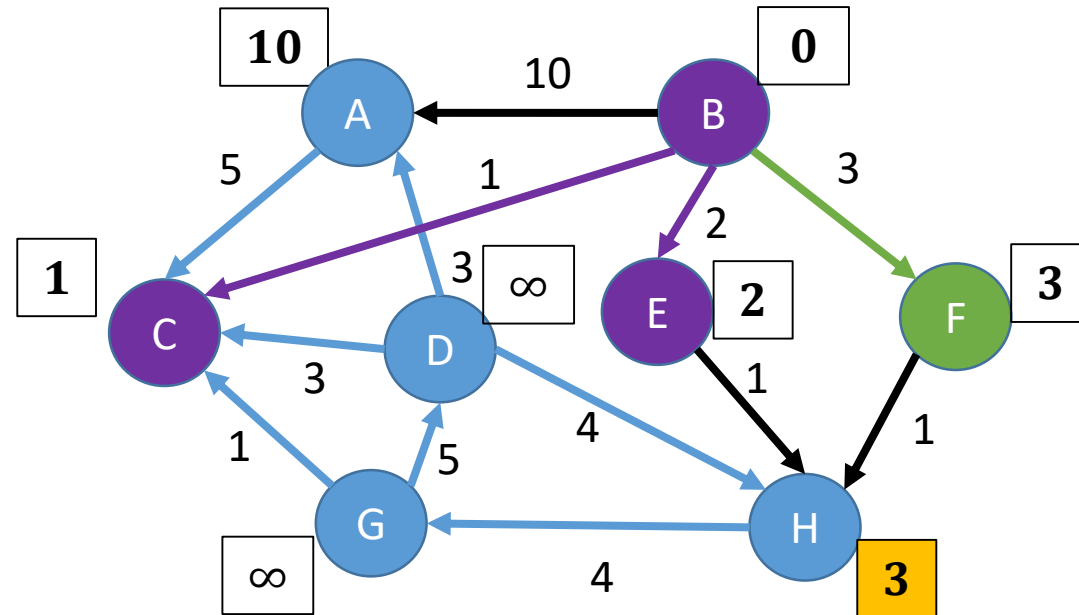
Priority Queue/Min-heap



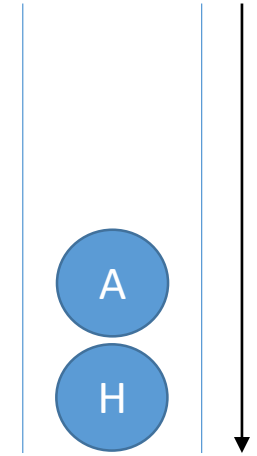
A cost = 10  
B cost = 0  
C cost = 1  
D cost =  $\infty$   
E cost = 2  
F cost = 3  
G cost =  $\infty$   
H cost = 3

# Dijkstra's Algorithm

- We'll now look at the nodes adjacent to F.
  - Cost of each is F's cost + cost of the edge
  - $3 + 1 = 4$  (NOT LESS THAN THE CURRENT COST OF H)



Priority Queue/Min-heap



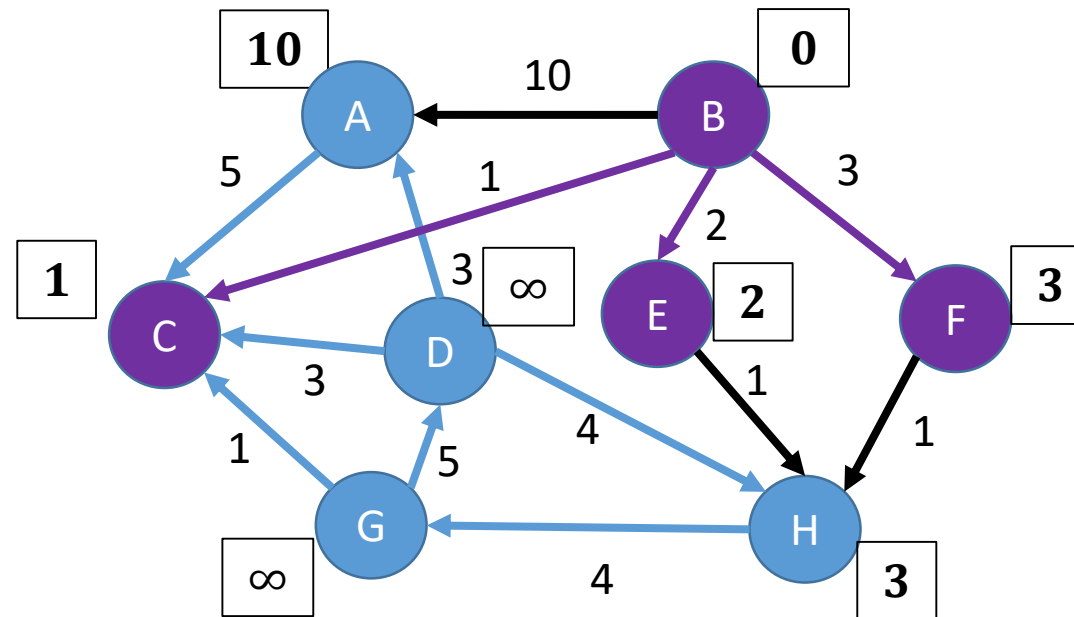
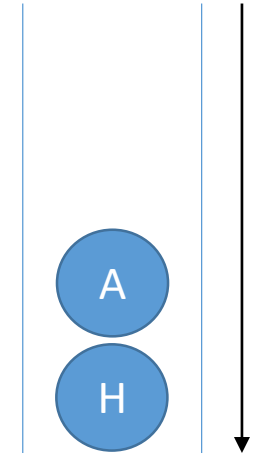
A cost = 10  
B cost = 0  
C cost = 1  
D cost =  $\infty$   
E cost = 2  
F cost = 3  
G cost =  $\infty$   
H cost = 3



# Dijkstra's Algorithm

- We are finished with F

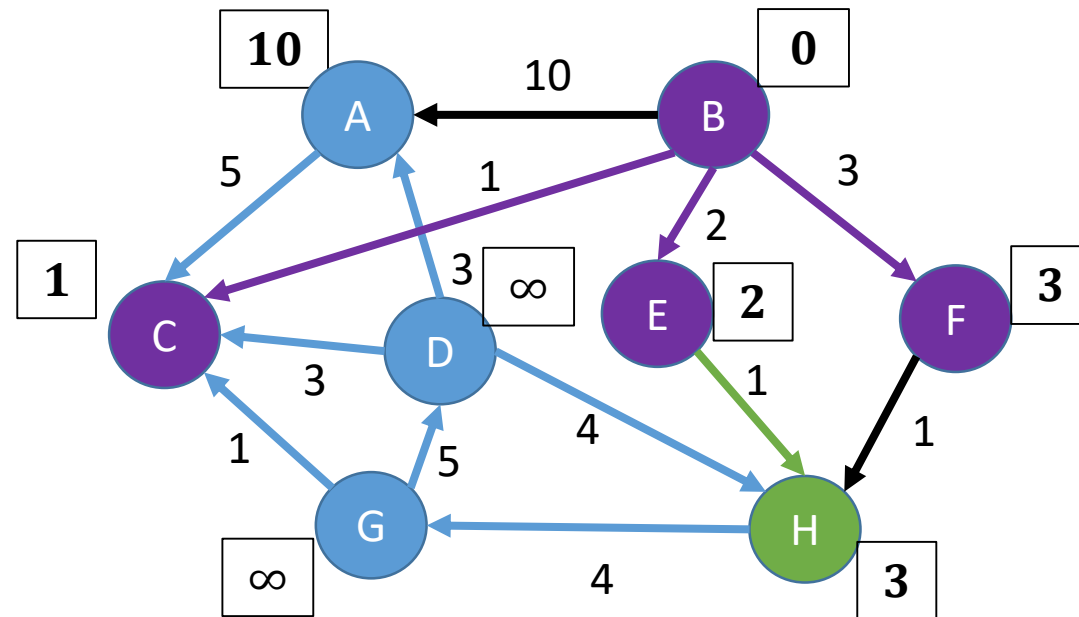
Priority Queue/Min-heap



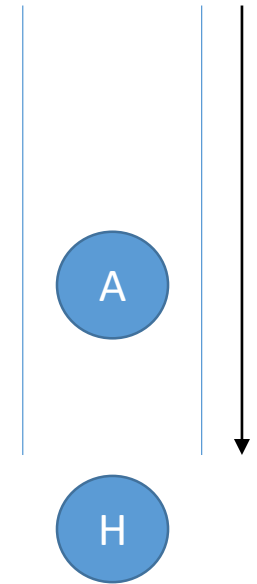
A cost = 10  
B cost = 0  
C cost = 1  
D cost =  $\infty$   
E cost = 2  
F cost = 3  
G cost =  $\infty$   
H cost = 3

# Dijkstra's Algorithm

- We move on to the next node in the priority queue



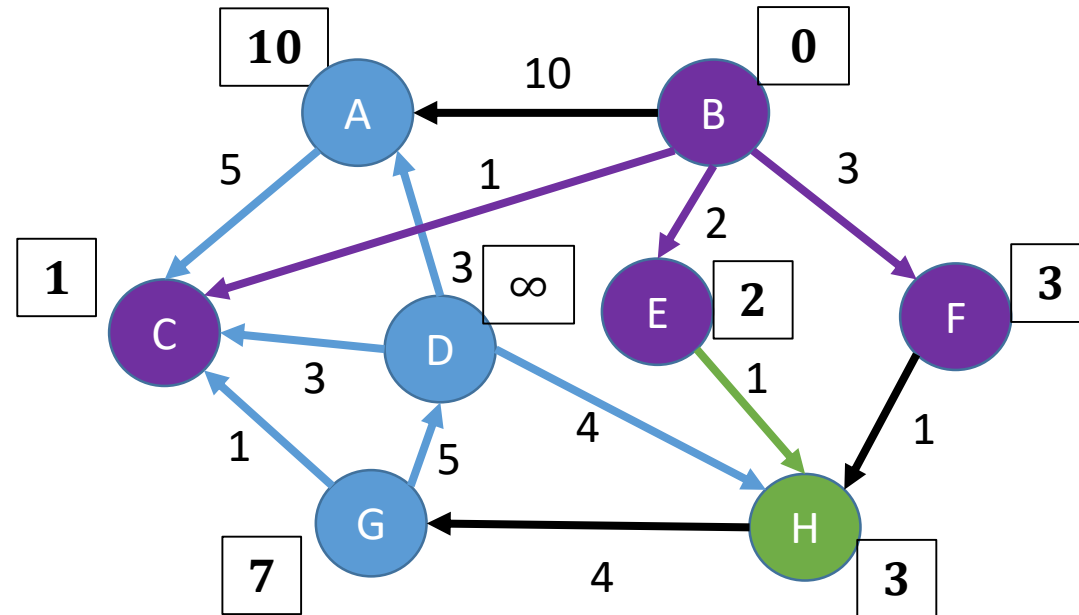
Priority Queue/Min-heap



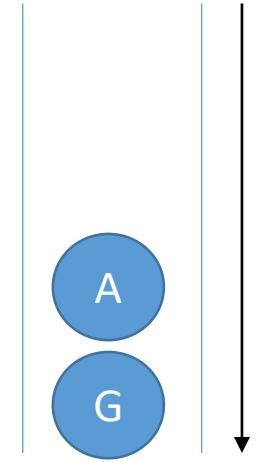
A cost = 10  
B cost = 0  
C cost = 1  
D cost =  $\infty$   
E cost = 2  
F cost = 3  
G cost =  $\infty$   
H cost = 3

# Dijkstra's Algorithm

- We'll now look at the nodes adjacent to H.
  - Cost of each is H's cost + cost of the edge



Priority Queue/Min-heap

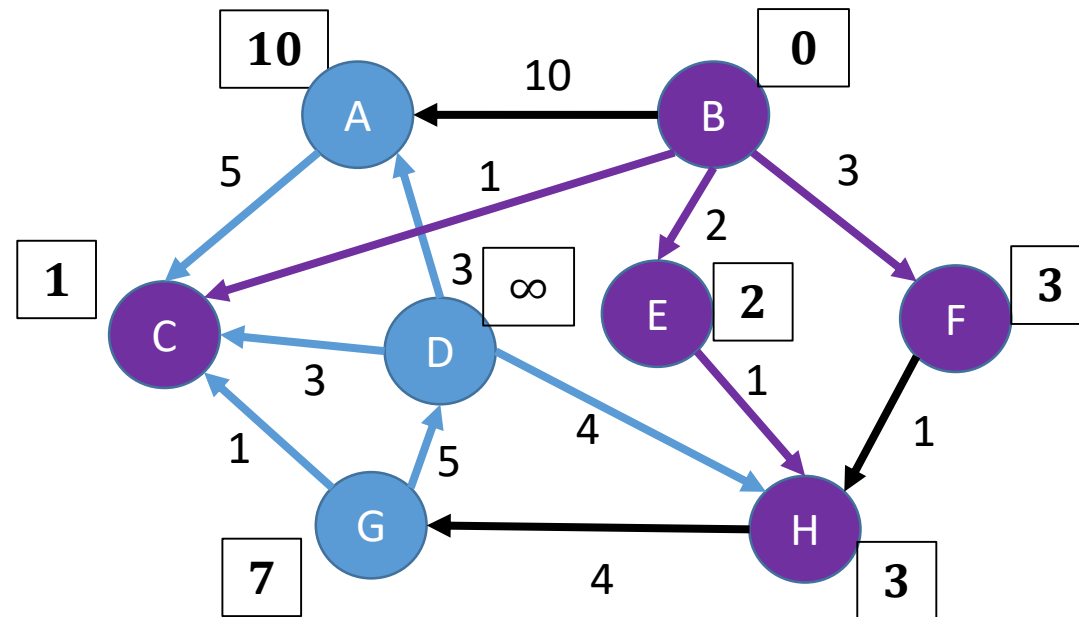
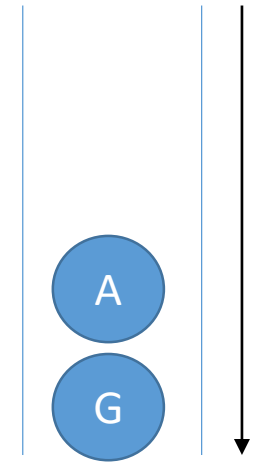


A cost = 10  
B cost = 0  
C cost = 1  
D cost =  $\infty$   
E cost = 2  
F cost = 3  
G cost = 3 + 4 = 7  
H cost = 3

# Dijkstra's Algorithm

- We are finished with H.

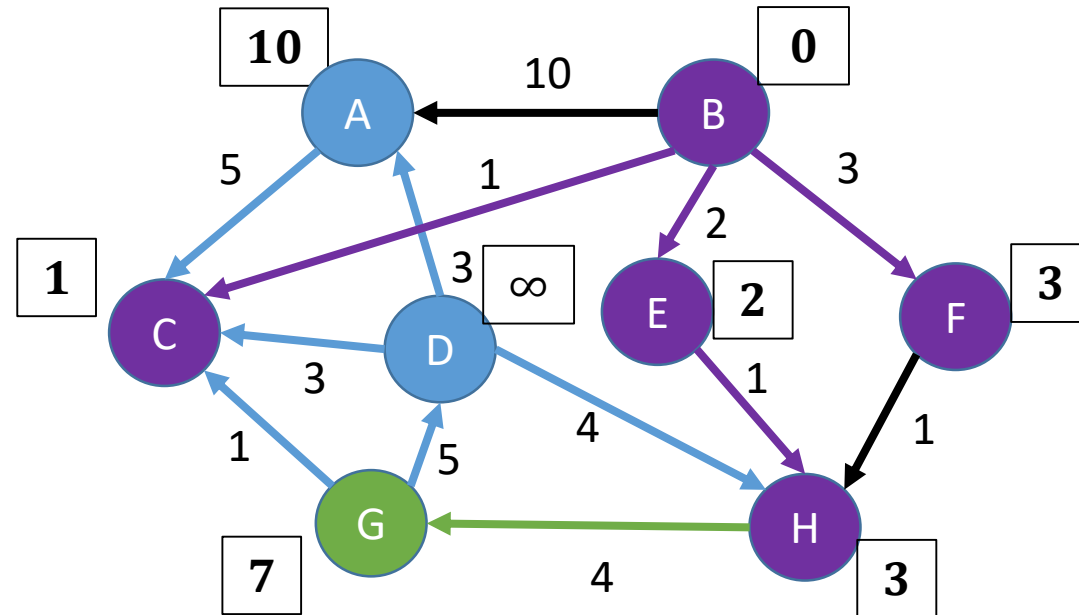
Priority Queue/Min-heap



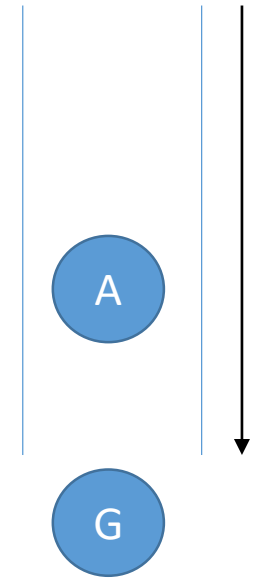
A cost = 10  
B cost = 0  
C cost = 1  
D cost =  $\infty$   
E cost = 2  
F cost = 3  
G cost = 7  
H cost = 3

# Dijkstra's Algorithm

- We move on to the next node in the priority queue



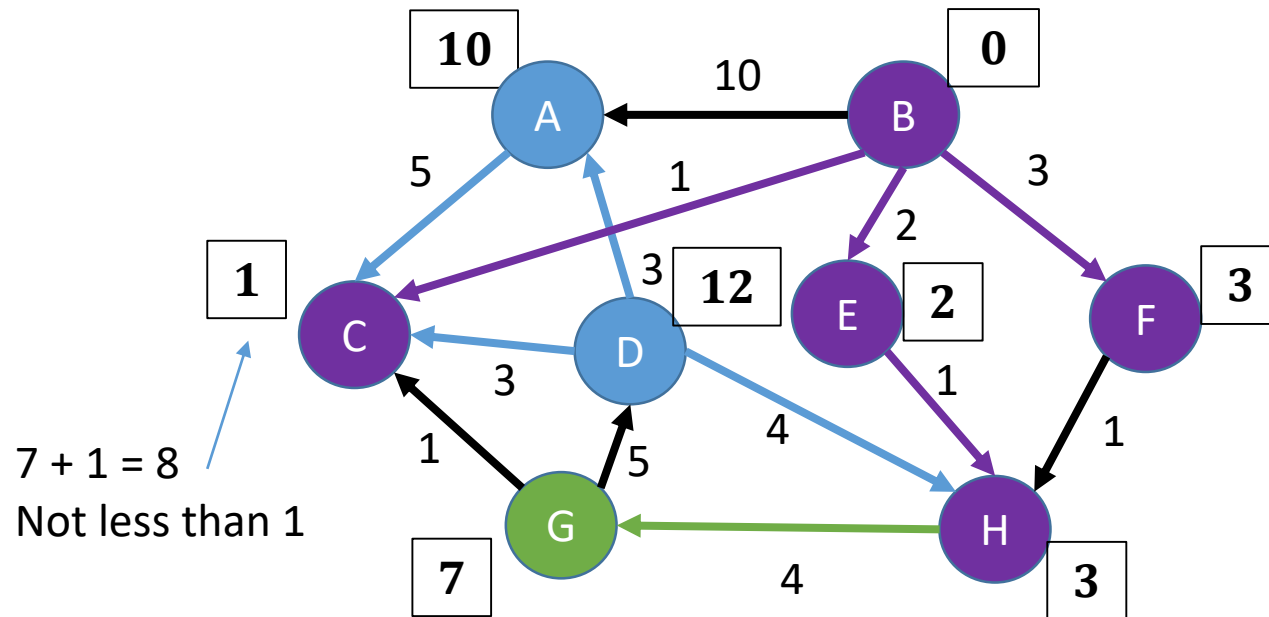
Priority Queue/Min-heap



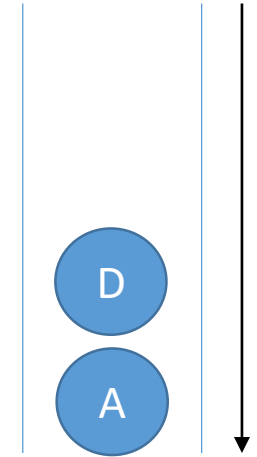
A cost = 10  
B cost = 0  
C cost = 1  
D cost =  $\infty$   
E cost = 2  
F cost = 3  
G cost = 7  
H cost = 3

# Dijkstra's Algorithm

- We'll now look at the nodes adjacent to G.
  - Cost of each is G's cost + cost of the edge



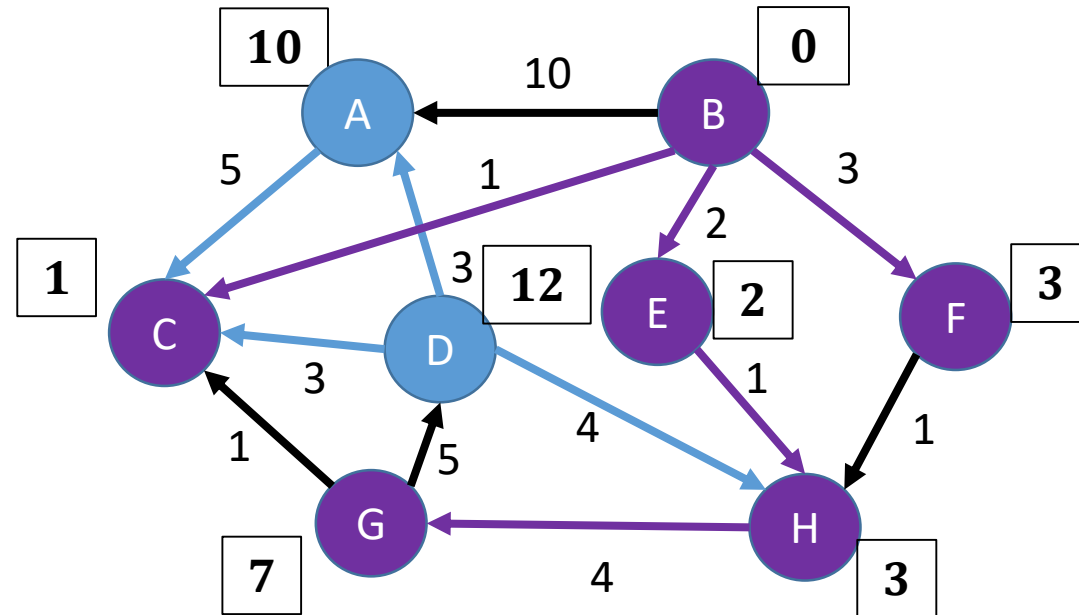
Priority Queue/Min-heap



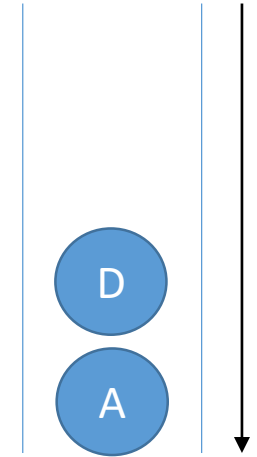
A cost = 10  
B cost = 0  
C cost = 1  
D cost = 7 + 5 = 12  
E cost = 2  
F cost = 3  
G cost = 7  
H cost = 3

# Dijkstra's Algorithm

- We are finished with G



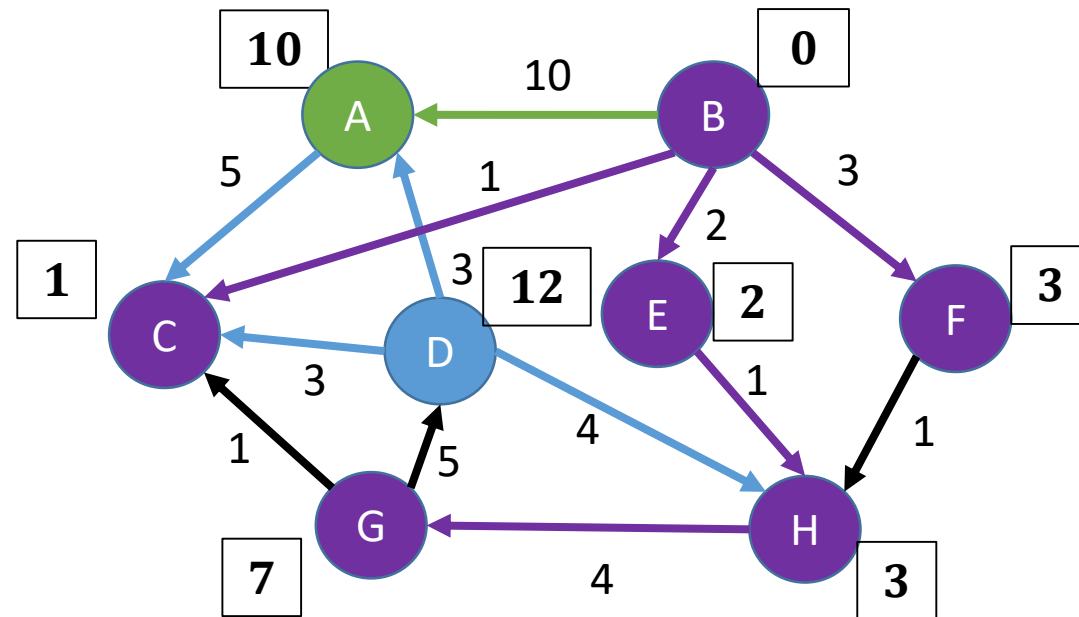
Priority Queue/Min-heap



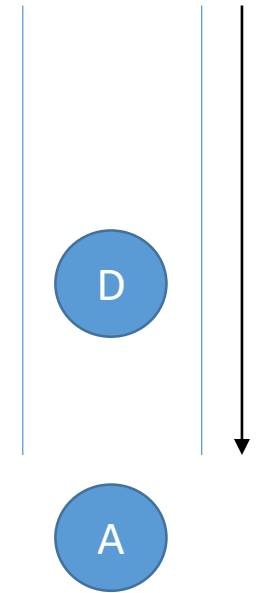
A cost = 10  
B cost = 0  
C cost = 1  
D cost = 12  
E cost = 2  
F cost = 3  
G cost = 7  
H cost = 3

# Dijkstra's Algorithm

- We move on to the next node in the priority queue



Priority Queue/Min-heap

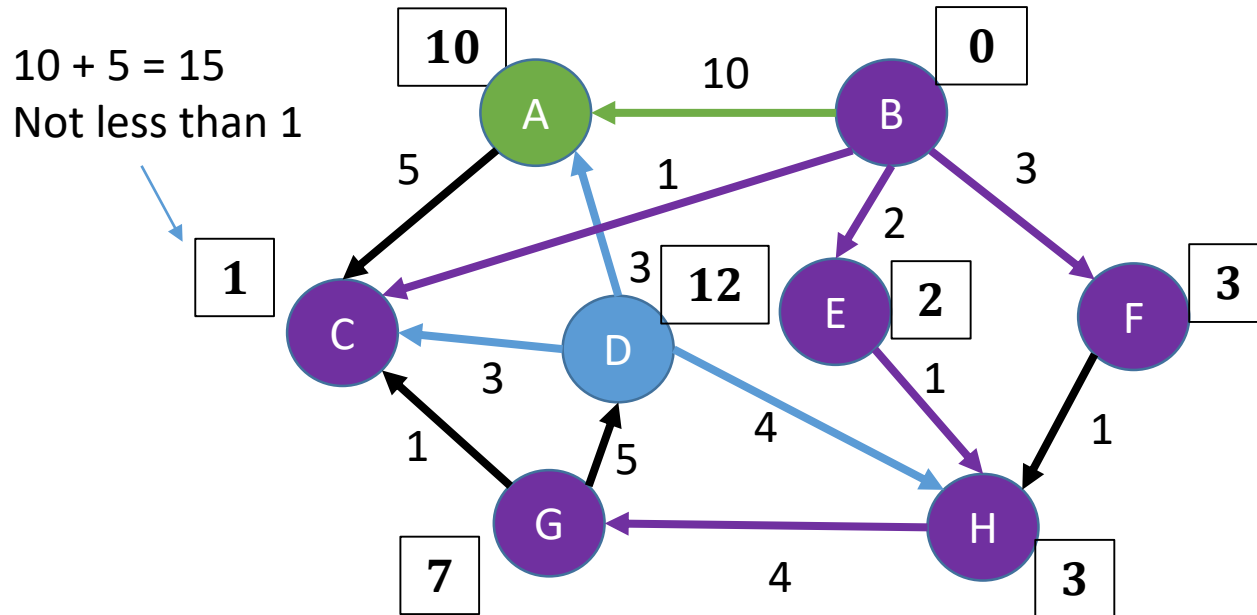


A cost = 10  
B cost = 0  
C cost = 1  
D cost = 12  
E cost = 2  
F cost = 3  
G cost = 7  
H cost = 3

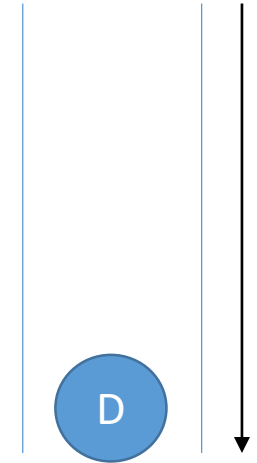


# Dijkstra's Algorithm

- We'll now look at the nodes adjacent to A.
  - Cost of each is A's cost + cost of the edge



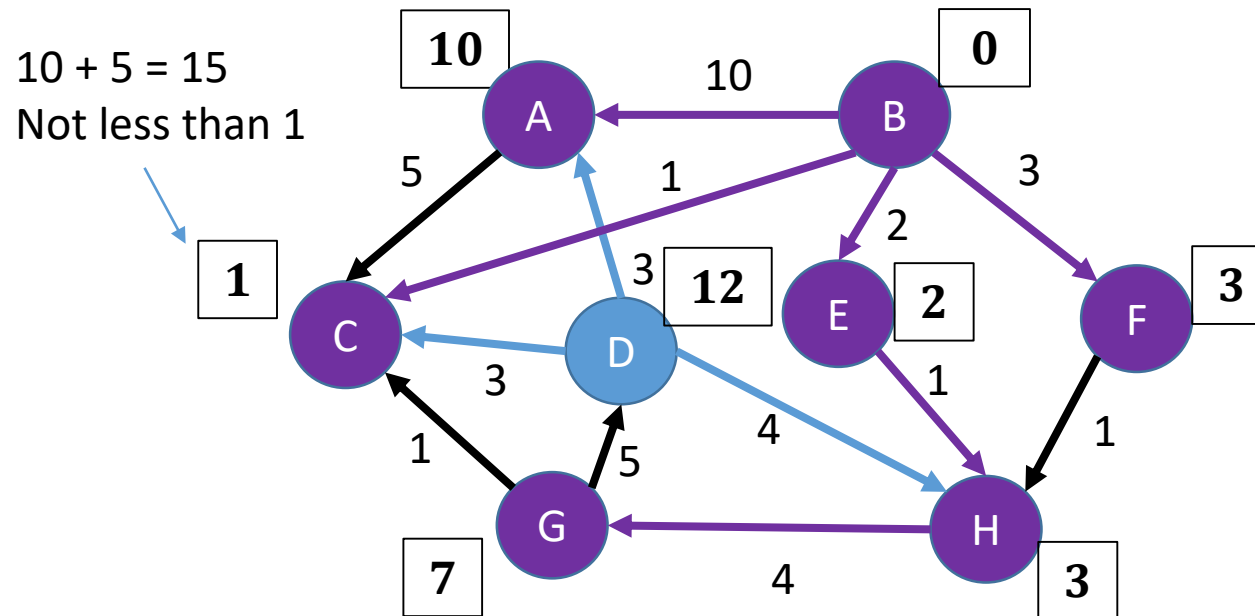
Priority Queue/Min-heap



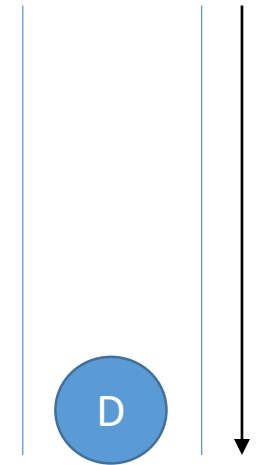
A cost = 10  
B cost = 0  
C cost = 1  
D cost = 12  
E cost = 2  
F cost = 3  
G cost = 7  
H cost = 3

# Dijkstra's Algorithm

- We are finished with A



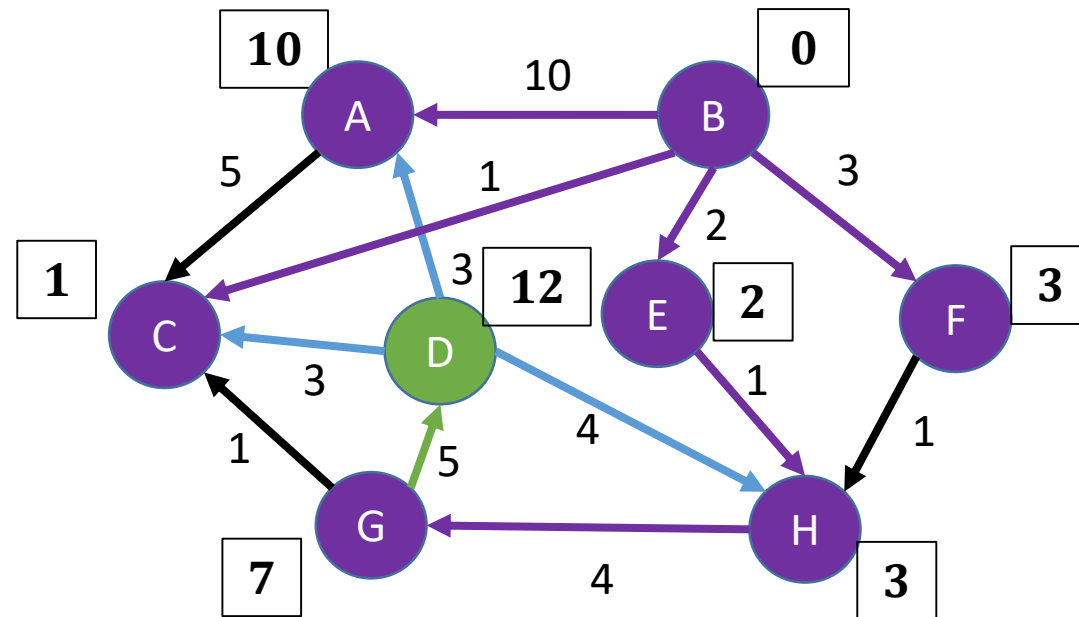
Priority Queue/Min-heap



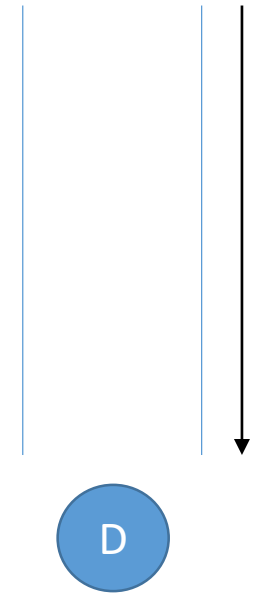
A cost = 10  
B cost = 0  
C cost = 1  
D cost = 12  
E cost = 2  
F cost = 3  
G cost = 7  
H cost = 3

# Dijkstra's Algorithm

- We move on to the next node in the priority queue



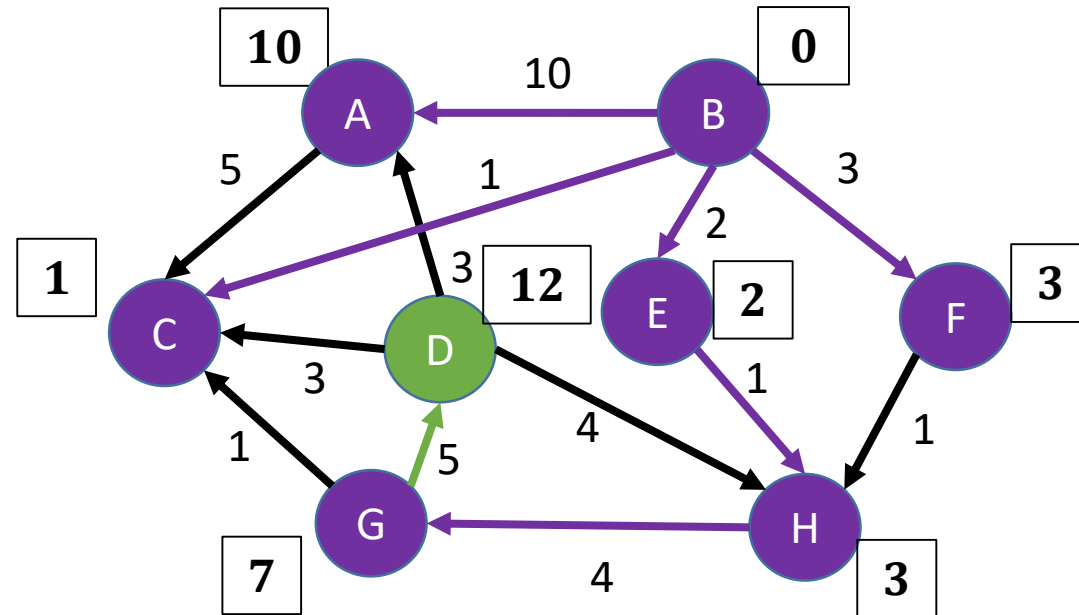
Priority Queue/Min-heap



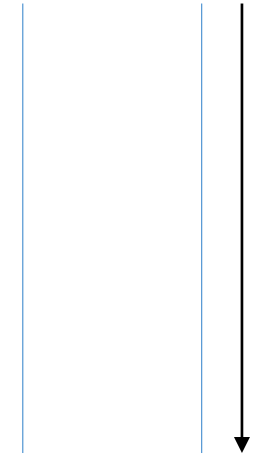
A cost = 10  
B cost = 0  
C cost = 1  
D cost = 12  
E cost = 2  
F cost = 3  
G cost = 7  
H cost = 3

# Dijkstra's Algorithm

- We'll now look at the nodes adjacent to D.
  - Cost of each is D's cost + cost of the edge



Priority Queue/Min-heap

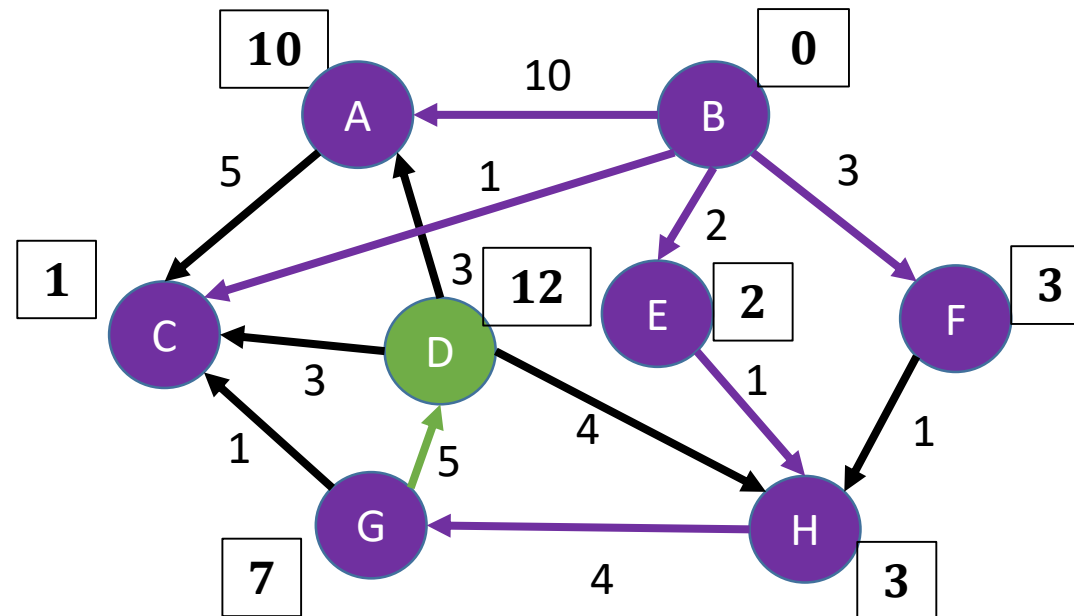
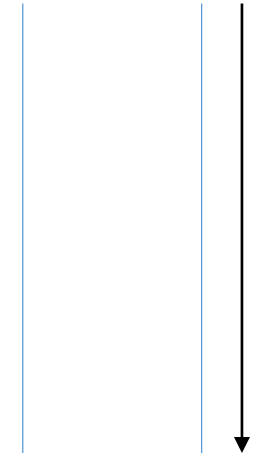


A cost = 10  
B cost = 0  
C cost = 1  
D cost = 12  
E cost = 2  
F cost = 3  
G cost = 7  
H cost = 3

# Dijkstra's Algorithm

- $D \rightarrow A = 12 + 3 = 15$  (Not less than 10)
- $D \rightarrow C = 12 + 3 = 15$  (Not less than 1)
- $D \rightarrow H = 12 + 4 = 16$  (Not less than 3)

Priority Queue/Min-heap

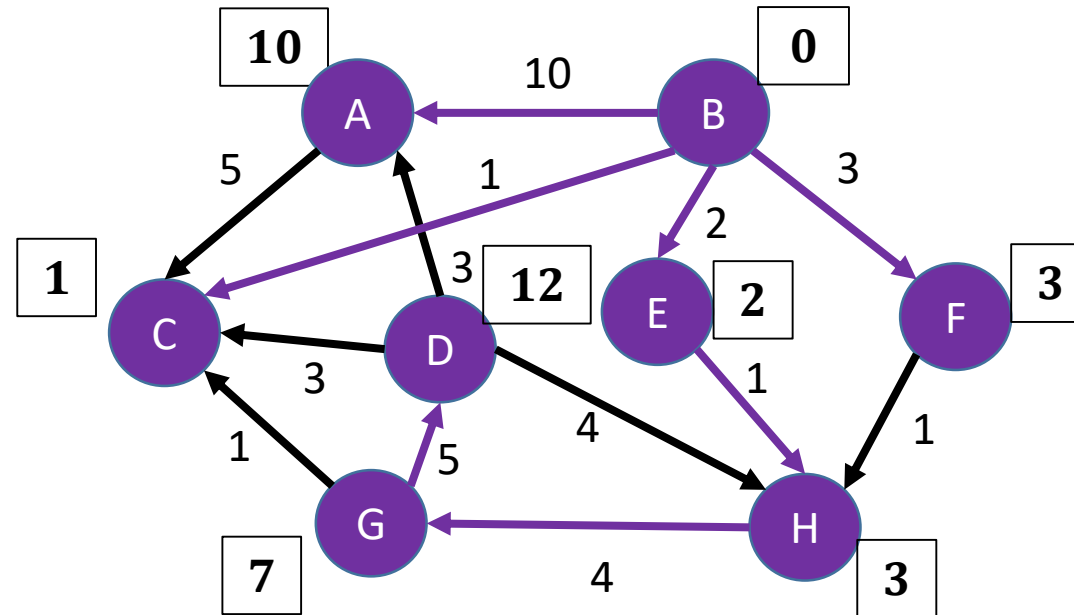
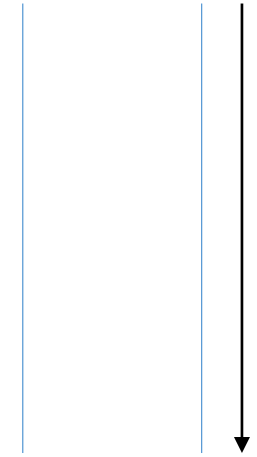


A cost = 10  
B cost = 0  
C cost = 1  
D cost = 12  
E cost = 2  
F cost = 3  
G cost = 7  
H cost = 3

# Dijkstra's Algorithm

- We are finished with D

Priority Queue/Min-heap

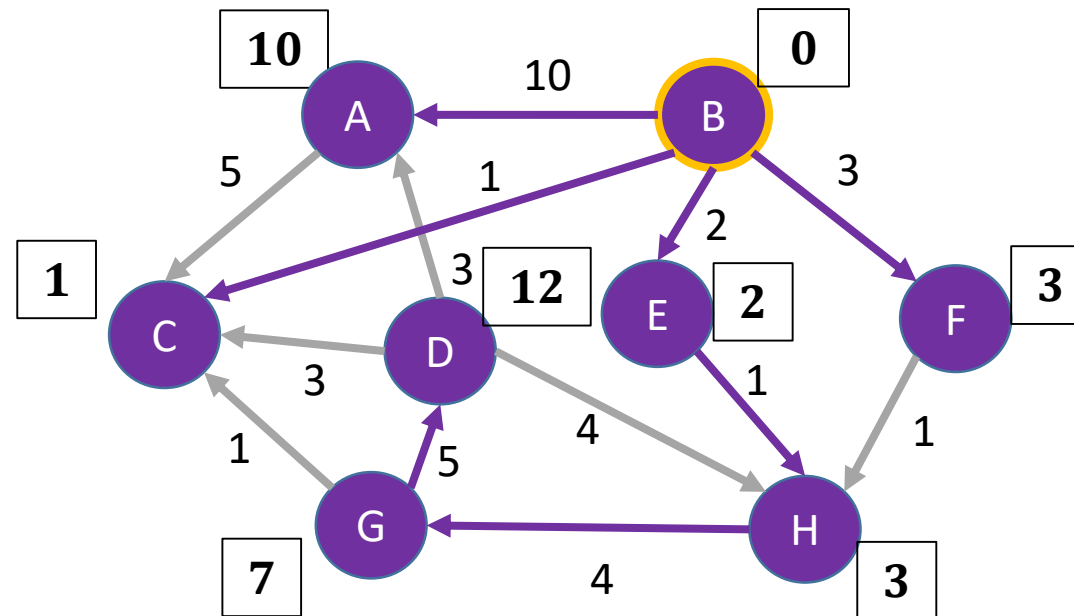
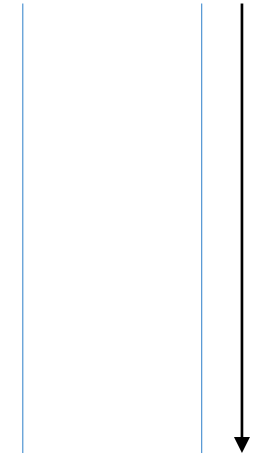


A cost = 10  
B cost = 0  
C cost = 1  
D cost = 12  
E cost = 2  
F cost = 3  
G cost = 7  
H cost = 3

# Dijkstra's Algorithm

- Priority Queue is empty
  - Algorithm is complete

Priority Queue/Min-heap



**Least cost from B to all other nodes:**

A cost = 10  
B cost = 0  
C cost = 1  
D cost = 12  
E cost = 2  
F cost = 3  
G cost = 7  
H cost = 3