

Graphs II

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Lecture Topics

- Graph Traversals
 - Depth-First
 - Breadth-First
- Finding distance with breadth-first
- Weighted Graphs
 - Dijkstra's Algorithm

Graph Traversal

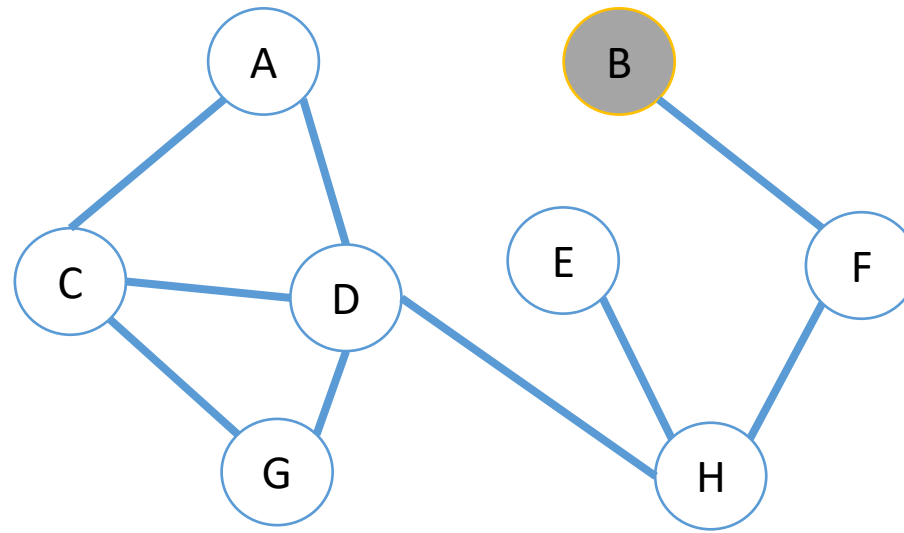
- Graphs can be traversed using depth-first or breadth-first traversals.
- Similar to traversing a tree, but we can begin at any node in the graph.
- Depth-first traversal uses a *stack*
- Breadth-first traversal uses a *queue*

Depth-First Traversal

- Depth-first traversals work by coloring each node
 - White: Node has not been seen previously
 - Gray: Node has been seen previously
 - Black: Node has been visited and we are done with it
- This prevents getting stuck in a cycle, if one existed
 - The traversal is the same for bi-directional graphs or digraphs

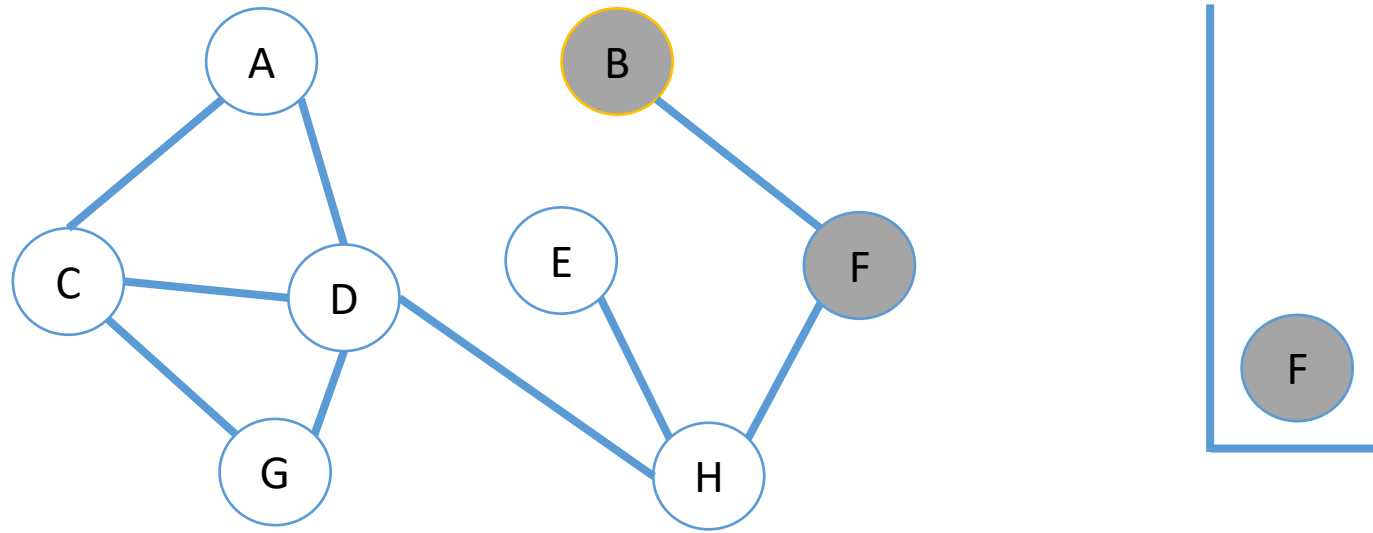
Depth-First Traversal

- Starting at B
- Order visited: N/A



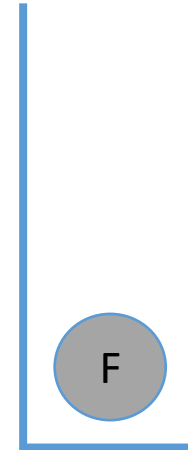
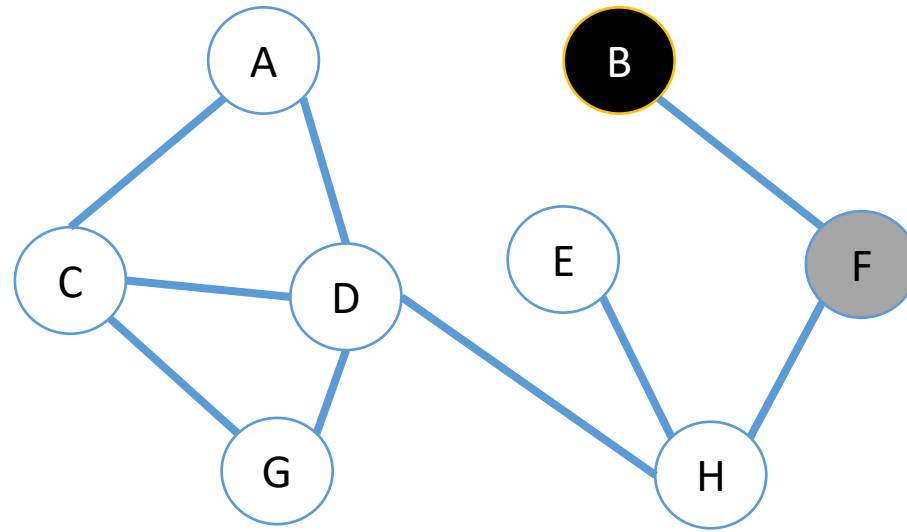
Depth-First Traversal

- Adds its adjacent white node(s) to the stack (turning the white node(s) gray)
- Order visited: N/A



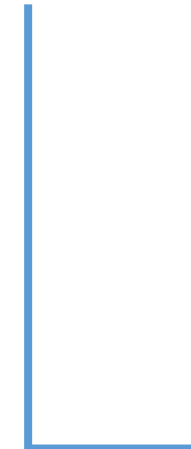
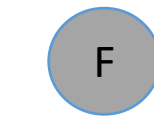
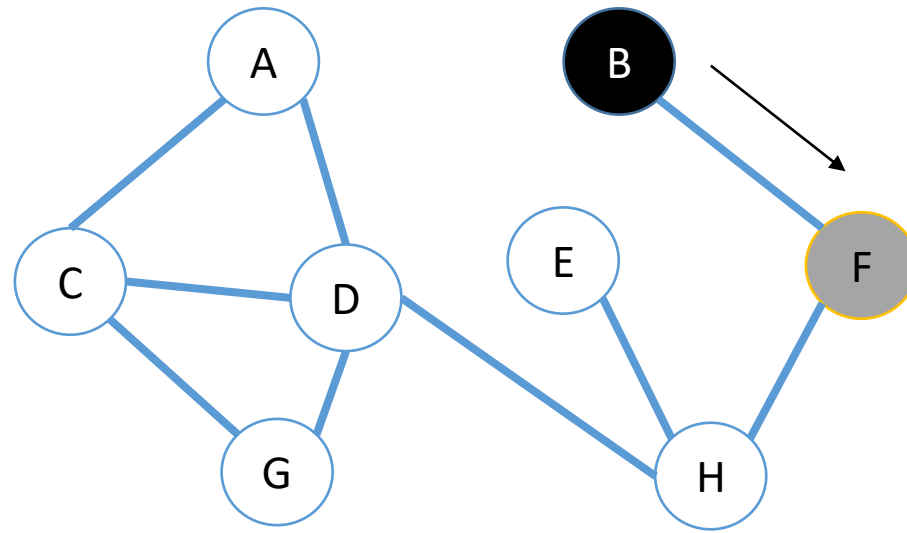
Depth-First Traversal

- B was visited
 - Changed to black
- Order visited: B



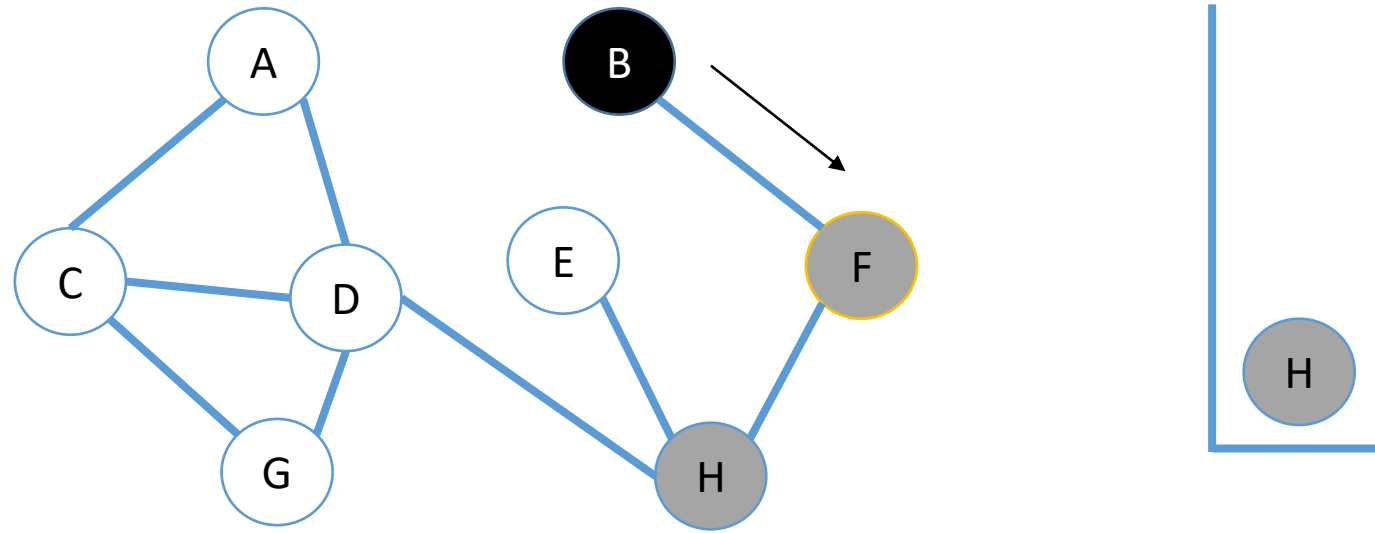
Depth-First Traversal

- Next node popped from stack
- Order visited: B



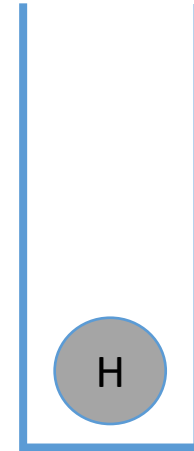
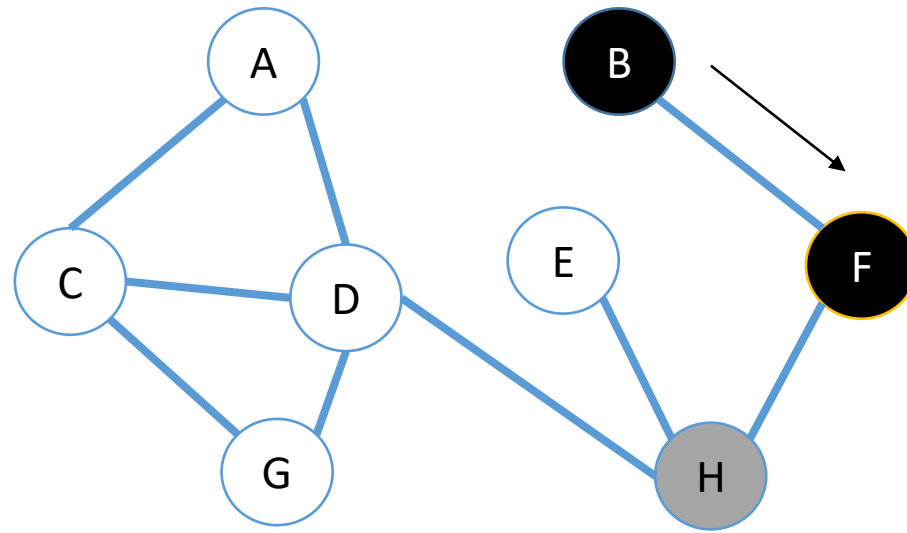
Depth-First Traversal

- Adds its adjacent white nodes to the stack (turning them gray)
- Order visited: B



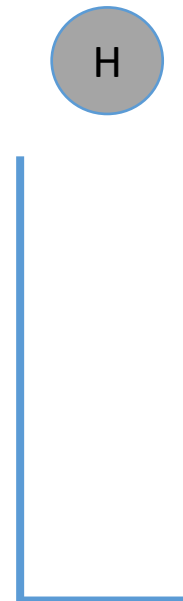
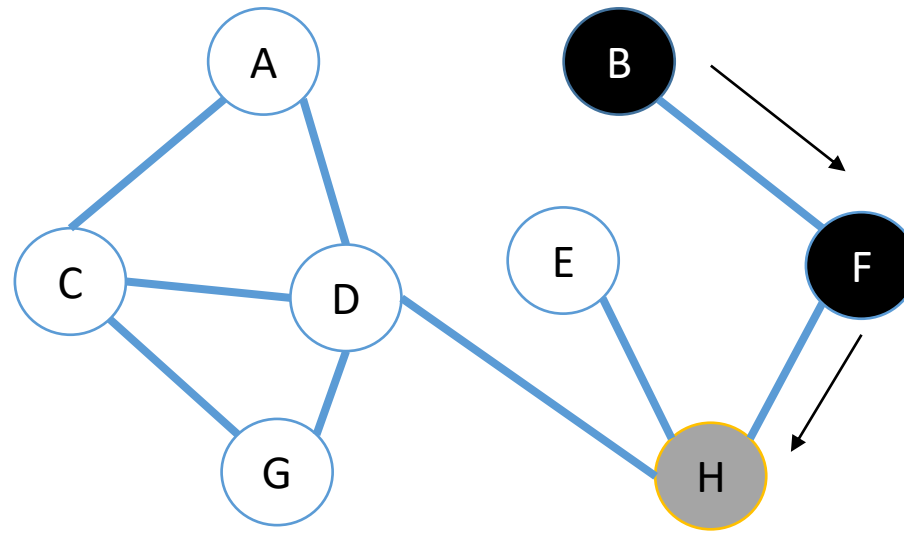
Depth-First Traversal

- F was visited
- Order visited: B, F



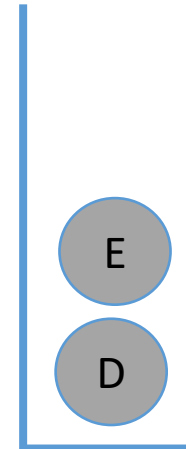
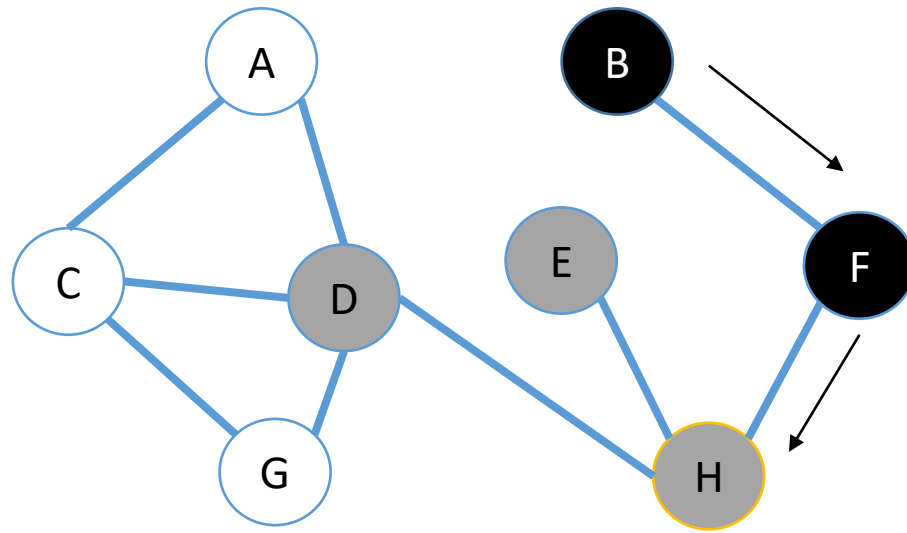
Depth-First Traversal

- Next node popped from stack
- Order visited: B, F



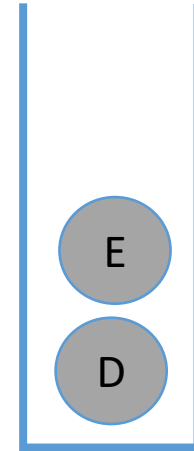
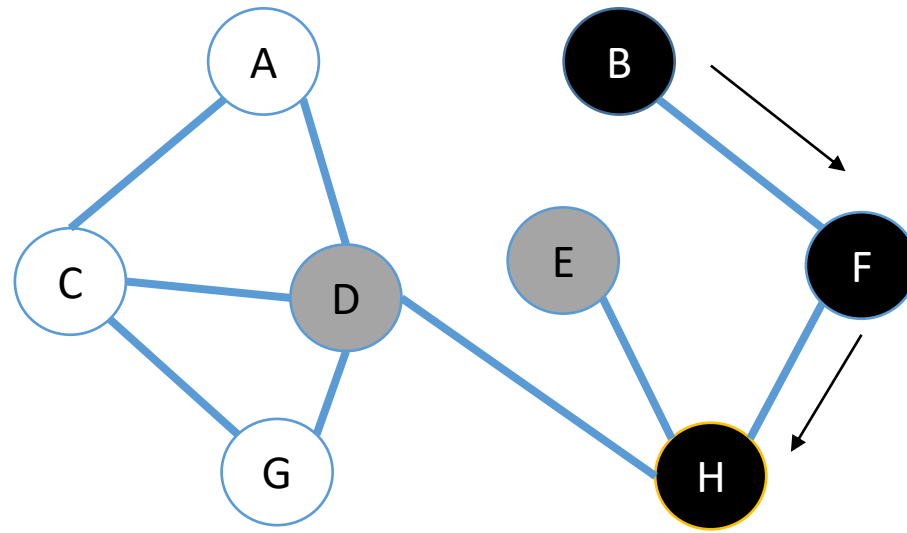
Depth-First Traversal

- Adds its adjacent white nodes to the stack
- Order visited: B, F



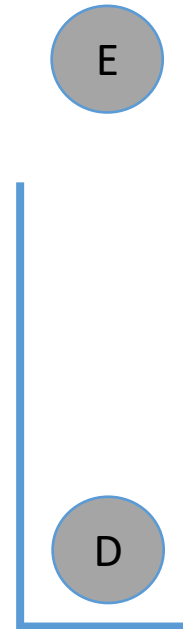
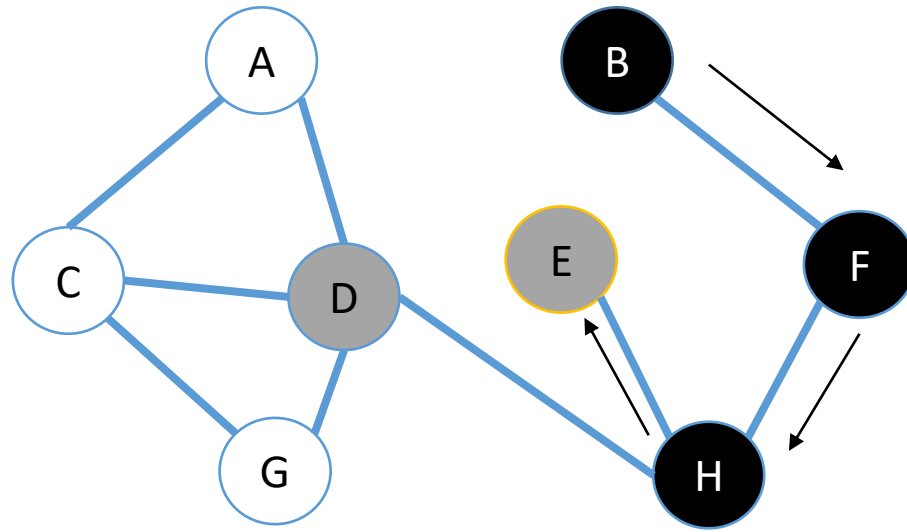
Depth-First Traversal

- H was visited
- Order visited: B, F, H



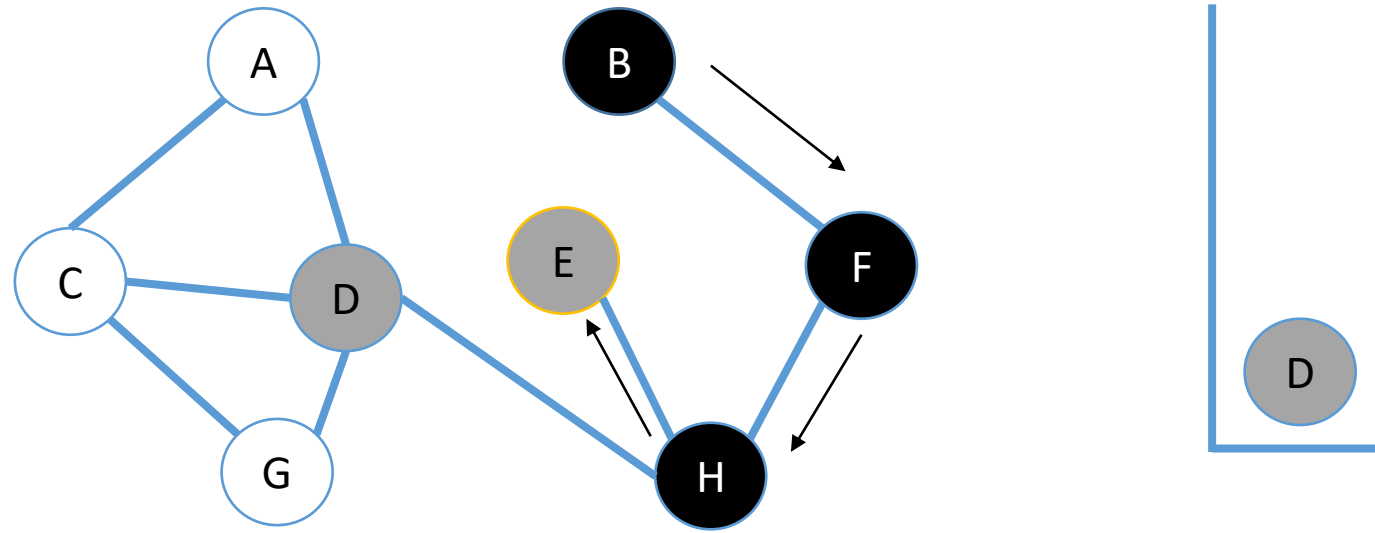
Depth-First Traversal

- Next node popped from stack
- Order visited: B, F, H



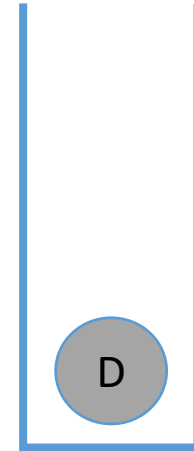
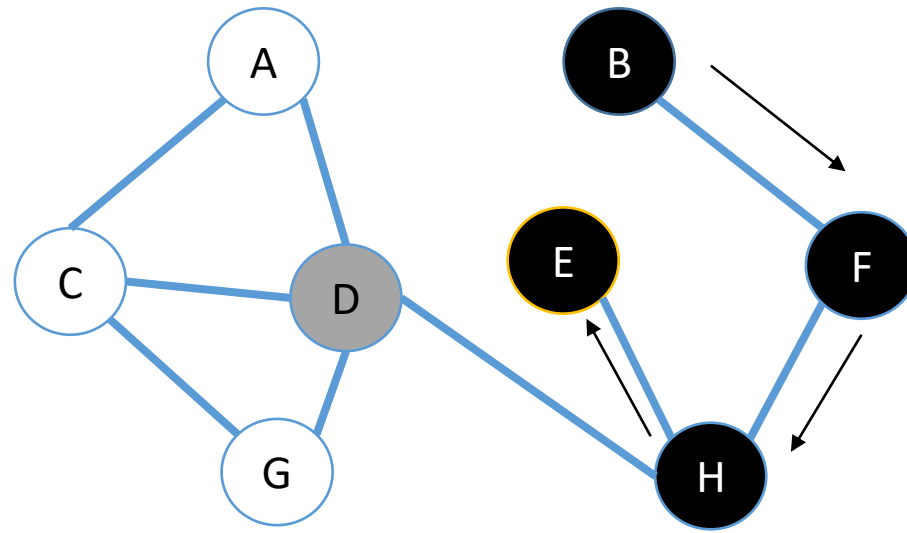
Depth-First Traversal

- Add its adjacent white nodes to the stack (there are none)
- Order visited: B, F, H



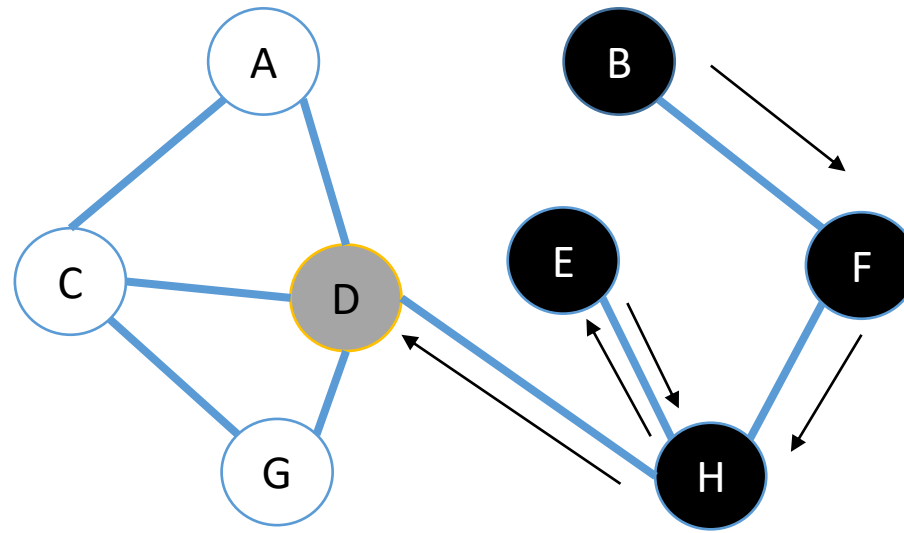
Depth-First Traversal

- E was visited
- Order visited: B, F, H, E



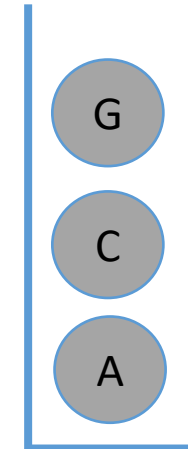
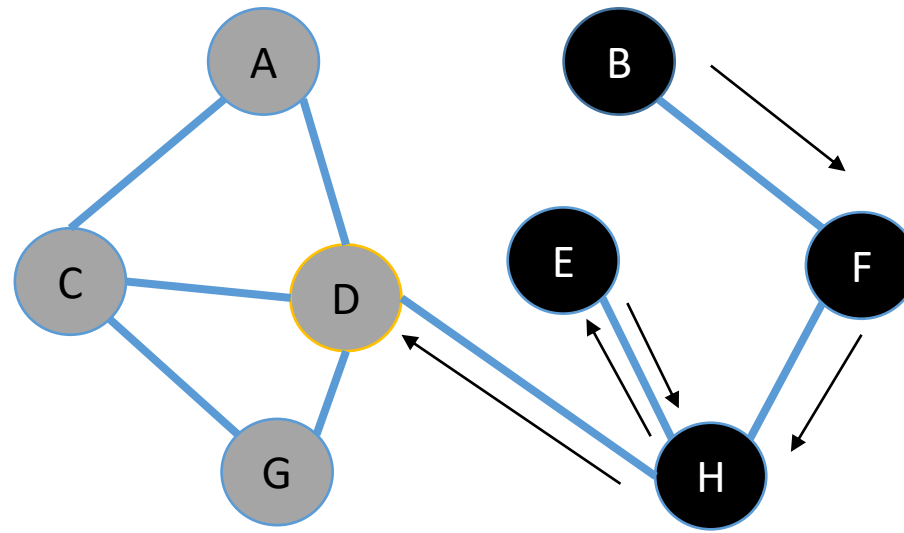
Depth-First Traversal

- Next node popped from stack
- Order visited: B, F, H, E



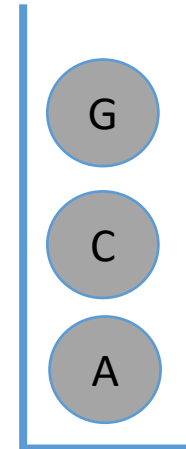
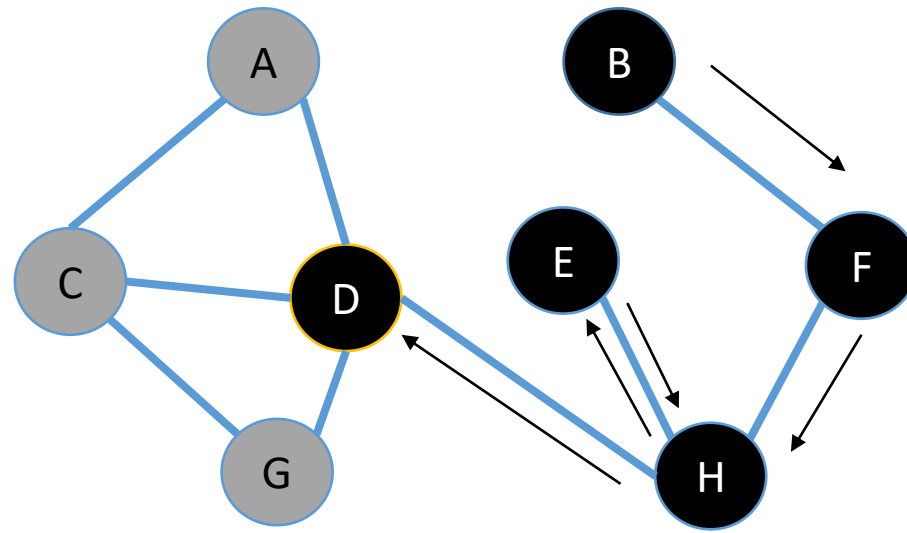
Depth-First Traversal

- Add its adjacent white nodes to the stack
- Order visited: B, F, H, E



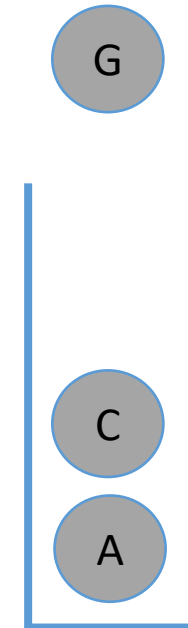
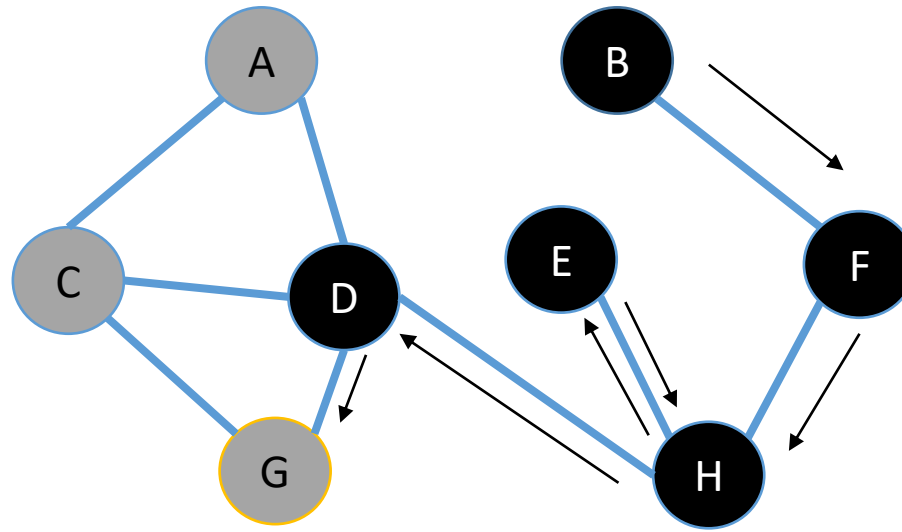
Depth-First Traversal

- D was visited
- Order visited: B, F, H, E, D



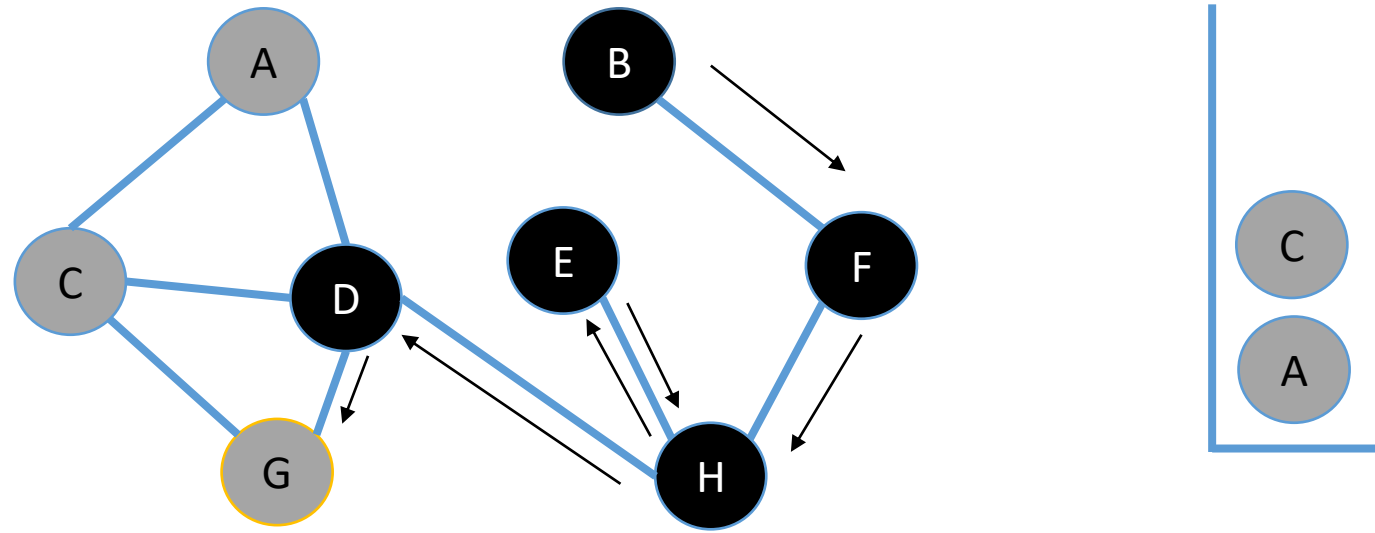
Depth-First Traversal

- Next node popped from the stack
- Order visited: B, F, H, E, D



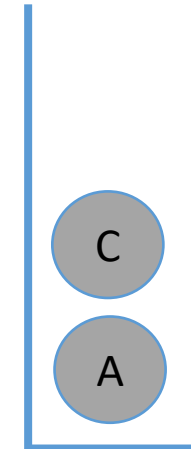
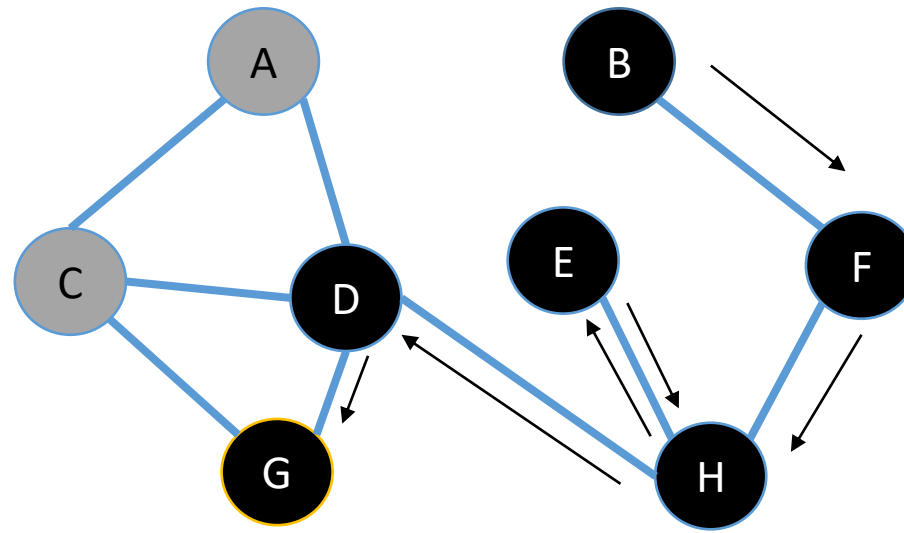
Depth-First Traversal

- Add its adjacent white nodes to the stack (there are none)
- Order visited: B, F, H, E, D



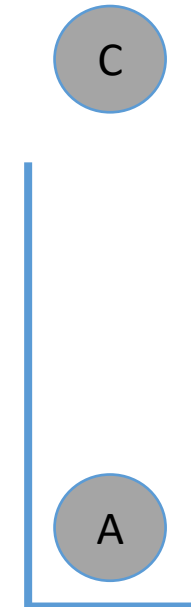
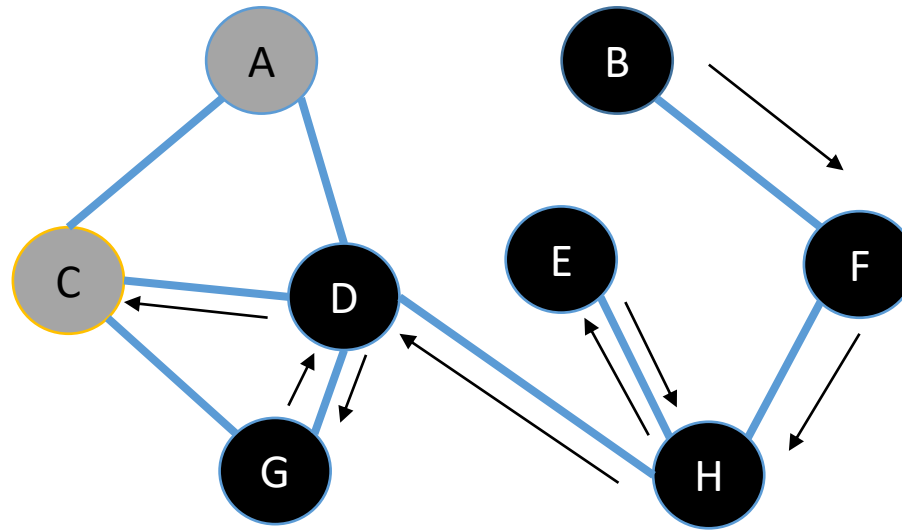
Depth-First Traversal

- G has been visited
- Order visited: B, F, H, E, D, G



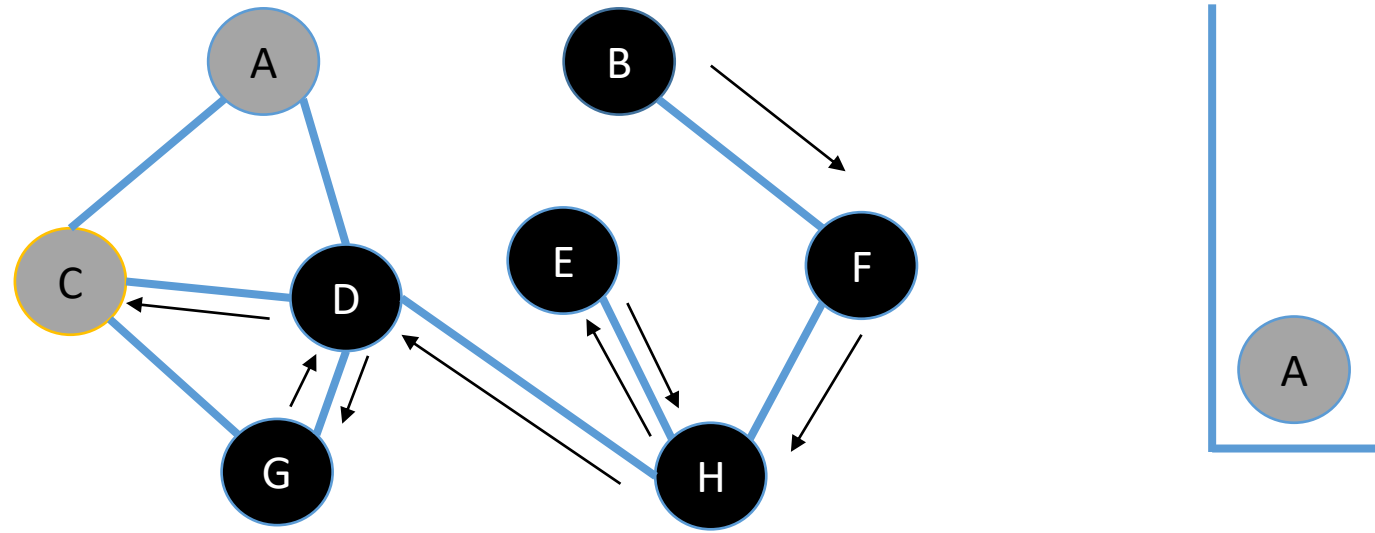
Depth-First Traversal

- Next node popped from the stack
- Order visited: B, F, H, E, D, G



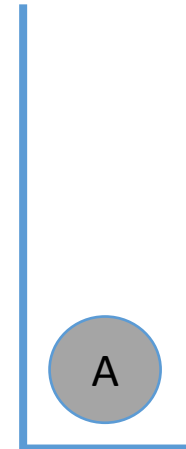
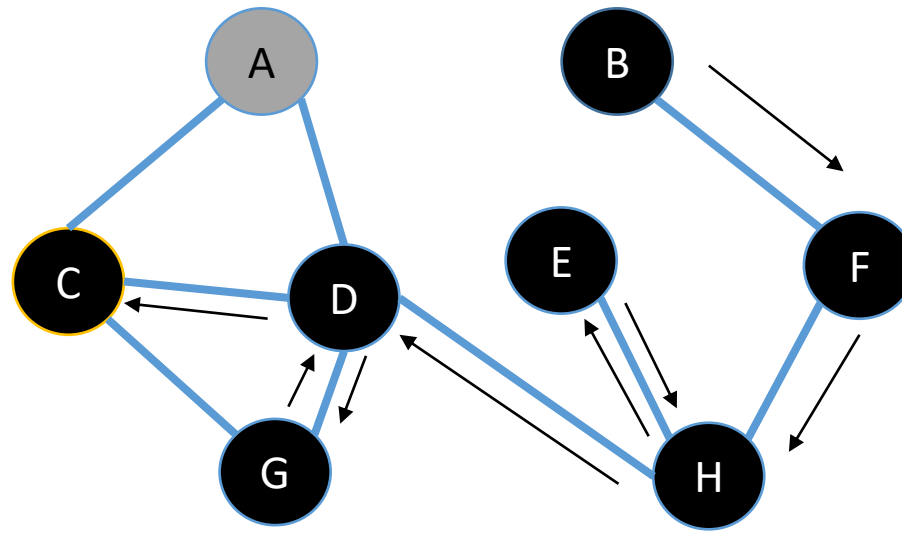
Depth-First Traversal

- Add its adjacent white nodes to the stack (there are none)
- Order visited: B, F, H, E, D, G



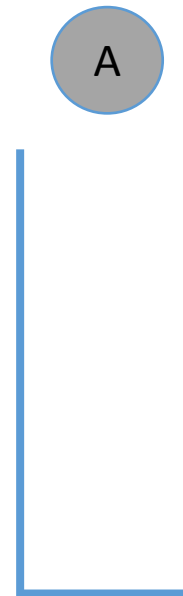
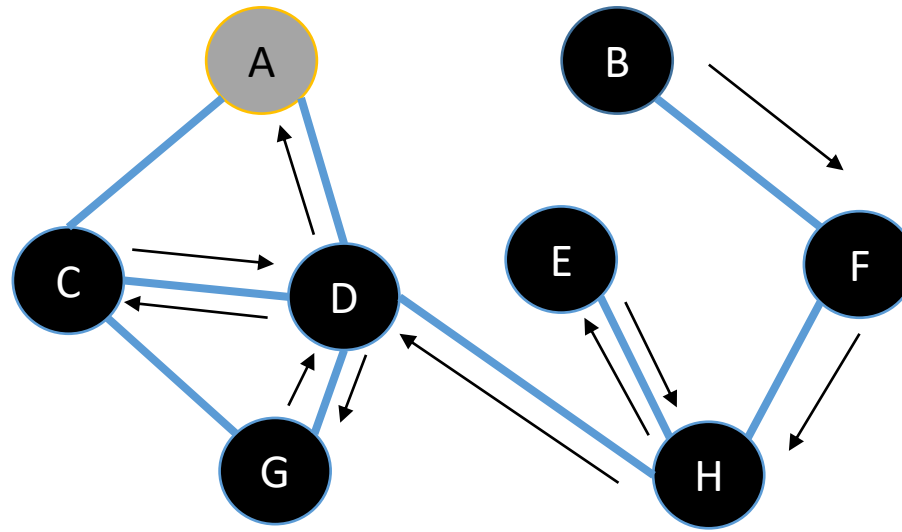
Depth-First Traversal

- C has been visited
- Order visited: B, F, H, E, D, G, C



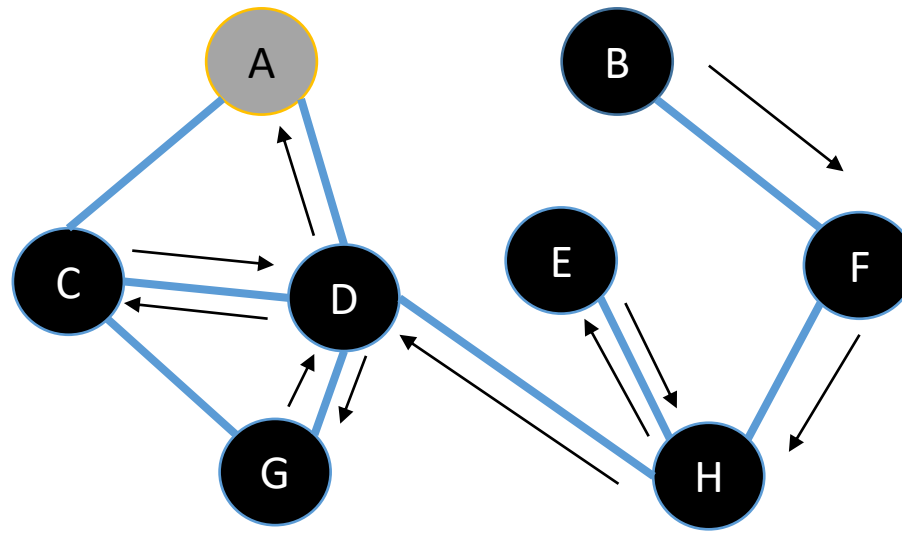
Depth-First Traversal

- Next node is popped from the stack
- Order visited: B, F, H, E, D, G, C



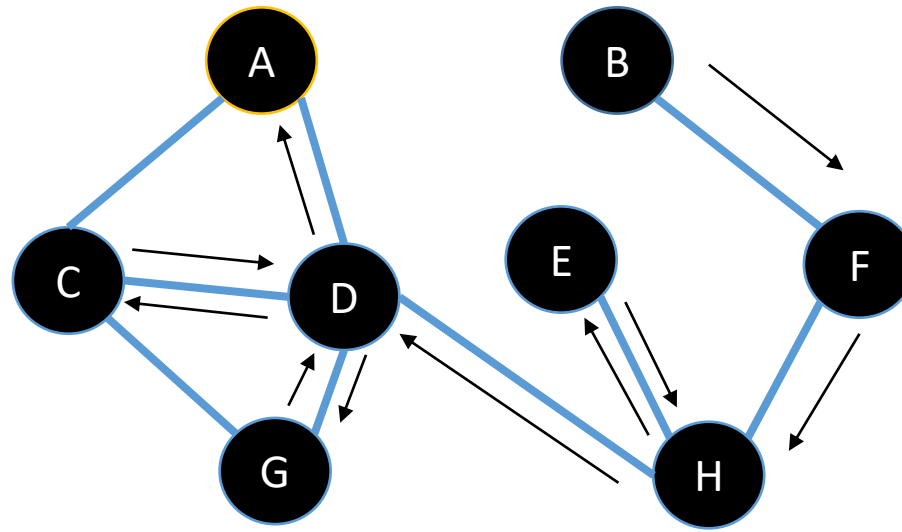
Depth-First Traversal

- Add its adjacent white nodes to the stack (there are none)
- Order visited: B, F, H, E, D, G, C



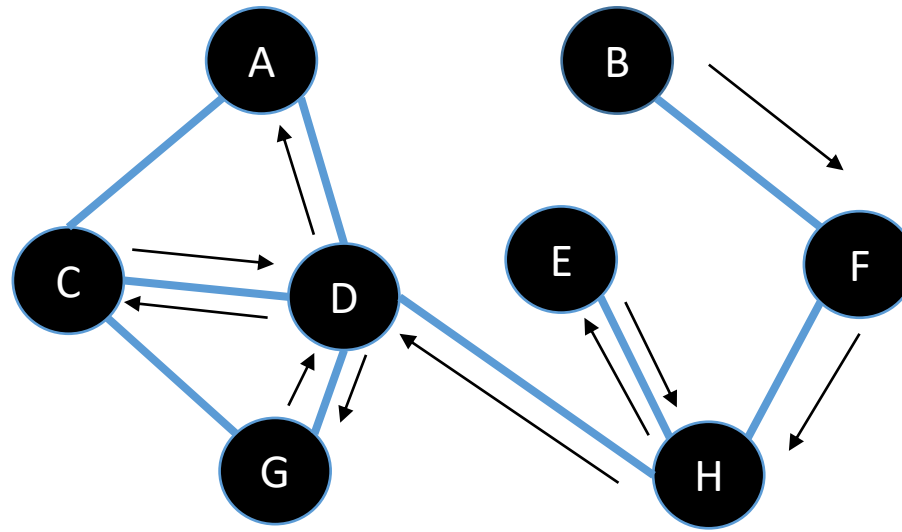
Depth-First Traversal

- A has been visited
- Order visited: B, F, H, E, D, G, C, A



Depth-First Traversal

- Next node is popped from the stack (there are none)
 - Traversal is complete when the stack is empty
- Order visited: **B, F, H, E, D, G, C, A**

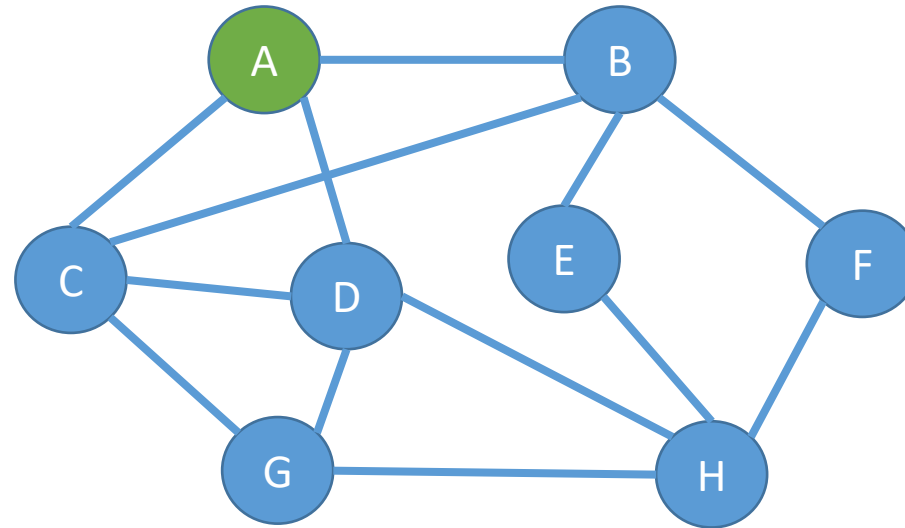


Breadth-First Traversal

- Breadth-first traversal of a graph is similar to the breadth-first traversal of a tree
- Uses a queue to manage the next nodes to visit
- Maintains an array of “seen” nodes to prevent getting stuck in a cycle, if one existed
 - The traversal process is the same for bi-directional graphs or digraphs

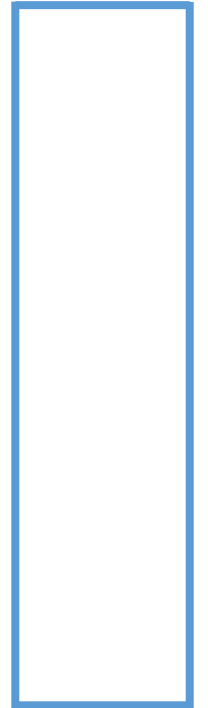
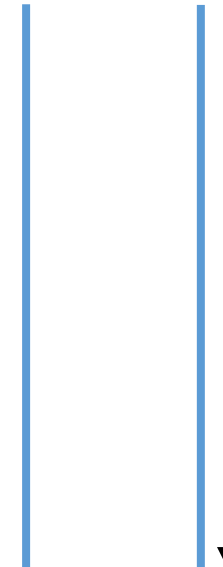
Breadth-First Traversal

- Starting at A
- Order visited: N/A



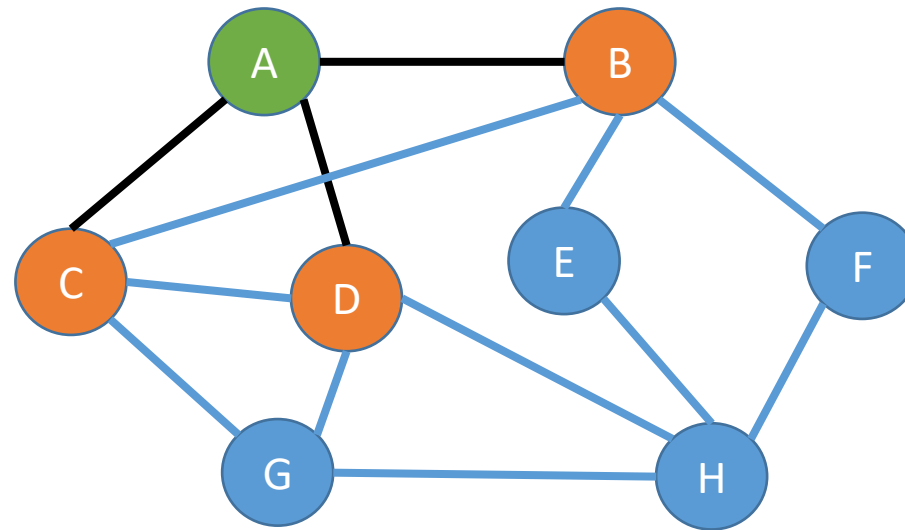
Queue

Nodes Seen

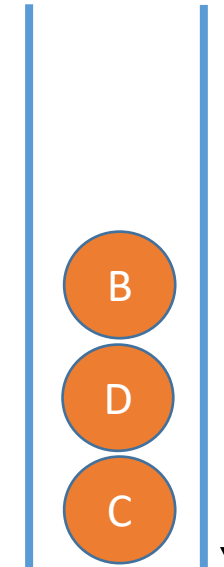


Breadth-First Traversal

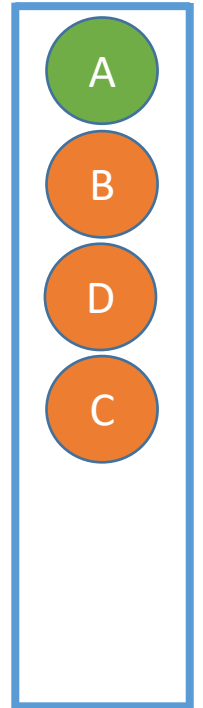
- A's unseen neighbors added to queue
 - A and unseen neighbors are added to nodes seen
- Order visited: A



Queue

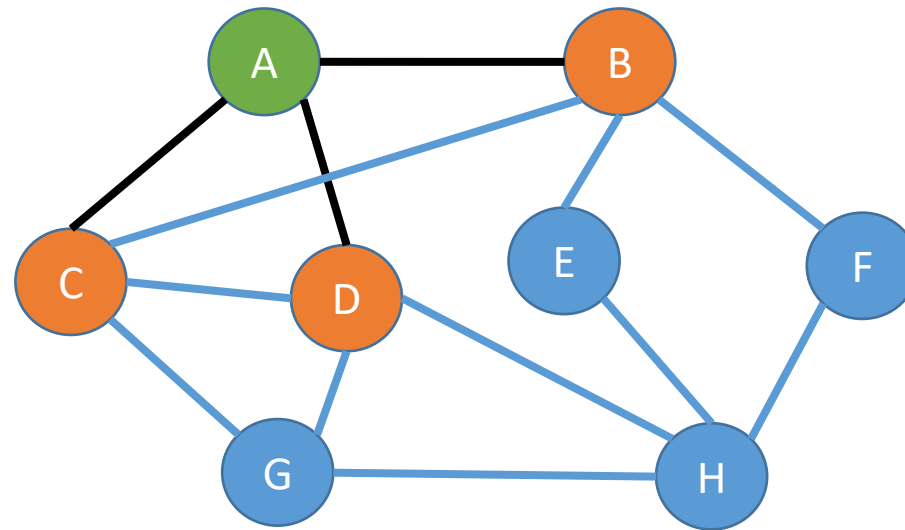


Nodes Seen

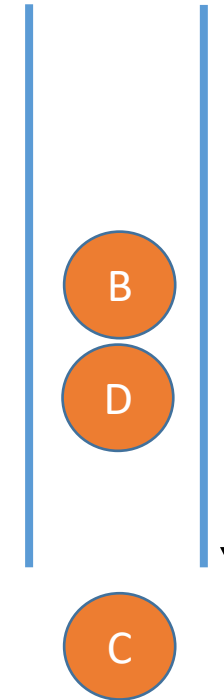


Breadth-First Traversal

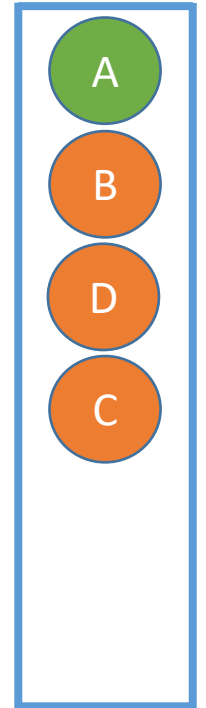
- C popped from the queue
- Order visited: A



Queue

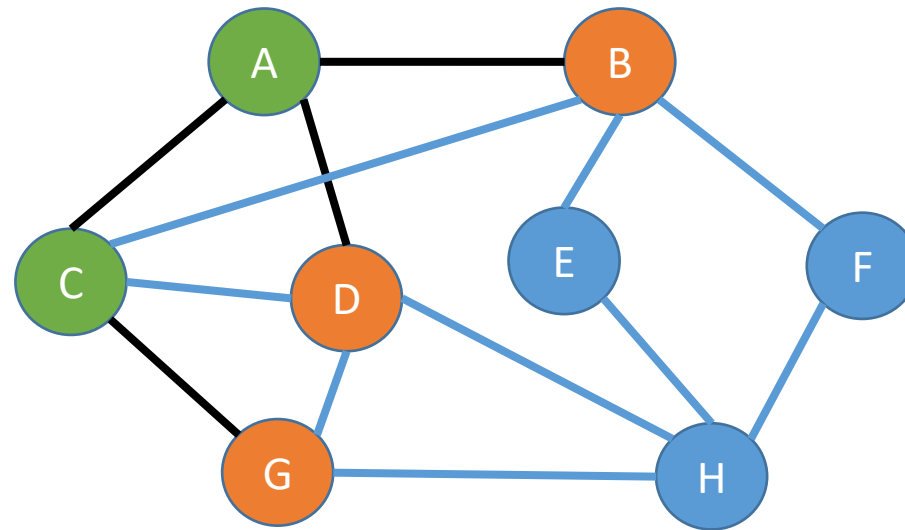


Nodes Seen

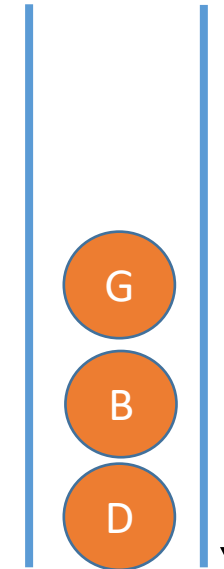


Breadth-First Traversal

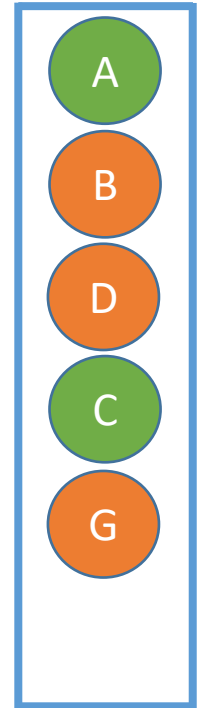
- C's unseen neighbors added to queue
 - Unseen neighbors are added to nodes seen
- Order visited: A, C



Queue

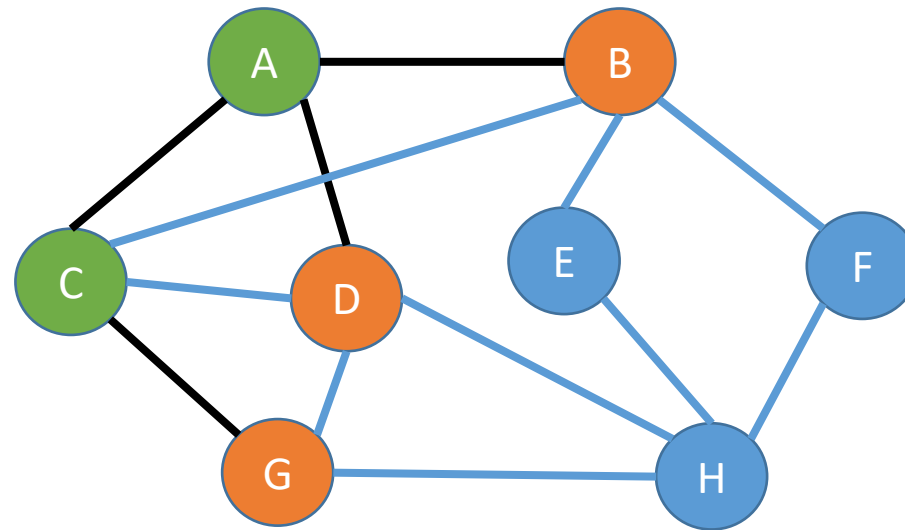


Nodes Seen

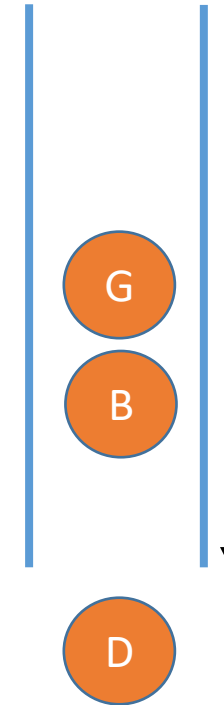


Breadth-First Traversal

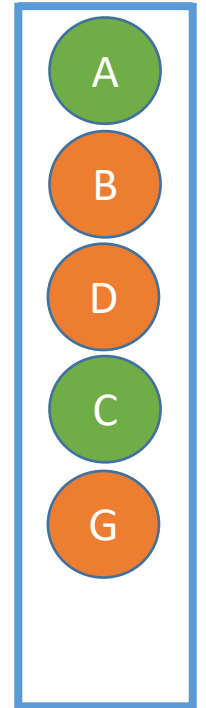
- D popped from the queue
- Order visited: A, C



Queue

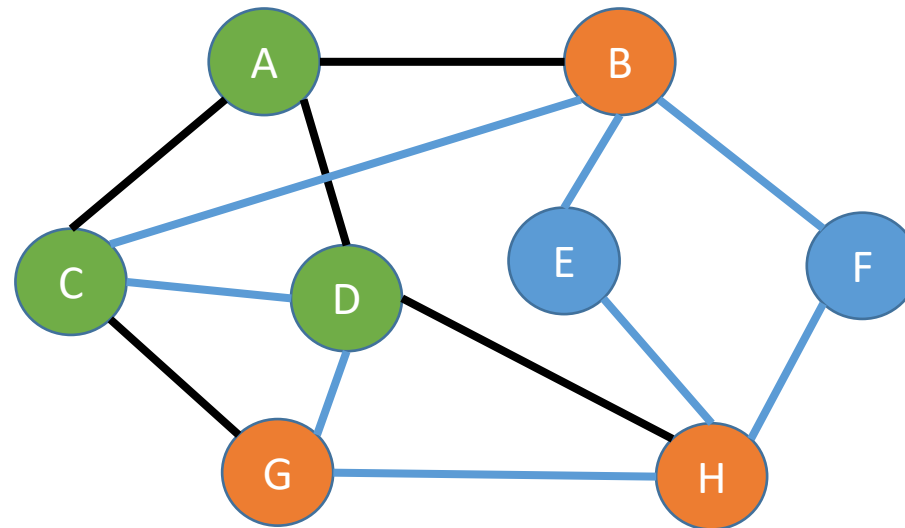


Nodes Seen

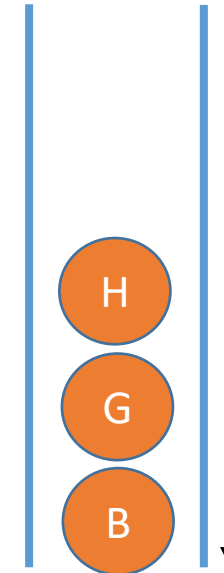


Breadth-First Traversal

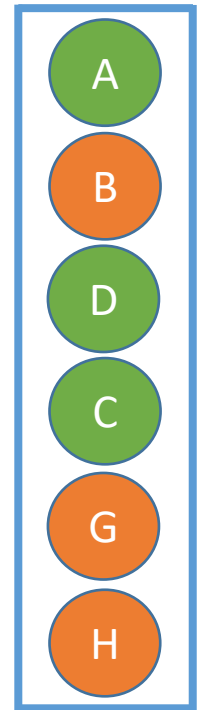
- D's unseen neighbors added to queue
 - Unseen neighbors are added to nodes seen
- Order visited: A, C, D



Queue

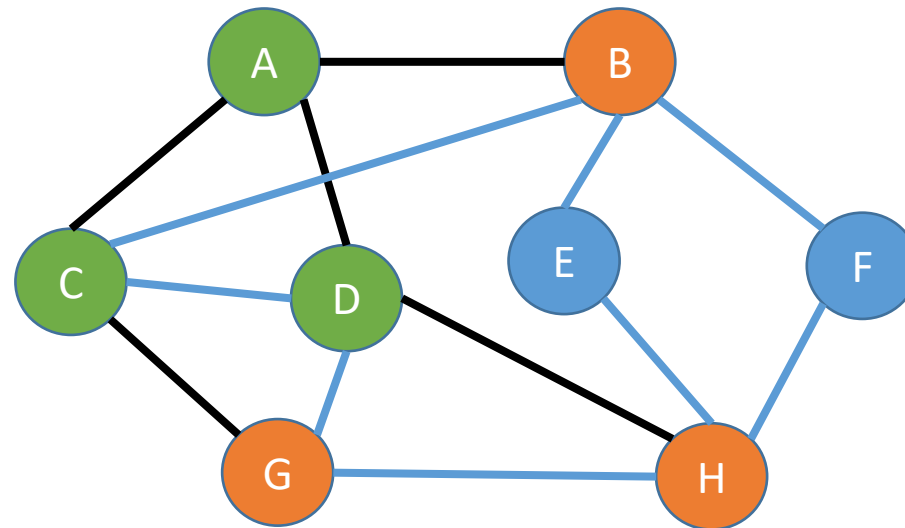


Nodes Seen

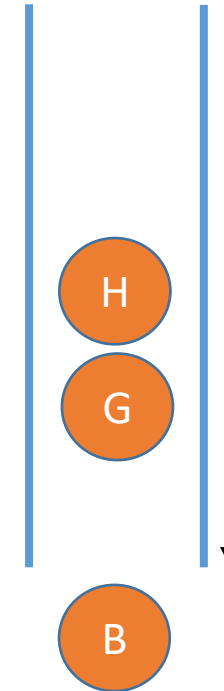


Breadth-First Traversal

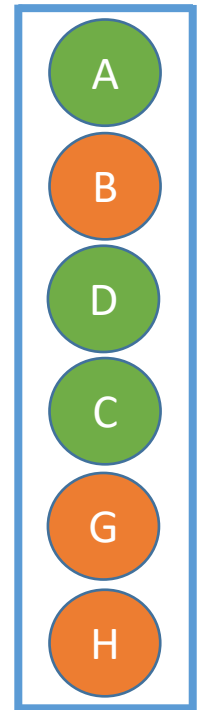
- B popped from the queue
- Order visited: A, C, D



Queue

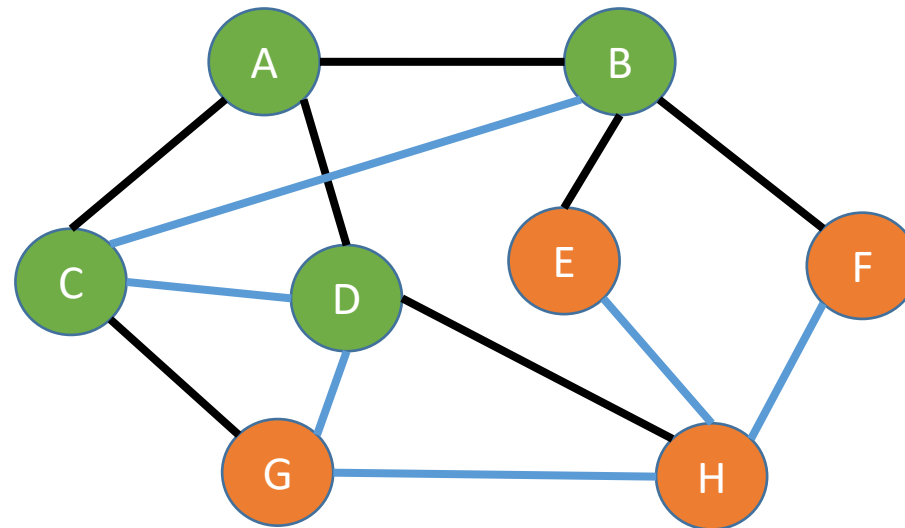


Nodes Seen

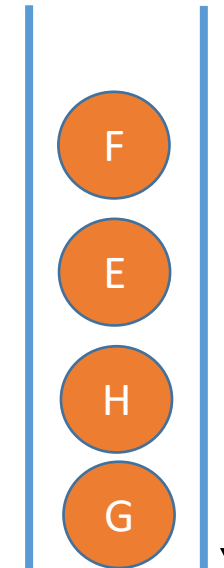


Breadth-First Traversal

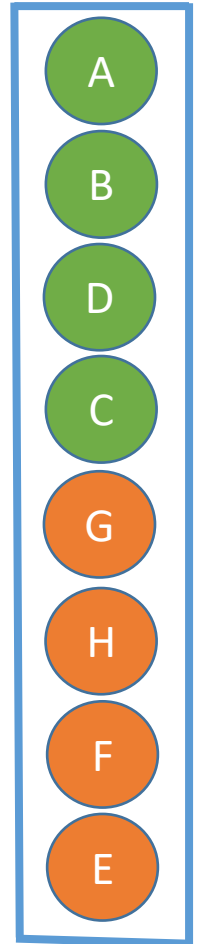
- B's unseen neighbors added to queue
 - Unseen neighbors are added to nodes seen
- Order visited: A, C, D, B



Queue

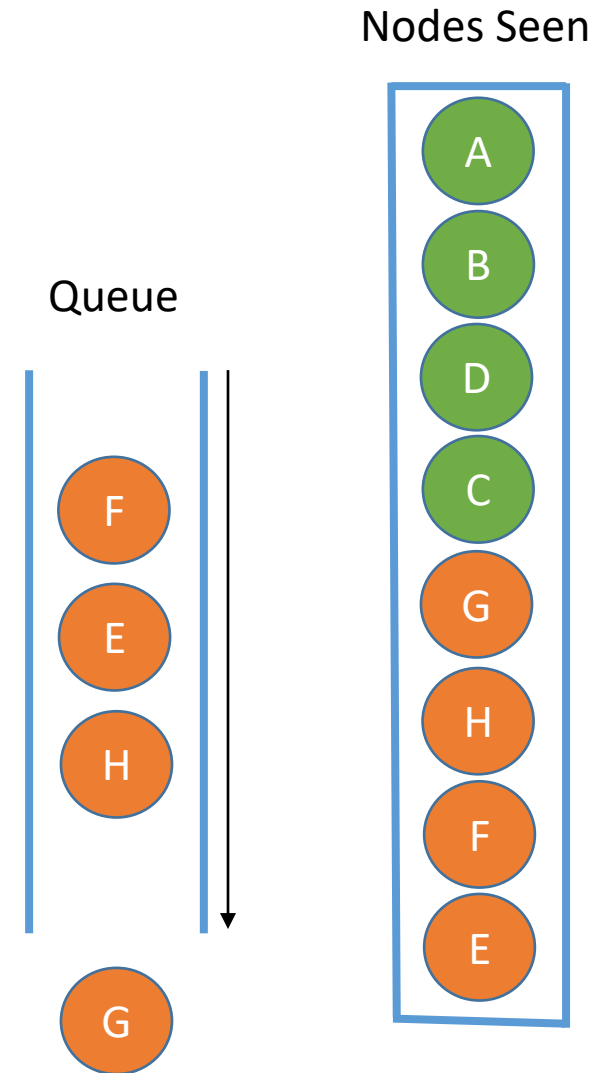
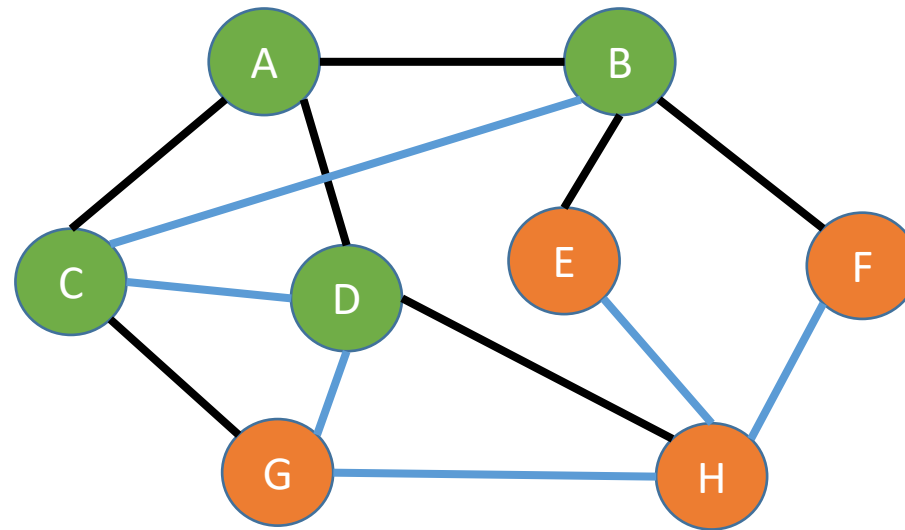


Nodes Seen



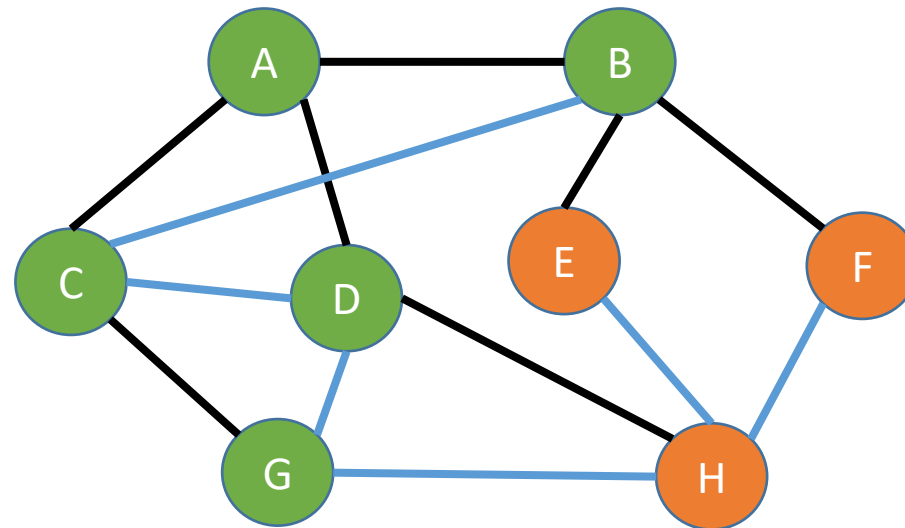
Breadth-First Traversal

- G popped from the queue
- Order visited: A, C, D, B

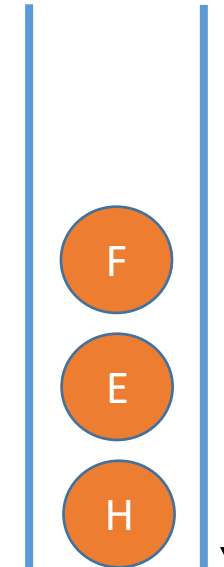


Breadth-First Traversal

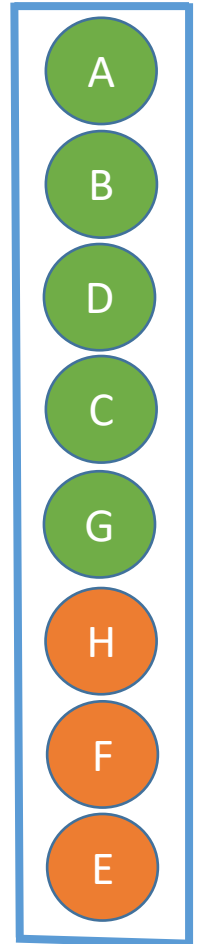
- G's unseen neighbors added to queue (none)
 - Unseen neighbors are added to nodes seen
- Order visited: A, C, D, B, G



Queue

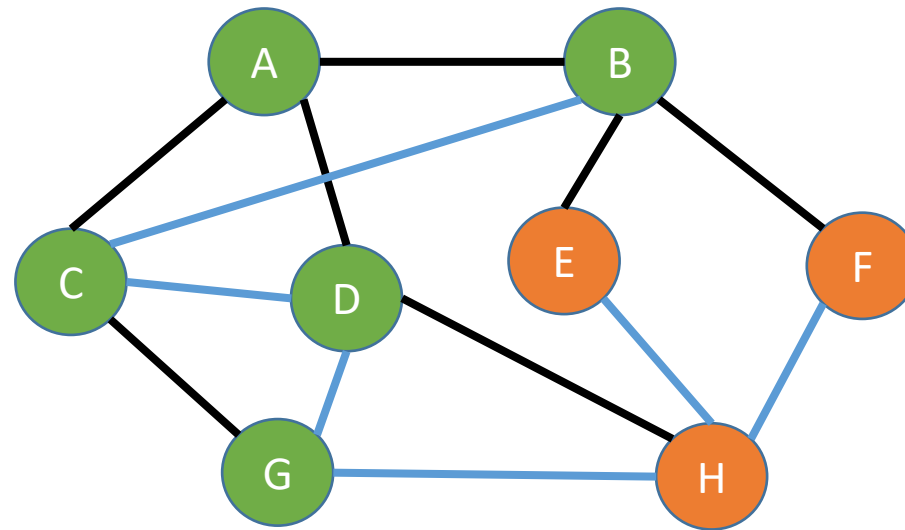


Nodes Seen

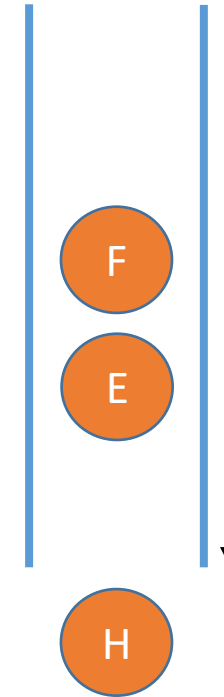


Breadth-First Traversal

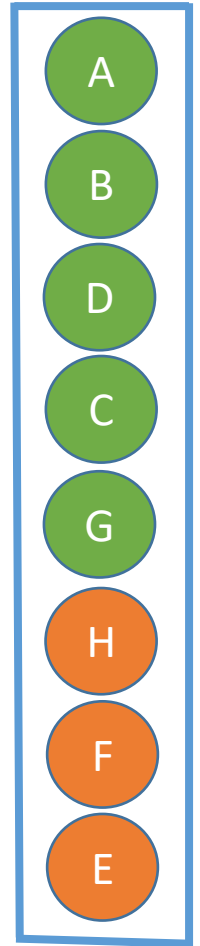
- H popped from the queue
- Order visited: A, C, D, B, G



Queue

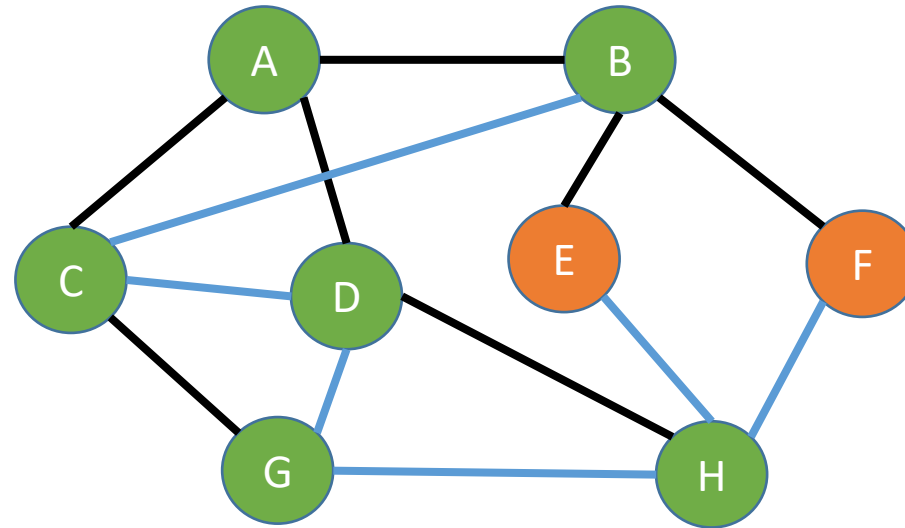


Nodes Seen

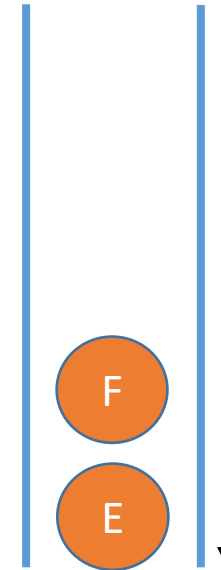


Breadth-First Traversal

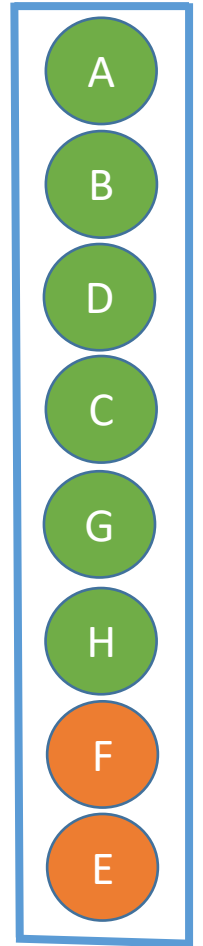
- H's unseen neighbors added to queue (none)
 - Unseen neighbors are added to nodes seen
- Order visited: A, C, D, B, G, H



Queue

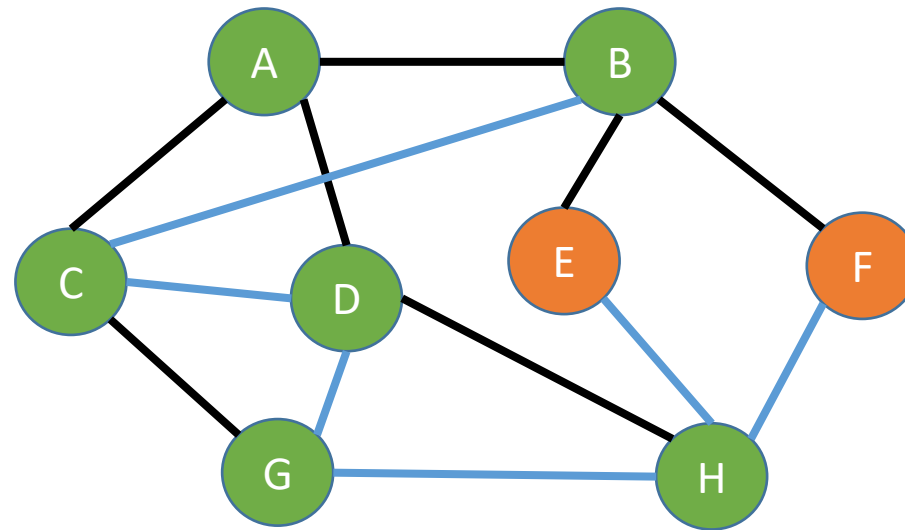


Nodes Seen

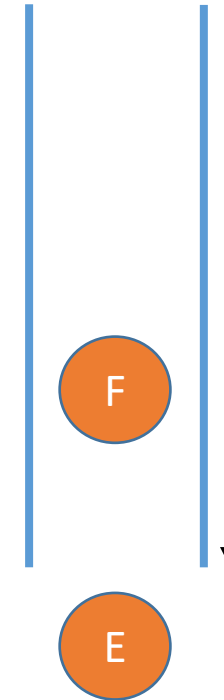


Breadth-First Traversal

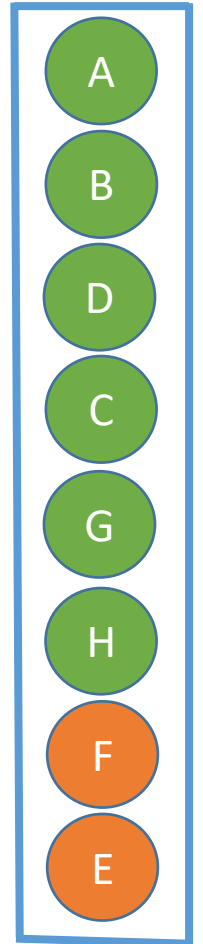
- E popped from the queue
- Order visited: A, C, D, B, G, H



Queue

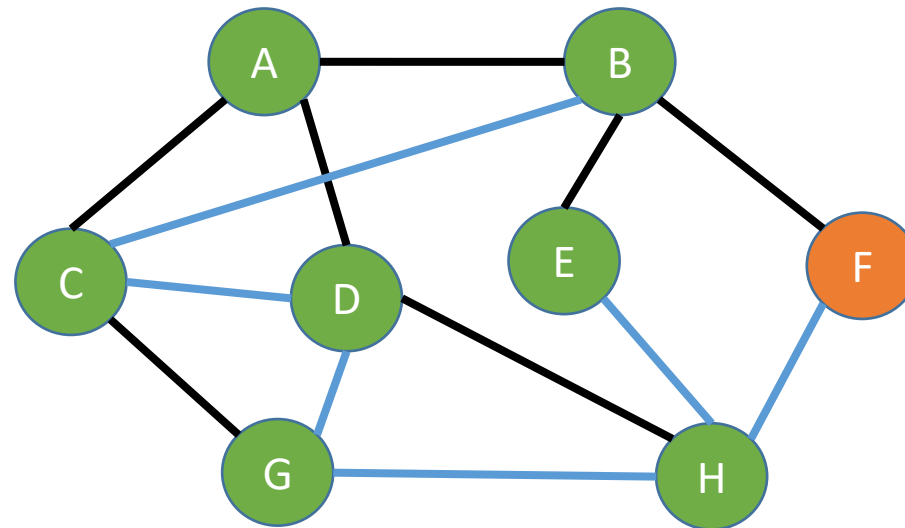


Nodes Seen

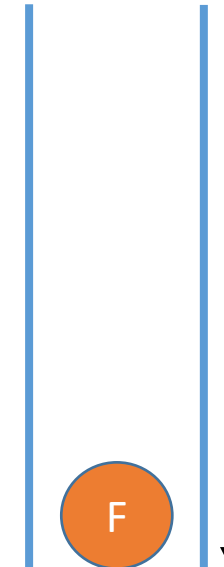


Breadth-First Traversal

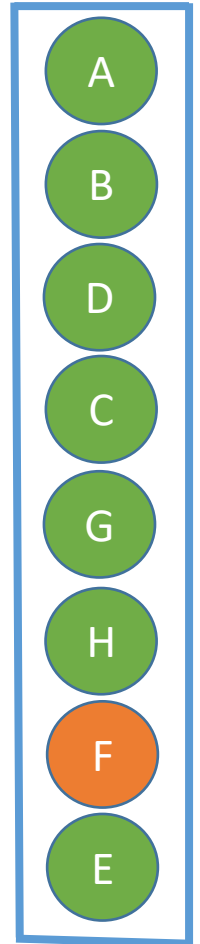
- E's unseen neighbors added to queue (none)
 - Unseen neighbors are added to nodes seen
- Order visited: A, C, D, B, G, H, E



Queue

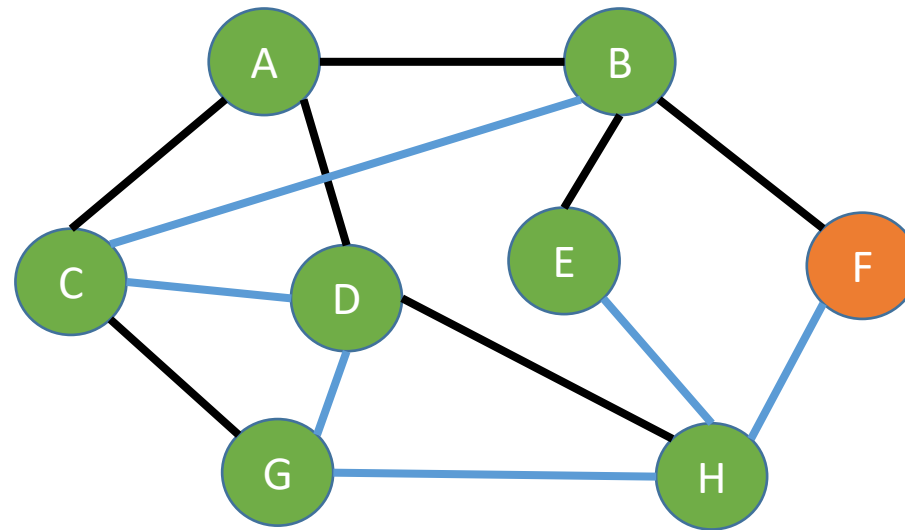


Nodes Seen

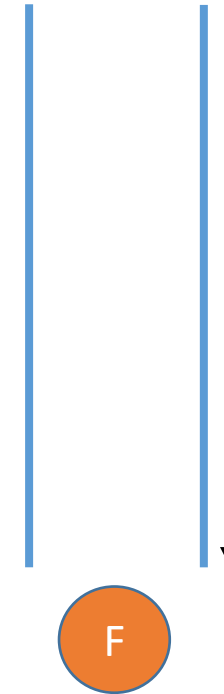


Breadth-First Traversal

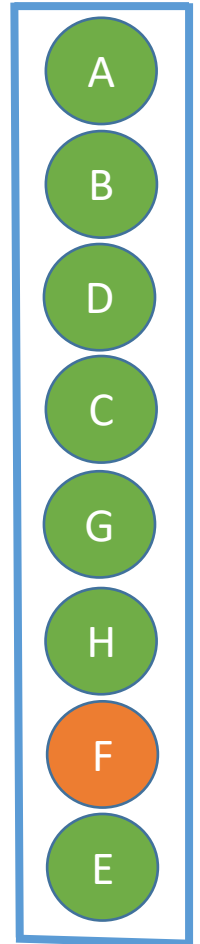
- F popped from queue
- Order visited: A, C, D, B, G, H, E



Queue

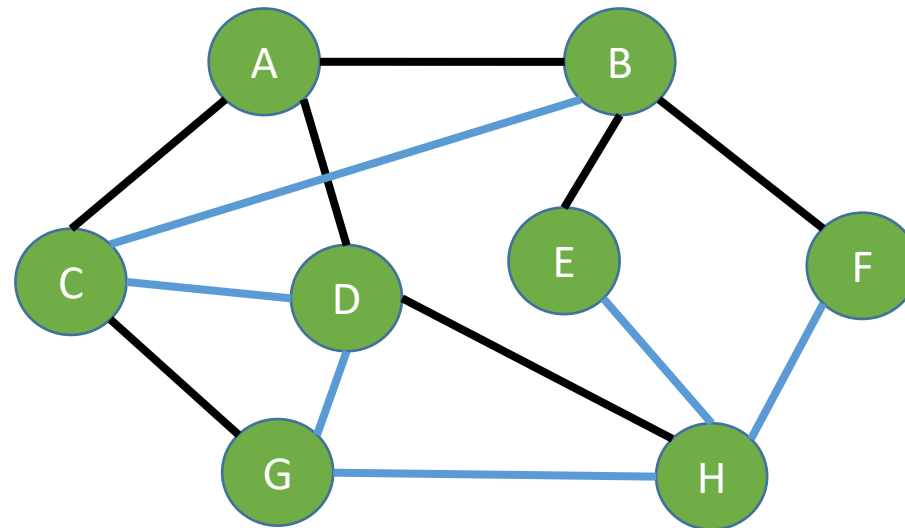


Nodes Seen

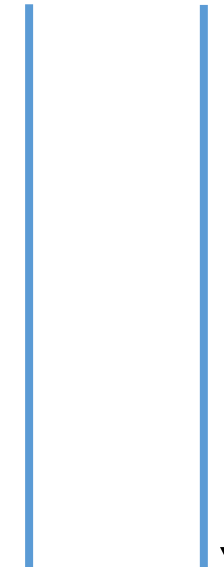


Breadth-First Traversal

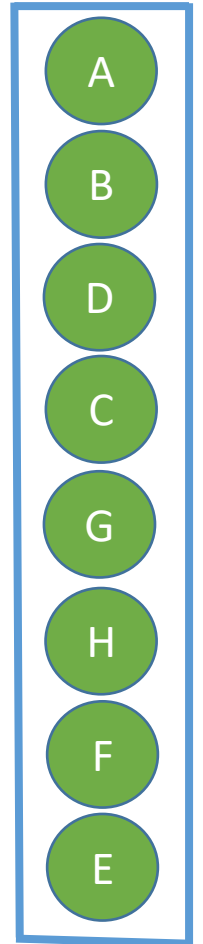
- F's unseen neighbors added to queue (none)
 - Unseen neighbors are added to nodes seen
- Order visited: A, C, D, B, G, H, E, F



Queue

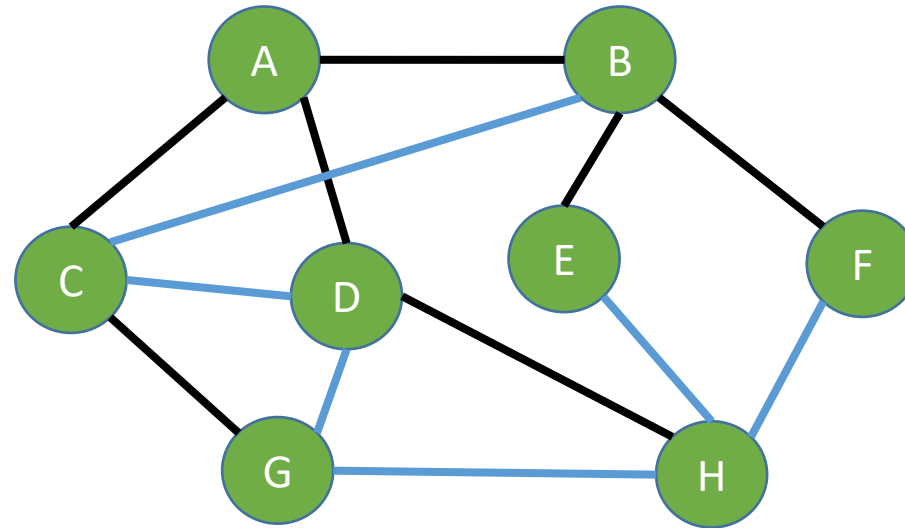


Nodes Seen



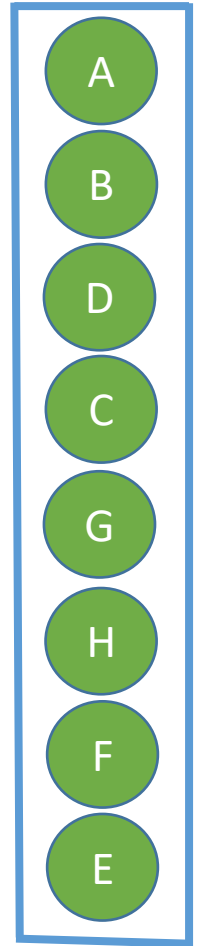
Breadth-First Traversal

- Queue is empty
 - Traversal is complete
- Order visited: **A, C, D, B, G, H, E, F**



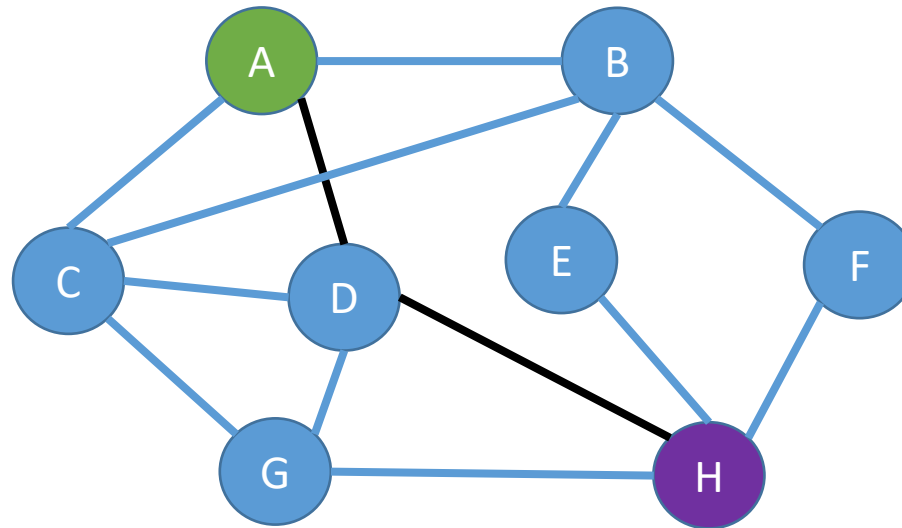
Queue

Nodes Seen



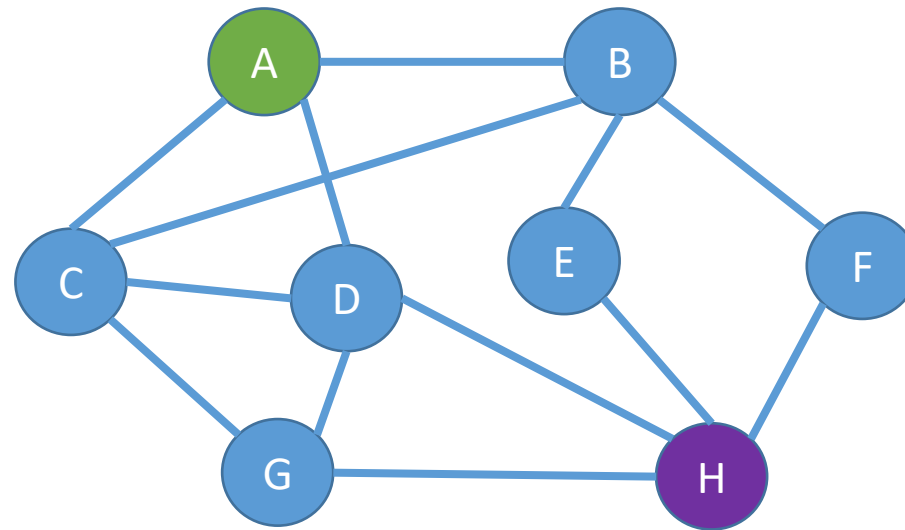
Finding Distance with Breadth-First

- Based on the last example, its easy to see a breadth-first traversal can be used to find the distance (shortest path) between two nodes
- As an example, we'll use a breadth-first traversal to find the distance between A and H



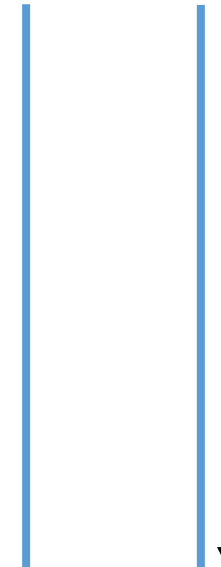
Finding Distance with Breadth-First

- Starting at A
- Order visited: N/A



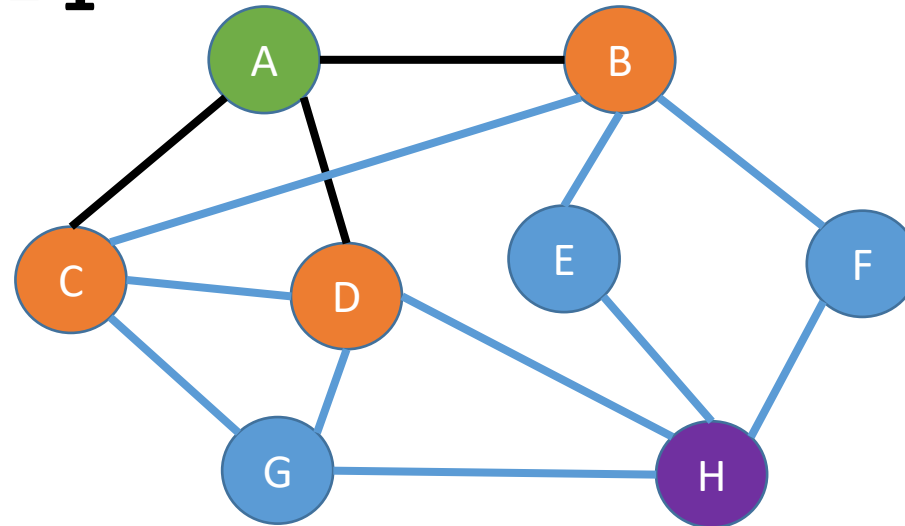
Queue

Nodes Seen

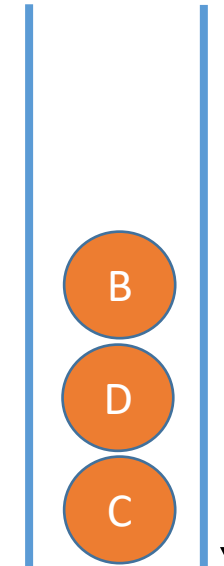


Finding Distance with Breadth-First

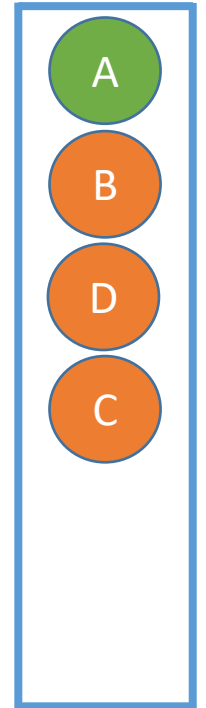
- A's unseen neighbors added to queue
 - A and unseen neighbors are added to nodes seen
- Order visited: A
- **DEEPEST LEVEL = 1**



Queue

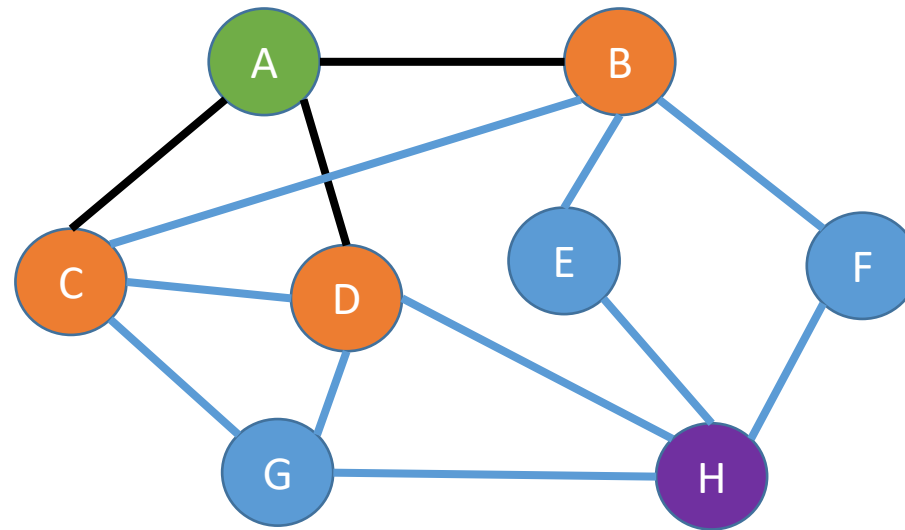


Nodes Seen

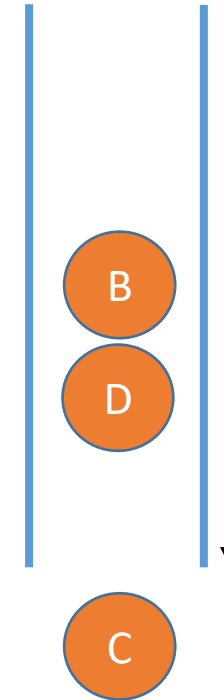


Finding Distance with Breadth-First

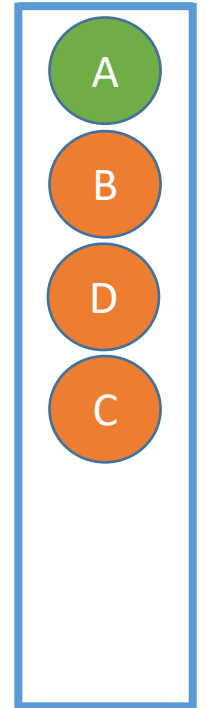
- C popped from the queue
- Order visited: A
- **DEEPEST LEVEL = 1**



Queue

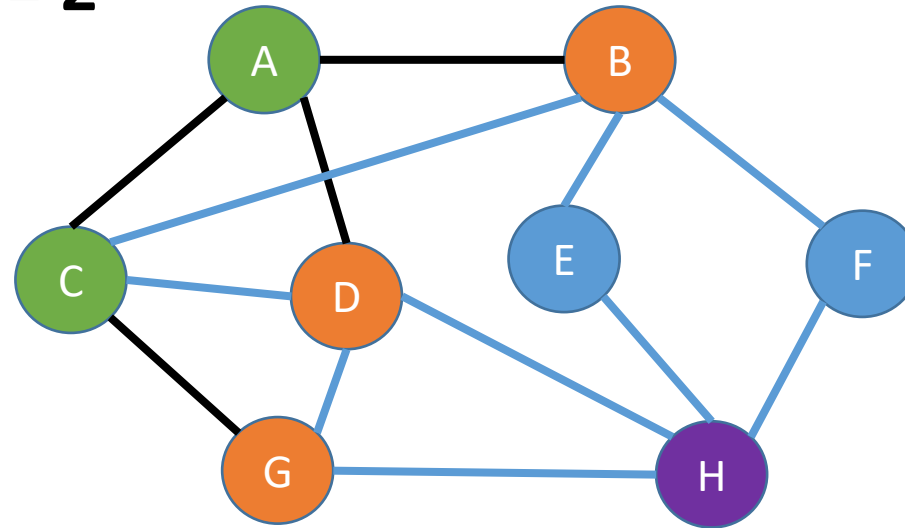


Nodes Seen

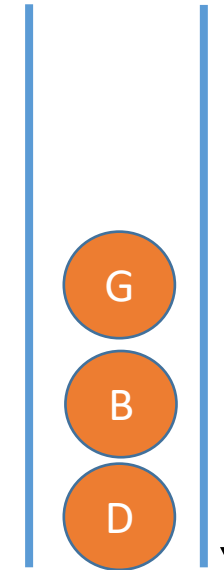


Finding Distance with Breadth-First

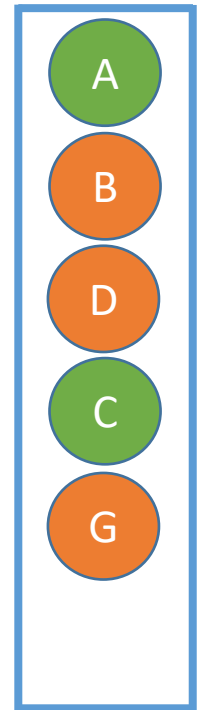
- C's unseen neighbors added to queue
 - Unseen neighbors are added to nodes seen
- Order visited: A, C
- **DEEPEST LEVEL = 2**



Queue

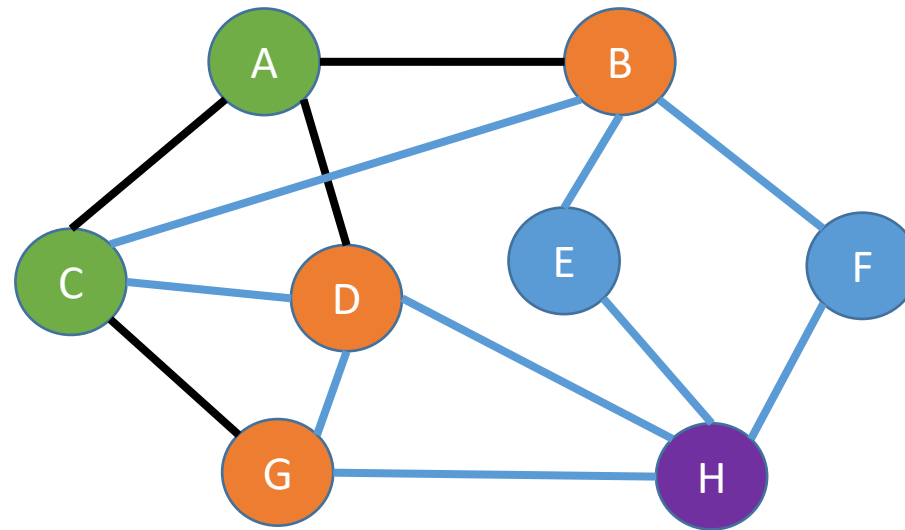


Nodes Seen

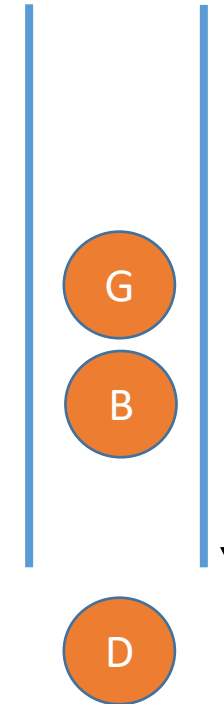


Finding Distance with Breadth-First

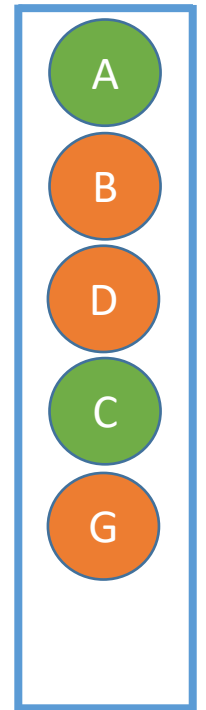
- D popped from the queue
- Order visited: A, C
- **DEEPEST LEVEL = 2**



Queue

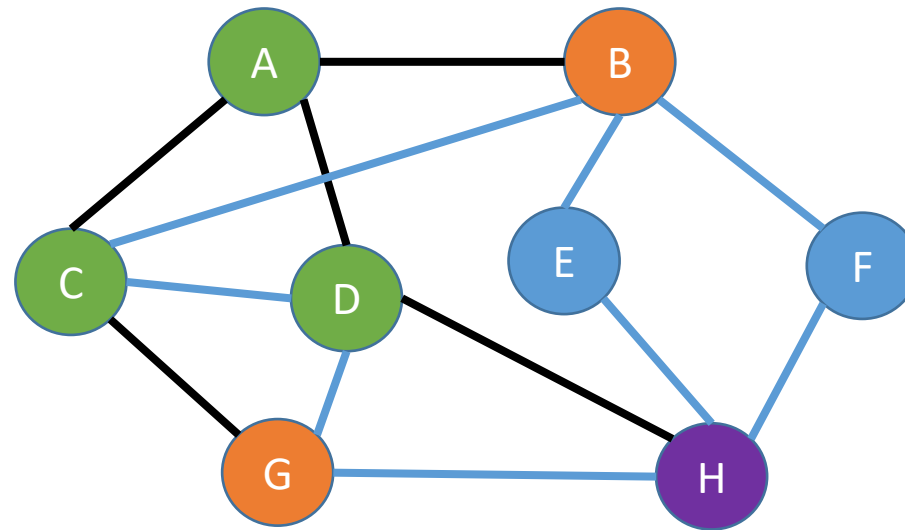


Nodes Seen

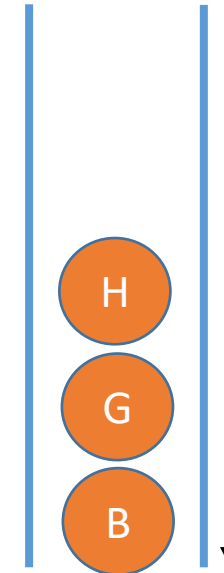


Finding Distance with Breadth-First

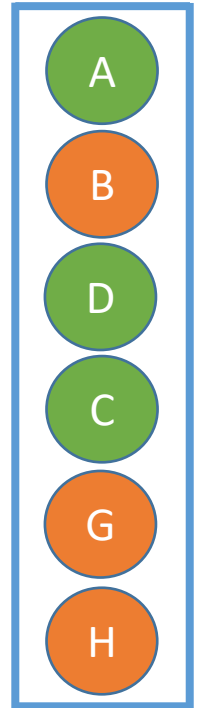
- D's unseen neighbors added to queue
 - H was reached
 - Traversal stops
- **DEEPEST LEVEL = 2**
 - **DISTANCE = 2**



Queue

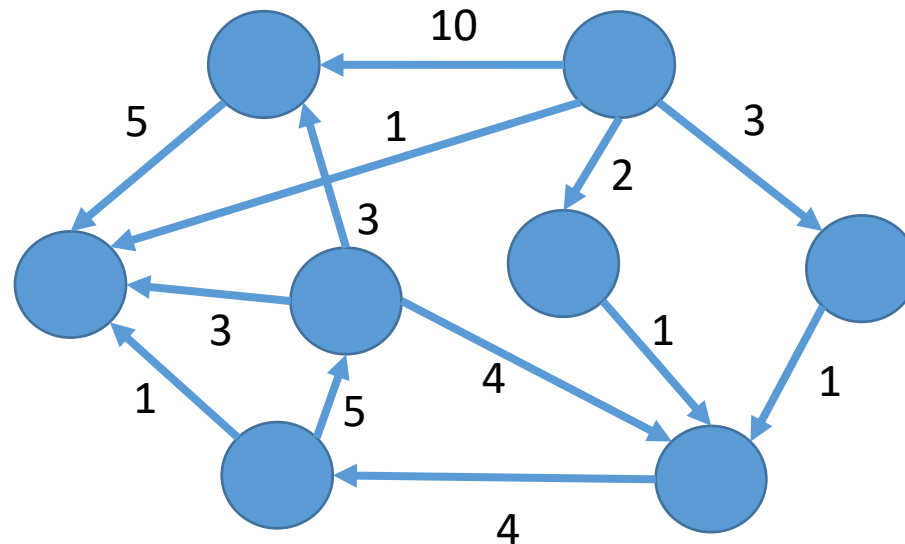


Nodes Seen



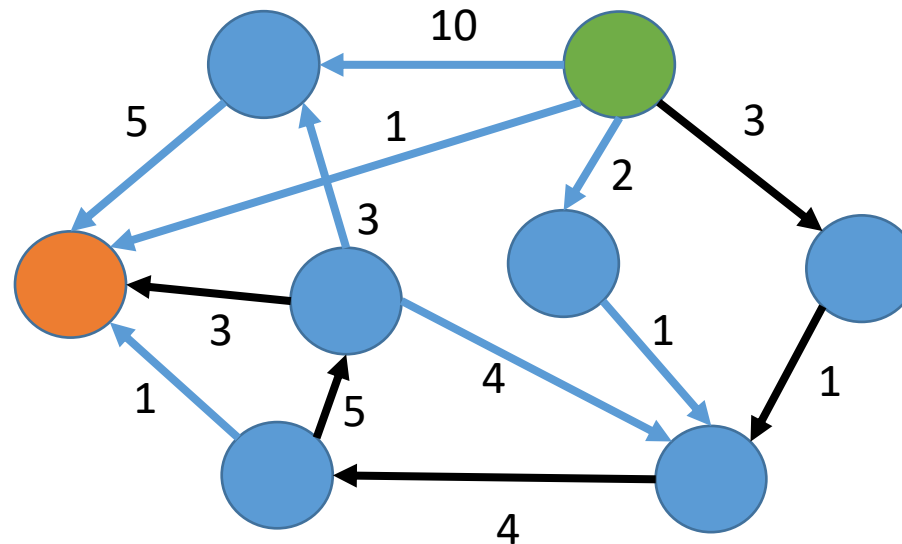
Weighted Graphs

- A **weighted graph** is a graph where each edge has a weight or cost.
 - Weighted graphs can be undirected or directed

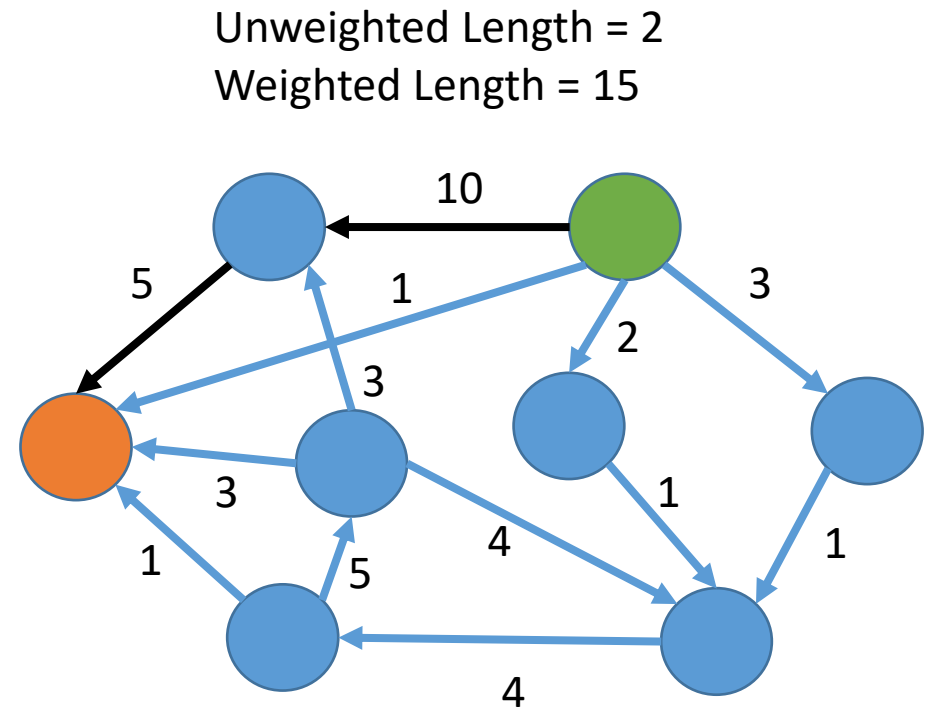
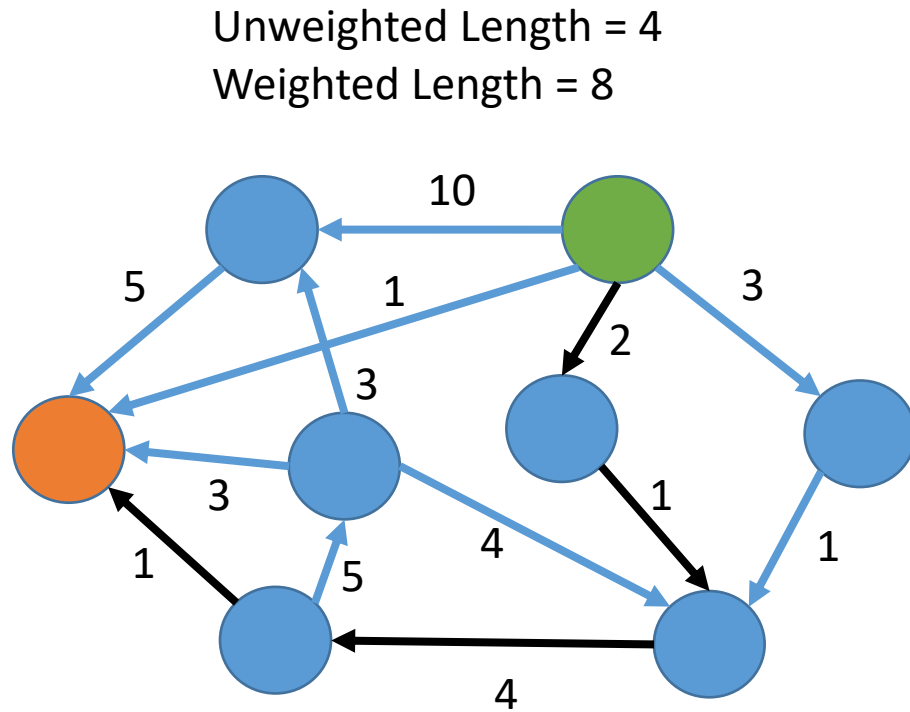


Weighted Graphs

- The path length of a weighted graph is the sum of the edge costs.
 - $3+1+4+5+3 = \mathbf{16}$



Weighted Graphs



- The first path is less costly than the second path, despite it being twice as long

Weighted Graphs

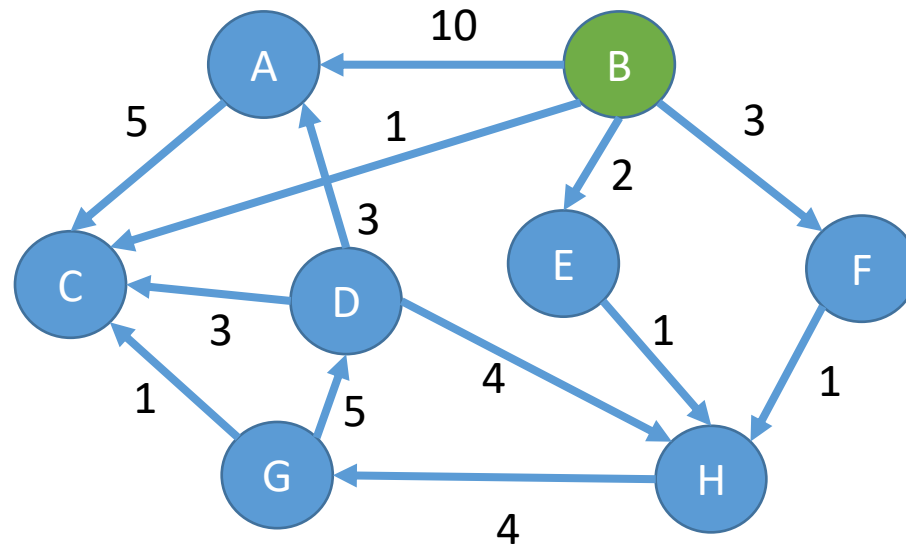
- A breadth-first traversal would not be useful for finding the path with the least cost
- Other algorithms are used to find the path with the least cost between two nodes
- The most well known is Dijkstra's Algorithm

Dijkstra's Algorithm

- This algorithm finds the shortest path between one node and **every other node** in a graph.
- *For each vertex*
 - The algorithm determines the vertex's distance (shortest/least costly path) from the starting vertex
 - The algorithm determines the vertex's predecessor pointer- the previous vertex with the shortest (or least costly) path from the starting vertex
- Can be used on:
 - Bi-directional and digraphs
 - Weighted graphs and unweighted graphs

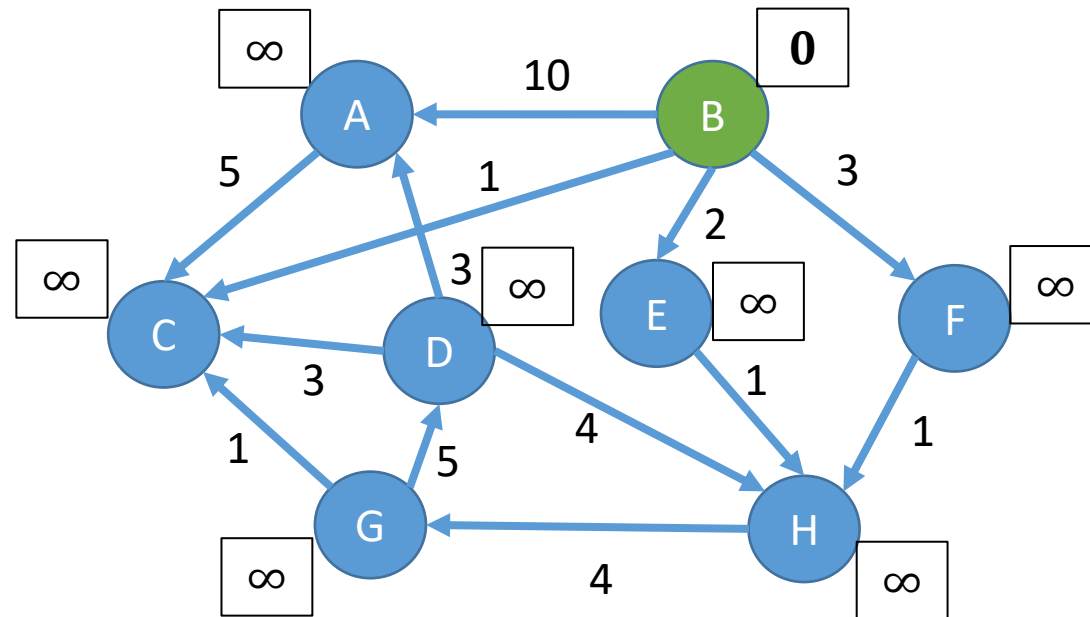
Dijkstra's Algorithm

- We can start with any node, but we'll start with node B since a path exists from B to all other nodes



Dijkstra's Algorithm

- We'll remember the cost from each node back to node B.
- B has a cost of 0; All other nodes are assumed to have a cost of infinity
 - Ensures the path found will be less than that

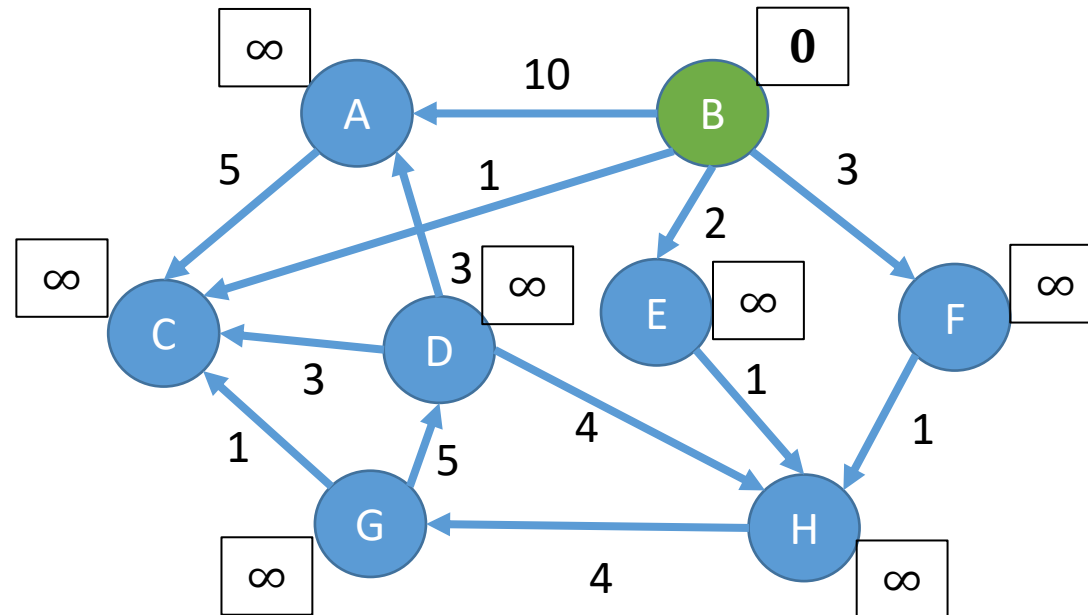
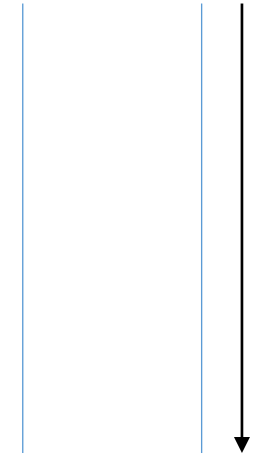


A cost = ∞
B cost = 0
C cost = ∞
D cost = ∞
E cost = ∞
F cost = ∞
G cost = ∞
H cost = ∞

Dijkstra's Algorithm

- A priority queue or min-heap is used to prioritize nodes with lower costs to be visited first

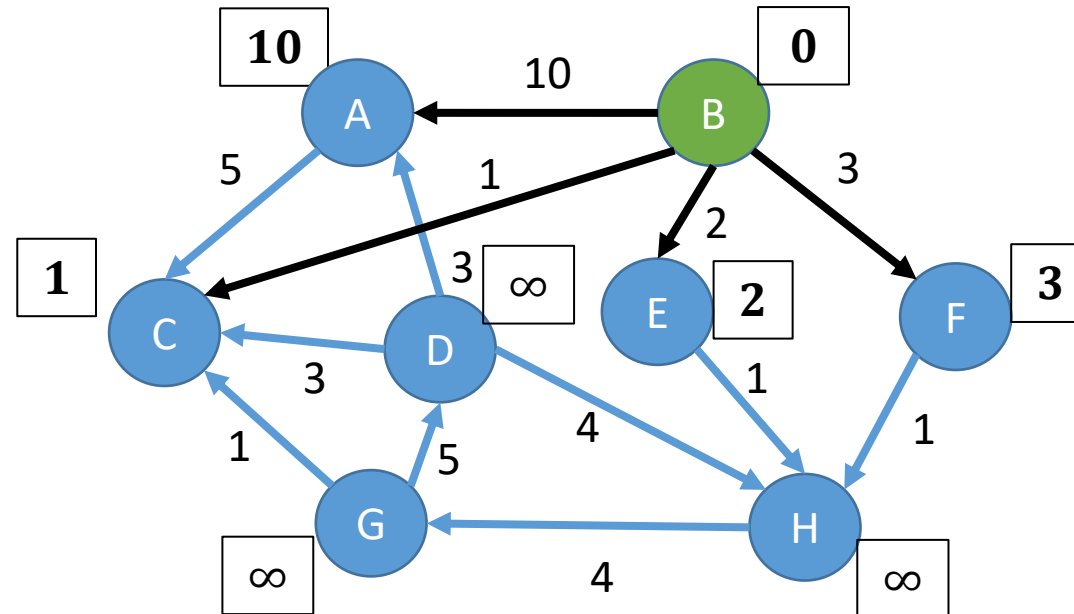
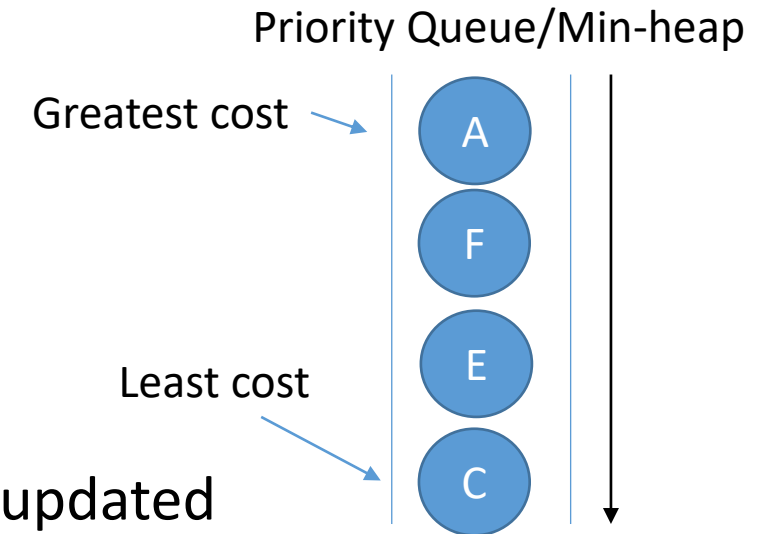
Priority Queue/Min-heap



A cost = ∞
B cost = 0
C cost = ∞
D cost = ∞
E cost = ∞
F cost = ∞
G cost = ∞
H cost = ∞

Dijkstra's Algorithm

- We'll now look at the nodes adjacent to B.
- Cost of each is B's cost + cost of the edge
 - All of which are less than infinity so the costs are updated

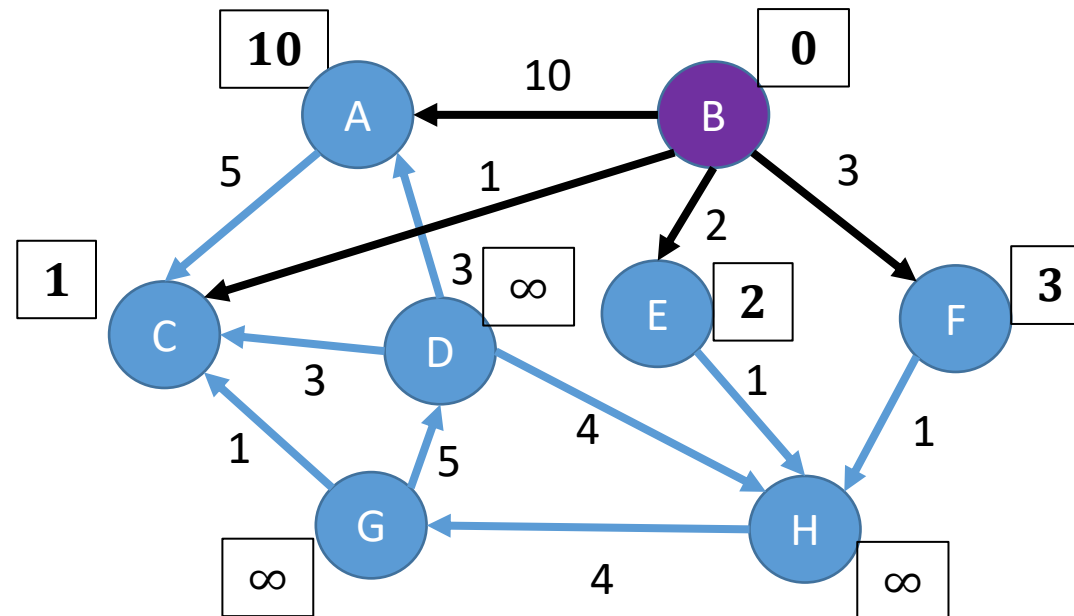
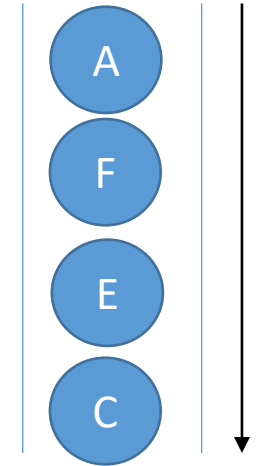


A cost = $0 + 10 = 10$
B cost = 0
C cost = $0 + 1 = 1$
D cost = ∞
E cost = $0 + 2 = 2$
F cost = $0 + 3 = 3$
G cost = ∞
H cost = ∞

Dijkstra's Algorithm

- We are finished with B

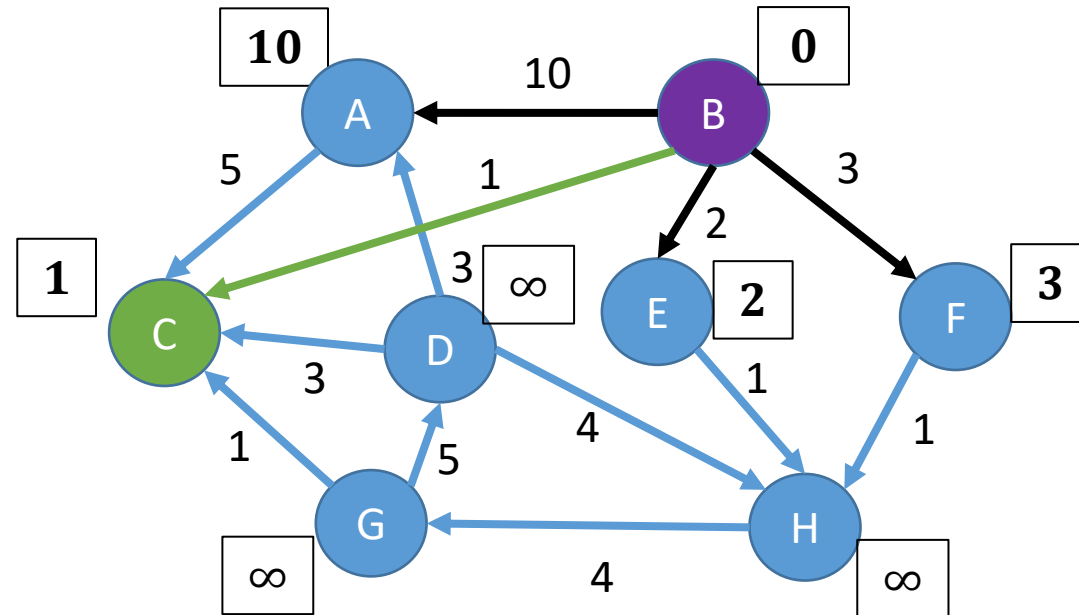
Priority Queue/Min-heap



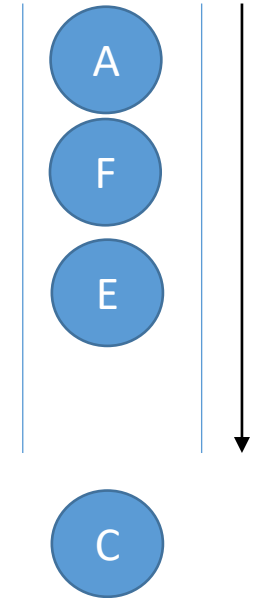
A cost = 10
B cost = 0
C cost = 1
D cost = ∞
E cost = 2
F cost = 3
G cost = ∞
H cost = ∞

Dijkstra's Algorithm

- We move on to the next node in the priority queue



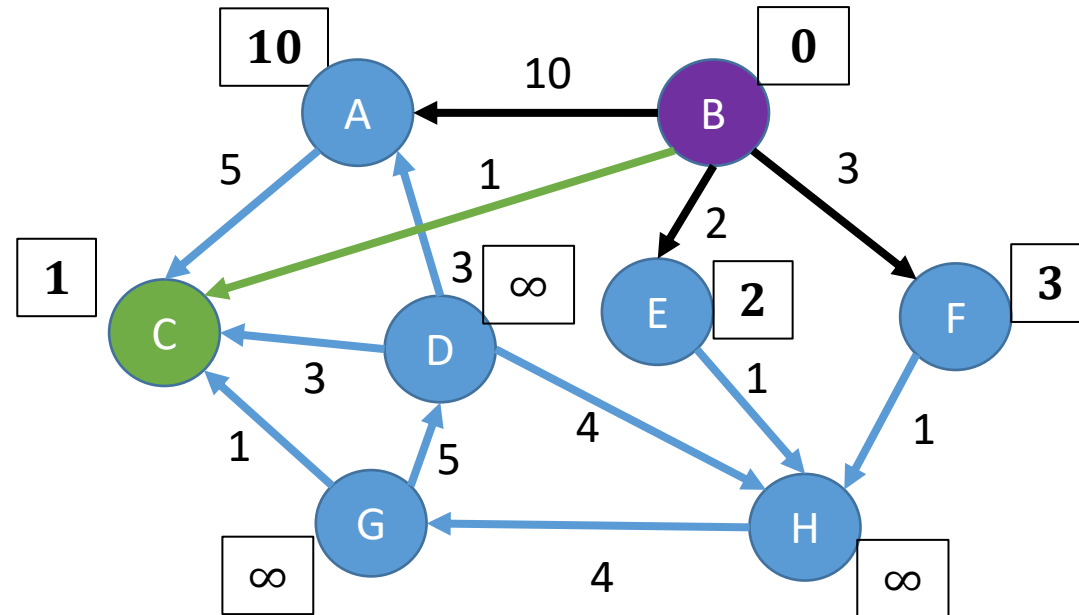
Priority Queue/Min-heap



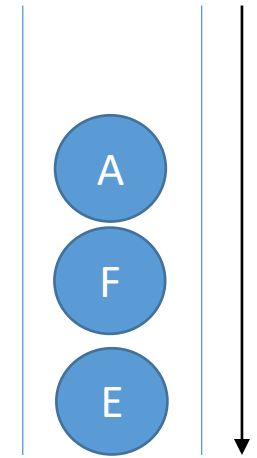
A cost = 10
B cost = 0
C cost = 1
D cost = ∞
E cost = 2
F cost = 3
G cost = ∞
H cost = ∞

Dijkstra's Algorithm

- We'll now look at the nodes adjacent to C.
 - There are none



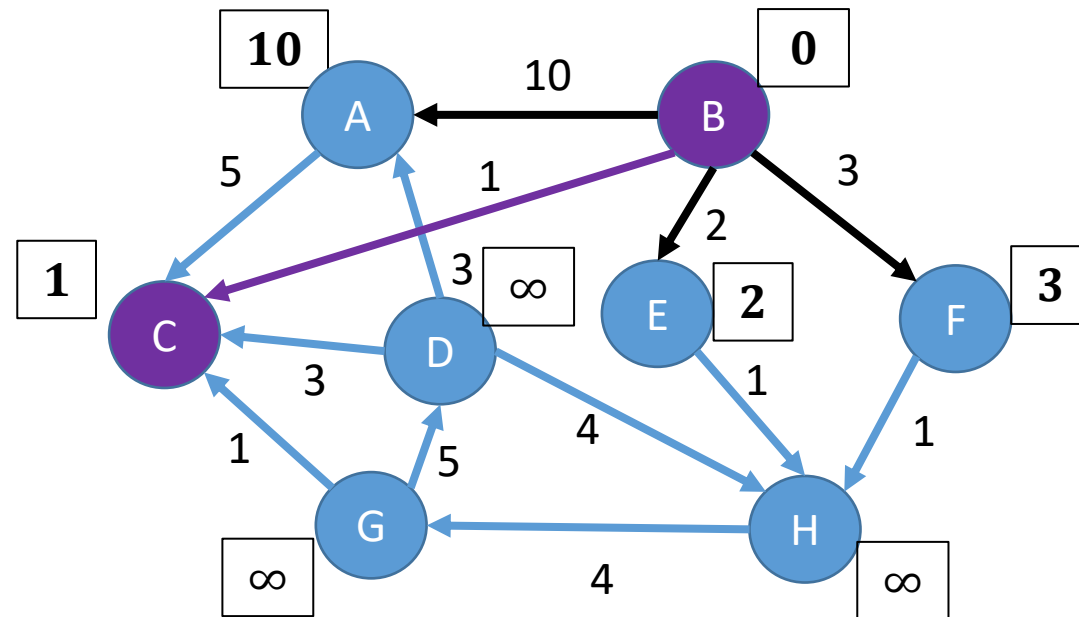
Priority Queue/Min-heap



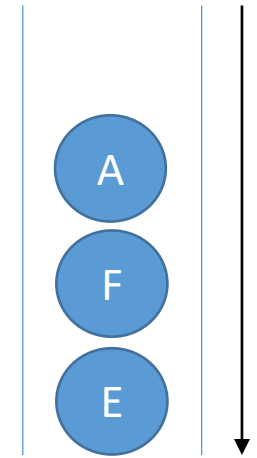
A cost = 10
B cost = 0
C cost = 1
D cost = ∞
E cost = 2
F cost = 3
G cost = ∞
H cost = ∞

Dijkstra's Algorithm

- We are finished with C



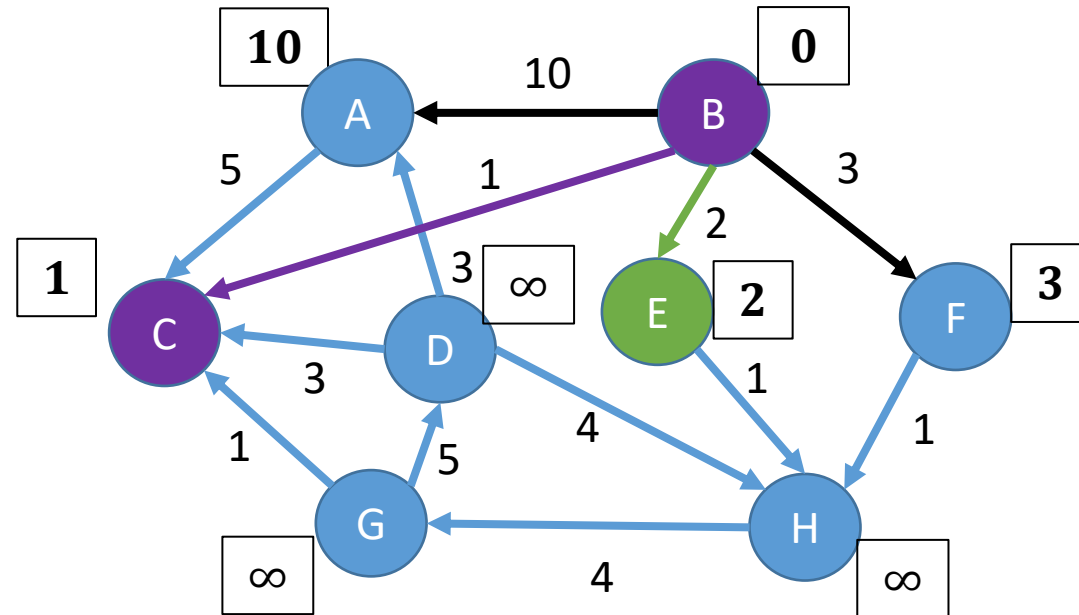
Priority Queue/Min-heap



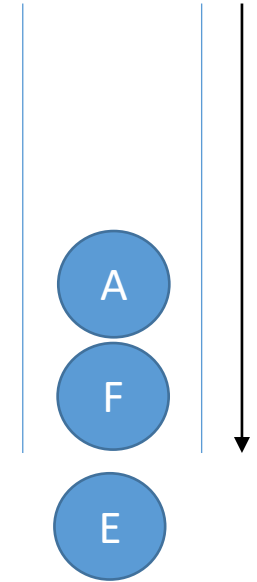
A cost = 10
B cost = 0
C cost = 1
D cost = ∞
E cost = 2
F cost = 3
G cost = ∞
H cost = ∞

Dijkstra's Algorithm

- We move on to the next node in the priority queue



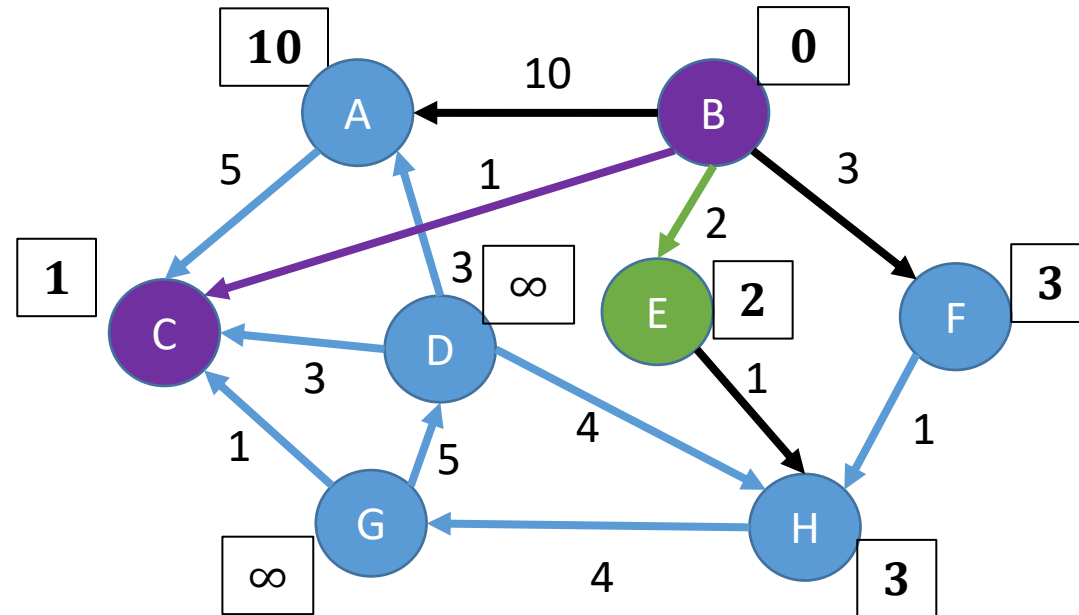
Priority Queue/Min-heap



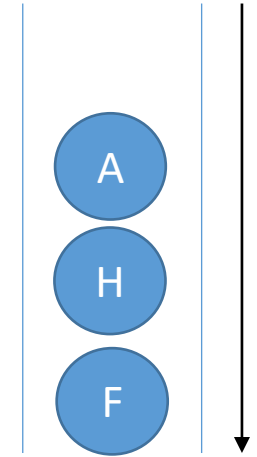
A cost = 10
B cost = 0
C cost = 1
D cost = ∞
E cost = 2
F cost = 3
G cost = ∞
H cost = ∞

Dijkstra's Algorithm

- We'll now look at the nodes adjacent to E.
 - Cost of each is E's cost + cost of the edge



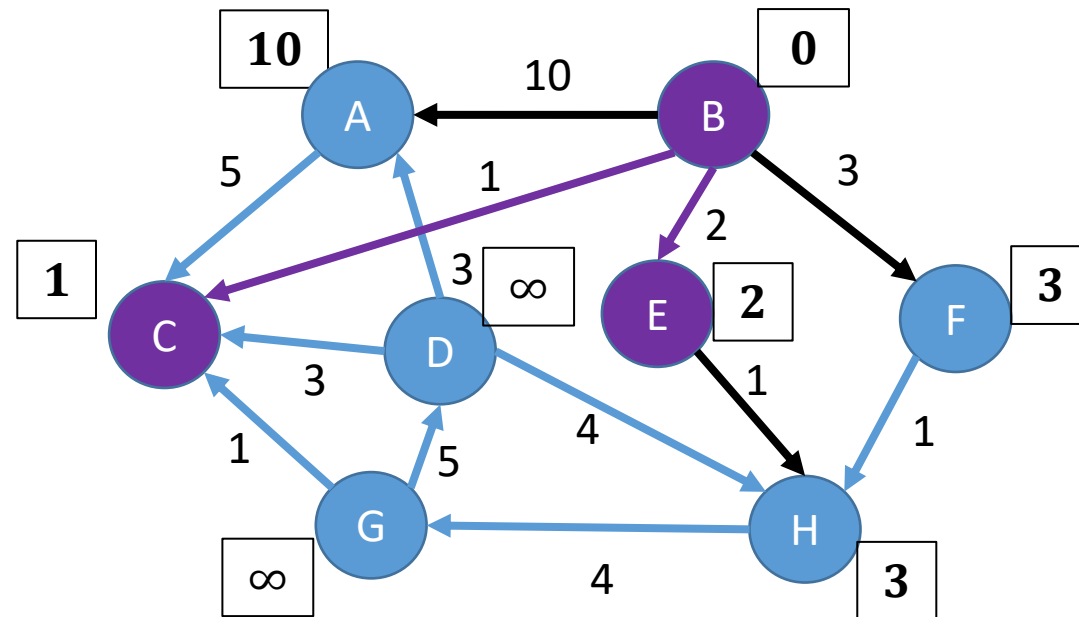
Priority Queue/Min-heap



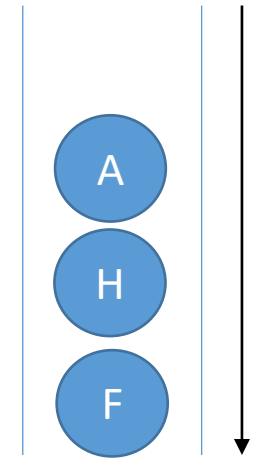
A cost = 10
B cost = 0
C cost = 1
D cost = ∞
E cost = 2
F cost = 3
G cost = ∞
H cost = $2 + 1 = 3$

Dijkstra's Algorithm

- We are finished with E.



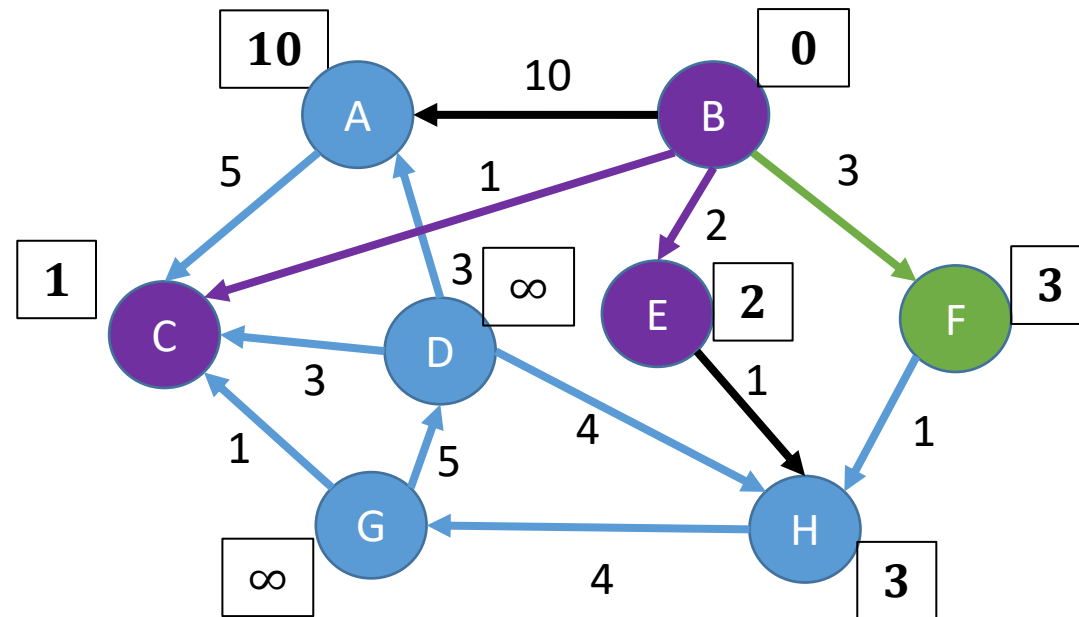
Priority Queue/Min-heap



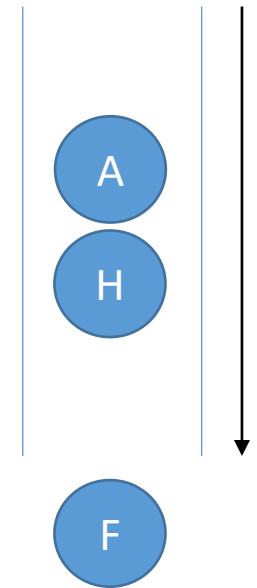
A cost = 10
B cost = 0
C cost = 1
D cost = ∞
E cost = 2
F cost = 3
G cost = ∞
H cost = 3

Dijkstra's Algorithm

- We move on to the next node in the priority queue



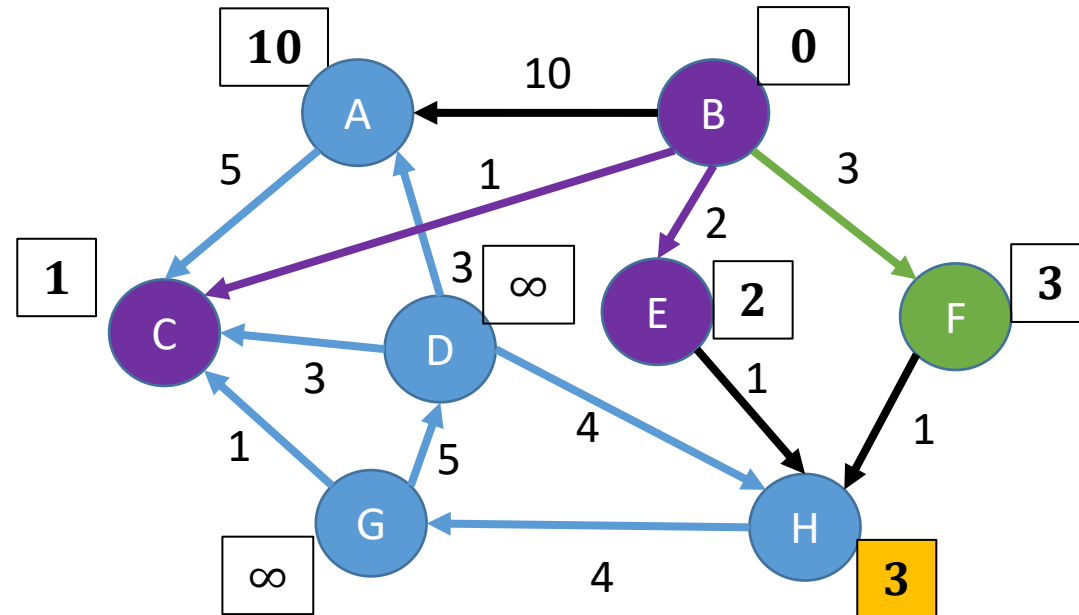
Priority Queue/Min-heap



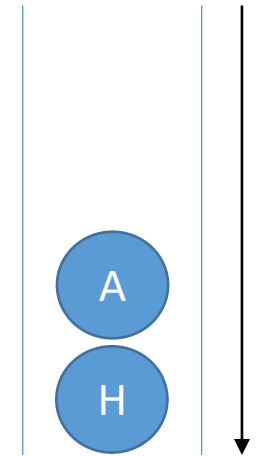
A cost = 10
B cost = 0
C cost = 1
D cost = ∞
E cost = 2
F cost = 3
G cost = ∞
H cost = 3

Dijkstra's Algorithm

- We'll now look at the nodes adjacent to F.
 - Cost of each is F's cost + cost of the edge
 - $3 + 1 = 4$ (NOT LESS THAN THE CURRENT COST OF H)



Priority Queue/Min-heap

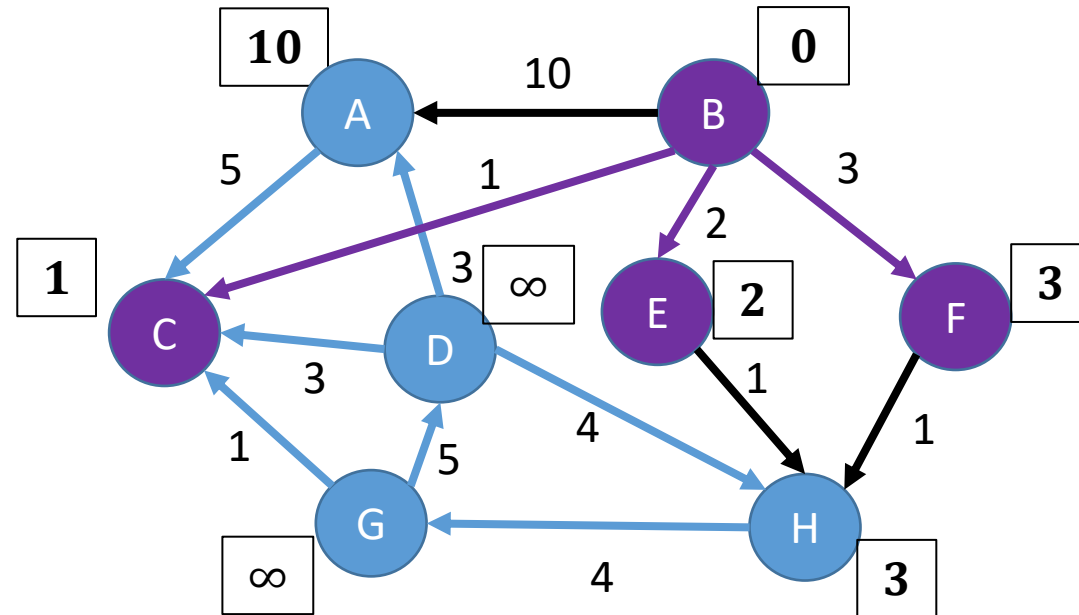
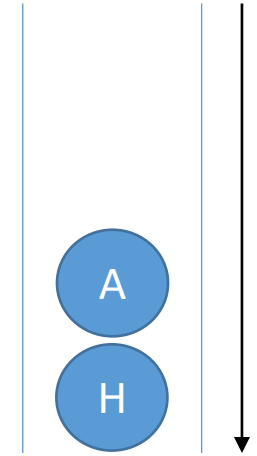


A cost = 10
B cost = 0
C cost = 1
D cost = ∞
E cost = 2
F cost = 3
G cost = ∞
H cost = 3

Dijkstra's Algorithm

- We are finished with F

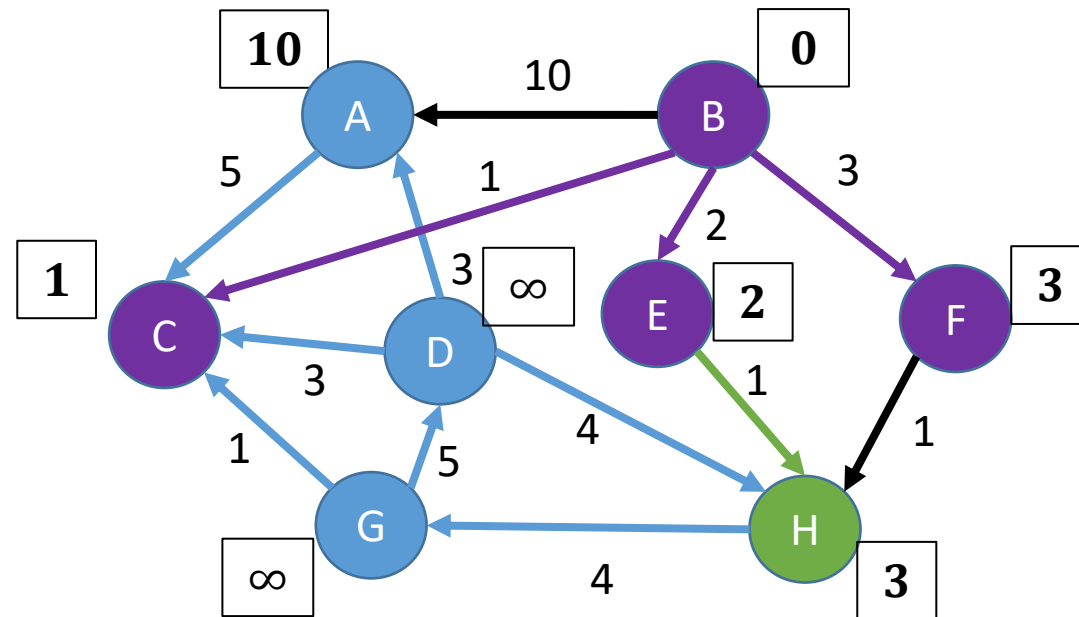
Priority Queue/Min-heap



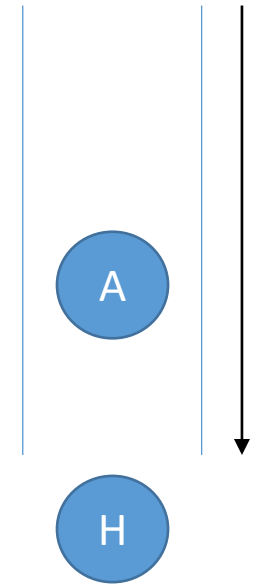
A cost = 10
B cost = 0
C cost = 1
D cost = ∞
E cost = 2
F cost = 3
G cost = ∞
H cost = 3

Dijkstra's Algorithm

- We move on to the next node in the priority queue



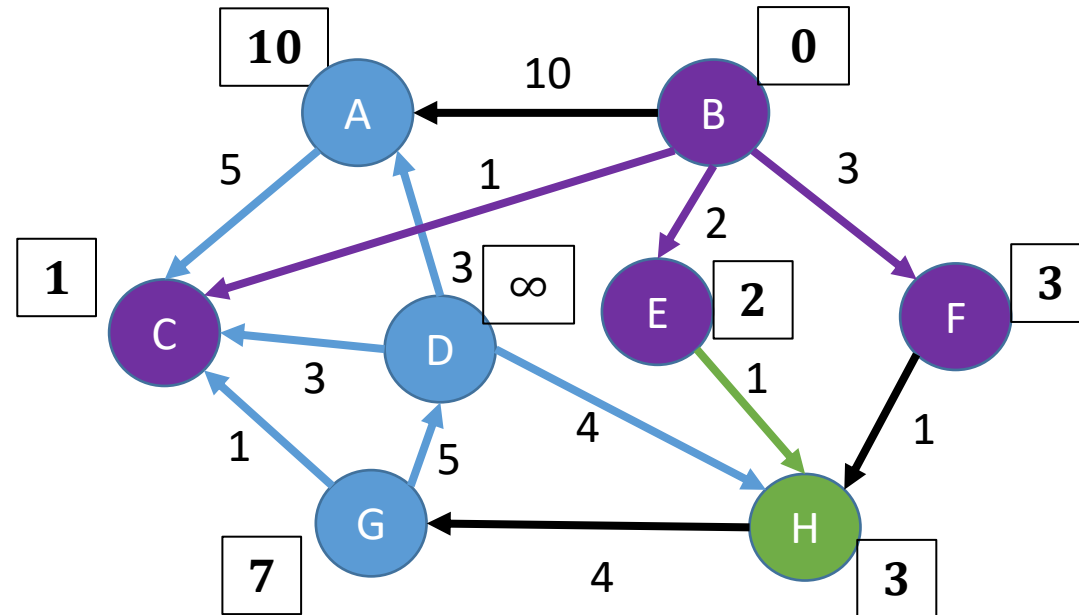
Priority Queue/Min-heap



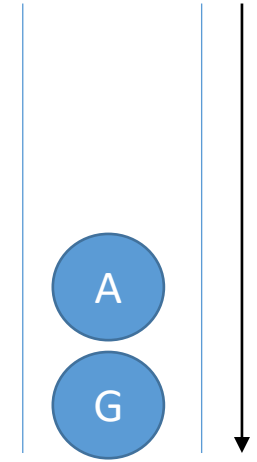
A cost = 10
B cost = 0
C cost = 1
D cost = ∞
E cost = 2
F cost = 3
G cost = ∞
H cost = 3

Dijkstra's Algorithm

- We'll now look at the nodes adjacent to H.
 - Cost of each is H's cost + cost of the edge



Priority Queue/Min-heap

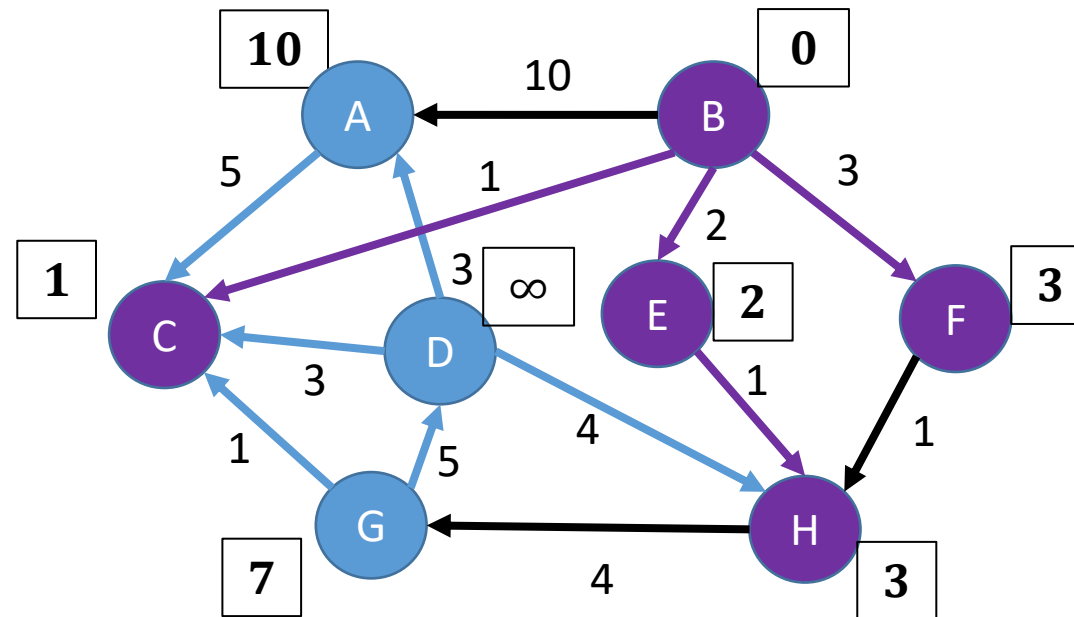
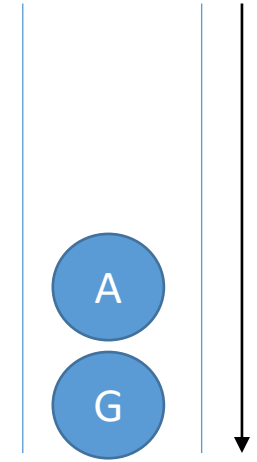


A cost = 10
B cost = 0
C cost = 1
D cost = ∞
E cost = 2
F cost = 3
G cost = 3 + 4 = 7
H cost = 3

Dijkstra's Algorithm

- We are finished with H.

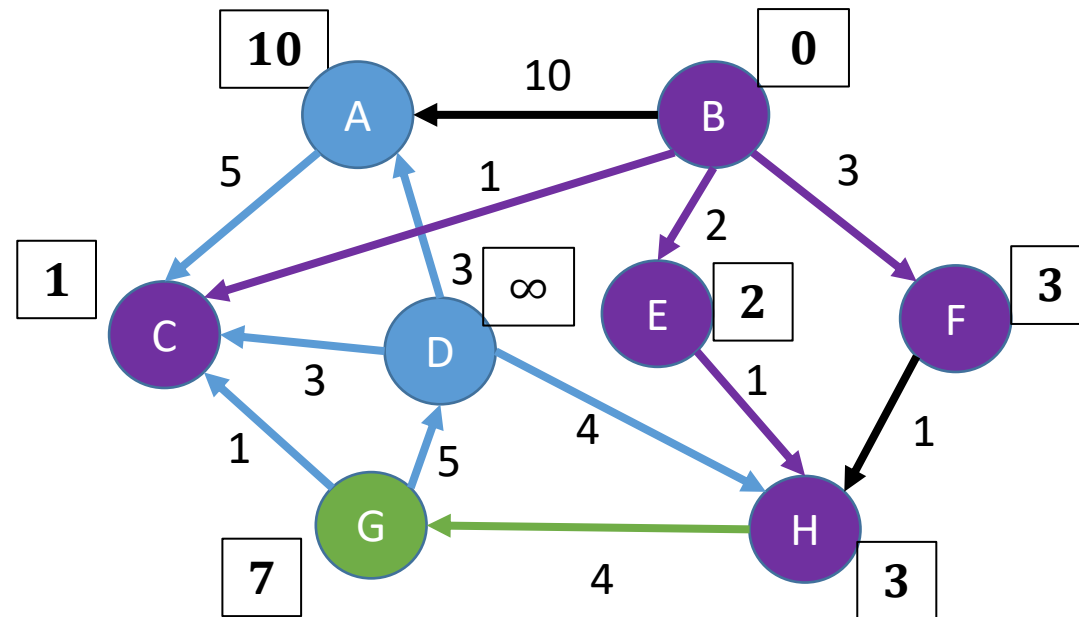
Priority Queue/Min-heap



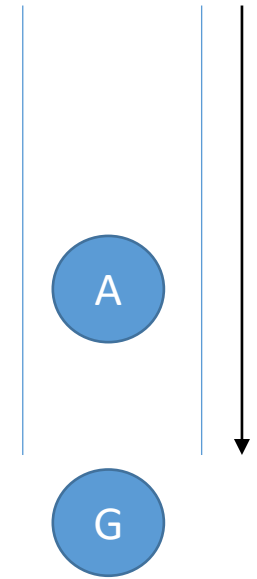
A cost = 10
B cost = 0
C cost = 1
D cost = ∞
E cost = 2
F cost = 3
G cost = 7
H cost = 3

Dijkstra's Algorithm

- We move on to the next node in the priority queue



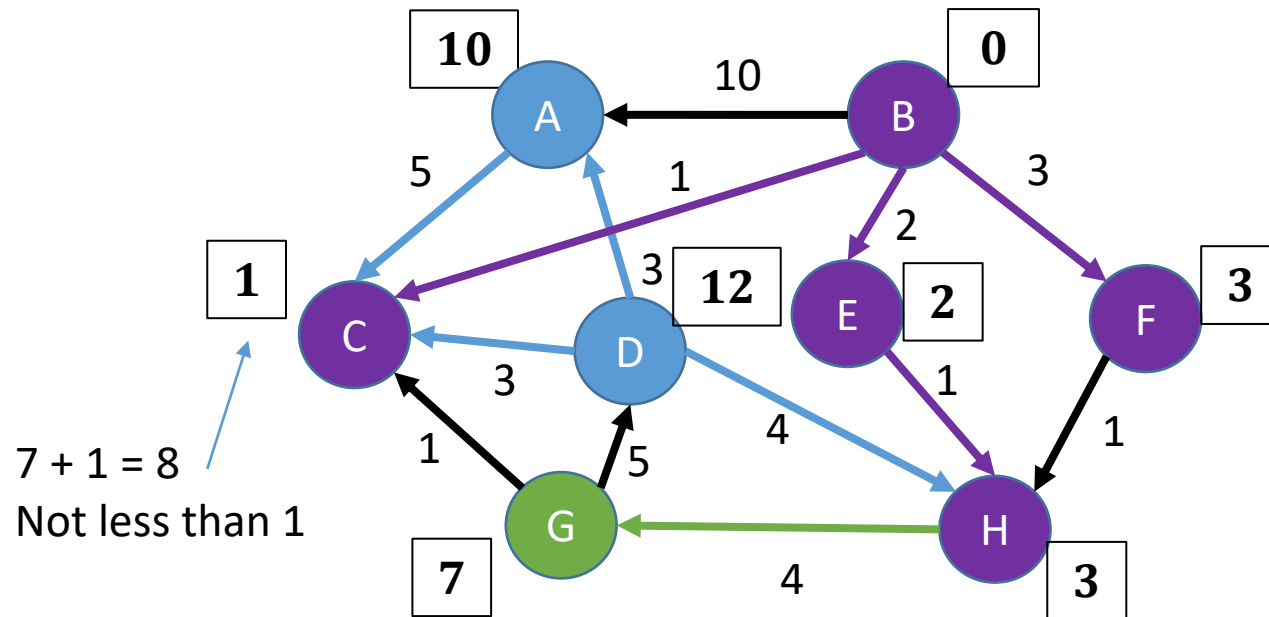
Priority Queue/Min-heap



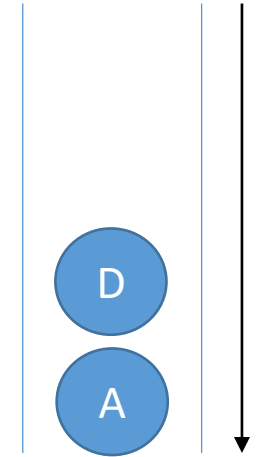
A cost = 10
B cost = 0
C cost = 1
D cost = ∞
E cost = 2
F cost = 3
G cost = 7
H cost = 3

Dijkstra's Algorithm

- We'll now look at the nodes adjacent to G.
 - Cost of each is G's cost + cost of the edge



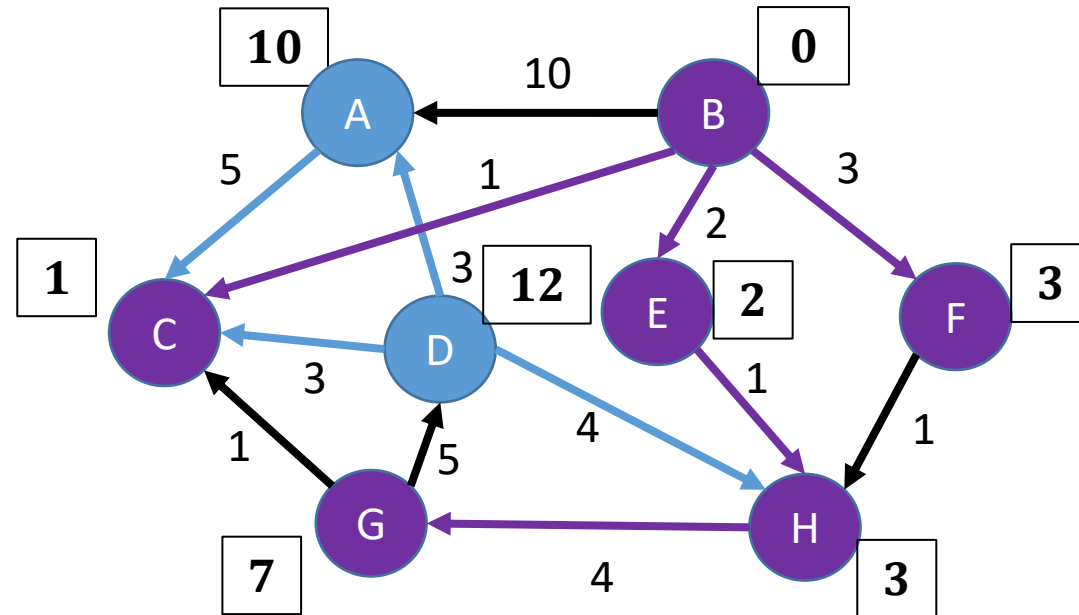
Priority Queue/Min-heap



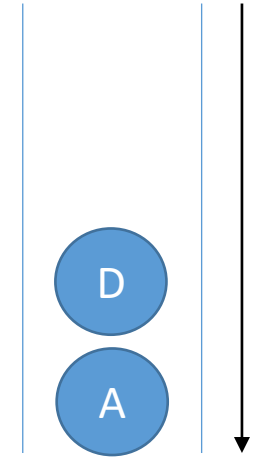
A cost = 10
B cost = 0
C cost = 1
D cost = $7 + 5 = 12$
E cost = 2
F cost = 3
G cost = 7
H cost = 3

Dijkstra's Algorithm

- We are finished with G



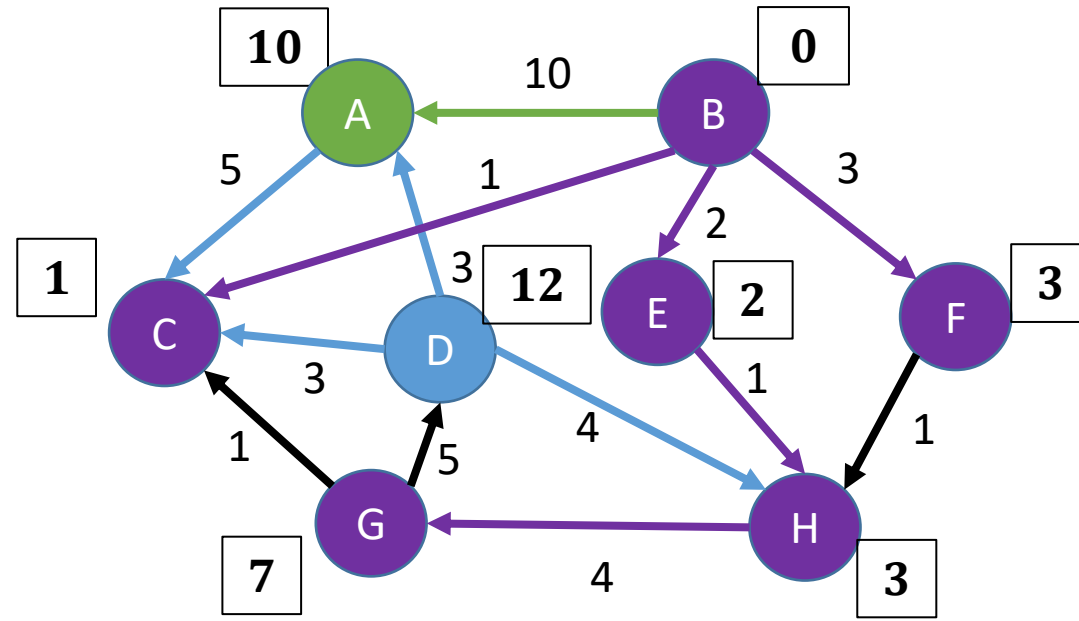
Priority Queue/Min-heap



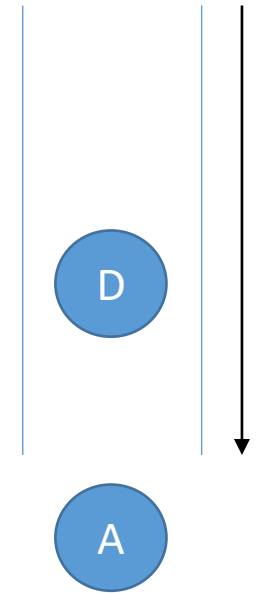
A cost = 10
B cost = 0
C cost = 1
D cost = 12
E cost = 2
F cost = 3
G cost = 7
H cost = 3

Dijkstra's Algorithm

- We move on to the next node in the priority queue



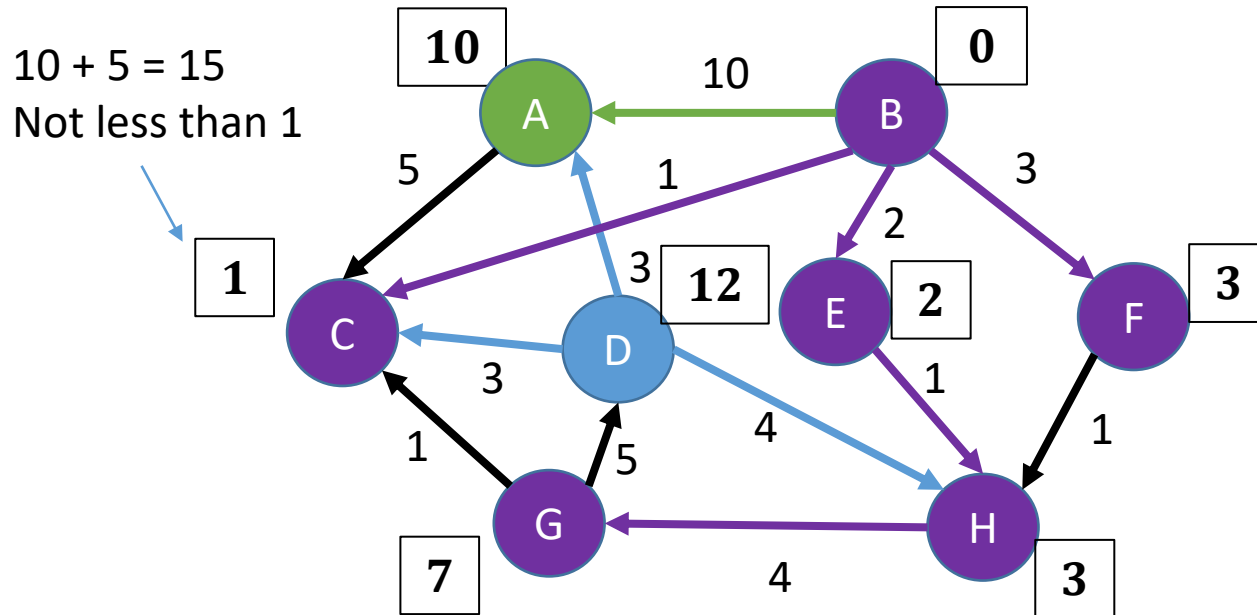
Priority Queue/Min-heap



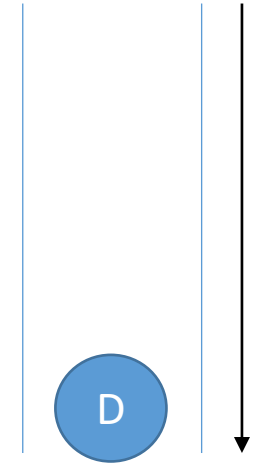
A cost = 10
B cost = 0
C cost = 1
D cost = 12
E cost = 2
F cost = 3
G cost = 7
H cost = 3

Dijkstra's Algorithm

- We'll now look at the nodes adjacent to A.
 - Cost of each is A's cost + cost of the edge



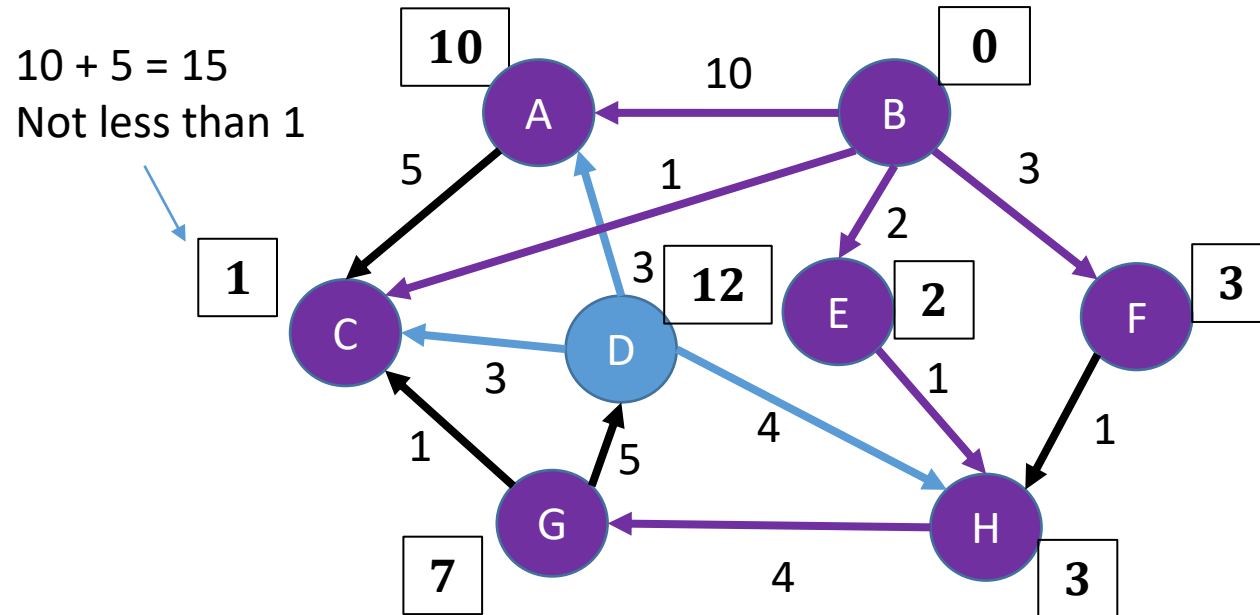
Priority Queue/Min-heap



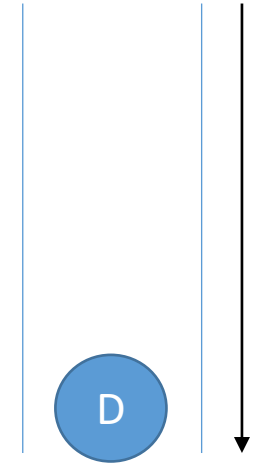
A cost = 10
B cost = 0
C cost = 1
D cost = 12
E cost = 2
F cost = 3
G cost = 7
H cost = 3

Dijkstra's Algorithm

- We are finished with A



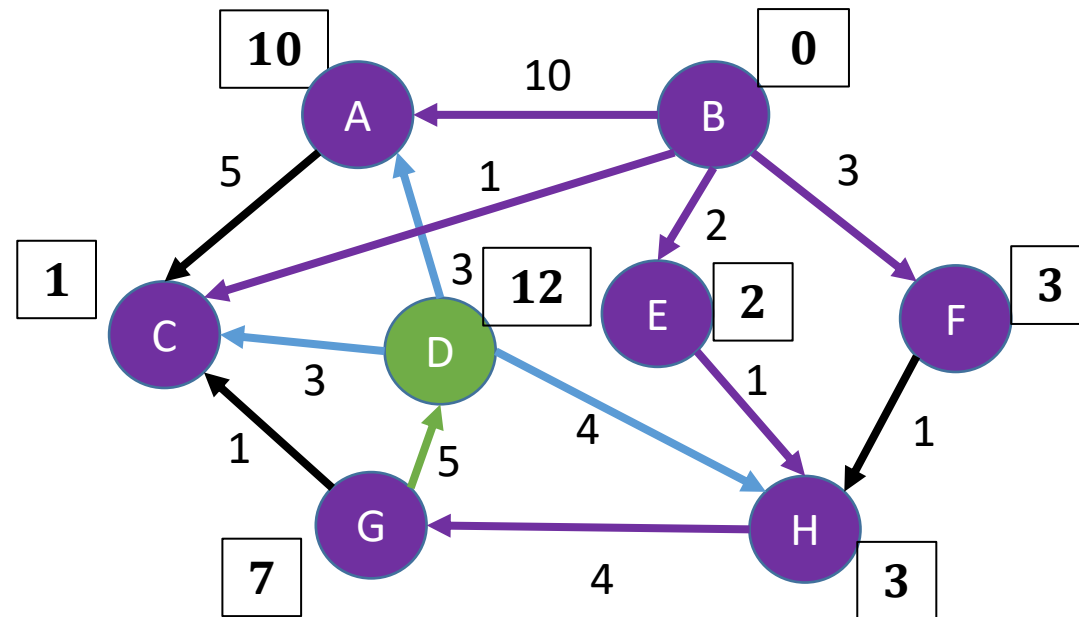
Priority Queue/Min-heap



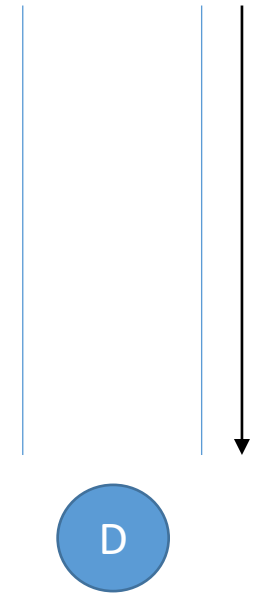
A cost = 10
B cost = 0
C cost = 1
D cost = 12
E cost = 2
F cost = 3
G cost = 7
H cost = 3

Dijkstra's Algorithm

- We move on to the next node in the priority queue



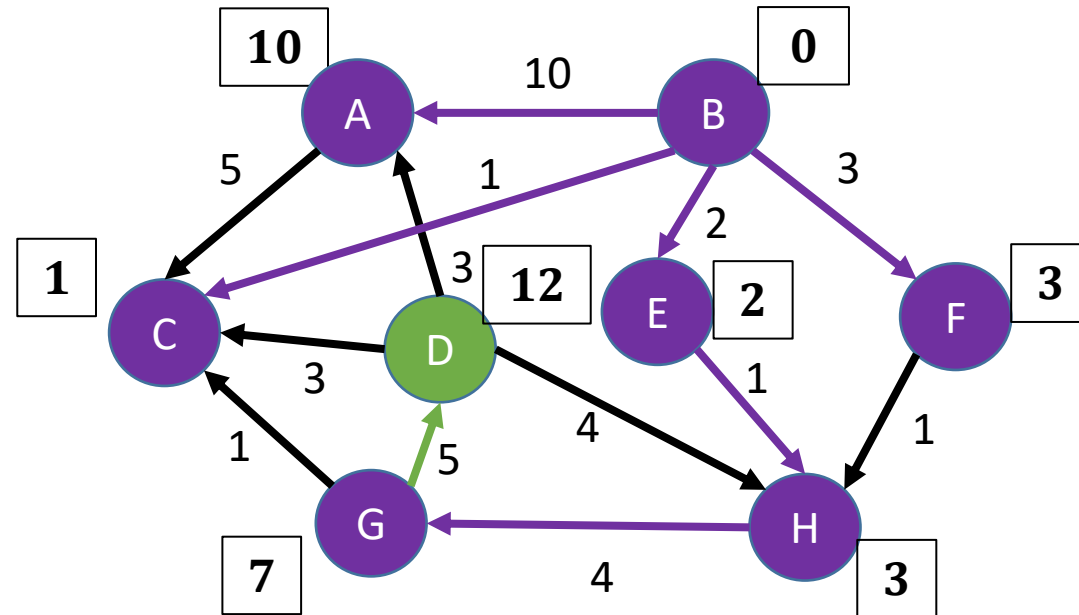
Priority Queue/Min-heap



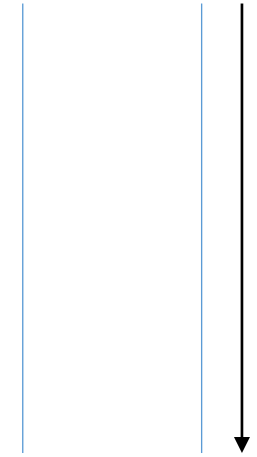
A cost = 10
B cost = 0
C cost = 1
D cost = 12
E cost = 2
F cost = 3
G cost = 7
H cost = 3

Dijkstra's Algorithm

- We'll now look at the nodes adjacent to D.
 - Cost of each is D's cost + cost of the edge



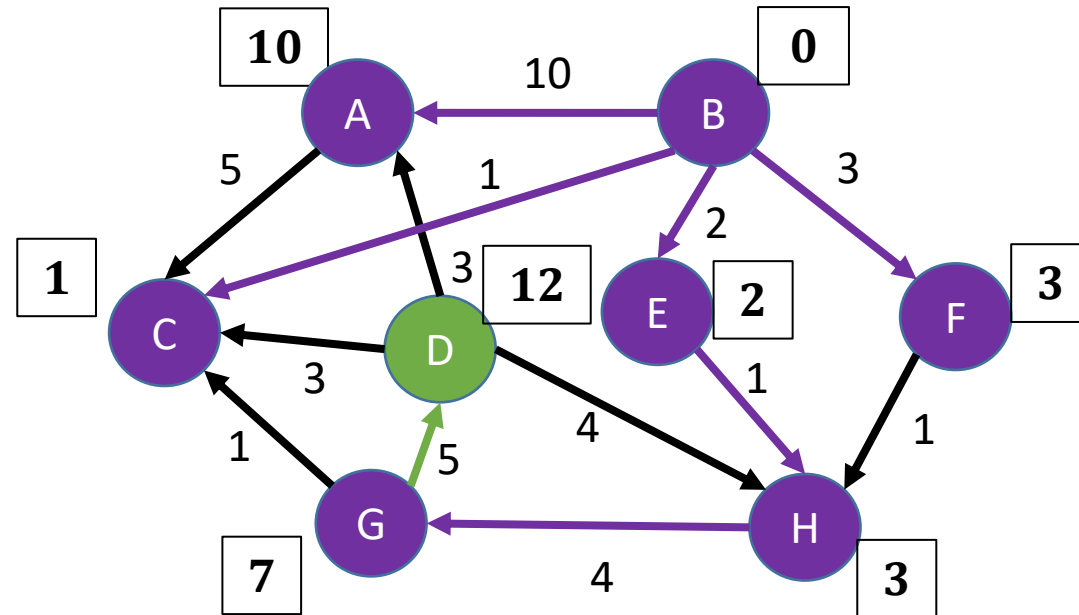
Priority Queue/Min-heap



A cost = 10
B cost = 0
C cost = 1
D cost = 12
E cost = 2
F cost = 3
G cost = 7
H cost = 3

Dijkstra's Algorithm

- $D \rightarrow A = 12 + 3 = 15$ (Not less than 10)
- $D \rightarrow C = 12 + 3 = 15$ (Not less than 1)
- $D \rightarrow H = 12 + 4 = 16$ (Not less than 3)

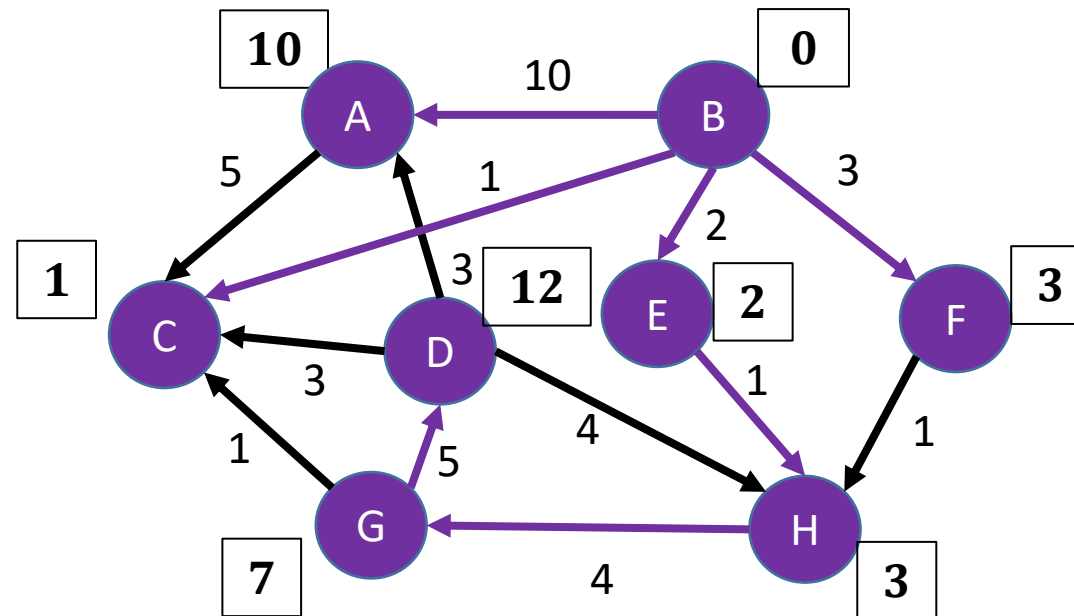
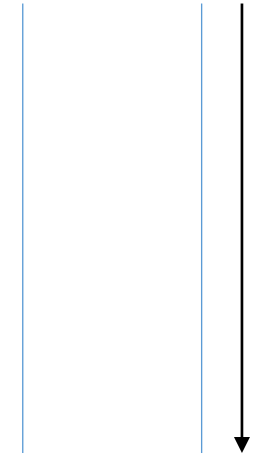


A cost = 10
 B cost = 0
 C cost = 1
 D cost = 12
 E cost = 2
 F cost = 3
 G cost = 7
 H cost = 3

Dijkstra's Algorithm

- We are finished with D

Priority Queue/Min-heap

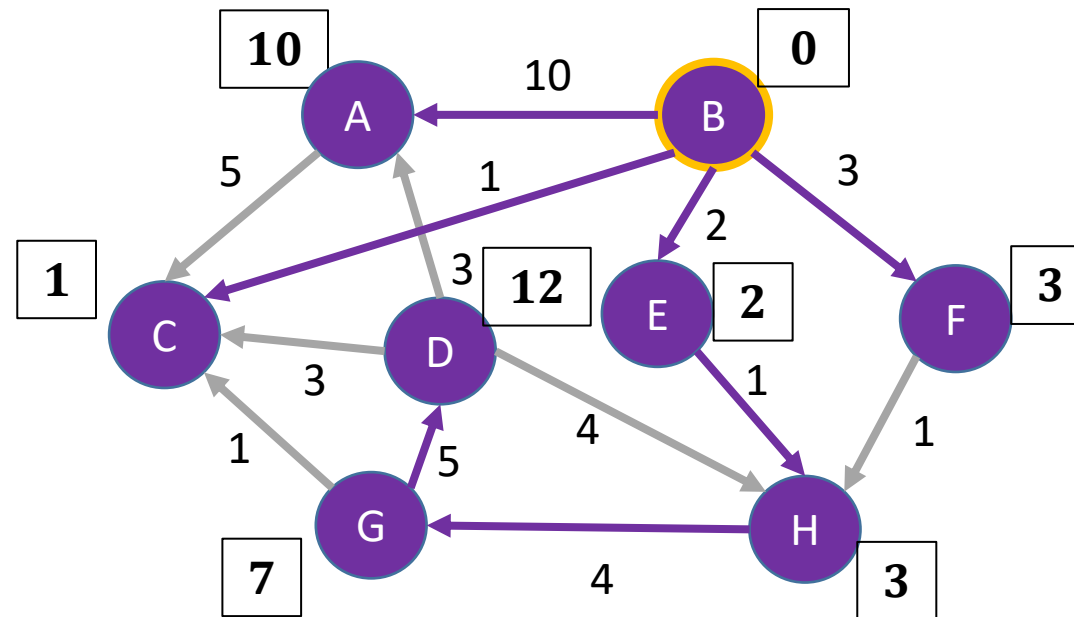
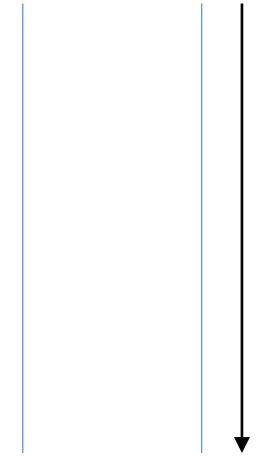


A cost = 10
B cost = 0
C cost = 1
D cost = 12
E cost = 2
F cost = 3
G cost = 7
H cost = 3

Dijkstra's Algorithm

- Priority Queue is empty
 - Algorithm is complete

Priority Queue/Min-heap



Least cost from B to all other nodes:

A cost = 10
B cost = 0
C cost = 1
D cost = 12
E cost = 2
F cost = 3
G cost = 7
H cost = 3