## CSC 139: Operating Systems Principles Second Quiz, Fall 2020 Friday, April 17th, 2020 Section 4

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1011

Student Name:		KCY	Student Number:				
	If in the avoic	RUE or FALS lance version of nave a deadloc	of the Banker's algorith	<b>m</b> , we fail to find a	safe sequence		oints] will FALSE
2.		ultiple processe ave a deadlock	es that are holding some	resources and wai	ting for other re	sources, we TRUE	FALSE
3.	A deadlock cannot happen unless mutual exclusion is required				resources.	TRUE	FALSE
4.		possible to red inating any prod	cover from a deadlock by cess.	preempting some	processes from	some resour	ces FALSE
	The worst-cas	se running time	ere is only <u>one</u> correct for finding a safe seque nd <i>q</i> resource types is: (C.)O(p <sup>2</sup> q)		s algorithm for e. O(pq)	[60 poi deadlock det f. O(p²c	tection on
V	<ul> <li>Which of the following statements is (are) true about cycles in the Resource Allocation Graph (RAG)?</li> <li>a. Having a cycle in the RAG is always a necessary condition for deadlocks.</li> <li>b. Having a cycle in the RAG is always a sufficient condition for deadlocks.</li> <li>c. Having a cycle in the RAG necessarily implies a deadlock if all resource types have multiple instances</li> <li>d. Having a cycle in the RAG necessarily implies a deadlock if all resource types have single instances.</li> <li>e. Both a and c are true.</li> <li>g. a, b and d are true.</li> </ul>						
3.	Which of the following is a <b>necessary</b> condition for a. Some resources are preemptive  C. Some resources are non-preemptive  e. Both a and b are true.			deadlocks:  b. All resources are preemptive  d. All resources are non-preemptive  f. Both c and d are true.			
4.	Po Po Po Po Po Po Po What's the m	rent Allocation 2 3 1	a system that has one re Current Request 4 8 2 6 or the <u>total</u> number of ins 9 d. 10	stances that will ma	Avail 2 3 2 3 3 3 3 3 3 3 3 3 3 3 3 3 3 3 3		?
5.	R <sub>2</sub> and there an instance of and P <sub>3</sub> is hold a. P <sub>2</sub> and P <sub>3</sub> b P <sub>1</sub> and P <sub>2</sub> c. P <sub>1</sub> and P <sub>3</sub>	are <u>two</u> instand of R <sub>2</sub> and an ins ding an instance are in a deadle are in a deadle	esses $P_1$ , $P_2$ and $P_3$ and es of each of $R_1$ and $R_3$ , tance of $R_3$ , $P_2$ is holding of $R_3$ and requesting alock but $P_1$ is not in a dealock but $P_3$ is not in a dealock but $P_2$ is not in a dealock and deadlock.	resource types R <sub>1</sub> , P <sub>1</sub> is currently hold go an instance of R <sub>2</sub> in instance of R <sub>1</sub> . Wadlock. adlock.	R₂ and R₃. The ding an instance and requesting	ere is <u>one</u> ins e of R₁ and re two instance: it state of the	questing s of R <sub>1</sub> , system?
6.	Consider a sy	stem with proc	esses P <sub>1</sub> , P <sub>2</sub> and P <sub>3</sub> and	resource types R <sub>1</sub>	and R <sub>2</sub> . There a	are <u>two</u> instar	nces of

each resource type. If  $P_1$  is currently holding an instance of  $R_1$ ,  $P_2$  is holding an instance of  $R_2$ , and  $P_3$  is holding an instance of  $R_2$ , which of the following sequences of requests will **necessarily** cause a deadlock?

- a. P<sub>1</sub> requests one instance of R<sub>2</sub>, P<sub>2</sub> requests one instance of R<sub>1</sub> and P<sub>3</sub> requests one instance of R<sub>1</sub>
  b. P<sub>1</sub> requests one instance of R<sub>2</sub>, P<sub>2</sub> requests one instance of R<sub>1</sub> and P<sub>3</sub> requests two instances of R<sub>1</sub>
  c. P<sub>1</sub> requests one instance of R<sub>2</sub>, P<sub>3</sub> requests one instance of R<sub>1</sub> and P<sub>2</sub> requests two instances of R<sub>1</sub>
  d. P<sub>1</sub> requests one instance of R<sub>2</sub>, P<sub>2</sub> requests two instances of R<sub>1</sub> and P<sub>3</sub> requests two instances of R<sub>1</sub>
- f. c and d are correct e. b and d are correct

g. b, c and d are correct



