

Improved Search for Integral, Impossible Differential and Zero-Correlation Attacks

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Motivation and Our Contributions



Motivation

- ✔ Providing a tool to find **complete** integral, and ID/ZC attacks



Contributions

- ✔ Improving the CP-based methods to find ID/ZC, and integral distinguishers
- ✔ Introducing a CP model for the partial-sum technique for the first time
- ✔ Improving distinguishers of Ascon, QARMAv2, and ForkSKINNY (25 Dists.)
- ✔ Improving key recovery attacks of SKINNY, and ForkSKINNY (24 Attacks)

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Part of Our Results Regarding Distinguishing Attacks

Cipher	#Rounds	Dist.	Data complexity	Ref.
QARMAv2-64	5	Integral	-	[Ava+23]
QARMAv2-64 ($\mathcal{T} = 1$)	7 / 8 / 9	Integral	2^8 / 2^{16} / 2^{44}	This work
QARMAv2-64 ($\mathcal{T} = 2$)	8 / 9 / 10	Integral	2^8 / 2^{16} / 2^{44}	This work
QARMAv2-128($\mathcal{T} = 2$)	10 / 11 / 12	Integral	2^{16} / 2^{44} / 2^{96}	This work
ForkSKINNY-64-192	16	Integral	2^{72}	[Niu+21]
ForkSKINNY-64-192	17	Integral	2^{60}	This work
ForkSKINNY-64-192	16	ID	-	[HSE23]
ForkSKINNY-64-192	21	ID	-	This work
ForkSKINNY-128-256	14	Integral	2^{56}	[HSE23]
ForkSKINNY-128-256	15	Integral	2^{56}	This work

Part of Our Results Regarding Key Recovery Attacks

Cipher	#R	Time	Data	Mem.	Attack	Setting / Model	Ref.
SKINNY-64-64	17	2^{59}	$2^{58.79}$	2^{40}	ID	STK / CP	[HSE23]
	18	$2^{53.58}$	$2^{53.58}$	2^{48}	Int	60,SK / CP,CT	This work
SKINNY-128-128	17	$2^{116.51}$	$2^{116.37}$	2^{80}	ID	STK / CP	[HSE23]
	18	$2^{105.58}$	$2^{105.58}$	2^{96}	Int	120,SK / CP,CT	This work
SKINNY-128-384	26	2^{344}	2^{121}	2^{340}	Int	360,SK / CP,CT	[HSE23]
	26	2^{331}	2^{122}	2^{328}	Int	360,SK / CP,CT	This work
ForkSKINNY-128-256	26	$2^{250.30}$	2^{127}	2^{160}	ID	256,RTK / CP	[BDL20]
	26	$2^{238.50}$	$2^{128.60}$	$2^{175.60}$	ID	256,RTK / CP	This paper

Outline

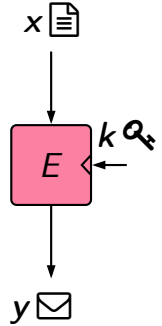
- 1 Background and the Research Gap
- 2 Search For Distinguishers
- 3 Our New Word-Wise Method for Finding Distinguishers
- 4 Our New Bit-Wise Method for Finding Distinguishers
- 5 Our Unified CP Model for Key-Recovery
- 6 Contributions and Future Works

Background and the Research Gap



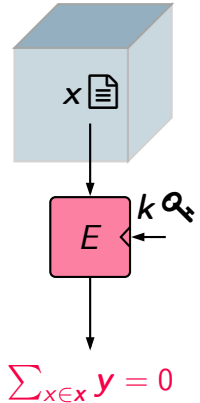
Integral, ID, and ZC Distinguishers

- Integral attack [Lai94; DKR97]
- Impossible-differential attack [BBS99; Knu98]
- Zero-correlation attack [BR14]



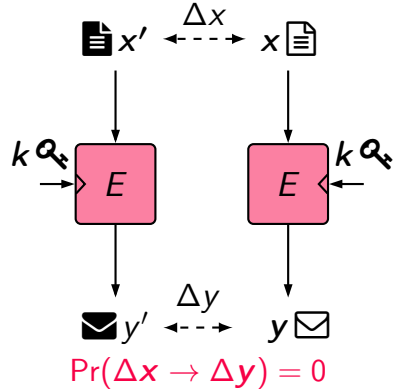
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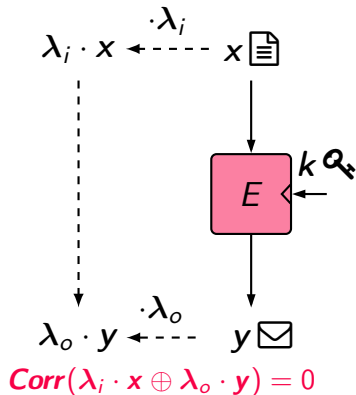
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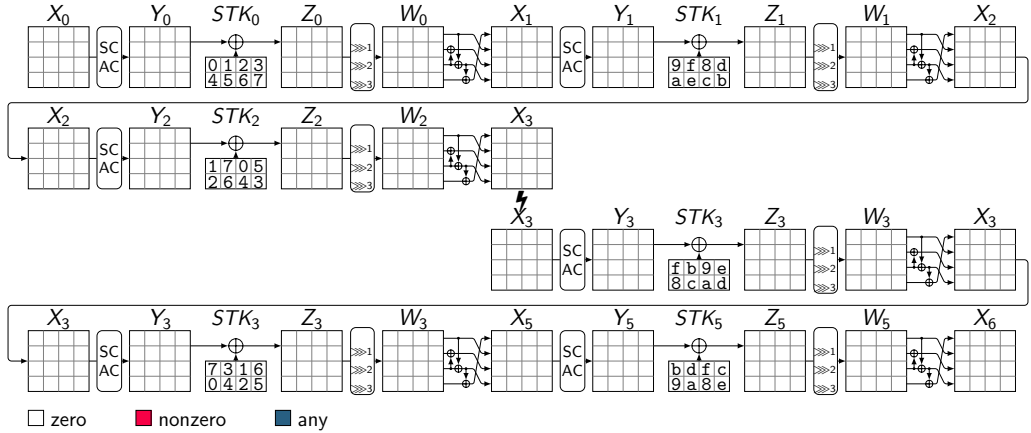
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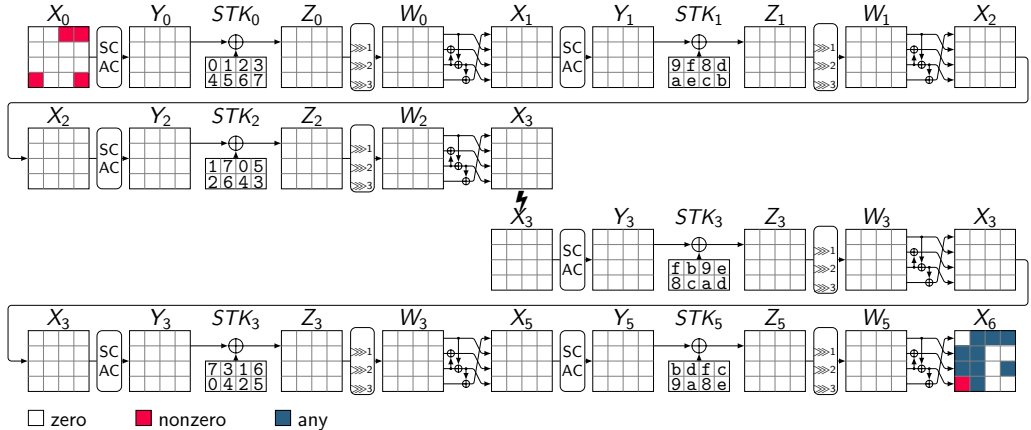
Miss-in-the-Middle Technique [BBS99]

- Find two differences (linear masks) that propagate forward and backward with probability one and contradict each other in the middle



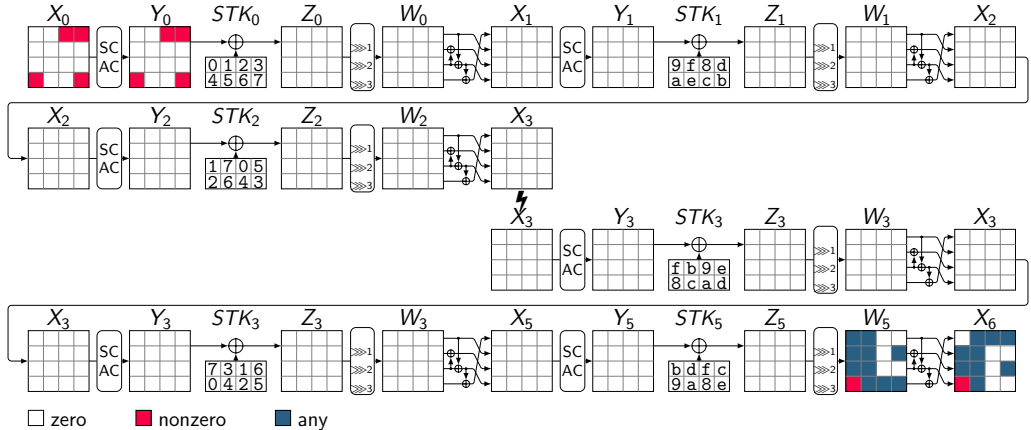
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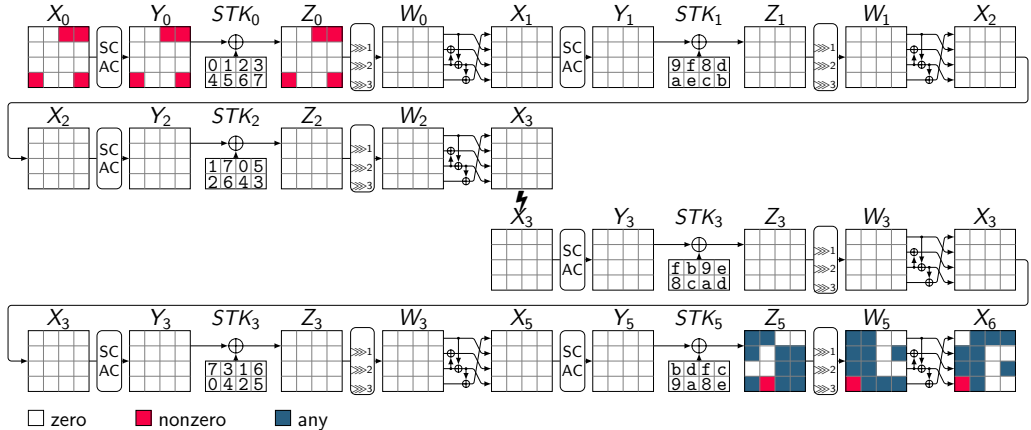
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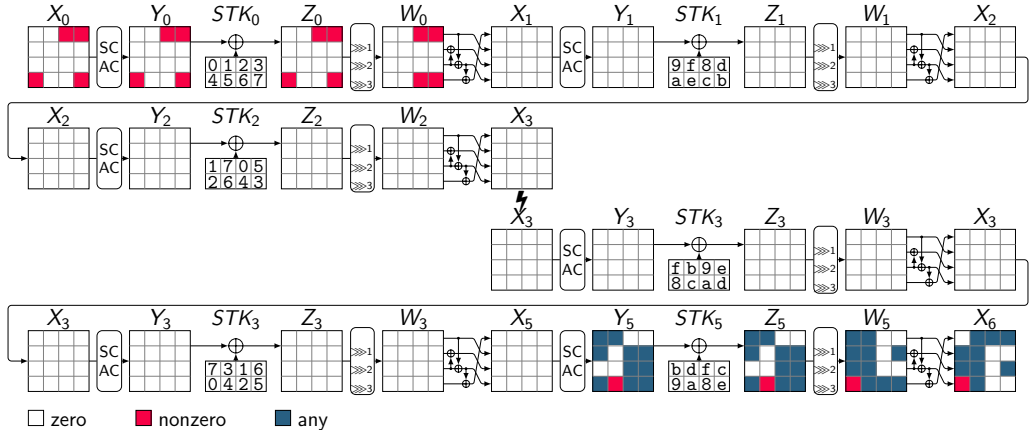
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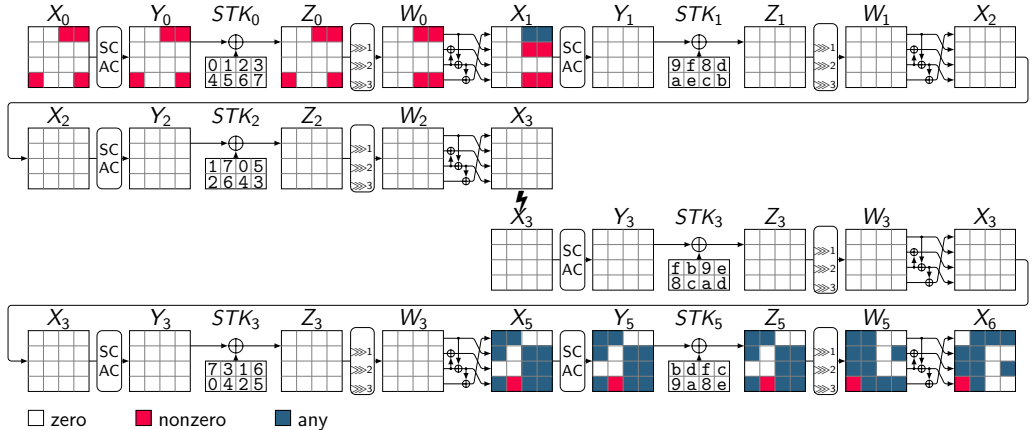
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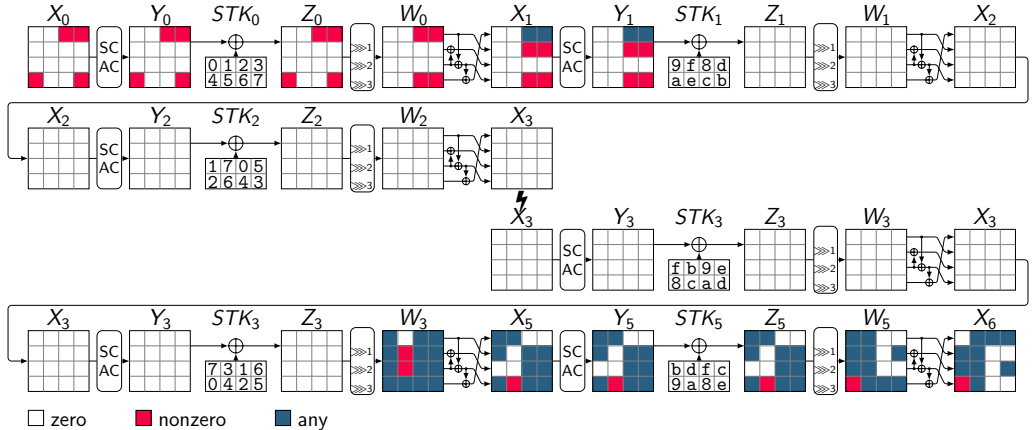
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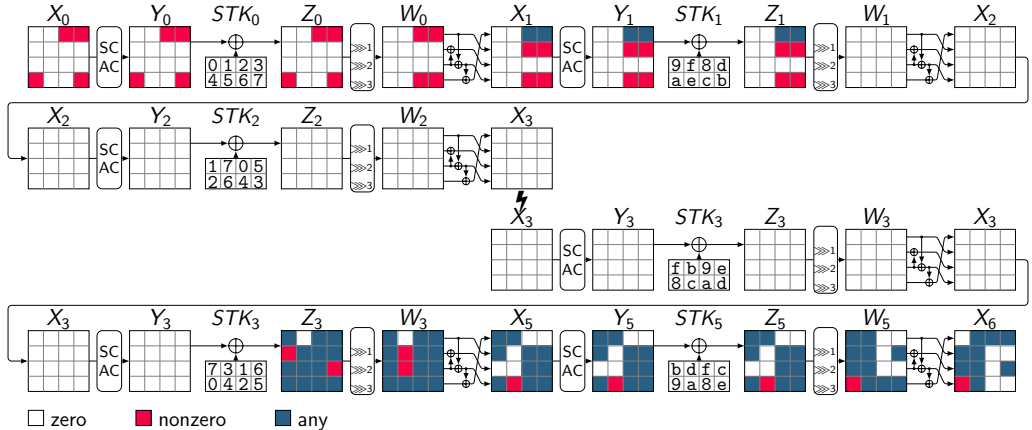
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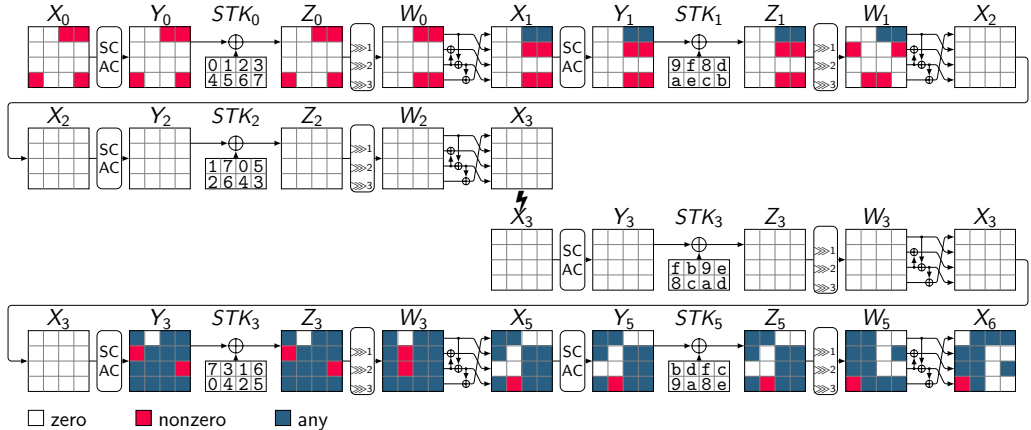
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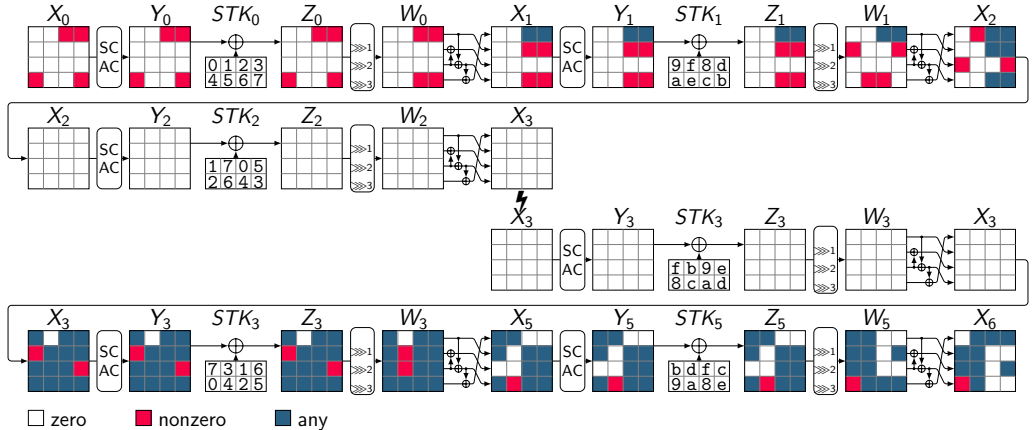
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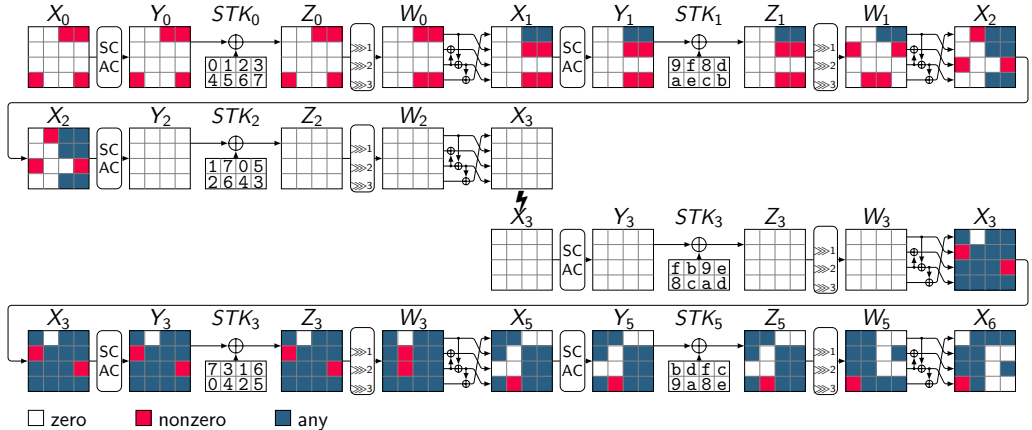
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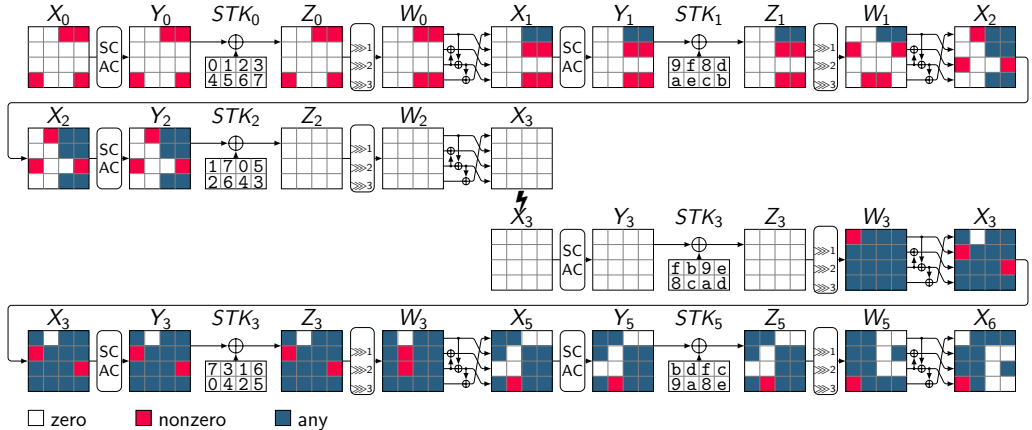
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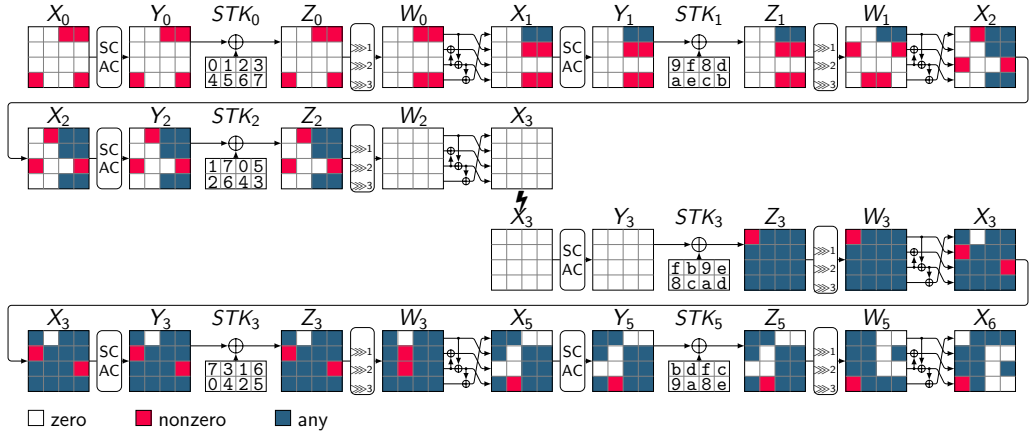
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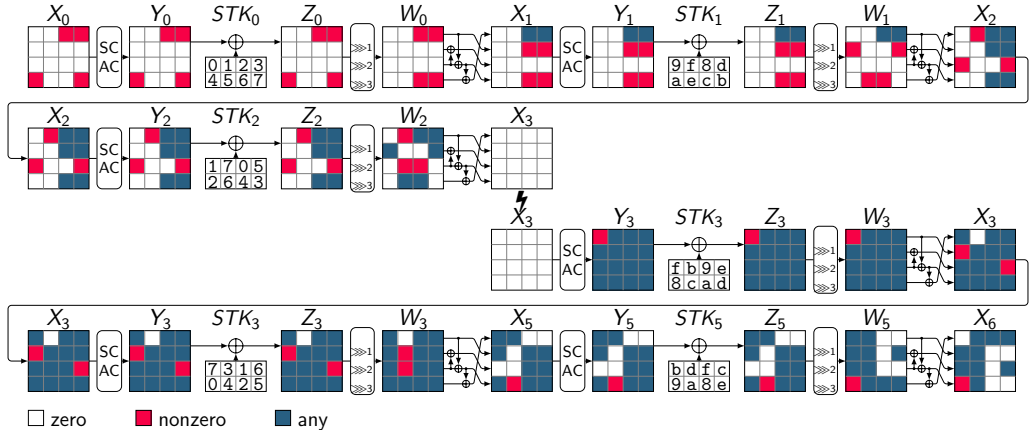
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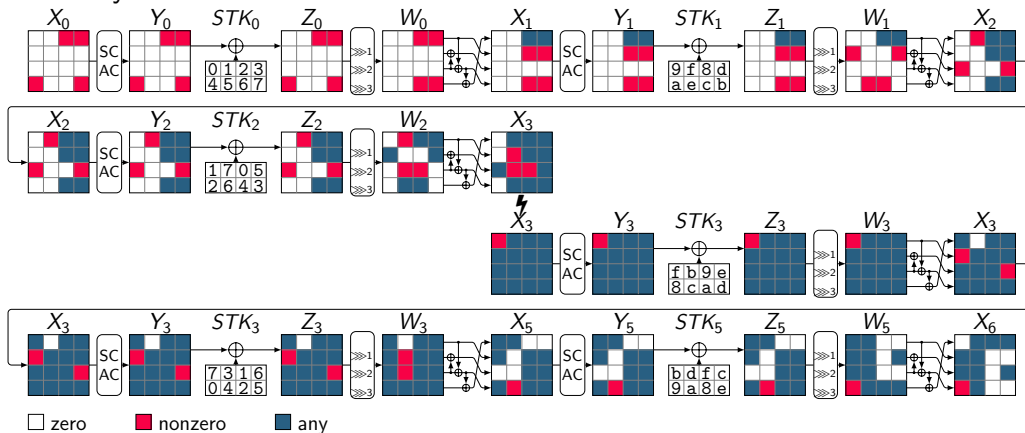
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Relation Between ZC and Integral Distinguishers

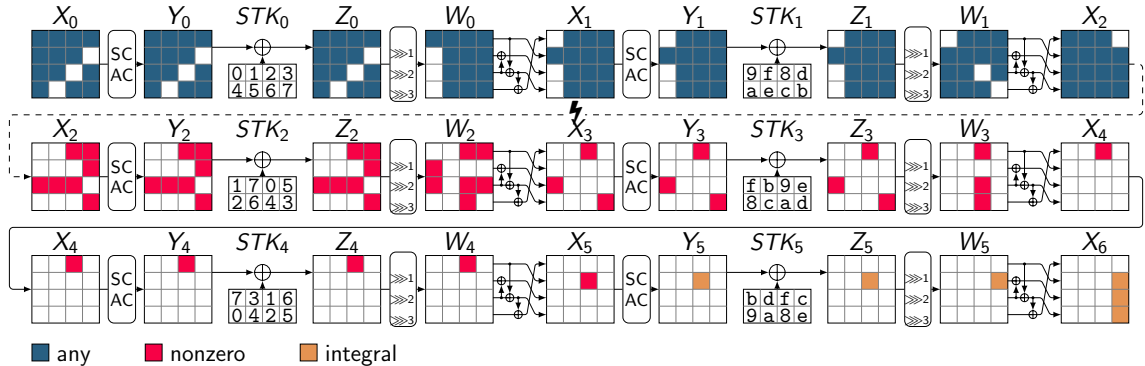
- Any ZC distinguisher can be converted to an integral distinguisher [Sun+15].

Link Between ZC and Integral Distinguishers [Sun+15]

Let $F : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^n$ be a vectorial Boolean function. Assume A is a subspace of \mathbb{F}_2^n and $\beta \in \mathbb{F}_2^n \setminus \{0\}$ such that (α, β) is a ZC approximation for any $\alpha \in A$. Then, for any $\lambda \in \mathbb{F}_2^n$, $\langle \beta, F(x + \lambda) \rangle$ is balanced over the set

$$A^\perp = \{x \in \mathbb{F}_2^n \mid \forall \alpha \in A : \langle \alpha, x \rangle = 0\}.$$

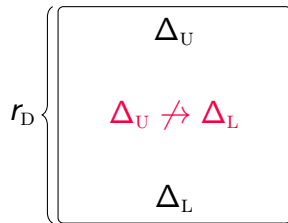
Example: Conversion of ZC Distinguisher to Integral Distinguisher



- $X_0[7, 10, 13]$ takes all possible values and the remaining cells take a fixed value
- $X_6[7] \oplus X_6[11] \oplus X_6[15]$ is balanced

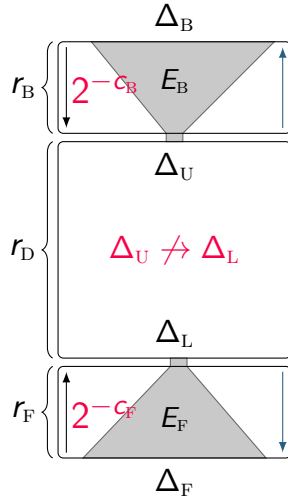
ID, ZC, and Integral Key Recovery

- Common technique for ID key recovery:
 - Early abort technique [Lu+08]
- Common technique for ZC/Integral key recovery:
 - Partial-sum technique [Fer+00]



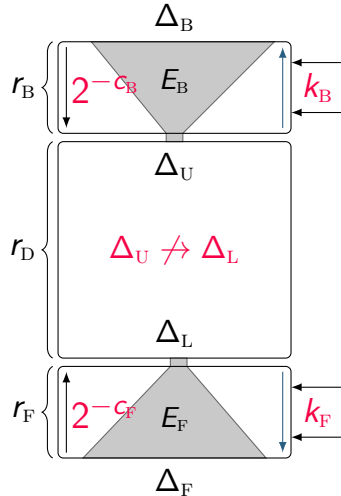
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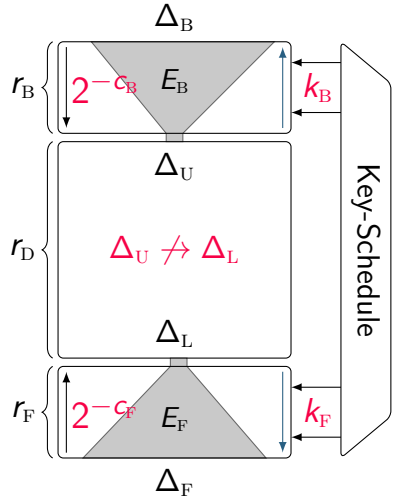
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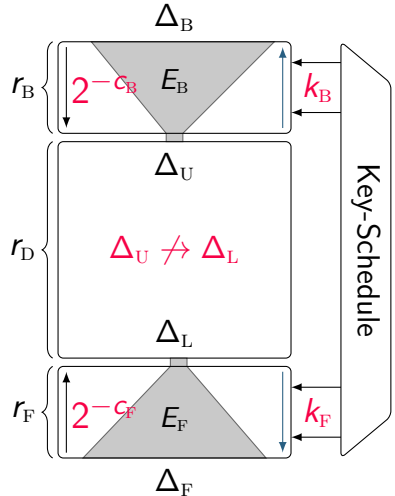
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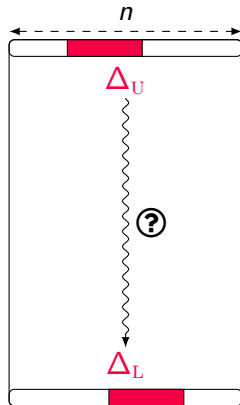
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Previous Tools for ID/ZC, and Integral Attacks

- Tools based on dedicated algorithms:
 - CRYPTO 2016 (\mathcal{DC} -MITM, ID) [DF16]
- Tools based on general purpose solvers:
 - Eprint 2016 (ID) [Cui+16]
 - ASIACRYPT 2016 (Integral) [Xia+16]
 - EUROCRYPT 2017 (ID, ZC) [ST17]
 - ToSC 2017 (ID, ZC) [Sun+17]
 - ToSC 2020 (ID, ZC) [Sun+20]



Search for Distinguishers



Our Previous Method to Search Distinguishers [HSE23]

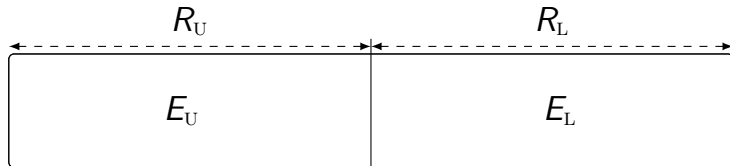
E

✓ $CSP_U(\Delta_U, \Delta'_U)$

✓ $CSP_L(\Delta_L, \Delta'_L)$

✓ $CSP_M(\Delta'_U, \Delta'_L)$

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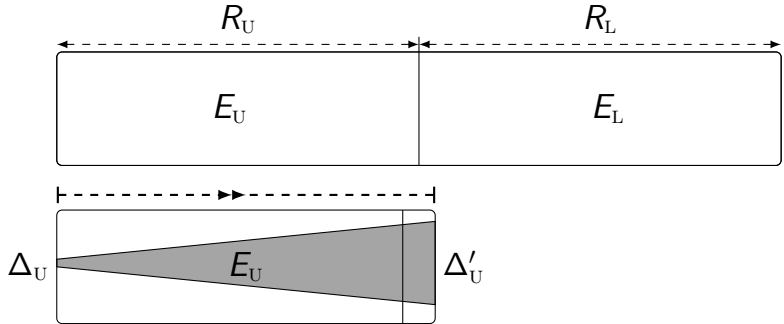
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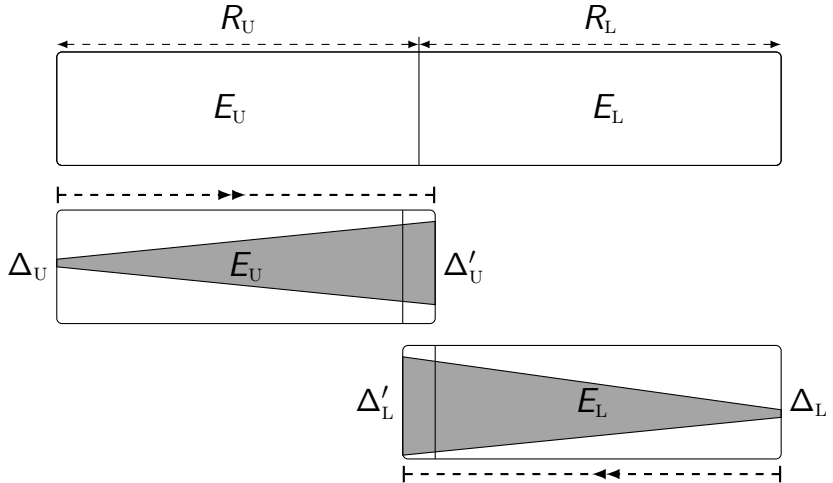


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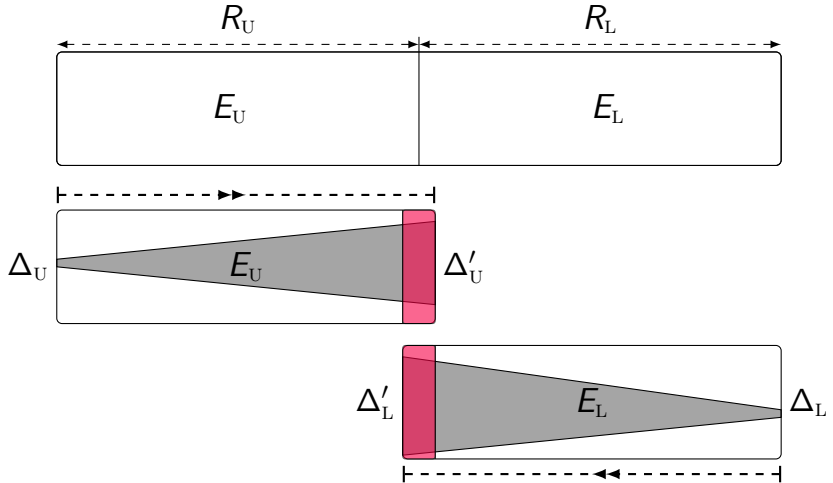


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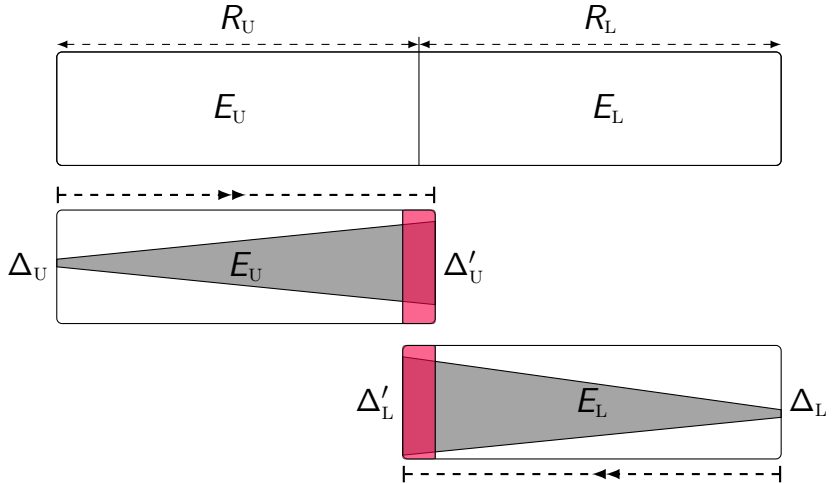
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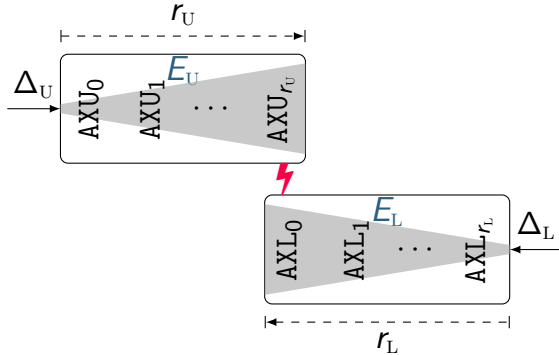


Our New Word-Wise Method for Finding Distinguishers

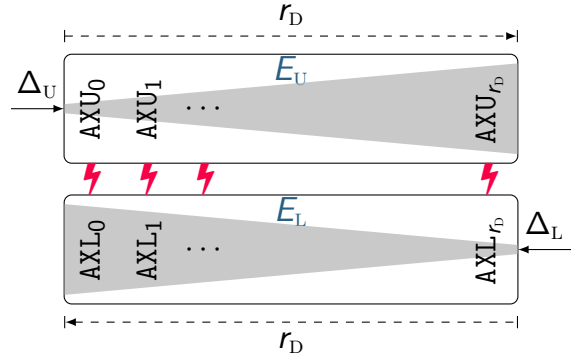


Relax the Limit of Fixing the Contradiction's Location

 Find ID distinguisher for $r_D (= r_U + r_L)$ rounds



Modeling the distinguishers in [HSE23].



Our modeling of the distinguishers.

Our New Bit-Wise Method for Finding Distinguishers



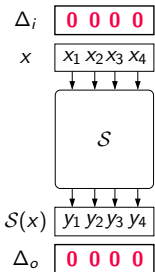
Deterministic Bit-Wise Differential Trails (a.k.a. Undisturbed Bits [Tez14])

x	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
$S(x)$	4	0	a	7	b	e	1	d	9	f	6	8	5	2	c	3

$\Delta_i \setminus \Delta_o$	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
0	16	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	2	2	2	2	0	0	0	0	2	2	2	2
2	0	2	0	2	0	0	0	4	0	2	2	0	0	0	2	2
3	0	2	0	2	0	0	4	0	0	2	2	0	0	0	2	2
4	0	0	0	0	0	0	0	0	0	0	4	4	2	2	2	2
5	0	0	0	0	2	2	2	2	0	0	4	4	0	0	0	0
6	0	2	0	2	0	4	0	0	0	2	2	0	2	2	0	0
7	0	2	0	2	4	0	0	0	0	2	2	0	2	2	0	0
8	0	0	0	0	0	0	0	0	0	0	0	0	4	4	4	4
9	0	4	4	0	0	0	0	0	0	4	0	4	0	0	0	0
a	0	0	2	2	2	0	0	2	4	0	0	0	0	2	0	2
b	0	0	2	2	0	2	2	0	4	0	0	0	2	0	2	0
c	0	4	4	0	2	2	2	2	0	0	0	0	0	0	0	0
d	0	0	0	0	2	2	2	2	0	4	0	4	0	0	0	0
e	0	0	2	2	0	2	2	0	4	0	0	0	0	2	0	2
f	0	0	2	2	2	0	0	2	4	0	0	0	2	0	2	0

$$\Delta_i = (0, 0, 0, 0) \xrightarrow{S} \Delta_o = (0, 0, 0, 0)$$

$$\Delta_i \neq (0, 0, 0, 0) \xrightarrow{S} \Delta_o \neq (0, 0, 0, 0)$$



Deterministic Bit-Wise Differential Trails (a.k.a. Undisturbed Bits [Tez14])

x	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
S(x)	4	0	a	7	b	e	1	d	9	f	6	8	5	2	c	3

$\Delta_i \backslash \Delta_o$		0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
Δ_i	0	16	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	1	0	0	0	0	2	2	2	2	0	0	0	0	2	2	2	2
x	2	0	2	0	2	0	0	0	4	0	2	2	0	0	0	2	2
	3	0	2	0	2	0	0	4	0	0	2	2	0	0	0	2	2
$S(x)$	4	0	0	0	0	0	0	0	0	0	4	4	2	2	2	2	2
	5	0	0	0	0	2	2	2	2	0	0	4	4	0	0	0	0
Δ_o	6	0	2	0	2	0	4	0	0	0	2	2	0	2	2	0	0
	7	0	2	0	2	4	0	0	0	0	2	2	0	2	2	0	0
	8	0	0	0	0	0	0	0	0	0	0	0	4	4	4	4	4
	9	0	4	4	0	0	0	0	0	4	0	4	0	0	0	0	0
	a	0	0	2	2	2	0	0	2	4	0	0	0	0	2	0	2
	b	0	0	2	2	0	2	2	0	4	0	0	0	2	0	2	0
	c	0	4	4	0	2	2	2	2	0	0	0	0	0	0	0	0
	d	0	0	0	0	2	2	2	2	0	4	0	4	0	0	0	0
	e	0	0	2	2	0	2	2	0	4	0	0	0	0	2	0	2
	f	0	0	2	2	2	0	0	2	4	0	0	0	2	0	2	0

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$\Delta_i \neq (0, 0, 0, 0) \xrightarrow{S} \Delta_o \neq (0, 0, 0, 0)$

$\Delta_i = (0, 0, 0, 1) \xrightarrow{S} \Delta_o = (?, 1, ?, ?)$

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Δ_i	0	16	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	1	0	0	0	0	2	2	2	2	0	0	0	0	2	2	2	2
	2	0	2	0	2	0	0	0	4	0	2	2	0	0	0	2	2
	3	0	2	0	2	0	0	4	0	0	2	2	0	0	0	2	2
x	4	0	0	0	0	0	0	0	0	0	0	4	4	2	2	2	2
	5	0	0	0	0	2	2	2	2	0	0	4	4	0	0	0	0
	6	0	2	0	2	0	4	0	0	0	2	2	0	2	2	0	0
	7	0	2	0	2	4	0	0	0	0	2	2	0	2	2	0	0
S	8	0	0	0	0	0	0	0	0	0	0	0	0	4	4	4	4
	9	0	4	4	0	0	0	0	0	4	0	4	0	0	0	0	0
	a	0	0	2	2	2	0	0	2	4	0	0	0	0	2	0	2
	b	0	0	2	2	0	2	2	0	4	0	0	0	2	0	2	0
$S(x)$	c	0	4	4	0	2	2	2	2	0	0	0	0	0	0	0	0
	d	0	0	0	0	2	2	2	2	0	4	0	4	0	0	0	0
	e	0	0	2	2	0	2	2	0	4	0	0	0	0	2	0	2
	f	0	0	2	2	2	0	0	2	4	0	0	0	2	0	2	0
Δ_o		1	?	?	?												

$\Delta_i = (0, 0, 0, 0) \xrightarrow{S} \Delta_o = (0, 0, 0, 0)$

$\Delta_i \neq (0, 0, 0, 0) \xrightarrow{S} \Delta_o \neq (0, 0, 0, 0)$

$\Delta_i = (0, 0, 0, 1) \xrightarrow{S} \Delta_o = (?, 1, ?, ?)$

$\Delta_i = (0, 1, 0, 0) \xrightarrow{S} \Delta_o = (1, ?, ?, ?)$

Deterministic Bit-Wise Differential Trails (a.k.a. Undisturbed Bits [Tez14])

x	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
$S(x)$	4	0	a	7	b	e	1	d	9	f	6	8	5	2	c	3

$\Delta_i \backslash \Delta_o$	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
Δ_i	0	16	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	1	0	0	0	0	2	2	2	2	0	0	0	0	2	2	2
	2	0	2	0	2	0	0	0	4	0	2	2	0	0	0	2
	3	0	2	0	2	0	0	4	0	0	2	2	0	0	0	2
x	4	0	0	0	0	0	0	0	0	0	4	4	2	2	2	2
	5	0	0	0	0	2	2	2	2	0	0	4	4	0	0	0
	6	0	2	0	2	0	4	0	0	0	2	2	0	2	2	0
	7	0	2	0	2	4	0	0	0	0	2	2	0	2	2	0
	8	0	0	0	0	0	0	0	0	0	0	0	4	4	4	4
	9	0	4	4	0	0	0	0	0	4	0	4	0	0	0	0
	a	0	0	2	2	2	0	0	2	4	0	0	0	0	2	0
	b	0	0	2	2	0	2	2	0	4	0	0	0	2	0	2
	c	0	4	4	0	2	2	2	2	0	0	0	0	0	0	0
	d	0	0	0	0	2	2	2	2	0	4	0	4	0	0	0
	e	0	0	2	2	0	2	2	0	4	0	0	0	2	0	2
	f	0	0	2	2	2	0	0	2	4	0	0	0	2	0	2

$$\Delta_i = (0, 0, 0, 0) \xrightarrow{S} \Delta_o = (0, 0, 0, 0)$$

$$\Delta_i \neq (0, 0, 0, 0) \xrightarrow{S} \Delta_o \neq (0, 0, 0, 0)$$

$$\Delta_i = (0, 0, 0, 1) \xrightarrow{S} \Delta_o = (?, 1, ?, ?)$$

$$\Delta_i = (0, 1, 0, 0) \xrightarrow{S} \Delta_o = (1, ?, ?, ?)$$

$$\Delta_i = (1, 0, 0, 0) \xrightarrow{S} \Delta_o = (1, 1, ?, ?)$$

Deterministic Bit-Wise Differential Trails (a.k.a. Undisturbed Bits [Tez14])

x	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
$S(x)$	4	0	a	7	b	e	1	d	9	f	6	8	5	2	c	3

$\Delta_i \backslash \Delta_o$	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
Δ_i	0	16	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	1	0	0	0	0	2	2	2	2	0	0	0	0	2	2	2
	2	0	2	0	2	0	0	0	4	0	2	2	0	0	0	2
	3	0	2	0	2	0	0	4	0	0	2	2	0	0	0	2
x	4	0	0	0	0	0	0	0	0	0	4	4	2	2	2	2
	5	0	0	0	0	2	2	2	2	0	0	4	4	0	0	0
	6	0	2	0	2	0	4	0	0	0	2	2	0	2	2	0
	7	0	2	0	2	4	0	0	0	0	2	2	0	2	2	0
	8	0	0	0	0	0	0	0	0	0	0	0	4	4	4	4
	9	0	4	4	0	0	0	0	0	4	0	4	0	0	0	0
	a	0	0	2	2	2	0	0	2	4	0	0	0	0	2	0
	b	0	0	2	2	0	2	2	0	4	0	0	0	2	0	2
	c	0	4	4	0	2	2	2	2	0	0	0	0	0	0	0
	d	0	0	0	0	2	2	2	2	0	4	0	4	0	0	0
	e	0	0	2	2	0	2	2	0	4	0	0	0	2	0	2
	f	0	0	2	2	2	0	0	2	4	0	0	0	2	0	2

$$\Delta_i = (0, 0, 0, 0) \xrightarrow{S} \Delta_o = (0, 0, 0, 0)$$

$$\Delta_i \neq (0, 0, 0, 0) \xrightarrow{S} \Delta_o \neq (0, 0, 0, 0)$$

$$\Delta_i = (0, 0, 0, 1) \xrightarrow{S} \Delta_o = (?, 1, ?, ?)$$

$$\Delta_i = (0, 1, 0, 0) \xrightarrow{S} \Delta_o = (1, ?, ?, ?)$$

$$\Delta_i = (1, 0, 0, 0) \xrightarrow{S} \Delta_o = (1, 1, ?, ?)$$

$$\Delta_i = (1, 0, 0, 1) \xrightarrow{S} \Delta_o = (?, 0, ?, ?)$$

Deterministic Bit-Wise Differential Trails (a.k.a. Undisturbed Bits [Tez14])

x	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
$S(x)$	4	0	a	7	b	e	1	d	9	f	6	8	5	2	c	3

$\Delta_i \backslash \Delta_o$	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
Δ_i	0	16	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	1	0	0	0	0	2	2	2	2	0	0	0	0	2	2	2
	2	0	2	0	2	0	0	0	4	0	2	2	0	0	0	2
	3	0	2	0	2	0	0	4	0	0	2	2	0	0	0	2
x	4	0	0	0	0	0	0	0	0	0	4	4	2	2	2	2
	5	0	0	0	0	2	2	2	2	0	0	4	4	0	0	0
	6	0	2	0	2	0	4	0	0	0	2	2	0	2	2	0
	7	0	2	0	2	4	0	0	0	0	2	2	0	2	2	0
	8	0	0	0	0	0	0	0	0	0	0	0	4	4	4	4
	9	0	4	4	0	0	0	0	0	4	0	4	0	0	0	0
	a	0	0	2	2	2	0	0	2	4	0	0	0	0	2	0
	b	0	0	2	2	0	2	2	0	4	0	0	0	2	0	2
	c	0	4	4	0	2	2	2	2	0	0	0	0	0	0	0
	d	0	0	0	0	2	2	2	2	0	4	0	4	0	0	0
	e	0	0	2	2	0	2	2	0	4	0	0	0	2	0	2
	f	0	0	2	2	2	0	0	2	4	0	0	0	2	0	2

$$\Delta_i = (0, 0, 0, 0) \xrightarrow{S} \Delta_o = (0, 0, 0, 0)$$

$$\Delta_i \neq (0, 0, 0, 0) \xrightarrow{S} \Delta_o \neq (0, 0, 0, 0)$$

$$\Delta_i = (0, 0, 0, 1) \xrightarrow{S} \Delta_o = (?, 1, ?, ?)$$

$$\Delta_i = (0, 1, 0, 0) \xrightarrow{S} \Delta_o = (1, ?, ?, ?)$$

$$\Delta_i = (1, 0, 0, 0) \xrightarrow{S} \Delta_o = (1, 1, ?, ?)$$

$$\Delta_i = (1, 0, 0, 1) \xrightarrow{S} \Delta_o = (?, 0, ?, ?)$$

$$\Delta_i = (1, 1, 0, 0) \xrightarrow{S} \Delta_o = (0, ?, ?, ?)$$

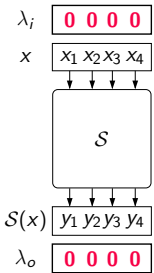
Deterministic Bit-Wise Linear Trails

x	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
$S(x)$	4	0	a	7	b	e	1	d	9	f	6	8	5	2	c	3

$\lambda_i \setminus \lambda_o$	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
0	16	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	4	-4	0	-8	-4	-4	0	0	4	-4	-8	0	4	4
2	0	0	0	0	0	0	0	0	0	8	8	0	0	8	-8	0
3	0	-8	4	4	0	0	-4	4	0	0	-4	4	-8	0	-4	-4
4	0	4	0	4	0	4	8	-4	0	4	0	4	-8	-4	0	4
5	0	4	-4	-8	0	-4	-4	0	0	4	-4	8	0	-4	-4	0
6	0	-4	8	4	0	-4	0	-4	0	4	0	4	8	-4	0	4
7	0	4	4	0	0	-4	4	-8	0	-4	-4	0	0	4	-4	-8
8	0	0	0	0	0	0	0	0	0	0	8	8	0	0	8	-8
9	0	0	-4	4	8	0	-4	-4	0	0	4	-4	0	-8	-4	-4
a	0	8	0	8	0	-8	0	8	0	0	0	0	0	0	0	0
b	0	0	-4	4	-8	0	-4	-4	0	8	-4	-4	0	0	4	-4
c	0	4	0	4	0	4	-8	-4	8	-4	0	4	0	4	0	4
d	0	4	4	0	-8	4	-4	0	-8	-4	4	0	0	-4	-4	0
e	0	4	8	-4	0	4	0	4	8	4	0	-4	0	-4	0	-4
f	0	-4	-4	0	-8	-4	4	0	8	-4	4	0	0	-4	-4	0

$$\lambda_i = (0, 0, 0, 0) \xrightarrow{S} \lambda_o = (0, 0, 0, 0)$$

$$\lambda_i \neq (0, 0, 0, 0) \xrightarrow{S} \lambda_o \neq (0, 0, 0, 0)$$



Deterministic Bit-Wise Linear Trails

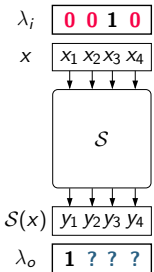
x	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
$S(x)$	4	0	a	7	b	e	1	d	9	f	6	8	5	2	c	3

$\lambda_i \setminus \lambda_o$	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
0	16	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	4	-4	0	-8	-4	-4	0	0	4	-4	-8	0	4	4
2	0	0	0	0	0	0	0	0	0	8	8	0	0	8	-8	0
3	0	-8	4	4	0	0	-4	4	0	0	-4	4	-8	0	-4	-4
4	0	4	0	4	0	4	8	-4	0	4	0	4	-8	-4	0	4
5	0	4	-4	-8	0	-4	-4	0	0	4	-4	8	0	-4	-4	0
6	0	-4	8	4	0	-4	0	-4	0	4	0	4	8	-4	0	4
7	0	4	4	0	0	-4	4	-8	0	-4	-4	0	0	4	-4	-8
8	0	0	0	0	0	0	0	0	0	0	8	8	0	0	8	-8
9	0	0	-4	4	8	0	-4	-4	0	0	4	-4	0	-8	-4	-4
a	0	8	0	8	0	-8	0	8	0	0	0	0	0	0	0	0
b	0	0	-4	4	-8	0	-4	-4	0	8	-4	-4	0	0	4	-4
c	0	4	0	4	0	4	-8	-4	8	-4	0	4	0	4	0	4
d	0	4	4	0	-8	4	-4	0	-8	-4	4	0	0	-4	-4	0
e	0	4	8	-4	0	4	0	4	8	4	0	-4	0	-4	0	-4
f	0	-4	-4	0	-8	-4	4	0	8	-4	4	0	0	-4	-4	0

$$\lambda_i = (0, 0, 0, 0) \xrightarrow{S} \lambda_o = (0, 0, 0, 0)$$

$$\lambda_i \neq (0, 0, 0, 0) \xrightarrow{S} \lambda_o \neq (0, 0, 0, 0)$$

$$\lambda_i = (0, 0, 1, 0) \xrightarrow{S} \lambda_o = (1, ?, ?, ?)$$



Deterministic Bit-Wise Linear Trails

x	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
$S(x)$	4	0	a	7	b	e	1	d	9	f	6	8	5	2	c	3

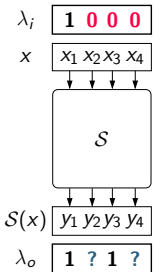
$\lambda_i \setminus \lambda_o$	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
0	16	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	4	-4	0	-8	-4	-4	0	0	4	-4	-8	0	4	4
2	0	0	0	0	0	0	0	0	0	8	8	0	0	8	-8	0
3	0	-8	4	4	0	0	-4	4	0	0	-4	4	-8	0	-4	-4
4	0	4	0	4	0	4	8	-4	0	4	0	4	-8	-4	0	4
5	0	4	-4	-8	0	-4	-4	0	0	4	-4	8	0	-4	-4	0
6	0	-4	8	4	0	-4	0	-4	0	4	0	4	8	-4	0	4
7	0	4	4	0	0	-4	4	-8	0	-4	-4	0	0	4	-4	-8
8	0	0	0	0	0	0	0	0	0	0	8	8	0	0	8	-8
9	0	0	-4	4	8	0	-4	-4	0	0	4	-4	0	-8	-4	-4
a	0	8	0	8	0	-8	0	8	0	0	0	0	0	0	0	0
b	0	0	-4	4	-8	0	-4	-4	0	8	-4	-4	0	0	4	-4
c	0	4	0	4	0	4	-8	-4	8	-4	0	4	0	4	0	4
d	0	4	4	0	-8	4	-4	0	-8	-4	4	0	0	-4	-4	0
e	0	4	8	-4	0	4	0	4	8	4	0	-4	0	-4	0	-4
f	0	-4	-4	0	-8	-4	4	0	8	-4	4	0	0	-4	-4	0

$$\lambda_i = (0, 0, 0, 0) \xrightarrow{S} \lambda_o = (0, 0, 0, 0)$$

$$\lambda_i \neq (0, 0, 0, 0) \xrightarrow{S} \lambda_o \neq (0, 0, 0, 0)$$

$$\lambda_i = (0, 0, 1, 0) \xrightarrow{S} \lambda_o = (1, ?, ?, ?)$$

$$\lambda_i = (1, 0, 0, 0) \xrightarrow{S} \lambda_o = (1, ?, 1, ?)$$



Deterministic Bit-Wise Linear Trails

x	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
$S(x)$	4	0	a	7	b	e	1	d	9	f	6	8	5	2	c	3

$\lambda_i \setminus \lambda_o$	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
0	16	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	4	-4	0	-8	-4	-4	0	0	4	-4	-8	0	4	4
2	0	0	0	0	0	0	0	0	0	8	8	0	0	8	-8	0
3	0	-8	4	4	0	0	-4	4	0	0	-4	4	-8	0	-4	-4
4	0	4	0	4	0	4	8	-4	0	4	0	4	-8	-4	0	4
5	0	4	-4	-8	0	-4	-4	0	0	4	-4	8	0	-4	-4	0
6	0	-4	8	4	0	-4	0	-4	0	4	0	4	8	-4	0	4
7	0	4	4	0	0	-4	4	-8	0	-4	-4	0	0	4	-4	-8
8	0	0	0	0	0	0	0	0	0	0	8	8	0	0	8	-8
9	0	0	-4	4	8	0	-4	-4	0	0	4	-4	0	-8	-4	-4
a	0	8	0	8	0	-8	0	8	0	0	0	0	0	0	0	0
b	0	0	-4	4	-8	0	-4	-4	0	8	-4	-4	0	0	4	-4
c	0	4	0	4	0	4	-8	-4	8	-4	0	4	0	4	0	4
d	0	4	4	0	-8	4	-4	0	-8	-4	4	0	0	-4	-4	0
e	0	4	8	-4	0	4	0	4	8	4	0	-4	0	-4	0	-4
f	0	-4	-4	0	-8	-4	4	0	8	-4	4	0	0	-4	-4	0

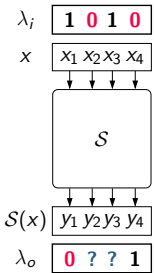
$$\lambda_i = (0, 0, 0, 0) \xrightarrow{S} \lambda_o = (0, 0, 0, 0)$$

$$\lambda_i \neq (0, 0, 0, 0) \xrightarrow{S} \lambda_o \neq (0, 0, 0, 0)$$

$$\lambda_i = (0, 0, 1, 0) \xrightarrow{S} \lambda_o = (1, ?, ?, ?)$$

$$\lambda_i = (1, 0, 0, 0) \xrightarrow{S} \lambda_o = (1, ?, 1, ?)$$

$$\lambda_i = (1, 0, 1, 0) \xrightarrow{S} \lambda_o = (0, ?, ?, 1)$$



CP Model for Deterministic Bit-Wise Trails - I

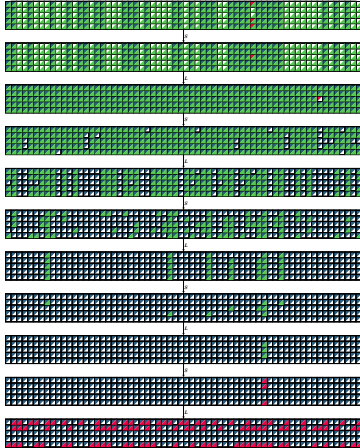
- For each bit position, we define an integer variable with domain $\{0, 1, -1\}$.
- Define CP constraints to model the propagation of deterministic bit-wise trails.



S-box

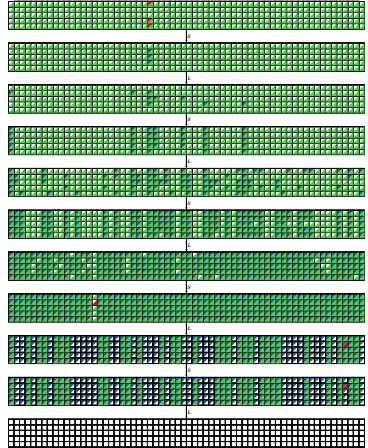
Assume that $x[i], y[i]$ are integer variables with domain $\{-1, 0, 1\}$ to encode the input and output differences at the i -th bit position, respectively. The valid deterministic differential transitions satisfy the following:


$$\left\{ \begin{array}{l} \text{if}(x[0] = 0 \wedge x[1] = 0 \wedge x[2] = 0 \wedge x[3] = 0) \text{ then } (y[0] = 0 \wedge y[1] = 0 \wedge y[2] = 0 \wedge y[3] = 0) \\ \text{elseif}(x[0] = 0 \wedge x[1] = 0 \wedge x[2] = 0 \wedge x[3] = 1) \text{ then } (y[0] = -1 \wedge y[1] = 1 \wedge y[2] = -1 \wedge y[3] = -1) \\ \text{elseif}(x[0] = 0 \wedge x[1] = 1 \wedge x[2] = 0 \wedge x[3] = 0) \text{ then } (y[0] = 1 \wedge y[1] = -1 \wedge y[2] = -1 \wedge y[3] = -1) \\ \text{elseif}(x[0] = 1 \wedge x[1] = 0 \wedge x[2] = 0 \wedge x[3] = 0) \text{ then } (y[0] = 1 \wedge y[1] = 1 \wedge y[2] = -1 \wedge y[3] = -1) \\ \text{elseif}(x[0] = 1 \wedge x[1] = 0 \wedge x[2] = 0 \wedge x[3] = 1) \text{ then } (y[0] = -1 \wedge y[1] = 0 \wedge y[2] = -1 \wedge y[3] = -1) \\ \text{elseif}(x[0] = 1 \wedge x[1] = 1 \wedge x[2] = 0 \wedge x[3] = 0) \text{ then } (y[0] = 0 \wedge y[1] = -1 \wedge y[2] = -1 \wedge y[3] = -1) \\ \text{else}(y[0] = -1 \wedge y[1] = -1 \wedge y[2] = -1 \wedge y[3] = -1) \text{ endif;} \end{array} \right.$$

Example: ID/ZC Distinguishers for 5 Rounds of Ascon



2^{155} ZC Distinguishers (upper/lower nonzero: /)



2^{155} ID Distinguishers (upper/lower unknown: /)

The Advantages of Our Method to Search for Distinguishers

- ✓ Based on satisfiability of the CP model
- ✓ Any feasible solutions of our CP model is a distinguisher
- ✓ We do not fix the input/output of distinguisher
- 💎 Extendable to a unified model for key-recovery
 - ✓ Enables us to find a distinguisher optimized for key-recovery
 - ✓ Enables us to consider key-recovery techniques:
 - ✓ MitM
 - ✓ Key bridging
 - ✓ *Partial-sum technique*

Our Unified CP Model for Partial-Sum Key-Recovery



Naive Approach v.s. Partial-Sum Technique



Naive approach:

☑ $x = F(k, c)$

☑ $T = N \cdot 2^{|k|}$



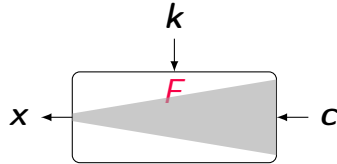
Partial-sum technique:

☑ $x_1 = f_1(k_1, x_0), x_2 = f_2(k_2, x_1), \dots, x = f_n(k_n, x_{n-1})$

☑ $x_0 = c, N_0 = N, N_i < N$

☑ $T = \sum_{i=1}^n \frac{N_{i-1}}{n} \cdot 2^{|k_1| + \dots + |k_i|} < \sum_{i=1}^n \frac{N}{n} \cdot 2^{|k|}$

☑ $T < N \cdot 2^{|k|}$



Naive Approach v.s. Partial-Sum Technique



Naive approach:

✔ $x = F(k, c)$

✔ $T = N \cdot 2^{|k|}$



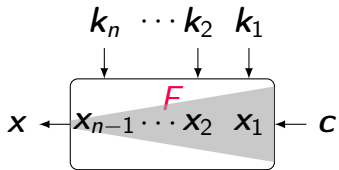
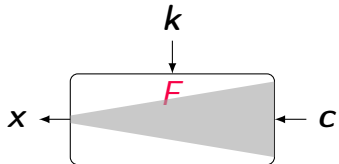
Partial-sum technique:

✔ $x_1 = f_1(k_1, x_0), x_2 = f_2(k_2, x_1), \dots, x = f_n(k_n, x_{n-1})$

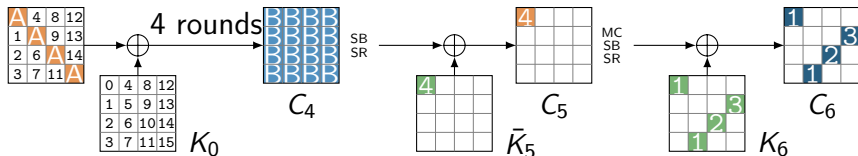
✔ $x_0 = c, N_0 = N, N_i < N$

✔ $T = \sum_{i=1}^n \frac{N_{i-1}}{n} \cdot 2^{|k_1| + \dots + |k_i|} < \sum_{i=1}^n \frac{N}{n} \cdot 2^{|k|}$

✔ $T < N \cdot 2^{|k|}$



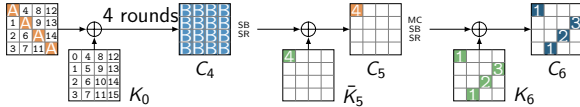
Example: Partial-Sum Integral Key Recovery for AES [Fer+00]



$$C_4[0] = \mathcal{S}^{-1} (\bar{K}_5[0] \oplus 0E \cdot \mathcal{S}^{-1} (C_6[0] \oplus K_6[0]) \oplus 09 \cdot \mathcal{S}^{-1} (C_6[7] \oplus K_6[7]) \\ \oplus 0D \cdot \mathcal{S}^{-1} (C_6[10] \oplus K_6[10]) \oplus 0B \cdot \mathcal{S}^{-1} (C_6[13] \oplus K_6[13]))$$

- Time complexity of naive key recovery: $6 \times 2^{32} \times 2^{40} \approx 2^{74.58}$

Partial-sum Technique for Integral Key Recovery [Fer+00]



- Guess $K_6[0, 7]$ and derive $\mathcal{S}_0 (C_6[0] \oplus K_6[0]) \oplus \mathcal{S}_1 (C_6[7] \oplus K_6[7])$
- Guess $K_6[10]$ and derive $\mathcal{S}_2 (C_6[10] \oplus K_6[10])$
- Guess $K_6[13]$ and derive $\mathcal{S}_3 (C_6[13] \oplus K_6[13])$
- Guess $\bar{K}_5[0]$ and derive $C_4[0]$
- Time complexity: $6 \times 4 \times 2^{48} \approx 2^{52}$ S-box lookups

Step 1: Key = 2^{16}

Data = 2^{32}

Time = 2^{48}

Step 2: Key = 2^{24}

Data = 2^{24}

Time = 2^{48}

Step 3: Key = 2^{32}

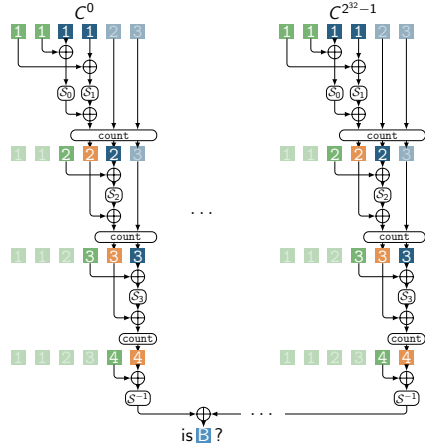
Data = 2^{16}

Time = 2^{48}

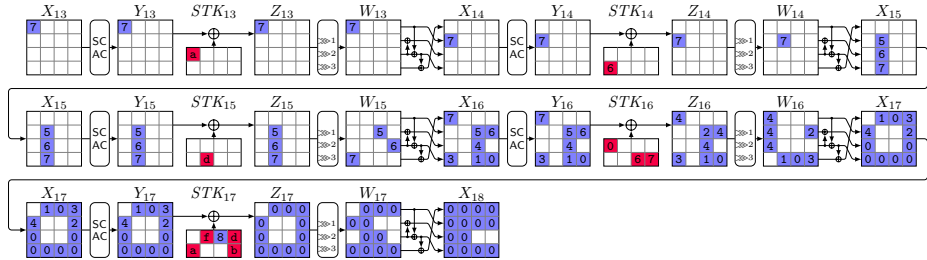
Step 4: Key = 2^{40}

Data = 2^8

Time = 2^{48}



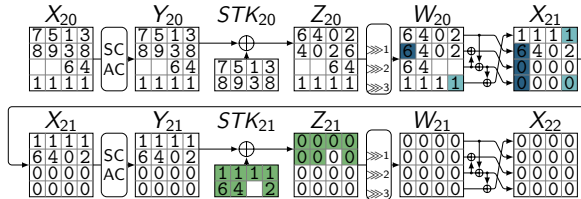
Our CP Model for Partial-Sum Technique - I



Step	Guessed	$K \times D = \text{Mem}$	Time	Stored Texts
0	–	$2^0 \times 2^{40} = 2^{40}$	$2^{40-5.2}$	$X_{17}[1, 3, 4, 7]; X_{17}[8, 11, 12, 13, 15]; X_{16}[15]$
1	$STK_{17}[1]$	$2^4 \times 2^{36} = 2^{40}$	$2^{44-7.2}$	$Z_{17}[3, 4, 7]; X_{17}[8, 11, 12, 15]; X_{16}[14, 15]$
2	$STK_{17}[7]$	$2^8 \times 2^{32} = 2^{40}$	$2^{44-8.2}$	$Z_{17}[3, 4]; X_{17}[8, 12, 15]; Z_{16}[6]; X_{16}[14, 15]$
3	$STK_{17}[3]$	$2^{12} \times 2^{28} = 2^{40}$	$2^{44-7.2}$	$Z_{17}[4]; X_{17}[8, 12]; Z_{16}[6]; X_{16}[12, 14, 15]$
4	$STK_{17}[4]$	$2^{16} \times 2^{28} = 2^{44}$	$2^{44-7.2}$	$Z_{16}[0, 6, 7]; X_{16}[10, 12, 14, 15]$
5	$STK_{16}[6]$	$2^{20} \times 2^{20} = 2^{40}$	$2^{48-7.2}$	$Z_{16}[0, 7]; X_{16}[12, 15]; X_{15}[5]$
6	$STK_{16}[7]$	$2^{24} \times 2^{16} = 2^{40}$	$2^{44-7.2}$	$Z_{16}[0]; X_{16}[12]; X_{15}[5, 9]$
7	$STK_{16}[0]$	$2^{28} \times 2^4 = 2^{32}$	$2^{44-6.2}$	$X_{13}[0]$
Σ		2^{44}	$2^{41.32}$	

Our CP Model for Partial-Sum Technique - II

- Assume that in each step we guess at least one cell of the involved keys.
- We define the number of steps s which is less than the number of involved key cells.
- For each cell we define an integer variable with domain $\{0, \dots, s\}$.
- We define some constraints to compute the step number of deriving each cell.

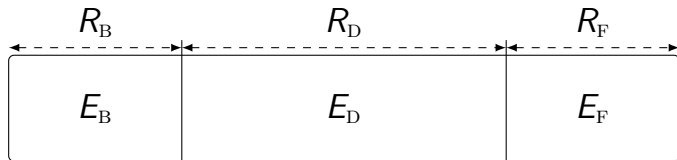


Our Unified Model for Finding Integral Attack

- Our CP model for finding complete integral attack includes the following modules:
 - Model the distinguisher part
 - Model the meet-in-the-middle technique
 - Model the involved cells in key recovery
 - Model the step assignment
 - Model the tweakey schedule (key-bridging)
 - Model the time/memory complexity evaluation
- Objective function: minimize the total time complexity

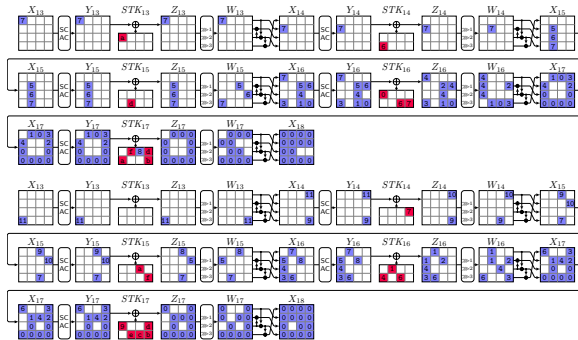
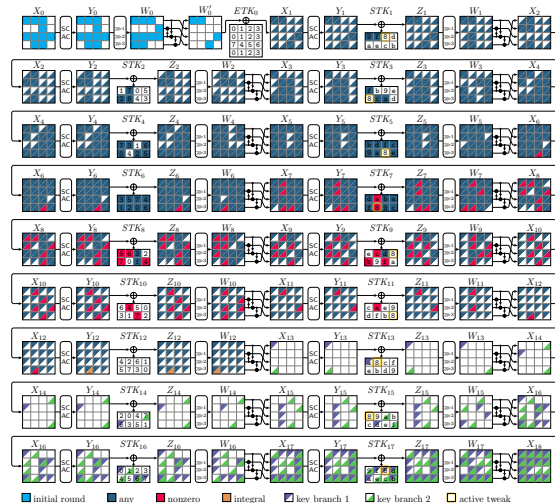
Usage of Our Tool

```
python3 attack.py -RB 1 -RD 12 -RF 5
```



- ✓ We use MiniZinc [Net+07] to create our CP models
- ✓ We use Gurobi [Gur22] and OrTools [PF] as the CP solvers
- 📖 Our tool can find the results in a few seconds running on a regular laptop

Example: 18-round Integral Attack on SKINNY- n - n



Contributions and Future Works



Contributions and Future Works

■ Contributions

- Improving unified models for finding complete ID/ZC/integral attacks
- Introducing a CP model for the partial-sum technique for the first time
- Found improved attacks for SKINNY, and ForskSKINNY, and QARMAv2

■ Future works

- A** Extending our distinguisher models for ID/ZC to find indirect contradictions
- A** Extending our tools to AndRX and ARX ciphers, e.g., Simeck, and SPECK.
- A** Extending our approach to division property or monomial prediction techniques
- A** Improving the key-recovery part of our CP models for ZC attacks

🐙: <https://github.com/hadipourh/zeroplus>

📄: <https://ia.cr/2023/1701>

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