SCIENCE
PASSION
TECHNOLOGY

Practical Multiple Persistent Fault Analysis

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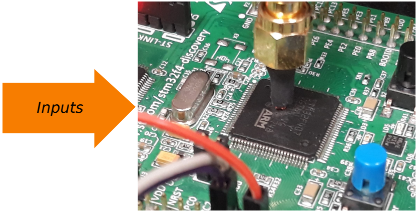
Outline

Introduction and the Research Gap

Fault Attacks

- **Fault attack**: An active side-channel attack [BDL97]:
  - **Fault injection**: Disturb the operation of a cryptographic device
  - **Fault analysis**: Analyze the erroneous outputs to retrieve the secret key



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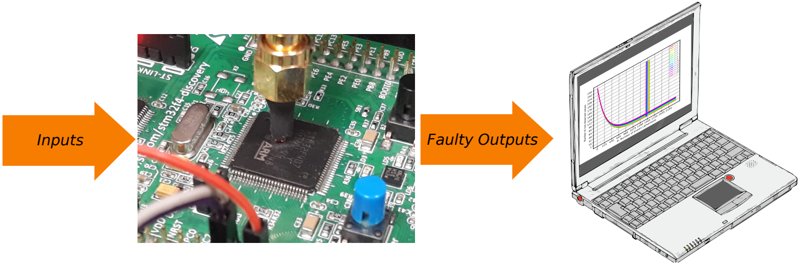

Persistent Fault Attack (PFA)

PFA fault model [Zha+18]:

- The injected faults are persistent until the reset of the device
- The injected faults typically alter the stored algorithm constants
- We can inject the faults before the encryption
- We can collect multiple faulty ciphertexts

x	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
$\mathcal{S}(x)$	6	4	c	5	0	7	2	e	1	f	3	d	8	a	9	b
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Core Idea of PFA

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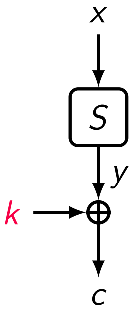

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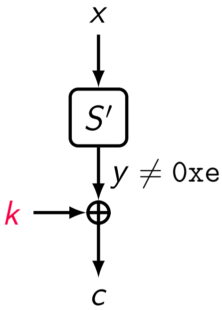

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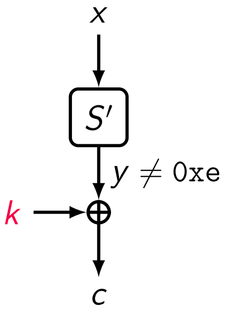

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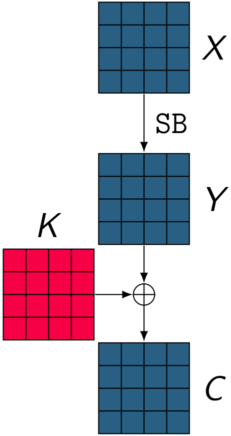

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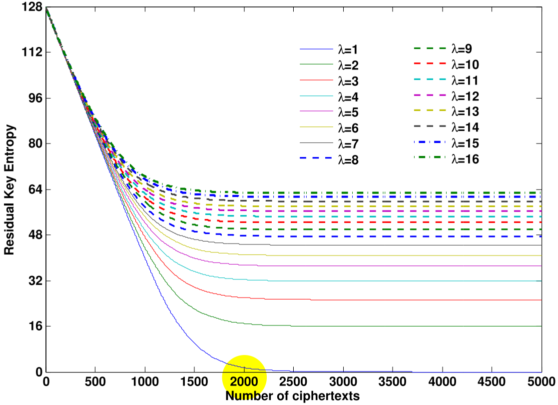
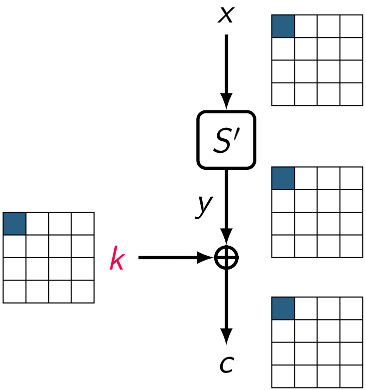
Limits of the Original PFA

PFA requires about 2000 faulty ciphertexts per key [Zha+18]





$$N \gtrapprox 2000 \Rightarrow |K| = \lambda^{16}$$

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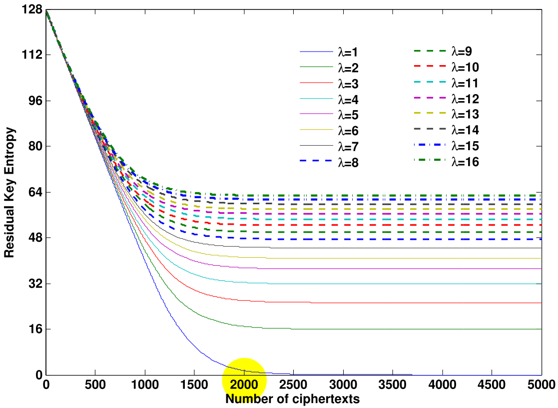
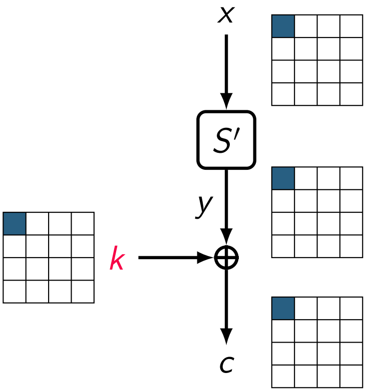
Limits of the Original PFA

PFA is very time consuming for multiple faults [Zha+18]





$\lambda = 12 \Rightarrow |K| = 12^{16} \approx \mathbf{2}^{57.36}$

More Limits of PFA and Its Enhanced Versions

- The location of the injected fault is supposed to be known
- For multiple fault injections:
  - We need a known plaintext/ciphertext pair to detect the correct key
- PFA only exploits the fault leakage in the last round
- Enhanced PFA (EPFA) [Xu+21] exploits the fault leakage in multiple rounds
- However, EPFA is not clear about exploiting multiple faults in deeper rounds
- Morever, EPFA still relies on the assumption of knowing the fault location

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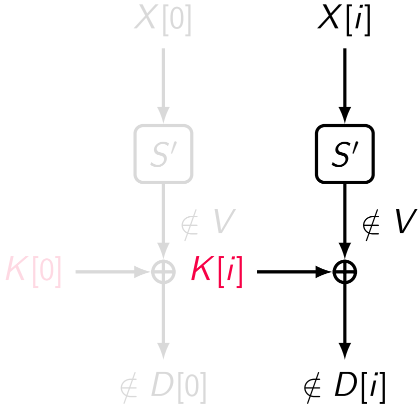
Our Framework for PFA With Multiple Faults

Core Idea

- V: Impossible values in the output of faulty S-box
- $D[i]$: Impossible values in the ith word of ciphertext
- $D[i] = V \oplus K[i]$ for all $i \in \{1, \ldots, 15\}$
- $V = K[0] \oplus D[0]$
- $D[i] = (K[0] \oplus K[i]) \oplus D[0]$
- $\delta[i] = K[0] \oplus K[i]$
- We can derive $\delta[i]$ from $(D[0], D[i])$



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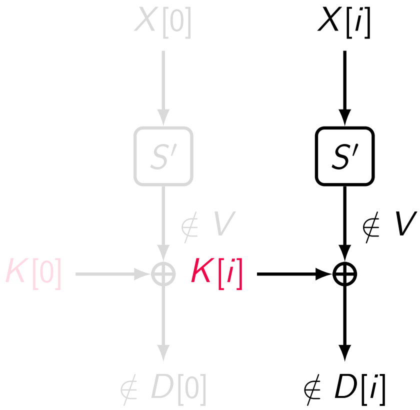

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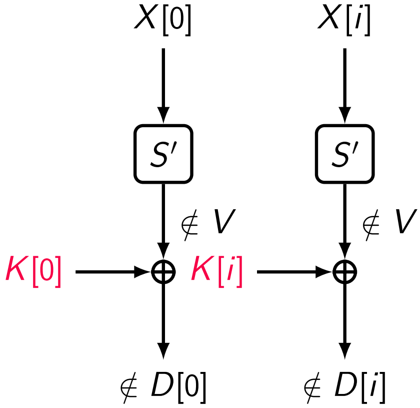

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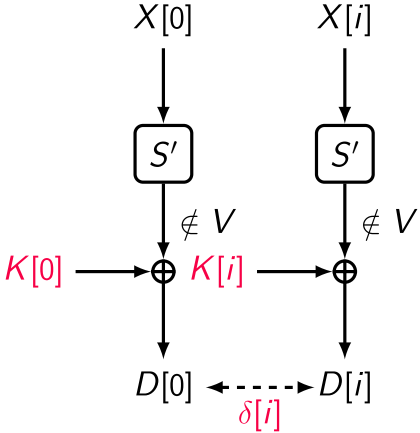
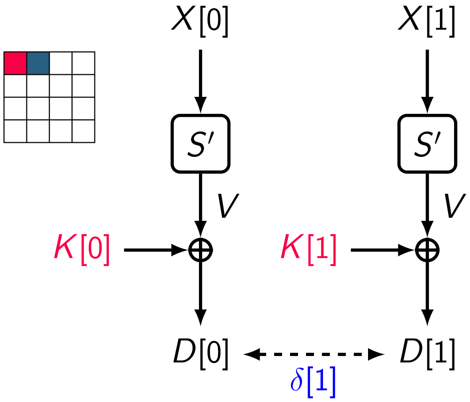


Reducing the Number of Key Candidates to 2^8



$\delta[1] = K[0] \oplus K[1]$

Guess $K[0]$ and determine $K[i]$ for all $i \in \{1, \ldots, 15\}$. So, $|K| = 2^8$ for AES!

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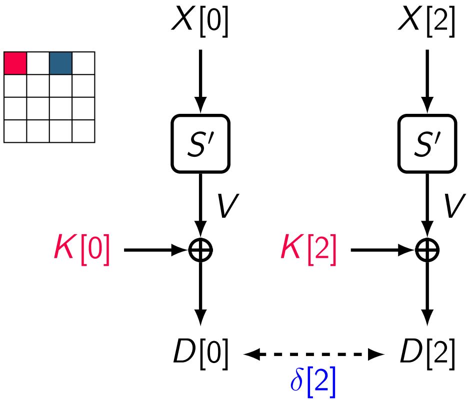

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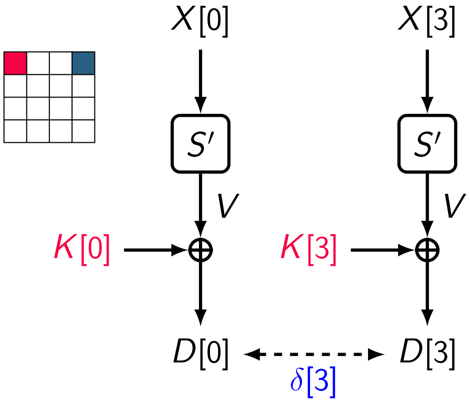

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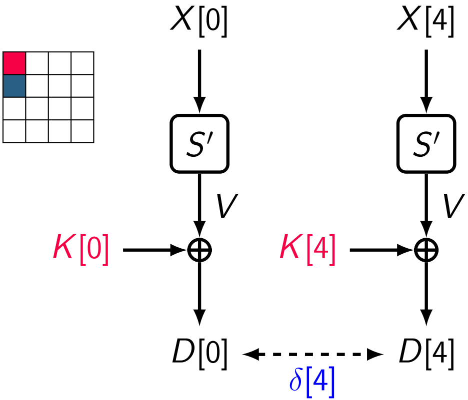

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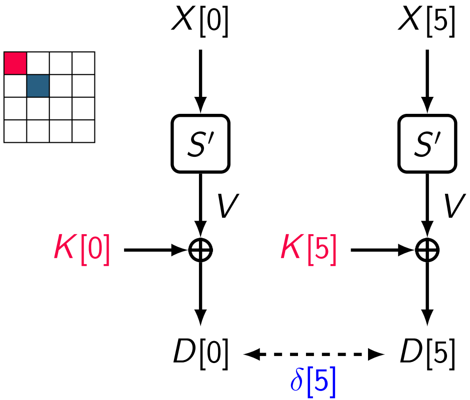

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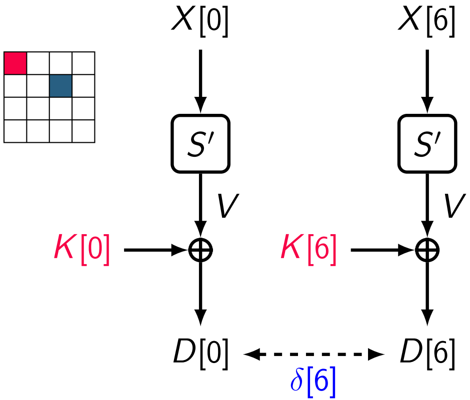

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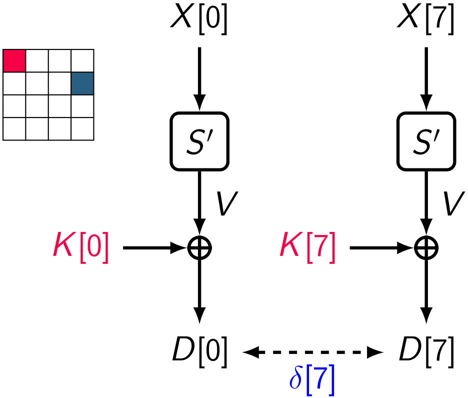

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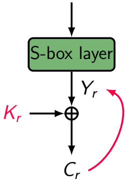
A Generic Key Recovery Framework

Going Deeper Into the Decryption Rounds

- For each key, compute the impossible values of S-box ($K \Rightarrow V$)
- Go deeper into the decryption to filter more wrong keys
- Challenge: the faulty S-box is not invertible
- We use the correct S-box for decryption
- We consider the wrong key assumption



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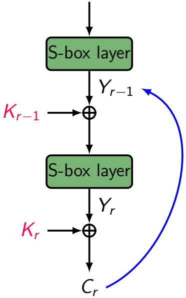

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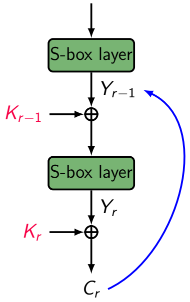

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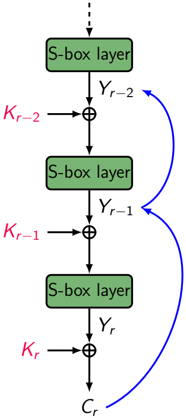

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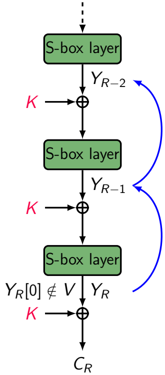
Our Key-recovery Framework

Input: Key candidates

Output: Master key

1 **for** *each key candidate* K **do**
2     $V \leftarrow K[0] \oplus D[0]$;
3     $\texttt{cnt}[K, V] \leftarrow 0$;
4   **foreach** *faulty ciphertext* **do**
5   **for** $r = R - 1, \ldots, 1$ **do**
6   Compute Y_r;
7   **foreach** *cell of* Y_r*, i.e.,* $Y_r[j]$ **do**
8   **if** $Y_r[j] \in V$ **then**
9   Go to line 4
10             $\texttt{cnt}[K, V] \leftarrow \texttt{cnt}[K, V] + 1$;
11 **return** key with maximum $\texttt{cnt}[K, V]$;



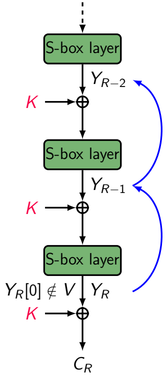
$$p = \left(1 - \frac{|V|}{256}\right)^{16}, \quad \texttt{cnt}_w = N \sum_{r=1}^{R-1} p^r, \quad \texttt{cnt}_c = N \sum_{r=1}^{R-1} p^r +$$

Hadi Soleimany, Nasour Bagheri, **Hosein Hadipour**, Prasanna Ravi, Shivam Bhasin, Sara Mansouri

Our Key-recovery Framework

Input: Key candidates
Output: Master key

1 **for** *each key candidate* K **do**
2 $\quad V \leftarrow K[0] \oplus D[0]$;
3 $\quad \texttt{cnt}[K, V] \leftarrow 0$;
4 $\quad$ **foreach** *faulty ciphertext* **do**
5 $\quad\quad$ **for** $r = R - 1, \ldots, 1$ **do**
6 $\quad\quad\quad$ Compute Y_r;
7 $\quad\quad\quad$ **foreach** *cell of* Y_r*, i.e.,* $Y_r[j]$ **do**
8 $\quad\quad\quad\quad$ **if** $Y_r[j] \in V$ **then**
9 $\quad\quad\quad\quad\quad$ Go to line 4
10 $\quad\quad\quad$ $\texttt{cnt}[K, V] \leftarrow \texttt{cnt}[K, V] + 1$;
11 **return** key with maximum $\texttt{cnt}[K, V]$;



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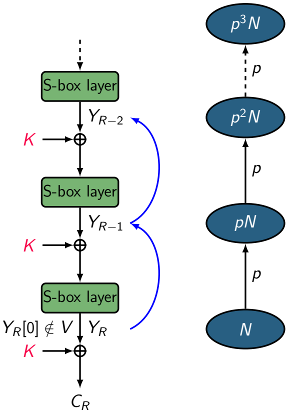

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Hadi Soleimany, Nasour Bagheri, **Hosein Hadipour**, Prasanna Ravi, Shivam Bhasin, Sara Mansouri
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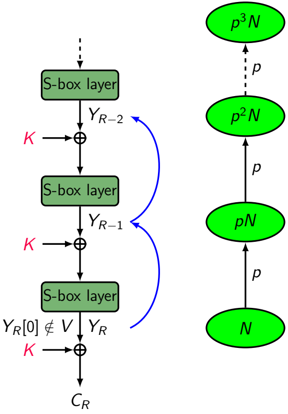

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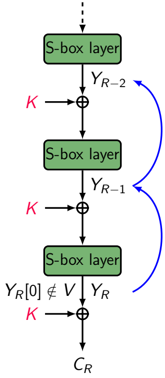
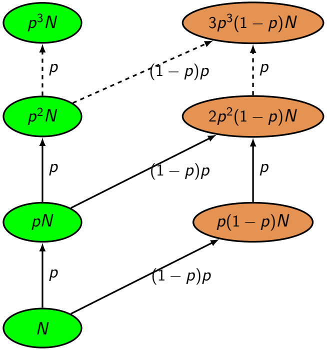
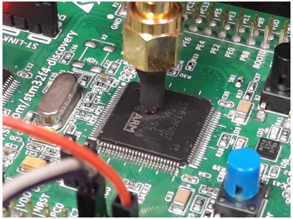
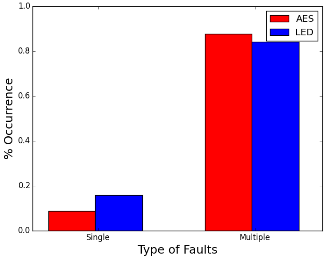




$$p = \left(1 - \frac{|V|}{256}\right)^{16}, \quad \boxed{\texttt{cnt}_w = N\sum_{r=1}^{R-1} p^r}, \quad \texttt{cnt}_c = \boxed{N\sum_{r=1}^{R-1} p^r} + \boxed{N\sum_{r=1}^{R-1} rp^r(1-p)}$$

Experimental Verification





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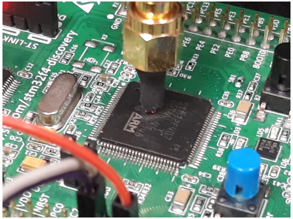

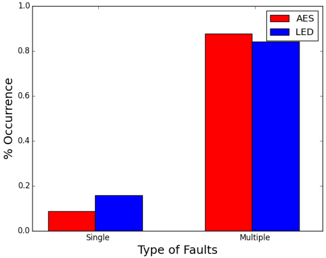

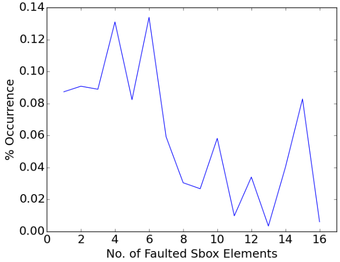

$$\lambda = 6, \ N = 1526, \ |K| = 256$$

Exp: $\texttt{cnt}_w = 3197.91, \ \texttt{cnt}_c = 6086.93$

The: $\texttt{cnt}_w = 3197.89, \ \texttt{cnt}_c = 6983.73$

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Conclusion

Hadi Soleimany, Nasour Bagheri, **Hosein Hadipour**, Prasanna Ravi, Shivam Bhasin, Sara Mansouri

Our Main Contributions

- We removed the assumption of knowing the fault location in PFA
- Our new technique decreases the number of key candidates by a factor of $\approx 2^{50}$
- We exploit the fault leakages in deeper rounds (until the first round)
- Our new technique reduces the number of required ciphertexts (refer to our paper)

Thanks for your attention!

https://github.com/hadipourh/faultyaes

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