



Applications to AES, Ascon, CLEFIA, SKINNY, PRESENT, KNOT, TWINE, WARP,

LBlock, Simeck, and SERPENT

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Research Gap and Our Contributions

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 - **②** How to formulate the correlation for more than one S-box layer?
 - ❷ How to (efficiently) find good DL distinguishers?
- Contributions
 - igspace Generalizing the DLCT framework [Bar+19] to handle multiple rounds
 - Introducing an efficient method to search for DL distinguishers applicable to
 - Strongly aligned SPN primitives: AES, SKINNY
 - Weakly aligned SPN primitives: Ascon, SERPENT, KNOT, PRESENT
 - Feistel structures: CLEFIA, TWINE, LBlock, LBlock-s, WARP
 - AndRX designs: Simeck

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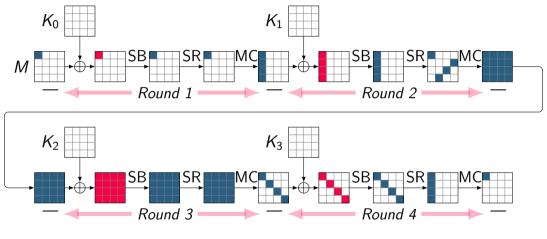
Outline

- Some Motivating Examples
- 2 Boomerang Analysis
- 3 Differential-Linear Cryptanalysis
- 4 Generalized DLCT Framework
- 5 Differential-Linear Switches and Deterministic Trails
- 6 Automatic Tools to Search for DL Distinguishers
- Contributions and Future Works

Some Motivating Examples

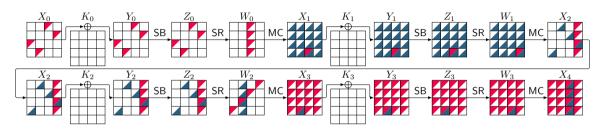


Security of AES Against Differential/Linear Attacks



$$\mathbb{P}_{4 \text{ rounds}} \leq 2^{-150}, \ \mathbb{C}_{4 \text{ rounds}}^2 \leq 2^{-150}$$

A 4-round DL Distinguisher for AES



$$r_u = 1, r_m = 3, r_\ell = 0, \ p = 2^{-24.00}, \ r = 2^{-7.66}, q^2 = 1, \ prq^2 = 2^{-31.66}$$

 ΔX_0 00005200000000f58f000000007b0000 ΓX_4 0032000000ab00000066000000980000

2^{63.32} v.s. 2¹⁵⁰

Boomerang Analysis



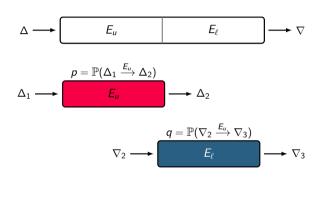
$$\Delta \longrightarrow \left[E : \mathbb{F}_2^n \to \mathbb{F}_2^n \right] \longrightarrow \nabla$$

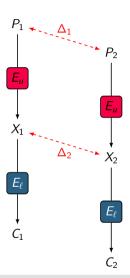
$$0 < \mathbb{P}(\Delta \xrightarrow{E} \nabla) \ll 2^{-n}$$

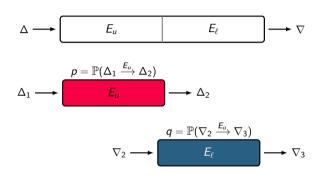
$$\Delta \longrightarrow \boxed{E_u \qquad E_\ell} \longrightarrow \nabla$$

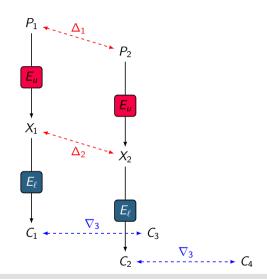
$$D_{\ell} \longrightarrow D_{\ell} \longrightarrow D_{\ell}$$

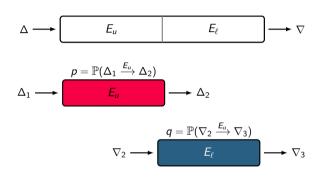
$$D_{\ell}$$

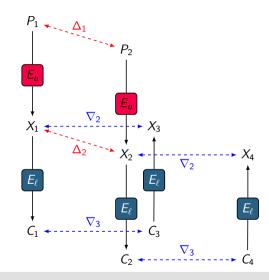


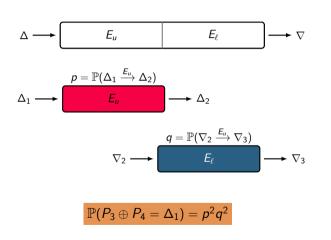


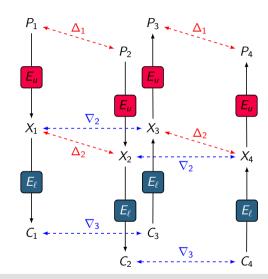




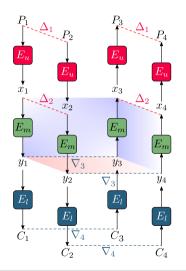






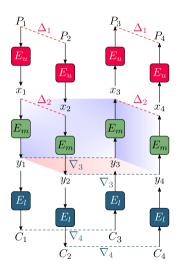


Sandwiching the Differentials! [DKS10; DKS14]



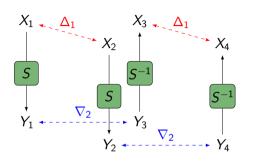


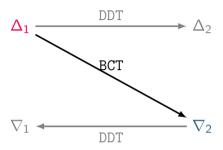
Sandwiching the Differentials! [DKS10; DKS14]



$$\mathbb{P}(P_3 \oplus P_4 = \Delta_1) \approx p^2 \times r \times q^2$$
$$r = \mathbb{P}(\Delta_2 \rightleftharpoons \nabla_3)$$

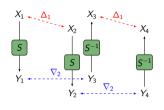
Boomerang Connectivity Table (BCT) [Cid+18]





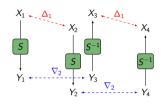
$$\mathrm{BCT}(\underline{\Delta_1},\nabla_2) \coloneqq \#\{X \in \mathbb{F}_2^n \,|\, S^{-1}\left(S(X) \oplus \nabla_2\right) \oplus S^{-1}(S(X \oplus \underline{\Delta_1}) \oplus \nabla_2) = \underline{\Delta_1}\}$$

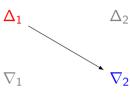
$$\mathbb{P}(\Delta_1 \rightleftarrows \nabla_2) = 2^{-n} \cdot \mathrm{BCT}(\Delta_1, \nabla_2)$$



$$\Delta_1 \longrightarrow \Delta_2$$

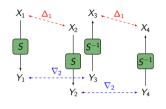
$$\nabla_1 \longleftarrow \nabla_2$$

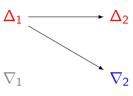




[DDV20; SQH19]

[Bou+20; DDV20]





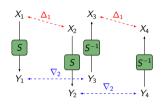
- $\mathcal{X}_{\mathrm{BCT}}(\Delta_1, \nabla_2) = \{x : S^{-1}(S(x) \oplus \nabla_2) \oplus S^{-1}(S(x \oplus \Delta_1) \oplus \nabla_2) = \Delta_1\}, \ \mathrm{BCT}(\Delta_1, \nabla_2) = \#\mathcal{X}_{\mathrm{BCT}}(\Delta_1, \nabla_2) = \#\mathcal{X}_{\mathrm{BCT}}(\Delta_1,$

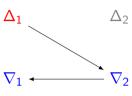
[WP19]

LBCT $(\Delta_1, \nabla_1, \nabla_2) = \#\{x : x \in \mathcal{X}_{BCT}(\Delta_1, \nabla_2) \cap \mathcal{X}_{DDT}(\nabla_1, \nabla_2)\}$

[DDV20; SQH19]

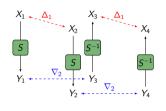
[Bou+20; DDV20]

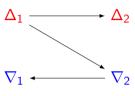




- $\mathcal{X}_{\mathtt{DDT}}(\Delta_1, \Delta_2) = \{x : S(x) \oplus S(x \oplus \Delta_1) = \Delta_2\}, \quad \mathtt{DDT}(\Delta_1, \Delta_2) = \#\mathcal{X}_{\mathtt{DDT}}(\Delta_1, \Delta_2)$
- $\mathcal{X}_{\mathrm{BCT}}(\Delta_1, \nabla_2) = \{x : S^{-1}(S(x) \oplus \nabla_2) \oplus S^{-1}(S(x \oplus \Delta_1) \oplus \nabla_2) = \Delta_1\}, \ \mathrm{BCT}(\Delta_1, \nabla_2) = \#\mathcal{X}_{\mathrm{BCT}}(\Delta_1, \nabla_2) = \#\mathcal{X}_{\mathrm{BCT}}(\Delta_1,$

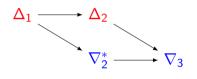
[DDV20; SQH19]



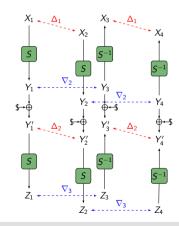


- $\mathcal{X}_{\mathrm{BCT}}(\Delta_1, \nabla_2) = \{x : S^{-1}(S(x) \oplus \nabla_2) \oplus S^{-1}(S(x \oplus \Delta_1) \oplus \nabla_2) = \Delta_1\}, \ \mathrm{BCT}(\Delta_1, \nabla_2) = \#\mathcal{X}_{\mathrm{BCT}}(\Delta_1, \nabla_2)$

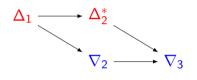
Double Boomerang Connectivity Table (DBCT) [HB21]



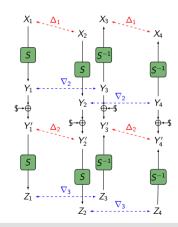
- igotagraphi DBCT $^{\vdash}(\Delta_1, \Delta_2, \nabla_3) = \sum_{\nabla_2} \mathtt{UBCT}(\Delta_1, \Delta_2, \nabla_2) \cdot \mathtt{LBCT}(\Delta_2, \nabla_2, \nabla_3)$
- igotimes DBCT $(\Delta_1, \nabla_3) = \sum_{\Delta_2}$ DBCT $^{\vdash}(\Delta_1, \Delta_2, \nabla_3) = \sum_{\nabla_2}$ DBCT $^{\dashv}(\Delta_1, \nabla_2, \nabla_3)$



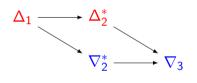
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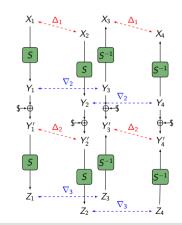
- $oldsymbol{oldsymbol{eta}}$ DBCT $^{\vdash}(\Delta_1,\Delta_2,
 abla_3) = \sum_{
 abla_2}$ UBCT $(\Delta_1,\Delta_2,
 abla_2) \cdot$ LBCT $(\Delta_2,
 abla_2,
 abla_3)$
- igotimes DBCT $(\Delta_1, \nabla_3) = \sum_{\Delta_2} ext{DBCT}^{\vdash}(\Delta_1, \Delta_2, \nabla_3) = \sum_{\nabla_2} ext{DBCT}^{\dashv}(\Delta_1, \nabla_2, \nabla_3)$



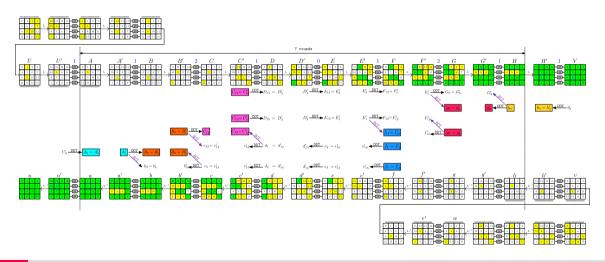
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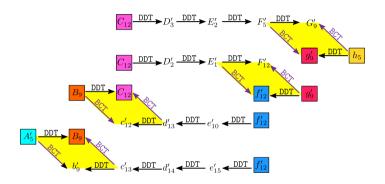
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Application of GBCT [HB21]



Application of GBCT [HB21]



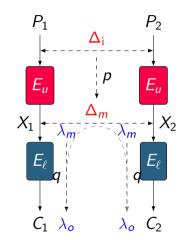
$$\begin{split} \text{DBCT}_{\text{total}} &= \text{DBCT}^{\vdash}(A_5, B_9, c_5) \cdot \text{DBCT}^{\vdash}(B_9, C_{12}, d_1) \cdot \text{DBCT}^{\dashv}(E_1', f_{12}', g_9') \cdot \text{DBCT}^{\dashv}(F_5', g_9', h_5) \\ \text{Pr}_{\text{total}} &= \text{Pr}(d_1 \xleftarrow{2 \text{ DDT}} f_{12}') \cdot \text{Pr}(c_5 \xleftarrow{3 \text{ DDT}} f_{12}') \cdot \text{Pr}(C_{12} \xrightarrow{2 \text{ DDT}} E_1') \cdot \text{Pr}(C_{12} \xrightarrow{3 \text{ DDT}} F_5') \\ r &= 2^{-8 \cdot n} \cdot \sum_{B_9} \sum_{C_{12}} \sum_{g_9'} \sum_{f_{12}'} \sum_{c_5} \sum_{d_1} \sum_{E_1'} \sum_{F_5'} \text{DBCT}_{\text{total}} \cdot \text{Pr}_{\text{total}}. \end{split}$$

Differential-Linear Cryptanalysis



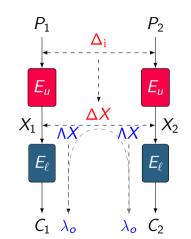
Differential-Linear (DL) Attack [LH94]

- $\blacksquare \quad \mathbb{P}(\Delta_{\mathrm{i}} \xrightarrow{E_u} \Delta_m) = p$
- $\qquad \mathbb{C}(\lambda_m \xrightarrow{E_\ell} \lambda_\mathrm{o}) = q$
- Assumptions $(\Delta X = X_1 \oplus X_2)$:
 - 1. E_u , and E_ℓ are statistically independent
 - 2. $\mathbb{P}(\lambda_m \cdot \Delta X = 0) = 1/2$ when $\Delta X \neq \Delta_m$
- $\mathbb{C} \left(\lambda_0 \cdot C_1 \oplus \lambda_0 \cdot C_2 \right) = (-1)^{\lambda_m \cdot \Delta_m} \cdot pq^2 = \pm pq^2$



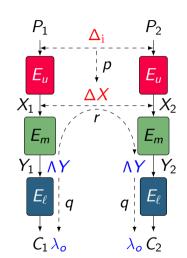
Differential-Linear Attack Revisited [BLN14; BLN17]

- Assumptions:
 - 1. E_u , and E_ℓ are statistically independent
- $\mathbb{C}(\lambda_{o}\cdot C_{1}\oplus\lambda_{o}\cdot C_{2})=\sum_{\Delta X,\Lambda X}\mathbb{C}(\Lambda X\cdot \Delta X)\cdot \mathbb{C}^{2}(\Lambda X,\lambda_{o})$



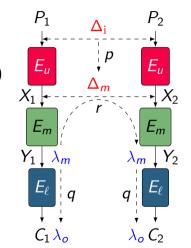
Sandwich Framework for DL Attack [DKS14; Bar+19]

- $\blacksquare \quad \mathbb{R}(\Delta X, \Lambda Y) = \mathbb{C}(\Lambda Y \cdot E_m(X) \oplus \Lambda Y \cdot E_m(X \oplus \Delta X))$
- $\qquad \mathbb{C}(\lambda_{o} \cdot \Delta C) = \sum_{\Delta X, \Lambda Y} \mathbb{P}(\Delta_{i}, \Delta X) \cdot \mathbb{R}(\Delta X, \Lambda Y) \cdot \mathbb{C}^{2}(\Lambda Y, \lambda_{o})$
- $\mathbb{R}(\Delta_m, \lambda_m) = r$
- $\mathbb{C}(\lambda_0 \cdot \Delta C) \approx prg^2$



Sandwich Framework for DL Attack [DKS14; Bar+19]

- $\blacksquare \quad \mathbb{R}(\Delta X, \Lambda Y) = \mathbb{C}(\Lambda Y \cdot E_m(X) \oplus \Lambda Y \cdot E_m(X \oplus \Delta X))$
- $\qquad \mathbb{C}(\lambda_{\mathrm{o}} \cdot \Delta C) = \sum_{\Delta X, \Lambda Y} \mathbb{P}(\Delta_{\mathrm{i}}, \Delta X) \cdot \mathbb{R}(\Delta X, \Lambda Y) \cdot \mathbb{C}^{2}(\Lambda Y, \lambda_{\mathrm{o}})$
- $\blacksquare \quad \mathbb{P}(\Delta_{\mathrm{i}} \xrightarrow{E_u} \Delta_m) = p$
- \blacksquare $\mathbb{R}(\Delta_m, \lambda_m) = r$
- $\blacksquare \quad \mathbb{C}(\lambda_m \xrightarrow{E_\ell} \lambda_{\mathrm{o}}) = q$
- $\mathbb{C}(\lambda_o \cdot \Delta C) \approx prq^2$



Differential-Linear Connectivity Table (DLCT)

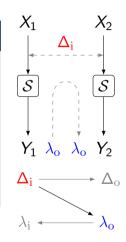
Differential-Linear Connectivity Table (DLCT) [Bar+19]

For a vectorial Boolean function $S: \mathbb{F}_2^n \to \mathbb{F}_2^m$, the DLCT of S is a $2^n \times 2^m$ table whose rows correspond to the input difference Δ_i to S and whose columns correspond to the output mask λ_o of S. The entry at index (Δ_i, λ_o) is

$$\mathtt{DLCT}(\Delta_{\mathrm{i}}, \lambda_{\mathrm{o}}) = |\mathtt{DLCT}_0(\Delta_{\mathrm{i}}, \lambda_{\mathrm{o}})| - |\mathtt{DLCT}_1(\Delta_{\mathrm{i}}, \lambda_{\mathrm{o}})|,$$

$$\text{where } \mathtt{DLCT}_b(\Delta_{\mathrm{i}},\lambda_{\mathrm{o}}) = \{x \in \mathbb{F}_2^n: \ \lambda_{\mathrm{o}} \cdot S(x) \oplus \lambda_{\mathrm{o}} \cdot S(x \oplus \Delta_{\mathrm{i}}) = b\}.$$

$$\mathbb{C}_{ exttt{DLCT}}\left(\Delta_{\mathrm{i}},\lambda_{\mathrm{o}}
ight)=2^{-n}\cdot exttt{DLCT}\left(\Delta_{\mathrm{i}},\lambda_{\mathrm{o}}
ight)$$



Generalized DLCT Framework



Upper Differential-Linear Connectivity Table (UDLCT)

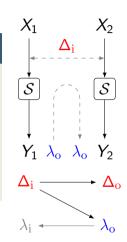
Upper Differential-Linear Connectivity Table (UDLCT)

For a vectorial Boolean function $S: \mathbb{F}_2^n \to \mathbb{F}_2^m$, the UDLCT of S is a $2^n \times 2^n \times 2^m$ table. The entry at index $(\Delta_i, \Delta_o, \lambda_o)$ is

$$\mathtt{UDLCT}(\Delta_{\mathrm{i}}, \Delta_{\mathrm{o}}, \lambda_{\mathrm{o}}) = |\mathtt{UDLCT}_0(\Delta_{\mathrm{i}}, \Delta_{\mathrm{o}}, \lambda_{\mathrm{o}})| - |\mathtt{UDLCT}_1(\Delta_{\mathrm{i}}, \Delta_{\mathrm{o}}, \lambda_{\mathrm{o}})|,$$

where
$$\mathrm{UDLCT}_b(\Delta_{\mathrm{i}}, \Delta_{\mathrm{o}}, \lambda_{\mathrm{o}}) = \{x \in \mathbb{F}_2^n : S(x) \oplus S(x \oplus \Delta_{\mathrm{i}}) = \Delta_{\mathrm{o}} \text{ and } \lambda_{\mathrm{o}} \cdot \Delta_{\mathrm{o}} = b\}.$$

$$\mathbb{C}_{\mathtt{UDLCT}}\left(\Delta_{\mathrm{i}}, \Delta_{\mathrm{o}}, \lambda_{\mathrm{o}}\right) = 2^{-n} \cdot \mathtt{UDLCT}\left(\Delta_{\mathrm{i}}, \Delta_{\mathrm{o}}, \lambda_{\mathrm{o}}\right)$$



Lower Differential-Linear Connectivity Table (LDLCT)

Lower Differential-Linear Connectivity Table (LDLCT)

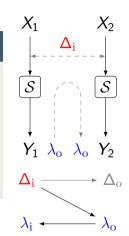
For a vectorial Boolean function $S: \mathbb{F}_2^n \to \mathbb{F}_2^m$, the LDLCT of S is a $2^n \times 2^m \times 2^m$ table. The entry at index $(\Delta_i, \lambda_i, \lambda_o)$ is

$$\texttt{LDLCT}(\Delta_i, \lambda_i, \lambda_o) = |\texttt{LDLCT}_0(\Delta_i, \lambda_i, \lambda_o)| - |\texttt{LDLCT}_1(\Delta_i, \lambda_i, \lambda_o)|,$$

where

$$\mathtt{LDLCT}_b(\Delta_{\mathbf{i}}, \lambda_{\mathbf{i}}, \lambda_{\mathbf{o}}) = \{x \in \mathbb{F}_2^n: \ \lambda_{\mathbf{i}} \cdot \Delta_{\mathbf{i}} \oplus \lambda_{\mathbf{o}} \cdot S(x) \oplus \lambda_{\mathbf{o}} \cdot S(x \oplus \Delta_{\mathbf{i}}) = b\}.$$

$$\mathbb{C}_{\mathtt{LDLCT}}\left(\Delta_{\mathrm{i}}, \lambda_{\mathrm{i}}, \lambda_{\mathrm{o}}\right) = 2^{-n} \cdot \mathtt{LDLCT}\left(\Delta_{\mathrm{i}}, \lambda_{\mathrm{i}}, \lambda_{\mathrm{o}}\right)$$



Extended Differential-Linear Connectivity Table (EDLCT)

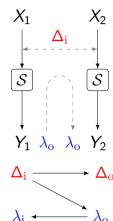
Extended Differential-Linear Connectivity Table (EDLCT)

For a vectorial Boolean function $S: \mathbb{F}_2^n \to \mathbb{F}_2^m$, the EDLCT of S is a $2^n \times 2^n \times 2^m \times 2^m$ table. The entry at index $(\Delta_i, \Delta_o, \lambda_i, \lambda_o)$ is

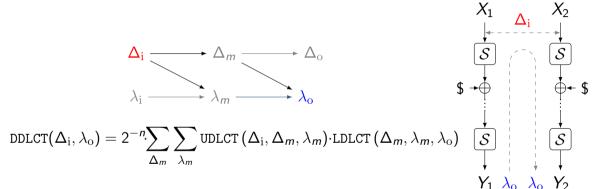
$$\mathtt{EDLCT}(\Delta_{\mathrm{i}}, \Delta_{\mathrm{o}}, \lambda_{\mathrm{i}}, \lambda_{\mathrm{o}}) \!=\! |\mathtt{EDLCT}_{0}(\Delta_{\mathrm{i}}, \! \Delta_{\mathrm{o}}, \! \lambda_{\mathrm{i}}, \! \lambda_{\mathrm{o}})| - |\mathtt{EDLCT}_{1}(\Delta_{\mathrm{i}}, \! \Delta_{\mathrm{o}}, \! \lambda_{\mathrm{i}}, \! \lambda_{\mathrm{o}})|,$$

where
$$\mathrm{EDLCT}_b(\Delta_\mathrm{i}, \Delta_\mathrm{o}, \lambda_\mathrm{i}, \lambda_\mathrm{o}) = \{x \in \mathbb{F}_2^n : S(x) \oplus S(x \oplus \Delta_\mathrm{i}) = \Delta_\mathrm{o} \text{ and } \lambda_\mathrm{i} \cdot \Delta_\mathrm{i} \oplus \lambda_\mathrm{o} \cdot \Delta_\mathrm{o} = b\}.$$

$$\mathbb{C}_{\mathtt{EDLCT}}\left(\Delta_{\mathrm{i}}, \Delta_{\mathrm{o}}, \lambda_{\mathrm{i}}, \lambda_{\mathrm{o}}\right) = 2^{-n} \cdot \mathtt{EDLCT}\left(\Delta_{\mathrm{i}}, \Delta_{\mathrm{o}}, \lambda_{\mathrm{i}}, \lambda_{\mathrm{o}}\right)$$

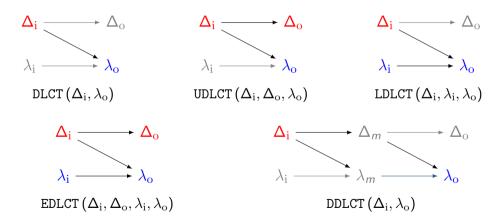


Double Differential-Linear Connectivity Table (DDLCT)

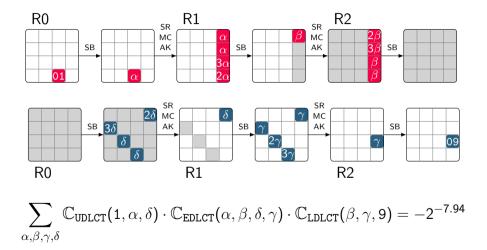


Generalized DLCT Framework (GBCT)

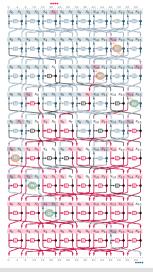
How to formulate the correlation for more than 1 round?



Application of the Generalized DLCT Tables - AES (- differential - linear)



Application of the Generalized DLCT Tables - TWINE (- differential - linear)



$$\begin{split} \mathbb{C}(\Delta_{\mathrm{i}},\lambda_{\mathrm{o}}) &= \sum_{\Delta_{m}} \mathbb{P}_{\mathtt{DDT}}(\Delta_{\mathrm{i}},\Delta_{m}) \cdot \mathbb{C}_{\mathtt{DDLCT}}\left(\Delta_{m},\lambda_{\mathrm{o}}\right) \\ &= \sum_{\lambda_{m}} \mathbb{C}_{\mathtt{DDLCT}}\left(\Delta_{\mathrm{i}},\lambda_{m}\right) \cdot \mathbb{C}_{\mathtt{LAT}}^{2}\left(\lambda_{m},\lambda_{\mathrm{o}}\right). \\ \mathbb{C}_{tot}(\Delta_{\mathrm{i}},\lambda_{\mathrm{o}}) &= \mathbb{C}^{2}(\Delta_{\mathrm{i}},\lambda_{\mathrm{o}}). \end{split}$$

$Input/Output\ Differences/Linear-mask$	Formula	Exp. Correlation
$\Delta_{ m i}(\Delta_{ m i},\lambda_{ m o})=$ (0xb4,0x67)	$-2^{-7.66}$	$-2^{-7.64}$
$(\Delta_{\mathrm{i}},\lambda_{\mathrm{o}})=(0\mathtt{x}\mathtt{02},0\mathtt{x}\mathtt{02})$	$-2^{-7.92}$	$-2^{-7.93}$
$(\Delta_{\mathrm{i}}, \lambda_{\mathrm{o}}) = (\texttt{0x55}, \texttt{0x55})$	$-2^{-7.99}$	$-2^{-7.98}$
$(\Delta_{ m i},\lambda_{ m o})=({ t Oxbf},{ t Oxef})$	$-2^{-8.05}$	$-2^{-8.06}$
$(\Delta_{ m i},\lambda_{ m o})=$ (0xfe,0x06)	$-2^{-8.26}$	$-2^{-8.25}$
$(\Delta_{\mathrm{i}},\lambda_{\mathrm{o}})=$ (0x4b,0x1a)	$-2^{-8.43}$	$-2^{-8.44}$

Differential-Linear Switches and Deterministic Trails



Cell-Wise and Bit-Wise Switches

X	0	1	2	3	4	5	6	7	8	9	а	b	С	d	е	f
S(x)	4	0	а	7	b	е	1	d	9	f	6	8	5	2	С	3

$\Delta \setminus \lambda$	0	1	2	3	4	5	6	7	8	9	a	b	С	d	е	f
0	16	16	16	16	16	16	16	16	16	16	16	16	16	16	16	16
1	16	0	0	0	-16	0	0	0	0	0	0	0	0	0	0	0
2	16	-8	-8	0	0	0	8	-8	0	-8	0	8	0	0	0	0
3	16	0	-8	-8	0	-8	8	0	0	0	0	0	0	-8	0	8
4	16	0	-8	0	0	0	-8	0	-16	0	8	0	0	0	8	0
5	16	0	-8	0	0	0	-8	0	0	0	8	0	-16	0	8	0
6	16	-8	8	-8	0	0	-8	0	0	-8	0	0	0	0	0	8
7	16	0	8	0	0	-8	-8	-8	0	0	0	8	0	-8	0	0
8	16	0	0	0	-16	0	0	0	-16	0	0	0	16	0	0	0
9	16	-8	0	-8	16	-8	0	-8	0	8	0	-8	0	8	0	-8
a	16	0	0	8	0	8	0	0	0	0	-8	0	0	-8	-8	-8
b	16	8	0	0	0	0	0	8	0	-8	-8	-8	0	0	-8	0
С	16	0	0	-8	0	0	0	-8	16	0	0	-8	0	0	0	-8
d	16	-8	0	0	0	-8	0	0	0	8	0	0	-16	8	0	0
е	16	0	0	0	0	8	0	8	0	0	-8	-8	0	-8	-8	0
f	16	8	0	8	0	0	0	0	0	-8	-8	0	0	0	-8	-8

- Cell-wise switches:
 - $ext{DLCT}(\Delta_{\mathrm{i}},0) = ext{DLCT}(0,\lambda_{\mathrm{o}}) = 2^n ext{ for a} \ \Delta_{\mathrm{i}},\lambda_{\mathrm{o}}$
- Bit-wise switches:

$$exttt{DLCT}(\Delta_{ ext{i}},\lambda_{ ext{o}})=\pm 2^n ext{ for } \Delta_{ ext{i}},\lambda_{ ext{o}}
eq 0$$

• Example: $\mathbb{C}(9,4) = \frac{16}{16}$

Cell-Wise and Bit-Wise Switches

x	0	1	2	3	4	5	6	7	8	9	а	b	С	d	е	f
S(x)	4	0	а	7	b	е	1	d	9	f	6	8	5	2	С	3

$\Delta \setminus \lambda$	0	1	2	3	4	5	6	7	8	9	a	b	С	d	е	f
0	16	16	16	16	16	16	16	16	16	16	16	16	16	16	16	16
1	16	0	0	0	-16	0	0	0	0	0	0	0	0	0	0	0
2	16	-8	-8	0	0	0	8	-8	0	-8	0	8	0	0	0	0
3	16	0	-8	-8	0	-8	8	0	0	0	0	0	0	-8	0	8
4	16	0	-8	0	0	0	-8	0	-16	0	8	0	0	0	8	0
5	16	0	-8	0	0	0	-8	0	0	0	8	0	-16	0	8	0
6	16	-8	8	-8	0	0	-8	0	0	-8	0	0	0	0	0	8
7	16	0	8	0	0	-8	-8	-8	0	0	0	8	0	-8	0	0
8	16	0	0	0	-16	0	0	0	-16	0	0	0	16	0	0	0
9	16	-8	0	-8	16	-8	0	-8	0	8	0	-8	0	8	0	-8
a	16	0	0	8	0	8	0	0	0	0	-8	0	0	-8	-8	-8
b	16	8	0	0	0	0	0	8	0	-8	-8	-8	0	0	-8	0
С	16	0	0	-8	0	0	0	-8	16	0	0	-8	0	0	0	-8
d	16	-8	0	0	0	-8	0	0	0	8	0	0	-16	8	0	0
е	16	0	0	0	0	8	0	8	0	0	-8	-8	0	-8	-8	0
f	16	8	0	8	0	0	0	0	0	-8	-8	0	0	0	-8	-8

- Cell-wise switches: $\text{DLCT}(\Delta_{\mathrm{i}},0) = \text{DLCT}(0,\lambda_{\mathrm{o}}) = 2^n \text{ for all } \Delta_{\mathrm{i}},\lambda_{\mathrm{o}}$
 - DLCT $(\Delta_i, \lambda_o) = \pm 2^n$ for $\Delta_i, \lambda_o \neq 0$
 - Example: $\mathbb{C}(9,4) = \frac{16}{16}$

Cell-Wise and Bit-Wise Switches

x	0	1	2	3	4	5	6	7	8	9	а	b	С	d	е	f
S(x)	4	0	а	7	b	е	1	d	9	f	6	8	5	2	С	3

$\Delta \setminus \lambda$	0	1	2	3	4	5	6	7	8	9	a	b	С	d	е	f
0	16	16	16	16	16	16	16	16	16	16	16	16	16	16	16	16
1	16	0	0	0	-16	0	0	0	0	0	0	0	0	0	0	0
2	16	-8	-8	0	0	0	8	-8	0	-8	0	8	0	0	0	0
3	16	0	-8	-8	0	-8	8	0	0	0	0	0	0	-8	0	8
4	16	0	-8	0	0	0	-8	0	-16	0	8	0	0	0	8	0
5	16	0	-8	0	0	0	-8	0	0	0	8	0	-16	0	8	0
6	16	-8	8	-8	0	0	-8	0	0	-8	0	0	0	0	0	8
7	16	0	8	0	0	-8	-8	-8	0	0	0	8	0	-8	0	0
8	16	0	0	0	-16	0	0	0	-16	0	0	0	16	0	0	0
9	16	-8	0	-8	16	-8	0	-8	0	8	0	-8	0	8	0	-8
a	16	0	0	8	0	8	0	0	0	0	-8	0	0	-8	-8	-8
b	16	8	0	0	0	0	0	8	0	-8	-8	-8	0	0	-8	0
С	16	0	0	-8	0	0	0	-8	16	0	0	-8	0	0	0	-8
d	16	-8	0	0	0	-8	0	0	0	8	0	0	-16	8	0	0
е	16	0	0	0	0	8	0	8	0	0	-8	-8	0	-8	-8	0
f	16	8	0	8	0	0	0	0	0	-8	-8	0	0	0	-8	-8

- Cell-wise switches: $ext{DLCT}(\Delta_{\mathrm{i}},0) = ext{DLCT}(0,\lambda_{\mathrm{o}}) = 2^n ext{ for all } \Delta_{\mathrm{i}},\lambda_{\mathrm{o}}$
 - Bit-wise switches: $\mathtt{DLCT}(\Delta_{\mathrm{i}},\lambda_{\mathrm{o}}) = \pm 2^{n} \text{ for } \Delta_{\mathrm{i}},\lambda_{\mathrm{o}} \neq 0$
 - Example: $\mathbb{C}(9,4) = \frac{16}{16}$

Deterministic Bit-Wise Differential Trails (Forward)

X	0	1	2	3	4	5	6	7	8	9	а	b	С	d	е	f
S(x)	4	0	а	7	b	е	1	d	9	f	6	8	5	2	С	3

$\Delta_i \setminus \Delta_o$	0	1	2	3	4	5	6	7	8	9	a	b	С	d	е	f
0	16	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	2	2	2	2	0	0	0	0	2	2	2	2
2	0	2	0	2	0	0	0	4	0	2	2	0	0	0	2	2
3	0	2	0	2	0	0	4	0	0	2	2	0	0	0	2	2
4	0	0	0	0	0	0	0	0	0	0	4	4	2	2	2	2
5	0	0	0	0	2	2	2	2	0	0	4	4	0	0	0	0
6	0	2	0	2	0	4	0	0	0	2	2	0	2	2	0	0
7	0	2	0	2	4	0	0	0	0	2	2	0	2	2	0	0
8	0	0	0	0	0	0	0	0	0	0	0	0	4	4	4	4
9	0	4	4	0	0	0	0	0	0	4	0	4	0	0	0	0
a	0	0	2	2	2	0	0	2	4	0	0	0	0	2	0	2
b	0	0	2	2	0	2	2	0	4	0	0	0	2	0	2	0
С	0	4	4	0	2	2	2	2	0	0	0	0	0	0	0	0
d	0	0	0	0	2	2	2	2	0	4	0	4	0	0	0	0
е	0	0	2	2	0	2	2	0	4	0	0	0	0	2	0	2
f	0	0	2	2	2	0	0	2	4	0	0	0	2	0	2	0

$$\Delta_{i} = (0,0,0,0) \xrightarrow{S} \Delta_{o} = (0,0,0,0)$$

$$\Delta_{i} = (0,0,0,1) \xrightarrow{S} \Delta_{o} = (?,1,?,?)$$

$$\Delta_{i} = (0,1,0,0) \xrightarrow{S} \Delta_{o} = (1,?,?,?)$$

$$\Delta_{i} = (1,0,0,0) \xrightarrow{S} \Delta_{o} = (1,1,?,?)$$

$$\Delta_{i} = (1,0,0,1) \xrightarrow{S} \Delta_{o} = (?,0,?,?)$$

$$\Delta_{i} = (1,1,0,0) \xrightarrow{S} \Delta_{o} = (0,?,?,?)$$

Deterministic Bit-Wise Linear Trails (Backward)

X	0	1	2	3	4	5	6	7	8	9	а	b	С	d	е	f
S(x)	4	0	а	7	b	е	1	d	9	f	6	8	5	2	С	3

$\lambda_i \setminus \lambda_o$	0	1	2	3	4	5	6	7	8	9	a	b	С	d	е	f
0	16	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	4	-4	0	-8	-4	-4	0	0	4	-4	-8	0	4	4
2	0	0	0	0	0	0	0	0	0	8	8	0	0	8	-8	0
3	0	-8	4	4	0	0	-4	4	0	0	-4	4	-8	0	-4	-4
4	0	4	0	4	0	4	8	-4	0	4	0	4	-8	-4	0	4
5	0	4	-4	-8	0	-4	-4	0	0	4	-4	8	0	-4	-4	0
6	0	-4	8	4	0	-4	0	-4	0	4	0	4	8	-4	0	4
7	0	4	4	0	0	-4	4	-8	0	-4	-4	0	0	4	-4	-8
8	0	0	0	0	0	0	0	0	0	0	8	8	0	0	8	-8
9	0	0	-4	4	8	0	-4	-4	0	0	4	-4	0	-8	-4	-4
a	0	8	0	8	0	-8	0	8	0	0	0	0	0	0	0	0
b	0	0	-4	4	-8	0	-4	-4	0	8	-4	-4	0	0	4	-4
С	0	4	0	4	0	4	-8	-4	8	-4	0	4	0	4	0	4
d	0	4	4	0	-8	4	-4	0	-8	-4	4	0	0	-4	-4	0
е	0	4	8	-4	0	4	0	4	8	4	0	-4	0	-4	0	-4
f	0	-4	-4	0	-8	-4	4	0	8	-4	4	0	0	-4	-4	0

$$\lambda_{i} = (1, ?, ?, 1) \stackrel{S}{\leftarrow} \lambda_{o} = (0, 1, 0, 0)$$

$$\lambda_{i} = (1, 1, ?, ?) \stackrel{S}{\leftarrow} \lambda_{o} = (1, 0, 0, 0)$$

$$\lambda_{i} = (0, ?, ?, ?) \stackrel{S}{\leftarrow} \lambda_{o} = (1, 1, 0, 0)$$

Bit-Wise Switches and Deterministic Trails

x																
S(x)	4	0	a	7	b	e	1	d	9	f	6	8	5	2	С	3

$\Delta \setminus \lambda$	0	1	2	3	4	5	6	7	8	9	a	b	С	d	е	f
0	16	16	16	16	16	16	16	16	16	16	16	16	16	16	16	16
1	16	0	0	0	-16	0	0	0	0	0	0	0	0	0	0	0
2	16	-8	-8	0	0	0	8	-8	0	-8	0	8	0	0	0	0
3	16	0	-8	-8	0	-8	8	0	0	0	0	0	0	-8	0	8
4	16	0	-8	0	0	0	-8	0	-16	0	8	0	0	0	8	0
5	16	0	-8	0	0	0	-8	0	0	0	8	0	-16	0	8	0
6	16	-8	8	-8	0	0	-8	0	0	-8	0	0	0	0	0	8
7	16	0	8	0	0	-8	-8	-8	0	0	0	8	0	-8	0	0
8	16	0	0	0	-16	0	0	0	-16	0	0	0	16	0	0	0
9	16	-8	0	-8	16	-8	0	-8	0	8	0	-8	0	8	0	-8
a	16	0	0	8	0	8	0	0	0	0	-8	0	0	-8	-8	-8
b	16	8	0	0	0	0	0	8	0	-8	-8	-8	0	0	-8	0
С	16	0	0	-8	0	0	0	-8	16	0	0	-8	0	0	0	-8
d	16	-8	0	0	0	-8	0	0	0	8	0	0	-16	8	0	0
е	16	0	0	0	0	8	0	8	0	0	-8	-8	0	-8	-8	0
f	16	8	0	8	0	0	0	0	0	-8	-8	0	0	0	-8	-8

$$\Delta_{\mathrm{i}} = (0,0,0,1) \xrightarrow{S} \Delta_{\mathrm{o}} = (?,1,?,?)$$

$$\Delta_{\mathrm{i}} = (0, 1, 0, 0) \xrightarrow{S} \Delta_{\mathrm{o}} = (1, ?, ?, ?)$$

$$\Delta_{\mathrm{i}} = (1,0,0,0) \xrightarrow{S} \Delta_{\mathrm{o}} = (1,1,?,?)$$

$$\Delta_{\mathrm{i}} = (1,0,0,1) \stackrel{\mathcal{S}}{\rightarrow} \Delta_{\mathrm{o}} = (?,0,?,?)$$

$$\Delta_{\rm i} = (1,1,0,0) \xrightarrow{S} \Delta_{\rm o} = (0,?,?,?)$$

$$\lambda_{\rm i} = (1,?,?,1) \stackrel{S}{\leftarrow} \lambda_{\rm o} = (0,1,0,0)$$

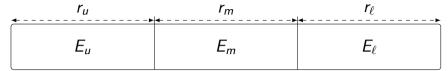
$$\lambda_{\mathrm{i}} = (1, 1, ?, ?) \stackrel{S}{\leftarrow} \lambda_{\mathrm{o}} = (1, 0, 0, 0)$$

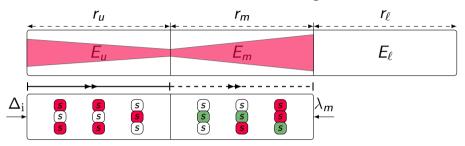
$$\lambda_{\rm i} = (0,?,?,?) \stackrel{S}{\leftarrow} \lambda_{\rm o} = (1,1,0,0)$$

Automatic Tools to Search for DL Distinguishers

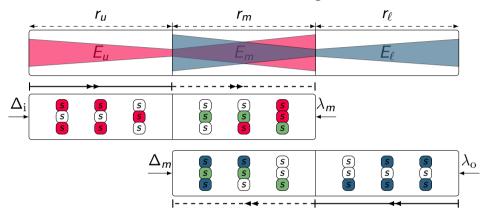


E

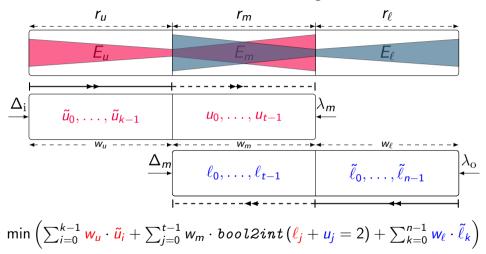




- differentially active S-box
 linearly active S-box
 common active S-box



- differentially active S-box
 linearly active S-box
 common active S-box

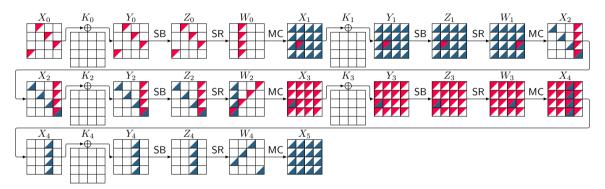


Usage of Our Tool

python3 attack.py -RU 6 -RM 10 -RL 6

<i>r_u</i>	r _m	r _ℓ
E _u	E _m	E_ℓ

Example: A 5-round DL Distinguisher for AES



$$r_0 = 1, r_m = 3, r_1 = 1, p = 2^{-24.00}, r = 2^{-7.66}, q^2 = 2^{-24.00}, prq^2 = 2^{-55.66}$$

Example: Distinguishers for up to 17 Rounds of TWINE

Comparing the data complexity of best boomerang and DL distinguishers

# Rounds	Boomerang [HNE22]	Differential-Linear	Gain
5	1	1	1
7	2 ^{3.20}	1	$2^{3.20}$
13	2 ^{34.32}	$2^{27.16}$	$2^{7.16}$
14	2 ^{42.25}	$2^{31.28}$	$2^{10.97}$
15	2 ^{51.03}	2 ^{38.98}	$2^{12.05}$
16	2 ^{58.04}	2 ^{47.28}	$2^{10.76}$
17	-	2 ^{59.24}	-

Example: Distinguishers for up to 17 Rounds of LBlock

Comparing the data complexity of best boomerang and DL distinguishers

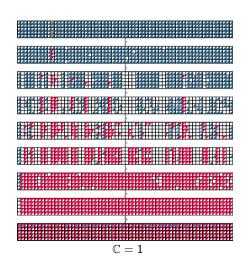
# Rounds	Boomerang [HNE22]	Differential-Linear	Gain
5	1	1	1
7	$2^{2.97}$	1	$2^{2.97}$
13	2 ^{30.28}	2 ^{23.78}	$2^{6.50}$
14	2 ^{38.86}	2 ^{30.34}	$2^{8.52}$
15	2 ^{46.90}	2 ^{38.26}	2 ^{8.64}
16	2 ^{57.16}	2 ^{46.26}	$2^{10.90}$
17	-	2 ^{58.30}	-

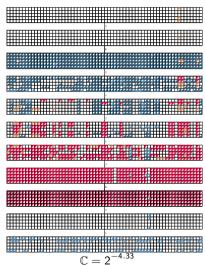
Example: Distinguishers for up to 8 Rounds of CLEFIA

Comparing the data complexity of best boomerang and DL distinguishers

# Rounds	Boomerang [HNE22]	Differential-Linear	Gain
3	1	1	1
4	$2^{6.32}$	1	$2^{6.32}$
5	$2^{12.26}$	$2^{5.36}$	$2^{6.90}$
6	2 ^{22.45}	$2^{14.14}$	$2^{8.31}$
7	2 ^{32.67}	$2^{23.50}$	$2^{9.17}$
8	2 ^{76.03}	2 ^{66.86}	$2^{9.17}$

Application to Ascon-p(active difference unknown difference active mask unknown mask)





Application to SERPENT

■ □: Experimentally verified

Cipher	#R	\mathbb{C}		Ref.
	3	$2^{-0.68}$	√	This work
	4	$2^{-12.75}$		[DIK08]
	4	$2^{-5.54}$	\checkmark	This work
CEDDENT	5	$2^{-16.75}$		[DIK08]
SERPENT	5	$2^{-11.10}$	\checkmark	This work
	8	$2^{-39.18}$		This work
	9	$2^{-56.50}$		[DIK08]
	9	$2^{-50.95}$		This work

Application to Simeck

■ □: Experimentally verified

Cipher	#R	\mathbb{C}		Ref.
	7	1	✓	This work
Simeck-32	14	$2^{-16.63}$		[ZWH24]
	14	$2^{-13.92}$	\checkmark	This work

Cipher	#R	\mathbb{C}		Ref.
	8	1	√	This work
	17	$2^{-22.37}$		[ZWH24]
	17	$2^{-13.89}$	\checkmark	This work
Simeck-48	18	$2^{-24.75}$		[ZWH24]
	18	$2^{-15.89}$		This work
	19	$2^{-17.89}$		This work
	20	$2^{-21.89}$		This work

Cipher	#R	\mathbb{C}		Ref.
	10	1	√	This work
	24	$2^{-38.13}$		[ZWH24]
C' - 1 C4	24	$2^{-25.14}$		This work
Simeck-64	25	$2^{-41.04}$		[ZWH24]
	25	$2^{-27.14}$		This work
	26	$2^{-30.35}$		This work

Contributions and Future Works



Contributions and Future Works

- Contributions
 - We generalized the DLCT framework from one S-box layer to multiple rounds
 - We proposed an automatic tool for finding optimum DL distinguishers
 - We applied our tool to almost any design paradigm
- Future works
 - A Extedning the application of our tool to other primitives, e.g., ARX
 - A Extending our tool to a unified model for finding complete attack (key recovery)
 - : https://ia.cr/2024/255

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Properties of Generalized DLCT Tables - I

- DLCT $(\Delta_{i}, \lambda_{o}) = \sum_{\Delta_{o}} \text{UDLCT}(\Delta_{i}, \Delta_{o}, \lambda_{o})$
- $\quad \quad \text{UDLCT}(\Delta_{\mathrm{i}}, \Delta_{\mathrm{o}}, \lambda_{\mathrm{o}}) = (-1)^{\Delta_{\mathrm{o}} \cdot \lambda_{\mathrm{o}}} \text{DDT}(\Delta_{\mathrm{i}}, \Delta_{\mathrm{o}})$
- $\qquad \texttt{LDLCT}(\Delta_{\mathrm{i}}, \lambda_{\mathrm{i}}, \lambda_{\mathrm{o}}) = (-1)^{\Delta_{\mathrm{i}} \cdot \lambda_{\mathrm{i}}} \texttt{DLCT}(\Delta_{\mathrm{i}}, \lambda_{\mathrm{o}})$
- $\qquad \text{EDLCT}(\Delta_{\mathrm{i}}, \Delta_{\mathrm{o}}, \lambda_{\mathrm{i}}, \lambda_{\mathrm{o}}) = (-1)^{\lambda_{\mathrm{i}} \cdot \Delta_{\mathrm{i}} \oplus \lambda_{\mathrm{o}} \cdot \Delta_{\mathrm{o}}} \text{DDT}(\Delta_{\mathrm{i}}, \Delta_{\mathrm{o}})$
- LDLCT $(\Delta_{i}, \lambda_{i}, \lambda_{o}) = \sum_{\Delta_{o}} \text{EDLCT}(\Delta_{i}, \Delta_{o}, \lambda_{i}, \lambda_{o})$
- $\sum_{\Lambda_i} \text{LDLCT}(\Delta_i, \lambda_i, \lambda_o) = \text{LAT}^2(\lambda_i, \lambda_o)$

Properties of Generalized DLCT Tables - II

 $\qquad \mathtt{DDLCT}(\Delta_{\mathrm{i}}, \lambda_{\mathrm{o}}) = \textstyle \sum_{\Delta_{m}} \sum_{\lambda_{m}} \mathtt{UDLCT}\left(\Delta_{\mathrm{i}}, \Delta_{m}, \lambda_{m}\right) \cdot \mathtt{LDLCT}\left(\Delta_{m}, \lambda_{m}, \lambda_{\mathrm{o}}\right)$

$$egin{aligned} exttt{DDLCT}(\Delta_{\mathrm{i}}, \lambda_{\mathrm{o}}) &= \sum_{\Delta_m} exttt{DDT}(\Delta_{\mathrm{i}}, \Delta_m) \cdot exttt{DLCT}(\Delta_m, \lambda_{\mathrm{o}}) \ &= 2^{-n} \sum_{\lambda_m} exttt{DLCT}(\Delta_{\mathrm{i}}, \lambda_m) \cdot exttt{LAT}^2(\lambda_m, \lambda_{\mathrm{o}}). \end{aligned}$$