Network layer: "data plane" roadmap

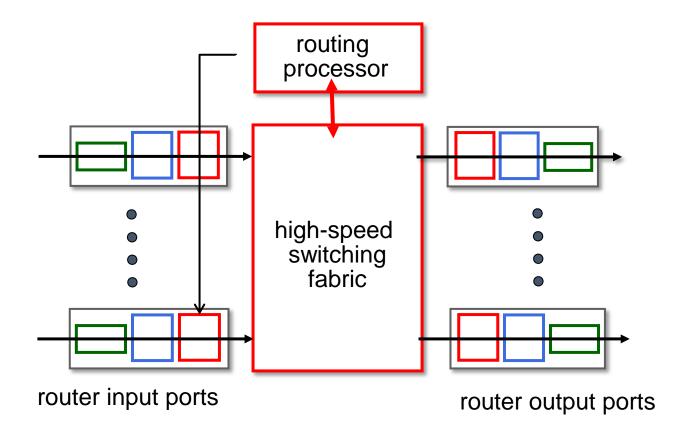
- Network layer: overview
 - data plane
 - control plane
- What's inside a router
 - input ports, switching, output ports
 - buffer management, scheduling
- IP: the Internet Protocol
 - datagram format
 - addressing
 - network address translation
 - IPv6



- Generalized Forwarding, SDN
 - Match+action
 - OpenFlow: match+action in action
- Middleboxes

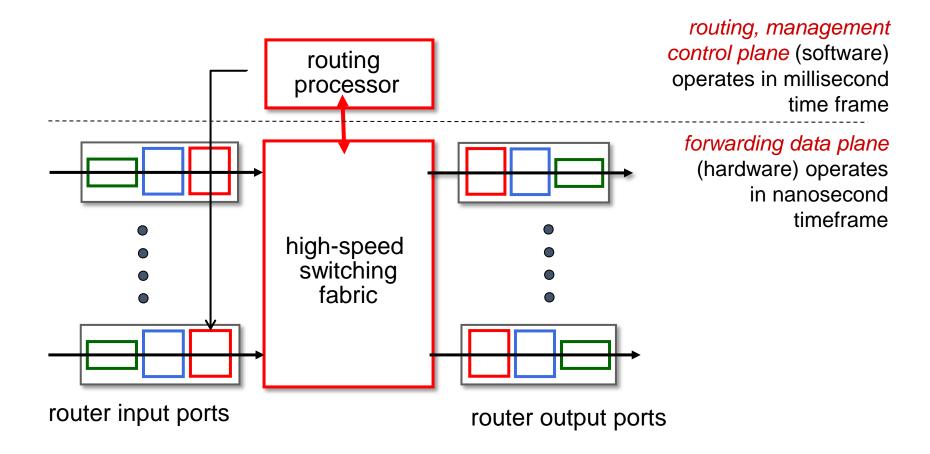
Router architecture overview

high-level view of generic router architecture:

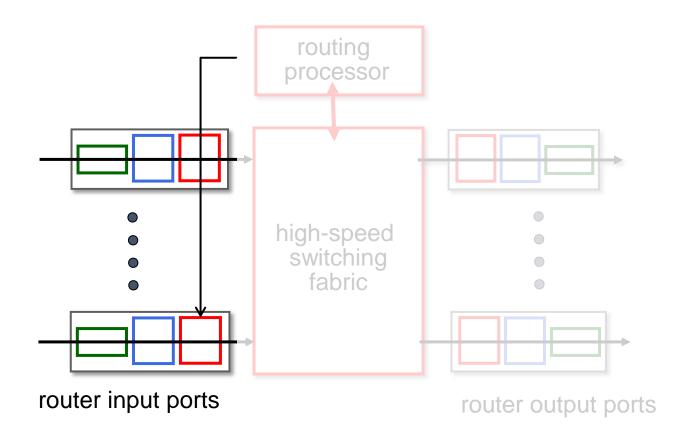


Router architecture overview

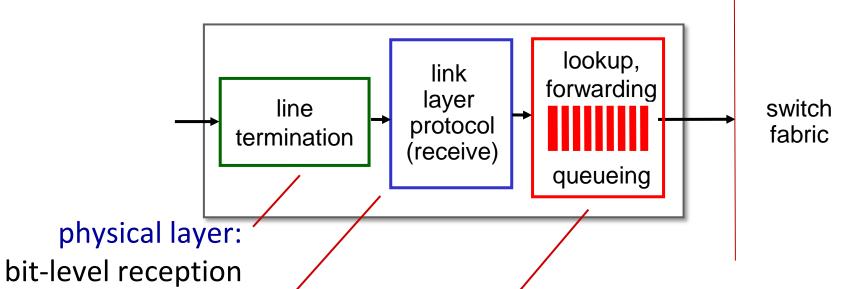
high-level view of generic router architecture:



1) Input ports



Input port functions



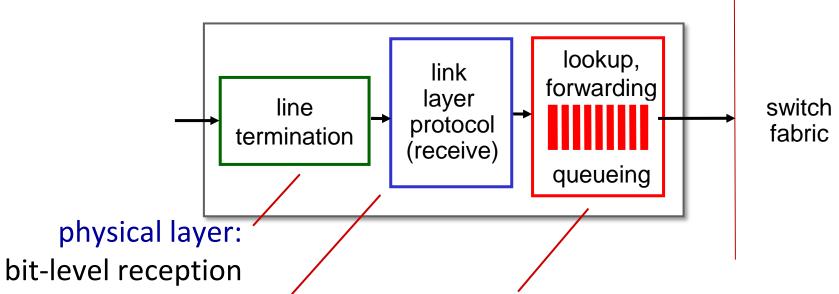
link layer:

e.g., Ethernet (chapter 6)

decentralized switching:

- using header field values, lookup output port using forwarding table in input port memory ("match plus action")
- goal: complete input port processing at 'line speed'
- input port queuing: if datagrams arrive faster than forwarding rate into switch fabric

Input port functions



link layer:

e.g., Ethernet (chapter 6)

decentralized switching:

- using header field values, lookup output port using forwarding table in input port memory ("match plus action")
- destination-based forwarding: forward based only on destination IP address (traditional)
- generalized forwarding: forward based on any set of header field values

Destination-based forwarding

forwarding table				
Destination Address Range	Link Interface			
11001000 00010111 000 <mark>10000 00000000000</mark>	0			
11001000 00010111 000 <mark>11000 00000000000</mark>	1			
11001000 00010111 000 <mark>11000 11111111</mark> 11001000 00010111 000 <mark>11001 00000000000</mark>	2			
11001000 00010111 000 <mark>11111 11111111</mark>	_			
otherwise	3			

Destination-based forwarding

forwarding table	
Destination Address Range	Link Interface
11001000 00010111 000 <mark>10000 00000000000</mark>	Q: but what happens if ranges don't divide up so
11001000 00010111 000 <mark>10000 00000</mark> 100 through 11001000 00010111 000 <mark>10000 00000111</mark>	nicely?
11001000 00010111 000 <mark>11000 00000000000</mark>	1
11001000 00010111 000 <mark>11001 00000000000</mark>	2
otherwise	3

longest prefix match

when looking for forwarding table entry for given destination address, use *longest* address prefix that matches destination address.

Destination A	Link interface			
11001000	00010111	00010***	*****	0
11001000	00010111	00011000	*****	1
11001000	00010111	00011***	*****	2
otherwise				3

examples:

which interface?	10100001	00010110	00010111	11001000
which interface?	10101010	00011000	00010111	11001000

longest prefix match

when looking for forwarding table entry for given destination address, use *longest* address prefix that matches destination address.

Destination Address Range				Link interface
11001000	00010111	00010***	*****	0
11001000	000.0111	00011000	*****	1
11001000	match! 1	00011***	*****	2
otherwise				3

examples

11001000 00010111 00010 110 10100001 which interface?
11001000 00010111 00011000 10101010 which interface?

longest prefix match

when looking for forwarding table entry for given destination address, use *longest* address prefix that matches destination address.

Destination Address Range				Link interface
11001000	00010111	00010***	*****	0
11001000	00010111	00011000	*****	1
11001000	00010111	00011 * * *	*****	2
otherwise				3

examples:

longest prefix match

11001000

when looking for forwarding table entry for given destination address, use *longest* address prefix that matches destination address.

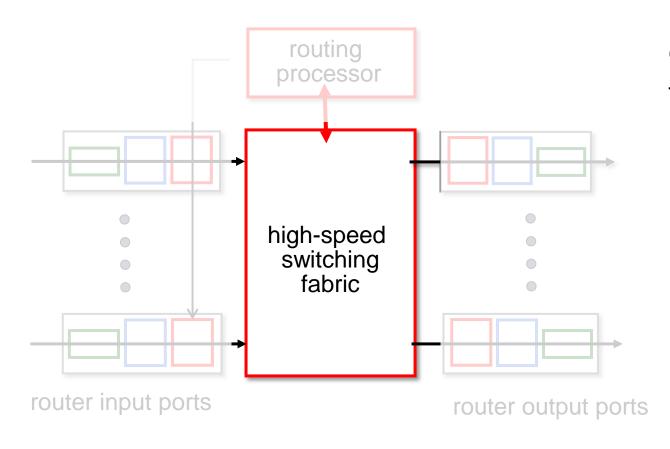
Doctination Address Pango				Link interface
Destination Address Range				LITIK II ILOTTAGO
11001000	00010111	00010***	*****	0
11001000	00010111	00011000	*****	1
11001000	0000111	00011***	*****	2
otherwise	match!			3
				1 . 1
11001000	000 0111	00010110	10100001	which interface?

00011000

examples:

- we'll see why longest prefix matching is used shortly, when we study addressing
- longest prefix matching: often performed using ternary content addressable memories (TCAMs)
 - content addressable: present address to TCAM: retrieve address in one clock cycle, regardless of table size
 - Cisco Catalyst: ~1M routing table entries in TCAM

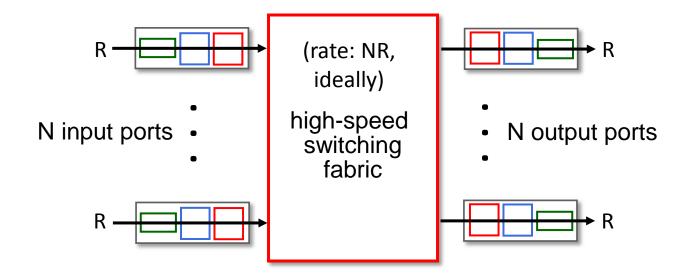
2)Switching Fabric



Connects input ports to the output ports (& the routing processor)

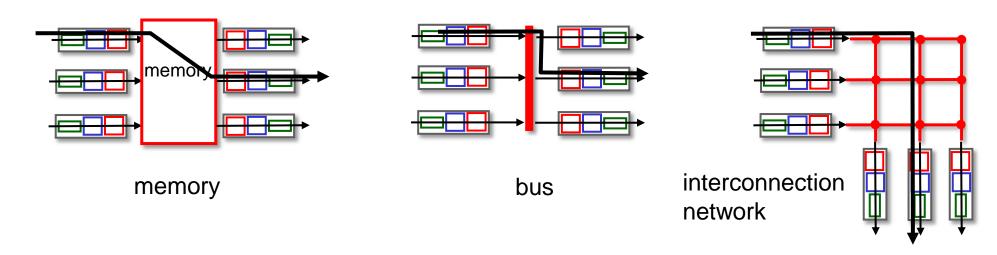
Switching fabrics

- transfer packet from input link to appropriate output link
- switching rate: rate at which packets can be transfer from inputs to outputs
 - often measured as multiple of input/output line rate
 - N inputs: switching rate N times line rate desirable



Switching fabrics

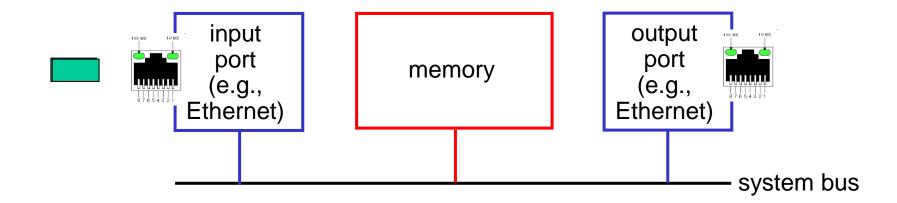
- transfer packet from input link to appropriate output link
- switching rate: rate at which packets can be transfer from inputs to outputs
 - often measured as multiple of input/output line rate
 - N inputs: switching rate N times line rate desirable
- three major types of switching fabrics:



Switching via memory

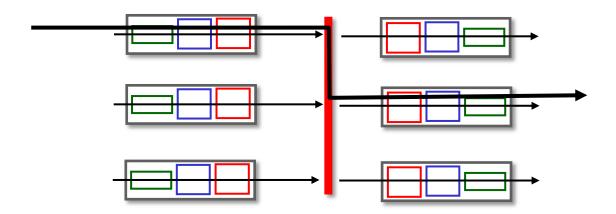
first generation routers:

- traditional computers with switching under direct control of CPU
- packet copied to system's memory
- speed limited by memory bandwidth (2 bus crossings per datagram)



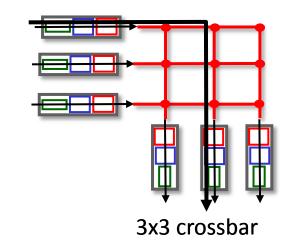
Switching via a bus

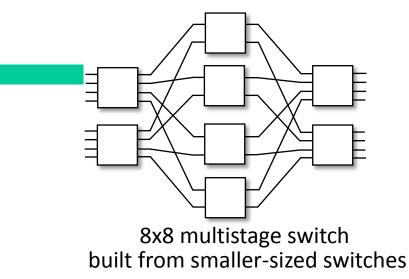
- datagram from input port memory to output port memory via a shared bus
- bus contention: switching speed limited by bus bandwidth
- 32 Gbps bus, Cisco 5600: sufficient speed for access routers



Switching via interconnection network

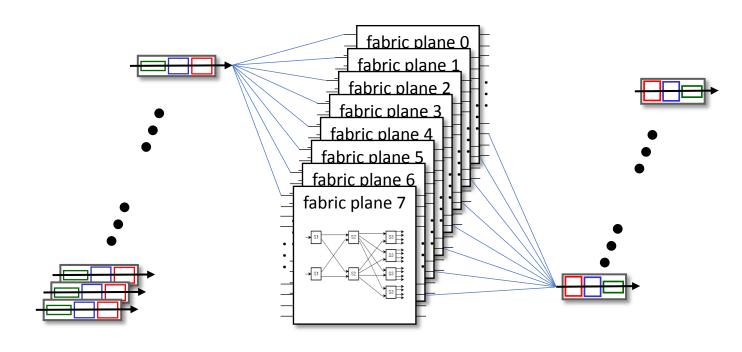
- Crossbar, Clos networks, other interconnection nets initially developed to connect processors in multiprocessor
- multistage switch: nxn switch from multiple stages of smaller switches
- exploiting parallelism:
 - fragment datagram into fixed length cells on entry
 - switch cells through the fabric, reassemble datagram at exit





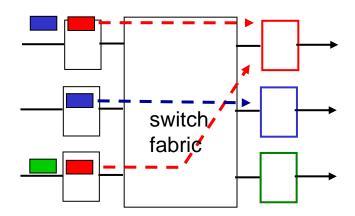
Switching via interconnection network

- scaling, using multiple switching "planes" in parallel:
 - speedup, scaleup via parallelism
- Cisco CRS router:
 - basic unit: 8 switching planes
 - each plane: 3-stage interconnection network
 - up to 100's Tbps switching capacity

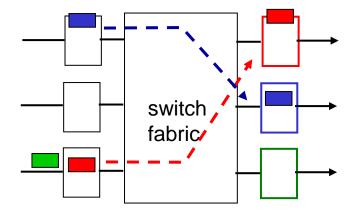


Input port queuing

- If switch fabric slower than input ports combined -> queueing may occur at input queues
 - queueing delay and loss due to input buffer overflow!
- Head-of-the-Line (HOL) blocking: queued datagram at front of queue prevents others in queue from moving forward

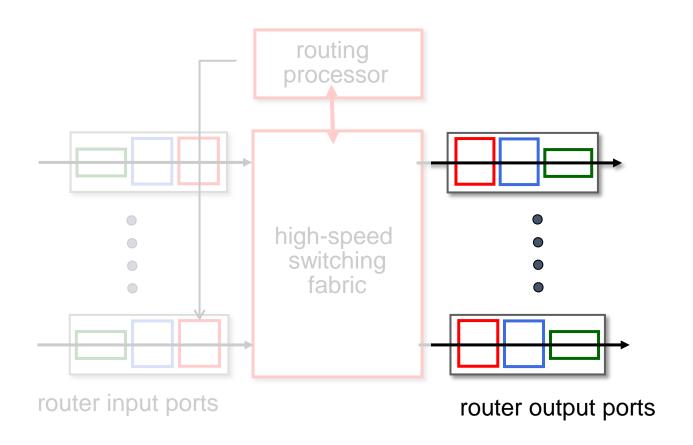


output port contention: only one red datagram can be transferred. lower red packet is *blocked*

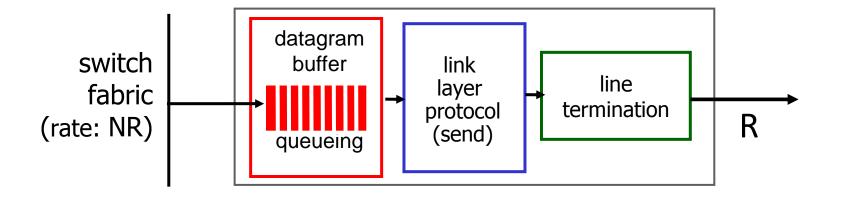


one packet time later: green packet experiences HOL blocking

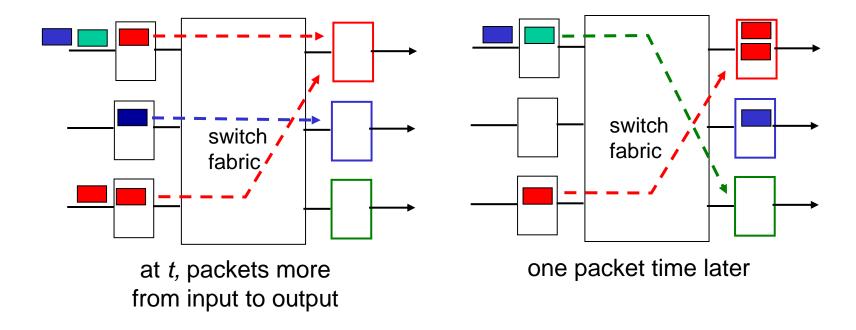
3) Output ports



Output port

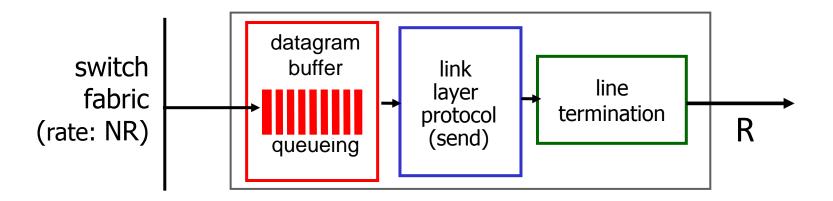


Output port queuing



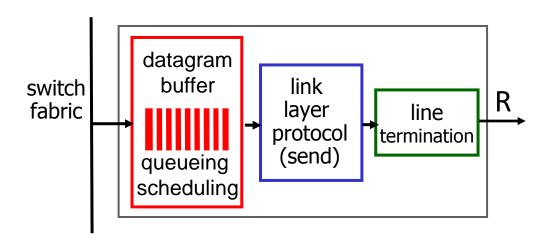
- buffering when arrival rate via switch exceeds output line speed
- queueing (delay) and loss due to output port buffer overflow!

Buffer Management

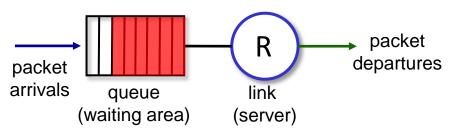


- Drop policy: which datagrams to drop if no free buffers?
- Scheduling discipline which datagram to choose among queued datagrams for transmission?

Drop Policies



Abstraction: queue



- drop: which packet to add, drop when buffers are full
 - tail drop: drop arriving packet
 - priority: drop/remove on priority basis

 marking: which packets to mark to signal congestion (ECN, RED)

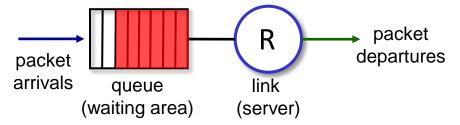
DQ

Packet Scheduling

packet scheduling: deciding which packet to send next on link

- first come, first served
- priority
- round robin
- weighted fair queueing

Abstraction: queue



Packet Scheduling: FCFS

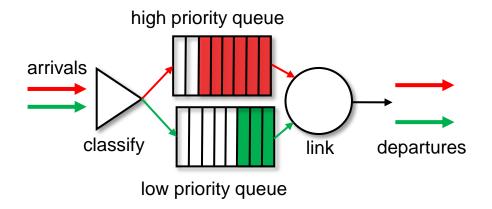
FCFS: packets transmitted in order of arrival to output port

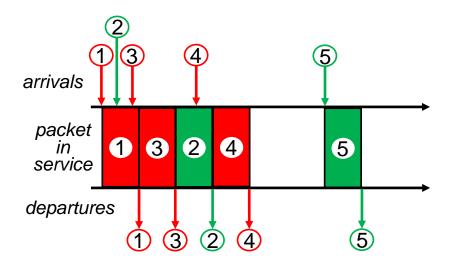
- also known as: First-in-first-out (FIFO)
- real world examples?

Scheduling policies: priority

Priority scheduling:

- arriving traffic classified, queued by class
 - any header fields can be used for classification
- send packet from highest priority queue that has buffered packets
 - FCFS within priority class

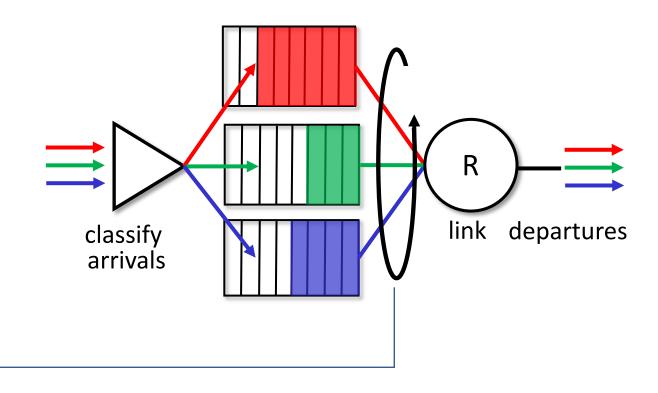




Scheduling policies: round robin

Round Robin (RR) scheduling:

- arriving traffic classified, queued by class
 - any header fields can be used for classification
- server cyclically, repeatedly scans class queues, sending one complete packet from each class (if available) in turn



Scheduling policies: weighted fair queueing

Weighted Fair Queuing (WFQ):

- generalized Round Robin
- each class, i, has weight, w_i, and gets weighted amount of service in each cycle:

$$\frac{w_i}{\sum_j w_j}$$

minimum bandwidth guarantee (per-traffic-class)

