

$D \rightarrow TL$	$L.a = T.type$	inherited syn	(Ug)
$T \rightarrow int$	$T.type = integer$	integer name (syn)	
$T \rightarrow float$	$T.type = float$		
$L \rightarrow L1, id$	$L1.a = L.a$	inherited	
	$addType(id.entry, L.a)$		
$L \rightarrow id$	$addType(id.entry, L.a)$		

$$T \rightarrow F \{ T.inh = F.val \} T' \{ T.val = T'.syn \}$$

$$T' \rightarrow *F \{ T'.inh = T.inh * F.val \} T' \{ T'.syn = T'.syn \}$$

$$T' \rightarrow \epsilon \{ T'.syn = T.inh \}$$

$$F \rightarrow digit \{ F.val = digit.lexval \}$$

↳ top down & given

non-terminal → recursive descent

↳ synthesized attr → inherited attr

val T ( ) { ①

val val\_F, val\_T'\_inh, val\_T'\_syn, val\_T

val\_F = F();

val\_T'\_inh = val\_F

val\_T'\_syn = T'(val\_T'\_inh);

val\_T = val\_T'\_syn; return val\_T;

②

val T'(val\_T'\_inh) {  
 val val\_F, val\_T'\_inh, val\_T'\_syn;  
 if (current == '\*') {  
 match('\*');  
 current++;  
 val\_F = F();  
 val\_T'\_inh = val\_T'\_inh \* val\_F;  
 val\_T'\_syn = T'(val\_T'\_inh);  
 val\_T'\_syn = val\_T'\_syn;  
 return val\_T'\_syn; } else  
 { return val\_T'\_inh; }

val F ( ) { ③  
 val val\_F;  
 val\_F = find(digit, lexval);  
 return val\_F; }





$S \rightarrow L.L \mid L$

$L \rightarrow LB \mid B$

$B \rightarrow 0 \mid 1$

attr isLeft  
will inherited  
for all non-terminals  
right to dot  
for the right len  
will

production	قواعد
$S \rightarrow L.L$	$L_1.isLeft = true$ $L_2.isLeft = false$ $S.val = L_1.val + L_2.val$
$S \rightarrow L$	$L.isLeft = true$ $S.val = L.val$
$L \rightarrow LB$	$L_1.isLeft = L.isLeft$ $L.len = L_1.len + 1$ $L.val = L.isLeft ?$ $L_1.val * 2 + B.val :$ $L_1.val + B.val * 2^{L.len}$
$L \rightarrow B$	$L.len = 1$ $L.val = L.isLeft ?$ $B.val : B.val / 2$
$B \rightarrow 0$	$B.val = 0$
$B \rightarrow 1$	$B.val = 1$

L-attributed نویس  
SDD

U101

(a)  $S \rightarrow \text{if } (C) S_1 \text{ else } S_2$

$L_1 = \text{new}()$

$C.\text{false} = L_1$

$S_1.\text{next} = S.\text{next}$

$S.\text{code} = C.\text{code} \parallel S_1.\text{code} \parallel \text{label} \parallel L_1 \parallel S_2.\text{code}$

☆ کد C،  $S_1$ ،  $S_2$ ، ایز بر

(b)  $S \rightarrow \text{do } S_1 \text{ while } (C)$

$L_1 = \text{new}()$

$C.\text{true} = L_1$

$S.\text{code} = \text{label} \parallel L_1 \parallel S_1.\text{code} \parallel C.\text{code}$

(c)  $S \rightarrow 'x' L 'y'; L \rightarrow LS \mid \epsilon$

$S \rightarrow \text{if}(B) S_1$

$B \rightarrow B_1 \parallel B_2$

$B \rightarrow B_1 \&\& B_2$

$B \rightarrow \text{true}$

$B \rightarrow \text{false}$

①

$S \rightarrow \text{if}(B) m S_1 \{ \text{backpatch}(B.\text{true list}, m.\text{instr});$   
 $S.\text{nextlist} = \text{merge}(B.\text{false list}, S_1.\text{nextlist});$

②

$B \rightarrow B_1 \parallel m B_2 \{ \text{backpatch}(B_1.\text{false list}, m.\text{instr});$   
 $B.\text{true list} = \text{merge}(B_1.\text{true list}, B_2.\text{true list});$   
 $B.\text{false list} = B_2.\text{false list};$

③  $B \rightarrow B_1 \&\& m B_2$

$\{ \text{backpatch}(B_1.\text{true list}, m.\text{instr});$   
 $B.\text{true list} = B_2.\text{true list};$   
 $B.\text{false list} = \text{merge}(B_1.\text{false list}, B_2.\text{false list});$

④  $B \rightarrow \text{true} \quad \{ B.\text{true list} = \text{makelist}(\text{nextinstr});$   
 $\text{gen}('goto -'); \}$

⑤  $B \rightarrow \text{false} \quad \{ B.\text{false list} = \text{makelist}(\text{nextinstr});$   
 $\text{gen}('goto -'); \}$

⑥  $m \rightarrow \epsilon \quad \{ m.\text{instr} = \text{nextinstr}; \}$

UP true  $\hookrightarrow$   $\cup$  true  $\in B_1$   $\&\&$   $\cup$  true  $\in B_2$  ②  $\cup$  true  
UP false  $\hookrightarrow$   $\cup$  false  $\in B_1$   $\&\&$   $\cup$  false  $\in B_2$   $\&\&$

