Operating Systems

Isfahan University of Technology Electrical and Computer Engineering Department 1400-1 semester

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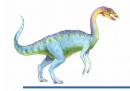
Session 16: Classical Problems of Synchronization



Classical Problems of Synchronization

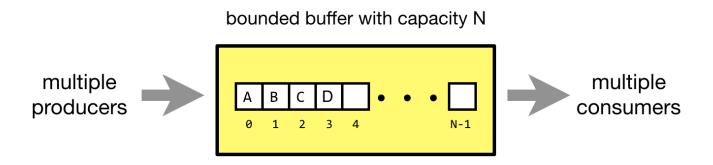
- Classical problems used to test newly-proposed synchronization schemes
 - Bounded-Buffer Problem
 - Readers and Writers Problem
 - Dining-Philosophers Problem





Bounded-Buffer Problem

- buffers, each can hold one item
- Semaphore mutex initialized to the value 1
- Semaphore **full** initialized to the value 0
- Semaphore empty initialized to the value n





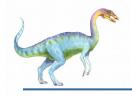


The structure of the producer process

```
while (true) {
    ...
    /* produce an item in next_produced */
    ...
    /* add next produced to the buffer */
    ...
```

}

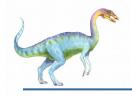




The structure of the producer process

```
while (true) {
     /* produce an item in next produced */
   wait(empty);
   wait(mutex);
     /* add next produced to the buffer */
   signal(mutex);
   signal(full);
```

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The structure of the consumer process

```
while (true) {
   /* remove an item from buffer to next consumed */
     /* consume the item in next consumed */
```

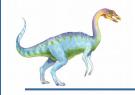




The structure of the consumer process

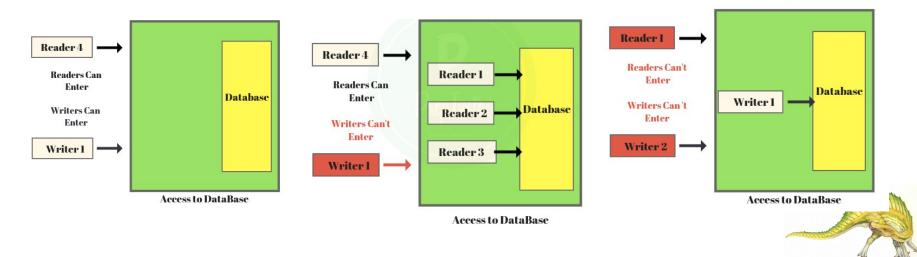
```
while (true) {
   wait(full);
   wait(mutex);
   /* remove an item from buffer to next consumed */
   signal(mutex);
   signal(empty);
     /* consume the item in next consumed */
```

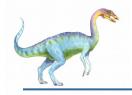




Readers-Writers Problem

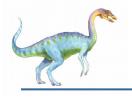
- A data set is shared among a number of concurrent processes
 - Readers only read the data set; they do not perform any updates
 - Writers can both read and write
- Problem allow multiple readers to read at the same time
 - Only one single writer can access the shared data at the same time
- Several variations of how readers and writers are considered all involve some form of priorities





- Shared Data
 - Data set
 - Semaphore rw mutex initialized to 1
 - Semaphore mutex initialized to 1
 - Integer read_count initialized to 0





The structure of a writer process

```
while (true) {
```

}





The structure of a writer process

```
while (true) {
          wait(rw_mutex);
          ...
     /* writing is performed */
          ...
          signal(rw_mutex);
}
```





The structure of a reader process

```
while (true) {
```

}

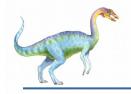




The structure of a reader process

```
while (true) {
       wait(mutex);
       read count++;
       if (read count == 1) /* first reader */
        wait(rw mutex);
            signal(mutex);
       /* reading is performed */
       wait(mutex);
       read count--;
       if (read count == 0) /* last reader */
           signal(rw mutex);
       signal(mutex);
```





Readers-Writers Problem Variations

- The solution in previous slide can result in a situation where a writer process never writes. It is referred to as the "First reader-writer" problem.
 - Once a reader is ready to read, no "newly arrived writer" is allowed to read.
- The "Second reader-writer" problem is a variation the first reader-writer problem that state:
 - Once a writer is ready to write, no "newly arrived reader" is allowed to read.
- Both the first and second may result in starvation. leading to even more variations
- Problem is solved on some systems by kernel providing reader-writer locks

