

Introduction to Software Testing (2nd edition) Chapter 7.5

Graph Coverage for Specifications

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Design Specifications

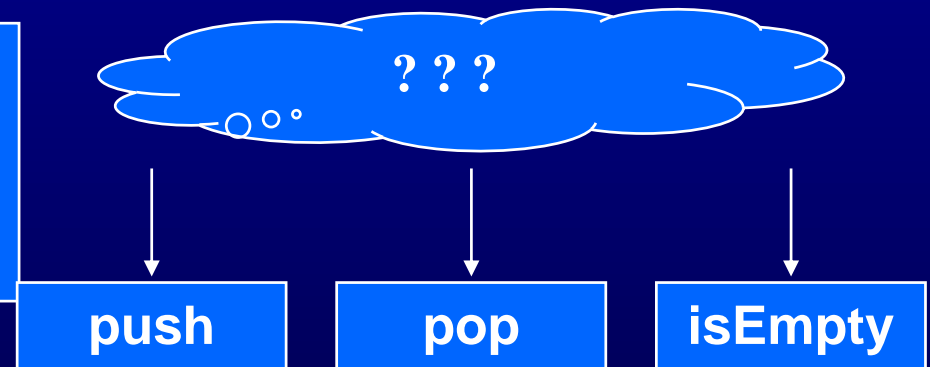
- A **design specification** describes **aspects** of **what behavior software should exhibit**
- A **design specification** **may** or **may not** **reflect** the **implementation**
 - More accurately – the **implementation** may **not** exactly reflect the **spec**
 - Design specifications are often called **models of the software**
- **Two types of descriptions** are used in this chapter
 1. **Sequencing constraints on class methods**
 2. **State behavior descriptions of software**

Sequencing Constraints

- Sequencing constraints are **rules** that **impose constraints** on the **order** in which **methods may be called**
- They can be encoded as **preconditions** or **other specifications**
- Section 7.4 said that classes often have methods that do not call each other

Class stack

```
public void push (Object o)
public Object pop ( )
public boolean isEmpty ( )
```



- Tests can be created for these classes **as sequences of method calls**
- Sequencing constraints give an **easy** and **effective** way to **choose which sequences to use**

Sequencing Constraints Overview

- Sequencing constraints might be
 - Expressed explicitly
 - Expressed implicitly
 - Not expressed at all
- Testers should derive them if they do not exist
 - Look at existing design documents
 - Look at requirements documents
 - Ask the developers
 - Last choice : Look at the implementation
- If they don't exist, expect to find more faults !
- Share with designers before designing tests
- Sequencing constraints do not capture all behavior

Queue Example

```
public int deQueue()  
{  
    // Pre: At least one element must be on the queue.  
    ... ..  
public enQueue (int e)  
{  
    // Post: e is on the end of the queue.
```

- **Sequencing constraints** are **implicitly** embedded in the **pre** and **postconditions**
 - **enQueue ()** must be called before **deQueue ()**
- **Does not include** the requirement that we must have at least as many **enQueue ()** calls as **deQueue ()** calls
 - Can be handled by **state behavior techniques**

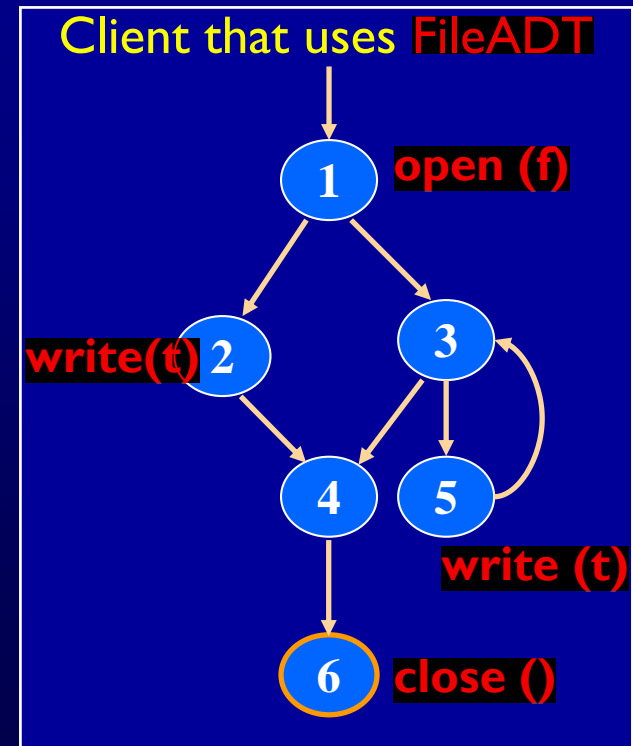
File ADT Example

class FileADT has **three methods**:

- **open (String fName)** // Opens file with name fName
- **close ()** // Closes the file and makes it unavailable
- **write (String textLine)** // Writes a line of text to the file

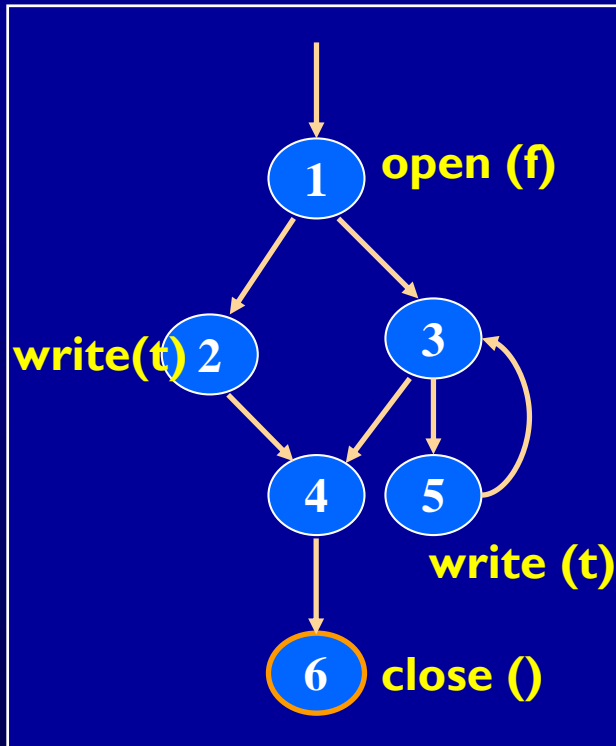
Valid sequencing constraints on FileADT:

1. An **open (f)** must be executed **before every write (t)**
2. An **open (f)** must be executed **before every close ()**
3. A **write (f)** may not be executed **after a close ()** unless there is an **open (f)** in between
4. A **write (t)** **should** be executed **before every close ()**



Static Checking

Is there a path that violates any of the sequencing constraints ?

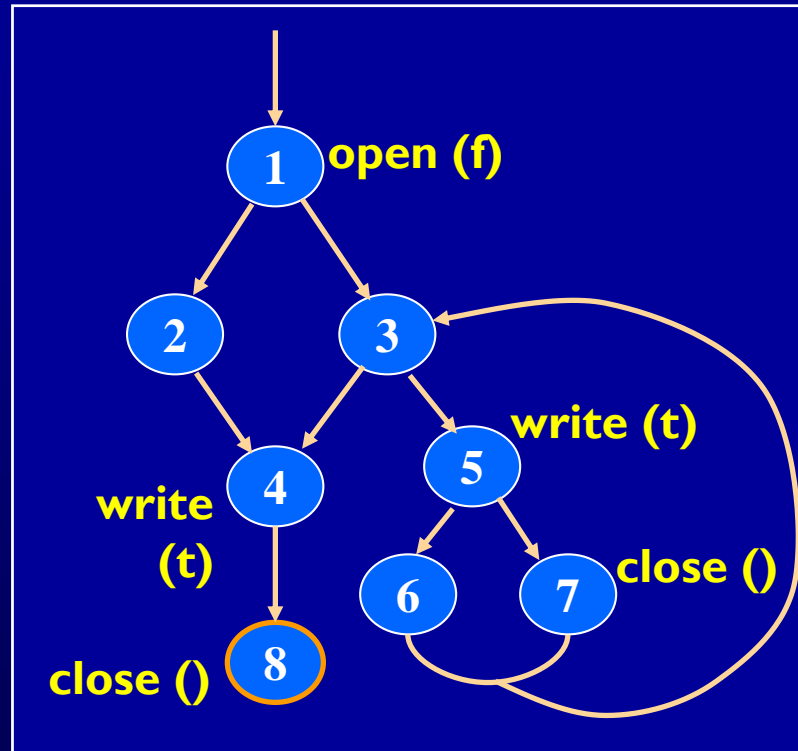


- Is there a path to a write() that **does not go through an open()** ?
- Is there **a path to a close()** that does not go through an open() ?
- Is there a path from a close() to a write()?
- Is there a path from an open() to a close() that **does not go through a write()** ? (“write-clear” path)

[1, 3, 4, 6] – ADT use anomaly!

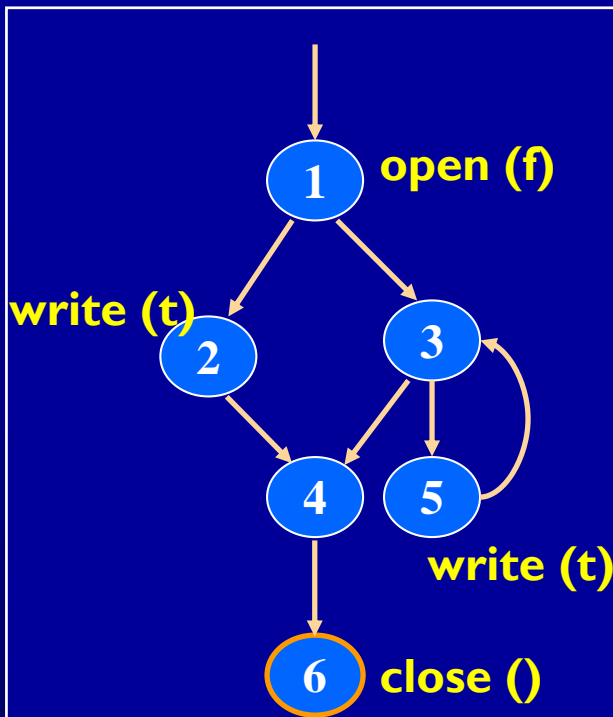
Static Checking

Consider the following graph :



[7, 3, 4] – close () before write () !

Generating Test Requirements



[1, 3, 4, 6] – ADT use anomaly!

- But it is possible that the **logic** of the **program** does **not allow** the **pair of edges [1, 3, 4]**
- That is – the **loop body** must be taken at least once
- Determining this is **undecidable** – so static methods are not enough

- Use the sequencing constraints to generate **test requirements**
- The goal is to **violate every sequencing constraint**

Test Requirements for FileADT

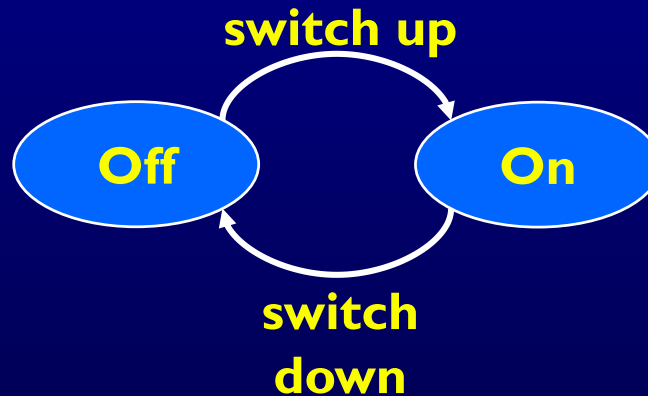
Apply to all programs that use FileADT

1. Cover every path from the start node to every node that contains a write() such that the path does not go through a node containing an open()
2. Cover every path from the start node to every node that contains a close() such that the path does not go through a node containing an open()
3. Cover every path from every node that contains a close() to every node that contains a write()
4. Cover every path from every node that contains an open() to every node that contains a close() such that the path does not go through a node containing a write()

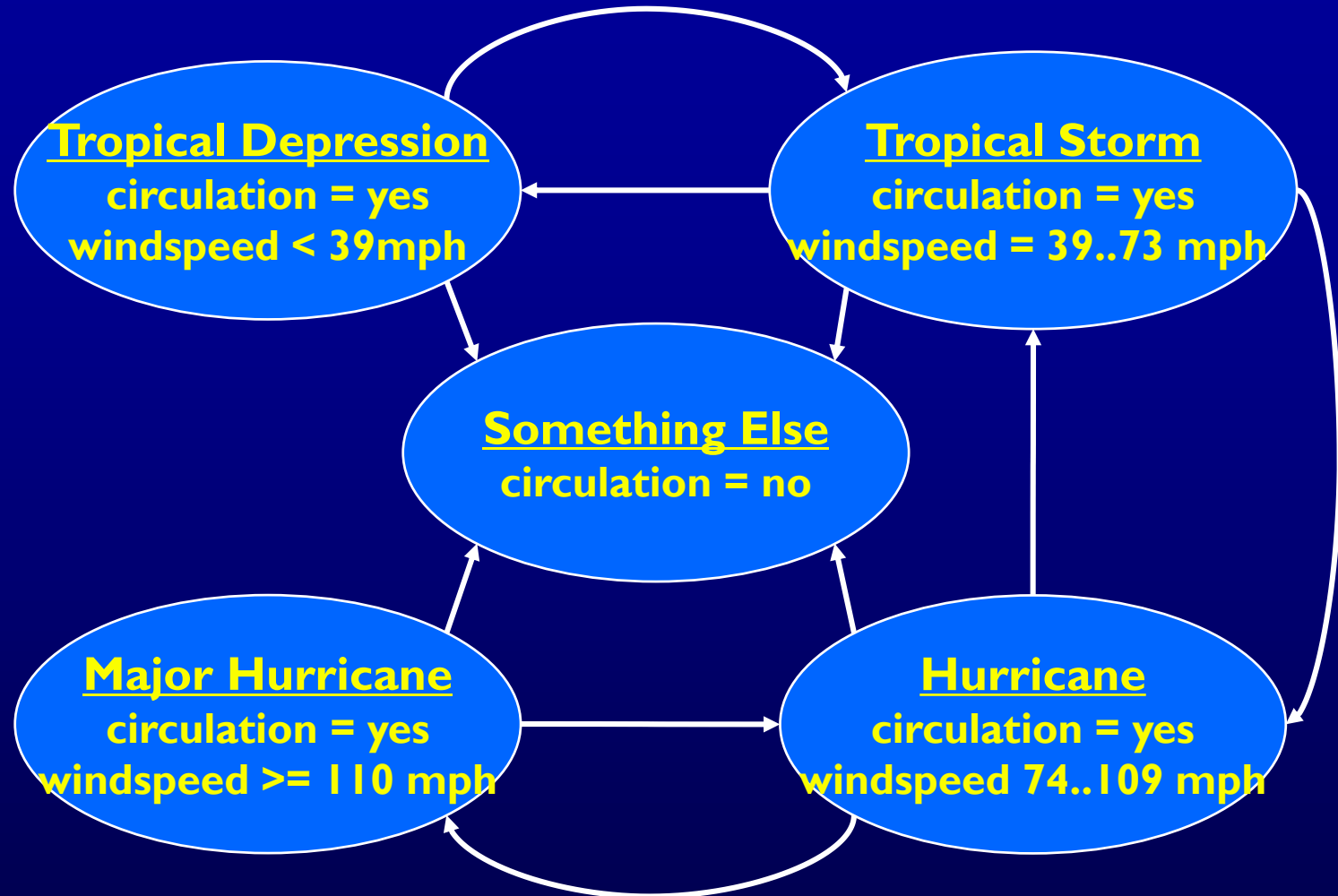
- If program is correct, all test requirements will be infeasible
- Any tests created will almost definitely find faults

Testing State Behavior

- A finite state machine (FSM) is a graph that describes how software variables are modified during execution
- **Nodes** : **States**, representing sets of values for key variables
- **Edges** : **Transitions**, possible changes in the state



Finite State Machine—Two Variables



Other variables may exist but **not be part of state**

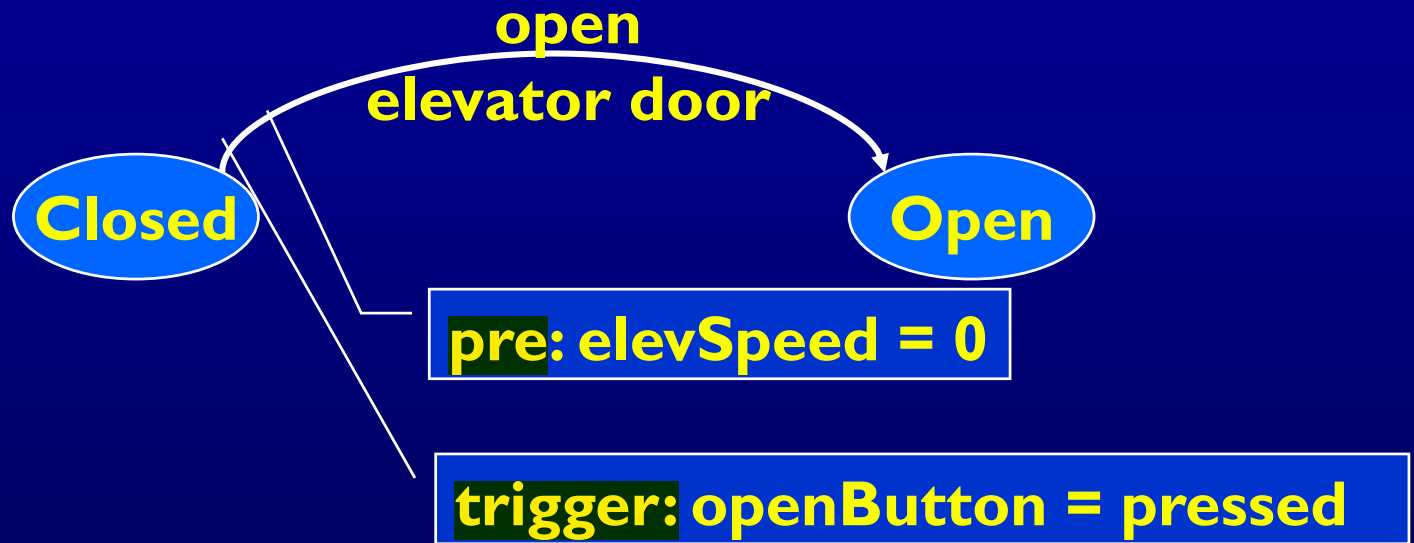
Finite State Machines are Common

- **FSMs** can **accurately model** many kinds of software
 - **Embedded** and **control software** (think electronic gadgets)
 - **Abstract data types**
 - **Compilers** and **operating systems**
 - **Web** applications
- **Creating** FSMs can **help** find **software problems**
- Numerous **languages** for **expressing** FSMs
 - **UML statecharts**
 - **Automata**
 - **State tables (SCR)**
 - **Petri nets**
- **Limitation** : FSMs are **not always practical** for programs that have **lots of states** (for example, **GUIs**)

Annotations on FSMs

- FSMs can be annotated with **different types of actions**
 - **Actions on transitions**
 - **Entry actions to nodes**
 - **Exit actions on nodes**
- Actions can express **changes to variables** or **conditions** on variables
- These slides use the basics:
 - **Preconditions (guards)** : **conditions that must be true** for transitions to be taken
 - **Triggering events** : **changes to variables** that cause transitions to be taken
- This is close to the **UML Statecharts**, but not exactly the same

Example Annotations



Covering FSMs

- **Node coverage** : execute **every state** (*state coverage*)
- **Edge coverage** : execute **every transition** (*transition coverage*)
- **Edge-pair coverage** : execute **every pair of transitions** (*transition-pair*)
- **Data flow**:
 - **Nodes** often do **not include** defs or uses of variables
 - **Defs** of variables in **triggers** are used immediately (the next state)
 - **Defs** and **uses** are usually computed for **guards**, or states are extended
 - **FSMs** typically only **model a subset of the variables**
- **Generating FSMs** is often **harder** than covering them ...

Deriving FSMs

- With some projects, an **FSM** (such as a statechart) was created **during design**
 - **Tester** should check to see if the **FSM is still current** with respect to the **implementation**
- If not, it is **very helpful** for the **tester to derive the FSM**
- Strategies for **deriving FSMs** from a program:
 1. **Combining control flow graphs** (*wrong*)
 2. Using the **software structure** (*wrong*)
 3. **Modeling state variables**
- Example based on a digital watch ...
 - Class Watch **uses** class Time

Class Watch

class Watch

```
// Constant values for the button (inputs)
private static final int NEXT = 0;
private static final int UP = 1;
private static final int DOWN = 2;
// Constant values for the state
private static final int TIME = 5;
private static final int STOPWATCH = 6;
private static final int ALARM = 7;
// Primary state variable
private int mode = TIME;
// Three separate times, one for each state
private Time watch, stopwatch, alarm;

public Watch () // Constructor
public void doTransition (int button) // Handles inputs
public String toString () // Converts values
```

class Time (inner class)

```
private int hour = 0;
private int minute = 0;

public void changeTime (int button)
public String toString ()
```

// Takes the appropriate transition when a button is pushed.

```
public void doTransition (int button)
{
    switch ( mode )
    {
        case TIME:
            if (button == NEXT)
                mode = STOPWATCH;
            else
                watch.changeTime (button);
            break;
        case STOPWATCH:
            if (button == NEXT)
                mode = ALARM;
            else
                stopwatch.changeTime (button);
            break;
        case ALARM:
            if (button == NEXT)
                mode = TIME;
            else
                alarm.changeTime (button);
            break;
        default:
            break;
    }
} // end doTransition()
```

// Increases or decreases the time.

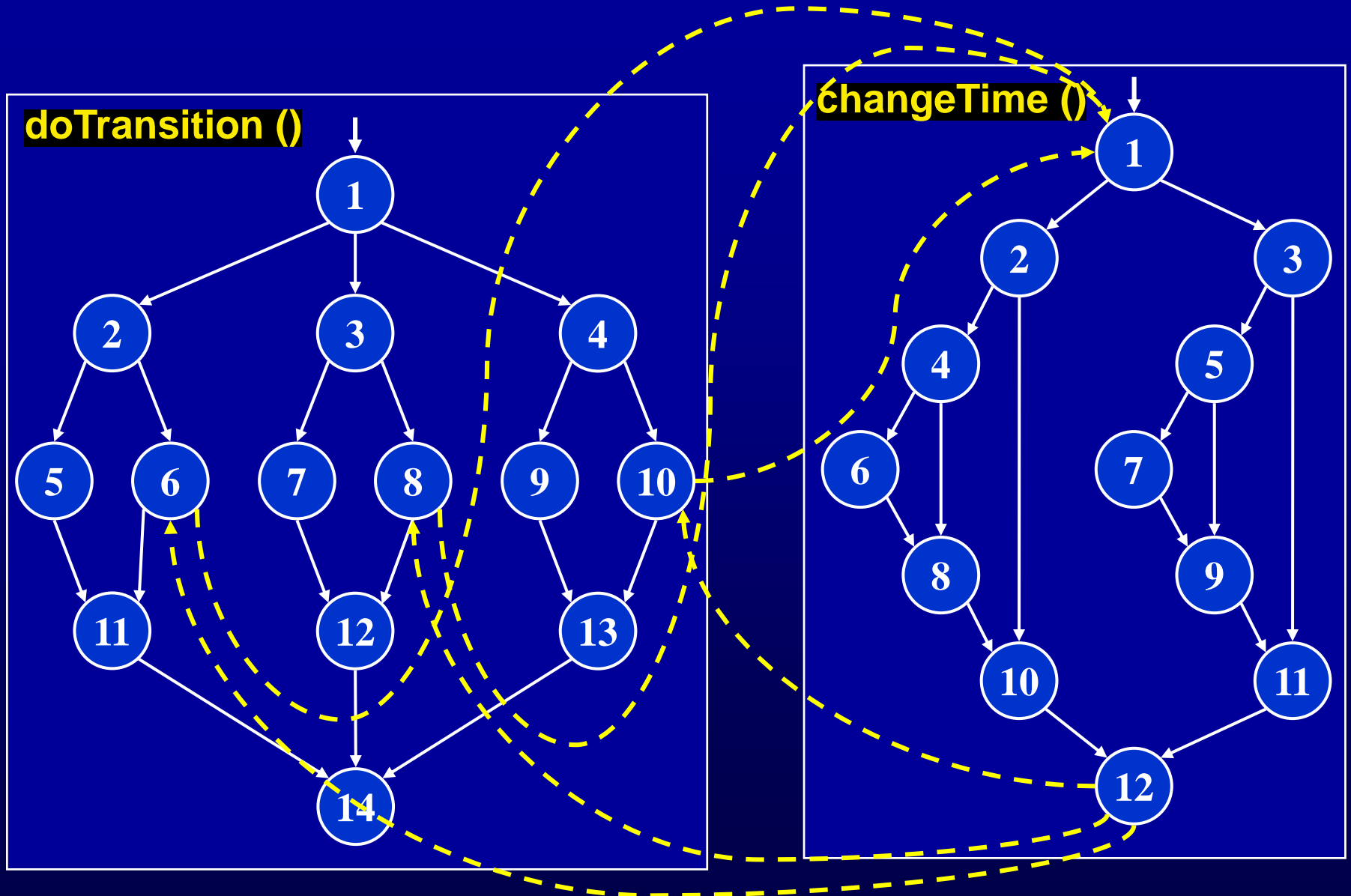
// Rolls around when necessary.

```
public void changeTime (int button)
{
    if (button == UP)
    {
        minute += 1;
        if (minute >= 60)
        {
            minute = 0;
            hour += 1;
            if (hour > 12)
                hour = 1;
        }
    }
    else if (button == DOWN)
    {
        minute -= 1;
        if (minute < 0)
        {
            minute = 59;
            hour -= 1;
            if (hour <= 0)
                hour = 12;
        }
    }
} // end changeTime()
```

1. Combining Control Flow Graphs

- The **first instinct** for inexperienced developers is to **draw CFGs** and **link them together**
- This is really **not an FSM**
- Several **problems**
 - Methods must return to correct callsites—implicit **nondeterminism**
 - **Implementation** must be **available** before graph can be built
 - This graph **does not scale up**
- Watch example ...

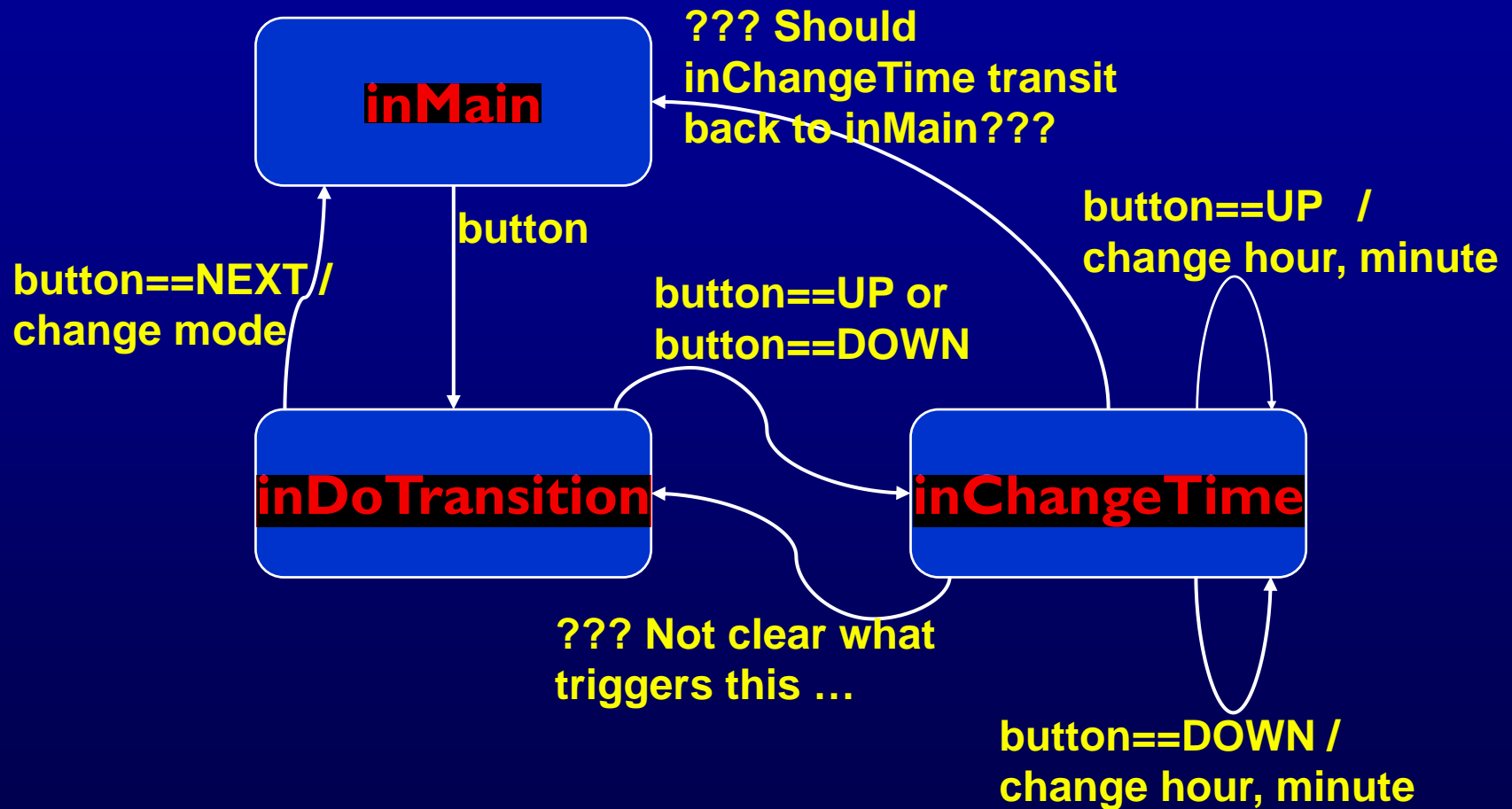
CFGs for Watch



2. Using the Software Structure

- A more experienced programmer may map methods to states
- These are really not states
- Problems
 - Subjective—different testers get different graphs
 - Requires in-depth knowledge of implementation
 - Detailed design must be present
- Watch example ...

Software Structure for Watch



3. Modeling State Variables

- More mechanical
- State variables are usually defined early
- First identify all state variables, then choose which are relevant
- In theory, every combination of values for the state variables defines a different state
- In practice, we must identify ranges, or sets of values, that are all in one state
- Some states may not be feasible

State Variables in Watch

Constants

- ~~NEXT, UP, DOWN~~
- ~~TIME, STOPWATCH, ALARM~~

Not relevant,
really just values

Non-constant variables in class Watch

- `int mode` (values: `TIME`, `STOPWATCH`, `ALARM`)
- `Time watch`, `stopwatch`, `alarm`

State Variables in Time

Non-constant variables in class Time

- int **hour** (values: 1..12)
- int **minute** (values: 0 .. 59)

12 X 60 values is 720 states

Clearly, that is too many

Combine values into **ranges of similar values** :

- **hour** : 1..11, 12
- **minute** : 0, 1..59

Clumsy ... Not sequential ...

let's combine hour and minute ...

Four states : (1..11, 0); (12, 0); (1..11, 1..59); (12, 1..59)

Time : 12:00, 12:01..12:59, 01:00 .. 11:59

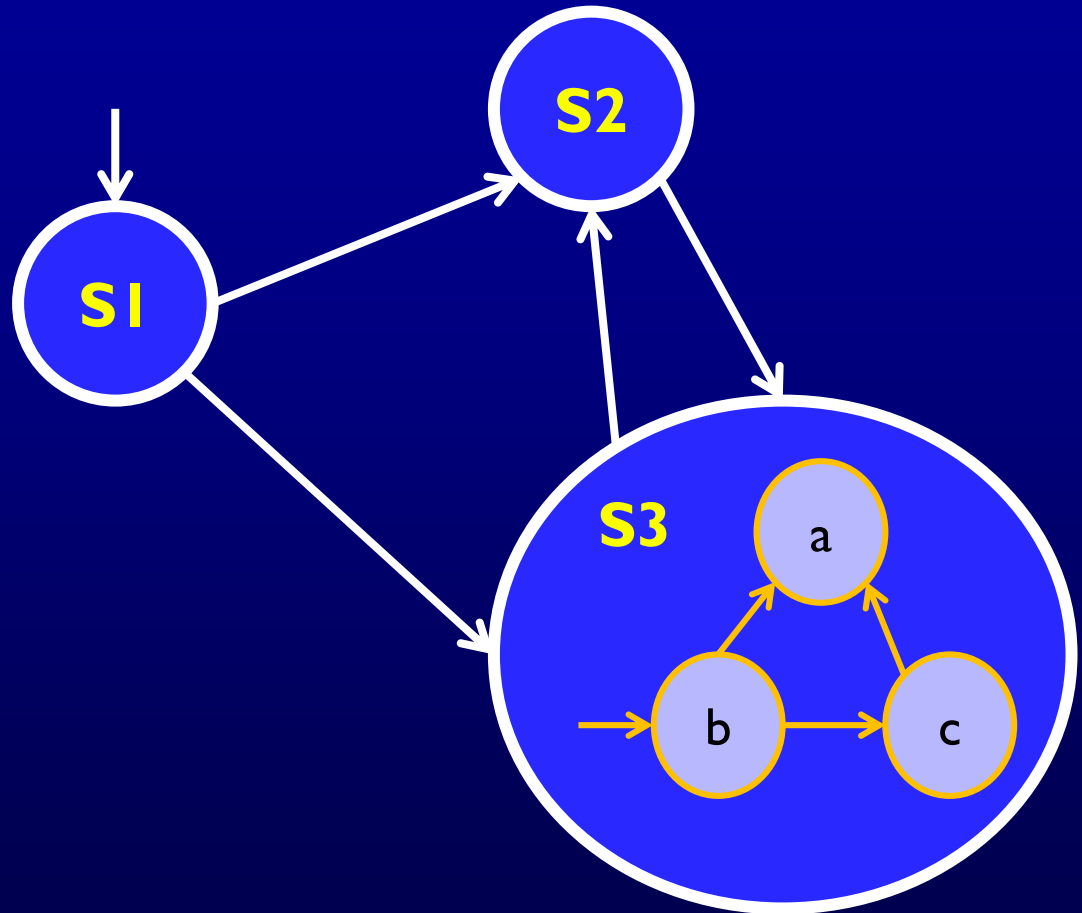
These require lots of
thought and
semantic domain
knowledge of the
program

Hierarchical FSMs

One **FSM** is **contained** within the other

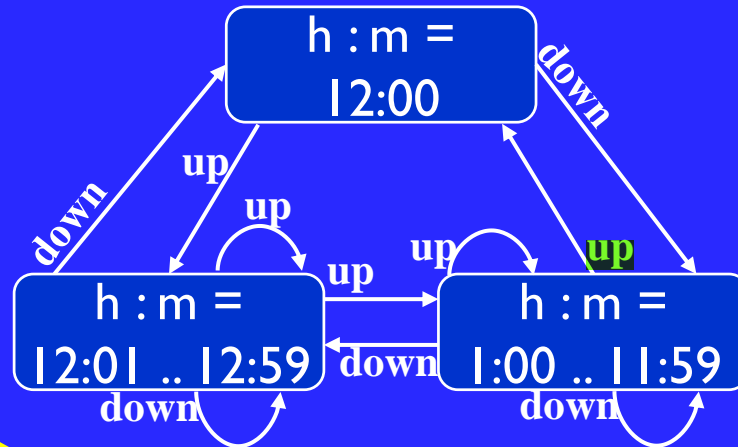
Class Watch uses
class Time

How can we model
two classes—one
that uses another?

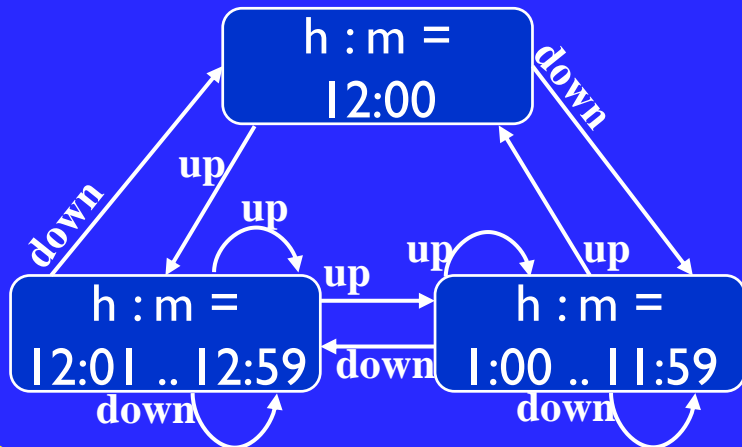


Watch / Time Hierarchical FSM

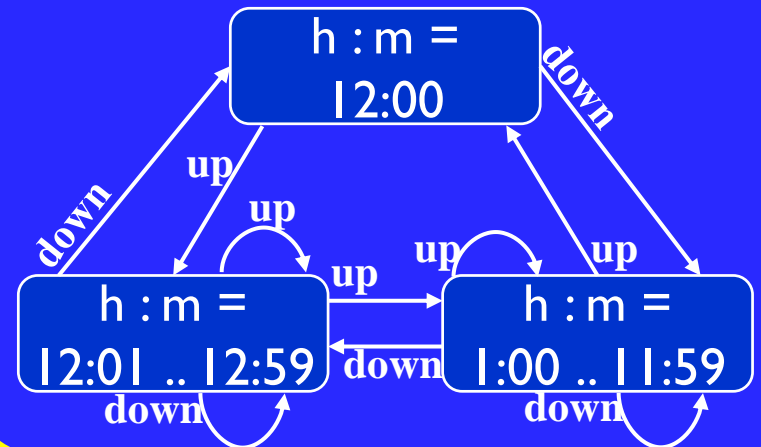
mode = TIME



mode = STOPWATCH



mode = ALARM



Summary–Tradeoffs in Applying Graph Coverage Criteria to FSMs

- **Two advantages**
 1. **Tests** can be **designed before implementation**
 2. **Analyzing FSMs** is much **easier** than analyzing source
- **Three disadvantages**
 1. Some **implementation** decisions are **not modeled in the FSM**
 2. There is some **variation in the results** because of the subjective nature of **deriving FSMs**
 3. **Tests** have to be “**mapped**” to **actual inputs** to the program – the **names** that appear in the **FSM** may not be the same as the **names** in the **program**