

# شبکه های کامپیوتری ۲

جلسه ۱۳ فصل ۹

**Multimedia Applications**

دانشگاه صنعتی اصفهان  
دانشکده مهندسی برق و کامپیوتر

# Chapter 9

## Multimedia

## Networking

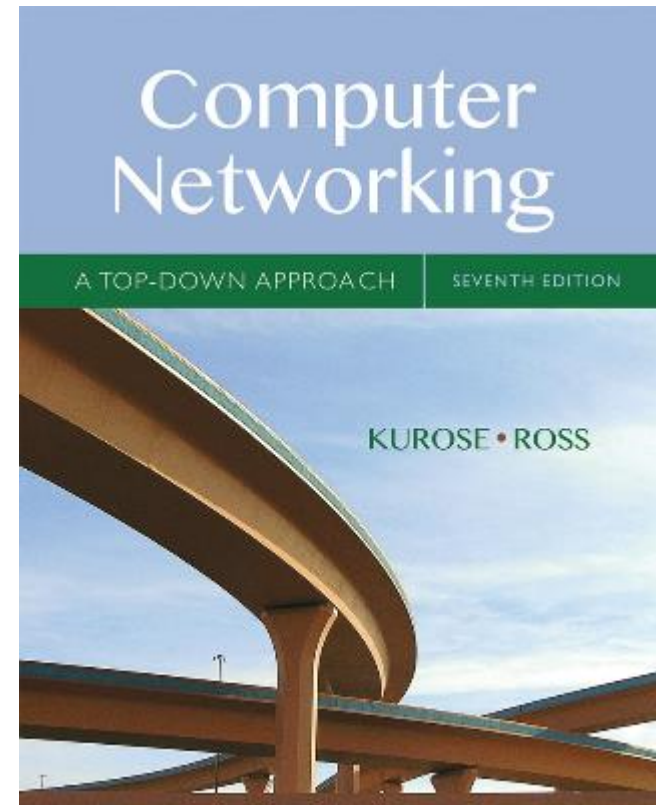
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## Computer Networking: A Top Down Approach

7<sup>th</sup> edition

Jim Kurose, Keith Ross

Pearson/Addison Wesley

April 2016

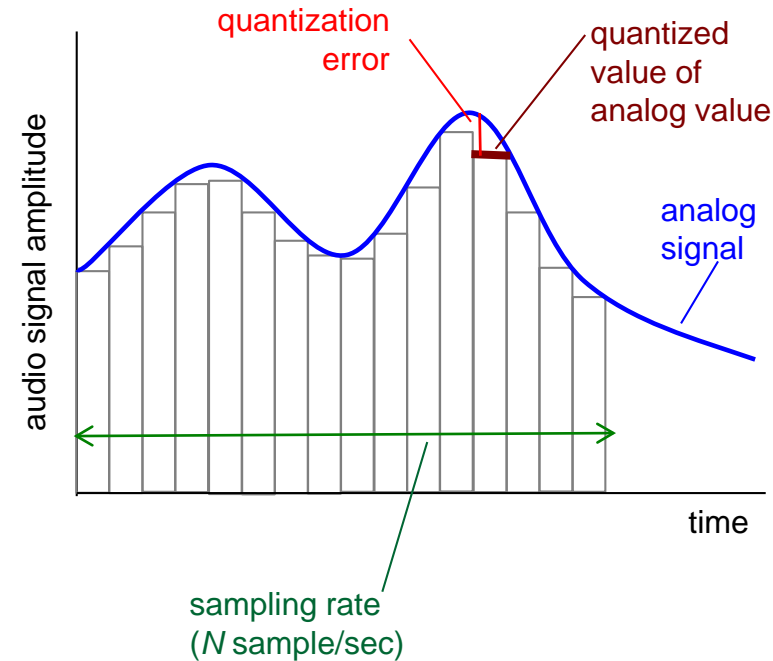
# Multimedia networking: outline

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- 9.1 multimedia networking applications
- 9.2 streaming *stored* video
- 9.3 voice-over-IP
- 9.4 protocols for *real-time* conversational applications
- 9.5 network support for multimedia

# Multimedia: audio

- analog audio signal sampled at constant rate
  - telephone: 8,000 samples/sec
  - CD music: 44,100 samples/sec
- each sample quantized, i.e., rounded
  - e.g.,  $2^8=256$  possible quantized values
  - each quantized value represented by bits, e.g., 8 bits for 256 values

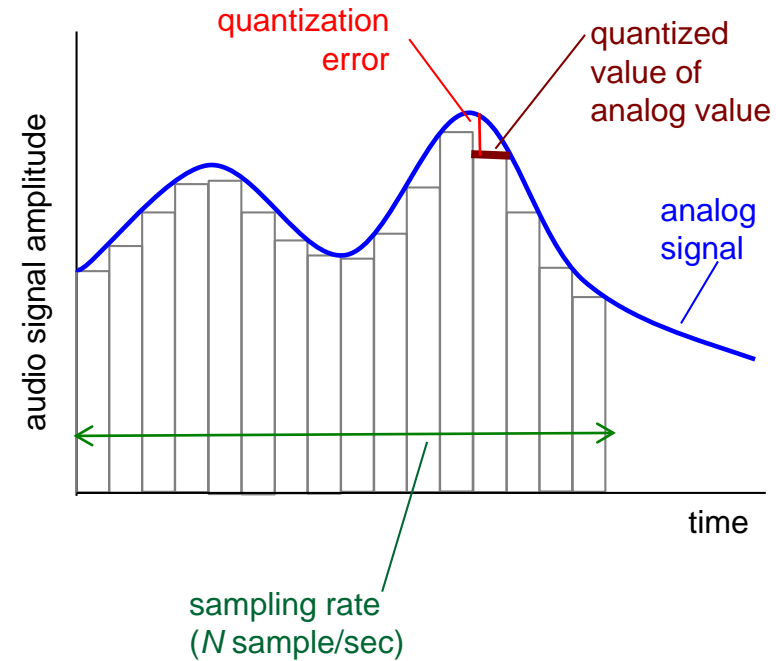


# Multimedia: audio

- example: 8,000 samples/sec, 256 quantized values: 64,000 bps
- receiver converts bits back to analog signal:
  - some quality reduction

## example rates

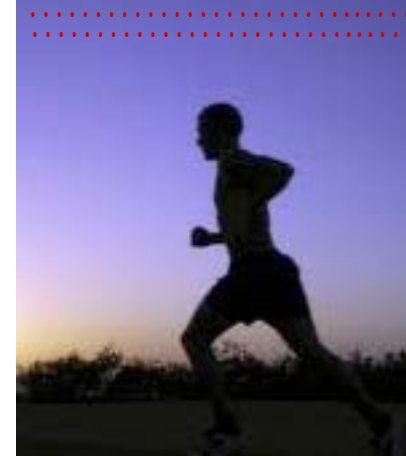
- CD: 1.411 Mbps
- MP3: 96, 128, 160 kbps
- Internet telephony: 5.3 kbps and up



# Multimedia: video

- video: sequence of images displayed at constant rate
  - e.g., 24 images/sec
- digital image: array of pixels
  - each pixel represented by bits
- coding: use redundancy *within* and *between* images to decrease # bits used to encode image
  - spatial (within image)
  - temporal (from one image to next)

*spatial coding example:* instead of sending  $N$  values of same color (all purple), send only two values: color value (*purple*) and number of repeated values ( $N$ )



frame  $i$

*temporal coding example:* instead of sending complete frame at  $i+1$ , send only differences from frame  $i$

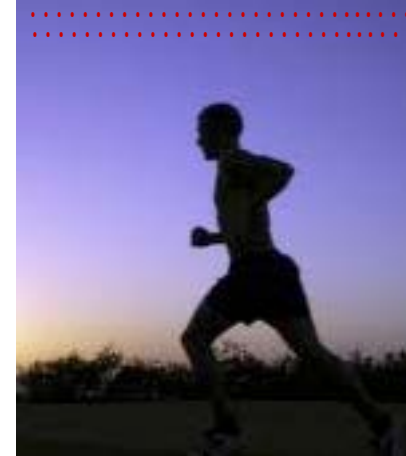


frame  $i+1$

# Multimedia: video

- **CBR: (constant bit rate):**  
video encoding rate fixed
- **VBR: (variable bit rate):**  
video encoding rate changes  
as amount of spatial,  
temporal coding changes
- **examples:**
  - MPEG I (CD-ROM) 1.5 Mbps
  - MPEG2 (DVD) 3-6 Mbps
  - MPEG4 (often used in Internet, < 1 Mbps)

*spatial coding example:* instead of sending  $N$  values of same color (all purple), send only two values: color value (*purple*) and number of repeated values ( $N$ )



frame  $i$

*temporal coding example:* instead of sending complete frame at  $i+1$ , send only differences from frame  $i$



frame  $i+1$

# Multimedia networking: 3 application types

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- *streaming, stored* audio, video
  - *streaming*: can begin playout before downloading entire file
  - *stored (at server)*: can transmit faster than audio/video will be rendered (implies storing/buffering at client)
  - e.g., YouTube, Netflix, Hulu
- *conversational* voice/video over IP
  - interactive nature of human-to-human conversation limits delay tolerance
  - e.g., Skype
- *streaming live* audio, video
  - e.g., live sporting event (futbol)



# Multimedia networking: outline

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9.1 multimedia networking applications

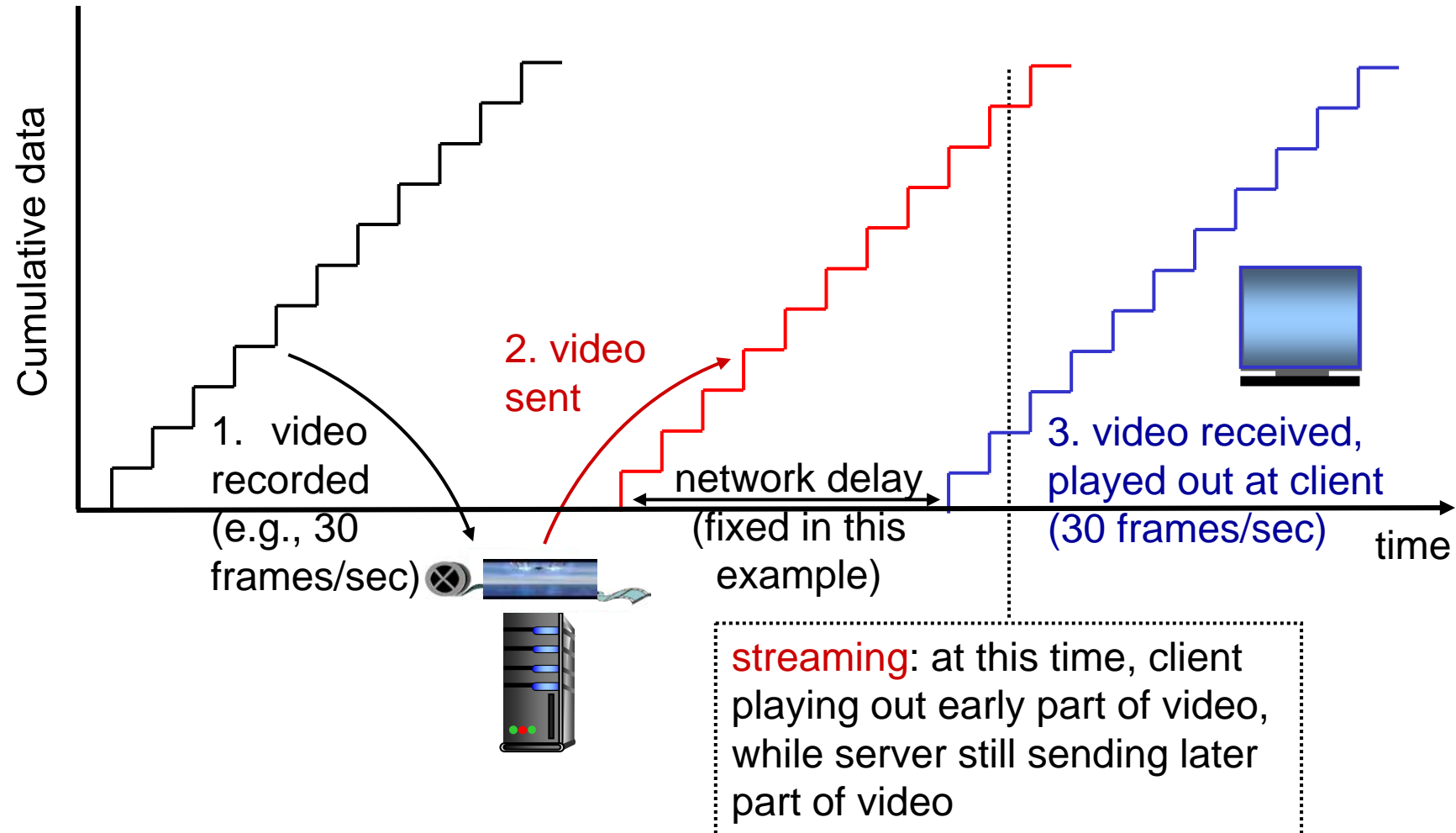
9.2 streaming *stored* video

9.3 voice-over-IP

9.4 protocols for *real-time* conversational applications

9.5 network support for multimedia

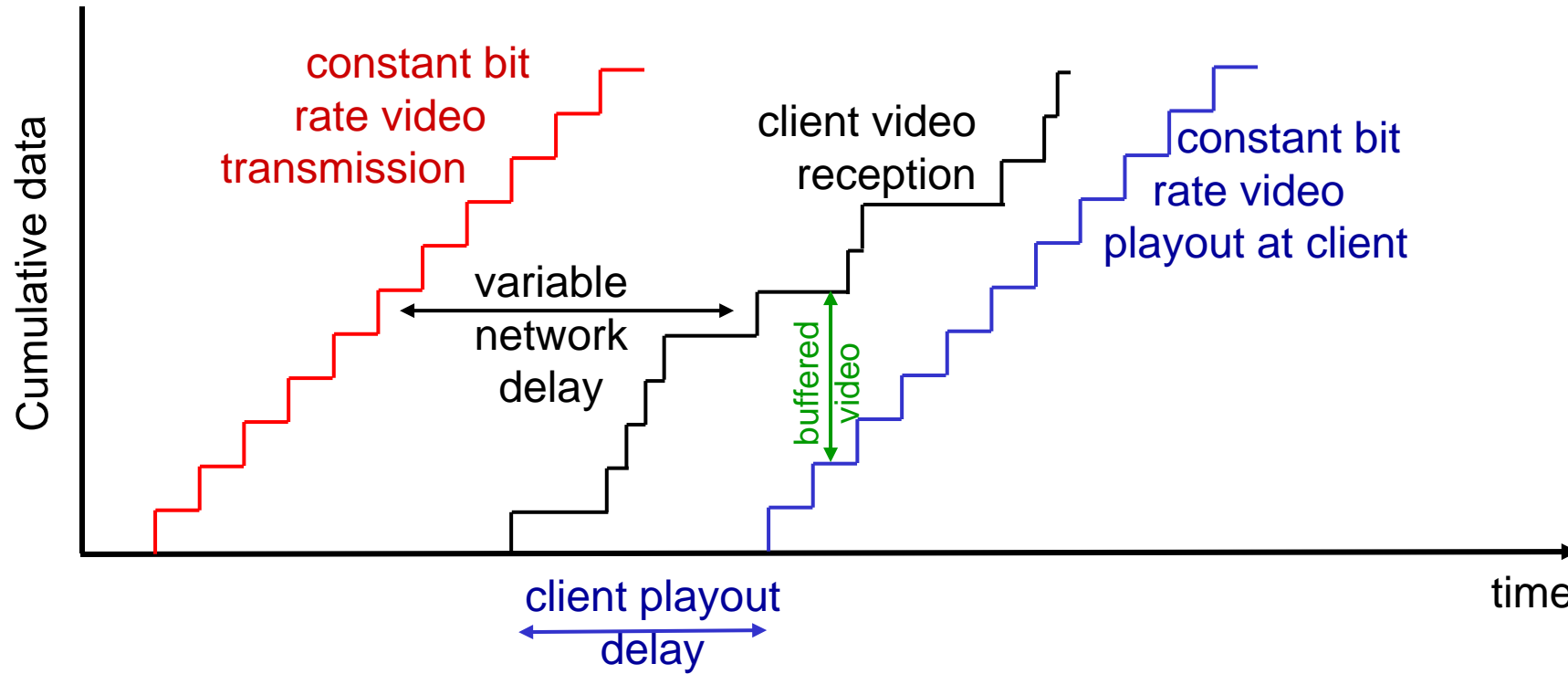
# Streaming stored video:



# Streaming stored video: challenges

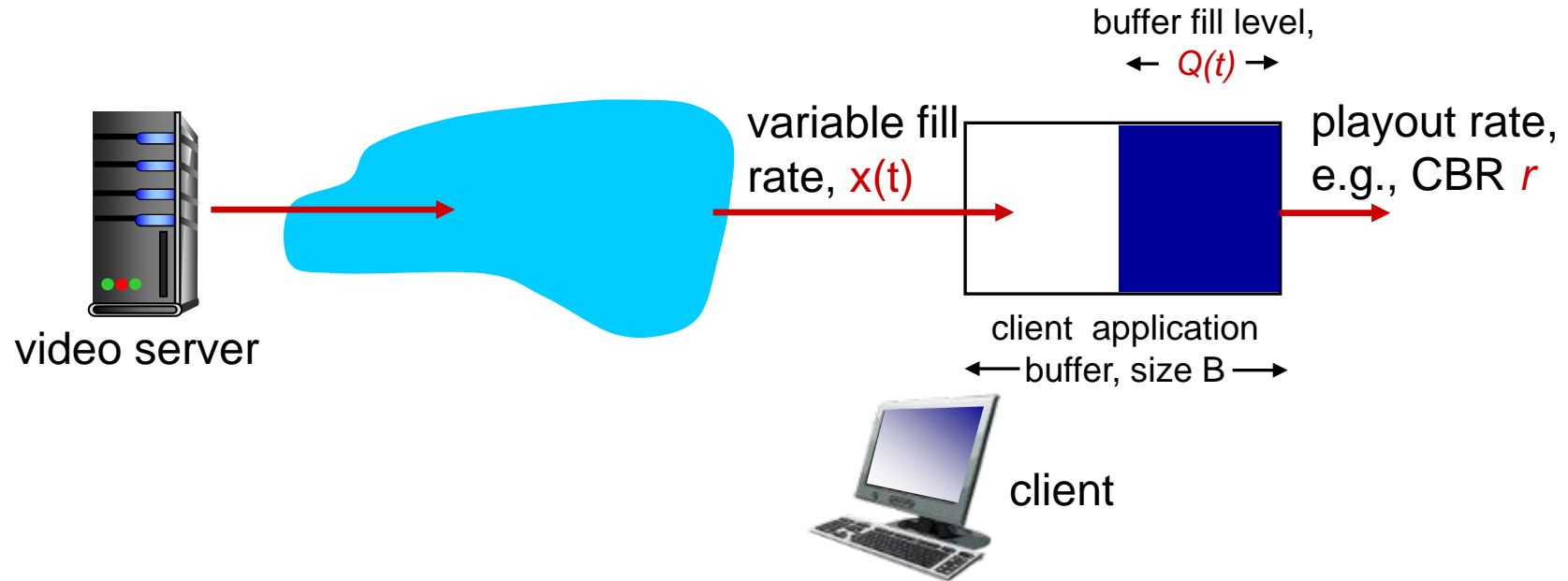
- **continuous playout constraint**: once client playout begins, playback must match original timing
  - ... but **network delays are variable** (jitter), so will need **client-side buffer** to match playout requirements
- other challenges:
  - client interactivity: pause, fast-forward, rewind, jump through video
  - video packets may be lost, retransmitted

# Streaming stored video: revisited

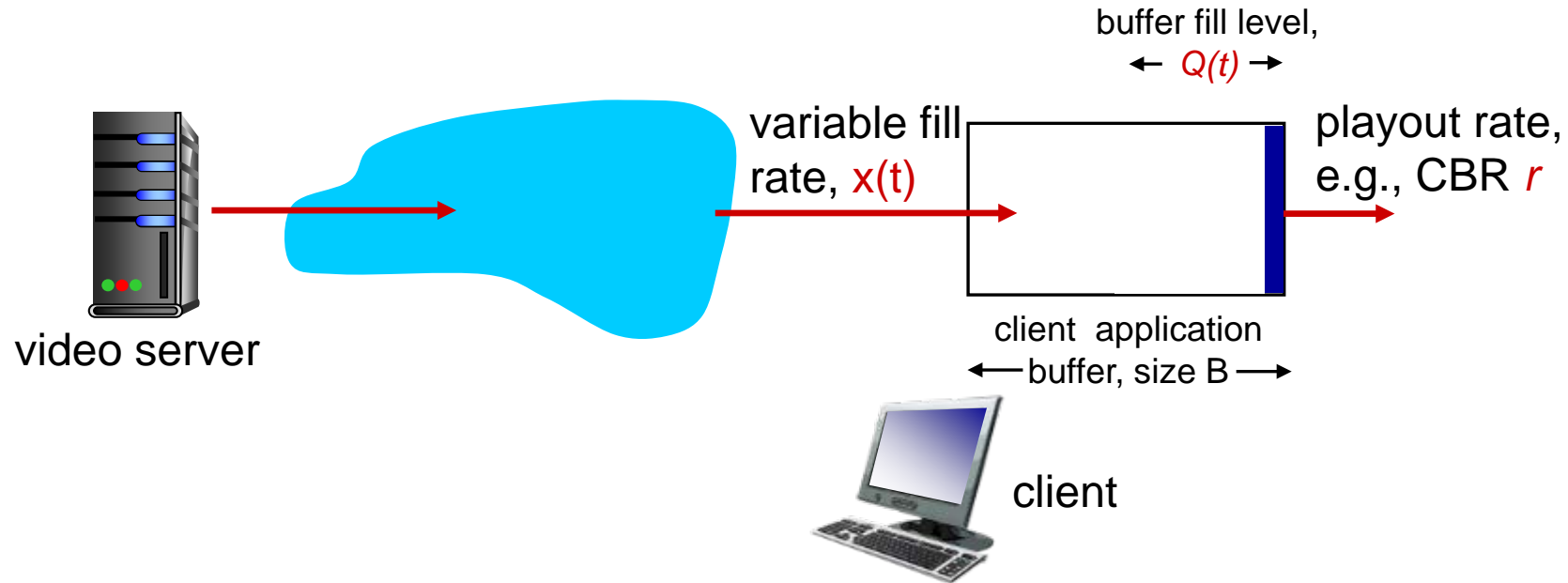


- *client-side buffering and playout delay*: compensate for network-added delay, delay jitter

# Client-side buffering, playout

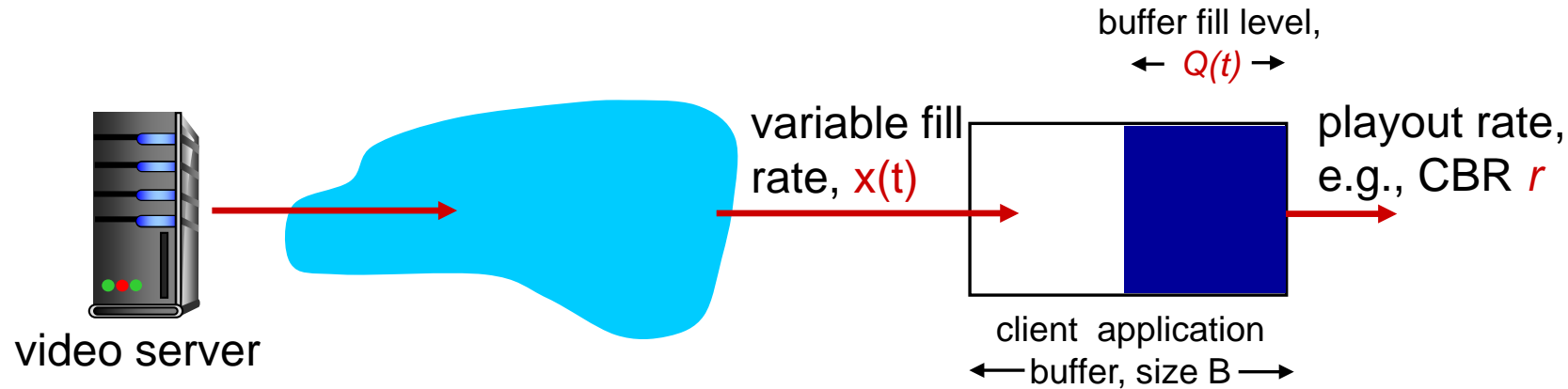


# Client-side buffering, playout



1. Initial fill of buffer until playout begins at  $t_p$
2. playout begins at  $t_p$ ,
3. buffer fill level varies over time as fill rate  $x(t)$  varies and playout rate  $r$  is constant

# Client-side buffering, playout



*playout buffering: average fill rate ( $\bar{x}$ ), playout rate ( $r$ ):*

- $\bar{x} < r$ : buffer eventually empties (causing freezing of video playout until buffer again fills)
- $\bar{x} > r$ : buffer will not empty, provided initial playout delay is large enough to absorb variability in  $x(t)$ 
  - *initial playout delay tradeoff*: buffer starvation less likely with larger delay, but larger delay until user begins watching

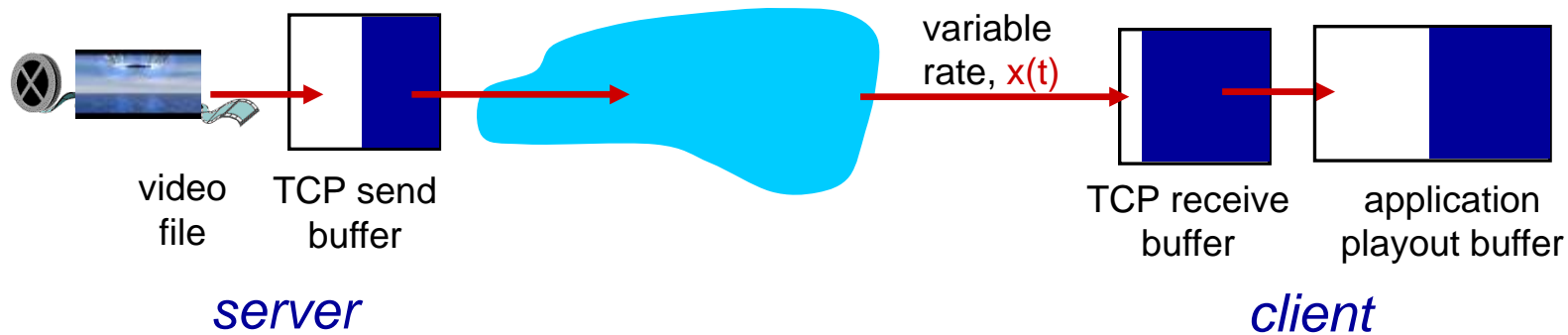
# Streaming multimedia: UDP

- server sends at rate appropriate for client
  - often: send rate = encoding rate = constant rate
  - transmission rate can be oblivious to congestion levels
- short playout delay (2-5 seconds) to remove network jitter
- error recovery: application-level, time permitting
- RTP [RFC 2326]: multimedia payload types
- UDP may *not* go through firewalls



# Streaming multimedia: HTTP

- multimedia file retrieved via HTTP GET
- send at maximum possible rate under TCP



- fill rate fluctuates due to TCP congestion control, retransmissions (in-order delivery)
- larger playout delay: smooth TCP delivery rate
- HTTP/TCP passes more easily through firewalls

# Multimedia networking: outline

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9.1 multimedia networking applications

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9.3 voice-over-IP

9.4 protocols for *real-time* conversational applications

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# Voice-over-IP (VoIP)

- *VoIP end-end-delay requirement*: needed to maintain “conversational” aspect
  - higher delays noticeable, impair interactivity
  - < 150 msec: good
  - > 400 msec bad
  - includes application-level (packetization, playout), network delays
- *session initialization*: how does callee advertise IP address, port number, encoding algorithms?
- *value-added services*: call forwarding, screening, recording
- *emergency services*: 911

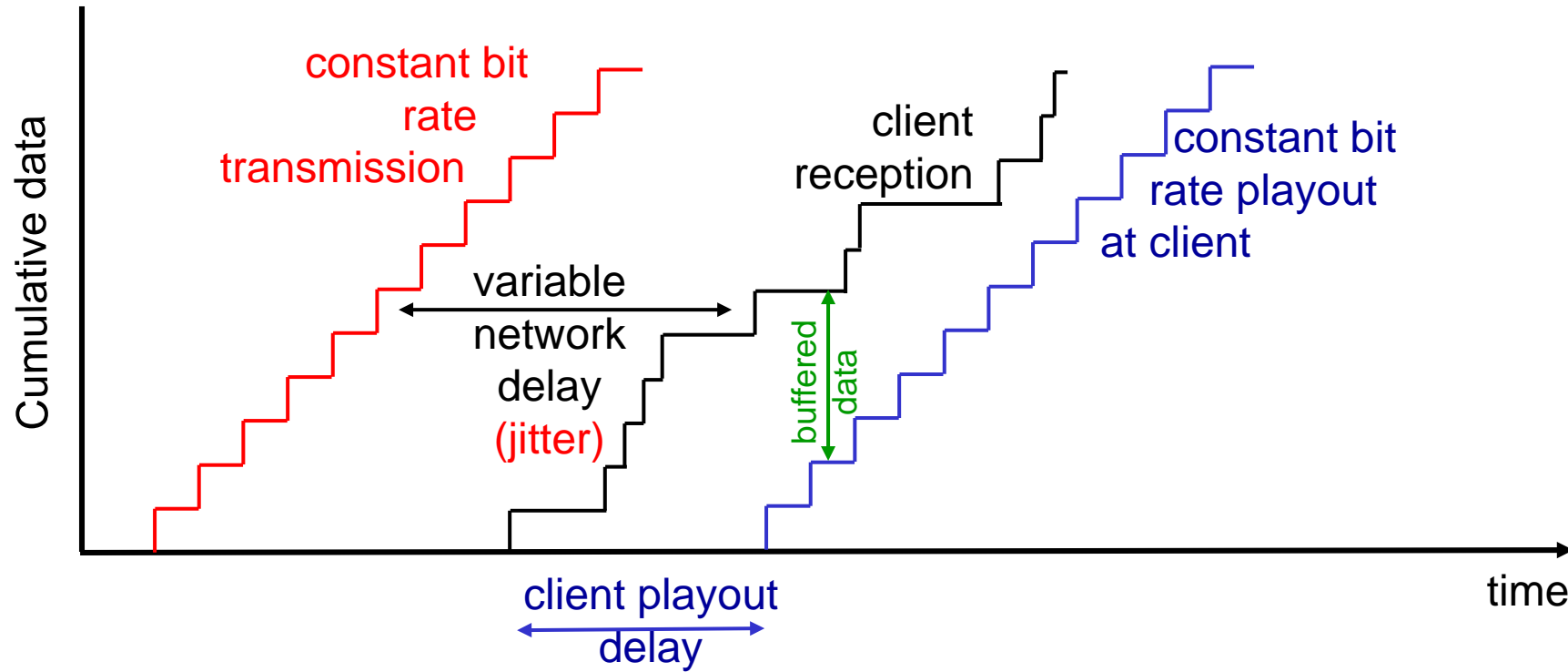
# VoIP characteristics

- speaker's audio: alternating talk spurts, silent periods.
  - 64 kbps during talk spurt
  - pkts generated only during talk spurts
  - 20 msec chunks at 8 Kbytes/sec: 160 bytes of data
- application-layer header added to each chunk
- chunk+header encapsulated into UDP or TCP segment
- application sends segment into socket every 20 msec during talkspurt

# VoIP: packet loss, delay

- *network loss*: IP datagram lost due to network congestion (router buffer overflow)
- *delay loss*: IP datagram arrives too late for playout at receiver
  - delays: processing, queueing in network; end-system (sender, receiver) delays
  - typical maximum tolerable delay: 400 ms
- *loss tolerance*: depending on voice encoding, loss concealment, packet loss rates between 1% and 10% can be tolerated

# Delay jitter



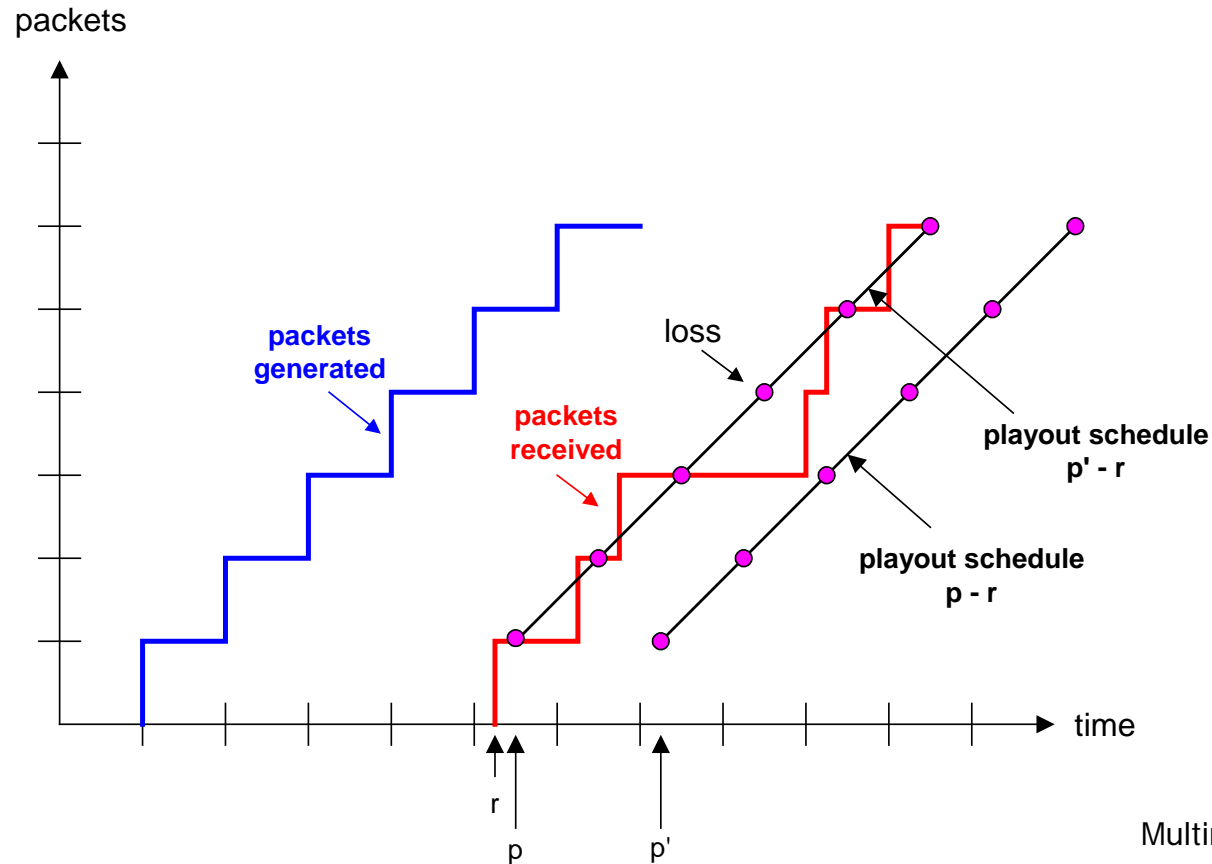
- end-to-end delays of two consecutive packets: difference can be more or less than 20 msec (transmission time difference)

# VoIP: fixed playout delay

- receiver attempts to playout each chunk exactly  $q$  msecs after chunk was generated.
  - chunk has time stamp  $t$ : play out chunk at  $t+q$
  - chunk arrives after  $t+q$ : data arrives too late for playout: data “lost”
- tradeoff in choosing  $q$ :
  - *large  $q$ : less packet loss*
  - *small  $q$ : better interactive experience*

# VoIP: fixed playout delay

- sender generates packets every 20 msec during talk spurt.
- first packet received at time  $r$
- first playout schedule: begins at  $p$
- second playout schedule: begins at  $p'$





# Adaptive playout delay (I)

- **goal:** low playout delay, low late loss rate
- **approach:** adaptive playout delay adjustment:
  - estimate network delay, adjust playout delay at beginning of each talk spurt
  - silent periods compressed and elongated
  - chunks still played out every 20 msec during talk spurt
- adaptively estimate packet delay: (EWMA - exponentially weighted moving average, **recall TCP RTT estimate**):

$$d_i = (1-\alpha)d_{i-1} + \alpha (r_i - t_i)$$

*delay estimate after ith packet*      *small constant, e.g. 0.1*      *time received - time sent (timestamp)*  
*measured delay of ith packet*

# Adaptive playout delay (2)

- also useful to estimate average deviation of delay,  $v_i$ :

$$v_i = (1-\beta)v_{i-1} + \beta |r_i - t_i - d_i|$$

- estimates  $d_i$ ,  $v_i$  calculated for every received packet, but used only at start of talk spurt
- for first packet in talk spurt, playout time is:

$$\text{playout-time}_i = t_i + d_i + Kv_i$$

- remaining packets in talkspurt are played out periodically

# Adaptive playout delay (3)

Q: How does receiver determine whether packet is first in a talkspurt?

- if no loss, receiver looks at successive timestamps
  - difference of successive stamps  $> 20$  msec --> talk spurt begins.
- with loss possible, receiver must look at both time stamps and sequence numbers
  - difference of successive stamps  $> 20$  msec *and* sequence numbers without gaps --> talk spurt begins.

# VoiP: recovery from packet loss (I)

*Challenge:* recover from packet loss given small tolerable delay between original transmission and playout

- each ACK/NAK takes  $\sim$  one RTT
- alternative: *Forward Error Correction (FEC)*
  - send enough bits to allow recovery without retransmission (recall two-dimensional parity in Ch. 5)

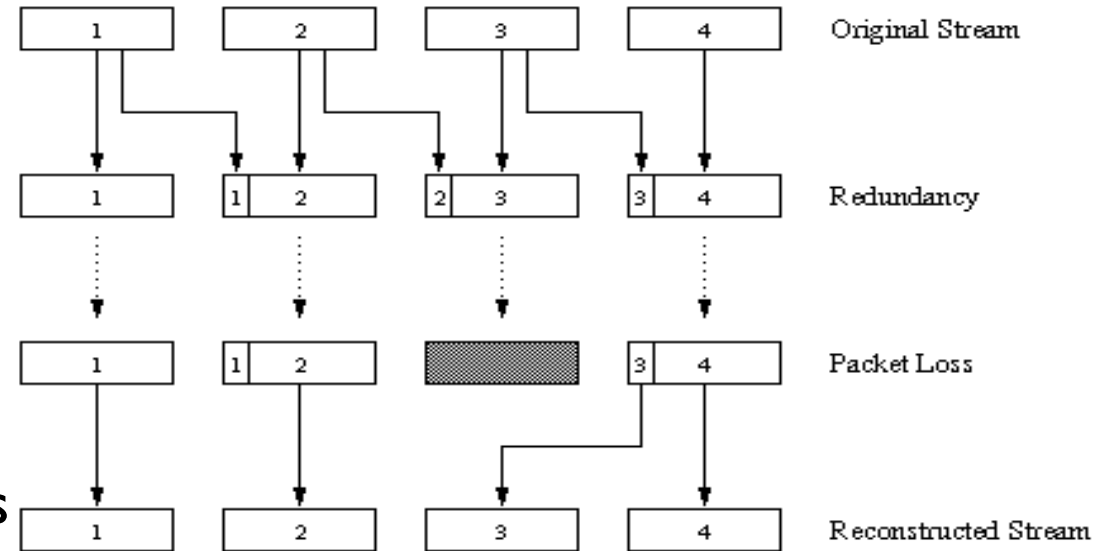
## *simple FEC*

- for every group of  $n$  chunks, create redundant chunk by exclusive OR-ing  $n$  original chunks
- send  $n+1$  chunks, increasing bandwidth by factor  $1/n$
- can reconstruct original  $n$  chunks if at most one lost chunk from  $n+1$  chunks, with playout delay

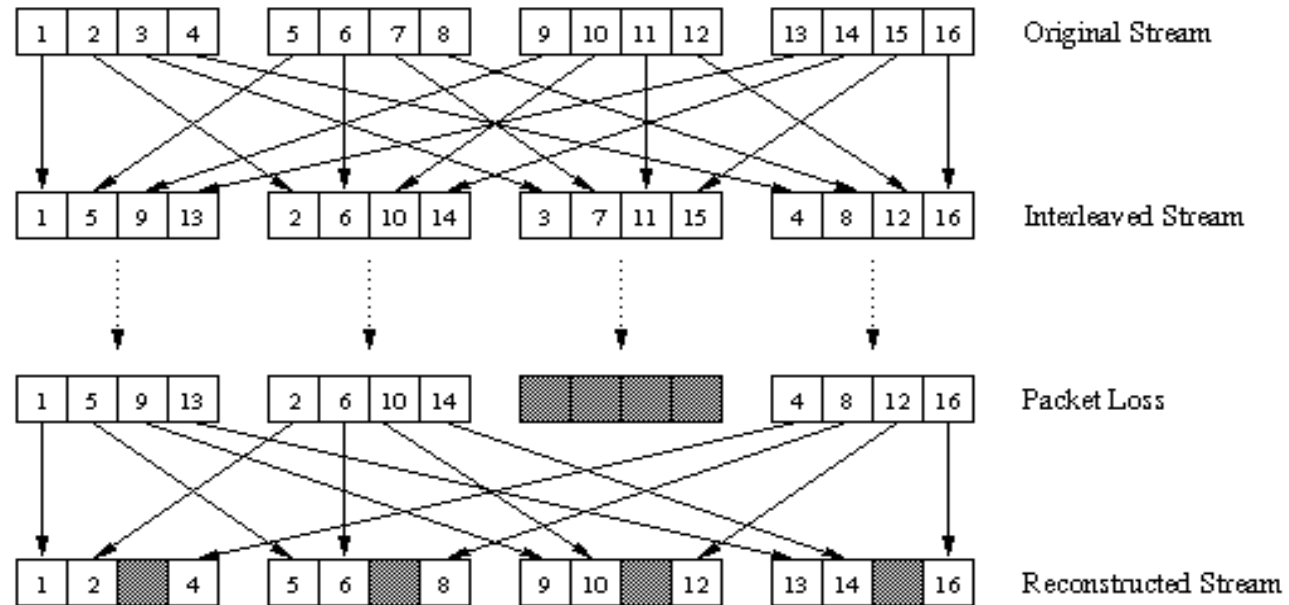
# VoiP: recovery from packet loss (2)

## another FEC scheme:

- “piggyback lower quality stream”
- send lower resolution audio stream as redundant information
- e.g., nominal stream PCM at 64 kbps and redundant stream GSM at 13 kbps
- non-consecutive loss: receiver can conceal loss
- generalization: can also append (n-1)st and (n-2)nd low-bit rate chunk



# VoiP: recovery from packet loss (3)



## *interleaving to conceal loss:*

- audio chunks divided into smaller units, e.g. four 5 msec units per 20 msec audio chunk
- packet contains small units from different chunks
- if packet lost, still have *most* of every original chunk
- no redundancy overhead, but increases playout delay