Let us define an optimal solution $O = \{L_{x1}, L_{x2}, L_{x3}, \ldots, L_{xn}\}$ and our greedy algorithm solution $G = \{L_{y1}, L_{y2}, L_{y3}, \ldots, L_{ym}\}$. Assuming that these two solutions are not the same, this means that there are elements in O that might not be in O and hence the orders of O and O might be different or there are elements in O in a different order than O. We can **exchange** the elements between the sets to make them similar to each other. If for example, the optimal algorithm stopped at a stop for O mins, we'd see my greedy algorithm stopping for O mins at O different stops. Both algorithms will give a total time of O mins and hence can be considered optimal for the first case. The second case is not possible as different orderings can not take place because O and O can't go in a reverse order. Hence we can conclude that O is also the optimal algorithm.