Hive queries

- Select statements
- Joins
- Built-in functions
- Create table as select (CTAS)
- Writing query results

HiveQL

- Basic syntax
- Joining Datasets
- Built-in Functions
- Simplifying queries using views
- Saving queries to tables
- Explaining the query

An Introduction to HiveQL

- HiveQL is Hive's query language
 - Based on a subset of SQL-92, plus Hive-specific extensions
- Some limitations compared to 'standard' SQL
 - Some features are not supported
 - Others are only partially implemented

Only non-materialized views

Joins are limited (only equality)

- HiveQL also has some features not offered in SQL
 - Maps, array and structs
 - Built-in table partitioning
 - Multi-table inserts

Selecting Data from Hive Tables

- The SELECT statement retrieves data from Hive tables
 - Can specify an ordered list of individual columns

```
hive> SELECT cust_id, fname, lname FROM customers;
```

- An asterisk matches all columns in the table

```
hive> SELECT * FROM customers;
```

Limiting and Sorting Query Results

• The LIMIT clause sets the maximum number of rows returned

```
hive> SELECT fname, lname FROM customers LIMIT 10;
```

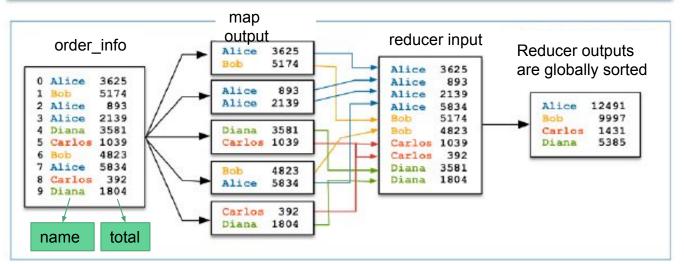
- Caution: no guarantee regarding which 10 results are returned
 - Use ORDER BY for top-N queries
 - The field(s) you ORDER BY must be selected

```
hive> SELECT cust_id, fname, lname FROM customers ORDER BY cust_id DESC LIMIT 10;
```

Sorting: use ORDER BY for complete sort

- As in SQL, ORDER BY sorts specified fields in HiveQL
 - Consider the result from the following query

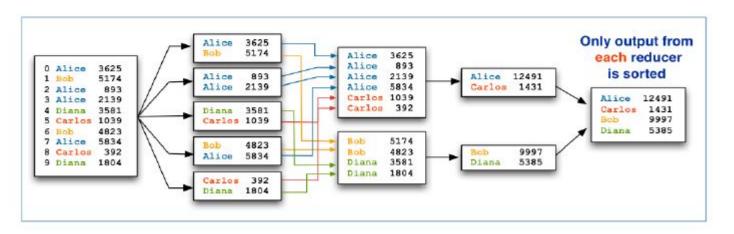
```
hive> SELECT name, SUM(total)
FROM order_info GROUP BY name
ORDER BY name;
```



Sorting: use SORT BY for partial sort

- HiveQL also supports partial ordering via SORT BY
 - Offers much better performance if global order isn't required

```
hive> SELECT name, SUM(total)
FROM order_info GROUP BY name
SORT BY name;
```



Using a WHERE Clause to Restrict Results

- •WHERE clauses restrict rows to those matching specified criteria
 - String comparisons are case-sensitive

```
hive> SELECT * FROM orders WHERE order_id=1287;
```

```
hive> SELECT * FROM customers WHERE state
IN ('CA', 'OR', 'WA', 'NV', 'AZ');
```

You can combine expressions using AND or OR

```
hive> SELECT * FROM customers
WHERE fname LIKE 'Ann%'
AND (city='Seattle' OR city='Portland');
```

Table Aliases

Table aliases can help simplify complex queries

```
hive> SELECT o.order_date, c.fname, c.lname
   FROM customers c JOIN orders o
   ON c.cust_id = o.cust_id
   WHERE c.zipcode='94306';
```

Combining Query Results with UNION ALL

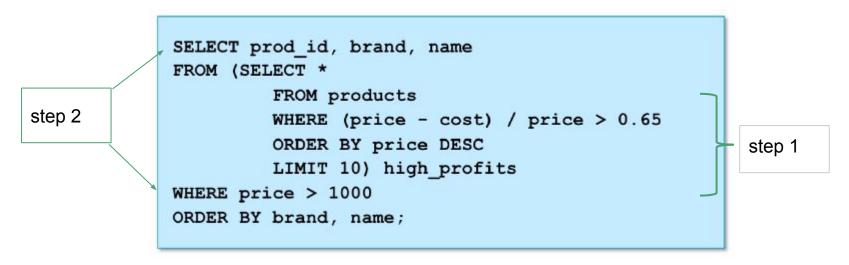
- Unifies output from SELECTs into a single result set
 - The name, order, and types of columns in each query must match
 - Hive only supports UNION ALL

```
SELECT emp_id, fname, lname, salary
FROM employees
WHERE state='CA' AND salary > 75000
UNION ALL
SELECT emp_id, fname, lname, salary
FROM employees
WHERE state != 'CA' AND salary > 50000;
```

UNION ALL can also be used with subqueries

Subqueries in Hive

Hive supports subqueries in the FROM clause of the SELECT statement



- Hive allows arbitrary levels of subqueries
 - Each subquery must be named (like high profits above)

Performance Patterns in HiveQL (1)

The fastest queries involve only metadata

```
DESCRIBE customers;
```

The next fastest simply read from HDFS

```
SELECT * FROM customers;
```

Then a query that requires a map-only job

```
SELECT * FROM customers WHERE zipcode = 94305;
```

Performance Patterns in HiveQL (2)

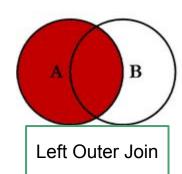
The next slowest type of query requires both Map and Reduce phases

```
SELECT COUNT(cust_id) FROM customers
WHERE zipcode=94305;
```

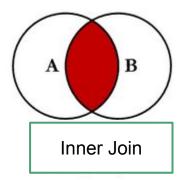
The slowest queries require multiple MapReduce jobs

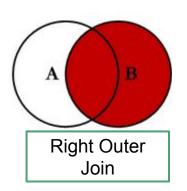
```
SELECT zipcode, COUNT(cust_id) AS num FROM customers
GROUP BY zipcode
ORDER BY num DESC
LIMIT 10;
```

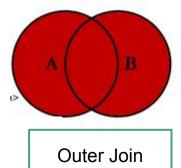
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SQL JOINS







Review: Joins in MapReduce

The rundown on joins (1)

Basic idea for Map-side joins:

- Load one set of data into memory, stored in a hash table
 - Key of the hash table is the join key
- Map over the other set of data, and perform a lookup on the hash table using the join key
- If the join key is found, you have a successful join
 - Otherwise, do nothing

The rundown on joins (2)

- For a Reduce-side join, the basic concept is:
 - Map over both data sets
 - Emit a (key, value) pair for each record
 - Key is the join key, value is the entire record
 - In the Reducer, do the actual join
 - Because of the Shuffle and Sort, values with the same key are brought together

Joins in Hive use MapReduce joins

Joins in Hive

- Joining disparate data sets is a common operation in Hive
- Hive supports several types of joins
 - Inner joins
 - Outer joins (left, right, and full)
 - Cross joins (supported in Hive 0.10 and later)
 - Left semi joins
- Only equality conditions are allowed in joins
 - Valid: customers.cust_id = orders.cust_id
 - Invalid: customers.cust_id <> orders.cust_id
 - Outputs records where the specified key is found in each table

Join Syntax

Hive requires the following syntax for joins

```
SELECT c.cust_id, name, total
FROM customers c
JOIN orders o ON (c.cust_id = o.cust_id);
```

- The above example is an inner join
 - Can replace JOIN with another type (e.g. RIGHT OUTER JOIN)

Since Hive 0.13, implicit joins are supported

```
SELECT *
FROM table1 t1, table2 t2, table3 t3
WHERE t1.id = t2.id AND t2.id = t3.id AND t1.zipcode = '02535';
```

Left Semi Joins (1)

- A less common type of join is the LEFT SEMI JOIN
 - They are a special (and efficient) type of inner join
 - They behave more like a filter than a join
- Left semi joins include additional criteria in the ON clause
 - Only unique records that match these criteria are returned
 - Fields listed in SELECT are limited to the left-side table

```
SELECT c.cust_id
FROM customers c
LEFT SEMI JOIN orders o
ON (c.cust_id = o.cust_id
   AND YEAR(o.order_date) = '2012');
```

Left Semi Joins (2)

Traditional left semi-join

```
SELECT c.customer_id FROM customers c
LEFT SEMI JOIN orders o
ON (c.customer_id = o.order_customer_id)
AND (o.order_date > 2012);
```

Recently added to Hive

```
SELECT c.customer_id FROM customers c
WHERE EXISTS(
        SELECT * FROM orders o
        WHERE (c.customer_id = o.order_customer_id)
        AND (o.order_date > 2012)
);
```

Join performance

- Most tables to be joined are buffered into memory.
- By default, only the last table named (rightmost table) is streamed.
- To change the table to be streamed use a /* hint */

```
SELECT /*+ STREAMTABLE(a)*/ a.val, b.val, c.val FROM a
JOIN b ON (a.key=b.key1)
JOIN c on (c.key=b.key1)
```

Join optimizations

- Use Map side joins work when one table fits in memory:
 - table is loaded into memory as a hash table
 - the larger table is scanned

- "star-schemas" are optimized for you
 - Star: event (sales), dimensions (time, buyer info, and store)

https://cwiki.apache.org/confluence/display/Hive/LanguageManual+JoinOptimization

Example of a star-schema chained-join

```
Select count(*) cnt
From sales sl
     join demographics demo on (sl.demo_key = demo.key)
     join time t on (sl.time_key = t.key)
     join store s on (sl.store key = s.key)
Where
     t.hour = 8
     t.minute >= 30
     demo.num_children = 2
order by cnt;
```

Optimization: Attempts to load all dimensions into hashmaps and use the sales table as Mapper input.

Resume here: Lecture 8 - Advanced Hive

Question from last week(1)

In addition to SORT BY:

- DISTRIBUTE BY which governs how Hive distributes records among reducers: that is, it customizes the Partitioner
- syntax: SELECT col1, col2 FROM t1 DISTRIBUTE BY col1

Example: distribute by dept -> each dept has a partition -> within each dept, sort by dept and emp_id

SELECT employee, dept FROM staff DISTRIBUTE BY dept SORT BY emp_id ASC, dept DESC

Question from last week(2)

In addition to SORT BY:

- CLUSTER BY combines SORT BY and DISTRIBUTE BY
- syntax: SELECT col1, col2 FROM t1 CLUSTER BY col1 ASC

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Hive Functions

- Hive offers dozens of built-in functions
 - Many are identical to those found in SQL
 - Others are Hive-specific
- Example function invocation
 - Function names are not case-sensitive

```
hive> SELECT CONCAT(fname, ' ', lname) AS fullname FROM customers;
```

To see information about a function

```
hive> DESCRIBE FUNCTION UPPER;
UPPER(str) - Returns str with all characters
changed to uppercase
```

Example Built-in Functions (1)

These functions operate on numeric values

Function Description	Example Invocation	Input	Output
Rounds to specified # of decimals	ROUND(total_price, 2)	23.492	23.49
Returns nearest integer above	CEIL(total_price)	23.492	24
Returns nearest integer below	FLOOR(total_price)	23.492	23
Return absolute value	ABS(temperature)	-49	49
Returns square root	SQRT(area)	64	8
Returns a random number	RAND()		0.584977

Example Built-in Functions (2)

These functions operate on timestamp values

Function Description	Example Invocation	Input	Output
Convert to UNIX format	UNIX_TIMESTAMP(order_dt)	2013-06-14 16:51:05	1371243065
Convert to string format	FROM_UNIXTIME (mod_time)	1371243065	2013-06-14 16:51:05
Extract date portion	TO_DATE(order_dt)	2013-06-14 16:51:05	2013-06-14
Extract year portion	YEAR(order_dt)	2013-06-14 16:51:05	2013
Returns # of days between dates	DATEDIFF(order_dt, ship_dt)	2013-06-14, 2013-06-17	3

Example: using date field (order_date) in orders

```
hive> select order_date from orders;
hive> select from unixtime(order date) from orders;
```

In your homework you will create a new table, with a "fixed" date-time field so that you can then select as follows:

```
hive> select year(from_unixtime(order_date)) from orders_corrected;
```

Example Built-in Functions (3)

Here are some other interesting functions

Function Description	Example Invocation	Input	Output
Converts to uppercase	UPPER(fname)	Bob	BOB
Extract portion of string	SUBSTRING(name, 0, 2)	Alice	Al
Selectively return value	IF(price > 1000, 'A', 'B')	1500	А
Convert to another type	CAST (weight as INT)	3.581	3
Returns size of array or map	SIZE(array_field)	N/A	6

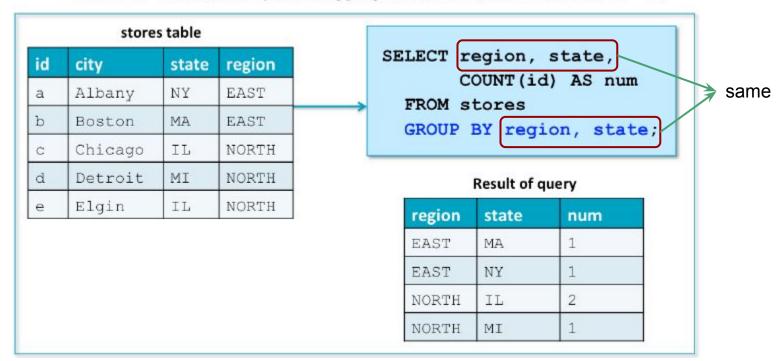
Built-in Aggregate Functions

Hive offers many aggregate functions, including

Function Description	Example Invocation	
Count all rows	COUNT (*)	
Count all rows where field is not null	COUNT (fname)	
Count all rows where field is unique and not null	COUNT (DISTINCT fname)	
Returns the largest value	MAX(salary)	
Returns the smallest value	MIN(salary)	
Adds all supplied values and returns result	SUM(price)	
Returns the average of all supplied values	AVG(salary)	

Group By

- •GROUP BY groups selected data by one or more columns
 - Caution: Columns not part of aggregation must be listed in GROUP BY



Another example of Group By

INSERT OVERWRITE TABLE sectors

SELECT sector, count(DISTINCT ticker) as sectCount, sum(cap)

FROM companyInfo

GROUP BY sector;

https://cwiki.apache.org/confluence/display/Hive/LanguageManual+GroupBy

Remember

You can combine several functions in a single select

You should combine multiple functions in a single select

- let the optimizer work for you

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Simplifying Complex Queries

- Complex queries can become cumbersome
 - Imagine typing this several times for different orders

```
SELECT o.order_id, order_date, p.prod_id, brand, name
FROM orders o
JOIN order_details d
ON (o.order_id = d.order_id)
JOIN products p
ON (d.prod_id = p.prod_id)
WHERE o.order_id=6584288;
```

Creating Views

- Views in Hive are conceptually like a table, but backed by a query
 - You cannot directly add data to a view

```
CREATE VIEW order_info AS

SELECT o.order_id, order_date, p.prod_id, brand, name

FROM orders o

JOIN order_details d

ON (o.order_id = d.order_id)

JOIN products p

ON (d.prod_id = p.prod_id);
```

Our query is now greatly simplified

```
hive> SELECT * FROM order_info WHERE order_id=6584288;
```

Inspecting and Removing Views

Use DESCRIBE FORMATTED to see underlying query

```
hive> DESCRIBE FORMATTED order_info;
```

Use DROP VIEW to remove a view

```
hive> DROP VIEW order_info;
```

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Saving Query Output to a Table

- SELECT statements display their results on screen
- To send results to a Hive table, use INSERT OVERWRITE TABLE
 - Destination table must already exist
 - Existing contents will be deleted

```
hive> INSERT OVERWRITE TABLE ny_customers

SELECT * FROM customers

WHERE state = 'NY';
```

INSERT INTO TABLE adds records without first deleting existing data

```
hive> INSERT INTO TABLE ny_customers

SELECT * FROM customers

WHERE state = 'NJ' OR state = 'CT';
```

Creating Tables Based on Existing Data

- Hive supports creating a table based on a SELECT statement
 - Often know as 'Create Table As Select' (CTAS)

```
CREATE TABLE ny_customers AS

SELECT cust_id, fname, lname FROM customers

WHERE state = 'NY';
```

- Column definitions are derived from the existing table
- Column names are inherited from the existing names
 - Use aliases in the SELECT statement to specify new names

Writing Output to a Filesystem

You can save output to a file in HDFS

```
hive> INSERT OVERWRITE DIRECTORY '/dualcore/ny/'
SELECT * FROM customers
WHERE state = 'NY';
```

Add LOCAL to store results to local disk instead

```
hive> INSERT OVERWRITE LOCAL DIRECTORY '/home/bob/ny/'
SELECT * FROM customers
WHERE state = 'NY';
```

Both produce text files delimited by Ctrl-A characters

Writing Output to HDFS, Specifying Format

- To write the files to HDFS with a user-specified format:
 - Create an external table in the required format
 - Use INSERT OVERWRITE TABLE

```
hive> CREATE EXTERNAL TABLE ny customers
        (cust id INT,
         fname STRING,
         lname STRING)
        ROW FORMAT DELIMITED
        FIELDS TERMINATED BY ','
        STORED AS TEXTFILE
        LOCATION '/dualcore/nydata';
hive> INSERT OVERWRITE TABLE ny customers
        SELECT cust id, fname, lname
        FROM customers WHERE state = 'NY';
```

Optimization: Multi-Table Insert (1)

We just saw that you can save output to an HDFS file

```
INSERT OVERWRITE DIRECTORY 'ny_customers'
SELECT cust_id, fname, lname
FROM customers WHERE state = 'NY';
```

This query could also be written as follows

```
FROM customers c

INSERT OVERWRITE DIRECTORY 'ny_customers'

SELECT cust_id, fname, lname WHERE state='NY';
```

Optimization: Multi-Table Insert (2)

- We sometimes need to extract data to multiple tables
 - Hive SELECT queries can take a long time to complete
- Hive allows us to do this with a single query
 - Much more efficient than using multiple queries
- The following example demonstrates multi-table insert
 - Result is two directories in HDFS

```
FROM customers c

INSERT OVERWRITE DIRECTORY 'ny_names'

SELECT fname, lname WHERE state = 'NY';

INSERT OVERWRITE DIRECTORY 'ny_count'

SELECT count(DISTINCT cust_id)

WHERE state = 'NY';
```

Optimization: Multi-Table Insert (3)

The following query produces the same result

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Viewing the Execution Plan

- How can you tell how Hive will execute a query?
 - Does it read only metadata?
 - Can it return data directly from HDFS?
 - Will it require a reduce phase or multiple MapReduce jobs?
- Prefix your query with EXPLAIN to view Hive's execution plan

```
hive> EXPLAIN SELECT * FROM customers;
```

- The output of EXPLAIN can be very long and complex
 - Fully understanding it requires in-depth knowledge of MapReduce
 - We will cover the basics here...

EXPLAIN: understanding the Query Plan (1)

The query plan contains three main sections

- Abstract syntax tree details how Hive parsed query (excerpt below)
- The stage dependencies and plans are more useful to most users

```
hive> EXPLAIN CREATE TABLE cust by zip AS
      SELECT zipcode, COUNT(cust id) AS num
      FROM customers GROUP BY zipcode;
ABSTRACT SYNTAX TREE:
    (TOK CREATETABLE (TOK TABNAME cust by zip)
STAGE DEPENDENCIES:
    ... (excerpt shown on next slide)
STAGE PLANS:
        (excerpt shown on upcoming slide)
```

EXPLAIN: understanding the Query Plan (2)

- Our query has three stages
- Dependencies define order
 - Stage-1 runs first
 - Stage-0 runs next
 - Stage-2 runs last

```
ABSTRACT SYNTAX TREE:
 ... (shown on previous slide)
STAGE DEPENDENCIES:
  Stage-1 is a root stage
  Stage-0 depends on stages: Stage-1
  Stage-2 depends on stages: Stage-0
STAGE PLANS:
 ... (shown on next slide)
```

EXPLAIN: understanding the Query Plan (3)

Stage-1: MapReduce job

Map phase

- Read customers table
- Selects zipcode and cust_id columns

Reduce phase

- Group by zipcode
- Count cust_id

```
STAGE PLANS:
  Stage: Stage-1
    Map Reduce
      Alias -> Map Operator Tree:
        TableScan
          alias: customers
            Select Operator
                zipcode, cust id
      Reduce Operator Tree:
        Group By Operator
          aggregations:
            expr: count(cust id)
          keys:
            expr: zipcode
```

EXPLAIN: understanding the Query Plan (4)

Stage-0: HDFS action

 Move previous stage's output to Hive's warehouse directory

Stage-2: Metastore action

- Create new table
- Has two columns

```
STAGE PLANS:
   Stage: Stage-1 (covered earlier)
   Stage: Stage-0
     Move Operator
       files:
         hdfs directory: true
         destination: (HDFS path...)
    Stage: Stage-2
      Create Table Operator:
        Create Table
         columns: zipcode string,
                  num bigint
         name: cust by zip
```

Bibliography

The following offer more information on topics discussed in this section

- HiveQL Language Manual
 - https://cwiki.apache.org/confluence/display/Hive/LanguageManual
- Hive Built-in Functions
 - https://cwiki.apache.org/confluence/display/Hive/LanguageManual+UDF

Hive transactions

As of Hive version **0.14.0**: INSERT...VALUES, UPDATE, and DELETE are now available with full ACID support.

INSERT ... VALUES Syntax:

```
INSERT INTO TABLE tablename [PARTITION (partcol1[=val1], partcol2[=val2] ...)] VALUES values_row [, values_row ...]
```

Where values_row is: (value [, value ...]) where a value is either null or any valid SQL literal

UPDATE Syntax:

```
UPDATE tablename SET column = value [, column = value ...] [WHERE expression]
```

DELETE Syntax:

DELETE FROM tablename [WHERE expression]

Hive transactions - references

Additionally, from the Hive Transactions doc:

If a table is to be used in ACID writes (insert, update, delete) then the table property "transactional" must be set on that table, starting with Hive 0.14.0. Without this value, inserts will be done in the old style; updates and deletes will be prohibited.

Hive DML reference:

https://cwiki.apache.org/confluence/display/Hive/LanguageManual+DML

Hive Transactions reference:

https://cwiki.apache.org/confluence/display/Hive/Hive+Transactions