

Contents

1	準備	2			
1.1	Caps Lock を Control に変更	2			
1.2	init.el	2			
1.3	tpl.cpp	2			
1.4	get input	2			
1.5	alias	2			
2	文字列	2			
2.1	マッチング	2			
2.1.1	複数文字列マッチング (Aho-Corasick 法)	2			
2.2	Suffix Array	3			
2.3	回文長 (Manacher)	3			
3	グラフ	3			
3.1	強連結成分分解	3			
3.1.1	関節点	3			
3.1.2	橋	4			
3.1.3	強連結成分分解	4			
3.1.4	無向中国人郵便配達問題	5			
3.1.5	全点対間最短路 (Johnson)	5			
3.1.6	無向グラフの全域最小カット	5			
3.2	フロー	5			
3.2.1	最大流	5			
3.2.2	二部マッチング	6			
3.2.3	最小費用流	6			
3.2.4	Gomory-Hu 木	7			
3.3	木	7			
3.3.1	木の直径	7			
3.3.2	最小全域木	7			
3.3.3	最小全域有向木	8			
3.3.4	最小シュタイナー木	8			
3.3.5	木の同型性判定	8			
3.3.6	HL 分解	9			
3.4	彩色数	9			
3.4.1	包除原理	9			
3.4.2	極大独立集合	10			
3.5		10			
4	数学	10			
4.1	整数	10			
4.1.1	剰余	10			
4.1.2	剰余積	11			
4.1.3	離散対数問題	11			
4.1.4	カタラン数	11			
4.1.5	乱数 (xor shift)	11			
4.1.6	確率的素数判定 (Miller-Rabin 法)	11			
4.2	多項式	12			
4.2.1	FFT(complex)	12			
			4.2.2	FFT(modulo)	12
			4.2.3	積 (FMT)	12
			4.2.4	逆元 (FMT)	12
			4.2.5	平方根 (FMT)	12
			4.3	行列	13
			4.3.1	線形方程式の解 (Givens 消去法)	13
			4.4	割り当て問題	13
			4.4.1	ハンガリアン法	13
5	幾何	14			
5.1	点	14			
5.2	直線と線分	14			
5.3	円	14			
5.4	多角形	15			
5.4.1	凸包	15			
5.4.2	最近点对	15			
5.4.3	点-多角形包含判定	16			
5.4.4	凸多角形の共通部分	16			
5.4.5	凸多角形の直径	16			
5.4.6	ドロネー三角形分割 (逐次添加法)	16			
6	データ構造	17			
6.1	Union-Find 木	17			
6.2	Meldable Heap	17			
6.3	Binary-Indexed-Tree	17			
6.4	Segment Tree	18			
6.5	Sparse table	18			
6.6	RBST	18			
6.7	永続 RBST	18			
6.8	赤黒木	19			
6.9	永続赤黒木	20			
6.10	wavelet 行列	21			
	6.10.1 完備辞書	21			
	6.10.2 wavelet 行列	22			
7	その他	23			
7.1	ビジュアライザ	23			

1 準備

1.1 Caps Lock を Control に変更

2 つ

1. 変更

```
1 setxkbmap -option ctrl:nocaps;
```

元に戻す

```
1 setxkbmap -option;
```

2. 上でダメな場合

```
1 xmodmap -e 'remove Lock = Caps_Lock';
2 xmodmap -e 'add Control = Caps_Lock';
3 xmodmap -e 'keysym Caps_Lock = Control_L';
```

1.2 init.el

linum は emacs24 のみ

```
1 (keyboard-translate ?\C-h ?\C-?)
2 (global-linum-mode t)
3 (setq linum-format "%4d ")
```

1.3 tpl.cpp

```
1 #include <bits/stdc++.h>
2 using namespace std;
3
4 #define rep(i,n) repi(i,0,n)
5 #define repi(i,a,b) for(int i=int(a);i<int(b);++i)
6 #define repit(it,u) for(auto it=begin(u);it!=end(u);++it)
7 #define all(u) begin(u),end(u)
8 #define uniq(u) (u).erase(unique(all(u)),end(u))
9 #define ll long
10 #define long int64_t
11 #define mp make_pair
12 #define pb push_back
13 #define eb emplace_back
14
15 bool input()
16 {
17     return true;
18 }
19
20 void solve()
21 {
22 }
23 }
24
25 int main()
26 {
27     cin.tie(0);
```

```
28 ios_base::sync_with_stdio(false);
29
30 while (input()) solve();
31 }
```

1.4 get input

```
1 wget -r http://(url of sample input)
```

1.5 alias

```
1 alias g++='g++ -g -O2 -std=gnu++0x -Wl,-stack_size,64000000';
2 alias emacs='emacs -nw';
```

2 文字列

2.1 マッチング

2.1.1 複数文字列マッチング (Aho-Corasick 法)

$O(N + M)$

```
1 const int C = 128;
2
3 struct pma_node {
4     pma_node *next[C]; // use next[0] as failure link
5     vector<int> match;
6     pma_node() { fill(next, next + C, (pma_node *) NULL); }
7     ~pma_node() { rep(i, C) if (next[i] != NULL) delete next[i]; }
8 };
9
10 pma_node *construct_pma(const vector<string>& pat) {
11     pma_node *const root = new pma_node();
12     root->next[0] = root;
13     // construct trie
14     rep(i, pat.size()) {
15         const string& s = pat[i];
16         pma_node *now = root;
17         for (const char c : s) {
18             if (now->next[int(c)] == NULL) now->next[int(c)] = new pma_node();
19             now = now->next[int(c)];
20         }
21         now->match.pb(i);
22     }
23     // make failure links by BFS
24     queue<pma_node *> q;
25     repi(i, 1, C) {
26         if (root->next[i] == NULL) root->next[i] = root;
27         else {
28             root->next[i]->next[0] = root;
29             q.push(root->next[i]);
30         }
31     }
32     while (not q.empty()) {
33         auto now = q.front();
34         q.pop();
35         repi(i, 1, C) if (now->next[i] != NULL) {
```

```

36         auto next = now->next[0];
37         while (next->next[i] == NULL) next = next->next[0];
38         now->next[i]->next[0] = next->next[i];
39         vector<int> tmp;
40         set_union(all(now->next[i]->match), all(next->next[i]->match), back_inserter
41                 (tmp));
42         now->next[i]->match = tmp;
43         q.push(now->next[i]);
44     }
45     return root;
46 }
47
48 void match(pma_node*& now, const string s, vector<int>& ret) {
49     for (const char c : s) {
50         while (now->next[int(c)] == NULL) now = now->next[0];
51         now = now->next[int(c)];
52         for (const int e : now->match) ret[e] = true;
53     }
54 }

```

2.2 Suffix Array

find_string(): $O(|T| \log |S|)$

S 中に T が含まれないなら -1, 含まれるならその先頭。

LCS(): $O(|S| + |T|)$

最長共通部分文字列。(先頭, 長さ) を返す。

```

1 // verify
2 // sa: http://www.spoj.com/problems/SARRAY/
3 // lcp: http://www.spoj.com/problems/SUBLEX/
4
5 int n, k;
6 vector<int> rnk, tmp, sa, lcp;
7
8 bool compare_sa(int i, int j) {
9     if(rnk[i] != rnk[j]) return rnk[i] < rnk[j];
10    else {
11        int ri = i + k <= n ? rnk[i+k] : -1;
12        int rj = j + k <= n ? rnk[j+k] : -1;
13        return ri < rj;
14    }
15 }
16
17 void construct_sa(const string &s) {
18     n = s.size();
19     rnk.assign(n+1, 0);
20     tmp.assign(n+1, 0);
21     sa.assign(n+1, 0);
22     lcp.assign(n+1, 0);
23     for(int i = 0; i <= n; i++) {
24         sa[i] = i;
25         rnk[i] = i < n ? s[i] : -1;
26     }
27     for(k = 1; k <= n; k*=2) {
28         sort(sa.begin(), sa.end(), compare_sa);
29         tmp[sa[0]] = 0;
30         for(int i = 1; i <= n; i++) {
31             tmp[sa[i]] = tmp[sa[i-1]] + (compare_sa(sa[i-1], sa[i]) ? 1 : 0);
32         }
33         for(int i = 0; i <= n; i++) {
34             rnk[i] = tmp[i];
35         }
36     }
37 }

```

```

36     }
37 }
38
39 void construct_lcp(const string &s) {
40     for(int i = 0; i <= n; i++) rnk[sa[i]] = i;
41     int h = 0;
42     lcp[0] = 0;
43     for(int i = 0; i < n; i++) {
44         int j = sa[rnk[i] - 1];
45         if(h > 0) h--;
46         for(; j + h < n && i + h < n; h++) {
47             if(s[j+h] != s[i+h]) break;
48         }
49         lcp[rnk[i] - 1] = h;
50     }
51 }

```

2.3 回文長 (Manacher)

$O(N)$

各文字を中心とした時の回文の最長の半径。

偶数長の回文はダミーを挟むことで求められている。

```

1 vector<int> manacher(const string &s) {
2     int n = s.size()*2;
3     vector<int> rad.assign(n, 0);
4     for (int i = 0, j = 0, k; i < n; i += k, j = max(j-k, 0)) {
5         while (i-j >= 0 && i+j+1 < n && s[(i-j)/2] == s[(i+j+1)/2]) ++j;
6         rad[i] = j;
7         for (k = 1; i-k >= 0 && rad[i]-k >= 0 && rad[i-k] != rad[i]-k; ++k)
8             rad[i+k] = min(rad[i-k], rad[i]-k);
9     }
10    return rad;
11 }

```

3 グラフ

```

1 struct edge {
2     int to; long w;
3     edge(int to, long w) : to(to), w(w) {}
4 };
5 typedef vector<vector<edge>> graph;
6
7 graph rev(const graph& G) {
8     const int n = G.size();
9     graph ret(n);
10    rep(i, n) for (const auto& e : G[i]) {
11        ret[e.to].eb(i, e.w);
12    }
13    return ret;
14 }

```

3.1 強連結成分分解

3.1.1 関節点

$O(E)$

ある関節点 u がグラフを k 個に分割するとき art には $k-1$ 個の u が含まれる. 不要な場合は `unique` を忘れないこと.

```
1  typedef vector<vector<int> > graph;
2
3  class articulation {
4      const int n;
5      graph G;
6      int cnt;
7      vector<int> num, low, art;
8      void dfs(int v) {
9          num[v] = low[v] = ++cnt;
10         for (int nv : G[v]) {
11             if (num[nv] == 0) {
12                 dfs(nv);
13                 low[v] = min(low[v], low[nv]);
14                 if ((num[v] == 1 and num[nv] != 2) or
15                     (num[v] != 1 and low[nv] >= num[v])) {
16                     art[v] = true;
17                 }
18             } else {
19                 low[v] = min(low[v], num[nv]);
20             }
21         }
22     }
23 public:
24     articulation(const graph& G) : n(G.size()), G(G), cnt(0), num(n), low(n), art(n) {
25         rep(i, n) if (num[i] == 0) dfs(i);
26     }
27     vector<int> get() {
28         return art;
29     }
30 };
```

3.1.2 橋

$O(V + E)$

```
1  typedef vector<vector<int> > graph;
2
3  class bridge {
4      const int n;
5      graph G;
6      int cnt;
7      vector<int> num, low, in;
8      stack<int> stk;
9      vector<pair<int, int> > brid;
10     vector<vector<int> > comp;
11     void dfs(int v, int p) {
12         num[v] = low[v] = ++cnt;
13         stk.push(v), in[v] = true;
14         for (const int nv : G[v]) {
15             if (num[nv] == 0) {
16                 dfs(nv, v);
17                 low[v] = min(low[v], low[nv]);
18             } else if (nv != p and in[nv]) {
19                 low[v] = min(low[v], num[nv]);
20             }
21         }
22         if (low[v] == num[v]) {
23             if (p != n) brid.eb(min(v, p), max(v, p));
24             comp.eb();
25             int w;
26             do {
```

```
27         w = stk.top();
28         stk.pop(), in[w] = false;
29         comp.back().pb(w);
30     } while (w != v);
31 }
32 }
33 public:
34     bridge(const graph& G) : n(G.size()), G(G), cnt(0), num(n), low(n), in(n) {
35         rep(i, n) if (num[i] == 0) dfs(i, n);
36     }
37     vector<pair<int, int> > get() {
38         return brid;
39     }
40     vector<vector<int> > components() {
41         return comp;
42     }
43 };
```

3.1.3 強連結成分分解

$O(V + E)$

```
1  typedef vector<vector<int> > graph;
2
3  class scc {
4      const int n;
5      graph G;
6      int cnt;
7      vector<int> num, low, in;
8      stack<int> stk;
9      vector<vector<int> > comp;
10     void dfs(int v) {
11         num[v] = low[v] = ++cnt;
12         stk.push(v), in[v] = true;
13         for (const int nv : G[v]) {
14             if (num[nv] == 0) {
15                 dfs(nv);
16                 low[v] = min(low[v], low[nv]);
17             } else if (in[nv]) {
18                 low[v] = min(low[v], num[nv]);
19             }
20         }
21         if (low[v] == num[v]) {
22             comp.eb();
23             int w;
24             do {
25                 w = stk.top();
26                 stk.pop(), in[w] = false;
27                 comp.back().pb(w);
28             } while (w != v);
29         }
30     }
31 public:
32     scc(const graph& G) : n(G.size()), G(G), cnt(0), num(n), low(n), in(n) {
33         rep(i, n) if (num[i] == 0) dfs(i);
34     }
35     vector<vector<int> > components() {
36         return comp;
37     }
38 };
```

3.1.4 無向中国人郵便配達問題

$O(m \log n + o^2 2^o)$, $-O2$ で $o \leq 18$ 程度が限界

```
1 long chinesePostman(const graph &g) {
2     long total = 0;
3     vector<int> odds;
4     rep(u, g.size()) {
5         for(auto &e: g[u]) total += e.w;
6         if (g[u].size() % 2) odds.push_back(u);
7     }
8     total /= 2;
9     int n = odds.size(), N = 1 << n;
10    int w[n][n]; // make odd vertices graph
11    rep(u,n) {
12        int s = odds[u]; // dijkstra's shortest path
13        vector<int> dist(g.size(), 1e9); dist[s] = 0;
14        vector<int> prev(g.size(), -2);
15        priority_queue<edge> Q;
16        Q.push( edge(-1, s, 0) );
17        while (!Q.empty()) {
18            edge e = Q.top(); Q.pop();
19            if (prev[e.to] != -2) continue;
20            prev[e.to] = e.src;
21            for(auto &f: g[e.to]) {
22                if (dist[f->to] > e.w+f->w) {
23                    dist[f->to] = e.w+f->w;
24                    Q.push(edge(f->src, f->to, e.w+f->w));
25                }
26            }
27        }
28        rep(v,n) w[u][v] = dist[odds[v]];
29    }
30    long best[N]; // DP for general matching
31    rep(S,N) best[S] = INF;
32    best[0] = 0;
33
34    for (int S = 0; S < N; ++S)
35        for (int i = 0; i < n; ++i)
36            if (!(S&(1<<i)))
37                for (int j = i+1; j < n; ++j)
38                    if (!(S&(1<<j)))
39                        best[S|(1<<i)|(1<<j)] = min(best[S|(1<<i)|(1<<j)], best[S]+w[i][j]);
40
41    return total + best[N-1];
42 }
```

3.1.5 全点对間最短路 (Johnson)

$O(\max(V, E) \log V, V^2)$

```
1 bool shortest_path(const graph &g, vector<vector<int>> &dist, vector<vector<int>> &
2     prev) {
3     int n = g.size();
4     vector<int> h(n+1);
5     rep(k,n) rep(i,n) for(auto &e: g[i]) {
6         if (h[e.to] > h[e.from] + e->w) {
7             h[e.to] = h[e.from] + e->w;
8             if (k == n-1) return false; // negative cycle
9         }
10    }
11    dist.assign(n, vector<int>(n, 1e9));
12    prev.assign(n, vector<int>(n, -2));
13    rep(s, n) {
```

```
13     priority_queue<edge> q;
14     q.push(edge(s, s, 0));
15     while (!q.empty()) {
16         edge e = q.top(); q.pop();
17         if (prev[s][e.dst] != -2) continue;
18         prev[s][e.to] = e.from;
19         for(auto &f: g[e.to]) {
20             if (dist[s][f.to] > e.w + f->w) {
21                 dist[s][f.to] = e.w + f->w;
22                 q.push(edge(f->from, f.to, e.w + f->w));
23             }
24         }
25     }
26     rep(u, n) dist[s][u] += h[u] - h[s];
27 }
28 }
29
30 vector<int> build_path(const vector<vector<int>> &prev, int s, int t) {
31     vector<int> path;
32     for (int u = t; u >= 0; u = prev[s][u])
33         path.push_back(u);
34     reverse(begin(path), end(path));
35     return path;
36 }
```

3.1.6 無向グラフの全域最小カット

$O(V^3)$

```
1 int minimum_cut(const graph &g) {
2     int n = g.size();
3     vector< vector<int>> > h(n, vector<int>(n)); // make adj. matrix
4     rep(u,n) for(auto &e: g[u]) h[e.src][e.dst] += e.weight;
5     vector<int> V(n); rep(u, n) V[u] = u;
6
7     int cut = 1e9;
8     for(int m = n; m > 1; m--) {
9         vector<int> ws(m, 0);
10        int u, v;
11        int w;
12        rep(k, m) {
13            u = v; v = max_element(ws.begin(), ws.end())-ws.begin();
14            w = ws[v]; ws[v] = -1;
15            rep(i, m) if (ws[i] >= 0) ws[i] += h[V[v]][V[i]];
16        }
17        rep(i, m) {
18            h[V[i]][V[u]] += h[V[i]][V[v]];
19            h[V[u]][V[i]] += h[V[v]][V[i]];
20        }
21        V.erase(V.begin()+v);
22        cut = min(cut, w);
23    }
24    return cut;
25 }
```

3.2 フロー

3.2.1 最大流

$O(EV^2)$

```

1  const int inf = 1e9;
2  struct edge {
3      int to, cap, rev;
4      edge(int to, int cap, int rev) : to(to), cap(cap), rev(rev) {}
5  };
6  typedef vector<vector<edge>> > graph;
7
8  void add_edge(graph& G, int from, int to, int cap) {
9      G[from].eb(to, cap, G[to].size());
10     G[to].eb(from, 0, G[from].size() - 1);
11 }
12
13 class max_flow {
14     const int n;
15     graph& G;
16     vector<int> level, iter;
17     void bfs(int s, int t) {
18         level.assign(n, -1);
19         queue<int> q;
20         level[s] = 0, q.push(s);
21         while (not q.empty()) {
22             const int v = q.front();
23             q.pop();
24             if (v == t) return;
25             for (const auto& e : G[v]) {
26                 if (e.cap > 0 and level[e.to] < 0) {
27                     level[e.to] = level[v] + 1;
28                     q.push(e.to);
29                 }
30             }
31         }
32     }
33     int dfs(int v, int t, int f) {
34         if (v == t) return f;
35         for (int& i = iter[v]; i < (int) G[v].size(); ++i) {
36             edge& e = G[v][i];
37             if (e.cap > 0 and level[v] < level[e.to]) {
38                 const int d = dfs(e.to, t, min(f, e.cap));
39                 if (d > 0) {
40                     e.cap -= d, G[e.to][e.rev].cap += d;
41                     return d;
42                 }
43             }
44         }
45         return 0;
46     }
47 public:
48     max_flow(graph& G) : n(G.size()), G(G) {}
49     int calc(int s, int t) {
50         int ret = 0, d;
51         while (bfs(s, t), level[t] >= 0) {
52             iter.assign(n, 0);
53             while ((d = dfs(s, t, inf)) > 0) ret += d;
54         }
55         return ret;
56     }
57 };

```

3.2.2 二部マッチング

$O(EV)$

```

1  int V;
2  vector<int> G[MAX_V];

```

```

3  int match[MAX_V];
4  bool used[MAX_V];
5
6  void add_edge(int u, int v){
7      G[u].push_back(v);
8      G[v].push_back(u);
9  }
10
11 bool dfs(int v){
12     used[v] = 1;
13     rep(i, G[v].size()){
14         int u = G[v][i], w = match[u];
15         if (w < 0 || !used[w] && dfs(w)){
16             match[v] = u;
17             match[u] = v;
18             return 1;
19         }
20     }
21     return 0;
22 }
23
24 int bi_matching(){
25     int res = 0;
26     memset(match, -1, sizeof(match));
27     rep(v, V) if (match[v] < 0){
28         memset(used, 0, sizeof(used));
29         if (dfs(v)) res++;
30     }
31     return res;
32 }

```

3.2.3 最小費用流

$O(FE \log V)$

```

1  const int inf = 1e9;
2  struct edge {
3      int to, cap, cost, rev;
4      edge(int to, int cap, int cost, int rev) : to(to), cap(cap), cost(cost), rev(rev) {}
5  };
6  typedef vector<vector<edge>> > graph;
7
8  void add_edge(graph& G, int from, int to, int cap, int cost) {
9      G[from].eb(to, cap, cost, G[to].size());
10     G[to].eb(from, 0, -cost, G[from].size() - 1);
11 }
12
13 int min_cost_flow(graph& G, int s, int t, int f) {
14     const int n = G.size();
15     struct state {
16         int v, d;
17         state(int v, int d) : v(v), d(d) {}
18         bool operator <(const state& t) const { return d > t.d; }
19     };
20
21     int ret = 0;
22     vector<int> h(n, 0), dist, prev(n), prev_e(n);
23     while (f > 0) {
24         dist.assign(n, inf);
25         priority_queue<state> q;
26         dist[s] = 0, q.emplace(s, 0);
27         while (not q.empty()) {
28             const int v = q.top().v;
29             const int d = q.top().d;
30             q.pop();

```

```

31         if (dist[v] < d) continue;
32         rep(i, G[v].size()) {
33             const edge& e = G[v][i];
34             if (e.cap > 0 and dist[e.to] > dist[v] + e.cost + h[v] - h[e.to]) {
35                 dist[e.to] = dist[v] + e.cost + h[v] - h[e.to];
36                 prev[e.to] = v, prev_e[e.to] = i;
37                 q.emplace(e.to, dist[e.to]);
38             }
39         }
40     }
41     if (dist[t] == inf) return -1;
42     rep(i, n) h[i] += dist[i];
43
44     int d = f;
45     for (int v = t; v != s; v = prev[v]) {
46         d = min(d, G[prev[v]][prev_e[v]].cap);
47     }
48     f -= d, ret += d * h[t];
49     for (int v = t; v != s; v = prev[v]) {
50         edge& e = G[prev[v]][prev_e[v]];
51         e.cap -= d, G[v][e.rev].cap += d;
52     }
53 }
54 return ret;
55 }

```

3.2.4 Gomory-Hu 木

$O(VMAXFLOW)$

```

1  #define RESIDUE(s,t) (capacity[s][t]-flow[s][t])
2  graph cutTree(const graph &g) {
3      int n = g.size();
4      Matrix capacity(n, Array(n)), flow(n, Array(n));
5      rep(u,n) for(auto &e: g[u]) capacity[e.from][e.to] += e.w;
6
7      vector<int> p(n), prev;
8      vector<int> w(n);
9      for (int s = 1; s < n; ++s) {
10         int t = p[s]; // max-flow(s, t)
11         rep(i,n) rep(j,n) flow[i][j] = 0;
12         int total = 0;
13         while (1) {
14             queue<int> Q; Q.push(s);
15             prev.assign(n, -1); prev[s] = s;
16             while (!Q.empty() && prev[t] < 0) {
17                 int u = Q.front(); Q.pop();
18                 for(auto &e: g[u]) if (prev[e.to] < 0 && RESIDUE(u, e.to) > 0) {
19                     prev[e.to] = u;
20                     Q.push(e.to);
21                 }
22             }
23             if (prev[t] < 0) goto esc;
24             int inc = 1e9;
25             for (int j = t; prev[j] != j; j = prev[j])
26                 inc = min(inc, RESIDUE(prev[j], j));
27             for (int j = t; prev[j] != j; j = prev[j])
28                 flow[prev[j]][j] += inc, flow[j][prev[j]] -= inc;
29             total += inc;
30         }
31     esc: w[s] = total; // make tree
32     rep(u, n) if (u != s && prev[u] != -1 && p[u] == t)
33         p[u] = s;
34     if (prev[p[t]] != -1)
35         p[s] = p[t], p[t] = s, w[s] = w[t], w[t] = total;

```

```

36     }
37     graph T(n); // (s, p[s]) is a tree edge of weight w[s]
38     rep(s, n) if (s != p[s]) {
39         T[ s ].push_back( Edge(s, p[s], w[s]) );
40         T[p[s]].push_back( Edge(p[s], s, w[s]) );
41     }
42     return T;
43 }
44
45 // Gomory-Hu tree を用いた最大流  $O(n)$ 
46 int max_flow(const graph &T, int u, int t, int p = -1, int w = 1e9) {
47     if (u == t) return w;
48     int d = 1e9;
49     for(auto &e: T[u]) if (e.to != p)
50         d = min(d, max_flow(T, e.to, t, u, min(w, e.w)));
51     return d;
52 }

```

3.3 木

3.3.1 木の直径

ある点 (どこでもよい) から一番遠い点 a を求める. 点 a から一番遠い点までの距離がその木の直径になる.

3.3.2 最小全域木

```

1  struct mst_edge {
2      int u, v; long w;
3      mst_edge(int u, int v, long w) : u(u), v(v), w(w) {}
4      bool operator <(const mst_edge& t) const { return w < t.w; }
5      bool operator >(const mst_edge& t) const { return w > t.w; }
6  };
7
8  graph kruskal(const graph& G) {
9      const int n = G.size();
10     vector<mst_edge> E;
11     rep(i, n) for (const auto& e : G[i]) {
12         if (i < e.to) E.eb(i, e.to, e.w);
13     }
14     sort(all(E));
15
16     graph T(n);
17     disjoint_set uf(n);
18     for (const auto& e : E) {
19         if (not uf.same(e.u, e.v)) {
20             T[e.u].eb(e.v, e.w);
21             T[e.v].eb(e.u, e.w);
22             uf.merge(e.u, e.v);
23         }
24     }
25     return T;
26 }
27
28 graph prim(const vector<vector<long>> &A, int s = 0) {
29     const int n = A.size();
30     graph T(n);
31     vector<int> done(n);
32     priority_queue<mst_edge, vector<mst_edge>, greater<mst_edge>> q;
33     q.emplace(-1, s, 0);
34     while (not q.empty()) {
35         const auto e = q.top();

```

```

36     q.pop();
37     if (done[e.v]) continue;
38     done[e.v] = 1;
39     if (e.u >= 0) {
40         T[e.u].eb(e.v, e.w);
41         T[e.v].eb(e.u, e.w);
42     }
43     rep(i, n) if (not done[i]) {
44         q.emplace(e.v, i, A[e.v][i]);
45     }
46 }
47 return T;
48 }

```

3.3.3 最小全域有向木

$O(VE)$

```

1 void visit(Graph &h, int v, int s, int r,
2     vector<int> &no, vector< vector<int> > &comp,
3     vector<int> &prev, vector< vector<int> > &next, vector<int> &mcost,
4     vector<int> &mark, int &cost, bool &found) {
5     const int n = h.size();
6     if (mark[v]) {
7         vector<int> temp = no;
8         found = true;
9         do {
10             cost += mcost[v];
11             v = prev[v];
12             if (v != s) {
13                 while (comp[v].size() > 0) {
14                     no[comp[v].back()] = s;
15                     comp[s].push_back(comp[v].back());
16                     comp[v].pop_back();
17                 }
18             }
19             while (v != s);
20             for(auto &j: comp[s]) if (j != r) for(auto &e: h[j])
21                 if (no[e.from] != s) e.w -= mcost[temp[j]];
22         }
23         mark[v] = true;
24         for(auto &i: next[v]) if (no[i] != no[v] && prev[no[i]] == v)
25             if (!mark[no[i]] || i == s)
26                 visit(h, i, s, r, no, comp, prev, next, mcost, mark, cost, found);
27     }
28     int minimum_spanning_arborescence(const graph &g, int r) {
29         const int n = g.size();
30         graph h(n);
31         rep(u,n) for(auto &e: g[u]) h[e.to].push_back(e);
32
33         vector<int> no(n);
34         vector< vector<int> > comp(n);
35         rep(u, n) comp[u].push_back(no[u] = u);
36
37         for (int cost = 0; ; ) {
38             vector<int> prev(n, -1);
39             vector<int> mcost(n, INF);
40
41             rep(j,n) if (j != r) for(auto &e: g[j])
42                 if (no[e.from] != no[j])
43                     if (e.w < mcost[no[j]])
44                         mcost[no[j]] = e.w, prev[no[j]] = no[e.from];
45
46             vector< vector<int> > next(n);
47             rep(u,n) if (prev[u] >= 0)

```

```

48         next[prev[u]].push_back(u);
49
50         bool stop = true;
51         vector<int> mark(n);
52         rep(u,n) if (u != r && !mark[u] && !comp[u].empty()) {
53             bool found = false;
54             visit(h, u, u, r, no, comp, prev, next, mcost, mark, cost, found);
55             if (found) stop = false;
56         }
57         if (stop) {
58             rep(u,n) if (prev[u] >= 0) cost += mcost[u];
59             return cost;
60         }
61     }
62 }

```

3.3.4 最小シュタイナー木

$O(4^{|T|}V)$

g は無向グラフの隣接行列. T は使いたい頂点の集合.

```

1 int minimum_steiner_tree(vi &T, vvi &g){
2     int n = g.size(), t = T.size();
3     if(t <= 1) return 0;
4     vvi d(g); // all-pair shortest
5     rep(k,n)rep(i,n)rep(j,n) //Warshall Floyd
6         d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
7
8     int opt[1 << t][n];
9     rep(S,1<<t) rep(x,n)
10         opt[S][x] = INF;
11
12     rep(p,t) rep(q,n) // trivial case
13         opt[1 << p][q] = d[T[p]][q];
14
15     repi(S,1,1<<t){ // DP step
16         if(!(S & (S-1))) continue;
17         rep(p,n) rep(E,S)
18             if((E | S) == S)
19                 opt[S][p] = min(opt[S][p], opt[E][p] + opt[S-E][p]);
20         rep(p,n) rep(q,n)
21             opt[S][p] = min(opt[S][p], opt[S][q] + d[p][q]);
22     }
23
24     int ans = INF;
25     rep(S,1<<t) rep(q,n)
26         ans = min(ans, opt[S][q] + opt[((1<<t)-1)-S][q]);
27     return ans;
28 }

```

3.3.5 木の同型性判定

順序付き $O(n)$

順序なし $O(n \log n)$

```

1 // ordered
2 struct node {
3     vector<node*> child;
4 };
5 bool otreeIsomorphism(node *n, node *m) {
6     if (n->child.size() != m->child.size()) return false;

```



```

7     rep(i, n->child.size())
8         if (!otreeIsomorphism(n->child[i], m->child[i])) return false;
9     return true;
10 }
11
12 // not ordered
13 struct node {
14     vector<node*> child;
15     vector<int> code;
16 };
17 void code(node *n) {
18     int size = 1;
19     vector< pair<vector<int>, int> > codes;
20     rep(i, n->child.size()) {
21         code(n->child[i]);
22         codes.push_back( make_pair(n->child[i]->code, i) );
23         size += codes[i].first[0];
24     }
25     sort(codes.rbegin(), codes.rend()); // !reverse
26     n->code.push_back(size);
27     for (int i = 0; i < n->child.size(); ++i) {
28         swap(n->child[i], n->child[ codes[i].second ]);
29         n->code.insert(n->code.end(),
30             codes[i].first.begin(), codes[i].first.end());
31     }
32 }
33 bool utreeIsomorphism(node *n, node *m) {
34     code(n); code(m); return n->code == m->code;
35 }

```

3.3.6 HL 分解

```

1 namespace HLD {
2     const int N = 200010;
3     vector<vector<int>> chains, childs;
4     int V, dep[N], par[N], heavy[N], head[N], chain[N], id[N], size[N], q[N];
5
6     void calc_heavy() {
7         int root = -1;
8         childs.assign(V, vector<int>());
9         for(int v = 0; v < V; v++) {
10             size[v] = 0;
11             heavy[v] = -1;
12             if(par[v] < 0) root = v;
13             else childs[par[v]].push_back(v);
14         }
15         int l = 0, r = 0;
16         q[r++] = root;
17         while(l < r) {
18             int v = q[l++];
19             for(auto &w: childs[v]) {
20                 if(w == par[v]) continue;
21                 dep[w] = dep[v] + 1;
22                 q[r++] = w;
23             }
24         }
25         reverse(q, q+V);
26         for(int i = 1; i < V; i++) {
27             int v = q[i], &u = par[v];
28             size[u] += ++size[v];
29             if(heavy[u] == -1 or size[v] > size[heavy[u]]) heavy[u] = v;
30         }
31     }
32     void calc_chain() {

```

```

33     chains.clear();
34     int idx = 0;
35     for (int v = 0; v < V; v++) {
36         if(par[v] < 0 or heavy[par[v]] != v) {
37             chains.push_back(vector<int>());
38             for (int w = v; w != -1; w = heavy[w]) {
39                 chain[w] = idx;
40                 head[w] = v;
41                 id[w] = chains.back().size();
42                 chains.back().push_back(w);
43             }
44             idx++;
45         }
46     }
47 }
48 void make_par(const vector<vector<int>> &g, int root = 0) {
49     memset(par, -1, sizeof(par));
50     par[root] = 0;
51     int l = 0, r = 0;
52     q[r++] = root;
53     while(l < r) {
54         int v = q[l++];
55         for(const int &w: g[v]) if(par[w] < 0) q[r++] = w, par[w] = v;
56     }
57     par[root] = -1;
58 }
59 void build(const vector<vector<int>> &g, int root = 0) {
60     V = g.size();
61     make_par(g, root);
62     calc_heavy();
63     calc_chain();
64 }
65 int lca(int u, int v) {
66     while (chain[u] != chain[v]) {
67         if (dep[head[u]] > dep[head[v]]) swap(u, v);
68         v = par[head[v]];
69     }
70     return dep[u] < dep[v]? u: v;
71 }
72 }

```

3.4 彩色数

3.4.1 包除原理

$O(2^V V)$

$N[i] := i$ と隣接する頂点の集合 (i も含む)

```

1     const int MAX_V=16;
2     const int mod = 10009;
3     int N[MAX_V], I[1<<MAX_V], V;
4     inline int mpow(int a, int k){ return k==0? 1: k%2? a*mpow(a,k-1)%mod: mpow(a*a%mod,k/2);}
5
6     bool can(int k){
7         int res = 0;
8         rep(S, 1<<V){
9             if(__builtin_popcountll(S)%2) res -= mpow(I[S], k);
10            else res += mpow(I[S],k);
11        }
12        return (res%mod+mod)%mod;
13    }
14
15    int color_number(){

```

```

16     memset(I, 0, sizeof(I));
17     I[0] = 1;
18     repi(S, 1, 1 << V) {
19         int v = 0;
20         while (!S & (1 << v)) v++;
21         I[S] = I[S - (1 << v)] + I[S & (~N[v])];
22     }
23     int lb = 0, ub = V, mid;
24     while (ub - lb > 1) {
25         mid = (lb + ub) / 2;
26         if (can(mid)) ub = mid;
27         else lb = mid;
28     }
29     return ub;
30 }

```

3.4.2 極大独立集合

```

1  typedef vector<vector<int>> graph;
2  class maximal_indsets {
3      const int n;
4      const graph& G;
5      vector<vector<int>> ret;
6      vector<int> cur, exists, deg, block;
7      void erase(int v) {
8          if (exists[v]) {
9              exists[v] = false;
10             for (int nv : G[v]) --deg[nv];
11         }
12     }
13     void restore(int v) {
14         exists[v] = true;
15         for (int nv : G[v]) ++deg[nv];
16     }
17     void select(int v) {
18         cur.push_back(v);
19         ++block[v], erase(v);
20         for (int nv : G[v]) ++block[nv], erase(nv);
21     }
22     void unselect(int v) {
23         cur.pop_back();
24         --block[v], restore(v);
25         for (int nv : G[v]) {
26             if (--block[nv] == 0) restore(nv);
27         }
28     }
29     void dfs() {
30         int mn = n, v = -1;
31         rep(u, n) if (exists[u]) {
32             if (deg[u] < mn) mn = deg[u], v = u;
33         }
34         if (v == -1) {
35             ret.push_back(cur);
36         } else {
37             select(v), dfs(), unselect(v);
38             for (int nv : G[v]) {
39                 if (exists[nv]) select(nv), dfs(), unselect(nv);
40             }
41         }
42     }
43 public:
44     maximal_indsets(const graph& G) : n(G.size()), G(G), exists(n, true), deg(n), block(
45         n) {
46         rep(v, n) deg[v] = G[v].size();
47     }
48 }

```

```

46     dfs();
47 }
48     const vector<vector<int>>& get() const { return ret; }
49 };

```

3.5

4 数学

4.1 整数

4.1.1 剩余

```

1  // (x, y) s.t. a x + b y = gcd(a, b)
2  long extgcd(long a, long b, long& x, long& y) {
3      long g = a; x = 1, y = 0;
4      if (b != 0) g = extgcd(b, a % b, y, x), y -= (a / b) * x;
5      return g;
6  }
7
8  // repi(i, 2, n) mod_inv[i] = mod_inv[m % i] * (m - m / i) % m
9  long mod_inv(long a, long m) {
10     long x, y;
11     if (extgcd(a, m, x, y) != 1) return 0;
12     return (x % m + m) % m;
13 }
14
15 // a mod p where n! = a p^e in O(log_p n)
16 long mod_fact(long n, long p, long& e) {
17     const int P = 1000010;
18     static long fact[P] = {1};
19     static bool done = false;
20     if (not done) {
21         repi(i, 1, P) fact[i] = fact[i - 1] * i % p;
22         done = true;
23     }
24     e = 0;
25     if (n == 0) return 1;
26     long ret = mod_fact(n / p, p, e);
27     e += n / p;
28     if (n / p % 2) return ret * (p - fact[n % p]) % p;
29     return ret * fact[n % p] % p;
30 }
31
32 // nCk mod p
33 long mod_binom(long n, long k, long p) {
34     if (k < 0 or n < k) return 0;
35     long e1, e2, e3;
36     long a1 = mod_fact(n, p, e1);
37     long a2 = mod_fact(k, p, e2);
38     long a3 = mod_fact(n - k, p, e3);
39     if (e1 > e2 + e3) return 0;
40     return a1 * mod_inv(a2 * a3 % p, p) % p;
41 }
42
43 // a^b mod m
44 long mod_pow(long a, long b, long m) {
45     long ret = 1;
46     do {
47         if (b & 1) ret = ret * a % m;
48         a = a * a % m;
49     } while (b >= 1);
50     return ret;
51 }

```

```
51 }
```

4.1.2 剰余積

```
1 inline long mod_mul(long a, long b, long m) {
2     long ret = a * b - m * long(roundl((long double)(a) * b / m));
3     return ret < 0 ? ret + m : ret;
4 }
5 long mod_powl(long a, long b, long m) {
6     long ret = 1;
7     do {
8         if (b & 1) ret = mod_mul(ret, a, m);
9         a = mod_mul(a, a, m);
10    } while (b >>= 1);
11    return ret;
12 }
```

4.1.3 離散対数問題

```
1 long discrete_log(long a, long m) {
2     if (a == 0) return -1;
3     long b = long(sqrt(m)) + 1, t = 1;
4     unordered_map<long, long> mem;
5     for (int i = 0; i < b; ++i) {
6         mem[t] = i;
7         t = t * a % m;
8         if (t == 1) return i + 1;
9     }
10    long u = t;
11    for (int i = b; i < m; i += b) {
12        if (mem.find(mod_inverse(u, m)) != mem.end()) {
13            return mem[mod_inverse(u, m)] + i;
14        }
15        u = u * t % m;
16    }
17    return -1;
18 }
```

4.1.4 カタラン数

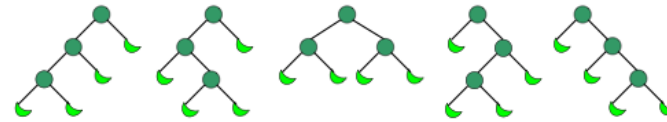
$n \leq 16$ 程度が限度. $n \geq 1$ について以下が成り立つ.

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n-1}$$

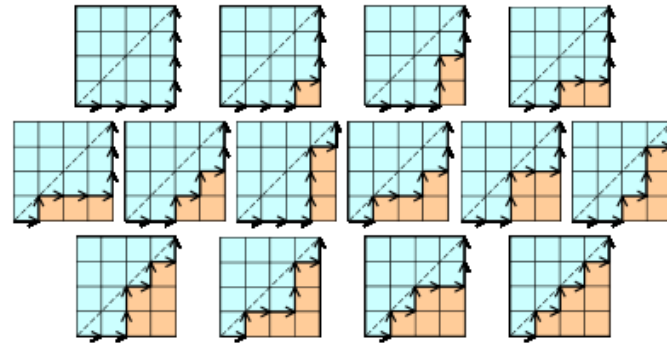
n が十分大きいとき, カタラン数は以下に近似できる.

$$C_n \approx \frac{4^n}{n^{3/2} \sqrt{\pi}}$$

() を正しく並べる方法, 二分木, 格子状の経路の数え上げ, 平面グラフの交差などに使われる.
 $C_3 = 5$



$C_4 = 14$



4.1.5 乱数 (xor shift)

周期は $2^{128} - 1$

```
1 unsigned xorshift() {
2     static unsigned x = 123456789;
3     static unsigned y = 362436069;
4     static unsigned z = 521288629;
5     static unsigned w = 88675123;
6     unsigned t;
7     t = x ^ cb ^ 86 (x << 11);
8     x = y; y = z; z = w;
9     return w = (w ^ cb ^ 86 (w >> 19)) ^ cb ^ 86 (t ^ cb ^ 86 (t >> 8));
10 }
```

4.1.6 確率的素数判定 (Miller-Rabin 法)

$O(k \log^3 n)$

合成数を素数と判定する確率は最大で 4^{-k}

```
1 bool suspect(long a, int s, long d, long n) {
2     long x = mod_pow(a, d, n); // use mod_powl instead for large n
3     if (x == 1) return true;
4     for (int r = 0; r < s; ++r) {
5         if (x == n - 1) return true;
6         x = x * x % n; // use mod_mul instead for large n
7     }
8     return false;
9 }
10 // {2,7,61,-1} is for n < 4759123141 (= 2^32)
11 // {2,3,5,7,11,13,17,19,23,-1} is for n < 10^16 (at least)
12 bool is_prime(long n) {
13     if (n <= 1 || (n > 2 && n % 2 == 0)) return false;
14     int test[] = {2,3,5,7,11,13,17,19,23,-1};
15     long d = n - 1, s = 0;
```

```

16 while (d % 2 == 0) ++s, d /= 2;
17 for (int i = 0; test[i] < n && test[i] != -1; ++i)
18     if (!suspect(test[i], s, d, n)) return false;
19 return true;
20 }

```

4.2 多項式

FFT は基本定数重めなので TLE に注意する。

4.2.1 FFT(complex)

$O(N \log N)$

複素数を用いた FFT. 変換する vector のサイズは 2 の冪乗にすること.

```

1 typedef complex<double> cd;
2 vector<cd> fft(vector<cd> f, bool inv){
3     int n, N = f.size();
4     for(n=0;;n++) if(N == (1<n)) break;
5     rep(m,N){
6         int m2 = 0;
7         rep(i,n) if(m&(1<i)) m2 |= (1<(n-1-i));
8         if(m < m2) swap(f[m], f[m2]);
9     }
10
11     for(int t=1;t<N;t*=2){
12         double theta = acos(-1.0) / t;
13         cd w(cos(theta), sin(theta));
14         if(inv) w = cd(cos(theta), -sin(theta));
15         for(int i=0;i<N;i+=2*t){
16             cd power(1.0, 0.0);
17             rep(j,t){
18                 cd tmp1 = f[i+j] + f[i+t+j] * power;
19                 cd tmp2 = f[i+j] - f[i+t+j] * power;
20                 f[i+j] = tmp1;
21                 f[i+t+j] = tmp2;
22                 power = power * w;
23             }
24         }
25     }
26     if(inv) rep(i,N) f[i] /= N;
27     return f;
28 }

```

4.2.2 FFT(modulo)

$O(N \log N)$

剰余環を用いた FFT(FMT). 変換する vector のサイズは 2 の冪乗にすること. mod は $a * 2^e + 1$ の形.

```

1 #include "number_theory.cpp"
2
3 const int mod = 7*17*(1<23)+1;
4 vector<int> fmt(vector<int> f, bool inv){
5     int e, N = f.size();
6     // assert((N&(N-1))==0 and "f.size() must be power of 2");
7     for(e=0;;e++) if(N == (1<e)) break;
8     rep(m,N){
9         int m2 = 0;
10        rep(i,e) if(m&(1<i)) m2 |= (1<(e-1-i));

```

```

11        if(m < m2) swap(f[m], f[m2]);
12    }
13    for(int t=1; t<N; t*=2){
14        int r = pow_mod(3,(mod-1)/(t*2),mod);
15        if(inv) r = mod_inverse(r,mod);
16        for(int i=0; i<N; i+=2*t){
17            int power = 1;
18            rep(j,t){
19                int x = f[i+j], y = 1LL*f[i+t+j]*power%mod;
20                f[i+j] = (x+y)%mod;
21                f[i+t+j] = (x-y+mod)%mod;
22                power = 1LL*power*r%mod;
23            }
24        }
25    }
26    if(inv) for(int i=0,ni=mod_inv(N,mod);i<N;i++) f[i] = 1LL*f[i]*ni%mod;
27    return f;
28 }

```

4.2.3 積 (FMT)

$O(N \log N)$

`poly_mul()` が必要.

```

1 vector<int> poly_mul(vector<int> f, vector<int> g){
2     int N = max(f.size(),g.size())*2;
3     f.resize(N); g.resize(N);
4     f = fmt(f,0); g = fmt(g,0);
5     rep(i,N) f[i] = 1LL*f[i]*g[i]%mod;
6     f = fmt(f,1);
7     return f;
8 }

```

4.2.4 逆元 (FMT)

$O(N \log N)$

`extgcd()`, `mod_inverse()`, `poly_mul()`, `fmt()` が必要.

```

1 vector<int> poly_inv(const vector<int> &f){
2     int N = f.size();
3     vector<int> r(1,mod_inv(f[0],mod));
4     for(int k = 2; k <= N; k <= 1){
5         vector<int> nr = poly_mul(poly_mul(r,r), vector<int>(f.begin(),f.begin()+k));
6         nr.resize(k);
7         rep(i,k/2) {
8             nr[i] = (2*r[i]-nr[i]+mod)%mod;
9             nr[i+k/2] = (mod-nr[i+k/2])%mod;
10        }
11        r = nr;
12    }
13    return r;
14 }

```

4.2.5 平方根 (FMT)

$O(N \log N)$

`extgcd()`, `mod_inverse()`, `poly_inv()`, `poly_mul()`, `fmt()` が必要.

```

1  const int inv2 = (mod+1)/2;
2  vector<int> poly_sqrt(const vector<int> &f) {
3      int N = f.size();
4      vector<int> s(1,1); // s[0] = sqrt(f[0])
5      for(int k = 2; k <= N; k <= 1) {
6          s.resize(k);
7          vector<int> ns = poly_mul(poly_inv(s), vector<int>(f.begin(),f.begin()+k));
8          ns.resize(k);
9          rep(i,k) s[i] = 1LL*(s[i]+ns[i])*inv2%mod;
10     }
11     return s;
12 }

```

4.3 行列

```

1  typedef double number;
2  typedef vector<number> vec;
3  typedef vector<vec> mat;
4
5  vec mul(const mat& A, const vec& x) {
6      const int n = A.size();
7      vec b(n);
8      rep(i, n) rep(j, A[0].size()) {
9          b[i] = A[i][j] * x[j];
10     }
11     return b;
12 }
13
14 mat mul(const mat& A, const mat& B) {
15     const int n = A.size();
16     const int o = A[0].size();
17     const int m = B[0].size();
18     mat C(n, vec(m));
19     rep(i, n) rep(k, o) rep(j, m) {
20         C[i][j] += A[i][k] * B[k][j];
21     }
22     return C;
23 }
24
25 mat pow(mat A, long m) {
26     const int n = A.size();
27     mat B(n, vec(n));
28     rep(i, n) B[i][i] = 1;
29     do {
30         if (m & 1) B = mul(B, A);
31         A = mul(A, A);
32     } while (m >>= 1);
33     return B;
34 }
35
36 const number eps = 1e-4;
37
38 // determinant; O(n^3)
39 number det(mat A) {
40     int n = A.size();
41     number D = 1;
42     rep(i,n){
43         int pivot = i;
44         repi(j,i+1,n)
45             if (abs(A[j][i]) > abs(A[pivot][i])) pivot = j;
46         swap(A[pivot], A[i]);
47         D *= A[i][i] * (i != pivot ? -1 : 1);
48         if (abs(A[i][i]) < eps) break;

```

```

49         repi(j,i+1,n)
50             for(int k=n-1;k>=i;--k)
51                 A[j][k] -= A[i][k] * A[j][i] / A[i][i];
52     }
53     return D;
54 }
55
56 // rank; O(n^3)
57 int rank(mat A) {
58     int n = A.size(), m = A[0].size(), r = 0;
59     for(int i = 0; i < m and r < n; i++){
60         int pivot = r;
61         repi(j,r+1,n)
62             if (abs(A[j][i]) > abs(A[pivot][i])) pivot = j;
63         swap(A[pivot], A[r]);
64         if (abs(A[r][i]) < eps) continue;
65         for(int k=m-1;k>=i;--k)
66             A[r][k] /= A[r][i];
67         repi(j,r+1,n) repi(k,i,m)
68             A[j][k] -= A[r][k] * A[j][i];
69         ++r;
70     }
71     return r;
72 }

```

4.3.1 線形方程式の解 (Givens 消去法)

$O(N^3)$

```

1  // Givens elimination; O(n^3)
2
3  typedef double number;
4  typedef vector<vector<number>> > matrix;
5
6  inline double my_hypot(double x, double y) { return sqrt(x * x + y * y); }
7  inline void givens_rotate(number& x, number& y, number c, number s) {
8      number u = c * x + s * y, v = -s * x + c * y;
9      x = u, y = v;
10 }
11 vector<number> givens(matrix A, vector<number> b) {
12     const int n = b.size();
13     rep(i, n) repi(j, i + 1, n) {
14         const number r = my_hypot(A[i][i], A[j][i]);
15         const number c = A[i][i] / r, s = A[j][i] / r;
16         givens_rotate(b[i], b[j], c, s);
17         repi(k, i, n) givens_rotate(A[i][k], A[j][k], c, s);
18     }
19     for (int i = n - 1; i >= 0; --i) {
20         repi(j, i + 1, n) b[i] -= A[i][j] * b[j];
21         b[i] /= A[i][i];
22     }
23     return b;
24 }

```

4.4 割り当て問題

4.4.1 ハンガリアン法

$O(N^2)$

```

1  int hungarian(const vector<vector<int>>> &a) {
2      int n = a.size(), p, q;

```

```

3   vector<int> fx(n, inf), fy(n, 0), x(n, -1), y(n, -1);
4   rep(i,n) rep(j,n) fx[i] = max(fx[i], a[i][j]);
5
6   for (int i = 0; i < n; ) {
7       vector<int> t(n, -1), s(n+1, i);
8       for (p = q = 0; p <= q && x[i] < 0; ++p)
9           for (int k = s[p], j = 0; j < n && x[i] < 0; ++j)
10              if (fx[k] + fy[j] == a[k][j] && t[j] < 0) {
11                  s[++q] = y[j], t[j] = k;
12                  if (s[q] < 0)
13                      for (p = j; p >= 0; j = p)
14                          y[j] = k = t[j], p = x[k], x[k] = j;
15              }
16              if (x[i] < 0) {
17                  int d = inf;
18                  rep(k,q+1) rep(j,n) if (t[j] < 0) d = min(d, fx[s[k]] + fy[j] - a[s[k]][j]);
19                  rep(j,n) fy[j] += (t[j] < 0 ? 0 : d);
20                  rep(k,q+1) fx[s[k]] -= d;
21              } else i++;
22          }
23          int ret = 0;
24          rep(i,n) ret += a[i][x[i]];
25          return ret;
26      }

```

5 幾何

```

1 // constants and eps-considered operators
2 const double eps = 1e-8; // choose carefully!
3 const double pi = acos(-1.0);
4
5 inline bool lt(double a, double b) { return a < b - eps; }
6 inline bool gt(double a, double b) { return lt(b, a); }
7 inline bool le(double a, double b) { return !lt(b, a); }
8 inline bool ge(double a, double b) { return !lt(a, b); }
9 inline bool ne(double a, double b) { return lt(a, b) or lt(b, a); }
10 inline bool eq(double a, double b) { return !ne(a, b); }

```

5.1 点

```

1 typedef complex<double> point;
2 inline double dot (point a, point b) { return real(conj(a) * b); }
3 inline double cross(point a, point b) { return imag(conj(a) * b); }
4 /*
5  * Here is what ccw(a, b, c) returns:
6  *
7  *      1
8  *  -----
9  *  2 |a  0  b| -2
10 *  -----
11 *      -1
12 *
13 * Note: we can implement intersectPS(p, s) as !ccw(s.a, s.b, p).
14 */
15 int ccw(point a, point b, point c) {
16     b -= a, c -= a;
17     if (cross(b, c) > eps) return +1;
18     if (cross(b, c) < eps) return -1;
19     if (dot(b, c) < eps) return +2; // c -- a -- b
20     if (lt(norm(b), norm(c))) return -2; // a -- b -- c

```

```

21     return 0;
22 }

```

5.2 直線と線分

```

1 struct line {
2     point a, b;
3     line(point a, point b) : a(a), b(b) {}
4 };
5
6 bool intersectLS(const line& l, const line& s) {
7     return ccw(l.a, l.b, s.a) * ccw(l.a, l.b, s.b) <= 0;
8 }
9 bool intersectSS(const line& s, const line& t) {
10    return intersectLS(s, t) and intersectLS(t, s);
11 }
12 bool intersectLL(const line& l, const line& m) {
13     return ne(cross(l.b - l.a, m.b - m.a), 0.0) // not parallel
14         or eq(cross(l.b - l.a, m.a - l.a), 0.0); // overlap
15 }
16 point crosspointLL(const line& l, const line& m) {
17     double A = cross(l.b - l.a, m.b - m.a);
18     double B = cross(l.b - l.a, m.a - l.a);
19     if (eq(A, 0.0) and eq(B, 0.0)) return m.a; // overlap
20     assert(ne(A, 0.0)); // not parallel
21     return m.a - B / A * (m.b - m.a);
22 }
23 point proj(const line& l, point p) {
24     double t = dot(l.b - l.a, p - l.a) / norm(l.b - l.a);
25     return l.a + t * (l.b - l.a);
26 }
27 point reflection(const line& l, point p) { return 2.0 * proj(l, p) - p; }
28
29 double distanceLP(const line& l, point p) { return abs(proj(l, p) - p); }
30 double distanceLL(const line& l, const line& m) {
31     return intersectLL(l, m) ? 0.0 : distanceLP(l, m.a);
32 }
33 double distanceLS(const line& l, const line& s) {
34     return intersectLS(l, s) ? 0.0 : min(distanceLP(l, s.a), distanceLP(l, s.b));
35 }
36 double distancePS(point p, const line& s) {
37     point h = proj(s, p);
38     return ccw(s.a, s.b, h) ? min(abs(s.a - p), abs(s.b - p)) : abs(h - p);
39 }
40 double distanceSS(const line& s, const line& t) {
41     if (intersectSS(s, t)) return 0.0;
42     return min(min(distancePS(s.a, t), distancePS(s.b, t)),
43               min(distancePS(t.a, s), distancePS(t.b, s)));
44 }

```

5.3 円

```

1 struct circle {
2     point o; double r;
3     circle(point o, double r) : o(o), r(r) {}
4 };
5
6 bool intersectCL(const circle& c, const line& l) {
7     return le(norm(proj(l, c.o) - c.o), c.r * c.r);
8 }

```

```

9 int intersectCS(const circle& c, const line& s) {
10     if (not intersectCL(c, s)) return 0;
11     double a = abs(s.a - c.o);
12     double b = abs(s.b - c.o);
13     if (lt(a, c.r) and lt(b, c.r)) return 0;
14     if (lt(a, c.r) or lt(b, c.r)) return 1;
15     return ccw(s.a, s.b, proj(s, c.o)) ? 0 : 2;
16 }
17 bool intersectCC(const circle& c, const circle& d) {
18     double dist = abs(d.o - c.o);
19     return le(abs(c.r - d.r), dist) and le(dist, c.r + d.r);
20 }
21 line crosspointCL(const circle& c, const line& l) {
22     point h = proj(l, c.o);
23     double a = sqrt(c.r * c.r - norm(h - c.o));
24     point d = a * (l.b - l.a) / abs(l.b - l.a);
25     return line(h - d, h + d);
26 }
27 line crosspointCC(const circle& c, const circle& d) {
28     double dist = abs(d.o - c.o), th = arg(d.o - c.o);
29     double ph = acos((c.r * c.r + dist * dist - d.r * d.r) / (2.0 * c.r * dist));
30     return line(c.o + polar(c.r, th - ph), c.o + polar(c.r, th + ph));
31 }
32
33 line tangent(const circle& c, double th) {
34     point h = c.o + polar(c.r, th);
35     point d = polar(c.r, th) * point(0, 1);
36     return line(h - d, h + d);
37 }
38 vector<line> common_tangents(const circle& c, const circle& d) {
39     vector<line> ret;
40     double dist = abs(d.o - c.o), th = arg(d.o - c.o);
41     if (abs(c.r - d.r) < dist) { // outer
42         double ph = acos((c.r - d.r) / dist);
43         ret.pb(tangent(c, th - ph));
44         ret.pb(tangent(c, th + ph));
45     }
46     if (abs(c.r + d.r) < dist) { // inner
47         double ph = acos((c.r + d.r) / dist);
48         ret.pb(tangent(c, th - ph));
49         ret.pb(tangent(c, th + ph));
50     }
51     return ret;
52 }
53 pair<circle, circle> tangent_circles(const line& l, const line& m, double r) {
54     double th = arg(m.b - m.a) - arg(l.b - l.a);
55     double ph = (arg(m.b - m.a) + arg(l.b - l.a)) / 2.0;
56     point p = crosspointLL(l, m);
57     point d = polar(r / sin(th / 2.0), ph);
58     return mp(circle(p - d, r), circle(p + d, r));
59 }
60 line bisector(point a, point b);
61 circle circum_circle(point a, point b, point c) {
62     point o = crosspointLL(bisector(a, b), bisector(a, c));
63     return circle(o, abs(a - o));
64 }

```

5.4 多角形

```

1 typedef vector<point> polygon;
2
3 double area(const polygon& g) {
4     double ret = 0.0;
5     int j = g.size() - 1;

```

```

6     rep(i, g.size()) {
7         ret += cross(g[j], g[i]), j = i;
8     }
9     return ret / 2.0;
10 }
11 point centroid(const polygon& g) {
12     if (g.size() == 1) return g[0];
13     if (g.size() == 2) return (g[0] + g[1]) / 2.0;
14     point ret = 0.0;
15     int j = g.size() - 1;
16     rep(i, g.size()) {
17         ret += cross(g[j], g[i]) * (g[j] + g[i]), j = i;
18     }
19     return ret / area(g) / 6.0;
20 }
21 line bisector(point a, point b) {
22     point m = (a + b) / 2.0;
23     return line(m, m + (b - a) * point(0, 1));
24 }
25 polygon convex_cut(const polygon& g, const line& l) {
26     polygon ret;
27     int j = g.size() - 1;
28     rep(i, g.size()) {
29         if (ccw(l.a, l.b, g[j]) != -1) ret.pb(g[j]);
30         if (intersectLS(l, line(g[j], g[i]))) ret.pb(crosspointLL(l, line(g[j], g[i])));
31         j = i;
32     }
33     return ret;
34 }
35 polygon voronoi_cell(polygon g, const vector<point>& v, int k) {
36     rep(i, v.size()) if (i != k) {
37         g = convex_cut(g, bisector(v[i], v[k]));
38     }
39     return g;
40 }

```

5.4.1 凸包

```

1 namespace std {
2     bool operator <(const point& a, const point& b) {
3         return ne(real(a), real(b)) ? lt(real(a), real(b)) : lt(imag(a), imag(b));
4     }
5 }
6
7 polygon convex_hull(vector<point> v) {
8     const int n = v.size();
9     sort(all(v));
10    polygon ret(2 * n);
11    int k = 0;
12    for (int i = 0; i < n; ret[k++] = v[i++]) {
13        while (k >= 2 and ccw(ret[k - 2], ret[k - 1], v[i]) <= 0) --k;
14    }
15    for (int i = n - 2, t = k + 1; i >= 0; ret[k++] = v[i--]) {
16        while (k >= t and ccw(ret[k - 2], ret[k - 1], v[i]) <= 0) --k;
17    }
18    ret.resize(k - 1);
19    return ret;
20 }

```

5.4.2 最近点对

だいたい $O(n \log n)$, 最悪縦 1 列に並んでる場合 $O(n^2)$

```

1 pair<point,point> closest_pair(vector<point> p) {
2     int n = p.size(), s = 0, t = 1, m = 2, S[n];
3     S[0] = 0, S[1] = 1;
4     sort(all(p)); // "p < q" <=> "p.x < q.x"
5     double d = norm(p[s]-p[t]);
6     for (int i = 2; i < n; S[m++] = i++) rep(j, m) {
7         if (norm(p[S[j]]-p[i])<d) d = norm(p[s = S[j]]-p[t = i]);
8         if (real(p[S[j]]) < real(p[i]) - d) S[j--] = S[--m];
9     }
10    return make_pair(p[s], p[t]);
11 }

```

5.4.3 点-多角形包含判定

$O(n)$

```

1 enum { OUT, ON, IN };
2 int contains(const polygon& P, const point& p) {
3     bool in = false;
4     for (int i = 0; i < (int)P.size(); ++i) {
5         point a = P[i] - p, b = P[(i+1)%P.size()] - p;
6         if (imag(a) > imag(b)) swap(a, b);
7         if (imag(a) <= 0 && 0 < imag(b) && cross(a, b) < 0) in = !in;
8         if (cross(a, b) == 0 && dot(a, b) <= 0) return ON;
9     }
10    return in ? IN : OUT;
11 }

```

5.4.4 凸多角形の共通部分

$O(n+m)$

```

1 bool intersect_lpt(const point& a, const point& b,
2                   const point& c, const point& d, point &r) {
3     number D = cross(b - a, d - c);
4     if (eq(D,0)) return false;
5     number t = cross(c - a, d - c) / D;
6     number s = -cross(a - c, b - a) / D;
7     r = a + t * (b - a);
8     return ge(t, 0) && le(t, 1) && ge(s, 0) && le(s, 1);
9 }
10 polygon convex_intersect(const polygon &P, const polygon &Q) {
11     const int n = P.size(), m = Q.size();
12     int a = 0, b = 0, aa = 0, ba = 0;
13     enum { Pin, Qin, Unknown } in = Unknown;
14     polygon R;
15     do {
16         int a1 = (a+n-1) % n, b1 = (b+m-1) % m;
17         number C = cross(P[a] - P[a1], Q[b] - Q[b1]);
18         number A = cross(P[a1] - Q[b], P[a] - Q[b]);
19         number B = cross(Q[b1] - P[a], Q[b] - P[a]);
20         point r;
21         if (intersect_lpt(P[a1], P[a], Q[b1], Q[b], r)) {
22             if (in == Unknown) aa = ba = 0;
23             R.push_back(r);
24             in = B > 0 ? Pin : A > 0 ? Qin : in;
25         }
26         if (C == 0 && B == 0 && A == 0) {
27             if (in == Pin) { b = (b + 1) % m; ++ba; }
28             else { a = (a + 1) % n; ++aa; }
29         } else if (C >= 0) {

```

```

30         if (A > 0) { if (in == Pin) R.push_back(P[a]); a = (a+1)%n; ++aa; }
31         else { if (in == Qin) R.push_back(Q[b]); b = (b+1)%m; ++ba; }
32     } else {
33         if (B > 0) { if (in == Qin) R.push_back(Q[b]); b = (b+1)%m; ++ba; }
34         else { if (in == Pin) R.push_back(P[a]); a = (a+1)%n; ++aa; }
35     }
36 } while ( (aa < n || ba < m) && aa < 2*n && ba < 2*m );
37 if (in == Unknown) {
38     if (convex_contains(Q, P[0])) return P;
39     if (convex_contains(P, Q[0])) return Q;
40 }
41 return R;
42 }

```

5.4.5 凸多角形の直径

$O(n)$

```

1 inline double diff(const vector<point> &P, const int &i) { return (P[(i+1)%P.size()] - P[i]); }
2 number convex_diameter(const polygon &pt) {
3     const int n = pt.size();
4     int is = 0, js = 0;
5     for (int i = 1; i < n; ++i) {
6         if (imag(pt[i]) > imag(pt[is])) is = i;
7         if (imag(pt[i]) < imag(pt[js])) js = i;
8     }
9     number maxd = norm(pt[is]-pt[js]);
10
11     int i, maxi, j, maxj;
12     i = maxi = is;
13     j = maxj = js;
14     do {
15         if (cross(diff(pt,i), diff(pt,j)) >= 0) j = (j+1) % n;
16         else i = (i+1) % n;
17         if (norm(pt[i]-pt[j]) > maxd) {
18             maxd = norm(pt[i]-pt[j]);
19             maxi = i; maxj = j;
20         }
21     } while (i != is || j != js);
22     return maxd; /* farthest pair is (maxi, maxj). */
23 }

```

5.4.6 ドロネー三角形分割 (逐次添加法)

$O(n^2)$

```

1 bool incircle(point a, point b, point c, point p) {
2     a -= p; b -= p; c -= p;
3     return norm(a) * cross(b, c)
4         + norm(b) * cross(c, a)
5         + norm(c) * cross(a, b) >= 0;
6     // < : inside, = cocircular, > outside
7 }
8 #define SET_TRIANGLE(i, j, r) \
9     E[i].insert(j); em[i][j] = r; \
10    E[j].insert(r); em[j][r] = i; \
11    E[r].insert(i); em[r][i] = j; \
12    S.push(pair<int,int>(i, j));
13 #define REMOVE_EDGE(i, j) \
14     E[i].erase(j); em[i][j] = -1; \
15     E[j].erase(i); em[j][i] = -1;

```



```

16 #define DECOMPOSE_ON(i,j,k,r) { \
17     int m = em[j][i]; REMOVE_EDGE(j,i); \
18     SET_TRIANGLE(i,m,r); SET_TRIANGLE(m,j,r); \
19     SET_TRIANGLE(j,k,r); SET_TRIANGLE(k,i,r); }
20 #define DECOMPOSE_IN(i,j,k,r) { \
21     SET_TRIANGLE(i,j,r); SET_TRIANGLE(j,k,r); \
22     SET_TRIANGLE(k,i,r); }
23 #define FLIP_EDGE(i,j) { \
24     int k = em[j][i]; REMOVE_EDGE(i,j); \
25     SET_TRIANGLE(i,k,r); SET_TRIANGLE(k,j,r); }
26 #define IS_LEGAL(i, j) \
27     (em[i][j] < 0 || em[j][i] < 0 || \
28     !incircle(P[i],P[j],P[em[i][j]],P[em[j][i]]))
29 double Delaunay(vector<point> P) {
30     const int n = P.size();
31     P.push_back( point(-inf,-inf) );
32     P.push_back( point(+inf,-inf) );
33     P.push_back( point( 0 ,+inf) );
34     int em[n+3][n+3]; memset(em, -1, sizeof(em));
35     set<int> E[n+3];
36     stack< pair<int,int> > S;
37     SET_TRIANGLE(n+0, n+1, n+2);
38     for (int r = 0; r < n; ++r) {
39         int i = n, j = n+1, k;
40         while (1) {
41             k = em[i][j];
42             if (ccw(P[i], P[em[i][j]], P[r]) == +1) j = k;
43             else if (ccw(P[j], P[em[i][j]], P[r]) == -1) i = k;
44             else break;
45         }
46         if (ccw(P[i], P[j], P[r]) != +1) { DECOMPOSE_ON(i,j,k,r); }
47         else if (ccw(P[j], P[k], P[r]) != +1) { DECOMPOSE_ON(j,k,i,r); }
48         else if (ccw(P[k], P[i], P[r]) != +1) { DECOMPOSE_ON(k,i,j,r); }
49         else { DECOMPOSE_IN(i,j,k,r); }
50         while (!S.empty()) {
51             int u = S.top().first, v = S.top().second; S.pop();
52             if (!IS_LEGAL(u, v)) FLIP_EDGE(u, v);
53         }
54     }
55     double minarg = 1e5;
56     for (int a = 0; a < n; ++a) {
57         for(auto &b: E[a]) {
58             int c = em[a][b];
59             if (b < n && c < n) {
60                 point p = P[a] - P[b], q = P[c] - P[b];
61                 minarg = min(minarg, acos(dot(p,q)/abs(p)/abs(q)));
62             }
63         }
64     }
65     return minarg;
66 }

```

6 データ構造

6.1 Union-Find 木

```

1 class disjoint_set {
2     vector<int> p;
3 public:
4     disjoint_set(int n) : p(n, -1) {}
5     int root(int i) { return p[i] >= 0 ? p[i] = root(p[i]) : i; }
6     bool same(int i, int j) { return root(i) == root(j); }
7     int size(int i) { return -p[root(i)]; }

```

```

8     void merge(int i, int j) {
9         i = root(i), j = root(j);
10        if (i == j) return;
11        if (p[i] > p[j]) swap(i, j);
12        p[i] += p[j], p[j] = i;
13    }
14 };

```

6.2 Meldable Heap

```

1 template <class T>
2 class meldable_heap {
3     struct node {
4         node *l = NULL, *r = NULL;
5         T val;
6         node(const T& val) : val(val) {}
7         ~node() { delete l, delete r; }
8     };
9     node *meld(node *a, node *b) {
10        if (!a) return b;
11        if (!b) return a;
12        if (a->val > b->val) swap(a, b);
13        a->r = meld(a->r, b);
14        swap(a->l, a->r);
15        return a;
16    }
17    node *root = NULL;
18    meldable_heap(node *root) : root(root) {}
19 public:
20    meldable_heap() {}
21    bool empty() const { return !root; }
22    const T& top() const { return root->val; }
23    void meld(const meldable_heap<T>&& t) { root = meld(root, t.root); }
24    void push(const T& val) { root = meld(root, new node(val)); }
25    void pop() {
26        node *t = root;
27        root = meld(t->l, t->r);
28        t->l = t->r = NULL;
29        delete t;
30    }
31 };

```

6.3 Binary-Indexed-Tree

0-indexed

```

1 template<class T> struct bit {
2     int n;
3     vector<T> dat;
4
5     bit(int n) : n(n) { dat.assign(n,0); }
6     // sum [0,i)
7     T sum(int i){
8         int ret = 0;
9         for(--i; i>=0; i=(i&(i+1))-1) ret += dat[i];
10        return ret;
11    }
12    // sum [i,j)
13    T sum(int i, int j){ return sum(j) - sum(i); }
14    // add x to i
15    void add(int i, T x){ for(; i < n; i|=i+1) dat[i] += x; }

```

```
16 };
```

6.4 Segment Tree

区間 add と RMQ ができる.

```
1 template<class T> struct segtree {
2     int N;
3     vector<T> dat, sum;
4     segtree(int n) {
5         N = 1;
6         while(N < n) N <= 1;
7         dat.assign(2*N-1,0);
8         sum.assign(2*N-1,0);
9     }
10    void add(int a, int b, T x) { add(a,b,x,0,0,N);}
11    T add(int a, int b, T x, int k, int l, int r) {
12        if(b <= l or r <= a) return dat[k];
13        if(a <= l and r <= b) {
14            sum[k] += x;
15            return dat[k] += x;
16        }
17        int m = (l+r)/2;
18        return dat[k] = min(add(a,b,x,2*k+1,l,m),add(a,b,x,2*k+2,m,r))+sum[k];
19    }
20    T minimum(int a, int b) { return minimum(a,b,0,0,N);}
21    T minimum(int a, int b, int k, int l, int r) {
22        if(b <= l or r <= a) return le9;
23        if(a <= l and r <= b) return dat[k];
24        int m = (l+r)/2;
25        return min(minimum(a,b,2*k+1,l,m),minimum(a,b,2*k+2,m,r))+sum[k];
26    }
27 };
```

6.5 Sparse table

```
1 const int N = 200010;
2 const int K = 18;
3 int st[K][N];
4 void construct(int *a, int n) {
5     copy_n(a, n, st[0]);
6     rep1(k, 1, K) {
7         for (int i = 0; i+(1<<k) <= n; ++i) {
8             st[k][i] = min(st[k-1][i], st[k-1][i+(1<<(k-1))]);
9         }
10    }
11 }
12 int query(int a, int b) {
13     int k = 31-__builtin_clz(b-a);
14     return min(st[k][a], st[k][b-(1<<k)]);
15 }
```

6.6 RBST

```
1 struct node {
2     long val, sum;
3     size_t size = 1;
4     node *left = NULL, *right = NULL;
```

```
5     node(long val) : val(val), sum(val) {}
6     ~node() { delete left, delete right; }
7 };
8 inline long sum(node *u) { return u ? u->sum : 0; }
9 inline size_t size(node *u) { return u ? u->size : 0; }
10 inline node *pull(node *u) {
11     u->sum = u->val + sum(u->left) + sum(u->right);
12     u->size = 1 + size(u->left) + size(u->right);
13     return u;
14 }
15 node *merge(node *u, node *v) {
16     if (!u) return v;
17     if (!v) return u;
18     if (rand() * long(size(u) + size(v)) < long(size(u)) * RAND_MAX) {
19         u->right = merge(u->right, v);
20         return pull(u);
21     } else {
22         v->left = merge(u, v->left);
23         return pull(v);
24     }
25 }
26 pair<node*,node*> split(node *u, size_t k) {
27     if (!u or k == 0) return {NULL, u};
28     if (k == size(u)) return {u, NULL};
29     if (size(u->left) >= k) {
30         auto p = split(u->left, k);
31         u->left = p.second;
32         return {p.first, pull(u)};
33     } else {
34         auto p = split(u->right, k - size(u->left) - 1);
35         u->right = p.first;
36         return {pull(u), p.second};
37     }
38 }
39 template <class ForwardIterator>
40 node *construct_from(ForwardIterator first, ForwardIterator last) {
41     if (first == last) return NULL;
42     auto mid = next(first, (last - first) / 2);
43     node *u = new node(*mid);
44     u->left = construct_from(first, mid);
45     u->right = construct_from(next(mid), last);
46     return pull(u);
47 }
```

6.7 永続 RBST

```
1 template <class T, size_t N>
2 struct mempool {
3     static T buf[N], *head;
4     static size_t cnt() { return head - buf; }
5     static void clear() { head = buf; }
6     void *operator new(size_t _ __attribute__((unused))) { return head++; }
7     void operator delete(void *_ __attribute__((unused))) {}
8 };
9 template <class T, size_t N> T mempool<T, N>::buf[N];
10 template <class T, size_t N> T *mempool<T, N>::head = mempool<T, N>::buf;
11
12 struct node;
13 long sum(node *u);
14 size_t size(node *u);
15 struct node : mempool<node, M> {
16     const long val = 0, sum = 0, lazy = 0;
17     const size_t size = 1;
18     node *const left = NULL, *const right = NULL;
```

```

19     node() {}
20     node(long val) : val(val), sum(val) {}
21     node(long val, long lazy, node *left, node *right)
22         : val(val),
23           sum(val + ::sum(left) + ::sum(right)),
24           lazy(lazy),
25           size(1 + ::size(left) + ::size(right)),
26           left(left),
27           right(right) {}
28 };
29 inline long sum(node *u) { return u ? u->sum + u->lazy * u->size : 0; }
30 inline size_t size(node *u) { return u ? u->size : 0; }
31 inline node *add(node *u, long x) { return u ? new node(u->val, u->lazy + x, u->left, u
->right) : NULL; }
32 node *merge(node *u, node *v) {
33     if (!u) return v;
34     if (!v) return u;
35     if (rand() * long(size(u) + size(v)) < long(size(u)) * RAND_MAX) {
36         return new node(u->val + u->lazy, 0, add(u->left, u->lazy), merge(add(u->right,
u->lazy), v));
37     } else {
38         return new node(v->val + v->lazy, 0, merge(u, add(v->left, v->lazy)), add(v->
right, v->lazy));
39     }
40 }
41 pair<node *, node *> split(node *u, size_t k) {
42     if (!u or k == 0) return {NULL, u};
43     if (k == size(u)) return {u, NULL};
44     if (size(u->left) >= k) {
45         auto p = split(add(u->left, u->lazy), k);
46         return {p.first, new node(u->val + u->lazy, 0, p.second, add(u->right, u->lazy
))};
47     } else {
48         auto p = split(add(u->right, u->lazy), k - size(u->left) - 1);
49         return {new node(u->val + u->lazy, 0, add(u->left, u->lazy), p.first), p.second
};
50     }
51 }
52 template <class OutputIterator>
53 OutputIterator dump(OutputIterator it, const node *u, long lazy = 0) {
54     if (!u) return it;
55     lazy += u->lazy;
56     it = dump(it, u->left, lazy);
57     *it++ = u->val + lazy;
58     return dump(it, u->right, lazy);
59 }
60 template <class ForwardIterator>
61 node *construct_from(ForwardIterator first, ForwardIterator last) {
62     if (first == last) return NULL;
63     auto mid = next(first, (last - first) / 2);
64     return new node(*mid, 0, construct_from(first, mid), construct_from(next(mid), last
));
65 }

```

6.8 赤黒木

```

1  template<class T> class rbtree {
2      enum COL { BLACK, RED,};
3      struct node {
4          T val, lazy, min_val;
5          int color, rnk, size;
6          node *left, *right;
7          // if !left then this node is leaf
8          node(){}

```

```

9      node(T v) : val(v), min_val(v), color(BLACK), rnk(0), size(1) {
10          lazy = 0;
11          left = right = NULL;
12      }
13      node(node *l, node *r, int c) : color(c) {
14          lazy = 0;
15          left = l;
16          right = r;
17          update();
18      }
19      void update() {
20          eval();
21          if(left) {
22              rnk = max(left->rnk+(left->color==BLACK),
23                        right->rnk+(right->color==BLACK));
24              size = left->size+right->size;
25              left->eval(); right->eval();
26              min_val = min(left->min_val, right->min_val);
27          }
28      }
29      void eval() {
30          min_val += lazy;
31          if(!left) val += lazy;
32          else {
33              left->lazy += lazy;
34              right->lazy += lazy;
35          }
36          lazy = 0;
37      }
38 };
39
40 node *new_node(T v) { return new node(v);}
41 node *new_node(node *l, node *r, int c) { return new node(l,r,c);}
42 node *rotate(node *v, int d) {
43     node *w = d? v->right: v->left;
44     if(d) {
45         v->right = w->left;
46         w->left = v;
47         v->right->update();
48     }
49     else {
50         v->left = w->right;
51         w->right = v;
52         v->left->update();
53     }
54     v->update(); w->update();
55     v->color = RED;
56     w->color = BLACK;
57     return w;
58 }
59 node *merge_sub(node *u, node *v) {
60     u->eval(); v->eval();
61     if(u->rnk < v->rnk) {
62         node *w = merge_sub(u,v->left);
63         v->left = w;
64         v->update();
65         if(v->color == BLACK and w->color == RED and w->left->color == RED) {
66             if(v->right->color == BLACK) return rotate(v,0);
67             else {
68                 v->color = RED;
69                 v->left->color = v->right->color = BLACK;
70                 return v;
71             }
72         }
73         else return v;
74     }
75     else if(u->rnk > v->rnk) {

```

```

76     node *w = merge_sub(u->right,v);
77     u->right = w;
78     u->update();
79     if(u->color == BLACK and w->color == RED and w->right->color == RED) {
80         if(u->left->color == BLACK) return rotate(u,l);
81         else {
82             u->color = RED;
83             u->left->color = u->right->color = BLACK;
84             return u;
85         }
86     }
87     else return u;
88 }
89 else return new_node(u,v,RED);
90 }
91 node *insert(node *v, int k) {
92     auto p = split(root,k);
93     return root = merge(merge(p.first,v),p.second);
94 }
95 void add(node *v, int res, T val) {
96     if(res < 1) return;
97     v->eval();
98     if(v->size == res) {
99         v->lazy += val;
100         return;
101     }
102     add(v->left, min(v->left->size, res), val);
103     add(v->right, res-v->left->size, val);
104     v->update();
105 }
106 T get(node *v, int k) {
107     v->eval();
108     if(!v->left) return v->val;
109     if(v->left->size > k) return get(v->left, k);
110     return get(v->right, k-v->left->size);
111 }
112 T minimum(node *v, int l, int r) {
113     if(r-l < 1) return inf;
114     v->eval();
115     if(v->size == r-l) return v->min_val;
116     return min(minimum(v->left, l, min(r, v->left->size)),
117               minimum(v->right, l-min(l, v->left->size), r-v->left->size));
118 }
119 T inf;
120 public:
121     node *root;
122     rbtree() {
123         inf = (((1LL<<(sizeof(T)*8-2))-1)<<1)+1;
124         root = NULL;
125     }
126     void clear() { delete root; root = NULL; }
127     node *build(const vector<T> &vs) {
128         if(!vs.size()) return root = NULL;
129         if((int)vs.size() == 1) return root = new_node(vs[0]);
130         int m = vs.size()/2;
131         return root = merge(build(vector<T>(begin(vs),begin(vs)+m)),
132                             build(vector<T>(begin(vs)+m,end(vs))));
133     }
134     int size() { return root? root->size: 0; }
135     node *push_back(T val) { return root = merge(root,new_node(val)); }
136     node *push_front(T val) { return root = merge(new_node(val),root); }
137     node *merge(node *u, node *v) {
138         if(!u) return v;
139         if(!v) return u;
140         u = merge_sub(u,v);
141         u->color = BLACK;
142

```

```

143     return u;
144 }
145 pair<node*,node*> split(node *v, int k) {
146     if(!k) return pair<node*,node*>(NULL,v);
147     if(k == v->size) return pair<node*,node*>(v,NULL);
148     v->eval();
149     if(k < v->left->size) {
150         auto p = split(v->left,k);
151         return pair<node*,node*>(p.first,merge(p.second,v->right));
152     }
153     else if(k > v->left->size) {
154         auto p = split(v->right,k-v->left->size);
155         return pair<node*,node*>(merge(v->left,p.first),p.second);
156     }
157     else return pair<node*,node*>(v->left,v->right);
158 }
159
160 node *insert(int k, T val) { return insert(new_node(val),k); }
161 node *erase(int k) {
162     auto p = split(root,k+1);
163     return root = merge(split(p.first,k).first, p.second);
164 }
165 void add(int l, int r, T val) { add(root, r, val); add(root, l, -val); }
166 T get(int k) { return get(root, k); }
167 T minimum(int l, int r) { return minimum(root, l, r); }
168 T operator[](const int &i) { return get(i); }
169 };

```

6.9 永続赤黒木

```

1 //const int MAX = 15000000, BOUND = 14000000;
2 template<class T> class prbtree {
3 public:
4     enum COL { BLACK, RED,};
5     struct node {
6         T val;
7         int color;
8         int rnk, size;
9         node *left, *right;
10
11     node(){}
12     node(T v) : val(v), color(BLACK), rnk(0), size(1) {
13         left = right = NULL;
14     }
15     node(node *l, node *r, int c) : color(c) {
16         left = l;
17         right = r;
18         rnk = max((!l? l->rnk+(l->color==BLACK): 0),
19                 (r? r->rnk+(r->color==BLACK): 0));
20         size = !l and !r? 1: !l? r->size: !r? r->size: l->size+r->size;
21     }
22 };
23
24 node *root;
25 //     node nodes[MAX];
26 //     int called;
27
28 prbtree() {
29     root = NULL;
30     // called = 0;
31 }
32
33 prbtree(T val) {
34     root = new_node(val);

```

```

35 // called = 0;
36 }
37
38 // node *new_node(T v) { return &(nodes[called++] = node(v));}
39 // node *new_node(node *l, node *r, int c) { return &(nodes[called++] = node(l,r,c
40 ));}
41 node *new_node(T v) { return new node(v);}
42 node *new_node(node *l, node *r, int c) { return new node(l,r,c);}
43
44 node *merge_sub(node *u, node *v) {
45     if(u->rnk < v->rnk) {
46         node *w = merge_sub(u,v->left);
47         if(v->color == BLACK and w->color == RED and w->left->color == RED){
48             if(v->right->color == BLACK) return new_node(w->left,new_node(w->right,
49                 v->right,RED),BLACK);
50             else return new_node(new_node(w->left,w->right,BLACK),new_node(v->right
51                 ->left,v->right->right,BLACK),RED);
52         }
53         else return new_node(w,v->right,v->color);
54     }
55     else if(u->rnk > v->rnk) {
56         node *w = merge_sub(u->right,v);
57         if(u->color == BLACK and w->color == RED and w->right->color == RED){
58             if(u->left->color == BLACK) return new_node(new_node(u->left,w->left,
59                 RED),w->right,BLACK);
60             else return new_node(new_node(u->left->left,u->left->right,BLACK),
61                 new_node(w->left,w->right,BLACK),RED);
62         }
63         else return new_node(u->left,w,u->color);
64     }
65     else return new_node(u,v,RED);
66 }
67
68 node *merge(node *u, node *v) {
69     if(!u) return v;
70     if(!v) return u;
71     u = merge_sub(u,v);
72     if(u->color == RED) return new_node(u->left,u->right,BLACK);
73     return u;
74 }
75
76 pair<node*,node*> split(node *v, int k) {
77     if(!k) return pair<node*,node*>(NULL,v);
78     if(k == v->size) return pair<node*,node*>(v,NULL);
79     if(k < v->left->size) {
80         auto p = split(v->left,k);
81         return pair<node*,node*>(p.first,merge(p.second,v->right));
82     }
83     else if(k > v->left->size) {
84         auto p = split(v->right,k-v->left->size);
85         return pair<node*,node*>(merge(v->left,p.first),p.second);
86     }
87     else return pair<node*,node*>(v->left,v->right);
88 }
89
90 node *build(const vector<T> &vs) {
91     if(!vs.size()) return NULL;
92     if((int)vs.size() == 1) return new_node(vs[0]);
93     int m = vs.size()/2;
94     return merge(build(vector<T>(begin(vs),begin(vs)+m)), build(vector<T>(begin(vs)+
95         m,end(vs))));
96 }
97
98 int size() { return root->size;}
99
100 void get(vector<T> &vs) { get(root,vs);}
101 void get(node *v, vector<T> &vs) {

```

```

96     if(!v->left and !v->right) vs.push_back(v->val);
97     else {
98         if(v->left) get(v->left,vs);
99         if(v->right) get(v->right,vs);
100     }
101 }
102
103 node *push_back(T val) {
104     node *v = new_node(val);
105     return root = merge(root,v);
106 }
107
108 // insert leaf at k
109 node *insert(int k, T val) {
110     return insert(new_node(val), k);
111 }
112
113 // insert tree v at k
114 node *insert(node *v, int k) {
115     auto p = split(root,k);
116     return root = merge(merge(p.first,v),p.second);
117 }
118
119 // copy [l,r)
120 node *copy(int l, int r) {
121     return split(split(root, l).second, r-l).first;
122 }
123
124 // copy and insert [l,r) at k
125 node *copy_paste(int l, int r, int k) {
126     return insert(copy(l,r),k);
127 }

```

6.10 wavelet 行列

N := 列の長さ

M := 最大値

6.10.1 完備辞書

function	計算量
count	$O(1)$
select	$O(\log N)$

```

1 template<int N> class FID {
2     static const int bucket = 512, block = 16;
3     static char popcount[];
4     int n, B[N/bucket+10];
5     unsigned short bs[N/block+10], b[N/block+10];
6
7 public:
8     FID(){}
9     FID(int n, bool s[]) : n(n) {
10         if(!popcount[i]) for (int i = 0; i < (1<<block); i++) popcount[i] =
11             __builtin_popcount(i);
12
13         bs[0] = B[0] = b[0] = 0;
14         for (int i = 0; i < n; i++) {

```

```

14         if(i%block == 0) {
15             bs[i/block+1] = 0;
16             if(i%bucket == 0) {
17                 B[i/bucket+1] = B[i/bucket];
18                 b[i/block+1] = b[i/block] = 0;
19             }
20             else b[i/block+1] = b[i/block];
21         }
22         bs[i/block] |= short(s[i])<<(i%block);
23         b[i/block+1] += s[i];
24         B[i/bucket+1] += s[i];
25     }
26     if(n%bucket == 0) b[n/block] = 0;
27 }
28
29 // number of val in [0,r), 0(1)
30 int count(bool val, int r) { return val? B[r/bucket]+b[r/block]+popcount[bs[r/block
    ]&((1<<(r%block))-1)]: r-count(1,r); }
31 // number of val in [l,r), 0(1)
32 int count(bool val, int l, int r) { return count(val,r)-count(val,l); }
33 // position of ith in val, 0-indexed, 0(log n)
34 int select(bool val, int i) {
35     if(i < 0 or count(val,n) <= i) return -1;
36     i++;
37     int lb = 0, ub = n, md;
38     while(ub-lb>1) {
39         md = (lb+ub)>>1;
40         if(count(val,md) >= i) ub = md;
41         else lb = md;
42     }
43     return ub-1;
44 }
45 int select(bool val, int i, int l) { return select(val,i+count(val,l)); }
46 bool operator[](int i) { return bs[i/block]>>(i%block)&1; }
47 };
48 template<int N> char FID<N>::popcount[1<<FID<N>::block];

```

6.10.2 wavelet 行列

function	計算量	FID::count	FID::select
count	$O(\log M)$	o	
select	$O(\log N \log M)$	o	o
get	$O(\log M)$	o	
maximum	$O(\log M)$ or $O(k \log M)$	o	
kth_number	$O(\log M)$	o	
freq	$O(\log M)$	o	
freq_list	$O(k \log M)$	o	
get_rect	$O(k \log N \log M)$	o	o

```

1 template<class T, int N, int D> class wavelet {
2     int n, zs[D];
3     FID<N> dat[D];
4
5     void max_dfs(int d, int l, int r, int &k, T val, vector<T> &vs) {
6         if(l >= r or !k) return;
7         if(d == D) {
8             while(l++ < r and k > 0) vs.push_back(val), k--;
9             return;

```

```

10     }
11     int lc = dat[d].count(1,l), rc = dat[d].count(1,r);
12     // if min, change this order
13     max_dfs(d+1, lc+zs[d], rc+zs[d], k, 1ULL<<(D-d-1)|val,vs);
14     max_dfs(d+1, l-lc, r-rc, k, val, vs);
15 }
16
17 T max_dfs(int d, int l, int r, T val, T a, T b) {
18     if(r-l <= 0 or val >= b) return -1;
19     if(d == D) return val>=a? val: -1;
20     int lc = dat[d].count(1,l), rc = dat[d].count(1,r);
21     T ret = max_dfs(d+1, lc+zs[d], rc+zs[d], 1ULL<<(D-d-1)|val, a, b);
22     if(~ret) return ret;
23     return max_dfs(d+1, l-lc, r-rc, val, a, b);
24 }
25
26 int freq_dfs(int d, int l, int r, T val, T a, T b) {
27     if(l == r) return 0;
28     if(d == D) return (a <= val and val < b)? r-l: 0;
29     T nv = 1ULL<<(D-d-1)|val, nnv = ((1ULL<<(D-d-1))-1)|nv;
30     if(nnv < a or b <= val) return 0;
31     if(a <= val and nnv < b) return r-l;
32     int lc = dat[d].count(1,l), rc = dat[d].count(1,r);
33     return freq_dfs(d+1,l-lc,r-rc,val,a,b)+
34         freq_dfs(d+1,lc+zs[d],rc+zs[d],nv,a,b);
35 }
36
37 void list_dfs(int d, int l, int r, T val, T a, T b, vector<pair<T,int>> &vs) {
38     if(val >= b or r-l <= 0) return;
39     if(d == D) {
40         if(a <= val) vs.push_back(make_pair(val,r-l));
41         return;
42     }
43     T nv = val|(1LL<<(D-d-1)), nnv = nv|(((1LL<<(D-d-1))-1));
44     if(nnv < a) return;
45     int lc = dat[d].count(1,l), rc = dat[d].count(1,r);
46     list_dfs(d+1,l-lc,r-rc,val,a,b,vs);
47     list_dfs(d+1,lc+zs[d],rc+zs[d],nv,a,b,vs);
48 }
49 public:
50     wavelet(int n, T seq[]) : n(n) {
51         T f[N], l[N], r[N];
52         bool b[N];
53         memcpy(f, seq, sizeof(T)*n);
54         for (int d = 0; d < D; d++) {
55             int lh = 0, rh = 0;
56             for (int i = 0; i < n; i++) {
57                 bool k = (f[i]>>(D-d-1))&1;
58                 if(k) r[rh++] = f[i];
59                 else l[lh++] = f[i];
60                 b[i] = k;
61             }
62             dat[d] = FID<N>(n,b);
63             zs[d] = lh;
64             swap(l,f);
65             memcpy(f+lh, r, rh*sizeof(T));
66         }
67     }
68
69     T get(int i) {
70         T ret = 0;
71         bool b;
72         for (int d = 0; d < D; d++) {
73             ret <<= 1;
74             b = dat[d][i];
75             ret |= b;
76             i = dat[d].count(b,i)+b*zs[d];

```

```

77     }
78     return ret;
79 }
80 T operator[](int i) { return get(i); }
81
82 int count(T val, int l, int r) {
83     for (int d = 0; d < D; d++) {
84         bool b = (val>>(D-d-1))&1;
85         l = dat[d].count(b,l)+b*zs[d];
86         r = dat[d].count(b,r)+b*zs[d];
87     }
88     return r-l;
89 }
90 int count(T val, int r) { return count(val,0,r); }
91
92 int select(T val, int k) {
93     int ls[D], rs[D], l = 0, r = n;
94     for (int d = 0; d < D; d++) {
95         ls[d] = l; rs[d] = r;
96         bool b = val>>(D-d-1)&1;
97         l = dat[d].count(b,l)+b*zs[d];
98         r = dat[d].count(b,r)+b*zs[d];
99     }
100    for (int d = D-1; d >= 0; d--) {
101        bool b = val>>(D-d-1)&1;
102        k = dat[d].select(b,k,ls[d]);
103        if(k >= rs[d] or k < 0) return -1;
104        k -= ls[d];
105    }
106    return k;
107 }
108 int select(T val, int k, int l) { return select(val,k+count(val,l)); }
109
110 vector<T> maximum(int l, int r, int k) {
111     if (r-l < k) k = r-l;
112     if(k < 0) return {};
113     vector<T> ret;
114     max_dfs(0,l,r,k,0,ret);
115     return ret;
116 }
117
118 T maximum(int l, int r, T a, T b) { return max_dfs(0,l,r,0,a,b); }
119
120 // k is 0-indexed
121 T kth_number(int l, int r, int k) {
122     if(r-l <= k or k < 0) return -1;
123     T ret = 0;
124     for (int d = 0; d < D; d++) {
125         int lc = dat[d].count(1,l), rc = dat[d].count(1,r);
126         if(rc-lc > k) {
127             l = lc+zs[d];
128             r = rc+zs[d];
129             ret |= 1ULL<<(D-d-1);
130         }
131         else {
132             k -= rc-lc;
133             l = lc;
134             r = rc;
135         }
136     }
137     return ret;
138 }
139
140 vector<pair<T,int>> freq_list(int l, int r, T a, T b) {
141     vector<pair<T,int>> ret;
142     list_dfs(0,l,r,0,a,b,ret);
143     return ret;

```

```

144     }
145
146     vector<pair<int,T>> get_rect(int l, int r, T a, T b) {
147         vector<pair<T,int>> res = freq_list(l,r,a,b);
148         vector<pair<int,T>> ret;
149         for(auto &e: res)
150             for (int i = 0; i < e.second; i++)
151                 ret.push_back(make_pair(select(e.first,i,l), e.first));
152         return ret;
153     }
154     // number of elements in [l,r) in [a,b), O(D)
155     int freq(int l, int r, T a, T b) { return freq_dfs(0,l,r,0,a,b); }
156 };

```

7 その他

7.1 ビジュアライザ

```

1 <script>
2 function line(x,y,a,b){c.b();c.moveTo(x,y);c.lineTo(a,b);c.s();}
3 function circle(x,y,r){c.b();c.arc(x,y,r,0,7,0);c.s();}
4 window.onload=function(){d=document;d.i=d.getElementById;
5 c=d.i('c').getContext('2d');c.b=c.beginPath;c.s=c.stroke;
6 d.i('s').src='data.js?';};
7 </script>
8 <body><canvas id="c" width="500" height="500" style="border:1px solid #000;"></canvas>
9 <script id="s"></script></body>

```