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1 準備

1.1 init.el

linum は emacs24 のみ

```
; key
(keyboard-translate ?\C-h ?\C-?)
(global-set-key "\M-g" 'goto-line)

; tab
(setq-default indent-tabs-mode nil)
(setq-default tab-width 4)
(setq indent-line-function 'insert-tab)

; line number
(global-linum-mode t)
(setq linum-format "%4d ")
```

1.2 tpl.cpp

3

```
#include <bits/stdc++.h>
   using namespace std;
   #define rep(i,a) for(int i = 0; i < (a); i++)
   #define repi(i,a,b) for(int i = (a); i < (b); i++)
   #define repd(i,a,b) for(int i = (a); i >= (b); i--)
   #define repit(i,a) for(__typeof((a).begin()) i = (a).begin(); i != (a).end(); i++)
   #define all(u) (u).begin(),(u).end()
   #define rall(u) (u).rbegin(),(u).rend()
   #define UNIQUE(u) (u).erase(unique(all(u)),(u).end())
   #define pb push_back
   #define mp make_pair
   const int INF = 1e9;
   const double EPS = 1e-8;
   const double PI = acos(-1.0);
   typedef long long 11;
   typedef vector<int> vi;
   typedef vector<vi> vvi;
   typedef pair<int,int> pii;
22
   int main(){
23
```

2 文字列

2.1 Aho-Corasick 法

O(N+M)

```
return res:
10
11
   // patternからパターンマッチングオートマトンの生成
12
   PMA *buildPMA(vector<string> pattern){
13
       PMA *root = new PMA, *now;
14
       root->next[0] = root;
15
       rep(i, patter.size()){
16
17
            now = root;
            rep(j, pattern[i].size()){
18
                if(now->next[(int)pattern[i][j]] == 0)
19
20
                    now->next[(int)pattern[i][j]] = new PMA;
                now = now->next[(int)pattern[i][j]];
21
22
23
            now->matched.push_back(i);
24
       queue < PMA*> que;
25
       repi(i,1,256){
26
            if(!root->next[i]) root->next[i] = root;
27
28
                root->next[i]->next[0] = root;
29
30
                que.push(root->next[i]);
31
32
       while(!que.empty()){
33
34
            now = que.front(); que.pop();
            repi(i,1,256){
35
36
               if(now->next[i]){
                    PMA *next = now->next[0];
37
38
                    while(!next->next[i]) next = next->next[0];
39
                    now->next[i]->next[0] = next->next[i];
                    now->next[i]->matched = set_union(now->next[i]->matched, next->next[i]->
40
                         matched);
                    que.push(now->next[i]);
41
42
43
44
45
       return root;
46
47
   void match(PMA* &pma, const string s, vector<int> &res){
       rep(i,s.size()){
48
49
            int c = s[i];
            while(!pma->next[c])
50
                pma = pma->next[0];
51
52
            pma = pma -> next[c]:
            rep(j,pma->matched.size())
53
                res[pma->matched[i]] = true;
54
55
56
```

3 グラフ

3.1 最大流

 $O(EV^2)$

```
struct edge{int to, cap, rev;};
vector<edge> G[MAX];
int level[MAX], itr[MAX];

void add_edge(int from, int to, int cap){
    G[from].push_back((edge){to, cap, int(G[to].size())});
    G[to].push_back((edge){from, 0, int(G[from].size()-1)});
}
```

```
void bfs(int s, int t){
10
        memset(level, -1, sizeof(level));
11
        queue<int> que; que.push(s);
12
        level[s] = 0;
13
        while(!que.empty()){
14
            int v = que.front(); que.pop();
15
            if(v == t) return;
16
            for(int i = 0; i < G[v].size(); i++){</pre>
17
                 edge &e = G[v][i];
18
                if(e.cap <= 0 or level[e.to] != -1) continue;</pre>
19
20
                que.push(e.to);
                level[e.to] = level[v]+1;
21
22
23
        }
   }
24
25
    int dfs(int v, int t, int f){
26
        if(v == t) return f;
27
        for(int &i = itr[v] ; i < G[v].size(); i++){</pre>
28
29
            edge &e = G[v][i];
30
            if(level[e.to] <= level[v] or e.cap <= 0) continue;</pre>
            int d = dfs(e.to, t, min(f, e.cap));
31
32
            if(d > 0)
33
                e.cap -= d;
34
                G[e.to][e.rev].cap += d;
35
                return d;
36
        }
37
38
        return 0;
39
40
    int max_flow(int s, int t){
        int flow = 0, f;
43
        while(1){
44
            bfs(s, t);
            if(level[t] == -1) return flow;
45
46
            memset(itr, 0, sizeof(itr));
47
            while((f = dfs(s, t, INF)) > 0) flow += f;
48
49
```

3.2 最小費用流

 $O(FE \log V)$

```
struct edge{ int to, cap, cost, rev;};
   int V;
   vector<edge> G[MAX_V];
   int h[MAX_V];
    int dist[MAX_V];
    int prevv[MAX_V], preve[MAX_V];
    void add_edge(int from, int to, int cap, int cost){
       G[from].push_back((edge){to, cap, cost, int(G[to].size())});
       G[to].push_back((edge){from, 0, -cost, int(G[from].size() - 1)});
11
12
13
    int min_cost_flow(int s, int t, int f){
14
       int res = 0;
15
       fill(h, h + V, 0);
16
17
       while(f > 0){
            priority_queue<pii, vector<pii>, greater<pii> > que;
18
            fill(dist, dist + V, inf);
19
            dist[s] = 0;
20
21
            que.push(pii(0, s));
```

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```
while(!que.empty()){
22
                pii p = que.top(); que.pop();
23
                int v = p.second;
24
                if(dist[v] < p.first) continue;</pre>
25
                rep(i,G[v].size()){
26
                    edge &e = G[v][i];
27
                    if(e.cap > 0 \&\& dist[e.to] > dist[v] + e.cost + h[v] - h[e.to]){
28
29
                        dist[e.to] = dist[v] + e.cost + h[v] - h[e.to];
                        prevv[e.to] = v;
30
                        preve[e.to] = i;
31
                        que.push(pii(dist[e.to], e.to));
32
33
                }
34
35
            if(dist[t] == inf) return -1;
36
            rep(v,V) h[v] += dist[v];
37
            int d = f;
38
            for(int v = t; v != s; v = prevv[v])
                d = min(d, G[prevv[v]][preve[v]].cap);
            f -= d;
42
            res += d * h[t];
            for(int v = t; v != s; v = prevv[v]){
43
                edge &e = G[prevv[v]][preve[v]];
                e.cap -= d;
45
                G[v][e.rev].cap += d;
47
48
49
       return res;
50
```

3.3 最小シュタイナー木

 $O(4^{|T|}V)$

g に無向グラフの隣接行列を入れる.

T に使いたい頂点の集合を入れる.

```
int minimum_steiner_tree(vi &T, vvi &g){
       int n = g.size(), t = T.size();
       if(t <= 1) return 0;
       vvi d(g); // all-pair shortest
       rep(k,n)rep(i,n)rep(j,n) //Warshall Floyd
            d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
       int opt[1 << t][n];</pre>
       rep(S,1<<t) rep(x,n)
            opt[S][x] = INF;
11
       rep(p,t) rep(q,n) // trivial case
12
            opt[1 << p][q] = d[T[p]][q];
13
14
       repi(S,1,1<<t){ // DP step
15
16
            if(!(S & (S-1))) continue;
            rep(p,n) rep(E,S)
17
18
                if((E \mid S) == S)
                    opt[S][p] = min(opt[S][p], opt[E][p] + opt[S-E][p]);
            rep(p,n) rep(q,n)
20
21
                opt[S][p] = min(opt[S][p], opt[S][q] + d[p][q]);
22
23
24
       int ans = INF;
       rep(S,1 << t) rep(q,n)
25
            ans = min(ans, opt[S][q] + opt[((1<<t)-1)-S][q]);
26
27
       return ans;
28 }
```

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