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1 準備

1.1 init.el

linum は emacs24 のみ

```
; key
(keyboard-translate ?\C-h ?\C-?)
(global-set-key "\M-g" 'goto-line)

; tab
(setq-default indent-tabs-mode nil)
(setq-default tab-width 4)
(setq indent-line-function 'insert-tab)

; line number
(global-linum-mode t)
(setq linum-format "%4d ")
```

1.2 tpl.cpp

```
#include <bits/stdc++.h>
using namespace std;
4 #define rep(i,n) repi(i,0,n)
   #define repi(i,a,b) for(int i=int(a);i<int(b);++i)</pre>
   #define repit(it,u) for(auto it=begin(u);it!=end(u);++it)
   #define all(u) begin(u),end(u)
   #define uniq(u) (u).erase(unique(all(u)),end(u))
   #define 11 long
#define long int64_t
   #define mp make_pair
12
   #define pb push_back
   #define eb emplace_back
14
   bool input()
15
16
17
       return true;
18
19
20
   void solve()
21
22
23
24
25
   int main()
26
27
        cin.tie(0);
28
        ios_base::sync_with_stdio(false);
29
30
        while (input()) solve();
31
```

2 文字列

2.1 マッチング

2.1.1 複数文字列マッチング (Aho-Corasick 法)

O(N+M)

```
const int C = 128:
2
   struct pma_node {
       pma_node *next[C]; // use next[0] as failure link
       vector<int> match;
       pma_node() { fill(next, next + C, (pma_node *) NULL); }
        "pma_node() { rep(i, C) if (next[i] != NULL) delete next[i]; }
   pma node *construct pma(const vector<string>& pat) {
10
       pma_node *const root = new pma_node();
11
       root->next[0] = root:
12
13
       // construct trie
       rep(i, pat.size()) {
14
            const string& s = pat[i];
15
            pma_node *now = root;
16
17
            for (const char c : s) {
                if (now->next[int(c)] == NULL) now->next[int(c)] = new pma_node();
18
                now = now->next[int(c)];
19
20
            now->match.pb(i);
21
22
       // make failure links by BFS
23
       queue<pma_node *> q;
24
       repi(i, 1, C) {
25
            if (root->next[i] == NULL) root->next[i] = root;
26
27
                root->next[i]->next[0] = root;
28
                q.push(root->next[i]);
29
30
31
       while (not q.empty()) {
32
            auto now = q.front();
33
            q.pop();
34
            repi(i, 1, C) if (now->next[i] != NULL) {
35
                auto next = now->next[0];
36
37
                while (next->next[i] == NULL) next = next->next[0];
                now->next[i]->next[0] = next->next[i];
38
                vector<int> tmp;
39
                set_union(all(now->next[i]->match), all(next->next[i]->match), back_inserter
40
                now->next[i]->match = tmp;
42
                q.push(now->next[i]);
43
44
45
       return root;
46
47
   void match(pma_node*& now, const string s, vector<int>& ret) {
48
49
       for (const char c : s) {
            while (now->next[int(c)] == NULL) now = now->next[0];
50
51
            now = now->next[int(c)];
52
            for (const int e : now->match) ret[e] = true;
53
   }
54
```

2.2 Suffix Array

find_string(): $O(|T|\log|S|)$ S 中に T が含まれないなら-1, 含まれるならその先頭. LCS(): O(|S+T|)最長共通部分文字列. (先頭、長さ) を返す.

```
const int MAX N = 1000000:
   int n. k:
   int rnk[MAX_N+1], tmp[MAX_N+1], sa[MAX_N+1], lcp[MAX_N+1];
   bool compare_sa(int i, int j) {
     if(rnk[i] != rnk[j]) return rnk[i] < rnk[j];</pre>
       int ri = i + k \leq n ? rnk[i+k] : -1;
       int rj = j + k <= n ? rnk[j+k] : -1;
10
       return ri < ri:
     }
11
   }
12
13
14
   void construct_sa(string S, int *sa) {
     n = S.length():
15
     for(int i = 0; i \le n; i++) {
16
17
       sa[i] = i;
       rnk[i] = i < n ? S[i] : -1;
18
19
     for (k = 1; k \le n; k*=2) {
20
21
       sort(sa, sa+n+1, compare_sa);
        tmp[sa[0]] = 0;
22
       for(int i = 1; i <= n; i++) {
23
         tmp[sa[i]] = tmp[sa[i-1]] + (compare_sa(sa[i-1], sa[i]) ? 1 : 0);
24
25
26
       for(int i = 0; i \le n; i++) {
         rnk[i] = tmp[i];
27
28
29
   }
30
31
   void construct_lcp(string S, int *sa, int *lcp) {
     int n = S.length();
     for(int i = 0; i \le n; i++) rnk[sa[i]] = i;
34
     int h = 0;
35
36
     lcp[0] = 0;
37
     for(int i = 0; i < n; i++) {
       int j = sa[rnk[i] - 1];
38
39
       if(h > 0) h--;
       for(; j + h < n && i + h < n; h++) {</pre>
         if(S[j+h] != S[i+h]) break;
42
43
       lcp[rnk[i] - 1] = h;
44
45
   //==========================//
   // 文字列検索(蟻本p338 改) O(|T|log|S|)
   // S中にTが含まれないなら -1. 含まれるならその先頭
   int find_string(string S, int *sa, string T) {
     int a = 0, b = S.length();
52
     while(b - a > 1) {
53
       int c = (a + b) / 2;
       if(S.compare(sa[c], T.length(), T) < 0) a = c;
54
55
       else b = c:
56
57
     return (S.compare(sa[b], T.length(), T) == 0)?sa[b]:-1;
58
59
60
   // 最長共通部分文字列(蟻本p341 改) construct_sa以外はO(|S+T|)
61
   // (先頭, 長さ)を返す
62
   pair<int, int> LCS(string S, string T) {
63
     int sl = S.length();
     S += ' \setminus 0' + T;
64
65
     construct_sa(S, sa);
     construct_lcp(S, sa, lcp);
```

```
int len = 0, pos = -1;
for(int i = 0; i < S.length(); i++) {
   if(((sa[i] < sl) != (sa[i+1] < sl)) && (len < lcp[i])) {
      len = lcp[i];
      pos = sa[i];
   }
}
return make_pair(pos, len);
}</pre>
```

3 グラフ

```
struct edge {
2.
       int to; long w;
3
       edge(int to, long w) : to(to), w(w) {}
4
5
   typedef vector<vector<edge> > graph;
   graph rev(const graph& G) {
       const int n = G.size();
       graph ret(n);
10
       rep(i, n) for (const auto& e : G[i]) {
11
            ret[e.to].eb(i, e.w);
12
13
       return ret;
14
```

3.1 強連結成分分解

3.1.1 関節点

O(E)

ある関節点 u がグラフを k 個に分割するとき art には k-1 個の u が含まれる. 不要な場合は unique を忘れないこと.

```
typedef vector<vector<int> > graph;
2
   class articulation {
       const int n;
       graph G;
       int cnt;
       vector<int> num, low, art;
       void dfs(int v) {
            num[v] = low[v] = ++cnt;
            for (int nv : G[v]) {
10
11
                if (num[nv] == 0) {
                    dfs(nv);
12
13
                    low[v] = min(low[v], low[nv]);
14
                    if ((num[v] == 1 and num[nv] != 2) or
                        (num[v] != 1 and low[nv] >= num[v])) {
15
                        art[v] = true;
16
17
                } else {
18
                    low[v] = min(low[v], num[nv]);
19
20
21
22
   public:
23
24
        articulation(const graph& G) : n(G.size()), G(G), cnt(0), num(n), low(n), art(n) {
```

3.1.2 橋

O(V+E)

```
typedef vector<vector<int> > graph;
   class bridge {
        const int n;
        graph G;
        int cnt;
        vector<int> num, low, in;
        stack<int> stk;
        vector<pair<int, int> > brid;
        vector<vector<int> > comp;
10
        void dfs(int v, int p) {
11
            num[v] = low[v] = ++cnt;
12
            stk.push(v), in[v] = true;
13
14
            for (const int nv : G[v]) {
15
                if (num[nv] == 0) {
                    dfs(nv, v);
16
17
                    low[v] = min(low[v], low[nv]);
18
                } else if (nv != p and in[nv]) {
19
                    low[v] = min(low[v], num[nv]);
20
21
            if (low[v] == num[v]) {
22
23
                if (p != n) brid.eb(min(v, p), max(v, p));
24
                comp.eb();
25
                int w;
                do {
26
                    w = stk.top();
28
                    stk.pop(), in[w] = false;
                    comp.back().pb(w);
29
30
                } while (w != v);
31
32
   public:
33
        bridge(const graph& G) : n(G.size()), G(G), cnt(0), num(n), low(n), in(n) {
34
35
            rep(i, n) if (num[i] == 0) dfs(i, n);
36
        vector<pair<int, int> > get() {
37
38
            return brid:
39
40
        vector<vector<int> > components() {
41
            return comp;
42
43
   };
```

3.1.3 強連結成分分解

O(V+E)

```
typedef vector<vector<int> > graph;
class scc {
```

```
const int n;
       graph G;
       int cnt:
       vector<int> num, low, in;
        stack<int> stk;
       vector<vector<int> > comp;
       void dfs(int v) {
10
            num[v] = low[v] = ++cnt;
11
12
            stk.push(v), in[v] = true;
            for (const int nv : G[v]) {
13
                if (num[nv] == 0) {
14
                    dfs(nv);
15
                    low[v] = min(low[v], low[nv]);
16
                } else if (in[nv]) {
17
                    low[v] = min(low[v], num[nv]);
18
19
20
21
            if (low[v] == num[v]) {
                comp.eb();
22
23
                int w;
                do {
24
25
                    w = stk.top();
                    stk.pop(), in[w] = false;
26
27
                    comp.back().pb(w);
28
                } while (w != v);
29
30
31
   public:
32
       scc(const graph& G) : n(G.size()), G(G), cnt(0), num(n), low(n), in(n) {
33
            rep(i, n) if (num[i] == 0) dfs(i);
34
35
        vector<vector<int> > components() {
36
            return comp;
37
38
   };
```

3.2.2 二部マッチング 3.2 フロー

3.2.1 最大流

```
O(EV^2)
```

```
const int inf = 1e9;
   struct edge {
       int to, cap, rev;
       edge(int to, int cap, int rev) : to(to), cap(cap), rev(rev) {}
   typedef vector<vector<edge> > graph;
   void add_edge(graph& G, int from, int to, int cap) {
       G[from].eb(to, cap, G[to].size());
       G[to].eb(from, 0, G[from].size() - 1);
10
11
   }
12
   class max_flow {
13
14
       const int n:
       graph& G;
15
       vector<int> level, iter;
16
       void bfs(int s, int t) {
17
18
            level.assign(n, -1);
19
            queue<int> q;
20
            level[s] = 0, q.push(s);
            while (not q.empty()) {
21
22
                const int v = q.front();
```

23

24

25

26

27

28

29

30

31

32

33

34

35

36

37

38

39

40

41

42

43

44

45

46

47

48

49

50

51

52

53

54

55

56

57 }; q.pop();

}

}

}

}

public:

return 0;

if (v == t) return;

int dfs(int v, int t, int f) {

edge& e = G[v][i];

 $if (d > 0) {$

return d;

max_flow(graph& G) : n(G.size()), G(G) {}

while (bfs(s, t), level[t] \geq 0) {

iter.assign(n, 0);

if (v == t) return f;

}

int calc(int s, int t) {

int ret = 0, d;

return ret;

for (const auto& e : G[v]) {

q.push(e.to);

if (e.cap > 0 and level[e.to] < 0) {</pre>

for (int& i = iter[v]; i < (int) G[v].size(); ++i) {</pre>

if (e.cap > 0 and level[v] < level[e.to]) {</pre>

while ((d = dfs(s, t, inf)) > 0) ret += d;

const int d = dfs(e.to, t, min(f, e.cap));

e.cap -= d, G[e.to][e.rev].cap += d;

level[e.to] = level[v] + 1;

```
O(EV)
   int V:
   vector<int> G[MAX_V];
   int match[MAX_V];
   bool used[MAX_V];
   void add_edge(int u, int v){
        G[u].push_back(v);
        G[v].push_back(u);
10
   bool dfs(int v){
11
12
        used[v] = 1;
13
        rep(i,G[v].size()){
14
            int u = G[v][i], w = match[u];
15
            if(w < 0 || !used[w] && dfs(w)){
16
                match[v] = u;
17
                match[u] = v:
                return 1;
18
19
           }
20
21
        return 0;
22
   }
23
   int bi_matching(){
24
        int res = 0;
25
```

```
26     memset(match, -1, sizeof(match));
27     rep(v,V) if(match[v] < 0){
28          memset(used, 0, sizeof(used));
29          if(dfs(v)) res++;
30     }
31     return res;
32  }</pre>
```

3.2.3 最小費用流

$O(FE \log V)$

```
const int inf = 1e9;
   struct edge {
2
       int to, cap, cost, rev;
        edge(int to, int cap, int cost, int rev): to(to), cap(cap), cost(cost), rev(rev) {}
5
   typedef vector<vector<edge> > graph;
   void add_edge(graph& G, int from, int to, int cap, int cost) {
       G[from].eb(to, cap, cost, G[to].size());
       G[to].eb(from, 0, -cost, G[from].size() - 1);
10
11
12
   int min_cost_flow(graph& G, int s, int t, int f) {
13
       const int n = G.size();
14
       struct state {
15
           int v. d:
16
            state(int v, int d) : v(v), d(d) {}
17
            bool operator <(const state& t) const { return d > t.d; }
18
19
       };
20
21
       int ret = 0;
       vector<int> h(n, 0), dist, prev(n), prev_e(n);
22
       while (f > 0) {
23
24
            dist.assign(n, inf);
            priority_queue<state> q;
25
            dist[s] = 0, q.emplace(s, 0);
26
27
            while (not q.empty()) {
                const int v = q.top().v;
28
29
                const int d = q.top().d;
30
                q.pop();
31
                if (dist[v] <= d) continue;</pre>
                rep(i, G[v].size()) {
                    const edge& e = G[v][i];
33
                    if (e.cap > 0 \text{ and } dist[e.to] > dist[v] + e.cost + h[v] - h[e.to]) {
34
35
                        dist[e.to] = dist[v] + e.cost + h[v] - h[e.to];
                        prev[e.to] = v, prev_e[e.to] = i;
36
37
                        q.emplace(e.to, dist[e.to]);
38
                }
39
            if (dist[t] == inf) return -1;
41
            rep(i, n) h[i] += dist[i];
43
44
45
            for (int v = t: v != s: v = prev[v]) {
                d = min(d, G[prev[v]][prev_e[v]].cap);
46
47
            f -= d. ret += d * h[t]:
48
            for (int v = t; v != s; v = prev[v]) {
49
                edge& e = G[prev[v]][prev_e[v]];
50
51
                e.cap -= d, G[v][e.rev].cap += d;
52
53
```

```
54 return ret;
55 }
```

3.3 木

3.3.1 木の直径

ある点(どこでもよい)から一番遠い点 a を求める. 点 a から一番遠い点までの距離がその木の直径になる.

3.3.2 最小全域木

```
#include "disjoint set.cpp"
   #include "graph.cpp"
   struct mst edge {
       int u, v; long w;
        mst_edge(int u, int v, long w) : u(u), v(v), w(w) {}
        bool operator <(const mst_edge& t) const { return w < t.w; }</pre>
7
        bool operator >(const mst_edge& t) const { return w > t.w; }
9
   };
10
   graph kruskal(const graph& G) {
11
        const int n = G.size();
12
        vector<mst_edge> E;
13
        rep(i, n) for (const auto& e : G[i]) {
14
            if (i < e.to) E.eb(i, e.to, e.w);</pre>
15
16
        sort(all(E));
17
18
        graph T(n);
19
        disjoint_set uf(n);
20
        for (const auto& e : E) {
21
            if (not uf.same(e.u, e.v)) {
22
23
                T[e.u].eb(e.v, e.w);
24
                T[e.v].eb(e.u, e.w);
25
                uf.merge(e.u, e.v);
26
27
28
        return T;
29
   graph prim(const vector<vector<long> >& A. int s = 0) {
        const int n = A.size();
        graph T(n);
33
34
        vector<int> done(n):
35
        priority_queue<mst_edge, vector<mst_edge>, greater<mst_edge> > g;
        q.emplace(-1, s, 0);
36
37
        while (not q.empty()) {
            const auto e = q.top();
38
39
            q.pop();
40
            if (done[e.v]) continue:
            done[e.v] = 1;
41
42
            if (e.u >= 0) {
43
                T[e.u].eb(e.v, e.w);
44
                T[e.v].eb(e.u, e.w);
45
            rep(i, n) if (not done[i]) {
46
47
                q.emplace(e.v, i, A[e.v][i]);
48
49
50
        return T;
```

1 }

3.3.3 最小シュタイナー木

 $O(4^{|T|}V)$

g は無向グラフの隣接行列. T は使いたい頂点の集合.

```
int minimum_steiner_tree(vi &T, vvi &g){
       int n = g.size(), t = T.size();
       if(t <= 1) return 0;
       vvi d(g); // all-pair shortest
       rep(k,n)rep(i,n)rep(j,n) //Warshall Floyd
            d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
       int opt[1 << t][n];</pre>
       rep(S,1 << t) rep(x,n)
10
            opt[S][x] = INF;
11
12
       rep(p,t) rep(q,n) // trivial case
13
            opt[1 << p][q] = d[T[p]][q];
14
       repi(S,1,1<<t){ // DP step
15
            if(!(S & (S-1))) continue;
16
           rep(p,n) rep(E,S)
17
                if((E \mid S) == S)
18
                    opt[S][p] = min(opt[S][p], opt[E][p] + opt[S-E][p]);
19
20
            rep(p,n) rep(q,n)
21
                opt[S][p] = min(opt[S][p], opt[S][q] + d[p][q]);
       }
22
23
       int ans = INF;
24
       rep(S,1 << t) rep(q,n)
25
            ans = min(ans, opt[S][q] + opt[((1<<t)-1)-S][q]);
26
27
        return ans;
28 }
```

3.4 包除原理

3.4.1 彩色数

 $O(2^VV)$

N[i] := i と隣接する頂点の集合 (i も含む)

```
const int MAX_V=16;
   const int mod = 10009;
   int N[MAX_V], I[1<<MAX_V], V;</pre>
   inline int mpow(int a, int k){ return k==0? 1: k%2? a*mpow(a,k-1)%mod: mpow(a*a%mod,k
        /2);}
   bool can(int k){
       int res = 0:
       rep(S, 1<<V){
            if( builtin popcountl1(S)%2) res -= mpow(I[S], k):
            else res += mpow(I[S],k);
10
11
       return (res%mod+mod)%mod:
12
13
14
15
   int color_number(){
       memset(I, 0, sizeof(I));
16
17
       I[0] = 1;
```

```
18
        repi(S,1,1<<V){
19
            int v = 0;
20
            while(!(S&(1<<v))) v++;
21
            I[S] = I[S-(1 << v)] + I[S&(~N[v])];
22
23
        int 1b = 0, ub = V, mid;
24
        while(ub-lb>1){
25
            mid = (1b+ub)/2:
26
            if(can(mid)) ub = mid;
            else lb = mid;
27
28
29
       return ub;
30
```

4 数学

4.1 整数

4.1.1 剰余

```
// (x, y) s.t. a x + b y = gcd(a, b)
   long extgcd(long a, long b, long& x, long& y) {
        long g = a; x = 1, y = 0;
        if (b != 0) g = extgcd(b, a % b, y, x), y -= (a / b) * x;
5
       return g;
   // repi(i, 2, n) mod_inv[i] = mod_inv[m % i] * (m - m / i) % m
   long mod_inv(long a, long m) {
        long x, y;
11
        if (extgcd(a, m, x, y) != 1) return 0;
        return (x % m + m) % m;
12
13
14
   // a mod p where n! = a p^e in O(log_p n)
   long mod_fact(long n, long p, long& e) {
       const int P = 1000010;
        static long fact[P] = {1};
19
        static bool done = false:
       if (not done) {
20
21
            repi(i, 1, P) fact[i] = fact[i - 1] * i % p;
22
            done = true:
23
24
        e = 0:
25
        if (n == 0) return 1;
        long ret = mod_fact(n / p, p, e);
26
        e += n / p;
28
        if (n / p % 2) return ret * (p - fact[n % p]) % p;
        return ret * fact[n % p] % p;
29
30
31
   // nCk mod p
   long mod_binom(long n, long k, long p) {
       if (k < 0 \text{ or } n < k) \text{ return } 0;
35
       long e1. e2. e3:
       long a1 = mod_fact(n, p, e1);
37
       long a2 = mod_fact(k, p, e2);
        long a3 = mod_fact(n - k, p, e3);
        if (e1 > e2 + e3) return 0;
39
        return a1 * mod_inv(a2 * a3 % p, p) % p;
41
   }
43 // a^b mod m
```

```
44  long mod_pow(long a, long b, long m) {
45    long ret = 1;
46    if (b & 1) ret = ret * a % m;
48    a = a * a % m;
49    } while (b >>= 1);
50    return ret;
51 }
```

4.1.2 カタラン数

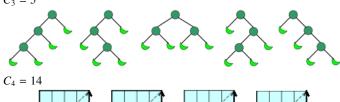
n < 16 程度が限度. n > 1 について以下が成り立つ.

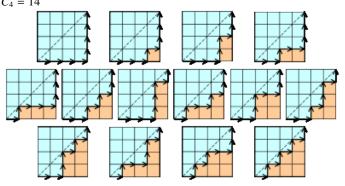
$$C_n = \frac{1}{n+1} {2n \choose n}$$
$$= {2n \choose n} - {2n \choose n-1}$$

n が十分大きいとき、カタラン数は以下に近似できる.

$$C_n = \frac{4^n}{n^{3/2} \sqrt{\pi}}$$

() を正しく並べる方法, 二分木, 格子状の経路の数え上げ, 平面グラフの交差などに使われる. $C_3=5$





4.2 多項式

FFT は基本定数重めなので TLE に注意する.

4.2.1 FFT(complex)

 $O(N \log N)$

複素数を用いた FFT. 変換する vector のサイズは2の冪乗にすること.

```
typedef complex<double> cd;
   vector<cd> fft(vector<cd> f, bool inv){
        int n, N = f.size();
        for(n=0;;n++) if(N == (1 << n)) break;
        rep(m,N){
            int m2 = 0:
            rep(i,n) if(m&(1<<i)) m2 |= (1<<(n-1-i));
            if(m < m2) swap(f[m], f[m2]);</pre>
10
        for(int t=1;t<N;t*=2){</pre>
11
12
            double theta = acos(-1.0) / t;
13
            cd w(cos(theta), sin(theta));
14
            if(inv) w = cd(cos(theta), -sin(theta));
            for(int i=0;i<N;i+=2*t){</pre>
15
                 cd power(1.0, 0.0);
                 rep(j,t){
17
                     cd tmp1 = f[i+j] + f[i+t+j] * power;
18
19
                     cd tmp2 = f[i+j] - f[i+t+j] * power;
                     f[i+j] = tmp1;
20
21
                     f[i+t+j] = tmp2;
22
                     power = power * w;
23
           }
24
25
26
        if(inv) rep(i,N) f[i] /= N;
27
        return f:
28
```

4.2.2 FFT(modulo)

 $O(N \log N)$

剰余環を用いた FFT(FMT). 変換する vector のサイズは 2 の冪乗にすること. mod は $a*2^e+1$ の形.

```
#include "number_theory.cpp"
   const int mod = 7*17*(1<<23)+1;
   vector<int> fmt(vector<int> f, bool inv){
        int e, N = f.size();
        // assert((N&(N-1))==0 and "f.size() must be power of 2");
        for(e=0;;e++) if(N == (1 << e)) break;
            int m2 = 0;
            rep(i,e) if(m&(1<< i)) m2 |= (1<<(e-1-i));
            if(m < m2) swap(f[m], f[m2]);</pre>
11
12
13
        for(int t=1; t<N; t*=2){</pre>
            int r = pow_mod(3, (mod-1)/(t*2), mod);
14
15
            if(inv) r = mod inverse(r.mod):
            for(int i=0; i<N; i+=2*t){
16
17
                int power = 1;
18
                rep(j,t){
                     int x = f[i+j], y = 1LL*f[i+t+j]*power%mod;
19
                     f[i+j] = (x+y)\%mod;
20
21
                    f[i+t+j] = (x-y+mod)%mod;
22
                    power = 1LL*power*r%mod;
                }
23
```

4.2.3 積 (FMT)

 $O(N \log N)$ poly_mul() が必要.

```
vector<int> poly_mul(vector<int> f, vector<int> g){
   int N = max(f.size(),g.size())*2;
   f.resize(N); g.resize(N);
   f = fmt(f,0); g = fmt(g,0);
   rep(i,N) f[i] = 1LL*f[i]*g[i]%mod;
   f = fmt(f,1);
   return f;
}
```

4.2.4 逆元 (FMT)

 $O(N \log N)$

extgcd(), mod_inverse(), poly_mul(), fmt() が必要.

```
vector<int> poly_inv(const vector<int> &f){
       int N = f.size();
2.
       vector<int> r(1,mod_inv(f[0],mod));
       for(int k = 2; k \le N; k \le 1)
            vector<int> nr = poly_mul(poly_mul(r,r), vector<int>(f.begin(),f.begin()+k));
            nr.resize(k);
            rep(i,k/2) {
                nr[i] = (2*r[i]-nr[i]+mod)%mod;
                nr[i+k/2] = (mod-nr[i+k/2])%mod;
10
           r = nr;
11
12
13
       return r;
14
```

4.2.5 平方根 (FMT)

O(NlogN)

extgcd(), mod_inverse(), poly_inv(), poly_mul(), fmt() が必要.

```
const int inv2 = (mod+1)/2;
vector<int> poly_sqrt(const vector<int> &f) {
   int N = f.size();
   vector<int> s(1,1); // s[0] = sqrt(f[0])
   for(int k = 2; k <= N; k <<= 1) {
        s.resize(k);
        vector<int> ns = poly_mul(poly_inv(s), vector<int>(f.begin(),f.begin()+k));
        ns.resize(k);
        rep(i,k) s[i] = 1LL*(s[i]+ns[i])*inv2%mod;
   }
   return s;
}
```

4.3 行列

```
typedef double number;
    typedef vector<number> vec;
    typedef vector<vec> mat;
   vec mul(const mat& A, const vec& x) {
        const int n = A.size();
        vec b(n):
        rep(i, n) rep(j, A[0].size()) {
           b[i] = A[i][j] * x[j];
10
11
        return b;
12
13
   mat mul(const mat& A, const mat& B) {
14
15
        const int n = A.size();
        const int o = A[0].size();
16
        const int m = B[0].size();
17
        mat C(n, vec(m));
18
19
        rep(i, n) rep(k, o) rep(j, m) {
20
            C[i][j] += A[i][k] * B[k][j];
21
22
        return C;
23
   }
24
25
   mat pow(mat A, long m) {
26
        const int n = A.size();
27
        mat B(n, vec(n));
        rep(i, n) B[i][i] = 1;
28
29
        do {
30
            if (m \& 1) B = mul(B, A);
31
            A = mul(A, A);
32
        } while (m >>= 1);
33
        return B;
34
35
   const number eps = 1e-4;
37
   // determinant; 0(n^3)
   number det(mat A) {
        int n = A.size();
41
        number D = 1;
42
        rep(i,n){
43
            int pivot = i;
44
            repi(j,i+1,n)
45
                if (abs(A[j][i]) > abs(A[pivot][i])) pivot = j;
46
            swap(A[pivot], A[i]);
47
            D *= A[i][i] * (i != pivot ? -1 : 1);
48
            if (abs(A[i][i]) < eps) break;</pre>
49
            repi(j,i+1,n)
                 for (int k=n-1; k>=i; --k)
50
51
                    A[j][k] -= A[i][k] * A[j][i] / A[i][i];
52
53
        return D;
54
   }
55
   // rank: 0(n^3)
   int rank(mat A) {
57
        int n = A.size(), m = A[0].size(), r = 0;
58
        for(int i = 0; i < m and r < n; i++){
59
60
            int pivot = r;
61
            repi(j,r+1,n)
62
                 if (abs(A[j][i]) > abs(A[pivot][i])) pivot = j;
63
            swap(A[pivot], A[r]);
            if (abs(A[r][i]) < eps) continue;</pre>
```

4.3.1 線形方程式の解 (Givens 消去法)

```
O(N^3)
   // Givens elimination; O(n^3)
   typedef double number;
   typedef vector<vector<number> > matrix;
   inline double my_hypot(double x, double y) { return sqrt(x * x + y * y); }
   inline void givens_rotate(number& x, number& y, number c, number s) {
       number u = c * x + s * y, v = -s * x + c * y;
       x = u, y = v;
10
   vector<number> givens(matrix A. vector<number> b) {
11
       const int n = b.size();
12
13
       rep(i, n) repi(j, i + 1, n) {
            const number r = my_hypot(A[i][i], A[j][i]);
14
15
            const number c = A[i][i] / r, s = A[i][i] / r;
            givens_rotate(b[i], b[j], c, s);
16
17
           repi(k, i + 1, n) givens_rotate(A[i][k], A[j][k], c, s);
18
       for (int i = n - 1: i >= 0: --i) {
19
           repi(j, i + 1, n) b[i] -= A[i][j] * b[j];
20
21
           b[i] /= A[i][i];
22
23
       return b;
24 }
```

5 幾何

```
// constants and eps-considered operators
   const double eps = 1e-8; // choose carefully!
   const double pi = acos(-1.0);
   inline bool lt(double a, double b) { return a < b - eps; }</pre>
   inline bool gt(double a, double b) { return lt(b, a); }
   inline bool le(double a, double b) { return !lt(b, a); }
   inline bool ge(double a, double b) { return !lt(a, b); }
   inline bool ne(double a. double b) { return lt(a. b) or lt(b. a): }
   inline bool eq(double a, double b) { return !ne(a, b); }
11
12
   // points and lines
13
14
   typedef complex<double> point:
15
16
   inline double dot (point a. point b) { return real(coni(a) * b): }
17
18
   inline double cross(point a, point b) { return imag(conj(a) * b); }
20 | struct line {
```

```
point a, b;
21
22
       line(point a, point b) : a(a), b(b) {}
23
   };
24
25
    * Here is what ccw(a, b, c) returns:
26
27
               1
28
29
         2 | a 0 b | -2
30
31
            - 1
32
33
    * Note: we can implement intersectPS(p, s) as !ccw(s.a, s.b, p).
34
35
36
   int ccw(point a, point b, point c) {
       b -= a, c -= a;
37
       if (cross(b, c) > eps) return +1;
38
       if (cross(b, c) < eps)</pre>
                               return -1:
39
                                 return +2; // c -- a -- b
40
       if (dot(b, c) < eps)</pre>
       if (lt(norm(b), norm(c))) return -2; // a -- b -- c
41
       return 0:
42
43
   bool intersectLS(const line& 1, const line& s) {
45
       return ccw(1.a, 1.b, s.a) * ccw(1.a, 1.b, s.b) <= 0;
46
   bool intersectSS(const line& s, const line& t) {
47
       return intersectLS(s, t) and intersectLS(t, s);
48
49
50
   bool intersectLL(const line& 1, const line& m) {
       return ne(cross(l.b - l.a, m.b - m.a), 0.0) // not parallel
52
           or eq(cross(l.b - l.a, m.a - l.a), 0.0); // overlap
53
   point crosspointLL(const line& 1, const line& m) {
       double A = cross(l.b - l.a. m.b - m.a):
        double B = cross(1.b - 1.a, m.a - 1.a);
       if (eq(A, 0.0)) and eq(B, 0.0) return m.a; // overlap
       assert(ne(A, 0.0));
                                                   // not parallel
59
       return m.a - B / A * (m.b - m.a);
60
   point proj(const line& 1, point p) {
        double t = dot(1.b - 1.a, p - 1.a) / norm(1.b - 1.a);
       return 1.a + t * (1.b - 1.a);
64
   point reflection(const line& 1, point p) { return 2.0 * proj(1, p) - p; }
65
   // distances (for shortest path)
67
   double distanceLP(const line& 1. point p) { return abs(proi(1. p) - p): }
   double distanceLL(const line& 1, const line& m) {
       return intersectLL(1, m) ? 0.0 : distanceLP(1, m.a);
71
72
   double distanceLS(const line& 1, const line& s) {
73
       return intersectLS(1, s) ? 0.0 : min(distanceLP(1, s.a), distanceLP(1, s.b));
74
75
76
   double distancePS(point p, const line& s) {
77
       point h = proj(s, p);
       return ccw(s.a, s.b, h)? min(abs(s.a - p), abs(s.b - p)): abs(h - p);
78
79
   double distanceSS(const line& s. const line& t) {
       if (intersectSS(s, t)) return 0.0;
82
        return min(min(distancePS(s.a, t), distancePS(s.b, t)),
                   min(distancePS(t.a, s), distancePS(t.b, s)));
83
84
85
86
   // circles
87
```

```
88
    struct circle {
        point o; double r;
89
90
        circle(point o, double r) : o(o), r(r) {}
91
    };
92
    bool intersectCL(const circle& c, const line& 1) {
93
        return le(norm(proj(1, c.o) - c.o), c.r * c.r);
94
95
96
    int intersectCS(const circle& c, const line& s) {
        if (not intersectCL(c, s)) return 0;
97
        double a = abs(s.a - c.o);
98
        double b = abs(s.b - c.o);
99
        if (lt(a, c.r) and lt(b, c.r)) return 0;
100
        if (lt(a, c.r) or lt(b, c.r)) return 1;
101
        return ccw(s.a, s.b, proj(s, c.o)) ? 0 : 2;
102
103
    bool intersectCC(const circle& c, const circle& d) {
104
        double dist = abs(d.o - c.o);
105
        return le(abs(c.r - d.r), dist) and le(dist, c.r + d.r);
106
107
    line crosspointCL(const circle& c, const line& l) {
108
        point h = proj(1, c.o);
109
        double a = sqrt(c.r * c.r - norm(h - c.o));
110
        point d = a * (1.b - 1.a) / abs(1.b - 1.a);
111
112
        return line(h - d, h + d);
113
    line crosspointCC(const circle& c, const circle& d) {
114
        double dist = abs(d.o - c.o), th = arg(d.o - c.o);
115
        double ph = acos((c.r * c.r + dist * dist - d.r * d.r) / (2.0 * c.r * dist));
116
117
        return line(c.o + polar(c.r, th - ph), c.o + polar(c.r, th + ph));
118
119
    line tangent(const circle& c, double th) {
120
121
        point h = c.o + polar(c.r, th);
        point d = polar(c.r, th) * point(0, 1);
122
123
        return line(h - d, h + d);
124
    vector<line> common_tangents(const circle& c, const circle& d) {
125
126
        vector<line> ret;
127
        double dist = abs(d.o - c.o), th = arg(d.o - c.o);
128
        if (abs(c.r - d.r) < dist) { // outer</pre>
            double ph = acos((c.r - d.r) / dist);
129
            ret.pb(tangent(c, th - ph));
130
            ret.pb(tangent(c, th + ph));
131
132
133
        if (abs(c.r + d.r) < dist) { // inner}
            double ph = acos((c.r + d.r) / dist);
134
135
            ret.pb(tangent(c, th - ph));
136
            ret.pb(tangent(c, th + ph));
137
        return ret;
138
139
    pair<circle, circle> tangent_circles(const line& 1, const line& m, double r) {
140
        double th = arg(m.b - m.a) - arg(1.b - 1.a);
141
        double ph = (arg(m.b - m.a) + arg(1.b - 1.a)) / 2.0;
142
143
        point p = crosspointLL(1, m);
144
        point d = polar(r / sin(th / 2.0), ph);
        return mp(circle(p - d, r), circle(p + d, r));
145
146
147
    line bisector(point a. point b):
    circle circum_circle(point a, point b, point c) {
148
149
        point o = crosspointLL(bisector(a, b), bisector(a, c));
        return circle(o, abs(a - o));
150
151
    }
152
    // polygons
153
154
```

```
155
   typedef vector<point> polygon;
156
    double area(const polygon& g) {
157
        double ret = 0.0;
158
159
        int j = g.size() - 1;
160
        rep(i, g.size()) {
            ret += cross(g[j], g[i]), j = i;
161
162
163
        return ret / 2.0;
164
165
    point centroid(const polygon& g) {
        if (g.size() == 1) return g[0];
166
        if (g.size() == 2) return (g[0] + g[1]) / 2.0;
167
        point ret = 0.0;
168
        int j = g.size() - 1;
169
170
        rep(i, g.size()) {
171
            ret += cross(g[j], g[i]) * (g[j] + g[i]), j = i;
172
173
        return ret / area(g) / 6.0;
174
   line bisector(point a, point b) {
175
        point m = (a + b) / 2.0;
176
        return line(m, m + (b - a) * point(0, 1));
177
178
179
    polygon convex_cut(const polygon& g, const line& l) {
180
        polygon ret;
181
         int j = g.size() - 1;
182
        rep(i, g.size()) {
183
             if (ccw(l.a, l.b, g[j]) != -1) ret.pb(g[j]);
184
             if (intersectLS(1, line(g[j], g[i]))) ret.pb(crosspointLL(1, line(g[j], g[i])));
185
            j = i;
186
        return ret:
187
188
    polygon voronoi_cell(polygon g, const vector<point>& v, int k) {
189
190
        rep(i, v.size()) if (i != k) {
            q = convex_cut(q, bisector(v[i], v[k]));
191
192
193
        return q;
194
```

5.1 凸包

```
#include "geometry.cpp"
    namespace std {
        bool operator <(const point& a, const point& b) {</pre>
            return ne(real(a), real(b)) ? lt(real(a), real(b)) : lt(imag(a), imag(b));
   }
    polygon convex_hull(vector<point> v) {
        const int n = v.size();
11
        sort(all(v)):
        polygon ret(2 * n);
12
13
        int k = 0:
        for (int i = 0; i < n; ret[k++] = v[i++]) {
14
15
            while (k \ge 2 \text{ and } ccw(ret[k - 2], ret[k - 1], v[i]) \le 0) --k;
16
        for (int i = n - 2, t = k + 1; i >= 0; ret[k++] = v[i--]) {
17
            while (k \ge t \text{ and } ccw(ret[k - 2]. ret[k - 1]. v[i]) \le 0) --k:
18
19
        ret.resize(k - 1);
20
21
        return ret:
```

2 }

6 データ構造

6.1 Union-Find 木

```
#include "macro.cpp"
   class disjoint_set {
       vector<int> p;
       int root(int i) { return p[i] >= 0 ? p[i] = root(p[i]) : i; }
       disjoint_set(int n) : p(n, -1) {}
       bool same(int i, int j) { return root(i) == root(j); }
       int size(int i) { return -p[root(i)]; }
       void merge(int i, int j) {
10
11
           i = root(i), j = root(j);
           if (i == j) return;
12
           if (p[i] > p[j]) swap(i, j);
13
           p[i] += p[j], p[j] = i;
14
15
   };
16
```

6.2 赤黒木

```
template < class T> class rbtree {
   public:
        enum COL { BLACK, RED,};
       struct node {
           T val;
            int color;
            int rnk, size;
            node *left, *right;
            node(T v) : val(v), color(BLACK), rnk(0), size(1) {
11
12
                left = right = NULL;
13
14
            node(node *1, node *r, int c) : color(c) {
15
                left = 1:
16
                right = r;
                update();
17
18
            void update() {
19
                rnk = max((left? left->rnk+(left->color==BLACK): 0),
20
21
                          (right? right->rnk+(right->color==BLACK): 0));
                size = (left? left->size: 0)+(right? right->size: 0)+(!left and !right);
22
23
24
       };
25
26
       node *root:
27
       rbtree() { root = NULL;}
28
       rbtree(T val) { root = new_node(val);}
29
30
       node *new node(T v) { return new node(v):}
31
       node *new_node(node *1, node *r, int c) { return new node(1,r,c);}
32
33
34
       node *right_rotate(node *v) {
```

```
node *w = v->left;
    v \rightarrow left = w \rightarrow right;
    w->right = v:
    v->left->update();
    v->update();
    w->right->update();
    v \rightarrow color = RED;
    w->color = BLACK:
    return w;
node *left_rotate(node *v) {
    node *w = v->right;
    v->right = w->left;
    w \rightarrow left = v;
    v->right->update();
    v->update();
    w->left->update();
    v \rightarrow color = RED;
    w->color = BLACK;
    return w:
}
node *merge_sub(node *u, node *v) {
    if(u->rnk < v->rnk) {
         node *w = merge_sub(u,v->left);
         v \rightarrow left = w;
         v->update();
         if(v->color == BLACK and w->color == RED and w->left->color == RED) {
             if(v->right->color == BLACK) return right_rotate(v);
             else {
                 v \rightarrow color = RED;
                 v->right->color = BLACK;
                 w->color = BLACK;
                  return v:
             }
         else return v;
    else if(u->rnk > v->rnk) {
         node *w = merge_sub(u->right,v);
        u \rightarrow right = w;
         u->update();
         if(u->color == BLACK and w->color == RED and w->right->color == RED) {
             if(u->left->color == BLACK) return left_rotate(u);
             else {
                 u \rightarrow color = RED;
                 u->left->color = BLACK;
                 w->color = BLACK:
                 return u;
             }
         else return u;
    else return new_node(u,v,RED);
node *merge(node *u, node *v) {
    if(!u) return v;
    if(!v) return u:
    u = merge_sub(u,v);
    u \rightarrow color = BLACK;
    return u;
pair<node*, node*> split(node *v, int k) {
    if(!k) return pair<node*,node*>(NULL,v);
```

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```
102
            if(k == v->size) return pair<node*,node*>(v,NULL);
            if(k < v->left->size) {
103
                 auto p = split(v->left.k):
104
                 return pair<node*,node*>(p.first,merge(p.second,v->right));
105
106
             else if(k > v->left->size) {
107
                 auto p = split(v->right,k-v->left->size);
108
                 return pair<node*,node*>(merge(v->left,p.first),p.second);
109
110
             else return pair<node*,node*>(v->left,v->right);
111
        }
112
113
        // insert val at k
114
        node *insert(T val, int k) { return insert(new_node(val),k);}
115
        // insert tree v at k
116
        node *insert(node *v, int k) {
117
            auto p = split(root,k);
118
119
             return root = merge(merge(p.first,v),p.second);
        }
120
121
        // delete at k
122
        node *erase(int k) {
123
            auto p = split(root,k+1);
124
            return root = merge(split(p.first,k).first, p.second);
125
126
        }
127
128
        node *build(const vector<T> &vs) {
            if(!vs.size()) return NULL;
129
130
            if((int)vs.size() == 1) return new_node(vs[0]);
131
            int m = vs.size()/2;
             return merge(build(vector<T>(begin(vs), begin(vs)+m)),
132
133
                          build(vector<T>(begin(vs)+m,end(vs))));
134
135
        int size() { return root->size;}
136
137
        void get(vector<T> &vs) { get(root,vs);}
138
        void get(node *v, vector<T> &vs) {
139
140
            if(!v->left and !v->right) vs.push_back(v->val);
141
142
                 if(v->left) get(v->left,vs);
                 if(v->right) get(v->right,vs);
143
144
145
        }
146
147
        node *push_back(T val) {
             node *v = new_node(val);
148
149
             return root = merge(root,v);
150
151
    };
```

6.3 永続赤黒木

```
//const int MAX = 15000000, BOUND = 14000000;
template<class T> class prbtree {
public:
    enum COL { BLACK, RED,};
    struct node {
        T val;
        int color;
        int rnk, size;
        node *left, *right;
        node(){}
```

```
node(T v) : val(v), color(BLACK), rnk(0), size(1) {
        left = right = NULL;
    node(node *1, node *r, int c) : color(c) {
        left = 1:
        right = r;
        rnk = max((1? 1->rnk+(1->color==BLACK): 0),
                  (r? r->rnk+(r->color==BLACK): 0));
        size = !1 and !r? 1: !1? r->size: !r? r->size: 1->size+r->size;
   }
};
node *root;
          node nodes[MAX]:
//
          int called:
prbtree() {
    root = NULL;
    // called = 0;
prbtree(T val) {
    root = new_node(val);
    // called = 0;
// node *new_node(T v) { return &(nodes[called++] = node(v));}
// node *new_node(node *1, node *r, int c) { return &(nodes[called++] = node(1,r,c
node *new_node(T v) { return new node(v);}
node *new_node(node *1, node *r, int c) { return new node(1,r,c);}
node *merge_sub(node *u, node *v) {
   if(u->rnk < v->rnk) {
        node *w = merge_sub(u,v->left);
        if(v->color == BLACK and w->color == RED and w->left->color == RED){
            if(v->right->color == BLACK) return new_node(w->left,new_node(w->right,
                 v->right.RED).BLACK):
            else return new_node(new_node(w->left,w->right,BLACK),new_node(v->right
                 ->left,v->right->right,BLACK),RED);
        else return new_node(w,v->right,v->color);
    else if(u->rnk > v->rnk) {
        node *w = merge_sub(u->right,v);
        if(u->color == BLACK and w->color == RED and w->right->color == RED){
            if(u->left->color == BLACK) return new_node(new_node(u->left,w->left,
                 RED),w->right,BLACK);
            else return new node(new node(u->left->left.u->left->right.BLACK).
                 new_node(w->left,w->right,BLACK),RED);
        else return new_node(u->left,w,u->color);
    else return new node(u.v.RED):
}
node *merge(node *u, node *v) {
    if(!u) return v;
   if(!v) return u;
   u = merge sub(u.v):
   if(u->color == RED) return new_node(u->left,u->right,BLACK);
    return u:
}
pair<node*,node*> split(node *v, int k) {
    if(!k) return pair<node*,node*>(NULL,v);
    if(k == v->size) return pair<node*,node*>(v,NULL);
```

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```
74
            if(k < v->left->size) {
                auto p = split(v->left,k);
75
76
                return pair<node*,node*>(p.first,merge(p.second,v->right));
77
            else if(k > v->left->size) {
78
                auto p = split(v->right,k-v->left->size);
79
                return pair<node*,node*>(merge(v->left,p.first),p.second);
80
81
82
            else return pair<node*,node*>(v->left,v->right);
        }
83
84
        node *build(const vector<T> &vs) {
85
            if(!vs.size()) return NULL;
86
            if((int)vs.size() == 1) return new_node(vs[0]);
87
            int m = vs.size()/2;
88
            return merge(build(vector<T>(begin(vs),begin(vs)+m)), build(vector<T>(begin(vs)+
89
                 m,end(vs)));
90
        }
91
        int size() { return root->size;}
92
93
        void get(vector<T> &vs) { get(root,vs);}
94
        void get(node *v, vector<T> &vs) {
95
            if(!v->left and !v->right) vs.push_back(v->val);
96
97
                if(v->left) get(v->left,vs);
98
99
                if(v->right) get(v->right, vs);
100
        }
101
102
        node *push_back(T val) {
103
104
            node *v = new_node(val);
105
            return root = merge(root, v);
106
107
        // insert leaf at k
108
        node *insert(int k, T val) {
109
110
            return insert(new_node(val), k);
111
112
113
        // insert tree v at k
114
        node *insert(node *v, int k) {
115
            auto p = split(root,k);
            return root = merge(merge(p.first,v),p.second);
116
        }
117
118
        // copy [1,r)
119
120
        node *copy(int 1, int r) {
121
            return split(split(root, 1).second, r-1).first;
122
123
        // copy and insert [1,r) at k
        node *copy_paste(int 1, int r, int k) {
124
125
            return insert(copy(l,r),k);
126
127 };
```