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# 1 準備

## 1.1 Caps Lock を Control に変更

2つ

1. 変更

```
setxkbmap -option ctrl:nocaps;
```

元に戻す

```
setxkbmap -option;
```

2. 上でダメな場合

```
xmodmap -e 'remove Lock = Caps_Lock';
xmodmap -e 'add Control = Caps_Lock';
xmodmap -e 'keysym Caps_Lock = Control_L';
```

#### 1.2 init.el

linum は emacs24 のみ

```
(keyboard-translate ?\C-h ?\C-?)
(global-linum-mode t)
(setq linum-format "%4d ")
```

## 1.3 tpl.cpp

```
#include <bits/stdc++.h>
   using namespace std;
   #define rep(i,n) repi(i,0,n)
   #define repi(i,a,b) for(int i=(int)(a);i<(int)(b);++i)
   #define all(u) begin(u), end(u)
   #define long int64_t
   #define mp make_pair
    #define pb push_back
    void input() {
10
11
12
   void solve() {
13
14
15
   int main() {
16
17
        cin.tie(0);
        ios_base::sync_with_stdio(false);
18
       input(); // multiple testcases?
19
        solve();
20
21
```

## 1.4 get input

```
wget -r http://(url of sample input)
```

## 1.5 alias

```
alias g++='g++ -g -02 -std=gnu++0x -Wl,-stack_size,64000000';
alias emacs='emacs -nw';
```

## 2 文字列

## 2.1 マッチング

## 2.1.1 複数文字列マッチング (Aho-Corasick 法)

O(N+M)

```
const int C = 128;
   struct pma_node {
        pma_node *next[C]; // use next[0] as failure link
        vector<int> match;
        pma_node() { fill(next, next + C, (pma_node *) NULL); }
        "pma_node() { rep(i, C) if (next[i] != NULL) delete next[i]; }
   pma_node *construct_pma(const vector<string>& pat) {
        pma_node *const root = new pma_node();
        root->next[0] = root:
10
        // construct trie
11
        rep(i, pat.size()) {
12
            const string& s = pat[i];
13
            pma_node *now = root;
14
            for (const char c : s) {
15
                if (now->next[int(c)] == NULL) now->next[int(c)] = new pma_node();
16
                now = now->next[int(c)];
17
18
            now->match.pb(i);
19
20
        // make failure links with BFS
21
        queue < pma_node *> q;
22
        repi(i, 1, C) {
23
            if (root->next[i] == NULL) root->next[i] = root;
24
            else {
25
26
                root->next[i]->next[0] = root;
27
                q.push(root->next[i]);
28
29
        while (not q.empty()) {
30
            auto now = q.front();
31
32
            q.pop();
            repi(i, 1, C) if (now->next[i] != NULL) {
33
34
                auto next = now->next[0];
35
                while (next->next[i] == NULL) next = next->next[0];
                now->next[i]->next[0] = next->next[i];
36
                vector<int> tmp;
37
                set_union(all(now->next[i]->match), all(next->next[i]->match), back_inserter
38
                     (tmp));
                now->next[i]->match = tmp;
39
                q.push(now->next[i]);
40
```

```
42
       }
       return root;
43
44
45
   void match(pma_node*& now, const string s, vector<int>& ret) {
       for (const char c : s) {
46
            while (now->next[int(c)] == NULL) now = now->next[0];
47
48
            now = now->next[int(c)];
            for (const int e : now->match) ret[e] = true;
49
50
51
   }
```

## 2.2 Suffix Array

```
find_string(): O(|T|\log|S|)
S 中に T が含まれないなら-1, 含まれるならその先頭.
LCS(): O(|S+T|)
最長共通部分文字列. (先頭, 長さ) を返す.
```

```
// verified: http://www.spoj.com/problems/{SARRAY,SUBLEX}/
   int n, k;
   vector<int> rnk, tmp, sa, lcp;
   bool compare_sa(int i, int j) {
       if (rnk[i] != rnk[j]) return rnk[i] < rnk[j];</pre>
            int ri = i+k <= n ? rnk[i+k] : -1;</pre>
            int rj = j+k \le n ? rnk[j+k] : -1;
            return ri < rj;
10
       }
11
   void construct_sa(const string &s) {
12
       n = s.size():
13
14
        rnk.assign(n+1, 0);
15
        tmp.assign(n+1, 0);
        sa.assign(n+1, 0);
17
        lcp.assign(n+1, 0);
        rep(i,n+1) {
19
            sa[i] = i;
            rnk[i] = i < n ? s[i] : -1;
20
21
        for (k = 1; k \le n; k *= 2) {
22
23
            sort(sa.begin(), sa.end(), compare_sa);
24
            tmp[sa[0]] = 0:
25
            repi(i,1,n+1) tmp[sa[i]] = tmp[sa[i-1]] + (compare_sa(sa[i-1], sa[i]) ? 1 : 0);
26
            rep(i,n+1) rnk[i] = tmp[i];
27
28
    void construct lcp(const string &s) {
29
        rep(i,n+1) rnk[sa[i]] = i;
30
31
        int h = lcp[0] = 0;
32
        rep(i,n) {
33
            int j = sa[rnk[i] - 1];
34
            if (h > 0) h--:
            for (; j+h < n and i+h < n; h++) {
35
                if (s[j+h] != s[i+h]) break;
36
37
            lcp[rnk[i] - 1] = h;
38
39
40
   }
```

# 2.3 Z-algorithm

```
s, s[i:] の最長共通部分列の長さ
```

```
vector<int> lcp0(const string& s) {
        const int n = s.length();
2
        vector<int> ret(n);
3
        ret[0] = n;
        for (int i = 1, j = 0, k; i < n; ) {
            while (i+j < n \text{ and } s[i+j] == s[j]) ++j;
            ret[i] = j;
            if (j == 0) { ++i; continue; }
            for (k = 1; i+k < n \text{ and } k+ret[k] < j; ++k) {
                 ret[i+k] = ret[k];
11
12
            i += k, j -= k;
13
14
        return ret;
15
```

## 2.4 回文長 (Manacher)

O(N)

各文字を中心とした時の回文の長さ.

```
偶数長の回文はダミーを挟むことで求められている.
```

```
vector<int> manacher(const string &s) {
    int n = s.size()*2;
    vector<int> rad(n,0);
    for (int i = 0, j = 0, k; i < n; i += k, j = max(j-k, 0)) {
        while (i-j >= 0 && i+j+1 < n && s[(i-j)/2] == s[(i+j+1)/2]) ++j;
        rad[i] = j;
        for (k = 1; i-k >= 0 && rad[i]-k >= 0 && rad[i-k] != rad[i]-k; ++k)
            rad[i+k] = min(rad[i-k], rad[i]-k);
    }
    return rad;
}
```

# 3 グラフ

## 3.1 強連結成分分解

### 3.1.1 関節点

O(E)

ある関節点 u がグラフを k 個に分割するとき art には k-1 個の u が含まれる. 不要な場合は unique を忘れないこと.

```
struct articulation {
        const int n; graph G;
        int cnt;
        vector<int> num, low, is_art;
        void dfs(int v) {
            num[v] = low[v] = ++cnt;
            for (int nv : G[v]) {
                if (num[nv] == 0) {
                     dfs(nv);
                     low[v] = min(low[v], low[nv]);
                    if ((num[v] == 1 \text{ and } num[nv] != 2) \text{ or}
11
                         (num[v] != 1 and low[nv] >= num[v])) {
12
                         is art[v] = true:
13
14
```

#### 3.1.2 橋

12

13

14

15

16

17

18

19

20

21

22

23

24

25

26

27

28

29

30

31

32 };

O(V+E)

```
struct bridge {
    const int n; graph G;
    int cnt;

vector<int> num, low, in;
    stack<int> stk;

vector<pair<int,int> brid;

vector<vector<int> > comp;

void dfs(int v, int p) {
    num[v] = low[v] = ++cnt;
    stk.push(v), in[v] = true;
    for (const int nv : G[v]) {
```

low[v] = min(low[v], low[nv]);

low[v] = min(low[v], num[nv]);

if (p != n) brid.eb(min(v, p), max(v, p));

} else if (nv != p and in[nv]) {

stk.pop(), in[w] = false;

if (num[nv] == 0) {

dfs(nv, v);

if (low[v] == num[v]) {

} while (w != v);

w = stk.top();

comp.back().pb(w);

rep(i, n) if (num[i] == 0) dfs(i, n);

comp.eb();

int w; do {

```
3.1.3 強連結成分分解
```

```
O(V + E)

struct scc {
    const int n;
    graph G;
    int cnt;
    vector<int> num, low, in;
    stack<int> stk;
    vector<vector<int> comp;
    void dfs(int v) {
        num[v] = low[v] = ++cnt;
        stk.push(v), in[v] = true;
    }
}
```

bridge(const graph& G) : n(G.size()), G(G), cnt(0), num(n), low(n), in(n) {

```
for (const int nv : G[v]) {
11
                if (num[nv] == 0) {
12
13
                    dfs(nv);
14
                    low[v] = min(low[v], low[nv]);
15
                } else if (in[nv]) {
                    low[v] = min(low[v], num[nv]);
17
18
            if (low[v] == num[v]) {
19
                comp.eb();
20
21
                int w; do {
22
                    w = stk.top();
                    stk.pop(), in[w] = false;
23
24
                    comp.back().pb(w);
25
                } while (w != v);
           }
26
27
        scc(const graph& G) : n(G.size()), G(G), cnt(0), num(n), low(n), in(n) {
28
            rep(i, n) if (num[i] == 0) dfs(i);
29
30
   };
31
```

## 3.1.4 無向中国人郵便配達問題

 $O(om \log n + o^2 2^o)$ , -O2 で  $o \le 18$  程度が限界

```
long chinesePostman(const graph &g) {
        long total = 0;
2
        vector<int> odds;
3
        rep(u, g.size()) {
            for(auto &e: g[u]) total += e.w;
            if (g[u].size() % 2) odds.push_back(u);
        total /= 2;
        int n = odds.size(), N = 1 << n;</pre>
        int w[n][n]; // make odd vertices graph
10
11
        rep(u,n) {
            int s = odds[u]; // dijkstra's shortest path
12
13
            vector<int> dist(g.size(), 1e9); dist[s] = 0;
14
            vector<int> prev(g.size(), -2);
            priority_queue<edge> Q;
15
            Q.push( edge(-1, s, 0) );
            while (!Q.empty()) {
17
18
                 edge e = Q.top(); Q.pop();
                if (prev[e.to] != -2) continue;
19
                prev[e.to] = e.src;
20
                for(auto &f: g[e.to]) {
22
                    if (dist[f->to] > e.w+f->w) {
23
                         dist[f->to] = e.w+f->w;
                         Q.push(edge(f->src, f->to, e.w+f->w));
24
25
26
27
28
            rep(v,n) w[u][v] = dist[odds[v]];
29
30
        long best[N]: // DP for general matching
31
        rep(S,N) best[S] = INF;
32
        best[0] = 0;
33
34
        for (int S = 0; S < N; ++S)
            for (int i = 0; i < n; ++i)
35
36
                if (!(S&(1<<i)))
                     for (int j = i+1; j < n; ++j)
37
38
                         if (!(S&(1<<j)))</pre>
```

### 3.1.5 全点対間最短路 (Johnson)

 $O(max(VE \log V, V^2))$ 

```
bool shortest_path(const graph &g, vector<vector<int> > &dist, vector<vector<int> > &
       int n = g.size();
       vector<int> h(n+1);
       rep(k,n) rep(i,n) for(auto &e: g[i]) {
            if (h[e.to] > h[e.from] + e->w) {
                h[e.to] = h[e.from] + e->w;
                if (k == n-1) return false; // negative cycle
       dist.assign(n, vector<int>(n, 1e9));
10
11
       prev.assign(n, vector<int>(n, -2));
12
       rep(s, n) {
13
            priority_queue<edge> q;
            q.push(edge(s, s, 0));
14
            while (!q.empty()) {
15
                edge e = q.top(); q.pop();
16
                if (prev[s][e.dst] != -2) continue;
17
18
                prev[s][e.to] = e.from;
                for(auto &f:g[e.to]) {
19
                    if (dist[s][f.to] > e.w + f->w) {
20
                        dist[s][f.to] = e.w + f->w;
21
22
                        q.push(edge(f-.from, f.to, e.w + f->w));
23
24
25
            rep(u, n) dist[s][u] += h[u] - h[s];
26
27
28
29
   vector<int> build_path(const vector<vector<int> >& prev, int s, int t) {
30
       vector<int> path;
31
32
       for (int u = t; u >= 0; u = prev[s][u])
33
            path.push_back(u);
34
        reverse(begin(path), end(path));
35
       return path;
36
```

### 3.1.6 無向グラフの全域最小カット

 $O(V^3)$ 

```
int minimum_cut(const graph &g) {
    int n = g.size();
    vector< vector<int>> h(n, vector<int>(n)); // make adj. matrix
    rep(u,n) for(auto &e: g[u]) h[e.src][e.dst] += e.weight;
    vector<int> V(n); rep(u, n) V[u] = u;

int cut = 1e9;
    for(int m = n; m > 1; m--) {
        vector<int> ws(m, 0);
        int u, v;
    int w;
}
```

```
12
                u = v; v = max_element(ws.begin(), ws.end())-ws.begin();
13
14
                w = ws[v]; ws[v] = -1;
15
                rep(i, m) if (ws[i] \geq 0) ws[i] += h[V[v]][V[i]];
16
            rep(i. m) {
17
18
                h[V[i]][V[u]] += h[V[i]][V[v]];
                h[V[u]][V[i]] += h[V[v]][V[i]];
19
20
21
            V.erase(V.begin()+v);
22
            cut = min(cut, w);
23
24
        return cut;
25
```

#### 3.2 フロー

### 3.2.1 最大流

 $O(EV^2)$ 

```
const int inf = 1e9:
   struct edge {
        int to, cap, rev;
        edge(int to, int cap, int rev) : to(to), cap(cap), rev(rev) {}
   typedef vector<vector<edge> > graph;
   void add_edge(graph& G, int from, int to, int cap) {
        G[from].eb(to, cap, G[to].size());
        G[to].eb(from, 0, G[from].size() - 1);
10
   struct max_flow {
11
        const int n; graph& G;
12
        vector<int> level, iter;
13
        void bfs(int s, int t) {
14
            level.assign(n, -1);
15
            queue<int> q;
16
            level[s] = 0, q.push(s);
17
            while (not q.empty()) {
18
19
                const int v = q.front();
                q.pop();
20
                if (v == t) return;
21
                for (const auto& e : G[v]) {
22
23
                     if (e.cap > 0 and level[e.to] < 0) {</pre>
                         level[e.to] = level[v] + 1;
24
                         q.push(e.to);
25
26
27
28
            }
29
        int dfs(int v, int t, int f) {
30
            if (v == t) return f:
31
            for (int& i = iter[v]; i < (int) G[v].size(); ++i) {</pre>
32
33
                 edge& e = G[v][i];
                if (e.cap > 0 and level[v] < level[e.to]) {</pre>
34
35
                     const int d = dfs(e.to, t, min(f, e.cap));
36
                     if (d > 0) {
37
                         e.cap -= d, G[e.to][e.rev].cap += d;
                         return d;
38
39
                    }
40
            }
41
42
            return 0:
43
```

```
max_flow(graph& G) : n(G.size()), G(G) {}
44
        int calc(int s, int t) {
45
46
            int ret = 0, d;
47
            while (bfs(s, t), level[t] \geq 0) {
                iter.assign(n, 0);
48
                while ((d = dfs(s, t, inf)) > 0) ret += d;
49
50
51
            return ret;
52
   };
53
```

#### 3.2.2 二部マッチング

```
O(EV)
```

```
int V;
   vector<int> G[MAX_V];
   int match[MAX_V];
    bool used[MAX_V];
    void add_edge(int u, int v){
        G[u].push_back(v);
        G[v].push_back(u);
9
10
11
   bool dfs(int v){
        used[v] = 1;
12
13
        rep(i,G[v].size()){
            int u = G[v][i], w = match[u];
14
15
            if(w < 0 \mid | !used[w] && dfs(w)){
                match[v] = u;
16
17
                match[u] = v;
                return 1;
18
19
20
21
        return 0;
22
23
24
    int bi_matching(){
       int res = 0;
25
        memset(match, -1, sizeof(match));
26
27
        rep(v,V) if (match[v] < 0){
28
            memset(used, 0, sizeof(used));
29
            if(dfs(v)) res++;
30
31
        return res;
32 }
```

#### 3.2.3 一般グラフの最大マッチング

```
O(V^3)
```

```
#define rep(i,n) repi(i,0,n)
#define repi(i,a,b) for(int i=(int)(a);i<(int)(b);++i)

#define even(x) (mu[x] == x or phi[mu[x]] != mu[x])
#define out(x) (mu[x] != x and phi[mu[x]] == mu[x] and phi[x] == x)
int maximum_matching(const vector<vector<int>>& G, vector<pair<int,int>>& ret) {
    const int n = G.size();
    vector<int>> mu(n), phi(n), rho(n), done(n);
    rep(v, n) mu[v] = phi[v] = rho[v] = v;
    for (int x = -1; ; ) {
```

```
11
12
                for (x = 0; x < n \text{ and } (done[x] \text{ or } !even(x)); ++x);
13
                if (x == n) break;
14
15
            int y = -1;
            for (int v : G[x]) if (out(v) or (even(v) and rho[v] != rho[x])) y = v;
16
17
            if (y == -1) {
                done[x] = true, x = -1;
18
            } else if (out(y)) {
19
                phi[y] = x;
20
21
            } else {
22
                vector<int> dx(n, -2), dy(n, -2); // x \% 2 --> x >= 0
                 for (int k = 0, w = x; dx[w] < 0; w = k % 2 ? mu[w] : phi[w]) <math>dx[w] = k++;
23
24
                 for (int k = 0, w = y; dy[w] < 0; w = k % 2 ? mu[w] : phi[w]) <math>dy[w] = k++;
25
                bool disjoint = true;
26
                rep(v, n) if (dx[v] >= 0 and dy[v] > 0) disjoint = false;
                if (disjoint) {
27
                    rep(v, n) if (dx[v] \% 2) mu[phi[v]] = v, mu[v] = phi[v];
28
                    rep(v, n) if (dy[v] % 2) mu[phi[v]] = v, mu[v] = phi[v];
29
                    mu[x] = y; mu[y] = x; x = -1;
30
                    rep(v, n) phi[v] = rho[v] = v, done[v] = false;
31
                } else {
32
33
                    int r = x, d = n;
                    rep(v, n) if (dx[v] >= 0 and dy[v] >= 0 and rho[v] == v and d > dx[v])
34
                          = dx[v], r = v;
35
                     rep(v, n) if (dx[v] \le d and dx[v] \% 2 and rho[phi[v]] != r) phi[phi[v]]
                     rep(v, n) if (dy[v] \le d and dy[v] % 2 and rho[phi[v]] != r) phi[phi[v]]
                           = v:
                     if (rho[x] != r) phi[x] = y;
38
                    if (rho[y] != r) phi[y] = x;
39
                    rep(v, n) if (dx[rho[v]] >= 0 or dy[rho[v]] >= 0) rho[v] = r;
40
41
            }
42
43
        ret.clear();
44
        rep(v, n) if (v < mu[v]) ret.emplace_back(v, mu[v]);</pre>
45
        return ret.size();
46
```

#### 3.2.4 最小費用流

#### $O(FE \log V)$

```
const int inf = 1e9;
   struct edge {
        int to, cap, cost, rev;
        edge(int to, int cap, int cost, int rev) : to(to), cap(cap), cost(cost), rev(rev) {}
    typedef vector<vector<edge> > graph;
   void add_edge(graph& G, int from, int to, int cap, int cost) {
        G[from].eb(to, cap, cost, G[to].size());
        G[to].eb(from, 0, -cost, G[from].size() - 1);
10
   int min_cost_flow(graph& G, int s, int t, int f) {
11
12
        const int n = G.size();
13
        struct state {
14
15
            state(int v, int d) : v(v), d(d) {}
            bool operator <(const state& t) const { return d > t.d; }
16
17
18
        int ret = 0;
        vector<int> h(n, 0), dist, prev(n), prev_e(n);
19
20
        while (f > 0) {
21
            dist.assign(n, inf);
```

```
priority queue < state > q:
22
            dist[s] = 0, q.emplace(s, 0);
23
24
            while (not q.empty()) {
25
                const int v = q.top().v;
                const int d = q.top().d;
26
27
                q.pop();
28
                if (dist[v] < d) continue;</pre>
                rep(i, G[v].size()) {
29
30
                     const edge& e = G[v][i];
                     if (e.cap > 0 \text{ and } dist[e.to] > dist[v] + e.cost + h[v] - h[e.to]) {
31
                         dist[e.to] = dist[v] + e.cost + h[v] - h[e.to];
32
33
                         prev[e.to] = v, prev_e[e.to] = i;
                         q.emplace(e.to, dist[e.to]);
34
35
                }
36
37
            if (dist[t] == inf) return -1:
38
            rep(i, n) h[i] += dist[i];
39
            int d = f;
40
            for (int v = t; v != s; v = prev[v]) {
41
                d = min(d, G[prev[v]][prev_e[v]].cap);
42
43
            f -= d, ret += d * h[t];
44
            for (int v = t; v != s; v = prev[v]) {
45
                edge& e = G[prev[v]][prev_e[v]];
46
47
                e.cap -= d, G[v][e.rev].cap <math>+= d;
48
49
50
        return ret;
51
```

### 3.2.5 Gomory-Hu 木

## O(VMAXFLOW)

```
#define RESIDUE(s,t) (capacity[s][t]-flow[s][t])
   graph cutTree(const graph &g) {
        int n = g.size();
       Matrix capacity(n, Array(n)), flow(n, Array(n));
       rep(u,n) for(auto &e: g[u]) capacity[e.from][e.to] += e.w;
       vector<int> p(n), prev;
       vector<int> w(n);
       for (int s = 1; s < n; ++s) {
10
            int t = p[s]; // max-flow(s, t)
            rep(i,n) rep(j,n) flow[i][j] = 0;
11
            int total = 0;
12
            while (1) {
13
14
                queue<int> Q; Q.push(s);
15
                prev.assign(n, -1); prev[s] = s;
                while (!Q.empty() && prev[t] < 0) {</pre>
16
17
                    int u = Q.front(); Q.pop();
                    for(auto &e: g[u]) if (prev[e.to] < 0 && RESIDUE(u, e.to) > 0) {
18
19
                        prev[e.to] = u;
20
                        Q.push(e.to);
21
22
23
                if (prev[t] < 0) goto esc;</pre>
24
                int inc = 1e9;
                for (int j = t; prev[j] != j; j = prev[j])
25
26
                    inc = min(inc, RESIDUE(prev[j], j));
27
                for (int j = t; prev[j] != j; j = prev[j])
28
                    flow[prev[j]][j] += inc, flow[j][prev[j]] -= inc;
                total += inc:
29
30
```

```
esc:w[s] = total: // make tree
31
            rep(u, n) if (u != s \&\& prev[u] != -1 \&\& p[u] == t)
32
33
                p[u] = s;
34
            if (prev[p[t]] != -1)
35
               p[s] = p[t], p[t] = s, w[s] = w[t], w[t] = total;
36
37
        graph T(n); // (s, p[s]) is a tree edge of weight w[s]
        rep(s, n) if (s != p[s]) {
38
           T[ s ].push_back( Edge(s, p[s], w[s]) );
39
40
           T[p[s]].push_back( Edge(p[s], s, w[s]) );
41
42
        return T:
43
44
   // Gomory-Hu tree を用いた最大流 O(n)
45
   int max_flow(const graph &T, int u, int t, int p = -1, int w = 1e9) {
        if (u == t) return w:
        int d = 1e9;
48
        for(auto &e: T[u]) if (e.to != p)
49
            d = min(d, max_flow(T, e.to, t, u, min(w, e.w)));
50
51
        return d;
52
```

- 3.3 木
- 3.3.1 木の直径: double sweep
- 3.3.2 最小全域木

```
struct uedge {
        int u, v; long w;
        uedge(int u, int v, long w) : u(u), v(v), w(w) {}
        bool operator <(const uedge& t) const { return w < t.w; }</pre>
        bool operator >(const uedge& t) const { return w > t.w; }
   graph kruskal(const graph& G) {
        const int n = G.size();
        vector<uedge> E;
        rep(i, n) for (const auto& e : G[i]) {
            if (i < e.to) E.eb(i, e.to, e.w);</pre>
11
12
        sort(all(E));
13
        graph T(n);
14
        disjoint_set uf(n);
15
        for (const auto& e : E) {
16
            if (not uf.same(e.u, e.v)) {
17
18
                T[e.u].eb(e.v, e.w);
19
                T[e.v].eb(e.u, e.w);
20
                uf.merge(e.u, e.v);
21
22
23
        return T;
24
   graph prim(const vector<vector<long> >& A, int s = 0) {
26
        const int n = A.size();
27
        graph T(n);
28
        vector<int> done(n);
        priority_queue<uedge, vector<uedge>, greater<uedge> > q;
29
30
        q.emplace(-1, s, 0);
        while (not q.empty()) {
31
32
            const auto e = q.top(); q.pop();
            if (done[e.v]) continue:
33
34
            done[e.v] = 1;
```

```
if (e.u >= 0) {
35
                T[e.u].eb(e.v, e.w);
36
37
                T[e.v].eb(e.u, e.w);
38
            rep(i, n) if (not done[i]) {
39
                q.emplace(e.v, i, A[e.v][i]);
40
41
42
43
        return T;
44
```

## 3.3.3 最小全域有向木

O(VE)

```
void visit(Graph &h, int v, int s, int r,
               vector<int> &no, vector< vector<int> > &comp,
2
               vector<int> &prev, vector< vector<int> > &next, vector<int> &mcost,
               vector<int> &mark, int &cost, bool &found) {
       const int n = h.size();
       if (mark[v]) {
            vector<int> temp = no;
            found = true:
            do {
10
                cost += mcost[v];
                v = prev[v];
11
                if (v != s) {
12
                    while (comp[v].size() > 0) {
13
                        no[comp[v].back()] = s;
14
                        comp[s].push_back(comp[v].back());
15
                        comp[v].pop_back();
16
17
18
            } while (v != s);
19
            for(auto &j: comp[s]) if (j != r) for(auto &e: h[j])
20
                if (no[e.from] != s) e.w -= mcost[temp[j]];
21
22
23
       mark[v] = true;
       for(auto &i: next[v]) if (no[i] != no[v] && prev[no[i]] == v)
24
            if (!mark[no[i]] || i == s)
25
26
                visit(h, i, s, r, no, comp, prev, next, mcost, mark, cost, found);
27
   int minimum_spanning_arborescence(const graph &g, int r) {
28
       const int n = g.size();
29
30
       graph h(n);
       rep(u,n) for(auto &e: g[u]) h[e.to].push_back(e);
31
32
33
       vector<int> no(n);
34
       vector < vector < int > > comp(n);
35
       rep(u, n) comp[u].push_back(no[u] = u);
36
37
        for (int cost = 0; ;) {
            vector<int> prev(n, -1);
38
39
            vector<int> mcost(n, INF);
40
            rep(j,n) if (j != r) for(auto &e: g[j])
41
42
                if (no[e.from] != no[i])
                    if (e.w < mcost[no[j]])</pre>
43
44
                        mcost[no[j]] = e.w, prev[no[j]] = no[e.from];
45
46
            vector< vector<int> > next(n);
            rep(u,n) if (prev[u] >= 0)
47
                next[prev[u]].push_back(u);
48
49
50
            bool stop = true;
```

```
vector<int> mark(n):
51
            rep(u,n) if (u != r && !mark[u] && !comp[u].empty()) {
52
53
                bool found = false;
54
                visit(h, u, u, r, no, comp, prev, next, mcost, mark, cost, found);
55
                if (found) stop = false;
56
57
            if (stop) {
                rep(u,n) if (prev[u] >= 0) cost += mcost[u];
58
59
                return cost;
60
61
62
   }
```

### 3.3.4 最小シュタイナー木

 $O(4^{|T|}V)$ 

g は無向グラフの隣接行列. T は使いたい頂点の集合.

```
int minimum_steiner_tree(vi &T, vvi &g){
        int n = q.size(), t = T.size();
        if(t <= 1) return 0;
        vvi d(q); // all-pair shortest
        rep(k,n)rep(i,n)rep(j,n) //Warshall Floyd
            d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
        int opt[1 << t][n];</pre>
        rep(S,1 << t) rep(x,n)
            opt[S][x] = INF;
10
11
12
        rep(p,t) rep(q,n) // trivial case
13
            opt[1 << p][q] = d[T[p]][q];</pre>
14
15
        repi(S,1,1<<t){ // DP step
16
            if(!(S & (S-1))) continue;
17
            rep(p,n) rep(E,S)
18
                if((E \mid S) == S)
                    opt[S][p] = min(opt[S][p], opt[E][p] + opt[S-E][p]);
19
20
            rep(p,n) rep(q,n)
21
                opt[S][p] = min(opt[S][p], opt[S][q] + d[p][q]);
22
23
24
        int ans = INF;
        rep(S,1 << t) rep(q,n)
25
            ans = min(ans, opt[S][q] + opt[((1<<t)-1)-S][q]);
26
27
        return ans;
28
```

### 3.3.5 木の同型性判定

順序付き O(n), 順序なし  $O(n \log n)$ 

```
// ordered
struct node {
    vector<node*> child;
};
bool otreeIsomorphism(node *n, node *m) {
    if (n->child.size() != m->child.size()) return false;
    rep(i, n->child.size())
        if (!otreeIsomorphism(n->child[i], m->child[i])) return false;
    return true;
}
```

```
11
    // not ordered
12
   struct node {
13
        vector<node *> child;
14
        vector<int> code;
15
   };
16
17
   void code(node *n) {
       int size = 1;
18
       vector< pair<vector<int>, int> > codes;
19
       rep(i, n->child.size()) {
20
            code(n->child[i]);
21
            codes.push_back( make_pair(n->child[i]->code, i) );
22
            size += codes[i].first[0];
23
24
       sort(codes.rbegin(), codes.rend()); // !reverse
25
       n->code.push_back(size);
26
       for (int i = 0; i < n->child.size(); ++i) {
27
            swap(n->child[i], n->child[ codes[i].second ]);
28
            n->code.insert(n->code.end(),
29
                            codes[i].first.begin(), codes[i].first.end());
30
31
32
   bool utreeIsomorphism(node *n, node *m) {
33
       code(n); code(m); return n->code == m->code;
34
35
```

#### 3.3.6 HL 分解

```
namespace HLD {
   const int N = 200010;
   vector<vector<int>> chains, childs;
   int V, dep[N], par[N], heavy[N], head[N], chain[N], id[N], size[N], q[N];
   void calc_heavy() {
       int root = -1;
       childs.assign(V, vector<int>());
       for(int v = 0; v < V; v++) {
            size[v] = 0;
10
11
            heavy[v] = -1;
            if(par[v] < 0) root = v;
12
            else childs[par[v]].push_back(v);
13
14
       int 1 = 0, r = 0;
15
       q[r++] = root;
16
       while(1 < r)  {
17
            int v = q[1++];
18
            for(auto &w: childs[v]) {
19
20
                if(w == par[v]) continue;
                dep[w] = dep[v]+1;
21
                q[r++] = w;
22
23
24
25
        reverse(q,q+V);
26
        for(int i = 1; i < V; i++) {
27
            int v = q[i], &u = par[v];
28
            size[u] += ++size[v];
            if(heavy[u] == -1 or size[v] > size[heavy[u]]) heavy[u] = v;
29
30
31
   void calc_chain() {
32
       chains.clear();
33
34
       int idx = 0;
        for (int v = 0; v < V; v++) {
35
36
            if(par[v] < 0 or heavy[par[v]] != v) {</pre>
```

```
chains.push back(vector<int>()):
37
                 for (int w = v; w != -1; w = heavy[w]) {
38
39
                     chain[w] = idx;
40
                    head[w] = v;
                    id[w] = chains.back().size();
41
                     chains.back().push_back(w);
42
43
                idx++;
44
45
46
47
   void make_par(const vector<vector<int>> &g, int root = 0) {
48
        memset(par,-1,sizeof(par));
49
50
        par[root] = 0;
        int 1 = 0, r = 0;
51
        q[r++] = root;
52
        while(1 < r)  {
53
54
            int v = q[1++];
55
            for(const int &w: g[v]) if(par[w] < 0) q[r++] = w, par[w] = v;
56
57
        par[root] = -1;
58
59
   void build(const vector<vector<int>> &g, int root = 0) {
        V = g.size();
60
        make_par(g,root);
61
        calc_heavy();
62
63
        calc_chain();
64
65
   int lca(int u, int v) {
        while (chain[u] != chain[v]) {
66
            if (dep[head[u]] > dep[head[v]]) swap(u,v);
67
68
            v = par[head[v]];
69
70
        return dep[u] < dep[v]? u: v;
71
72
```

#### 3.3.7 重心分解

```
const int N = 100010;
   int level[N], par[N], done[N];
   vector<int> bfs(int s) {
        vector<int> ret;
        queue < int > que;
        que.push(s), par[s] = -1;
        while (not que.empty()) {
            int v = que.front(); que.pop();
            ret.push_back(v);
10
            done[v] = true;
            for (int u : G[v]) {
12
                if (level[u] == 0 and not done[u]) {
                    que.push(u), par[u] = v;
14
15
16
17
        return ret;
18
19
   int size[N], ch[N];
   void update(int v) {
20
21
        size[v] = 1, ch[v] = 0;
22
        for (int u : G[v]) {
23
            if (u != par[v] and level[u] == 0) {
                size[v] += size[u]:
24
25
                ch[v] = max(ch[v], size[u]);
```

```
27
28
29
    void decompomposite() {
        auto ord = bfs(0);
30
        rep(i, 26) {
31
            fill_n(done, n, 0);
32
            for (int v : ord) {
33
                if (level[v] == 0 and not done[v]) {
34
                     auto sub = bfs(v);
35
36
                     reverse(all(sub));
                    for (int u : sub) update(u);
37
                     int whole = size[v], petal = ch[v];
38
                     for (bool flag = true; flag; ) {
39
                         flag = false;
40
                         for (int c : G[v]) {
41
                             if (level[c] == 0) {
42
                                 int tmp = max(ch[c], whole - size[c]);
43
                                 if (petal > tmp) {
44
                                     v = c, petal = tmp;
45
                                      flag = true;
46
                                      break:
47
48
49
                         }
50
51
52
                     // v is a centroid
53
                     level[v] = i + 1;
54
55
56
57
```

## 3.4 彩色数

#### 3.4.1 包除原理

 $O(2^VV)$ 

N[i] := i と隣接する頂点の集合 (i も含む)

```
const int MAX_V=16;
   const int mod = 10009;
   int N[MAX_V], I[1<<MAX_V], V;
   inline int mpow(int a, int k){ return k==0? 1: k%2? a*mpow(a,k-1)%mod: mpow(a*a%mod,k
        /2);}
   bool can(int k){
       int res = 0;
       rep(S, 1<<V){
            if(__builtin_popcountl1(S)%2) res -= mpow(I[S], k);
10
            else res += mpow(I[S],k);
11
12
       return (res%mod+mod)%mod;
13
14
15
    int color_number(){
16
       memset(I, 0, sizeof(I));
17
       I[0] = 1;
       repi(S,1,1<<V){
18
19
            int v = 0;
            while(!(S&(1<<v))) v++;
20
21
            I[S] = I[S-(1 << v)] + I[S&(~N[v])];
22
23
       int lb = 0, ub = V, mid;
```

#### 3.4.2 極大独立集合

```
typedef vector<vector<int>> graph;
   class maximal_indsets {
        const int n;
        const graph& G;
        vector<vector<int>> ret;
        vector<int> cur, exists, deg, block;
7
        void erase(int v) {
            if (exists[v]) {
                exists[v] = false;
                for (int nv : G[v]) --deg[nv];
10
11
12
13
        void restore(int v) {
14
            exists[v] = true;
15
            for (int nv : G[v]) ++deg[nv];
16
17
        void select(int v) {
18
            cur.push_back(v);
19
            ++block[v], erase(v);
20
            for (int nv : G[v]) ++block[nv], erase(nv);
21
22
        void unselect(int v) {
23
            cur.pop back():
            --block[v], restore(v);
24
25
            for (int nv : G[v]) {
                if (--block[nv] == 0) restore(nv);
26
27
           }
28
        void dfs() {
29
30
            int mn = n, v = -1;
31
            rep(u, n) if (exists[u]) {
32
                if (deg[u] < mn) mn = deg[u], v = u;
33
            if (v == -1) {
34
                ret.push_back(cur);
35
36
           } else {
37
                select(v), dfs(), unselect(v);
38
                for (int nv : G[v]) {
39
                    if (exists[nv]) select(nv), dfs(), unselect(nv);
40
41
           }
42
   public:
43
        maximal_indsets(const graph& G) : n(G.size()), G(G), exists(n, true), deg(n), block(
45
            rep(v, n) deg[v] = G[v].size();
            dfs();
46
47
        const vector<vector<int>>& get() const { return ret; }
48
   };
```

## 4 数学

## 4.1 整数

### 4.1.1 剰余

```
// (x, y) s.t. a x + b y = gcd(a, b)
   long extgcd(long a, long b, long& x, long& y) {
        long q = a; x = 1, y = 0;
       if (b != 0) g = extgcd(b, a % b, y, x), y -= <math>(a / b) * x;
   // inv[1] = 1; repi(i,2,n) inv[i] = inv[p%i] * (p - p/i) % p;
   long mod_inv(long a, long m) {
       long x, y;
       if (extgcd(a, m, x, y) != 1) return 0;
10
       return (x % m + m) % m;
11
12 }
13
   // a mod p where n! = a p^e in O(log_p n)
   long mod_fact(long n, long p, long& e) {
14
       const int P = 1000010:
15
        static long fac[P] = {1};
16
        for (static int once = 1: once: --once) {
17
            repi(i,1,P) fac[i] = fac[i-1] * i % p;
18
       }
19
20
       e = 0:
21
       if (n == 0) return 1;
       long ret = mod_fact(n/p, p, e);
22
23
       e += n/p;
       return ret * (n/p%2 ? p - fac[n%p] : fac[n%p]) % p;
24
25
   long mod_binom(long n, long k, long p) {
26
       if (k < 0 \text{ or } n < k) \text{ return } 0;
27
28
       long e1, e2, e3;
       long a1 = mod_fact(n, p, e1);
29
       long a2 = mod_fact(k, p, e2);
30
       long a3 = mod_fact(n - k, p, e3);
31
       if (e1 > e2 + e3) return 0;
32
       return a1 * mod_inv(a2 * a3 % p, p) % p;
33
34
   long mod_pow(long a, long b, long m) {
35
       long ret = 1;
36
       do {
37
            if (b & 1) ret = ret * a % m;
38
            a = a * a % m;
39
       } while (b >>= 1);
40
41
       return ret;
42
   inline long mod_mul(long a, long b, long m) {
43
       long ret = a * b - m * long(roundl((long double)(a) * b / m));
44
        return ret < 0 ? ret + m : ret;</pre>
45
46
```

#### 4.1.2 離散対数

```
long discrete_log(long a, long m) {
    if (a == 0) return -1;
    long b = sqrt(m)+1, t = 1;
    unordered_map<long,long> mem;
    rep(i, b) {
        mem[t] = i;
        t = t * a % m;
    }
}
```

```
if (t == 1) return i+1:
8
10
        long u = t;
11
        for (int i = b; i < m; i += b) {
            if (mem.find(mod_inv(u, m)) != mem.end()) {
12
                return mem[mod_inv(u, m)] + i;
13
14
           u = u * t % m;
15
16
       }
17
        return -1:
18
```

#### 4.1.3 カタラン数

() を正しく並べる方法, 二分木, 格子状の経路の数え上げ, 平面グラフの交差などに使われる.  $C_{16}=35357670$  が限界?

$$C_n = \frac{1}{n+1} {2n \choose n} = {2n \choose n} - {2n \choose n-1} \approx \frac{4^n}{n^{3/2} \sqrt{\pi}}$$

## 4.1.4 乱数 (xor shift)

周期は 2128 - 1

```
unsigned xorshift() {
    static unsigned x = 123456789;
    static unsigned y = 362436069;
    static unsigned z = 521288629;
    static unsigned w = 88675123;
    unsigned t;
    t = x ^cb^86 (x << 11);
    x = y; y = z; z = w;
    return w = (w ^cb^86 (w >> 19)) ^cb^86 (t ^cb^86 (t >> 8));
}
```

#### 4.1.5 確率的素数判定 (Miller-Rabin 法)

 $O(k \log^3 n)$ 合成数を素数と判定する確率は最大で  $4^{-k}$ 

```
bool suspect(long a, int s, long d, long n) {
        long x = mod_pow(a, d, n); // use mod_powl instead for large n
        if (x == 1) return true;
        for (int r = 0; r < s; ++r) {
           if (x == n - 1) return true;
            x = x * x % n; // use mod_mul instead for large n
       return false;
   // {2,7,61,-1}
                                   is for n < 4759123141 (= 2^32)
11
   // {2.3.5.7.11.13.17.19.23.-1} is for n < 10^16 (at least)
   bool is_prime(long n) {
12
13
        if (n \le 1 \mid | (n > 2 \&\& n \% 2 == 0)) return false;
        int test[] = \{2,3,5,7,11,13,17,19,23,-1\};
14
15
        long d = n - 1, s = 0;
        while (d \% 2 == 0) ++s, d /= 2;
16
        for (int i = 0; test[i] < n && test[i] != -1; ++i)
17
            if (!suspect(test[i]. s. d. n)) return false:
18
19
        return true:
```

20 }

## 4.2 多項式

FFT は基本定数重めなので TLE に注意する.

#### 4.2.1 FFT(complex)

 $O(N \log N)$ 

複素数を用いた FFT. 変換する vector のサイズは2の冪乗にすること.

```
typedef complex < double > cd;
    vector<cd> fft(vector<cd> f. bool inv){
        int n, N = f.size();
        for(n=0;;n++) if(N == (1 << n)) break;
        rep(m,N){
            int m2 = 0;
            rep(i,n) if(m&(1<<i)) m2 |= (1<<(n-1-i));
            if(m < m2) swap(f[m], f[m2]);</pre>
       }
10
        for(int t=1;t<N;t*=2){
11
            double theta = acos(-1.0) / t;
12
13
            cd w(cos(theta), sin(theta));
            if(inv) w = cd(cos(theta), -sin(theta));
14
            for(int i=0;i<N;i+=2*t){</pre>
15
                cd power(1.0, 0.0);
16
17
                rep(j,t){
                     cd tmp1 = f[i+j] + f[i+t+j] * power;
18
                     cd tmp2 = f[i+j] - f[i+t+j] * power;
19
                     f[i+j] = tmp1;
20
                     f[i+t+j] = tmp2;
21
                     power = power * w;
22
23
24
25
26
        if(inv) rep(i,N) f[i] /= N;
27
        return f:
28
```

#### 4.2.2 FFT(modulo)

 $O(N \log N)$ 

剰余環を用いた FFT(FMT). 変換する vector のサイズは 2 の冪乗にすること. mod は  $a*2^e+1$  の形.

```
#include "number_theory.cpp"
    const int mod = 7*17*(1<<23)+1;
   vector<int> fmt(vector<int> f, bool inv){
       int e, N = f.size();
        // assert((N&(N-1))==0 and "f.size() must be power of 2");
       for(e=0;;e++) if(N == (1 << e)) break;
       rep(m,N){
10
            rep(i,e) if(m&(1<<i)) m2 |= (1<<(e-1-i));
            if(m < m2) swap(f[m], f[m2]);</pre>
11
12
        for(int t=1: t<N: t*=2){
13
14
            int r = pow_mod(3, (mod-1)/(t*2), mod);
```

```
if(inv) r = mod inverse(r.mod):
15
            for(int i=0; i<N; i+=2*t){
16
17
                int power = 1;
18
                rep(j,t){
19
                     int x = f[i+j], y = 1LL*f[i+t+j]*power*mod;
                     f[i+j] = (x+y)\%mod;
20
                    f[i+t+j] = (x-y+mod)\%mod;
21
                    power = 1LL*power*r%mod;
22
23
24
           }
25
26
        if(inv) for(int i=0.ni=mod inv(N.mod):i<N:i++) f[i] = 1LL*f[i]*ni%mod:
27
        return f;
28
```

### 4.2.3 積 (FMT)

O(N log N). fmt() が必要.

```
vector<int> poly_mul(vector<int> f, vector<int> g){
   int N = max(f.size(),g.size())*2;
   f.resize(N); g.resize(N);
   f = fmt(f,0); g = fmt(g,0);
   rep(i,N) f[i] = 1LL*f[i]*g[i]%mod;
   f = fmt(f,1);
   return f;
}
```

#### 4.2.4 逆元 (FMT)

 $O(N \log N)$ . extgcd(), mod\_inverse(), poly\_mul(), fmt() が必要.

```
vector<int> poly_inv(const vector<int> &f){
       int N = f.size();
2
3
       vector<int> r(1,mod_inv(f[0],mod));
       for (int k = 2; k \le N; k \le 1)
           vector<int> nr = poly_mul(poly_mul(r,r), vector<int>(f.begin(),f.begin()+k));
           nr.resize(k);
           rep(i,k/2) {
               nr[i] = (2*r[i]-nr[i]+mod)%mod;
               nr[i+k/2] = (mod-nr[i+k/2])%mod;
10
11
           r = nr;
12
13
       return r;
```

### 4.2.5 平方根 (FMT)

O(NlogN). extgcd(), mod\_inverse(), poly\_inv(), poly\_mul(), fmt() が必要.

```
const int inv2 = (mod+1)/2;
vector<int> poly_sqrt(const vector<int> &f) {
   int N = f.size();
   vector<int> s(1,1); // s[0] = sqrt(f[0])
   for(int k = 2; k <= N; k <<= 1) {
        s.resize(k);
        vector<int> ns = poly_mul(poly_inv(s), vector<int>(f.begin(),f.begin()+k));
        ns.resize(k);
}
```

## 4.3 行列

```
typedef double number;
   typedef vector<number> vec;
   typedef vector<vec> mat;
   vec mul(const mat& A, const vec& x) {
       const int n = A.size();
       vec b(n):
       rep(i, n) rep(j, A[0].size()) {
            b[i] = A[i][j] * x[j];
       return b:
10
11
12
   mat mul(const mat& A, const mat& B) {
       const int n = A.size();
13
       const int o = A[0].size();
14
       const int m = B[0].size();
15
16
       mat C(n, vec(m));
17
       rep(i, n) rep(k, o) rep(j, m) {
            C[i][j] += A[i][k] * B[k][j];
18
       3
19
       return C;
20
21
   mat pow(mat A, long m) {
22
       const int n = A.size();
23
       mat B(n, vec(n));
24
25
       rep(i, n) B[i][i] = 1;
26
            if (m \& 1) B = mul(B, A);
27
            A = mul(A, A);
28
       } while (m >>= 1);
29
30
       return B;
31
32
   const number eps = 1e-4;
33
    // determinant; 0(n^3)
   number det(mat A) {
34
35
       int n = A.size();
       number D = 1;
36
37
       rep(i,n){
            int pivot = i;
38
39
            repi(j,i+1,n)
                if (abs(A[j][i]) > abs(A[pivot][i])) pivot = j;
40
41
            swap(A[pivot], A[i]);
42
            D *= A[i][i] * (i != pivot ? -1 : 1);
43
            if (abs(A[i][i]) < eps) break;</pre>
44
            repi(j,i+1,n)
                for(int k=n-1; k>=i;--k)
45
46
                    A[j][k] -= A[i][k] * A[j][i] / A[i][i];
47
48
       return D;
49
   // rank; 0(n^3)
50
51
   int rank(mat A) {
        int n = A.size(), m = A[0].size(), r = 0;
52
53
        for(int i = 0; i < m and r < n; i++){
54
            int pivot = r;
55
            repi(j,r+1,n)
                if (abs(A[j][i]) > abs(A[pivot][i])) pivot = j;
56
            swap(A[pivot], A[r]);
57
```

```
58
            if (abs(A[r][i]) < eps) continue:
59
            for (int k=m-1; k>=i; --k)
60
                 A[r][k] /= A[r][i];
61
            repi(j,r+1,n) repi(k,i,m)
62
                A[j][k] -= A[r][k] * A[j][i];
63
64
        }
65
        return r;
66
```

### 4.3.1 線形方程式の解 (Givens 消去法)

 $O(N^3)$ 

```
typedef double number;
   typedef vector<vector<number> > matrix;
   inline double my_hypot(double x, double y) { return sqrt(x * x + y * y); }
   inline void givens_rotate(number& x, number& y, number c, number s) {
        number u = c * x + s * y, v = -s * x + c * y;
        x = u, y = v;
7
   vector<number> givens(matrix A, vector<number> b) {
        const int n = b.size();
        rep(i, n) repi(j, i + 1, n) {
10
11
            const number r = my_hypot(A[i][i], A[j][i]);
            const number c = A[i][i] / r, s = A[j][i] / r;
12
13
            givens_rotate(b[i], b[j], c, s);
            repi(k, i, n) givens_rotate(A[i][k], A[j][k], c, s);
14
15
16
        for (int i = n - 1; i >= 0; --i) {
17
            repi(j, i + 1, n) b[i] -= A[i][j] * b[j];
           b[i] /= A[i][i];
18
19
20
       return b;
21
```

## 4.4 割り当て問題

#### **4.4.1** ハンガリアン法

 $O(N^2)$ 

```
int hungarian(const vector<vector<int>> &a) {
        int n = a.size(), p, q;
2
        vector<int> fx(n, inf), fy(n, 0), x(n, -1), y(n, -1);
        rep(i,n) rep(j,n) fx[i] = max(fx[i], a[i][j]);
        for (int i = 0; i < n; ) {
            vector\langle int \rangle t(n, -1), s(n+1, i);
            for (p = q = 0; p \le q \&\& x[i] < 0; ++p)
                 for (int k = s[p], j = 0; j < n && x[i] < 0; ++j)
                    if (fx[k] + fy[j] == a[k][j] && t[j] < 0) {
                         s[++q] = y[j], t[j] = k;
11
                         if (s[q] < 0)
12
                             for (p = j; p >= 0; j = p)
                                 y[j] = k = t[j], p = x[k], x[k] = j;
13
14
            if (x[i] < 0) {
15
16
                int d = inf;
                rep(k,q+1) \ rep(j,n) \ if \ (t[j] < 0) \ d = min(d, fx[s[k]] + fy[j] - a[s[k]][j]);
17
                rep(j,n) fy[j] += (t[j] < 0 ? 0 : d);
18
19
                rep(k,q+1) fx[s[k]] -= d;
20
            } else i++;
```

### 5 幾何

```
// constants and eps-considered operators
const double eps = 1e-8; // choose carefully!
const double pi = acos(-1.0);

inline bool lt(double a, double b) { return a < b - eps; }
inline bool gt(double a, double b) { return lt(b, a); }
inline bool le(double a, double b) { return !lt(b, a); }
inline bool ge(double a, double b) { return !lt(a, b); }
inline bool ne(double a, double b) { return !lt(a, b) or lt(b, a); }
inline bool eq(double a, double b) { return !ne(a, b); }
```

## 5.1 点

```
typedef complex<double> point;
   inline double dot (point a, point b) { return real(conj(a) * b); }
   inline double cross(point a, point b) { return imag(conj(a) * b); }
       Here is what ccw(a, b, c) returns:
             1
       _____
        2 |a 0 b| -2
10
            - 1
11
12
    * Note: we can implement intersectPS(p. s) as !ccw(s.a. s.b. p).
13
14
   int ccw(point a, point b, point c) {
15
16
       b -= a, c -= a;
17
       if (cross(b, c) > eps)
                              return +1:
       if (cross(b, c) < eps)</pre>
18
                              return -1;
       if (dot(b, c) < eps)</pre>
                                return +2; // c -- a -- b
19
       if (lt(norm(b), norm(c))) return -2; // a -- b -- c
20
21
       return 0:
22 }
```

# 5.2 直線と線分

```
struct line {
   point a, b;
   line(point a, point b) : a(a), b(b) {}

bool intersectLS(const line& 1, const line& s) {
   return ccw(l.a, l.b, s.a) * ccw(l.a, l.b, s.b) <= 0;
}

bool intersectSS(const line& s, const line& t) {
   return intersectLS(s, t) and intersectLS(t, s);</pre>
```

```
bool intersectLL(const line& 1, const line& m) {
12
13
       return ne(cross(1.b - 1.a, m.b - m.a), 0.0) // not parallel
14
           or eq(cross(l.b - l.a, m.a - l.a), 0.0); // overlap
15
   point crosspointLL(const line& 1, const line& m) {
16
17
        double A = cross(1.b - 1.a, m.b - m.a);
        double B = cross(1.b - 1.a, m.a - 1.a);
18
19
        if (eq(A, 0.0) \text{ and } eq(B, 0.0)) return m.a; // overlap
                                                    // not parallel
20
        assert(ne(A, 0.0));
21
       return m.a - B / A * (m.b - m.a);
22
   point proj(const line& 1, point p) {
23
        double t = dot(1.b - 1.a, p - 1.a) / norm(1.b - 1.a);
       return 1.a + t * (1.b - 1.a);
25
26
   point reflection(const line& 1, point p) { return 2.0 * proj(1, p) - p; }
27
    double distanceLP(const line& 1, point p) { return abs(proj(1, p) - p); }
   double distanceLL(const line& 1, const line& m) {
       return intersectLL(1, m) ? 0.0 : distanceLP(1, m.a);
31
32
33
   double distanceLS(const line& 1, const line& s) {
        return intersectLS(1, s) ? 0.0 : min(distanceLP(1, s.a), distanceLP(1, s.b));
34
35
36
   double distancePS(point p, const line& s) {
37
        point h = proj(s, p);
        return ccw(s.a, s.b, h) ? min(abs(s.a - p), abs(s.b - p)) : abs(h - p);
38
39
   double distanceSS(const line& s, const line& t) {
       if (intersectSS(s, t)) return 0.0;
        return min(min(distancePS(s.a, t), distancePS(s.b, t)),
42
                   min(distancePS(t.a, s), distancePS(t.b, s)));
43
44
```

## 5.3 円

```
struct circle {
        point o; double r;
        circle(point o, double r) : o(o), r(r) {}
   };
   bool intersectCL(const circle& c, const line& 1) {
       return le(norm(proj(1, c.o) - c.o), c.r * c.r);
   int intersectCS(const circle& c, const line& s) {
10
        if (not intersectCL(c, s)) return 0;
11
        double a = abs(s.a - c.o);
        double b = abs(s.b - c.o);
        if (lt(a, c.r) and lt(b, c.r)) return 0;
13
14
        if (lt(a, c.r) or lt(b, c.r)) return 1;
        return ccw(s.a, s.b, proj(s, c.o)) ? 0 : 2;
15
16
   bool intersectCC(const circle& c, const circle& d) {
        double dist = abs(d.o - c.o);
19
        return le(abs(c.r - d.r), dist) and le(dist, c.r + d.r);
20
21
   line crosspointCL(const circle& c, const line& 1) {
        point h = proj(1, c.o);
22
23
        double a = sqrt(c.r * c.r - norm(h - c.o));
        point d = a * (1.b - 1.a) / abs(1.b - 1.a);
24
       return line(h - d, h + d);
27 | line crosspointCC(const circle& c, const circle& d) {
```

```
double dist = abs(d.o - c.o). th = arg(d.o - c.o):
28
        double ph = acos((c.r * c.r + dist * dist - d.r * d.r) / (2.0 * c.r * dist));
29
30
       return line(c.o + polar(c.r, th - ph), c.o + polar(c.r, th + ph));
31
32
   line tangent(const circle& c. double th) {
33
34
       point h = c.o + polar(c.r, th);
       point d = polar(c.r, th) * point(0, 1);
35
36
       return line(h - d, h + d);
37
38
   vector<line> common_tangents(const circle& c, const circle& d) {
       vector<line> ret:
39
        double dist = abs(d.o - c.o), th = arg(d.o - c.o);
40
       if (abs(c.r - d.r) < dist) { // outer</pre>
41
            double ph = acos((c.r - d.r) / dist);
42
43
            ret.pb(tangent(c, th - ph));
            ret.pb(tangent(c, th + ph));
44
45
       if (abs(c.r + d.r) < dist) { // inner</pre>
46
            double ph = acos((c.r + d.r) / dist);
47
            ret.pb(tangent(c, th - ph));
48
            ret.pb(tangent(c, th + ph));
49
50
51
       return ret:
52
   pair<circle, circle> tangent_circles(const line& 1, const line& m, double r) {
53
54
       double th = arg(m.b - m.a) - arg(l.b - l.a);
       double ph = (arg(m.b - m.a) + arg(1.b - 1.a)) / 2.0;
55
56
       point p = crosspointLL(1, m);
       point d = polar(r / sin(th / 2.0), ph);
57
       return mp(circle(p - d, r), circle(p + d, r));
58
59
   line bisector(point a, point b);
60
61
   circle circum_circle(point a, point b, point c) {
62
       point o = crosspointLL(bisector(a, b), bisector(a, c));
        return circle(o, abs(a - o));
63
64
```

## 5.4 多角形

```
typedef vector<point> polygon;
   double area(const polygon& g) {
       double ret = 0.0;
       int j = g.size() - 1;
       rep(i, g.size()) {
            ret += cross(g[j], g[i]), j = i;
       return ret / 2.0;
10
11
   point centroid(const polygon& g) {
12
       if (g.size() == 1) return g[0];
13
       if (g.size() == 2) return (g[0] + g[1]) / 2.0;
       point ret = 0.0:
14
15
       int j = g.size() - 1;
16
       rep(i, q.size()) {
            ret += cross(g[j], g[i]) * (g[j] + g[i]), j = i;
17
18
19
       return ret / area(g) / 6.0;
20
   line bisector(point a, point b) {
21
22
       point m = (a + b) / 2.0;
23
       return line(m. m + (b - a) * point(0. 1)):
24 }
```

```
polygon convex cut(const polygon& g. const line& 1) {
26
        polygon ret;
27
        int j = q.size() - 1;
28
        rep(i, g.size()) {
29
            if (ccw(l.a, l.b, g[j]) != -1) ret.pb(g[j]);
            if (intersectLS(1, line(g[j], g[i]))) ret.pb(crosspointLL(1, line(g[j], g[i])));
30
31
           j = i;
32
33
        return ret;
34
35
   polygon voronoi_cell(polygon g, const vector<point>& v, int k) {
        rep(i, v.size()) if (i != k) {
           g = convex_cut(g, bisector(v[i], v[k]));
37
38
39
        return g;
40
```

## 5.4.1 凸包

```
namespace std {
        bool operator <(const point& a, const point& b) {</pre>
            return ne(real(a), real(b)) ? lt(real(a), real(b)) : lt(imag(a), imag(b));
   }
   polygon convex_hull(vector<point> v) {
        const int n = v.size();
        sort(all(v));
        polygon ret(2 * n);
        int k = 0;
11
12
        for (int i = 0; i < n; ret[k++] = v[i++]) {
13
            while (k \ge 2 \text{ and } ccw(ret[k - 2], ret[k - 1], v[i]) \le 0) --k;
14
15
        for (int i = n - 2, t = k + 1; i >= 0; ret[k++] = v[i--]) {
16
            while (k \ge t \text{ and } ccw(ret[k - 2], ret[k - 1], v[i]) \le 0) --k;
17
18
        ret.resize(k - 1);
19
        return ret;
```

## 5.4.2 最近点対

だいたい  $O(n \log n)$ , 最悪縦 1 列に並んでる場合  $O(n^2)$ 

```
pair<point, point> closest_pair(vector<point> p) {
    int n = p.size(), s = 0, t = 1, m = 2, S[n];
    S[0] = 0, S[1] = 1;
    sort(all(p)); // "p < q" <=> "p.x < q.x"
    double d = norm(p[s]-p[t]);
    for (int i = 2; i < n; S[m++] = i++) rep(j, m) {
        if (norm(p[S[j]]-p[i]) < d) d = norm(p[s = S[j]]-p[t = i]);
        if (real(p[S[j]]) < real(p[i]) - d) S[j--] = S[--m];
    }
    return make_pair(p[s], p[t]);
}</pre>
```

## 5.4.3 点-多角形包含判定

O(n)

```
enum { OUT, ON, IN };
int contains(const polygon& P, const point& p) {
   bool in = false;
   for (int i = 0; i < (int)P.size(); ++i) {
       point a = P[i] - p, b = P[(i+1)%P.size()] - p;
       if (imag(a) > imag(b)) swap(a, b);
       if (imag(a) <= 0 && 0 < imag(b) && cross(a, b) < 0) in = !in;
       if (cross(a, b) == 0 && dot(a, b) <= 0) return ON;
   }
   return in ? IN : OUT;
}</pre>
```

#### 5.4.4 凸多角形の共通部分

```
O(n+m)
```

```
bool intersect_1pt(const point& a, const point& b,
2
                       const point& c, const point& d, point &r) {
       number D = cross(b - a, d - c);
3
       if (eq(D,0)) return false;
       number t = cross(c - a. d - c) / D:
       number s = -cross(a - c, b - a) / D;
       r = a + t * (b - a);
       return ge(t, 0) && le(t, 1) && ge(s, 0) && le(s, 1);
9
10
   polygon convex_intersect(const polygon &P, const polygon &Q) {
11
       const int n = P.size(), m = Q.size();
       int a = 0, b = 0, aa = 0, ba = 0;
12
13
       enum { Pin, Qin, Unknown } in = Unknown;
14
       polygon R;
15
       do {
16
            int a1 = (a+n-1) % n, b1 = (b+m-1) % m;
            number C = cross(P[a] - P[a1], Q[b] - Q[b1]);
17
18
            number A = cross(P[a1] - Q[b], P[a] - Q[b]);
            number B = cross(Q[b1] - P[a], Q[b] - P[a]);
19
20
21
            if (intersect_1pt(P[a1], P[a], Q[b1], Q[b], r)) {
22
                if (in == Unknown) aa = ba = 0;
                R.push_back( r );
23
                in = B > 0 ? Pin : A > 0 ? Oin : in:
24
25
26
            if (C == 0 \&\& B == 0 \&\& A == 0) 
               if (in == Pin) { b = (b + 1) \% m; ++ba; }
27
               else
                               \{ a = (a + 1) \% m; ++aa; \}
28
            } else if (C >= 0) {
29
30
               if (A > 0) { if (in == Pin) R.push_back(P[a]); a = (a+1)%n; ++aa; }
                           { if (in == Qin) R.push_back(Q[b]); b = (b+1)%m; ++ba; }
31
           } else {
32
               if (B > 0) { if (in == Qin) R.push_back(Q[b]); b = (b+1)%m; ++ba; }
33
                           { if (in == Pin) R.push_back(P[a]); a = (a+1)%n; ++aa; }
34
35
       } while ( (aa < n || ba < m) && aa < 2*n && ba < 2*m );
36
37
       if (in == Unknown) {
            if (convex_contains(Q, P[0])) return P;
38
            if (convex_contains(P, Q[0])) return Q;
39
40
41
       return R;
42 }
```

## 5.4.5 凸多角形の直径

O(n)

```
inline double diff(const vector<point> &P, const int &i) { return (P[(i+1)%P.size()] - P
    number convex_diameter(const polygon &pt) {
        const int n = pt.size();
3
        int is = 0, js = 0;
        for (int i = 1; i < n; ++i) {
5
            if (imag(pt[i]) > imag(pt[is])) is = i;
7
            if (imag(pt[i]) < imag(pt[js])) js = i;</pre>
8
        number maxd = norm(pt[is]-pt[js]);
9
10
11
        int i, maxi, j, maxj;
12
        i = maxi = is:
13
       j = maxj = js;
14
        do {
15
            if (cross(diff(pt,i), diff(pt,j)) >= 0) j = (j+1) % n;
16
            else i = (i+1) \% n;
17
            if (norm(pt[i]-pt[j]) > maxd) {
18
                maxd = norm(pt[i]-pt[j]);
19
                maxi = i; maxj = j;
20
21
        } while (i != is || j != js);
22
        return maxd; /* farthest pair is (maxi, maxj). */
23
```

### 5.4.6 ドロネー三角形分割(逐次添加法)

 $O(n^2)$ 

```
bool incircle(point a, point b, point c, point p) {
        a -= p; b -= p; c -= p;
        return norm(a) * cross(b, c)
             + norm(b) * cross(c, a)
             + norm(c) * cross(a, b) >= 0;
        // < : inside, = cocircular, > outside
7
   #define SET_TRIANGLE(i, j, r) \
        E[i].insert(j); em[i][j] = r; \
10
        E[j].insert(r); em[j][r] = i; \setminus
        E[r].insert(i); em[r][i] = j; \
12
        S.push(pair<int,int>(i, j));
13
    #define REMOVE_EDGE(i, j) \
        E[i].erase(j); em[i][j] = -1; \setminus
        E[j].erase(i); em[j][i] = -1;
    #define DECOMPOSE_ON(i,j,k,r) { \
            int m = em[j][i]; REMOVE_EDGE(j,i); \
            SET_TRIANGLE(i,m,r); SET_TRIANGLE(m,j,r); \
            SET_TRIANGLE(j,k,r); SET_TRIANGLE(k,i,r); }
    #define DECOMPOSE_IN(i,j,k,r) { \
            SET_TRIANGLE(i,j,r); SET_TRIANGLE(j,k,r); \
22
            SET_TRIANGLE(k,i,r); }
    #define FLIP_EDGE(i,j) { \
            int k = em[j][i]; REMOVE_EDGE(i,j); \
24
25
            SET_TRIANGLE(i,k,r); SET_TRIANGLE(k,j,r); }
    #define IS_LEGAL(i, j) \
27
        (em[i][j] < 0 \mid | em[j][i] < 0 \mid | \setminus
         !incircle(P[i],P[j],P[em[i][j]],P[em[j][i]]))
29
    double Delaunay(vector<point> P) {
        const int n = P.size();
30
31
        P.push_back( point(-inf,-inf) );
        P.push_back( point(+inf,-inf) );
32
        P.push_back( point( 0 ,+inf) );
33
        int em[n+3][n+3]: memset(em. -1, sizeof(em)):
34
        set < int > E[n+3];
```

```
stack< pair<int.int> > S:
        SET_TRIANGLE(n+0, n+1, n+2);
37
38
        for (int r = 0; r < n; ++r) {
39
            int i = n, j = n+1, k;
40
            while (1) {
                k = em[i][j];
41
42
                        (ccw(P[i], P[em[i][j]], P[r]) == +1) j = k;
                else if (ccw(P[j], P[em[i][j]], P[r]) == -1) i = k;
43
44
                else break:
45
46
                    (ccw(P[i], P[j], P[r]) != +1) { DECOMPOSE_ON(i,j,k,r); }
47
            else if (ccw(P[j], P[k], P[r]) != +1) \{ DECOMPOSE_ON(j,k,i,r); \}
            else if (ccw(P[k], P[i], P[r]) != +1) \{ DECOMPOSE_ON(k,i,j,r); \}
48
                                                    { DECOMPOSE_IN(i,j,k,r); }
49
            while (!S.empty()) {
50
                int u = S.top().first, v = S.top().second; S.pop();
51
                if (!IS_LEGAL(u, v)) FLIP_EDGE(u, v);
52
53
54
       double minarg = 1e5;
55
       for (int a = 0; a < n; ++a) {
56
            for(auto &b: E[a]) {
57
                int c = em[a][b];
58
                if (b < n && c < n) {
59
                    point p = P[a] - P[b], q = P[c] - P[b];
60
61
                    minarg = min(minarg, acos(dot(p,q)/abs(p)/abs(q)));
62
63
64
65
       return minarg;
66
```

# 6 データ構造

## 6.1 Union-Find 木

```
class disjoint_set {
       vector<int> p;
   public:
       disjoint_set(int n) : p(n, -1) {}
       int root(int i) { return p[i] >= 0 ? p[i] = root(p[i]) : i; }
       bool same(int i, int j) { return root(i) == root(j); }
       int size(int i) { return -p[root(i)]; }
       void merge(int i, int j) {
           i = root(i), j = root(j);
           if (i == j) return;
11
           if (p[i] > p[j]) swap(i, j);
12
           p[i] += p[j], p[j] = i;
13
14
   };
```

# 6.2 Meldable Heap

```
template <class T>
class meldable_heap {
    struct node {
        node *1 = NULL, *r = NULL;
        T val;
        node(const T& val) : val(val) {}
        "node() { delete 1, delete r; }
}
```

```
node *meld(node *a, node *b) {
10
            if (!a) return b;
11
            if (!b) return a:
12
            if (a->val > b->val) swap(a, b);
13
            a \rightarrow r = meld(a \rightarrow r, b);
14
            swap(a->1, a->r);
            return a;
15
16
17
        node *root = NULL;
        meldable_heap(node *root) : root(root) {}
18
19
    public:
        meldable_heap() {}
20
        bool empty() const { return !root; }
21
        const T& top() const { return root->val; }
22
        void meld(const meldable_heap<T>&& t) { root = meld(root, t.root); }
23
        void push(const T& val) { root = meld(root, new node(val)); }
24
        void pop() {
25
26
            node *t = root;
27
            root = meld(t->1, t->r);
            t.1 = t.r = NULL;
28
            delete t:
29
30
   };
31
```

## 6.3 Binary-Indexed-Tree

0-indexed

```
template < class T> struct bit {
        int n;
       vector<T> dat;
       bit(int n) : n(n) { dat.assign(n,0); }
       // sum [0,i)
       T sum(int i){
           int ret = 0;
           for(--i; i>=0; i=(i&(i+1))-1) ret += bit[i];
10
           return ret;
11
       // sum [i,j)
12
13
       T sum(int i, int j){ return sum(j) - sum(i);}
14
       // add x to i
15
        void add(int i, T x){ for(; i < n; i|=i+1) bit[i] += x;}
16
```

## 6.4 Segment Tree

区間 add と RMQ ができる.

```
template < class T> struct segtree {
        int N:
        vector<T> dat, sum;
        segtree(int n) {
            N = 1:
            while (N < n) N <<= 1;
            dat.assign(2*N-1,0);
7
            sum.assign(2*N-1,0);
9
        void add(int a, int b, T x) { add(a,b,x,0,0,N);}
10
        T add(int a. int b. T x. int k. int l. int r) {
11
12
            if(b <= l or r <= a) return dat[k];</pre>
```

```
if(a \le 1 \text{ and } r \le b) {
13
14
                 sum[k] += x;
15
                 return dat[k] += x;
16
            int m = (1+r)/2;
17
            return dat[k] = min(add(a,b,x,2*k+1,1,m),add(a,b,x,2*k+2,m,r))+sum[k];
18
19
        T minimum(int a, int b) { return minimum(a,b,0,0,N);}
20
        T minimum(int a, int b, int k, int l, int r) {
21
            if(b \leq 1 or r \leq a) return 1e9;
22
23
            if(a <= l and r <= b) return dat[k];</pre>
24
            int m = (1+r)/2:
            return min(minimum(a,b,2*k+1,1,m),minimum(a,b,2*k+2,m,r))+sum[k];
25
26
27
   };
```

## 6.5 Range Tree (Simple)

```
vector<pair<int,int> > ps;
   struct node {
        int a, b;
       node *1, *r;
       vector<int> ys, bit;
       node(int a, int b) : a(a), b(b), l(NULL), r(NULL) {}
   inline int leftmost(node *v) { return ps[v->a].first; }
   inline int rightmost(node *v) { return ps[v->b - 1].first; }
   node *construct(int a, int b) {
10
       if (a == b) return NULL;
11
       node *ret = new node(a, b);
12
       if (a == b-1) {
13
            ret->ys.push_back(ps[a].second);
14
            ret->bit.push_back(0);
15
       } else {
16
            int m = (a+b)/2;
17
            if ((ret->l = construct(a, m))) {
18
19
                ret->ys.insert(ret->ys.end(), all(ret->l->ys));
20
            if ((ret->r = construct(m, b))) {
21
22
                ret->ys.insert(ret->ys.end(), all(ret->r->ys));
23
            sort(all(ret->ys));
24
            ret->bit.resize(b-a):
25
26
27
       return ret;
28
   void insert(node *v, int x, int y) {
29
30
       if (!v) return;
31
       if (make_pair(x, y) < ps[v->a]) return;
32
       if (make_pair(x, y) > ps[v->b - 1]) return;
33
       int k = lower_bound(all(v->ys), y) - v->ys.begin();
        for (; k < (int)v->bit.size(); k |= k+1) {
34
35
            ++v->bit[k];
36
37
        insert(v->1, x, y);
38
       insert(v->r, x, y);
39
40
   int query(node *v, int x, int y) {
       if (!v or leftmost(v) > x) return 0;
41
42
        if (rightmost(v) <= x) {</pre>
            int ret = 0, k = upper_bound(all(v->ys), y) - v->ys.begin();
43
            for (; k; k &= k - 1) {
44
                ret += v->bit[k-1]:
45
46
```

## 6.6 Sparse table

```
const int N = 200010;
   const int K = 18;
   int st[K][N]:
   void construct(int *a, int n) {
        copy_n(a, n, st[0]);
        repi(k. 1. K) {
            for (int i = 0; i+(1 << k) <= n; ++i) {
                st[k][i] = min(st[k-1][i], st[k-1][i+(1<<(k-1))]);
9
10
11
12
   int query(int a, int b) {
13
        int k = 31-__builtin_clz(b-a);
        return min(st[k][a], st[k][b-(1<<k)]);</pre>
14
15
   }
```

### **6.7 RBST**

```
struct node {
        long val, sum;
        size_t size = 1;
        node *left = NULL, *right = NULL;
        node(long val) : val(val), sum(val) {}
        ~node() { delete left, delete right; }
   inline long sum(node *u) { return u ? u->sum : 0; }
   inline size_t size(node *u) { return u ? u->size : 0; }
   inline node *pull(node *u) {
        u->sum = u->val + sum(u->left) + sum(u->right);
12
        u \rightarrow size = 1 + size(u \rightarrow left) + size(u \rightarrow right);
        return u;
13
14
   node *merge(node *u, node *v) {
15
        if (!u) return v;
        if (!v) return u;
        if (rand() * long(size(u) + size(v)) < long(size(u)) * RAND_MAX) {</pre>
18
            u->right = merge(u->right, v);
19
20
            return pull(u);
21
        } else {
            v->left = merge(u, v->left);
22
23
            return pull(v);
24
25
   pair<node*,node*> split(node *u, size_t k) {
        if (!u or k == 0) return {NULL, u};
28
        if (k == size(u)) return {u, NULL};
29
        if (size(u->left) >= k) {
30
            auto p = split(u->left, k);
            u \rightarrow left = p.second;
31
32
            return {p.first, pull(u)};
33
            auto p = split(u->right, k - size(u->left) - 1);
34
            u->right = p.first:
35
36
            return {pull(u), p.second};
```

```
38
39
   template <class ForwardIterator>
   node *construct_from(ForwardIterator first, ForwardIterator last) {
       if (first == last) return NULL;
41
       auto mid = next(first. (last - first) / 2):
42
43
       node *u = new node(*mid):
       u->left = construct_from(first, mid);
44
       u->right = construct_from(next(mid), last);
45
       return pull(u);
46
47
```

## 6.8 永続 RBST

```
template <class T, size_t N>
   struct mempool {
       static T buf[N]. *head:
       static size_t cnt() { return head - buf; }
       static void clear() { head = buf; }
       void *operator new(size_t _ __attribute__((unused))) { return head++; }
       void operator delete(void *_ __attribute__((unused))) {}
    template <class T, size_t N> T mempool<T, N>::buf[N];
   template <class T, size_t N> T *mempool<T, N>::head = mempool<T, N>::buf;
   struct node;
12
   long sum(node *u):
13
   size_t size(node *u);
14
   struct node : mempool<node, M> {
15
       const long val = 0, sum = 0, lazy = 0;
16
       const size_t size = 1;
17
       node *const left = NULL, *const right = NULL;
18
19
       node(long val) : val(val), sum(val) {}
20
       node(long val, long lazy, node *left, node *right)
21
            : val(val),
22
              sum(val + ::sum(left) + ::sum(right)),
23
              lazy(lazy),
24
              size(1 + ::size(left) + ::size(right)),
25
26
              left(left),
              right(right) {}
27
28
   inline long sum(node *u) { return u ? u->sum + u->lazy * u->size : 0; }
29
   inline size_t size(node *u) { return u ? u->size : 0; }
   inline node *add(node *u, long x) { return u ? new node(u->val, u->lazy + x, u->left, u
31
         ->right) : NULL; }
   node *merge(node *u, node *v) {
32
33
       if (!u) return v;
       if (!v) return u;
34
       if (rand() * long(size(u) + size(v)) < long(size(u)) * RAND_MAX) {</pre>
35
            return new node(u->val + u->lazy, 0, add(u->left, u->lazy), merge(add(u->right,
36
                 u \rightarrow lazy, v);
37
38
            return new node(v->val + v->lazy, 0, merge(u, add(v->left, v->lazy)), add(v->
                 right, v->lazy));
39
40
41
   pair<node *, node *> split(node *u, size_t k) {
       if (!u or k == 0) return {NULL, u};
42
43
        if (k == size(u)) return {u, NULL};
       if (size(u->left) >= k) {
44
            auto p = split(add(u->left, u->lazy), k);
45
            return {p.first. new node(u->val + u->lazv. 0. p.second. add(u->right. u->lazv
46
                ))};
```

```
47
            auto p = split(add(u->right, u->lazy), k - size(u->left) - 1);
48
            return {new node(u->val + u->lazy, 0, add(u->left, u->lazy), p.first), p.second
49
                3:
50
51
   }
52
   template <class OutputIterator>
   OutputIterator dump(OutputIterator it, const node *u, long lazy = 0) {
        if (!u) return it;
        lazy += u->lazy;
55
56
        it = dump(it, u->left, lazy);
57
        *it++ = u->val + lazy;
        return dump(it, u->right, lazy);
58
59
   template <class ForwardIterator>
   node *construct_from(ForwardIterator first, ForwardIterator last) {
        if (first == last) return NULL:
        auto mid = next(first, (last - first) / 2);
63
        return new node(*mid, 0, construct_from(first, mid), construct_from(next(mid), last
64
65
```

## 6.9 赤黒木

```
template < class T> class rbtree {
        enum COL { BLACK, RED,};
        struct node {
            T val, lazy, min_val;
            int color, rnk, size;
            node *left, *right;
            // if !left then this node is leaf
            node(){}
            node(T v) : val(v), min_val(v), color(BLACK), rnk(0), size(1) {
                lazy = 0;
10
                left = right = NULL;
11
12
13
            node(node *1, node *r, int c) : color(c) {
                lazv = 0:
14
                left = 1;
15
                right = r;
                update();
17
18
            void update() {
                eval();
                if(left) {
                     rnk = max(left->rnk+(left->color==BLACK),
                               right -> rnk + (right -> color == BLACK));
23
24
                     size = left->size+right->size;
25
                    left->eval(); right->eval();
                    min_val = min(left->min_val, right->min_val);
27
            void eval() {
29
                min val += lazv:
                if(!left) val += lazy;
31
32
                    left->lazy += lazy;
33
34
                     right -> lazy += lazy;
35
36
                lazy = 0;
37
        };
38
39
        node *new_node(T v) { return new node(v);}
```

```
node *rotate(node *v, int d) {
42
43
             node *w = d? v->right: v->left;
44
             if(d) {
                  v->right = w->left;
45
                  w \rightarrow left = v:
46
47
                  v->right->update();
48
49
              else {
                  v \rightarrow left = w \rightarrow right;
50
51
                  w->right = v;
52
                  v->left->update():
53
             v->update(); w->update();
54
             v \rightarrow color = RED:
55
             w->color = BLACK;
56
57
             return w:
58
         node *merge_sub(node *u, node *v) {
59
             u->eval(): v->eval():
60
             if(u->rnk < v->rnk) {
61
                  node *w = merge_sub(u,v->left);
62
                  v \rightarrow left = w;
63
                  v->update();
64
                  if(v->color == BLACK and w->color == RED and w->left->color == RED) {
65
                      if(v->right->color == BLACK) return rotate(v,0);
66
67
                      else {
                           v \rightarrow color = RED;
68
                           v->left->color = v->right->color = BLACK;
69
70
                           return v;
                      }
71
72
                  else return v;
73
74
              else if(u->rnk > v->rnk) {
75
                  node *w = merge_sub(u->right,v);
76
77
                  u -> right = w;
                  u->update();
78
                  if(u->color == BLACK and w->color == RED and w->right->color == RED) {
79
                      if(u->left->color == BLACK) return rotate(u,1);
80
81
                      else {
82
                           u \rightarrow color = RED:
83
                           u->left->color = u->right->color = BLACK;
                           return u;
84
                      }
85
86
                  else return u;
87
88
              else return new_node(u,v,RED);
89
90
91
         node *insert(node *v, int k) {
              auto p = split(root,k);
92
93
             return root = merge(merge(p.first.v).p.second):
94
95
         void add(node *v, int res, T val) {
             if(res < 1) return;</pre>
96
97
             v->eval();
             if(v->size == res) {
98
99
                  v \rightarrow lazy += val;
100
101
              add(v->left, min(v->left->size, res), val);
102
              add(v->right. res-v->left->size. val):
103
104
             v->update();
105
         T get(node *v, int k) {
106
107
             v->eval();
```

node \*new node(node \*1. node \*r. int c) { return new node(1.r.c):}

41

```
if(!v->left) return v->val:
108
             if(v->left->size > k) return get(v->left, k);
109
110
             return get(v->right, k-v->left->size);
111
112
        T minimum(node *v, int 1, int r) {
113
             if(r-1 < 1) return inf:
114
             v->eval():
             if(v->size == r-1) return v->min_val;
115
116
             return min(minimum(v->left, 1, min(r, v->left->size)),
                        minimum(v->right, l-min(l, v->left->size), r-v->left->size));
117
118
        T inf:
119
    public:
120
121
122
         node *root:
         rbtree() {
123
             \inf = (((1LL << (sizeof(T)*8-2))-1) << 1)+1;
124
125
             root = NULL;
126
         void clear() { delete root; root = NULL;}
127
         node *build(const vector<T> &vs) {
128
             if(!vs.size()) return root = NULL;
129
130
             if((int)vs.size() == 1) return root = new_node(vs[0]);
             int m = vs.size()/2;
131
132
             return root = merge(build(vector<T>(begin(vs),begin(vs)+m)),
133
                                  build(vector<T>(begin(vs)+m,end(vs))));
134
         int size() { return root? root->size: 0;}
135
         node *push_back(T val) { return root = merge(root,new_node(val));}
136
         node *push_front(T val) { return root = merge(new_node(val),root);}
137
         node *merge(node *u, node *v) {
138
             if(!u) return v;
139
             if(!v) return u;
140
141
             u = merge_sub(u,v);
142
             u \rightarrow color = BLACK;
143
             return u;
144
         pair<node*, node*> split(node *v, int k) {
145
146
             if(!k) return pair<node*,node*>(NULL,v);
147
             if(k == v->size) return pair<node*,node*>(v,NULL);
148
             v->eval();
149
             if(k < v->left->size) {
150
                 auto p = split(v->left,k);
151
                 return pair<node*,node*>(p.first,merge(p.second,v->right));
152
             else if(k > v->left->size) {
153
                 auto p = split(v->right,k-v->left->size);
154
155
                 return pair<node*,node*>(merge(v->left,p.first),p.second);
156
157
             else return pair<node*,node*>(v->left,v->right);
158
159
160
         node *insert(int k. T val) { return insert(new node(val).k):}
         node *erase(int k) {
161
162
             auto p = split(root,k+1);
             return root = merge(split(p.first,k).first, p.second);
163
164
         void add(int 1. int r. T val) { add(root. r. val): add(root. 1. -val):}
165
166
        T get(int k) { return get(root, k);}
167
         T minimum(int 1, int r) { return minimum(root, 1, r);}
         T operator[](const int &i) { return get(i);}
168
169
```

#### 6.10 永続赤黒木

62

```
//const int MAX = 15000000, BOUND = 14000000;
   template < class T> class prbtree {
   public:
        enum COL { BLACK, RED,};
       struct node {
            T val;
            int color:
            int rnk. size:
            node *left, *right;
10
11
            node(){}
12
            node(T v) : val(v), color(BLACK), rnk(0), size(1) {
                left = right = NULL:
13
14
            node(node *1, node *r, int c) : color(c) {
15
16
                left = 1:
                right = r;
17
                rnk = max((1? 1->rnk+(1->color==BLACK): 0).
18
19
                           (r? r->rnk+(r->color==BLACK): 0));
                size = !1 and !r? 1: !1? r->size: !r? r->size: 1->size+r->size;
20
21
       };
22
23
24
       node *root:
                  node nodes[MAX]:
25
       //
26
                  int called:
27
28
       prbtree() {
            root = NULL:
29
            // called = 0;
30
       }
31
32
       prbtree(T val) {
33
            root = new_node(val);
34
35
            // called = 0;
36
37
38
        // node *new_node(T v) { return &(nodes[called++] = node(v));}
39
        // node *new_node(node *1, node *r, int c) { return &(nodes[called++] = node(1,r,c
             )):}
       node *new_node(T v) { return new node(v);}
40
       node *new_node(node *1, node *r, int c) { return new node(1,r,c);}
41
42
       node *merge_sub(node *u, node *v) {
43
            if(u->rnk < v->rnk) {
44
45
                node *w = merge_sub(u,v->left);
46
                if(v->color == BLACK and w->color == RED and w->left->color == RED){
                    if(v->right->color == BLACK) return new_node(w->left,new_node(w->right,
47
                         v->right, RED), BLACK);
                    else return new_node(new_node(w->left,w->right,BLACK),new_node(v->right
48
                         ->left, v->right->right, BLACK), RED);
49
50
                else return new_node(w,v->right,v->color);
51
52
            else if(u \rightarrow rnk > v \rightarrow rnk) {
53
                node *w = merge_sub(u->right,v);
                if(u->color == BLACK and w->color == RED and w->right->color == RED){
54
                    if(u->left->color == BLACK) return new_node(new_node(u->left,w->left,
55
                         RED).w->right.BLACK):
56
                    else return new_node(new_node(u->left->left,u->left->right,BLACK),
                         new_node(w->left,w->right,BLACK),RED);
57
58
                else return new_node(u->left,w,u->color);
59
60
            else return new node(u.v.RED):
61
```

```
node *merge(node *u, node *v) {
63
64
             if(!u) return v;
65
             if(!v) return u:
            u = merge\_sub(u,v);
66
67
             if(u->color == RED) return new_node(u->left,u->right,BLACK);
             return u;
68
        }
69
70
71
        pair<node*,node*> split(node *v, int k) {
72
             if(!k) return pair<node*,node*>(NULL,v);
73
             if(k == v->size) return pair<node*, node*>(v, NULL);
            if(k < v->left->size) {
74
75
                 auto p = split(v->left,k);
                 return pair<node*,node*>(p.first,merge(p.second,v->right));
76
77
             else if(k > v->left->size) {
78
                 auto p = split(v->right,k-v->left->size);
79
80
                 return pair<node*,node*>(merge(v->left,p.first),p.second);
81
             else return pair<node*,node*>(v->left,v->right);
82
        }
83
84
        node *build(const vector<T> &vs) {
85
             if(!vs.size()) return NULL;
86
             if((int)vs.size() == 1) return new_node(vs[0]);
87
88
             int m = vs.size()/2;
             return merge(build(vector<T>(begin(vs),begin(vs)+m)), build(vector<T>(begin(vs)+
89
                 m, end(vs))));
90
91
92
         int size() { return root->size;}
93
         void get(vector<T> &vs) { get(root,vs);}
94
95
         void get(node *v, vector<T> &vs) {
             if(!v->left and !v->right) vs.push_back(v->val);
96
97
             else {
98
                 if(v->left) get(v->left,vs);
99
                 if(v->right) get(v->right, vs);
100
        }
101
102
103
         node *push_back(T val) {
104
             node *v = new_node(val);
105
             return root = merge(root, v);
106
107
108
         // insert leaf at k
         node *insert(int k, T val) {
109
110
             return insert(new_node(val), k);
111
112
113
         // insert tree v at k
        node *insert(node *v, int k) {
114
115
             auto p = split(root,k);
             return root = merge(merge(p.first,v),p.second);
116
117
118
119
         // copy [1,r)
120
        node *copy(int 1, int r) {
             return split(split(root, 1).second, r-1).first;
121
122
         // copy and insert [1,r) at k
123
124
        node *copy_paste(int 1, int r, int k) {
             return insert(copy(l,r),k);
125
126
127
   };
```

## 6.11 wavelet 行列

N := 列の長さ M := 最大値

### 6.11.1 完備辞書

function	計算量
count	O(1)
select	$O(\log N)$

```
template<int N> class FID {
       static const int bucket = 512, block = 16;
       static char popcount[];
       int n, B[N/bucket+10];
       unsigned short bs[N/block+10], b[N/block+10];
   public:
       FID(){}
       FID(int n, bool s[]) : n(n) {
            if(!popcount[1]) for (int i = 0; i < (1 << block); i++) popcount[i] =
10
                 __builtin_popcount(i);
11
            bs[0] = B[0] = b[0] = 0;
12
            for (int i = 0; i < n; i++) {
13
                if(i%block == 0) {
14
                    bs[i/block+1] = 0;
15
                    if(i%bucket == 0) {
16
                        B[i/bucket+1] = B[i/bucket];
17
                        b[i/block+1] = b[i/block] = 0;
18
19
                    else b[i/block+1] = b[i/block];
20
21
                bs[i/block] |= short(s[i])<<(i%block);</pre>
22
                b[i/block+1] += s[i];
23
24
                B[i/bucket+1] += s[i];
25
            if(n%bucket == 0) b[n/block] = 0;
26
27
28
        // number of val in [0,r), O(1)
29
       int count(bool val, int r) { return val? B[r/bucket]+b[r/block]+popcount[bs[r/block
30
            ]&((1<<(r%block))-1)]: r-count(1,r); }
31
        // number of val in [1,r), 0(1)
32
       int count(bool val, int 1, int r) { return count(val,r)-count(val,l); }
        // position of ith in val, 0-indexed, 0(log n)
33
34
       int select(bool val, int i) {
35
            if(i < 0 or count(val,n) <= i) return -1;
36
37
            int 1b = 0, ub = n, md;
            while(ub-lb>1) {
38
39
                md = (1b+ub)>>1;
40
                if(count(val,md) >= i) ub = md;
41
                else lb = md;
42
43
            return ub-1;
44
45
        int select(bool val, int i, int l) { return select(val,i+count(val,l)); }
        bool operator[](int i) { return bs[i/block]>>(i%block)&1: }
46
47 };
```

```
48 | template<int N> char FID<N>::popcount[1<<FID<N>::block];
```

#### **6.11.2** wavelet 行列

function	計算量	FID::count	FID::select
count	$O(\log M)$	О	
select	$O(\log N \log M)$	0	0
get	$O(\log M)$	0	
maximum	$O(\log M)$ or $O(k \log M)$	0	
kth_number	$O(\log M)$	О	
freq	$O(\log M)$	0	
freq_list	$O(k \log M)$	0	
get_rect	$O(k \log N \log M)$	0	О

```
template < class T, int N, int D> class wavelet {
         int n, zs[D];
        FID<N> dat[D];
         void max_dfs(int d, int 1, int r, int &k, T val, vector<T> &vs) {
             if(1 >= r or !k) return;
             if(d == D) {
                 while (1++ < r \text{ and } k > 0) vs.push_back(val), k--;
10
             int lc = dat[d].count(1,1), rc = dat[d].count(1,r);
11
             // if min, change this order
12
13
             \max_{dfs(d+1, lc+zs[d], rc+zs[d], k, lULL << (D-d-1)|val,vs)}
             max_dfs(d+1, l-lc, r-rc, k, val, vs);
14
15
16
        T max_dfs(int d, int l, int r, T val, T a, T b) {
17
             if(r-1 \le 0 \text{ or val} >= b) \text{ return } -1;
18
             if(d == D) return val>=a? val: -1;
19
             int lc = dat[d].count(1,1), rc = dat[d].count(1,r);
20
21
             T ret = \max_{d} ds(d+1, lc+zs[d], rc+zs[d], 1ULL << (D-d-1) | val, a, b);
             if(~ret) return ret;
22
             return max_dfs(d+1, l-lc, r-rc, val, a, b);
23
24
25
        int freq_dfs(int d, int l, int r, T val, T a, T b) {
             if(1 == r) return 0;
27
28
             if(d == D) return (a <= val and val < b)? r-1: 0;</pre>
29
             T \text{ nv} = 1ULL << (D-d-1) | val, nnv = ((1ULL << (D-d-1)) -1) | nv;
             if(nnv < a or b <= val) return 0;</pre>
30
31
             if(a <= val and nnv < b) return r-l;</pre>
             int lc = dat[d].count(1,1), rc = dat[d].count(1,r);
32
33
             return freq_dfs(d+1,1-lc,r-rc,val,a,b)+
34
                     freq_dfs(d+1,lc+zs[d],rc+zs[d],nv,a,b);
35
36
37
         void list_dfs(int d, int l, int r, T val, T a, T b, vector<pair<T,int>> &vs) {
38
             if(val >= b or r-1 <= 0) return;
             if(d == D) {
39
40
                 if(a <= val) vs.push_back(make_pair(val,r-1));</pre>
41
42
             T \text{ nv} = \text{val} | (1LL << (D-d-1)) \cdot \text{nnv} = \text{nv} | (((1LL << (D-d-1))-1)) :
43
44
             if(nnv < a) return;
```

```
int lc = dat[d].count(1.1), rc = dat[d].count(1.r):
45
             list_dfs(d+1,1-lc,r-rc,val,a,b,vs);
46
47
             list_dfs(d+1,lc+zs[d],rc+zs[d],nv,a,b,vs);
48
49
    public:
        wavelet(int n, T seq[]) : n(n) {
50
51
             T f[N], l[N], r[N];
             bool b[N];
52
             memcpy(f, seq, sizeof(T)*n);
53
             for (int d = 0; d < D; d++) {
54
55
                 int 1h = 0, rh = 0;
                 for (int i = 0: i < n: i++) {
56
                      bool k = (f[i] >> (D-d-1))&1;
57
58
                      if(k) r[rh++] = f[i];
                      else 1[lh++] = f[i];
59
                      b[i] = k;
60
61
                 dat[d] = FID < N > (n,b);
62
                 zs[d] = lh;
63
                 swap(1,f);
64
                 memcpy(f+lh, r, rh*sizeof(T));
65
66
67
        }
68
        T get(int i) {
69
             T ret = 0:
70
71
             bool b;
             for (int d = 0; d < D; d++) {
72
                 ret <<= 1;
73
                 b = dat[d][i];
74
                 ret |= b:
75
76
                 i = dat[d].count(b,i)+b*zs[d];
77
78
             return ret;
79
        T operator[](int i) { return get(i); }
80
81
        int count(T val, int l, int r) {
82
83
             for (int d = 0; d < D; d++) {
                 bool b = (val >> (D-d-1))&1;
84
                 1 = dat[d].count(b,1)+b*zs[d];
85
86
                 r = dat[d].count(b,r)+b*zs[d];
87
88
             return r-1;
89
        int count(T val, int r) { return count(val,0,r); }
90
91
        int select(T val, int k) {
92
             int ls[D], rs[D], l = 0, r = n;
93
94
             for (int d = 0; d < D; d++) {
95
                 ls[d] = 1; rs[d] = r;
                 bool b = val >> (D-d-1)\&1;
96
97
                 1 = dat[d].count(b.1)+b*zs[d]:
                 r = dat[d].count(b,r)+b*zs[d];
98
99
             for (int d = D-1; d >= 0; d--) {
100
                 bool b = val >> (D-d-1)\&1;
101
                 k = dat[d].select(b.k.ls[d]):
102
103
                 if (k >= rs[d] \text{ or } k < 0) \text{ return } -1;
104
                 k \rightarrow [d];
105
             return k;
106
107
        int select(T val, int k, int l) { return select(val,k+count(val,l)); }
108
109
        vector<T> maximum(int 1, int r, int k) {
110
111
             if (r-1 < k) k = r-1;
```

```
if(k < 0) return {}:
112
             vector<T> ret;
113
114
             max_dfs(0,1,r,k,0,ret);
115
             return ret:
116
117
118
         T \max(int 1, int r, T a, T b) \{ return \max_dfs(0,1,r,0,a,b); \}
119
120
         // k is 0-indexed
         T kth_number(int 1, int r, int k) {
121
122
             if(r-1 \le k \text{ or } k < 0) \text{ return } -1;
123
             T ret = 0:
             for (int d = 0; d < D; d++) {
124
125
                  int lc = dat[d].count(1,1), rc = dat[d].count(1,r);
                  if(rc-lc > k) {
126
                     l = lc+zs[d];
127
                     r = rc + zs[d]:
128
                      ret |= 1ULL << (D-d-1);
129
130
                  else {
131
                     k -= rc-lc;
132
                     1 -= 1c:
133
                     r -= rc;
134
135
             }
136
137
             return ret;
138
139
         vector<pair<T,int>> freq_list(int 1, int r, T a, T b) {
140
             vector<pair<T,int>> ret;
141
             list_dfs(0,1,r,0,a,b,ret);
142
143
             return ret;
        }
144
145
146
         vector<pair<int,T>> get_rect(int 1, int r, T a, T b) {
             vector<pair<T,int>> res = freq_list(l,r,a,b);
147
148
             vector<pair<int,T>> ret;
149
             for(auto &e: res)
150
                  for (int i = 0; i < e.second; i++)
                     ret.push_back(make_pair(select(e.first,i,l), e.first));
151
152
             return ret;
153
154
         // number of elements in [l,r) in [a,b), O(D)
155
         int freq(int 1, int r, T a, T b) { return freq_dfs(0,1,r,0,a,b); }
    };
156
```

# 7 その他

# 7.1 ビジュアライザ

```
cscript>
function line(x,y,a,b){c.b();c.moveTo(x,y);c.lineTo(a,b);c.s();}

function circle(x,y,r){c.b();c.arc(x,y,r,0,7,0);c.s();}

window.onload=function(){d=document;d.i=d.getElementById;
c=d.i('c').getContext('2d');c.b=c.beginPath;c.s=c.stroke;
d.i('s').src='data.js?';};

c/script>

cbody><canvas id="c" width="500" height="500" style="border:1px solid #000;"></canvas>

cscript id="s"></script>
```