



# CHAPTER 6

## Basic SQL

# Chapter 6 Outline

- SQL Data Definition and Data Types
- Specifying Constraints in SQL
- Basic Retrieval Queries in SQL
- INSERT, DELETE, and UPDATE Statements in SQL
- Additional Features of SQL

# Basic SQL

## ■ SQL language

- Considered one of the major reasons for the commercial success of relational databases

## ■ SQL

- The origin of SQL is relational predicate calculus called tuple calculus (see Ch.8) which was proposed initially as the language SQUARE.
- **SQL** Actually comes from the word “SEQUEL” which was the original term used in the paper: “SEQUEL TO SQUARE” by Chamberlin and Boyce. IBM could not copyright that term, so they abbreviated to SQL and copyrighted the term SQL.
- Now popularly known as “Structured Query language”.
- SQL is an informal or practical rendering of the relational data model with syntax

# SQL Data Definition, Data Types, Standards

- Terminology: *relational*
  - **Table**, **row**, and **column** used for relational model terms relation, tuple, and attribute
- CREATE statement
  - Main SQL command for data definition
- The language has features for : Data definition, Data Manipulation, Transaction control (Transact-SQL, Ch. 20), Indexing (Ch.17), Security specification (Grant and Revoke- see Ch.30), Active databases (Ch.26), Multi-media (Ch.26), Distributed databases (Ch.23) etc.

# SQL Standards

- **SQL-86** : ANSI Standard
- **SQL1(SQL-89)** : Integrity constraint
- **SQL2(1992)**
  - Data types extensions (varchar, bit, date, time, interval), Various JOIN operators, Catalogs, Domains, Derived tables in FROM clause, Assertions, Temporary tables, Referential integrity options, Schema manipulation language, Dynamic SQL, Scrollable cursors
- **SQL3(1999-)**
  - ~~Object Relational extensions~~(User-defined data types, Reference types, Collection types, Large object support (LOBs), Table hierarchies), Triggers, Stored procedures and User-defined functions, Recursive queries, SQL procedural constructs, Savepoints, Update through unions and joins
- **SQL4(2003)**
  - Bug fixes and enhancement to all 8 parts of SQL3, One new part (SQL/XML)
- **SQL5(2008)**
  - 3D-Data, XML <--> Tables, Fuzzy String Search, Distributed Request

# Schema and Catalog Concepts in SQL (1)

- We cover the basic standard SQL syntax – there are variations in existing RDBMS systems
- **SQL schema**
  - Identified by a schema name
  - Includes an authorization identifier and descriptors for each element
- **Schema elements** include
  - Tables, constraints, views, domains, and other constructs
- Each statement in SQL ends with a **semicolon**

# Schema and Catalog Concepts in SQL (2)

- CREATE SCHEMA statement
  - CREATE SCHEMA COMPANY AUTHORIZATION 'Jsmith' ;
- **Catalog**
  - Named collection of schemas in an SQL environment
- SQL also has the concept of a cluster of catalogs.



# The CREATE TABLE Command in SQL (1)

- Specifying a new relation
  - Provide name of table
  - Specify attributes, their types and initial constraints
- Can optionally specify schema:
  - `CREATE TABLE COMPANY.EMPLOYEE ...`  
or
  - `CREATE TABLE EMPLOYEE ...`

# The CREATE TABLE Command in SQL (2)

- **Base tables (base relations)** real table
  - Relation and its tuples are actually created and stored as a file by the DBMS
- **Virtual relations (views)**
  - Created through the `CREATE VIEW` statement.  
Do not correspond to any physical file.

# COMPANY relational database schema (Fig. 5.7)

## EMPLOYEE

Fname	Minit	Lname	<u>Ssn</u>	Bdate	Address	Sex	Salary	Super_ssn	Dno
-------	-------	-------	------------	-------	---------	-----	--------	-----------	-----

## DEPARTMENT

Dname	<u>Dnumber</u>	Mgr_ssn	Mgr_start_date
-------	----------------	---------	----------------

## DEPT\_LOCATIONS

<u>Dnumber</u>	<u>Dlocation</u>
----------------	------------------

## PROJECT

Pname	<u>Pnumber</u>	Plocation	Dnum
-------	----------------	-----------	------

## WORKS\_ON

<u>Essn</u>	<u>Pno</u>	Hours
-------------	------------	-------

## DEPENDENT

<u>Essn</u>	<u>Dependent_name</u>	Sex	Bdate	Relationship
-------------	-----------------------	-----	-------	--------------

# One possible database state for the COMPANY relational database schema (Fig. 5.6) (1)

## EMPLOYEE

Fname	Minit	Lname	<u>Ssn</u>	Bdate	Address	Sex	Salary	Super_ssn	Dno
John	B	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	M	30000	333445555	5
Franklin	T	Wong	333445555	1955-12-08	638 Voss, Houston, TX	M	40000	888665555	5
Alicia	J	Zelaya	999887777	1968-01-19	3321 Castle, Spring, TX	F	25000	987654321	4
Jennifer	S	Wallace	987654321	1941-06-20	291 Berry, Bellaire, TX	F	43000	888665555	4
Ramesh	K	Narayan	666884444	1962-09-15	975 Fire Oak, Humble, TX	M	38000	333445555	5
Joyce	A	English	453453453	1972-07-31	5631 Rice, Houston, TX	F	25000	333445555	5
Ahmad	V	Jabbar	987987987	1969-03-29	980 Dallas, Houston, TX	M	25000	987654321	4
James	E	Borg	888665555	1937-11-10	450 Stone, Houston, TX	M	55000	NULL	1

## DEPARTMENT

Dname	<u>Dnumber</u>	Mgr_ssn	Mgr_start_date
Research	5	333445555	1988-05-22
Administration	4	987654321	1995-01-01
Headquarters	1	888665555	1981-06-19

## DEPT\_LOCATIONS

<u>Dnumber</u>	<u>Dlocation</u>
1	Houston
4	Stafford
5	Bellaire
5	Sugarland
5	Houston

# One possible database state for the COMPANY relational database schema (Fig. 5.6) (2)

**WORKS\_ON**

<u>Essn</u>	<u>Pno</u>	Hours
123456789	1	32.5
123456789	2	7.5
666884444	3	40.0
453453453	1	20.0
453453453	2	20.0
333445555	2	10.0
333445555	3	10.0
333445555	10	10.0
333445555	20	10.0
999887777	30	30.0
999887777	10	10.0
987987987	10	35.0
987987987	30	5.0
987654321	30	20.0
987654321	20	15.0
888665555	20	NULL

**PROJECT**

<u>Pname</u>	<u>Pnumber</u>	Plocation	Dnum
ProductX	1	Bellaire	5
ProductY	2	Sugarland	5
ProductZ	3	Houston	5
Computerization	10	Stafford	4
Reorganization	20	Houston	1
Newbenefits	30	Stafford	4

**DEPENDENT**

<u>Essn</u>	<u>Dependent_name</u>	Sex	Bdate	Relationship
333445555	Alice	F	1986-04-05	Daughter
333445555	Theodore	M	1983-10-25	Son
333445555	Joy	F	1958-05-03	Spouse
987654321	Abner	M	1942-02-28	Spouse
123456789	Michael	M	1988-01-04	Son
123456789	Alice	F	1988-12-30	Daughter
123456789	Elizabeth	F	1967-05-05	Spouse

# SQL CREATE TABLE data definition statements for defining the COMPANY schema from Figure 5.7 (Fig. 6.1) (1)

**CREATE TABLE EMPLOYEE**

( Fname	VARCHAR(15)	NOT NULL,
Minit	CHAR,	
Lname	VARCHAR(15)	NOT NULL,
Ssn	CHAR(9)	NOT NULL,
Bdate	DATE,	
Address	VARCHAR(30),	
Sex	CHAR,	
Salary	DECIMAL(10,2),	
Super_ssn	CHAR(9),	
Dno	INT	NOT NULL,

**PRIMARY KEY (Ssn),**

**CREATE TABLE DEPARTMENT**

( Dname	VARCHAR(15)	NOT NULL,
Dnumber	INT	NOT NULL,
Mgr_ssn	CHAR(9)	NOT NULL,
Mgr_start_date	DATE,	

**PRIMARY KEY (Dnumber),**

**UNIQUE (Dname),**  **key candidate**

~~**FOREIGN KEY (Mgr\_ssn)**~~ **REFERENCES EMPLOYEE(Ssn) );**

**CREATE TABLE DEPT\_LOCATIONS**

( Dnumber	INT	NOT NULL,
Dlocation	VARCHAR(15)	NOT NULL,

**PRIMARY KEY (Dnumber, Dlocation),**

**FOREIGN KEY (Dnumber) REFERENCES DEPARTMENT(Dnumber) );**

*continued on next slide*

# SQL CREATE TABLE data definition statements for defining the COMPANY schema from Figure 5.7 (Fig. 6.1) (2)

**CREATE TABLE PROJECT**

( Pname	VARCHAR(15)	NOT NULL,
Pnumber	INT	NOT NULL,
Plocation	VARCHAR(15),	
Dnum	INT	NOT NULL,

**PRIMARY KEY** (Pnumber),

**UNIQUE** (Pname),

**FOREIGN KEY** (Dnum) **REFERENCES** DEPARTMENT(Dnumber) );

**CREATE TABLE WORKS\_ON**

( Essn	CHAR(9)	NOT NULL,
Pno	INT	NOT NULL,
Hours	DECIMAL(3,1)	NOT NULL,

**PRIMARY KEY** (Essn, Pno),

**FOREIGN KEY** (Essn) **REFERENCES** EMPLOYEE(Ssn),

**FOREIGN KEY** (Pno) **REFERENCES** PROJECT(Pnumber) );

**CREATE TABLE DEPENDENT**

( Essn	CHAR(9)	NOT NULL,
Dependent_name	VARCHAR(15)	NOT NULL,
Sex	CHAR,	
Bdate	DATE,	
Relationship	VARCHAR(8),	

**PRIMARY KEY** (Essn, Dependent\_name),

**FOREIGN KEY** (Essn) **REFERENCES** EMPLOYEE(Ssn) );

# Attribute Data Types and Domains in SQL (1)

- **Basic data types**

- **Numeric data types**

- Integer numbers: INTEGER, INT, and SMALLINT
    - Floating-point (real) numbers: FLOAT or REAL, and DOUBLE PRECISION

- **Character-string data types**

- Fixed length: CHAR(*n*), CHARACTER(*n*)
    - Varying length: VARCHAR(*n*), CHAR VARYING(*n*), CHARACTER VARYING(*n*)



# Attribute Data Types and Domains in SQL (2)

- **Bit-string** data types
  - Fixed length: `BIT (n)`
  - Varying length: `BIT VARYING (n)`
- **Boolean** data type
  - Values of `TRUE` or `FALSE` or `NULL`
- **DATE** data type
  - Ten positions
  - Components are `YEAR`, `MONTH`, and `DAY` in the form `YYYY-MM-DD`
  - Multiple mapping functions available in RDBMSs to change date formats

# Attribute Data Types and Domains in SQL (3)

- Additional data types

- **Timestamp** data type '2008-09-27 09:12:47.648302'

Includes the `DATE` and `TIME` fields

- Plus a minimum of six positions for decimal fractions of seconds

- Optional `WITH TIME ZONE` qualifier

- **INTERVAL** data type

- Specifies a relative value that can be used to increment or decrement an absolute value of a date, time, or timestamp

- **DATE, TIME, Timestamp, INTERVAL** data types can be cast or converted to string formats for comparison.

# Attribute Data Types and Domains in SQL (4)

## ■ Domain

- Name used with the attribute specification
- Makes it easier to change the data type for a domain that is used by numerous attributes
- Improves schema readability
- Example:

like typedef

- `CREATE DOMAIN SSN_TYPE AS CHAR(9);`

## ■ TYPE

- User Defined Types (UDTs) are supported for object-oriented applications. (See Ch.12) Uses the command: `CREATE TYPE`

# Specifying Constraints in SQL

## Basic constraints:

- Relational Model has 3 basic constraint types that are supported in SQL:
  - **Key** constraint: A primary key value cannot be duplicated
  - **Entity Integrity** Constraint: A primary key value cannot be null
  - **Referential integrity** constraints : The “foreign key” must have a value that is already present as a primary key, or may be null.

# Specifying Attribute Constraints

Other Restrictions on attribute domains:

- Default value of an attribute
  - **DEFAULT** <value>
  - NULL is not permitted for a particular attribute (NOT NULL)
- **CHECK** clause
  - Dnumber ~~/INT/~~ ~~NOT NULL/~~ ~~CHECK~~ (Dnumber > 0 AND Dnumber < 21);

# Specifying Key and Referential Integrity Constraints (1)

- **PRIMARY KEY** clause

- Specifies one or more attributes that make up the primary key of a relation

- Dnumber /INT /PRIMARY KEY;

- **UNIQUE** clause

- Specifies alternate (secondary) keys (called CANDIDATE keys in the relational model).

- Dname /VARCHAR(15) /UNIQUE;

# Specifying Key and Referential Integrity Constraints (2)

- **FOREIGN KEY** clause

- 1 ■ Default operation: reject update on violation
- 2 ■ Attach **referential triggered action** clause
  - Options include SET NULL, CASCADE, and SET DEFAULT
  - Action taken by the DBMS for SET NULL or SET DEFAULT is the same for both ON DELETE and ON UPDATE
  - CASCADE option suitable for “relationship” relations

# Giving Names to Constraints

- Using the Keyword **CONSTRAINT**
  - Name a constraint
  - Useful for later altering



# Default attribute values and referential integrity triggered action specification (Fig. 6.2)

```
CREATE TABLE EMPLOYEE
(
    ...,
    Dno          INT          NOT NULL          DEFAULT 1,
    CONSTRAINT EMPPK
        PRIMARY KEY (Ssn),
    CONSTRAINT EMPSUPERFK
        FOREIGN KEY (Super_ssn) REFERENCES EMPLOYEE(Ssn)
            ON DELETE SET NULL          ON UPDATE CASCADE,
    CONSTRAINT EMPDEPTFK
        FOREIGN KEY (Dno) REFERENCES DEPARTMENT(Dnumber)
            ON DELETE SET DEFAULT       ON UPDATE CASCADE);
CREATE TABLE DEPARTMENT
(
    ...,
    Mgr_ssn CHAR(9)          NOT NULL          DEFAULT '888665555',
    ...,
    CONSTRAINT DEPTPK
        PRIMARY KEY (Dnumber),
    CONSTRAINT DEPTSK
        UNIQUE (Dname),
    CONSTRAINT DEPTMGRFK
        FOREIGN KEY (Mgr_ssn) REFERENCES EMPLOYEE(Ssn)
            ON DELETE SET DEFAULT       ON UPDATE CASCADE);
CREATE TABLE DEPT_LOCATIONS
(
    ...,
    PRIMARY KEY (Dnumber, Dlocation),
    FOREIGN KEY (Dnumber) REFERENCES DEPARTMENT(Dnumber)
        ON DELETE CASCADE              ON UPDATE CASCADE);
```

연쇄적으로  
업데이트

# Specifying Constraints on Tuples Using CHECK

- Additional Constraints on individual tuples within a relation are also possible using CHECK
- CHECK clauses at the end of a CREATE TABLE statement
  - Apply to each tuple individually
  - CHECK (Dept\_create\_date <= Mgr\_start\_date);

# Basic Retrieval Queries in SQL

- **SELECT statement**
    - One basic statement for retrieving information from a database
  - SQL allows a table to have two or more tuples that are identical in all their attribute values
    - Unlike relational model (relational model is strictly set-theory based)
    - Multiset or bag behavior
    - Tuple-id may be used as a key
- Handwritten notes: "same" with an arrow pointing to "identical", and a red box around "identical".*

# The SELECT-FROM-WHERE

## Structure of Basic SQL Queries (1)

### ■ Basic form of the SELECT statement:

```
SELECT    <attribute list>  
FROM      <table list>  
WHERE     <condition>;
```

where

- <attribute list> is a list of attribute names whose values are to be retrieved by the query.
- <table list> is a list of the relation names required to process the query.
- <condition> is a conditional (Boolean) expression that identifies the tuples to be retrieved by the query.

# The SELECT-FROM-WHERE Structure of Basic SQL Queries (2)

- Logical comparison operators
  - =, <, <=, >, >=, and <>
- **Projection attributes**
  - Attributes whose values are to be retrieved
- **Selection condition**
  - Boolean condition that must be true for any retrieved tuple. Selection conditions include join conditions (see Ch.8) when multiple relations are involved.

# Basic Retrieval Queries (1)

<u>Bdate</u>	<u>Address</u>
1965-01-09	731 Fondren, Houston, TX

<u>Fname</u>	<u>Lname</u>	<u>Address</u>
John	Smith	731 Fondren, Houston, TX
Franklin	Wong	638 Voss, Houston, TX
Ramesh	Narayan	975 Fire Oak, Humble, TX
Joyce	English	5631 Rice, Houston, TX

**Query 0.** Retrieve the birth date and address of the employee(s) whose name is 'John B. Smith'.

**Q0:**     **SELECT**     Bdate, Address  
          **FROM**     EMPLOYEE  
          **WHERE**    Fname='John' **AND** Minit='B' **AND** Lname='Smith';

**Query 1.** Retrieve the name and address of all employees who work for the 'Research' department.

**Q1:**     **SELECT**     Fname, Lname, Address  
          **FROM**     EMPLOYEE, DEPARTMENT  
          **WHERE**    Dname='Research' **AND** Dnumber=Dno;

# Basic Retrieval Queries (2)

(c)

<u>Pnumber</u>	<u>Dnum</u>	<u>Lname</u>	<u>Address</u>	<u>Bdate</u>
10	4	Wallace	291Berry, Bellaire, TX	1941-06-20
30	4	Wallace	291Berry, Bellaire, TX	1941-06-20

**Query 2.** For every project located in 'Stafford', list the project number, the controlling department number, and the department manager's last name, address, and birth date.

**Q2:**      **SELECT**      Pnumber, Dnum, Lname, Address, Bdate  
             **FROM**        PROJECT, DEPARTMENT, EMPLOYEE  
             **WHERE**      Dnum=Dnumber **AND** Mgr\_ssn=Ssn **AND**  
                         Plocation='Stafford';

# Ambiguous Attribute Names

- Same name can be used for two (or more) attributes in different relations
  - As long as the attributes are in different relations
  - Must **qualify** the attribute name with the relation name to prevent ambiguity

```
Q1A:  SELECT  Fname, EMPLOYEE.Name, Address
        FROM    EMPLOYEE, DEPARTMENT
        WHERE   DEPARTMENT.Name='Research' AND
                DEPARTMENT.Dnumber=EMPLOYEE.Dnumber;
```



# Aliasing, and Renaming and Tuple Variables (1)

## ■ Aliases or tuple variables

- Declare alternative relation names E and S to refer to the EMPLOYEE relation twice in a query:

**Query 8.** For each employee, retrieve the employee's first and last name and the first and last name of his or her immediate supervisor.

```
SELECT E.Fname, E.Lname, S.Fname, S.Lname  
FROM EMPLOYEE AS E, EMPLOYEE AS S  
WHERE E.Super_ssn=S.Ssn;
```

- Recommended practice to abbreviate names and to prefix same or similar attribute from multiple tables.

# Aliasing, Renaming and Tuple Variables (2)

- The attribute names can also be renamed

```
EMPLOYEE AS E(Fn, Mi, Ln, Ssn, Bd,  
  Addr, Sex, Sal, Sssn, Dno)
```

- Note that the relation EMPLOYEE now has a variable name E which corresponds to a tuple variable
- The “AS” may be dropped in most SQL implementations

# Unspecified WHERE Clause and Use of the Asterisk (1)

- Missing WHERE clause
  - Indicates no condition on tuple selection
- Effect is a CROSS PRODUCT
  - Result is all possible tuple combinations (or the Algebra operation of Cartesian Product– see Ch.8)

Queries 9 and 10. Select all EMPLOYEE Ssns (Q9) and all combinations of EMPLOYEE Ssn and DEPARTMENT Dname (Q10) in the database.

Q9:     SELECT     Ssn  
          FROM     EMPLOYEE;

Q10:    SELECT     Ssn, Dname  
          FROM     EMPLOYEE, DEPARTMENT;

모든 조합 반환

# Unspecified WHERE Clause and Use of the Asterisk (\*)

- Specify an asterisk (\*)
  - Retrieve all the attribute values of the selected tuples
  - The \* can be prefixed by the relation name; e.g., EMPLOYEE \*

Q1C:    SELECT    \*  
         FROM     EMPLOYEE  
         WHERE    Dno=5;

Q1D:    SELECT    \*  
         FROM     EMPLOYEE, DEPARTMENT  
         WHERE    Dname='Research' AND Dno=Dnumber;

Q10A:   SELECT    \*  
         FROM     EMPLOYEE, DEPARTMENT;

# Tables as Sets in SQL (1)

- SQL does not automatically eliminate duplicate tuples in query results
- For aggregate operations (See sec 7.1.7) duplicates must be accounted for
- Use the keyword **DISTINCT** in the SELECT clause
  - Only distinct tuples should remain in the result

Query 11. Retrieve the salary of every employee (Q11) and all distinct salary values (Q11A).

Q11:     SELECT     ~~ALL~~ Salary     can be omitted  
          FROM       EMPLOYEE;

Q11A:    SELECT     DISTINCT Salary  
          FROM       EMPLOYEE;

# Tables as Sets in SQL (2)

## ■ Set operations

- **UNION, EXCEPT (difference), INTERSECT**
- Corresponding multiset operations: UNION ALL, EXCEPT ALL, INTERSECT ALL)
- Type compatibility is needed for these operations to be valid

**Query 4.** Make a list of all project numbers for projects that involve an employee whose last name is 'Smith', either as a worker or as a manager of the department that controls the project.

```
Q4A: (SELECT   DISTINCT Pnumber
      FROM     PROJECT, DEPARTMENT, EMPLOYEE
      WHERE    Dnum=Dnumber AND Mgr_ssn=Ssn
              AND Lname='Smith' )

      UNION

      ( SELECT   DISTINCT Pnumber
        FROM     PROJECT, WORKS_ON, EMPLOYEE
        WHERE    Pnumber=Pno AND Essn=Ssn
              AND Lname='Smith' );
```

# Substring Pattern Matching and Arithmetic Operators

- **LIKE** comparison operator
  - Used for string **pattern matching**
  - % replaces an arbitrary number of zero or more characters
  - underscore (\_) replaces a single character
  - Examples: **WHERE** Address **LIKE** '%Houston,TX%';
  - **WHERE** Ssn **LIKE** ' \_ \_ 1 \_ \_ 8901';
- **BETWEEN** comparison operator

E.g., in Q14 :  $30000 \leq \text{Salary} \leq 40000$

**WHERE**(Salary **BETWEEN** 30000 **AND** 40000)  
**AND** Dno = 5;

# Arithmetic Operations

- Standard arithmetic operators:
  - Addition (+), subtraction (−), multiplication (\*), and division (/) may be included as a part of **SELECT**
- **Query 13.** Show the resulting salaries if every employee working on the 'ProductX' project is given a 10 percent raise.

아래 조건에 해당하는 e.salary를 찾고 거기에  
1.1을 곱한값을 보여줌

```
SELECT E.Fname, E.Lname, 1.1 * E.Salary AS Increased_sal  
FROM EMPLOYEE AS E, WORKS_ON AS W, PROJECT AS P  
WHERE E.Ssn=W.Essn AND W.Pno=P.Pnumber AND  
P.Pname='ProductX';
```



# Ordering of Query Results

- Use **ORDER BY** clause
  - Keyword DESC to see result in a descending order of values
  - Keyword ASC to specify ascending order explicitly
  - Typically placed at the end of the query

```
ORDER BY D.Dname DESC, E.Lname ASC,  
        E.Fname ASC
```

# Basic SQL Retrieval Query Block

```
SELECT    <attribute list>  
FROM      <table list>  
[ WHERE    <condition> ]  
[ ORDER BY <attribute list> ];
```

# INSERT, DELETE, and UPDATE Statements in SQL

- Three commands used to modify the database:
  - INSERT, DELETE, and UPDATE
- INSERT typically inserts a tuple (row) in a relation (table)
- UPDATE may update a number of tuples (rows) in a relation (table) that satisfy the condition
- DELETE may also update a number of tuples (rows) in a relation (table) that satisfy the condition

# INSERT

- In its simplest form, it is used to add one or more tuples to a relation
- Attribute values should be listed in the same order as the attributes were specified in the **CREATE TABLE** command
- Any integrity constraints as a part of the DDL specification are enforced

# The INSERT Command

- Specify the relation name and a list of values for the tuple. All values including nulls are supplied.

```
U1:    INSERT INTO    EMPLOYEE
        VALUES      ( 'Richard', 'K', 'Marini', '653298653', '1962-12-30', '98
                        Oak Forest, Katy, TX', 'M', 37000, '653298653', 4 );
```

- The variation below inserts multiple tuples where a new table is loaded values from the result of a query.

```
U3B:    INSERT INTO    WORKS_ON_INFO ( Emp_name, Proj_name,
                                         Hours_per_week )
        SELECT          E.Lname, P.Pname, W.Hours
        FROM             PROJECT P, WORKS_ON W, EMPLOYEE E
        WHERE            P.Pnumber=W.Pno AND W.Essn=E.Ssn;
```

# BULK LOADING OF TABLES

- Another variation of **INSERT** is used for bulk-loading of several tuples into tables
- A new table TNEW can be created with the same attributes as T and using LIKE and DATA in the syntax, it can be loaded with entire data.
- **EXAMPLE:**

```
CREATE TABLE D5EMPS LIKE EMPLOYEE
```

```
(SELECT E.*  
FROM EMPLOYEE AS E  
WHERE E.Dno=5)
```

```
WITH DATA;
```

“생성된 D5EMPS에 EMPLOYEE와 같은 애트리뷰트들이 생성되고  
검색된 전체 데이터가 적재된다.”

# DELETE

- Removes tuples from a relation
  - Includes a WHERE-clause to select the tuples to be deleted
  - Referential integrity should be enforced
  - Tuples are deleted from only *one table* at a time (unless CASCADE is specified on a referential integrity constraint)
  - A missing WHERE-clause specifies that *all tuples* in the relation are to be deleted; the table then becomes an empty table
  - The number of tuples deleted depends on the number of tuples in the relation that satisfy the WHERE-clause

# The DELETE Command

- Removes tuples from a relation
  - Includes a `WHERE` clause to select the tuples to be deleted. The number of tuples deleted will vary.

<b>U4A:</b>	<b>DELETE FROM</b>	<b>EMPLOYEE</b>
	<b>WHERE</b>	<b>Lname='Brown';</b>
<b>U4B:</b>	<b>DELETE FROM</b>	<b>EMPLOYEE</b>
	<b>WHERE</b>	<b>Ssn='123456789';</b>
<b>U4C:</b>	<b>DELETE FROM</b>	<b>EMPLOYEE</b>
	<b>WHERE</b>	<b>Dno=5;</b>
<b>U4D:</b>	<b>DELETE FROM</b>	<b>EMPLOYEE;</b>



# UPDATE

- Used to modify attribute values of one or more selected tuples
- A WHERE-clause selects the tuples to be modified
- An additional SET-clause specifies the attributes to be modified and their new values
- Each command modifies tuples *in the same relation*
- Referential integrity specified as part of DDL specification is enforced

# The UPDATE Command (1)

- Example: Change the location and controlling department number of project number 10 to 'Bellaire' and 5, respectively

U5:            UPDATE    PROJECT  
              SET        ① Plocation = 'Bellaire',  
                             ② Dnum = 5  
              WHERE     Pnumber = 10

# The UPDATE Command (2)

- Example: Give all employees in the 'Research' department a 10% raise in salary.

```
U6:UPDATE      EMPLOYEE
      SET      Salary = Salary *1.1
      WHERE    Dno IN (SELECT      Dnumber
                        FROM        DEPARTMENT
                        WHERE       Dname = 'Research')
```

해당 집합에 속하는지

- In this request, the modified Salary value depends on the original Salary value in each tuple
  - The reference to the Salary attribute on the right of = refers to the old Salary value before modification
  - The reference to the Salary attribute on the left of = refers to the new Salary value after modification

# Additional Features of SQL (1)

- Techniques for specifying complex retrieval queries (see Ch.7)
- Writing programs in various programming languages that include SQL statements: Embedded and dynamic SQL, SQL/CLI (Call Level Interface) and its predecessor ODBC, SQL/PSM (Persistent Stored Module) (See Ch.10)
- Set of commands for specifying physical database design parameters, file structures for relations, and access paths, e.g., CREATE INDEX

# Additional Features of SQL (2)

- Transaction control commands (Ch.20)
- Specifying the granting and revoking of privileges to users (Ch.30)
- Constructs for creating triggers (Ch.26)
- Enhanced relational systems known as object-relational define relations as classes. Abstract data types (called User Defined Types- UDTs) are supported with CREATE TYPE
- New technologies such as XML (Ch.13) and OLAP (Ch.29) are added to versions of SQL

# Summary

- SQL
  - A Comprehensive language for relational database management
  - Data definition, queries, updates, constraint specification, and view definition
- Covered :
  - Data definition commands for creating tables
  - Commands for constraint specification
  - Simple retrieval queries
  - Database update commands