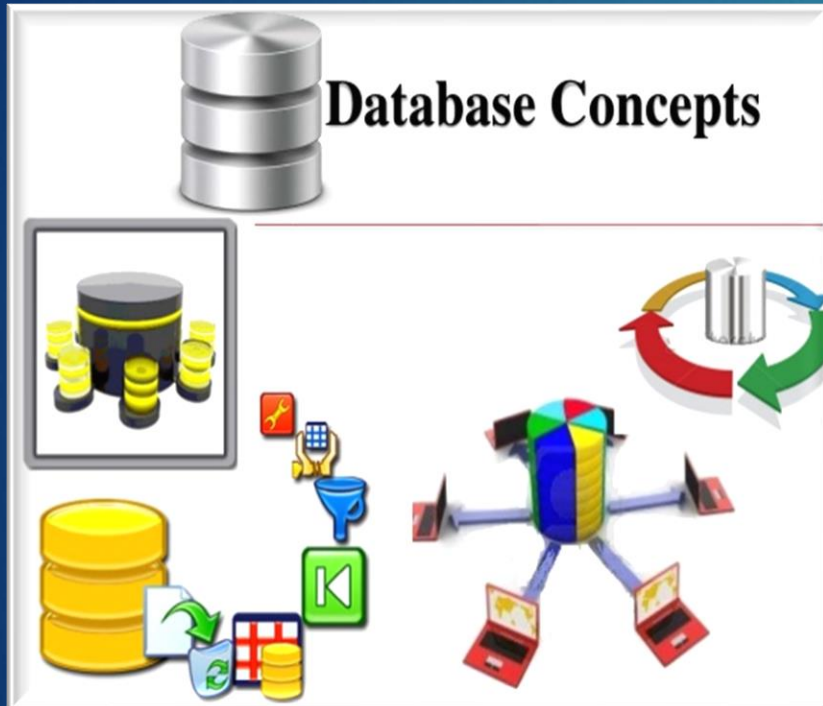


# Database Normalization



*Prof. K ADISHA (Ph. D)*

# Database Normalization



Introduction



Normalization



Database Anomaly



Functional Dependency



Normal Forms

*Prof. K. Adisesha*

# Introduction

## *Database Design Strategies:*

*Database design can be defined as a collection of tasks or processes that enhance the designing, development, implementation, and maintenance of enterprise data management system.*

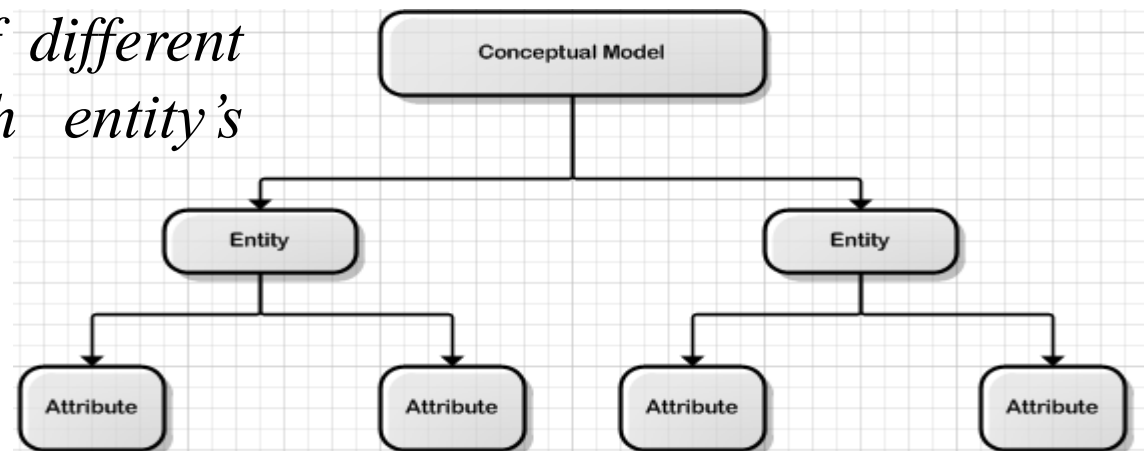
- *Designing a proper database reduces the maintenance cost thereby improving data consistency and the cost-effective measures are greatly influenced in terms of disk storage space.*
- *There are two approaches for developing any database:*
  - ❖ *The top-down method*
  - ❖ *The bottom-up method*

# Introduction

## *Top – down design method:*

*The top-down design method starts from the general and moves to the specific.*

- *It start with a general idea of what is needed for the system and then work your way down to the more specific details of how the system will interact.*
- *This process involves the identification of different entity types and the definition of each entity's attributes.*

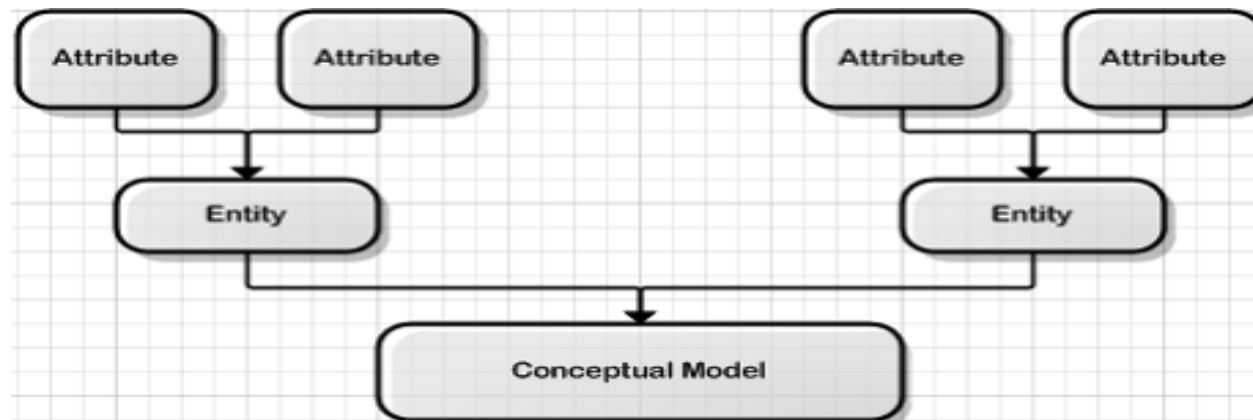


# Introduction

## *Bottom – up design method:*

*The bottom-up approach begins with the specific details and moves up to the general.*

- *This is done by first identifying the data elements (items) and then grouping them together in data sets.*
- *In other words, this method first identifies the attributes, and then groups them to form entities.*



# Introduction

*Normalization is the process of organizing the data and the attributes of a database.*

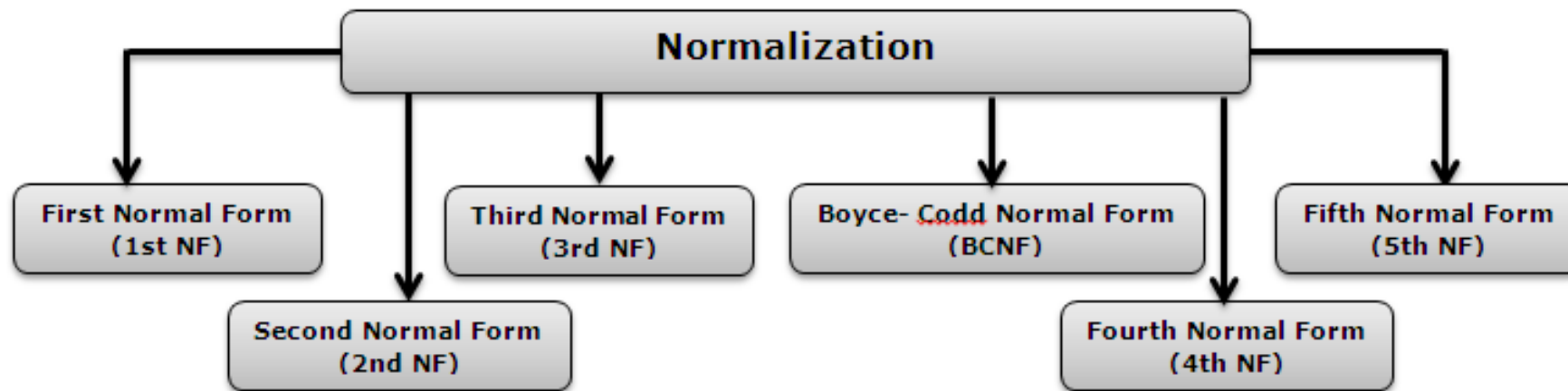
- *It is performed to reduce the data redundancy in a database and to ensure that data is stored logically.*
- *Data redundancy means having the same data but at multiple places.*
- *It is necessary to remove data redundancy because it causes anomalies in a database which makes it very hard for a database administrator to maintain it.*

# Introduction

## *Normalization:*

*Normalization is a step by step process of removing the different kinds of redundancy and anomaly one step at a time from the database.*

- E.F Codd developed for the relation data model in 1970.
- Normalization rules are divided into following normal form:





# Normalization

*Normalization is the process of organizing the data and the attributes of a database. The main reason for normalizing the relations is removing these anomalies.*

- *Normalization is the process of organizing the data in the database.*
- *Normalization is used to minimize the redundancy from a relation or set of relations. It is also used to eliminate undesirable characteristics like Insertion, Update, and Deletion Anomalies.*
- *Normalization divides the larger table into smaller and links them using relationships.*
- *The normal form is used to reduce redundancy from the database table.*



# Normalization

*The main reason for normalizing the relations is removing these anomalies.*

➤ ***Advantages of Normalization***

- ❖ *Normalization helps to minimize data redundancy.*
- ❖ *Greater overall database organization.*
- ❖ *Data consistency within the database.*
- ❖ *Much more flexible database design.*
- ❖ *Enforces the concept of relational integrity.*

➤ ***Disadvantages of Normalization***

- ❖ *It is very time-consuming and difficult to normalize relations of a higher degree.*
- ❖ *Careless decomposition may lead to a bad database design, leading to serious problems.*

# Normalization

## *Need for Normalization:*

*Normalization is used to reduce data redundancy. It provides a method to remove the following anomalies from the database and bring it to a more consistent state:*

- A **database anomaly** is a flaw in the database that occurs because of poor planning and redundancy.
  - ❖ **Insertion anomalies:** *This occurs when we are not able to insert data into a database because some attributes may be missing at the time of insertion.*
  - ❖ **Updation anomalies:** *This occurs when the same data items are repeated with the same values and are not linked to each other.*
  - ❖ **Deletion anomalies:** *This occurs when deleting one part of the data deletes the other necessary information from the database.*

# Normalization

## *Database Normal Forms.*

*The Theory of Data Normalization in MySQL server is still being developed further, there are discussions even on 6th Normal Form.*

➤ *However, in most practical applications, normalization achieves its best in 3rd Normal Form. The evolution of Normalization in SQL theories is illustrated below-*

- ❖ *1NF (First Normal Form)*
- ❖ *2NF (Second Normal Form)*
- ❖ *3NF (Third Normal Form)*
- ❖ *BCNF (Boyce-Codd Normal Form)*
- ❖ *4NF (Fourth Normal Form)*
- ❖ *5NF (Fifth Normal Form)*
- ❖ *6NF (Sixth Normal Form)*

1NF	2NF	3NF	BCNF
Atomic	Functional Dependency not Partial Dependency	Not Transitive Dependency	$X \rightarrow Y$ X is Super Key 0% redundancy (due to Functional Dependency)

# Normalization

## *Database Normal Forms.*

*To understand the above-mentioned normal forms, we first need to have an understanding of the functional dependencies.*

- *There are four types of normal forms that are usually used in relational databases*
  - ❖ **1NF:** *A relation is in 1NF if all its attributes have an atomic value.*
  - ❖ **2NF:** *A relation is in 2NF if it is in 1NF and all non-key attributes are fully functional dependent on the candidate key.*
  - ❖ **3NF:** *A relation is in 3NF if it is in 2NF and there is no transitive dependency.*
  - ❖ **BCNF:** *A relation is in BCNF if it is in 3NF and for every Functional Dependency, LHS is the super key.*

# Functional dependency

## *Functional dependency .*

*The functional dependency is a relationship that exists between two attributes. It typically exists between the primary key and non-key attribute within a table.*

- *The left side of FD is known as a determinant, the right side of the production is known as a dependent.*
- *It is denoted by  $X \rightarrow Y$ , where  $X$  is called a determinant and  $Y$  is called dependent.*
- *Example:*
  - ❖ *Assume we have an employee table with attributes:  $Emp\_Id$ ,  $Emp\_Name$ ,  $Emp\_Address$ .*
  - ❖ *Functional dependency can be written as:  $Emp\_Id \rightarrow Emp\_Name$*
  - ❖ *We can say that  $Emp\_Name$  is functionally dependent on  $Emp\_Id$ .*
  - ❖ *Here  $Emp\_Id$  attribute can uniquely identify the  $Emp\_Name$  attribute of employee table because if we know the  $Emp\_Id$ , we can tell that employee name associated with it.*

# Functional dependency

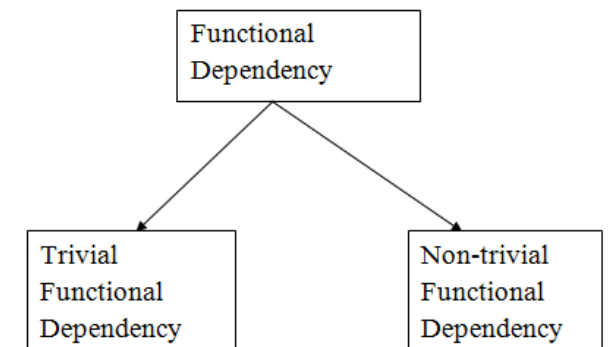
## *Functional dependency .*

*The functional dependency is a relationship that exists between two attributes. It typically exists between the primary key and non-key attribute within a table.*

### ➤ *Types of Functional dependency*

- ❖ *Trivial functional dependency*
- ❖ *Non-Trivial functional dependency*
- ❖ *Multivalued functional dependency*
- ❖ *Transitive functional dependency.*

roll_no	name	age
42	abc	17
43	pqr	18
44	xyz	18



- *Here, {roll\_no, name} → name is a trivial functional dependency, since the dependent name is a subset of determinant set {roll\_no, name}*
- *roll\_no → name is a non-trivial functional dependency, since the dependent name is not a subset of determinant roll\_no.*
- *If  $a \rightarrow \{b, c\}$  and there exists no functional dependency between b and c, then it is called a multivalued functional dependency.*
- *If  $a \rightarrow b$  &  $b \rightarrow c$ , then according to axiom of transitivity,  $a \rightarrow c$ . This is a transitive*



# Functional dependency

## *Functional dependency .*

*The functional dependency is a relationship that exists between two attributes. It typically exists between the primary key and non-key attribute within a table.*

### ➤ *Types of Functional dependency*

#### ❖ *Trivial functional dependency*

❖  *$A \rightarrow B$  has trivial functional dependency if  $B$  is a subset of  $A$ .*

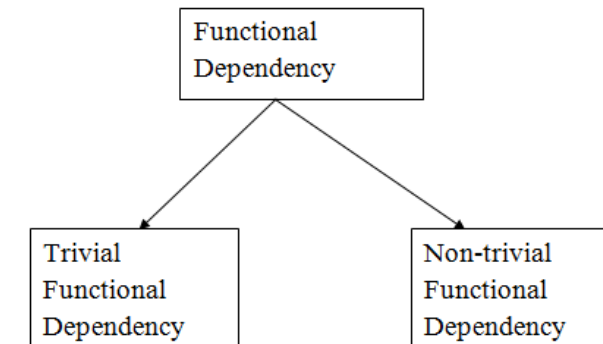
❖ *The following dependencies are also trivial like:  $A \rightarrow A$ ,  $B \rightarrow B$*

#### ❖ *Non-trivial functional dependency*

❖  *$A \rightarrow B$  has a non-trivial functional dependency if  $B$  is not a subset of  $A$ .*

❖ *When  $A$  intersection  $B$  is NULL, then  $A \rightarrow B$  is called as complete non-trivial.*

❖ *In Multivalued functional dependency: entities of the dependent set are not dependent on each other. i.e. If  $a \rightarrow \{b, c\}$  and there exists no functional dependency between  $b$  and  $c$ , then it is called a multivalued functional dependency.*





# Inference Rule (IR):

## *Inference Rule (IR):.*

*The Functional dependency has 6 types of inference rule:.*

- **Reflexive Rule (IR1):** In the reflexive rule, if  $Y$  is a subset of  $X$ , then  $X$  determines  $Y$ .

**If  $X \supseteq Y$  then  $X \rightarrow Y$**

- **Augmentation Rule (IR2):** The augmentation is also called as a partial dependency. In augmentation, if  $X$  determines  $Y$ , then  $XZ$  determines  $YZ$  for any  $Z$ . **If  $X \rightarrow Y$  then  $XZ \rightarrow YZ$**

- **Transitive Rule (IR3):** In the transitive rule, if  $X$  determines  $Y$  and  $Y$  determine  $Z$ , then  $X$  must also determine  $Z$ .

**If  $X \rightarrow Y$  and  $Y \rightarrow Z$  then  $X \rightarrow Z$**

- **Union Rule (IR4):** Union rule says, if  $X$  determines  $Y$  and  $X$  determines  $Z$ , then  $X$  must also determine  $Y$  and  $Z$ .

**If  $X \rightarrow Y$  and  $X \rightarrow Z$  then  $X \rightarrow YZ$**

- **Decomposition Rule (IR5):** Decomposition rule is also known as project rule. This Rule says, if  $X$  determines  $Y$  and  $Z$ , then  $X$  determines  $Y$  and  $X$  determines  $Z$  separately.

**If  $X \rightarrow YZ$  then  $X \rightarrow Y$  and  $X \rightarrow Z$**

- **Pseudo transitive Rule (IR6):** In Pseudo transitive Rule, if  $X$  determines  $Y$  and  $YZ$  determines  $W$ , then  $XZ$  determines  $W$ .

**If  $X \rightarrow Y$  and  $YZ \rightarrow W$  then  $XZ \rightarrow W$**

# Normal Form :

## *First Normal Form (1NF):*

**1NF:** *A relation is in 1NF if all its attributes have an atomic value.*

- *A relation will be 1NF if it contains an atomic value.*
- *It states that an attribute of a table cannot hold multiple values. It must hold only single-valued attribute.*
- *First normal form disallows the multi-valued attribute, composite attribute, and their combinations.*

S_ID	S_NAME	S_PHONE	S_COURSE
101	SUNNY	72XXXX9064 738XXXX238	BCA
102	PRAJWAL	857XXXX832	BBA



S_ID	S_NAME	S_PHONE	S_COURSE
101	SUNNY	72XXXX9064	BCA
101	SUNNY	738XXXX238	BCA
102	PRAJWAL	857XXXX832	BBA

# Normal Form :

## *Second Normal Form (2NF):*

**2NF:** *The first condition for the table to be in Second Normal Form is that the table has to be in First Normal Form. The table should not possess partial dependency.*

- *In the 2NF, relational must be in 1NF.*
- *In the second normal form, all non-key attributes are fully functional dependent on the primary key*

S_ID	S_NAME	S_COURSE
101	SUNNY	BCA
102	SHAILU	BCA
103	PRAJWAL	BBA



**Second Normal Form (2NF)**

S_ID	S_NAME	C_ID
101	SUNNY	C01
102	SHAILU	C01
103	PRAJWAL	CO2

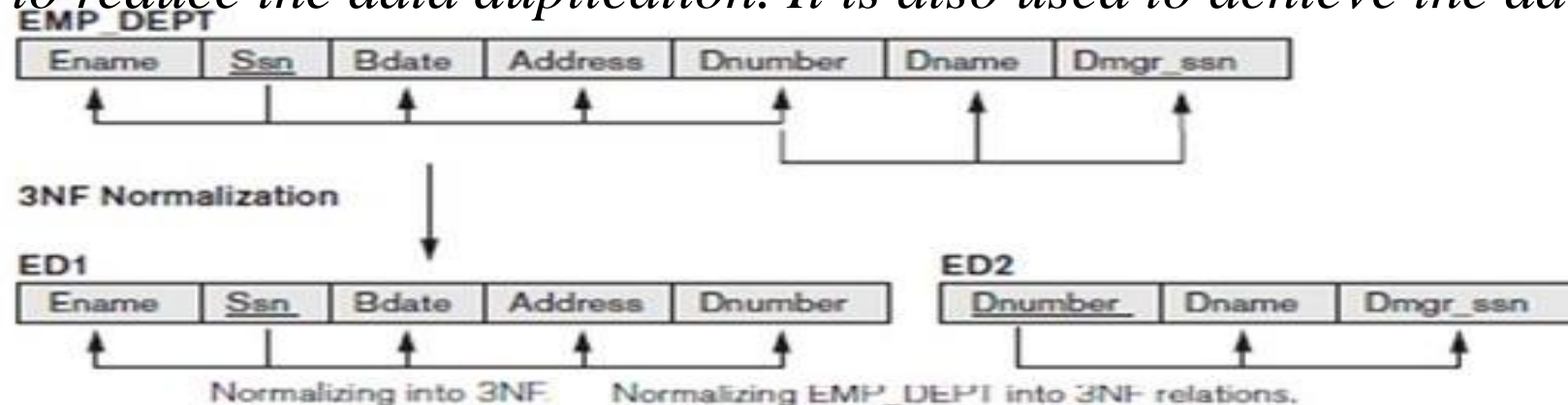
C_ID	S_COURSE
C01	BCA
C01	BCA
CO2	BBA

# Normal Form :

## *Third Normal Form (3NF):*

**3NF:** *The third Normal Form ensures the reduction of data duplication. It is also used to achieve data integrity.*

- *The third Normal Form ensures the reduction of data duplication. It is also used to achieve data integrity.*
- *A relation will be in 3NF if it is in 2NF and not contain any transitive partial dependency.*
- *3NF is used to reduce the data duplication. It is also used to achieve the data integrity.*



# Normal Form :

## *Third Normal Form (3NF):*

**3NF:** *A relation is in 3NF if it is in 2NF and there is no transitive dependency.*

- *The third Normal Form ensures the reduction of data duplication. It is also used to achieve data integrity.*
- *In the student table, stu\_id determines subid, and subid determines sub. Therefore, stu\_id determines sub via subid. This implies that the table possesses a transitive functional dependency, and it does not fulfill the third normal form criteria..*

stu_id	name	subid	sub	address
1	Arun	11	SQL	Delhi
2	Varun	12	Java	Bangalore
3	Harsh	13	C++	Delhi
4	Keshav	12	Java	Kochi



stu_id	name	subid	address
1	Arun	11	Delhi
2	Varun	12	Bangalore
3	Harsh	13	Delhi
4	Keshav	12	Kochi

subid	subject
11	SQL
12	java
13	C++
12	Java

# Normal Form

## *Boyce Codd normal form (BCNF):*

**BCNF:** A relation is in BCNF if it is in 3NF and for every Functional Dependency, LHS is the super key.

- BCNF is the advance version of 3NF. It is stricter than 3NF.
- A table is in BCNF if every functional dependency  $X \rightarrow Y$ ,  $X$  is the super key of the table.
- For BCNF, the table should be in 3NF, and for every FD, LHS is super key.
- To transform the table into the BCNF, you will divide the table into two parts. One table will hold stuid which already exists and the second table will hold a newly created column profid.

stuid	subject	professor
1	SQL	Prof. Mishra
2	Java	Prof. Anand
2	C++	Prof. Kanth
3	Java	Prof. James
4	DBMS	Prof. Lokesh

→

	stuid	profid
▶	1	101
	2	102
	2	103
	3	102
	4	104

	profid	subject	professor
▶	1	SQL	Prof. Mishra
	2	Java	Prof. Anand
	2	C++	Prof. Kanth
	3	Java	Prof. James
	4	DBMS	Prof. Lokesh



# Normal Form

## *Fourth normal form (4NF):*

**4NF:** A relation will be in 4NF if it is in Boyce Codd normal form and has no multi-valued dependency.

- For a dependency  $A \twoheadrightarrow B$ , if for a single value of  $A$ , multiple values of  $B$  exists, then the relation will be a multi-valued dependency.
- The given *STUDENT* table is in 3NF, but the *COURSE* and *HOBBY* are two independent entity. Hence, there is no relationship between *COURSE* and *HOBBY*.
- So to make the *Student* table into 4NF, we can decompose it into two tables *Course* & *Hobby* tables:

STU_ID	COURSE	HOBBY		STU_ID	COURSE	STU_ID	HOBBY
21	Computer	Dancing	→	21	Computer	21	Dancing
21	Math	Singing		21	Math	21	Singing
34	Chemistry	Dancing		34	Chemistry	34	Dancing
74	Biology	Cricket		74	Biology	74	Cricket



# Normal Form

## *Fifth normal form (5NF):*

**4NF:** A relation is in 5NF if it is in 4NF and not contains any join dependency and joining should be lossless.

- 5NF is satisfied when all the tables are broken into as many tables as possible in order to avoid redundancy.
- 5NF is also known as Project-join normal form (PJ/NF).
- So to make the SPC table into 4NF, we can decompose it into three tables SP, SC & PC tables:

Table: SPC

Suplier	Product	Customer
Ali	ABC	Nauman
Ahmad	KLM	Qasim
Haider	XYZ	Saim



Table: SP

Suplier	Product
Ali	ABC
Ahmad	KLM
Haider	XYZ

Table: SC

Suplier	Customer
Ali	Nauman
Ahmad	Qasim
Haider	Saim

Table: PC

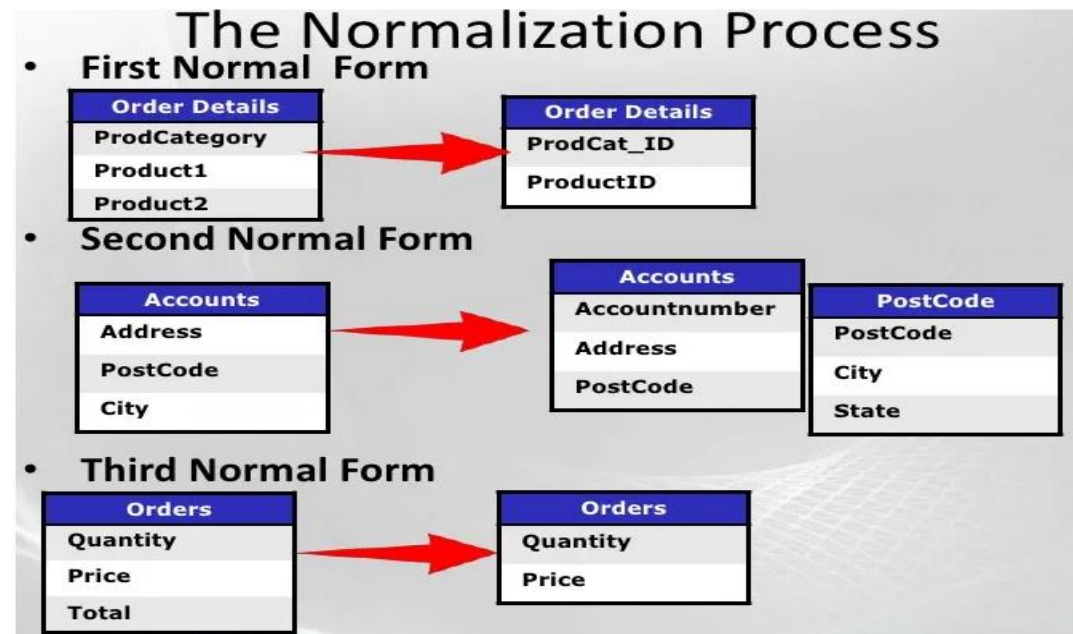
Product	Customer
ABC	Nauman
KLM	Qasim
XYZ	Saim

# Normal Form

## *Normalization:*

*Normalization is a step by step process of removing the different kinds of redundancy and anomaly one step at a time from the database.*

- 1<sup>st</sup> normal form
  - All attributes depend on **the key**
- 2<sup>nd</sup> normal form
  - All attributes depend on **the whole key**
- 3<sup>rd</sup> normal form
  - All attributes depend on **nothing but the key**



# Discussion

**Queries ?**  
**Prof. K. Adishesha**  
9449081542  
**Thank you**

