

Due Date:11-August-2024

20% penalty for 1 day late

40% penalty for 2 days late

Submission not allowed afterwards

CS2009: Design and Analysis of Algorithms (Fall 2024)

Assignment 1

Total Marks: 100

1. Show the steps merge sort uses to sort the following list of integers in the descending order (from the highest to the lowest / biggest to the smallest): [7.5 points]

**9, 19, 14, 29, 13, 4, 23, 10, 37**

Consider the following variation on Merge Sort, that instead of dividing input in half at each step of Merge Sort, you divide into **four** part, sort each part, and finally combine all of them using a **four-way merge** subroutine. What is the overall asymptotic running time of this algorithm? [7.5 points].

2. Prove  $3n^3 + 5n^2 + 25n = \Theta(n^3)$ . Determine the values of constant  $c$  and  $n_0$ . [5 Points]
3. Prove  $5 \log_2 \log_2 n + 4 \log^2 n = O(?)$ . Determine the values of constant  $c$  and  $n_0$ . [5 Points].
4. By using the frequency count method, find the time complexity of the following. [10 Points]

```
int p = 0;
for (i = n / 2; i > 1; i /= 6) {
    for (j = 2; j <= n; j *= 4) {
        for (k = 0; k <= j; k *= 3) {
            p = p + n / 2;
        }
    }
}
```

```
int p = 0;
for (i = n; i > 1; i = i / 2) {
    for (j = 2; j <= n; j += 1) {
        p = p + n;
        if (p > (n ^ i)) break;
    }
}
```

5. List the six functions below in non-decreasing asymptotic order of growth. Please give mathematical justification for the ordering you specify [5 Points]

$\lg \lg n^2$	$\log^2 n$	$n^{-2}$	$\lg(2^{\lg(n^2)})$	$n^{(1/2)}$	$n^2$
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6. Watch the video lecture on Big O, Big  $\Omega$  and Big  $\Theta$  notation from <http://www.youtube.com/watch?v=6Ol2JbwoJp0>. Write the summary of the lecture in your words. [10 Points].
7. For each of the following questions, indicate whether it is T (True) or F (False) and justify using some examples e.g. assuming a function? [20 Points]
  - For all positive  $n = \omega(\sqrt{n})$
  - For all positive  $f(n)$ ; if  $f(n) = O(g(n))$ , then  $g(n) = \Omega(f(n))$
  - Disprove that  $f(n) = O(f(n)^2)$  by showing a counterexample, show a function  $f(n)$  for which this does not hold.
  - Prove that the running time of an algorithm is  $\Theta(f(n))$  if and only if its worst-case running time is  $O(f(n))$  and its best-case running time is  $\Omega(f(n))$ .

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8. You have a computer and  $n$  processes with the processing time  $t_1, t_2, \dots, t_n$ . You have to pick the order in which to run the processes. Let  $p_j$  denote the  $j$ th process you run. Then, the completion time  $C_i$  for the process  $p_j$  is defined as (the sum of times for all the processes up till this one ends).

$$C_i = \sum_{j=1}^i t_{p_j}$$

Design 2 algorithms that give its time complexity **in polynomial time and logarithmic time respectively** that minimizes the average completion time  $\frac{1}{n} \sum_{i=1}^n C_i$