Review of loop invariants using Insertion Sort

Insertion Sort

- Problem: sort n numbers in A[1..n].
- Input: *n*, numbers in *A*
- Output: A in sorted order: $\forall i \in [2..n], A[i-1] \le A[i]$

```
for j=2 to length(A)
do key=A[j]
     i=j-1
   while i>0 and A[i]>key
     do A[i+1]=A[i]
     i--
     A[i+1]=key
```

Loop Invariant

• A **loop invariant** is a condition [among program variables] that is necessarily true immediately before and immediately after each iteration of a **loop**. (Note that this says nothing about its truth or falsity part way through an iteration.)

Loop Invariants

- Invariants statements about an algorithm that remain valid
- We must show three things about loop invariants:
 - Initialization statement is true before first iteration
 - **Maintenance** -if it is true before an iteration, *then* it remains true before the next iteration
 - **Termination** when loop terminates the invariant gives a useful property to show the correctness of the algorithm

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■ Initialization: j = 2, the invariant trivially holds because A[1] is a sorted array.

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```

■ **Maintenance**: the inner **while** loop finds the position i with A[i] <= key, and shifts A[j-1], A[j-2], ..., A[i+1] right by one position. Then key, formerly known as A[j], is placed in position i+1 so that $A[i] \le A[i+1] < A[i+2]$.

A[1...j-1] sorted + $A[j] \rightarrow A[1...j]$ sorted

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■ **Termination**: the loop terminates, when *j=n+1*. Then the invariant states: "A[1...n] consists of elements originally in A[1...n] but in sorted order."