# EECS-3311 – Lab – Analyzer

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# 1 Design Goals

```
require
    across 0 | .. | 4 as i all lab_completed(i.item) end
    read accompanying document: Eiffel101¹
ensure
    submitted on time
    no submission errors
rescue
    ask for help during scheduled labs
    attend office hours for TA William
```

This lab requires the implementation of the **visitor design pattern**. This Lab also makes use of the composite design pattern to parse expressions, check if their type is correct and evaluate them. The actual parser for expressions is provided.

#### 2 Resources

• Lecture recording of the **composite** and **visitor design patterns** from Section M can be found here (Lectures 15 and 16):

https://www.eecs.yorku.ca/~jackie/teaching/lectures/index.html#EECS3311 W19

• A tutorial series on the **composite** and **visitor design patterns** is available here:

https://www.eecs.yorku.ca/~jackie/teaching/tutorials/index.html#composite visitor

**Aside**: In the sequel, you are provided with a context free grammar for simple expressions, and the parsing is done for you via the Eiffel *gelex* and *geyacc* tools. You do not need to know how this parser works. However, if you are interested in learning more (not needed for this course), a simple tutorial and example is provided at (with anonymous login):

https://svn.eecs.yorku.ca/repos/sel-open/misc/tutorial/yacc-tutorial/docs/index.html

<sup>&</sup>lt;sup>1</sup> See: https://www.eecs.yorku.ca/~eiffel/eiffel101/Eiffel101.pdf

## 3 The Problem

In this lab you are required to *analyze* (i.e., evaluate, calculate type, and type check) any given string whose **syntax** conforms to the following context-free grammar (CFG) of a small expression language:

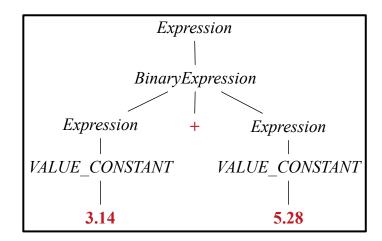
```
Expression ::=
   VALUE CONSTANT
  BinaryExpression
  UnaryExpression
  (Expression)
BinaryExpression ::=
                                    /* addition */
   Expression
                   Expression
  Expression
                                    /* subtraction */
                   Expression
                                    /* multiplication */
  Expression
                   Expression
                                    /* division */
                   Expression
  Expression
                                    /* equal to */
  Expression
                   Expression
                                    /* not equal to */
  Expression /= Expression
                                    /* greater than */
  Expression
                   Expression
                                    /* greater than or equal to */
  Expression >= Expression
  Expression < Expression
                                    /* less than */
                                    /* less than or equal to */
  Expression
              <= Expression
                                    /* conjunction */
  Expression \land Expression
  Expression
                   Expression
                                    /* disjunction */
               => Expression
                                    /* implication */
  Expression
                                    /* if-and-only-if */
               <=> Expression
  Expression
UnaryExpression ::=
                                     /* logical negation */
   not Expression
```

The above grammar consists of three rules, each of which with the form R := D: the left-hand side of := denotes the name of rule (e.g., BinaryExpression), whereas the right-hand side denotes the definition of rule. When a rule has multiple, alternative definitions, they are separated by a vertical bar (|).

Each italicized, capitalized word denotes a non-terminal (i.e., *Expression*, *BinaryExpression*, *UnaryExpression*). Each all-capital word (e.g., VALUE\_CONSTANT), as well as a symbol or word in bold red, denotes a terminal (e.g., **not**, (, \(\lambda\), etc.). The terminal VALUE\_CONSTANT denotes any valid floating-point number (e.g., 3, 3.14, etc.). Texts enclosed within /\* and \*/ denote comments that are not part of the grammar.

The above grammar can produce an infinite number of *syntactically-correct* strings. Terminals are base cases of the production, whereas non-terminal are recursive cases. For the purpose of this lab, we say that a given string is *syntactically correct* if that string can be produced by the above grammar. As examples:

- The string 3.14 can be produced (because *Expression* can be a VALUE CONSTANT).
- The string 3.14 + 5.28 can be produced (because Expression can be *Expression* + *Expression*, and the left and right *Expression* can each produce a VALUE\_CONSTANT). Here is the so called **abstract syntax tree** (AST) for this syntactically correct string:



- As exercises:
  - $\circ$  Why is (23.4 \* (34.04 + 28.9)) syntactically correct?
  - $\circ$  Why is (23.4 \* (not (34.04 + 28.9))) also syntactically correct?
  - $\circ$  Why is (23.4 \* (not (34.04 + 28.9) (34.28 / 19.29))) not syntactically correct?

In this lab, you are given a parser that is able to take as input a given string and determine if it is syntactically correct with respect to the above grammar. This parser is useful: we should not bother to analyze a random string which does not even conform to the expected syntax. If there is no syntax, then the input string is transformed into an **AST**, for which we implement using the **composite design pattern**, where the top deferred class is **EXPRESSION**. That is, a runtime object of type **EXPRESSION** has a tree structure similar to the above AST. See Figure 2.

Can we always evaluate a string that is *syntactically correct*? Maybe not. For example:

conforms to the above grammar but it is not *type-correct*: **not** is the unary operator for logical negation, but its operand (34.04 + 28.9) is not a Boolean expression.

In this lab exercise, you are given working implementation of classes that implement the above CFG using the **composite design pattern**. Your tasks will be to implement the visitor design pattern to support the following operations that may be performed upon any string which conforms to the above grammar:

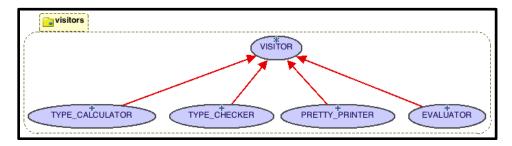
- Type Calculation
- Type Checking (which should make use of the type calculation)
- Evaluation

A working implementation for the **PRETTY\_PRINTER** visitor is given to you as the starting reference.

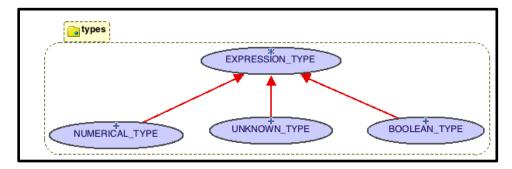
# 4 Design of Classes

Your starter project contains the following clusters:

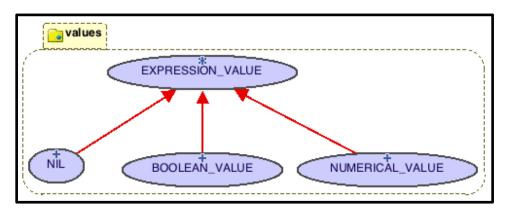
- parsing: classes in this cluster parse a given string and determine if it conforms to the above grammar of the small expression language. If the given string is syntactically correct, the parser constructs a COMPOSITE object (which is an AST of the given string). The parsing classes are provided, so you do not need to do the actual parsing.
- composite: classes in this cluster implement the composite design pattern for constructing objects that correspond to ASTs of string that are syntactically correct. **See Figure 2**. You will have to complete classes in this cluster, as these classes are used to represent simple and composite expressions.
- visitors: classes in this cluster implement the visitor design pattern for performing various operations (pretty printing, evaluation, type calculation, and type checking) upon the composite objects.



• types: classes in this cluster should be used by the TYPE\_CALCULATOR visitor class.



• values: classes in this cluster should be used by the EVALUATOR visitor class.



```
analyzer
       analyzer.ecf
      composite
          - addition.e
          - binary expression.e
          - composite.e
          - conjunction.e
           disjunction.e
          - division.e
          equals.e
          - expression.e
           greater than.e
           greater\_than\_or\_equals.e
          - iff.e
         - implication.e
          less than.e
          - less than or equals.e
          - multiplication.e
          - negation.e
          - not equals.e
         – subtraction.e
          - unary_expression.e
          - value_constant.e
      - parse
      - docs/
      - root
     └── root.e
      - tests
         instructor
         basic_tests.e
         student
        student_tests.e
      - types
         boolean_type.e
           expression_type.e
          - numerical_type.e
       — unknown_type.e
      - values
          -boolean value.e
          expression_value.e
         – nil.e
         - numerical value.e
      visitors
         - evaluator.e
        - pretty printer.e
        - type_calculator.e
         type_checker.e
        visitor.e
```

Figure 1 starter directory analyzer (you complete classes in red)

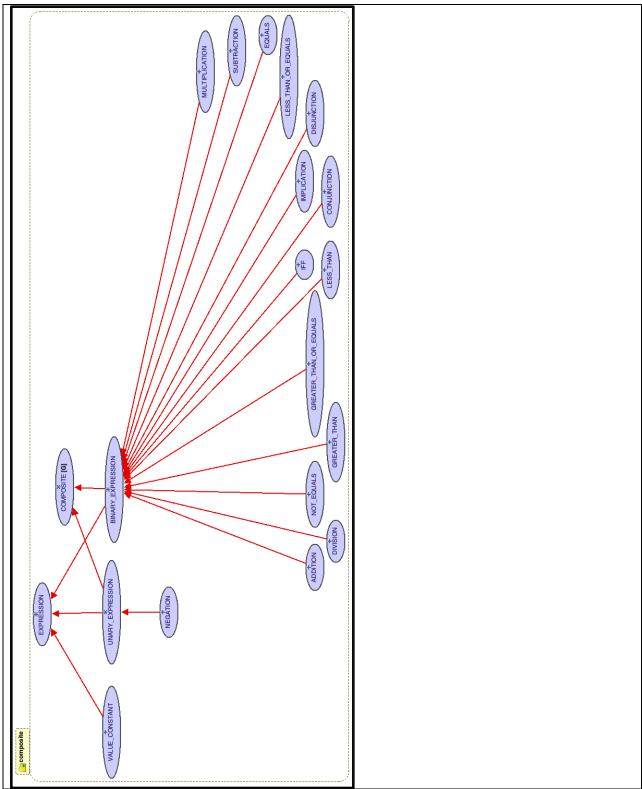


Figure 2 Expressions and Composite Expressions

# 5 Getting started

These instructions are for when you work on one of the EECS Linux Workstations or Servers (e.g *red*). You should not compile on *red* as it is a shared server; compile on an EECS workstation. See the course wiki for how to use the SEL-VM or your Laptop.

#### 5.1 Retrieve Lab5

#### > ~sel/retrieve/3311/lab5

This will provide you with a starter directory analyzer with the project structure as shown in Figure 1. Classes in **bold red** are those that you are required to modify.

# 5.2 Compile the Lab

You can now compile the Lab.

```
> estudio analyzer/analyzer.ecf &
```

where analyzer.ecf is the Eiffel configuration file for this Lab, with the ROOT class:

```
note
    description: "gelex-calculator application root class"
    date: "$Date$"
    revision: "$Revision$"

class
    ROOT|
inherit
    ES_SUITE

create
    make
    feature {NONE} -- Initialization

    make
    do
        add_test (create {BASIC_TESTS}.make)
        show_browser
        run_espec
    end
```

This initial starter project given to you compiles. Running Workbench System will give you the following test report:

	FAILED (13 failed & 2 passed out of 15)		
Case Type	Passed	Total	
Violation	0	0	
Boolean	2	15	
All Cases	2	15	
State	Contract Violation	Test Name	
Test1	BASIC_TESTS		
FAILED	NONE	test_01: print 3	
FAILED	NONE	test_02: print $((3 + 4) * 45 >= 64) \land (4 /= 34) \lor (not (5 = 65))$	
FAILED	Check assertion violated.	t_03: 3.23 evaluates to 3.23	
FAILED	Check assertion violated.	t_04a: (3.23 + 4) * 3 evaluates to 21.69	
FAILED	Check assertion violated.	$t_04b: (3.23 + 4) * 3 = 21.69$ evaluates to true	
FAILED	NONE	test_05: 3.23 has numerical type	
FAILED	NONE	test_06: 3.23 + 4 has numerical type	
PASSED	NONE	test_07: $(3 < 4) + (23 >= 34)$ has unknown type	
FAILED	NONE	test_08: $(3 < 4) \Rightarrow (23 \Rightarrow 34)$ has boolean type	
FAILED	NONE	test_09: (3 < 4) <=> (23 >= 34) has boolean type	
FAILED	NONE	test_10: 3.23 is type correct	
FAILED	NONE	test_11: 3.23 + 4 is type correct	
PASSED	NONE	test_12: $(3 < 4) + (23 >= 34)$ is not type correct	
FAILED	NONE	test_13: $(3 < 4) \Rightarrow (23 \Rightarrow 34)$ is type correct	
FAILED	NONE	test_14: (3 < 4) <=> (23 >= 34) is type correct	

Note Two tests (test\_07 and test\_12) are passed initially because the expected results (unknown type and false) happen to match the initial, default values of the corresponding VISITOR objects (TYPE\_CALCULATOR and TYPE\_CHECKER).

# 5.3 Understanding the API

Study carefully the tests given to you in the BASIC TESTS class. For example:

```
test_02: BOOLEAN
    local
        v: PRETTY_PRINTER
do
        comment ("test_02: print ((3 + 4) * 45 >= 64) /\ (4 /= 34) \/ (not (5 = 65))")
        parser.parse_string ("((3 + 4) * 45 >= 64) /\ (4 /= 34) \/ (not (5 = 65))")
        Result := parser.error_count = 0
        check Result end

        create v.make
        -- parser.expression returns a composite object of static type EXPRESSION
        -- In this case, its dynamic type is DISJUNCTION
        parser.expression.accept (v)
        Result := v.value ~ "(((((3+4)*45)>=64)/\((4/=34))\/((not(5=65)))")
        end
```

The call parser.parse\_string (...) takes a string and determines if it is syntactically correct (it is not if parser.error\_count is larger than zero). If there is no syntax error in the input string, then parser.expression returns an EXPRESSION object. To use the **visitor design pattern**, we call the accept feature on the EXPRESSION object:

```
parser.expression.accept (v)
```

Here you need to review how the **double dispatch** work in the visitor pattern, given that the dynamic type of parser.expression is DISJUNCTION and the dynamic type of v is PRETTY\_PRINTER.

# 6 What You Must Do

Here is what you are required to complete (see the **To Do** tags):

- composite cluster
  - The deferred feature accept is declared in the EXPRESSION class. You are required to implement that feature in all its effective descendant classes.
- visitors cluster
  - o Complete the classes here according to the **visitor design pattern** taught in class.
  - o Note that:
    - A working implementation is given to you for the **PRETTY PRINTER**.
    - For each of the three remaining visitors, an attribute value is declared with the type that corresponds to the visitor's expected operation:
      - For the EVALUATOR visitor, attribute value is with type EXPRESSION VALUE (for the result of evaluation).

For evaluating equality (=) between floating-point numbers, we have 0.001 tolerance. A helper query <code>is\_equal\_within</code> is included in this class for you to use.

- For the TYPE\_CALCULATOR visitor, attribute value is with type EXPRESSION\_TYPE (for the result of type calculation).
- For the TYPE\_CHECKER visitor, attribute value is declared with type BOOLEAN (for whether or not the string is type-correct).

## 7 To Submit

Before submitting, please consult the course wiki for final details as well as when the submission machinery is available.

To obtain a passing grade, your submission must compile and pass the tests we provided.

We will be grading the correctness of your submission with additional tests.

The main requirements for submission are as follows.

- 1. Add correct implementations as specified.
- 2. Add your own contracts (precondition, postcondition, and invariant) when appropriate.
- 3. Work incrementally one feature at a time. Run all regression tests before moving to the next feature. This will help to ensure that you have not added new bugs, and that the prior code you developed still executes correctly.
- 4. Add at least 4 tests of your own to STUDENT TESTS, i.e. do not just rely on our tests.
- 5. Don't make any changes to classes other than the ones specified.
- 6. Ensure that you get a green bar for all tests. Before running the tests, always freeze first.

You must make an electronic submission as described below.

- 1. On Prism (Linux), *eclean* your system, freeze it, and re-run all the tests to ensure that you get the green bar.
- 2. eclean your directory sorted-variants again to remove all EIFGENs.

Submit your Lab from the command line as follows:

```
submit 3311 Lab5 analyzer
```

You will be provided with some feedback. Examine your feedback carefully. Submit often and as many times as you like.

### Remember

- Your code must compile and execute on the departmental Linux system (Prism) under CentOS7. That is where it must work and that is where it will be compiled and tested for correctness.
- Equip each test t with a *comment* ("t: ...") clause to ensure that the ESpec testing framework and grading scripts process your tests properly. (Note that the colon ":" in test comments is mandatory.). An improper submission will not be given a passing grade.
- The directory structure of your folder sorted-variants must be a superset of Figure 1.