

Question 1:

- 1) First,  $EG \rightarrow EG$ , with  $E \rightarrow CA$  we have  
 $EG \rightarrow EGCA$ , with  $CG \rightarrow DI$  we have  
 $EG \rightarrow EGCADI$ ,  
Hence,  $EG \rightarrow I$  holds.
- 2) Candidate keys are:  
AGJ and EGJ
- 3) **1NF** since it only contains atomic attribute values.  
Not in 2NF since  $AJ \rightarrow C$  violates 2NF, the non-prime attribute C is partially dependent on key {AGJ}.
- 4) The minimal cover  $F_m = \{CG \rightarrow DI, AG \rightarrow B, ADI \rightarrow EH, E \rightarrow CA, AJ \rightarrow BE, EJ \rightarrow I\}$   
Or  $F_m = \{CG \rightarrow DI, AG \rightarrow B, ADI \rightarrow EH, E \rightarrow CA, AJ \rightarrow BEI\}$

5) **Not lossless**

	A	B	C	D	E	G	H	I	J
R1	b	b	a	b	a	a	b	a	b
R2	a	b	b	a	a	b	a	b	a
R3	b	a	b	b	a	a	a	b	b

First scan of F, we have:

	A	B	C	D	E	G	H	I	J
R1	a	a	a	b	a	a	b	a	b
R2	a	b	a	a	a	b	a	b	a
R3	a	a	a	b	a	a	a	a	b

Second scan of F, we have:

	A	B	C	D	E	G	H	I	J
R1	a	a	a	b	a	a	b	a	b
R2	a	b	a	a	a	b	a	b	a
R3	a	a	a	b	a	a	a	a	b

There is no change during the second scan. The result table doesn't contain a line with all columns marked with 'a'. Hence, the decomposition is not lossless.

- 6) R violates BCNF because of  $CG \rightarrow DI$ , hence R is decomposed to  
 $R1(C,D,G,I)$  and  $R2(A,B,C,E,G,H,J)$ ,  
 $R2$  violates BCNF because of  $AG \rightarrow B$ , hence  $R2$  is decomposed to  
 $R21(A,B,G)$  and  $R22(A,C,E,G,H,J)$ ,  
 $R22$  violates BCNF because of  $E \rightarrow AC$ , hence  $R22$  is decomposed to  
 $R221(A,C,E)$  and  $R222(E,G,H,J)$   
 $R222$  violates BCNF because of  $EG \rightarrow H$ , hence  $R222$  is decomposed to  
 $R2221(E,G,H)$  and  $R2222(E,G,J)$   
Hence, a BCNF decomposition of R is:  
 $R1(C,D,G,I)$

R21(A,B,G)

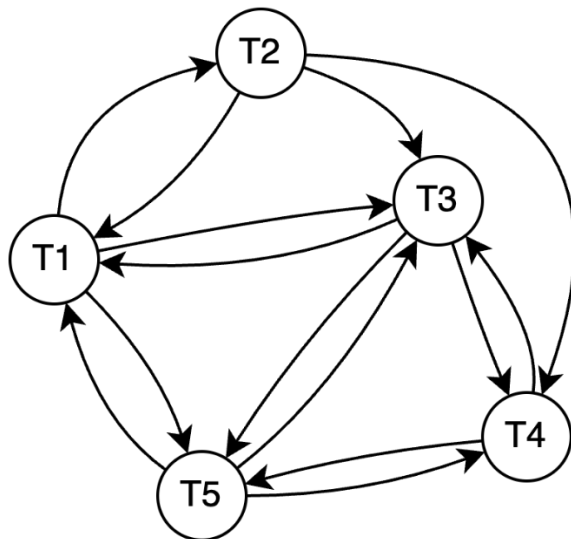
R221(A,C,E)

R2221(E,G,H)

R2222(E,G,J)

Question 2:

- 1) T1, T5: Redo  
T3, T4: Undo  
T2: Do Nothing
- 2) The transaction schedule is not conflict serializable because its precedence graphs is not acyclic:



- 3) One possible schedule, regardless of T2 and T3, that must cause deadlock is shown below; in the deadlock, T1 waits for T5 for Y, T5 waits for T3 for Z, and T3 waits for T1 for X.

T1	WL(X)	R(X)	...				WL(Y)	R(Y)	W(Y)	W(X)	...			UL(X)	UL(Y)					
T2	...																			
T3			WL(Z)	R(Z)	...				WL(X)	R(X)	W(Z)	W(X)	...			UL(Z)	UL(X)			
T4	...																			
T5					WL(Y)	W(Y)	...					WL(Z)	R(Z)	W(Z)	...				UL(Y)	UL(Z)

Question 3:

### 1) LRU Policy

	P1	P2	P3	P2	P4	P5	P6	P6	P3	P7	P2	P3
Buffer1	P1	P1	P1	P1	P4	P4	P4	P4	P3	P3	P3	P3
Buffer2		P2	P2	P2	P2	P2	P6	P6	P6	P6	P2	P2
Buffer3			P3	P3	P3	P5	P5	P5	P5	P7	P7	P7

Fault	F	F	F		F	F	F		F	F	F	
Hit				H				H				H

9 page faults, 3 hits

## 2) MRU Policy

	P1	P2	P3	P2	P4	P5	P6	P6	P3	P7	P2	P3
Buffer1	P1	P1	P1	P1	P1	P1	P1	P1	P1	P1	P1	P1
Buffer2		P2	P2	P2	P4	P5	P6	P6	P6	P6	P6	P6
Buffer3			P3	P3	P3	P3	P3	P3	P3	P7	P2	P3
Fault	F	F	F		F	F	F			F	F	F
Hit				H				H	H			

9 page faults, 3 hits

## 3) FIFO Policy

	P1	P2	P3	P2	P4	P5	P6	P6	P3	P7	P2	P3
Buffer1	P1	P1	P1	P1	P4	P4	P4	P4	P3	P3	P3	P3
Buffer2		P2	P2	P2	P2	P5	P5	P5	P5	P7	P7	P7
Buffer3			P3	P3	P3	P3	P6	P6	P6	P6	P2	P2
Fault	F	F	F		F	F	F		F	F	F	
Hit				H				H				H

9 page faults, 3 hits

## 4)

LRU,MRU, and FIFO all have the same page faults (9 faults and 3 hits), and they have the same performance.