Assignment 1

Complete the following problems and turn in a hardcopy of your solutions by the beginng of class on January 21^{st} (5:00 pm).

1. 4 points. Consider an extension of diML that provides support for references, arrays, and while loops (each defined as we did in class). Call this language diml-RAW. We wish to extend diml-RAW so that it supports for and foreach loops. The for loops should have the form for(e₁; e₂; e₃) { e₄ }, and should behave as for loops in C, except they should evaluate to false when the loop terminates (like while loops in diml-RAW). The foreach loops should have the form foreach(x in e₁) { e₂ } and operate as shown below (note that foreach loops should also evaluate to false when the loop terminates). Provide the static and dynamic semantics of both for and foreach loops, keeping consistent with the existing language semantics of diml-RAW.

The following function should, when given an integer array, add 1 to every element in the array.

```
fun f (xs : int array) : bool = foreach (x in xs) \{x := !x + 1;\}
```

2. **6 points.** Let's assume that you have already defined a single-step judgement (i.e., $e \to e'$) for some language L. Now, define the multi-step judgement $e \to^* e'$. Next, prove the following theorem, which states that the alpha-equivalence relationship is preserved by the multi-step relation.

```
\alpha-safety: ((e_1 =_{\alpha} e_2) \land (e_1 \to^* e'_1)) \Rightarrow \exists e'_2 : (e_2 \to^* e'_2 \land e'_1 =_{\alpha} e'_2)
```

In your proof of the alpha-safety theorem, you may assume that the following two lemmas hold:

```
\begin{array}{l} \alpha\text{-progress: } ((e_1 =_\alpha e_2) \wedge (e_1 \to e_1')) \Rightarrow \exists e_2' : e_2 \to e_2' \\ \alpha\text{-preservation: } ((e_1 =_\alpha e_2) \wedge (e_1 \to e_1') \wedge (e_2 \to e_2')) \Rightarrow e_1' =_\alpha e_2' \end{array}
```

3. 1.5 points extra credit. Prove that the α -progress and α -preservation lemmas hold for the call-by-name untyped lambda calculus.