

Learning quantum states and unitaries of bounded gate complexity

arXiv:2310.19882

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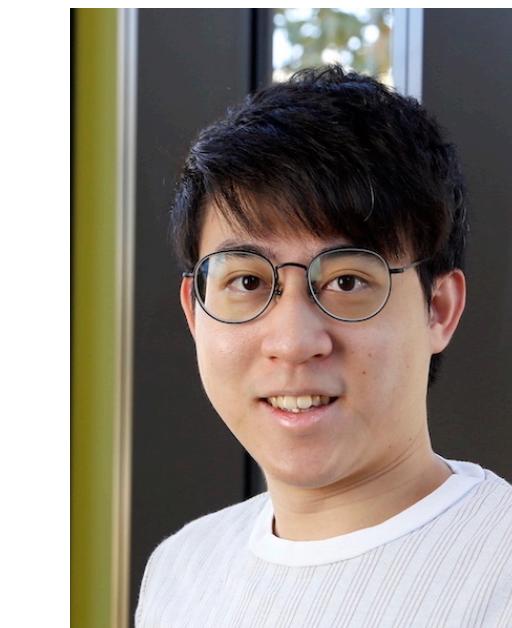
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Yihui Quek



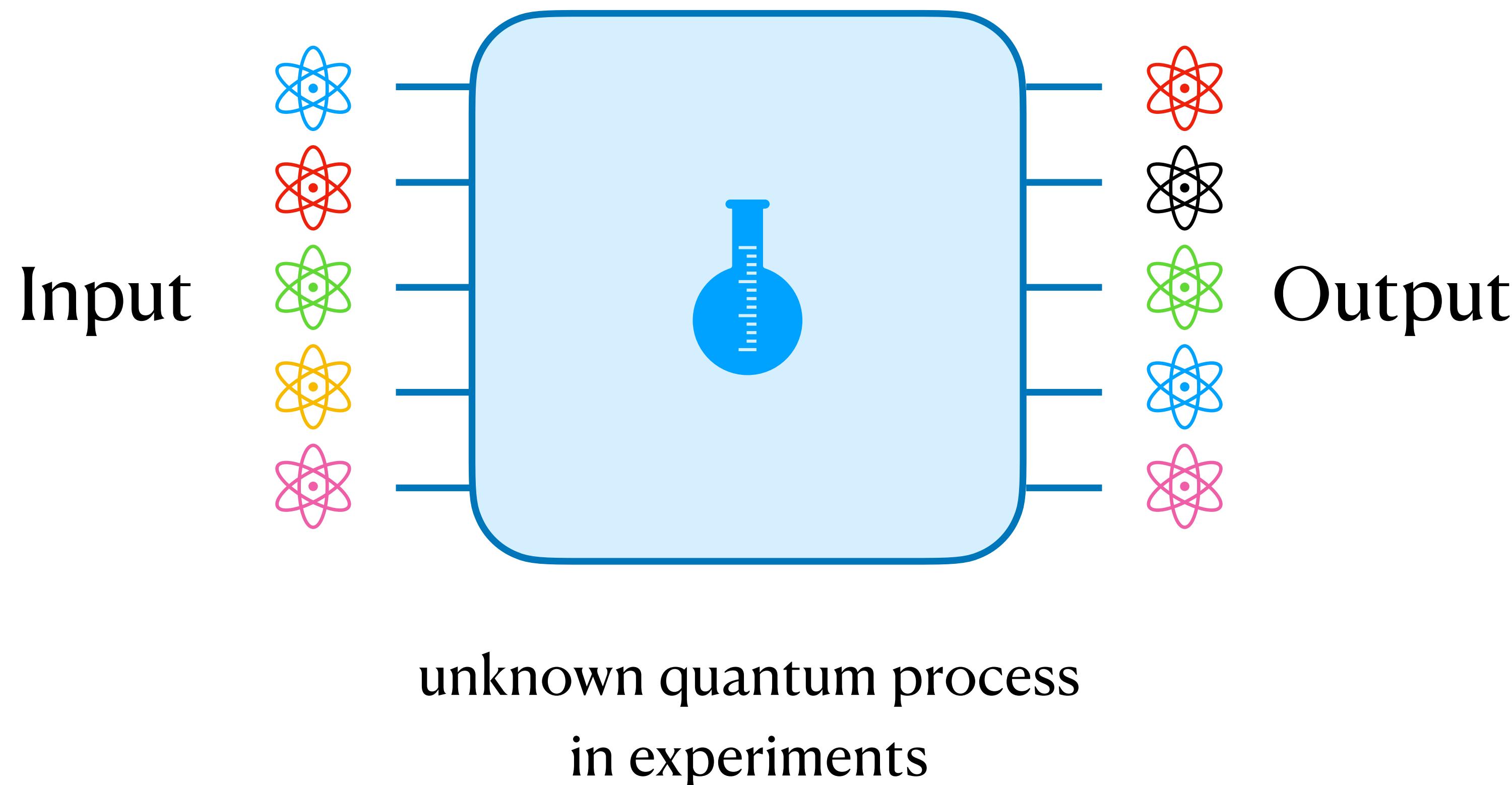
Hsin-Yuan Huang **Matthias Caro
(Robert)**



Motivation

How to learn from Nature?

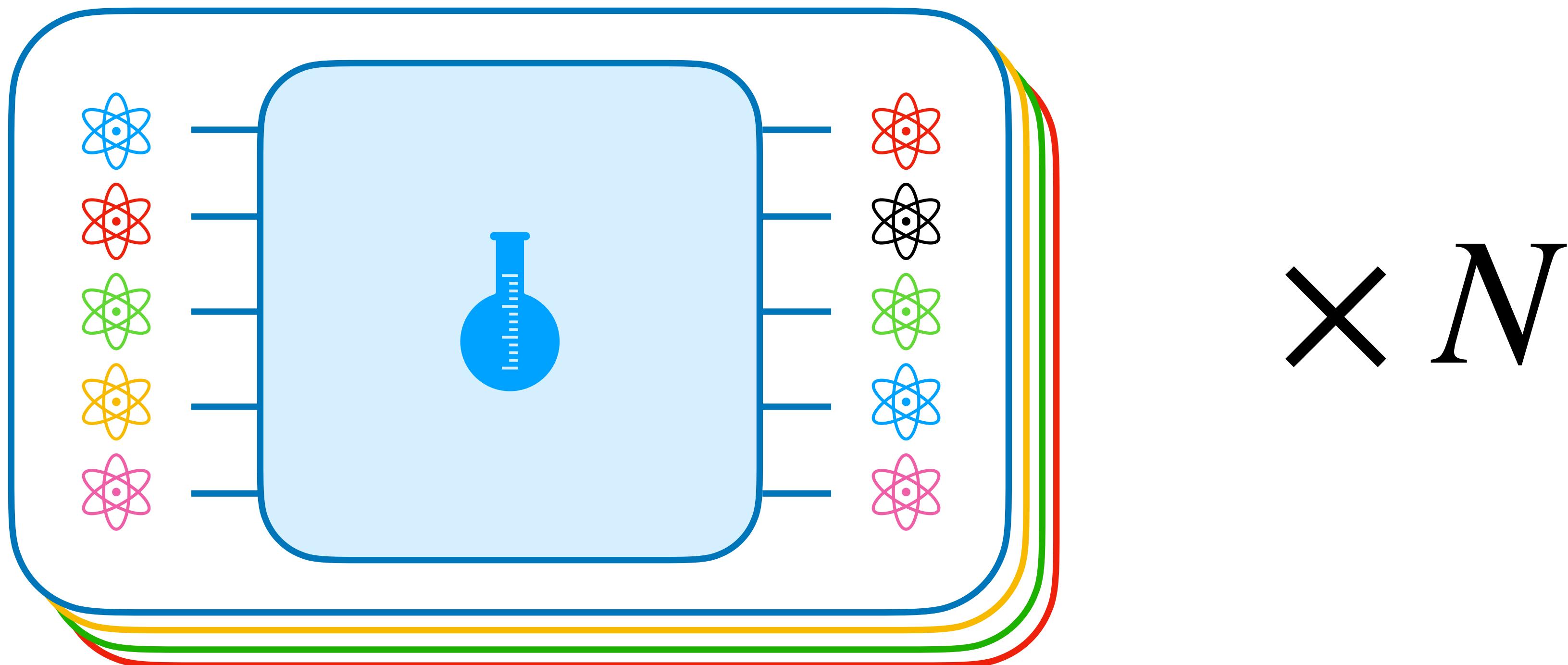
- Do experiments



Motivation

How to learn from Nature?

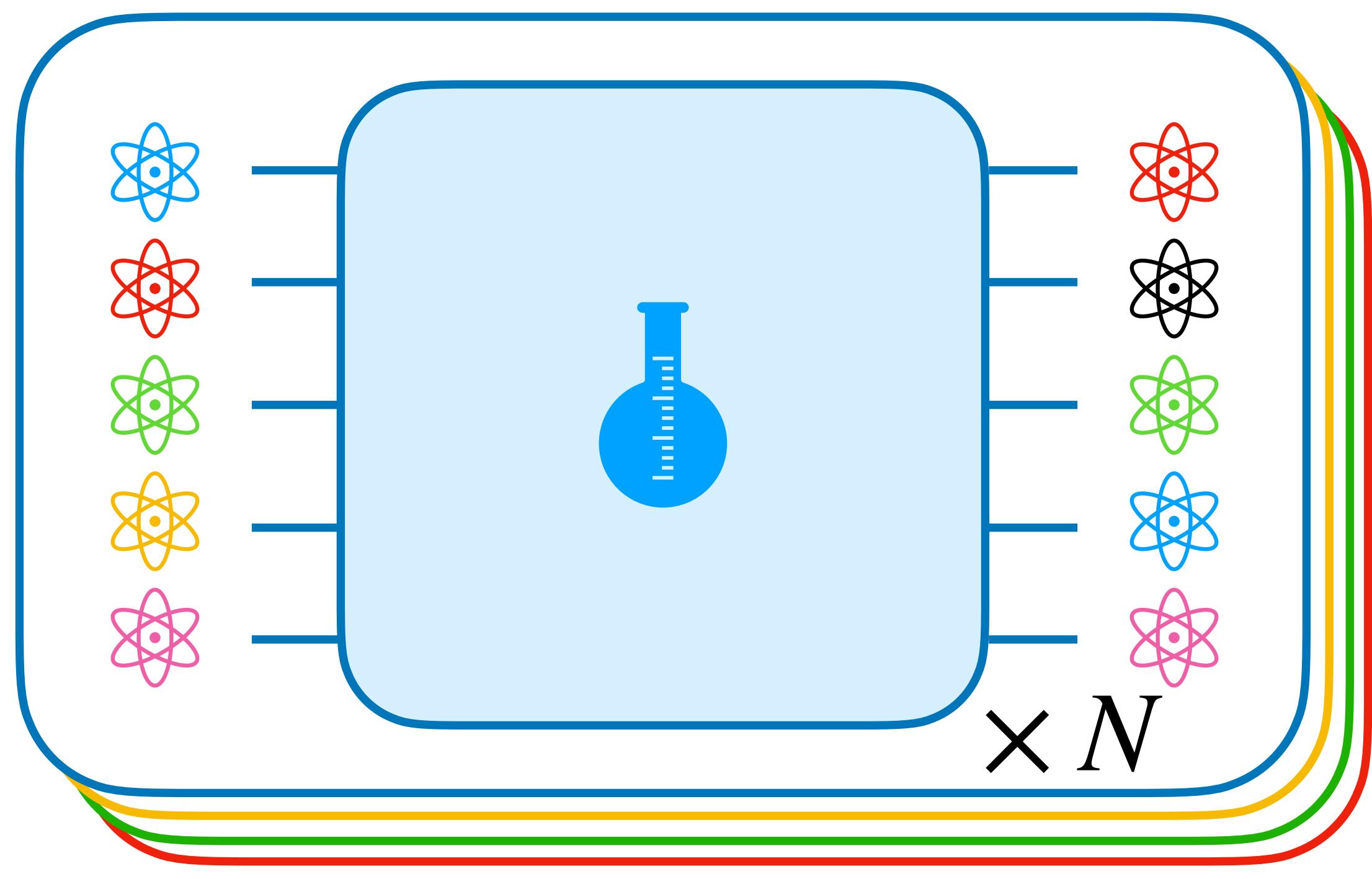
- Do experiments => Collect many samples



Motivation

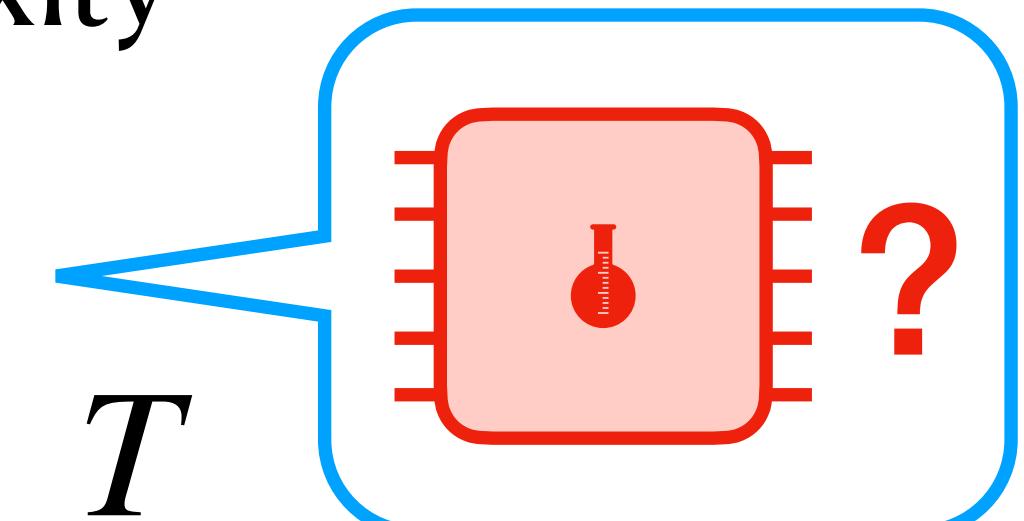
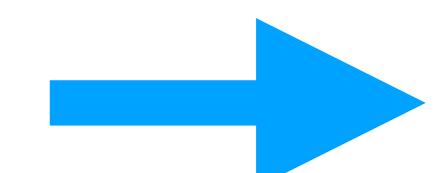
How to learn from Nature?

- Do experiments => Collect many samples => Try to learn the underlying mechanism



N : sample complexity

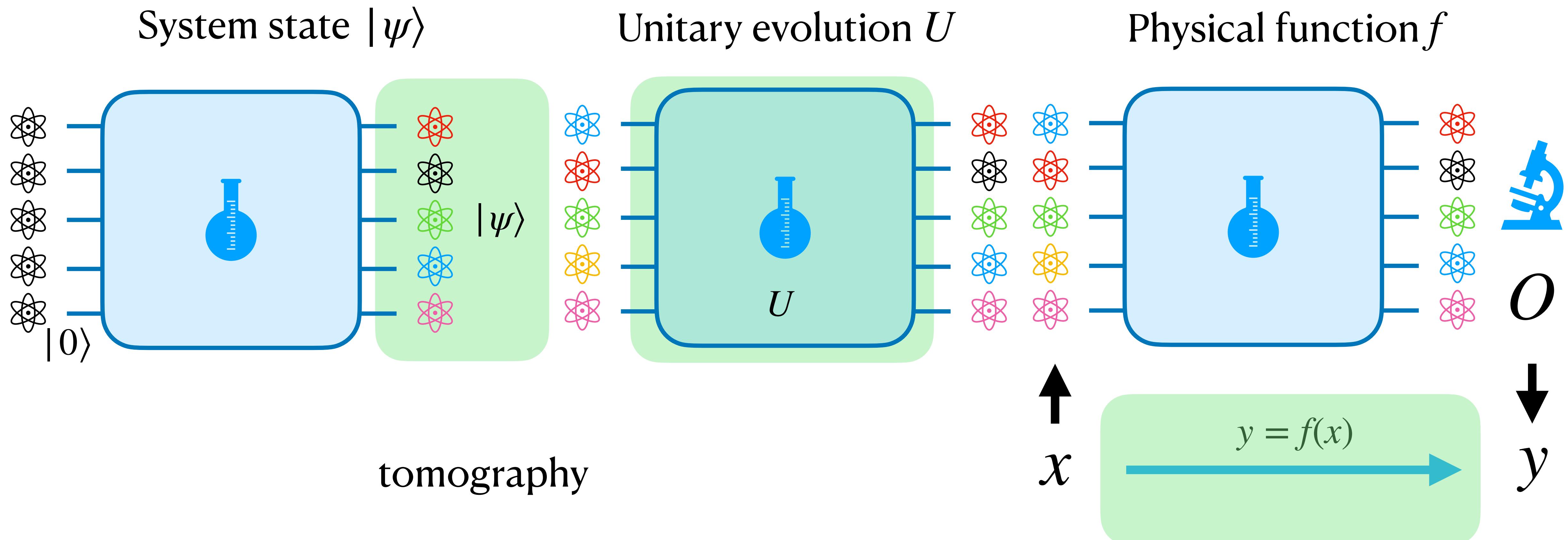
T : computational complexity



$$d(\text{blue box}, \text{red box}) \leq \epsilon \text{ w.h.p}$$

Motivation

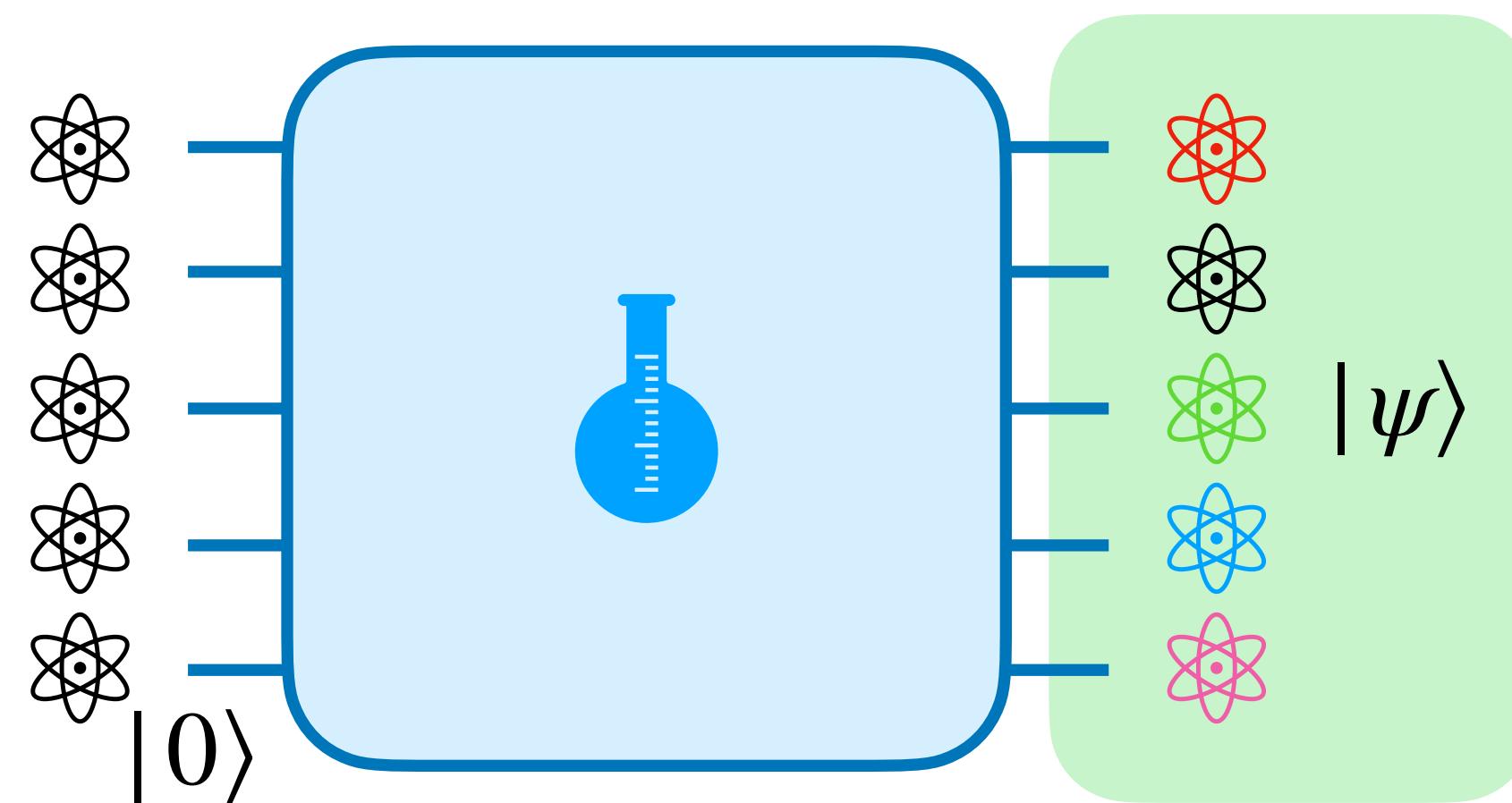
What to learn?



Motivation

Learning is hard in general!

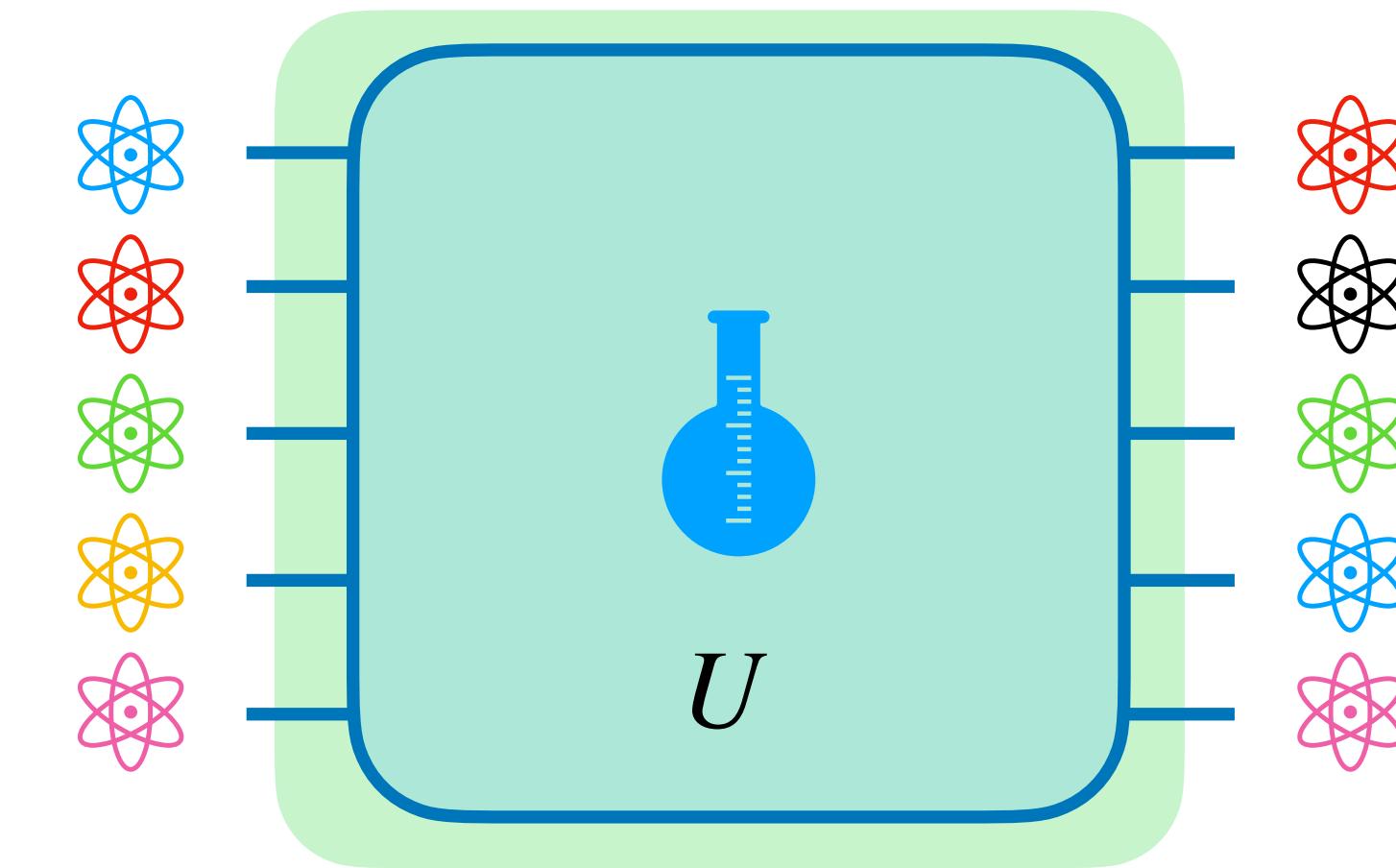
System state $|\psi\rangle$



$$N = \Theta(2^n)$$

How is learning even possible?

Unitary evolution U



$$N = \Theta(4^n)$$

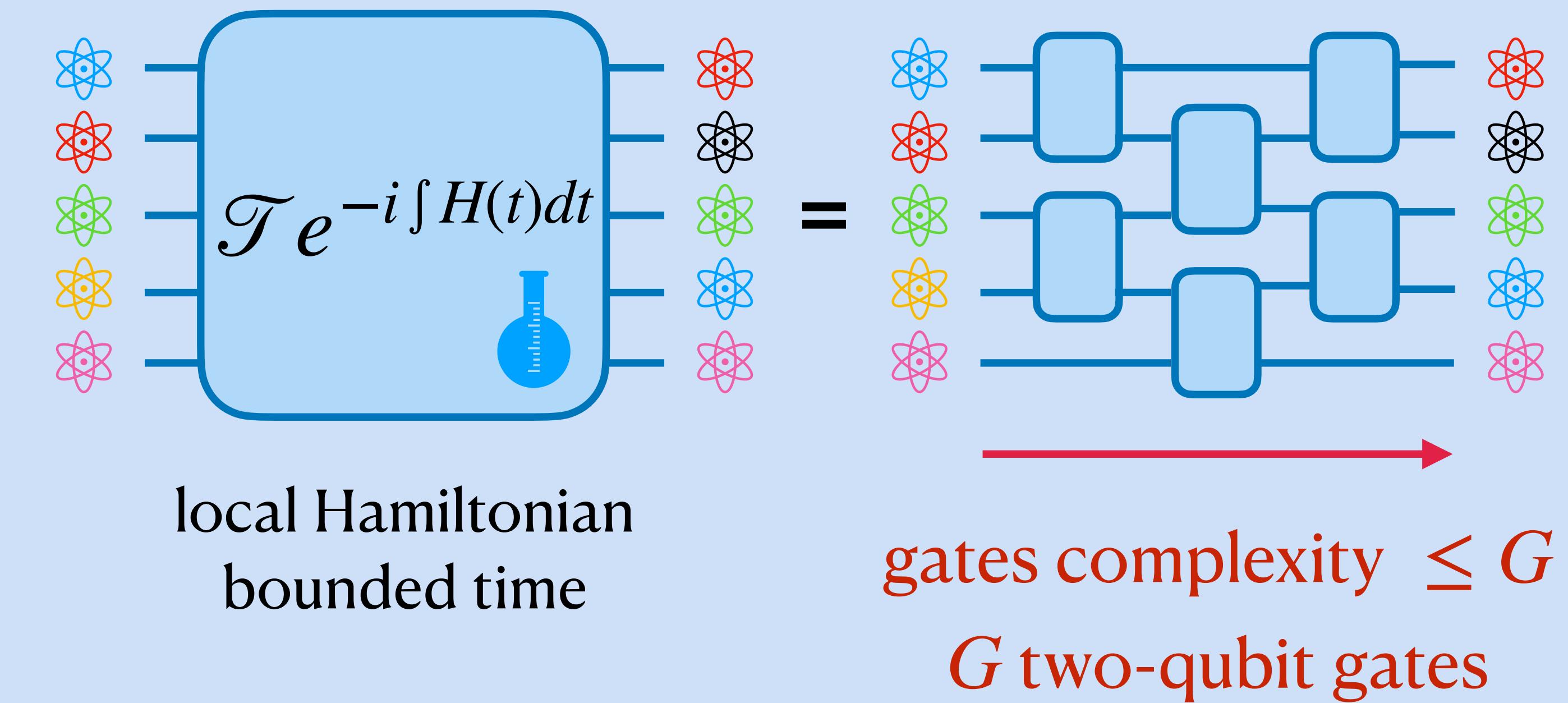
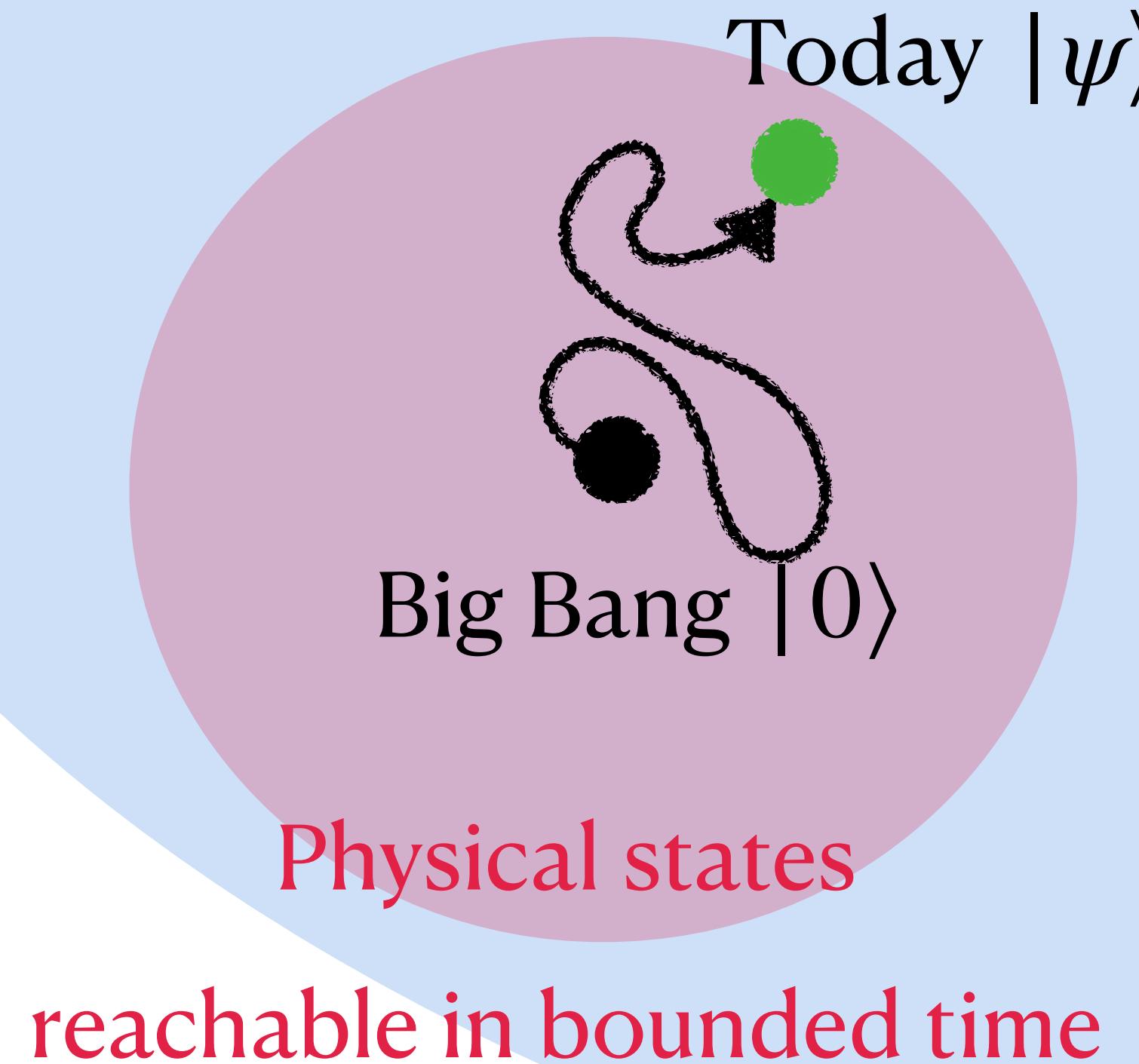
$n \sim 10^{23}!$

O'Donnell and Wright, STOC 2016

Haah, Kothari, O'Donnell, Tang, FOCS 2023

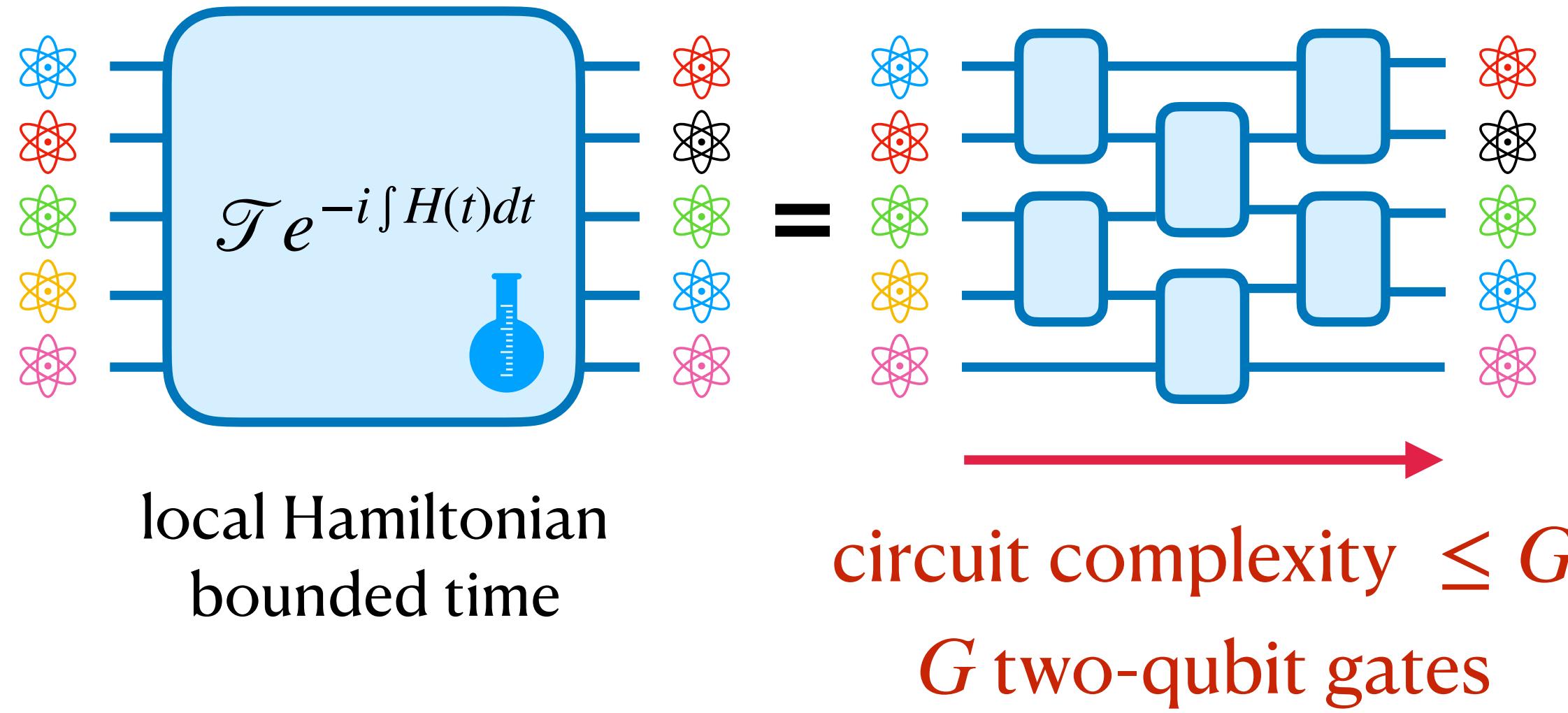
Illusive Hilbert space $\sim 2^n$

Motivation Physical constraints



In this work, we don't need “geometric” locality, nor discrete gate-set

Main Question



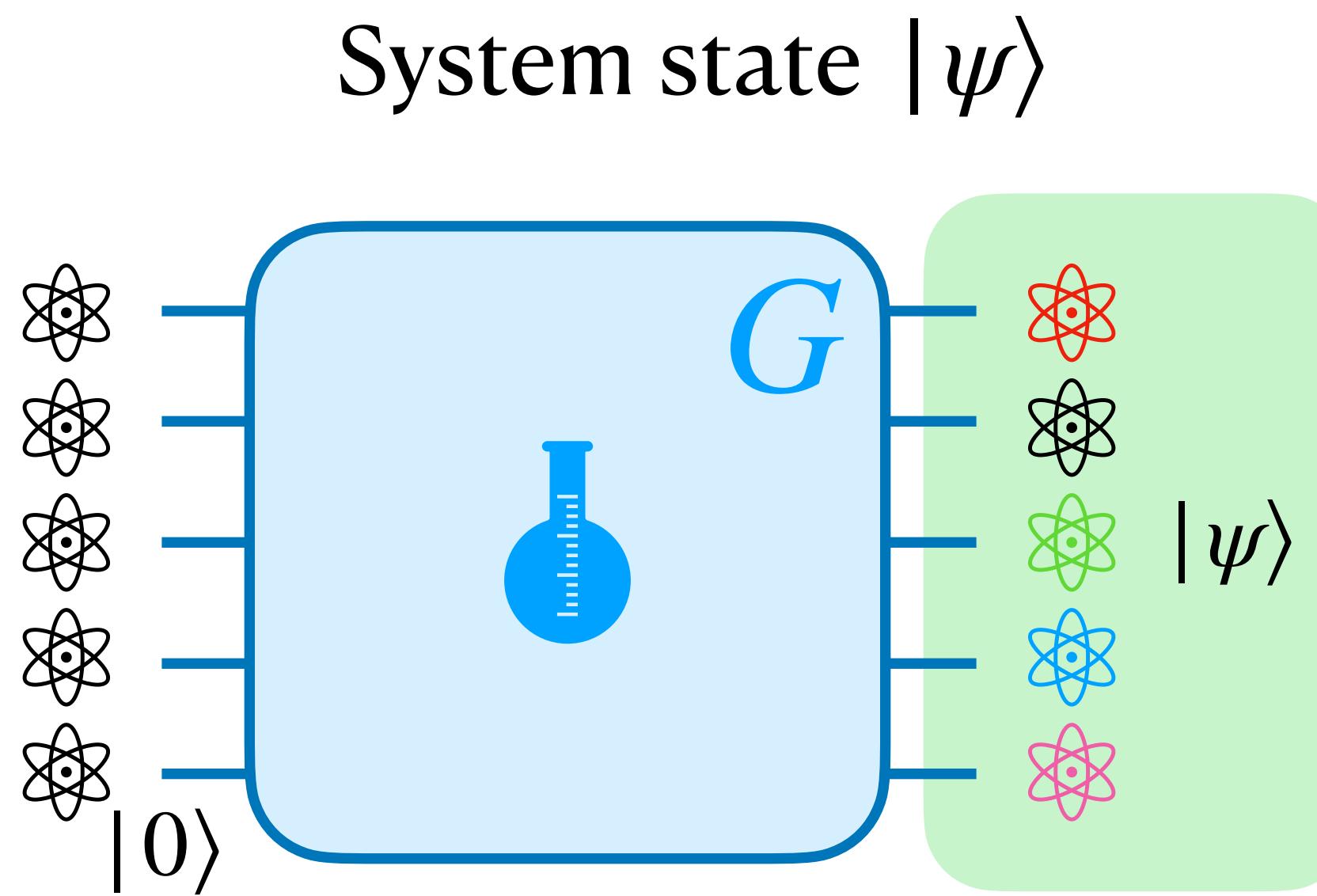
Can we efficiently learn states/unitaries of bounded gate complexity?

Applications on near-term quantum devices: limited G (both digital and analog)

Relating different notions of complexity: the complexity of learning and creating

Results

sample complexity



Learning $|\psi\rangle$ in ϵ trace distance
requires $N = \tilde{\Theta}\left(\frac{G}{\epsilon^2}\right)$ samples.

1. Complexity of learning = **complexity of creating** (information theoretically)
2. Completely independent of system size n . Can learn $n = 10000, G = 10$ states!
3. Non-adaptive/incoherent scheme is already optimal.

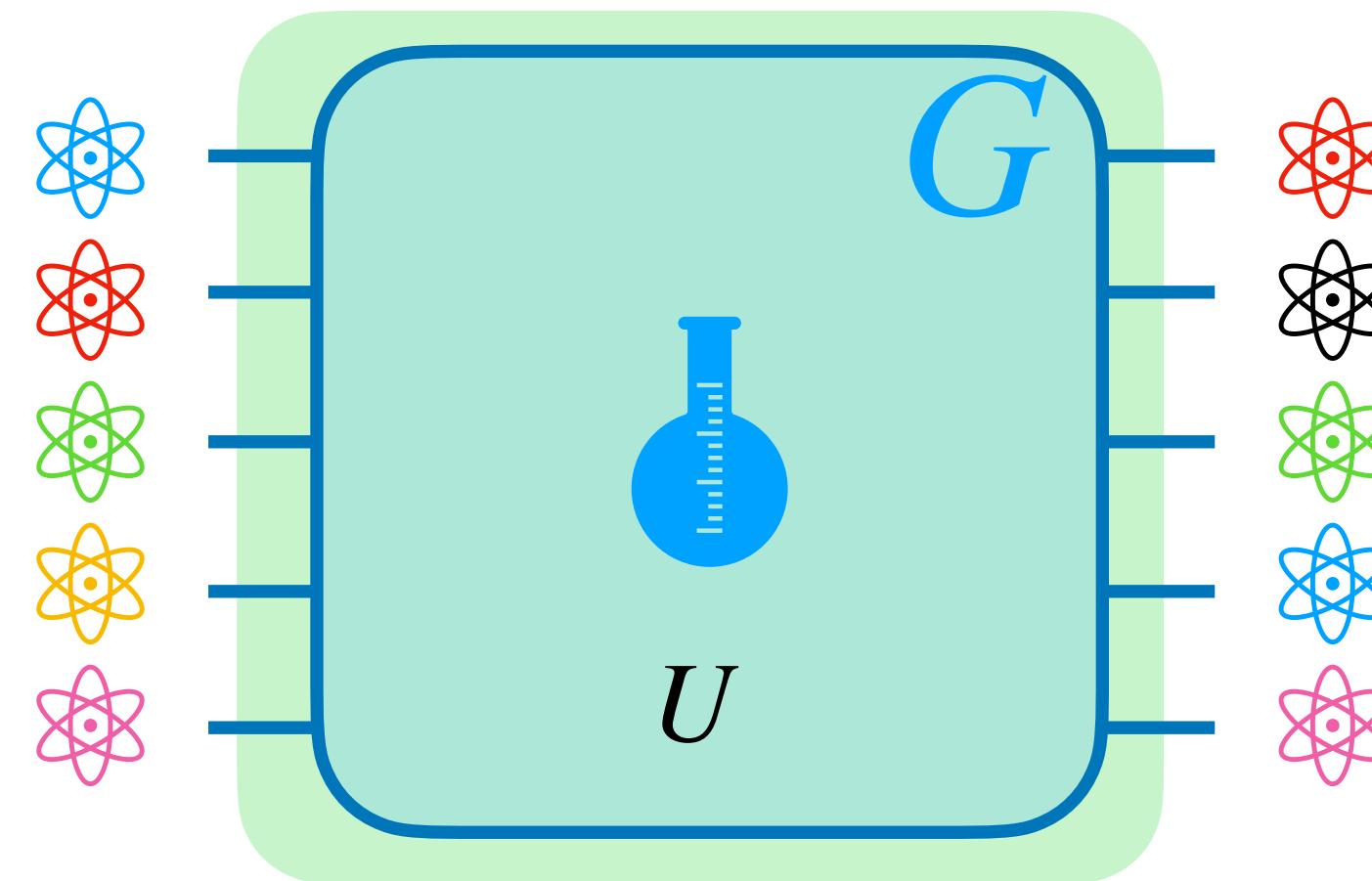
Results

sample complexity

= trace distance between Choi states

$$d_{\text{avg}}(U, V) = \sqrt{\mathbb{E}_{|\psi\rangle} [d_{\text{tr}}(U|\psi\rangle, V|\psi\rangle)^2]}$$

Unitary evolution U

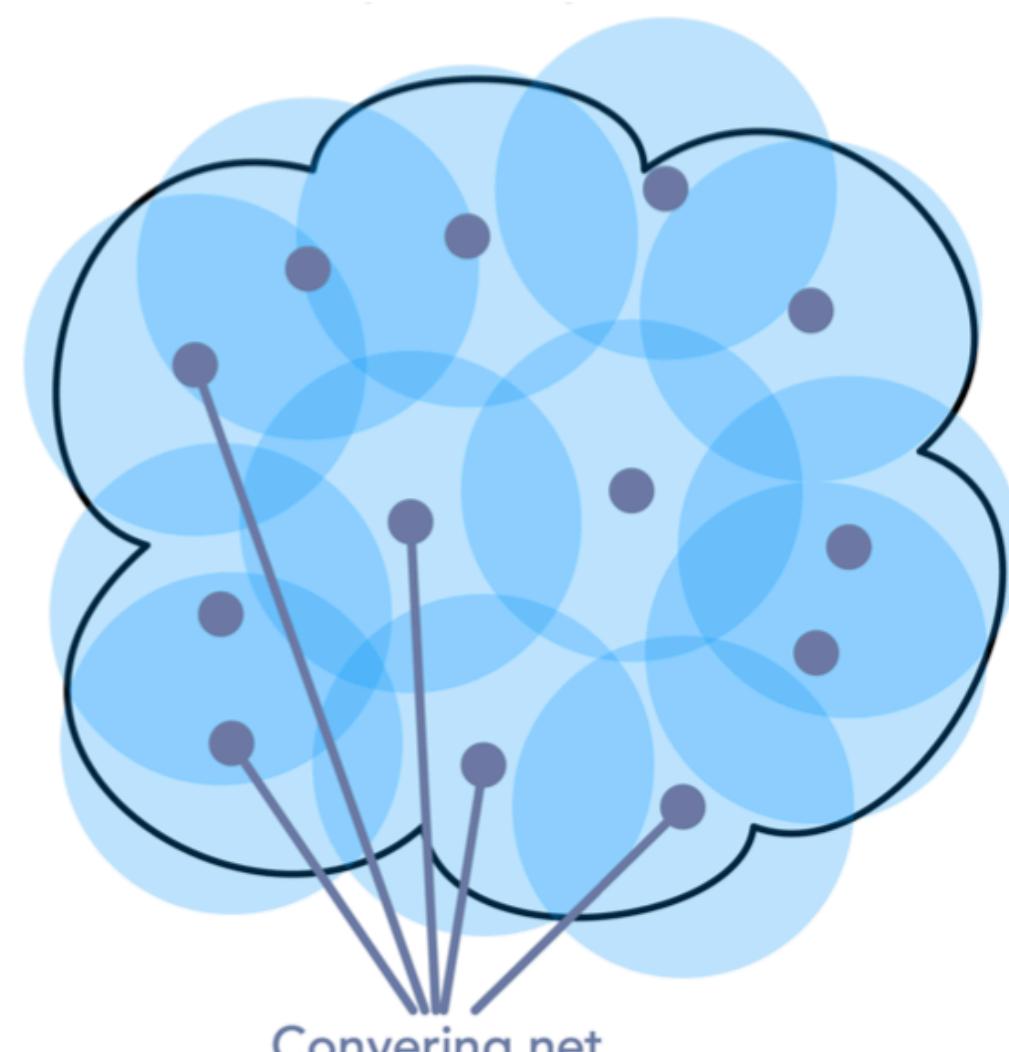
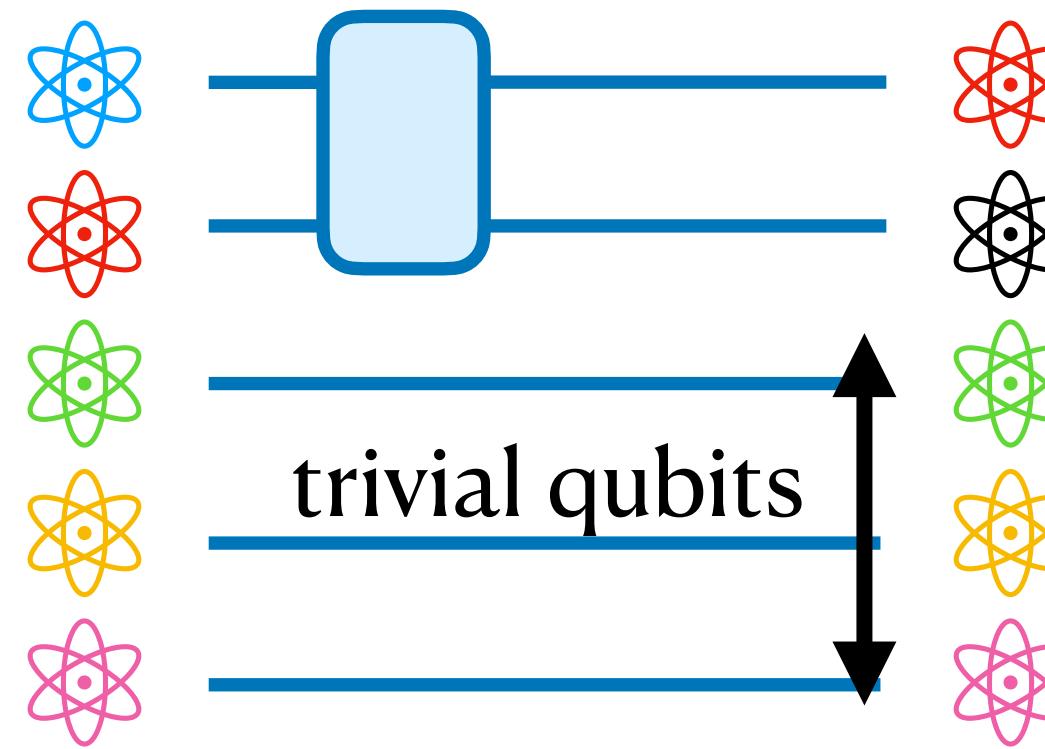


Learning U in average-case distance
requires $N = \tilde{\Theta}(G)$ queries.

1. Non-adaptive/incoherent query is already optimal (in G).
2. Learning in **worst-case distance** (diamond norm) requires $N = \exp(\Omega(\min\{G, n\}))/\epsilon$
3. ϵ -dependence: $\tilde{O}(\min\{1/\epsilon^2, \sqrt{2^n}/\epsilon\})$, $\Omega(1/\epsilon)$, **Heisenberg scaling** open Grover

Proof sketch

Upper bound: the learning algorithm



1. **Junta learning:** measure to identify non-trivial qubits, remove n -dependence
2. **Hypothesis selection:** construct a covering net \mathcal{N} that covers the set of G -gate states with ϵ -balls
$$\log |\mathcal{N}| = \tilde{\Theta}(G)$$
3. Find the best candidate by estimating all distances

$$\text{with classical shadow } N = O\left(\frac{\log |\mathcal{N}|}{\epsilon^2}\right) \leq \tilde{O}\left(\frac{G}{\epsilon^2}\right)$$

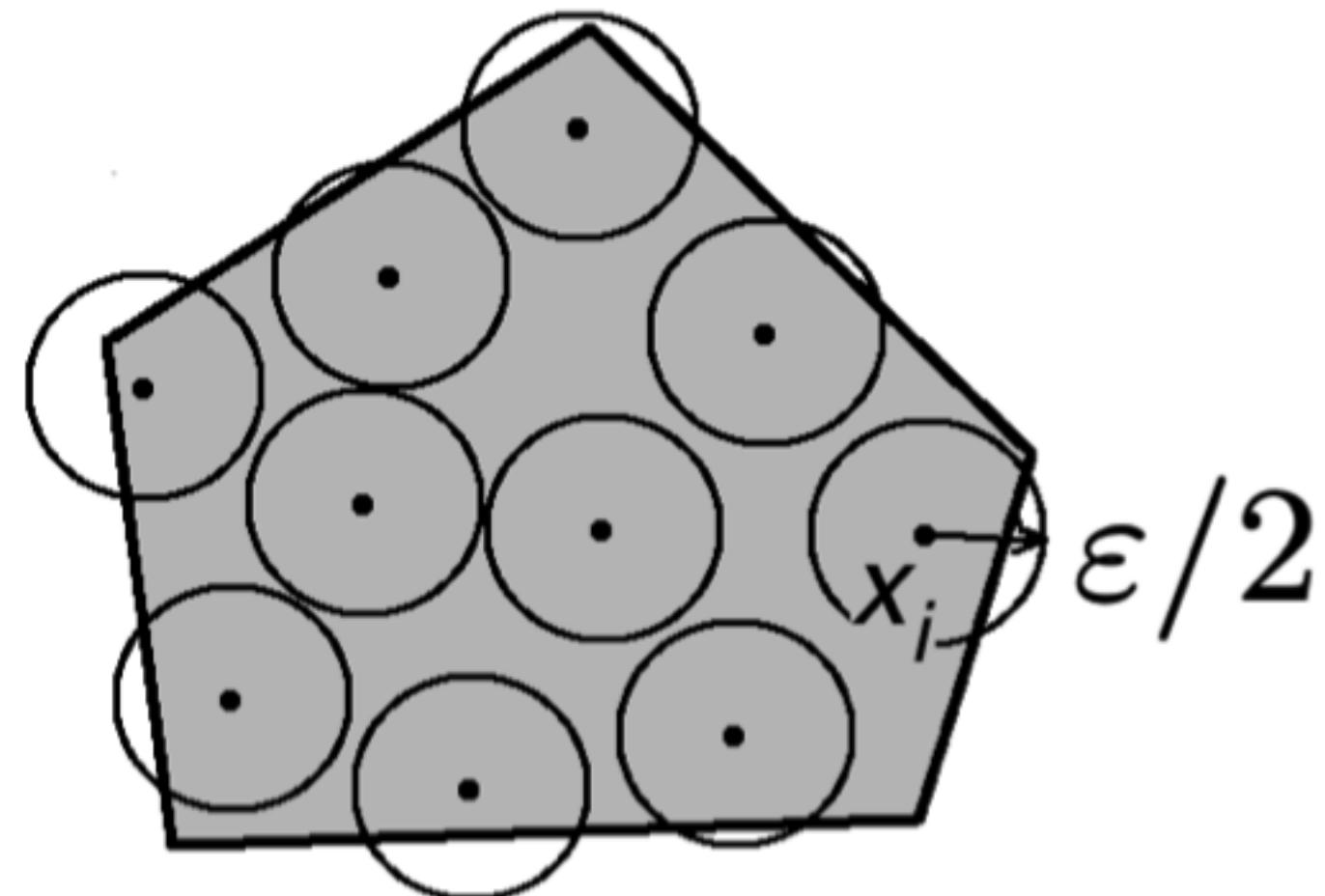
Unitary: $O(G \min\{1/\epsilon^2, \sqrt{2^n}/\epsilon\})$

via Choi states + quantum phase estimation

Proof sketch

Lower bound: information theory

1. Learning \Rightarrow distinguish elements of a packing net \mathcal{P}



2. Fano's inequality: distinguishing requires $\Omega(\log |\mathcal{P}|)$

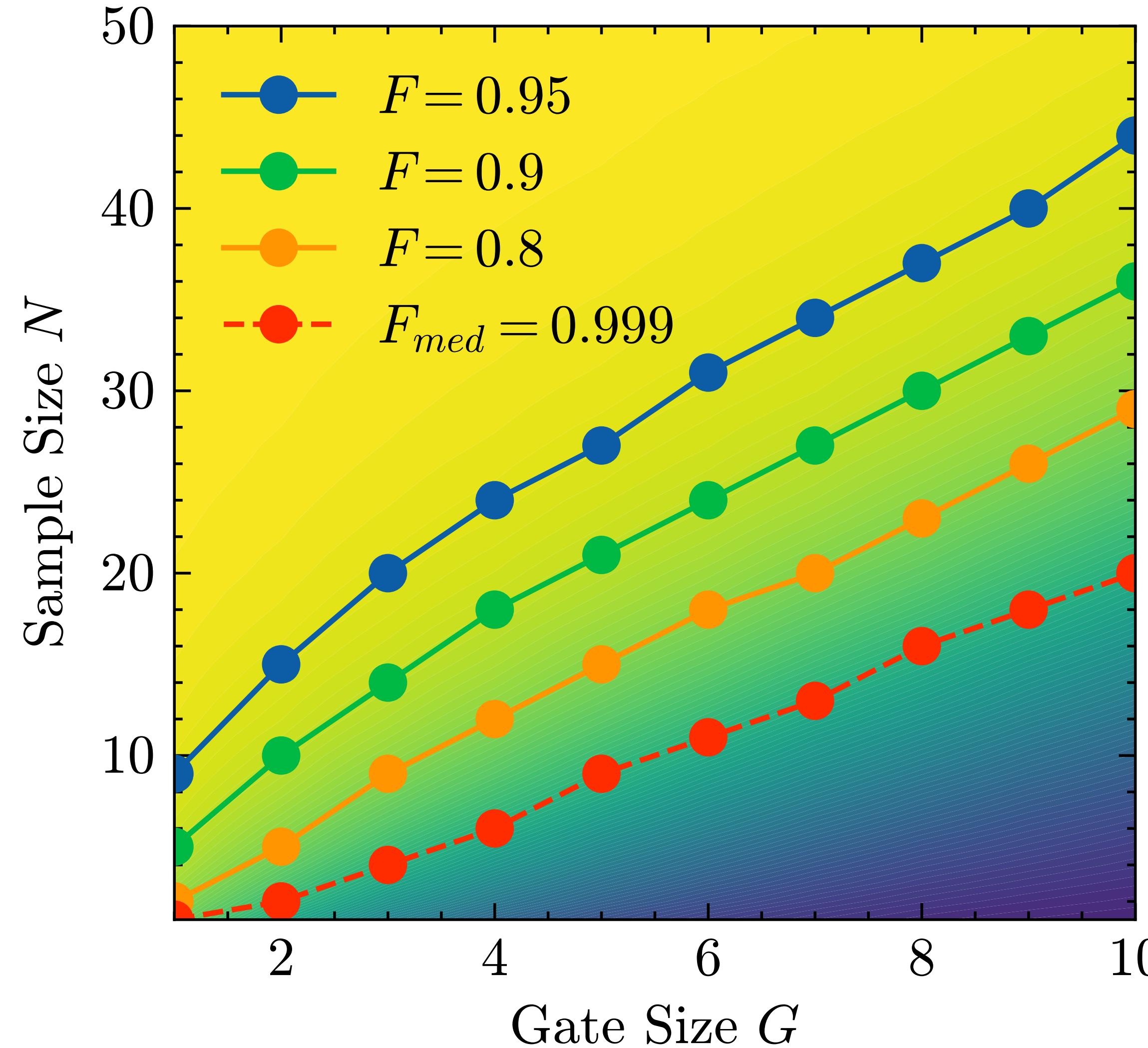
bits of info

3. Holevo's theorem: each sample gives $\tilde{O}(\epsilon^2)$ bits of

$$\text{info} \Rightarrow N \geq \Omega\left(\frac{\log |\mathcal{P}|}{\epsilon^2}\right) = \tilde{\Omega}\left(\frac{G}{\epsilon^2}\right)$$

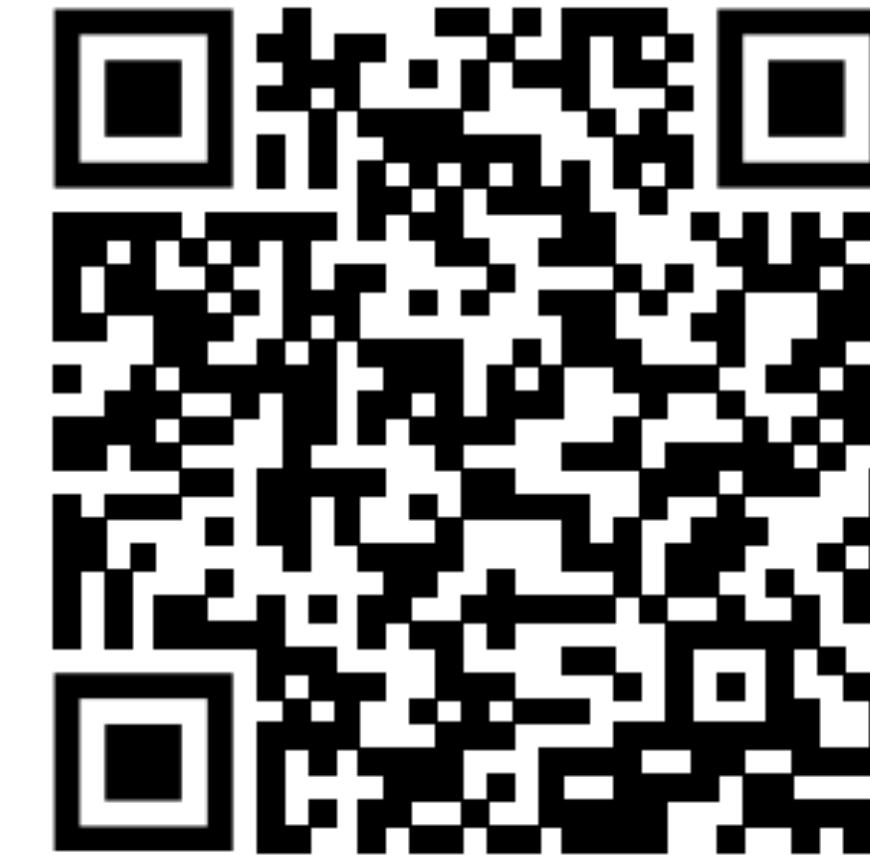
$\log |\mathcal{P}| \approx \log |\mathcal{N}|$: a general way to prove matching sample complexity bound

Numerical experiments



Learning random G -gate
states on $n = 10000$ qubits

Sample linear in G , runtime e^G

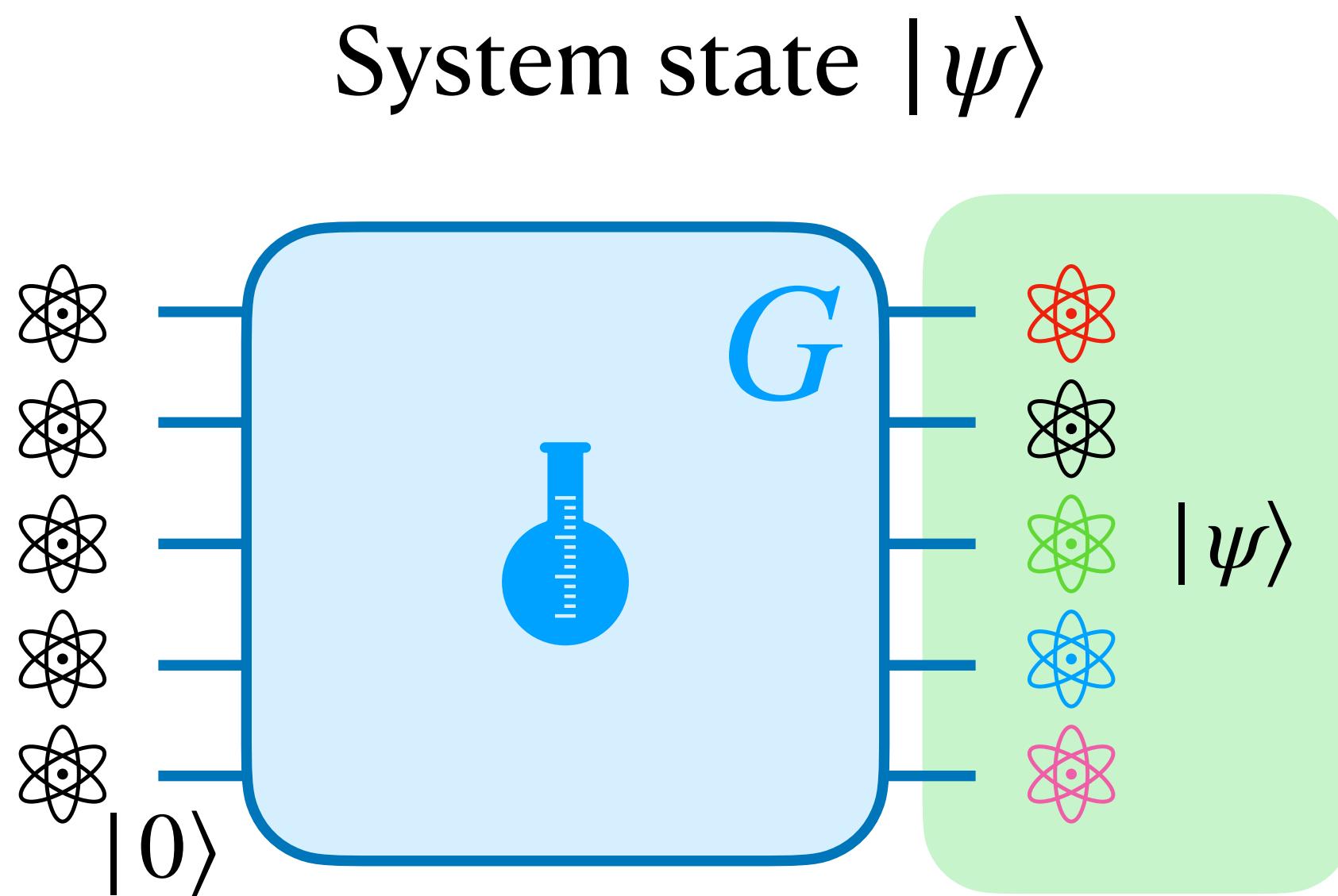


GitHub

[haimengzhao/bounded-gate-tomography](https://github.com/aimengzhao/bounded-gate-tomography)

Results

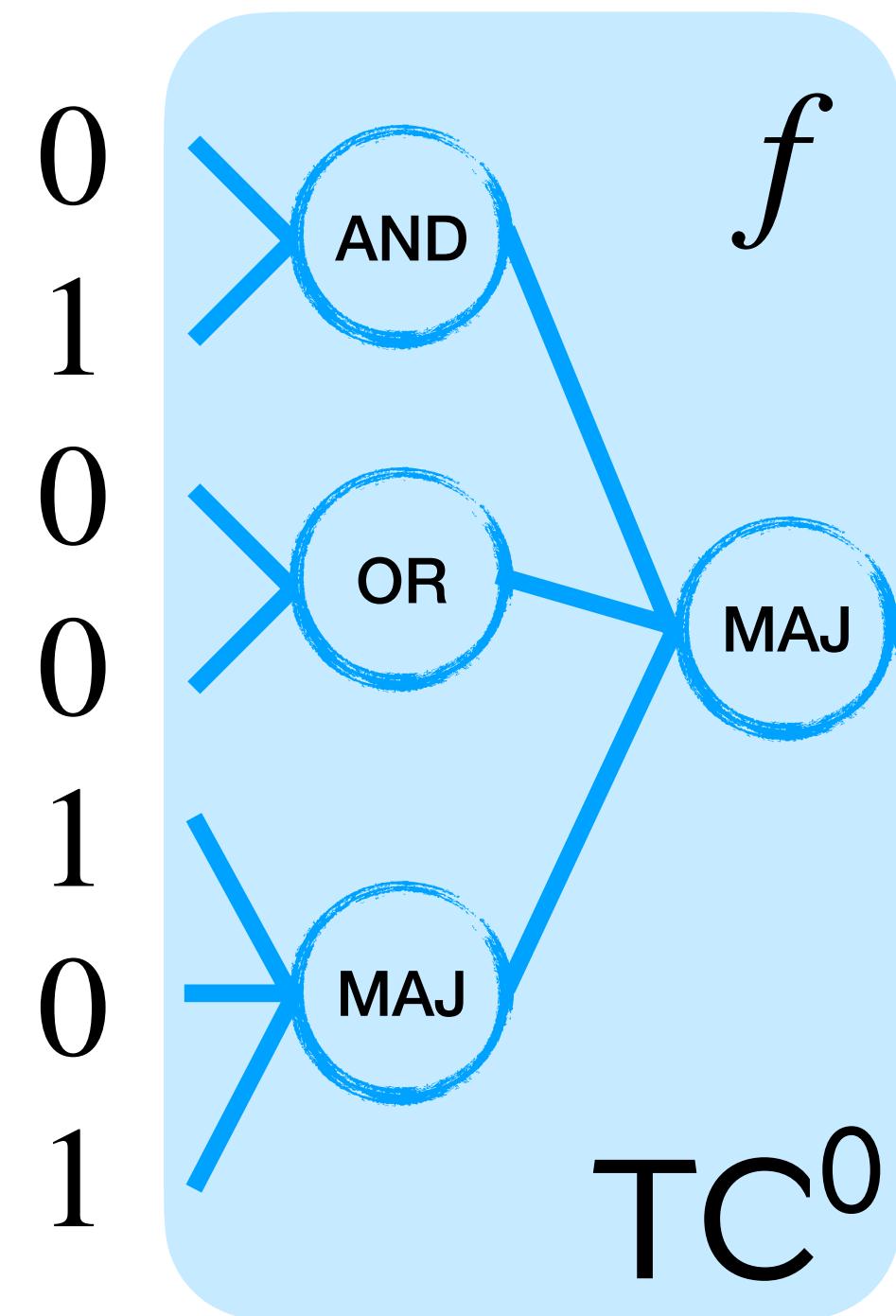
computational complexity



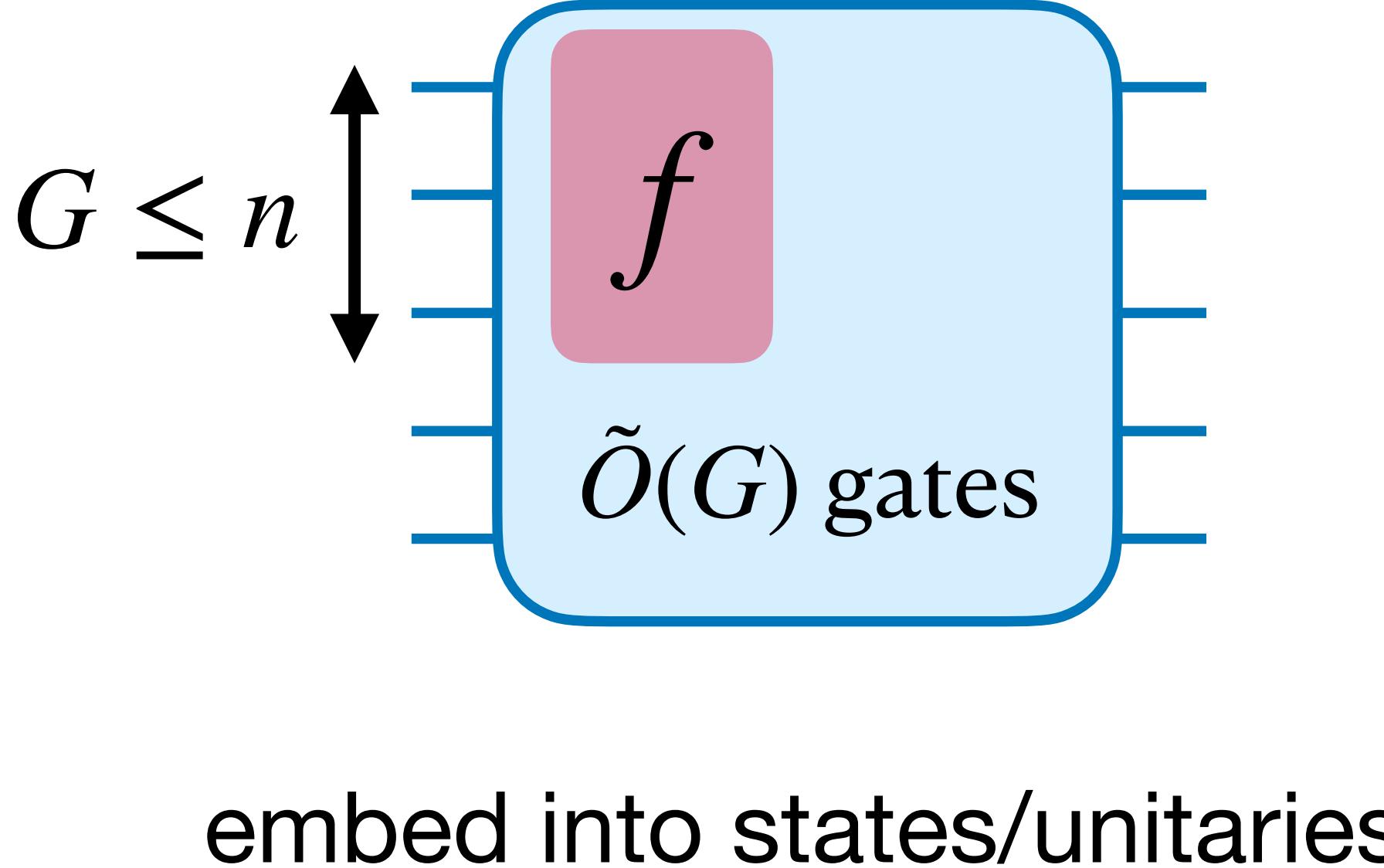
Learning $|\psi\rangle$ in ϵ trace distance
requires $T = \exp(\Omega(\min\{G, n\}))$ time,
if RingLWE is sub-exponential hard.

1. Complexity of learning = $e^{\text{complexity of creating}}$ (computationally), efficient $\log n$ gates
2. Even for **quantum learners**: RingLWE is expected to be hard for quantum computers.
3. **Worst-case** statement: efficiency possible with additional assumptions.
4. Same for average-case unitary learning

Proof sketch computational complexity



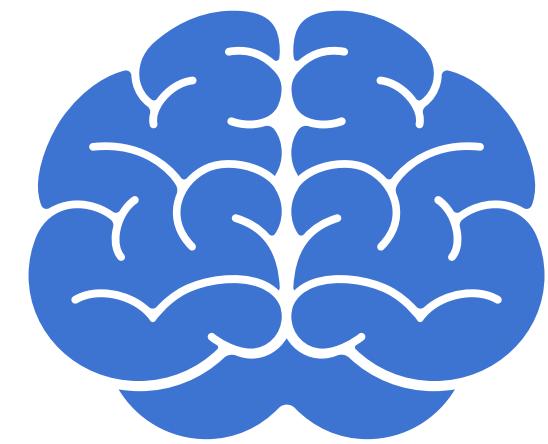
pseudorandom functions



Learning breaks PRF/PRS
requires $T = e^{\Omega(G)}$ time

$$G \leq n$$

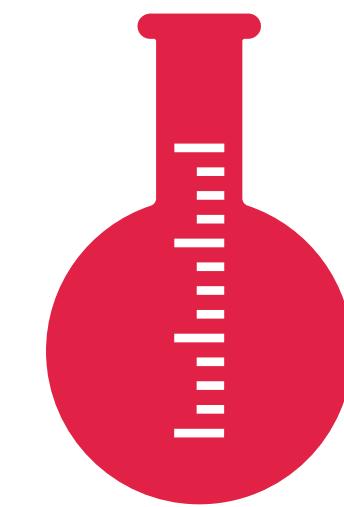
Message



sample size

$$N \approx \log T \approx G \approx t$$

compute time

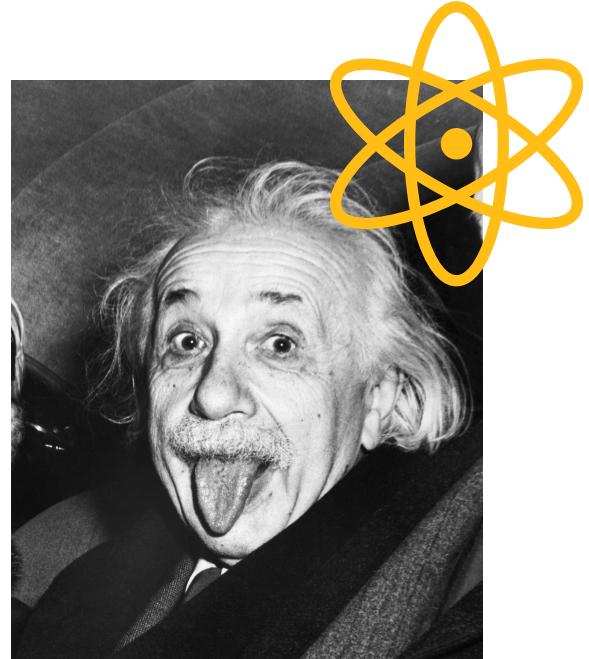


gate complexity

evolution time

(Brown-Susskind conjecture)

Learning physical states/unitaries is information-theoretically easy,
but computationally hard!

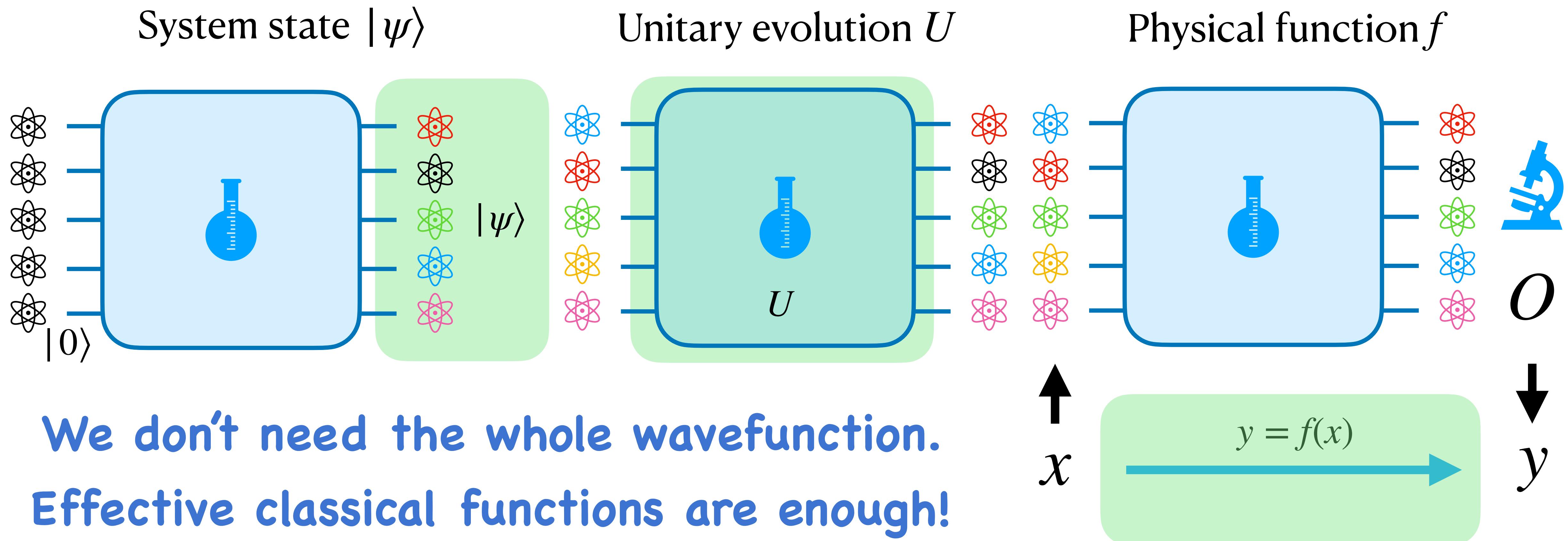


In scientific discovery, a few samples might already be enough,
but coming up with a theory requires some real genius!

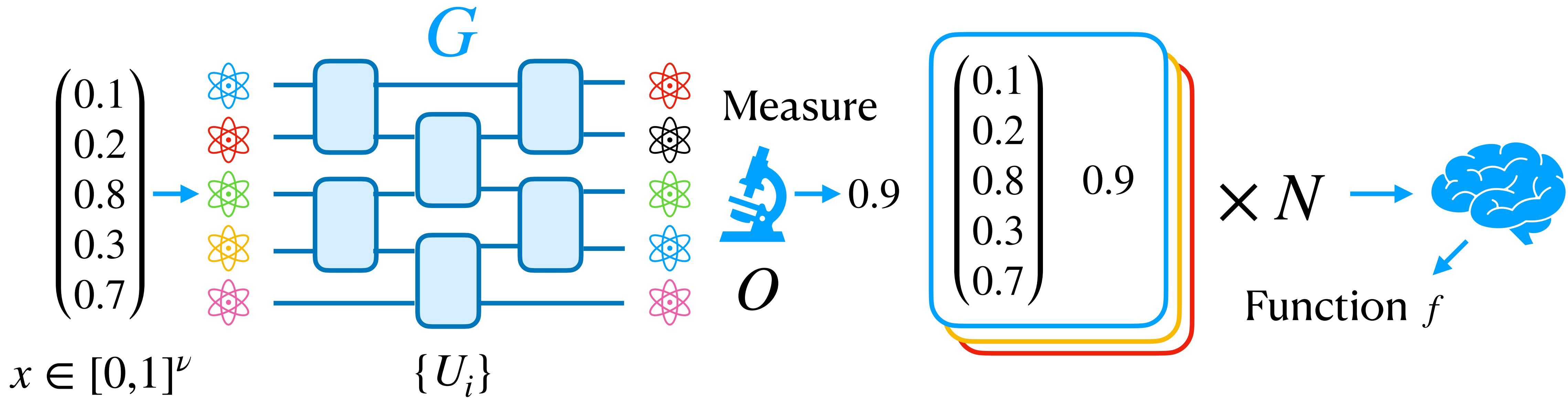
Reminder:
Doing math proofs needs
no data but is NP-hard!

Motivation

What to learn?



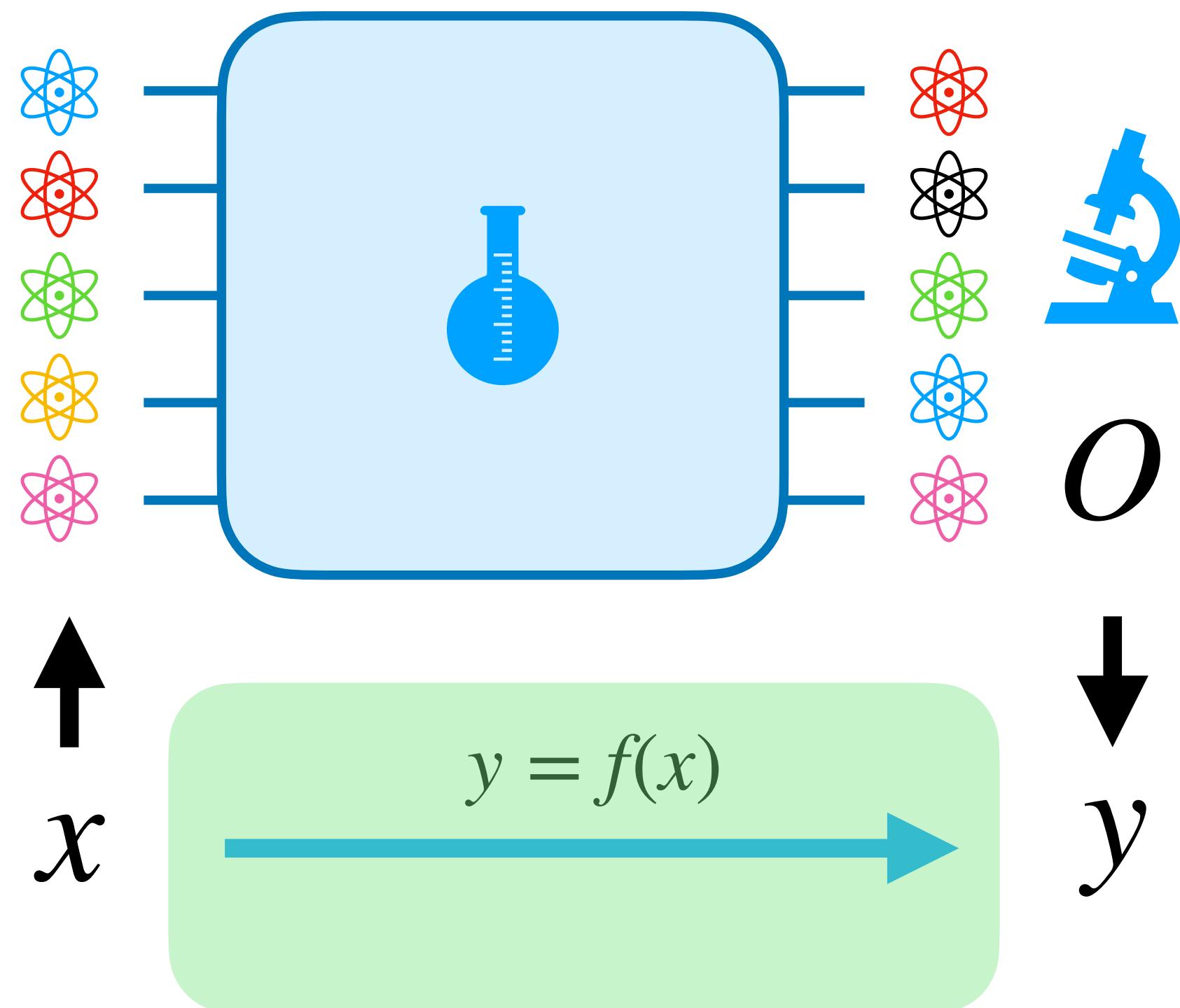
Physical functions



$$f(x) = \langle 0^n | U^\dagger(x, \{U_i\}) O U(x, \{U_i\}) | 0^n \rangle$$

What kind of functions are physically implementable?

Physical function f



Results

To implement/learn 1-Lip 1-bounded functions in ϵ infinite norm,
 $G, N = \tilde{\Omega}(1/\epsilon^\nu)$ gates/samples are needed

1. Certain well-behaved function class is not physical! $G = \Omega(\exp \nu)$
2. For these functions, no quantum advantage! $\tilde{\Theta}(1/\epsilon^\nu)$ with classical ReLU NN.
3. More restricted function classes: possible physicality/advantage
4. Complement universal approximation theorems of QML

Fourier integrable $O(1/\epsilon^2)$
Gonon and Jacquier 2023

Pérez-Salinas et al., PRA 2021
Schuld, Sweke, Meyer, PRA 2021
Manzano, Dechant, Tura, Dunjko 2023

Proof Sketch

$$U_i |\psi\rangle = \begin{pmatrix} U_i^{11} & U_i^{12} \\ U_i^{21} & U_i^{22} \end{pmatrix} \begin{pmatrix} \psi_1 \\ \psi_2 \end{pmatrix} = \begin{pmatrix} U_i^{11}\psi_1 + U_i^{12}\psi_2 \\ U_i^{21}\psi_1 + U_i^{22}\psi_2 \end{pmatrix}$$

$$f(x) = \langle 0^n | U^\dagger(x, \{U_i\}) O U(x, \{U_i\}) | 0^n \rangle$$

$f(x)$ is a polynomial of $\{U_i\}$ with degree at most $2G$

\Rightarrow fat-shattering/pseudo dimension $\leq \tilde{O}(G)$

Approximating 1-Lip 1-bounded functions require $\Omega(1/\epsilon^\nu)$

$\Rightarrow G, N \geq \tilde{\Omega}(1/\epsilon^\nu)$

Summary and Outlook

- Learning complexity = circuit complexity (information theoretically)
- Learning complexity = $e^{\text{circuit complexity}}$ (computationally)
- Possible efficient learning for more restricted states/unitaries? Clifford+T, MPS, Shallow Circuits, Fermion, Boson, etc.
- What physical properties does learning complexity relate to? Geometry? Phase? Thermalization? Entanglement?
- Worst-case unitary learning requires exponential samples. Du, Hsieh, Tao, arXiv last week
- Initiate the study of physical functions. Learning complexity? Q signal processing?
- Information-theoretic quantum no-free-lunch theorem (not covered)
- Technical: matching bounds, Heisenberg scaling, mixed states/channels, etc.