

# Heapsort

- $\Theta(n \lg n)$  like merge sort
- In-place like insertion sort
- Uses a data structure - **heap**

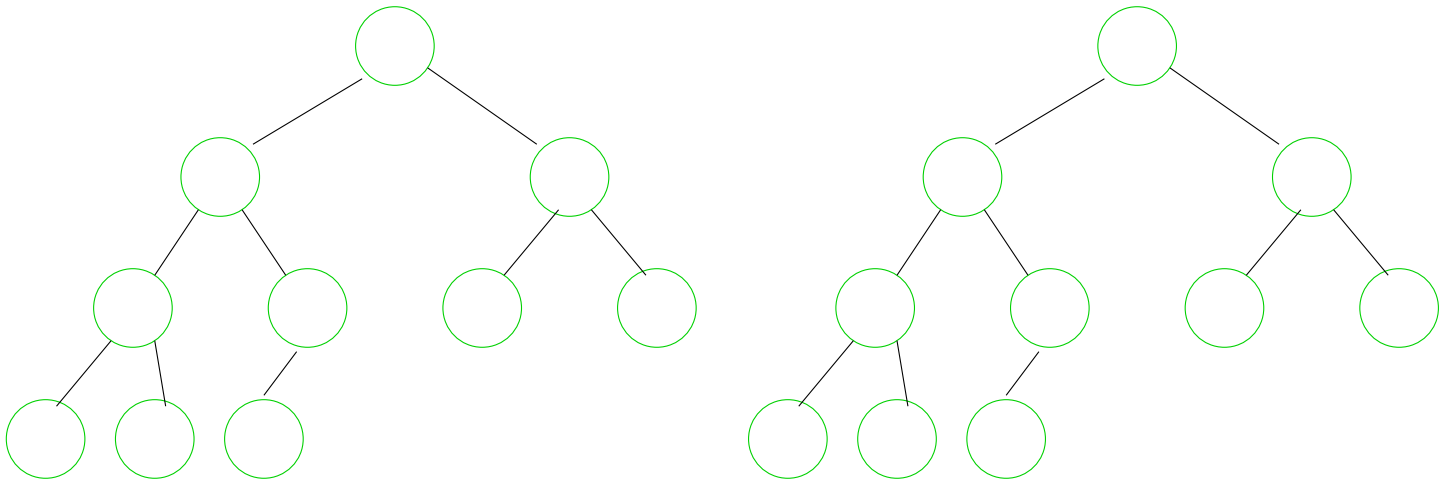
# HEAP

## Structure:

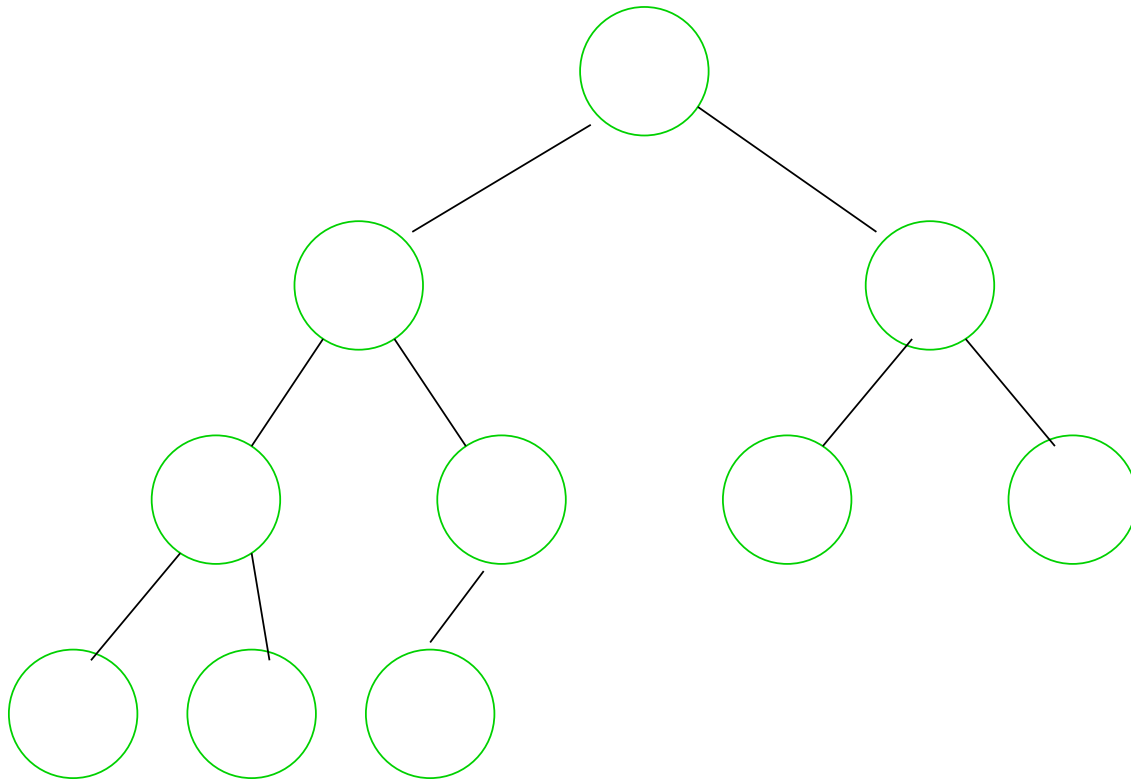
- labeled binary tree
- all levels full, except possibly last
- nodes in last level as far left as possible

## Max-heap property:

- For all nodes  $x$ ,  $\text{label}(x) \geq$  labels of children of  $x$ .



# Implementation



10	9	7	8	4	6	1	3	5	2
1	2	3	4	5	6	7	8	9	10

If node  $x$  is in location  $i$

$$\text{lchild}(x) \longrightarrow 2i$$

$$\text{rchild}(x) \longrightarrow 2i + 1$$

$$\text{parent}(x) \longrightarrow \lfloor i/2 \rfloor$$

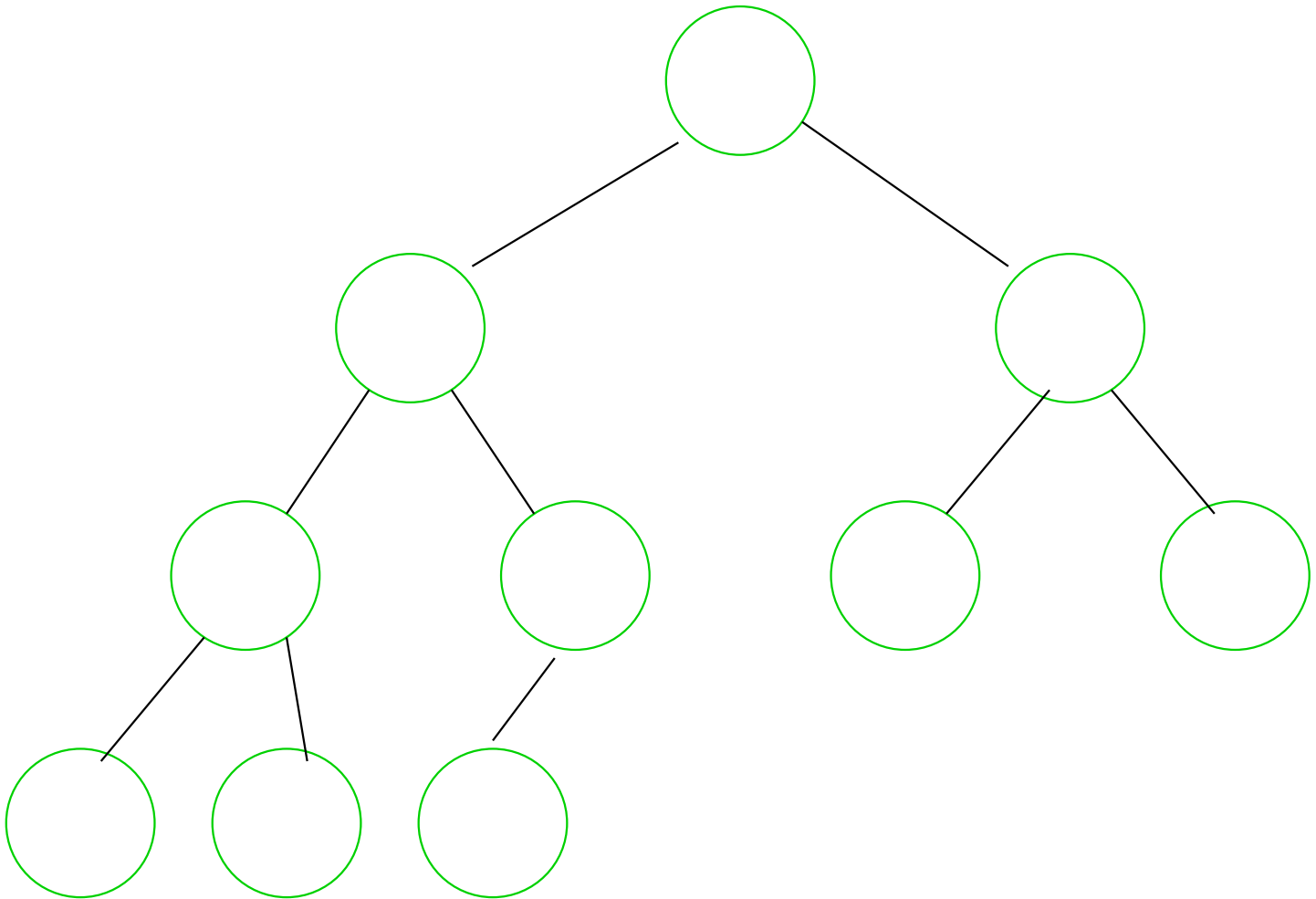
# How to construct heap?

## How to use heap?

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### Heapify:

“Fixes” tree which is heap except possibly at root.



(Root value filters down.)

## Heapify( $x$ )

▷ Assumes tree rooted at  $x$  is heap except possibly at root

**if**  $x$  is not a leaf **then**

    let  $y$  be child of  $x$  with maximum value

**if**  $\text{value}(x) < \text{value}(y)$  **then**

        swap values of  $x$  and  $y$

        Heapify( $y$ )

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Max. # of compares for Heapify( $x$ ) in tree with  $n$  nodes?

(2 \* height of  $x$  in tree)

# Height $h$ of heap with $n$ nodes

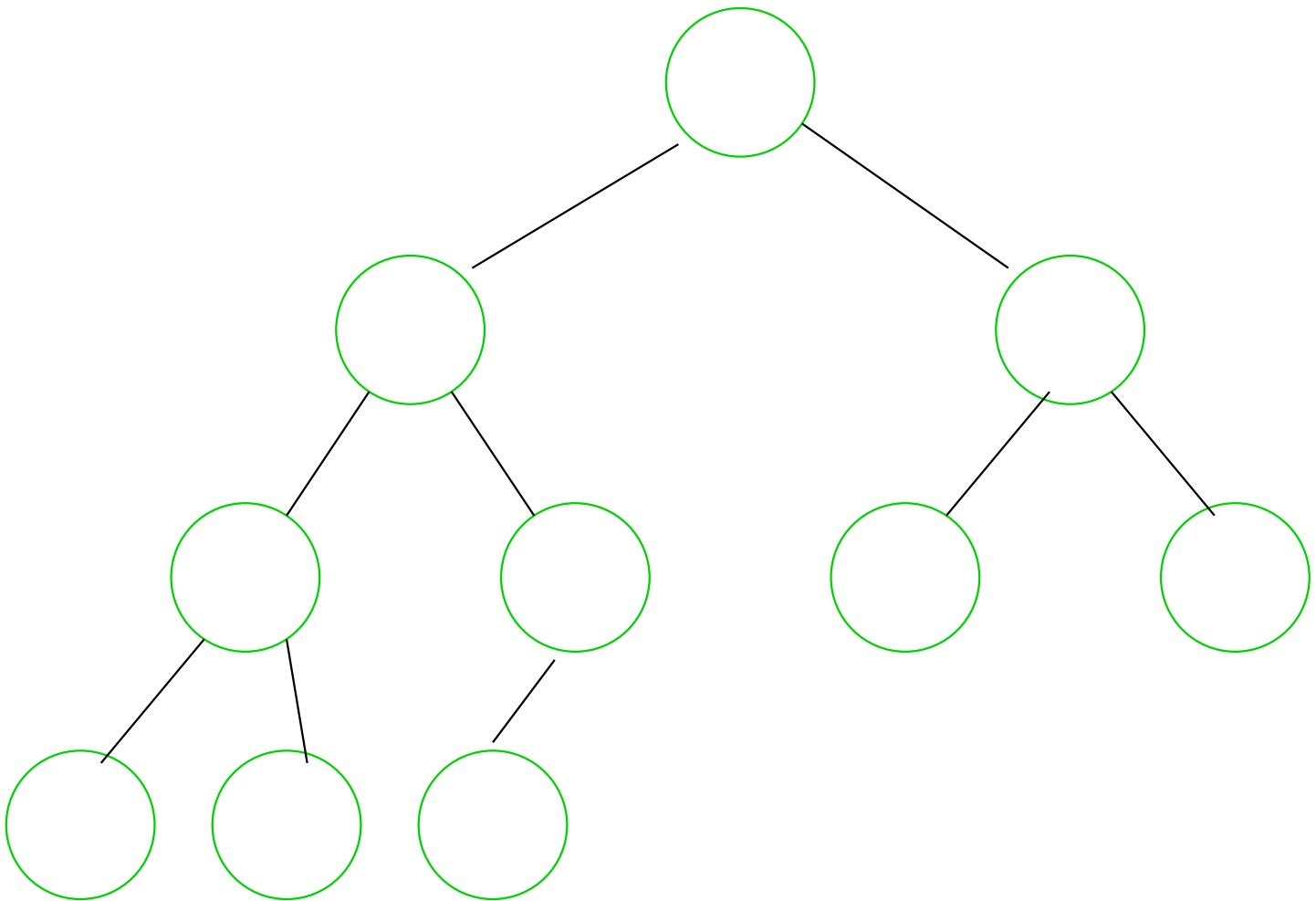
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$$1 + \sum_{i=0}^{h-1} 2^i \leq n \leq \sum_{i=0}^h 2^i$$

$$2^h \leq n \leq 2^{h+1} - 1$$

# How to organize $n$ elements into a heap?

1. Place elements arbitrarily in tree.
2. **Buildheap**: Heapify( $x$ ) as  $x$  runs through all nodes, from bottom to top level, right-to-left across levels.



# Array implementation

Place elements in  $A[1..n]$ , then:

## Buildheap

**for**  $i \leftarrow n$  **downto** 1 **do** (\*)

**Heapify**( $i$ )

(\*) could start at  $\lfloor n/2 \rfloor$

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Total # of compares:

$$\leq \sum_{i=1}^n (2 * \text{height of node } i)$$

Upper Bound:

$$\sum_{i=1}^n 2 * \lfloor \lg n \rfloor = 2n \lfloor \lg n \rfloor$$

## Better Bound?



# Better Bound

Let  $h = \lfloor \lg n \rfloor$ .

$$\sum_{i=1}^n (2^{\text{height of node } i})$$

$$= \sum_{j=0}^h 2^j * \# \text{ of nodes of height } j$$

$$\leq 2 \sum_{j=0}^h j * 2^{h-j}$$

$$= 2^{h+1} \sum_{j=0}^h j * (1/2)^j$$

$$\leq 2^{h+1} \sum_{j=0}^{\infty} j * (1/2)^j = 2^{h+1} \frac{1/2}{(1 - 1/2)^2}$$

$$= 2^{h+2} = 4 * 2^h \leq 4n$$

**linear!**

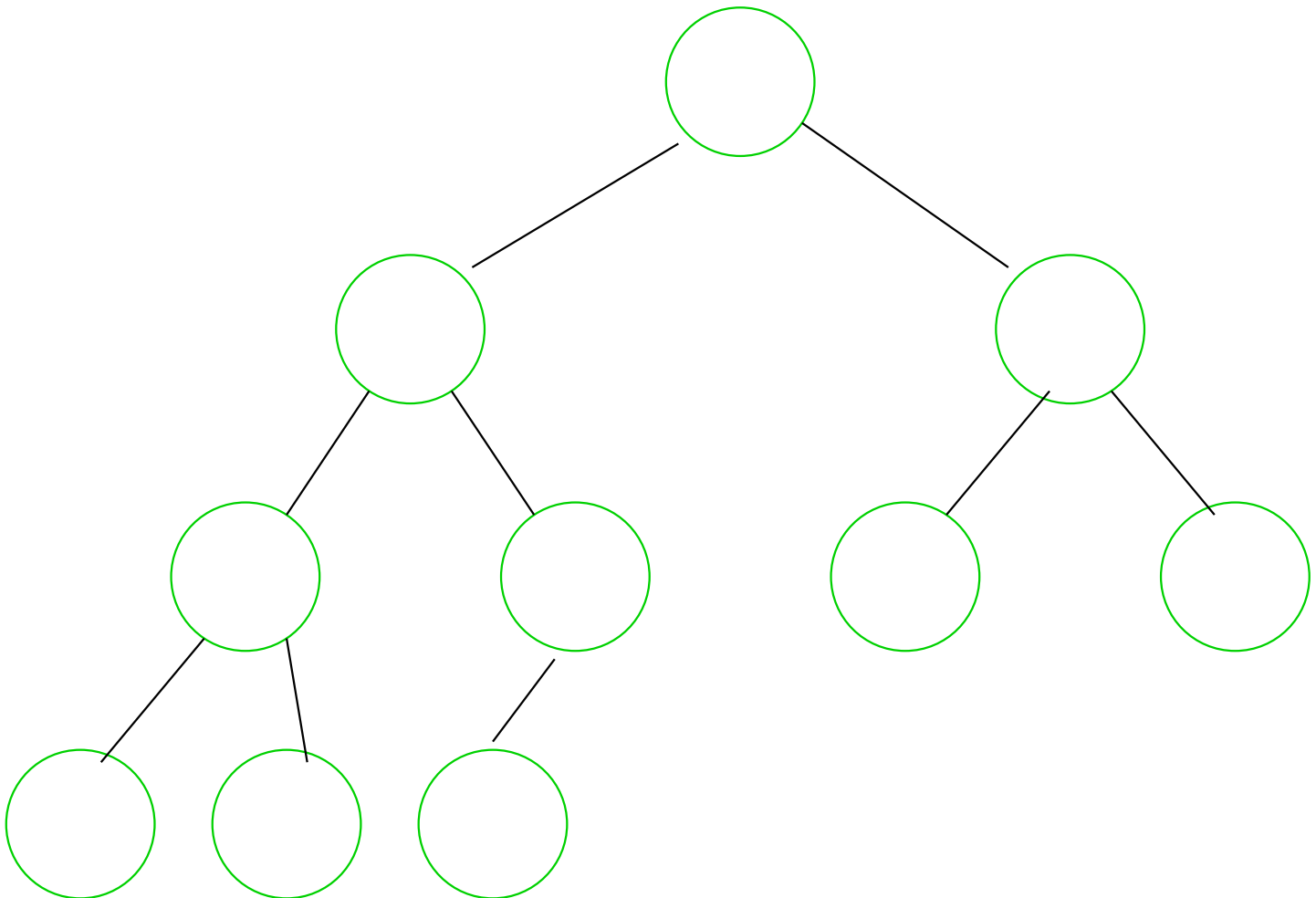
# Heapsort

1. Builheap

2. **while** heap not empty **do**

- Remove maximum element (root)
- Move last element at bottom level to root
- Heapify

(# of compares  $\leq 4n + 2n \lg n$ )



# Priority Queue

Abstract data structure

**Objects:** Set  $S$  of records, each with a key from a totally ordered set.

**Operations:**

- **Maximum( $S$ ):** returns element of  $S$  with maximum key
- **Extract-Max( $S$ ):** returns and deletes element of  $S$  with maximum key
- **Insert( $S, x$ ):** adds element  $x$  to  $S$
- **Increase-Key( $S, x, k$ ):** increases the key value of  $x$  in  $S$  to  $k$ ; assumes  $k \geq$  current key of  $x$  and that a pointer is given to  $x$ .

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Implement with **HEAP**

- how to implement each operation?
- w.c. time for each operation?