

# Efficient Lossy Trapdoor Functions based on Subgroup Membership Assumptions

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## 1 Introduction

## 2 Our Contribution

- $SMA \implies LTDF$
- Concrete Examples

## 3 Conclusion

# Outline

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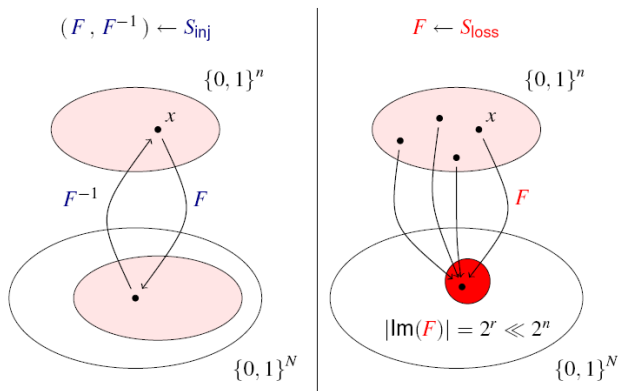
# Lossy Trapdoor Function (LTDF)

Peikert and Waters proposed the LTDF in STOC 2008.

$$DDH, LWE \rightarrow \textcolor{red}{LTDF} \rightarrow \begin{cases} TDF, \text{ Hard Core}; \\ OT; \\ CR \text{ Hash}; \\ CCA, \dots \end{cases}$$

# Lossy Trapdoor Function [PW'08]

From Peikert's slides



$$F \stackrel{c}{\approx} F$$

# Definition of LTDF

## Injective model

- $(s, t) \leftarrow S_{inj}(1^n);$
- $F_{ltdf}(s, \cdot) : \{0, 1\}^m \rightarrow \{0, 1\}^*$
- $F_{ltdf}^{-1}(t, F_{ltdf}(s, x)) = x.$

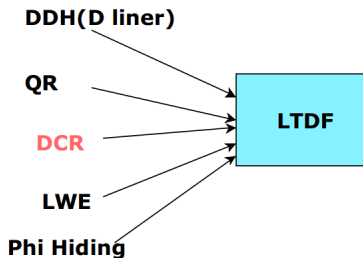
## Lossy with $l$ bits

- $s \leftarrow S_{loss}(1^n);$
- $F_{ltdf}(s, \cdot) : \{0, 1\}^m \rightarrow \{0, 1\}^*$
- $F_{ltdf}(s, \cdot)$  has size at most  $2^{m-l};$

$$\{s : s \leftarrow S_{lossy}\} \stackrel{c}{\approx} \{s : (s, t) \leftarrow S_{inj}\}.$$

# Constructions of LTDF

- DDH or  $d$ -liner  
[PW'08],[FGKRS'10], [Wee12];
- QR assumption  
[FGKRS'10],[JL'13], [Wee12]
- DCR assumption [BFO'08],  
[FGKRS'10], [Wee12]
- LWE assumption  
[PW'08],[Wee12]
- $\Phi$ -Hiding [KOS'10].



The DCR based construction is one of the most efficient constructions.

# DCR Assumption [Pai99, Dam01]

## Definition

Let  $N = pq$  for  $p = 2p' + 1$ ,  $q = 2q' + 1$  and  $s \geq 2$

$$P := \{a = x^{N^{s-1}} \bmod N^s \mid x \in \mathbb{Z}_N^*\},$$

$$M := \{a = (1 + N)^y x^{N^{s-1}} \bmod N^s \mid x \in \mathbb{Z}_N^*, y \in \mathbb{Z}_{N^{s-1}}\}.$$

$$\{a \leftarrow P\} \stackrel{c}{\approx} \{a \leftarrow M\}$$

①  $N^{s-1}$ -th residuosity is a **subgroup** with order  $2p'q' \approx N/2$ .

② For  $a$  in  $M$ ,

$$a^{2p'q'} = 1 + y2p'q'N \bmod N^s.$$



# DCR Based LTDF

For input  $m \in [0, N^{s-1}]$ , the two function models follow:

*Injective model*

$$\{(1 + N)x^{N^{s-1}}\}^m$$

*Lossy model*

$$\{x^{N^{s-1}}\}^m$$

- $G_{N^{s-1}} = H \times K = \langle (1 + N) \rangle \times \{x^{N^{s-1}}\}$
- $s \geq 3$  in order to make enough lossiness.

# Motivation

*General Subgroup membership assumption*  $\xrightarrow{?}$  *LTDF*

$$\text{mod } N^3 \xrightarrow{?} \text{mod } N^2 \xrightarrow{?} \text{mod } N$$

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# Our Contribution

*Subgroup membership assumption + 2 Properties  $\xrightarrow{\checkmark}$  LTDF*

$$\text{mod } N^3 \xrightarrow{\checkmark} \text{mod } N^2 \xrightarrow{\checkmark} \text{mod } N$$

*Shrinking the subgroup or Enlarging the quotient group.*

# Subgroup Membership Assumption [Gjøsteen 05]

## Definition (SMA)

Let  $G$  be a finite cyclic group.

$$G = \langle g \rangle = G/K \times K = G/K \times \langle h \rangle$$

The subgroup membership assumption  $SM_{(G,K)}$  asserts that,

$$\{x, x \leftarrow K\} \stackrel{c}{\approx} \{x, x \leftarrow G \setminus K\}.$$

$$\mathbb{Z}_{N^s}^* = \langle (1 + N) \times \{x^{N^{s-1}}\} \rangle.$$

## 2 Properties

- ①  $SDL_{(G,K,g)}$  is easy with a trapdoor  $t$ ;
- ②  $|G/K| \gg |K|$ . (Lossy property)

Definition (Subgroup Discrete Logarithm Problem [Gjøsteen 05])

If  $\varphi : G \rightarrow G/K$  is the canonical epimorphism, then  $SDL_{(G,K,g)}$  is:

To compute  $\log_{\varphi(g)}(\varphi(x))$  for  $x \leftarrow G$ .

$$(1 + N)^y z^{N^{s-1}} \rightarrow y?$$

# Generic construction

Let  $(G, K, g, h, t)$  be an instance of  $SM_{(G,K)}$  with 2 properties. For  $m \in [0, |G/K|]$ , the two models follow,

## *Injective model*

- 1  $a = gh^r$  for  $r \leq |K|$  and  $t=t$ ;
- 2  $F_{ltdf}(a, m) = a^m = [gh^r]^m$
- 3 Recover  $m$  by solving  $SDL_{(G,K,g)}$  with  $t$ .

## *Lossy model*

- 1  $a = h^r$  for  $r \leq |K|$ ;
- 2  $F_{ltdf}(a, m) = a^m = [h^r]^m$
- 3  $|F_{ltdf}(a, \cdot)| < |K|$  as  $F_{ltdf}(a, \cdot)$  falls into  $K$  ;

# SMA $\Rightarrow$ LTDF

## Theorem (1 in page 240)

*If the  $SM_{G,K}$  with two above properties holds, This is an  $(\log |G/K|, \log |G/K| - \log |K|)$  LTDF.*



# DCR& QR based LTDF

Let  $N = pq$  with  $p = 2^k p' + 1, q = 2^k q' + 1$ .

- For  $y \in QNR_N$ , let  $G = \langle (1 + N)y^N \rangle$  with order  $N2^k p' q'$ ;
- For  $h_1 \in \mathbb{Z}_N^*$ , let  $K = \langle h_1^{2^k N} \rangle$  with order  $p' q'$ .

Theorem (3 in page 243)

$$DCR\&QR \Rightarrow SM_{(G,K)}.$$

## Extended $p$ -subgroup based LTDF

Let  $N = p^2q$  with  $p = 2p' + 1, q = 2q' + 1$ , For  $y \in \mathbb{Z}_N^*$ , Let  $h = y^{2N^2}$

- Let  $G = \langle (1 + N)h \rangle$  with order  $Np'q'$ ;
- Let  $K = \langle h \rangle$  with order  $p'q'$ .

$SM_{(G,K)}$  is a generalization of  $p$  subgroup in [OU98]

# Decisional RSA [Groth 05] based LTDF

Let  $N = pq$  with  $p = 2p'r_p + 1, q = 2q'r_q + 1$ , Let  $r_p, r_q$  be  $B$ -smooth with  $t$  distinct prime factors and  $l \approx \log B$ .

For  $x \in \mathbb{Z}_N^*$ , let  $h = x^{2r_p r_q}$  and  $g \leftarrow QR_N$ .

- Let  $G = \langle g \rangle$  with order larger than  $p'q'2^{(t-d)(l-1)}$ ;
- Let  $K = \langle h \rangle$  with order  $p'q'$ .

This  $SM_{(G,K)}$  assumption is the Decisional RSA assumption in [Groth 05].

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# Comparison with previous constructions

Assumption	Input size	Lossiness	Index size	Efficiency
DDH	$n$	$n -  \mathbb{G} $	$n^2 \mathbb{G}$	$n^2$ Multi
LWE	$n$	$cn$	$n(d+w)\mathbb{Z}_q$	$n(d+w)$ Multi
d-linear	$n$	$n - d \mathbb{G} $	$n^2 \mathbb{G}$	$n^2$ Multi
QR	$\log N$	1	$\mathbb{Z}_N^*$	1 Multi
DDH& QR	$n$	$n - \log N$	$(\frac{n}{k})^2 \mathbb{Z}_N^*$	$(\frac{n}{k})^2$ Multi
$\Phi$ -hiding	$\log N$	$\log e$	$\mathbb{Z}_N^*$	$\log e \log N$
DCR	$2 \log N$	$\log N$	$\mathbb{Z}_{N^3}^*$	$3 \log x \log N$
QR & DCR	$\frac{9}{8} \log N$	$\frac{3}{8} \log N$	$\mathbb{Z}_{N^2}^*$	$2 \log x \log N$
E $p$ -sub	$\log N$	$\frac{1}{3} \log N$	$\mathbb{Z}_{N^2}^*$	$2 \log x \log N$
D RSA	$l_x$	$l_x - l_{p'} - l_{q'}$	$\mathbb{Z}_N^*$	$\log x \log N$

$$l_x \leq 698, l_{p'} = l_{q'} = 160$$

# Conclusion

We present a generic construction of LTDFs from subgroup membership assumptions.

We give three efficient constructions based on

- 1 DCR & QR;
- 2 Extended  $p$  Subgroup;
- 3 Decisional RSA.

**Thank you**