Curs 9

Taskuri

Future-Promise

Executie asincrona

Executori

Apeluri asincrone

Future

Promise

Istoric

- Termenul promise a fost propus de catre Daniel P. Friedman si David Wise in 1976;
- ~ aceeasi perioada Peter Hibbard I-a denumit eventual;
- conceptul future a fost introdus in 1977 intr-un articol scris de catre Henry Baker si Carl Hewitt .
- Future si promise isi au originea in programarea functionala si paradigmele conexe (progr. logica)
- Scop: decuplarea unei valori (a future) de ceea ce o calculeaza
 - Permite calcul flexibil si paralelizabil
- Folosirea in progr. Paralela si distribuita a aparut ulterior mai intai pentru
 - reducerea latentei de comunicatie (round trips).
 apoi
 - in programele asincrone.

Promise pipelining

Barbara Liskov and Liuba Shrira in 1988 Mark S. Miller, Dean Tribble and Rob Jellinghaus 1989

```
Conventional RPC

t3 := (x.a()).c(y.b())

Echivalent cu

t1 := x.a();

t2 := y.b();

t3 := t1.c(t2); //executie dupa ce t1 si t2 se termina

Apel remote atunci este nevoie de 3 round-trip.

(a,b,c se executa remote)

Folosind futures

("Dataflow" with Promises)

t3 := (x <- a()) <- c(y <- b())

Echivalent cu

t1 := x <- a()

t2 := y <- b()
```

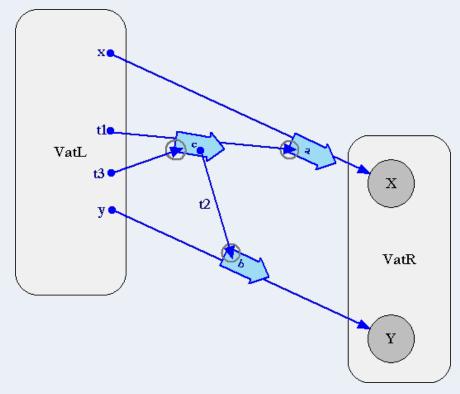
t3 := t1 < -c(t2)

Daca x, y, t1, si t2 sunt localizate pe aceeasi masina remote atunci se poate rezolva in 1 round-trip. - O cerere trimisa si un raspuns necesar!

Promise Pipelining

From http://www.erights.org/elib/distrib/pipeline.html

The key is that later messages can be sent specifying promises for the results of earlier messages as recipients or arguments, despite the fact that these promises are unresolved at the time these messages are sent. This allows to stream out all three messages at the same time, possibly in the same packet.



Evaluare

- call by future
 - non-deterministic: valoarea se va calcula candva intre momentul crearii variabilei future si momentul cand aceasta se va folosi
 - eager evaluation: imediat ce future a fost creata
 - lazy evaluation, doar atunci cand e folosita
 - Odata ce valoarea a fost atribuita nu se mai recalculeaza atunci cand se refoloseste.
- lazy future: calculul valorii incepe prima oara cand aceasta este ceruta (folosita)
 - in C++11
 - Politica de apel std::launch::deferred ->La apelul std::async.

Future and Promise

- the two sides of an asynchronous operation:
- consumer/caller vs. producer/implementor
- a caller of an asynchronous task will get a Future as a handle to the computation's result
- Future handles the computation's result
 - e.g. call get()
- The implementor must return a Future
 - it is responsible for completing that future as soon as the computation is done.

Blocking vs non-blocking semantic

- Accesare sincrona->
 - De exemplu prin transmiterea unui mesaj (se asteapta pana la primirea mesajului)
- Accesare sincrona-> posibilitati:
 - Accesul blocheaza threadul curent /procesul pana cand se calculeaza valoarea (eventual timeout).
 - Accesul sincronizat produce o eroare (aruncare exceptie)
 - Se poate obtine fie succes daca valoarea este deja calculata sau se transmite eroare daca nu este inca calculata -> poate introduce race conditions.
- in C++11, un thread care are nevoie de valoarea unei *future* se poate bloca pana cand se calculeaza (wait() ori get()). Eventual timeout.
 - Daca future a aparut prin apelul de tip std::async atunci un apel wait cu blocare poate produce invocare sincrona a functiei care calculeaza rezultatul.

C++11

- future
 - promise
 - async
 - packaged_task

Future

- (1) future from a packaged_task
- (2) future from an async()
- (3) future from a promise

packaged_task

- std::packaged_task object = wraps a callable object
 - can be wrapped in a std::function object,
 - passed to a std::thread as the thread function,
 - passed to another function that requires a callable object,
 - invoked directly.

async

- Depinde de implementare daca std::async porneste un nou thread sau daca taskul se va executa sincron atunci cand se cere valoarea pt future.
 - std::launch::deferred se amana pana cand se apeleaza fie wait() fie get() si se va rula in threadul curent (care poate sa nu fie cel care a apelat async)
 - std::launch::async se ruleaza in thread separat. (lazy evaluation)

async

- async
 - Executa o functia f asyncron
 - posibil in alt thread si
 - returneaza un obiect std::future care va contine rezultatul

std::promise

 Furnizeaza un mecanism de a stoca o valoare sau o exceptie care va fi apoi obtinuta asincron via un obiect <u>std::future</u> care a fost creat prin obiectul <u>promise</u>.

Actiuni:

- make ready: se stocheaza rezultatul in 'shared state'.
 - Deblocheaza threadurile care asteapta actualizarea unui obiect future asociat cu 'shared state'.
- release: se elibereaza referinta la 'shared state'.
- abandon: shared state = ready +
 - exception of type <u>std::future_error</u> with error code std::future_errc::broken_promise

std::promise

promise furnizeaza un obiect future.

- Se furnizeaza si un mecanism de transfer de informatie intre threaduri
 - T1-> wait()
 - T2-> set_value() => future ready.
- d.p.d.v al threadului care asteapta nu e important de unde a aparut informatia.

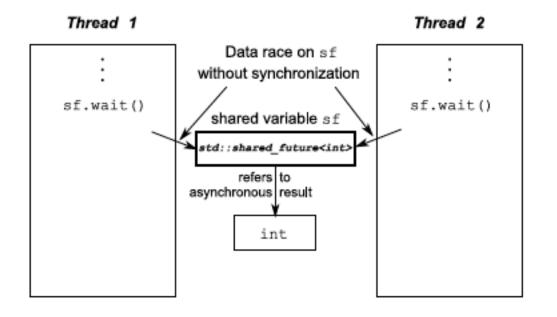
Exemple

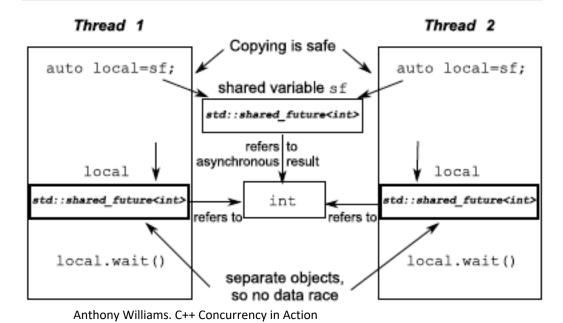
```
//(1) future from a packaged_task
  std::packaged_task<int()> task( []() { return 7; } ); // wrap the function
  std::future<int> f1 = task.get future(); // get a future
  std::thread(std::move(task)).detach(); // launch on a thread
  //(2) future from an async()
  std::future<int> f2 = std::async(std::launch::async, [](){ return 8; });
  // (3) future from a promise
  std::promise<int> p;
  std::future<int> f3 = p.get future();
  std::thread( [&p]{ p.set_value_at_thread_exit(9); }).detach();
  f1.wait();
  f2.wait();
  f3.wait();
```

std::shared_future

PROBLEMA: daca se acceseaza un obiect std::future din mai multe threaduri fara sincronizare aditionala => data race.

- std::future modeleaza unique ownership
 - doar un thread poate sa preia valoarea
- std::shared_future permite accesarea din mai multe threaduri
- std::future este moveable (ownership can be transferred between instances)
- std::shared_future este copyable (mai multe obiecte pot referi aceeasi stare asociata).





std::shared_future, member functions on an individual object are still unsynchronized.

- To avoid data races when accessing a single object from multiple threads, you must protect accesses with a lock.
- The preferred way to use it would be to take a copy of the object instead and have each thread access its own copy.
- Accesses to the shared
 asynchronous state from
 multiple threads are safe if
 each thread accesses that state
 through its own
 std::shared_future object.

Java

- task (Runnable vs. Callable)
- Future
- Executor
- CompletableFuture

Task

- Task = activitate independenta
- Nu depinde de :
 - starea,
 - resultatul, ori
 - 'side effects'

ale altor taskuri

=> Concurenta /Paralelism

Ref: Java Concurrency in Practice

Curs 9 - PPD -

Exemplul 1

Aplicatii client server-> task = cerere client

```
class SingleThreadWebServer {
  public static void main(String[] args) throws IOException {
     ServerSocket socket = new ServerSocket(80);
     while (true) {
          Socket connection = socket.accept();
          handleRequest(connection);
     }
}
....
}
```

Analiza

- procesare cerere =
 - socket I/O (read the request + write the response) -> se poate bloca
 - file I/O or make database requests-> se poate bloca
 - Procesare efectiva

Single-threaded => ineficient

- Timp mare de raspuns
- Utilizare ineficienta CPU

Exemplu 2

```
class ThreadPerTaskWebServer {
     public static void main(String[] args) throws IOException {
          ServerSocket socket = new ServerSocket(80);
             while (true) {
              final Socket connection = socket.accept();
                  Runnable task = new Runnable() {
                       public void run() {
                         handleRequest(connection);
               new Thread(task).start();
```

Dezavantaje ale nelimitarii numarului de threaduri create

Thread lifecycle overhead

Creare threaduri

Resource consumption

- Threadurile active consuma resursele sistemului (memorie)
- Multe threaduri inactive bloacheaza spatiu de memorie -> probleme garbage collector
- Multe threaduri => probleme cu CPU-> costuri de performanta

Stability

- exista o limita a nr de threaduri care se pot crea (depinde de platforma)
- -> OutOfMemoryError.

Executori

- Task = unitate logica
- Thread -> un mecanism care poate executa taskurile asincron
- Interfata/obiect Executor
 - Mecanism de decuplare a submiterii unui task de executia lui
 - Suport pentru monitorizarea executiei
 - Se bazeaza pe sablonul producator-consumator

```
public interface Executor {
    void execute(Runnable command);
}
```

Exemplu 3

```
class TaskExecutionWebServer {
    private static final int NTHREADS = 50;
    private static final Executor exec= Executors.newFixedThreadPool(NTHREADS);
    public static void main(String[] args) throws IOException {
         ServerSocket socket = new ServerSocket(80);
         while (true) {
               final Socket connection = socket.accept();
               Runnable task = new Runnable() {
                    public void run() {
                      handleRequest(connection);
               };
               exec.execute(task);
```

Adaptare – task per thread

```
public class ThreadPerTaskExecutor implements Executor {
   public void execute(Runnable r) {
        new Thread(r).start();
   };
public class WithinThreadExecutor implements Executor {
   public void execute(Runnable r) {
        r.run();
   };
//Executor care executa taskurile direct in threadul apelant (synchronously).
```

Execution policy

"what, where, when, how" pentru executia taskurilor

- = instrument de management al resurselor
- In ce thread se executa un anumit task?
- In ce ordine se aleg taskurile pentru executie (FIFO, LIFO, priority)?
- Cate taskuri se pot executa concurent?
- Cate taskuri se pot adauga in coada de executie?
- Daca sistemul este supracincarcat *overloaded*, care task se va alege pentru anulare si cum se notifica aplicatia care l-a trimis?
- Ce actiuni trebuie sa fie facute inainte si dupa executia unui task?

Thread pool

- Un executor care gestioneaza un set omogen de threaduri = worker threads
- Se foloseste
 - work queue pentru stocare task-uri
- Worker thread =>
 - cerere task din work queue,
 - Executie task
 - Intoarcere in starea de asteptare task.

Variante

newFixedThreadPool.

A fixed-size thread pool creates threads as tasks are submitted, up to the maximum pool size, and then attempts to keep the pool size constant (adding new threads if a thread dies due to an unexpected Exception).

newCachedThreadPool.

A cached thread pool has more flexibility to reap idle threads when the current size of the pool exceeds the demand for processing, and to add new threads when demand increases, but places no bounds on the size of the pool.

newSingleThreadExecutor.

A single-threaded executor creates a single worker thread to process tasks, replacing it if it dies unexpectedly. Tasks are guaranteed to be processed sequentially according to the order imposed by the task queue (FIFO, LIFO, priority order).

newScheduledThreadPool.

A fixed-size thread pool that supports delayed and periodic task execution, similar to Timer.

Runnable vs. Callable

- abstract computational tasks:
 - Runnable
 - Callable
 - Return a value
- Task
- Start
- [eventually] terminates
- Task Lifecycle:
 - created
 - submitted
 - started
 - completed
- Anulare (cancel)
 - Taskurile submise dar nepornite se pot anula
 - Taskurile pornite se pot anulare doar daca raspund la intreruperi
 - Taskurile terminate nu sunt influentate de 'cancel'.

Interfetele Callable si Future

```
public interface Callable<V> {
     V call() throws Exception;
public interface Future<V> {
    boolean cancel(boolean mayInterruptIfRunning);
    boolean isCancelled();
    boolean isDone();
    V get() throws InterruptedException, ExecutionException,
     CancellationException;
    V get(long timeout, TimeUnit unit) throws InterruptedException,
     ExecutionException, CancellationException, TimeoutException;
```

Java

Future

FutureTask -> A cancellable asynchronous computation.

CompletableFuture

apel direct fara executor

```
public class Test {
      public static class AfisareMesaj implements Callable<String>{
            private String msg;
            public AfisareMesaj(String m){
                  msg = m;
            public String call(){
                   String threadName = Thread.currentThread().getName();
                  // System.out.println(msg +" "+ threadName);
                   return msg + " "+threadName;
public static void main(String a[]){
      FutureTask<String> fs = new FutureTask<String>(new AfisareMesaj("TEST"));
     fs.run();
     try {
            System.out.println(fs.get());
      } catch (InterruptedException | ExecutionException e2) {
                  e2.printStackTrace();
```

Suma numere consecutive – afisare rezultate

```
public class MyRunnable implements Runnable {
    private final long countUntil;
    MyRunnable(long countUntil) {
         this.countUntil = countUntil;
    @Override
    public void run() {
         long sum = 0;
         for (long i = 1; i < countUntil; i++) {
             sum += i;
         System.out.println(sum);
//global_variable = sum;
```

```
import java.util.concurrent.ExecutorService;
import java.util.concurrent.Executors;
public class Main {
    private static final int NTHREADS = 10;
    public static void main(String[] args) {
         ExecutorService executor = Executors.newFixedThreadPool(NTHREADS);
         for (int i = 0; i < 500; i++) {
             Runnable worker = new MyRunnable(1000000L + i);
             executor.execute(worker);
         // This will make the executor accept no new threads
         // and finish all existing threads in the queue
         executor.shutdown();
         // Wait until all threads are finish
         executor.awaitTermination();
         System.out.println("Finished all threads");
```

Exemplu: Futures & Callable Suma de numere consecutive – acumulare

```
import java.util.concurrent.Callable;
public class MyCallable implements Callable<Long> {
    @Override
    public Long call() throws Exception {
        long sum = 0;
        for (long i = 0; i \le 100; i++) {
             sum += i;
        return sum;
```

```
import java.util.ArrayList;
import java.util.List;
import java.util.concurrent.Callable;
import java.util.concurrent.ExecutionException;
import java.util.concurrent.ExecutorService;
import java.util.concurrent.Executors;
import java.util.concurrent.Future;
public class CallableFutures {
    private static final int NTHREADS = 10;
    private static final int MAX = 200;
 public static void main(String[] args) {
  ExecutorService executor =
      Executors.newFixedThreadPool(NTHREDS);
  List<Future<Long>> list =
             new ArrayList<Future<Long>>();
  for (int i = 0; i < MAX; i++) {
    Callable<Long> worker = new MyCallable();
    Future<Long> submit =
                   executor.submit(worker);
    list.add(submit);
```

```
long sum = 0;
        System.out.println(list.size());
        // now retrieve the result
        for (Future<Long> future : list) {
            try {
                 sum += future.get();
            } catch (InterruptedException e) {
                 e.printStackTrace();
            } catch (ExecutionException e) {
                 e.printStackTrace();
        System.out.println(sum);
        executor.shutdown();
```

CompletableFuture (from docs...)

- A Future that may be explicitly completed (setting its value and status), and may be used as a CompletionStage, supporting dependent functions and actions that trigger upon its completion.
- When two or more threads attempt to complete or cancel a CompletableFuture, only one of them succeeds.
- CompletableFuture implements interface CompletionStage with the following policies:
 - Actions supplied for dependent completions of non-async methods may be performed by the thread that completes the current CompletableFuture, or by any other caller of a completion method.
 - All async methods without an explicit Executor argument are performed using the ForkJoinPool.commonPool() (unless it does not support a parallelism level of at least two, in which case, a new Thread is created to run each task).
 - To simplify monitoring, debugging, and tracking, all generated asynchronous tasks are instances of the marker interface CompletableFuture.AsynchronousCompletionTask.
 - All CompletionStage methods are implemented independently of other public methods, so the behavior of one method is not impacted by overrides of others in subclasses.

runAsync

supplyAsync

Variants of runAsync() and supplyAsync()

```
static CompletableFuture<Void> runAsync(Runnable runnable)
static CompletableFuture<Void> runAsync(Runnable runnable, Executor executor)
static <U> CompletableFuture<U> supplyAsync(Supplier<U> supplier)
static <U> CompletableFuture<U> supplyAsync(Supplier<U> supplier, Executor executor)
```

ForkJoinPool.commonPool()

thenApply

```
// Create a CompletableFuture
CompletableFuture<String> whatsYourNameFuture = CompletableFuture.supplyAsync(
 () -> { try { TimeUnit.SECONDS.sleep(1); }
       return "Ana";}
);
// Attach a callback to the Future using thenApply()
CompletableFuture<String> greetingFuture = whatsYourNameFuture.thenApply(
     name -> { return "Hello " + name;}
);
// Block and get the result of the future.
System.out.println(greetingFuture.get());
```

.thenApply - takes a Function<T,R> as an argument

.thenApply(.....).thenApply(

thenApply() variants

class CompletableFuture<T>

methods:

<U> CompletableFuture<U> thenApply(Function<? super T,? extends U> fn)

<U> CompletableFuture<U> thenApplyAsync(Function<? super T,? extends U> fn)

<u> CompletableFuture<u> thenApplyAsync(Function<? super T,? extends U> fn, Executor executor)

thenAccept

- thenAccept() takes a Consumer<T> and returns CompletableFuture<Void>.
- It has access to the result of the CompletableFuture on which it is attached.