

SDL - Buffer Overflows



Buffer Overflows Overview

Buffer Overflow: Occurs when data is written into a fixed-length buffer and the size of that data exceeds the capacity of the receiving buffer

- Primary Risks: Corrupt data, crash programs and control execution flow
- Common in native applications (C/C++)
 - Rare, but still possible in managed code (.NET, Java)
- Cause is failing to validate input
- Can occur on stacks and heaps

Review of Application Stack Frames (detailed)



```
int function_B(int a, int b)
      push EBP
   intox, MBP. ESP // local variable
                                                                                                                     0xB FF FF000
                                                                    0xBFFFF000
                                                                                                                   (small addresses)
                                                                   (small addresses)
   v = \cdot b * b;
                                                                                                                                             int y
                                                                                                                                             int x
   return (x + y);
                                                                                                                                          saved EBP
                                                                                 function B()'s stack frame
                                                                                                                                   saved EIP (ret addr in A)
                                                                                                                                                                           function A()'s stack frame
                                                                                                                                    0x01 (B's arg 1 = int a)
int function A(int p, int q) int function A(int p, int q) push EBP-
                                                                                                                                    0x01 (B's arg 2 = int b)
                                                                                                                                             int c
   into; EBP, ESP // local variables
                                                                                                                                                                               local parameters
                                                                                                                                          saved EBP
                                                                                 function A()'s stack frame
            ESP, 44h
                                                                                                                                  saved EIP (ret addr in main)
                                                                                                                                                                   Ą
   c = p * q * function B(p, p);
                                                                                                                                    0x01 (A's arg 1 = int p)
                                                                                                                  Stack grows toward lo memory addresses
                                                                                                                                                                             saved frame pointer
                                                                                                                                    0x02 (A's arg 2 = int q)
   push 1 return c;
                                                                                                                                            int res
                                                                                                                                                                               return address
      push 1
                                                                                                                                          saved EBP
      call function B
                                                                                    main()'s stack frame
                                                                                                                                     saved EIP (ret addr)
                                                                                                                                                                              calling parameters
                                                                                                                                            int argc
                                                                    0xBFFFFF00
                                                                   (big addresses)
int main (int arge, char **argv, char **env)
                                                                                                                                         char **argv
                                                                                                                    03BFFFFF00
                                                                                                                   (big addresses)
                                                                                                                                          char **env
{ push 2
   pnshres;
   call function A
res = function_A(1, 2);
   return res;
                                                                                        Stack base
                                                                                                                                         Stack base
```



Stack-Based Buffer Overflows

Primary Risk: Ability to overwrite control structures

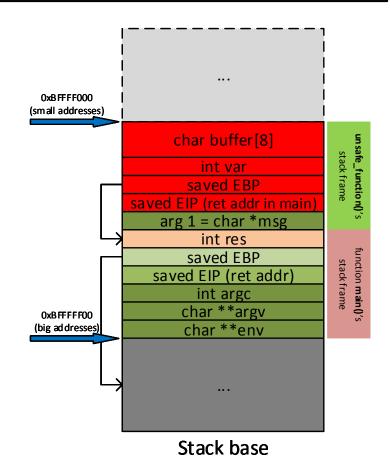
```
/* UNSAFE Function */
                                              SAMPLE INPUTS (STR VALUES):
                                              1. "Kevin"
void UnsafeFunction(char * str)
                                              2. "A" repeated 40 times
    char Buffer[32];
    /* Copy str into Buffer */
    strcpy(Buffer,str);
            Buffer (32 Bytes)
                                                              Address
                                                                          Function
Malicious Payload or Machine Instructions
                                                              of Buffer
                                                                         parameters
```



Stack-Based Buffer Overflows (details)

Primary Risk: Ability to overwrite control structures

```
int unsafe function(char *msg)
                    // local variables
  int var;
  char buffer[8];
  var = 10;
  strcpy(buffer, msg);
  return var;
int main(int argc, char **argv, char **env)
  int res;
  /* Buffer overflow for "strlen(argv[1]) >= 8"
  res = unsafe function(argv[1]);
  return res;
```





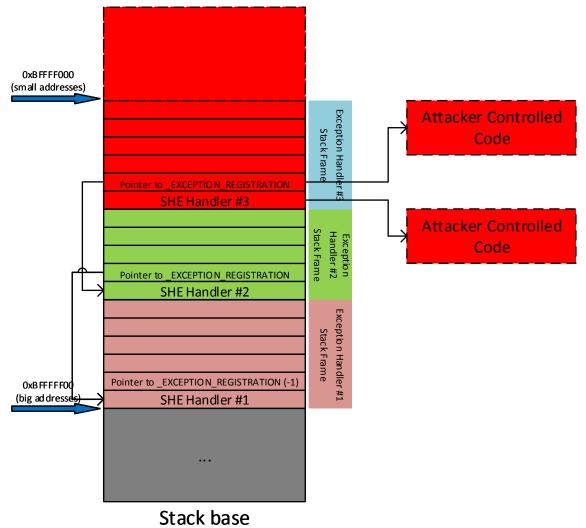
Heap-Based Buffer Overflows

Primary Risk: Ability to write arbitrary 4 byte DWORD anywhere in memory (return address, pointers, etc.)

```
/* UNSAFE Function */
                                                Pseudo-code For Chunk Freeing:
void UnsafeFunction(char * str)
                                                NextChunk =
                                                                AAAA
                                                PreviousChunk =
                                                                    AAAA
   /* Allocate 32 bytes heap space */
                                                           ->BP =
                                                  AAAA
                                                                      AAAA
   char * Buffer = (char *)malloc(32);
                                                PreviousChunk->FP = NextChunk
   /* Copy str into Buffer */
   strcpy(Buffer,str);
           AA(32 times)AA
                                             Data
FP
     BP
                                                              FP
                                                                    BP
                                                                         Data
                                                                 Chunk #1
      Chunk #3
                                      Chunk #2
```

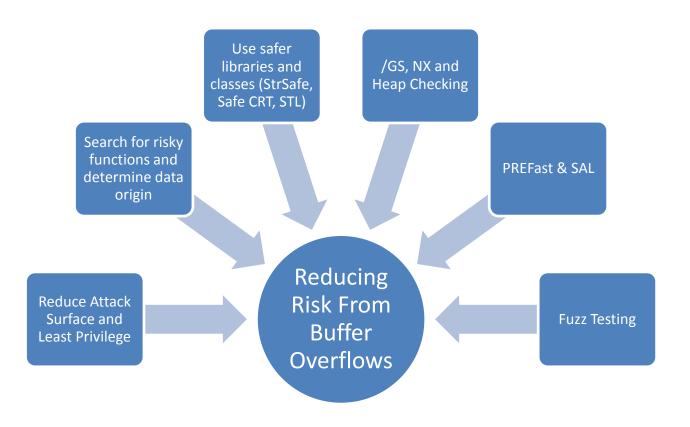


Structured Exception Handling (SEH)



Reducing Exposure to Buffer Overflows with the Microsoft SDL





 Presentation content is available for all of these topics

Platform Protection From Buffer Overflows



- Modern day operating systems and processors have built-in buffer overflow protection
 - Address Space Layout Randomization (ASLR)
 - Data Execution Protection (DEP)
- However none of these are "silver bullets"
 - Denial of Service (DoS) attacks usually not prevented
 - More subtle attacks could still be performed
 - Developers still need to follow security best practices
 - Developers should always apply the Microsoft SDL

Bitdefender

CWE Buffer-Overflow Related

- CWE-119: Improper Restriction of Operations within the Bounds of a Memory Buffer
- CWE-120: Buffer Copy without Checking Size of Input ('Classic Buffer Overflow')
 - Rank 3 in the Top 25
- CWE-121: Stack-based Buffer Overflow
- CWE-122: Heap-based Buffer Overflow
- CWE-124: Buffer Underwrite ('Buffer Underflow')
- CWE-125: Out-of-bounds Read
- CWE-131: Incorrect Calculation of Buffer Size (!)
 - Rank 20 in the Top 25
- CWE-170: Improper Null Termination
- CWE-190: Integer Overflow (!)
 - Rank 24 in the Top 25
- CWE-193: Off-by-one Error
- CWE-805: Buffer Access with Incorrect Length Value
- ...

Real-Life Examples



- First well-known Internet worm: Morris finger worm (1988)
- Common Vulnerabilities and Exposures (https://cve.mitre.org/find/index.html)
 - Searching string "buffer overflow" → "About 639 results" (actually few thousands)
- Vulnerability Notes Database (https://www.kb.cert.org/vuls/)
 - Searching string "buffer overflow" → "About 240 results"
- Examples
 - CVE-2015-0235 GHOST: glibc gethostbyname buffer overflow
 - CVE-2014-0001 Buffer overflow in client/mysql.cc in Oracle MySQL and MariaDB before 5.5.35
 - CVE-2014-0182 Heap-based buffer overflow in the virtio_load function in hw/virtio/virtio.c in QEMU before 1.7.2
 - CVE-2014-0498 Stack-based buffer overflow in Adobe Flash Player before 11.7.700.269
 - CVE-2014-0513 Stack-based buffer overflow in Adobe Illustrator CS6 before 16.0.5
 - CVE-2014-8271 Tianocore UEFI implementation reclaim function vulnerable to buffer overflow
 - CVE-2013-0002 Buffer overflow in the Windows Forms (aka WinForms) component in Microsoft .NET Framework
 - CVE-2005-3267 Integer overflow in Skype client ... leads to a resultant heap-based buffer overflow

- ..

Conclusions



- Buffer Overflow
 - classical, well known, still present ("oldie but goldie")
 - due to
 - the usage of unsafe function and non-validated user input
 - logic (calculation) errors
- Recommendations for Code Developers
 - Do not use unsafe (string) functions
 - Use the right compiler/linker options
 - Check allocation size calculations
 - Do make size checking
 - Take care of automatic type casting and possible integer overflows
 - use size_t (when possible) for allocation size variables
 - take care at casts from signed to unsigned
 - **–**
- Recommendations for Code Reviewers
 - Check for user input and trace it through the application
 - Check for unsafe functions
 - Check for allocation size calculations