

TÁMOP-4.2.2.C-11/1/KONV-2012-0004

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# PETRI NET BASED TRAJECTORY OPTIMIZATION

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- 1. Introduction**
- 2. The CEGAR approach on Petri nets**
- 3. Trajectory optimization using CEGAR**
- 4. Evaluation**
- 5. Conclusions**

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# INTRODUCTION

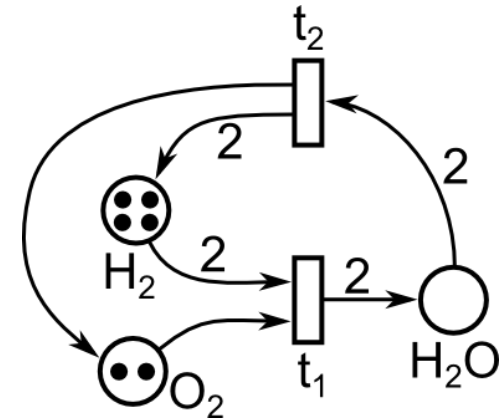
## Petri Nets

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- **Information systems are becoming more complex**
  - Modeling and automatic analysis is important

- **Modeling: Petri Nets**

- Widely used modeling formalism
  - Asynchronous, distributed, parallel, non-deterministic systems
- Behavior: possible states and transitions
- Optimization problems
  - Optimal trajectory from the initial state to a given goal state
  - Reachability analysis



# INTRODUCTION

## Reachability analysis

- **Reachability analysis**

- Checks, if a given state is reachable from the initial state
- $m_1 \in R(PN, m_0) \rightarrow$  „Is  $m_1$  reachable from  $m_0$  in the Petri net  $PN$ ?”
- Drawback: complexity

- **Complexity**

- State space can be large or infinite
- Reachability is decidable, but at least EXPSPACE-hard
- No upper bound is known
- A possible solution is to use abstraction

# INTRODUCTION

## The CEGAR approach

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- **CounterExample Guided Abstraction Refinement**
  - General approach
    - Can handle large or infinite state spaces
  - Works on an abstraction of the original model
    - Less detailed state space
    - Finite, smaller representation
  - Abstraction refinement is required
    - An action in the abstract model may not be realizable in the original model
    - Refine the abstraction using the information from the explored part of the state space
  - H. Wimmel, K. Wolf
    - Applying CEGAR to the Petri Net State Equation (2011)

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# CEGAR APPROACH ON PETRI NETS

## Initial abstraction

- Abstraction of Petri nets: state equation**

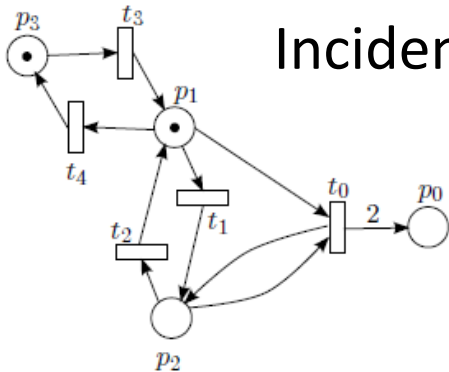
Initial state

$$m_0 + Cx = m_1$$

Target state

Firing count  
of transitions  
(*unknown*)

Incidence matrix



$$\begin{bmatrix} 2 & 0 & 0 & 0 & 0 \\ -1 & -1 & 1 & 1 & -1 \\ 0 & 1 & -1 & 0 & 0 \\ 0 & 0 & 0 & -1 & 1 \end{bmatrix}$$

State  
equationInitial  
abstractionReachability  
problem



# CEGAR APPROACH ON PETRI NETS

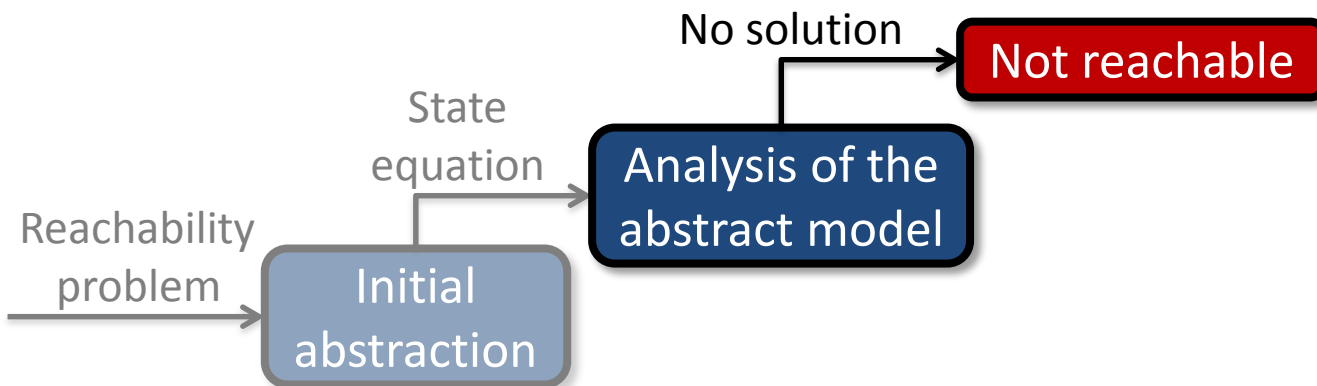
## Analysis of the abstract model

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- Solving the state equation for the firing count of transitions
- Integer Linear Programming problem
- Necessary, but not sufficient criterion for reachability

$$m_0 + C\mathbf{x} = m_1$$

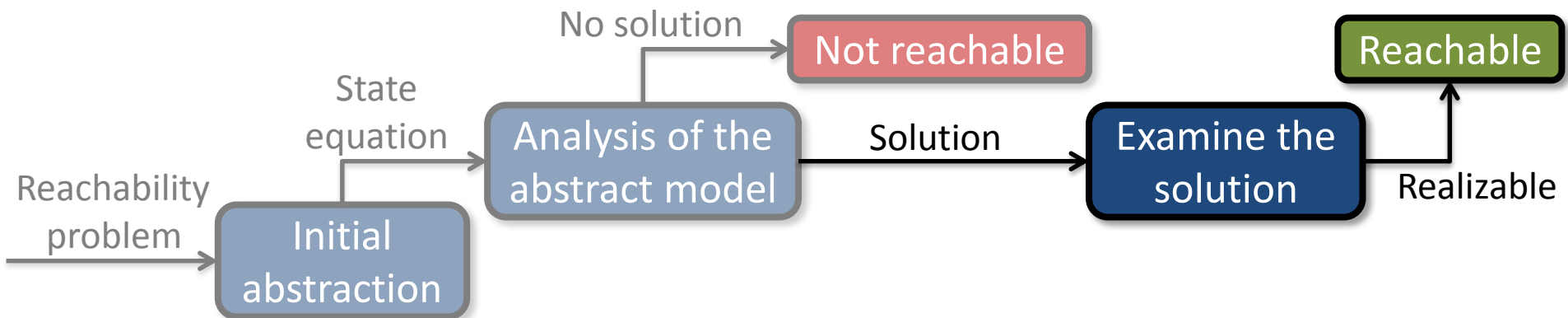
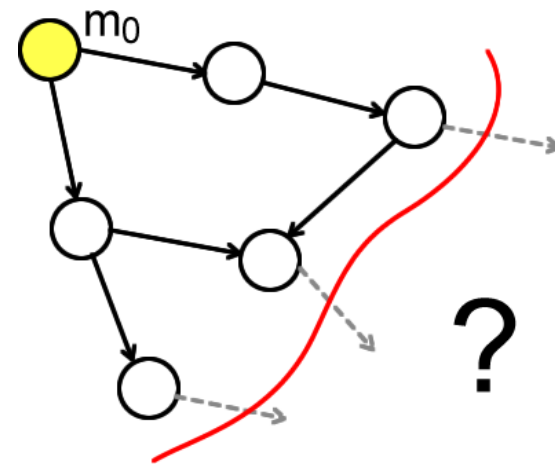
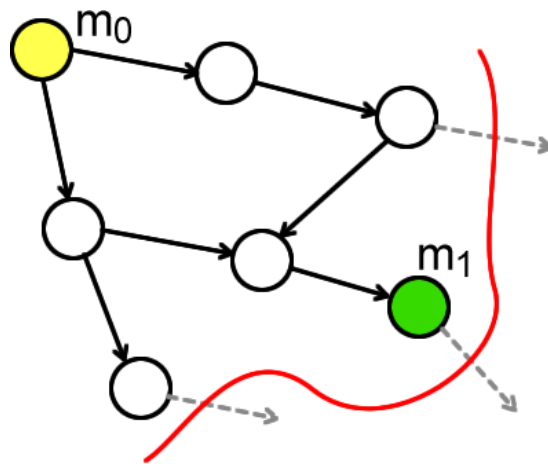


# CEGAR APPROACH ON PETRI NETS

## Examining the solution

- **Bounded exploration of the state space**

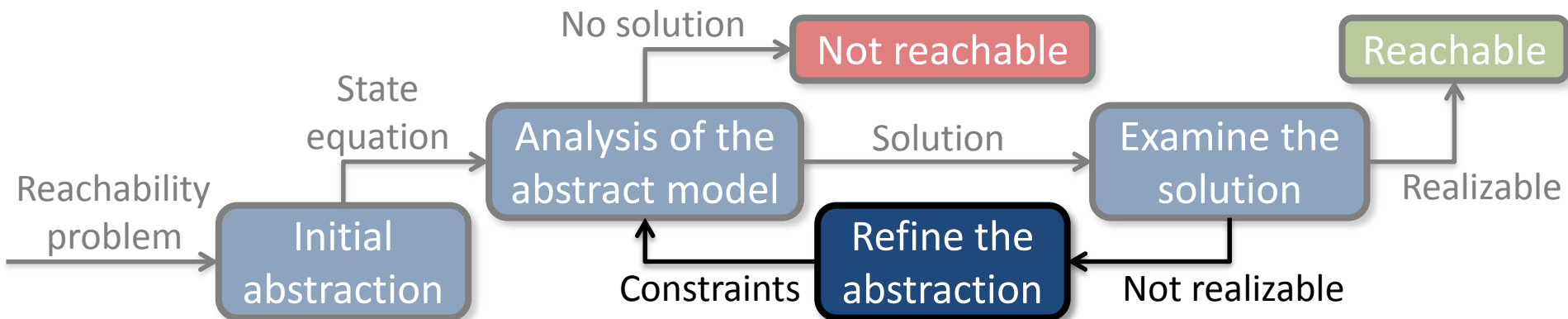
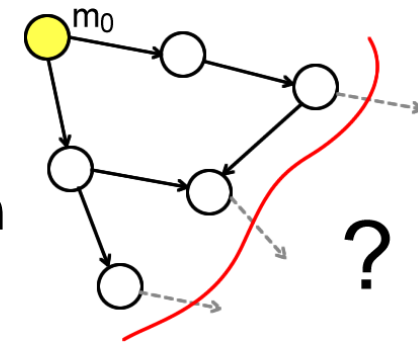
- Try to fire the transitions in some order



# CEGAR APPROACH ON PETRI NETS

## Abstraction refinement

- **Exclude the counterexample without losing any realizable solution**
- **Constraints can be added to the state equation**
  - The state equation may become infeasible
  - A new solution can be obtained
- **Traversing the solution space instead of the state space**



# CEGAR APPROACH ON PETRI NETS

## Solution space

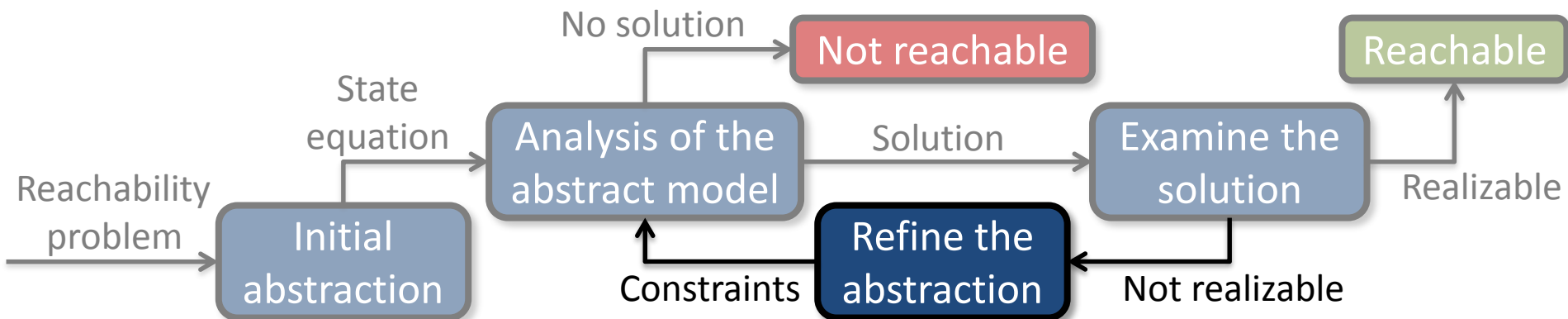
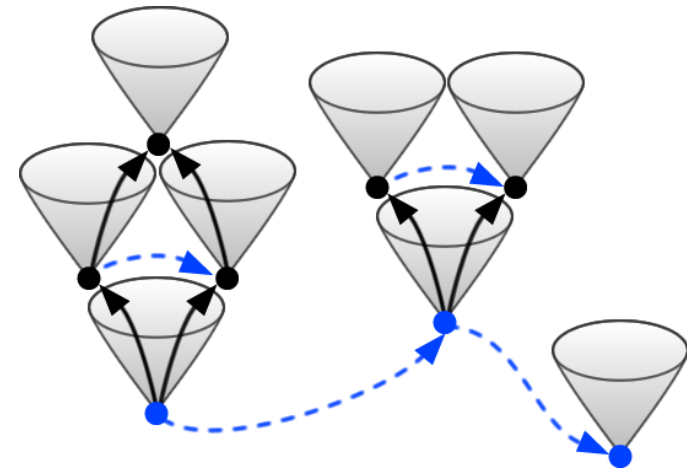
- **Semi-linear space**

- Base solutions
- T-invariants
  - Solutions of the homogenous part  $Cy = 0$
  - Possible cycles in the Petri Net

- **Two types of constraints**

- Jump: switch between base solutions
- Increment: reach non-base solutions

$$m_0 + C\mathbf{x} = m_1$$



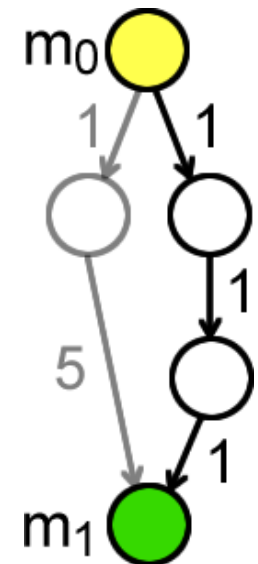
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- **Our previous work**
  - Analyzing the algorithm
    - Correctness
    - Completeness
  - Extending the set of decidable problems
  - New optimizations
- **Current work**
  - Trajectory optimization using CEGAR
    - Assigning costs to transitions
    - New strategy for the solution space traversal

# TRAJECTORY OPTIMIZATION

## Assigning costs to transitions

- **Core of the CEGAR approach: state equation**
  - ILP problem
  - ILP solver minimizes a function over the variables
  - Variables are transitions in our case
- **Original algorithm**
  - Verification purpose → *Is there a solution or not?*
  - Equal cost for each transition → shortest trajectories
- **Our new approach**
  - Optimization purpose → *What is the optimal solution?*
  - Arbitrary cost for transitions
  - ILP solver minimizes using the given cost










- **Traversing the solution space of the state equation**
- **Original algorithm**
  - Verification purpose → *Is there a solution or not?*
  - Fast convergence → DFS
- **Our new approach**
  - Optimization purpose → *What is the optimal solution?*
  - Store the solutions in a sorted queue
  - Continue with the one with the lowest cost



Input: Reachability problem  $m_1 \in R(PN, m_0)$  and cost function  $z$

Output: Trajectory  $\sigma$  or „*Not reachable*”

1.  $C \leftarrow$  incidence matrix of  $PN$  
2.  $Q \leftarrow \text{SolveLP}(m_0, m_1, C, z, \emptyset)$  
3. while  $Q \neq \emptyset$  do
4. |  $x \leftarrow$  solution from  $Q$  minimizing  $z \cdot x$
5. | if  $x$  is realizable then stop and output  $\sigma$  for  $x$  
6. | else 
  7. | | foreach jump and increment constraint  $c'$  do
  8. | | |  $Q \leftarrow \text{SolveLP}(m_0, m_1, C, z, \{\text{constraints of } x\} \cup \{c'\})$
  9. | | end foreach
10. | end else  
11. end while
12. Output „*Not reachable*” 

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- **Implementation**

- PetriDotNet framework
  - Modeling and analysis of Petri nets
  - Supports add-ins

- **Measurements**

- Traveling salesman problem
  - Graph traversal optimization
  - NP-complete



Number of nodes	Runtime (s)
4	0,04
6	0,14
8	0,66
9	0,90
10	1,95
11	9,49
12	24,57
13	1067,00

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- **New approach for the optimal trajectory problem**
  - Translation to the reachability of Petri nets
  - Solving reachability using CEGAR
    - Handle transition costs
    - New strategy for solution space traversal
- **Implementation and evaluation**
- **Possible future direction**
  - Optimization of continuous systems

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# THANK YOU FOR YOUR ATTENTION! *QUESTIONS?*

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*<https://inf.mit.bme.hu/en/research/tools/petridotnet>*

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