

# Exploitation of dynamic symmetries for solving SAT problems

Doctorat de Sorbonne Université

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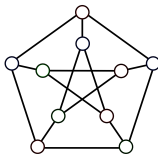
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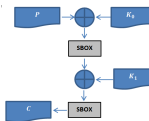


# Motivation

Graph coloring



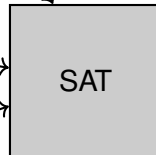
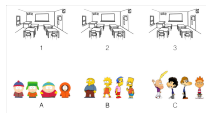
Cryptanalysis



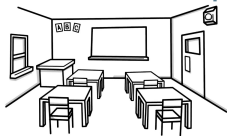
Hardware model checking



Planning



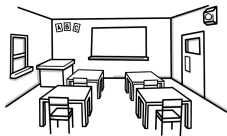
## SAT: an example (1/2)



1



2



3



A



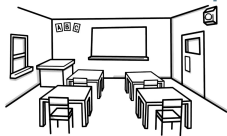
B



C

Is it possible to attribute each group to a unique classroom?

## SAT: an example (1/2)



1  
↑



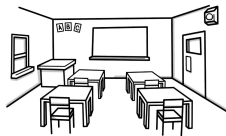
A



2  
↑



B



3  
↑

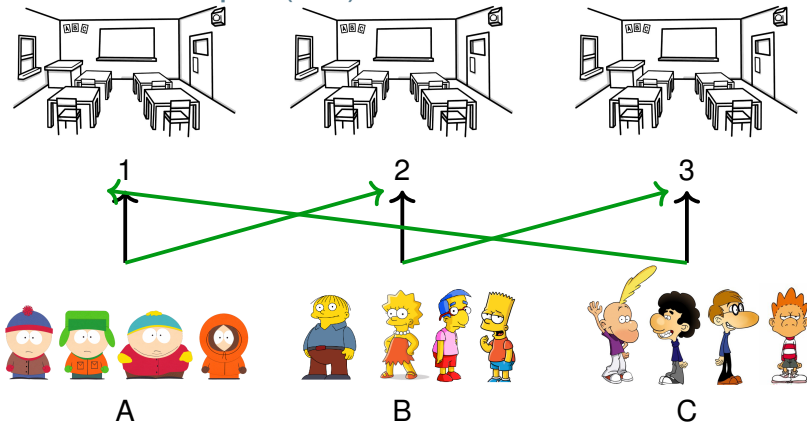


C

Is it possible to attribute each group to a unique classroom?

YES!  $\text{SAT } \alpha = (A, 1), (B, 2), (C, 3)$

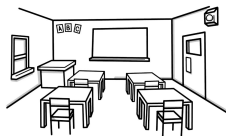
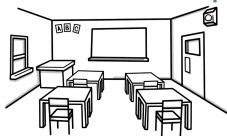
## SAT: an example (1/2)



Is it possible to attribute each group to a unique classroom?

YES!  $\text{SAT } \alpha = (A, 1), (B, 2), (C, 3)$   
Many solutions  $\alpha' = (A, 2), (B, 3), (C, 1)$

# SAT: an example (1/2)



A



B



C

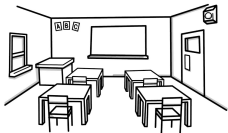
Is it possible to attribute each group to a unique classroom?

YES!  $\text{SAT } \alpha = (A, 1), (B, 2), (C, 3)$

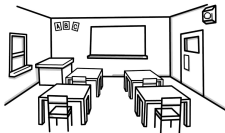
Many solutions  $\alpha' = (A, 2), (B, 3), (C, 1)$

$\alpha'' = \dots$

## SAT: an example (2/2)



1



2



A



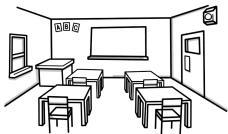
B



C

Is it possible to attribute each group to a unique classroom?

## SAT: an example (2/2)



1



2



A



B



C

Is it possible to attribute each group to a unique classroom?

No! UNSAT



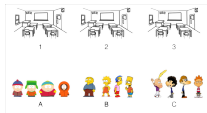
# Encoding the problem

$$\begin{array}{ccc} \overbrace{(A, 1)}^{x_1} & \overbrace{(A, 2)}^{x_2} & \overbrace{(A, 3)}^{x_3} \\ (B, 1) & (B, 2) & (B, 3) \\ (C, 1) & (C, 2) & (C, 3) \end{array}$$

$$(x_1 \vee x_2 \vee x_3) \wedge$$

$$(x_4 \vee x_5 \vee x_6) \wedge$$

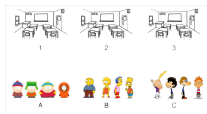
$$(x_7 \vee x_8 \vee x_9) \wedge$$



# Encoding the problem

$$\begin{array}{c} \overbrace{(A, 1)}^{x_1} \overbrace{(A, 2)}^{x_2} \overbrace{(A, 3)}^{x_3} \\ (B, 1)(B, 2)(B, 3) \\ (C, 1)(C, 2)(C, 3) \end{array}$$

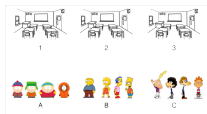
$$\begin{array}{l} \neg(A, 1) \neg(B, 1) \\ \neg(A, 1) \neg(C, 1) \\ \neg(B, 1) \neg(C, 1) \end{array}$$



$$\begin{array}{l} (x_1 \vee x_2 \vee x_3) \wedge \\ (x_4 \vee x_5 \vee x_6) \wedge \\ (x_7 \vee x_8 \vee x_9) \wedge \end{array}$$

$$\begin{array}{l} (\neg x_1 \vee \neg x_4) \wedge \\ (\neg x_1 \vee \neg x_7) \wedge \\ (\neg x_4 \vee \neg x_7) \wedge \end{array}$$

# Encoding the problem



$$\begin{matrix} x_1 & x_2 & x_3 \\ \overbrace{(A, 1)} & \overbrace{(A, 2)} & \overbrace{(A, 3)} \\ \overbrace{(B, 1)} & \overbrace{(B, 2)} & \overbrace{(B, 3)} \\ \overbrace{(C, 1)} & \overbrace{(C, 2)} & \overbrace{(C, 3)} \end{matrix}$$

$$\neg(A, 1) \neg(B, 1)$$

$$\neg(A, 1) \neg(C, 1)$$

$$\neg(B, 1) \neg(C, 1)$$

$$\neg(A, 2) \neg(B, 2)$$

$$\neg(A, 2) \neg(C, 2)$$

$$\neg(B, 2) \neg(C, 2)$$

$$\neg(A, 3) \neg(B, 3)$$

$$\neg(A, 3) \neg(C, 3)$$

$$\neg(B, 3) \neg(C, 3)$$

$$(x_1 \vee x_2 \vee x_3) \wedge$$

$$(x_4 \vee x_5 \vee x_6) \wedge$$

$$(x_7 \vee x_8 \vee x_9) \wedge$$

$$(\neg x_1 \vee \neg x_4) \wedge$$

$$(\neg x_1 \vee \neg x_7) \wedge$$

$$(\neg x_4 \vee \neg x_7) \wedge$$

$$(\neg x_2 \vee \neg x_5) \wedge$$

$$(\neg x_2 \vee \neg x_8) \wedge$$

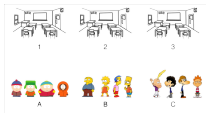
$$(\neg x_5 \vee \neg x_8) \wedge$$

$$(\neg x_3 \vee \neg x_6) \wedge$$

$$(\neg x_3 \vee \neg x_9) \wedge$$

$$(\neg x_6 \vee \neg x_9)$$

# Encoding the problem



$$\begin{matrix} x_1 & x_2 & x_3 \\ \overbrace{(A, 1)} & \overbrace{(A, 2)} & \overbrace{(A, 3)} \\ \overbrace{(B, 1)} & \overbrace{(B, 2)} & \overbrace{(B, 3)} \\ \overbrace{(C, 1)} & \overbrace{(C, 2)} & \overbrace{(C, 3)} \end{matrix}$$

$$\neg(A, 1) \neg(B, 1)$$

$$\neg(A, 1) \neg(C, 1)$$

$$\neg(B, 1) \neg(C, 1)$$

$$\neg(A, 2) \neg(B, 2)$$

$$\neg(A, 2) \neg(C, 2)$$

$$\neg(B, 2) \neg(C, 2)$$

$$\neg(A, 3) \neg(B, 3)$$

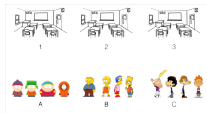
$$\neg(A, 3) \neg(C, 3)$$

$$\neg(B, 3) \neg(C, 3)$$

Clause

$$\begin{aligned} & (x_1 \vee x_2 \vee x_3) \wedge \\ & (x_4 \vee x_5 \vee x_6) \wedge \\ & (x_7 \vee x_8 \vee x_9) \wedge \\ & (\neg x_1 \vee \neg x_4) \wedge \\ & (\neg x_1 \vee \neg x_7) \wedge \\ & (\neg x_4 \vee \neg x_7) \wedge \\ & (\neg x_2 \vee \neg x_5) \wedge \\ & (\neg x_2 \vee \neg x_8) \wedge \\ & (\neg x_5 \vee \neg x_8) \wedge \\ & (\neg x_3 \vee \neg x_6) \wedge \\ & (\neg x_3 \vee \neg x_9) \wedge \\ & (\neg x_6 \vee \neg x_9) \end{aligned}$$

# Encoding the problem



$$\begin{matrix} x_1 & x_2 & x_3 \\ \overbrace{(A, 1)} & \overbrace{(A, 2)} & \overbrace{(A, 3)} \\ \overbrace{(B, 1)} & \overbrace{(B, 2)} & \overbrace{(B, 3)} \\ \overbrace{(C, 1)} & \overbrace{(C, 2)} & \overbrace{(C, 3)} \end{matrix}$$

$$\neg(A, 1) \neg(B, 1)$$

$$\neg(A, 1) \neg(C, 1)$$

$$\neg(B, 1) \neg(C, 1)$$

$$\neg(A, 2) \neg(B, 2)$$

$$\neg(A, 2) \neg(C, 2)$$

$$\neg(B, 2) \neg(C, 2)$$

$$\neg(A, 3) \neg(B, 3)$$

$$\neg(A, 3) \neg(C, 3)$$

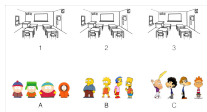
$$\neg(B, 3) \neg(C, 3)$$

Clause

$$\begin{aligned} & (x_1 \vee x_2 \vee x_3) \wedge \\ & (x_4 \vee x_5 \vee x_6) \wedge \\ & (x_7 \vee x_8 \vee x_9) \wedge \\ & (\neg x_1 \vee \neg x_4) \wedge \\ & (\neg x_1 \vee \neg x_7) \wedge \\ & (\neg x_4 \vee \neg x_7) \wedge \\ & (\neg x_2 \vee \neg x_5) \wedge \\ & (\neg x_2 \vee \neg x_8) \wedge \\ & (\neg x_5 \vee \neg x_8) \wedge \\ & (\neg x_3 \vee \neg x_6) \wedge \\ & (\neg x_3 \vee \neg x_9) \wedge \\ & (\neg x_6 \vee \neg x_9) \end{aligned}$$

Conjunctive Normal Form (CNF)

# Encoding the problem



$$\begin{matrix} x_1 & x_2 & x_3 \\ \overbrace{(A, 1)} & \overbrace{(A, 2)} & \overbrace{(A, 3)} \\ \overbrace{(B, 1)} & \overbrace{(B, 2)} & \overbrace{(B, 3)} \\ \overbrace{(C, 1)} & \overbrace{(C, 2)} & \overbrace{(C, 3)} \end{matrix}$$

$$\neg(A, 1) \neg(B, 1)$$

$$\neg(A, 1) \neg(C, 1)$$

$$\neg(B, 1) \neg(C, 1)$$

$$\neg(A, 2) \neg(B, 2)$$

$$\neg(A, 2) \neg(C, 2)$$

$$\neg(B, 2) \neg(C, 2)$$

$$\neg(A, 3) \neg(B, 3)$$

$$\neg(A, 3) \neg(C, 3)$$

$$\neg(B, 3) \neg(C, 3)$$

Clause represented as a set:

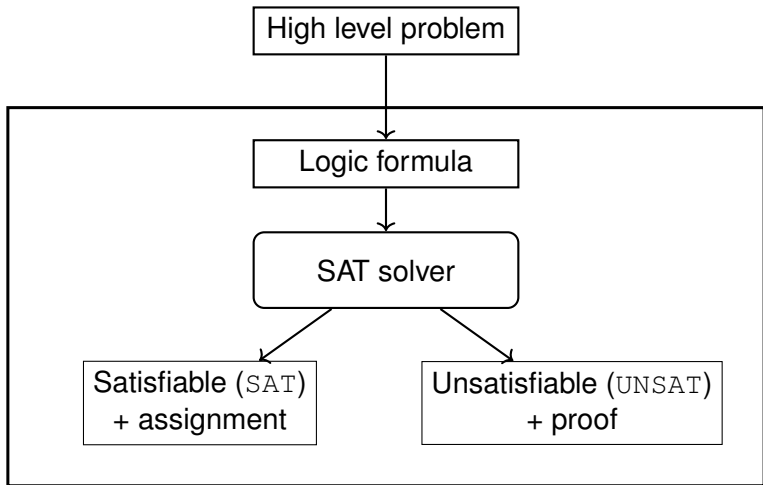
$$(x_1 \vee x_2 \vee x_3) \rightarrow \{x_1, x_2, x_3\}$$

Clause

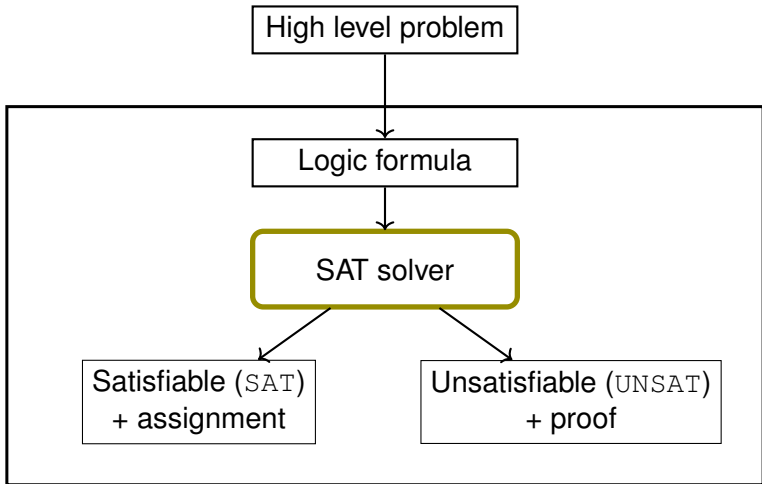
$$\begin{aligned} & (x_1 \vee x_2 \vee x_3) \wedge \\ & (x_4 \vee x_5 \vee x_6) \wedge \\ & (x_7 \vee x_8 \vee x_9) \wedge \\ & (\neg x_1 \vee \neg x_4) \wedge \\ & (\neg x_1 \vee \neg x_7) \wedge \\ & (\neg x_4 \vee \neg x_7) \wedge \\ & (\neg x_2 \vee \neg x_5) \wedge \\ & (\neg x_2 \vee \neg x_8) \wedge \\ & (\neg x_5 \vee \neg x_8) \wedge \\ & (\neg x_3 \vee \neg x_6) \wedge \\ & (\neg x_3 \vee \neg x_9) \wedge \\ & (\neg x_6 \vee \neg x_9) \end{aligned}$$

Conjunctive Normal Form (CNF)

# SAT design



# SAT design





# SAT Solving

Solving SAT formula is known to be **NP-complete** [Coo71]

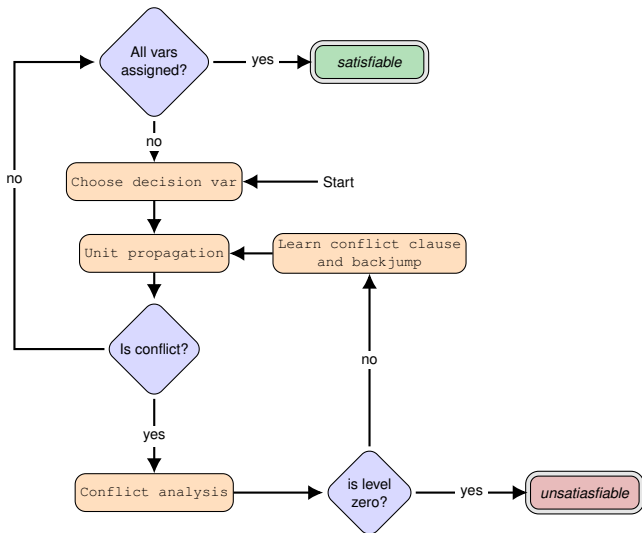
Good performance in practice:

- Handle large problem (million variables and clauses)
- International SAT competition each year on academic and industrial problems

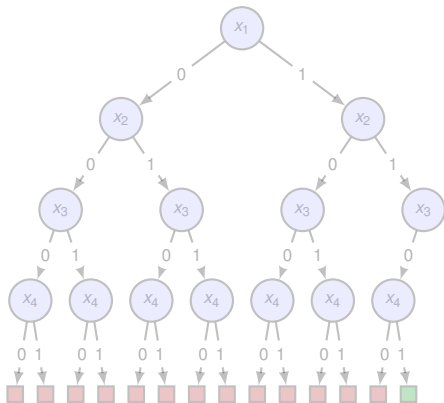
Enumerative algorithms:

- Davis, Putnam, Logemann, and Loveland (DPLL) [DLL62]
  - Boolean Constraint Propagation (BCP)
- Conflict Driven Clause Learning (CDCL) [MSS99]
  - Derived from DPLL
  - Clause learning

# CDCL in detail



# CDCL in action



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

$$\omega_5 = \{x_2, x_4\}$$

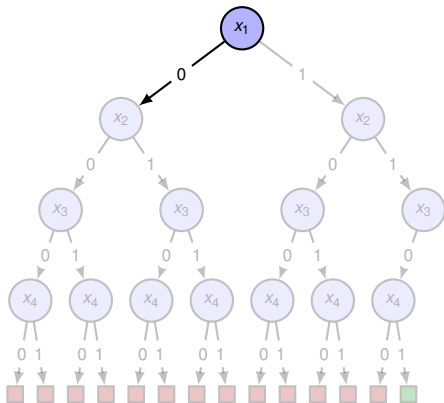
$$\omega_6 = \{x_3, x_4\}$$

$$\alpha = \{\}$$

# CDCL in action

Choose decision var

Unit Propagation



$$\omega_1 = \{\mathbf{x}_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{\mathbf{x}_1, \neg x_4\}$$

$$\omega_3 = \{\mathbf{x}_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

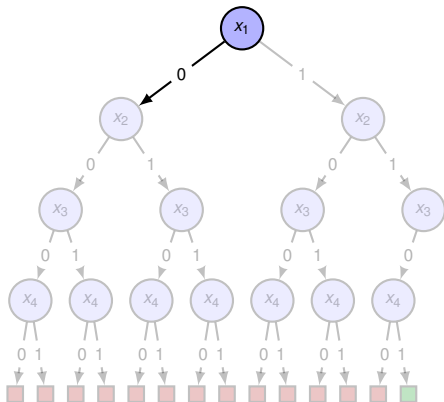
$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

$$\alpha = \{\neg x_1\}$$

# CDCL in action

Conflict Analysis



$$\omega_1 = \{\mathbf{x}_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{\mathbf{x}_1, \neg \mathbf{x}_4\}$$

$$\omega_3 = \{\mathbf{x}_1, \mathbf{x}_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

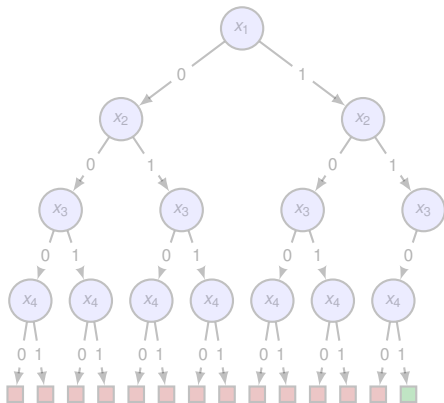
$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

$$\alpha = \{\neg x_1\}$$

# CDCL in action

Learn conflict clause  
and backjump



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

$$\omega_5 = \{x_2, x_4\}$$

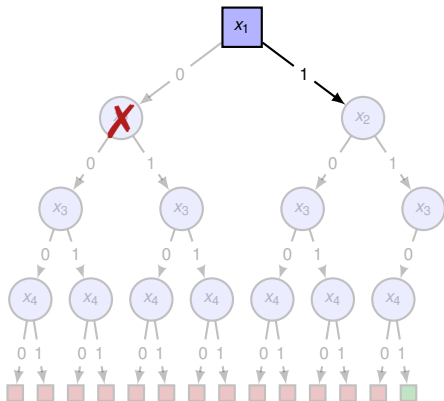
$$\omega_6 = \{x_3, x_4\}$$

$$\omega_7 = \{x_1\}$$

$$\alpha = \{\}$$

# CDCL in action

Unit Propagation



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

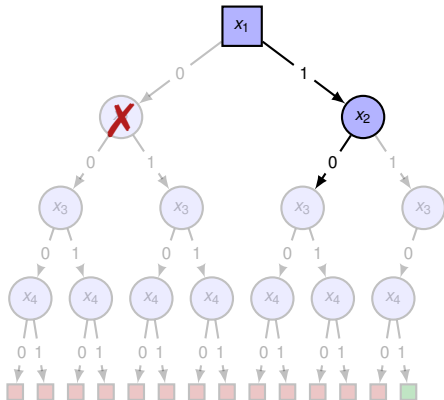
$$\omega_7 = \{x_1\}$$

$$\alpha = \{x_1\}$$

# CDCL in action

Choose decision var

Unit Propagation



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

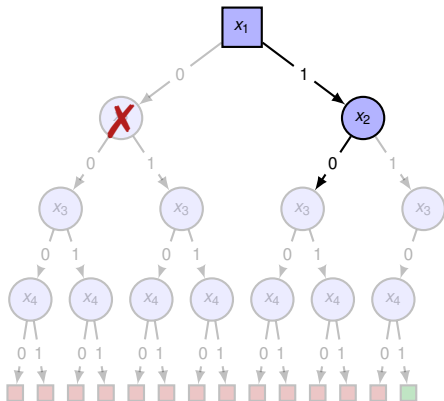
$$\omega_7 = \{x_1\}$$

$$\alpha = \{x_1, \neg x_2\}$$



# CDCL in action

Conflict Analysis



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

$$\omega_5 = \{x_2, x_4\}$$

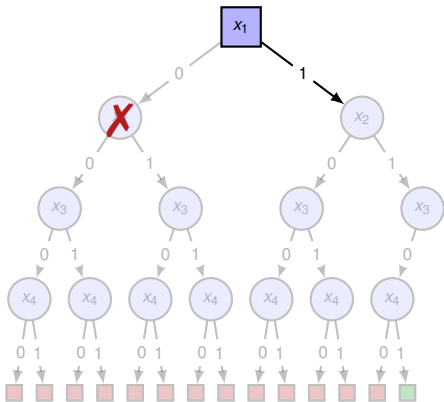
$$\omega_6 = \{x_3, x_4\}$$

$$\omega_7 = \{x_1\}$$

$$\alpha = \{x_1, \neg x_2\}$$

## CDCL in action

Learn conflict clause  
and backjump



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

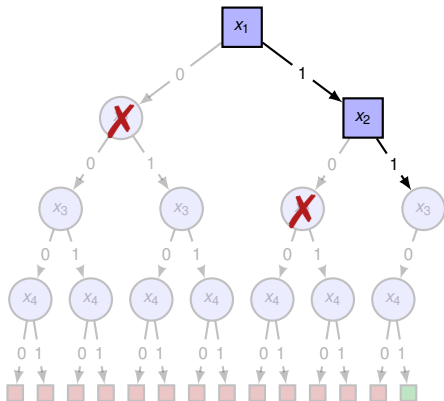
$$\omega_7 = \{x_1\}$$

$$\omega_8 = \{x_2\}$$

$$\alpha = \{x_1\}$$

# CDCL in action

Unit Propagation



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

$$\omega_7 = \{x_1\}$$

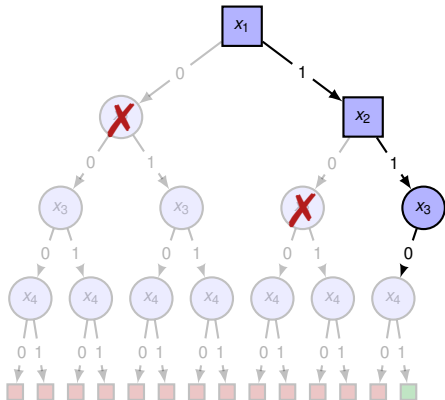
$$\omega_8 = \{x_2\}$$

$$\alpha = \{x_1, x_2\}$$

# CDCL in action

Choose decision var

Unit Propagation



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

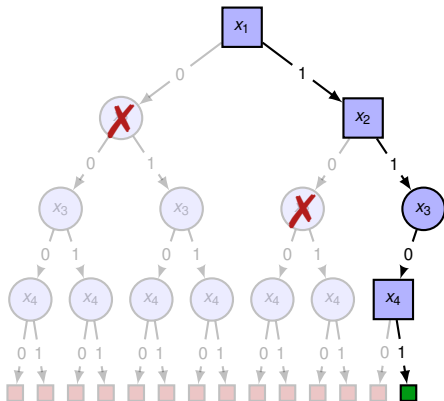
$$\omega_7 = \{x_1\}$$

$$\omega_8 = \{x_2\}$$

$$\alpha = \{x_1, x_2, \neg x_3\}$$

# CDCL in action

Unit Propagation



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

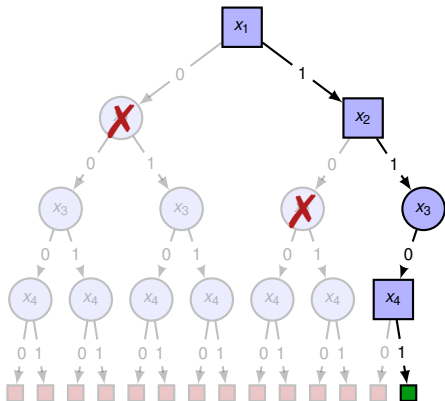
$$\omega_7 = \{x_1\}$$

$$\omega_8 = \{x_2\}$$

$$\alpha = \{x_1, x_2, \neg x_3, x_4\}$$

# CDCL in action

satisfiable



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

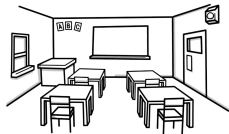
$$\omega_7 = \{x_1\}$$

$$\omega_8 = \{x_2\}$$

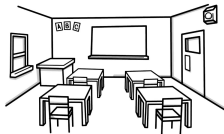
$$\alpha = \{x_1, x_2, \neg x_3, x_4\}$$

# SAT and symmetries

# Presence of symmetries



1



2



A



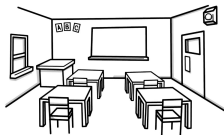
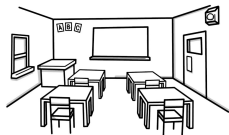
B



C



# Presence of symmetries



A

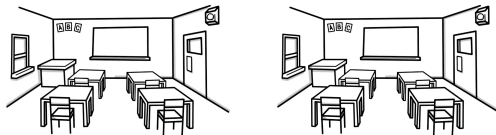


B



C

# Presence of symmetries



1 ← → 2



A

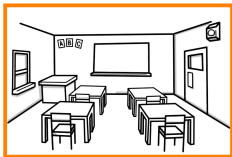


B

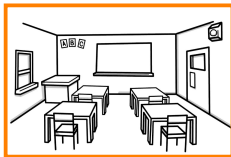


C

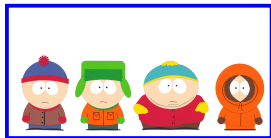
# Presence of symmetries



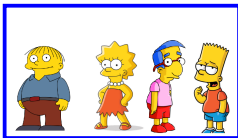
1



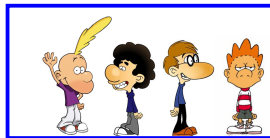
2



A

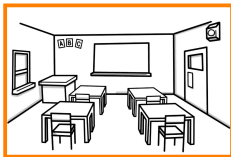


B

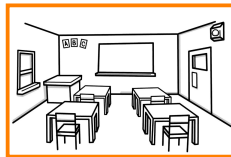


C

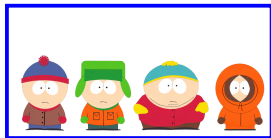
# Presence of symmetries



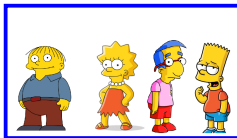
$\neq 2$



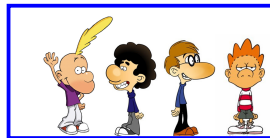
$\neq 1$



A



B C



$\in B$

## Syntactic symmetry

A symmetry (permutation)  $g$  is a bijective function (on variables) that leaves the formula  $\varphi$  invariant

## Syntactic symmetry

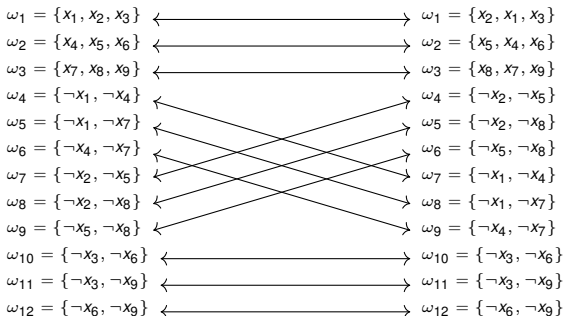
A symmetry (permutation)  $g$  is a bijective function (on variables) that leaves the formula  $\varphi$  invariant

$$g = \left( \begin{array}{cccccccccc} x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 & x_8 & x_9 \\ x_2 & x_1 & x_3 & x_5 & x_4 & x_6 & x_8 & x_7 & x_9 \end{array} \right) \rightarrow (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$

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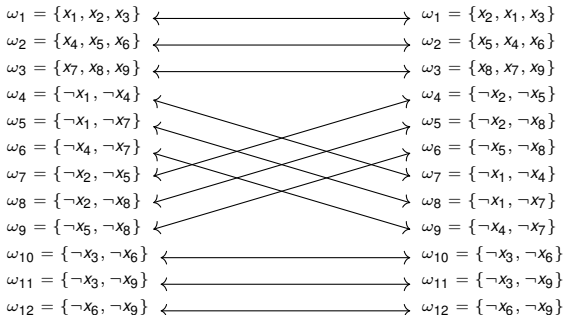
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## Equi-satisfiability

$$\alpha \models \varphi \Leftrightarrow g.\alpha \models \varphi$$



# Computing symmetries of a SAT problem

*CNF formula*

$$\begin{aligned} & (x_1 \vee x_2 \vee x_3) \wedge (x_4 \vee x_5 \vee x_6) \wedge (x_7 \vee x_8 \vee x_9) \\ & \wedge (\neg x_1 \vee \neg x_4) \wedge (\neg x_1 \vee \neg x_7) \wedge (\neg x_4 \vee \neg x_7) \\ & \wedge (\neg x_2 \vee \neg x_5) \wedge (\neg x_2 \vee \neg x_8) \wedge (\neg x_5 \vee \neg x_8) \\ & \wedge (\neg x_3 \vee \neg x_6) \wedge (\neg x_3 \vee \neg x_9) \wedge (\neg x_6 \vee \neg x_9) \end{aligned}$$

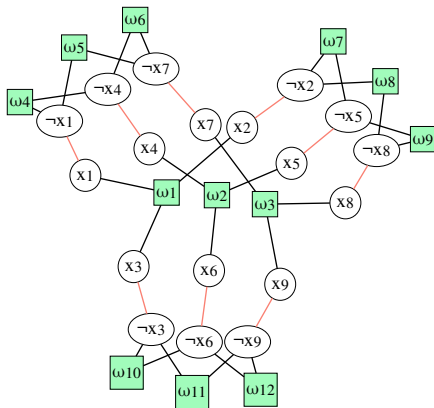
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colored graph



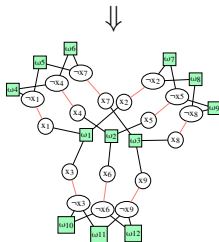
# Computing symmetries of a SAT problem

CNF formula

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colored graph



graph automorphism



(bliss, saucy, ...)

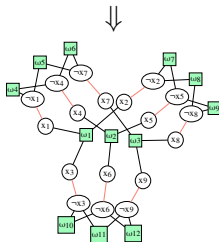
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colored graph



⇓  
graph automorphism

⇓  
(bliss, saucy, ...)

⇓  
set of symmetries

⇓

$$\begin{aligned} g_1 &= (x_2 \ x_3)(x_5 \ x_6)(x_8 \ x_9) \\ g_2 &= (x_4 \ x_7)(x_5 \ x_8)(x_6 \ x_9) \\ g_3 &= (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8) \\ g_4 &= (x_1 \ x_4)(x_2 \ x_5)(x_3 \ x_6) \end{aligned}$$

The set of symmetries of a formula is a group noted  $G$

Exploitation of symmetries

Static symmetry breaking

# Orbit

Orbit of an assignment  $\alpha$  for a group  $G$ :

$$G.\alpha = \{g.\alpha \mid g \in G\}$$

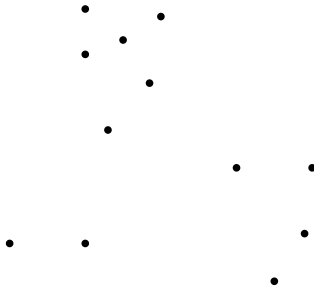
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Example:

- full assignment



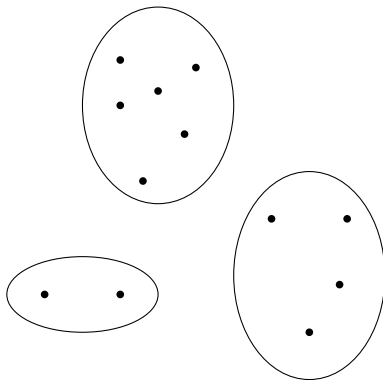
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Example:

- full assignment
- orbit





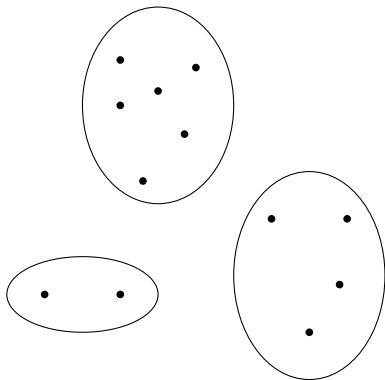
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Equivalence relation with respect to SAT:

- Either  $G.\alpha$  contains no solution
- Or all elements of  $G.\alpha$  are solutions

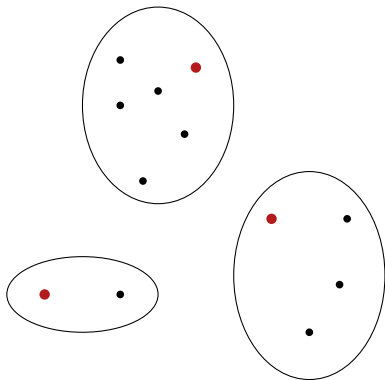
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- orbit
- representative



Equivalence relation with respect to SAT:

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## Compare assignment

Define an ordering relation to compare assignment ( $\prec$ )

- Total ordering on variables
- Minimum value:  $F < T$  or  $T < F$

Allow only minimal value (lex-leader)

Forbids other assignments in each orbit

→ Add Symmetry breaking predicates (SBP)

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Example:

$$x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8; F < T$$

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$

	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$x_7$	$x_8$
$g.\alpha$	F	T	F	F	F	F	F	F
$\alpha$	T	F	F	F	F	F	F	F

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	$\prec$							
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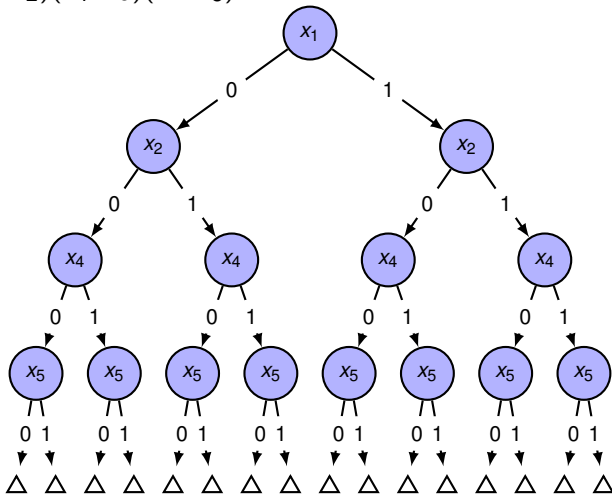
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	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$x_7$	$x_8$
$g.\alpha$	F	T	F	F	F	F	F	F
	$\prec$							
$\alpha$	T	F	F	F	F	F	F	F

Generate SBP  $\omega = \{\neg x_1, x_2\}$

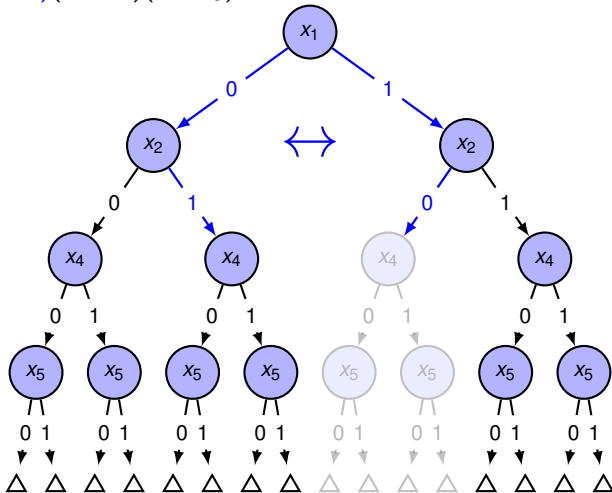
# Using symmetries to prune the search space

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$



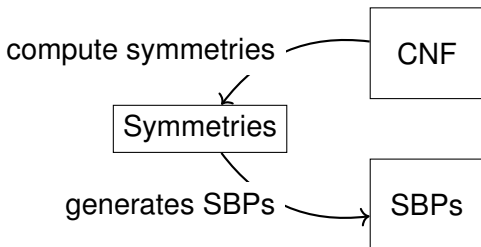
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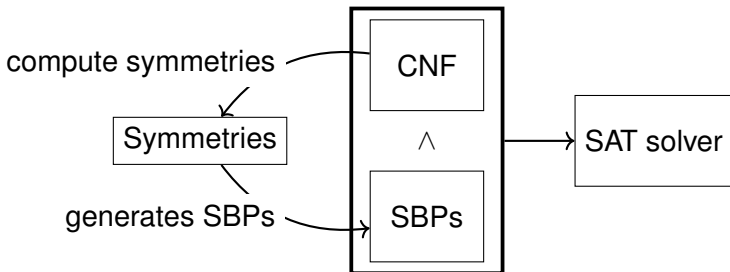
# Static symmetry breaking



State-of-the-art approaches:

- Shatter [ASM06]
- BreakID [DBBD16]
- ...

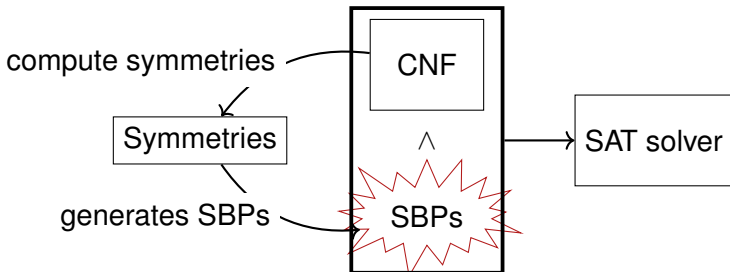
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# Static symmetry breaking



State-of-the-art approaches:

- Shatter [ASM06]
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The solver can "explode" instead of being helped

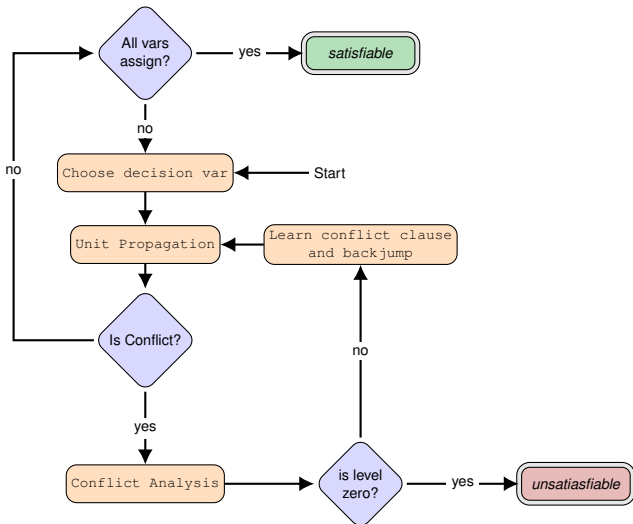
First contribution

CDCL[sym] Introducing Effective Symmetry  
Breaking in SAT Solving

TACAS'18 [MBCK18]

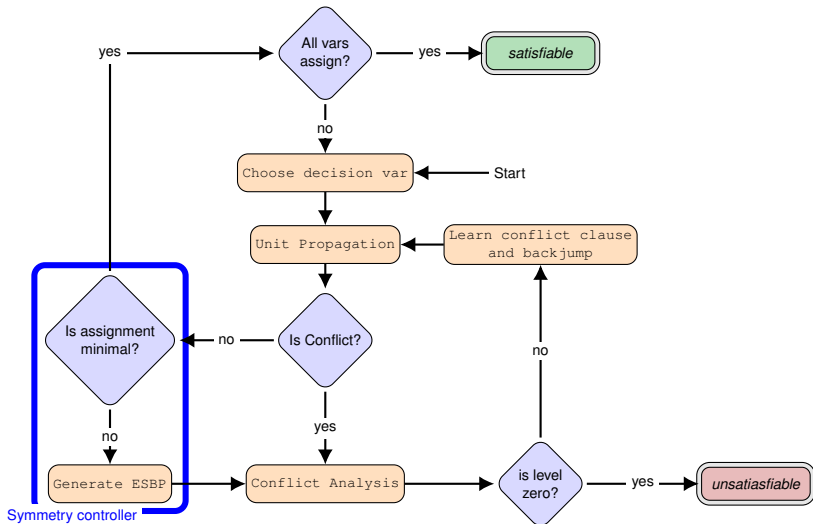
# CDCL[Sym]

Compute and inject SBP **opportunistically**, during the solving



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Compute and inject SBP **opportunistically**, during the solving



# Symmetry status

- reducer:  $g.\alpha \prec \alpha$
- inactive:  $\alpha \prec g.\alpha$
- active: *not enough information*

## Efficient implementation of symmetry status

Keep track the smallest unassigned variable  $x$  :

- ①  $\alpha(g.x) \leq \alpha(x)$ , then  $g$  is reducer  $\Rightarrow$  Effective SBP (ESBP)
- ②  $\alpha(x) \leq \alpha(g.x)$ , then  $g$  is inactive  $\Rightarrow g$  cannot reduce  $\alpha$
- ③  $\alpha(g.x)$  or  $\alpha(x)$  is unassigned then  $g$  is active

Update whenever variables are assigned / unassigned

## Example

$$x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8; \textcolor{red}{F} < \textcolor{green}{T}$$

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$

	$\downarrow$							
	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$x_7$	$x_8$
$g.\alpha$	U	U	U	U	U	U	U	U
$\alpha$	U	U	U	U	U	U	U	U

status of permutation  $g$ : active



## Example

$$x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8; \textcolor{red}{F} < \textcolor{green}{T}$$

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$

	$\downarrow$	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$x_7$	$x_8$
$g.\alpha$		$\textcolor{red}{F}$	U	U	U	U	U	U	U
$\alpha$		U	$\textcolor{red}{F}$	U	U	U	U	U	U

status of permutation  $g$ : active

## Example

$$x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8; \textcolor{red}{F} < \textcolor{green}{T}$$

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	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$x_7$	$x_8$
$g.\alpha$	$\textcolor{red}{F}$	U	U	U	$\textcolor{red}{F}$	U	U	U
$\alpha$	U	$\textcolor{red}{F}$	U	$\textcolor{red}{F}$	U	U	U	U

status of permutation  $g$ : active

# Example

$$x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8; \text{F} < \text{T}$$

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	↓							
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$g.\alpha$	F	T	U	U	F	U	U	U
				↗				
$\alpha$	T	F	U	F	U	U	U	U

status of permutation g: reducer

# Example

$$x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8; \text{F} < \text{T}$$

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$

	↓							
	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$x_7$	$x_8$
$g.\alpha$	F	T	U	U	F	U	U	U
	↖							
$\alpha$	T	F	U	F	U	U	U	U

status of permutation g: reducer

Generate ESBP  $\omega = \{\neg x_1, x_2\}$

# CDCL[Sym] implementation

- C++ Implementation
  - Packaged as a library **cosy**<sup>1</sup> (Controller of Symmetry)
  - Lightweight
  - Fast update
  - Low memory consumption
  - Follows symmetry status
- 
- Works with any enumerative SAT solver
  - Can be integrated easily
- e.g. +3% LOC on MiniSAT.

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<sup>1</sup><https://github.com/lip6/cosy>

# Experiments

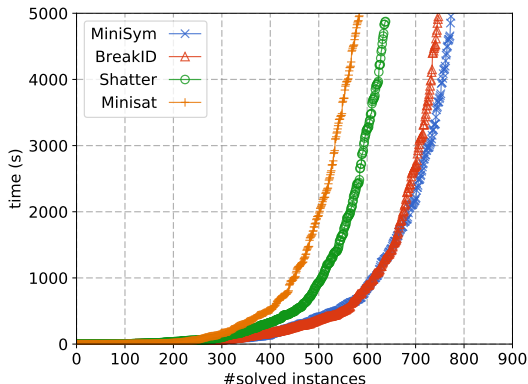
## Benchmark:

- from SAT contests 2012 – 2017
- filter: `bliss` finds symmetries in 1000 seconds
- 36 % of instances, 1 350/3 700

## Setup:

- four tools
  - MiniSat (no symmetry, baseline)
  - MiniSat + BreakID (SOTA SAT solver using symmetries)
  - MiniSat + Shatter (SOTA SAT solver using symmetries)
  - **MiniSym** = MiniSat + CDCL[Sym] (our approach)
- 5000 seconds timeout, 8GB memory
- includes time to compute symmetries (except for MiniSat)

# Experimental results



Solver	PAR-2	ALL	SAT	UNSAT
MiniSAT	2243h	586	325	261
Shatter	2088h	640	316	324
BreakID	1790h	749	334	415
MiniSym	1735h	775	336	439

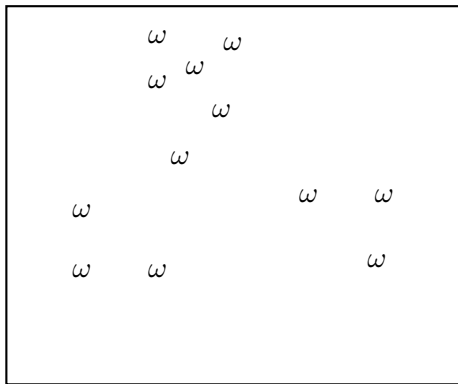
Exploitation of symmetries

Dynamic symmetry breaking



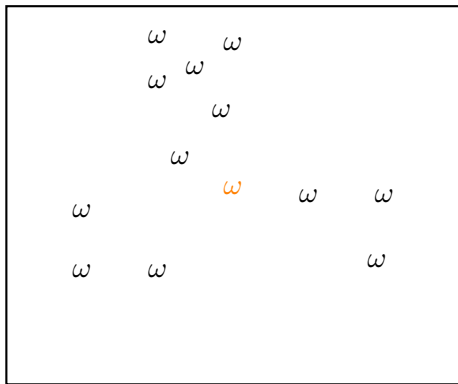
# Learn symmetrical clauses

□ formula  
ω clause



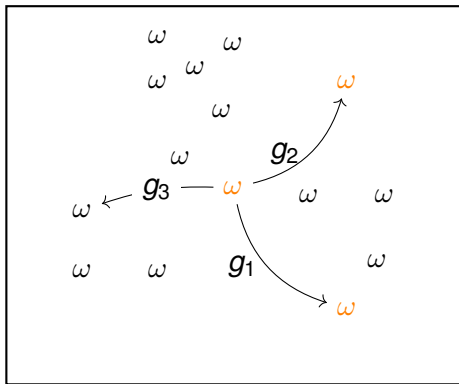
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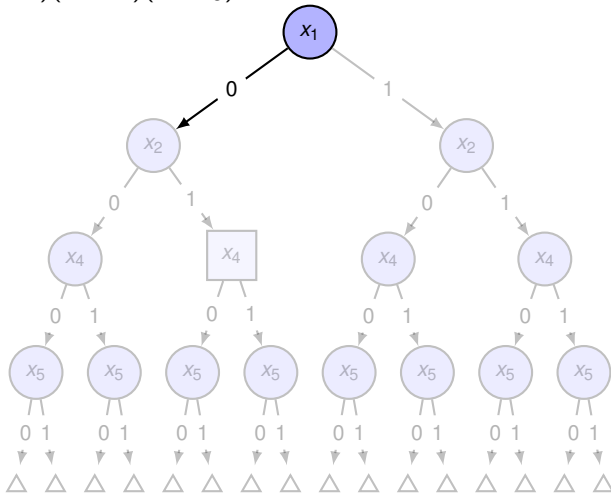
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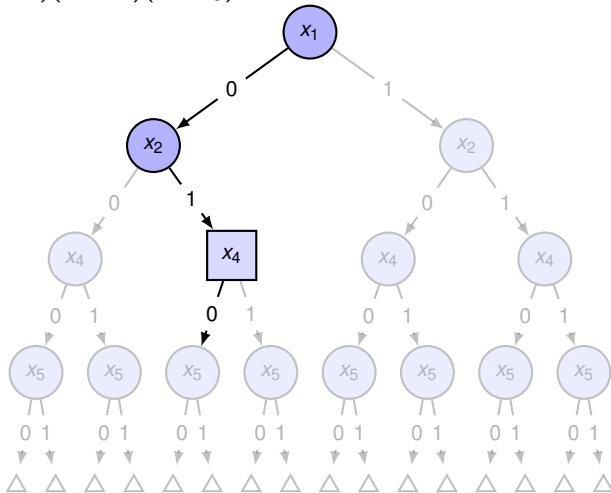
# Using symmetries to accelerate the tree traversal

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$



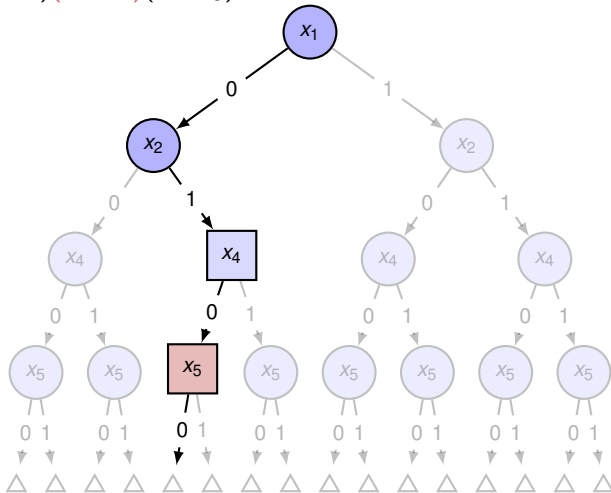
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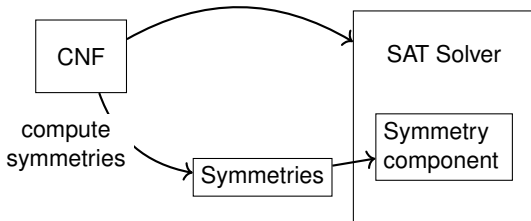
# Using symmetries to accelerate the tree traversal

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$



Use symmetries to deduce symmetrical facts.

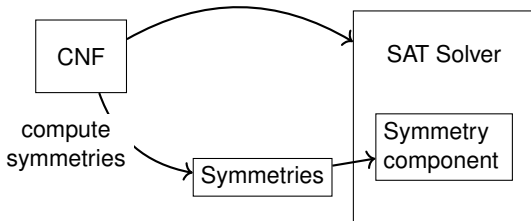
# Dynamic Symmetry Breaking



State-of-the-art:

- Symmchaff [Sab05]
- Symmetry Propagation (SP) [DBdC<sup>+</sup>12]
- Symmetry Learning Scheme (SLS) [BNOS10]
- Symmetry Explanation Learning (SEL) [DBB17]

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Cannot handle some instances solved by static approach



Second contribution

Composing Symmetry Propagation and  
Effective Symmetry Breaking for SAT Solving

NFM'19 [MBK19]

# ESBP + SP

Compose the symmetry propagation and the ESBP

*prune the decision tree while accelerating its traversal*

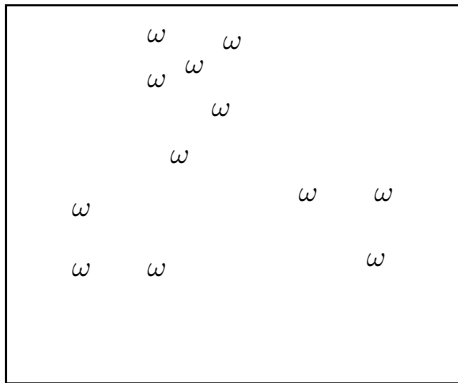
Problems:

- ESBP breaks symmetries (incrementally)
- SP considers the manipulated symmetries valid all time

In a hybrid approach, SP must be able to identify  
**valid symmetries**

# Local symmetry

- formula
- $\omega$  clause
- $\omega$  learnt clause

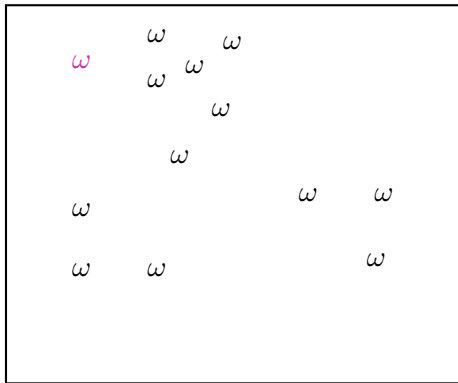


Local symmetries:

$$\omega \leftarrow \{g_1, g_2, g_3\}$$

# Local symmetry

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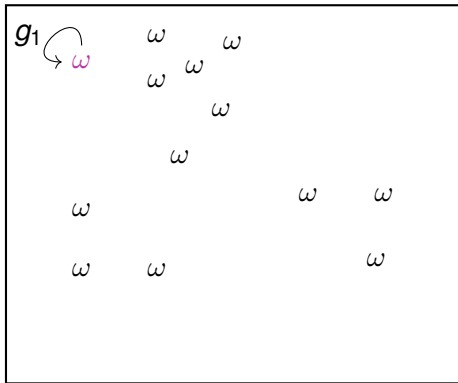


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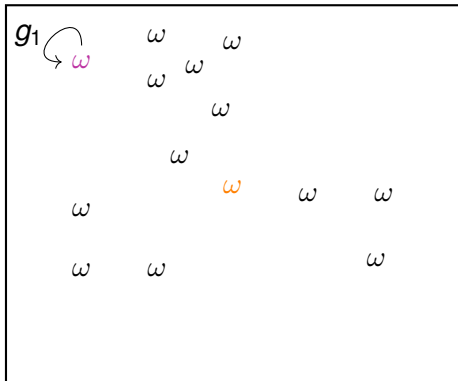
$$\omega \leftarrow \{g_1, g_2, g_3\}$$

$$\omega \leftarrow \{g_1\}$$

- Compute valid local symmetries
- On the fly
- At minimal cost

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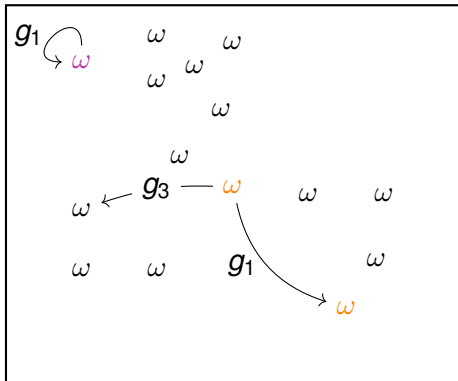
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**Inductive construction**

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- Compute valid local symmetries
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**Inductive construction**

# Experimental results

## Benchmark:

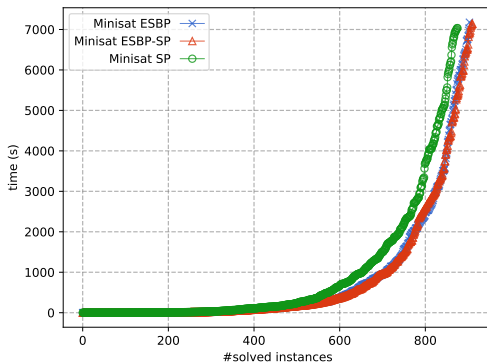
- from SAT contests 2012 – 2018
- retain only instances for which `bliss` finds significant symmetries in 1000 seconds
- 1400 symmetric instances (out of 4000)

## Setup:

- three tools
  - MiniSat SP (Minisat with Symmetry Propagation)
  - MiniSat ESBP (Minisat with CDCL[Sym])
  - **Minisat ESBP-SP** (our approach)
- 7200 seconds timeout



# Experimental results



Solver	PAR-2	ALL	SAT	UNSAT
SP	1674h00	876	406	470
ESBP	1578h30	904	416	488
<b>ESBP-SP</b>	<b>1570h15</b>	<b>911</b>	<b>420</b>	<b>491</b>

# Conclusion & Perspectives

## Conclusion

- A new dynamic symmetry breaking approach
  - Generation of SBP on the fly
  - Package as a library cosy usable with any CDCL solver
- A new hybrid approach (ESBP-SP)
  - Take advantage of static and dynamic approach
  - Combination of CDCL[Sym] with SEL

## Perspectives

- Symmetries and parallel SAT solver
- Exploitation of partial symmetries

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**Thanks !**



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# CDCL in action TODO



$$\omega_1 = \{x_1, x_2, x_3\}$$

$$\omega_2 = \{x_4, x_5, x_6\}$$

$$\omega_3 = \{\neg x_1, \neg x_5\}$$

$$\omega_4 = \{\neg x_2, \neg x_4\}$$

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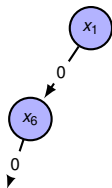
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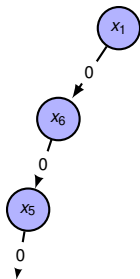
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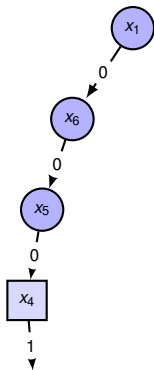
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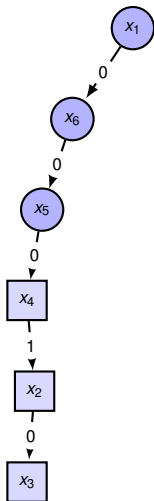
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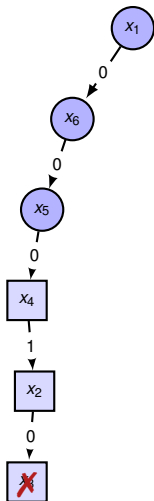
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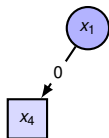
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# Weakly active symmetries

## Logical consequence

When  $\omega$  is satisfied in all satisfying assignments of  $\varphi$ , we say that  $\omega$  is a logical consequence of  $\varphi$ , and we denote this by  $\varphi \vdash \omega$ .

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## Symmetry propagation

Let  $\sigma$  a weakly active symmetry, then

$$\varphi \cup \alpha \vdash \{I\} \Leftrightarrow \varphi \cup \alpha \vdash \sigma.\{I\}$$



# Local symmetries

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Let  $\varphi$  be a formula. We define  $L_{\omega, \varphi}$ , the set of *local symmetries* for a clause  $\omega$ , and with respect to a formula  $\varphi$ , as follows:

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$$L_{\omega, \varphi} = \{\sigma \in \mathfrak{S} \mid \varphi \vdash \sigma.\omega\}$$

We can state that:

$$\bigcap_{\omega \in \varphi} L_{\omega, \varphi} \subseteq G.$$

# Computing local symmetries

Formula can be decomposed as :  $\varphi = \varphi_o \cup \varphi_e \cup \varphi_d$  where

- $\varphi_o$  is the set of the original clauses
- $\varphi_e$  is the set of ESBPs
- $\varphi_d$  is the set of deduced clauses.

## Local symmetries

- $\omega \in \varphi_o, L_{\omega, \varphi} \supseteq G$
- $\omega \in \varphi_e, L_{\omega, \varphi} \supseteq \text{Stab}(\omega) = \{\sigma \in G \mid \omega = \sigma.\omega\}$
- $\omega \in \varphi_d, L_{\omega, \varphi} \supseteq \left( \bigcap_{\omega' \in \varphi_1} L_{\omega', \varphi} \right) \cup \text{Stab}(\omega)$

where  $\varphi_1$  is the set of clauses that derives  $\omega$ .

# Generates symmetry breaking predicates (SBP)

- Define lexicographic order
  - Define total order on variables
  - Define minimal value
- Forbid non minimal assignment for each orbit

Example:

$$x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8; \textcolor{red}{F} < \textcolor{green}{T}$$

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$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$\dots$	lex-leader	SBP

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	T	F	—	—	—	$\dots$	✗	

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	$\textcolor{green}{T}$	$\textcolor{red}{F}$	—	—	—	$\dots$	✗	
$O_2$	$\textcolor{red}{F}$	$\textcolor{red}{F}$	—	$\textcolor{red}{F}$	$\textcolor{green}{T}$	$\dots$	✓	

# Generates symmetry breaking predicates (SBP)

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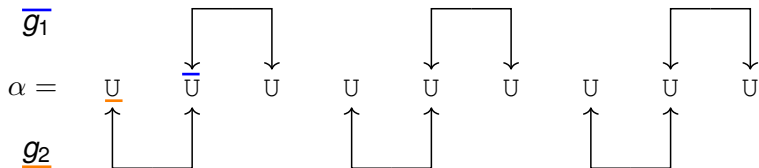
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$\dots$								

# Example

$$g_1 = (x_2 \ x_3)(x_5 \ x_6)(x_8 \ x_9)$$

$$g_2 = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$

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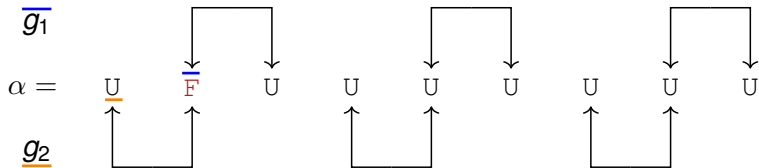


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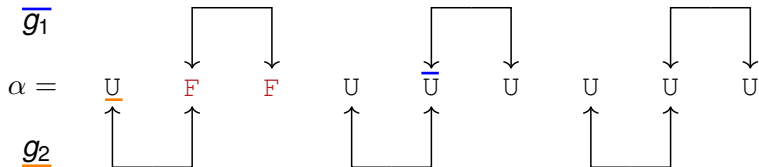


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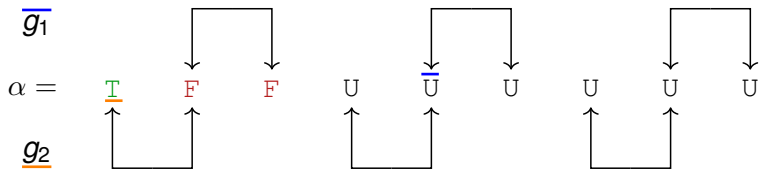


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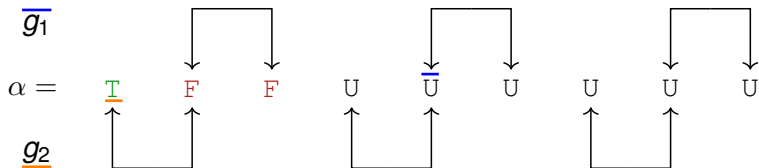


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$g_2$  generates ESBP  $\omega = \{\neg x_1, x_2\}$

# Example

- 1 reducer:  $\alpha(g.x) \leq \alpha(x)$
- 2 inactive:  $\alpha(x) \leq \alpha(g.x)$
- 3 active:  $\alpha(g.x)$  or  $\alpha(x)$  is unassigned

$$x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8 ; \text{ F } < \text{ T }$$

$$g_1 = \begin{array}{ccc} (x_2 & x_3) & (x_5 & x_6) & (x_8 & x_9) \end{array} \left| \begin{array}{l} x = x_2 \\ g.x = x_3 \\ \text{active} \end{array} \right.$$

↑

$$g_2 = \begin{array}{ccc} (x_1 & x_2) & (x_4 & x_5) & (x_7 & x_8) \end{array} \left| \begin{array}{l} x = x_1 \\ g.x = x_2 \\ \text{active} \end{array} \right.$$

↑

...

$$\alpha = \{ \quad \quad \quad \}$$

# Example

- 1 reducer:  $\alpha(g.x) \leq \alpha(x)$
- 2 inactive:  $\alpha(x) \leq \alpha(g.x)$
- 3 active:  $\alpha(g.x)$  or  $\alpha(x)$  is unassigned

$$x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8 ; \text{ F } < \text{ T }$$

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...

$$\alpha = \{ \neg x_2 \quad \}$$



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...

$$\alpha = \{\neg x_2, \neg x_3, x_1\}$$

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...

$$\alpha = \{\neg x_2, \neg x_3, x_1\}$$

$$g_2 \text{ generates } \omega = \{\neg x_1, x_2\}$$

# Encoding the problem

$(A, 1)(A, 2)(A, 3)$

$(B, 1)(B, 2)(B, 3)$

$(C, 1)(C, 2)(C, 3)$

$x_1 \vee x_2 \vee x_3$

$x_4 \vee x_5 \vee x_6$

$x_7 \vee x_8 \vee x_9$

$\neg(A, 1)\neg(B, 1)$

$\neg(A, 1)\neg(C, 1)$

$\neg(B, 1)\neg(C, 1)$

$\neg x_1 \vee \neg x_4$

$\neg x_1 \vee \neg x_7$

$\neg x_4 \vee \neg x_7$

$\neg(A, 2)\neg(B, 2)$

$\neg(A, 2)\neg(C, 2)$

$\neg(B, 2)\neg(C, 2)$

$\neg x_2 \vee \neg x_5$

$\neg x_2 \vee \neg x_8$

$\neg x_5 \vee \neg x_8$

$\neg(A, 3)\neg(B, 3)$

$\neg(A, 3)\neg(C, 3)$

$\neg(B, 3)\neg(C, 3)$

$\neg x_3 \vee \neg x_6$

$\neg x_3 \vee \neg x_9$

$\neg x_6 \vee \neg x_9$