

Exploitation des symétries dynamiques pour la résolution des problèmes SAT

Thèse de doctorat de Sorbonne Université

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Motivation

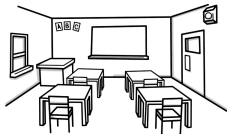
Applications:

- planning
- scheduling
- combinatorial design
- system analysis (model-checking)
- etc

Outline

- ① SAT
 - SAT solving
 - Symmetry
 - Challenges
- ② Tackling the explosion in the static symmetry breaking approach
 -
 -
 - Experimental results
- ③ Taking the maximum benefits from static and dynamic approaches
 - a
 - b
 - Experimental results

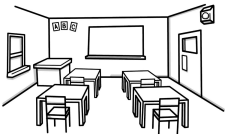
SAT an example



1
↑



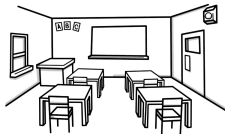
A



2
↑



B



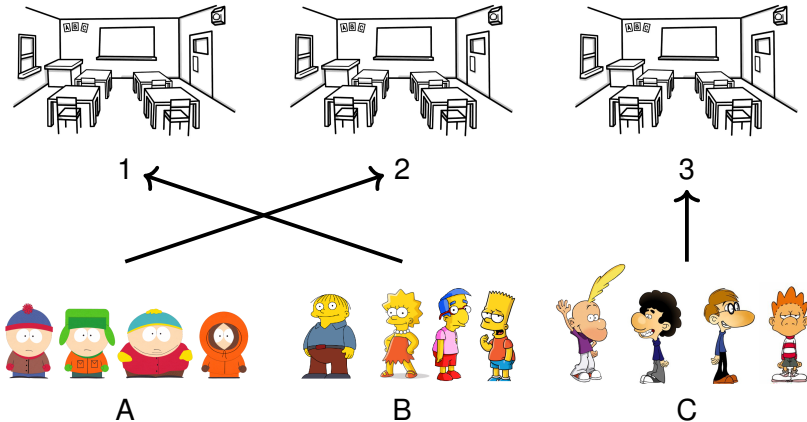
3
↑



C

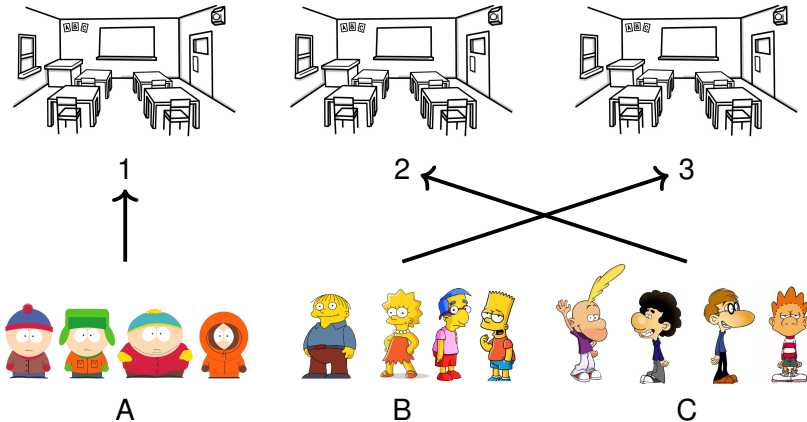
Attribute each group to a class room

SAT an example



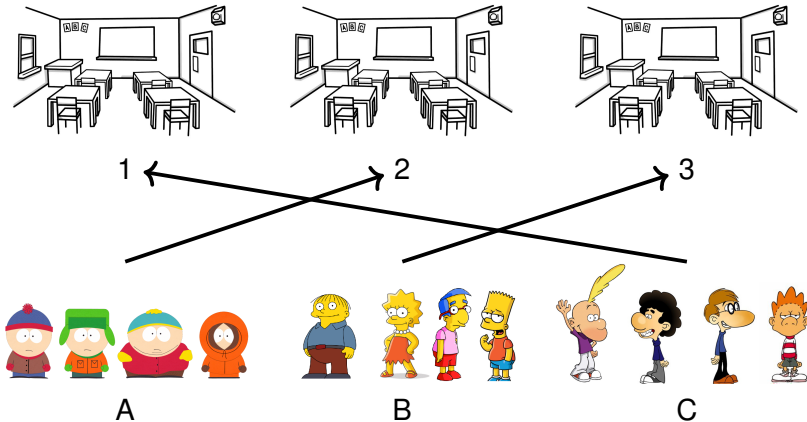
Attribute each group to a class room

SAT an example



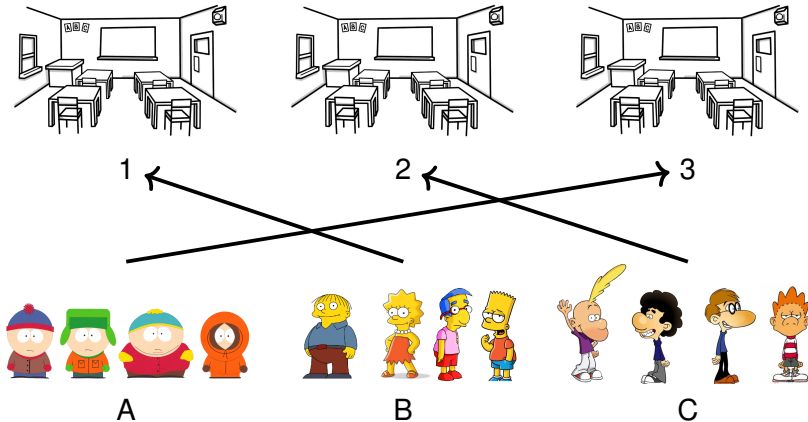
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SAT an example



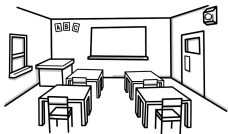
Attribute each group to a class room

SAT an example

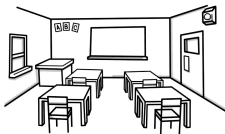


Attribute each group to a class room

SAT an example 2



1



2



A



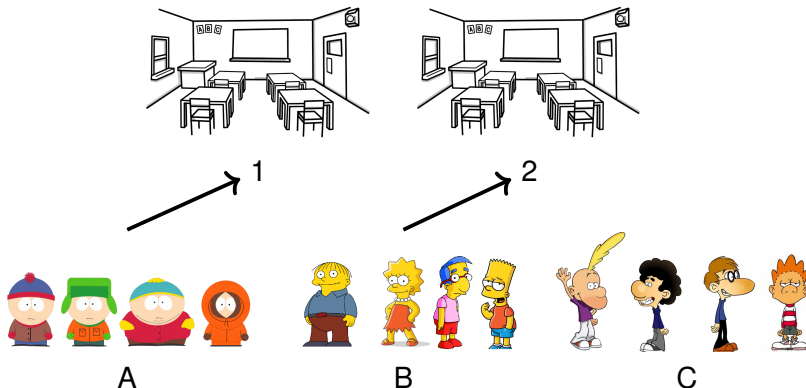
B



C

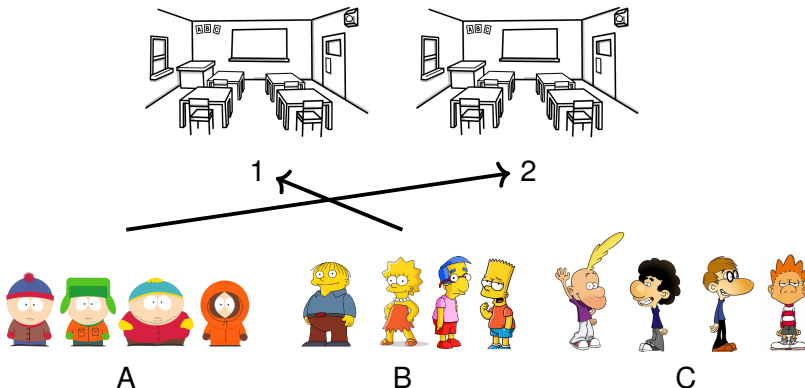
Attribute each group to a class room

SAT an example 2



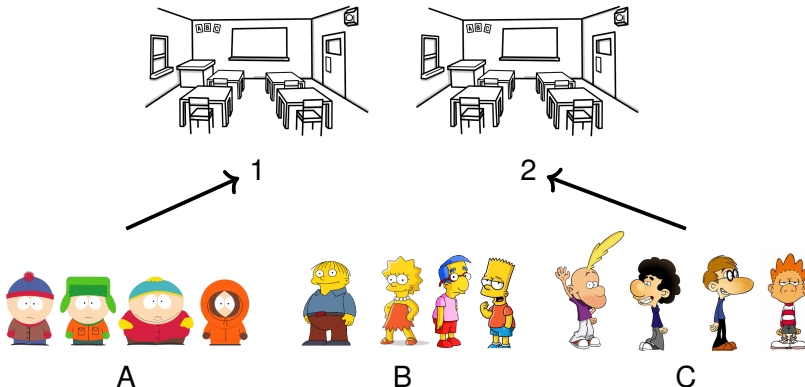
Attribute each group to a class room

SAT an example 2



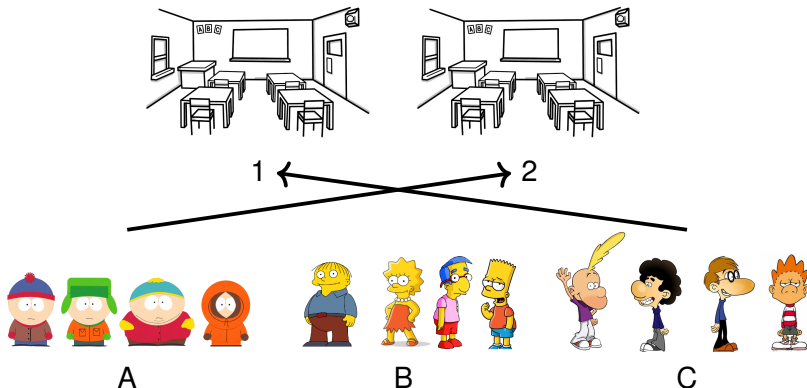
Attribute each group to a class room

SAT an example 2



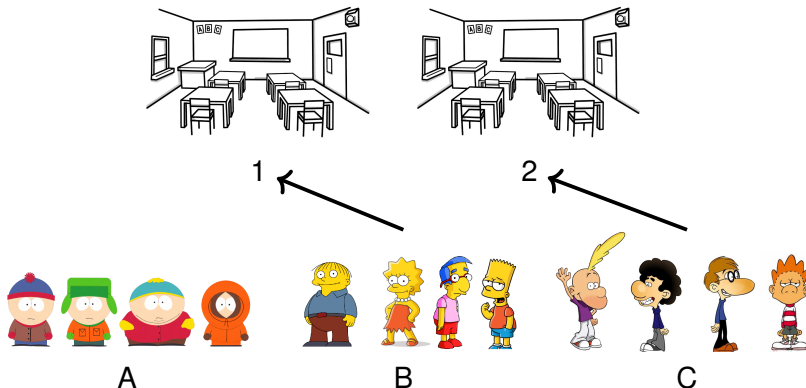
Attribute each group to a class room

SAT an example 2



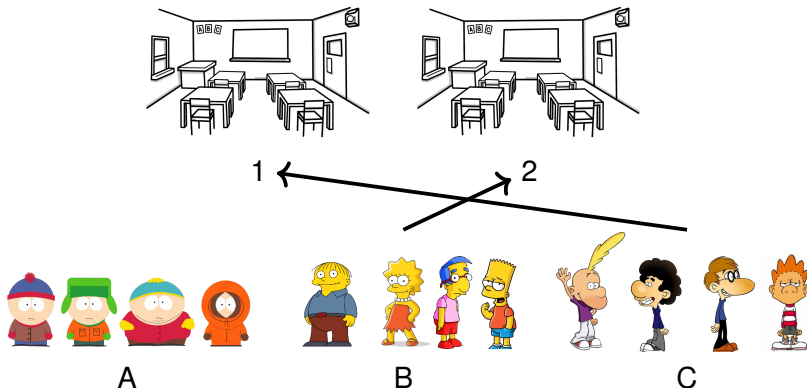
Attribute each group to a class room

SAT an example 2



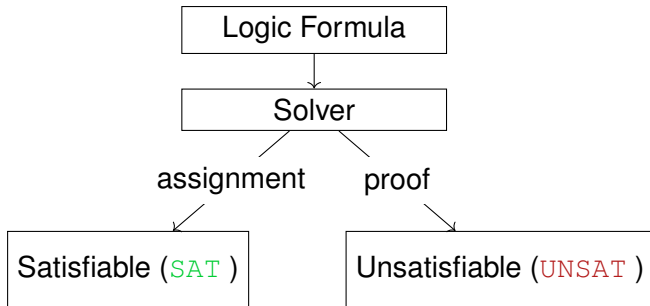
Attribute each group to a class room

SAT an example 2

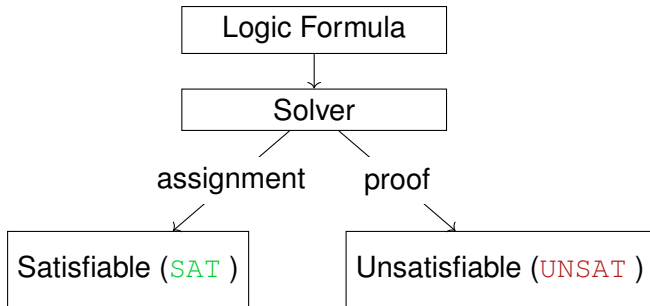


Attribute each group to a class room

SAT



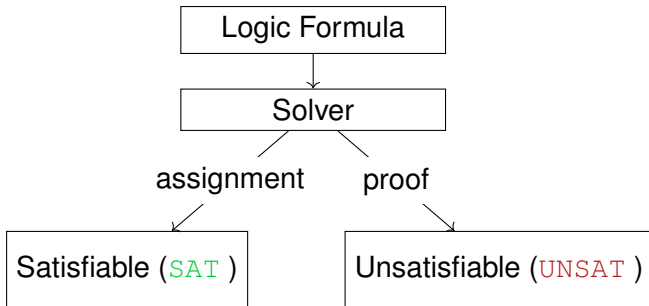
SAT



$$\underbrace{(x_1 \vee x_2 \vee \neg x_3)}$$

Clause with literals $x_1, x_2, \neg x_3$

SAT



$$\overbrace{(x_1 \vee x_2 \vee \neg x_3) \wedge (\neg x_1 \vee \neg x_2) \wedge (x_2 \vee \neg x_4)}^{\text{Formula (CNF)}}$$

Clause

Computing symmetries of a SAT problem

CNF formula

$$\begin{aligned} & (x_1 \vee x_2 \vee x_3) \wedge (x_4 \vee x_5 \vee x_6) \wedge (x_7 \vee x_8 \vee x_9) \\ & \wedge (\neg x_1 \vee \neg x_4) \wedge (\neg x_1 \vee \neg x_7) \wedge (\neg x_4 \vee \neg x_7) \\ & \wedge (\neg x_2 \vee \neg x_5) \wedge (\neg x_2 \vee \neg x_8) \wedge (\neg x_5 \vee \neg x_8) \\ & \wedge (\neg x_3 \vee \neg x_6) \wedge (\neg x_3 \vee \neg x_9) \wedge (\neg x_6 \vee \neg x_9) \end{aligned}$$

¹<http://www.tcs.hut.fi/Software/bliss/>

²<http://vlsicad.eecs.umich.edu/BK/SAUCY/>

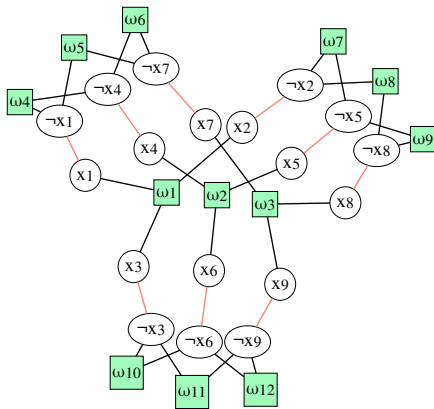
Computing symmetries of a SAT problem

CNF formula

$$\begin{aligned} & (x_1 \vee x_2 \vee x_3) \wedge (x_4 \vee x_5 \vee x_6) \wedge (x_7 \vee x_8 \vee x_9) \\ & \wedge (\neg x_1 \vee \neg x_4) \wedge (\neg x_1 \vee \neg x_7) \wedge (\neg x_4 \vee \neg x_7) \\ & \wedge (\neg x_2 \vee \neg x_5) \wedge (\neg x_2 \vee \neg x_8) \wedge (\neg x_5 \vee \neg x_8) \\ & \wedge (\neg x_3 \vee \neg x_6) \wedge (\neg x_3 \vee \neg x_9) \wedge (\neg x_6 \vee \neg x_9) \end{aligned}$$



colored graph



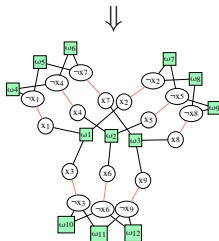
Computing symmetries of a SAT problem

CNF formula

$$\begin{aligned} &(x_1 \vee x_2 \vee x_3) \wedge (x_4 \vee x_5 \vee x_6) \wedge (x_7 \vee x_8 \vee x_9) \\ &\wedge (\neg x_1 \vee \neg x_4) \wedge (\neg x_1 \vee \neg x_7) \wedge (\neg x_4 \vee \neg x_7) \\ &\wedge (\neg x_2 \vee \neg x_5) \wedge (\neg x_2 \vee \neg x_8) \wedge (\neg x_5 \vee \neg x_8) \\ &\wedge (\neg x_3 \vee \neg x_6) \wedge (\neg x_3 \vee \neg x_9) \wedge (\neg x_6 \vee \neg x_9) \end{aligned}$$



colored graph



graph automorphism



(bliss¹ or saucy²)

¹<http://www.tcs.hut.fi/Software/bliss/>

²<http://vlsicad.eecs.umich.edu/BK/SAUCY/>

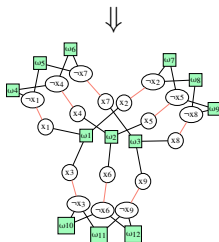
Computing symmetries of a SAT problem

CNF formula

$$\begin{aligned} & (x_1 \vee x_2 \vee x_3) \wedge (x_4 \vee x_5 \vee x_6) \wedge (x_7 \vee x_8 \vee x_9) \\ & \wedge (\neg x_1 \vee \neg x_4) \wedge (\neg x_1 \vee \neg x_7) \wedge (\neg x_4 \vee \neg x_7) \\ & \wedge (\neg x_2 \vee \neg x_5) \wedge (\neg x_2 \vee \neg x_8) \wedge (\neg x_5 \vee \neg x_8) \\ & \wedge (\neg x_3 \vee \neg x_6) \wedge (\neg x_3 \vee \neg x_9) \wedge (\neg x_6 \vee \neg x_9) \end{aligned}$$



colored graph



⇓
graph automorphism

⇓
set of symmetries

⇓
(bliss¹ or saucy²)

⇓

$$\begin{aligned} g_1 &= (x_2 \ x_3)(x_5 \ x_6)(x_8 \ x_9) \\ g_2 &= (x_4 \ x_7)(x_5 \ x_8)(x_6 \ x_9) \\ g_3 &= (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8) \\ g_4 &= (x_1 \ x_4)(x_2 \ x_5)(x_3 \ x_6) \end{aligned}$$

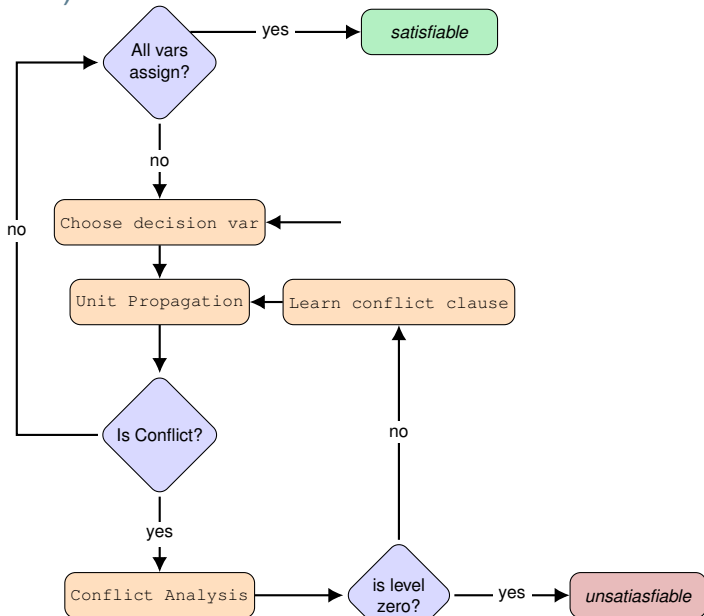
¹<http://www.tcs.hut.fi/Software/bliss/>

²<http://vlsicad.eecs.umich.edu/BK/SAUCY/>

SAT

example solving arbre

Conflict Driven Clause Learning Algorithm (CDCL)



Tree

