

# Exploitation of dynamic symmetries for solving SAT problems

Thèse de doctorat de Sorbonne Université

Hakan METIN

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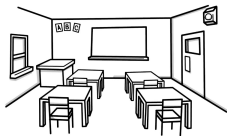


# Motivation

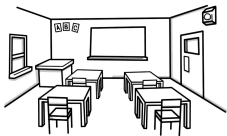
Boolean SATisfiability is widely used in different domains

- Artificial intelligence (planning [KS<sup>+</sup>92], ...)
- Bioinformatics (haplotype inference [LMS06], ...)
- Security (cryptanalysis [MM00], ...)
- Computationally hard problems (ramsey numbers, graph coloring, ...)
- Formal methods, (bounded model checking [BCCZ99], ...)

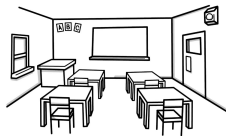
# SAT: an example



1



2



3



A



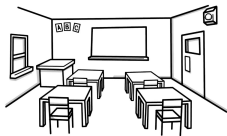
B



C

Is it possible to attribute each group to a classroom?

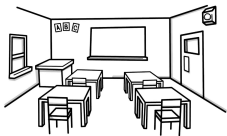
# SAT: an example



1  
↑



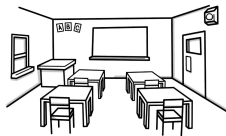
A



2  
↑



B



3  
↑

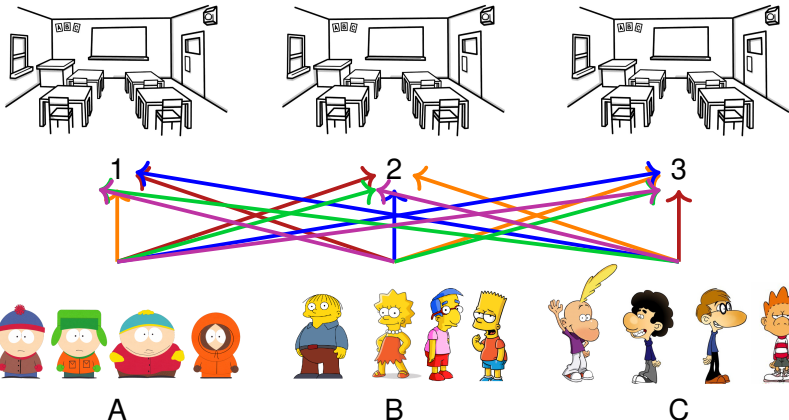


C

Is it possible to attribute each group to a classroom?

YES! SAT

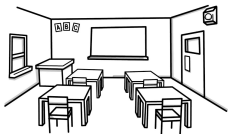
# SAT: an example



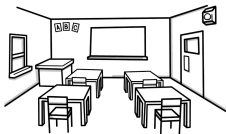
Is it possible to attribute each group to a classroom?

YES! SAT Many solutions

# An UNSAT example



1



2



A



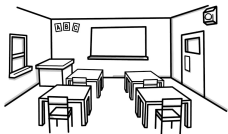
B



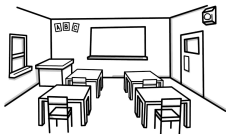
C

Is it possible to attribute each group to a classroom?

# An UNSAT example



1



2



A



B

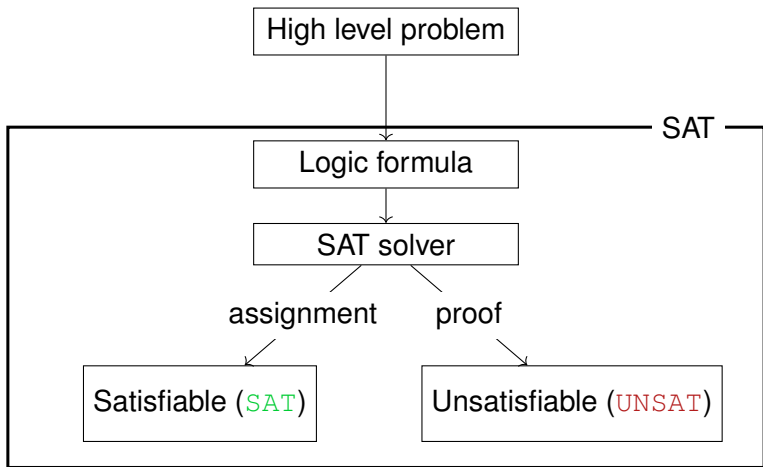


C

Is it possible to attribute each group to a classroom?

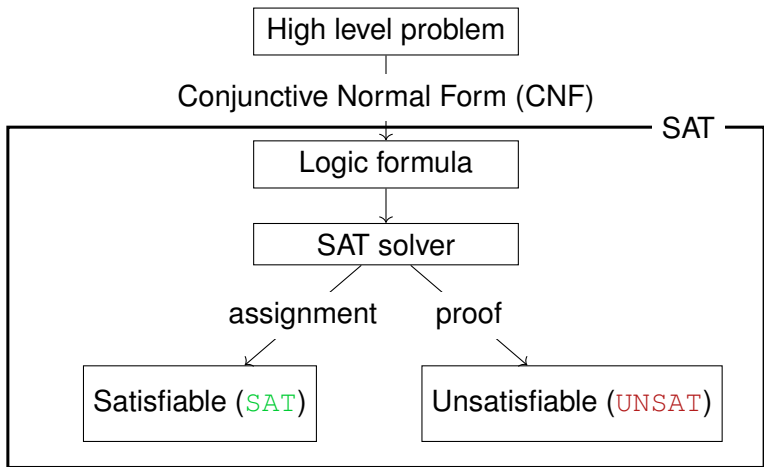
No! UNSAT

# Boolean SATisfiability problem





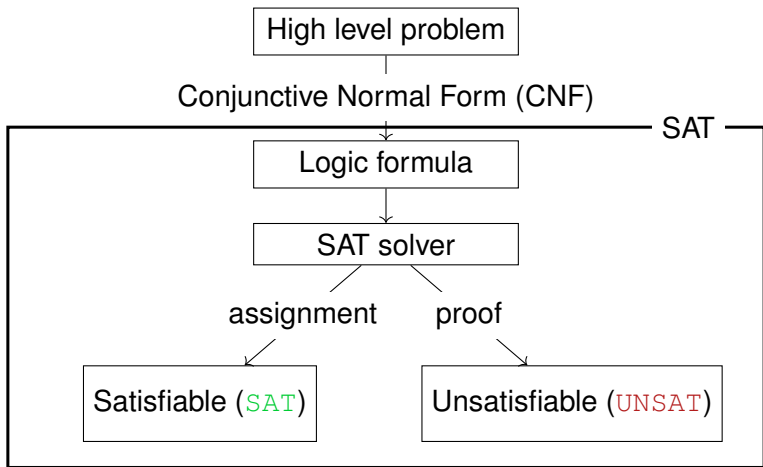
# Boolean SATisfiability problem



CNF representation:

$$\underbrace{(x_1 \vee x_2 \vee \neg x_3)}_{\text{Clause with literals } x_1, x_2, \neg x_3}$$

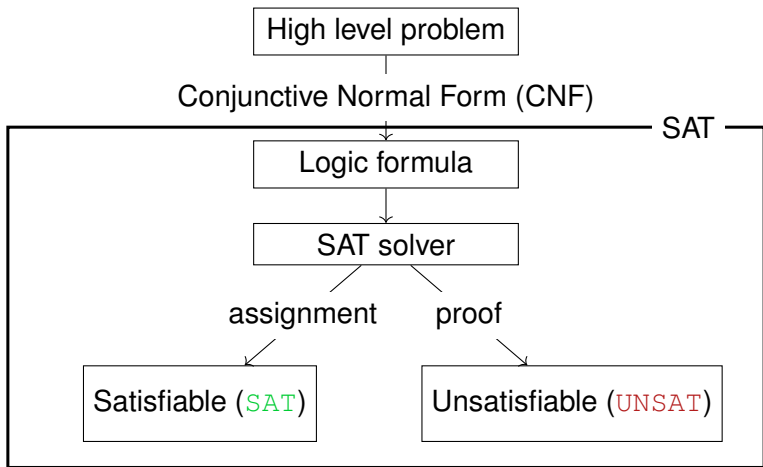
# Boolean SATisfiability problem



CNF representation:

$$\underbrace{(x_1 \vee x_2 \vee \neg x_3)}_{\text{Clause}} \wedge \underbrace{(\neg x_1 \vee \neg x_2) \wedge (x_2 \vee \neg x_4)}_{\text{Formula (CNF)}}$$

# Boolean SATisfiability problem



Clause representation as a set:

$$(x_1 \vee x_2 \vee \neg x_3) \rightarrow \{x_1, x_2, \neg x_3\}$$

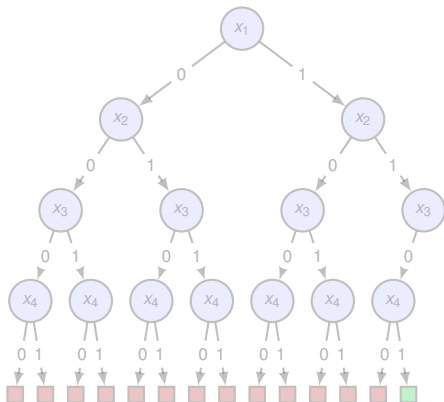
# SAT Solving

Solving SAT formula is known to be **NP-complete** [Coo71]

Enumerative algorithms:

- Davis, Putnam, Logemann, and Loveland (DPLL) [DLL62]
  - Boolean Constraint Propagation (BCP)
- Conflict Driven Clause Learning (CDCL) [MSS99]
  - Derived from DPLL
  - Clause learning

# CDCL in action



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

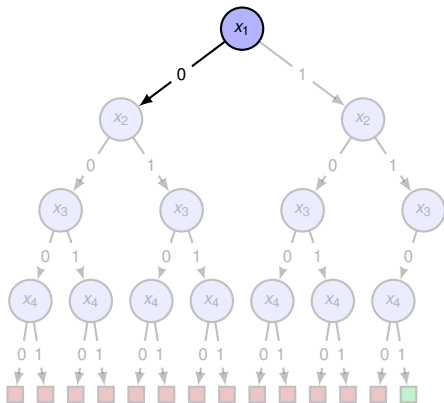
$$\omega_3 = \{x_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

# CDCL in action



$$\omega_1 = \{\mathbf{x}_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{\mathbf{x}_1, \neg x_4\}$$

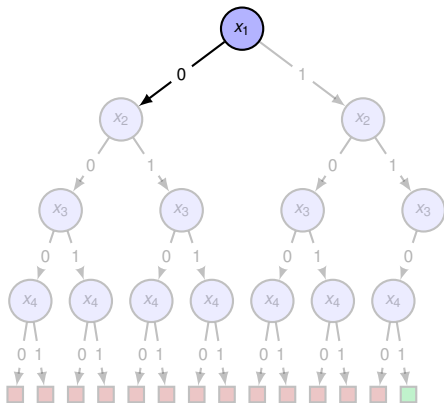
$$\omega_3 = \{\mathbf{x}_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

# CDCL in action



$$\omega_1 = \{\textcolor{red}{x}_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{\textcolor{red}{x}_1, \neg \textcolor{blue}{x}_4\}$$

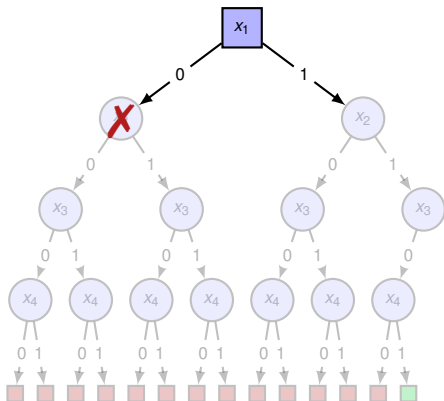
$$\omega_3 = \{\textcolor{red}{x}_1, \textcolor{blue}{x}_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

# CDCL in action



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

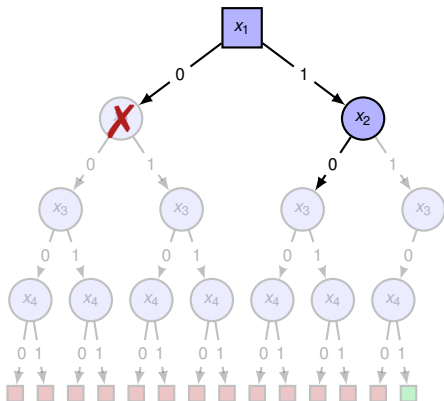
$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

$$\omega_7 = \{x_1\}$$



# CDCL in action



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

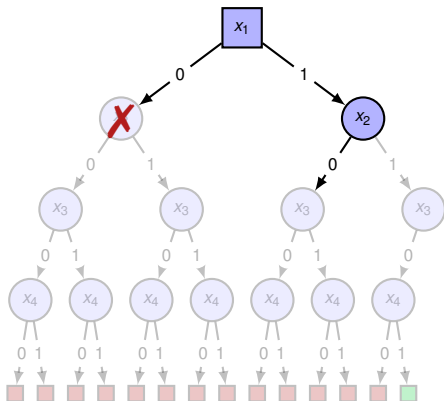
$$\omega_4 = \{x_2, \neg x_4\}$$

$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

$$\omega_7 = \{x_1\}$$

# CDCL in action



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

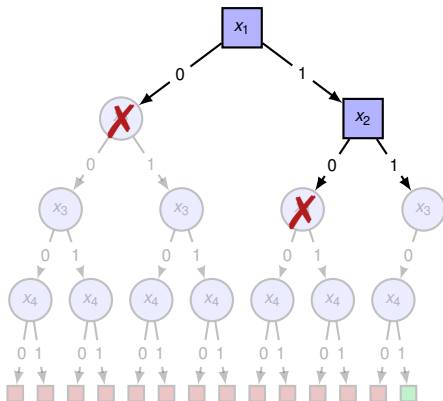
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$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

$$\omega_7 = \{x_1\}$$

# CDCL in action



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

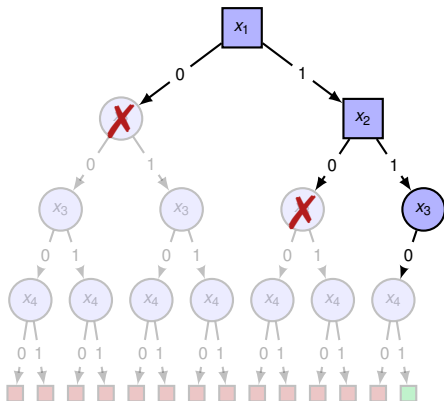
$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

$$\omega_7 = \{x_1\}$$

$$\omega_8 = \{x_2\}$$

# CDCL in action



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

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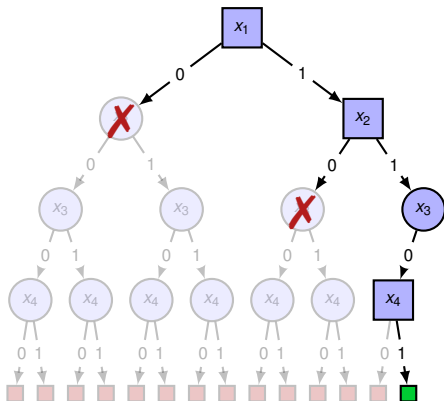
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$$\omega_7 = \{x_1\}$$

$$\omega_8 = \{x_2\}$$

# CDCL in action



$$\omega_1 = \{x_1, x_2, x_3, x_4\}$$

$$\omega_2 = \{x_1, \neg x_4\}$$

$$\omega_3 = \{x_1, x_4\}$$

$$\omega_4 = \{x_2, \neg x_4\}$$

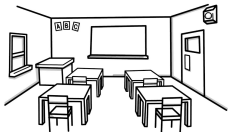
$$\omega_5 = \{x_2, x_4\}$$

$$\omega_6 = \{x_3, x_4\}$$

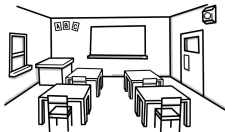
$$\omega_7 = \{x_1\}$$

$$\omega_8 = \{x_2\}$$

# An UNSAT example



1



2



A



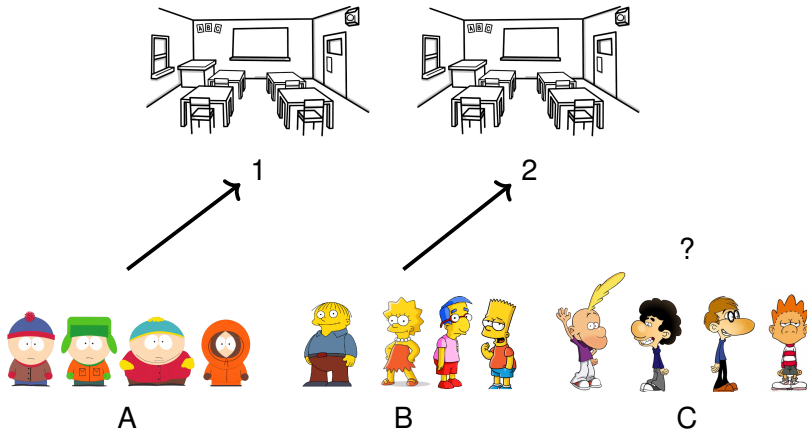
B



C

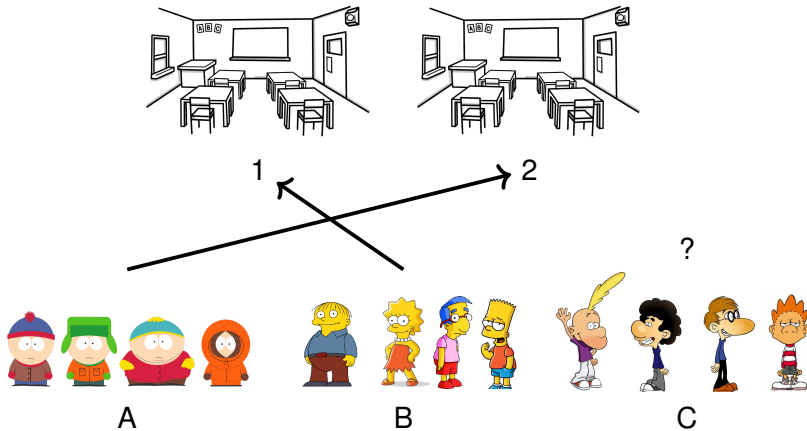
Is it possible to attribute each group to a classroom?

# An UNSAT example



Is it possible to attribute each group to a classroom?

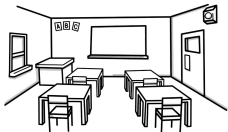
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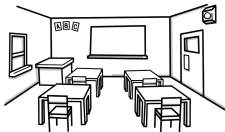
Is it possible to attribute each group to a classroom?



# An UNSAT example



1



2

...



A



B

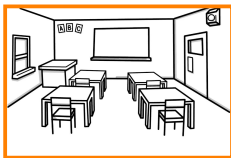


C

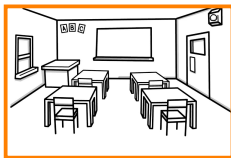
Is it possible to attribute each group to a classroom?

No!

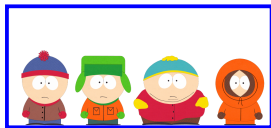
# An UNSAT example



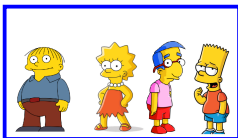
1



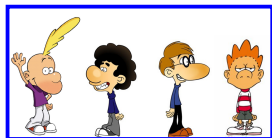
2



A



B



C

Is it possible to attribute each group to a classroom?

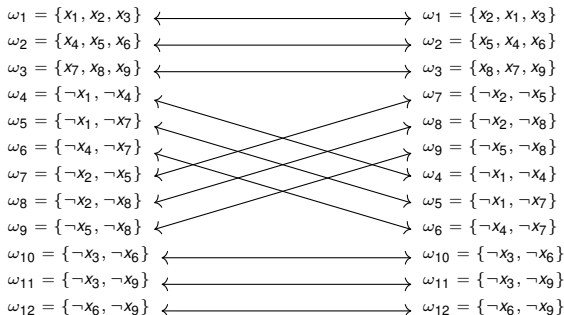
No!

Presence of symmetries hinders the performance of the solver

# Symmetry

A symmetry (permutation)  $g$  is a bijective function (on variables) that leaves the formula invariant

$$g = \begin{pmatrix} x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 & x_8 & x_9 \\ x_2 & x_1 & x_3 & x_5 & x_4 & x_6 & x_8 & x_7 & x_9 \end{pmatrix} \rightarrow (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$



The set of symmetries of a formula is a group noted  $G$

# Computing symmetries of a SAT problem

*CNF formula*

$$\begin{aligned} & (x_1 \vee x_2 \vee x_3) \wedge (x_4 \vee x_5 \vee x_6) \wedge (x_7 \vee x_8 \vee x_9) \\ & \wedge (\neg x_1 \vee \neg x_4) \wedge (\neg x_1 \vee \neg x_7) \wedge (\neg x_4 \vee \neg x_7) \\ & \wedge (\neg x_2 \vee \neg x_5) \wedge (\neg x_2 \vee \neg x_8) \wedge (\neg x_5 \vee \neg x_8) \\ & \wedge (\neg x_3 \vee \neg x_6) \wedge (\neg x_3 \vee \neg x_9) \wedge (\neg x_6 \vee \neg x_9) \end{aligned}$$

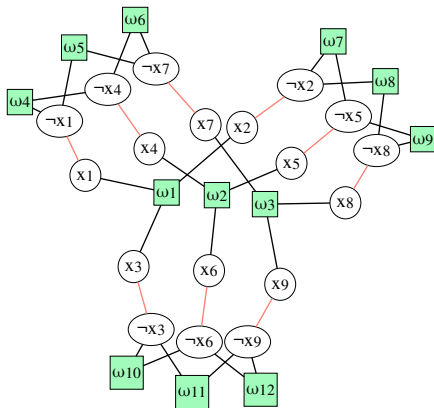
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colored graph



# Computing symmetries of a SAT problem

*CNF formula*

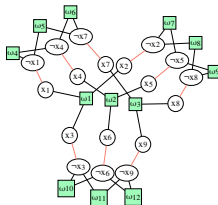


colored graph



graph automorphism

$$\begin{aligned} & (x_1 \vee x_2 \vee x_3) \wedge (x_4 \vee x_5 \vee x_6) \wedge (x_7 \vee x_8 \vee x_9) \\ & \wedge (\neg x_1 \vee \neg x_4) \wedge (\neg x_1 \vee \neg x_7) \wedge (\neg x_4 \vee \neg x_7) \\ & \wedge (\neg x_2 \vee \neg x_5) \wedge (\neg x_2 \vee \neg x_8) \wedge (\neg x_5 \vee \neg x_8) \\ & \wedge (\neg x_3 \vee \neg x_6) \wedge (\neg x_3 \vee \neg x_9) \wedge (\neg x_6 \vee \neg x_9) \end{aligned}$$



(bliss, saucy, ...)

# Computing symmetries of a SAT problem

*CNF formula*



colored graph

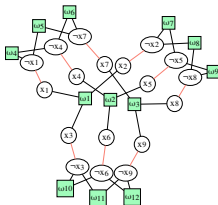


graph automorphism



set of symmetries

$$\begin{aligned} & (x_1 \vee x_2 \vee x_3) \wedge (x_4 \vee x_5 \vee x_6) \wedge (x_7 \vee x_8 \vee x_9) \\ & \wedge (\neg x_1 \vee \neg x_4) \wedge (\neg x_1 \vee \neg x_7) \wedge (\neg x_4 \vee \neg x_7) \\ & \wedge (\neg x_2 \vee \neg x_5) \wedge (\neg x_2 \vee \neg x_8) \wedge (\neg x_5 \vee \neg x_8) \\ & \wedge (\neg x_3 \vee \neg x_6) \wedge (\neg x_3 \vee \neg x_9) \wedge (\neg x_6 \vee \neg x_9) \end{aligned}$$



(bliss, saucy, ...)



$$\begin{aligned} g_1 &= (x_2 \ x_3)(x_5 \ x_6)(x_8 \ x_9) \\ g_2 &= (x_4 \ x_7)(x_5 \ x_8)(x_6 \ x_9) \\ g_3 &= (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8) \\ g_4 &= (x_1 \ x_4)(x_2 \ x_5)(x_3 \ x_6) \end{aligned}$$

# Orbit

Orbit of an assignment  $\alpha = G.\alpha = \{g.\alpha \mid g \in G\}$

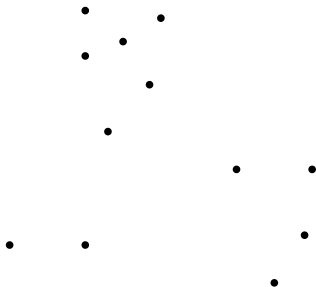


# Orbit

Orbit of an assignment  $\alpha = G.\alpha = \{g.\alpha \mid g \in G\}$

Example:

- full assignment

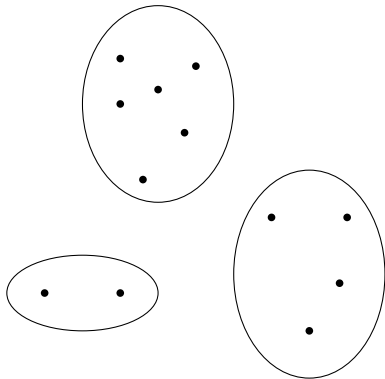


# Orbit

Orbit of an assignment  $\alpha = G.\alpha = \{g.\alpha \mid g \in G\}$

Example:

- full assignment
- orbit

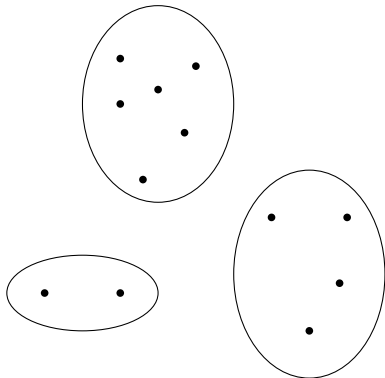


# Orbit

Orbit of an assignment  $\alpha = G.\alpha = \{g.\alpha \mid g \in G\}$

Example:

- full assignment
- orbit

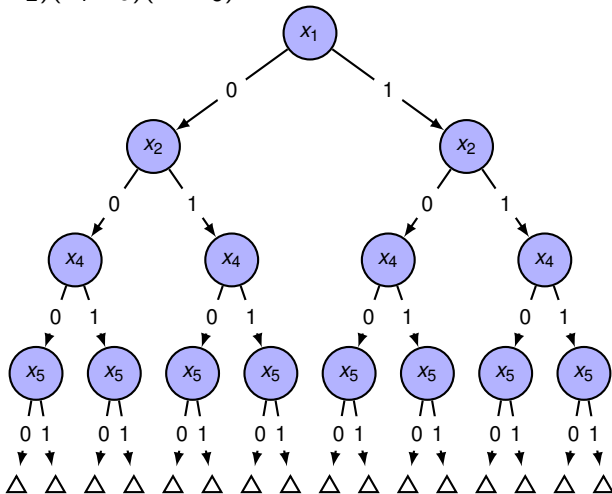


Equivalence relation with respect to SAT:

- Either  $G.\alpha$  contains no solution
- Or all elements of  $G.\alpha$  are solutions

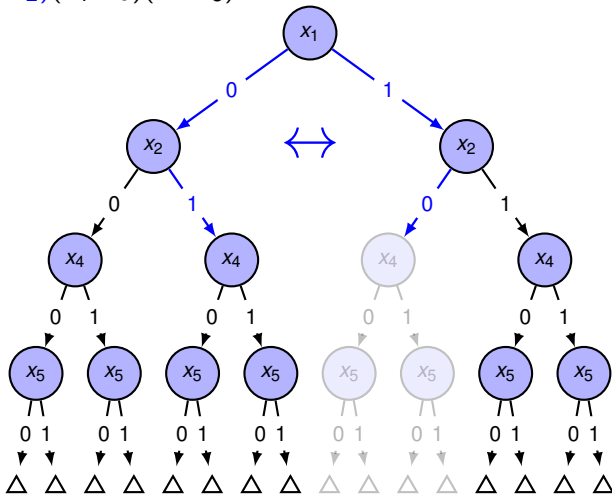
# Using symmetries to prune the search space

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$

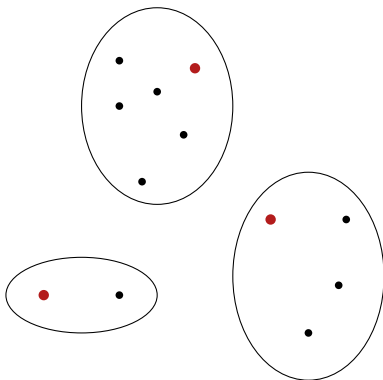


# Using symmetries to prune the search space

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$



# Representative assignment



- full assignment
- orbit
- representative assignment

# Generates symmetry breaking predicates (SBP)

- Define lexicographic order
  - Define total order on variables
  - Define minimal value
- Forbid non minimal assignment for each orbit

Example:

$$x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8; \textcolor{red}{F} < \textcolor{green}{T}$$

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$

$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$\dots$	lex-leader	SBP

# Generates symmetry breaking predicates (SBP)

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Example:

$$x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8; \textcolor{red}{F} < \textcolor{green}{T}$$

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$

	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$\dots$	lex-leader	SBP
$O_1$	$\textcolor{red}{F}$	$\textcolor{green}{T}$	—	—	—	$\dots$	✓	



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Example:

$$x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8; \text{F} < \text{T}$$

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$

	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$\dots$	lex-leader	SBP
$O_1$	F	T	—	—	—	$\dots$	✓	$\rightarrow \neg x_1 \vee x_2$
	T	F	—	—	—	$\dots$	✗	

# Generates symmetry breaking predicates (SBP)

- Define lexicographic order
  - Define total order on variables
  - Define minimal value
- Forbid non minimal assignment for each orbit

Example:

$$x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8; \text{F} < \text{T}$$

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$

	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$\dots$	lex-leader	SBP
$O_1$	F	T	—	—	—	$\dots$	✓	$\rightarrow \neg x_1 \vee x_2$
	T	F	—	—	—	$\dots$	✗	
$O_2$	F	F	—	F	T	$\dots$	✓	

# Generates symmetry breaking predicates (SBP)

- Define lexicographic order
  - Define total order on variables
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Example:

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	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$\dots$	lex-leader	SBP
$O_1$	F	T	—	—	—	$\dots$	✓	$\rightarrow \neg x_1 \vee x_2$
	T	F	—	—	—	$\dots$	✗	
$O_2$	F	F	—	F	T	$\dots$	✓	$\rightarrow x_1 \vee x_2 \vee \neg x_4 \vee x_5$
	F	F	—	T	F	$\dots$	✗	

# Generates symmetry breaking predicates (SBP)

- Define lexicographic order
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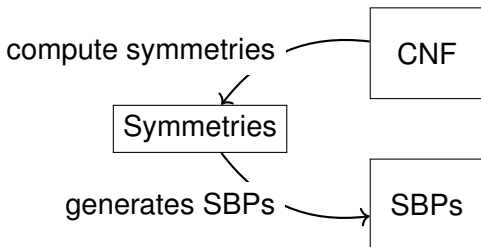
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$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$

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$O_1$	F	T	—	—	—	$\cdots$	✓	$\rightarrow \neg x_1 \vee x_2$
	T	F	—	—	—	$\cdots$	✗	
$O_2$	F	F	—	F	T	$\cdots$	✓	$\rightarrow x_1 \vee x_2 \vee \neg x_4 \vee x_5$
	F	F	—	T	F	$\cdots$	✗	
$\dots$								

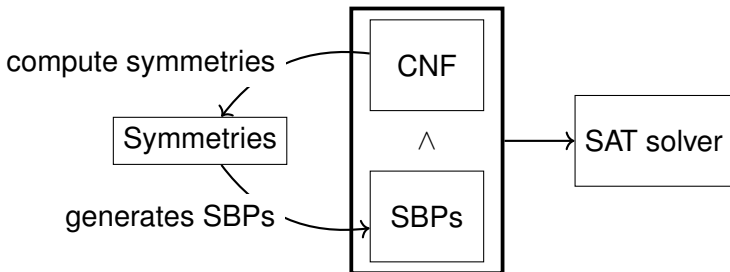
# Static symmetry breaking



Different approaches:

- Shatter [ASM06]
- BreakID [DBBD16]
- ...

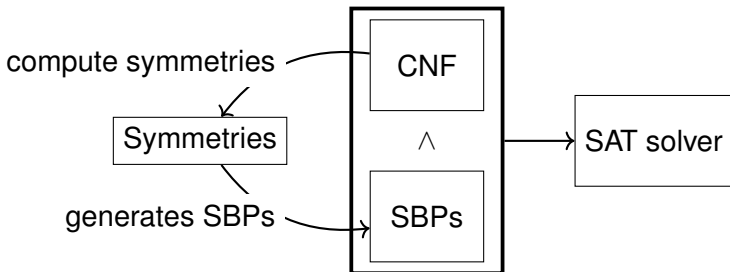
# Static symmetry breaking



Different approaches:

- Shatter [ASM06]
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# Static symmetry breaking



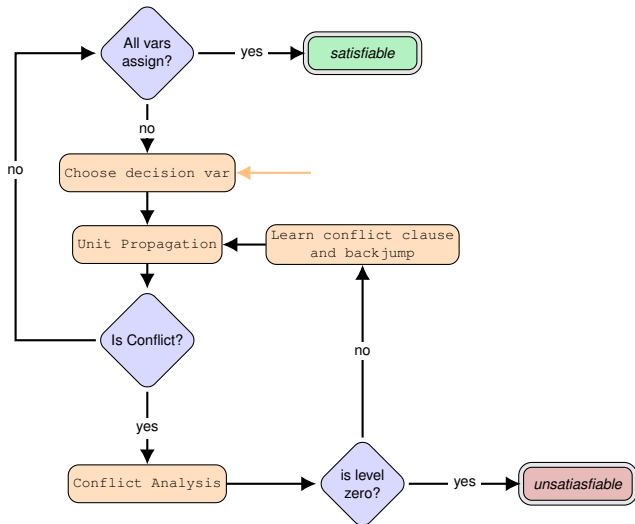
Different approaches:

- Shatter [ASM06]
- BreakID [DBBD16]
- ...

The solver can "explode" instead of being helped

# Our contribution CDCL[Sym] [MBCK18]

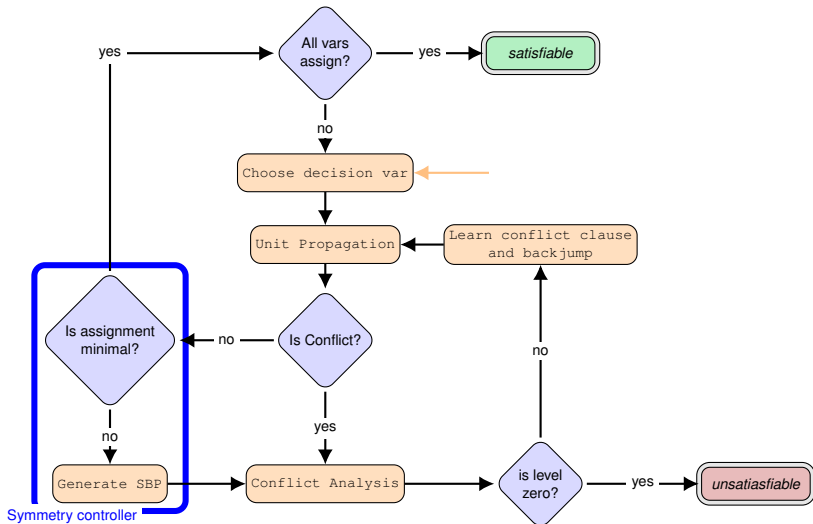
Compute and inject SBP *opportunistically*, during the solving





# Our contribution CDCL[Sym] [MBCK18]

Compute and inject SBP **opportunistically**, during the solving



# Symmetry status

- reducer:  $g.\alpha \prec \alpha$
- inactive:  $\alpha \prec g.\alpha$
- active: *not enough information*

## Efficient implementation of symmetry status

Keep track the smallest unassigned variable  $x$  :

- ①  $\alpha(g.x) \leq \alpha(x)$ , then  $g$  is `reducer`  $\Rightarrow$  Effective SBP (ESBP)
- ②  $\alpha(x) \leq \alpha(g.x)$ , then  $g$  is `inactive`  $\Rightarrow g$  cannot reduce  $\alpha$
- ③  $\alpha(g.x)$  or  $\alpha(x)$  is unassigned then  $g$  is `active`

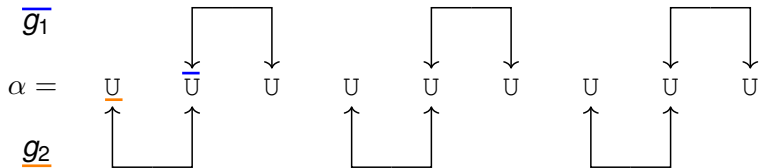
Update whenever variables are assigned / unassigned

# Example

$$g_1 = (x_2 \ x_3)(x_5 \ x_6)(x_8 \ x_9)$$

$$g_2 = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$

$$\textcolor{red}{F} < \textcolor{green}{T} \quad x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8 \leq x_9$$

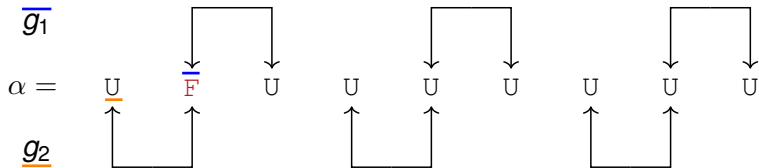


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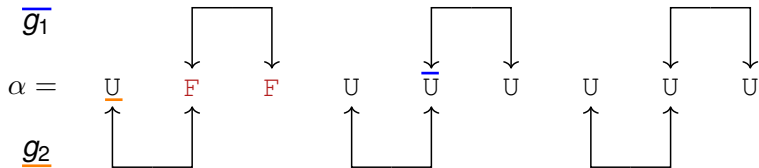


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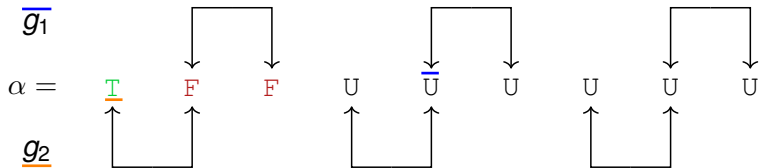


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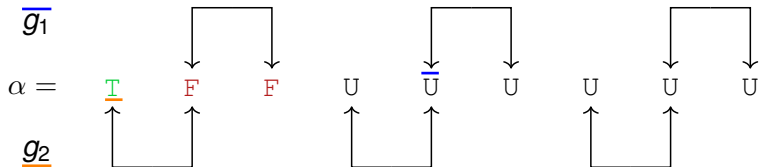


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$g_2$  generates ESBP  $\omega = \{\neg x_1, x_2\}$

# Example

- 1 reducer:  $\alpha(g.x) \leq \alpha(x)$
- 2 inactive:  $\alpha(x) \leq \alpha(g.x)$
- 3 active:  $\alpha(g.x)$  or  $\alpha(x)$  is unassigned

$$x_1 \leq x_2 \leq x_3 \leq x_4 \leq x_5 \leq x_6 \leq x_7 \leq x_8 ; \text{ F } < \text{ T }$$

$$g_1 = \begin{array}{ccc} (x_2 & x_3) & (x_5 & x_6) & (x_8 & x_9) \end{array} \left| \begin{array}{l} x = x_2 \\ g.x = x_3 \\ \text{active} \end{array} \right.$$

↑

$$g_2 = \begin{array}{ccc} (x_1 & x_2) & (x_4 & x_5) & (x_7 & x_8) \end{array} \left| \begin{array}{l} x = x_1 \\ g.x = x_2 \\ \text{active} \end{array} \right.$$

↑

...

$$\alpha = \{ \quad \quad \quad \}$$



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...

$$\alpha = \{ \neg x_2 \quad \}$$

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$$g_2 = \begin{array}{cc|cc} (x_1 & x_2) & (x_4 & x_5) & (x_7 & x_8) \\ \uparrow & & & & & \end{array} \quad \left| \quad x = x_1 \quad \begin{array}{l} g.x = x_2 \\ \text{reducer} \end{array} \right.$$

...

$$\alpha = \{\neg x_2, \neg x_3, x_1\}$$

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...

$$\alpha = \{\neg x_2, \neg x_3, x_1\}$$

$$g_2 \text{ generates } \omega = \{\neg x_1, x_2\}$$

# CDCL[Sym] Implementation

- Packaged as a library **cosy**<sup>1</sup>
  - Lightweight
  - Fast Update and low memory
  - Follows symmetry status
- 
- Works with any enumerative SAT solver

---

<sup>1</sup><https://github.com/lip6/cosy>

# Experiments

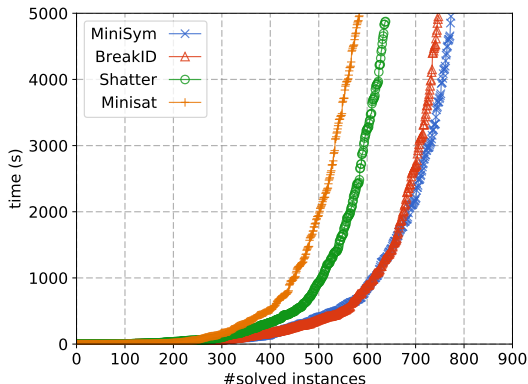
## Benchmark:

- from SAT contests 2012 – 2017
- filter: `bliss` finds significant symmetries in 1000s
- 36 % of instances, 1 350/3 700

## Setup:

- four tools
  - MiniSat (no symmetry, baseline)
  - MiniSat + BreakID (SOTA SAT solver using symmetries)
  - MiniSat + Shatter (SOTA SAT solver using symmetries)
  - **MiniSym** = MiniSat + CDCL[Sym] (our approach)
- 5000s timeout, 8GB memory
- includes time to compute symmetries (except for MiniSat)

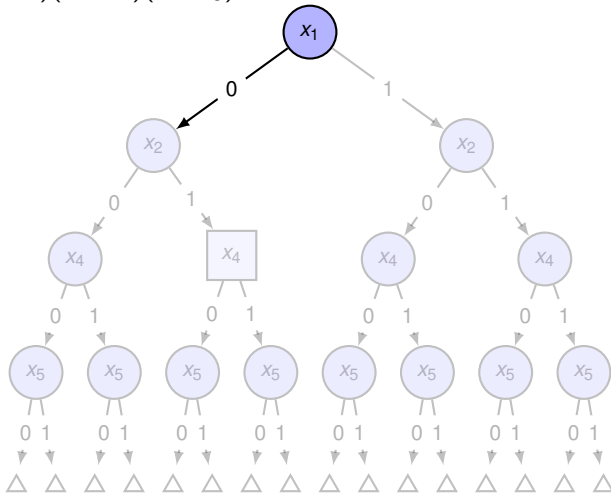
# Experimental results



Solver	PAR-2	ALL	SAT	UNSAT
MiniSAT	2243h	586	325	261
Shatter	2088h	640	316	324
BreakID	1790h	749	334	415
MiniSym	<b>1735h</b>	<b>775</b>	<b>336</b>	<b>439</b>

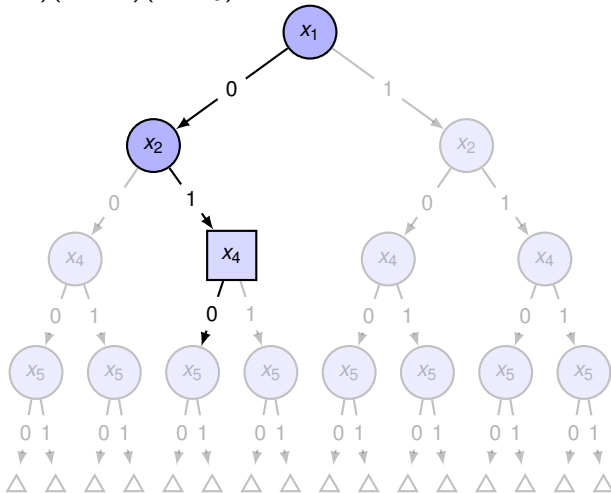
# Using symmetries to accelerate the tree traversal

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$



# Using symmetries to accelerate the tree traversal

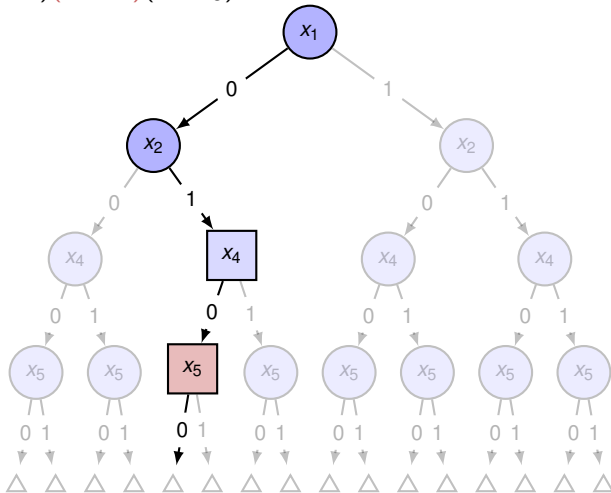
$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$





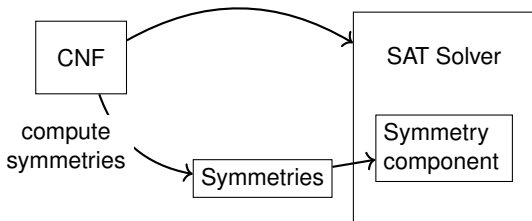
# Using symmetries to accelerate the tree traversal

$$g = (x_1 \ x_2)(x_4 \ x_5)(x_7 \ x_8)$$



Use symmetries to deduce symmetrical facts.

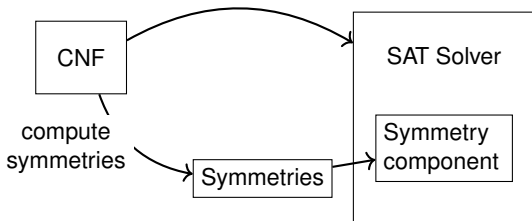
# Dynamic Symmetry Breaking



Different approaches:

- Symmchaff [Sab05]
- Symmetry Propagation (SP) [DBdC<sup>+</sup>12]
- Symmetry Learning Scheme (SLS) [BNOS10]
- Symmetry Explanation Learning (SEL) [DBB17]
- ...

# Dynamic Symmetry Breaking



Different approaches:

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- Symmetry Explanation Learning (SEL) [DBB17]
- ...

Cannot handle some instances solved by static approach

## ESBP + SP [MBK19]

Compose the symmetry propagation and the ESBP

*prune the decision tree while accelerating its traversal*

Problems:

- ESBP breaks symmetries (incrementally)
- SP considers the manipulated symmetries valid all time

In a hybrid approach, SP must be able to identify  
**valid symmetries**

# Local symmetries

Formula  $\leftarrow$  (Symmetries)

$\omega_1 \leftarrow$  (Local symmetries)

$\omega_2 \leftarrow$  (Local symmetries)

$\omega_3 \leftarrow$  (Local symmetries)

$\omega_4 \leftarrow$  (Local symmetries)

Macro level

$\rightarrow$

Micro level

# Local symmetries

Formula  $\leftarrow$  (Symmetries)

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$\omega_3 \leftarrow$  (Local symmetries)

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$\omega_5$

Macro level

$\rightarrow$

Micro level

# Local symmetries

Formula  $\leftarrow$  (Symmetries)

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Macro level

$\rightarrow$

Micro level

Compute valid local symmetries on-the-fly at a minimal cost.

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Formula  $\leftarrow$  (Symmetries)

$\omega_1 \leftarrow$  (Local symmetries)

$\omega_2 \leftarrow$  (Local symmetries)

$\omega_3 \leftarrow$  (Local symmetries)

$\omega_4 \leftarrow$  (Local symmetries)

$\omega_5$

Macro level

$\rightarrow$

Micro level

Compute valid local symmetries on-the-fly at a minimal cost.

- Inductive construction of the valid symmetries
- During the solving
- At a minimal cost



# Experimental results

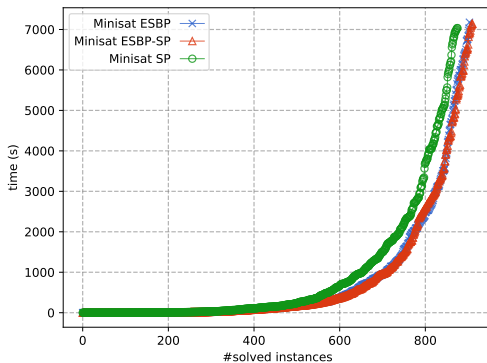
## Benchmark:

- from SAT contests 2012 – 2018
- retain only instances for which `bliss` finds significant symmetries in 1000s
- 1400 symmetric instances (out of 4000)

## Setup:

- three tools
  - MiniSat SP (Minisat with Symmetry Propagation)
  - MiniSat ESBP (Minisat with CDCL[Sym])
  - **Minisat ESBP-SP** (our approach)
- 7200s timeout

# Experimental results



Solver	PAR-2	ALL	SAT	UNSAT
SP	1674h00	876	406	470
ESBP	1578h30	904	416	488
ESBP-SP	<b>1570h15</b>	<b>911</b>	<b>420</b>	<b>491</b>

# Conclusion

- A new dynamic symmetry breaking approach
  - Generation of SBP on the fly
  - Package as a library cosy usable with any CDCL solver
  - Overcomes drawbacks of the existing approaches
- A new hybrid approach (ESBP-SP)
  - Take advantage of static and dynamic approach
  - Introduce local symmetries

# Perspectives

- Combination of CDCL[Sym] with other dynamic symmetry breaking approach
- Exploitation of partial symmetries

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- Combination of CDCL[Sym] with other dynamic symmetry breaking approach
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**Thanks !**



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# CDCL in action TODO



$$\omega_1 = \{x_1, x_2, x_3\}$$

$$\omega_2 = \{x_4, x_5, x_6\}$$

$$\omega_3 = \{\neg x_1, \neg x_5\}$$

$$\omega_4 = \{\neg x_2, \neg x_4\}$$

$$\omega_5 = \{\neg x_3, \neg x_4\}$$

$$\omega_6 = \{\neg x_3, \neg x_6\}$$

# CDCL in action TODO



$$\omega_1 = \{\mathbf{x}_1, x_2, x_3\}$$

$$\omega_2 = \{x_4, x_5, x_6\}$$

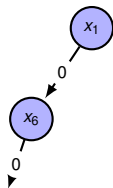
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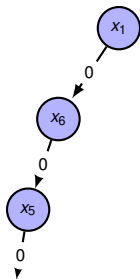
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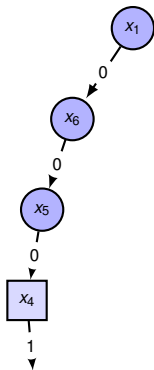
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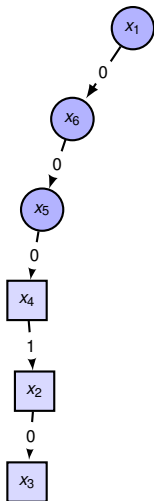
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$$\omega_5 = \{\neg \mathbf{x}_3, \neg \mathbf{x}_4\}$$

$$\omega_6 = \{\neg x_3, \neg \mathbf{x}_6\}$$

# CDCL in action TODO



$$\omega_1 = \{x_1, x_2, x_3\}$$

$$\omega_2 = \{x_4, x_5, x_6\}$$

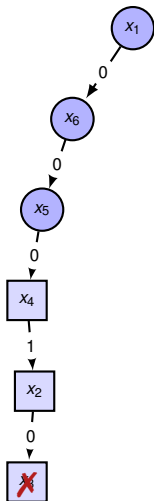
$$\omega_3 = \{\neg x_1, \neg x_5\}$$

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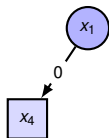
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$$\omega_7 = \{x_1, \neg x_4\}$$



# Weakly active symmetries

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When  $\omega$  is satisfied in all satisfying assignments of  $\varphi$ , we say that  $\omega$  is a logical consequence of  $\varphi$ , and we denote this by  $\varphi \vdash \omega$ .

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## Symmetry propagation

Let  $\sigma$  a weakly active symmetry, then

$$\varphi \cup \alpha \vdash \{I\} \Leftrightarrow \varphi \cup \alpha \vdash \sigma.\{I\}$$

# Local symmetries

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## Local Symmetries

Let  $\varphi$  be a formula. We define  $L_{\omega, \varphi}$ , the set of *local symmetries* for a clause  $\omega$ , and with respect to a formula  $\varphi$ , as follows:

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We can state that:

$$\bigcap_{\omega \in \varphi} L_{\omega, \varphi} \subseteq G.$$

# Computing local symmetries

Formula can be decomposed as :  $\varphi = \varphi_o \cup \varphi_e \cup \varphi_d$  where

- $\varphi_o$  is the set of the original clauses
- $\varphi_e$  is the set of ESBPs
- $\varphi_d$  is the set of deduced clauses.

## Local symmetries

- $\omega \in \varphi_o, L_{\omega, \varphi} \supseteq G$
- $\omega \in \varphi_e, L_{\omega, \varphi} \supseteq \text{Stab}(\omega) = \{\sigma \in G \mid \omega = \sigma.\omega\}$
- $\omega \in \varphi_d, L_{\omega, \varphi} \supseteq \left( \bigcap_{\omega' \in \varphi_1} L_{\omega', \varphi} \right) \cup \text{Stab}(\omega)$

where  $\varphi_1$  is the set of clauses that derives  $\omega$ .