```
\mathbb{M}(A) \to ((A \to \mathbb{R}) \to \mathbb{R})
                     = LO interpretation
            \llbracket - \rrbracket
                      : \mathbb{M}(A,T)
             m
                     = do\{h; a \leftarrow w; M\}
                                                                                                                                           input to disint
             m
                      = \mathbf{do}\{a \leftarrow \mu; t \leftarrow k \, a; \operatorname{return}(a, t)\}
                                                                                                                                            m to \hat{m} by algebraic manipulation
             \hat{m}
                      : \mathbb{M}(A)
                               {lebesgue,counting,dirac}
                                                                                                                                            in concrete code, \mu is determined by A
                                                                                                                                            (type directed)
               k
                     : A \to \mathbb{M}(T)
                      = the interesting thing
                     = \lambda f. \llbracket \mu \rrbracket (\lambda a. \llbracket k \, a \rrbracket (\lambda t. f(a, t)))
                                                                                                                                            defin of \llbracket \_ \rrbracket. Given \hat{m}, solve for k
          \llbracket \hat{m} \rrbracket
                     = do{a \leftarrow \mu; \langle \triangleleft w \, a \, \bar{M} \, h}
                                           \bar{M} = \lambda h'.\mathbf{do}\{h'; M\}
      where
                                  < \triangleleft____ = constrain outcome ???
Target:
          [\![\hat{m}]\!] = \lambda f. \int_{a=(-\infty,\infty)} [\![k\,a]\!] (\lambda t. f(a,t)) \wedge da
Assuming \mu = \text{lebesgue}
                    = \lambda f. \int_{a=(-\infty,\infty)} \llbracket \langle \triangleleft w \, a \, \bar{M} \, h \rrbracket(f) \wedge da
          \llbracket \hat{m} \rrbracket
                                [\![\bar{M}h']\!] = [\![h']\!][[\![M]\!](f)]
      where
                                                                                                                                                                                                                             (a)
General form
                           \lambda f. \llbracket h \rrbracket \lceil \llbracket w \rrbracket (\lambda a. \llbracket M \rrbracket (f)) \rceil
          \llbracket \hat{m} \rrbracket =
                                 \llbracket h \rrbracket is
                                                    an expression with holes, such that substituting a PL (patently linear) -in-f
      where
                                                     expression gives a PL-in-f expression.
                                                    a \in \text{free-vars}(\llbracket M \rrbracket)
assuming
                                 free-vars(h) \subseteq free-vars(\llbracket w \rrbracket) \cup free-vars(\llbracket M \rrbracket)
Then case split on w:
     (1) w = lebesgue
   [\![\hat{m}]\!](f) = [\![h]\!] \left[ \int_{a=(-\infty,\infty)} [\![M]\!](f) \wedge \mathrm{d}a \right]
                      = \int_{a=(-\infty,\infty)} \llbracket h \rrbracket \llbracket \llbracket M \rrbracket (f) \rrbracket \wedge \mathrm{d}a
                                                                                                                                            because of linearity/Tonelli/h is nice
                                                                                                                                            enough
Thus define \langle \neg | \text{lebesgue } a \, \bar{M} \, h' = \bar{M} h' \text{ where we used } \text{eqn}_{(a)}
     ② w = \operatorname{Normal} \mu \sigma
                                                                                                                                            PDF is \hat{N}_x
    \llbracket \hat{m} \rrbracket (f) \quad = \quad \llbracket h \rrbracket \big[ \int_{a=(-\infty,\infty)} \hat{N}_a \cdot \llbracket M \rrbracket (f) \wedge \mathrm{d}a \big] 
                     = \int_{a=(-\infty,\infty)} \llbracket h \rrbracket \left[ \hat{N}_a \cdot \llbracket M \rrbracket(f) \right] \wedge da
                                                                                                                                            linearity/Tonelli/h nice
                     = \int_{a=(-\infty,\infty)} \bar{M}\Big(\llbracket h \rrbracket \big[ \hat{N}_a \cdot \big[ \big] \big] \Big) \wedge da
                                                                                                                                            ???
```

Thus define $\langle \triangleleft (\text{Normal } \mu \, \sigma) \, a \, \bar{M} \, h' = \bar{M}(h'; \text{factor } \hat{N}_{\mu,\sigma})$

(3) w = counting $[\![\hat{m}]\!](f) = [\![h]\!] [\sum_{a=(-\infty,\infty)} [\![M]\!](f)]$ $= \sum_{a=(-\infty,\infty)} \llbracket h \rrbracket \llbracket M \rrbracket (f) \rrbracket$ linearity/switch/h nice

But we "knew" that the type of w is $\mathbb{M}(\mathbb{Z})$ and we are in $\mathbb{M}(\mathbb{R})$

 $h \vdash d : \mathbb{R}$ (4) $w = \operatorname{Dirac} d$

```
[\hat{m}](f) = [h][M](f)|_{a=d}
```

d is a constant, likely impossible?

d is a var bound in h, see POPL backup slides. case return x. see also outline + proofs (???)case 2:

case 3:
$$d = g x$$
 where x bound in h
 $< < [Dirac $(g x)] a c (h_1; x \leftarrow w_1; h_2) = < < w_1 (g^{-1} a) ? ?$$

worry: g^{-1} may involve (????)

$$\textcircled{5}$$
 $w = \operatorname{return}|e|$

$$= \int_{b=(-\infty,\infty)} [\![\langle \neg \langle \text{return } e \rangle b \ M |_{a=|b|} \ h]\!](f) \wedge db$$
 by induction

$$= \int_{b=(-\infty,\infty)} \left[\!\!\left[< \triangleleft \left(\text{return } e \right) b \, \bar{M} \, h \right]\!\!\right] (f) \Big|_{a=|b|} \, \wedge \, \mathrm{d}b \qquad \text{substitution commutative}$$

=
$$\int_{b=(-\infty,0)} [\![\langle \neg (\text{return } e) \, b \, \bar{M} \, h]\!](f) |_{a=-b} \wedge db$$

+
$$\int_{b=(0, \infty)} [\langle \langle (\text{return } e) \, b \, \bar{M} \, h]] (f) \Big|_{a=b} \wedge db$$

$$= \int_{a=(0,\infty)} \left[\!\!\left[< \triangleleft \left(\text{return } e \right) b \, \bar{M} \, h \right]\!\!\right] (f) \Big|_{b=-a} \wedge \mathrm{d}a \qquad \text{change of variables from } b \text{ to } -a$$

+
$$\int_{b=(0,\infty)} \left[\langle \langle (\text{return } e) \, b \, \bar{M} \, h \right] (f) \Big|_{a=b} \wedge db$$

$$= \int_{a=(0,\infty)} [\![\langle \neg (\text{return } e) - a \, \bar{M} \, h]\!](f) \wedge da$$

+
$$\int_{a=(0,\infty)} \llbracket < \lhd (\text{return } e) \quad a \, \bar{M} \, h \rrbracket (f) \wedge da$$
 \hat{M} is b-free

$$= \int_{a=(0,\infty)} \llbracket \cdots \oplus \cdots \rrbracket (f) \wedge da$$

$$= \int_{a=(-\infty,\infty)} \mathbb{1}_{[0,\infty)}(a) \cdot \llbracket \cdots \rrbracket(f) \wedge da$$

=
$$\int_{a=(-\infty,\infty)} [[guard [0,\infty) a; \cdots]](f) \wedge da$$

And now onto $\frac{d}{da}!$ Definition 2. Let:

= $\mathbf{do}\{x \leftarrow \mu; y \leftarrow k x; \operatorname{return}(x, y)\}$ where $\mu = \text{lebesgue}$

Show that: