

## Activity # 3

Key: 00111010, 01010101, 1100001  
00111100, 10100000, 00011110  
10111111, 00001011

Plain Text: 11010000 11110000 10101010  
11001010 11100001 10111101  
10101111 11110010

PC - 1

|    |    |    |    |    |    |    |
|----|----|----|----|----|----|----|
| 57 | 49 | 41 | 33 | 25 | 17 | 9  |
| 1  | 38 | 30 | 42 | 34 | 26 | 18 |
| 10 | 2  | 59 | 51 | 43 | 35 | 27 |
| 19 | 44 | 3  | 60 | 52 | 44 | 36 |
| 63 | 55 | 47 | 39 | 31 | 23 | 15 |
| 7  | 62 | 54 | 46 | 38 | 30 | 22 |
| 14 | 6  | 61 | 53 | 45 | 37 | 29 |
| 21 | 13 | 5  | 28 | 20 | 12 | 4  |

We use above table PC-1  
to convert 64 bit key  
into 56 bit.

Permute key is  
following:

$k^+$  = 0101010 0000001 1001011  
1010110 1110000 1011010  
1011101 0011011

For  $C_0$  &  $D_0$  divide  
 $k^+$  into half.

$C_0$  = 0101010 0000001 1001011  
1010110

$D_0$  = 11100000 1011010 1011101  
0011011

| Iteration No | No. of left shift |
|--------------|-------------------|
| 1            | 1                 |
| 2            | 1                 |
| 3            | 2                 |
| 4            | 2                 |
| 5            | 2                 |
| 6            | 2                 |
| 7            | 2                 |
| 8            | 2                 |
| 9            | 1                 |
| 10           | 2                 |
| 11           | 2                 |
| 12           | 2                 |
| 13           | 2                 |
| 14           | 2                 |
| 15           | 2                 |
| 16           | 1                 |

We will generate keys

$$C_2 = 010101000000110010111010110$$

$$D_2 = 1110000101101010111010011011$$

$$C_1 = 1010100000001100101110101100$$

$$D_1 = 1100001011010101110100110111$$

$$C_2 = ?$$

$$D_2 = ?$$

$$C_{15} = ?$$

$$D_{15} = ?$$

Now will convert 56 bit  
into 48 bit by  
using following PC-2  
table.

## PC - 2

|    |    |    |    |    |    |
|----|----|----|----|----|----|
| 14 | 17 | 11 | 24 | 1  | 5  |
| 3  | 28 | 15 | 6  | 21 | 10 |
| 23 | 19 | 12 | 4  | 26 | 8  |
| 16 | 7  | 27 | 20 | 13 | 2  |
| 41 | 52 | 31 | 37 | 47 | 55 |
| 30 | 40 | 51 | 45 | 33 | 48 |
| 44 | 49 | 39 | 56 | 34 | 53 |
| 46 | 42 | 50 | 38 | 29 | 32 |

For  $C_1 D_1$ ,

|         |        |        |        |
|---------|--------|--------|--------|
| $k_1 =$ | 110011 | 100010 | 110010 |
|         | 000110 | 010101 | 111101 |
|         | 100100 | 110010 |        |

Similarly generate keys for

$C_2 D_2 = k_2$  till  $C_{16} D_{16} = k_{16}$

## Initial Round:

Now we will take plain text convert it into another by using following table IP

IP

|    |    |    |    |    |    |    |   |
|----|----|----|----|----|----|----|---|
| 58 | 50 | 42 | 34 | 26 | 18 | 10 | 2 |
| 60 | 52 | 44 | 36 | 28 | 20 | 12 | 4 |
| 62 | 54 | 46 | 38 | 30 | 22 | 14 | 6 |
| 64 | 56 | 48 | 40 | 32 | 24 | 16 | 8 |
| 57 | 49 | 41 | 33 | 25 | 17 | 9  | 1 |
| 59 | 51 | 43 | 35 | 27 | 19 | 11 | 3 |
| 61 | 53 | 45 | 37 | 29 | 21 | 13 | 5 |
| 63 | 55 | 47 | 39 | 31 | 23 | 15 | 7 |

plain text = 11010000 11100000

~~IP~~ = 10101010 11001010 11100001

10111101 10101111 11110010

IP = 10011011 10100011 01100000  
01110000 11111111 11110110  
01101100 11001100

Divide IP into two  
half  $L_0$  and  $R_0$

$$L_0 = 10011011 \quad 10100011 \quad 01100000 \\ 01110000$$

$$R_0 = 11111111 \quad 1110110 \quad 01101100 \\ 11001100$$

Round 1

$$L_n = R_{n-1}$$

$$R_n = L_{n-1} + f(R_{n-1}, K_n)$$

put  $n=1$

$$L_1 = R_0$$

$$R_1 = L_0 + f(R_0, K_1)$$

$(R_o, K_1)$

$K_1$  is 18 bit &  $R_o$  is 32 bit  
first we convert  $R_o$  into  
8 bit, using following  
E-bit selection Table.

E-bit Selection Table

|    |    |    |    |    |    |
|----|----|----|----|----|----|
| 32 | 1  | 2  | 3  | 4  | 5  |
| 4  | 5  | 6  | 7  | 8  | 9  |
| 8  | 9  | 10 | 4  | 12 | 13 |
| 12 | 13 | 14 | 15 | 16 | 17 |
| 16 | 17 | 18 | 19 | 20 | 21 |
| 20 | 21 | 22 | 23 | 24 | 25 |
| 24 | 25 | 26 | 27 | 28 | 29 |
| 28 | 29 | 30 | 31 | 32 | 1  |

$R_o$ , 1111111 11110110 01101100  
11001100

$E(R_o)$  = 011111 1111H 111110 101100  
001101 011001 011001 011001

$$K_1 \oplus E(R_0) = \begin{array}{cc} 101100 & 011101 \\ 001100 & 101010 \\ 011000 & 100100 \\ 111101 & 101011 \end{array}$$

No we will convert this into 32 bit again  
 We divide 48 bit into 6 bit chunk of 8 and named them  $B_1, B_2, \dots, B_8$

$$S_1(B_1), S_2(B_2), \dots, S_8(B_8)$$

$$\begin{aligned} S_1 &= 101100 & 10 \Rightarrow 2 \text{ col} \\ &= 0010 & 0110 = 6 \text{ Row} \end{aligned}$$

$$\begin{aligned} S_2 &= 011101 & 01 \Rightarrow 1 \text{ col} \\ &= 1011 & 1110 \Rightarrow 14 \text{ Row} \end{aligned}$$

$$\begin{aligned} S_3 &= 001100 & 00 \Rightarrow 0 \text{ col} \\ &= 1111 & 0110 \Rightarrow 6 \text{ Row} \end{aligned}$$

$$S_4 = \begin{matrix} 1 & 0 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 1 \end{matrix}$$

$10 \Rightarrow 2$  Col  
 $0101 \Rightarrow 5$  Row

$$S_5 = \begin{matrix} 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 \end{matrix}$$

$00 \Rightarrow 0$  Col  
 $1100 \Rightarrow 12$  Row

$$S_6 = \begin{matrix} 1 & 0 & 0 & 1 & 0 & 0 \\ 1 & 1 & 1 & 1 \end{matrix}$$

$10 \Rightarrow 2$  Col  
 $0010 \Rightarrow 2$  Row

$$S_7 = \begin{matrix} 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \end{matrix}$$

$11 \Rightarrow 3$  Col  
 $110 \Rightarrow 14$  Row

$$S_8 = \begin{matrix} 1 & 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 \end{matrix}$$

$11 \Rightarrow 3$  Col  
 $0101 \Rightarrow 5$  Row

$$= \begin{matrix} 0 & 0 & 1 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 1 & 1 & 0 & 1 \end{matrix}$$
  
$$\begin{matrix} 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 1 & 0 & 1 & 0 \end{matrix}$$

Now we will perform  
the permutation again

P

|    |    |    |    |
|----|----|----|----|
| 16 | 7  | 20 | 21 |
| 29 | 12 | 28 | 17 |
| 1  | 15 | 23 | 26 |
| 5  | 18 | 31 | 10 |
| 2  | 8  | 24 | 14 |
| 32 | 27 | 3  | 9  |
| 19 | 13 | 30 | 6  |
| 22 | 11 | 4  | 25 |

$f_2 = 1111 \quad 1111 \quad 0110 \quad 1111 \quad 0110$   
 $0111 \quad 0100 \quad 1100$

~~R<sub>1</sub> = f~~

$$R_1 = R_0 \oplus f(R_0, k_1)$$

$$\begin{array}{r} 100110111010001101000 \\ 00011100000 \oplus 111111110110 \\ 11110110011101001100 \end{array}$$

$$\begin{array}{r} R_1 = 011001001100110000000 \\ 11100111100 \end{array}$$

$L_1 = 1111 \ 1111 \ 1111 \ 0110 \ 0011 \ 01100$

$1100 \ 1100$

Round  $\Rightarrow 2$

$L_2 = R_1$

$R_2 = L_1 + f(R_1, K_2)$

$K_2 = 000100 \ 01000 \ 011100$

$100111 \ 100111 \ 001100$

$10111 \ 001110$

$R_1 = 01100100110000110000$

$1110011100$

$E(R_1) = 001100 \ 001001 \ 011001$

$011000 \ 000000 \ 001100$

$100111 \ 111000$

$$K_2 \oplus L(R_1) \rightarrow \begin{matrix} 001000 & 011001 & 000101 \\ 111111 & 100111 & 000010 \\ 001000 & 110110 & \end{matrix}$$

By applying  $S_d(B_1)$  ...  $S_8(B_2)$

We have

$$\begin{matrix} 0010 & 011000 & 0001100111 \\ 00011111101 & \end{matrix}$$

Now we perform permutation,

$$F_2 = \begin{matrix} 0110 & 1010 & 0101 & 0100 \\ 0011 & 1110 & 1111 & 0001 \end{matrix}$$

$$R_2 = L \oplus F(R_1, K_2)$$

$$L_2 = \begin{matrix} 0110010011001100000000 \\ 11100111100 \end{matrix}$$

$$R_2 = \begin{matrix} 10010101101000100101001 \\ 000111101 \end{matrix}$$

Now Reverse Order.

$R_2 L_2 =$  10010101101000100101  
0010000111010110010011  
00000000 11100111100

Inverse permutation.

Use the following table

|    |    |    |    |    |    |    |    |
|----|----|----|----|----|----|----|----|
| 40 | 8  | 48 | 16 | 52 | 24 | 64 | 32 |
| 39 | 7  | 47 | 15 | 53 | 23 | 63 | 31 |
| 38 | 6  | 46 | 14 | 54 | 22 | 62 | 30 |
| 37 | 5  | 45 | 13 | 53 | 21 | 61 | 29 |
| 36 | 31 | 49 | 12 | 52 | 20 | 60 | 28 |
| 35 | 3  | 43 | 11 | 51 | 19 | 59 | 27 |
| 34 | 2  | 42 | 10 | 50 | 18 | 58 | 26 |
| 33 | 1  | 41 | 9  | 49 | 17 | 57 | 25 |

$IP^{-1} =$  01001001 00011100 11101011  
00100011 01000111 10010011  
10100100 01110000

Result. 2 73 28 235 35 71  
147 112