1)											
16/1											1
A-104											
16/1	////////	/////////	16/1						2 (1		
A-104		,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,	B-120								
16/1	////////	/////////	16/1		16/1	///////					
A-104			B-120		C-136						
16/1	////////	//////////	16/1		16/1	///////	32/1			//////////////////////////////////////	////////
A-104			B-120		C-136	İ	D-152				
16/1	////////	/////////	16/0		16/1	////////	32/1			///////////////////////////////////////	////////
A-104			B-120		C-136		D-152				
32/0					16/1	///////	32/1			///////////////////////////////////////	////////
104					C-136		D-152	A			1
32/0				V C	16/1	////////	32/0				
104					C-136		152				
Best fit and first	t fit are both the	same:									
8 slash 1	24/0				16/1	///////	32/0				I I
E-100	108				C-136						1

////////	////////	10 slash 1								
		B-112								
////////	/////////	10 slash 1		////////						
		B-112								
////////	////////	10 slash 1		////////	10 slash 1	///////				
		B-112			C-122					
////////	////////	10 slash 1		////////	10 slash 1	///////	////////	10 slash 1		
		B-112		-	C-122			D-132		
////////	/////////	10 slash 1		////////	10 slash 1	///////	///////	10 slash 0		
		B-112			C-122			132		
////////	/////////	10 slash 0			10 slash 1	///////	////////	10 slash 0		
		112		A TABLE TO SERVICE	C-122			132		
e the same:										
////////	/////////	6 slash 1	4 slash 0		10 slash 1	///////	////////	10 slash 0		ľ
		E-112	F-118		C-122			132		
////////	////////	6 slash 1	4 slash 0		10 slash 0			10 slash 0		
		E-112	F-118		122			132		
////////	////////	6 slash 1	4 slash 0		10 slash 1		///////	10 slash 0		
		E-112	F-118		122		111111111111111111111111111111111111111	132		
	//////////////////////////////////////		//////// /////////////////////////////			B-112				

4)													
p	////////	////////											
A-103					4								
24/1	////////	///////	12 slash 1	///////////////////////////////////////	////////	+							
A-103			B-127		11111111								
24/1	////////	////////	12 slash 1	///////////////////////////////////////	////////	24 slash 1		////////	/ ///////	////////	1		
A-103			B-127			C-139							
24/1	////////	////////	12 slash 1	///////////////////////////////////////	////////	24 slash 1		////////	////////	////////	12 slash 1	////////	///////
A-103	(5337211)		B-127	2000		C-139					D-163	10000000000	
First fit and best fit are the same:												7	8
24/1	////////	////////	12 slash 0			24 slash 1		////////	/////////	///////////////////////////////////////	12 slash 1	////////	/////////
A-103			127			C-139					D-163		
24/1		////////	6 slash 1	6 slash 0		24 slash 1		////////	/////////	////////	12 slash 1	////////	///////
A-103			E-127	133		C-139					D-163		
24/1	////////	////////	6 slash 1	6 slash 0		24 slash 1		////////	/ ////////	////////	12 slash 0		
A-103			E-127	133		C-139					163		
24/1	////////	////////	6 slash 1	42/0									
A-103	[1377]	-	E-127	133									
24/1	////////	////////	6 slash 1	15/1			27/0					*	
A-103		, ,	E-127	F-133			148						

- 1. Expanding data inside an allocated segment to fill an internal segment is possible, but since this internal fragmentation is small it is unlikely to be big enough to store any new memory, so it would likely just results in slower memory allocation.
- 2. It seems like the ideal size for blocks of memory to be allocated is as small as possible so that there is no internal fragmentation because you can allocate new memory to precisely the amount of memory needed
- 3. In some cases, you may be able to use program logic to put small pieces of data into the internal fragments.
- 4. Explicit free lists make it faster to allocate new memory as you only have to search through all the free blocks to find memory big enough to place the new data rather than searching through all blocks.