





# Advanced C++ for HPC: Making a SIMD library

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### **Outline**

# Why?

- Vectorization.
- SIMD intrinsics.
- Alternatives.

### How?

- Design constraints.
- Representing ABIs: type maps.
- CRTP for fun and profit.
- Front-end interface.





#### CPU SIMD instructions: do more with less.

```
g++-9.2 -fno-unroll-loops -03 -mno-sse

fldl 8(%rsp)
    xorl %eax, %eax
.L2:

fldl (%rsi,%rax)
    fmul %st(1), %st
    faddl (%rdx,%rax)
    fstpl (%rdi,%rax)
    addq $8, %rax
    cmpq $262144, %rax
    jne .L2
    fstp %st(0)
    Ret
```



#### CPU SIMD instructions: do more with less.

```
g++-9.2 -03 -ffp-contract=off -march=skylake-avx512

vbroadcastsd %xmm0, %ymm1
   xorl %eax, %eax
.L2:

vmulpd (%rsi,%rax), %ymm1, %ymm0
  vaddpd (%rdx,%rax), %ymm0, %ymm0
  vmovupd %ymm0, (%rdi,%rax)
  addq $32, %rax
  cmpq $262144, %rax
  jne .L2
  vzeroupper
  ret
```





#### CPU SIMD instructions: do more with less.

vbroadcast<mark>s</mark>d — fill target register with single double-precision value.

xorl %eax, %eax — set eax/rax to zero.

vmulpd — multiply packed SIMD double-precision values.

(%rsi, %rax) — indirect/indexed addressing: contents of address rsi+rax

xmmn: first 128 bits of nth SIMD register

ymmn: first 256 bits of nth SIMD register [AVX] zmmn: all 512 bits of nth SIMD register [AVX512]

```
vbroadcastsd %xmm0, %ymm1
xorl %eax, %eax

.L2:

vmulpd (%rsi,%rax), %ymm1, %ymm0
vaddpd (%rdx,%rax), %ymm0, %ymm0
vmovupd %ymm0, (%rdi,%rax)
addq $32, %rax
cmpq $262144, %rax
jne .L2
vzeroupper
ret
```





### Wait, why is this only 4-wide?

```
#define N 32768
                                                              q++-9.2 -03 -ffp-contract=off -march=skylake-avx512
// c = k*a + b
                                                                 vbroadcastsd %xmm0, %ymm1
void sma(double* __restrict c, double k,
                                                                 xorl %eax, %eax
         const double* a, const double* b)
                                                            .L2:
                                                                 vmulpd (%rsi,%rax), %ymm1, %ymm0
      for (unsigned i = 0; i < N; ++i) {
                                                                  vaddpd (%rdx,%rax), %ymm0, %ymm0
            c[i] = k*a[i] + b[i];
                                                                 vmovupd %ymm0, (%rdi,%rax)
                                                                                                256 bits: 4 x double
                                                                  addq $32, %rax
                                                                  cmpq $262144, %rax
                                                                  jne .L2
                                                                 vzeroupper
                                                                 ret
```



32 bytes per iteration: 4 x double

### How about clang?

```
clang++-9 -03 -fno-unroll-loops -march=skylake-avx512

vbroadcastsd %xmm0, %zmm0
    xorl %eax, %eax
.LBB0_1:
    vmulpd (%rsi,%rax,8), %zmm0, %zmm1
    vaddpd (%rdx,%rax,8), %zmm1, %zmm1
    vmovupd %zmm1, (%rdi,%rax,8)
    addq $8, %rax
    cmpq $32768, %rax
    jne .LBB0_1
    vzeroupper
    retq
```



## **Auto-vectorization is tricksy**

Compiler support in gcc and clang improving every release, but:

- No guarantees on what you'll get.
- 2. Might perform unsafe transformations without being asked (gcc, icc 😠).



### Sometimes very clever...

#### clang++-9 -03 -march=skylake-avx512

```
vxorpd %xmm0, %xmm0, %xmm0
       xorl %eax, %eax
.LBB0 1:
       vmovsd (%rdi,%rax,8), %xmm1
       vmovsd 8(%rdi,%rax,8), %xmm2
       vsubsd (%rsi,%rax,8), %xmm1, %xmm1
       vmulsd (%rdx,%rax,8), %xmm1, %xmm1
       vsubsd 8(%rsi,%rax,8), %xmm2, %xmm2
       vmulsd 8(%rdx,%rax,8), %xmm2, %xmm2
       vaddsd %xmm1, %xmm0, %xmm0
       vaddsd %xmm2, %xmm0, %xmm0
       vmovsd 16(%rdi,%rax,8), %xmm1
       vsubsd 16(%rsi,%rax,8), %xmm1, %xmm1
       vmulsd 16(%rdx,%rax,8), %xmm1, %xmm1
       vmovsd 24(%rdi,%rax,8), %xmm2
       vsubsd 24(%rsi,%rax,8), %xmm2, %xmm2
       vaddsd %xmm1, %xmm0, %xmm0
       vmulsd 24(%rdx,%rax,8), %xmm2, %xmm1
       vaddsd %xmm1, %xmm0, %xmm0
       addq $4, %rax
       cmpq $32768, %rax
       ine .LBB0 1
        reta
```





### Sometimes very clever... if given enough leeway.

```
clang++-9 -03 -ffast-math -march=skylake-avx512
       vxorpd %xmm0, %xmm0, %xmm0
       xorl %eax. %eax
       vxorpd %xmm1, %xmm1, %xmm1
       vxorpd %xmm2, %xmm2, %xmm2
      vxorpd %xmm3, %xmm3, %xmm3
.LBB0 1:
                                                   4x unroll
       vmovupd (%rdi,%rax,8), %zmm4
       vmovupd 64(%rdi,%rax,8), %zmm5
                                                8-wide fma
       vmovupd 128(%rdi,%rax,8), %zmm6
       vmovupd 192(%rdi,%rax,8), %zmm7
       vsubpd (%rsi,%rax,8), %zmm4, %zmm4
       vsubpd 64(%rsi,%rax,8), %zmm5, %zmm5
       vsubpd 128(%rsi,%rax,8), %zmm6, %zmm6
       vsubpd 192(%rsi,%rax,8), %zmm7, %zmm7
       vfmadd231pd (%rdx,%rax,8), %zmm4, %zmm0
       vfmadd231pd 64(%rdx,%rax,8), %zmm5, %zmm1
       vfmadd231pd 128(%rdx,%rax,8), %zmm6, %zmm2
       vfmadd231pd 192(%rdx,%rax,8), %zmm7, %zmm3
       addq $32, %rax
       cmpq $32768, %rax
       jne .LBB0 1
```

```
vaddpd %zmm0, %zmm1, %zmm0
vaddpd %zmm0, %zmm2, %zmm0
vaddpd %zmm0, %zmm3, %zmm0
vextractf64x4 $1, %zmm0, %ymm1
vaddpd %zmm1, %zmm0, %zmm0
vextractf128 $1, %ymm0, %xmm1
vaddpd %xmm1, %xmm0, %xmm0
vpermilpd $1, %xmm0, %xmm0
vzeroupper
retq
```



#### Transcendentals?

```
#include <cmath>
#define N 32768

// c = exp(a)
void vexp(double* __restrict c, const double* a)
{
    for (unsigned i = 0; i<N; ++i) {
        c[i] = std::exp(a[i]);
    }
}</pre>
```

```
clang++-9 -03 -fno-unroll-loops -ffast-math
            -march=skylake-avx512
# inside the loop:
vmovups (%r15,%rbx,8), %xmm0
vmovaps %xmm0, 48(%rsp)
vmovups 16(%r15,%rbx,8), %xmm0
vmovaps %xmm0, 16(%rsp)
vmovups 32(%r15,%rbx,8), %xmm0
vmovaps %xmm0, 64(%rsp)
vmovups 48(%r15,%rbx,8), %xmm0
vmovaps %xmm0, 32(%rsp)
vzeroupper
callq exp finite
vmovaps %xmm0, (%rsp)
vpermilpd $1, 32(%rsp), %xmm0
callq exp finite
# [ ... and another 6 calls to exp finite ]
# [ ... then pack everything into zmm0 and store ]
```



### **Auto-vectorization is tricksy**

Compiler support in gcc and clang improving every release, but:

- No guarantees on what you'll get.
- 2. Might perform unsafe transformations without being asked (gcc, icc 😠).
- 3. Often need to permit unsafe assumptions with -ffast-math anyway.
- 4. Even with -ffast-math, transcendentals may be serialized.



#### Scatter conflicts

```
clang++-9 -03 -fno-unroll-loops -ffast-math
                -march=skyLake-avx512
     xorl %eax, %eax
.LBB0 1:
     vmovupd (%rcx,%rax,8), %zmm0
     vmulpd (%rdx,%rax,8), %zmm0, %zmm0
     vpmovzxdq (%rsi,%rax,4), %zmm1
     kxnorw %k0, %k0, %k1
     vscatterqpd %zmm0, (%rdi,%zmm1,8) {%k1}
     addq $8, %rax
     cmpq $32768, %rax
     jne .LBB0 1
     vzeroupper
     retq
```





#### Scatter conflicts

Potential conflict: e.g. if index[2]==index[4]

```
clang++-9 -03 -fno-unroll-loops -ffast-math
                 -march=skyLake-avx512
      xorl %eax, %eax
.LBB0 1:
      vmovsd (%rdx,%rax,8), %xmm0
      vmovsd (%rcx,%rax,8), %xmm1
      mov1 (%rsi,%rax,4), %r8d
      vfmadd213sd (%rdi,%r8,8), %xmm0, %xmm1
      vmovsd %xmm1, (%rdi,%r8,8)
      incq %rax
      cmpq $32768, %rax
      jne .LBB0 1
      retq
          With -Rpass-analysis=loop-vectorize
      <source>:7:5: remark: loop not vectorized: unsafe
      dependent memory operations in loop. [...]
```





## **Auto-vectorization is tricksy**

Compiler support in gcc and clang improving every release, but:

- No guarantees on what you'll get.
- 2. Might perform unsafe transformations without being asked (gcc, icc 😠).
- 3. Often need to permit unsafe assumptions with -ffast-math anyway.
- 4. Even with -ffast-math, transcendentals may be serialized.
- 5. Can't express knowledge of potential conflicts or lack of conflicts in arithmetic syntax.



### **SIMD Intrinsics**

Compilers generally support a number of *intrinsic* built-in functions.

- Provide an interface to non-standard functionality.
- Compiler specific.

SIMD operations for a particular CPU architecture are provided by intrinsics in gcc, clang, and icc, and are (mostly) the same for each compiler.

- Most correspond to a single CPU SIMD operation.
- They aren't necessarily compiled exactly to that CPU SIMD operation.



### **SIMD Intrinsics**

#### Auto-vectorized

```
clang++-9 -03 -fno-unroll-loops -march=skylake-avx512
    vbroadcastsd %xmm0, %zmm1
    xorl %eax, %eax
.L2:
    vmovupd (%rsi,%rax), %zmm0
    vfmadd213pd (%rdx,%rax), %zmm1, %zmm0
    vmovupd %zmm0, (%rdi,%rax)
    addq $64, %rax
    cmpq $262144, %rax
    jne .L2
    vzeroupper
    ret
```





#### **SIMD Intrinsics**

### Explicit vectorization with intrinsics

clang++-9 -03 -fno-unroll-loops -march=skylake-avx512

```
vbroadcastsd %xmm0, %zmm1
    xorl %eax, %eax
.L2:

vmovupd (%rsi,%rax), %zmm0
    vfmadd213pd (%rdx,%rax), %zmm1, %zmm0
    vmovupd %zmm0, (%rdi,%rax)
    addq $64, %rax
    cmpq $262144, %rax
    jne .L2
    vzeroupper
    retq
```

Exactly the same generated code.





## Using SIMD intrinsics vs auto-vectorization

#### The Good

- Generated code will be close to what you write.
- Free to support full IEEE semantics or make numerical assumptions.
- Can supply own vectorized implementations of transcendental functions.
- Can be explicit about conflict assumptions in indirect memory operations.

#### The Bad

- Different code for every architecture: huge maintenance burden.
- Very, very non-standard C++.

### The Ugly

```
    a = _mm512_castsi512_pd(_mm512_and_epi64(_mm512_castpd_si512(x),
        _mm512_set1_epi64(0x7fffffffffffffff)))
```





#### **Alternatives**

- OpenMP SIMD extensions
- 2. The Vc library: <a href="mailto:github.com/VcDevel/Vc">github.com/VcDevel/Vc</a>
- 3. std::experimental::simd coming soon to a libstdc++ near you.

Proposed SIMD extensions described in working draft N4808 are based on a subset of the Vc library.

Or we can write our own.



# Write your own SIMD library, what are you crazy?

#### Excuses:

- Vc didn't do everything we wanted in our project.
- We really like writing ratpoly transcendental function approximations.\*
- It turned out to be useful as an example of C++ patterns.





<sup>\*</sup> This might be a slight overstatement.

- 1. Broadly follow the API given in the N4808 proposal, or a simpler version.
  - Lower burden for migration to a future standard library implementation.
  - Benefit from existing API design work.



Broadly follow the API given in the N4808 proposal, or a simpler version.

```
#include <experimental/simd>
#define N 32768
using double4 = std::experimental::fixed size simd<double, 4>;
using std::experimental::element aligned;
void sma(double* restrict c, double k, const double* a, const double* b) {
      for (unsigned i = 0; i < N; i += 4) {
             double4 va(a+i, element aligned);
             double4 vb(b+i, element aligned);
             auto vc = k*va + vb;
             vc.copy to(c+i, element aligned);
```





### You can run this code on Compiler Explorer right now.

```
gcc-9.2 -03 -march=skylake-avx512 -ffast-math -std=c++17
```

```
#include <experimental/simd><https://raw.githubusercontent.com/VcDevel/std-simd/compiler explorer/simd.h>
#define N 32768
using double4 = std::experimental::fixed size simd<double, 4>;
using std::experimental::element aligned;
void sma(double* restrict c, double k, const double* a, const double* b) {
      for (unsigned i = 0; i < N; i += 4) {
             double4 va(a+i, element aligned);
             double4 vb(b+i, element aligned);
             auto vc = k*va + vb;
             vc.copy to(c+i, element aligned);
```





- 1. Broadly follow the API given in the N4808 proposal, or a simpler version.
- Separate user-visible API from architecture-specific implementations.
  - Many SIMD back-ends can share a common front-end that handles all the syntactic sugar.
  - Reduces the development cost of writing a new back-end.



- 1. Broadly follow the API given in the N4808 proposal, or a simpler version.
- Separate user-visible API from architecture-specific implementations.

```
// user code
double4 a = b + c;
double* x = ...;
int4 index = ...;
indirect(x, index) = a;
```

```
// equivalent back-end code
double4_impl::vector_type A =
double4_impl::add(B, C);
double4_impl::scatter(tag<int4_impl
>{}, A, x, I);
```



- 1. Broadly follow the API given in the N4808 proposal, or a simpler version.
- 2. Separate user-visible API from architecture-specific implementations.
- 3. Decouple back-end functionality from any specific SIMD data representation.
  - Operations on the same data representation (e.g. \_\_m256) may have different implementations for different SIMD ISAs that are supported on the same platform (e.g. AVX2, AVX512).
  - Back-ends then should provide functions operating on SIMD data, rather than wrapping SIMD data and providing operations as methods.

- 1. Broadly follow the API given in the N4808 proposal, or a simpler version.
- 2. Separate user-visible API from architecture-specific implementations.
- 3. Decouple back-end functionality from any specific SIMD data representation.
- Supply fallback implementations for operations that are unimplemented or unsupported by the architecture.
  - If we can always move to and from an array representation, we can write generic fallback code, relying on the optimizer to remove redundant copying.





### Let's make a prototype!

- Two back-ends: AVX2 and a generic std::array-based back-end.
- Arithmetic operations: + and \* for SIMD double and int values.
- Store to/from memory.
- Element (lane) access.
- Indirect memory operations (gather/scatter).



Back-end design: ABI mapping

How will we represent a particular SIMD data-type and set of operations?

- N4808: an ABI tag describes a mapping from a data-parallel type to a particular width and binary representation.
- Simplify for tinysimd: make sizes explicit, and use an ABI tag to select a SIMD ISA implementation, too.
- Two explicit choices: avx2 and generic.





### Back-end design: ABI mapping

#### Convention:

- abi::abitag<T, N>::type is the abitag implementation type for N-wide vector of T, or void if unavailable.
- abi::generic<T, N>::type will map to a valid implementation based on std::array for every T and N.
- abi::default\_abi<T, N>::type will map to the 'best' available implementation for T and N.

```
#include <tinysimd/generic.h>
#include <tinysimd/avx2.h>
namespace tinysimd {
template <typename...>
struct first not void of; // (details elided)
namespace abi {
    template <typename T, unsigned N>
    struct default abi {
        using type = typename first not void of<
            typename avx2<T, N>::type,
            typename generic<T, N>::type
        >::type;
   };
```





### Back-end design: implementation API

What will an implementation class look like? Need a class which provides:

- A type for the SIMD representation can use a traits class.
- Functions operating on this representation these can be static member functions.

### What functions do we need for tinysimd?

- Copy to/from memory
- Lane access
- Scalar to vector

- Arithmetic: add, multiply
- Indirect reads (gather)
- Indirect writes (scatter)





### Back-end design: fallback implementations

If an implementation provides a traits class instance, and provides a copy\_to and copy\_from function for writing to and reading from memory, a fallback implementation class can perform every other operation with everyday C++.

```
// in tinysimd/fallback.h

template <typename I> struct simd_traits {
    static constexpr unsigned width = 0;
    using scalar_type = void;
    using vector_type = void;
};
```

```
// in tinysimd/avx2.h

struct avx2_double4;
template <> struct simd_traits<avx2_double4> {
    static constexpr unsigned width = 4;
    using scalar_type = std::int32_t;
    using vector_type = __m256d;
};
```





```
template <typename I> struct fallback {
      static constexpr unsigned width = simd traits<I>::width;
      using scalar type = typename simd traits<I>::scalar type;
      using vector_type = typename simd_traits<I>::vector_type;
      using store = scalar type[width];
      static vector type broadcast(scalar type x) {
             store a;
             std::fill(std::begin(a), std::end(a), x);
             return I::copy from(a);
      }
      static vector type add(vector type u, const vector type v) {
             store a, b, r;
             I::copy_to(u, a);
             I::copy to(v, b);
             for (unsigned i = 0; i<width; ++i) r[i] = a[i]+b[i];</pre>
             return I::copy from(r);
      // similarly: element(), set element(), mul() ...
};
```



### Back-end design: fallback implementations

An actual SIMD implementation class can use intrinsics, or call the fallback implementation. But:

- Forwarding methods add a lot of boilerplate.
- Adding a new back-end method requires adding this method to every implementation class.

Solution: just inherit from the fallback class!

```
struct avx2 double4 {
    static void copy_to(__m256d v, double* p) {
        mm256 storeu pd(p, v);
    static m256d copy_from(const double* p) {
        return _mm256_loadu_pd(p);
    static m256d broadcast(double v) {
        return _mm256_set1_pd(v);
    static __m256d add(__m256d a, __m256d b) {
        return fallback<avx2_double4>::add(a, b);
    // And mul, etc. ...
};
```





### Back-end design: CRTP

**CRTP** is the Curiously Recurring Template Pattern: a class X derives from a parameterized class B, with parameter X.

- Used for 'static polymorphism' (among other things).
- Often abused.

#### For us:

- Base class provides generic implementations; uses specialized methods from the derived class where it can.
- Missing methods in the derived class automatically are provided by the base.

```
struct avx2_double4: fallback<avx_double4> {
    static void copy_to(__m256d v, double* p) {
         mm256 storeu pd(p, v);
    static m256d copy from(const double* p) {
        return mm256 loadu pd(p);
    static m256d broadcast(double v) {
        return mm256 set1 pd(v);
    // Fallback methods for everything else just inherited.
};
```





#### **Aside: CRTP**

### Perils of CRTP implementation

Compiler won't catch errors that it would with dynamic types and overloads:

- A static method fallback<I>::foo() might call add() (invoking the fallback version) instead of I::add() (the possibly specialized version from I).
- A static method I::bar(...) may have a mismatch in its signature preventing
  it from specializing fallback<i>::bar(...).



#### **Aside: CRTP**

#### CRTP without static methods

Member functions in derived class can be invoked by base class by downcasting this.

- Base class can implement a common functionality based on operations implemented in the derived class reduces code duplication.
- Same sorts of caveats apply as in static case.

```
template <typename D> struct base {
    D* derived() {
        return static cast<D*>(this);
    void foo() {
        derived()->op bar();
        derived()->op baz();
};
struct derived: base<derived> {
    friend struct base<derived>;
private:
    void op bar() { /* ... */ }
    void op baz() { /* ... */ }
};
```





### Back-end design: generic implementation

Generic no-intrinsics back-end: just use fallback routines.

Hopefully compiler will elide all those copies?

- Sometimes it will.
- Sometimes the optimizer gets horribly confused if width>1.

```
struct generic: fallback<generic<T, N>> {
    using vector type = std::array<T, N>;
    static void copy to(vector type v, T* p) {
        std::copy(std::begin(v), std::end(v), p);
    static vector type copy from(const T* p) {
        vector type v;
        std::copy(p, p+N, std::begin(v));
        return v;
};
template <typename T, unsigned N>
struct simd traits<generic<T, N>> {
    static constexpr unsigned width = N;
    using scalar type = T;
    using vector type = std::array<T, N>;
};
```





#### Does it work?

With implementations filled out for avx2\_double4, do we get vectorized code?

```
#include <tinysimd/avx2.h>
#define N 32768
using namespace tinysimd;
using S = avx2 double4;
using vec = simd traits<S>::vector type;
constexpr unsigned width = simd_traits<S>::width;
void sfma(double* __restrict c, double k,
         const double* a, const double* b)
   for (unsigned i = 0; i<N; i += width) {</pre>
        vec va = S::copy from(a+i);
        vec vb = S::copy from(b+i);
        vec vc = S::fma(S::broadcast(k), va, vb);
        S::copy to(vc, c+i);
```

```
clang++-9 -03 -fno-unroll-loops -march=skylake-avx512
sfma(double*, double, double const*):
      vbroadcastsd %xmm0, %ymm0
      xorl %eax, %eax
.LBB0 1:
      vmovupd (%rsi,%rax,8), %ymm1
      vfmadd213pd (%rdx,%rax,8), %ymm0, %ymm1
      vmovupd %ymm1, (%rdi,%rax,8)
      addq $4, %rax
      cmpq $32768, %rax
      jb .LBB0 1
      vzeroupper
      retq
```





Back-end design: heterogeneous operations — gather and scatter

```
Scatter operation (indirect SIMD write): p[index[i]] = v[i]
Gather operation (indirect SIMD read): v[i] = p[index[i]]
```

Heterogeneous operation: SIMD value type and SIMD index type.

SIMD value implementation class knows how to interpret the value vector-type, but how will it know how to interpret the index vector-type?



Back-end design: heterogeneous operations — gather and scatter

A 'tag' type can represent the implementation class for the index vector, without requiring that the implementation class be complete.

```
template <typename X> struct tag {};
```

Back-end can then overload gather for different index implementations, that may share the same index *representation*.

• Example: suppose avx512\_int8 and avx512\_long8 both use a \_\_m512i representation. Then different gather overloads for tag<avx512\_int8> and tag<avx512\_long8> can use the corresponding intrinsic.





### Back-end design: heterogeneous operations — gather and scatter

In the fallback class, use the Index implementation provided by the tag to extract the indicies, and perform the loads one by one:

```
template <typename I> struct tag {};
template <typename I> struct fallback<I> {
    // ...
    template <typename IndexI>
    static vector type gather(tag<IndexI>, const scalar type* restrict p,
                              typename simd_traits<IndexI>::vector_type index) {
        using index store = typename simd_traits<IndexI>::scalar_type[width];
        index store j;
        IndexI::copy to(index, j);
        store a;
        for (unsigned i = 0; i<width; ++i) a[i] = p[j[i]];</pre>
        return I::copy from(a);
};
```





### Back-end design: heterogeneous operations — gather and scatter

In the implementation class:

- Use a using declaration to bring in the templated overloads from fallback.
- Provide specific implementations as required.

Extend similarly for scatter operations (and other heterogeneous operations such as value casting).





### Back-end design: gather demo

```
#include <tinysimd/avx2.h>
#define N 32768
using namespace tinysimd;
                                                                const*):
using d4impl = avx2 double4;
                                                                       xorl %eax. %eax
using d4 = simd_traits<d4impl>::vector_type;
                                                                 .LBB0 1:
using i4impl = avx2 int4;
using i4 = simd traits<i4impl>::vector type;
constexpr unsigned width = 4;
// c[i] = a[index[i]]*b[index[i]]
void mul_indexed(double* __restrict c, const int* index,
                 const double* a, const double* b)
    for (unsigned i = 0; i<N; i += width) {</pre>
        i4 vi = i4impl::copy from(index+i);
                                                                       addq $4, %rax
        d4 va = d4impl::gather(tag<i4impl>{}, a, vi);
                                                                       cmpq $32768, %rax
        d4 vb = d4impl::gather(tag<i4impl>{}, b, vi);
                                                                       jb .LBB0 1
        d4impl::copy to(d4impl::mul(va, vb), c+i);
                                                                       vzeroupper
                                                                       retq
```

```
clang++-9 -03 -fno-unroll-loops -march=skylake-avx512
mul indexed(double*, int const*, double const*, double
      vmovupd (%rsi,%rax,4), %xmm0
      vpcmpeqd %ymm1, %ymm1, %ymm1
      vxorpd %xmm2, %xmm2, %xmm2
      vgatherdpd %ymm1, (%rdx,%xmm0,8), %ymm2
      vpcmpeqd %ymm1, %ymm1, %ymm1
      vxorpd %xmm3, %xmm3, %xmm3
      vgatherdpd %ymm1, (%rcx,%xmm0,8), %ymm3
      vmulpd %ymm3, %ymm2, %ymm0
      vmovupd %ymm0, (%rdi,%rax,8)
```





#### Front-end

### Public API components:

- A wrapper simd\_wrap<I> around the implementation class I, containing a member of type simd\_traits<I>::vector\_type. Provides operator overloads.
- 2. A type map simd<V, N, Abi> to the wrapper simd\_wrap<I> where I is Abi<V, N>::type.
- 3. A function indirect that wraps a pointer and a SIMD index into an indirect expression object used to describe indexed operations.





#### Front-end: simd\_wrap

#### Constructors

```
template <typename I> struct simd_wrap {
private:
   static assert(!std::is void<I>::value, "unsupported");
   using vector type = typename simd traits<I>::vector type;
   vector type value ;
public:
   using scalar type = typename simd traits<I>::scalar type;
   static constexpr unsigned width = simd traits<I>::width;
   simd wrap() = default;
   simd wrap(const simd_wrap& other) = default;
   simd wrap(const scalar type& x):
       value (I::broadcast(x)) {}
   simd wrap(const scalar type (&a)[width]):
       value (I::copy_from(a)) {}
   explicit simd wrap(const scalar type *p):
       value (I::copv from(p)) {}
```

#### Assignment and writes to memory

```
simd_wrap& operator=(const simd_wrap& other) = default;
simd_wrap& operator=(const scalar_type& x) {
    value_ = I::broadcast(x);
    return *this;
}
void copy_to(scalar_type* p) const {
    I::copy_to(value_, p);
}
// ...
```





### Front-end: simd\_wrap arithmetic operations

```
template <typename I> struct simd wrap {
   // ...
private:
    static simd wrap wrap(const vector type& v) {
        simd wrap s;
        s.value = v
       return s;
public:
   friend simd wrap operator+(const simd wrap& a, const simd wrap& b) {
        return wrap(I::add(a.value , b.value ));
   friend simd wrap operator*(const simd_wrap& a, const simd_wrap& b) {
        return wrap(I::mul(a.value , b.value ));
   friend simd wrap fma(const simd wrap& a, const simd wrap& b, const simd wrap& c) {
        return wrap(I::fma(a.value , b.value , c.value ));
    }
    simd wrap& operator+=(const simd wrap& a) { value = I::add(value , a.value ); return *this; }
    simd wrap& operator*=(const simd wrap& a) { value = I::mul(value , a.value ); return *this; }
    // ...
};
```





### Front-end: simd\_wrap element access

If v is a SIMD vector value, how do we change the elements in this vector?

A proxy class can hold on to the SIMD value and the element index, and perform an update when it is assigned:

```
simd<double, 4> v = ...;
// v[3] is actually of type
// simd<double, 4>::element_proxy
v[3] = 2.3;
```

```
template <typename I> struct simd wrap { // ...
    struct element proxy{
        vector type* vptr;
        int i:
        element proxy operator=(scalar type x){
            I::set element(*vptr, i, x);
            return *this:
        operator scalar_type() const {
            return I::element(*vptr, i);
    };
    element proxy operator[](int i) {
        return element proxy{&value , i};
    scalar type operator[](int i) const {
        return I::element(value , i);
};
```





### Front-end: adding indirect operations

How to represent indexed, indirect access to memory? Unify syntax for gather and scatter operations?

```
simd<int, 4> indices = ...;

// gather: v[i] = p[indices[i]]
simd<double, 4> v = indirect(p, indices);

// scatter: p[indices[i]] = v[i]
indirect(p, indices) = v;
```

Make an indirect\_expression object that is returned by indirect(), and represents the pointer to memory and SIMD value of offsets.

```
template <typename I> struct simd wrap;
template <typename I, typename T>
struct indirect expression {
    using index type = typename simd traits<I>::vector type;
    T* p;
    index type index;
    indirect_expression(T* p, const index_type& index):
        p(p), index(index) {}
    template <typename J>
    indirect expression& operator=(const simd wrap<J>& a) {
        a.copy to(*this);
        return *this;
```





### Front-end: adding indirect operations

We can add indirect() to simd\_wrap as a friend function (for access to the private value\_):

```
template <typename I> struct simd_wrap {
    // ...
    template <
        typename Ptr,
        typename =
            std::enable_if_t<std::is_pointer<Ptr>::value>
        }
    friend auto indirect(Ptr p, const simd_wrap& index)
    {
        using V = std::remove_reference_t<decltype(*p)>;
        return indirect_expression<I, V>(p, index.value_);
    }
};
```

Then overload simd\_wrap constructor and copy\_to for indirect\_expression.

```
template <typename I> struct simd wrap {
    // ...
    template <typename J>
    simd wrap(indirect expression<J, const scalar type> pi):
        value (I::gather(tag<J>{}, pi.p, pi.index)) {}
    template <typename J>
    simd wrap(indirect expression<J, scalar type> pi):
        value (I::gather(tag<J>{}, pi.p, pi.index)) {}
    template <typename J>
    void copy to(indirect expression<J, scalar type> pi)
        const
        I::scatter(tag<J>{}, value , pi.p, pi.index);
};
```





# 

### Examples

```
clang++-9 -03 -fno-unroll-loops -march=skylake-avx512
sma(double*, double, double const*, double const*):
    vbroadcastsd %xmm0, %ymm0
    xorl %eax, %eax
.LBB0_1:
    vmulpd (%rsi,%rax,8), %ymm0, %ymm1
    vaddpd (%rdx,%rax,8), %ymm1, %ymm1
    vmovupd %ymm1, (%rdi,%rax,8)
    addq $4, %rax
    cmpq $32768, %rax
    jb .LBB0_1
    vzeroupper
    retq
```





### Examples

```
#include <tinysimd/simd.h>
using int4 = tinysimd::simd<int, 4>;
using double4 = tinysimd::simd<double, 4>;
// c[i] = a[index[i]] + b[index[i]]
void mul indexed(
    double* restrict c, const int* index,
    const double* a, const double* b)
    for (unsigned i = 0; i < N; i+=4) {
        int4 vi(index+i);
        double4 vc = double4(indirect(a, vi))*
                     double4(indirect(b, vi));
       vc.copy to(c+i);
```

```
clang++-9 -03 -fno-unroll-loops -march=skylake-avx512
permuted mul(double*, int const*, double const*,
double const*):
      xorl %eax, %eax
.LBB0 1:
      vmovupd (%rsi,%rax,4), %xmm0
      vpcmpeqd %ymm1, %ymm1, %ymm1
      vxorpd %xmm2, %xmm2, %xmm2
      vgatherdpd %ymm1, (%rdx,%xmm0,8), %ymm2
      vpcmpeqd %ymm1, %ymm1, %ymm1
      vxorpd %xmm3, %xmm3, %xmm3
      vgatherdpd %ymm1, (%rcx,%xmm0,8), %ymm3
      vmulpd %ymm3, %ymm2, %ymm0
      vmovupd %ymm0, (%rdi,%rax,8)
      addq $4, %rax
      cmpq $32768, %rax
      ib .LBB0 1
      vzeroupper
      retq
```





Missing features — masks and conditionals

where(
$$x>3$$
, a) = b + c

Comparisons produce a simd\_mask\_wrap, that also wraps a SIMD implementation, but provides logical operations.

The where function produces a where\_expression, representing a mask and a simd\_wrap lvalue. Assignments are translated to masked SIMD expressions and blend operations provided by the implementation class.



Missing features — conversion/casting

Add constructors to simd\_wrap<I> that call another heterogeneous operation I::cast\_from in the implementation class, with a tag for the source SIMD implementation.

A top-level templated function simd\_wrap<I> simd\_cast(simd\_wrap<J>) asserts width equality and wraps the conversion functions.



#### Resources

### Sample code

Git repo: <a href="https://github.com/eth-cscs/examples\_cpp">https://github.com/eth-cscs/examples\_cpp</a>
tinysimd code and examples in directory Code/tinysimd

#### x86-64 instructions and intrinsics

Compiler intrinsics: <a href="https://software.intel.com/sites/landingpage/IntrinsicsGuide/">https://software.intel.com/sites/landingpage/IntrinsicsGuide/</a>
Félix Cloutier's x86 instruction set reference: <a href="https://www.felixcloutier.com/x86/">https://www.felixcloutier.com/x86/</a>
Agner Fog's instruction tables: <a href="https://www.agner.org/optimize/#manuals">https://www.agner.org/optimize/#manuals</a>

### Compiler explorer

https://godbolt.org





Addendum — constrained indices

General gather/scatter is typically slow, even if supported in the architecture.

#### But:

- If we knew the indices were all the same:
  - gather is a single scalar load.
- If we knew the indices were all distinct:
  - could do an indirect compound add without conflict.





#### Addendum — constrained indices

# Allow user code to add a *constraint* to an indirect expression

Indices	Constraint
3, 3, 3, 3	constant
3, 4, 5, 6	contiguous
3, 5, 2, 7	independent
3, 5, 5, 8	monotonic
3, 5, 2, 5	none





#### Addendum — constrained indices

- Add a constraint field to indirect\_expression, defaulting to constraint::none.
- Add fallback implementations for gather and scatter that take a constraint, and dispatch accordingly.
- 3. Add a fallback implementation for a scatter-reduce operation, that implements the indirect += operation, dispatching on constraint.