Threshold Modes in Elliptic Curves

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<author>Phillip Hallam-Baker

<surname>Hallam-Baker

<initials>P. M.

<firstname>Phillip

<email>phill@hallambaker.com

<organization>ThresholdSecrets.com

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Threshold cryptography operation modes are described with application to the Ed25519, Ed448, X25519 and X448 Elliptic Curves. Threshold key generation allows generation of keypairs to be divided between two or more parties with verifiable security guaranties. Threshold decryption allows elliptic curve key agreement to be divided between two or more parties such that all the parties must co-operate to complete a private key agreement operation. The same primitives may be applied to improve resistance to side channel attacks. A Threshold signature scheme is described in a separate document.

Discussion of this draft should take place on the CFRG mailing list (cfrg@irtf.org), which is archived at <https://mailarchive.ietf.org/arch/browse/cfrg/>.

# Introduction

Public key cryptography provides greater functionality than symmetric key cryptography by introducing separate keys for separate roles. Knowledge of the public encryption key does not provide the ability to decrypt. Knowledge of the public signature verification key does not provide the ability to sign. Threshold cryptography extends the scope of traditional public key cryptography with further separation of roles by splitting the private key. This allows greater control of (e.g.) decryption operations by requiring the use of two decryption key shares rather than just the decryption key.

This document describes threshold modes for decryption and key generation for the elliptic curves described in <norm="RFC7748"/> and <norm="RFC8032"/>. Both schemes are interchangeable in their own right but require minor modifications to the underlying elliptic curve systems to encode the necessary information in the public (X25519/X448) or private key (Ed25519/Ed448).

In its most general form, threshold cryptography allows a private key to be divided between (*n*, *t*) shares such that *n* is the total number of shares created and *t* is the threshold number of shares required to perform the operation. For most applications however, the number of shares is the same as the threshold (*n* = *t*) and the most common case is (*n* = *t* = 2).

This document sets out the principles that support the definition threshold modes in elliptic curve Diffie Hellman system first using simple additive key sharing, an approach which is limited to the case (*n* = *t*). The use of Shamir secret sharing <info="Shamir79"/> is then described to support the case (*n* > *t*). For convenience, we refer to the non Shamir secret sharing case as 'direct key sharing'.

## Applications

Development of the threshold modes described in this document were motivated by the following applications.

### Cloud control of decryption

The security of data at rest is of increasing concern in enterprises and for individual users. Transport layer security such as TLS and SSH provide effective confidentiality controls for data in motion but not for data at rest.

Of particular concern is the case in which a large store of confidential data is held on a server. Encryption provides a simple and effective means of protecting the confidentiality of such data. But the real challenge is how to effect decryption of the data on demand for the parties authorized to access it.

Storing the decryption keys on the server that holds the data provides no real security advantage as any compromise that affects the server is likely to result in disclosure of the keys. Use of end-to-end security in which each document is encrypted under the public key of each person authorized to access it provides the desired security but introduces a complex key management problem and provides no effective means of revoking access after it is granted.

Threshold encryption allows a key service to control decryption of stored data without having the ability to decrypt. The data decryption key is split into two (or more) parts with a different split being created for each user. One decryption share is held at the server allowing it to control access to the data without being able to decrypt. The other decryption share is encrypted under the public encryption key of the corresponding user allowing them to decrypt the stored data but only with the co-operation of the key service.

### Device Onboarding

The term 'onboarding' is used to refer to the configuration of a device for use by a particular user or within a particular enterprise. One of the typical concerns of onboarding is to initialize the device with a set of public key pairs for various purposes and to issue credentials (e.g. PKIX certificates) to enable their use.

One of the concerns that arises in such cases is whether keys should be generated on the device on which they are to be used or on another device administering the onboarding process.

Both approaches are unsatisfactory. While generation of keys on the device on which they are to be used is generally preferred, experience has shown that many devices, particularly IoT devices use insufficiently random processes to generate such keys and this has led to numerous cases of duplicate and otherwise weak keys being discovered in running systems.

Generation of keys on an administration device and transferring them to the device on which they are to be used means that the administration device used for key generation represents a single point of failure. Compromise of this device or of the means used to install the keys will lead to compromise of the device.

Threshold key generation allows the advantages of both approaches to be realized avoiding dependence on either the target device or the administration device.

### Cryptographic co-processor

Most real-world compromises of cryptographic security systems involve the disclosure of a private key. Common means of disclosure being inadvertent inclusion in backup tapes, keys being stored on public file shares and (increasingly) keys being inadvertently uploaded to source code management systems.

A new and emerging set of key disclosure threats come from the recent generation of hardware vulnerabilities being discovered in CPUs including ROWHAMMER and SPECTRE.

The advantages of Hardware Security Modules (HSMs) for storing and using private keys are well established. An HSM allows a private key to be used in an isolated environment that is designed to be resistant to side channel attacks.

Regrettably, the 'black box' nature of HSMs means that their use introduces a new security concern - the possibility that the device itself has been compromised during manufacture or in the supply chain.

Threshold approaches allows a key exchange or signature operation to be divided between two private keys, one of which is generated by the application that is to use it and the other of which is tightly bound to a specific host such that it cannot be extracted.

### Side Channel Resistance

The same techniques that make threshold cryptography possible are the basis for Kocher's side-channel resistance technique <info="Kocher96"/>. Differential side channel attacks are a powerful tool capable of revealing a private key value that is used repeatedly in practically any algorithm. The claims made with respect to 'constant time' algorithms such as the Montgomery ladder depend upon the actual implementation performing operations in constant time.

# Definitions

This section presents the related specifications and standard, the terms that are used as terms of art within the documents and the terms used as requirements language.

## Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in <norm="RFC2119"/>.

## Defined Terms

The following terms are used as terms of art in this document and the accompanying specification <info="draft-hallambaker-threshold-sigs"/>.

Dealer

A party that coordinates the actions of a group of participants performing a threshold operation.

Multi-Encryption

The use of multiple decryption fields to allow a document encrypted under a session key to be decrypted by multiple parties under different decryption keys.

Multi-Encryption allows a document to be shared with multiple recipients but does not allow the decryption role to be divided between multiple parties.

Multi-Signatures

The use of multiple independently verifiable digital signatures to authenticate a document.

Multi-Signatures allow separation of the signing roles and thus achieve a threshold capability. But they are not true threshold signatures as the set of signers is visible to external parties.

Onboarding

The process by which an embedded device is provisioned for deployment in a local network.

Threshold Key Generation

An aggregate public, private key pair is constructed from a set of contributions such that the private key is a function of the private key of all the contributions.

A Threshold Key Generation function is auditable if and only if the public component of the aggregate key can be computed from the public keys of the contributions alone.

Threshold Decryption

Threshold decryption divides the decryption role between two or more parties.

Threshold Key Agreement

A bilateral key agreement exchange in which one or both sides present multiple public keys and the key agreement value is a function of all of them.

This approach allows a party to present multiple credentials in a single exchange, a de

Threshold Signatures

Threshold signatures divide the signature role between two or more parties in such a way that the parties and their roles is not visible to an external observer.

## Related Specifications

This document extends the elliptic curve cryptography systems described in <norm="RFC7748"/> and <norm="RFC8032"/> to provide threshold capabilities.

This work was originally motivated by the requirements of the Mathematical Mesh <info="draft-hallambaker-mesh-architecture"/>.

## Implementation Status

The implementation status of the reference code base is described in the companion document <info="draft-hallambaker-mesh-developer"/>.

# Threshold Cryptography in Diffie-Hellman

The threshold modes described in this specification are made possible by the fact that Diffie Hellman key agreement and elliptic curve variants thereof support properties we call the Key Combination Law and the Result Combination Law.

Let {*X*, *x*}, {*Y*, *y*}, {*A*, *a*} be {public, private} key pairs and r [.] S represent the Diffie Hellman operation applying the private key r to the public key S.

The Key Combination law states that we can define an operator [x] such that there is a keypair {*Z*, *z*} such that:

*Z* = *X* [x] *Y* and *z* = (*x* + *y*) mod *o* (where *o* is the order of the group)

The Result Combination Law states that we can define an operator [+] such that:

(*x* [.] *A*) [+] (*y* [.] *A*) = (*z* [.] *A*) = (*a* [.] *Z*)

It will be noted that each of these laws is interchangeable. The output of the key combination law to a Diffie Hellman key pair is a Diffie Hellman key pair and the output of the result combination law is a Diffie Hellman result. This allows modular and recursive application of these principles.

## Application to Diffie Hellman (not normative)

Diffie Hellman in a modular field provides a concise demonstration of the key combination and result combination laws <info="RFC2631"/>. The realization of the threshold schemes in a modular field is outside the scope of this document.

For the Diffie Hellman system in a modular field p, with exponent e:

* r [.] S = Sr mod p
* o = p-1
* *A* [x] *B* = *A* [.] *B* = *AB* mod *p*.

*Proof:*

Let z = x + y

By definition, X = ex mod p, Y = ey mod p, and Z = ez mod p.

Therefore,

Z = ez mod p

= ex+y mod p

= (exey) mod p

= ex mod p.ey mod p

= X.Y

Moreover, A = ea mod p,

Therefore,

(Ax mod p).(Ay mod p)

= (AxAy) mod p)

= (Ax+y) mod p)

= Az mod p

= eaz mod p

= (ez)a mod p

= Za mod p

Since eo mod p = 1, the same result holds for z = (x + y) mod o since ex+y+no = ex+y.eno = ex+y.1 = ex+y.

## Threshold decryption

Threshold decryption allows a decryption key to be divided into two or more parts, allowing these roles to be assigned to different parties. This capability can be employed within a machine to divide use of a private key between an implementation running in the user mode and a process running in a privileged mode that is bound to the machine. Alternatively, threshold cryptography can be employed to

The key combination law and result combination law provide a basis for threshold decryption.

### Direct Key Splitting

We begin by creating a base key pair { X, x }. The public component X is used to perform encryption operations by means of a key agreement against an ephemeral key in the usual fashion. The private component x may be used for decryption in the normal fashion or to provide the source material for a key splitting operation.

To split the base key into n shares { S1, s1 }, { S2, s2 }, … { Sn, sn }, we begin by generating the first n-1 private keys in the normal fashion. It is not necessary to generate the corresponding public keys as these are not required.

The private key of the final key share sn is given by:

*sn = (x - s1 - s2 - … - sn-1)* mod *o*

Thus, the base private key x is equal to the sum of the private key shares modulo the group order.

### Direct Decryption

To encrypt a document, we first generate an ephemeral key pair { Y, y }. The key agreement value exy is calculated from the base public key X = ex and the ephemeral private key y. A key derivation function is then used to obtain the session key to be used to encrypt the document under a symmetric cipher.

To decrypt a document using the threshold key shares, each key share holder first performs a Diffie Hellman operation using their private key on the ephemeral public key. The key shares are then combined using the result combination law to obtain the key exchange value from which the session key is recovered.

The key contribution ci for the holder of the ith key share is calculated as:

ci = Ysi

The key agreement value is thus

A = c1 . c2 . … . cn

This value is equal to the encryption key agreement value according to the group law.

## Direct threshold key generation

The key combination law allows an aggregate private key to be derived from private key contributions provided by two or more parties such that the corresponding aggregate public key may be derived from the public keys corresponding to the contributions. The resulting key generation mechanism is thus, auditable and interoperable.

### Device Provisioning

Auditable Threshold Key Generation provides a simple and efficient means of securely provisioning keys to devices. This is encountered in the IoT space as a concern in 'onboarding' and in the provisioning of unique keys for use with cryptographic applications (e.g. SSH, S/MIME, OpenPGP, etc.).

Consider the case in which Alice purchases an IoT connected Device D which requires a unique device key pair *{ X , x }* for its operation. The process of provisioning (aka 'onboarding') is performed using an administration device. Traditional key pair generation gives us three options:

* Generate and install a key pair during manufacture.
* Generate a new key pair during device provisioning.
* Generate a key pair on the administration device and transfer it to the device being provisioned.

The first approach has the obvious disadvantage that the manufacturer has knowledge of the private key. This represents a liability for both the user and the manufacturer. Less obvious is the fact that the second approach doesn't actually provide a solution unless the process of generating keys on the device is auditable. The third approach is auditable with respect to the device being provisioned but not with respect to the administration device being used for provisioning. If that device were to be compromised, so could every device it was used to provision.

Threshold key generation allows the administration device and the joining device being provisioned to jointly provision a key pair as follows:

* The joining device has public, private key pair *{ D, d }*.
* A provisioning request is made for the device which includes the joining device public key *D*.
* The administration device generates a key pair contribution *{ A, a }*.
* The administration device private key is transmitted to the Device by means of a secure channel.
* The joining device calculates the aggregate key pair *{ A [x] D, a+d }*.
* The administration device authorizes the joining device to participate in the local network using the public key *A [x] D*.

The Device key pair MAY be installed during manufacture or generated during provisioning or be derived from a combination of both using threshold key generation recursively. The provisioning request MAY be originated by the device or be generated by a purchasing system.

Note that the provisioning protocol does not require either party to authenticate the aggregate key pair. The protocol provides security by ensuring that both parties obtain the correct results if and only if the values each communicated to the other were correct.

Out of band authentication of the joining device public key *D* is sufficient to prevent Man-in-the-Middle attack. This may be achieved by means of a QR code printed on the device itself that discloses or provides a means of obtaining *D.*

[Note add serious warning that a party providing a private contribution MUST make sure they see the public key first. Otherwise a rogue key attack is possible. The Mesh protocols ensure that this is the case but other implementations may overlook this detail.]

### Key Rollover

Traditional key rollover protocols in PKIX and other PKIs following the Kohnfelder model, require a multi-step interaction between the key holder and the Certificate Authority.

Threshold key generation allows a Certificate Authority to implement key rollover with a single communication:

Consider the case in which the service host has a base key pair { B , b }. A CA that has knowledge of the Host public key B may generate a certificate for the time period *t* with a fresh key as follows:

* Generate a key pair contribution { At, at }.
* Generate and sign an end entity certificate Ct for the key B [x] At.
* Transmit Ct, at to the host by means of a secure channel

### Host Activation

Modern Internet service architectures frequently make use of short lived, ephemeral hosts running on virtualized machines. Provisioning cryptographic material in such environments is a significant challenge and especially so when the underlying hardware is a shared resource.

The key rollover approach described above provides a means of provisioning short lived credentials to ephemeral hosts that potentially avoids the need to build sensitive keys into the service image or configuration.

### Separation of Duties

Threshold key generation provides a means of separating administration of cryptographic keys between individuals. This allows two or more administrators to control access to a private key without having the ability to use it themselves. This approach is of particular utility when used in combination with threshold decryption. Alice and Bob can be granted the ability to create key contributions allowing a user to decrypt information without having the ability to decrypt themselves.

## Side Channel Resistance

Side-channel attacks, present a major concern in the implementation of public key cryptosystems. Of particular concern are the timing attacks identified by Paul Kocher <info="Kocher96"/> and related attacks in the power and emissions ranges. Performing repeated observations of the use of the same private key material provides an attacker with considerably greater opportunity to extract the private key material.

A simple but effective means of defeating such attacks is to split the private key value into two or more random shares for every private key operation and use the result combination law to recover the result.

The implementation of this approach is identical to that for Threshold Decryption except that instead of giving the key shares to different parties, they are kept by the party performing the private key operation.

While this approach doubles the number of private key operations required, the operations MAY be performed in parallel. Thus avoiding impact on the user experience.

# Shamir Secret Sharing

The direct threshold modes described above may be extended to support the case (*n* > *t*) through application of Shamir secret sharing and the use of the Lagrange basis to recover the secret value.

Shamir Secret Sharing makes use of the fact that a polynomial of degree t-1 is defined by t points on the curve. To share a secret *s*, we construct a polynomial of degree *t-1* in the modular field *L* where *L* > *s*.

*f*(*x*) = *s* + *a1*.*x* + *a2*.*x2* + … *at-1*.*xt-1* mod *L*

The shares *p1*, *p2* … *pn* are given by the values *f*(*x1*), *f*(*x2*), … ,*f*(*xn*) where *x1*, *x2* … *xn* are in the range 1 < *xi* < *L*.

We can use the Lagrange basis function to construct a set of coefficients l1, l2, … lt from a set of *t* shares p1, p2 … pt such that:

*s* = l1p1 + l2p2 + … + ltpt mod *L*

Thus, if we choose the sub-group order of the curve as the value of *L*, the formula used to recover the secret *s* is a sum of terms as was used for the case where *n* = *t*. We can thus apply a set of Lagrange coefficients provided by the dealer to the secret shares to extend the threshold operations to the case where *n* > *t*.

## Shamir secret share generation

To create *n* shares for the secret *s* with a threshold of *t*, the dealer constructs a polynomial of degree *t* in the modular field *L*, where *L* is the order of the curve sub-group:

f(x) = a0 + a1.x + a2.x2 + … at.xt-1 mod L

where

a0 = s

a1 … at are randomly chosen integers in the range 1 < ai < L

The values of the key shares are the values *f*(x1), *f*(x2), … ,*f*(xn). That is

pi = *f*(xi)

where

p1 … pn are the private key shares

x1 … xn are distinct integers in the range 1 < xi < L

The means of constructing distinct values x1 … xn is left to the implementation. It is not necessary for these values to be secret.

## Lagrange basis recovery

Given *n* shares (*x0*, *y0*), (*x1*, *y1*), … (*xn-1*, *yn-1*), The corresponding the Lagrange basis polynomials *l0*, *l1*, .. *ln-1* are given by:

lm = ((*x* - *xm0*) / (*xm* - x*m0*)) . ((*x* - *xm1*) / (*xm* - x*m1*)) . … . ((*x* - *xmn-2*) / (*xm* - *xmn-*2))

Where the values *m0*, *m1*, … *mn-2*, are the integers 0, 1, .. *n*-1, excluding the value *m*.

These can be used to compute *f(x)* as follows:

*f*(*x*) = *y0l0* + *y1l1* + … *yn-1ln-1*

Since it is only the value of *f(*0*)* that we are interested in, we compute the Lagrange basis for the value *x* = 0:

*lzm* = ((*xm1*) / (*x*m - *xm1*)) . ((*xm2*) / (*xm* - *xm2*))

Hence,

*a0* = *f*(*0*) = *y0lz0* + *y1lz1* + … *yn-1ln-1*

## Verifiable Secret Sharing

The secret share generation mechanism described above allows a private key to be split into *n* shares such that *t* shares are required for recovery. While this supports a wide variety of applications, there are many cases in which it is desirable for generation of the private key to be distributed.

Feldman’s Verifiable Secret Sharing (VSS) Scheme provides a mechanism that allows the participants in a distributed generation scheme to determine that their share has been correctly formed by means of a commitment.

To generate a commitment the dealer generates the polynomial *f*(*x*) as before and in addition selects a generator *g*

[TBS]

## Distributed Key Generation

[TBS]

# Application to Elliptic Curves

For elliptic curve cryptosystems, the operators [x] and [.] are point addition.

Implementing a robust Key Co-Generation for the Elliptic Curve Cryptography schemes described in <norm="RFC7748"/> and <norm="RFC8032"/> requires some additional considerations to be addressed.

* The secret scalar used in the EdDSA algorithm is calculated from the private key using a digest function. It is therefore necessary to specify the Key Co-Generation mechanism by reference to operations on the secret scalar values rather than operations on the private keys.
* The Montgomery Ladder traditionally used to perform X25519 and X448 point multiplication does not require implementation of a function to add two arbitrary points. While the steps required to create such a function are fully constrained by <norm="RFC7748"/>, the means of performing point addition is not.

## Implementation for Ed25519 and Ed448

<norm="RFC8032"/> provides all the cryptographic operations required to perform threshold operations and corresponding public keys.

The secret scalars used in <norm="RFC8032"/> private key operations are derived from a private key value using a cryptographic digest function. This encoding allows the inputs to a private key combination to be described but not the output. Contrawise, the encoding allows the inputs to a private key splitting operation to be described but not the output

It is therefore necessary to provide an alternative representation for the Ed25519 and Ed448 private keys. Moreover, the signature algorithm requires both a secret scalar and a prefix value as inputs.

Since threshold signatures are out of scope for this document and <norm="RFC8032"/> does not specify a key agreement mechanism, it suffices to specify the data formats required to encode private key values generated by means of threshold key generation.

### Ed25519

Let the inputs to the threshold key generation scheme be a set of 32 byte private key values *P1, P2 … Pn*. For each private key value *i* in turn:

1. Hash the 32-byte private key using SHA-512, storing the digest in a 64-octet large buffer, denoted *hi*. Let ni be the first 32 octets of hi and mi be the remaining 32 octets of hi.
2. Prune ni: The lowest three bits of the first octet are cleared, the highest bit of the last octet is cleared, and the second highest bit of the last octet is set.
3. Interpret the buffer as the little-endian integer, forming a secret scalar si.

The private key values are calculated as follows:

The aggregate secret scalar value *sa = s1 + s2 + … sn* mod *L*, where *L* is the order of the curve.

The aggregate prefix value is calculated by either

* Some function TBS on the values m1, m2, … mn, or
* Taking the SHA256 digest of sa.

The second approach is the simplest and the most robust. It does however mean that the prefix is a function of the secret scalar rather than both being functions of the same seed.

### Ed448

Let the inputs to the threshold key generation scheme be a set of 57 byte private key values *P1, P2 … Pn*. For each private key value *i* in turn:

1. Hash the 57-byte private key using SHAKE256(x, 114), storing the digest in a 114-octet large buffer, denoted *hi*. Let ni be the first 57 octets of hi and mi be the remaining 57 octets of hi.
2. Prune ni: The two least significant bits of the first octet are cleared, all eight bits the last octet are cleared, and the highest bit of the second to last octet is set.
3. Interpret the buffer as the little-endian integer, forming a secret scalar si.

The private key values are calculated as follows:

The aggregate secret scalar value *sa = s1 + s2 + … sn* mod *L*, where *L* is the order of the curve.

The aggregate prefix value is calculated by either

* Some function TBS on the values m1, m2, … mn, or
* Taking the SHAKE256(x, 57) digest of sa.

The second approach is the simplest and the most robust. It does however mean that the prefix is a function of the secret scalar rather than both being functions of the same seed.

## Implementation for X25519 and X448

<norm="RFC7748"/> defines all the cryptographic operations required to perform threshold key generation and threshold decryption but does not describe how to implement them.

The Montgomery curve described in <norm="RFC7748"/> allows for efficient scalar multiplication using arithmetic operations on a single coordinate. Point addition requires both coordinate values. It is thus necessary to provide an extended representation for point encoding and provide an algorithm for recovering both coordinates from a scalar multiplication operation and an algorithm for point addition.

The notation of <norm="RFC7748"/> is followed using {u, v} to represent the coordinates on the Montgomery curve and {x, y} for coordinates on the corresponding Edwards curve.

### Point Encoding

The relationship between the u and v coordinates is specified by the Montgomery Curve formula itself:

*v2* = *u3 + Au2 + u*

An algorithm for extracting a square root of a number in a finite field is specified in <norm="RFC8032">.

Since *v2* has a positive (*v*) and a negative solution (*-v*), it follows that *v2* mod p will have the solutions *v*, *p-v*. Furthermore, since *p* is odd, if *v* is odd, *p-v* must be even and vice versa. It is thus sufficient to record whether *v* is odd or even to enable recovery of the *v* coordinate from *u*.

### X25519 Point Encoding

The extended point encoding allowing the v coordinate to be recovered is as given in <norm="draft-ietf-lwig-curve-representations"/>

A curve point (u, v), with coordinates in the range 0 <= u,v < p, is coded as follows. First, encode the u-coordinate as a little-endian string of 57 octets. The final octet is always zero. To form the encoding of the point, copy the least significant bit of the v-coordinate to the most significant bit of the final octet.

### X448 Point Encoding

The extended point encoding allowing the v coordinate to be recovered is as given in <norm="draft-ietf-lwig-curve-representations"/>

A curve point (u, v), with coordinates in the range 0 <= u,v < p, is coded as follows. First, encode the u-coordinate as a little-endian string of 57 octets. The final octet is always zero. To form the encoding of the point, copy the least significant bit of the v-coordinate to the most significant bit of the final octet.

### Point Addition

The point addition formula for the Montgomery curve is defined as follows:

Let P1 = {u1, v1}, P2 = {u2, v2}, P3 = {u3, v3} = P1 + P2

By definition:

u3

= B(v2 - v1)2 / (u2 - u1)2 - A - u1 - u2

= B((u2v1 - u1v2)2 ) / u1u2 (u2 - u1)2

v3 = ((2u1 + u2 + A)(v2 - v1) / (u2 - u1)) - B (v2 - v1)3 / (u2 -u1)3 - v1

For curves X25519 and X448, B = 1 and so:

u3 = ((v2 - v1).(u2 - u1)-1)2 - A - u1 - u2

v3 = ((2u1 + u2 + A)(v2 - v1).(u2 - u1)-1) - ((v2 - v1).(u2 -u1)-1)3 - v1

This may be implemented using the following code:

<include=..\Source\MontyPointAdd.md>

Performing point addition thus requires that we have sufficient knowledge of the values v1, v2. At minimum whether one is odd and the other even or if both are the same.

### Montgomery Ladder with Coordinate Recovery

As originally described, the Montgomery Ladder only provides the u coordinate as output. López and Dahab <info="Lopez99"/> provided a formula for recovery of the v coordinate of the result for curves over binary fields. This result was then extended by Okeya and Sakurai <info="Okeya01"/> to prime field Montgomery curves such as X25519 and X448. The realization of this result described by Costello and Smith <info="Costello17"/> is applied here.

The scalar multiplication function specified in <norm="RFC7748"/> takes as input the scalar value k and the coordinate u1 of the point P1 = {u1, v1} to be multiplied. The return value in this case is u2 where P2 = {u2, v2} = k.P1.

To recover the coordinate v2 we require the values x\_2, z\_2, x\_3, z\_3 calculated in the penultimate step:

<include=..\Source\MontyLadder.md>

The values x\_2, z\_2 give the projective form of the u coordinate of the point P2 = {u2, v2} = k.P1 and the values x\_3, z\_3 give the projective form of the u coordinate of the point P3 = {u3, v3} = (k+1).P1 = P1 + k.P1 = P1 + P2.

Given the coordinates {u1, v1} of the point P1 and the u coordinates of the points P2, P1 + P2, the coordinate v2 MAY be recovered by trial and error as follows:

<include=..\Source\MontyGetVSqrt.md>

Alternatively, the following MAY be used to recover {u2, v2} without the need to extract the square root and using a single modular exponentiation operation to convert from the projective coordinates used in the calculation. As with the Montgomery ladder algorithm above, the expression has been modified to be consistent with the approach used in <norm="RFC7748"/> but any correct formula may be used.

<include=..\Source\MontyGetVFast.md>

# Test Vectors

<include=..\Examples\ExamplesThreshold.md>

# Security Considerations

All the security considerations of <norm="RFC7748"/> and <norm="RFC8032"/> apply and are hereby incorporated by reference.

## Complacency Risk

Separation of duties can lead to a reduction in overall security through complacency and lack of oversight.

Consider the case in which a role that was performed by A alone is split into two roles B and C. If B and C each do their job with the same diligence as A did alone, risk should be reduced but if B and C each decide they can be careless because security is the responsibility of the other, the risk of a breach may be substantially increased.

It is therefore important that each of the participants in a threshold scheme perform their responsibilities with the same degree of diligence as if they were the sole control and for those responsible for oversight to treat single point failures that do not lead to an actual breach with the same degree of concern as if a breach had occurred.

Use of threshold operation modes mitigates but does not eliminate security considerations relating to private key operations of the underlying algorithm. For example, use of threshold key generation to generate a composite keypair {b+c, B+C} from key contributions {b, B} and {c, C} produces a strong composite private key if either of the key contributions *a*, *b* are strong. But the composite key will be weak if neither contribution is strong.

## Rogue Key Attack

In general, threshold modes of operation provide a work factor that is at least as high as that of the work factor of the strongest private key share. The karmic tradeoff for this benefit is that the trustworthiness of a composite public key is that of the least trustworthy input.

For example, consider the case in which a client with keypair {c, C} generates an ephemeral keypair {e, E} for use in an authentication algorithm. We might decide to create an 'efficient' proof of knowledge of c and e by using the composite public key A = C+E to test for knowledge of both at the same time.

This approach fails because an attacker with knowledge of C can generate a keypair {a, A} and calculate the corresponding public key E = A-C. The attacker can then use the value a in the authentication protocol.

# IANA Considerations

This document requires no IANA actions (yet).

It will be necessary to define additional code points for the signed version of the X25519 and X448 public key and the threshold decryption final private keys.

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# Appendix A: Calculating Lagrange coefficients

The following C# code calculates the Lagrange coefficients used to recover the secret from a set of shares.

[TBS]