

Functional Languages

4190.310
Programming Languages
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Lecture 05

Reading assignments: Chapter 10

Historical Origins

- The imperative and functional models grew out of work undertaken Alan Turing, Alonzo Church, Stephen Kleene, Emil Post, etc. ~1930s
 - Different formalizations of the notion of an algorithm, or effective procedure, based on automata, symbolic manipulation, recursive function definitions, and combinatorics
- These results led Church to conjecture that any intuitively appealing model of computing would be equally powerful as well
 - This conjecture is known as Church's thesis

Turing Machine

- Turing's model of computing was the Turing machine
 - A sort of pushdown automaton using an unbounded storage “tape”
- The Turing machine computes in an imperative way, by changing the values in cells of its tape – like variables just as a high level imperative program computes by changing the values of variables

Lambda Calculus

- Church's model of computing is called the lambda calculus
 - Based on the notion of parameterized expressions (with each parameter introduced by an occurrence of the letter λ —hence the notation's name)
- Lambda calculus was the inspiration for functional programming
- One uses it to compute by substituting parameters into expressions, just as one computes in a high level functional program by passing arguments to functions

Functional Programming Concepts

- Functional languages such as Lisp, Scheme, FP, ML, Miranda, and Haskell are an attempt to realize Church's lambda calculus in practical form as a programming language
- The key idea: do everything by composing functions
 - No mutable state
 - No side effects

Necessary Features in Functional Programming

- Many of which are missing in some imperative languages
- 1st class and high-order functions
- Serious polymorphism
- Powerful list facilities
- Structured function returns
- Fully general aggregates
- Garbage collection

How do You Get Anything Done?

- Recursion (especially tail recursion) takes the place of iteration
- In general, you can get the effect of a series of assignments

x := 0 ...

x := expr1 ...

x := expr2 ...

- From $f_3(f_2(f_1(0)))$, where each f expects the value of x as an argument, f_1 returns expr1 , and f_2 returns expr2

Tail Recursion

- Tail call
 - Suppose function g calls function f
 - A call to f in the body of g is a tail call if g returns the result of calling f without any further computation
- Tail recursion
 - A function is tail recursive if all recursive calls in the body of f are tail calls to f

```
int g(int n) {                                int f(int n) {  
    if (n == 0)                                if (n == 0)  
        return f(x);                            return f(x);  
    else                                         else  
        return (f(x) + n);                      return f(n-1);  
}                                              }
```

Recursion

- Recursion even does a nifty job of replacing looping

```
x := 0; i := 1; j := 100;  
while i < j do  
    x := x + i*j; i := i + 1;  
    j := j - 1  
end while  
return x
```

- Becomes $f(0, 1, 100)$, where
 $f(x, i, j) == \text{if } i < j \text{ then } f(x+i*j, i+1, j-1) \text{ else } x$

Higher-order Functions

- Thinking about recursion as a direct, mechanical replacement for iteration, however, is the wrong way to look at things
 - One has to get used to thinking in a recursive style
- Even more important than recursion is the notion of higher-order functions
 - Take a function as argument, or return a function as a result
 - Great for building things

Variants of Lisp

- Pure (original) Lisp
- Interlisp, MacLisp, Emacs Lisp
- Common Lisp
- Scheme

Variants of Lisp (contd.)

- Pure Lisp is purely functional; all other Lisps have imperative features
- All early Lisps dynamically scoped
 - Not clear whether this was deliberate or if it happened by accident
- Scheme and Common Lisp statically scoped
 - Common Lisp provides dynamic scope as an option for explicitly-declared special functions
 - Common Lisp now THE standard Lisp
 - Very big; complicated (The Ada of functional programming)

Variants of Lisp (contd.)

- Scheme is a particularly elegant Lisp
- Other functional languages
 - ML
 - Miranda
 - Haskell
 - FP
- Haskell is the leading language for research in functional programming

Advantages of Functional Languages

- Lack of side effects makes programs easier to understand
- Lack of explicit evaluation order (in some languages) offers possibility of parallel evaluation (e.g. MultiLisp)
- Lack of side effects and explicit evaluation order simplifies some things for a compiler (provided you don't blow it in other ways)
- Programs are often surprisingly short
- Language can be extremely small and yet powerful

Problems of Functional Languages

- Difficult (but not impossible!) to implement efficiently on von Neumann machines
- Lots of copying of data through parameters
- (apparent) need to create a whole new array in order to change one element

Problems of Functional Languages (contd.)

- Heavy use of pointers (space/time and locality problem)
- Frequent procedure calls
- Heavy space use for recursion
- Requires garbage collection
- Requires a different mode of thinking by the programmer
- Difficult to integrate I/O into purely functional model